

Computational Methods and Applications

Scribe: II

Linear Systems and Interpolation

Equations of the form $Ax=b$, where

A is an $n \times n$ square matrix whose elements are a_{ij}
 x and b are column vectors of dimension n

$Ax=b$ can be written as:

$$\sum_{j=1}^n a_{ij} x_j = b_i \quad \forall i \in \{1, 2, \dots, n\}$$

b can also be interpreted as a linear combination of the

column matrix A weighted by vector x .

An example of pipe network was illustrated where we would write a system of linear eqn to compute the pressure at each node.
→ Direct numerical methods are not ideal for very large systems of linear equations

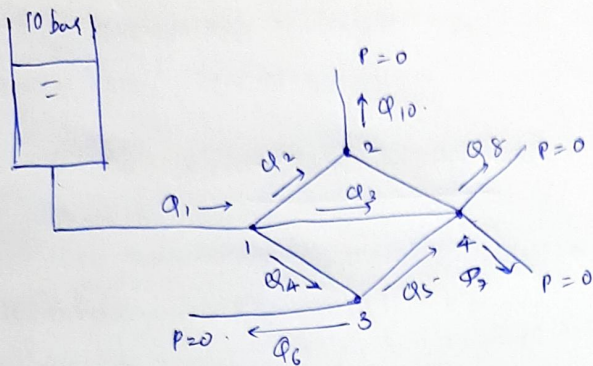
Iterative solution method

→ An iterative method for solution of linear ~~equation~~ system results in a sequence of vectors $\{x^{(k)}, k \geq 0\}$ of \mathbb{R}^n that converges to the exact solution x^* , that is

$$\lim_{k \rightarrow \infty} x^{(k)} = x^* \quad \text{for any given initial}$$

vector $x^{(0)} \in \mathbb{R}^n$.

Constructive



Res

$$Q_j = L_j \Delta P_j$$

Write a linear system of equations to compute the pressure at each node

$$Q_1 = 10 - P_1$$

$$Q_2 = P_1 - P_2$$

$$Q_3 = P_1 - P_4$$

$$Q_4 = P_1 - P_3$$

$$Q_5 = P_3 - P_4$$

$$Q_6 = P_3$$

$$Q_7 = P_4$$

$$Q_8 = P_4$$

$$Q_9 = P_2 - P_4$$

$$Q_{10} = P_2$$

$$10 - P_1 = P_1 - P_2 + P_1 - P_4 + P_1 - P_3 \quad \text{--- (1)}$$

$$P_1 - P_2 = P_2 + P_2 - P_4 \quad \text{--- (2)}$$

$$P_1 - P_3 = P_3 + P_3 - P_4 \quad \text{--- (3)}$$

$$P_1 - P_4 + P_2 - P_4 + P_3 - P_4 = 2P_4 \quad \text{--- (4)}$$

$$\text{from (1)} \Rightarrow 10 = 4P_1 - P_2 - P_4 - P_3$$

$$\text{(2)} \Rightarrow 0 = 3P_2 - P_4 - P_1$$

$$\text{(3)} \Rightarrow 0 = 3P_3 - P_4 - P_1$$

$$\text{(4)} \Rightarrow 0 = P_1 + P_2 + P_3 - 5P_4$$

Constructing iterative method

Select a suitable non singular matrix P , such that split matrix A

$$A = P - (P - A)$$

$$\text{Then } Px^* = b - (A - P)x^* \quad \text{--- (1)}$$

Correspondingly for $k \geq 0$

$$Px^{(k+1)} = b - (A - P)x^{(k)} \quad \text{--- (2)}$$

(2) - (1) gives.

$$x^{(k+1)} - x^* = (I - P^{-1}A)(x^{(k)} - x^*)$$

Convergence.

Let $e^{(k)} = x^{(k)} - x^*$ denote error at step k .

If $(I - P^{-1}A)$ is symmetric and positive definite (all eigenvalues positive)

$$\|e^{(k+1)}\|_2 = \|(I - P^{-1}A)e^{(k)}\|_2 \leq \rho(I - P^{-1}A)\|e^{(k)}\|_2$$

where $\rho(\cdot)$ is known as the spectral radius (maximum modulus of eigenvalues). If $\rho(\cdot) < 1$ there is convergence.

The Jacobi Method

If diagonal entries of A are non zero, we can set $P = D$ (diagonal matrix containing ~~entries~~ diagonal entries of A). Then we get

$$x_i^{(k+1)} = \frac{1}{a_{ii}} \left(b_i - \sum_{j=1, j \neq i}^n a_{ij} x_j^{(k)} \right) \quad \forall i \in \{1, 2, \dots, n\}$$

Proposition

If the matrix A is strictly diagonally dominant by row, then the Jacobi method converges

- It may converge otherwise also
- The dominant diagonal elements becomes the denominator and drives the iterations towards convergence.

The Gauss - Seidel method

Faster convergence could be achieved if the new $(k+1)$ components already available are used

$$x_i^{(k+1)} = \frac{1}{a_{ii}} \left(b_i - \sum_{j=1}^{i-1} a_{ij} x_j^{(k+1)} - \sum_{j=i+1}^n a_{ij} x_j^{(k)} \right)$$

$\forall i \in \{1, 2, \dots, n\}$

- There are no general results stating this method converges faster than Jacobi's

→ Python's `mumpy.linalg` module provide efficient low level implementations of standard linear algebra algorithms

Interpolation

In several applications we only know value of a function f at some given points $\{(x_i, y_i), i=0, 1, \dots, n\}$. How do we determine f ?

We figure out an approximate function \tilde{f} that satisfies

$$\tilde{f}(x_i) = y_i \quad \forall i \in \{0, 1, \dots, n\}$$

→ 1 cubic function exists passing through points $(0, 1)$, $(1, 4)$, $(1, 0)$ and $(2, 15)$

$$f(x) = x^3 + x^2 + x + 1$$

we got this solving

$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 2 & 4 & 8 \end{bmatrix} \begin{bmatrix} a \\ b \\ c \\ d \end{bmatrix} = \begin{bmatrix} 1 \\ 4 \\ 0 \\ 15 \end{bmatrix}$$

Complexity
 $O(n^3)$

Different kinds of interpolation.

polynomial interpolation

$$\tilde{f}(x) = a_0 + a_1x + a_2x^2 + \dots + a_nx^n.$$

trigonometric interpolation

$$\tilde{f}(x) = a_{-m}e^{-imx} + \dots + a_0 + \dots + a_me^{imx}$$

rational interpolation

$$\tilde{f}(x) = \frac{a_0 + a_1x + \dots + a_kx^k}{b_0 + b_1x + \dots + b_mx^m}.$$

Lagrangian polynomial interpolation

for $J = \{0, 1, \dots, n\}$, define

$$\psi_j(x) = \prod_{\substack{i=0, i \neq j}}^n \frac{x - x_i}{x_j - x_i}$$

Note that

$$\psi_j(x_k) = \begin{cases} \prod_{i=0, i \neq j}^n \frac{x_k - x_i}{x_j - x_i} = 1 & \text{if } k=j \\ 0 & \text{otherwise} \end{cases}$$

then required approximation is

$$\tilde{f}(x) = \sum_{j=0}^n y_j \psi_j(x)$$

eg: Function passing through the points $(0,1)$, $(1,4)$, $(-1,0)$ and $(2,15)$

$$\psi_0(x) = \frac{x^3 - 2x^2 - x + 2}{2}$$

$$\psi_1(x) = \frac{-x^3 + x^2 + 2x}{2}$$

$$\psi_2(x) = \frac{-x^3 + 3x^2 - 2x}{6}$$

$$\psi_3(x) = \frac{x^3 - x}{6}$$

$$\begin{aligned} \tilde{f}(x) &= \psi_0(x) + 4\psi_1(x) + 15\psi_2(x) \\ &= x^3 + x^2 + x + 1 \end{aligned}$$