

$$CR_{y}(\theta) = \frac{10}{20} \frac{10}{10} = \frac{10}{10} \frac{10}{10} = \frac{10}{10$$

Exercise 8.20: (Circuit model for amplitude damping) Show that the circuit in

Figure 8.13 models the amplitude damping quantum operation, with  $\sin^2(\theta/2) = \omega$ 

$$S_{in} \mapsto S_{in} \otimes |o\rangle\langle o| = S_{in}$$
 $\mapsto S_{in} \otimes |o\rangle\langle o| = S_{in}$ 
 $\mapsto (|o\rangle\langle o| \otimes I + |1\rangle\langle 1| \otimes R_{in}(\Theta)) (S_{in} \otimes |o\rangle\langle o|) (|o\rangle\langle o| \otimes I + |1\rangle\langle 1| \otimes R_{in}(\Theta)^{t})$ 
 $= |o\rangle\langle o|S_{in}|o\rangle\langle o| \otimes |o\rangle\langle o| + |o\rangle\langle o|S_{in}|o\rangle\langle o|R_{in}(O)\rangle\langle o|R_{in}(O)^{t}$ 
 $+ |o\rangle\langle o|S_{in}|o\rangle\langle o| \otimes R_{in}(O) |o\rangle\langle o|R_{in}(O)^{t}$ 
 $+ |o\rangle\langle o|S_{in}|o\rangle\langle o| \otimes R_{in}(O) |o\rangle\langle o|R_{in}(O)^{t}$ 
 $+ |o\rangle\langle o|S_{in}(O)\langle o|S_{in}(O)\rangle\langle o|R_{in}(O)\rangle\langle o|R$ 

Exercise 8.23: (Amplitude damping of dual-rail qubits) Suppose that a single qubit state is represented by using two qubits, as

$$|\psi\rangle = a|01\rangle + b|10\rangle. \tag{8.113}$$

Show that  $\mathcal{E}_{AD} \otimes \mathcal{E}_{AD}$  applied to this state gives a process which can be described by the operation elements

$$E_0^{\rm dr} = \sqrt{1 - \gamma} I \tag{8.114}$$

$$E_1^{\text{dr}} = \sqrt{\gamma} \left[ |00\rangle\langle 01| + |00\rangle\langle 10| \right] , \qquad (8.115)$$

that is, either nothing  $(E_0^{dr})$  happens to the qubit, or the qubit is transformed  $(E_1^{dr})$  into the state  $|00\rangle$ , which is orthogonal to  $|\psi\rangle$ . This is a simple error-detection code, and is also the basis for the robustness of the 'dual-rail' qubit discussed in Section 7.4.

Exercise 8.30:  $(T_2 \le T_1/2)$  The  $T_2$  phase coherence relaxation rate is just the exponential decay rate of the off-diagonal elements in the qubit density matrix, while  $T_1$  is the decay rate of the diagonal elements (see Equation (7.144)). Amplitude damping has both nonzero  $T_1$  and  $T_2$  rates; show that for amplitude damping  $T_2 = T_1/2$ . Also show that if amplitude and phase damping are both applied then  $T_2 \le T_1/2$ .

$$\begin{bmatrix} a & b \\ b^* & 1-a \end{bmatrix} \rightarrow \begin{bmatrix} (a-a_0)e^{-t/T_1} + a_0 & be^{-t/T_2} \\ b^*e^{-t/T_2} & (a_0-a)e^{-t/T_1} + 1-a_0 \end{bmatrix}, (7.144)$$

## ibm\_washington Exploratory

## IBM Quantum Washington machine specs

Details

127

Status: • Online

Version: 1.1.0

Your usage: --

Total pending jobs: 140 jobs

64

Qubits

Processor type ①: Eagle r1

QV

Basis gates: CX, ID, RZ, SX, X

850

CLOPS

Avg. CNOT Error: 1.000e+0

Avg. Readout Error: 2.826e-2

Avg. T1: 101.21 us

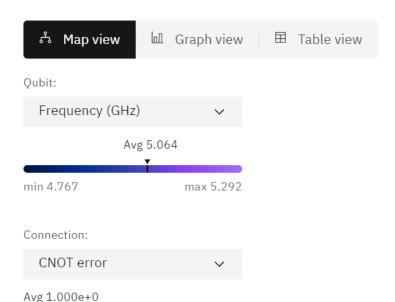
Avg. T2: 96.28 us

Providers with access: --

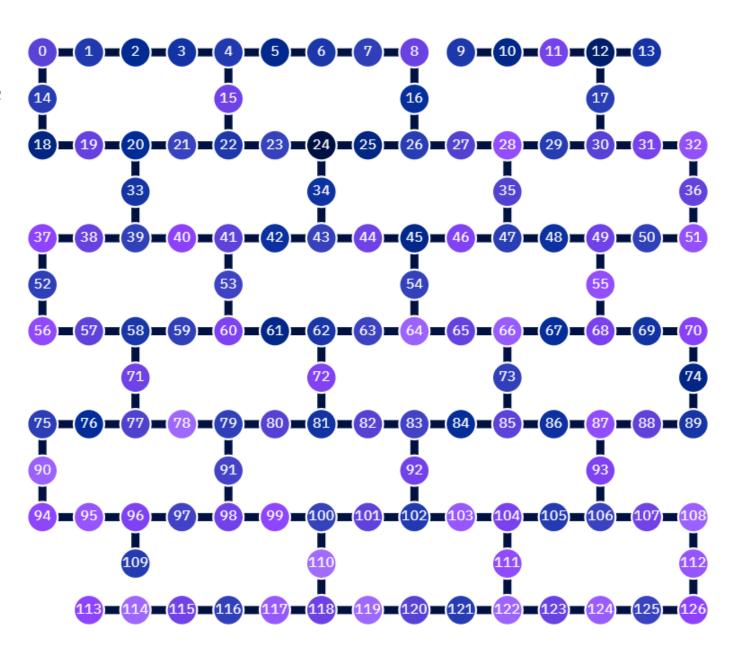
Supports Qiskit Runtime: Yes

## Calibration data

min 1.000e+0



max 1.000e+0



[1907.09528] Topological and subsystem codes on low-degree graphs with flag qubits (arxiv.org)

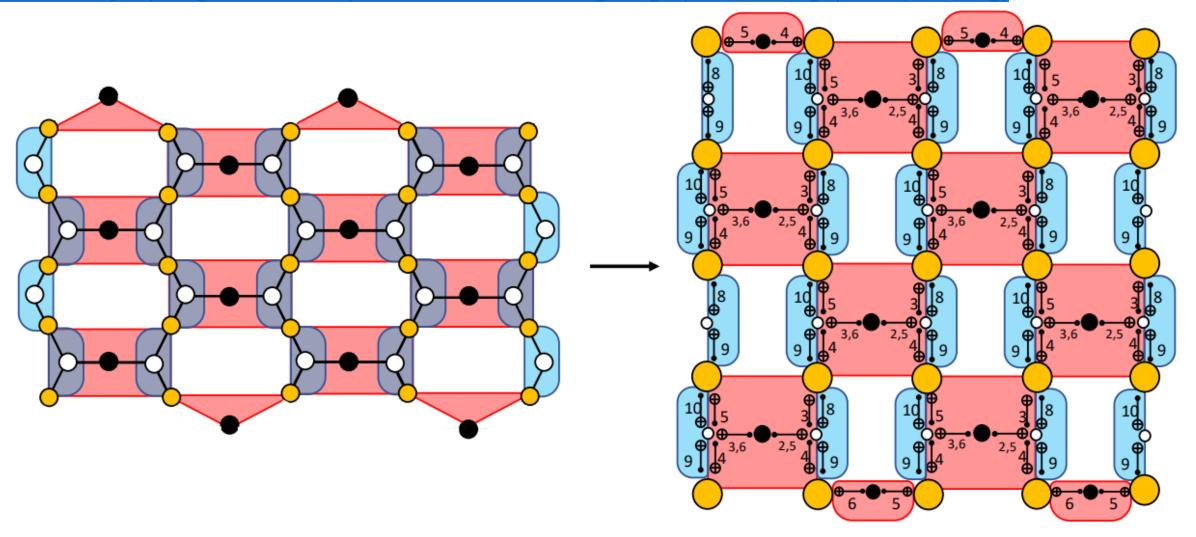


FIG. 2. The left of the figure corresponds to the actual layout of the d=5 heavy hexagon code which encodes one logical qubit. The data qubits are represented by yellow vertices, white vertices are the flag qubits and dark vertices represent the ancilla to measure the X-type gauge generators (red areas) and the Z-type gauge generators (blue areas). In the bulk, products of the two Z-type gauge generators at each white face forms a Z-type stabilizer. The right of the figure provides a circuit illustration of the heavy hexagon code with the scheduling of the CNOT gates used the measure the X-type and Z-type gauge generators.

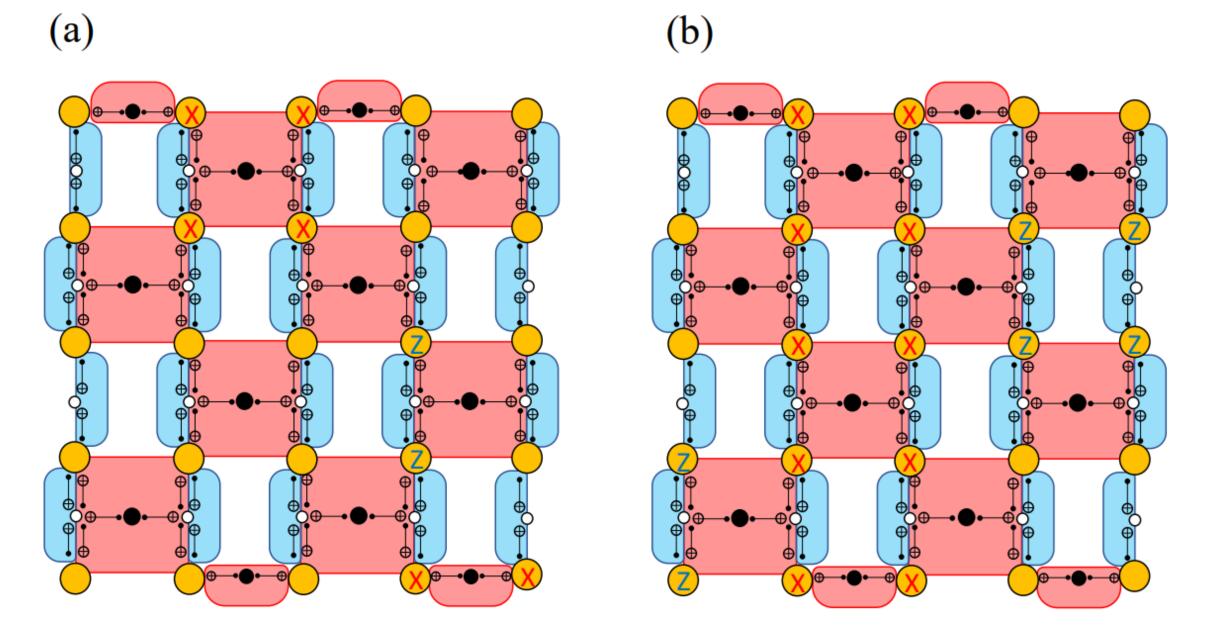


FIG. 3. (a) Gauge generators: weight-four X-type in the bulk, weight-two X-type on the upper and lower boundaries, and weight-two Z-type. (b) Stabilizer operators: a two-column vertical strip with X-type, weight-four Z-type in the bulk, and weight-two Z-type on the left and right boundaries.

The gauge group of the heavy hexagon code is

$$\mathcal{G} = \langle Z_{i,j} Z_{i+1,j}, \ X_{i,j} X_{i,j+1} X_{i+1,j} X_{i+1,j+1}, X_{i+1,j+1}, X_{i,2m-1} X_{1,2m}, \ X_{d,2m} X_{d,2m+1} \rangle$$

The stabilizer group which specifies the logical subspace  $\mathcal{H}_{\mathcal{L}}$  is the center of the gauge group or, explicitly,

$$S = \langle Z_{i,j} Z_{i,j+1} Z_{i+1,j} Z_{i+1,j+1}, Z_{2m,d} Z_{2m+1,d}, Z_{2m-1,1} Z_{2m,1}, \prod_{i} X_{i,j} X_{i,j+1} \rangle$$
(2)

In general, a distance d version of the code will have d data qubits along each row and each column of the hexagonal lattice so that the code parameters are given by  $[d^2, 1, d]$ . In addition, a distance d implementation of the code requires a total of  $\frac{d+1}{2}(d-1)$  syndrome measurement qubits and d(d-1) flag qubits. Hence the total number of qubits in the implementation of the code is  $\frac{5d^2-2d-1}{2}$ .

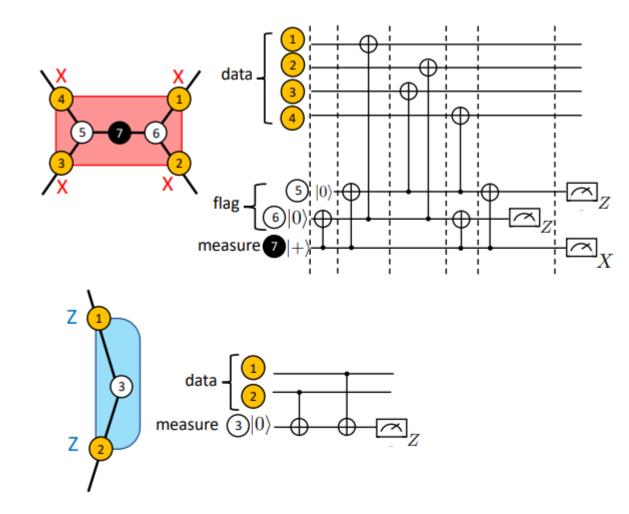
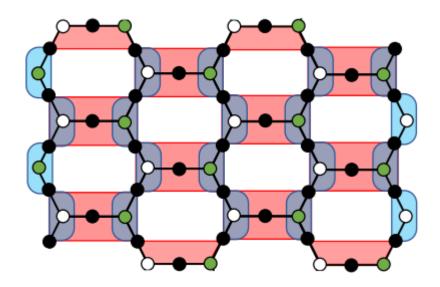
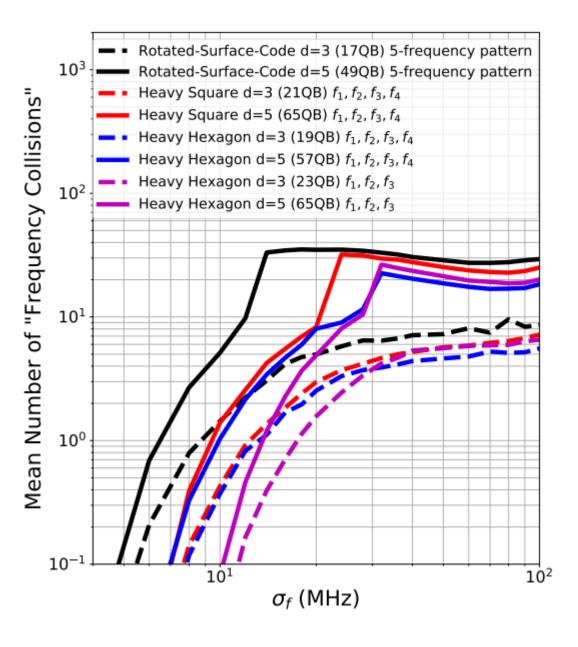


FIG. 4. Circuit to perform the X and Z-type parity measurements of the heavy hexagon code. Two flag qubits (white circles) are used to measure the weight-four X-type gauge generators.

control (f<sub>1</sub>)target (f<sub>2</sub>)target (f<sub>3</sub>)





## References

Quantum error correction intro: <u>arXiv:1907.11157</u>

Subsystem codes: <a href="mailto:arXiv:quant-ph/0508131">arXiv:quant-ph/0506023</a>

Surface codes: <u>arXiv:1208.0928</u>

Flag qubits: <u>arXiv:1705.02329</u>

Heavy hexagon code: arXiv:1907.09528