

Pseudocode examples

BFS

Queue implementation:

Algorithm 1: BFS(G,s)	
1	<i>/* G is hash table, the adjacency list of a graph */</i>
2	<i>/* s is a source vertex in G */</i>
1	$parents \leftarrow \{\}$ <i>/* empty hash table, $parents[v] = v$'s parent. */</i>
2	$dist \leftarrow \{\}$ <i>/* empty hash table, $dist[v] =$ distance from s. */</i>
3	$Q \leftarrow$ empty FIFO queue <i>/* keep track of active nodes */</i>
4	$Q.enqueue(s)$, $parents[s] = None$, $dist[s] = 0$ <i>/* initialization */</i>
5	while Q is not empty do
6	$u \leftarrow Q.dequeue()$;
7	for v in $G[u]$ do
8	<i>/* explore neighbors of active node u */</i>
8	if v not in $parents$ then
9	<i>/* v was so far undiscovered */</i>
9	$parents[v] = u$;
10	$dist[v] = dist[u] + 1$;
11	$Q.enqueue(v)$;
12	return $parents, dist$

Layers implementation:

Algorithm 2: BFS(G, s)

```
/*  $G$  is hash table, the adjacency list of a graph */
/*  $s$  is a source vertex in  $G$  */
1  $parents \leftarrow$  hash table /*  $parents[v] = v$ 's parent in BFS tree. */
2  $layers \leftarrow$  hash table /*  $layers[i] =$  nodes in layer  $i$  */
3  $dist \leftarrow$  hash table /* empty hash table,  $dist[v] =$  distance from  $s$ . */
4  $layers[0] \leftarrow s, parents[s] = None, dist[s] = 0, i \leftarrow 0$  /* initialization */
5 while  $layers[i]$  is not empty do
6    $layers[i + 1] \leftarrow []$ ;
7   for  $u$  in  $layers[i]$  do
8     for  $v$  in  $G[u]$  do
9       /* explore neighbors of active node  $u$  */
10      if  $v$  not in  $parents$  then
11        /*  $v$  was so far undiscovered */
12         $parents[v] = u;$ 
13         $dist[v] = dist[u] + 1;$ 
14         $layers[i + 1].add(v)$  /*  $v$  is at layer  $i+1$  */
15  $i \leftarrow i + 1;$ 
16 return  $parents, dist, layers$  (optional)
```

Algorithm 3: bipartiteBFS(G)

```
/*  $G$  is hash table, the adjacency list of a graph */
1  $s \leftarrow$  choose a vertex at random from  $V$  /* serves as the source */
2  $parents, dist = BFS(G, s);$ 
3  $C1, C2 \leftarrow []$  /* empty lists, contains the vertices of the two classes in the
   bipartition */
4 for  $v$  in  $dist$  do
5   /* assign nodes to  $C1, C2$  based on odd/even layers */
6   if  $dist[v] == odd$  then
7      $C1.add(v);$ 
8   else
9      $C2.add(v);$ 
10 return  $C1, C2$ 
```

DFS

Algorithm 4: DFSWrapper(G, s)

```
/*  $G$  adjacency list as nested hash table,  $s$  source node */
1  $parents \leftarrow$  empty hash table /* parents of nodes in the DFS tree (forest) */
2  $times \leftarrow$  empty hash table /* tuples of <discovery time, finish time> */
3  $time \leftarrow 0$  /* time counter */
4  $parents[s] \leftarrow None$ ;
5 DFS( $G, s$ );
6 return  $parents, times$ 
```

Algorithm 5: DFS(G, u)

```
1  $time \leftarrow time + 1$ ;
2  $times[u][0] = time$ ;
3 for  $v$  in  $G[u]$  do
    /* recursively explore  $u$ 's neighbors */
4     if  $v$  not in  $parents$  then
5          $parents[v] = u$ ;
6         DFS( $G, v$ )
7  $time \leftarrow time + 1$ ;
8  $times[u][1] = time$ ;
```

Dijkstra

Algorithm 6: Dijkstra(G, s)	
	<i>/* G adjacency list fo weighted directed graph. $G[u][v] = l(u, v)$, source s */</i>
1	$\pi \leftarrow \{ \}$ <i>/* hash table, current best list for v */</i>
2	$d \leftarrow \{ \}$ <i>/* hash table, distance of v */</i>
3	$parents \leftarrow \{ \}$ <i>/* parents in shortest paths tree */</i>
4	$Q \leftarrow PQ$ <i>/* priority queue to keep track of min π */</i>
5	$\pi[s] = 0;$
6	$Q.INSERT(< 0, s >);$
7	for $v \neq s$ in G do
8	$\pi[v] = \infty;$
9	$Q.INSERT(< \pi[v], v >);$
10	while Q <i>is not empty</i> do
11	$< \pi[u], u > \leftarrow \text{EXTRACT-MIN}(Q);$
12	$d[u] \leftarrow \pi[u];$
13	for v in $G[u]$ do
14	if $\pi[v] > d[u] + l(u, v)$ then
15	$\text{DECREASE-KEY}(< \pi[v], v >, < d[u] + l(u, v), v >);$
16	$\pi[v] \leftarrow d[u] + l(u, v);$
17	$parents[v] \leftarrow u;$
18	return $d, parents$

Prim

Algorithm 7: Prim(G)	
<pre> /* G adjacency list of weighted undirected graph $G[u][v] = w(u, v)$ */ 1 $parents \leftarrow$ hash table /* parents list in MST, contains current best edge */ 2 $wT \leftarrow 0$ /* total weight of MST, optional */ 3 $Tnodes \leftarrow$ empty set /* keep track of nodes fixed in MST */ 4 $Q \leftarrow$ empty priority queue /* $\langle weight, v \rangle$ least edge weight connecting v to MST */ */ 5 $s \leftarrow$ random node as source; 6 $parents[s] \leftarrow None, Q.INSERT(\langle 0, s \rangle);$ 7 for $v \neq s$ in G do 8 $Q.INSERT(\langle \infty, v \rangle)$ /* use ∞ for unreachable nodes. Can use really large int instead. */ 9 $parents[v] \leftarrow None;$ 10 while Q is not empty do 11 $\langle weight, u \rangle \leftarrow Q.EXTRACT - MIN()$ /* next lightest edge */ 12 $wT += weight;$ 13 $Tnodes.add(u)$ /* u is added to the MST and is fixed */ 14 for v in $G[u]$ do 15 if v not in $Tnodes$ AND $G[v][parents[v]] > G[v][u]$ then 16 $Q.DECREASE - KEY(\langle G[v][parent[v]], v \rangle, \langle G[v][u], v \rangle);$ 17 $parents[v] = u;$ 18 return $parents, wT$ </pre>	

Divide and Conquer

Algorithm 8: MergeSort(A, p, r)	
<pre> /* Sorts the subarray $A[p:r]$ in place */ 1 if $p == r$ then 2 return $A;$ 3 $q \leftarrow \lfloor \frac{p+r}{2} \rfloor;$ 4 $A[p : q] \leftarrow MergeSort(A, p, q);$ 5 $A[q + 1 : r] \leftarrow MergeSort(A, q + 1, r);$ 6 $A[p : r] \leftarrow Merge(A, p, q, r);$ 7 return A </pre>	

Algorithm 9: BinarySearch($A, p, r, query$)

```

    /* find the index of query in the subarray A[p:r] */
1   $q \leftarrow \lfloor \frac{p+r}{2} \rfloor$ ;
2   $middle \leftarrow A[q]$ ;
3  if  $query == middle$  then
4    return  $q$ ;
5  else if  $query < middle$  then
6    BinarySearch( $A, p, q, query$ );
7  else
8    BinarySearch( $A, q + 1, r, query$ );
9  return  $query$  not in  $A$ 

```

Algorithm 10: Karatsuba(a, b, n)

```

    /* multiply  $n$ -bit ints  $a$  and  $b$  */
1  if  $n == 1$  then
2    return  $ab$ ;
3   $m \leftarrow \lfloor \frac{n}{2} \rfloor$ ;
4   $A_1 \leftarrow \lfloor a/2^m \rfloor$  /* high  $n/2$  bits of  $a$  and  $b$  */
5   $B_1 \leftarrow \lfloor b/2^m \rfloor$  /* implemented by bit-shift */
6   $A_0 \leftarrow a \bmod 2^m$  /* lower  $n/2$  bits of  $a$  and  $b$  */
7   $B_0 \leftarrow b \bmod 2^m$ ;
8   $x \leftarrow \text{Karatsuba}(A_1, B_1, m)$ ;
9   $y \leftarrow \text{Karatsuba}(A_0, B_0, m)$ ;
10  $z \leftarrow \text{Karatsuba}(A_1 + A_0, B_1 + B_0, m)$ ;
11 return  $x^{2^m} + (z - x - y)2^m + y$ ;

```

DP algorithms**Algorithm 11:** SubsetSum(w_1, w_2, \dots, w_n, W)

```

1   $M \leftarrow (n + 1) \times (W + 1)$  /* table (matrix/2D array) */
2  for  $w = 0 \dots W$  do
3    /* set border cases */
4     $M[0][w] = 0$ ;
5  for  $j = 1 \dots n$  do
6    for  $w = 0 \dots W$  do
7      /* Apply recursive formula */
8       $M[j][w] = \max\{w_j + M[j - 1][w - w_j]; M[j - 1][w]\}$ ;
9  return  $M[n][W]$ 

```

Algorithm 12: SubsetSumSolution($M, w = [w_1, w_2, \dots, w_n], W$)

```

1  $S \leftarrow []$  /* set of opt weights */
2 while  $i > 0$  AND  $j > 0$  do
3   if  $M[i][j] > M[i-1][j]$  then
4     /* the case where  $w_i$  is chosen */
5      $S.append(w[i]);$ 
6      $i \leftarrow i - 1;$ 
7   else
8     /*  $w_i$  is not chosen */
9      $i \leftarrow i - 1;$ 
9 return  $S$ 

```