

On Multiphase-Linear Ranking Functions

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Contributions

- ▶ Equivalence of different classes of ranking function.
- ▶ Algorithms for converting between ranking functions.
- ▶ Complete solution for ranking functions on integer.
- ▶ Depth bound and iteration bound for $M\Phi RF$.

Single Path Linear Constraint Loop

Example

`while` $(x \geq -z)$ `do` $x' = x + y$, $y' = y + z$, $z' = z - 1$

`while` $(x_2 - x_1 \leq 0, x_1 + x_2 \geq 1)$ `do` $x_2' = x_2 - 2x_1 + 1$, $x_1' = x_1$

Definition (SLC)

while $(B\mathbf{x} \leq \mathbf{b})$ *do* $A \begin{pmatrix} \mathbf{x} \\ \mathbf{x}' \end{pmatrix} \leq \mathbf{c}$

$$A'' = \begin{pmatrix} B & 0 \\ A \end{pmatrix}$$

$$\mathbf{c}'' = \begin{pmatrix} \mathbf{b} \\ \mathbf{c} \end{pmatrix}$$

$$A''\mathbf{x}'' \leq \mathbf{c}''$$

Ranking Functions

Definition (Linear Ranking Function(LRF))

$f(x_1, \dots, x_n) = a_1x_1 + \dots a_nx_n + a_0$, such that

- ▶ $f(\mathbf{x}) \geq 0$ for any \mathbf{x} satisfies the loop constraints.
- ▶ $f(\mathbf{x}) - f(\mathbf{x}') \geq 1$ for any transition from \mathbf{x} to \mathbf{x}' .

Example

`while (x - 1 > 0)do x' = x - 1`

Its LRF: $f(x) = x - 1$

Example: Multiphase Ranking Function

Problem: LRF is not strong enough for all loops.

Example

`while ($x > -z$)do $x' = x + y, y' = y + z, z = z - 1$`

$f(x, y, z) = a_1x + a_2y + a_3z + b$

y cannot be used for non-existence of its lower bound.

$f(x, y, z) = x + z$

Problem?

Example: Multiphase Ranking Function

while $(x > -z)$ do $x' = x + y, y' = y + z, z = z - 1$

Attempt to use a ranking function that has several phases:

$\langle z + 1, y + 1, x \rangle$

x	y	z	$z + 1$	$y + 1$	x
1	1	1	2	2	1
2	2	0	1	3	2
4	2	-1	0	3	4
6	1	-2	-1	2	6
7	-1	-3	-2	0	7
6	-4	-4	-3	-3	6
2	-8	-5	-4	-7	2
-6	-13	-6	-5	-12	-6

Multiphase Ranking Function

Definition

Given a set of transitions $T \subseteq \mathbb{Q}^{2n}$, we say $\langle f_1, \dots, f_d \rangle$ is a multiphase ranking function for T if for every $\mathbf{x}'' \in T$, there is an index $i \in [1, d]$, s.t.

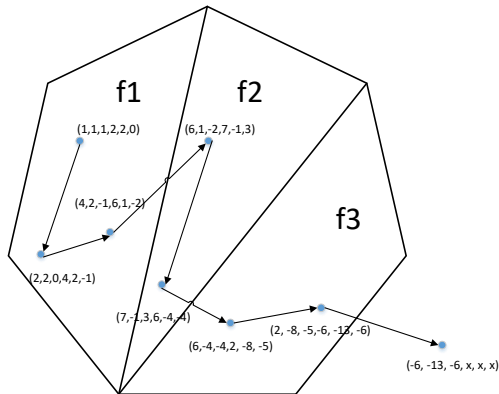
$$\begin{aligned} \forall j \leq i . \Delta f_j(\mathbf{x}'') &\geq 1, \\ f_i(\mathbf{x}) &\geq 0, \\ \forall j < i . f_j(\mathbf{x}) &\leq 0. \end{aligned}$$

We say that \mathbf{x}'' is ranked by f_i (for the minimal).

Example Revisit

`while (x > -z)do x' = x + y, y' = y + z, z = z - 1`

$$\begin{aligned} \forall j \leq i . \Delta f_j(\mathbf{x}'') &\geq 1, \\ f_i(\mathbf{x}) &\geq 0, \\ \forall j < i . f_j(\mathbf{x}) &\leq 0. \end{aligned}$$



Nested Ranking Function

`while ($x > -z$)do $x' = x + y, y' = y + z, z = z - 1$`

Loop condition: $x + z > 0$. We only want to use this constraint for the ranking function.

$\langle z + 1, y + 1, x + z \rangle$

Definition (Nested Ranking Function)

A tuple $\langle f_1, \dots, f_d \rangle$ is a nested ranking function for T if the following requirements are satisfied for all $\mathbf{x}'' \in T$

$$\begin{aligned} f_d(\mathbf{x}) &\geq 0 \\ (\Delta f_i(\mathbf{x}'') - 1) + f_{i-1}(\mathbf{x}) &\geq 0 \quad \text{for all } i = 1, \dots, d. \end{aligned}$$

Let $f_0 = 0$.

