

Operating System

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(玉泉教7-206)

周二第7,8节;周四第9,10节/教4-301

实验: 周四第7,8节/ 曹西-503

Lecture 1: Overview

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Outlines of the overview

- Also known as “Principles of Operating System”
- Background
 - The main **purpose** of this course
 - The **contents** and **schedule** of the course
 - The pre-requisite knowledge for the course
- Chapter 1 Introduction
- Chapter 2 Operating-System Structures





Main objectives

An introductory course for Operating System

- To learn the basics and internal design of operating systems
- To look at both the history and the state-of-the-art techniques used in operating systems
- To study the design **methodologies** applied in OS
- To **prepare to practice** the techniques in your future work and
- To **have fun !?**





These are NOT the main objectives

- To learn the cool Linux/RISCV commands (stuff like shell programming, pipes, scripts, ...)
- To write a new OS for my PC
- To learn how to design a big system
- To analyze the source code of an OS
 - <http://www.chromium.org/chromium-os>
 - <https://source.android.com>
 - <http://www.tinyos.net/>





Contents

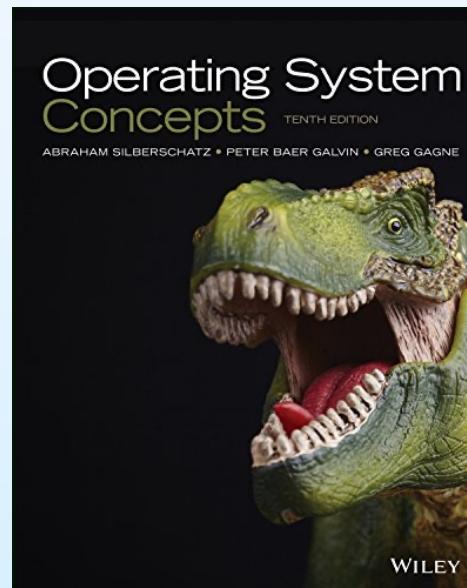
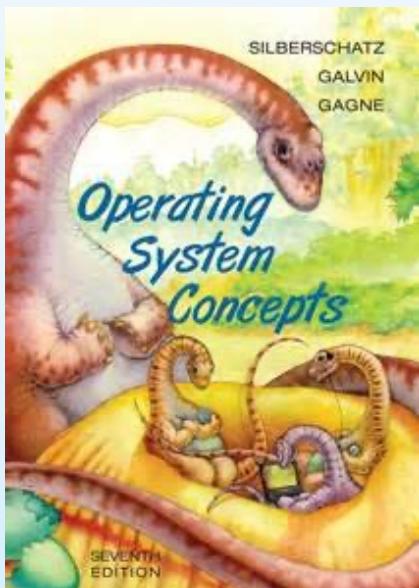
- Overview - 1 session
 - Intro
 - OS structure
- Process Management - 6 sessions
 - Processes
 - Threads
 - CPU scheduling
 - Process Synchronization
 - Deadlocks
- Memory Management- 4 sessions
 - Main memory
 - Virtual memory
- Storage Management-4 sessions
 - File-system interface
 - File-system implementation
 - Mass-storage structure
 - I/O systems
- Additional Topic – 2 session





Textbook & resources

- Main Textbook: **Operating System Concepts 7th/10th, Abrraham Silberschatz、Peter Galvin, Greg Gagne**
- 高等教育出版社 (第7/10版影印版)
- Course webpage: <http://course.zju.edu.cn>



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Pre-requisite Knowledge

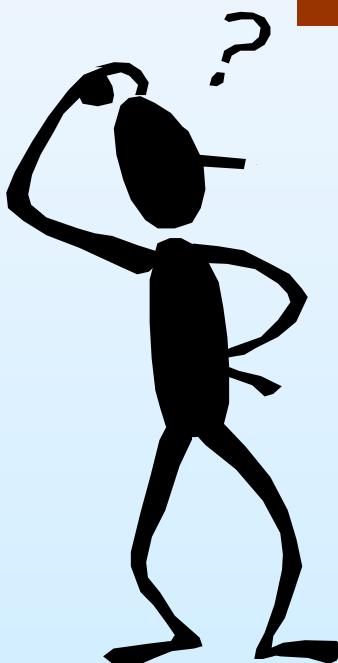
- Computer system architecture/organization
- Programming language C
- Assembling language
- Data structures





Final Grade

- Final Exam = 50%
- Homework = 5%
- Quiz = 5%
- Project Report = 20%
- Project Face-face check = 20%





Project of OS

■ Individual Project

- Lab0: Riscv64 Kernel Compilation 5%
- Lab1: Riscv64 Kernel Bootup and Time Interrupt 15%
- Lab2: Riscv64 Process Scheduling 15%

■ Group Project, 2 ppl/group

- Lab3: Riscv64 Virtual Memory 20%
- Lab4: Riscv64 User Space and System Call 15%
- Lab5: Riscv64 Fork and Page Fault 30%

■ Selective Project (bonus), group

- Lab6: File System 10%

Project Site : <https://yuque.zju.edu.cn/qz21wl/ernn3p/xw0d3r>





My suggestions on learning OS

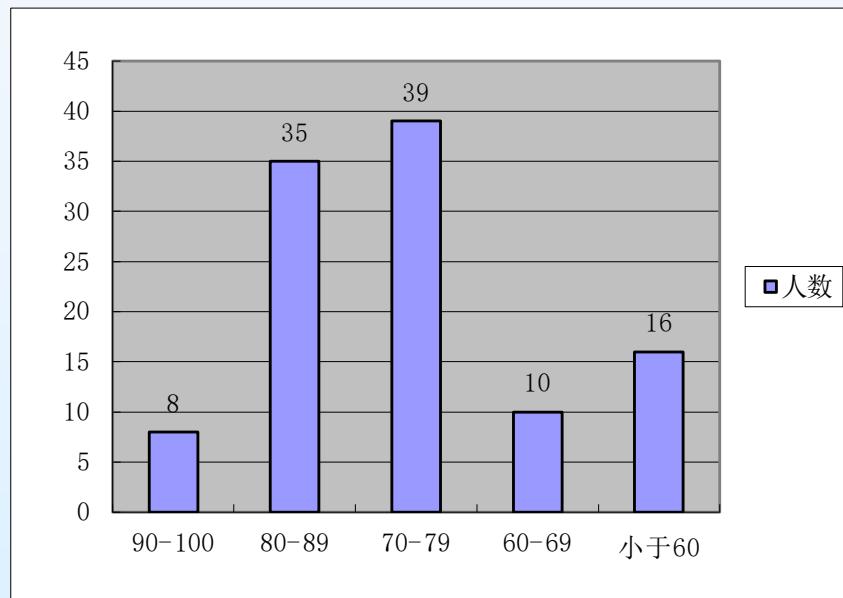
- Do a lot of readings before or after the lecture sessions, especially from the **English** textbook
- If you have any question, raise your hand!!!
- Do NOT refer to the so-called “standard” answers to the exercises, they contain many mistakes. Once you are caught using the wrong “standard” answers, penalty might be applicable.
- Hopefully more interactions during the lectures



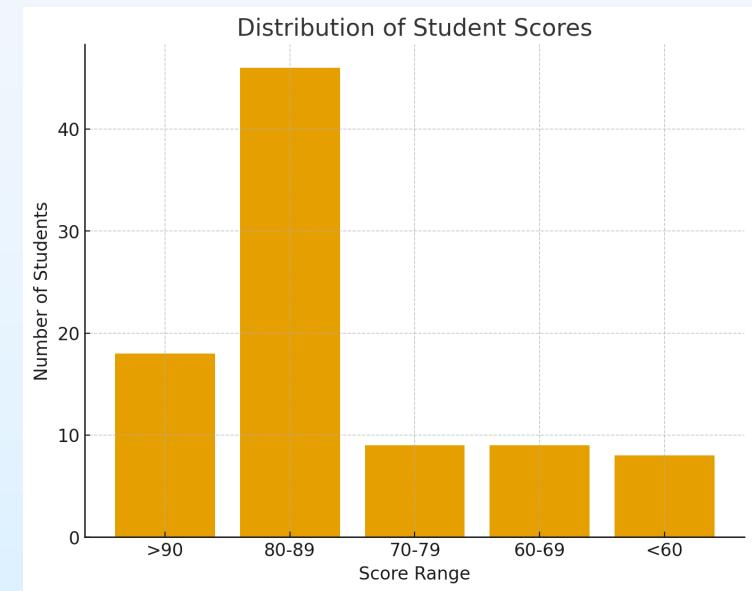


Previous Performance

Only the final exam

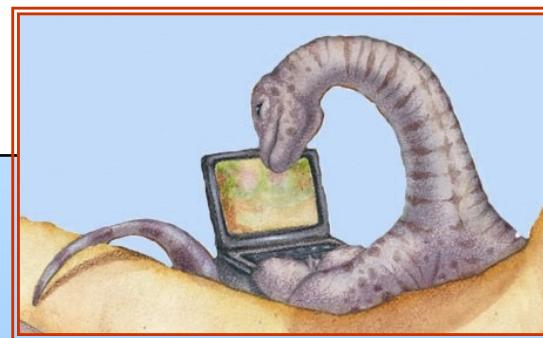


Combined with other grades



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Chapter 1 Introduction





Outlines for Chapter ONE

- What Operating Systems Do
- Computer-System Organization
- Computer-System Architecture

Overview

- Operating-System Structure
- Operating-System Operations

Structure

- Process Management
- Memory Management
- Storage Management

Main parts

- Protection and Security
- Distributed Systems
- Special-Purpose Systems
- Computing Environments





Objectives

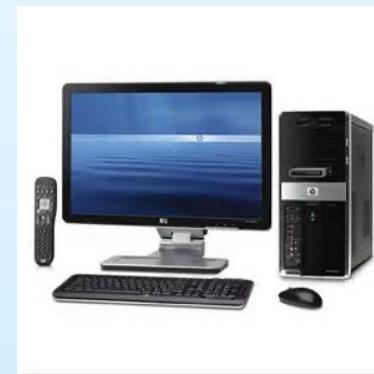
- To provide a grand tour of the major operating systems components
- To provide coverage of basic computer system organization





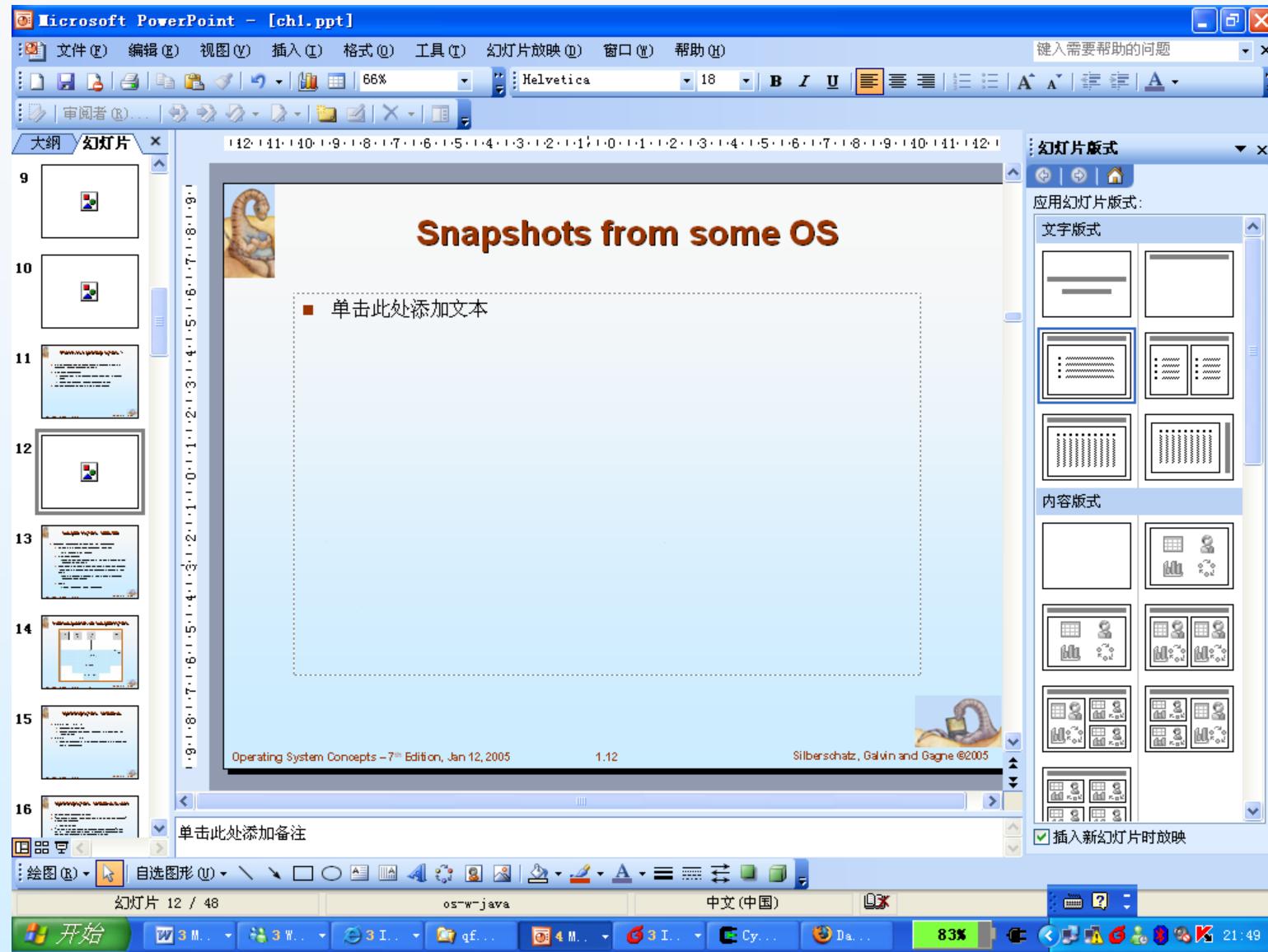
What is an Operating System?

- A program that acts as an intermediary between a user of a computer and the computer hardware.
- Operating system goals:
 - Execute user programs and make solving user problems easier.
 - Make the computer system convenient to use.
- Use the computer hardware in an efficient manner.





Snapshots from some OS



The screenshot shows a Microsoft PowerPoint presentation window. The title slide has the heading "Snapshots from some OS" and a placeholder "单击此处添加文本". The status bar at the bottom indicates the slide is 12 of 48, the file is "os-w-java", and the date is Jan 12, 2005. A note in the status bar also mentions "Silberschatz, Galvin and Gagne ©2005". The ribbon menu is visible, and the "幻灯片版式" (Slide Layout) pane is open, showing various layout options. The taskbar at the bottom shows multiple open windows, including a browser with several tabs.



开始

蔡耀明 

张婉茹
答复：今晚一起吃饭好吗？好呀！就去那家后面带露台的新餐厅吧？

邮件 8

足球训练
体育场
下午 5:30 - 下午 8:30

8 星期一

Internet Explorer

应用商店

必应 Bing



照片

地图

SkyDrive

游戏



桌面

消息

23°C
北京
多云 17° / 24°

天气



热气球节本周末开幕

音乐





From my linux box

```
dbgpucluster.di.comp.nus.edu.sg - dbgpu* - SSH Secure Shell
File Edit View Window Help
Quick Connect Profiles

[wusai@dbgpucluster ~]$ 
[wusai@dbgpucluster ~]$ cd knnsearch/
[wusai@dbgpucluster knnsearch]$ make
make: Nothing to be done for `default'.
[wusai@dbgpucluster knnsearch]$ ls
clib_port.h  db_int.h          dbstl_dbt.h      dbstl_resource_manager.h HyperplaneLSH.h IndexDB.o    main.cpp
data.db       dbstl_base_iterator.h dbstl_element_ref.h dbstl_set.h      HyperplaneLSH.o  knnSearch   main.o
db_config.h   dbstl_common.h     dbstl_exception.h dbstl_utility.h   index.db        lib        makefile
db_cxx.h      dbstl_container.h  dbstl_inner_utility.h dbstl_vector.h  IndexDB.cpp   libdb-6.1.so  siftgeo.bin
db.h         dbstl_dbc.h        dbstl_map.h      HyperplaneLSH.cpp IndexDB.h     libdb_cxx-6.1.so sqlite3.h
[wusai@dbgpucluster knnsearch]$ make clean
rm -f knnSearch *.o *
[wusai@dbgpucluster knnsearch]$ make
g++ -g -Wall -c main.cpp
main.cpp: In function `int main()'
main.cpp:22:37: warning: deprecated conversion from string constant to `char*'?[-Wwrite-strings]
  indexer.importAsLSH("./siftgeo.bin");
                           ^
main.cpp:42:26: warning: comparison between signed and unsigned integer expressions [-Wsign-compare]
  for(int i=0; i<val.size(); i++)
                  ^
g++ -g -Wall -c IndexDB.cpp
IndexDB.cpp: In member function `std::vector<int> IndexDB::getApproximateKNNByLSH(int*, int)'
IndexDB.cpp:209:28: warning: comparison between signed and unsigned integer expressions [-Wsign-compare]
  for(int j=0; j<next.size(); j++)
                  ^
IndexDB.cpp:216:19: warning: comparison between signed and unsigned integer expressions [-Wsign-compare]
  while(all.size()<K)
                  ^
IndexDB.cpp:222:29: warning: comparison between signed and unsigned integer expressions [-Wsign-compare]
  for(int j=0; j<next.size(); j++)
                  ^
IndexDB.cpp:229:29: warning: comparison between signed and unsigned integer expressions [-Wsign-compare]
  for(int j=0; j<next.size(); j++)
                  ^
IndexDB.cpp:260:17: warning: comparison between signed and unsigned integer expressions [-Wsign-compare]
  if(knn.size()<K)
                  ^
g++ -g -Wall -c HyperplaneLSH.cpp
g++ -g -Wall -L./lib -ldb-6.1 -ldb_cxx-6.1 -lpthread main.o IndexDB.o HyperplaneLSH.o -o knnSearch
[wusai@dbgpucluster knnsearch]$ 
```



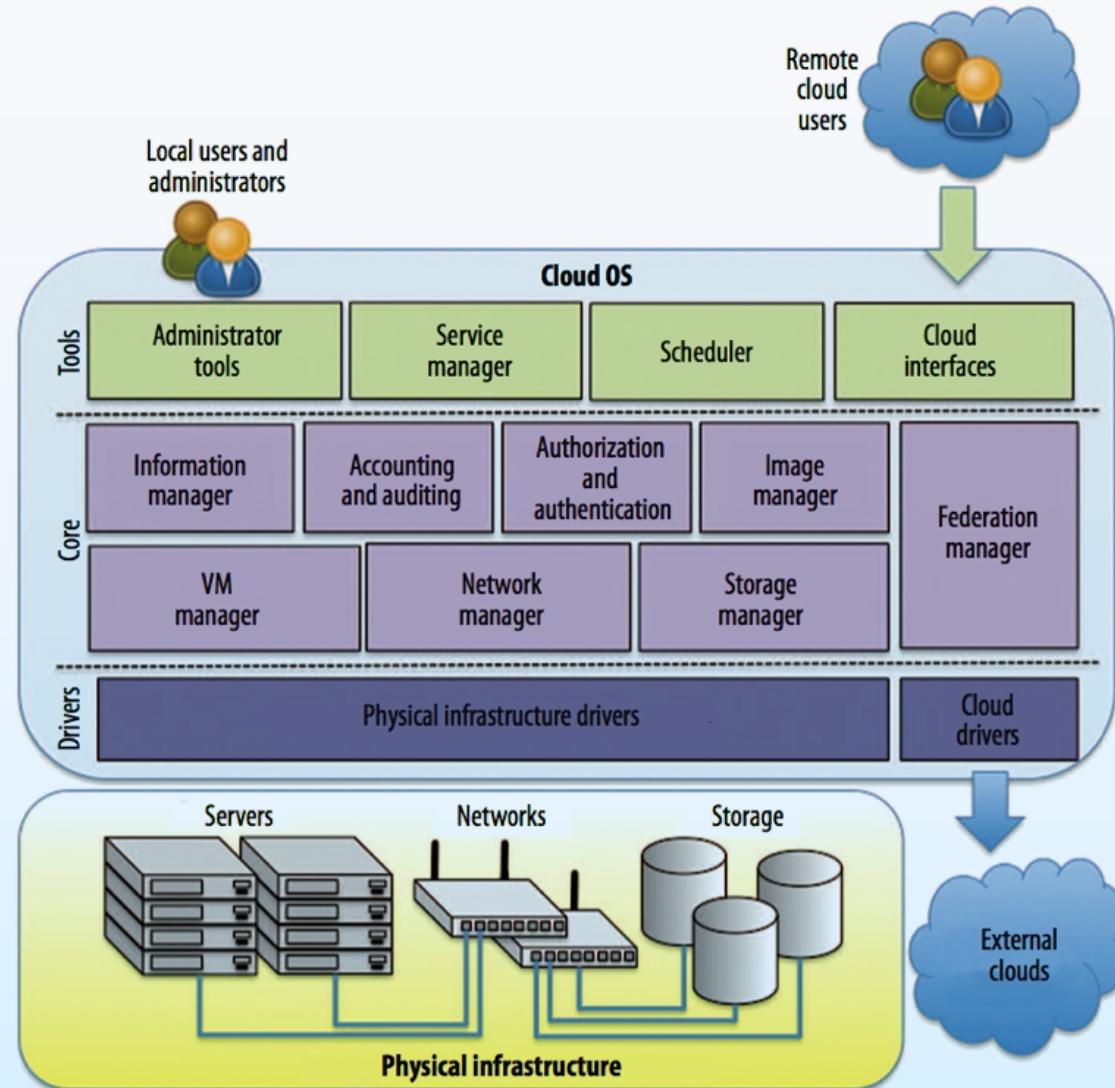


Smartphones & Tablets



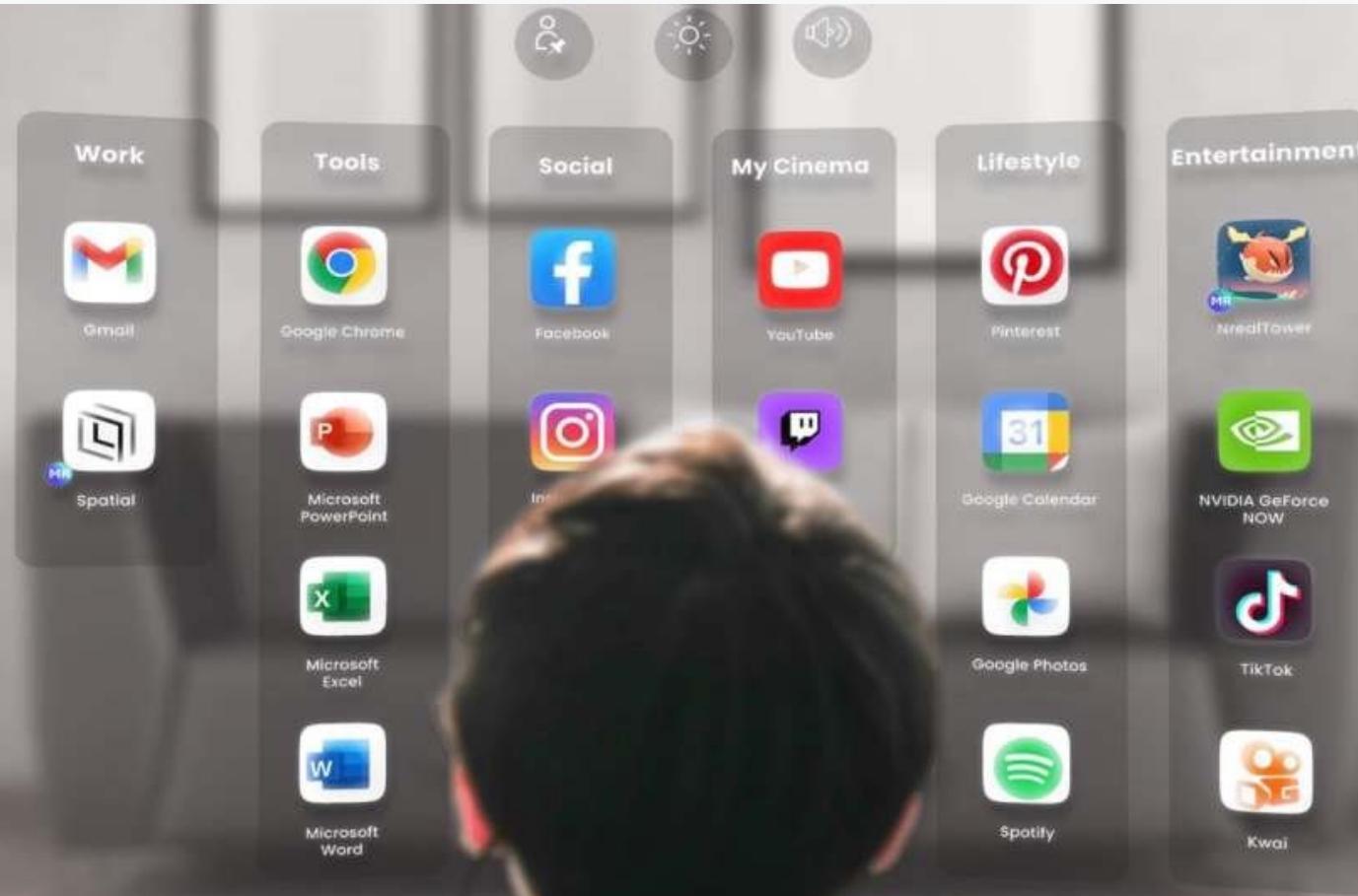


Cloud Operating System





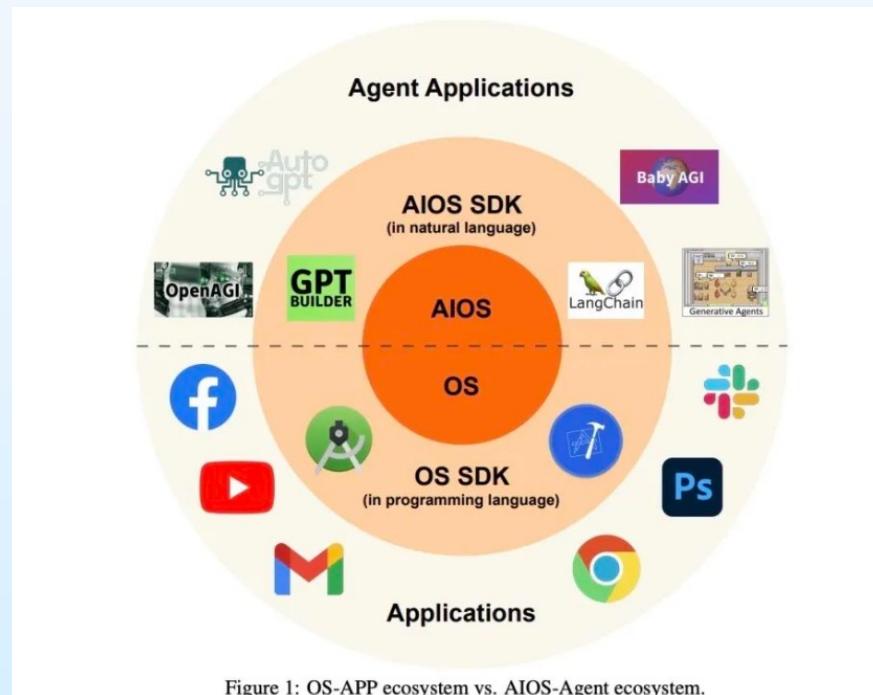
Metaverse Operating System





AI Operating System

- Merging LLM with OS turns the operating system into a conversational, adaptive, and self-optimizing platform. Instead of being a passive resource manager, the OS becomes an active collaborator





It used to look like this ...





Firmware vs Operating System

- Firmware: low-level code, stored in bios/cd
- OS: a system-level software





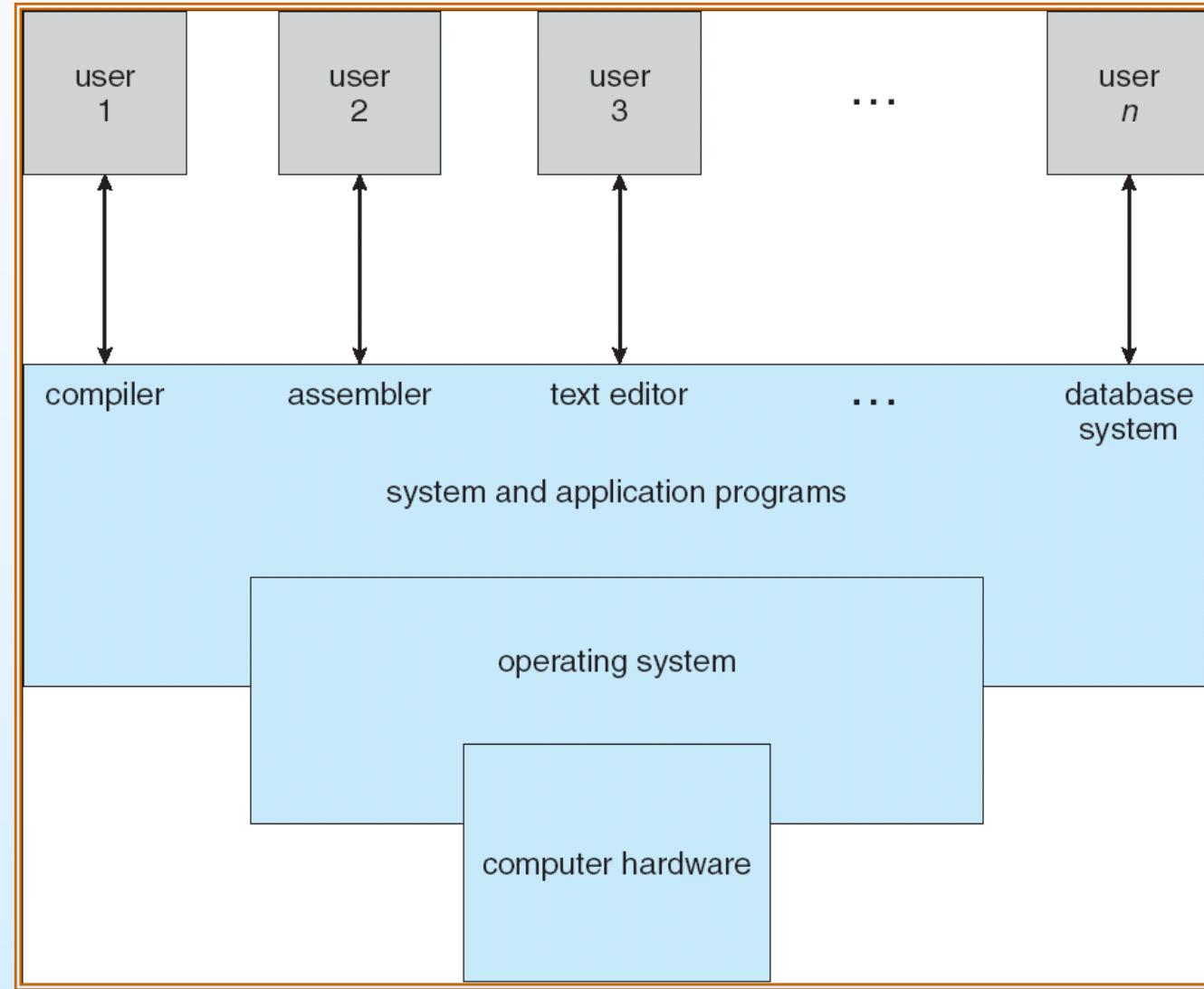
Computer System Structure

- Computer system can be divided into four components
 - Hardware – provides basic computing resources
 - ▶ CPU, memory, I/O devices
 - Operating system
 - ▶ Controls and coordinates use of hardware among various applications and users
 - System & application programs – define the ways in which the system resources are used to solve the computing problems of the users
 - ▶ Word processors, compilers, web browsers, database systems, video games
 - Users
 - ▶ People, machines, other computers





Four Components of a Computer System



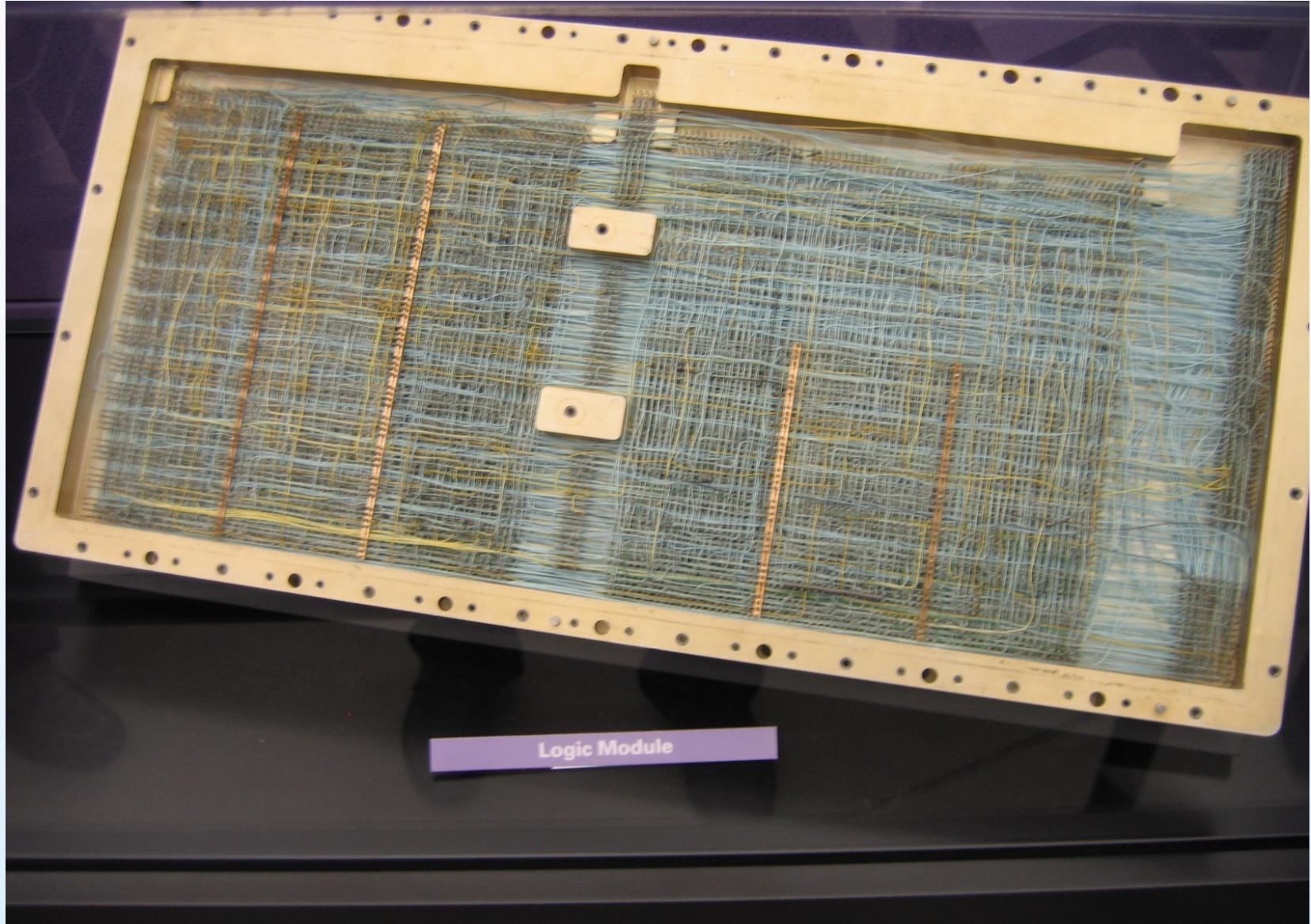


How hard each component has evolved!



The WISC computer developed by Univ of Wisconsin.





The Logic Module used by Apollo guidance computer, putting a Man on the moon with a computer that had 36KB of memory





My very first computer: 386

- A 386 running a win 32 system
- Mainly used for typing practice and gaming...





Operating System Definition

■ OS is a **resource allocator**

- Manages all resources
- Decides between conflicting requests for efficient and fair resource use

■ OS is a **control program**

- Controls execution of programs to prevent errors and improper use of the computer





Operating System Definition (Cont.)

- No universally accepted definition
- “Everything a vendor ships when you order an operating system” is good approximation
 - But varies wildly
- “The one program running at all times on the computer” is the **kernel**. Everything else is either a system program (ships with the operating system) or an application program





Operating System Definition (Cont.)

■ From deepseek

什么是操作系统



当然！这是一个关于“操作系统”的详细解释，从简单定义到深入功能，希望能让你全面理解。

一句话概括

操作系统（Operating System，简称 OS）是计算机中最基本、最核心的系统软件，它是硬件和应用程序之间的桥梁，负责管理计算机的硬件和软件资源，并为用户提供一个方便、高效、安全的操作环境。

你可以把它想象成一个公司的总经理或交通系统的指挥中心：

- 硬件（电脑的CPU、内存、硬盘等）就像公司的员工和资源（桌椅、打印机）、或者道路上的车辆和信号灯。
- 应用程序（浏览器、游戏、Word等）就像公司里各个部门的项目组，或者想要去不同目的地的司机。
- 操作系统（总经理/指挥中心）不直接做具体项目（不直接算数学题、画图），但它负责分配谁用什么资源、什么时候用、用多少，确保整个系统有条不紊地高效运行，避免混乱和冲突。





Computer Startup

- **bootstrap program** is loaded at power-up or reboot
 - Typically stored in ROM or EPROM, generally known as **firmware**
 - Initializes all aspects of system
 - Loads operating system kernel and starts execution

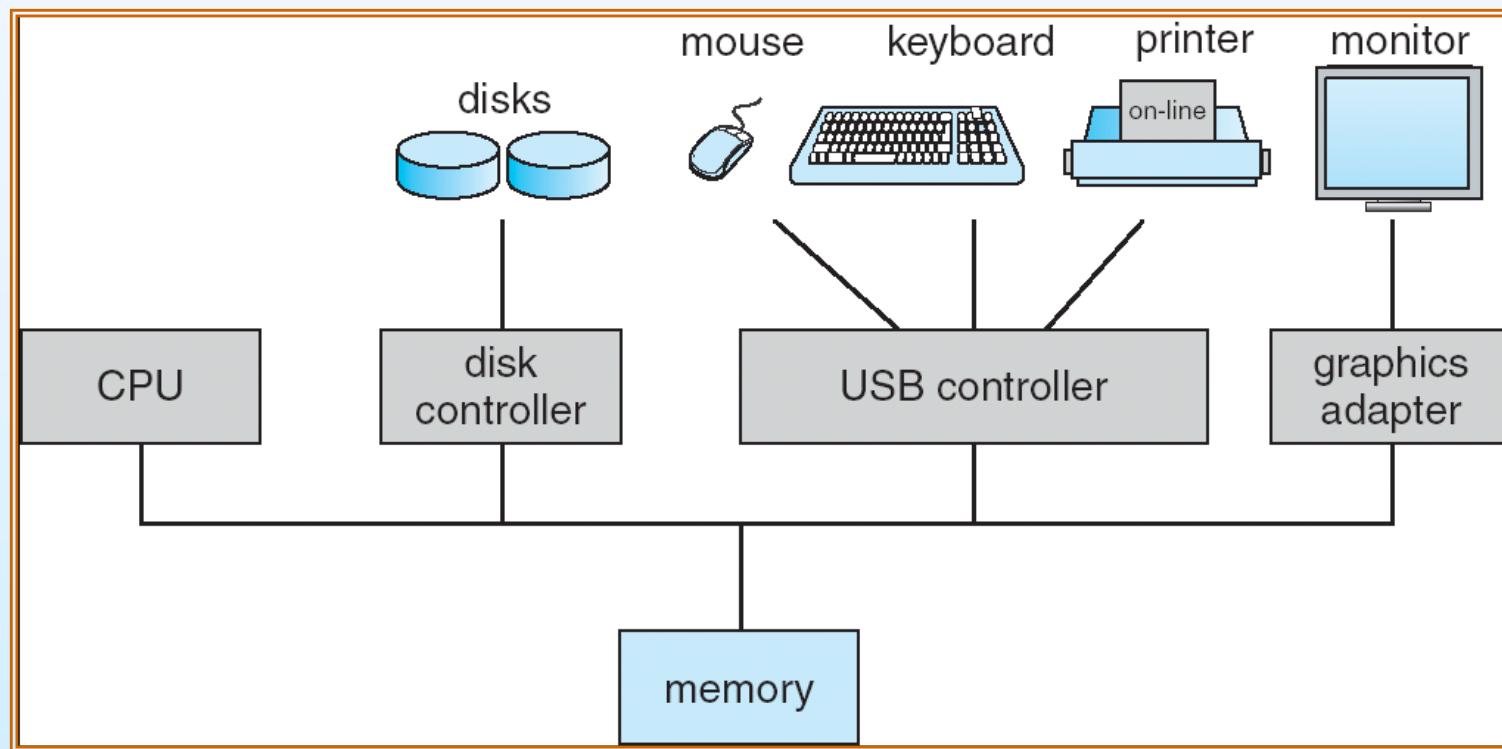




Computer System Organization

■ Computer-system operation

- One or more CPUs, device controllers connect through common bus providing access to shared memory
- Concurrent execution of CPUs and devices competing for memory cycles





Computer-System Operation

- I/O devices and the CPU can execute concurrently.
- Each device controller is in charge of a particular device type.
- Each device controller has a **local buffer**.
- CPU moves data from/to main memory to/from local buffers
- I/O is from the device to local buffer of controller.
- Device controller informs CPU that it has finished its operation by causing an *interrupt* (**via system bus**).





Common Functions of Interrupts

- Interrupt transfers control to the interrupt service routine generally, through the *interrupt vector*, which contains the addresses of all the service routines.
- Interrupt architecture must save the address of the interrupted instruction.
- Incoming interrupts are *disabled* while another interrupt is being processed to prevent a *lost interrupt*.
- A *trap* is a software-generated interrupt caused either by an *error* or a *user request* (the latter is often referred to as a *system call*).
- An operating system is *interrupt driven*.





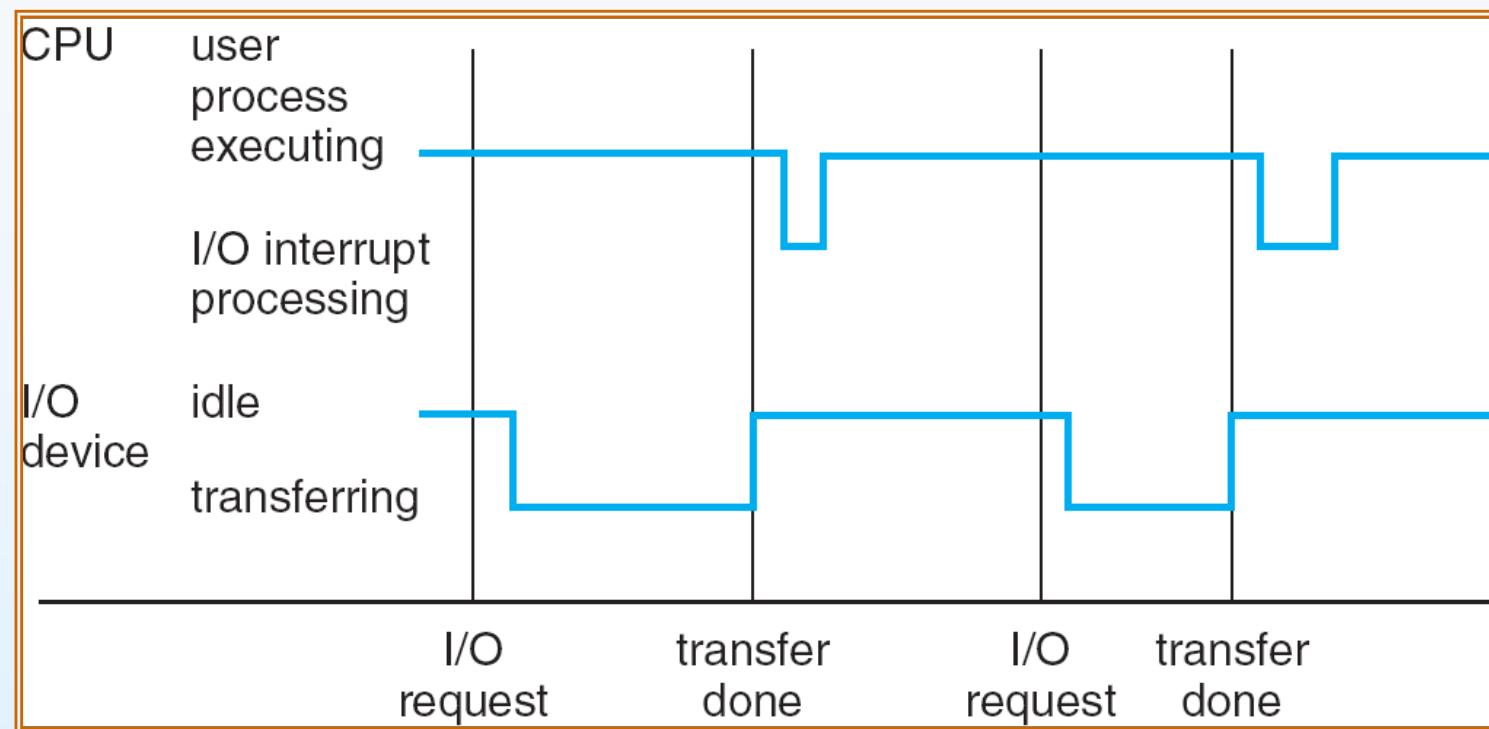
Interrupt Handling

- The operating system preserves the state of the CPU by storing **registers** and the **program counter**.
- Determines which type of interrupt has occurred:
 - *polling* by a generic routine
 - *vectored* interrupt system
- Separate segments of code determine what action should be taken for each type of interrupt





Interrupt Timeline





I/O Structure

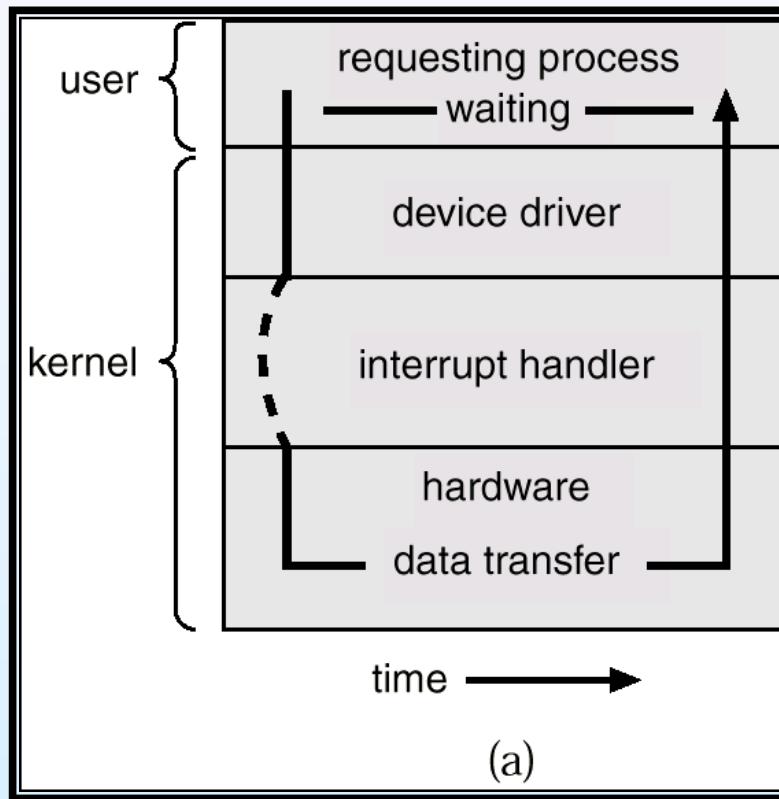
- After I/O starts, control returns to user program only upon I/O completion.
 - Wait instruction idles the CPU until the next interrupt
 - Wait loop (contention for memory access).
 - At most one I/O request is outstanding at a time, no simultaneous I/O processing.
- After I/O starts, control returns to user program without waiting for I/O completion.
 - *System call* – request to the operating system to allow user to wait for I/O completion.
 - *Device-status table* contains entry for each I/O device indicating its type, address, and state.
 - Operating system indexes into I/O device table to determine device status and to modify table entry to include interrupt.



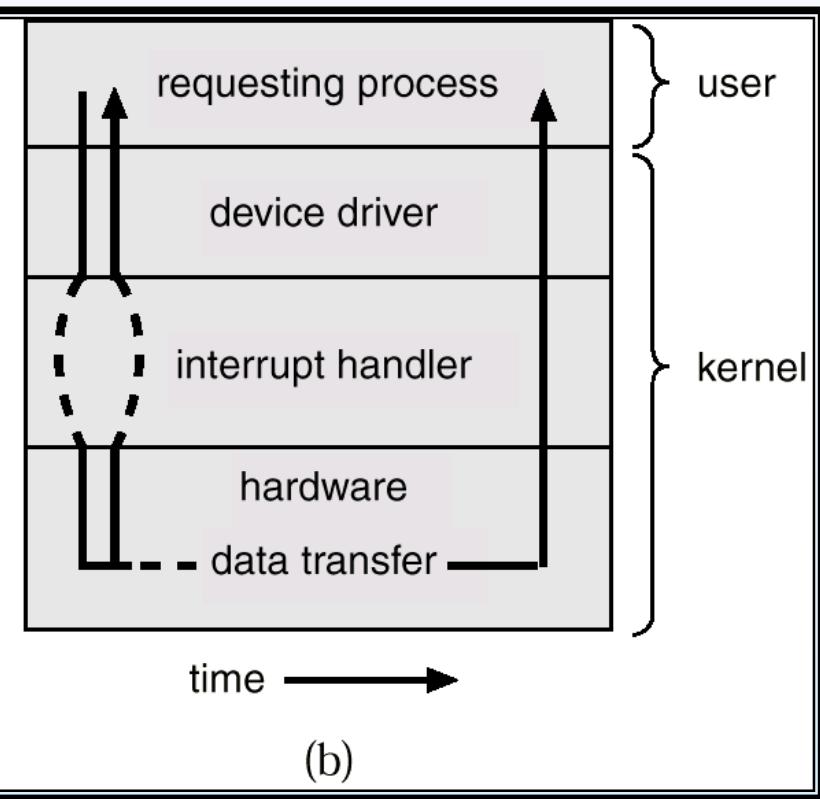


Two I/O Methods

Synchronous

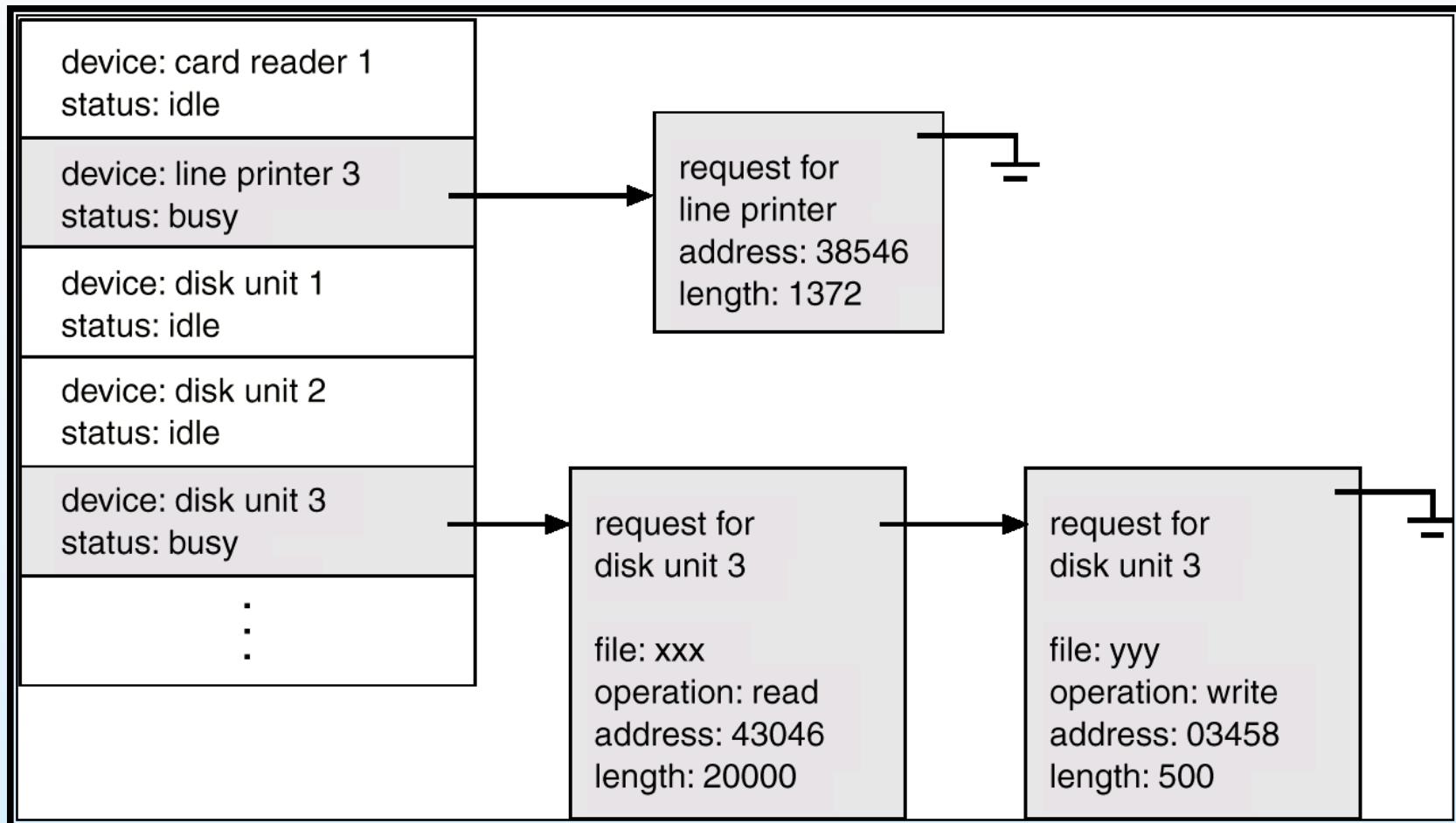


Asynchronous





Device-Status Table





Direct Memory Access Structure

- Used for high-speed I/O devices able to transmit information at close to memory speeds.
- Device controller transfers blocks of data from buffer storage directly to main memory **without** CPU intervention.
- Only one interrupt is generated **per block**, rather than the one interrupt per byte.





Storage Structure

- Main memory – only large storage media that the CPU can access directly.
- Secondary storage – extension of main memory that provides large nonvolatile storage capacity.
- Magnetic disks – rigid metal or glass platters covered with magnetic recording material
 - Disk surface is logically divided into *tracks*, which are subdivided into *sectors*.
 - The *disk controller* determines the logical interaction between the device and the computer.





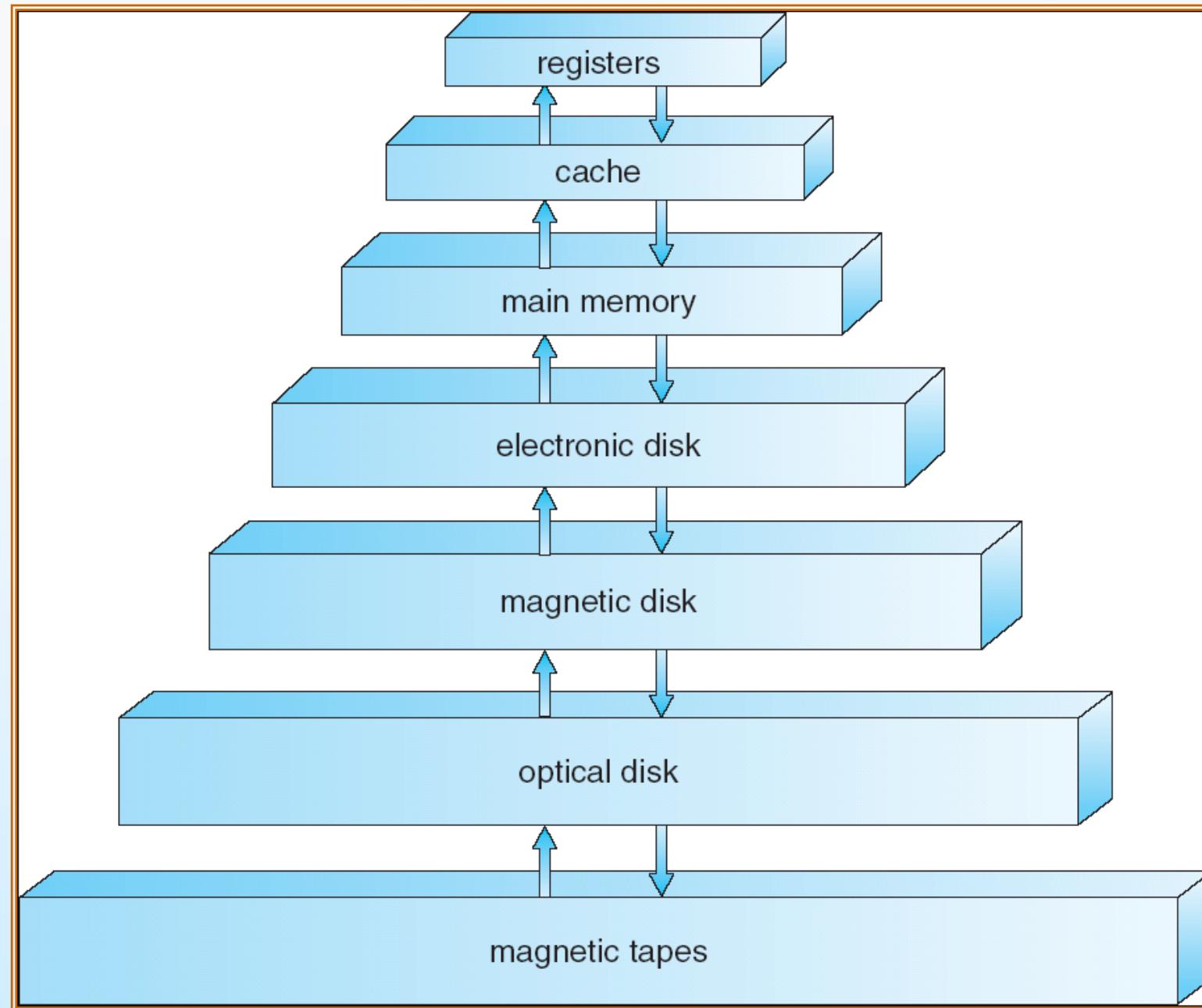
Storage Hierarchy

- Storage systems organized in hierarchy.
 - Speed
 - Cost
 - Volatility
- *Caching* – copying information into faster storage system; main memory can be viewed as a last *cache* for secondary storage.





Storage-Device Hierarchy





Caching

- Important principle, performed at many levels in a computer (in hardware, operating system, software)
- Information in use copied from slower to faster storage temporarily – **Speed mismatch**
- Faster storage (cache) checked first to determine if information is there
 - If it is, information used directly from the cache (fast)
 - If not, data copied to cache and used there
- Cache smaller than storage being cached
 - Cache management important design problem
 - Cache size and replacement policy





Performance of Various Levels of Storage

- Movement between levels of storage hierarchy can be explicit or implicit

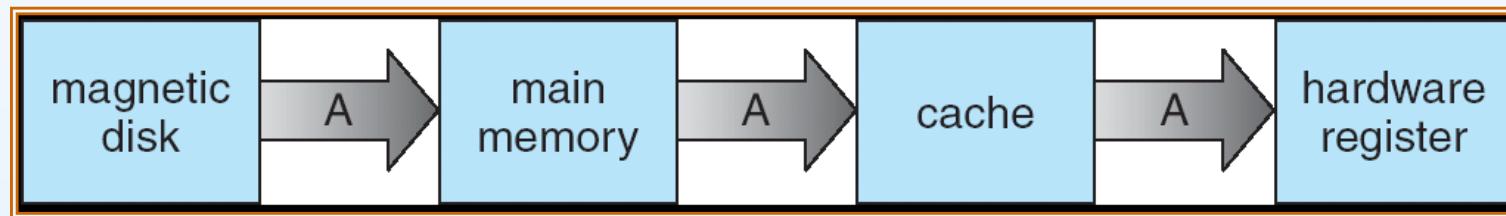
Level	1	2	3	4
Name	registers	cache	main memory	disk storage
Typical size	< 1 KB	> 16 MB	> 16 GB	> 100 GB
Implementation technology	custom memory with multiple ports, CMOS	on-chip or off-chip CMOS SRAM	CMOS DRAM	magnetic disk
Access time (ns)	0.25 – 0.5	0.5 – 25	80 – 250	5,000.000
Bandwidth (MB/sec)	20,000 – 100,000	5000 – 10,000	1000 – 5000	20 – 150
Managed by	compiler	hardware	operating system	operating system
Backed by	cache	main memory	disk	CD or tape





Migration of Integer A from Disk to Register

- **Multitasking** environments must be careful to use most recent value, no matter where it is stored in the storage hierarchy



- **Multiprocessor** environment must provide **cache coherency** in hardware such that all CPUs have the most recent value in their cache
- Distributed environment situation even more complex
 - Several copies of a datum can exist
 - Various solutions covered in Chapter 17





Operating System Structure

■ Multiprogramming needed for efficiency (**CPU utilization**)

- Single user cannot keep CPU and I/O devices busy at all times
- Multiprogramming organizes jobs (code and data) so CPU always has one to execute
- A subset of total jobs in system is kept in memory
- One job selected and run via **job scheduling**
- When it has to wait (for I/O for example), OS switches to another job

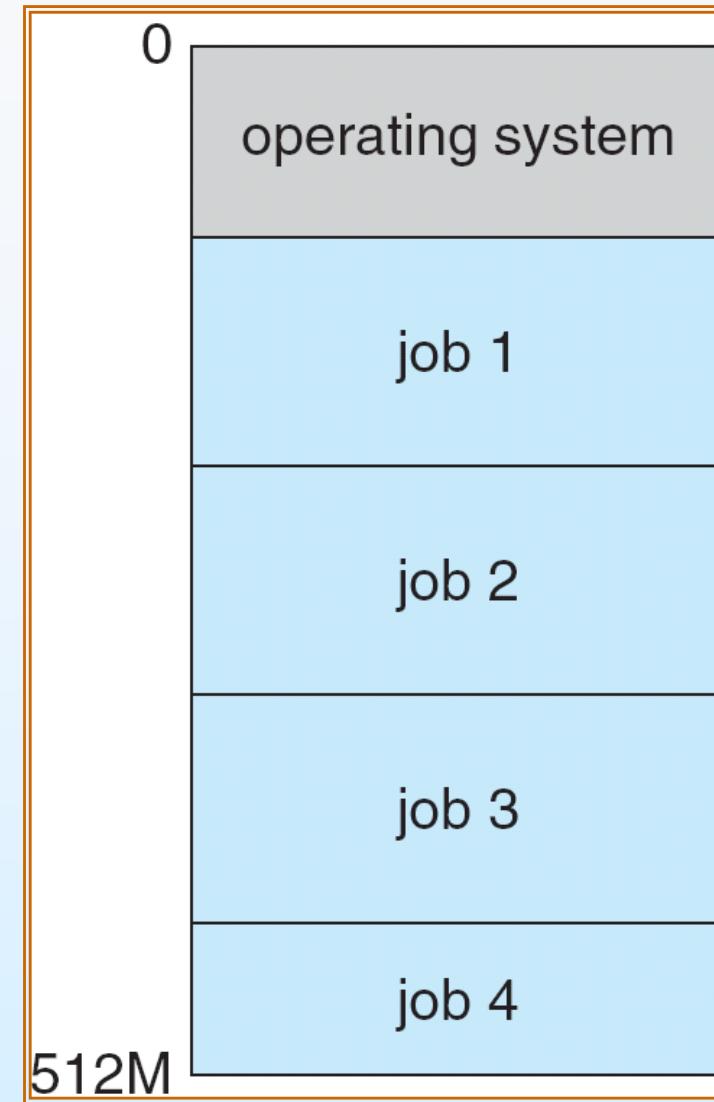
■ Timesharing (multitasking) is logical extension in which CPU switches jobs so frequently that users can interact with each job while it is running, creating **interactive** computing (**interactivity**)

- **Response time** should be < 1 second
- Each user has at least one program executing in memory
⇒ **process**
- If several jobs ready to run at the same time ⇒ **CPU scheduling**
- If processes don't fit in memory, **swapping** moves them in and out to run
- **Virtual memory** allows execution of processes not completely in memory





Memory Layout for Multiprogrammed System





Operating-System Operations

- Interrupt driven by hardware
- Software error or request creates **exception** or **trap**
 - Division by zero, request for operating system service
- Other process problems include **infinite loop**, processes **modifying** each other or the operating system

Thus we need protection:

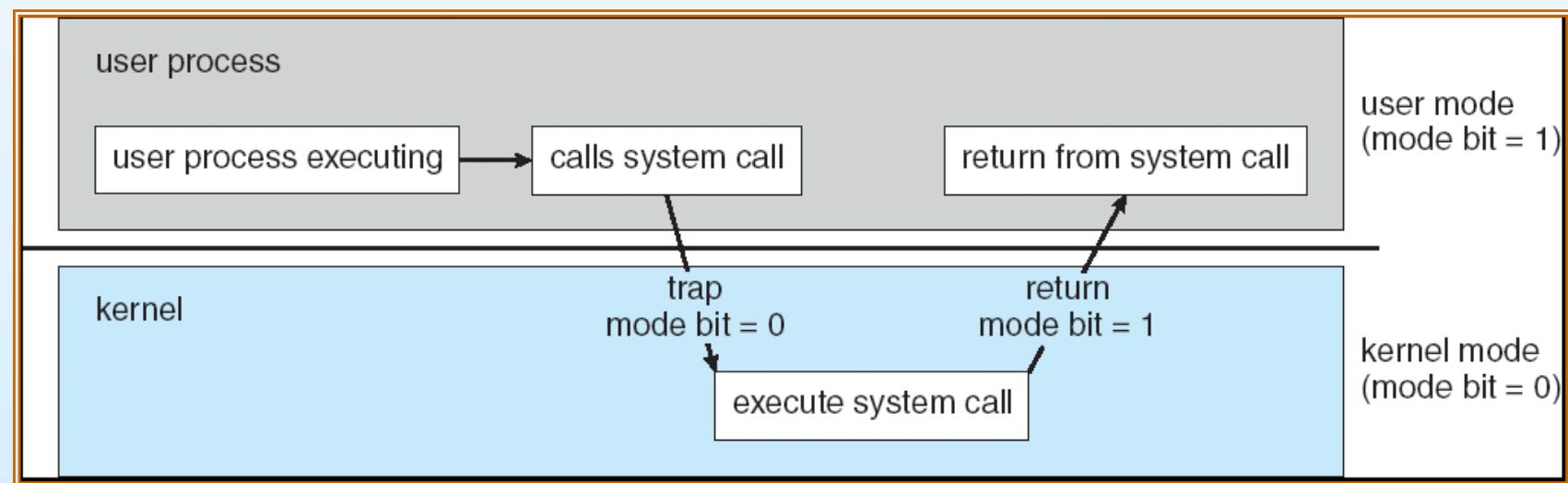
- **Dual-mode** operation allows OS to protect itself and other system components
 - **User mode** and **kernel mode**
 - **Mode bit** provided by hardware
 - ▶ Provides ability to distinguish when system is running user code or kernel code
 - ▶ Some instructions designated as **privileged**, only executable in kernel mode
 - ▶ System call changes mode to kernel, return from call resets it to user





Transition from User to Kernel Mode

- Timer to prevent infinite loop / process hogging resources
 - Set interrupt after specific period
 - Operating system decrements counter
 - When counter zero generate an interrupt
 - Set up before scheduling process to regain control or terminate program that exceeds allotted time





Process Management

- A process is a program in execution. It is a unit of work within the system. Program is a *passive entity*, process is an *active entity*.
- Process needs resources to accomplish its task
 - CPU, memory, I/O, files
 - Initialization data
- Process termination requires reclaim of any reusable resources
- Single-threaded process has one **program counter** specifying location of next instruction to execute
 - Process executes instructions sequentially, one at a time, until completion
- Multi-threaded process has one program counter per thread
- Typically system has many processes, some user, some operating system running concurrently on **one or more CPUs**
 - Concurrency by multiplexing the CPUs among the processes / threads





Process Management Activities

The operating system is responsible for the following activities in connection with process management:

- Creating and **deleting** both user and system processes
- Suspending and **resuming** processes
- Providing mechanisms for process **synchronization**
- Providing mechanisms for process **communication**
- Providing mechanisms for **deadlock** handling





Memory Management

- All **data** must be in memory before and after processing
- All **instructions** must be in memory in order to execute
- Memory management determines what is in memory when
 - Optimizing CPU utilization and computer response to users
- Memory management activities
 - Keeping track of which parts of memory are currently being used and by whom
 - Deciding which processes (or parts thereof) and data to move into and out of memory
 - Allocating and deallocating memory space as needed





Storage Management

- OS provides uniform, logical view of information storage
 - Abstracts physical properties to logical storage unit - **file**
 - Each medium is controlled by device (i.e., disk drive, tape drive)
 - ▶ Varying properties include access speed, capacity, data-transfer rate, access method (sequential or random)
- File-System management
 - Files usually organized into **directories**
 - **Access control** on most systems to determine who can access what
 - OS activities include
 - ▶ Creating and deleting files and directories
 - ▶ Primitives to manipulate files and dirs
 - ▶ Mapping files onto secondary storage
 - ▶ Backup files onto stable (non-volatile) storage media





Mass-Storage Management

- Usually disks used to store data that does not fit in main memory or data that must be kept for a “long” period of time.
- Proper management is of central importance
- Entire **speed** of computer operation hinges on disk subsystem and its algorithms
- OS activities
 - Free-space management
 - Storage allocation
 - Disk scheduling
- Some storage need not be fast
 - Tertiary storage includes optical storage, magnetic tape
 - Still must be managed
 - Varies between WORM (write-once, read-many-times) and RW (read-write)





I/O Subsystem

- One purpose of OS is to **hide peculiarities** of hardware devices from the user – *ease of usage & programming*
- I/O subsystem responsible for
 - Memory management of I/O including **buffering** (storing data temporarily while it is being transferred), **caching** (storing parts of data in faster storage for performance), **spooling** (the overlapping of output of one job with input of other jobs)
 - General device-driver interface
 - Drivers for specific hardware devices





Protection and Security

- **Protection** – any mechanism for controlling access of processes or users to resources defined by the OS
- **Security** – defense of the system against internal and external attacks
 - Huge range, including denial-of-service, worms, viruses, identity theft, theft of service
- Systems generally first distinguish among users, to determine who can do what
 - **User identities** (**user IDs**, security IDs) include name and associated number, one per user
 - User ID then associated with all files, processes of that user to determine **access control**
 - **Group identifier** (**group ID**) allows set of users to be defined and controls managed, then also associated with each process, file
 - **Privilege escalation** allows user to change to effective ID with more rights





Computing Environments

■ Traditional computer

- Blurring over time
- Office environment
 - ▶ PCs connected to a network, terminals attached to mainframe or minicomputers providing batch and timesharing
 - ▶ Now **portals** allowing networked and remote systems access to same resources
- Home networks
 - ▶ Used to be single system, then modems
 - ▶ Now firewalled, networked





Computing Environments (Cont.)

Mainframe Computing
1960s



Mini Computing
1970s



Personal Computing
1980s



Desktop Internet Computing
1990s



Mobile Internet Computing
2000s

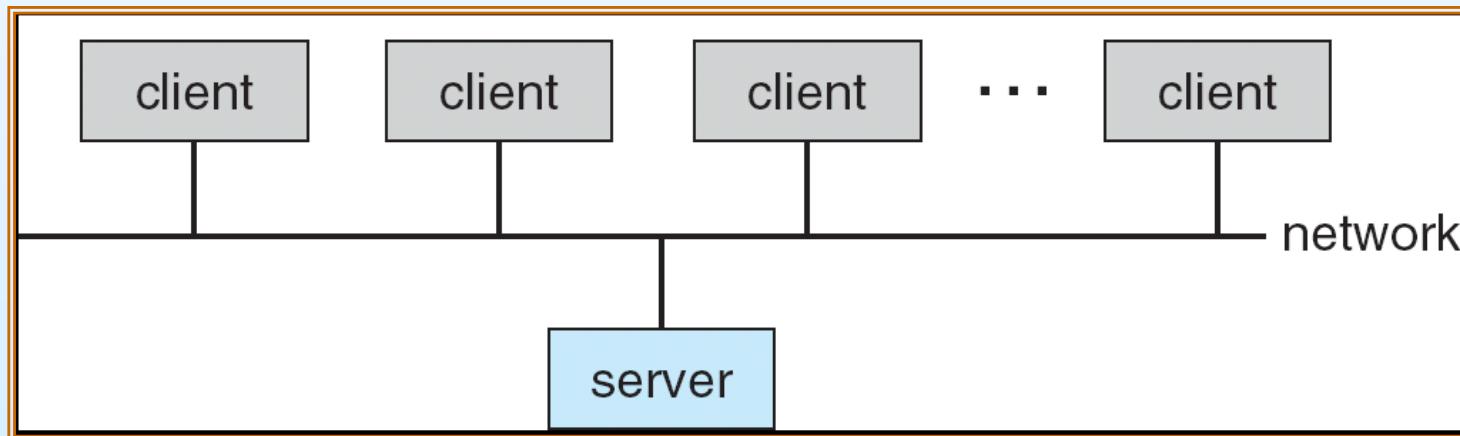




Computing Environments (Cont.)

■ Client-Server Computing

- Dumb terminals supplanted by smart PCs
- Many systems now **servers**, responding to requests generated by **clients**
 - ▶ **Compute-server** provides an interface to client to request services (i.e. database)
 - ▶ **File-server** provides interface for clients to store and retrieve files





Peer-to-Peer Computing

- Another model of distributed system
- P2P does not distinguish clients and servers
 - Instead all nodes are considered peers
 - May each act as client, server or both
 - Node must join P2P network
 - ▶ Registers its service with **central lookup service** on network, or
 - ▶ Broadcast request for service and respond to requests for service via ***discovery protocol***
 - Examples include *Napster* and *Gnutella*





Web-Based Computing

- Web has become ubiquitous
- PCs most prevalent devices
- More devices becoming networked to allow web access
- New category of devices to manage web traffic among similar servers: **load balancers**
- Use of operating systems like Windows 95, client-side, have evolved into Linux and Windows XP, which can be clients and servers
- Web based office software! <http://Docs.google.com>



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End of Chapter 1

