

# Linear Invariant Generation Using Non-Linear Constraint Solving

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CAV'03 and SAS'04

Report date: March 23, 2021

# Problem and Contribution

- ▶ Problem: Generation of linear invariant for linear transition system.
- ▶ Contribution: An exact method for finding the invariant which avoid the widening operator in the classical abstract interpretation.

# Transition System and Invariant

**Definition 1 (Transition System)** A *transition system*  $P : \langle V, L, l_0, \Theta, \mathcal{T} \rangle$  consists of a set of *variables*  $V$ , a set of *locations*  $L$ , an *initial location*  $l_0$ , an *initial assertion*  $\Theta$  over the variables  $V$ , and a set of *transitions*  $\mathcal{T}$ . Each transition  $\tau \in \mathcal{T}$  is a tuple  $\langle l, l', \rho_\tau \rangle$ , where  $l, l' \in L$  are the *pre* and *post locations*, and  $\rho_\tau$  is the *transition relation*, an assertion over  $V \cup V'$ , where  $V$  represents current-state variables and its primed version  $V'$  represents the next-state variables.

**Definition 2 (Inductive Assertion Map)** Given a program  $P$  with a cutset  $C$  and an assertion  $\eta_c(l)$ , for each cutpoint  $l$ , we say that  $\eta_c$  is an *inductive assertion map* for  $C$  if it satisfies the following conditions for all cutpoints  $l, l'$ :

**Initiation** For each basic path  $\pi$  from  $l_0$  to  $l$ ,  $\Theta \wedge \rho_\pi \models \eta_c(l)$ .

**Consecution** For each basic path  $\pi$  from  $l$  to  $l'$ ,  $\eta_c(l) \wedge \rho_\pi \models \eta_c(l')$ .

# Linear Constraint

# Farkas' Lemma

**Theorem 2.5** (Farkas' lemma). *The system  $Ax = b$  has a nonnegative solution if and only if there is no vector  $y$  satisfying  $y^T A \geq 0$  and  $y^T b < 0$ .*

Intuitive understanding of farkas lemma.

**Corollary 2.5b.** *Suppose that the system  $Ax \leq b$  has at least one solution. Then for every solution  $x$  of  $Ax \leq b$  one has  $c^T x \leq \delta$  if and only if there exists a vector  $y \geq 0$  such that  $y^T A = c^T$  and  $y^T b \leq \delta$ .*

# Farkas' Lemma

A better demonstration of the Corollary.

**Theorem 1 (Farkas' Lemma).** *Consider the following system of linear inequalities over real-valued variables  $x_1, \dots, x_n$ ,*

$$S : \begin{bmatrix} a_{11}x_1 + \dots + a_{1n}x_n + b_1 \leq 0 \\ \vdots \\ a_{m1}x_1 + \dots + a_{mn}x_n + b_m \leq 0 \end{bmatrix}$$

*When  $S$  is satisfiable, it entails a given linear inequality*

$$\psi : c_1x_1 + \dots + c_nx_n + d \leq 0$$

*if and only if there exist non-negative real numbers  $\lambda_0, \lambda_1, \dots, \lambda_m$ , such that*

$$c_1 = \sum_{i=1}^m \lambda_i a_{i1}, \quad \dots, \quad c_n = \sum_{i=1}^m \lambda_i a_{in}, \quad d = \left( \sum_{i=1}^m \lambda_i b_i \right) - \lambda_0$$

*Furthermore,  $S$  is unsatisfiable if and only if the inequality  $1 \leq 0$  can be derived as shown above.*

# Solving the Invariant of Transition System

# Quantifier Elimination

- ▶ Exact quantifier elimination.
- ▶ Under-approximate elimination approach