## ALGORITHMICS, graduate program 2025/2026 CRYPTOGRAPHY, 2025/2026, lab assignments list # 1, 09.10.2025 **Regular deadline:** 07.11 or 13.11.2025

Familiarize yourself with the MD5 hash function [1]. Your task in this list is to reproduce the collision attack by Wang et al. [2, 3].

- 1. (2pts) To reproduce the collision attack, you need a working MD5 implementation and its internals. Find an open source implementation, for instance MbedTLS [4] (you can choose any programming language you want). Familiarize yourself with the source code and identify inner workings, in particular F, G, H and I functions, A, B, C, D variables, etc.
- 2. (4pts) Based on the implementation from the previous step (feel free to copy or link/reference, as long as the license permits!) try to compute step-by-step the example from Table 2 in [3]. Ensure that the inputs result in a collision equal to the H from the table. Debug if not. In particular, check for byte order (endianness). In general, we need to check if  $MD5(MD5(IV, M_0), M_1) = MD5(MD5(IV, M_0'), M_1') = H$ . Create an interface that allows you to verify this for any choice of  $M_0, M_0'$  and  $M_1, M_1'$ .
- 3. (4pts) Having this interface, implement the (easier) second step of the attack (Step 2 from Sect. 4.5 in [3]). Create a procedure that, given  $M_0$  and  $M'_0$  on input, finds  $M_1$  and  $M'_1$  that result in a collision (the same hash value). For now, test with  $M_0$  and  $M'_0$  from Table 2.
- 4. (**OPTIONAL**, 5pts) Implement as much of the remainder of the attack as you want. Based on the effort, you can get up to 5 additional points (beyond the standard 10). In particular, you may want to:
  - Implement the first step of the attack (Step 1 from Sect. 4.5 in [3]) using the "Single-message Modification" method (Sect. 4.4). This step should take slightly longer to compute, so give yourself ample time for potential debugging.
  - Next, you may try to re-implement the second step of the attack for your selected message pair. Note that the differential pairs may be different.
  - Finally, generate a collision and test with an external software if you actually obtain an MD5 collision (use, for instance, the md5sum Linux command). Remember about potential encoding differences, byte order, and so on.
  - Add an option to use the "Multi-message Modification" method (Sect. 4.4)

## Hints

- 1. The exhaustive search stage is quite intensive. Make sure that the environment you choose (language, libraries, execution environment, etc) can execute a  $2^{40}$  search in a reasonable time (hours, not weeks).
- 2. Test as much as possible with the data from the paper. If you write a function that verifies if a message clears the requirements from Table 6 check if the messages from Table 2 pass those checks.
- 3. Focus first on correctness, then on performance. Performance is still critical in the main search loop, but leaving the loop running  $2^{40}$  times with a bug in the code can be painful.

- 4. Byte order: MD5 loads data in *Little Endian* byte order, so the order on disk will be different from the order in paper (or the "natural", MSB first order), i.e. the first 4 bytes of  $M_0$  (word 2dd31d1) are d1, 31, dd, 02. All operations in MD5 are performed on 32-bit words, so once the data is stored in that type, the order does not matter.
- 5. Section 4.4 Message Modification describes how to modify a randomly generated message to fulfill at least some of the differential requirements (from tables 4 and 6). The remaining requirements should still be tested. Not all of them have to be; in fact, you can just skip some of them and check if the results satisfy  $\Delta H_1 / \Delta H$ . Early detection allows for weeding out "bad" candidates.
- 6. *Single-message Modification*, as presented in the paper, may be a bit imprecise, or straight-up invalid (rotation *IS NOT* distributive over addition!). To better understand how to apply it, consider the following:
  - (a) In a regular MD5 round,  $a_{i+1} \leftarrow ((a_i + F(b_i, c_i, d_i) + m_i + k_i) \ll s_i) + b_i$ , c.f. [this image].
  - (b) Since we need some specific properties of  $a_{i+1}$  (let us call the modified version  $a'_{i+1}$ ), we want to modify  $m_i$  into  $m'_i$ , st.t.  $a'_{i+1} \leftarrow ((a_i + F(b_i, c_i, d_i) + m'_i + k_i) \ll s_i) + b_i$ .
  - (c) Now you have two equations with a single variable, which should be easy to solve by hand. Note that huge chunks of the two equations are common, but some parts need unwinding, like  $b_i$  which needs to be subtracted from  $a'_{i+1}$  before the rotation can be inverted. Recall that here you solve for  $m'_i$ , all the other variables are known.
- 7. The condition tables are long, and rewriting them by hand is very error-prone. Consider writing a script that can parse them into an easy-to-use form. For instance, translate

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## References

- [1] Ronald L. Rivest. *The MD5 Message-Digest Algorithm*. RFC 1321. Apr. 1992. DOI: 10.17487/RFC1321. URL: https://www.rfc-editor.org/info/rfc1321.
- [2] Xiaoyun Wang et al. "Collisions for Hash Functions MD4, MD5, HAVAL-128 and RIPEMD". In: *IACR Cryptol. ePrint Arch.* (2004), p. 199. URL: http://eprint.iacr.org/2004/199.
- [3] Xiaoyun Wang and Hongbo Yu. "How to Break MD5 and Other Hash Functions". In: Advances in Cryptology EUROCRYPT 2005, 24th Annual International Conference on the Theory and Applications of Cryptographic Techniques, Aarhus, Denmark, May 22-26, 2005, Proceedings. Ed. by Ronald Cramer. Vol. 3494. Lecture Notes in Computer Science. Springer, 2005, pp. 19–35. DOI: 10.1007/11426639\_2. URL: https://iacr.org/archive/eurocrypt2005/34940019/34940019.pdf.

[4]	TrustedFirmware and MbedTLS contributors. <i>MbedTLS MD5 implementation</i> . 2021. URL: https://github.com/Mbed-TLS/mbedtls/blob/3304f253d7aa4b5b18e772c455b2113f7af29ca5/library/md5.c.