

SVMRANKER: A General Termination Analysis Framework of Loop Programs via SVM Presentation & Demonstration

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Outline

1. Introduction to Ranking Functions
2. Overview of SVMRANKER
3. Demonstration of the Tool in Command Line.

Single Path Linear Constraint Loop

Example

while $(x \geq -z)$ **do** $x' = x + y$, $y' = y + z$, $z' = z - 1$

Let $B = (-1, 0, 1)$, $\mathbf{x} = (x, y, z)^T$, $\mathbf{b} = 0$.

Let $\mathbf{x}'' = (x, y, z, x', y', z')$,

$$A = \begin{bmatrix} 1 & 1 & 0 & -1 & 0 & 0 \\ 0 & 1 & 1 & 0 & -1 & 0 \\ 0 & 0 & 1 & 0 & 0 & -1 \end{bmatrix} \quad (1)$$

and $\mathbf{c} = (0, 0, 1)^T$

Definition (SLC)

while $(B\mathbf{x} \leq \mathbf{b})$ *do* $A \begin{pmatrix} \mathbf{x} \\ \mathbf{x}' \end{pmatrix} \leq \mathbf{c}$

$$A'' = \begin{pmatrix} B & 0 \\ A \end{pmatrix}$$

$$\mathbf{c}'' = \begin{pmatrix} \mathbf{b} \\ \mathbf{c} \end{pmatrix}$$

$$A''\mathbf{x}'' \leq \mathbf{c}''$$

Ranking Functions

Definition (Single Linear Ranking Function(LRF))

$f(x_1, \dots, x_n) = a_1 x_1 + \dots a_n x_n + a_0$, such that

- ▶ $f(\mathbf{x}) \geq 0$ for any \mathbf{x} satisfies the loop constraints.
- ▶ $f(\mathbf{x}) - f(\mathbf{x}') \geq 1$ for any transition from \mathbf{x} to \mathbf{x}' .

Example

while $(x - 1 > 0)$ **do** $x' = x - 5$

LRF: $f(x) = ax + b$.

- ▶ $ax + b \geq 0 \Rightarrow x \geq -\frac{b}{a} = 1$.
- ▶ $ax + b - (ax' + b) = a(x - x') = 5a \Rightarrow 5a \geq 1$

A possible SLRF: $f(x) = x - 1$

Limitation of SLRF

`while ($q > 0$)do $q' = q - y, y' = y + 1$`

Assume there is a LRF for this loop, say $f(q, y) = a_1 q + a_2 y + b$

$$f(q, y) - f(q', y') = a_1 y + a_2$$

Since y is not bounded, we cannot guarantee $\Delta f(q, y, q', y') > 0$

The loop does not has a SLRF, however, it does terminate.

We still wish to use q for ranking function, but to distinguish different “phase” of q base on either $y \geq 0$ or $y < 0$

Nested RF

Definition (Nested Ranking Function)

A tuple $\langle f_1, \dots, f_d \rangle$ is a nested ranking function for T if the following requirements are satisfied for all $\mathbf{x}'' \in T$

$$\begin{aligned} f_d(\mathbf{x}) &\geq 0 \\ (\Delta f_i(\mathbf{x}'') - 1) + f_{i-1}(\mathbf{x}) &\geq 0 \quad \text{for all } i = 1, \dots, d. \end{aligned}$$

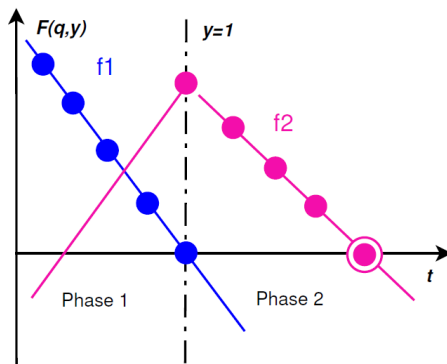
Let $f_0 = 0$.

Example: Nested RF

$$\begin{aligned} f_d(\mathbf{x}) &\geq 0 \\ (\Delta f_i(\mathbf{x}'') - 1) + f_{i-1}(\mathbf{x}) &\geq 0 \quad \text{for all } i = 1, \dots, d. \end{aligned}$$

while $(q > 0)$ do $q' = q - y, y' = y + 1$

- Above loop has Nested RF $\langle 1 - y, q + 1 \rangle$



Linear Loop Program

Definition

A linear loop program $\text{LOOP}(x, x')$ is a binary relation defined by a formula with the free variables x and x' of the form

$$\bigvee_{i \in I} (A_i(x) \leq b_i \wedge C_i(x') < d_i)$$

for some finite index set I .

Example

`while (q > 0){if (y > 0) : q' = q - y - 1; else : q' = q + y - 1}`

can be represented by

$$\begin{aligned} & (q > 0 \wedge y > 0 \wedge y' = y \wedge q' = q - y - 1) \\ \vee & (q > 0 \wedge y \leq 0 \wedge y' = y \wedge q' = q + y - 1) \end{aligned}$$

Limitation of Nested RF

Example

```
while ( $q > 0 \vee y > 0$ )  
{ if ( $y > 0$ ) :  $y' = y - 1$ ;  $q' = q$ ; else :  $q' = q - 1$  }
```

This program does not have a nested ranking function for we require $f_d \geq 0$ but the guard is $q > 0 \vee y > 0$.

However, this loop does terminate. Then we use a “multi-phase” ranking function $\langle y, q \rangle$ to prove the termination.

$$f_d(\mathbf{x}) \geq 0$$
$$(\Delta f_i(\mathbf{x}'') - 1) + f_{i-1}(\mathbf{x}) \geq 0 \quad \text{for all } i = 1, \dots, d.$$

Multiphase Ranking Function

Definition

Given a set of transitions $T \subseteq \mathbb{Q}^{2n}$, we say $\langle f_1, \dots, f_d \rangle$ is a multiphase ranking function for T if for every $\mathbf{x}'' \in T$, there is an index $i \in [1, d]$, s.t.

$$\begin{aligned} \forall j \leq i . \quad & \Delta f_j(\mathbf{x}'') \geq 1, \\ & f_i(\mathbf{x}) \geq 0, \\ \forall j < i . \quad & f_j(\mathbf{x}) \leq 0. \end{aligned}$$

We say that \mathbf{x}'' is ranked by f_i (for the minimal).

Example: Multiphase Ranking Function

while $(x > -z)$ **do** $x' = x + y, y' = y + z, z = z - 1$

Attempt to use a ranking function that has several phases:

$\langle z + 1, y + 1, x \rangle$

x	y	z	$z + 1$	$y + 1$	x
1	1	1	2	2	1
2	2	0	1	3	2
4	2	-1	0	3	4
6	1	-2	-1	2	6
7	-1	-3	-2	0	7
6	-4	-4	-3	-3	6
2	-8	-5	-4	-7	2
-6	-13	-6	-5	-12	-6

Example: Multiphase Ranking Function

while $(x > -z)$ **do** $x' = x + y, y' = y + z, z' = z - 1$
 $\langle z + 1, y + 1, x \rangle$

\mathbf{x}'' is ranked by f_k when $i = k$. In this example, $f_1(x, y, z) = z + 1$,
 $f_2(x, y, z) = y + 1$ and $f_3(x, y, z) = x$

$$\begin{aligned} \forall j \leq i. \Delta f_j(\mathbf{x}'') &\geq 1, \\ f_i(\mathbf{x}) &\geq 0, \\ \forall j < i. f_j(\mathbf{x}) &\leq 0. \end{aligned}$$

