GLOBAL AND MULTI-OBJECTIVE OPTIMIZATION

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INTRODUCTION TO GENETIC PROGRAMMING (PART 2)

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LINEAR GP

MOTIVATIONS

- Trees are not the only way of representing programs
- Also streams of instructions can represent programs
- Instead of LISP-like structure we now use assembly-like commands
- Linear GP: a linear stream of "assembly-like" instructions

AN EXAMPLE OF LINEAR GP

Register I

Register 2

Register3

Register 4

Registers of a virtual (or real!) machine

Add RI, R2, RI

Sub R3, R1, R4

Add R4,R3,R2

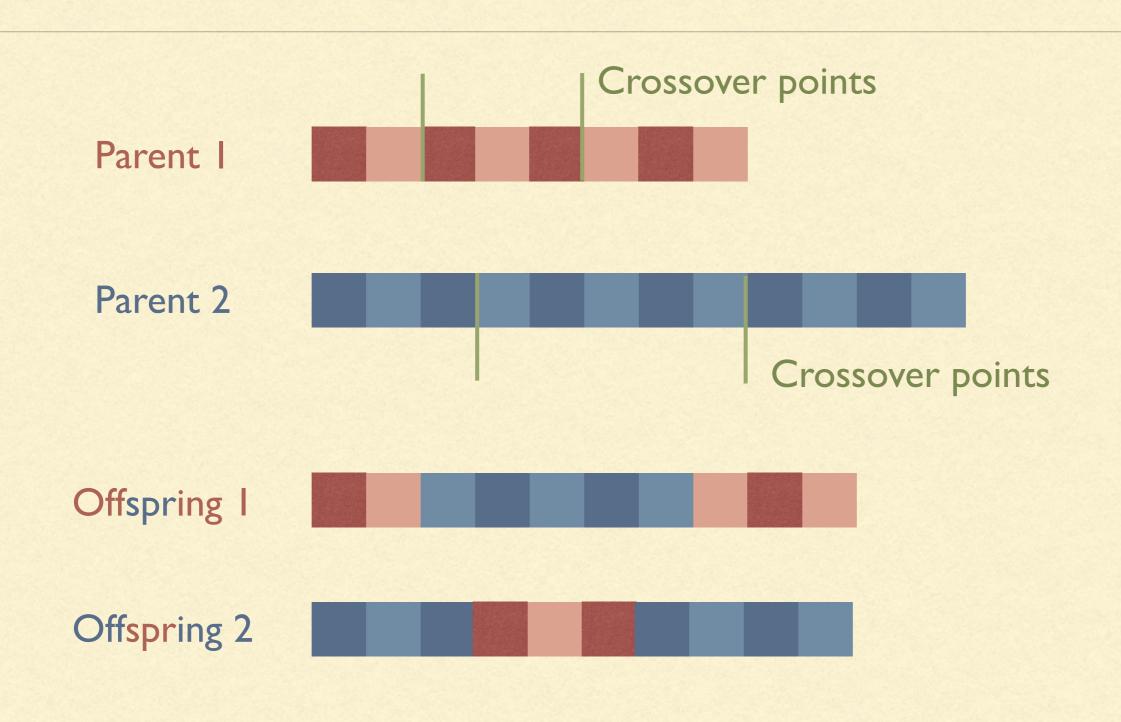
Mul RI,RI,R2

A Linear GP individual:
a list of instructions for the machine communicating via registers

LINEAR GP AND GA

- Linear GP seems pretty similar to standard GA
- Except that the individuals can be of non-fixed length
- An important difference is that we are evolving programs and we "execute" the individuals
- Most of the operators of GA can be used for Linear GP
- Two points crossover (with possibly different crossover points between the two individuals) is usually employed

TWO-POINTS CROSSOVER

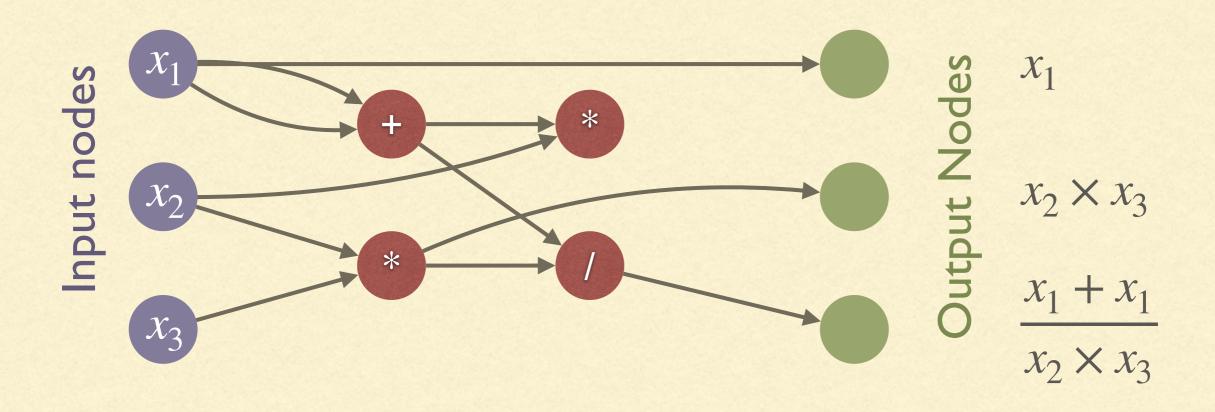


CARTESIAN GP

MOTIVATIONS

- It is possible to represent programs as circuits/graphs
- Cartesian GP represents individuals in this way (invented by Julian F. Miller)
- Naturally suited for problems with many inputs and many outputs (instead of using multiple trees)
- Has some tracts in common with linear GP (the representation for the circuit/graph is encoded in a linear way)

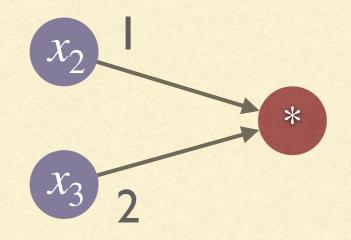
AN INDIVIDUAL OF CGP



But how is the individual encoded?

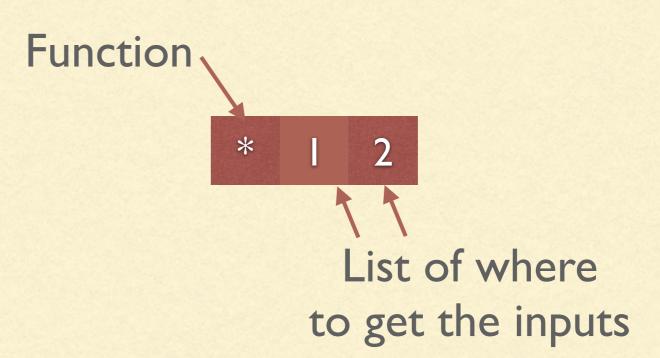
ENCODING IN CGP

Internal nodes



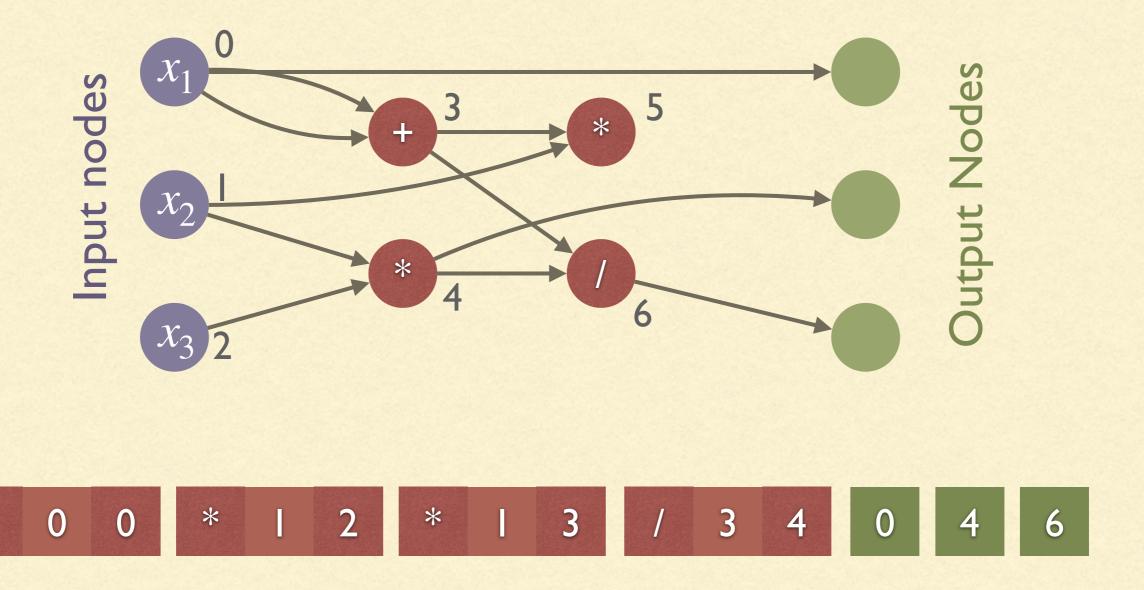
Output nodes



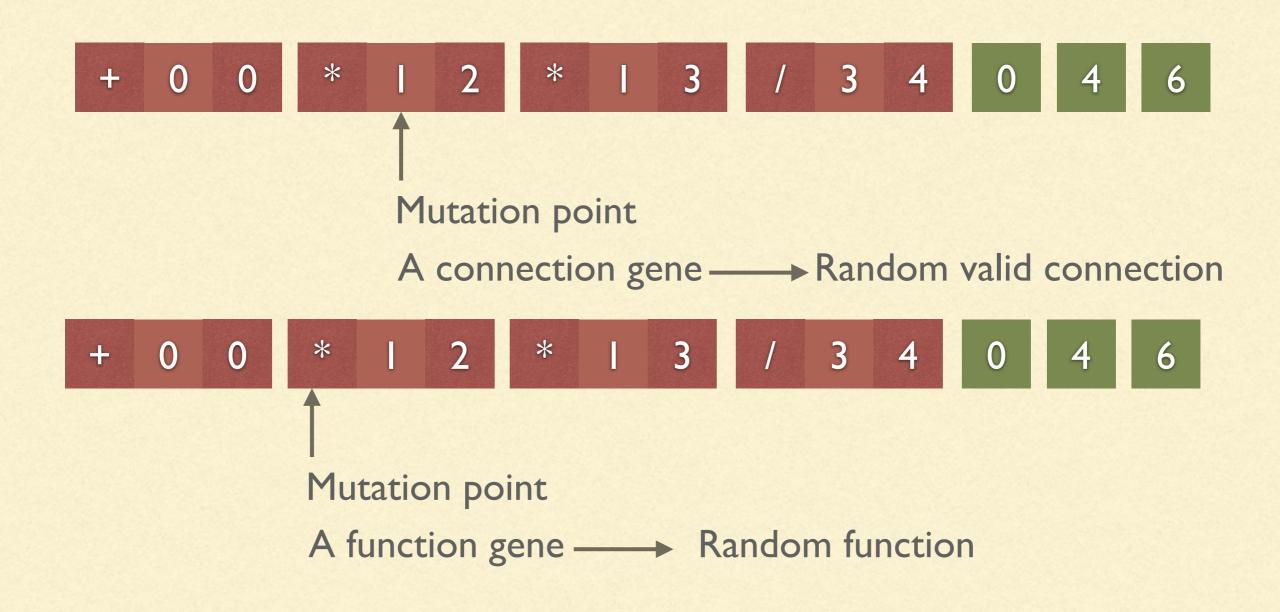


6 Where to get the output

AN INDIVIDUAL OF CGP



CGP MUTATION



GRAMMATICAL EVOLUTION



- GP trees allow to express constructs of high-level languages
- A linear representation is easier to manage
- Grammatical Evolution combines a linear genome with a treebased representation (in an indirect way)
- The idea is to define a grammar that is then used to interpret a linear genome as a derivation of that grammar

BACKUS NAUR FORM (BNF)

- Invented by John Backus and refined by Peter Naur
- Used to define context-free grammars, initially for the ALGOL programming language
- It expresses a language as a set of production rules
 - Each rule maps a non-terminal to a string of zero or more non-terminals and terminals

BACKUS NAUR FORM (BNF)

- A grammar in BNF is a quadruple (T, N, P, S) where:
 - T is a set of terminal symbols
 - N is a set of non-terminal symbols
 - P is a set of production rules
 - $S \in N$ is the axiom, i.e., the start symbol

Terminal symbols

$$T = \{+,0,1,2,3,4,5,6,7,8,9\}$$

Non-terminal symbols

$$N = \{S, C, D\}$$
 Axiom S

What kind of strings can be defined with this grammar?

Production Rules

$$S \rightarrow S + S \mid C$$

$$C \rightarrow D \mid DC$$

$$D \rightarrow 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$$

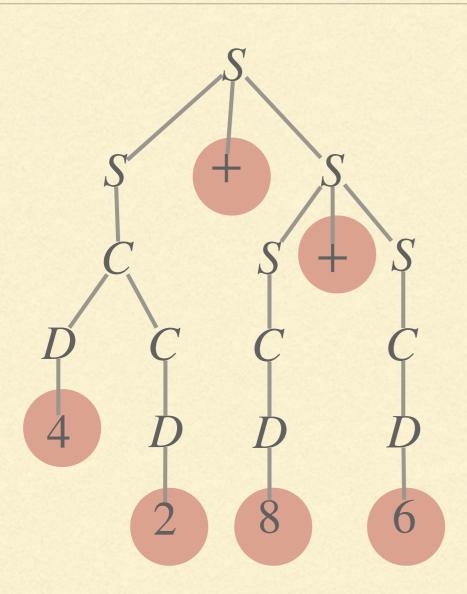
$$S \to S + S \to C + S \to DC + S \to 4C + S \to 4D + S$$

 $\to 42 + S \to 42 + S + S \to 42 + C + S \to 42 + D + S$
 $\to 42 + 8 + S \to 42 + 8 + C \to 42 + 8 + D \to 42 + 8 + 6$

Production Rules

$$S \to S + S \mid C$$

 $C \to D \mid DC$
 $D \to 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$



The expansion can also be seen as building a tree

All the non-terminal are internal nodes

The terminals are the leaves

The resulting string is obtained by reading all the leaves in order

- With the previous grammar we can, depending on the expansion that we select, generate all strings that represents the sum of integers numbers
- Notice that if we want to obtain the result of the sum we need to evaluate the string
- This means that if we can have a grammar expressing a program (with variables) and evaluate it on a specific input

FROM BNFTO GE

- What if we want to generate a specific string?
- We can keep track of which expansion to select among all the possible ones for a non-terminal symbol
- Something like:
 - Select the expansion number 1, then number 3, then number 7, then number 2, etc.
 - Seems like we can have a vector (1,3,7,2,...)

FROM BNFTO GE

- Once we have a vector of integers we can use GA-like evolution!
- The genome is a sequence of integers
- To evaluate an individual we need to generate a string via a grammar...
- ...and then evaluate the string
- We are sure that the string is a valid program because it was generated by a grammar!

TOWARDS GE

$$S \rightarrow S + S \mid C$$

$$C \to D \mid DC$$

We number the RHS of all production rules

But not all production rules have the same number of RHS!

$$0$$
 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | $D \rightarrow 0$ | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9

Should our genome be in $\{0,1\}^n$ or in $\{0,1,\ldots,9\}^n$?

TOWARDS GE

Should our genome be in $\{0,1\}^n$ or in $\{0,1,\ldots,9\}^n$?

$$S \to S + S \mid C$$

Neither, we can use any integer! Simply take it modulus the number of RHS

$$C \rightarrow D \mid DC$$

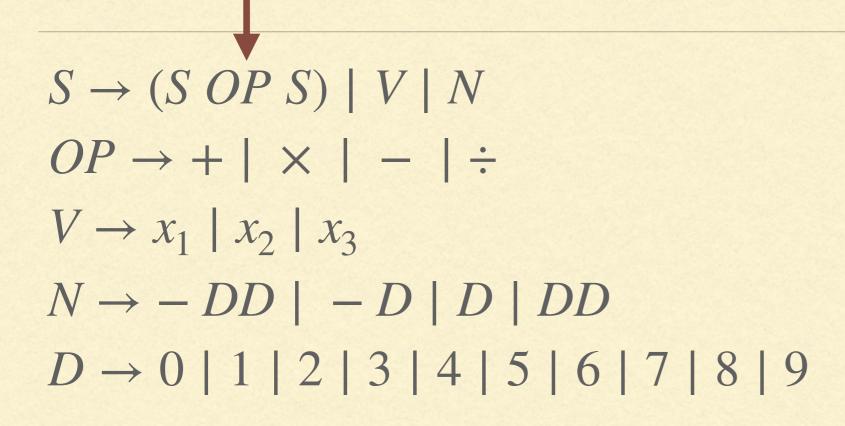
$$Mod 2 = 1$$

$$7548397623$$

$$S \to (S \ OP \ S) \mid V \mid N$$

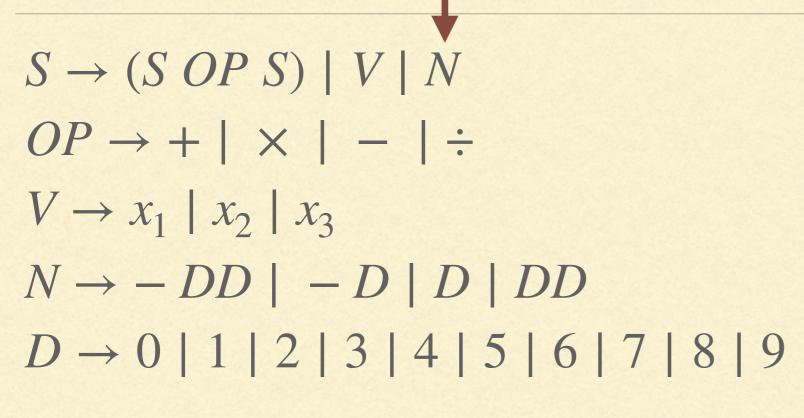
 $OP \to + \mid \times \mid - \mid \div$
 $V \to x_1 \mid x_2 \mid x_3$
 $N \to -DD \mid -D \mid D \mid DD$
 $D \to 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$





S OP S









N OP S

D OP S

$$S \to (S \ OP \ S) \mid V \mid N$$

$$OP \to + \mid \times \mid - \mid \div$$

$$V \to x_1 \mid x_2 \mid x_3$$

$$N \to -DD \mid -D \mid D \mid DD$$

$$D \to 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$$

$$\downarrow \qquad \qquad \downarrow$$

$$3 \quad 5 \quad 22 \quad 6 \quad 9 \quad 1 \quad 56 \quad 18 \quad 24 \quad 1$$

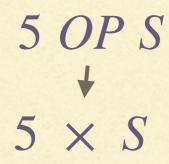
$$0 \quad 2 \quad 2 \quad 6$$

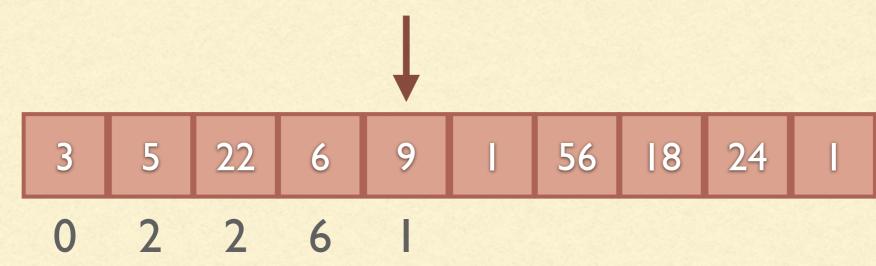
DOPS

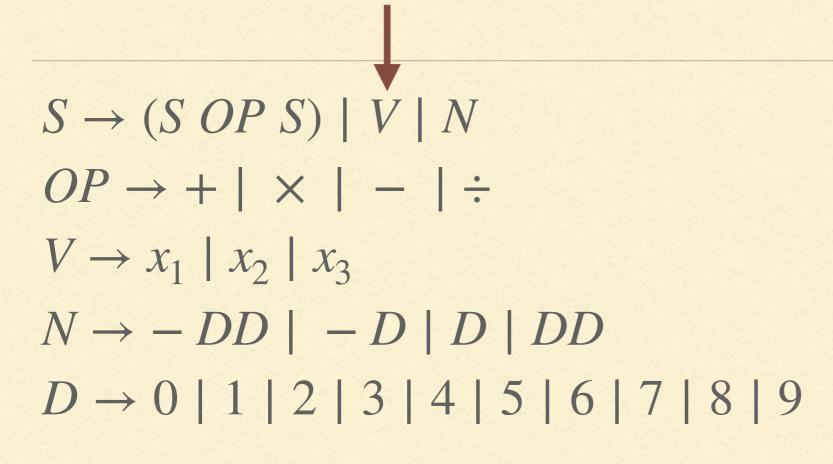
5 OP S

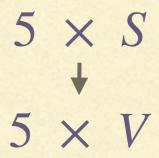
$$S \rightarrow (S \ OP \ S) \mid V \mid N$$

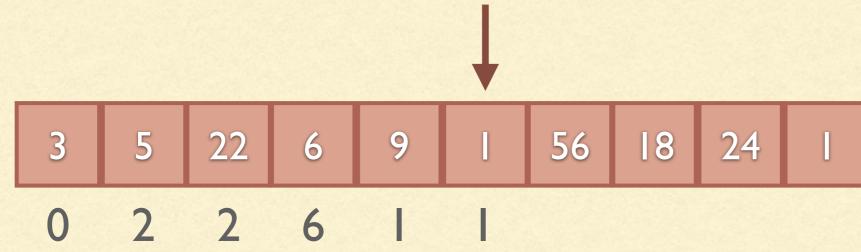
 $OP \rightarrow + \mid \times \mid - \mid \div$
 $V \rightarrow x_1 \mid x_2 \mid x_3$
 $N \rightarrow -DD \mid -D \mid D \mid DD$
 $D \rightarrow 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$





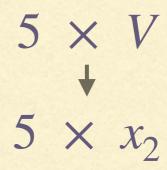


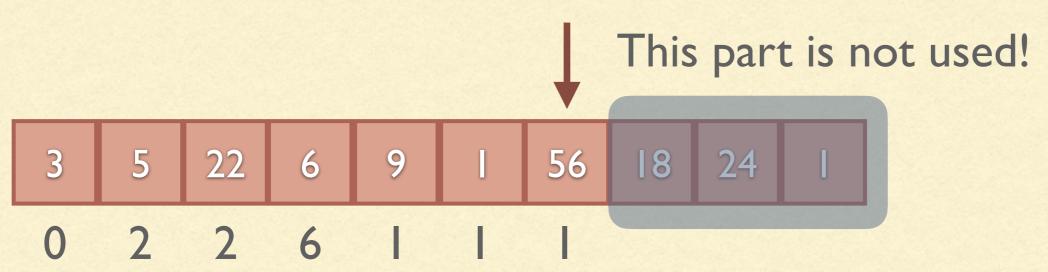




$$S \to (S \ OP \ S) \mid V \mid N$$

 $OP \to + \downarrow \times \mid - \mid \div$
 $V \to x_1 \mid x_2 \mid x_3$
 $N \to -DD \mid -D \mid D \mid DD$
 $D \to 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$





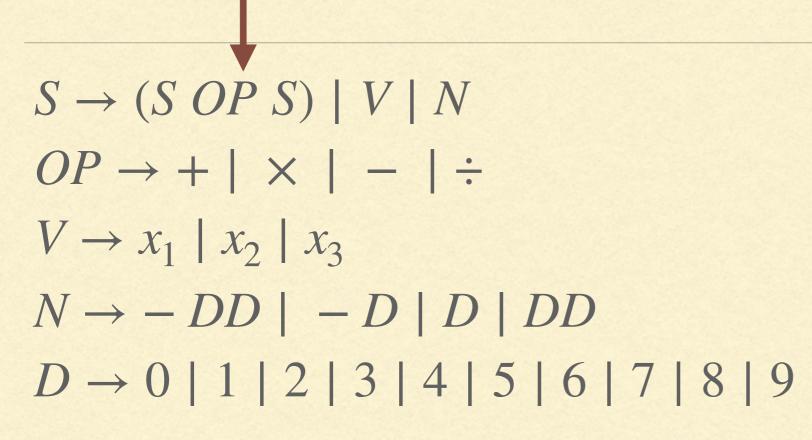
GE PECULIARITIES

- The individuals might not be all of the same length
- The effective length of an individual can be shorter than the length of the genome (previous example)...
- ...but it can also be longer (in this case we "wrap around" and continue from the beginning)
- We must have a maximum length for the expansion!

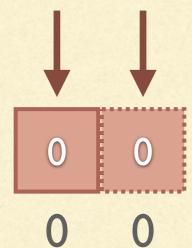
$$S \rightarrow (S \ OP \ S) \mid V \mid N$$

 $OP \rightarrow + \mid \times \mid - \mid \div$
 $V \rightarrow x_1 \mid x_2 \mid x_3$
 $N \rightarrow -DD \mid -D \mid D \mid DD$
 $D \rightarrow 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$

SOPS



$$S OP S$$
 \downarrow
 $(S OP S) OP S$



The expansion will never end!

We will simply assign a very bad fitness to this individual

GE PECULIARITIES

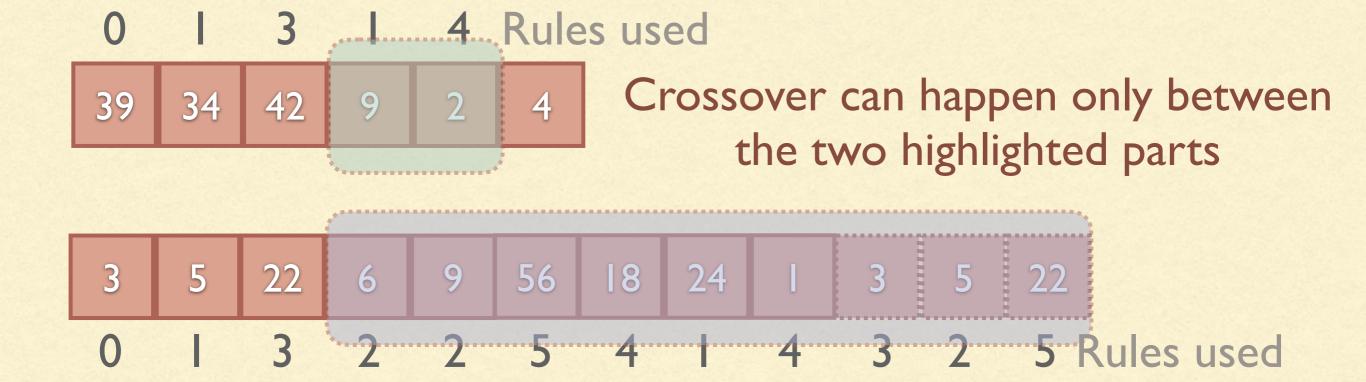
- Notice that there is no one-to-one mapping between genotype and the phenotype
- Multiple genotypes can be expanded to represent the same derivation tree/phenotype
- Notice that the "meaning" of a part of an individual depends from all the genes before it (since they define in which part of the derivation we are)

GE: OPERATORS

- We can use the same operators as GA for crossover and mutation
- We can also use specific operators for crossover:
 - Sensible crossover is a one-point crossover that uses the actual length of an individual
 - Ripple crossover
 - Homologous crossover

HOMOLOGOUS CROSSOVER

- In homologous crossover we align the history of rule activations in two individuals
- We perform a one-point crossover in the area where they differs



LIBRARIES

- Grammatical evolution is implemented in PonyGE2 library
- https://github.com/PonyGE/PonyGE2
- You need to clone/download the repository (you cannot use pip or anaconda for the installation)