

Report on

"Compiler Design Mini Project - C++"

Submitted in partial fulfillment of the requirements for Sem VI

Compiler Design Laboratory

Bachelor of Technology in Computer Science & Engineering

Submitted by:

Karan Kumar G PES1201801883 Sparsh Temani PES1201800284 Sumukh Aithal K PES1201801461

Under the guidance of

Madhura V

Assistant Professor PES University, Bengaluru

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DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING FACULTY OF ENGINEERING PES UNIVERSITY

(Established under Karnataka Act No. 16 of 2013) 100ft Ring Road, Bengaluru – 560 085, Karnataka, India

TABLE OF CONTENTS

Chapter No.	Title	Page No.
1.	INTRODUCTION (Mini-Compiler is built for which language. Provide sample input and output of your project)	03
2.	ARCHITECTURE OF LANGUAGE: • What all have you handled in terms of syntax and semantics for the chosen language.	27
3.	LITERATURE SURVEY (if any paper referred or link used)	28
4.	CONTEXT FREE GRAMMAR (which you used to implement your project)	29
5.	 DESIGN STRATEGY (used to implement the following) SYMBOL TABLE CREATION INTERMEDIATE CODE GENERATION CODE OPTIMIZATION ERROR HANDLING - strategies and solutions used in your Mini-Compiler implementation (in its scanner, parser, semantic analyzer, and code generator). 	46
6.	 IMPLEMENTATION DETAILS (TOOL AND DATA STRUCTURES USED in order to implement the following): SYMBOL TABLE CREATION INTERMEDIATE CODE GENERATION CODE OPTIMIZATION ERROR HANDLING - strategies and solutions used in your Mini-Compiler implementation (in its scanner, parser, semantic analyzer, and code generator). Provide instructions on how to build and run your program. 	49
7.	RESULTS AND possible shortcomings of your Mini-Compiler	51
8.	SNAPSHOTS (of different outputs)	51
9.	CONCLUSIONS	66
10.	FURTHER ENHANCEMENTS	66
REFERENC	ES/BIBLIOGRAPHY	67

1. Introduction

As part of the Compiler Design Project component, we have built a mini-compiler for C++ language. The compiler aims to cover the basic syntax, grammar, intermediate code generation and some optimisation techniques.

Sample input:

A complete C++ program consisting of different constructs like if, if-else and for along with basic syntax:

```
#include<iostream>
using namespace std;
/* Note this is a comment to check the correctness of the comment removal code in
Lex
  This comment spans multiple lines.
// This is a single comment line.....
int main()
  // Variable declaration check
  int a = 5;
  int b = 6;
  int c = a * b;
  int d = c+b;
  int e = a*b+d-b+a:
  int f = a*5 + 8-4 + 7*3 + d*e/1;
  int g = (4+5)*f+e-d/c+b;
  int k = c + 4;
  g = d+4;
  int sum = 0;
  // Simple for loop
  for(int i=0; i<10; i++)
     sum = sum + 1;
```

```
}
// Double nested for loop
for(int i=0;i<10;i++)
  for(int j=i; j<10; j++)
     sum = sum + 1;
// Simple if condition
if(a>b)
  g = c + 4;
// Simple if condition
if(sum==55)
  int k = 2;
// Multi-Nested if-else
if(a+b==c)
  if(b+c>=f)
     k = 2;
  else
     if(a-b>0)
       k = 5;
     else
       k = 4;
     g = 3;
```

```
b = g + 1;
  e = f - 1;
else if(b+c==d)
  a = g - 1;
  b = f*2;
else
  a = g-2;
int x = 5;
x+=6;
x-=a;
int mx = 4+5-3;
mx = -mx;
mx = +5;
// constant expression as if condition
if(1+2)
  k = 1;
// if-else inside a for loop
for(int j=1; j<10; j++)
  if(a+b>c)
    k = 1+2;
  if(1)
     x = 1;
  else
    x = -1;
// Different type of if condition
if((a+b>x)&&(c+d>b))
```

```
a = 100;
// Nested for with if inside
for(int q=10;q>0;--q)
  for(int p=0;p<5;++p)
     sum+=p+q;
     if(p+q>10)
       a = a - 1;
// Simple if construct
if(a)
  a = 20;
if((a+b>c))
  for(int z=0;z<6;z=z+5)
     if(z)
       a=1;
     else
       a+=1;
// Special cases of for loops
for(;a>10;)
  a = 1;
```

```
for(;;)
{
    sum-=1;
}

for(a=0;;a+=1)
{
    sum+=1;
}
```

Sample output:

i. Symbol table with required information

YMBOL TABLE	MBOL TABLE					
oken	Category	Туре	Line Number	Scope	Value String	
	Identifier	int	13	1	t82	
	Identifier	int	14	1	6	
	Identifier	int	15	1	t0	
	Identifier	int	16	1	t1	
	Identifier	int	17	1	t5	
	Identifier	int	18	1	t13	
	Identifier	int	19	1	t19	
	Identifier	int	25	2	0	
	Identifier	int	31	3	0	
	Identifier	int	33	4	Ō	
	Identifier	int	96	16	i	
	Identifier	int	20	1	t20	
	Identifier	int	48	6	2	
	Identifier	int	117	22	0	
	Identifier	int	115	21	10	
	Identifier	int	83	1	t47	
	Identifier	int	135	26	0	
9	temporary		0	1	a * b	
ĺ	temporary		Ö	ī	c + b	
2	temporary		0	ī	a * b	
- 3	temporary		0	ī	t2 + d	
í	temporary		0	ī	t3 - b	
5	temporary		0	ī	t4 + a	
5	temporary		0	1	a * 5	
7	temporary		0	1	t6 + 8	
, B	temporary		0	1	t7 - 4	
9	temporary		ö	1	7 * 3	
10	temporary		0	1	t8 + t9	
11	temporary		0	1	d * e	
20	temporary		Θ	i	c + 4	
12	temporary		9	1	t11 / 1	
21			8	1	d + 4	
30	temporary		8	1	u + 4 a > b	
13	temporary		9	1	t10 + t12	
13 22	temporary		0	2	i < 10	
22 31	temporary			5	1 < 10 c + 4	
31 40	temporary temporary		0 0	5 7	c + 4 f - 1	

t14	temporary		Θ	1	4 + 5
t32	temporary		ō	ī	sum == 55
t41	temporary		0	12	b + c
t50	temporary		0	1	- mx
t15			0	1	- IIIX t14 * f
	temporary				
t24	temporary		0	2	sum + 1
t33	temporary		0	1	a + b
t42	temporary		0	12	t41 == d
t51	temporary		0	1	+ 5
t60	temporary		0	1	t59 > x
t16	temporary		Θ	1	t15 + e
t25	temporary		ō	3	i < 10
t34	temporary		0	1	t33 == c
t43			0	13	
	temporary				g - 1
t52	temporary		0	1	1 + 2
t61	temporary		0	1	c + d
t70	temporary		0	22	p + q
t17	temporary		0	1	d / c
t35	temporary		0	7	b + c
t44	temporary		0	13	f * 2
t53	temporary		0	16	j < 10
t62			0	1	
	temporary				t61 > b
t71	temporary		0	22	t70 > 10
t80	temporary		0	29	a - 1
t18	temporary		0	1	t16 - t17
t27	temporary		0	4	j < 10
t36	temporary		0	7	t35 >= f
t45	temporary		Θ	14	g - 2
t63	temporary		0	i	t60 && t62
t72	temporary		0	23	a - 1
t81			0	30	a - 1 sum - 1
	temporary				
t19	temporary		0	1	t18 + b
t37	temporary		0	9	a - b
t46	temporary		0	1	x + 6
t55	temporary		0	16	a + b
t64	temporary		0	21	q > 0
t73	temporary		0	1	a + b
t82	temporary		0	31	a + 1
t29	temporary		0	4	sum + 1
			0	9	
t38	temporary				t37 > 0
t47	temporary		0	1	x - a
t56	temporary		0	16	t55 > c
t74	temporary		0	1	t73 > c
			-		
t83	temporary		0	31	sum + 1
t39	temporary		0	7 1	g + 1
t48	temporary		0	1	4 + 5
t57	temporary		0	17	1 + 2
t66	temporary		0	22	p < 5
t75	temporary		Ö	26	z < 6
t49			0	1	t48 - 3
	temporary				
t58	temporary		0	19	- 1
t76	temporary		0	26	z + 5
t59	temporary		0	1	a + b
t68	temporary		0	22	p + q
	*		0	27	a - 1
t77	temporary		0	22	sum + t68
	temporary temporary				
t77 t69	temporary				a + 1
t77 t69 t78	temporary temporary		0	28	a + 1
t77 t69 t78 t79	temporary temporary temporary	int	0 0	28 29	a > 10
t77 t69 t78 t79 mx	temporary temporary temporary Identifier	int	0 0 87	28 29 1	a > 10 t49
t77 t69 t78 t79 mx sum	temporary temporary temporary Identifier Identifier	int	0 0 87 23	28 29 1 1	a > 10 t49 t83
t77 t69 t78 t79 mx	temporary temporary temporary Identifier		0 0 87	28 29 1	a > 10 t49
t77 t69 t78 t79 mx sum	temporary temporary temporary Identifier Identifier	int	0 0 87 23	28 29 1 1	a > 10 t49 t83

ii. Intermediate code (both in Three Address Code and Quadruple format)

Three Address Code:

$$a = 5$$

 $b = 6$
 $t0 = a * b$
 $c = t0$
 $t1 = c + b$

d = t1t2 = a * bt3 = t2 + dt4 = t3 - bt5 = t4 + ae = t5t6 = a * 5t7 = t6 + 8t8 = t7 - 4t9 = 7 * 3t10 = t8 + t9t11 = d * et12 = t11 / 1t13 = t10 + t12f = t13t14 = 4 + 5t15 = t14 * ft16 = t15 + et17 = d/ct18 = t16 - t17t19 = t18 + bg = t19t20 = c + 4k = t20t21 = d + 4g = t21sum = 0i = 0L0: t22 = i < 10if t22 goto L1 goto L2 L3: t23 = i + 1i = t23goto L0 L1: t24 = sum + 1sum = t24goto L3 L2:

```
i = 0
L4:
t25 = i < 10
if t25 goto L5
goto L6
L7:
t26 = i + 1
i = t26
goto L4
L5:
j = i
L8:
t27 = j < 10
if t27 goto L9
goto L10
L11:
t28 = j + 1
j = t28
goto L8
L9:
t29 = sum + 1
sum = t29
goto L11
L10:
goto L7
L6:
t30 = a > b
if t30 goto L12
goto L13
L12:
t31 = c + 4
g = t31
goto L14
L13:
L14:
t32 = sum == 55
if t32 goto L15
goto L16
L15:
k = 2
goto L17
```

L16: L17: t33 = a + bt34 = t33 == cif t34 goto L18 goto L19 L18: t35 = b + ct36 = t35 >= fif t36 goto L20 goto L21 L20: k = 2goto L22 L21: t37 = a - bt38 = t37 > 0if t38 goto L23 goto L24 L23: k = 5goto L25 L24: k = 4L25: g = 3L22: t39 = g + 1b = t39t40 = f - 1e = t40goto L26 L19: t41 = b + ct42 = t41 == dif t42 goto L27 goto L28 L27: t43 = g - 1a = t43t44 = f * 2

```
b = t44
goto L29
L28:
t45 = g - 2
a = t45
L29:
L26:
x = 5
t46 = x + 6
x = t46
t47 = x - a
x = t47
t48 = 4 + 5
t49 = t48 - 3
mx = t49
t50 = -mx
mx = t50
t51 = +5
mx = t51
t52 = 1 + 2
if t52 goto L30
goto L31
L30:
k = 1
goto L32
L31:
L32:
j = 1
L33:
t53 = j < 10
if t53 goto L34
goto L35
L36:
t54 = j + 1
j = t54
goto L33
L34:
t55 = a + b
t56 = t55 > c
```

if t56 goto L37

goto L38

```
L37:
t57 = 1 + 2
k = t57
goto L39
L38:
L39:
if 1 goto L40
goto L41
L40:
x = 1
goto L42
L41:
t58 = -1
x = t58
L42:
goto L36
L35:
t59 = a + b
t60 = t59 > x
t61 = c + d
t62 = t61 > b
t63 = t60 \&\& t62
if t63 goto L43
goto L44
L43:
a = 100
goto L45
L44:
L45:
q = 10
L46:
t64 = q > 0
if t64 goto L47
goto L48
L49:
t65 = q - 1
q = t65
goto L46
L47:
p = 0
```

L50:

t66 = p < 5if t66 goto L51 goto L52 L53: t67 = p + 1p = t67goto L50 L51: t68 = p + qt69 = sum + t68sum = t69t70 = p + qt71 = t70 > 10if t71 goto L54 goto L55 L54: t72 = a - 1a = t72goto L56 L55: L56: goto L53 L52: goto L49 L48: if a goto L57 goto L58 L57: a = 20goto L59 L58: L59: t73 = a + bt74 = t73 > cif t74 goto L60 goto L61 L60: z = 0L62: t75 = z < 6if t75 goto L63

goto L64 L65:

t76 = z + 5

z = t76

goto L62

L63:

if z goto L66

goto L67

L66:

t77 = a - 1

a = t77

goto L68

L67:

t78 = a + 1

a = t78

L68:

goto L65

L64:

goto L69

L61:

L69:

L70:

t79 = a > 10

if t79 goto L71

goto L72

L73:

goto L70

L71:

t80 = a - 1

a = t80

goto L73

L72:

L74:

goto L75

L77:

goto L74

L75:

t81 = sum - 1

sum = t81

goto L77

L76:

a = 0 L78: goto L79 L81: t82 = a + 1 a = t82 goto L78 L79: t83 = sum + 1 sum = t83 goto L81 L80:

Quadruple Format:

QUADRUPLES

Op	arg1	arg2	res
=	5		a
=	6		b
*	a	b	t0
=	t0		c
+	c	b	t1
=	t1		d
*	a	b	t2
+	t2	d	t3
-	t3	b	t4
+	t4	a	t5
=	t5		e
*	a	5	t6
+	t6	8	t7
-	t7	4	t8
*	7	3	t9
+	t8	t9	t10
*	d	e	t11
/	t11	1	t12
+	t10	t12	t13
=	t13		f
+	4	5	t14
*	t14	f	t15
+	t15	e	t16
/	d	c	t17

_	t16	t17	t18
+	t18	b	t19
=	t19		g
+	c	4	t20
=	t20		k
+ = =	d	4	t21
=	t21		g
	0		sum
=	0		i
L0			label
<	i	10	t22
if	t22		L1
goto			L2
L3			label
+	i	1	t23
=	t23		i
goto			L0
L0			label
+	sum	1	t24
=	t24		sum
goto			L3
L2			label
=	0		i
L4			label
<	i	10	t25
if	t25		L5
goto			L6
L7			label
+	i	1	t26
=	t26		i
goto			L4
L4			label
=	i		j
L8			label
<	j	10	t27
if	t27		L9
goto			L10
L11			label
+	j	1	t28
=	t28		j
goto			L8

L8			label
+	G11100	1	
	sum	1	t29
=	t29		sum
goto			L11
L10			label
goto			L7
L6			label
>	a	b	t30
L12	t30		if
goto	•••		L13
L12			label
+	c	4	t31
=	t31		g
goto			L14
L13			label
L14			label
==	sum	55	t32
L15	t32		if
goto			L16
L15			label
_	2		
	2		k
goto			L17
L16			label
L17			label
+	a	b	t33
===	t33	c	t34
		C	
L18	t34		if
goto			L19
L18			label
+	b	c	t35
>=	t35	f	t36
L20	t36	1	if
	130		
goto			L21
L20			label
=	2		k
goto			L22
L21			label
L/ L 1	0	h	
-	a .27	b	t37
>	t37	0	t38
L23	t38		if
goto			L24
0-10			

L23			label
= ,	5		k
goto L24			L25 label
L2 4 =	4		k
L25	•		label
=	3		g
L22			label
+	g	1	t39
=	t39		b
-	f	1	t40
=	t40		e
goto			L26
L19	1		label
+	b	C	t41
	t41 t42	d	t42 if
L27 goto	142		L28
L27			label
L2 /	σ	1	t43
=	g t43	1	a
*	f	2	t44
=	t44		b
goto			L29
L28			label
-	g	2	t45
=	t45		a
L29			label
L26			label
=	5		X
+	X	6	t46
=	X		t46
_	X	a	t47
= +	x 4	5	t47 t48
_	4 t48	3	t49
- =	t49	5	mx
t50	mx		-
=	t50		mx
t51	5		+
=	t51		mx

+	1	2	t52
L30	t52		if
	132		
goto			L31
L30			label
=	1		k
goto			L32
L31			label
L32	1		label
=	1		j
L33			label
<	j 452	10	t53
if	t53		L34
goto			L35
L36			label
	:	1	
+	j	1	t54
=	t54		j
goto			L33
L33			label
+	a	b	t55
>	t55	c	t56
L37	t56	·	if
	130		
goto			L38
L37			label
+	1	2	t57
=	t57		k
goto			L39
L38			label
L39			label
	1		
L40	1		if
goto			L41
L40			label
=	1		X
goto			L42
L41			label
t58	1		14001
=			
	t58		X
L42			label
goto			L36
L35			label
+	a	b	t59
>	t59	X	t60

+	c	d	t61
>	t61	b	t62
&&	t60	t62	t63
L43	t63		if
goto			L44
L43			label
=	100		a
goto			L45
L44			label
L45			label
=	10		q
L46	10		label
>	q	0	t64
if	t64	Ü	L47
goto			L48
L49			label
_	q	1	t65
=	t65	_	q
goto			L46
L46			label
=	0		p
L50			label
<	p	5	t66
if	t66		L51
goto			L52
L53			label
+	p	1	t67
=	t67		p
goto			L50
L50			label
+	p	q	t68
+	sum	t68	t69
=	sum		t69
+	p	q	t70
>	t70	10	t71
L54	t71		if
goto			L55
L54			label
-	a	1	t72
=	t72		a
goto			L56

L55 L56 goto L52 goto L48 L57 goto L57	a		label L53 label L49 label if L58 label
=	20		a 1.50
goto L58			L59 label
L59			label
+	a	b	t73
>	t73	c	t74
L60	t74		if
goto			L61
L60			label
=	0		Z
L62		(label
< if	z t75	6	t75
goto	1/3		L63 L64
L65			label
+	Z	5	t76
=	t76		Z
goto			L62
L62			label
L66	Z		if
goto			L67
L66		1	label
_	a	1	t77
goto	a		t77 L68
L67			label
+	a	1	t78
=	a		t78
L68			label
goto			L65
L64			label
goto			L69

L61 L69			label
L70 >	0	10	label t79
if	a t79	10	L71
goto	• • • • • • • • • • • • • • • • • • • •		L72
L73			label
goto			L70
L70			label
-	a	1	t80
=	a		t80
goto			L73
L72			label
L74			label
goto			L75
L77			label
goto			L74
L74			label
-	sum	1	t81
=	sum		t81
goto			L77
L76			label
=	0		a
L78			label
goto			L79
L81			label
+	a	1	t82
=	a		t82
goto			L78
L78			label
+	sum	1	t83
=	sum		t83
goto			L81
L80			label

iii. Optimized Intermediate code (in Quadruple format)

QUADRUPLES

op	arg1	arg2	result	
=	5		a	
=	6		b	
=	30		t0	
=	36		t1	
=	66		t3	
=	60		t4	
=	65		t5	
=	25		t6	
=	33		t7	
=	29		t8	
=	21		t9	
=	50		t10	
=	2340		t11	
=	2340.	0	t12	
=	2390.	0	t13	
=	9		t14	
=	21510	0.0		t15
=	21575	5.0		t16
=	1.2		t17	
=	21573	3.8		t18
=	21579	8.6		t19
=	34		t20	
=	40		t21	
L0			label	
=	True		t22	
if	t22		L1	
goto			L2	
L3			label	
L0			label	
L2			label	
=	0		t23	
L7			label	
L4			label	
L11			label	
L8			label	
L10			label	
L6			label	
=	False		t30	
L12	t30		if	

L12		label
=	34	t21
L13		label
L14		label
L15	t30	if
L15	130	label
L13		label
L17	1.1	label
=	11	t33
L18	t30	if
L18		label
L20	t30	if
L20		label
L21		label
=	-1	t37
L23	t30	if
L23		label
L24		label
=	4	a
L25		label
=	3	t21
L22		label
=	4	b
=	2389.0	t40
L19		label
L27	t30	if
L27		label
=	4780.0	t44
L28	1700.0	label
L29		label
L26		label
=	6	
t50	6	mx
=	t50	mv
		mx
t51	5	+
=	t51	mx :c
L30	3	if
L30		label
L31		label
L32		label
L36		label

L33	4781.0	label t55
L37	t22	if
L37	3	label t28
L38	5	label
L39		label
L40	1	if
L40		label
L41	1	label
t58	1	- +20
_ L42	t58	t28 label
L35		label
>	4781.0t28	t60
&&	t60 t30	t63
L43	t63	if
L43		label
=	100	t28
L44		label
L45 =	10	label
– L49	10	q label
=	9	q
L46		label
L53		label
=	1	t28
L50		label
=	12	t69
L54	t30	if
L54	00	label
_ L55	99	t72 label
L55		label
L52		label
L48		label
L57	99	if
L57		label
=	20	t72
L58		label
L59		label

=	4800.0	t73
L60	t22	if
L60		label
L65		label
=	5	t72
L62		label
L66	5	if
L66		label
=	19	t77
L67		label
L68		label
L64		label
L61		label
L69		label
L73		label
L70		label
L72		label
L77		label
L74		label
=	2	t72
L76		label
L81		label
L78		label
L80		label

2. ARCHITECTURE OF LANGUAGE

The designed mini compiler aims to handle basic working syntax of C++ programs along with specific constructs.

In terms of syntax, the following cases have been handled:

- Single line and multiline comments
- Incomplete multiline comments resulting in error message generation
- Recognition of multiple keywords like return, void, class, public, private, protected, int, float, double, bool, if, else, for, cin, cout, printf, scanf, break, continue, exit, string, char, true, false, etc.
- Recognition of valid identifiers (limited to maximum 32 characters)

- Conversion of exponential notation floating point numbers (like 3.14E10) to standard decimal notation floating point numbers
- Preprocessor directives
- Functions (with prototype, declaration and definition)
- Single-line if construct
- Block if constructs
- Block if else construct
- Single line for construct
- Block for construct
- All arithmetic operators (+, -, *, /, %)
- All bitwise operators (&, |, ^)
- All logical and relational operators
- Multiple cases of Assignment expressions
- Jump statements

In terms of semantics, the following cases have been handled:

- Usage of undeclared variables
- Implicit type casting between primitive data types
- Incompatible type assignments
- Incompatible type operations
- Illegal redeclarations
- Occurence of break statements at only appropriate places (inside loop bodies)

3. LITERATURE SURVEY

- Lex and Yacc Anchor material provided by PES University Compiler Design Faculty
- Regular Expressions online website (https://www.regular-expressions.info/)
- ISO C++ reference (https://isocpp.org/wiki/faq/compiler-dependencies#yaccable-grammar)
- C++ yacc-able grammar reference (http://www.computing.surrey.ac.uk/research/dsrg/fog/CxxGrammar.y)
- Lecture slides on Intermediate Code generation and Optimisation techniques provided by PES University Compiler Design Faculty

4. CONTEXT FREE GRAMMAR

Following is the C++ context free grammar written by us in the yacc file:

```
START
     : INCLUDE BODY
     BODY
     | INCLUDE
INCLUDE
     : INCLUDE T HEADER_INCLUDE '<' T_HEADER_FILE '>'
     INCLUDE T HEADER INCLUDE T STRING LITERAL
     T HEADER INCLUDE '<' T HEADER FILE '>'
     T HEADER INCLUDE T STRING LITERAL
BODY
     : BODY BLOCK BODY
     | BODY BLOCK
BODY BLOCK
     : FUNCTION
     BLOCK
FUNCTION
     : FUNCTION_PROTOTYPE
     | FUNCTION DEFINITION
     | FUNCTION DECLARATION
FUNCTION PROTOTYPE
     : FUNCTION PREFIX TYPE LIST')' ';'
     | FUNCTION PREFIX ')' ';'
TYPE LIST
    : TYPE ',' TYPE LIST
```

```
TYPE
FUNCTION DEFINITION
      : FUNCTION PREFIX FUNCTION PARAMETER LIST ')' ';' {
            scope leave();
FUNCTION DECLARATION
      : FUNCTION PREFIX FUNCTION PARAMETER LIST')''{
STATEMENTS '}' {
            scope leave();
      | FUNCTION PREFIX ')' '{' STATEMENTS '}' {
            scope leave();
FUNCTION PARAMETER LIST
      : TYPE T IDENTIFIER ',' FUNCTION PARAMETER LIST {
            if (insert($2, "Identifier", $1, @2.last line, NULL) == NULL) {
                  printf("[Error] at line:%d - Function Parameter \"%s\" has
already been declared\n", @2.last line, $2);
      TYPE T IDENTIFIER '=' EXPRESSION ','
FUNCTION PARAMETER LIST {
            if (insert($2, "Identifier", $1, @2.last line, NULL) == NULL) {
                  printf("[Error] at line:%d - Function Parameter \"%s\" has
already been declared\n", @2.last line, $2);
    symbol table* element = lookup($2);
            strcpy(element->value, $4);
      | TYPE T | IDENTIFIER {
            if (insert($2, "Identifier", $1, @2.last line, NULL) == NULL) {
                  printf("[Error] at line:%d - Function Parameter \"%s\" has
already been declared\n", @2.last line, $2);
```

```
| TYPE T IDENTIFIER '=' EXPRESSION {
            if (insert($2, "Identifier", $1, @2.last line, NULL) == NULL) {
                  printf("[Error] at line:%d - Function Parameter \"%s\" has
already been declared\n", @2.last line, $2);
    symbol table* element = lookup($2);
            strcpy(element->value, $4);
FUNCTION PREFIX
      : TYPE T IDENTIFIER '(' {
            insert($2, "Function-Identifier", $1, @2.last line, NULL);
            scope enter();
BLOCK
      : BLOCK START STATEMENTS BLOCK END
BLOCK START
     :'{' {
    scope enter();
BLOCK END
     :'}' {
    scope leave();
STATEMENTS
      : STATEMENT STATEMENTS
      STATEMENT
SINGLE LINE IF
      : IF PREFIX LINE STATEMENT ';' {
            scope leave();
```

```
| IF_PREFIX ';' {
          scope leave();
     | IF PREFIX CONSTRUCT {
          scope leave();
      SINGLE LINE IF SINGLE LINE ELSE
      SINGLE LINE IF BLOCK ELSE
BLOCK IF
     : T CONSTRUCT IF '(' EXPRESSION ')' BLOCK
     BLOCK IF SINGLE LINE ELSE
     BLOCK IF BLOCK ELSE
IF PREFIX
     : T CONSTRUCT IF '('EXPRESSION')' {
          scope_enter();
SINGLE LINE ELSE
     : ELSE_PREFIX LINE_STATEMENT ';'{
          scope leave();
     | ELSE PREFIX ';'{
          scope leave();
     | ELSE PREFIX CONSTRUCT {
          scope_leave();
BLOCK ELSE
     : T CONSTRUCT ELSE BLOCK
ELSE PREFIX
     : T_CONSTRUCT_ELSE {
```

```
scope enter();
SINGLE LINE FOR
     : FOR PREFIX FOR INIT STATEMENT';'
FOR CONDITION STATEMENT'; FOR ACTION STATEMENT')'
LINE STATEMENT ';'{
           scope leave();
           is in construct = 1;
     FOR PREFIX FOR INIT STATEMENT ';'
FOR CONDITION STATEMENT ';' FOR ACTION STATEMENT ')' ';' {
           scope_leave();
           is in construct -= 1;
     FOR PREFIX FOR INIT STATEMENT';'
FOR_CONDITION_STATEMENT';' FOR_ACTION_STATEMENT')'
CONSTRUCT {
           scope leave();
           is_in_construct -= 1;
BLOCK FOR
     : FOR PREFIX FOR INIT STATEMENT';'
FOR CONDITION STATEMENT ';' FOR ACTION STATEMENT ')' '{'
STATEMENTS '}'{
           scope leave();
           is in construct -= 1;
FOR PREFIX
     : T CONSTRUCT FOR '(' {
           scope_enter();
           is in construct += 1;
FOR INIT STATEMENT
```

```
LINE STATEMENT
FOR CONDITION STATEMENT
     CONDITIONAL EXPRESSION
FOR ACTION STATEMENT
     | LINE STATEMENT
BITWISE OPERATOR
     : '&' {
   $\$ = strdup(\$1);
     $\$ = strdup(\$1);
    | '^' {
   $$ = strdup($1);
CONDITIONAL EXPRESSION
     : EXPRESSION LOGICAL_OPERATOR EXPRESSION_GRAMMAR {
   sprintf($$, "%s %s %s", $1, $2, $3);
   $ = strdup($);
     | EXPRESSION RELATIONAL OPERATOR EXPRESSION GRAMMAR {
   sprintf($$, "%s %s %s", $1, $2, $3);
   \$\$ = strdup(\$\$);
     EXPRESSION BITWISE OPERATOR EXPRESSION GRAMMAR {
   sprintf($$, "%s %s %s", $1, $2, $3);
   $ = strdup($);
```

```
ASSIGNMENT
      : T_IDENTIFIER ASSIGNMENT_OPERATOR EXPRESSION_GRAMMAR
            if (lookup(\$1) == NULL) 
                  printf("[Error] at line:%d - Undeclared Variable \"%s\" \n",
@1.last line, $1);
    sprintf($$, "%s %s %s", $1, $2, $3);
    $ = strdup($);
      T IDENTIFIER ASSIGNMENT OPERATOR ASSIGNMENT {
            if(lookup(\$1) == NULL) {
                  printf("[Error] at line:%d - Undeclared Variable \"%s\" \n",
@1.last line, $1);
    sprintf($$, "%s %s %s", $1, $2, $3);
    $ = strdup($);
      T IDENTIFIER '[' EXPRESSION ']' ASSIGNMENT OPERATOR
EXPRESSION GRAMMAR {
            if(lookup(\$1) == NULL) {
                  printf("[Error] at line:%d - Undeclared Variable \"%s\" \n",
@1.last line, $1);
    sprintf($$, "%s [ %s ] %s %s", $1, $3, $5, $6);
    $ = strdup($);
      T IDENTIFIER '[' EXPRESSION ']' ASSIGNMENT OPERATOR
ASSIGNMENT {
            if(lookup(\$1) == NULL) {
                  printf("[Error] at line:%d - Undeclared Variable \"%s\" \n",
@1.last line, $1);
    sprintf($$, "%s [ %s ] %s %s", $1, $3, $5, $6);
    $ = strdup($);
ASSIGNMENT OPERATOR
     : '=' {
    $$ = strdup($1);
```

```
T OP ADD ASSIGNMENT {
   $$ = strdup($1);
     T_OP_SUBTRACT_ASSIGNMENT {
   $$ = strdup($1);
     T OP MULTIPLY ASSIGNMENT {
   $$ = strdup($1);
     T OP DIVIDE ASSIGNMENT {
   $$ = strdup($1);
     T OP MOD ASSIGNMENT {
   $$ = strdup($1);
EXPRESSION
     : ASSIGNMENT {
   $ = strdup($1);
 }
     | CONDITIONAL EXPRESSION {
   $$ = strdup($1);
     EXPRESSION GRAMMAR {
   $$ = strdup($1);
EXPRESSION GRAMMAR
     : EXPRESSION_GRAMMAR '+' EXPRESSION_TERM {
   sprintf($$, "%s + %s", $1, $3);
   $ = strdup($);
     | EXPRESSION GRAMMAR '-' EXPRESSION TERM {
   sprintf($$, "%s - %s", $1, $3);
   $ = strdup($);
     EXPRESSION_TERM {
   $$ = strdup($1);
```

```
EXPRESSION TERM
      : EXPRESSION TERM '*' EXPRESSION F {
    sprintf($$, "%s * %s", $1, $3);
    $ = strdup($);
     EXPRESSION_TERM '/' EXPRESSION_F {
    sprintf($$, "%s / %s", $1, $3);
    $ = strdup($);
     EXPRESSION TERM '%' EXPRESSION F {
    sprintf($$, "%s %% %s", $1, $3);
    $ = strdup($);
     | EXPRESSION_F {
    $$ = strdup($1);
     | '!' EXPRESSION F {
    sprintf($$, "! %s", $2);
    $ = strdup($);
EXPRESSION F
     : IDENTIFIER_OR_LITERAL {
    $$ = strdup($1);
  }
     '('EXPRESSION')' {
    sprintf($$, "( %s )", $2);
    $ = strdup($);
     '+' EXPRESSION F {
    sprintf($$, "+ %s", $2);
    $ = strdup($);
     '-' EXPRESSION F {
    sprintf($$, "- %s", $2);
    $ = strdup($);
```

```
CONSTRUCT
     : SINGLE_LINE_CONSTRUCT
     | BLOCK_CONSTRUCT
BLOCK CONSTRUCT
     : BLOCK FOR
     BLOCK IF %prec IF PREC
SINGLE LINE CONSTRUCT
     : SINGLE LINE FOR
     | SINGLE LINE IF %prec IF PREC
STATEMENT
     : LINE STATEMENT ';'
     | CONSTRUCT
     BLOCK
JUMP STATEMENT
    : T JUMP BREAK {
          if (is in construct == 0)
               printf("[Error] at line:%d - \"break\" statement not within loop or
switch\n", @1.last_line);
     T JUMP_EXIT
     | T_JUMP_CONTINUE
LINE STATEMENT
     : VARIABLE DECLARATION
     | EXPRESSION
     COUT
     CIN
     RETURN
     JUMP STATEMENT
```

```
VARIABLE DECLARATION
      : VARIABLE DECLARATION TYPE VARIABLE LIST {
           strcpy(variable declaration type, "\0");
    sprintf($$, "%s %s", $1, $2);
    $$ = strdup($$);
VARIABLE DECLARATION TYPE
     : TYPE {
           strcpy(variable declaration type, $1);
    $$ = strdup($1);
VARIABLE LIST
     : VARIABLE DECLARATION IDENTIFIER ',' VARIABLE LIST {
    sprintf($$, "%s, %s", $1, $3);
    $ = strdup($);
     | VARIABLE DECLARATION IDENTIFIER '=' EXPRESSION ','
VARIABLE LIST {
           symbol table* element = lookup($1);
           sprintf(element->value, "%s", $3);
    sprintf(\$\$, "\%s = \%s, \%s", \$1, \$3, \$5);
    $$ = strdup($$);
     | VARIABLE DECLARATION IDENTIFIER {
    $ = strdup($1);
     | VARIABLE DECLARATION IDENTIFIER '=' EXPRESSION {
           symbol table* element = lookup($1);
           strcpy(element->value, $3);
    sprintf(\$\$, "\%s = \%s", \$1, \$3);
    $ = strdup($);
     ARRAY VARIABLE DECLARATION IDENTIFIER WITH SIZE ','
VARIABLE LIST {
```

```
sprintf($$, "%s, %s", $1, $3);
    $ = strdup($);
      ARRAY VARIABLE DECLARATION IDENTIFIER '=' ARRAY LIST ','
VARIABLE LIST {
    sprintf(\$\$, "\%s = \%s, \%s", \$1, \$3, \$5);
    $$ = strdup($$);
      ARRAY VARIABLE DECLARATION IDENTIFIER WITH SIZE {
    $$ = strdup($1);
      | ARRAY VARIABLE DECLARATION IDENTIFIER '=' ARRAY LIST {
    sprintf(\$\$, "\%s = \%s", \$1, \$3);
    $ = strdup($);
VARIABLE DECLARATION IDENTIFIER
      : T IDENTIFIER {
            if (insert($1, "Identifier", variable declaration type, @1.last line,
NULL) == NULL) {
                  printf("[Error] at line:%d - \"%s\" has already been declared\n",
@1.last line, $1);
    $$ = strdup($1);
ARRAY VARIABLE DECLARATION IDENTIFIER
      : T IDENTIFIER '[' ']' {
            if (insert($1, "Identifier-Array", variable declaration type,
@1.last line, NULL) == NULL) {
                  printf("[Error] at line:%d - \"%s\" has already been declared\n",
@1.last line, $1);
    sprintf($$, "%s []", $1);
    $ = strdup($);
      T IDENTIFIER '[' EXPRESSION ']' {
            if (insert($1, "Identifier-Array", variable_declaration_type,
@1.last line, NULL) == NULL) {
```

```
printf("[Error] at line:%d - \"%s\" has already been declared\n",
@1.last line, $1);
    sprintf($$, "%s [ %s ]", $1, $3);
    $ = strdup($);
      }
ARRAY VARIABLE DECLARATION IDENTIFIER WITH SIZE
      : T IDENTIFIER '[' EXPRESSION ']' {
            if (insert($1, "Identifier-Array", variable declaration type,
@1.last line, NULL) == NULL) {
                  printf("[Error] at line:%d - \"%s\" has already been declared\n",
@1.last_line, $1);
    sprintf($$, "%s [ %s ]", $1, $3);
    $ = strdup($);
ARRAY LIST
      : '{' LITERAL LIST '}' {
    sprintf($$, "{ %s }", $2);
    $ = strdup($);
      T STRING LITERAL {
    $ = strdup($1);
LITERAL LIST
      : IDENTIFIER_OR_LITERAL ',' LITERAL_LIST {
    sprintf($$, "%s , %s", $1, $3);
    $ = strdup($);
      | IDENTIFIER OR LITERAL {
    $$ = strdup($1);
```

COUT

```
: T IO COUT T IO INSERTION INSERTION LIST {
    sprintf($$, "%s %s %s", $1, $2, $3);
    $ = strdup($);
INSERTION LIST
      : EXPRESSION T IO INSERTION INSERTION LIST {
    sprintf($$, "%s %s %s", $1, $2, $3);
    $ = strdup($);
      | EXPRESSION {
    $$ = strdup($1);
CIN
      : T IO CINT IO EXTRACTION EXTRACTION LIST {
    sprintf($\$, "\%s \\%s \\%s", $1, $2, $3);
    $ = strdup($);
EXTRACTION LIST
      : T IDENTIFIER T IO EXTRACTION EXTRACTION LIST {
            if (lookup(\$1) == NULL) 
                  printf("[Error] at line:%d - Undeclared variable \"%s\" \n",
@1.last_line, $1);
    sprintf($$, "%s %s %s", $1, $2, $3);
    $ = strdup($);
      | T IDENTIFIER {
            if(lookup(\$1) == NULL) {
                  printf("[Error] at line:%d - Undeclared variable \"%s\" \n",
@1.last line, $1);
            \$\$ = strdup(\$1);
```

```
RETURN
     : T RETURN EXPRESSION {
    sprintf($$, "%s %s", $1, $2);
    $ = strdup($);
LOGICAL OPERATOR
     : T LOG OP AND {
    $ = strdup($1);
     T LOG OP OR {
    $$ = strdup($1);
RELATIONAL OPERATOR
     : T_REL_OP_EQUAL {
    $$ = strdup($1);
     | '>' {
    $$ = strdup($1);
     T REL OP GREATER THAN EQUAL {
    $$ = strdup($1);
     | '<' {
    $$ = strdup($1);
     |T REL OP LESS THAN EQUAL {
    $$ = strdup($1);
     | T_REL_OP_NOT_EQUAL {
    $$ = strdup($1);
IDENTIFIER OR LITERAL
     : T IDENTIFIER {
           if(lookup(\$1) == NULL) \{
```

```
printf("[Error] at line:%d - Undeclared variable \"%s\" \n",
@1.last line, $1);
             $\$ = strdup(\$1);
      T IDENTIFIER '(' ')' {
             if (lookup(\$1) == NULL) 
                    printf("[Error] at line:%d - Function \"%s\" not defined \n",
@1.last line, $1);
             sprintf($$, "%s()", $1);
    $$ = strdup($$);
      | T IDENTIFIER '(' LITERAL LIST ')' {
             if(lookup(\$1) == NULL) {
                   printf("[Error] at line:%d - Function \"%s\" not defined \n",
@1.last line, $1);
             sprintf($$, "%s ( %s )", $1, $3);
    $ = strdup($);
      T IDENTIFIER UNARY OPERATOR {
             if(lookup(\$1) == NULL) {
                    printf("[Error] at line:%d - Undeclared variable \"%s\" \n",
@1.last line, $1);
             sprintf($$, "%s %s", $1, $2);
    $$ = strdup($$);
      | UNARY OPERATOR T IDENTIFIER {
             if (lookup(\$2) == NULL) {
                    printf("[Error] at line:%d - Undeclared variable \"%s\" \n",
@2.last line, $2);
             sprintf($$, "%s %s", $1, $2);
    $ = strdup($);
      T IDENTIFIER '[' EXPRESSION ']' {
             if (lookup(\$1) == NULL) 
                   printf("[Error] at line:%d - Undeclared variable \"%s\" \n",
@1.last line, $1);
```

```
sprintf($$, "%s [ %s ]", $1, $3);
    $ = strdup($);
      }
      | UNARY OPERATOR T IDENTIFIER '[' EXPRESSION ']' {
            if(lookup(\$2) == NULL) {
                  printf("[Error] at line:%d - Undeclared variable \"%s\" \n",
@2.last_line, $2);
            sprintf($$, "%s %s [ %s ]", $1, $2, $4);
    $ = strdup($);
      | T IDENTIFIER '[' EXPRESSION ']' UNARY OPERATOR {
            if(lookup(\$1) == NULL) \{
                  printf("[Error] at line:%d - Undeclared variable \"%s\" \n",
@1.last line, $1);
            sprintf($$, "%s [ %s ] %s", $1, $3, $5);
    $ = strdup($);
      T CHAR LITERAL {
    $$ = strdup($1);
      T NUMBER LITERAL {
    $$ = strdup($1);
      T STRING LITERAL {
    $$ = strdup($1);
      T BOOL LITERAL {
    $$ = strdup($1);
UNARY OPERATOR
      : T OP INCREMENT {
    $$ = strdup($1);
      | T OP DECREMENT {
    $$ = strdup($1);
```

5. DESIGN STRATEGY

i) SYMBOL TABLE CREATION

The symbol table is implemented as a hash table for constant time (O(1)) access. Chaining is done to avoid collision. A unique hash function based on the name of the identifier is calculated and the location in the symbol table is determined based on that.

ii) INTERMEDIATE CODE GENERATION

A stack based approach is used to generate Three Address Code which is then converted to quadruple format as specified.

iii) CODE OPTIMIZATION

The quadruple data structure created by the generation of intermediate code is taken as input for the optimization phase. This phase performs the different optimizations and gives the resulting quadruple data structure as the output.

The following code optimizations have been performed:

i. Constant folding

- Expressions that can be evaluated at compile time as the arguments forming the expressions are constants, are evaluated and the resulting value is assigned to the appropriate variable.
- Algebraic identities are also constant folded by this optimization. Algebraic identities are equations that are always true regardless of the value assigned to the variables
- Example for algebraic identity constant folding: a + 0 = 0 + a = a

ii. Constant propagation

- If the value of the variable is a constant that is known at compile time, this value is propagated and substituted whenever this variable is encountered.
- This is usually followed by constant folding

iii. Common Subexpression Elimination

- An occurrence of an expression E is called a common subexpression, if E is previously computed and the values in E have not changed since the previous computation.
- All such future occurrences of the expression can be eliminated as there is no need to recompute the value of the expression
- The Variables that are assigned to these future occurrences of the expression are assigned to the temporary that holds the value of the original expression E

iv. Strength Reduction

- Here an expensive operation is replaced by a cheaper operation.
- The cost being talked about here is with respect to the evaluation of the expression by the underlying hardware
- Example:
 - \circ a*2 => a<<1
 - \circ a/2 => a>>1

iv) ERROR HANDLING

The following error handling strategies have been employed:

i) In Scanner

- Length of identifier greater than 32 characters (automatically truncated to first 32 characters after generating the error)
- Incomplete character literals (For eg: 'a)
- Incomplete string literals (For eg: "string)
- Non terminating comments are flagged as erroneous
- Any invalid characters not supported by the programming language are identified to be errors

ii) In Parser

- Redeclaration of variables
- Usage of undeclared variables
- Incorrect usage of break statements (must be used only inside loop bodies)
- Usage of undefined functions

iii) Semantic Analyzer

- Invalid operations on unsupported data types (Semantic error)
- Invalid assignment of incompatible types (Semantic error)

6. IMPLEMENTATION DETAILS

SYMBOL TABLE CREATION

```
typedef struct symbol_table_ {
    int line_number;
    char name[MAX_IDENTIFIER_SIZE];
    char type[MAX_IDENTIFIER_SIZE];
    char category[MAX_IDENTIFIER_SIZE];
    char value[MAX_IDENTIFIER_SIZE];
    int size;
    int scope;
} symbol_table;

typedef struct node {
    symbol_table *st;
    struct node *next;
} node_t* complete_symbol_table[SYMBOL_TABLE_SIZE];
```

This is the structure of the symbol table.

INTERMEDIATE CODE GENERATION

For each type of TAC which has to be generated, a separate function is defined. This function is called as part of the action when the grammar is matched.

For example, expr_code_gen function is called when a binary expression is matched. When an expression is matched, the stack now contains the top two elements as the left hand side and right hand side of the expression and thus the Three Address Code is generated. The corresponding quadruple is also generated.

```
Example:
```

```
void expr_code_gen(char *op)
{
    // create temp
    if(TAC)
    {
        char temp_var[20];
        sprintf(temp_var,"t%d",temp_id);
        temp_id+=1;
        printf("%s = %s %s %s\n",temp_var,stack[top-1],op,stack[top]);
        char value_temp[40000];
```

```
sprintf(value_temp,"%s %s %s",stack[top-1],op,stack[top]);
insert(temp_var,"temporary","", 0, value_temp);
//quadruple add
create_quad(stack[top-1],stack[top],temp_var,op);

top-=1;
strcpy(stack[top],temp_var);
}
```

CODE OPTIMIZATION

- Input to code optimizer program is the Intermediate code generated in Quadruple format
- A python program inputs this Intermediate code and stores it in a List data structure
- Logic is implemented using general if and for loops
- Output generated is also in Quadruple format displayed accordingly

ERROR HANDLING

Provide instructions on how to build and run your program.

- yyleng used to identify the length of identifiers, to flag very long identifiers
- Invalid characters are flagged using Regular expression rules
- Incomplete character literals are identified by use of Regex
- Incomplete string literals are identified by use of Regex
- Nonterminating comments are identified by use of Regex
- All parser errors are identified by looking up the Symbol table which stores all information about variables
- Semantic errors are identified by usage of flag variables to correspondingly identify matching types and operations

7. Results and possible shortcomings of your Mini-Compiler

- All functionalities defined in the Architecture of our mini-compiler have been implemented and works as expected
- Basic constructs of if, if-else and for with different variations have been implemented and works as expected
- Intermediate code generation works for all the syntax handled by the grammar defined in the vacc file
- Optimization processes are detailed in handling cases of Constant folding, Constant propagation, Common subexpression elimination and Strength Reduction
- This only being a mini-compiler, it cannot be used as a complete alternative to the standard C++ compiler. This compiler only handles restricted syntax and constructs due to time constraints
- Compiler is designed only to handle static operations and no run time operations

8. SNAPSHOTS (of different outputs)

Example 1:

Input:

```
#include <stdio.h>

int main()
{
    int a = 10;
    int b = 20;
    int c = a + b;
    float d = b / a;
    double e = a + b;
    bool f = a * 2 - b;
    char g = 'r';
    double h = 1 + 4.5 / 3 + 10;
```

Output:

Symbol Table:

oken	Category	Туре	Line Number	Scope	Value String
	Identifier	int	5	1	10
	Identifier	int	6	1	20
	Identifier	int	7	1	t0
	Identifier	float	8	1	t1
	Identifier	double	9	1	t2
	Identifier	bool	10	1	false
	Identifier	char	11	1	'r'
	Identifier	double	12	1	t7
)	temporary		Θ	1	a + b
	temporary		0	ī	b / a
2	temporary		Θ	1	a + b
8	temporary		Θ	1	a * 2
4	temporary		Θ	1	t3 - b
5	temporary		0	1	4.5 / 3
5	temporary		0	ī	1 + t5
,	temporary		0	ī	t6 + 10
ain	Function-Identifier	int	3	0	NA NA

Intermediate Code (in TAC):

a = 10

b = 20

t0 = a + b

c = t0

t1 = b / a

d = t1

t2 = a + b

e = t2

t3 = a * 2

t4 = t3 - b

f = t4

g = 'r'

t5 = 4.5 / 3

t6 = 1 + t5

t7 = t6 + 10

h = t7

Intermediate Code (in Quadruple format):

QUADRUPLES

Op arg1 arg2 res = 10 a = 20 b

b t0 a t0 c t1b a t1 d b t2 a t2 e 2 t3 a t3 t4 b f t4 'r' g 4.5 t5 3 t6 t5 1 t6 10 t7 t7 h

Optimised Intermediate Code:

With Common Subexpression Elimination and Strength Reduction:

op	arg1	arg2	result
=	10		0
			a
=	20		b
= + = /	a	b	t0
=	t0		c
/	b	a	t1
= << -	t1		d
<<	a	1	t3
-	t3	b	t4
=	t4		f
= = /	'r'		g
/	4.5	3	t5
+	1	t5	t6
+ =	t6	10	t7
=	t7		h

With Constant folding and Constant propagation:

op arg1 arg2 result

```
10
                    a
       20
                    b
=
       30
                    t0
=
       30
                     c
       2.0
                    t1
       2.0
                     d
                    t3
       20
       0
                    t4
                    f
       0
      'r'
                    g
       1.5
                    t5
=
       2.5
                    t6
      12.5
                    t7
       12.5
                    h
```

Example 2:

Input:

```
#include <stdio.h>
int main()
{
   int a = 10 + 5;
   int b = 5 * 3;
   double c;
   if(a > b)
   {
      c = 1.0;
   }
   else
   {
      c = 0.0;
   }
   if(a)
   {
      if(b)
      {
        int temp1 = 10;
   }
}
```

```
else if(c)
{
    int temp2 = 20;
}
else
{
    int temp3 = 1;
}
}
```

Output:

Symbol Table:

oken	Category	Туре	Line Number	Scope	Value String
	Identifier	int	5	1	t0
)	Identifier	int	6	1	t1
	Identifier	double		1	NA
:0	temporary		0	1	10 + 5
1	temporary		0	1	5 * 3
2	temporary		0	1	a > b
ain	Function-Identifier	int	3	0	NA
emp1	Identifier	int	20	5	10
emp2	Identifier	int	24	7	20
emp3	Identifier	int	28	8	1

Intermediate Code (in TAC):

```
t0 = 10 + 5

a = t0

t1 = 5 * 3

b = t1

t2 = a > b

if t2 goto L0

goto L1

L0:

c = 1

goto L2

L1:

c = 0

L2:

if a goto L3

goto L4
```

L3: if b goto L5 goto L6 L5: temp1 = 10goto L7 Ľ6: if c goto L8 goto L9 L8: temp2 = 20goto L10 L9: temp3 = 1L10: L7: goto L11 L4: L11:

Intermediate Code (in Quadruple format):

QUADRUPLES

arg2 Op arg1 res 5 t0 +10 t0= a 5 * 3 t1 t1 b = > a b t2 if L0 t2 L1 goto L0 label = 1 c L2 goto L1 label 0 = c label L2 if L3 a goto L4 L3 label L5 if b

goto		L6
L5		label
=	10	temp1
goto		L7 _
L6		label
L8	c	if
goto		L9
L8		label
=	20	temp2
goto		L10
L9		label
=	1	temp3
L10		label
L7		label
goto		L11
L4		label
L11		label

Optimised Intermediate Code:

With Common Subexpression Elimination and Strength Reduction:

op	arg1	arg2	result
+	10	5	t0
=	t0		a
*	5	3	t1
=	t1		b
>	a	b	t2
L0	t2		if
goto			L1
L0			label
=	1		c
L1			label
=	0		c
L2			label
L3	a		if
L3			label
L5	b		if
L5			label
=	10		temp1

L6		label
L8	c	if
L8		label
=	20	temp2
L9		label
L10		label
L7		label
L4		label
L11		label

With Constant folding and Constant propagation:

op	arg1	arg2	result
=	15		t0
=	15		a
=	15		t1
=	15		b
=	False		t2
L0	t2		if
goto			L1
L0			label
=	1		c
L1			label
=	0		c
L2			label
L3	15		if
L3			label
L5	15		if
L5			label
=	10		temp1
L6			label
L8	0		if
L8			label
=	20		temp2
L9			label
L10			label
L7			label
L4			label
L11			label

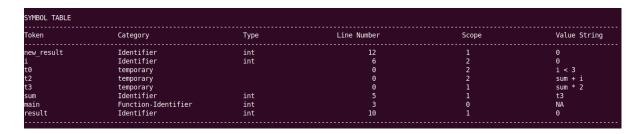
Example 3:

Input:

```
#include <stdio.h>
int main()
{
    int sum = 0;
    for(int i = 0; i < 3; i++)
    {
        sum += i;
    }
    int result = sum;
    sum *= 2;
    int new_result = sum;
}</pre>
```

Output:

Symbol Table:



Intermediate Code (in TAC):

```
sum = 0
i = 0
L0:
t0 = i < 3
if t0 goto L1 goto L2
```

```
L3:

t1 = i + 1

i = t1

goto L0

L1:

t2 = sum + i

sum = t2

goto L3

L2:

result = sum

t3 = sum * 2

sum = t3

new_result = sum
```

Intermediate Code (in Quadruple format):

QUADRUPLES

QUIDICITEE						
Op	arg1	arg2	res			
=	0		sum			
=	0		i			
L0			label			
<	i	3	t0			
if	t0		L1			
goto			L2			
L3			label			
+	i	1	t1			
=	t1		i			
goto			L0			
L0			label			
+	sum	i	t2			
=	sum		t2			
goto			L3			
L2			label			
=	sum		result			
*	sum	2	t3			
=	sum		t3			
=	sum		new_result			

Optimised Intermediate Code:

With Common Subexpression Elimination and Strength Reduction:

op	arg1	arg2	result
=	0		sum
<	sum	3	t0
if	t0		L1
goto			L2
L3			label
+	sum	1	t1
=	t1		sum
L0			label
+	sum	sum	t2
=	sum		t2
L2			label
<<	sum	1	t3
=	sum		new_result

With Constant folding and Constant propagation:

op	arg1	arg2	result
_	0		
=	0		sum
=	True		t0
if	t0		L1
goto			L2
L3			label
=	1		t1
=	1		sum
L0			label
=	2		t2
=	1		t2
L2			label
=	2		t3
=	1		new result

Example 4:

Input:

```
#include <stdio.h>
int main()
{
    int a = 0;
    double b = 1.5;
    int sum;
    if(b)
    {
        sum = 0;
        for(int i = 0;i<2;i++)
        {
            sum = sum + b;
        }
        if(sum > 3)
        {
            sum = 10;
        }
        else
        {
            sum = 0;
        }
    }
    int result = sum;
}
```

Output:

Symbol Table:

SYMBOL TABLE						
Token	Category	Туре	Line Number	Scope	Value String	
1	Identifier	int	5	1	0	
	Identifier	double	6	1	1.5	
	Identifier	int	10	3	Θ	
0	temporary		0	3	i < 2	
2	temporary		Θ	3	sum + b	
3	temporary		0	2	sum > 3	
um	Identifiér	int	9	2	Θ	
ain	Function-Identifier	int	3	0	NA	
esult	Identifier	int	23	1	Θ	

Intermediate Code (in TAC):

a = 0

b = 1.5

if b goto L0

goto L1

L0:

sum = 0

i = 0

L2:

t0 = i < 2

if t0 goto L3

goto L4

L5:

t1 = i + 1

i = t1

goto L2

L3:

t2 = sum + b

sum = t2

goto L5

L4:

t3 = sum > 3

if t3 goto L6

goto L7

L6:

sum = 10

goto L8

L7:

sum = 0

L8:

goto L9

Ľ1:

L9:

result = sum

Intermediate Code (in Quadruple format):

QUADRUPLES

1		. 0
b		if
		L1
		label
0		sum
0		i
		label
i	2	t0
t0		L3
		L4
		label
i	1	t1
t1		i
		L2
		label
sum	b	t2
t2		sum
		L5
		label
sum	3	t3
t3		if
		L7
		label
10		sum
		L8
		label
0		sum
		label
		L9
		label
		label
sum		result
	0 i t0 i t1 sum t2 sum t3	0 0 1 2 t0 1 1 1 1 1 1 2 sum b t2 sum 3 t3

Optimised Intermediate Code:

With Common Subexpression Elimination and Strength Reduction:

goto			L1
L0			label
=	0		i
<	i	2	t0
if	t0		L3
L5			label
+	i	1	t1
=	t1		i
L2			label
+	a	b	t2
=	t2		a
L4			label
>	a	3	t3
L6	t3		if
L6			label
=	10		a
L7			label
L8			label
L1			label
L9			label
=	i		result

With Constant folding and Constant propagation:

op	arg1	arg2	result
=	1.5		b
L0	1.5		if
goto			L1
L0			label
=	0		i
=	True		t0
if	t0		L3
L5			label
=	1		t1
=	1		i
L2			label
+	a	1.5	t2
=	t2		a
L4			label

>	a	3	t3
L6	t3		if
L6			label
=	10		a
L7			label
L8			label
L1			label
L9			label
=	1		result

9. CONCLUSIONS

This implementation of a C++ mini-compiler enables compilation of basic C++ programs covering basic constructs of if, if-else and for. Corresponding Intermediate code is generated by the front end of the compiler, in the form of Three Address Code, represented in Quadruple format. Multiple machine independent optimizations have been implemented which improves efficiency and execution times of the compiled programs.

This project has enabled us to experience a true hands-on approach to learning about Compiler Design. We have learnt a lot about different phases involved in the compilation process. We have learnt to write the Lexer to generate tokens, write the required context-free grammar for handling the language syntax and generate a parse tree, write modular code to generate the intermediate code and develop functions to convert the generated intermediate code into a more optimized version.

10. FURTHER ENHANCEMENTS

This being only a mini-compiler, it can be extended to cover more standard C++ functionalities like classes, STL functions, etc. This would require building upon our current context-free grammar.

We can also bring about further enhancements through optimizations. Loop level optimizations can be implemented to improve efficiency and execution times of our programs.

REFERENCES/BIBLIOGRAPHY

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- ISO C++ reference (https://isocpp.org/wiki/faq/compiler-dependencies#yaccable-grammar)
- C++ yacc-able grammar reference (http://www.computing.surrey.ac.uk/research/dsrg/fog/CxxGrammar.y)
- Lecture slides on Intermediate Code generation and Optimisation techniques provided by PES University Compiler Design Faculty