

## Assignment # 2

### Homework

Homework problems are a preparation for the quizzes. They are *not* graded. Please use the `mywpi` forum to post questions you have on these problems.

- 2.2, 2.4, 2.5 ( $C(x) = 1 + x + x^3$ ), 2.6, 2.8, 2.10, 3.3, 3.5, 3.7, 3.12

### Project

1. Problem 2.11 from the book.
2. Problem 3.6 from the book.
3. One important property that makes DES secure is the non-linearity of the S-boxes. In particular, a single change in the input should affect at least half the bits of the output. To verify this, we will check the *strict avalanche criterion* (SAC). It is computed as:

$$\sum_{x \in \mathbb{Z}_{2^n}} S(x) \oplus S(x \oplus 2^i) \geq (2^{(n-1)}, 2^{(n-1)}, \dots, 2^{(n-1)})$$

The S-box is a function that maps a  $n$ -bit input to an  $m$ -bit output  $S : \mathbb{Z}_{2^n} \rightarrow \mathbb{Z}_{2^m}$ , i.e. for DES  $n = 6, m = 4$ . Furthermore,  $(2^{(n-1)}, 2^{(n-1)}, 2^{(n-1)}, 2^{(n-1)})$  is the individual sum over each output bit position. In order for the SAC to be fulfilled, the value of each individual sum must be greater or equal to  $2^{(n-1)}$ .

*Hint:* The SAC is computed for each output bit separately, i.e. you should get 6x4 SAC coefficients: one per s-box output bit is  $m = 4$  per row, one row for each input bit position, giving  $n = 6$  rows, i.e. you are checking 6x4 values.

- (a) Implement a function that on the input  $x \in \mathbb{Z}_{2^6}$  produces an output  $y \in \mathbb{Z}_{2^4}$  according to the DES S-box  $S_4$  (i.e. implement the DES S-box  $S_4$ ).
- (b) Write a program that, based on the above implementation of  $S_4$  computes the SAC for  $i = 0 \dots 5$ . Is the SAC fulfilled by  $S_4$ ? Please turn in both the code and the output.

4. Implement an exhaustive key search for DES. Recover the key for the following plaintext-ciphertext pair (both given in hex notation):

```
pt = 48656c6c6f212121
ct = 1f6339383e8da6c4
```

Please turn in your working code together with the correct 64-bit representation of the key.

*Note:* while the key space of DES is way too small by now, it is still too large to be searched in reasonable time on a simple desktop PC. In order to facilitate the search, the first four bytes of the 64-bit key have been fixed to 0. i.e. the 64-bit key looks like this (in hex representation): 00000000XXXXXXX.

*Hint:* Do not try to implement DES yourself. Use an existing implementation of DES. Possible choices are:

- In **sage**, a DES engine is available in `Crypto.Cipher`

```
from Crypto.Cipher import DES
...
```

(Note that the server might quickly be overloaded if many people try to run an exhaustive search in parallel).

- For **Python** there is a versatile crypto library available at: <https://www.dlitz.net/software/pycrypto/>. It needs an additional install, but provides the same functionality
- Any **openssl** implementation will do. The mode to be selected should be `des_ede3_ecb`. The key needs to be replicated three times.

## Bonus Problem

Recover the DES key for the following plaintext-ciphertext pair (both given in hex notation):

```
pt = 48656c6c6f212121
ct = ddb92846141922b8
```

The key starts with zeros, just as the one before (but less). Please turn in your working code together with the correct 64-bit representation of the key and the runtime needed.

# Good Luck and Have Fun!