v1.20.4 < Course This course has already ended. The latest instance of the course can be found at: Concurrent Programming: 2023 **CS-E4110** Course materials « 3 Semaphore Blocking Queue -- Concurrent data structures based on lo... Course materials Your points CS-E4110 / Round 3 - Synchronization primitives / 4 Scala (JVM) Memory model and Data Races. / Submission 5 MyCourses Assignment description My submissions (5/5) \checkmark Scala (JVM) Memory model and Data Races. 1. Which are among the reasons a thread might not immediately see results of an operation in another thread? 1/1 A Compiler may reorder instructions otherwise than stated in the original source code. Out of thin air execution in JVM. Processors may process instructions in parallel or out of order. Caches may vary the order in which writes happen to main memory. Threads may choose to store results of computation in their local cache before committing them later. ✓ Correct!

2. The JMM defines a partial ordering called happens-before on all actions within a program. Choose the ones that belong to this happens-before rule?

- An unlock on a monitor lock happens-before every subsequent lock on that same monitor lock.
- ☐ A write to a volatile field happens-before every subsequent read of any field.
- **☑** A call to Thread.start on a thread happens-before every action in the started thread.
- **✓** Any action in a thread happens-before any other thread detects the thread has terminated.
- **☐** If a happens-before b, and b happens-before c, then a happens-before c.
- **✓** Correct!
- 3. Which are possible results of executing the following program under JVM?? 0/1

```
//shared variables
int x = 0, y = 0;
int a = 0, b = 0;
```

```
// thread 1 (t1)
                     // thread 2 (t2)
a = 1;
                      b = 1;
x = b;
                      y = a;
```

```
// results
t1.start(); t2.start();
t1.join(); t2.join();
println("("+x+","+y+")");
```

- \square (1,0).
- ☑ (0,1).
- \square (1,1).
- **☑** (0,0).

Imagine all possible reorderings and interleavings.

- **X** Incorrect
- 4. Which ones are true regarding the following program? 2/2

```
//shared variables
   int x = 0, y = 0;
    // thread 1 (t1)
                          // thread 2 (t2)
     r1 = x;
                          r2 = y;
                         if (r2 != 0) x = 1;
     if (r1 != 0) y = 1;
\square JMM allows the line r1 = x; to see the write of x = 1.
```

- The line r1 = x; being able to see the write of x = 1 is happens-before consistent according to JMM.
- **☐** The program is correctly synchronized according to JMM.
- \square The line $r_2 = y$; being able to see the write of y = 1 is sequentially consistent according to JMM.

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✓ Correct!

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Exercise info

Assignment category

Multiple choice questionnaires

Your submissions

5/5

Deadline

Monday, 22 November 2021, 14:00

Late submission deadline

Monday, 29 November 2021, 14:00 (-30%)

Total number of submitters

43

Submission info

Submitted on

Sunday, 14 November 2021, 16:21:20

Status

Ready

Grade

8 / 10

Submitters

Binh Nguyen