

Lemma 4.1.6

Clexma

February 26, 2019

Lemma 1 (4.1.6) *A flow f is a s - t path flow iff f satisfies the following conditions:*

- *If $s = t$ then*
 1. $\sum_{w \in \text{out}(v)} f(v, w) = \sum_{w \in \text{in}(v)} f(w, v)$ for all v
 2. $F(f)$ is a s - t support
- *If $s \neq t$ then*
 1. $\sum_{w \in \text{out}(v)} f(v, w) = \sum_{w \in \text{out}(v)} f(w, v)$ for all $v \in V - \{s, t\}$
 $\sum_{w \in \text{out}(s)} f(s, w) = \sum_{w \in \text{out}(s)} f(w, s) + 1$
 $\sum_{w \in \text{out}(t)} f(w, t) = \sum_{w \in \text{out}(t)} f(w, t) - 1$
 2. $F(f)$ is connected.

Proof 0.1 *Prove by induction on $n = \sum_{e \in E} f(e)$*

- \Rightarrow : *By definition direction of path flow the proof is obvious.*
- \Leftarrow :
 1. $s = t, F(f)$ is disconnected, we have the conclusion that f is a s - t path flow.
 - case 1 : there is a self loop from s to s .
 - case 2 : there is no self loop from s to s . In this case we delete an edge from v to s such that $v \neq s$.
 - * subcase 2.1 the resulting graph is still connected let the new flow be f' . By the induction hypothesis f' is a path flow from s to v .
 - * subcase 2.2 the resulting graph is not connected, in this case, there are exactly two connected components such in the resulting graph, we call them C_1 and C_2 and assume they contain s and v respectively.
 - sub-subcase 2.2.1: C_2 contains no edges.
 - sub-subcase 2.2.2: C_1 and C_2 both contain at least one edge.