

L19-20 Bi-partite Matching & Revision Lecture

G54AAD Advanced Algorithms and Data Structures

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Matching

Definition (Matching)

A **matching** of a graph $G = (V, E)$ is a subset $M \subset E$ of its edges such that for all distinct edges $\{u, v\} \in M$ and $\{x, y\} \in M$

- $\{u, v\} \cap \{x, y\} = \emptyset$

i.e., no two edges in M share a common vertex.

Example

Bi-Partite Graph

Definition (Bipartite Graph)

A graph $G = (V, E)$ is called **bi-partite** if there exists $S, T \subseteq V$ st.

- $S \cup T = V$ and $S \cap T = \emptyset$ (i.e., S and T is a partition of V), and
- for all $\{u, v\} \in E$, either
 - $u \in S$ and $v \in T$, or
 - $u \in T$ and $v \in S$

Example

Maximum Bi-partite Matching Problem

Problem

Given a bi-partite graph, find a matching with maximal cardinality.

Example (Allocating gates to flights)

Allocate gates to as many flights as possible within a time window.

Flights have gate constraints, e.g.

- some types of aircraft require certain gates
- airlines may have certain requirements

Solving Maximum Bi-partite Matching as Max Flow

Converting a bi-partite graph to a flow network

- ➊ Add an edge from s to all nodes in S
- ➋ Add an edge from t to all nodes in T
- ➌ Give all edges capacity 1.

Solving Maximum Bi-partite Matching as Max Flow

Converting a bi-partite graph to a flow network

- 1 Add an edge from s to all nodes in S
- 2 Add an edge from t to all nodes in T
- 3 Give all edges capacity 1.

Theorem (Integrality Theorem)

If all capacities in a flow network are integer-valued, then the Ford-Fulkerson method produces a maximum flow f where $f(u, v)$ is integer-valued for all pairs (u, v) .

Theorem

The bi-partite graph has a matching of size k if and only if the flow network has a flow of value k .

Outline of revision lecture

- Revision of some of the topics covered
- Example questions from each topic
- Revision material (past exam papers in G64ADS)
- Complete syllabus given on Moodle
(see reading column and lecture notes)

Assessment

Activity	% of Overall Mark
Exam (2 hr written)	100 %
Coursework (report)	25 %

Exam Format¹

G54AAD-E1

The University of Nottingham

SCHOOL OF COMPUTER SCIENCE

A LEVEL 4 MODULE, AUTUMN SEMESTER 2015-2016

G54AAD ADVANCED ALGORITHMS AND DATA STRUCTURES

Time allowed TWO hours

Candidates may complete the front cover of their answer book and sign their desk card but must NOT write anything else until the start of the examination period is announced.

**Answer Question 1 in Section A and
TWO questions in Section B.**

No calculators are permitted in this examination.

Dictionaries are not allowed with one exception. Those whose first language is not English may use a standard translation dictionary to translate between that language and English provided that neither language is the subject of this examination. Subject-specific translation directories are not permitted.

¹Subject to suggestions by external moderator.

General advice

- Format
 - Time allowed is 2 hours.
 - Section A (40%): Answer everything
 - Section B (60%): Answer two of three questions.
- Read the entire paper before answering (plan your time).
- Read questions carefully.
- Answers have to be justified.
- State any needed assumptions.

Basics assumed

- Solving equations, inequalities
- Proof by induction, contradiction etc.
- Properties of basic functions,
eg logarithms, exponential function, factorials etc.

L01-02 Asymptotic Notation & Recurrence relations

Asymptotic notation

- Definitions

Example questions

- Explicit: Classify functions according to asymptotic growth rate.
- Implicit: State answers to other problems in asymptotic notation.

L03-04 The Master Theorem

Recurrences and the Master theorem

- Statement
- Proof (disregarding round-off errors)
- Applicability (when can the theorem be used)

Example questions

- Explicit: Solve some given recurrences.
- Implicit: Runtime analysis of a divide and conquer algorithm, where the runtime can be expressed as a recurrence for which the theorem applies.

L05-06 Leftist Heaps

```
merge :: (Ord a) => LHeap a -> LHeap a -> LHeap a
merge x E = x
merge E y = y
merge h1@(B _ x1 left1 right1) h2@(B _ x2 _ _)
  | x1 <= x2 = B (1 + (rank right1)) x1 left right
  | otherwise = merge h2 h1
  where (left,right) = leftistOrder left1 (merge right1 h2)

leftistOrder :: LHeap a -> LHeap a -> (LHeap a, LHeap a)
leftistOrder h1 h2
  | (rank h1) <= (rank h2) = (h2,h1)
  | otherwise              = (h1,h2)
```

- Definition of heaps, min-heaps, max-heaps.
- Definition of leftist heaps, including leftist heap property
- Analysis of height of leftist heaps
- Implementation and analysis of operations (findMin, merge, insert, delMin)

Example question

- Explicit: Implementation of leftist heaps.
- Implicit: Use some of the ideas behind leftist heaps to analyse other data structures.

L07-08 Binomial Heaps

Binomial trees

- Definition
- Properties: eg rank vs size or depth

Binomial heaps

- Definition
- Properties: eg largest rank vs size etc.
- Implementation and analysis of basic operations `insTree`, `merge`, `findMin`, `deleteMin`

Example question:

- Find size, depth, and number of children in a binomial tree of rank r .

L09 Persistence

Red Black Trees

- Persistent vs ephemeral data structures
- Sharing of data

Example question

- Illustrate sharing of data from example code.

L10 Red-Black Trees

Red Black Trees

- Definition and properties
- Implementation and analysis of operations
eg member, insert (delete not required)

Example question

- Analysis of depth.

L11-12 Dynamic Programming

Dynamic Programming

- Applicable on problems with
 - optimal substructure
 - overlapping substructure
- Longest Common Subsequence
- Floyd-Warshall Algorithm

Example question

- Explicit: Define solution to problem recursively in terms of sub-solutions.
- Implicit: Recognise that a problem can be solved by dynamic programming.

L12-18 Maximum Flow

Maximum flow

- Definition and properties of flows
- Augmenting paths and Ford-Fulkerson method
- Cuts and Max-flow Min-cut Theorem
- Proof ideas

Example questions

- make use of basic properties of flows
- find residual networks and augmenting paths
- recognise that a problem is an instance of max-flow

L19-20 Edmonds-Karp Algorithm

Edmonds-Karp

- Instance of Ford-Fulkerson method
- Augmenting paths by Shortest paths
- Analysis of runtime

Example question

- write pseudo-code
- explain behind runtime analysis

L21 Bi-partite Matching

Bi-partite matching

- problem definition
- formulation as max-flow

Example question

- Show that some real-world problem can be described as bipartite matching.

Revision Material on Moodle

- G54AAD was a new module from 2014/2015, but overlaps with old module G64ADS.
- Three past exam papers available on Moodle
 - G64ADS regular exam 2012/2013
 - G64ADS regular exam 2013/2014
 - G54AAD regular exam 2014/2015
- Relevant exam questions from G64ADS are marked in red.

Thank you!