Example Method approach— method not accepting Parameters but doing 16/11/2 12 Methods in Java and Fundamentals of Stack and Heap area. 11 dust ance Variables Void add (1) O Methods: Methods in Java is Collection -> method of instructions that perform a specific - no input and no ortput nethod a=10; b=20; 1 Method have (=a+b) System.out. Printle(1); 1) input (Parameters) (3) Body (b) return type. Public Class Calculator & Public Static Void main (String [] args){ retumpre name (Parameter) Calculator (alc = new Calculator) ); 7, - Method Syntax Activity/Body. theologish (alc. add (); add Method) is called gualant of start little hay -> instance Mariables, properties, field are referrig The thirty and of for Execution · bottom the first in get Executed. -) Method, Activity, behaviour, does part are referring same. worthed to apply the farther will a

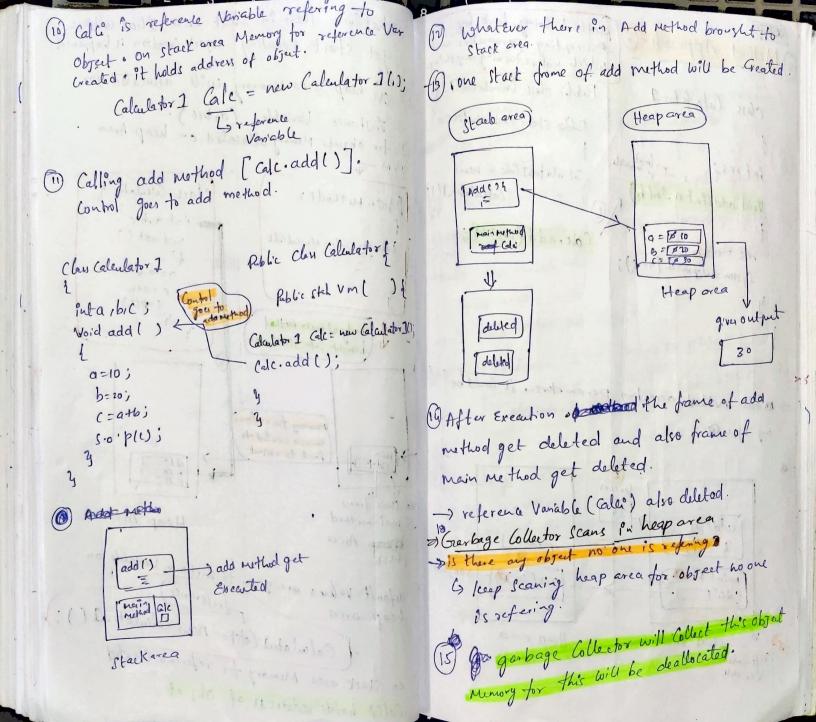
> Calc Managing object type of data of Execution 1) Jum got byte Code of Entire Wn Hen program.

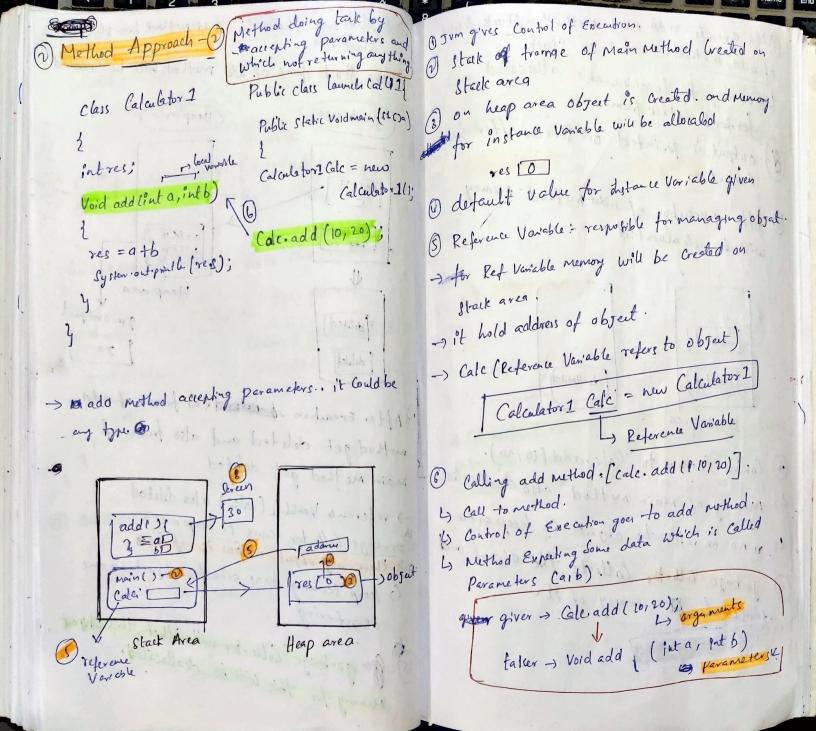
Jum calle Main nethod.

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Jum gives Control of Execution. -) Clan Calculator 1 got life after Creating objet Bon order Execute task drêide main nothod on Stack area one recorde will be Created -> have of the object is Calci Cartivation record of Main Method (or) Start 7 -> JUM Executes Tava program frame of main Method brought into Ly it arranges data area and Executers or it. Italk area. Public class Calculatory Pablic Class Chicaban (Stry 10) 75) E can Calculator I Public Static Void mail ){ Calculator 1 Cale = New Calculated) intabic; void add() Calc. add(); (or) Heap luntime stack area. Stack area a=10; b=20; (=atb; -> Any activity / task in program should not brought 3.0.b(1) in Stark area for Execution. -> whatever there in get Executed. stack area -> JVM got byte Gode of Entire written I sativation record (00) Steak frame of main method. Rogram

B what Ever there in Steek area that will get start Executing. Object Creation is happened. 6) on heap area Memory will allocated for Instance Variable (a,b,e). for objects Memory allocated on heap Area. JUM J Class Calculator 4 (Main Method () Calculator Cole = new (of cult void add () a=10,6=20) body of mein method brought into Itaeli erea memory for reference Variable Created to Point to object stall frame Heap Area. of main method. Area (4) default values are given to distance variable on Preference · Calculatori Calci = new Calculatori (); heap orea. on Stack area Memory for reference Created Calci holds address of object

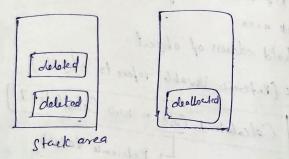




To Execute add mald method Stack frame of add method created on Stack. C) numby for local variable allocated

output is printed on Screen after Execution

(9) After Execution of add method Memory will be de alocated.



- (0) Control. Comes to Calc. add (10, 10).
- L) memoy of main method also deallocated.
- 4) a no on refer to object in heap area.
- (1) garbage collector Collects the object and de allocate to memorey of it:

teles or in the property of the property of

3 Method approach (3) Clas Calculator I Public Class Counch Calc 14 Public Static Vold main (Shy () int arbites;

(Peters) -> int.add ( a= 10; b = 20; res = atb;
return res

Calculator 1 Calc = new Calatator 10; [ Calc. add(); -> taking returned value ] inte=(alc.adde); 3.0.p(c); you prepare boston had

-> it is not accepting parameters. After Execution it is returning a value.

> If method accepting parameters we must give parameters it it is giving (an returning avalue it is capto in to ros take (01) not.

> Variable c is und to take value from add
method.

as experience where in strate, get the central

The goal botter blood of sop today

rough done is propert of police of

[ int c = Calc. add();

on stack area Memory for beal varidle are (4) Method Approach - (4) Created. A. Public class launch Calcal After Execution Memory of Add method Class Calculator & areland Main method get deblocated Rublic State Void main ( ) int add (intarint b) { Calculator 1 Calc = new garbage Collector Collects and object and 11 Calc. addl); Calculatori int reseatb; de allocate nemony of it. returning; # Possible ways of writing Method + Int res = Calc. add (10,20) 5.0. P(res); Void add ( Calc. add (); -> here Method accepting Parameters. and it returning a Value. Ly add Method b = 20. c=atb; 3 (Nowah for spe 5.0'p(c)' @ accepting parameter and not returning value intres void add (intarintb) Call add (10,20); adoren. res=atb; s.o'p (res); Stack frame Steel Heap area of main ll calciaddi); int add ()? orea nethod int C = calliadd 1); a=10) 6) Whatever there in Stack get Executed. S. o. ple) res = atb; 3 return res; 1 Object is created on heap area (3) Control gow to add method. These add int res = calc. add (10,20) int add (inta, intb) { method is brought to stack area and return res; 5.0.D (ras), get Executed.

O Possible ways of writing Method: Code Snippets (21/10/22) 1: 45:32 1) Method not accepts parameters and not return out x=0/1/2/3 Switch (X) any value defent soppade fault "); , fall through Void add ( Calcadd(); Caro: S.o.p("0"); a=10; b = 20; Ladd Method is Car 2; 5.0. p("2"); I fall through. C=a+b; Called. S.o. p(c); @ Method accepts parameters and not return any Value. V.int res; void mainl into jutb) { x=2 010 = 2 , wrappe (lay Booler bi=trus; Primite Cale, addl); 11 b3 = true; f (bitwise operator)

11 (1b1 f b2) 1 (b2 f b3) f b3) -> f > booler res = atb) s.op (res); 3) Method not accepts Parameters and but returns s.o.p("alpha"); if ((b) = falk) ( (b) \$ b3) 1(b) 162) ->f Value. 11 Calc. addl); int adal) 1 x . s.o.p("beta"); a=10; inte-Calvaddelj -> no outpute of p. b=20; res=atb; S.o.p(c) refum res; ((Xola4) 1 + 119) 10) (4) Method accepting parameter Put add (inta , int b) { int res = atb; intres = calciadd 192 return ry; S.o.P (res);