

Recall:

Notation: \mathbb{N} = natural numbers
 $\mathbb{N}_0 = \mathbb{N} \cup \{0\}$.

$\mathcal{I}_N = \{A \mid A \subseteq \mathbb{N}_0\}$
 $P_N : \mathcal{I}_N \rightarrow [0, 1]$ (P = \mathbb{P} , intuitive)
 (Uniform Probability)

Simple random walks on \mathbb{Z} .

Fix $N \in \mathbb{N}$.

$$\Omega_N = \{\omega = (w_1, \dots, w_N) \mid w_i \in \{-1, +1\}\}^N$$

Define for $1 \leq k \leq N$

$$X_k : \Omega_N \rightarrow \{-1, 1\} \text{ by } X_k(\omega) = w_k \\ (\text{step of random walk at time } k)$$

$$S_0(\omega) = 0 ; S_n(\omega) = \sum_{k=1}^n X_k(\omega) \quad 1 \leq n \leq N \\ (\text{position of walk at time } n)$$

(i.e. any sequence in $\{-1, 1\}^N$ has the same probability)

Definition: Sequence of random variables $\{S_n\}_{n=0}^N$ on probability space $(\Omega_N, \mathcal{I}_N, P_N)$ is called a simple symmetric random walk of length N starting at 0.

1.2 Stopping times:

Interpretation: S_n represents "amount of capital" of the player after n rounds

X_k = amount a player wins in round k .

* Expected "amount of capital" after n rounds
 $= E[S_n] = 0 \quad 0 \leq n \leq N$

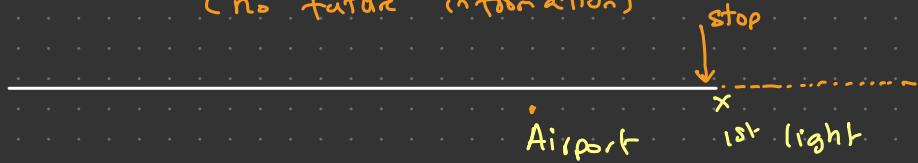
Question: Is it possible to stop the game in a favorable moment?

(clever stopping strategy \leadsto to a positive expected gain)

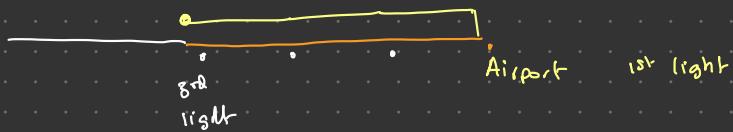
(no future). no inside trading { decision to stop may

Intuition :

- stop at the first light after the airport
(no future information)



- stop at the 3rd last light before the airport
(use future information)



Definition : An event $A \subseteq \Omega$ is observable until time n when it can be written as a union of basic events of the form

$$\{ \omega \in \Omega \mid \omega_1 = o_1, \dots, \omega_n = o_n \} \quad o_i \in d-1, i \in \{1, \dots, n\}$$

i.e. A - can be determined from the outcome of the 1st n tosses.

$A_n :=$ class of event A that can be observed by time n . [include \emptyset]

Notation: $A \subseteq \Omega$ (indicator of A)

$$\mathbb{1}_A(\omega) = \begin{cases} 1 & \omega \in A \\ 0 & \omega \notin A \end{cases}$$

Definition: A map $T: \Omega \rightarrow \{0, 1, \dots, N\} \cup \{\infty\}$ is called a stopping time if

$$\{T=n\} = \{\omega \in \Omega \mid T(\omega) = n\} \in \mathcal{A}_n$$
$$0 \leq n \leq N$$

i.e. $\{T=n\}$ is an event observable until time n .

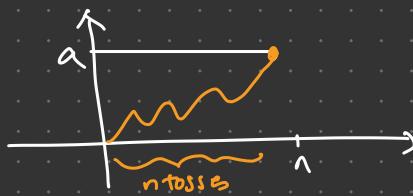
Example: For $a \in \mathbb{Z}$ let

$$\sigma_a: \Omega \rightarrow \{0, 1, \dots, N\} \cup \{\infty\}$$

$$\sigma_a(\omega) = \min \{k \in [0, N] \mid S_k(\omega) = a\}$$
$$\min(\emptyset) = \infty$$

Ex:- $\sigma_a(\cdot)$ is a stopping time

$\{n \in \mathbb{N} : \sigma_a(n) = n\} = \dots$ "only depends on last n tosses"



Theorem 1 :- For any stopping time

[Impossibility of a favorite stopping] $T: \Omega \rightarrow \{0, 1, \dots, N\}$

$$E[S_T] = 0$$

$S_T(\omega) \equiv S_{T(\omega)} \equiv$ outcome of the
trajectories ω at
stopping time $T(\omega)$

Proof :- T is a stopping time.

$$k=0, 1, 2, \dots, N \quad \{T \geq k\} \in A_{k-1} \quad \text{--- (1)}$$

$\left[\because \{T \geq k\} = \overline{\underbrace{\{T=1\} \cup \dots \cup \{T=k-1\}}^{\in A_k} \in A_k} \right]$
 $\in A_{k-1} \in A_k$ is
closed under complement

$$S_T = \sum_{k=1}^N X_k \mathbf{1}_{\{T \geq k\}} \quad \text{--- (2)}$$

$$\left[\because S_T = \sum_{k=1}^N S_k \mathbf{1}_{\{T \geq k\}}$$

$$= \sum_{k=1}^{N-1} S_k [\mathbf{1}_{\{T \geq k\}} - \mathbf{1}_{\{T \geq k+1\}}]$$

$$+ S_N \mathbf{1}_{\{T=N\}}$$

$$= \sum_{k=1}^{N-1} S_k \mathbf{1}_{\{T \geq k\}} - \sum_{k=1}^{N-1} S_k \mathbf{1}_{\{T \geq k+1\}} + S_N \mathbf{1}_{\{T=N\}}$$

$$= S_1 \mathbb{1}_{T \geq 1} + \sum_{k=2}^{n-1} S_k \mathbb{1}_{(T \geq k)} + S_n \mathbb{1}_{T=n}$$

$$- \sum_{k=1}^{n-1} S_k \mathbb{1}_{(T \geq k+1)}$$

$$= S_1 \mathbb{1}_{T \geq 1} + \sum_{m=1}^{n-1} S_{m+1} \mathbb{1}_{(T \geq m+1)} + S_n \mathbb{1}_{T=n}$$

$$- \sum_{k=1}^{n-1} S_k \mathbb{1}_{(T \geq k+1)}$$

$$= S_1 \mathbb{1}_{T \geq 1} + \sum_{l=1}^{n-1} (S_{l+1} - S_l) \mathbb{1}_{T \geq l+1} + (S_n - S_{n-1}) \mathbb{1}_{T=n}$$

\downarrow \downarrow \downarrow
 $\equiv \{T \geq n\}$ $x_l \mathbb{1}_{T \geq 1} + \sum_{l=1}^{n-1} x_l \mathbb{1}_{T \geq l+1} + x_n \mathbb{1}_{T=n}$

$$\mathbb{E}[S_T] \stackrel{(2)}{=} \mathbb{E}\left[\sum_{k=1}^n x_k \mathbb{1}_{T \geq k}\right]$$

linearity of \mathbb{E}

$$\mathbb{E}\left[\sum_{k=1}^n x_k \mathbb{1}_{T \geq k}\right] = \sum_{k=1}^n \mathbb{E}[x_k \mathbb{1}_{T \geq k}]$$

\downarrow $\underbrace{\mathbb{E} x_k}_{\text{depends on the } k\text{-th toss}} \downarrow$
 w_k " w_1, \dots, w_{n-1} "
 \downarrow

independence $\leftarrow \mathbb{E} x_k$

$$\sum_{k=1}^n \mathbb{E}[x_k] \mathbb{E}[\mathbb{1}_{(T \geq k)}]$$

\downarrow

$$= \sum_{k=1}^n \textcircled{O} \Pr(T > s_k) = 0$$

□

Game System: $v_1, \dots, v_n: \Omega_N \rightarrow \mathbb{R}$

are random variables such that

$$\{v_k = c\} \in A_{k-1} \quad c \in \mathbb{R}, \quad k=1, 2, \dots, N$$

$v_k \equiv$ amount of money you will place as a bet
 in the k^{th} round
 . Result in the k^{th} round $\equiv v_k X_k$

$$S_n^V = \sum_{k=1}^n v_k X_k$$

Theorem 2: For any v_1, v_2, \dots, v_n game system
 Then $E[S_n^V] = 0$

Proof: like before

$$E[S_n^V] = \sum_{k=1}^n E[v_k X_k] - \textcircled{#}$$

It is enough to show $E[v_k X_k] = 0$.

A_{k-1}

depends on
 k^{th} too

$\text{Range}_c(v_k) = \{c_1, c_2, \dots, c_m\}$

$$\Rightarrow X_{kc} v_k = \sum_{j=1}^m c_j X_{kc} \mathbb{1}_{(v_{jc} = c_j)}$$

$$E[X_{kc} v_k] = \sum_{j=1}^m c_j E[X_{kc} \mathbb{1}_{(v_{jc} = c_j)}]$$

\swarrow "w_k" \searrow "c₁, .. c_m"

$$= \sum_{j=1}^m c_j E[X_{kc}] P(v_{jc} = c_j)$$

$$\downarrow$$

$$= \dots 0 \dots = 0$$

↪ ⊗

Place ⊗ into # \Rightarrow to get result □

Impact of Theorem 2:

T: Ω → {0, 1, .., n} is a stopping time

$$\cdot V_k = \mathbb{1}_{\{T \geq k\}}$$

$$S_N = \sum_{k=1}^n X_k V_k = \sum_{k=1}^n X_k \mathbb{1}_{\{T \geq k\}}$$

$$\stackrel{\textcircled{2}}{=} S_T$$

[Theorem 2
 \Rightarrow Theorem 1]

$$v_k = s_{k-1} \underbrace{1}_{\textcircled{1}} \underbrace{\tau_{\geq k}}_{(w_1, \dots, w_{k-1})} \quad (\text{take this definition})$$

$\Rightarrow v_1, v_2, \dots, v_n$ is a game system.

$$v_k x_k = s_{k-1} x_k \underbrace{1}_{\tau_{\geq k}}$$

$$\begin{aligned} \text{Observe: } s_k^2 &= (s_{k-1} + x_k)^2 \\ &= s_{k-1}^2 + 2 \underbrace{s_{k-1} x_k}_{\frac{s_k^2 - s_{k-1}^2 - x_k^2}{2}} + x_k^2 \\ \Rightarrow v_k x_k &= \left(\frac{s_k^2 - s_{k-1}^2 - x_k^2}{2} \right) \underbrace{1}_{\tau_{\geq k}} \end{aligned}$$

$$\begin{aligned} s_n &= \sum_{k=1}^n v_k x_k \\ &= \sum_{k=1}^n \left(\frac{s_k^2 - s_{k-1}^2 - x_k^2}{2} \right) \underbrace{1}_{\tau_{\geq k}} \\ (x_{1c}^2 = 1) &= \sum_{k=1}^n \frac{s_k^2}{2} \underbrace{1}_{\tau_{\geq k}} - \sum_{k=1}^n \frac{s_{k-1}^2}{2} \underbrace{1}_{\tau_{\geq k}} \\ &\quad - \sum_{k=1}^n \frac{\underbrace{1}_{\tau_{\geq k}}}{2} \\ &= \sum_{k=1}^n \frac{s_k^2}{2} \underbrace{1}_{\tau_{\geq k}} - \frac{T}{2} \end{aligned}$$

$$S_n^V = \frac{(S_T^2 - T)}{2} - \textcircled{XXX}$$

Theorem 2 $\Rightarrow E[S_n^V] = 0$
 $\therefore B_2 \textcircled{XXX}$ for any $T: \Omega \rightarrow \{0, 1, \dots, n\}$

we have

$$0 = E[S_n^V] = E\left[\frac{S_T^2 - T}{2}\right]$$

$$\Rightarrow E[S_T^2] = E[T] - \textcircled{\#}$$

Theorem 1: $E[S_T] = 0$

$$\textcircled{\#} \quad \text{Var}[S_T] = E[T]$$

1.3 Ruin Problem:

Two players. Each with capital $\begin{cases} a - I \\ b - II \end{cases}$

Interpret

$S_n \equiv$ gain of player 1

Ruin (of player 1)

$$\{ \sigma_a < \sigma_b ; \sigma_a \leq n \}$$

where, $\sigma_u = \min \{ k : S_k = u \}$

Define: Ruin Probability

$$\gamma'_N = \mathbb{P}(\sigma_{-a} \leq \sigma_b, \sigma_{-a} \leq N)$$

$$T_N = \min (\sigma_{-a}, \sigma_b, N)$$

Ex: T_N is a stopping time

Theorem 1 $\Rightarrow E[S_{T_N}] = 0 \quad -\textcircled{3}$

$$T_N = \sigma_{-a} \Rightarrow S_{T_N} = -a$$

$$T_N = \sigma_b \Rightarrow S_{T_N} = b$$

$$T_N = N \Rightarrow S_{T_N} = S_N$$

- Use this in $\textcircled{3}$

$$0 = -a \mathbb{P}(T_N = \sigma_{-a}) + b \mathbb{P}(T_N = \sigma_b) + E[S_N | T_N = N]$$

As $T_N = \min (\sigma_{-a}, \sigma_b, N)$

$$0 = -a \mathbb{P}(\sigma_{-a} < \sigma_b \text{ & } \sigma_{-a} \leq N)$$

$$+ b \mathbb{P}(\sigma_b < \sigma_a \text{ and } \sigma_b \leq n) \\ + \mathbb{E}[S_n | \min(\sigma_a, \sigma_b) \geq n]$$

$$\Rightarrow O = -a \mathbb{P}_n(\text{Ruin of Player 1})$$

$$+ b \mathbb{P}_n(\text{Ruin of Player 2})$$

$$+ \mathbb{E}[S_n | \min(\sigma_a, \sigma_b) \geq n]$$

(4)

$$\gamma^1 = \lim_{n \rightarrow \infty} \mathbb{P}_N(\text{Ruin of Player 1}) \quad \left. \right\} \text{Ex-} \underline{\text{stinction}}$$

$$\gamma^2 = \lim_{n \rightarrow \infty} \mathbb{P}_N(\text{Ruin of Player 2})$$

$$(4) \quad \text{Ex} \quad O = -a \gamma^1 + b \gamma^2 + 0$$

$\lim_{n \rightarrow \infty}$

$$\Rightarrow -a \gamma^1 + b \gamma^2 \rightarrow - \quad (5)$$

$$\text{Ex} \quad \gamma^1 + \gamma^2 = 1 \quad \rightarrow (6)$$

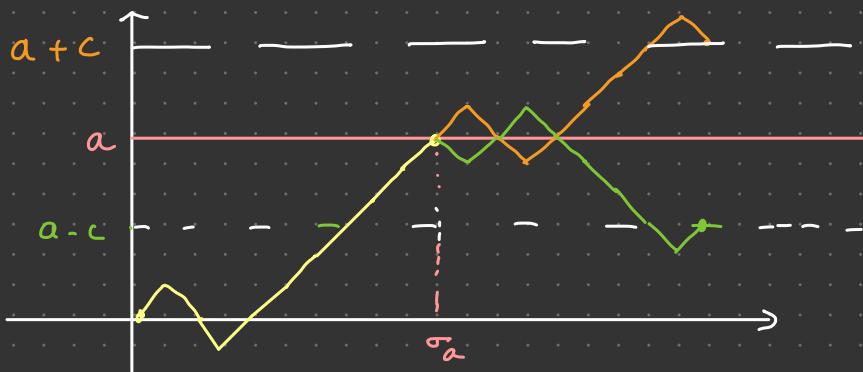
$$(5) \Leftarrow (6)$$

$$\gamma^1 = \frac{b}{a+b} \quad \gamma^2 = \frac{a}{a+b} !$$

□

1.4 Reflection Principle

Pictorial Concept : $a \in \mathbb{N}$ $c \in \mathbb{N} \cup \{0\}$



Ex: For every path of the random walk after σ_a :

Each orange path that lands at $a+c$

There is a "Reflected" path across $y=a$

green path that lands at $a-c$

$$|\{w \in \Omega : S_n = a+c\}|$$

$$= |\{w \in \Omega : \sigma_a \leq n, S_n = a+c\}|$$

$$= |\{w \in \Omega : \sigma_a \leq n, S_n = a-c\}|$$

Lemma 1 : $a, c \in \mathbb{N} \cup \{0\}$

$$\mathbb{P}(S_n = a+c) = \mathbb{P}(\sigma_a \leq n, S_n = a-c)$$

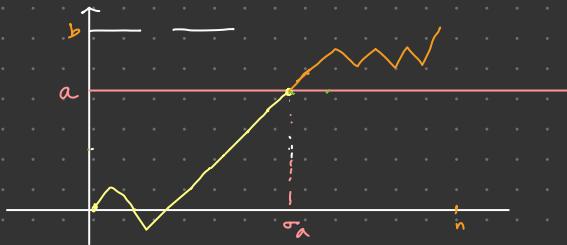
Theorem 2

$$\textcircled{a} \quad \mathbb{P}(\sigma_a \leq n) = \mathbb{P}(S_n \notin [-a, a-1])$$

$$\textcircled{b} \quad \mathbb{P}(\sigma_a = n) = \frac{1}{2} [\mathbb{P}(S_n = a-1) - \mathbb{P}(S_{n+1} = a+1)]$$

Proof :-

$$\mathbb{P}(\sigma_a \leq n) = \mathbb{P}(\sigma_a \leq n, \bigcup_{b \in \mathbb{Z}} \{S_n = b\})$$



$$\text{i.e. } \mathbb{P}(\sigma_a \leq n) = \sum_{b \in \mathbb{Z}} \mathbb{P}(\sigma_a \leq n, S_n = b)$$

$$= \sum_{\substack{b \in \mathbb{Z} \\ b \geq a}} \mathbb{P}(\sigma_a \leq n, S_n = b) + \sum_{\substack{b \in \mathbb{Z} \\ b < a}} \mathbb{P}(\sigma_a \leq n, S_n = b)$$

Apply Lemma 1.3.1
 $b = a - c$
 $(\Rightarrow a+c = 2a-b)$

Reflection Principle

$$= \sum_{\substack{b \in \mathbb{Z} \\ b \geq a}} \mathbb{P}(\sigma_a \leq n, S_n = b) + \sum_{\substack{b \in \mathbb{Z} \\ b < a}} \mathbb{P}(S_n = 2a-b)$$

(Ex)

$$= \sum_{\substack{b \in \mathbb{Z} \\ b \geq a}} \mathbb{P}(S_n = b) + \mathbb{P}(S_n > a)$$

$\{\sigma_a \leq n\} \subseteq \{S_n = b\}$

$$= \mathbb{P}(S_n \geq a) + \mathbb{P}(S_n > a)$$

$$\stackrel{\text{(Symmetry)}}{=} \mathbb{P}(S_n \geq a) + \mathbb{P}(S_n < -a)$$

$$= \mathbb{P}(S_n \notin [-a, a])$$

$$\textcircled{b} \quad \mathbb{P}(S_a = n) = \mathbb{P}(S_a \leq n) - \mathbb{P}(S_a \leq n-1)$$

$$\stackrel{\text{Ex}}{=} \frac{1}{2} [\mathbb{P}(S_{n-1} = a-1) - \mathbb{P}(S_{n-1} = a+1)]$$

(Apply \textcircled{a})
(\oplus)
(Symmetry)

D

Corollary 3 :- $a, n \in \mathbb{N}$

$$\mathbb{P}(S_a = n) = \frac{a}{n} \mathbb{P}(S_n = a)$$

Proof: Earlier class:

$$\mathbb{P}(S_{n-1} = a-1) = \frac{1}{2^{n-1}} \binom{n-1}{\frac{(n-1)+(a-1)}{2}}$$

$$\mathbb{P}(S_{n-1} = a+1) = \frac{1}{2^{n-1}} \binom{n-1}{\frac{n-1+a+1}{2}}$$

$$\mathbb{P}(S_n = a) = \frac{1}{2^n} \binom{n}{\frac{n+a}{2}}$$

$$\left(\begin{array}{l} \text{Do the} \\ \text{Combinatorics} \end{array} \right) \quad \binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}$$

$$\binom{n-1}{k} = \underbrace{\frac{n-1}{k}}_{k} \binom{n}{k}$$

$$\text{Thm 2} \quad \Rightarrow \quad \mathbb{P}(\sigma_a = n) = \frac{\alpha}{n} \mathbb{P}(S_n = a) \quad \square$$

$$\forall n \in \mathbb{N} \quad \sigma_a = \min \{n \geq 1 \mid S_n = a\}$$

(Chitting time of a) $a \in \mathbb{Z}/\{0\}$

$$(\text{Return time}) \quad \sigma_0 = \min \{n \geq 1 \mid S_n = 0\}$$

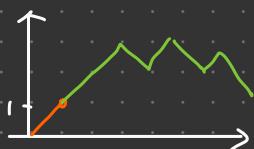
Lemma 4: (Escape from origin)

$$\mathbb{P}(\sigma_0 > 2n) = \mathbb{P}(S_{2n} = 0)$$

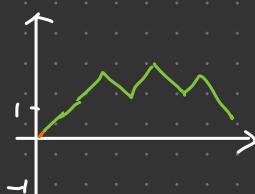
Proof:-

$$\mathbb{P}(\sigma_0 > 2n) = \mathbb{P}(S_1 \neq 0, S_2 \neq 0, \dots, S_{2n} \neq 0)$$

$$\xleftarrow[\text{(Reflection Principle)}]{\text{Eq}} = 2 \mathbb{P}(S_1 > 0, S_2 > 0, \dots, S_{2n} > 0)$$



$$= 2 \cdot \frac{\# \text{ paths start at } 1 \text{ & stay positive for } 2n-1 \text{ steps}}{2^{2n}}$$



$$= \frac{1}{2^{2n-1}} \cdot \frac{\# \text{ paths start at } 0 \text{ & stay above } -1 \text{ for } 2n-1 \text{ steps}}{2^{2n}}$$

(Reflection Principle)

$$\xleftarrow[\text{(Reflection Principle)}]{\text{Eq}} = \frac{1}{2^{2n-1}} \cdot \frac{\# \text{ paths start at } 0 \text{ & stay below } 1 \text{ for } 2n-1 \text{ steps}}{2^{2n}}$$

$$= \mathbb{P}(\sigma_1 > 2^{n-1})$$

$$= 1 - \mathbb{P}(\sigma_1 \leq 2^{n-1})$$

Theorem 2 $\hat{\Sigma} = 1 - \mathbb{P}(S_{2^{n-1}} \notin [-1, 0])$

with $a=1$

$$\stackrel{\text{Ex}}{=} \mathbb{P}(S_{2^{n-1}} = -1) = \mathbb{P}(S_{2^n} = 0) \quad \square$$