



Course materials Your points

This course has already ended.

The latest instance of the course can be found at: Principles of Algorithmic Techniques: 2023 Autumn

« 11.1 Implementing greedy algorithms

Course materials

Lecture and Exercise Set 7 - Randomized Algorithms 2 »

Programming Exercise 3 - Set Cover / 11.2 The set cover problem

The set cover problem

Exercise: Set cover

This exercise asks you to implement an algorithm that computes a low-cost set cover. More precisely, the input consists of a nonempty family $\mathcal{F} = \{S_0, S_1, \dots, S_{m-1}\}$ of nonempty subsets $S_j \subseteq \{0, 1, \dots, n-1\}$ for $j=0,1,\ldots,m-1$. A set $R\subseteq\{0,1,\ldots,m-1\}$ is a *set cover* if for all $x\in\{0,1,\ldots,n-1\}$ there exists a $j \in R$ with $x \in S_j$. The cost of a set cover R is c(R) = |R|. The optimum cost OPT of a set cover is the minimum cost of a set cover, where the minimum is taken over all possible set covers; if no set cover exists, the optimum is undefined. The algorithm must either

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1. output a set cover R with cost c(R) \leq H_q \cdot \text{OPT}, where q = \max_{j=0,1,\ldots,m-1} |S_j| and
H_q = 1 + 2 + \ldots + \frac{1}{q}, or
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2. correctly assert that no set cover exists.

For example, an implementation of the greedy set cover algorithm suffices for this purpose.

Your task in this exercise is to complete the subroutine

```
void solver(int n, int m, const int *p, const int *f, int &k, int *r)
```

which should compute the size k and the elements r of a set cover as described in the previous paragraph from the given input consisting of positive integers n and m, as well as the arrays p and f, whose format is as follows.

The array f concatenates the sets $S_0, S_1, \ldots, S_{m-1}$. That is, writing $S_i[0], S_i[1], \ldots, S_i[|S_i|-1]$ for the $|S_i|$ distinct elements of the set S_i , we have

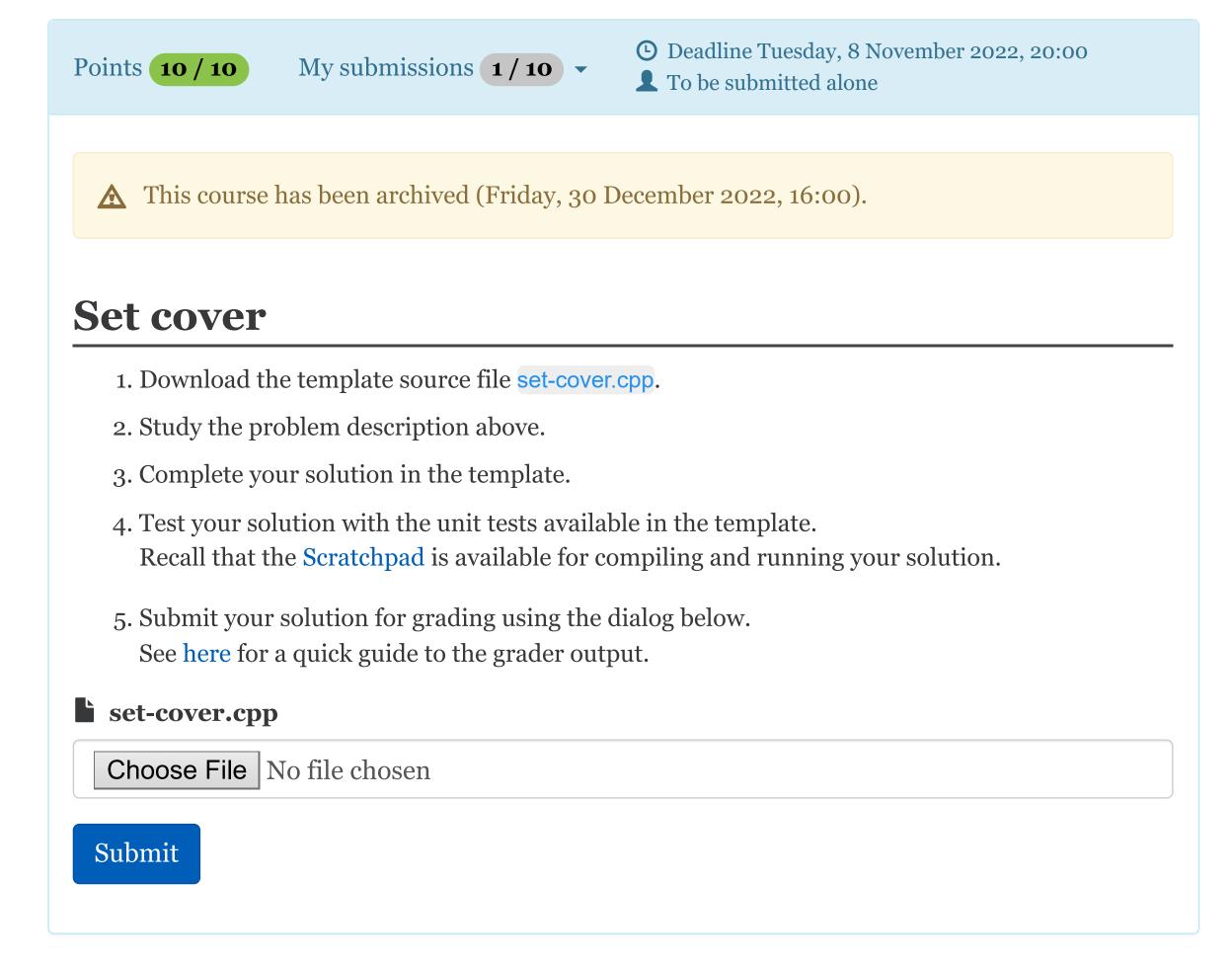
$$f = (S_0[0], S_0[1], \ldots, S_0[|S_0|-1], \ldots, S_{m-1}[0], S_{m-1}[1], \ldots, S_{m-1}[|S_{m-1}|-1]).$$

The array [p] satisfies $p[i] = \sum_{0 \le i \le i} |S_j|$ for all $i = 0, 1, \ldots, m$. Thus, the array [p] provides an index to the array f with $S_j = \{f[p[j]], f[p[j]+1], \dots, f[p[j+1]-1]\}$ for all $j = 0, 1, \dots, m-1$.

For example, when m = 2, $S_0 = \{4, 6, 8\}$, and $S_1 = \{7, 9\}$, we have f = (4, 6, 8, 7, 9) and p = (0, 3, 5).

The output of the subroutine should be as follows. To give as output a set cover $R = \{j_0, j_1, \dots, j_{k-1}\}$, set k equal to k and the element r[i] equal to j_i for all $i=0,1,\ldots,k-1$. When no set cover exists, set k equal to 0. You may assume that $1 \le n \le 1048576$, $0 \le m \le 2097152$, $0 \le p[m] \le 16777216$, $0 \le k \le n$, $q \leq 10$, and that the array r has capacity for at least n elements. To locate the subroutine quickly, you can search for "???" in the source file.

Grading. This exercise awards you up to 10 points in the course grading. The number of points awarded is the maximum points times the number of tests passed over the total number of tests, rounded up. To successfully complete a test, your implementation must use no more than 10 seconds of wall clock time and 1 GiB of memory. Each test will in general require the successful solution of one or more problem instances. In each batch of scaling tests, the first failed test will cause all subsequent tests in the batch to be skipped.



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