

# 基于 SVM 算法的循环程序的终止性证明

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# Verification

```
1 assume (True);  
2 int x = 1;  
3 while (x>0)  
4 {  
5     x++;  
6 }  
7 assertion (x<=0);
```

- Step 1: Can the program reach the assertion?
- Step 2: Does the assertion hold?

# Verification

```
1 assume( True );  
2 int x = 1;  
3 while( x>0 )  
4 {  
5     x++;  
6 }  
7 assertion( x<=0 );
```

- It is very important to check whether the program is **TERMINATING** in the verification problem.

# Bad news & Good news

- In theory,  
the termination problem of programs has been proven to be undecidable.
- In practice,  
just return "UNKNOWN" when we can not prove whether it is terminating.

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Goal:

For a given program, we try to avoid "UNKNOWN" result as much as possible.

# Background

- It attracts many researchers to work on the termination of programs.
- **Ultimate Automizer**, one of the leading tools in program analysis according to the outcomes of the **SV-COMP** competitions.
- The termination of **loops** is at the core of the termination analysis techniques used in **Ultimate Automizer**.

```
1 while(g_1) \\Loop 1
2 {
3     ....;
4     ....;
5     ....;
6 }
```

the termination of simple loop programs, i.e., **no nested loops**.

# Ranking function (RF)

- Informal definition

$$f : S \rightarrow \mathbb{R}$$

where  $S$  is the set of program states,  $\mathbb{R}$  is a well-founded ordered set.

- Decrease condition

$$\forall x, x' : f(x) - f(x') \geq \delta > 0$$

where  $\delta \in \mathbb{R}$ ,  $x$  is current program state,  $x'$  is the program state after updating.

- Bounded condition

$$\forall x : f(x) \geq C$$

where  $C \in \mathbb{R}$ .



# Ranking function (RF)

Thm.

A loop program  $P$  is terminating, if  $P$  has a ranking function  $f(x)$ .

# Synthesis of ranking functions

- Termination of programs is undecidable.
- Synthesis of ranking function is, however, decidable, given certain classes of ranking functions.

# Related works

- **Linear program & Single phase linear ranking function**
  - In 2002, Colón and Sipma synthesized linear ranking functions for linear-constraint loops.
  - In 2004, a complete and efficient solution was proposed by Podelski and Rybalchenko.
  - ...
- **Linear program & K phases ranking function**
  - In 2005, Bradley et al. showed how to synthesize lexicographic linear ranking functions (LLRFs).
  - In 2014, Ben-Amram and Genaim provided a complete polynomial-time solution for M $\Phi$ RFs with bounded depth.
  - ...
- **Polynomial program & Single phase ranking function**
  - In 2005, Cousot made use of parametric abstraction and SDP to compute ranking functions of loops.
  - In 2019, Yuan et al. proposed a ranking function detection method exploiting SVM .
  - ...

# Our contributions

- Linear program & Single phase linear ranking function
- Linear program & K phases ranking function
- Polynomial program & Single phase ranking function
- Polynomial program & K phases ranking function
  - Based on SVM to synthesize nested ranking functions.
  - Provide the comprehensive empirical evaluation.

# Preliminaries

## Def. (Loop Program)

A loop program  $\Omega(x, x')$  is a binary relation with free variables  $x$  and  $x'$ , where  $x$  is the current state, and  $x'$  is the next state.

```
1 while (x>0)
2 {
3   x++;
4 }
```

$$\Omega(x, x') \triangleq x > 0 \wedge x' = x + 1$$

## Def. ( $k$ -Nested Ranking Function)

Given a loop program  $\Omega$ , let  $k \in \mathbb{N}_{>0}$  and, for each  $i \in \{1, \dots, k\}$ ,  $f_i(\mathbf{x})$  be a polynomial or an algebraic fraction over the program variables  $\mathbf{x}$ . We call the  $k$ -tuple  $\langle f_1, f_2, \dots, f_k \rangle$  a  $k$ -nested ranking function of  $\Omega$  if the following condition holds for a set of parameters  $\{C_i \in \mathbb{R}_{>0} \mid 1 \leq i \leq k+1\}$ :

$$\forall (\mathbf{x}, \mathbf{x}') \in \Omega : \begin{cases} f_1(\mathbf{x}) - f_1(\mathbf{x}') \geq C_1 \\ f_2(\mathbf{x}) - f_2(\mathbf{x}') + f_1(\mathbf{x}) \geq C_2 \\ \vdots \\ f_k(\mathbf{x}) - f_k(\mathbf{x}') + f_{k-1}(\mathbf{x}) \geq C_k \\ f_k(\mathbf{x}) \geq C_{k+1} \end{cases}$$

# Example

```
1 int q,y;  
2 while(q>0)  
3 {  
4     q = q-y;  
5     y = y+1;  
6 }
```

- Single phase linear ranking function does not work.
- However, it has a 2-nested ranking function.

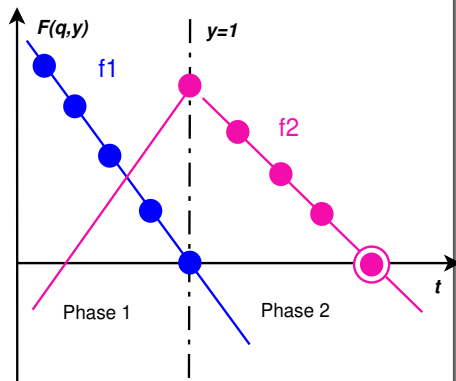
$$f_1(q,y) = 1 - y, f_2(q,y) = q + 1, C_1 = C_2 = C_3 = 1;$$

# Example

## Loop Program

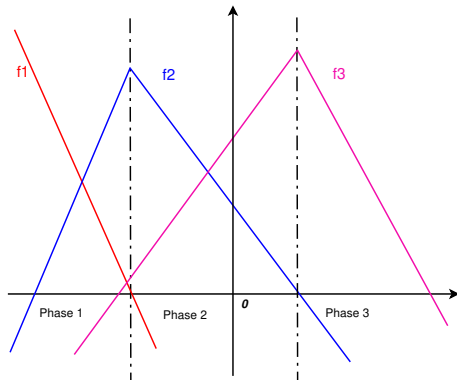
```
1 int q, y;  
2 while (q > 0)  
3 {  
4   q = q - y;  
5   y = y + 1;  
6 }
```

## Phases



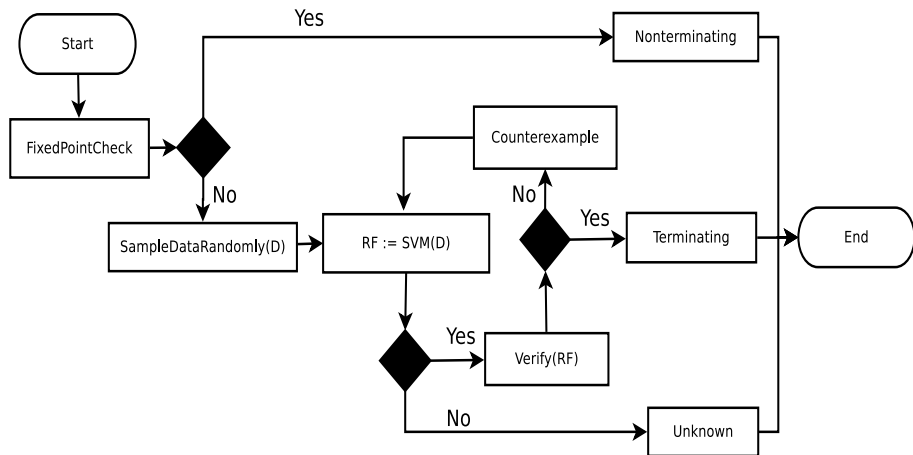


# Intuition



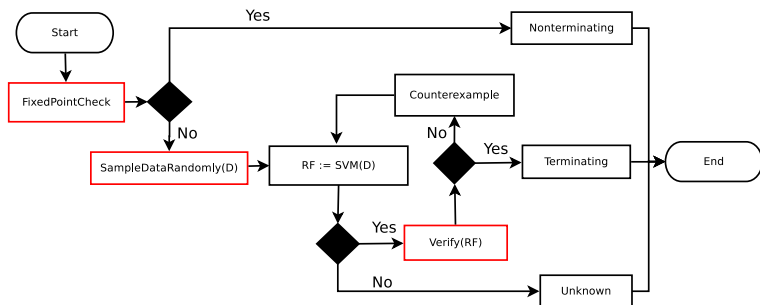
- We allow that the ranking function can increase in the previous phases.
- It will finally decrease in their own phase and the following phases.

# Overview of the Algorithm



# Problems

- How to check the fixed point?
- How to verify the correctness of learned ranking function?
- How to construct the data set?



# How to check the fixed point?

## Def. (Fixed point)

Given a loop program specified by  $\Omega$ , we say that  $\mathbf{x} \in \mathbb{R}^n$  is a fixed point of the loop if  $(\mathbf{x}, \mathbf{x}) \in \Omega$ .

# Example

```
1 while (x>0 && y>0)
2 {
3   x = x+y;
4   x = x-y;
5 }
```

$$\begin{aligned} X' &= (x', y') = ((x + y) - y, y) = (x, y) = X \\ \implies \forall X \in \mathbb{R}_{>0}^2, (X, X) &\in \Omega \\ \implies \forall X \in \mathbb{R}_{>0}^2, X &\text{ is a fixed point.} \end{aligned}$$

# How to verify the correctness of learned ranking functions?

- Decreased condition & Bounded condition
- SMT tools, like Z3, etc.

Check for single phase ranking function

$$\forall (x, x') \in \Omega : f(x) \geq C \wedge f(x) - f(x') \geq \delta$$

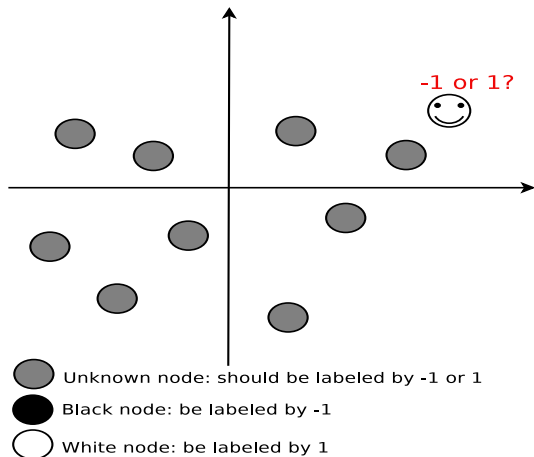
# How to verify the correctness of learned ranking functions?

Check for nested ranking functions

$$\forall (x, x') \in \Omega : \begin{cases} f_1(x) - f_1(x') \geq C_1 \\ f_2(x) - f_2(x') + f_1(x) \geq C_2 \\ \vdots \\ f_k(x) - f_k(x') + f_{k-1}(x) \geq C_k \\ f_k(x) \geq C_{k+1} \end{cases}$$

# How to construct the data set?

- Relation between the SVM & the sampled data set.
- Relation between the SVM & the ranking function.





# Ranking function

Let  $\langle f_1, \dots, f_k \rangle$  be a  $k$ -tuple representing a  $k$ -nested ranking functions;

$$f_j(x) = a_j^T \cdot U_j(x)$$

where  $a_j = (a_{j,1}, \dots, a_{j,s_j})$  is a real vector of coefficients and  $U_j(x) = (U_{j,1}(x), \dots, U_{j,s_j}(x))^T$  is an  $s_j$ -tuple with  $U_{j,i}(x) = \frac{q_{j,i}(x)}{p_{j,i}(x)}$ , where  $q_{j,i}(x), p_{j,i}(x) \in \mathbb{R}[x]$ , for  $i \in \{1, \dots, s_j\}$ .

# Ranking function

Exp.

$$f_j(x) = 3x^2 - 4xy + \frac{5y^3}{3x^3 + 2y + 1} + 7$$

$$f_j(x) = a_j^T \cdot U_j(x)$$

where  $a_j^T = (3, -4, 5, 7)$  and  $U_j(x) = (x^2, xy, \frac{y^3}{3x^3+2y+1}, 1)^T$

# Nested ranking function

$$\forall (x, x') \in \Omega : \begin{cases} f_1(x) - f_1(x') \geq C_1 \\ f_2(x) - f_2(x') + f_1(x) \geq C_2 \\ \vdots \\ f_k(x) - f_k(x') + f_{k-1}(x) \geq C_k \\ f_k(x) \geq C_{k+1} \end{cases}$$

$\Rightarrow$

$$\forall (x, x') \in \Omega : \begin{cases} a_1^T \cdot (U_1(x) - U_1(x')) \geq C_1 \\ a_2^T \cdot (U_2(x) - U_2(x')) + a_1^T \cdot U_1(x) \geq C_2 \\ \vdots \\ a_k^T \cdot (U_k(x) - U_k(x')) + a_{k-1}^T \cdot U_{k-1}(x) \geq C_k \\ a_k^T \cdot U_k(x) \geq C_{k+1} \end{cases}$$

# Substitute

$$\begin{aligned} G_1(x, x') \mapsto & \begin{pmatrix} 0 \\ \vdots \\ 0 \\ 0 \\ U_1(x) - U_1(x') \end{pmatrix}, \quad G_2(x, x') \mapsto \begin{pmatrix} 0 \\ \vdots \\ 0 \\ U_2(x) - U_2(x') \\ U_1(x) \end{pmatrix}, \quad \dots \\ \dots, \quad G_k(x, x') \mapsto & \begin{pmatrix} U_k(x) - U_k(x') \\ U_{k-1}(x) \\ 0 \\ \vdots \\ 0 \end{pmatrix}, \quad G_{k+1}(x, x') \mapsto \begin{pmatrix} U_k(x) \\ 0 \\ 0 \\ \vdots \\ 0 \end{pmatrix} \end{aligned}$$

$$\forall (x, x') \in \Omega : \left\{ \begin{array}{l} (a_k^T, \dots, a_1^T) \cdot G_1(x, x') \geq C_1 \\ (a_k^T, \dots, a_1^T) \cdot G_2(x, x') \geq C_2 \\ \vdots \\ (a_k^T, \dots, a_1^T) \cdot G_k(x, x') \geq C_k \\ (a_k^T, \dots, a_1^T) \cdot G_{k+1}(x, x') \geq C_{k+1} \end{array} \right. \quad (1)$$

where  $C_i > 0, i \in 1..k+1$ .

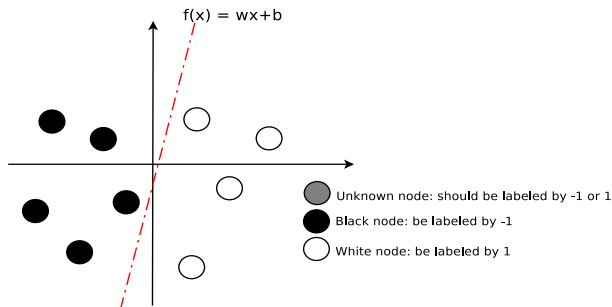
# SVM

- Positive examples:

$$f(x) = wx + b > 0$$

- Negative examples:

$$f(x) = wx + b < 0$$



# Positive examples

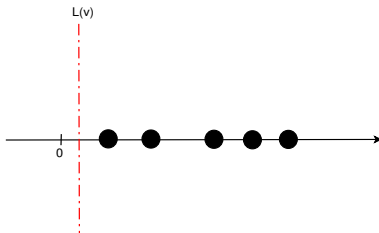
$$\forall (x, x') \in \Omega : \begin{cases} w \cdot G_1(x, x') + b > 0 \\ w \cdot G_2(x, x') + b > 0 \\ \vdots \\ w \cdot G_k(x, x') + b > 0 \\ w \cdot G_{k+1}(x, x') + b > 0 \end{cases}$$

where  $w = (a_k^T, \dots, a_1^T)$ ,  $b = 0$ .

Thm.

Given a loop specified by  $\Omega$ , it has nested polynomial ranking functions as defined in Eq. 1 if and only if there exists a hyperplane  $L(v)$  strictly separating the origin  $O \in \mathbb{R}^m$  from  $G(\Omega) \subseteq \mathbb{R}^m$ .

Proof.



□



# Example

```
1 int q,y;  
2 while(q>0)  
3 {  
4     q = q-y;  
5     y = y+1;  
6 }
```

---

$$f_1(q, y) = 1 - y, f_2(q, y) = q + 1, C_1 = C_2 = C_3 = 1;$$

# Example

## Template

$$f_1(q, y) = a_1^T \cdot U_1(q, y) = (a_{11}, a_{12}, a_{13}) \cdot (q, y, 1)$$

$$f_2(q, y) = a_2^T \cdot U_2(q, y) = (a_{21}, a_{22}, a_{23}) \cdot (q, y, 1)$$

---

## Constraint

$$\forall ((q, y), (q, y)') \in \Omega : \begin{cases} a_1^T \cdot (U_1(q, y) - U_1(q', y')) \geq C_1 \\ a_2^T \cdot (U_2(q, y) - U_2(q', y')) + a_1^T \cdot U_1(q, y) \geq C_2 \\ a_2^T \cdot U_2(q, y) \geq C_3 \end{cases}$$

$$\begin{aligned}G_1(q, y, q', y') &\mapsto \begin{pmatrix} 0 \\ U_1(q, y) - U_1(q', y') \end{pmatrix}, \\G_2(q, y, q', y') &\mapsto \begin{pmatrix} U_2(q, y) - U_2(q', y') \\ U_1(q, y) \end{pmatrix}, \\G_3(q, y, q', y') &\mapsto \begin{pmatrix} U_2(q, y) \\ 0 \end{pmatrix}\end{aligned}$$

# Constraint

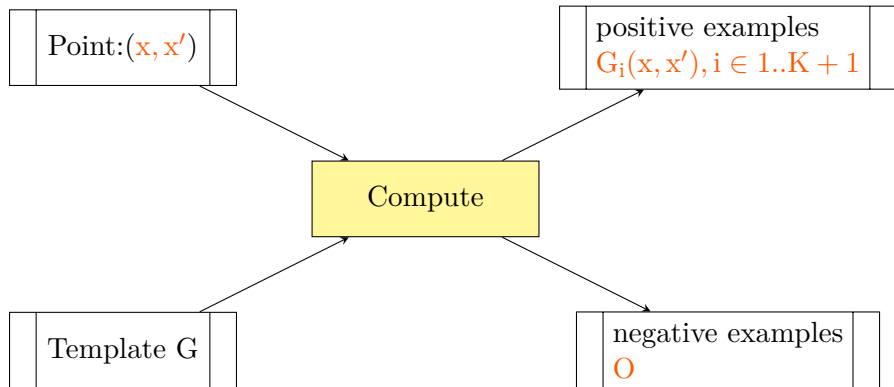
$$\forall (q, y, q', y') \in \Omega : \begin{cases} (a_2^T, a_1^T) \cdot G_1(q, y, q', y') \geq C_1 \\ (a_2^T, a_1^T) \cdot G_2(q, y, q', y') \geq C_2 \\ (a_2^T, a_1^T) \cdot G_3(q, y, q', y') \geq C_3 \end{cases}$$

where  $C_i > 0, i \in 1..3$ .

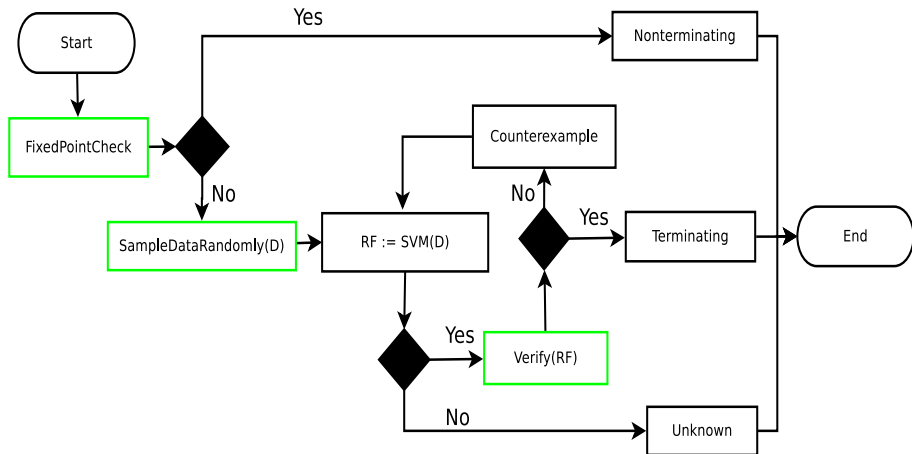
$$\forall (q, y, q', y') \in \Omega : \begin{cases} w \cdot G_1(q, y, q', y') + b \geq C_1 > 0 \\ w \cdot G_2(q, y, q', y') + b \geq C_2 > 0 \\ w \cdot G_3(q, y, q', y') + b \geq C_3 > 0 \end{cases}$$

where  $w = (a_2^T, a_1^T)$ ,  $b = 0$ .

# Data set



# Algorithm



# Algorithm

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**Algorithm 1:** The SVM-based algorithm for synthesizing  $k$ -nested ranking functions

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**Input:** Program  $\Omega$ , initial sample size  $n$ , functions template  $U = \langle U_1, \dots, U_k \rangle$

**Output:** The  $k$ -nested ranking coefficients  $(\mathbf{a}_k^T, \dots, \mathbf{a}_1^T)$  if  $\Omega$  is well-founded;  
“nonterminating” if  $\Omega$  contains fixed points; “unknown” otherwise

```
1 begin
2   if HASFIXEDPOINT( $\Omega$ ) then
3     | return (nonterminating, GETFIXEDPOINT( $\Omega$ ));
4    $D := \{(\mathbf{O}, -1)\}$ ;
5   for  $i := 1$  to  $n$  do
6     | Sample  $(\mathbf{x}, \mathbf{x}')$  randomly from  $\Omega$ ;
7     |  $D := D \cup \text{GETDATAPOINT}(U, (\mathbf{x}, \mathbf{x}'))$ ;
8   while true do
9     |  $svm := \text{SVM}(D)$ ;
10    | if  $svm = (\mathbf{a}_k^T, \dots, \mathbf{a}_1^T)$  then
11      | check := VERIFY( $svm, U, \Omega$ );
12      | if check = true then
13        | return (terminating,  $svm$ );
14      | else // check is a counterexample of the form  $(\mathbf{x}, \mathbf{x}')$ 
15        |  $D := D \cup \text{GETDATAPOINT}(U, \text{check})$ ;
16    | else // SVM failed to separate  $\mathbf{O}$  from the points from  $G(\Omega)$ 
17      | return unknown
```

---



# Experiment

- Implement algorithm in a prototype tool named **SVMRanker**.
- Compare with **LassoRanker**, which is the main part of the **Ultimate Automizer** for proving termination of loop programs.
- Linear program
  - All program cases are adapted from the programs in the official website of **LassoRanker**.
- Non-linear program
  - All cases are adapted from the programs in related papers which focus on the termination of non-linear programs.

# Result

	Terminating	Non-terminating	Unknown	Timeout
Dataset	65	69	-	-
SVMRanker	40	34	0	60
LassoRanker	24	37	73	0
Common cases	19	34	0	0

Tab.: Summary of the experiments

# Conclusion & Future Work

- Conclusion
  - Based on SVM to synthesize nested ranking functions.
    - Able to deal with the polynomial programs with nested ranking functions.
    - Show the relation between the nested ranking functions and the SVM algorithm.
  - Provide the comprehensive empirical evaluation.
    - Implement algorithm in a prototype tool named **SVMRanker**.
    - Solve many cases with the non-linear ranking function.
- Future Work
  - Try to apply the learning algorithm to prove the non-termination of programs.
  - Find more efficient SMT tools to verify the ranking function.

# Experience of Master Student

- Yi Li, **Xuechao Sun**, Yong Li, Andrea Turrini, Lijun Zhang: Synthesizing Nested Ranking Functions for Loop Programs via SVM. ICFEM 2019: 438-454.(CCF- C 类会议, 大会报告作者)
- Yong Li, **Xuechao Sun**, Andrea Turrini, Yu-Fang Chen, Junnan Xu: ROLL 1.0: -Regular Language Learning Library. TACAS (1) 2019: 365-371.(CCF-B 类会议, 工具短文, 大会报告作者)

Thanks! & Questions?