

# Title

## Subtitle

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## Problem of the Week (Monday, April 2, 2018 ~ Saturday, April 7, 2018)

- (a) Given an array  $A[0 \cdots n - 1]$ , to determine whether there is a value that *occurs more than  $\lfloor n/k \rfloor$  times* in  $\Theta(n \lg k)$  time and  $\Theta(k)$  extra space.
- (b) Prove that the *lower bound* of this problem is  $\Theta(n \lg k)$ .

# TIP#1

(Monday, April 2, 2018)

Take  $k = 2$ .

$\Theta(n)$  time &  $\Theta(1)$  space

# TIP#2

(Tuesday, April 3, 2018)

## 定义 ( $k$ -simplified Multiset)

Consider a multiset  $\mathcal{M}$ . A  *$k$ -simplified multiset* for  $\mathcal{M}$  is a multiset derived from  $\mathcal{M}$  by repeating *deleting  $k$  distinct elements* from it until no longer possible.

## 定理

*If a value occurs more than  $\lfloor n/k \rfloor$  times in  $\mathcal{M}$  of  $n$  elements, then it is in a  $k$ -simplified multiset for  $\mathcal{M}$ .*

Prove this theorem. Take  $k = 2$  again. Design an  $\Theta(n)$  algorithm for  $k = 2$ . Generalize it to an algorithm for general  $k$  (ignoring  $\Theta(n \lg k)$  for now).

# TIP#3

(Wednesday, April 4, 2018)

Today, you have an efficient data structure  $T$  for a multiset:

We denote a multiset by  $\mathcal{M} = \text{set}(v_i, c_i)$ , where  $c_i$  is the number of times  $v_i$  occurs in  $\mathcal{M}$ . The number of distinct values in  $\mathcal{M}$  is denoted by  $d = |\text{set} v_i|$ .

The data structure  $T$  for  $\mathcal{M}$  contains  $d$  nodes, each being a pair  $(v_i, c_i)$ .

It supports INSERT, DELETE, and SEARCH in  $\Theta(\lg d)$ .

Use this data structure (you are not required to implement it) in your algorithm to achieve the time complexity of  $\Theta(n \lg k)$ .

# TIP#4

(Thursday, April 5, 2018)

## 定义 ( $r$ -list)

Let  $r = \lfloor n/k \rfloor$ . An  $r$ -list is a list of  $n$  values such that each of the values  $0, 1, \dots, \lfloor n/r \rfloor - 1$  occurs  $r$  times and the value  $\lfloor n/r \rfloor$  occurs  $(n \bmod r)$  times.

## 引理

*There are at least  $(k/e)^n$  different  $r$ -lists.*

## 定理

*Executing these  $r$ -lists on a decision tree will follow different paths.*

# TIP#5

(Friday, April 6, 2018)

定义 ( $L = L_1 * L_2$ )

Given two lists  $L_1$  and  $L_2$ , a new list  $L = L_1 * L_2$  is defined as follows:

$$L[i] = \min \{L_1[i], L_2[i]\}, 0 \leq i < n$$

定理

If  $r$ -lists  $L_1$  and  $L_2$  are different, then  $L = L_1 * L_2$  *contains* a value which occurs more than  $r$  times.

Prove the theorem on Thursday *by contradiction*, with the help of the definition and theorem above.



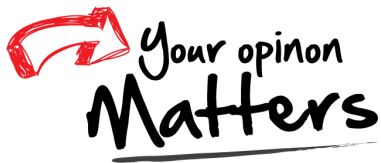
# TIP#6

(Saturday, April 7, 2018)

Paper (Clickable):

*Finding Repeated Elements* (J. Misra, David Gries, 1982)

Thank  
You!



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