Notes Made By

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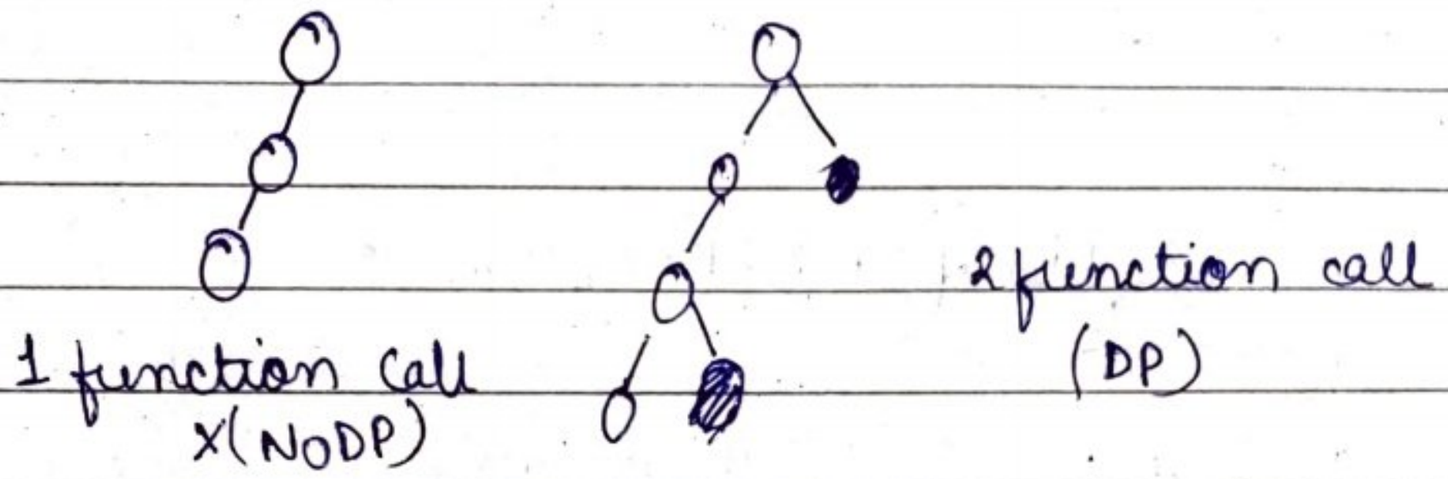
Dynamic Playlist (Aditya Verma)

DP = enhanced recursion

How to identify DP problem → (2 calls)

1) Where there is recursion, DP is used (for overlapping problem)
   a) Choice



   b) Optimal - Min, max, largest

How to write DP code?

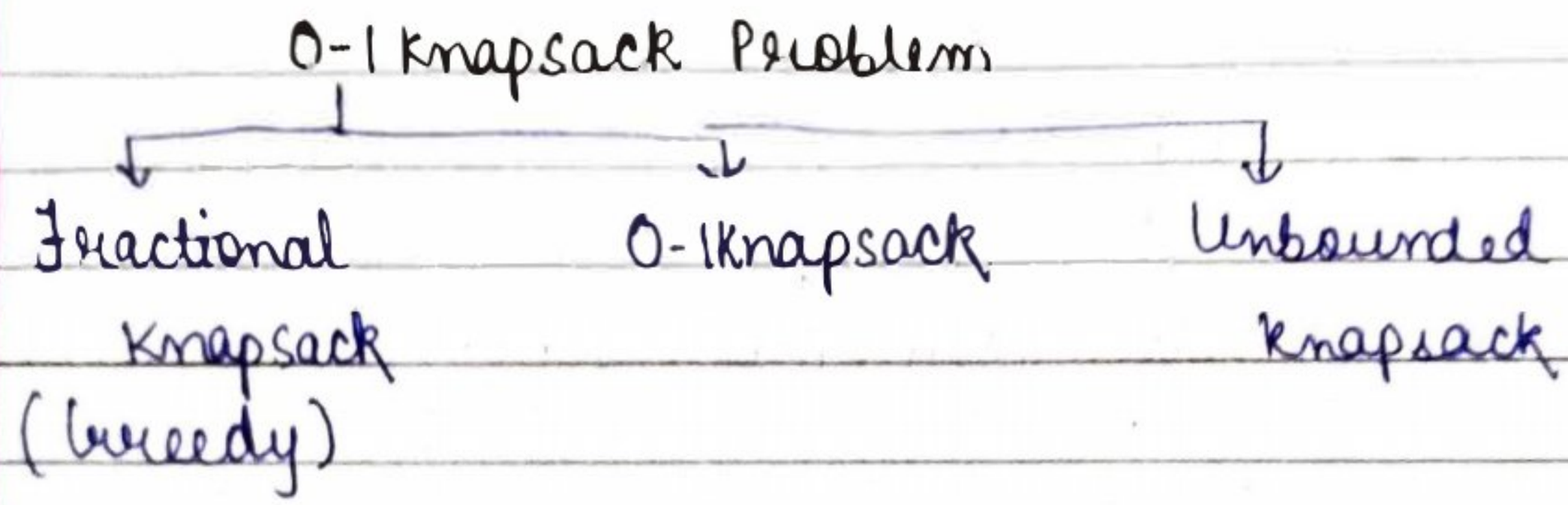
Recursive → memoization → Top down approach
solution

Questions on DP.

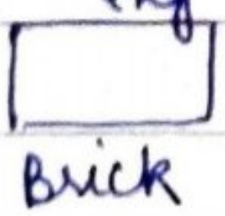
1) 0-1 Knapsack (6)
2) Unbounded knapsack (5)
3) Fibonacci (7)
4) LCS (15) (longest common subsequence)
5) LIS (10) (longest increasing subsequence)
6) Kadane's Algorithm (6)
7) Matrix chain Multiplication (7)
8) DP on trees (4)
9) DP on grid (14)
10) Others (5)

Types of knapsack

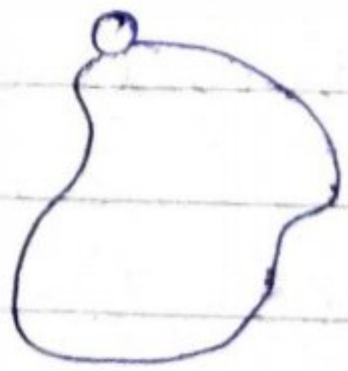
1. Subset sum
2. Equal sum partition
3. Count of subset sum
4. Minimum subset diff
5. Target sum.
6.



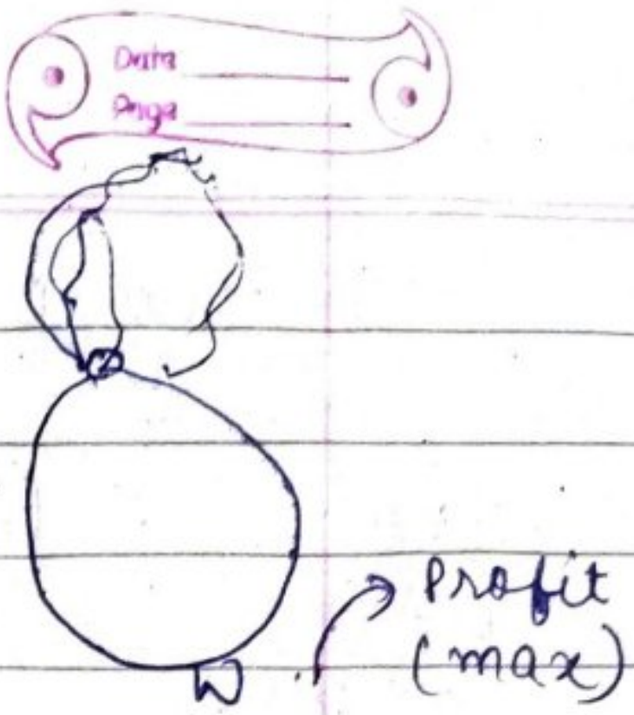
0-1 Knapsack → We are given some weight & some value. Then a max weight W. pick items so that the profit is maximum. And the weight has a given bound W.

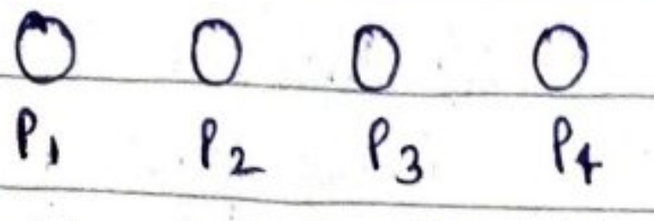


	I_1	I_2	I_3	I_4
wt[] =	1	3	4	5
val[] =	1	4	5	7

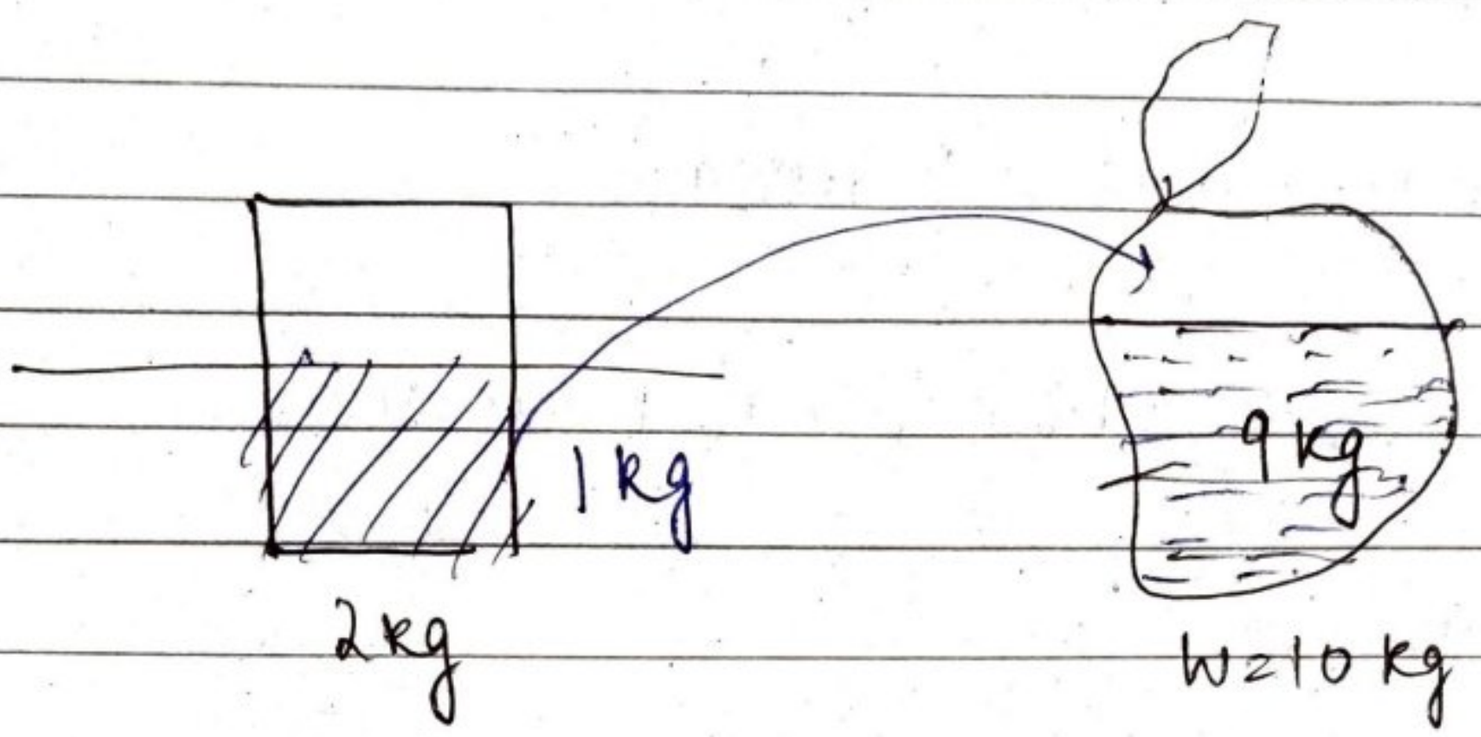


Max Profit = ?





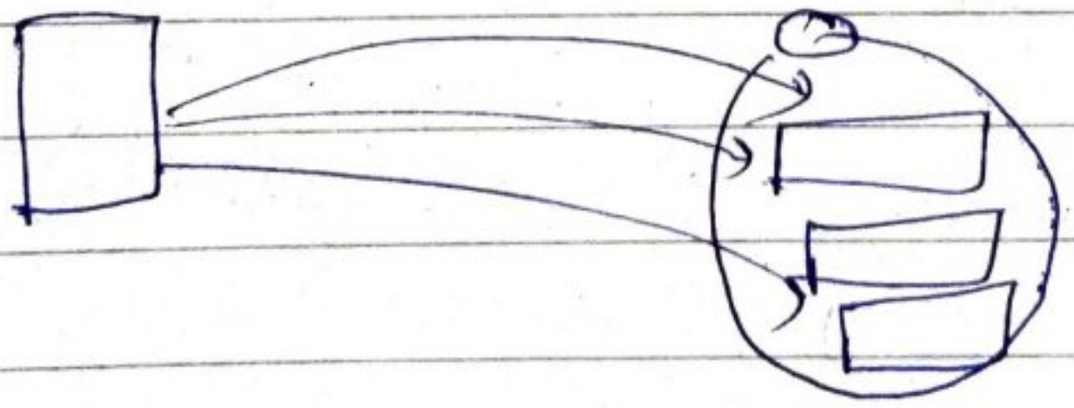
Fractional knapsack



Greedy approach

Unbounded Knapsack

Unlimited supply of every item



0-1 Knapsack

i) How to identify

wt[] : 1 3 4 5 W : 7kg

val[] : 1 4 5 7

O/p : max profit

1) Choice → (either put in knapsack) (either don't put in knapsack)

2) Optimal (max profit)

DP: Recursive → Memoization (DP) → Top down (DP)

DP → recursion + storage

0-1 knapsack Recursive

Identify

DP → Recursive → DP (top down)
→ DP (memoization)

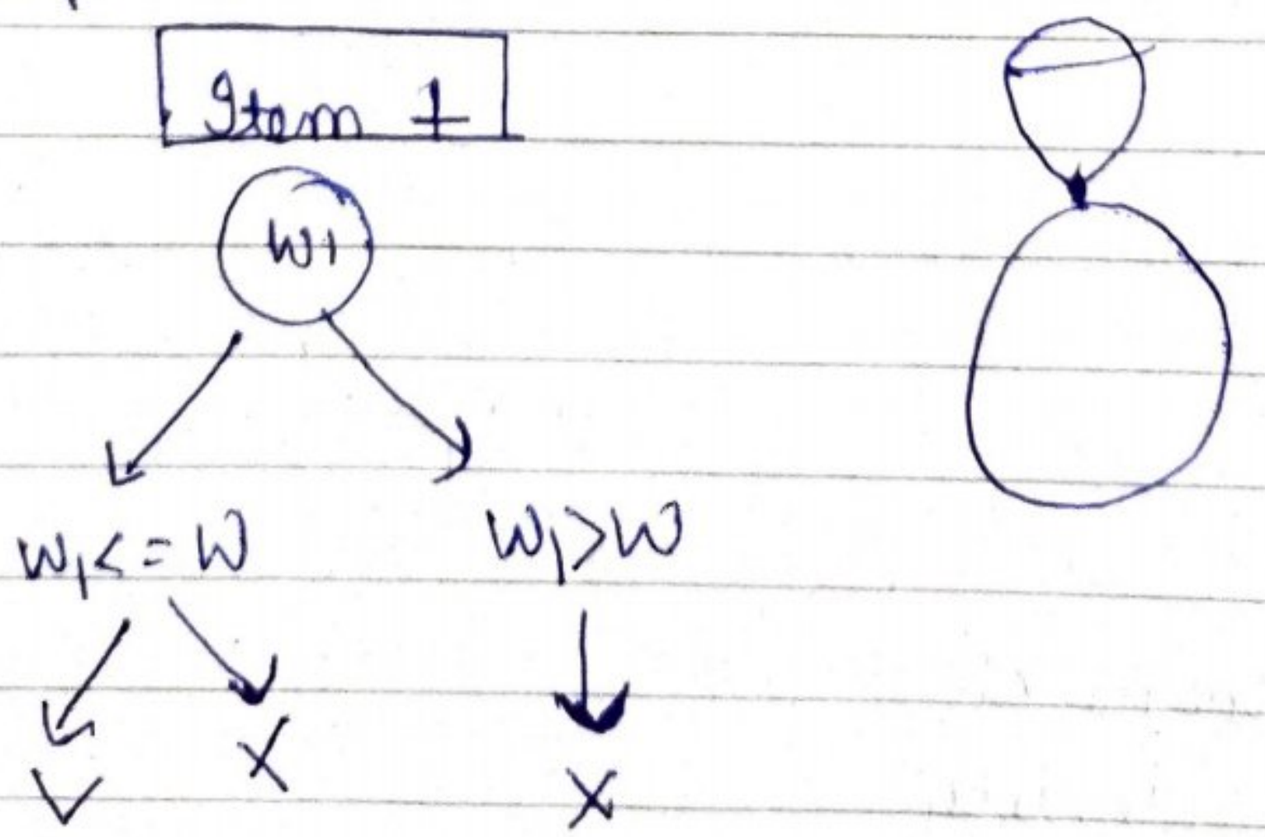
I/P wt[] = [1|3|4|5]

val[] = [1|4|5|7]

(Capacity of knapsack) W = 7 kg

O/P → Max Profit

Choice diagram

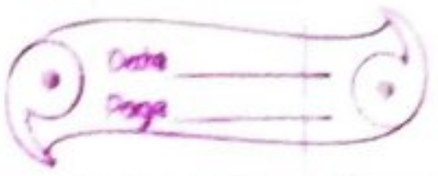


We have to return the max profit so return type would be int.

Base cond^n → think of the smallest valid i/p.

I/P wt[]: [ ][ ] → (n) → 0.
val[]: [ ][ ]

W : 10kg ⟶ 0 kg



```
int Knapsack ( int wt[], int val[], int W, int n) {
// base condition
  if (n == 0 || W == 0)
      return 0;
```

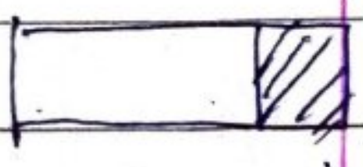
max profit

// choice diagram

```
if ( wt[n-1] <= W) {
    return max(val[n-1] + Knapsack ( wt, val, W - wt[n-1], n-1),
                Knapsack ( wt, val, W, n-1));
}
else if ( wt[n-1] > W) {
    return Knapsack ( wt, val, W, n-1);
}
```

weight less hai & including

weight less hai but don't include



01 Knapsack Memoization

Memoization = Recursive + 2 lines

t

n*m

val of n, m {

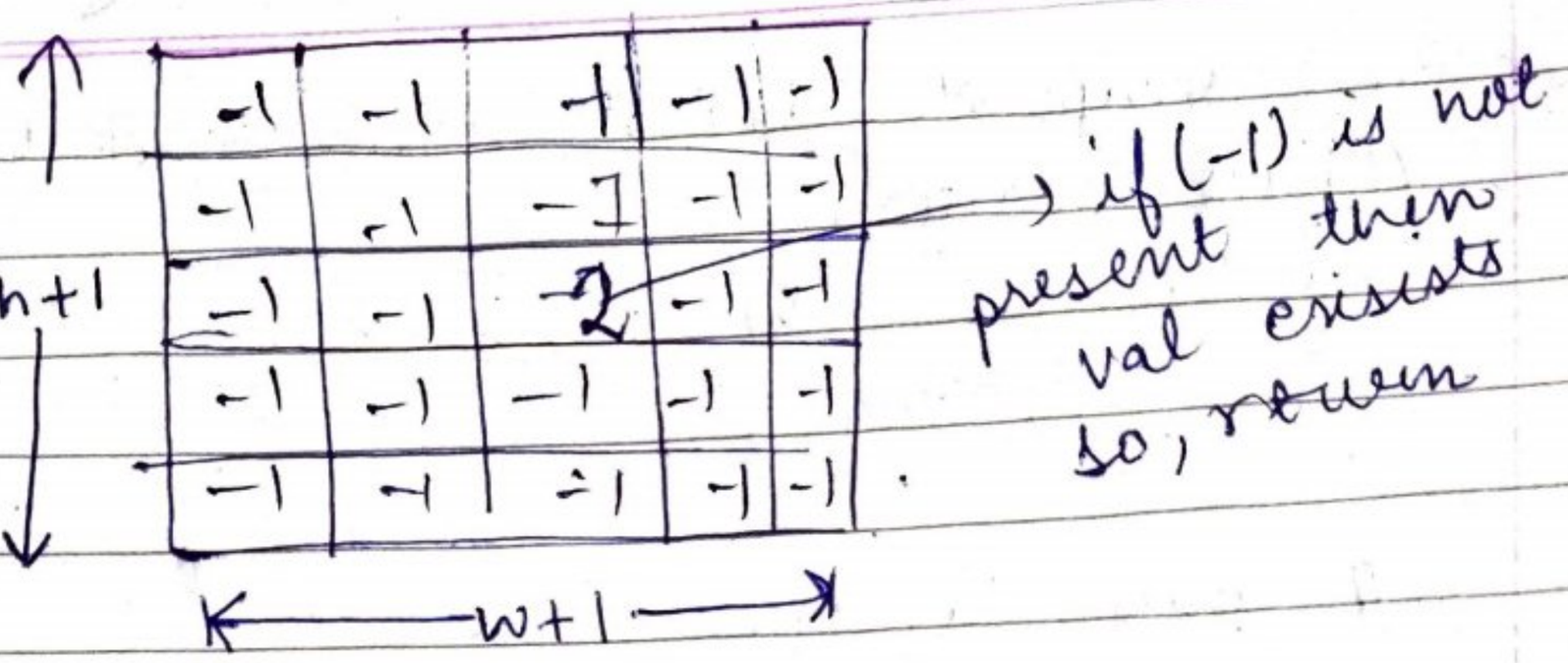
func (IP/doesn't change , IP/ change)

n-1

W - wt[n-1]

KnapSack (wt[], val[], n, W)

t [n+1] [W+1]



initialise this matrix with -1.

```
int t[n+1][w+1]
memset (t, -1, sizeof(t))
```

Changes

Now declare the matrix globally.

```
int static t[102][002]; constrainst
                          n<=100
memset (t, -1, sizeof t)   w<=1000

int knapsack (int wt[], int val[], int w, int n)
  {
    if (n==0 || w==0)
      return 0;
    if (t[n][w] != -1)
      return t[n][w];

if (wt[n-1] <= w)
return t[n][w] = max (val[n-1] + knapsack (wt, val, w-wt[n-1], n-1
, knapsack (wt, val, w, n-1));
```

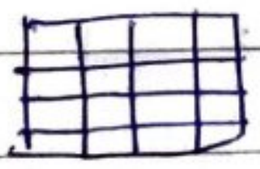
```
else if (wt[n-1] > W)
return t[n][W] = knapsack(wt, val, W, n-1);
}
```

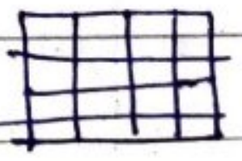
The complexity of top down & memoization remains same but the problem with memoization is the stack gets full. due to repeated func^{n} calls.

(0-1 Top Down Knapsack) → Real DP

Recursive → Memoize → Top down → 6 Problems

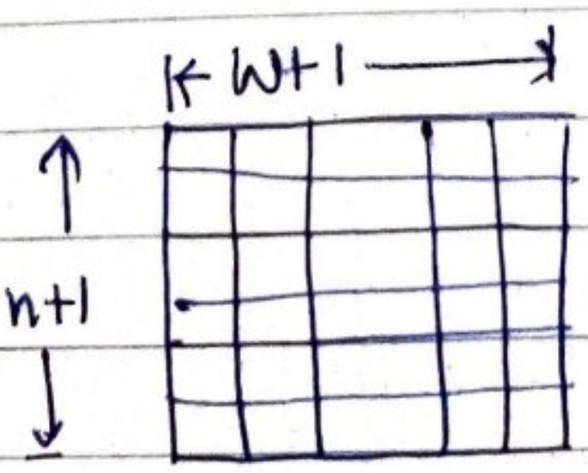
↓	↓	↓
BC + recursive calls	RC + table	Only Table





Recursive → Memoization — Initialising matrix with -1

Recursive → Top-down (totally omit the recursive call & use the table only)



We only make table for the i/p which is changing.

2 steps to make table for top down

Step1: Initialization

Step2: Recursive code $\xrightarrow{\text{changes}}$ Iterative code

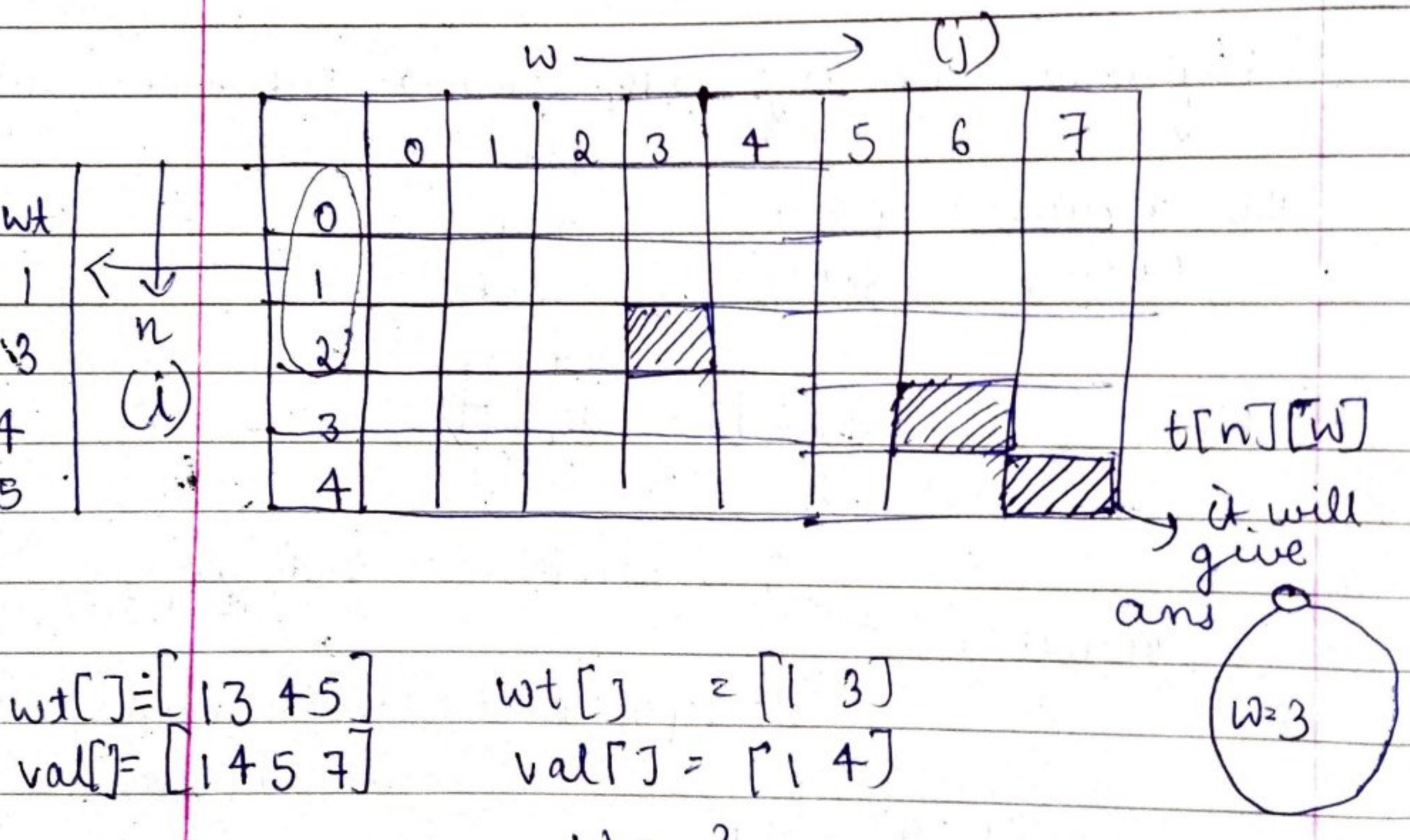
Step1: Initialize

W = 7

n = 4 (0 1 2 3)

wt[] = [1 3 4 5]

val[] = [1 4 5 7]



wt[] = [1 3 4 5]

val[] = [1 4 5 7]

wt[] = [1 3]

val[] = [1 4]

W = 3

wt[] = [1 3 4]

val[] = [1 4 5]

W = 6

RC + table ⟶ table

Base cond^n ⟶ Initialization

Base cond^n

RC

```
if (n==0 || W==0)
        return 0;
```

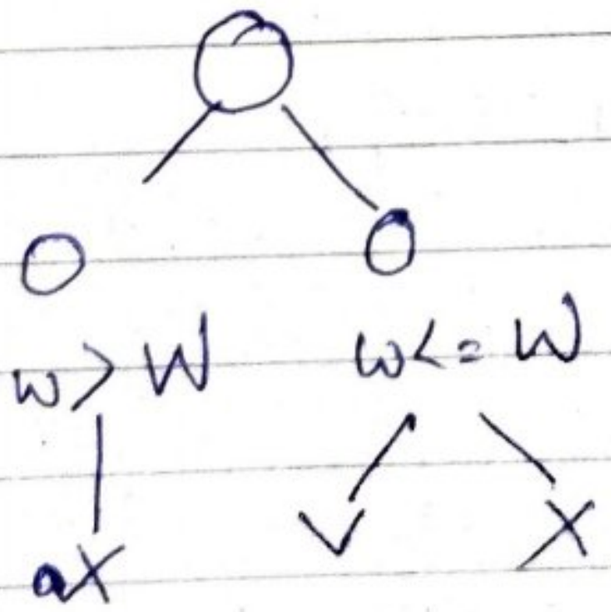
↓

table

	0	1	2	...
1	0	0	0	0
2	0			
3	0			
4	0			

```
for (int i=0; i<n+1; i++) {
    for (int j=0; j<w+1; j++)
    {
        if (i==0 || j==0)
            t[i][j] = 0;
    }
}
```

Choice diagram



n, w ⟶ i, j

7+

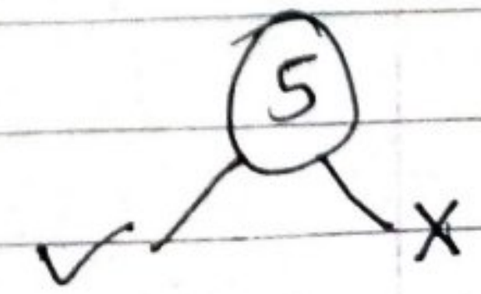
dp[i][j]

RC.

wt = 1 3 4 5

val = 1 4 5 7

W = 7



max(val[n-1] + dp[i-1][j-wt[i-1]], dp[i-1][j])

Recursive

```
if(wt[n-1]<= w)
return max(val[n-1]+Knapsack(wt,
        val, w-wt[n-1],n-1)
        knapsack(wt, val, w, n-1))
else if (wt[n-1] > w)
return knapsack(wt, val, w, n-1)
```

Top down

```
if (wt[n-1] <= w)
t[n][w] = max( val[n-1] +
          t[w-wt[n-1]]
          [n-1],
          t[n-1][w])
else
t[n][w] = t[n-1][w]
```

Top-down approach

```
int t[n+1][w+1];
for(int i=1; i<n+1; i++)
    for(int j=1; j<w+1; j++)
    {
      if (wt[i-1] <= j)
t[i][j] = max( val[i-1] + t[i-1][j-wt[i-1],
                t[i-1][j] )
      else
        t[i][j] = t[i-1][j]
    return t[n][w];
```

	0	1	2	3	
wt =		1	3	4	5
val =		1	4	5	7

W = 7, n = 4

W =	0	1	2	3	4	5	6	7
0								
1								
2								
3								
4								

pseudo code

i=1 → 4

j=1 → 7

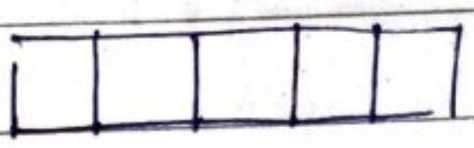
i=1, wt[i-1] <= 1 → true

1 <= 1 → true

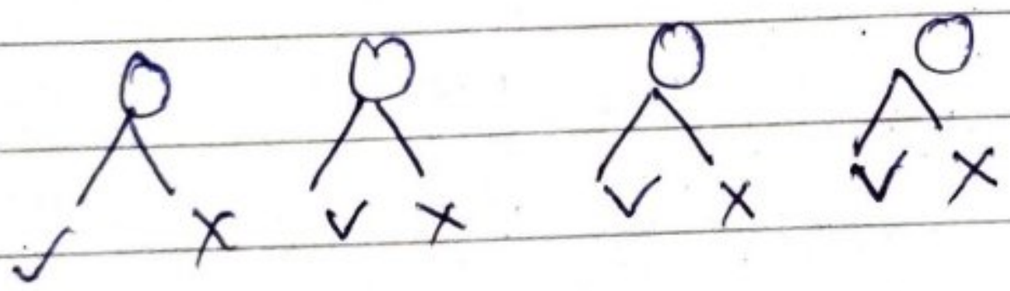
t[i][j] = max(1 + t[0][0], t[0][j]) = max(1, 0) = 1

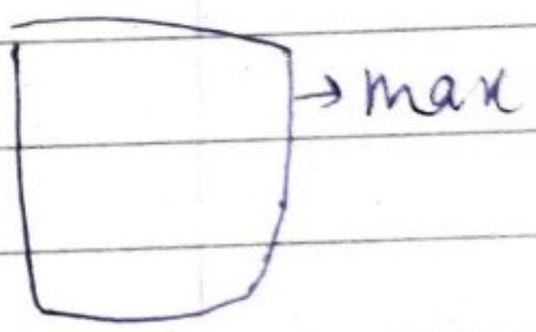
Identification of knapsack problem

1) Subset sum problem
2) Equal sum partition
3) Count of subset sum.
4) Minimum subset sum diff
5) Target sum
6) No of subset with a given diff.

Ip: Item array : 

W: Capacity





1. Subset Sum problem

arr[]: 2 3 7 8 10

Sum = ...

a) Problem statement
b) similarity with knapsack
c) Code variation

1) Problem statement : find if there is a subset present in an array with given sum.

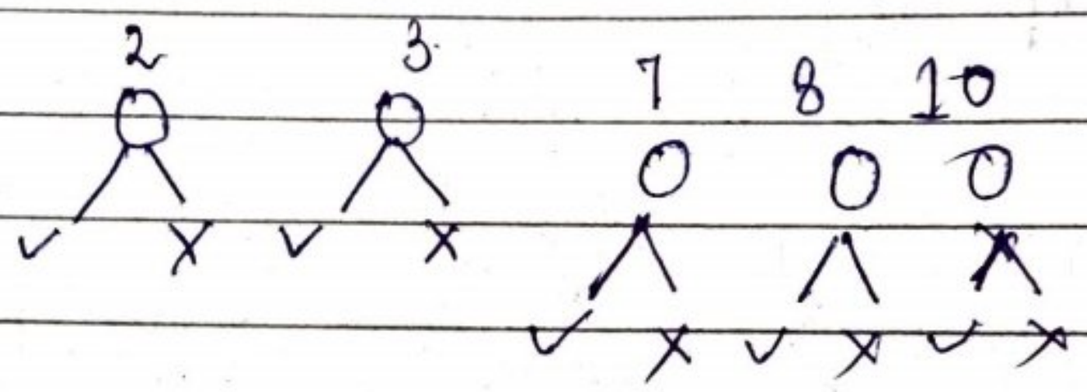
arr[]: 2 3 7 8 10

Sum : 11

2) Similarity

item array[] : → 2 3 7 8 10

weight : capacity → 11



3) Code variation

t[n+1] [w+1]

t[S+1] [11+1]

sum → t[6] [12]

Initialization :

	0	1	2	3	4	5	6	7	8	9	10	11
0	T	F	F	F	F	F	F	F	F	F	F	F
1	T											
2	T											
3	T											
4	T											
5	T											

↓ n

~~Q~~ → True/false

arr[]:

sum : 0.

arr[] : 1

sum = 0 '{}'

arr[] : 2

sum = 0 '{}'

arr[] :

Sum = 2

when array is empty sum can't be anything

arr[ ] : no elements

Sum : 1 → not possible

t[n+1][sum+1] n

Initialisation: for(int i = 1 ... n+1)

for(int j = 1 ... W+1)

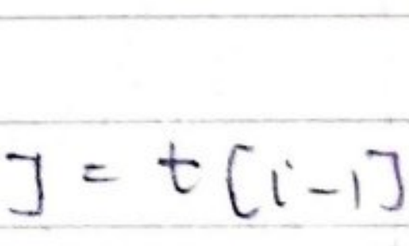
if (i == 0)

t[i][j] = false.

2 3 7 8 10

if (j == 0)

t[i][j] = True



knapsack → arr.

if (wt[i-1] <= j)

t[i][j] = max(val[i-1] + t[i-1][j - wt[i-1]], t[i-1][j])

there is no max in true or false.

else

t[i][j] = t[i-1][j]

i j

t[1][0] = t[0][0]

= t[1][-2] || t[0][0]

= T

Subset sum

if (arr[i-1] <= j)

t[i][j] = t[i][j - arr[i-1]]

||

t[i-1][j]

else

t[i][j] = t[i-1][j]

return t[n][sum];

2, 3, 8

t[i][j]

{3, 8} → true ✓

{3} → false ✗

2. Equal Sum Partition Problem

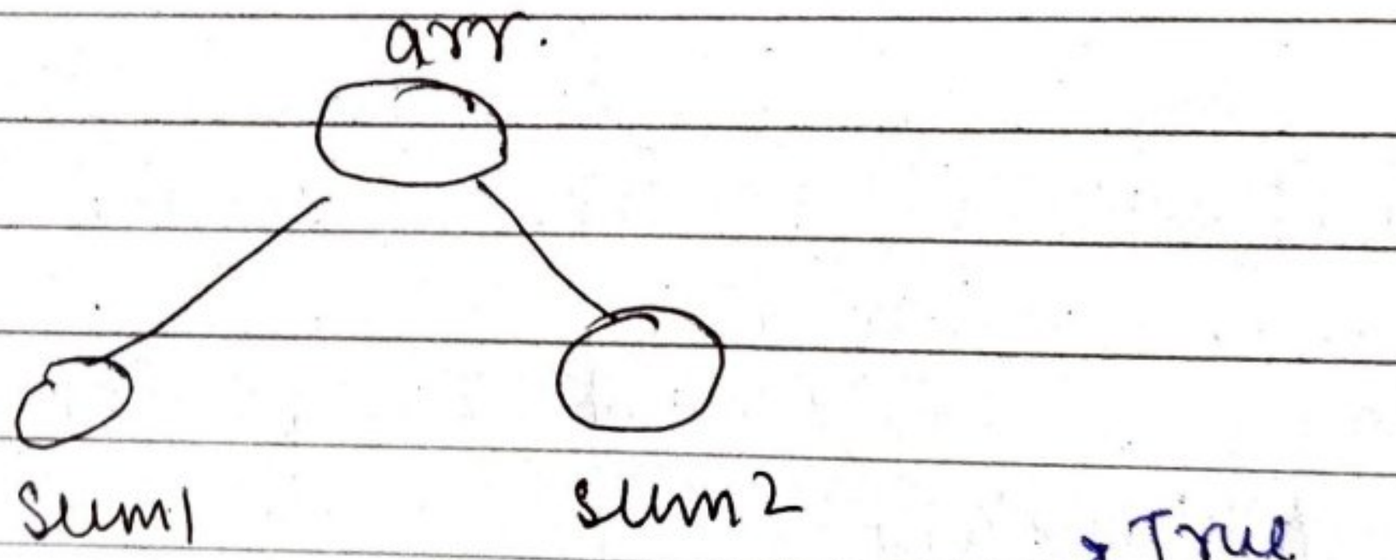
1) Problem statement
2) Subset sum similarity
3) Odd / Even significance
4) Code variation

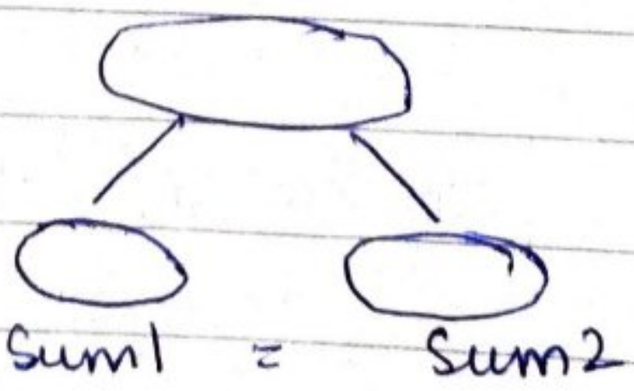
1) Problem statement

arr[] = {1, 5, 11, 5}

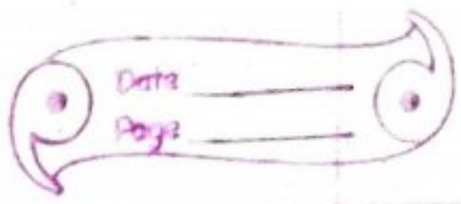
O/p : T/F

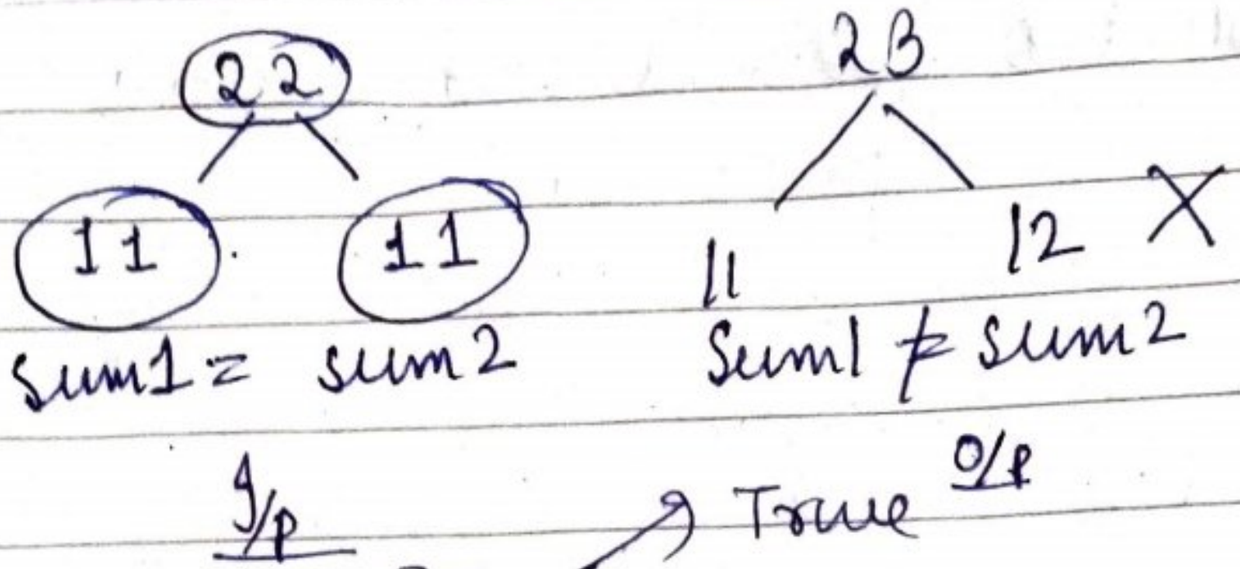
Is it possible to divide the array such that both the subset gives an equal sum





Sum1 = Sum2 → It is equal when the no is even. Sum of all the array elements should be even to change it into equal parts.





{1, 5, 11, 5}

Subset even (sum=22) Subset odd (false)

Subset =11 Subset 11

```
for (int i = 0; i<size; i++) {
        sum = sum + arr[i];
    }
if (sum % 2 != 0)        (sum is odd)
    return false
else if (sum % 2 == 0)
```

* → we need to find one subset with sum 11 the next subset would automatically be 11.

return subsetsum(arr, sum/2);

3. Code:

```
ip → arr[], n.
int sum = 0;
for (int i = 0; i<n; i++)
    sum += arr[i];
if (sum % 2 != 0)
    return false
else
    return subsetsum(arr, sum/2);
```

3. Count of Subsets sum with a given sum.

I/p:

arr[] = 2 3 5 6 8 10

Sum = 10.

Flow

1) Problem Statement
2) Similarity to Subset sum
3) Code variation.
   - Initialisation
   - Code
4) Return Type.

1) Problem statement

arr[] 2 3 5 6 8 10

Sum : 10 O/p = 3.

{2,8} → Yes/True (in subset sum)

{2,8} = 10
{5,2,3} = 10
{10} = 10
} Count = 3

2) Similarity

2, 8 | return count

Yes NO

3. Code variation

Subset sum

due to ← bool
T/F.

False → 0 (no of subset 0)

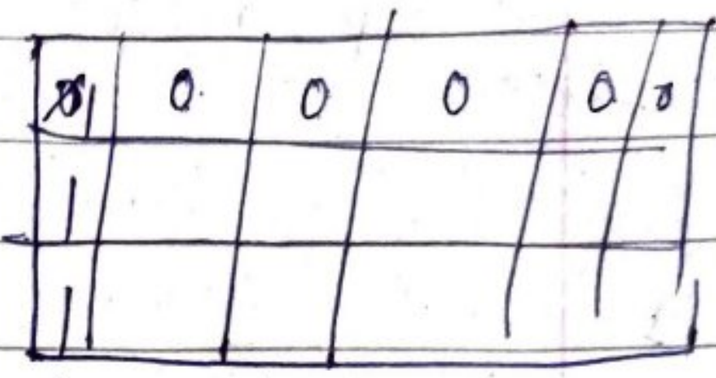
True → null subset (no of subset 1)

. . . . . .

if (arr[i-1] <= j)
dp[i][j] = dp[i-1][j] || t[i-1][j-arr[i-1]]

Else.
dp[i][j] = dp[i-1][j]

Count
int



We will add all the subset so ou would be changed to +. Ture & false case we can use bool ou.

if (arr[i-1] <= j)
dp[i][j] = dp[i-1][j] + dp[i-1][j-arr[i-1]]

else
dp[i][j] = dp[i-1][j]

if (arr[i-1] <= j)

dp[i][j] = dp[i-1][j-arr[i-1]] + dp[i-1][j]

else

dp[i][j] = dp[i-1][j]

arr[] = 2 3 5 6 8 10

Sum = 10

O/p = 3

1, 3, 2, 5 5

O/p → {3, 2}

{5}

arr[0] <= 1

i=1 j=1 2-2>=0

i=1 j=2

arr[0] <= 2

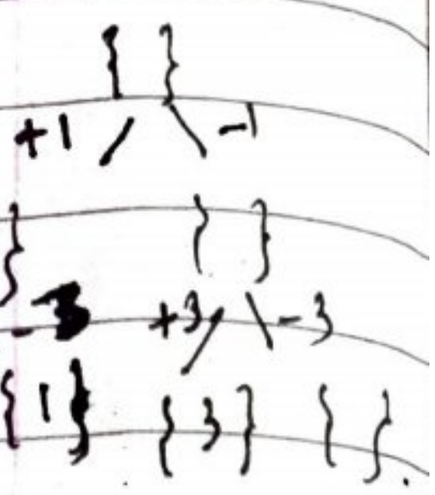
2 <= 2

dp[0][0]

dp[0][2]

dp[N+1][sum+1] = dp[6+1][10+1]

= dp[7][11]



Sum →

	0	1	2	3	4	5	6	7	8	9	10
0	1	0	0	0	0	0	0	0	0	0	0
1	1	0	-1								
2	1										
3	1										
4	1										
5	1										
6	1										

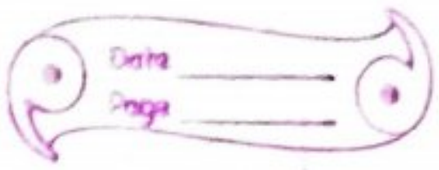
Minimum subset sum difference

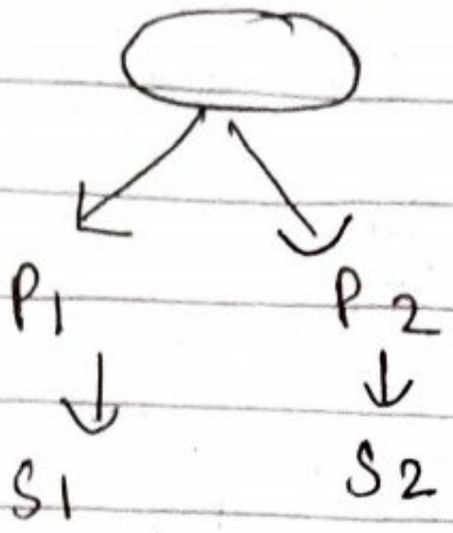
1) Problem statement
2) Similarity
3) Solve using its previous concept +

1) Problem statement

arr[]: 1 6 11 5

OP: 1





Equal sum : $S1 - S2 = 0$

Num^m Subset : $S1 - S2 = min$

$abs(S1 - S2) = min$ (should be min)

1, 6, 11, 5

{1, 6} {11, 5}

7 16

$16 - 7 = 9$ → minimize of find could it be done in a better way.

{1, 6, 5} {11}

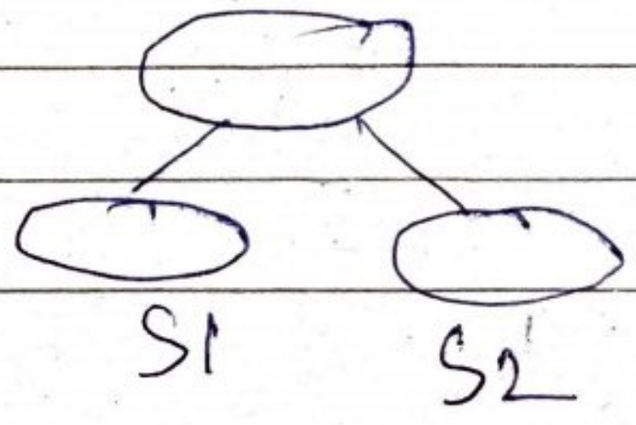
12 11

$12 - 11 = (1)$ → Can't be minimized further

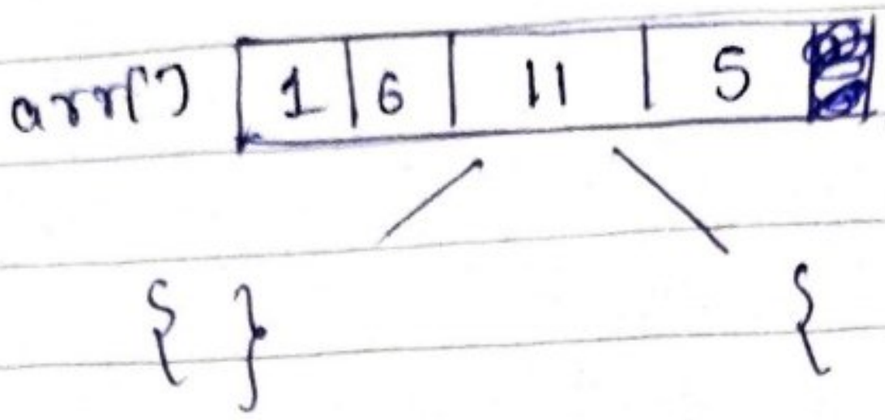
↳ O/P.

2) Similarity

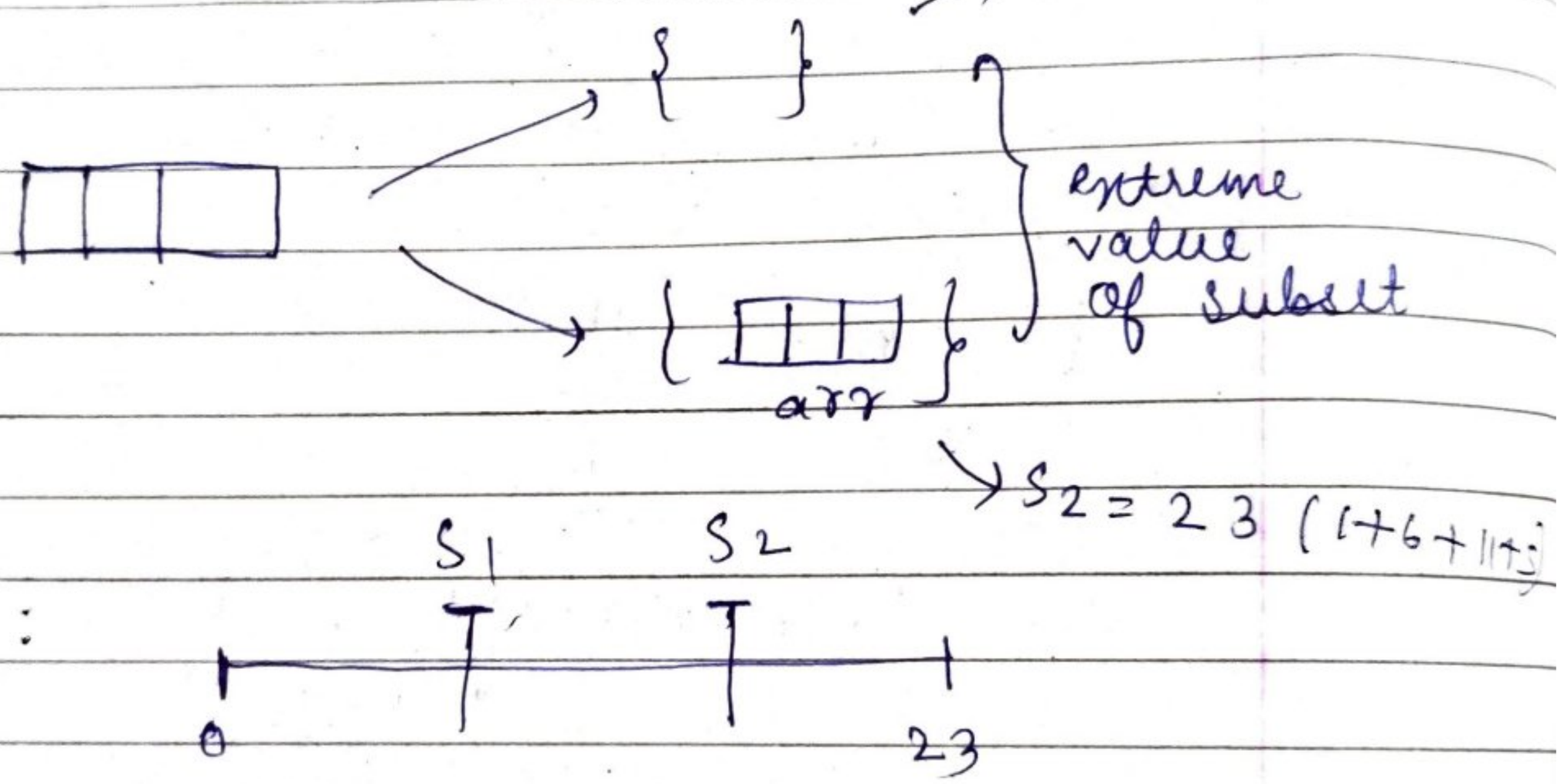
It is similar to equal sum partition.



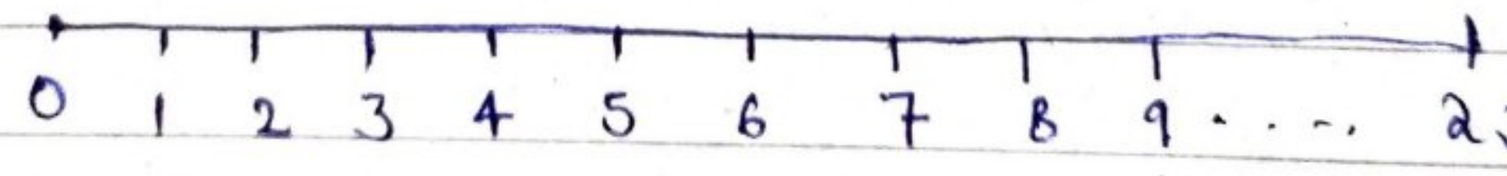
We need to find S1 & S2.



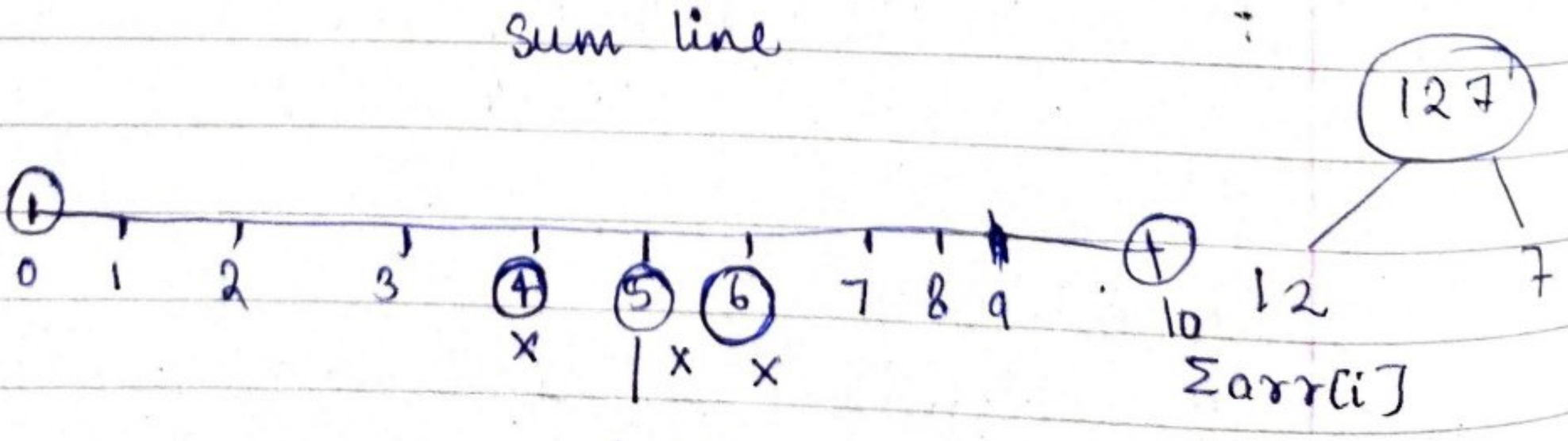
We can find the range of S1 & S2.



arr[]: {1, 2, 7}

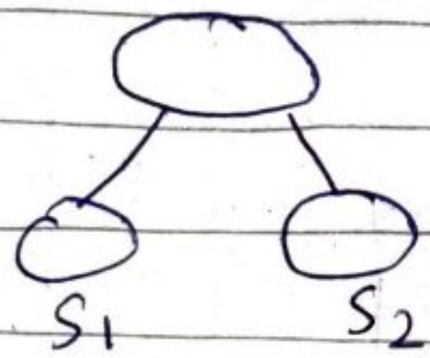


Sum line :



S1/S2 can't be 5

S1/S2 = {0, 1, 2, 3, 7, 8, 9}

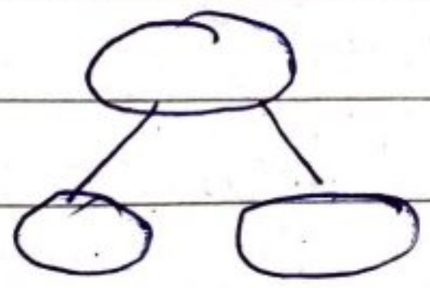


$S_1 + S_2 = \Sigma arr[i]$.

$S_1/S_2 = \{ 0, 1, 2, 3, \mid 7, 8, 9, 10 \}$.

S_1 (S_1) — Range-S_1 (S_2)

$S_1 - S_2$ = minimize



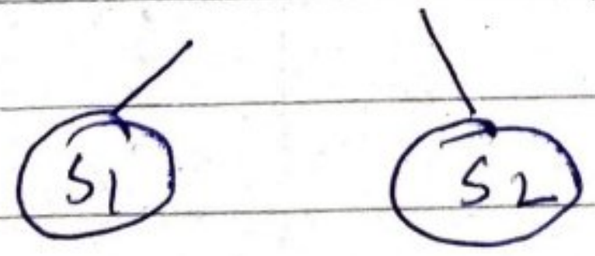
abs $(S_2 - S_1)$ or abs $(S_1 - S_2)$

abs $(S_2 - S_1)$ = minimize

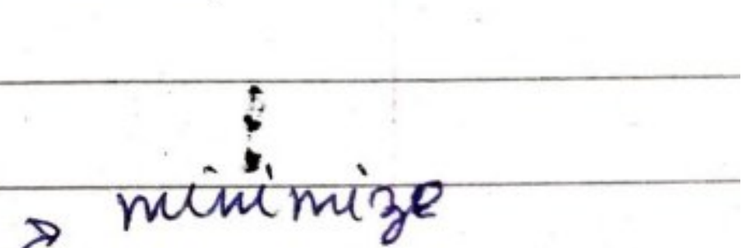
abs (Range-S1 -S1) = minimize. abs (Range - 2S1) ← minimize

Sum up

1 6 11 5

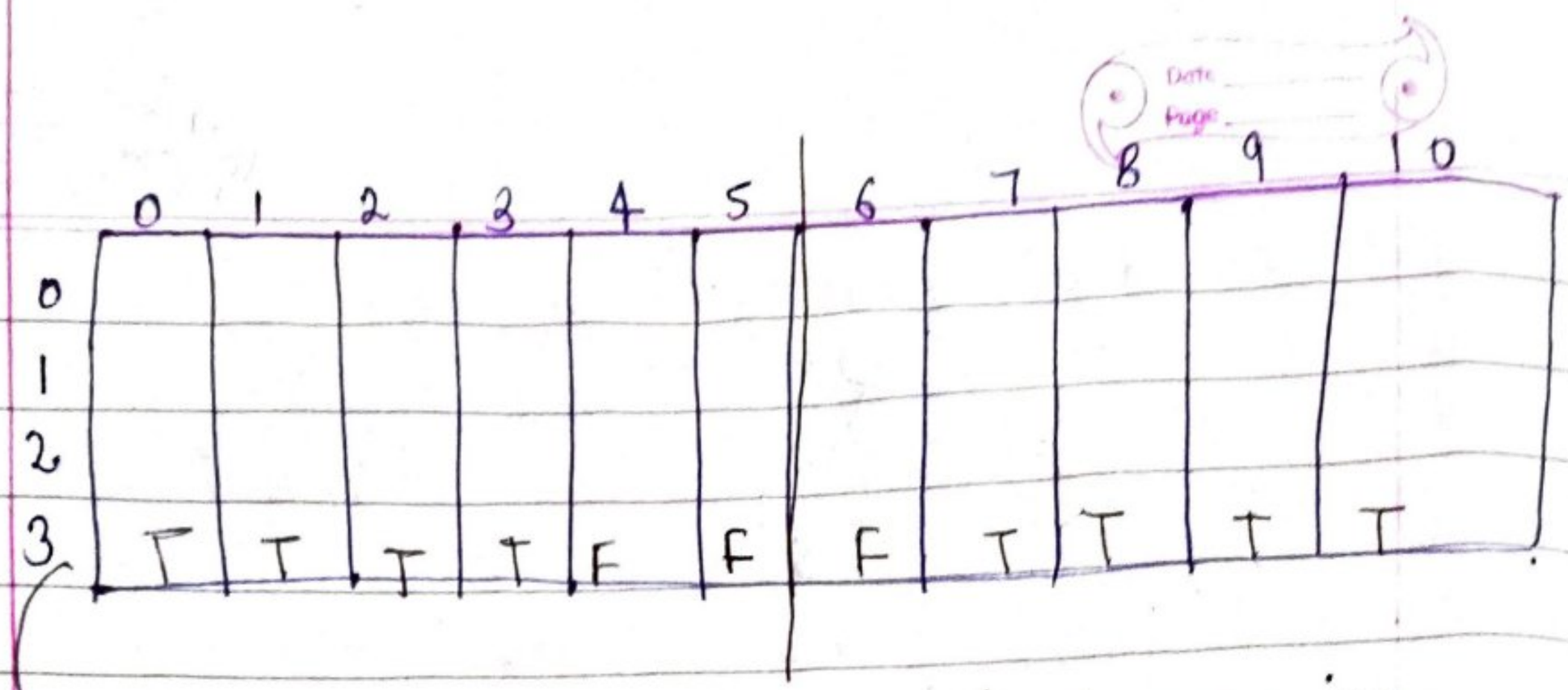


$S_1 - S_2$
$S_2 - S_1$ } minimize.



((Range - S1) - S1)

(Range - 2S1)

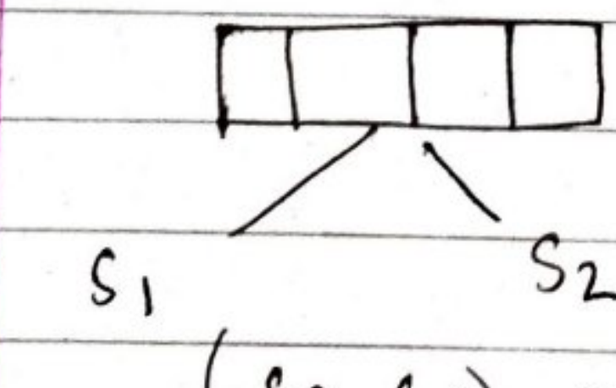


↳ we will push the last row in vector till half way.

0	1	2	3.

```
int min = INT_MAX;
for (int i = 0; i < v.size(); i++) {
    {
        mn = min (mn, Range - 2v[i])
    }
}
return mn;
```

Concept



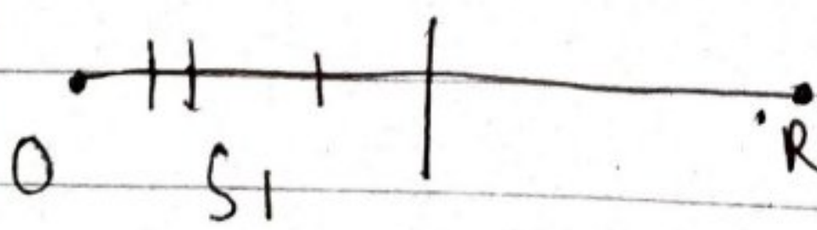
(S2 - S1) = min

↓

smaller

(S1) (Range - S1)

Range - 2S1 → Minm



Code

```
Subset sum( int arr[ ], int range)
{
```

last row filled in vec till halfway

0	0	0	

```
}
```

```
int mindiff (int arr[], int n) {
  int sum = 0;
    for (int i = 0; i < n; i++) {
        sum += arr[i];
    }
    bool dp[n+1][sum+1];
    for (int i = 0; i < n+1; i++) {
        for (int j = 0; j < sum+1; j++) {
        if (i == 0)    dp[i][j] = false;
        else if (j == 0)  dp[i][j] = true;
        }
    }
  for (int i = 1; i < n+1; i++) {
        for (int j = 1; j < sum+1; j++) {
        if (arr[i-1] <= j)
            dp[i][j] = dp[i-1][j-arr[i-1]] || dp[i-1][j];
        else
            dp[i][j] = dp[i-1][j];
        }
    }
  int diff = INT_MAX;
    for (int j = sum/2; j >= 0; j--) {
        if (dp[n][j] == true) {
            diff = min(sum - 2*j, diff)
        }
    }
    return diff;
}
```

Date
Page

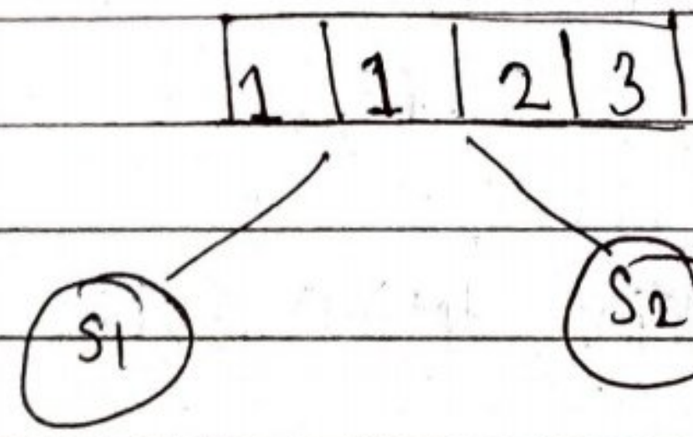
Count the number of subset with a given diff

1) Problem statement
2) Will try to reduce the actual statement
3) Solve it using already solved problem.

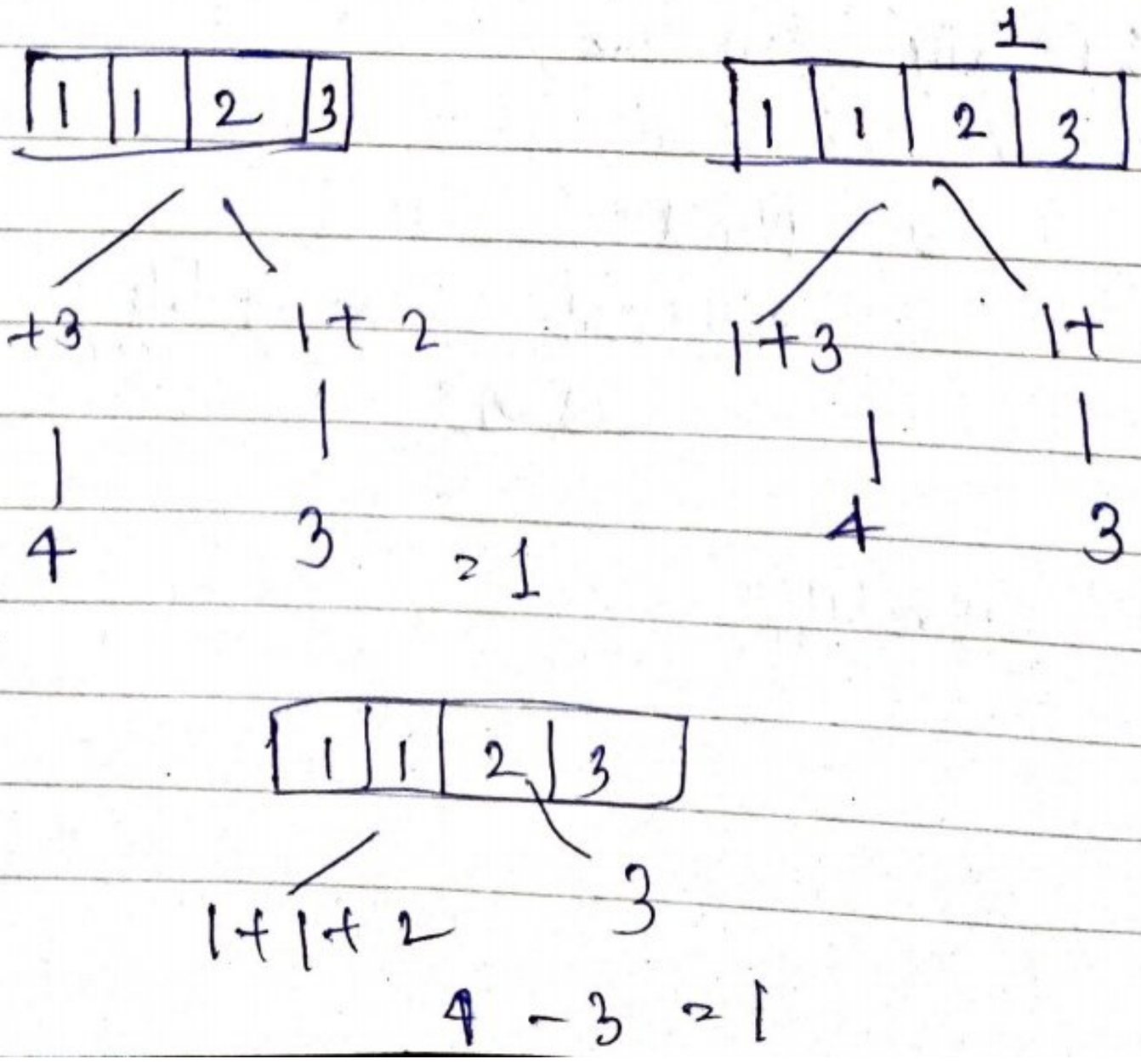
1) Problem statement

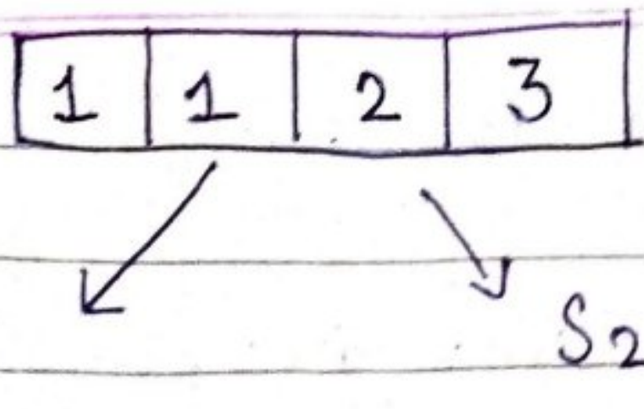
arr[] : | 1 | 1 | 2 | 3 |

Diff : 1



S1 - S2 = Diff → return the count of subset with diff^n.





(diff) S = Sum of arr

Sum(s1) - Sum(s2) = diff

Sum (s1) + Sum(s2) = S

$2S_1 = diff + Sum(arr)$

$S_1 = \left(\frac{diff + Sum(arr)}{2}\right)$

$= \frac{1+7}{2} = \frac{8}{2} = 4$

$S_1 = 4$ $S_2 = S_1 - diff$

$= 4 - 1$

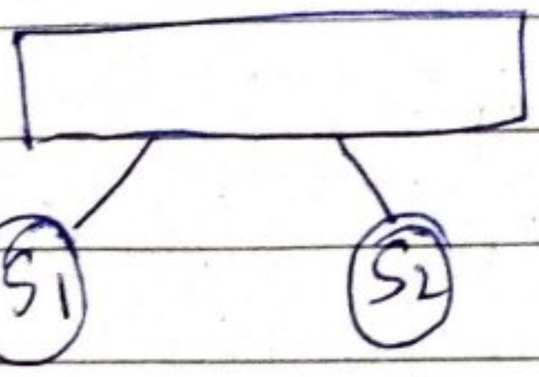
$= 3.$

Count = ?

- Find count when sum of S1 = 4

no of count of Subset with given diff ⟶ count of Subset sum.

Sum up.



$S_1 - S_2 = diff$

$S_1 + S_2 = sum(arr)$

$2S_1 = diff + sum(arr)$

$S_1 = \frac{diff + sum(arr)}{2}$.

int sum = diff + sum(arr)
2.

return countofsubset sum(arr, sum);

int countofsubsetwithdiff (int arr[], int n){
~~int sum=0 for(int i=0; i<n; i++){~~ int diff
~~sum=sum+arr[i];~~

int arrsum=0
for (int i=0; i<n; i++){
arrsum += arr[i];
}

int sum = 0;
sum = diff+arrsum
2.

return countofsubsetsum(arr, sum, n);
}

int countofsubsetsum(int arr[], int sum, int n){
int dp [n+1][sum+1];
for(int i=0; i<n+1; i++){
for(int j=0; j<sum+1; j++){

if (i==0) dp[i][j]=0;
if (j==0) dp[i][j] = 1;
}
}

```
for (int i = 1; i < n+1 ; i++){
    for(int j = 1; j < sum+1 ; j++) {

        if (arr[i-1] <= j ) {
dp[i][j] = dp[i-1][j - arr[i-1]] + dp[i-1][j];
        else {
        {dp[i][j] = dp[i-1][j];
        }
    }
}
return dp[n][sum];
}
```

Target Sum..

arr : | 1 | 1 | 2 | 3 |

Sum : 1

+/-

+c : [0, 0, 0, 0, 0, 0, 1]

target = 1

o/p : 3

Target value = 2

+1 -1 -2 + 3 = 1

-1 + 1 - 2 + 3 = 1 {0, 2}

+1 + 1 + 2 - 3 = 1

(+0, 2) (-0, 2)

count = 2

arr[] → | + | - | + | + |

{0, 0, 2}

{+0, +0, 2} {+0, -0, 2} {-0, -0,}

| | | | |

S1 S2

{-0, -0, 2}

S1 - S2 = diff

= 4 (count

So your answer is in the form of power of (2, no of zeroes)

+ - - +

1 1 2 3

+1+3 -1-2

(1+3) - (1+2)

S1 - S2 → count of subset with given diff^n.

~~0/1 Unbounded Knapsack~~

```
int findTargetSumWays(vector<int>& nums, int S) {
    int cnt = 0, sum = 0;
    int n = nums.size();
    for(int i = 0; i < nums.size(); i++) {
        sum = sum + nums[i];
        if(nums[i] == 0)
            cnt = cnt + 1;
    }
    S = abs(S);
    if(S > sum || (S + sum) % 2 != 0)
        return 0;
    int s = (S + sum) / 2;
    int dp[n+1][s+1];

    for(int i = 0; i < n+1; i++)
        for(int j = 0; j < sum+1; j++)
            if(i == 0) dp[i][j] = 0;
            if(j == 0) dp[i][j] = 1;
```

cnt the no of zeroes in subset

```
for (int i=1; i<n+1; i++) {
    for (int j=1; j<s+1; j++) {

        if(nums[i-1]==0)
            dp[i][j] = dp[i-1][j];

        else if (num[i-1]>j)
            dp[i][j] = dp[i-1][j];

        else
            dp[i][j] = dp[i-1][j-nums[i-1]] +
                       dp[i-1][j]

    }.
}.
    return pow(2,cnt) * dp[n][s];
}.
```

if 0 is there in the array dont take it

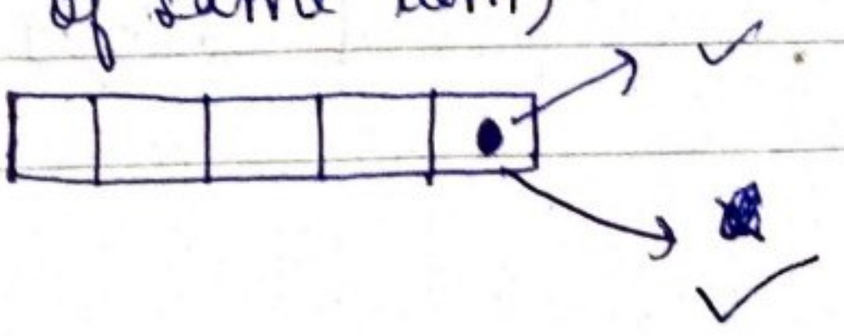
13. Unbounded Knapsack

Related Problems

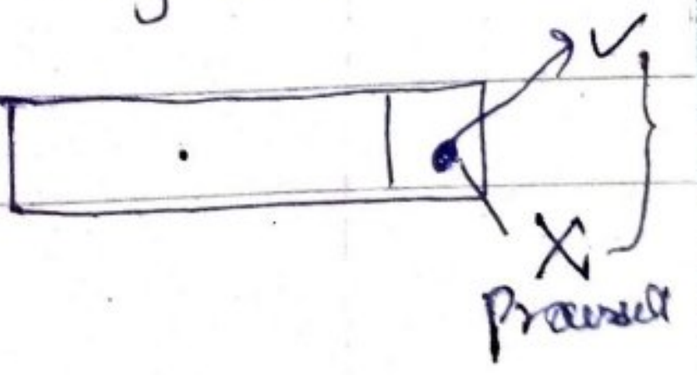
1) Rood cutting
2) Coin change I (Max no of ways)
3) Coin change II (Min no of ways)
4) Max^m. Ribbon cut

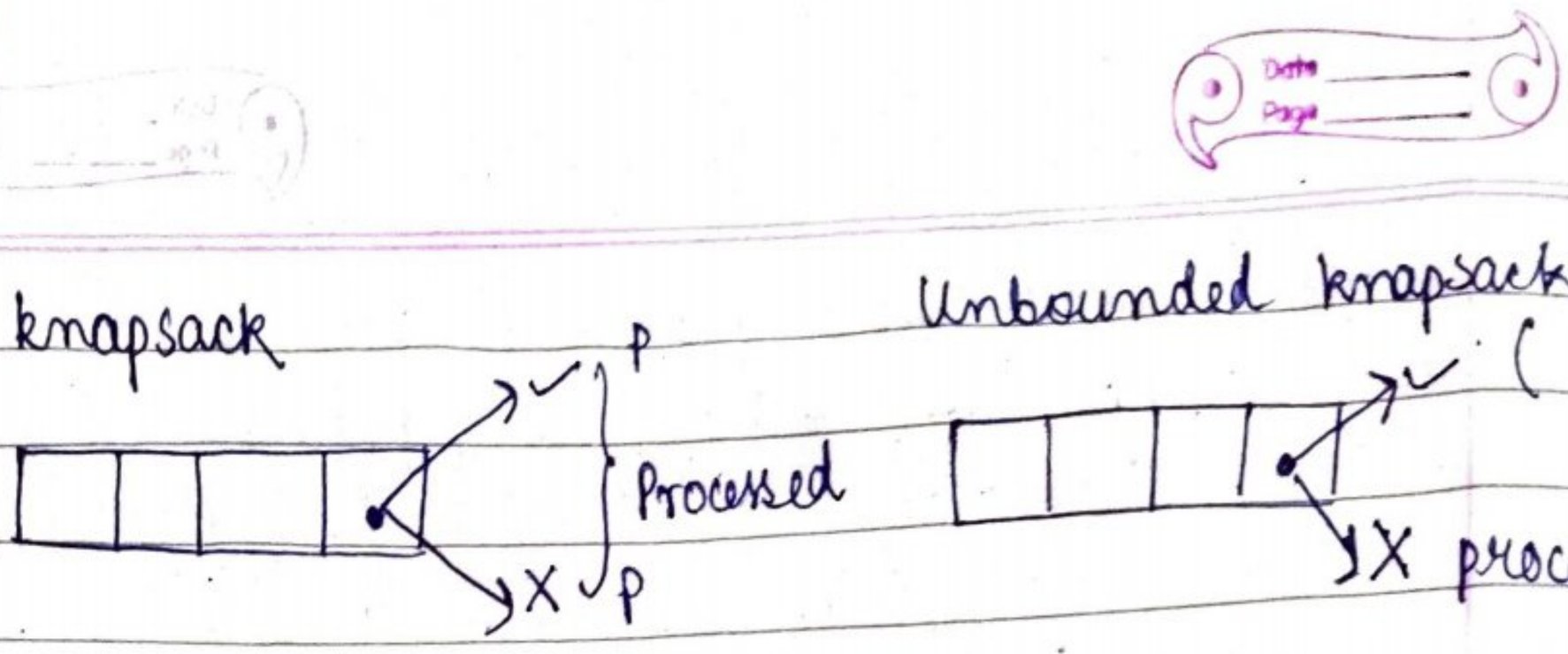
(Variations of unbounded knapsack)

Unbounded (multiple occurences of same item)



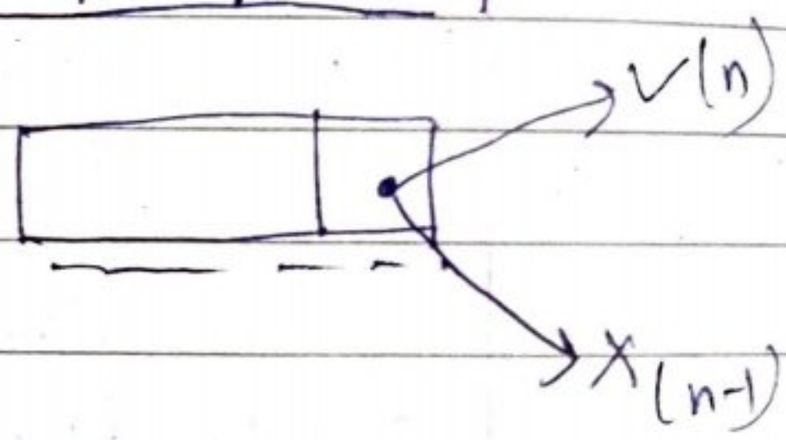
Knapsack (only one occurence)





Multiple occurrences
- ✓ (can revisit many times)
- X

Comparision betn.

Knapsack	Unbounded knapsack
0 1 2 3 4 0 0 0 0 0 0 1 0 2 0	0 1 2 3 4 0 0 0 0 0 0 1 0 2 0
t[n+1][w+1]	

if (wt[i-1] <= j){
t[i][j] = max(val[i-1] + t[i-1][j-wt[i-1]],
t[i-1][j]); }
else {
t[i][j] = t[i-1][j];
}.
	if (wt[i-1] <= j)
	t[i][j] = max(val[i-1] + t[i][j-wt[i-1]],
	t[i-1][j]);
	else
	t[i][j] = t[i-1][j]

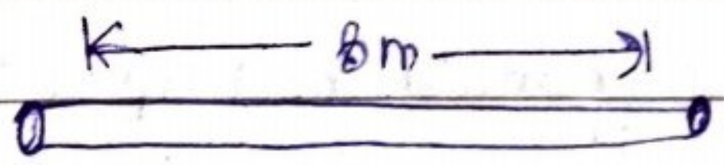
14. Road Cutting Problem

length [ ] : | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |

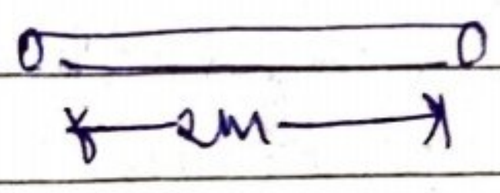
price [ ] : | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |

N : 8

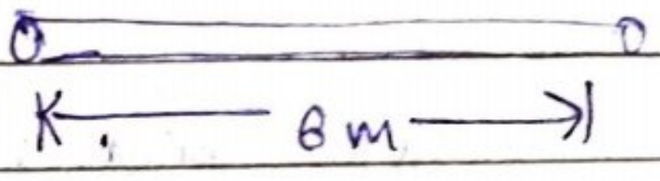
1) Problem statement



length of rod is given. We need to cut it in such a way the profit is max^m.



Rs 5



Rs 17

17
5
22

Total = 5+17
= Rs 22

knapsack

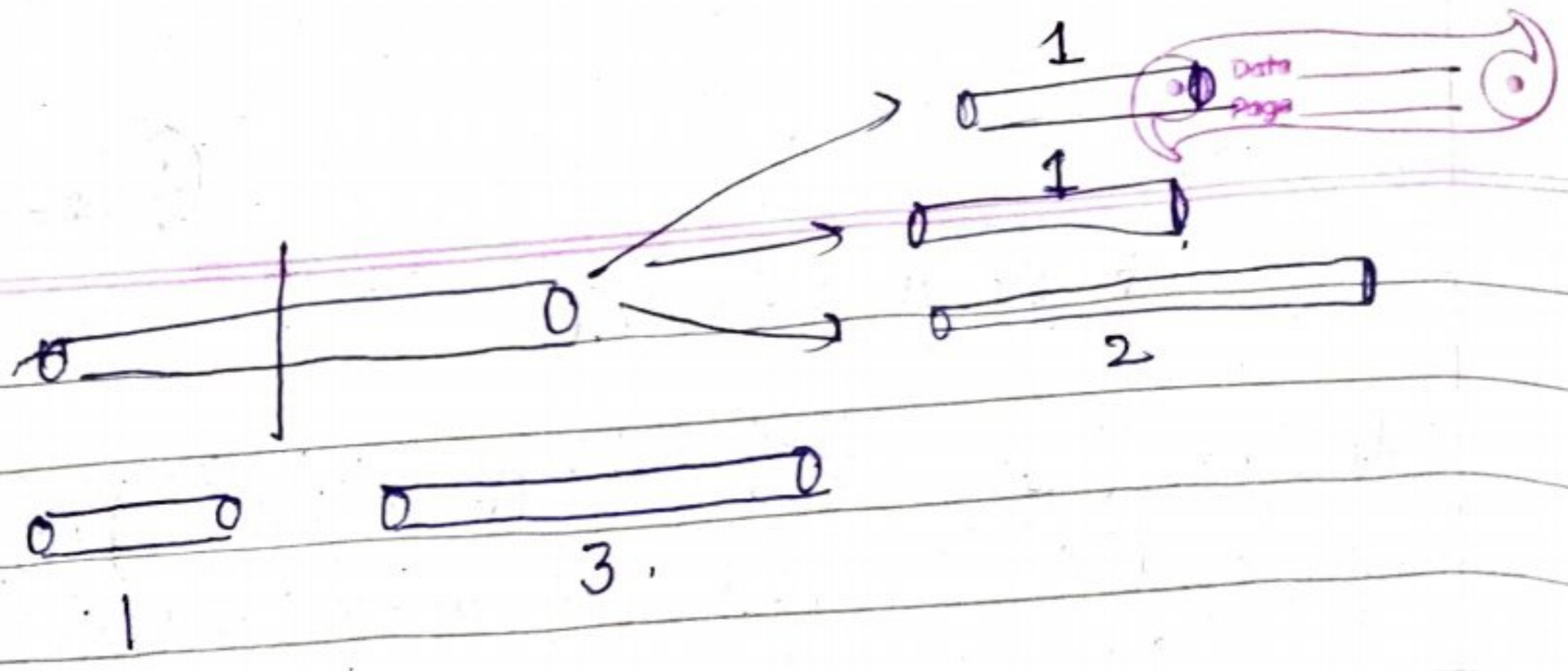
wt
val
W

length = 1 to N
Price = | | | | | |
N = 8

if length array is not give make it & fill it will 1 to N.

knapsack
price[ ]
val[ ]
wt[ ]
W

unbounded knapsack
price[ ]
length[ ]
N



Code variation

```
t[N+1] [N+1]
if (length[i-1] <= j)
dp[i][j] = max(price[i-1] + dp[i][j- length[i-1]],
                                    dp[i-1][j])
else
dp[i][j] = dp[i-1][j];
```

Sometimes arr size changes therefore we need to find the size of array.

dp[size+1][N+1].

Length →

↓ size	0	0	0	0	0	0
	0					
	0					

Coinchange Problem

Maxm no of ways

Minm no of coins

Max^m no of ways.

Problem statement

coin[ ]: | 1 | 2 | 3 | → (unlimited)

sum: 5.

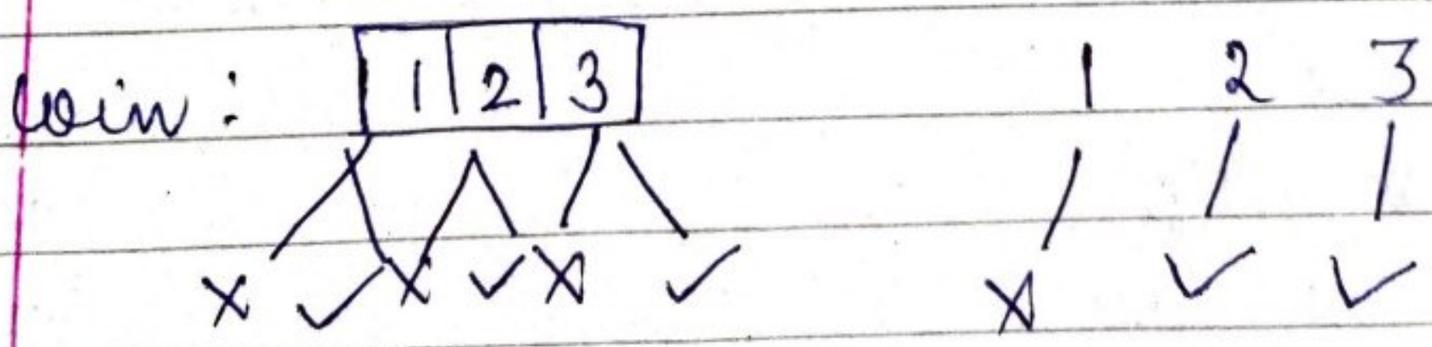
Use the coin array unlimited time & get the sum 5. Now find the max^m no of ways in which we can get 5.

	O/P
2 + 3	5.
1 + 2 + 2	5.
1 + 1 + 3	5.
1 + 1 + 1 + 1 + 1	5.
1 + 1 + 1 + 2	5

5 ways.

Q Why knapsack?

→ Every coin has a choice to include or not



item — wt[ ] ✓

item — val[ ] ✗ (when one array is given)

Matching

wt[] ⟶ coin[]
W ⟶ sum

1	2	3.

Sum 5 : 1 + 1 + 3

one item used many times

If taking one item many times does the sum allows then it is unbounded knapsack.

Subset sum

1	2	3	5

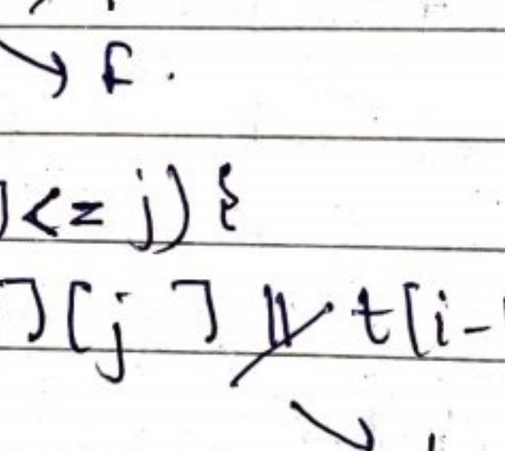
Sum : 8 → T / F.

Subset sum.

```
if (arr[i-1]<=j) {
t[i][j] = t[i-1][j] || t[i-1][j-arr[i-1]]
} else {
t[i][j] = t[i-1][j];
}
```

→ for count we remove || by +.

count / no of ways



We need to find maxim. no of ways.

Matching → Knapsack → 0-1
                    → Unbounded

wt → coin
W → sum.

if (coin[i-1] <= j)
  t[i][j] = t[i-1][j] + t[i][j - coin[i-1]].
else
  t[i][j] = t[i-1][j]

sum →

size ↓	0	1	2	3	. . .
0	1	0	0	0	0 .
1	1				
2	1				
3	1				

Coin change - II (Min^m no of coins)

coin[ ] = [1 | 2 | 3]
sum = 5

2+3 → 5 → 2 coins
1+2+2 → 5 → 3 coins
1+1+1+2 → 5 → 4 coins
1+1+3 → 5 → 3 coins
1+1+1+1+1 → 5. → 5 coins

} find min^m. no. of coins to make sum as 5.

coin[ ] = | 1 | 2 | 3 |

Sum = 5

O/p : 2 (2+3 = 5)

Initialisation: t[n+1][w+1]

↓

t[n+1][sum+1]

wt → coin[ ] = n

val → x

W → sum.

n=3
sum=5

Sum (S) →

	0	1	2	3	4	5	6
0		INT_MAX -1	INT_MAX	INT_MAX	INT_MAX	INT_MAX	INT_MAX
1	0						
2	0						
3	0						

→ 2nd Row initialize

Size (n) ↓

↓ code

sum = 3
coin[ ] = 1

Initialisation + Twist

coin[ ]: empty
sum : 1 } INT_MAX (∞ coins)

coin[ ] : empty
sum : 0 } INTMAX - 1

coin[ ] : 1
sum : 0 } 0 coins

coin[ ] : 12
sum : 0 } 0 coins

sum = 5.
arr = [3, 5, 2]

when
arr [3]
sum = 4

For quer

$\frac{3}{3} = 1$

$\frac{4}{3} = INT\_MAX$

j	0	1	2	3	↑	5
0		INT	MAX	-1		
i = 1		1				
2	0					
3	↓					
4	↓					

sum = 1] 1
coin = 1

sum = 2
coin = 1 (1)

$\frac{4}{3} = INT\_MAX$

$\frac{j}{arr[0]} = \frac{3}{3} = 1$

size will always remains 1 so we will use arr[0] as it is the second row

for second row

```
for(int i=1; j<sum+1; j++) {
    if (j % arr[0] == 0)
        t[i][j] = j / arr[0]
    else
        t[i][j] = INT_MAX-1
```

$\frac{j}{arr[0]} \rightarrow \frac{3}{3} = 1$

$\frac{4}{3} = INT\_MAX$

Code variation

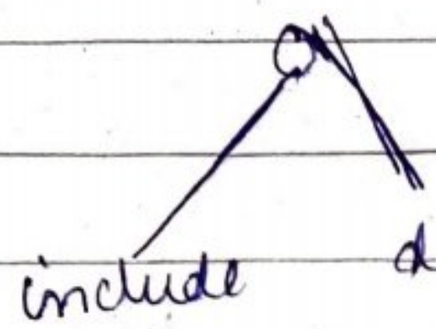
1st row → INT_MAX 1st col → 0

if (coin[i-1] <= j)

t[i][j] = min (t[i-1][j] , t[i][j-coin[i-1]] + 1)

else

t[i][j] = t[i-1][j]




```
int minCoins( int coins[], int M, int V) {
    // m = arrsize   v = sum.
    int dp [M+1][ V+1];
    for (int i=0; i<M+1; i++){
        for(int j=0; j<V+1; j++) {
            if (j==0) dp[i][j] = 0;
            if (i==0) dp[i][j] = INT_MAX-1;
        } }

    for(int i=1, j=1 ; j<V+1; j++) {
        if(j % coins[0] ==0) dp[i][j] = j/coins[0];
        else                 dp[i][j] = INT_MAX-1;

    for(int i = 2; i<M+1; i++) {
        for(int j = 1; j<V+1; j++) {
            if (coins[i-1]<=j)
                dp[i][j] = min(dp[i-1][j], 1+dp[i][j-coins[i-1]]
            else dp[i][j] = dp[i-1][j];
        }}
    if(dp[M][V] == INT_MAX-1) return -1;
    return dp[M][V];
}.
```

Longest Common Subsequence.

1) Longest common substring
2) Print LCS
3) Shortest common supersequence
4) Print SCS
5) Min^m no of insertion and deletion a→b
6) Largest repeating subsequence
7) Length of largest subsequence of a which is a substring is b.
8) Subsequence pattern matching
9) Count how many times a appear subsequence in b
10) Largest Palindromic subsequence
11) Largest palindromic substring
12) Count of palindromic substring
13) Minimum no of deletion in a string to make it a palindrome
14) Minimum no of insertion in a string to make it a palindrome

Longest Common Subsequence (Recursive)

Problem Statement

x : a b c d g h
y : a b e d f h r.

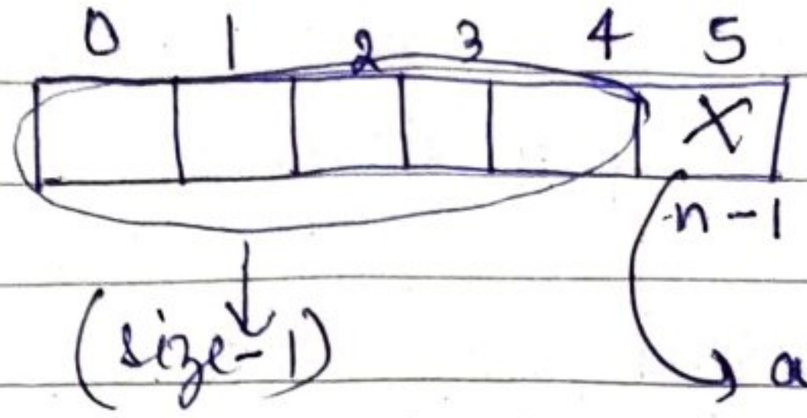
String x = abcdgh
String y = abedfhr

o/p → abdh
→ 4 (length of string)

Subsequence → longest common subsequence - abdh.
substring → longest common substring - ab

Recursive approach: Base cond^n + Choice diagram + i/p small

i/p (x, y)




```
func(x, y)
{
  func(x, y)
    ↓
  smaller x, y
}
```

Base cond^n: Think of the smallest valid input.

x: ⬭ → n

y: ⬭ → m.

n = 0 m = 0 LCS = 0
(empty string)

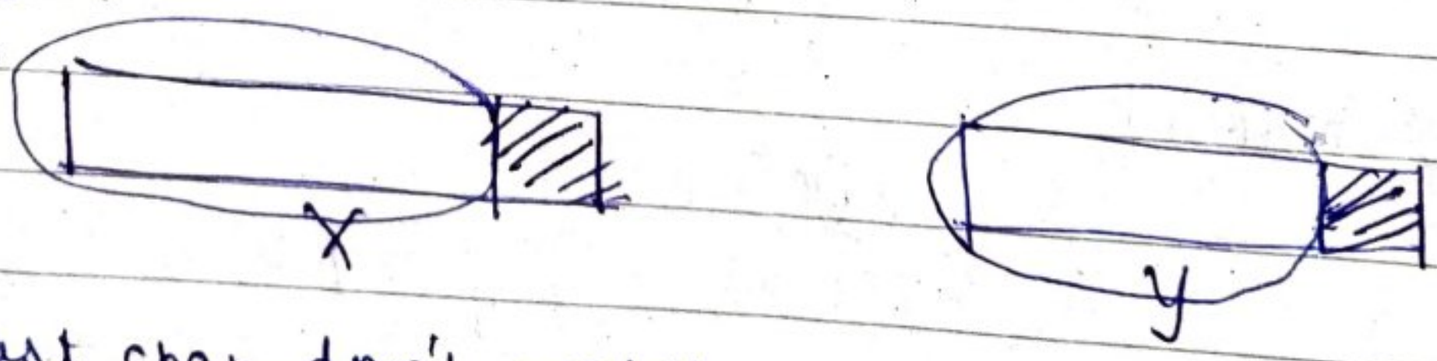
```
if(n == 0 || m == 0)
    return 0
```

Choice diagram:

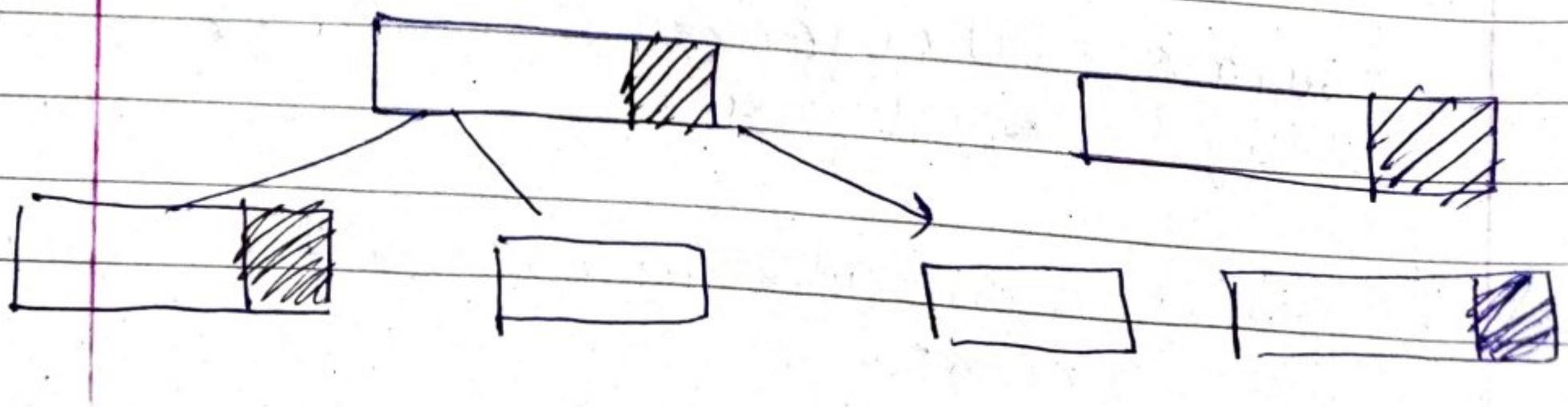
x : abcdgh

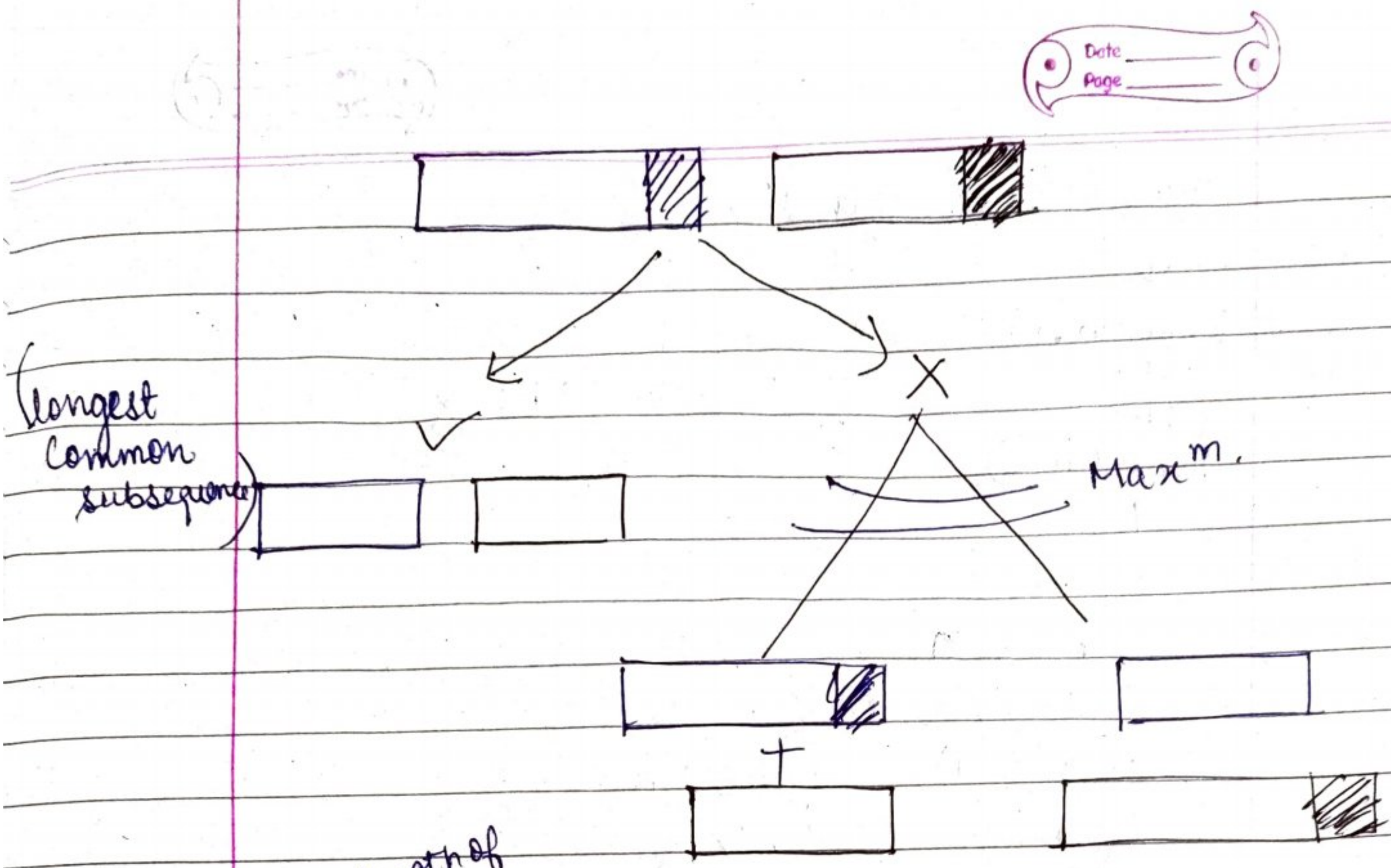
y : abedfhr

① when last char matches



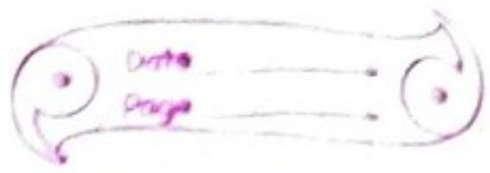
② when last char don't match






```
int LCS (string X, string Y, int n, int m) {

    if (n==0 || m==0) return 0;

    if (X[n-1] == Y[m-1])
        return 1+LCS(X, Y, n-1, m-1);

    else
        return max (LCS(X, Y, n, m-1),
                    LCS(X, Y, n-1, m));
}
```

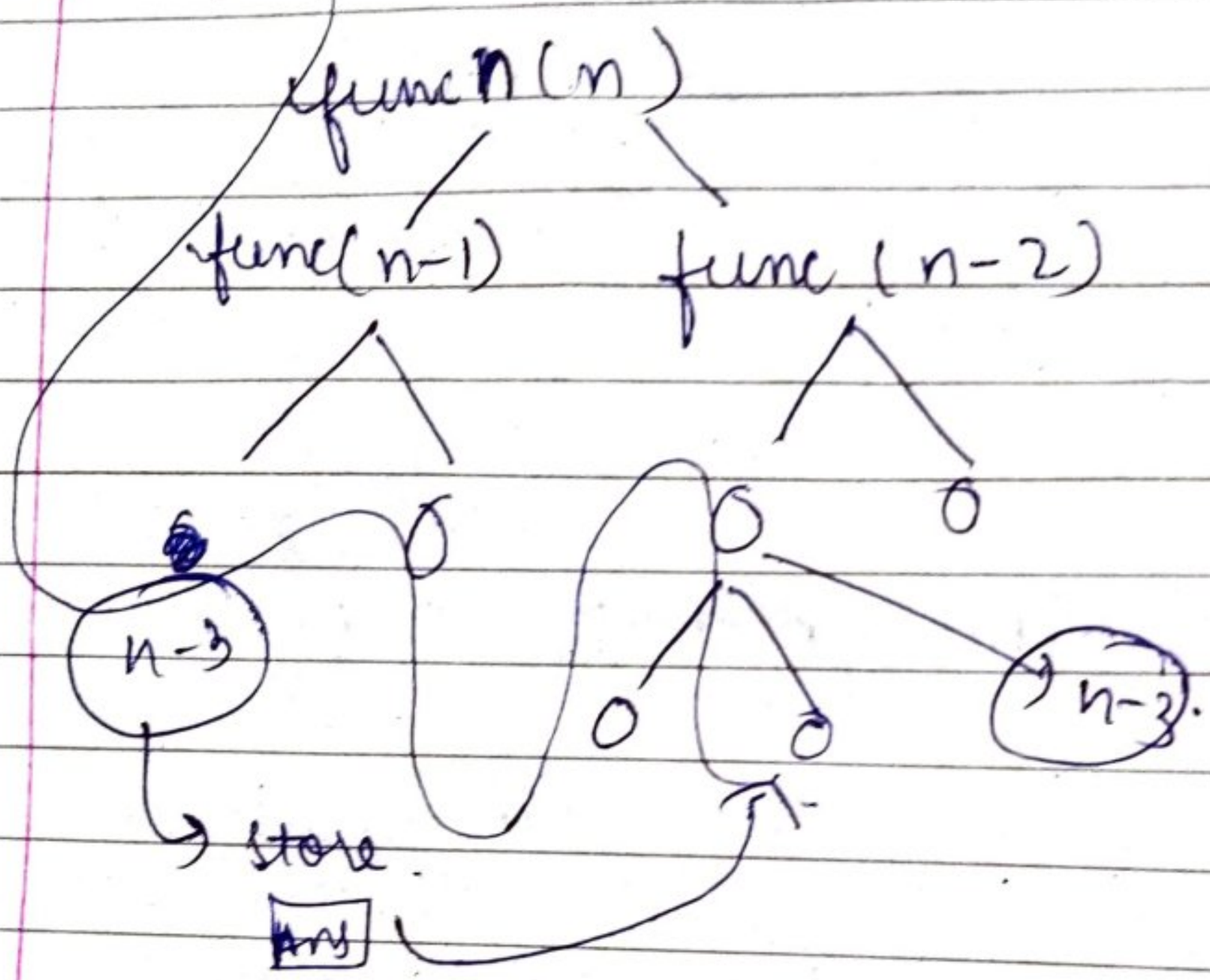
LCS memoization (bottom up approach)



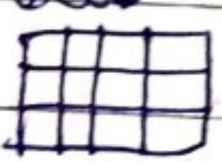
fact(n)
↓
func(n-1)
|
func(n-2)
|

```
func( )
{
func( )
}
```

→ DP X



R.C + table



"AXyT" "AyZx"

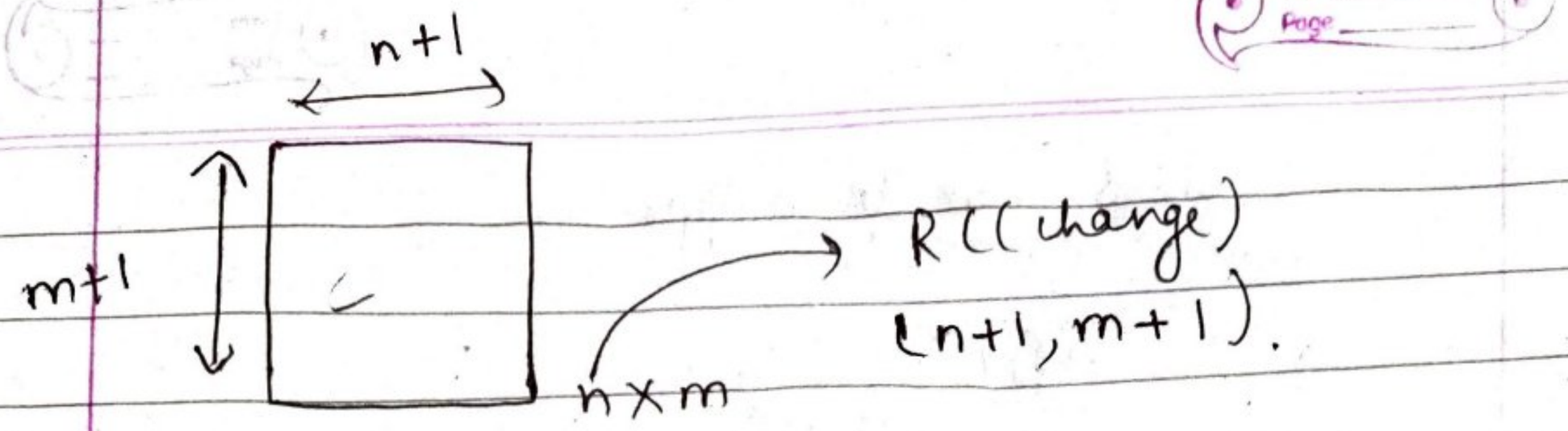
LCS(AXy, AyZx)

LCS(AXyT, AyZ)

LCS(AX AyZ) LCS(AXy, AyZ) LCS(AXy, AyT)

LCS(AXyT, Ay)

Sub - Problem repeating twice




```
int t[100][100];
```

-1	-1	-1	-1
-1	-1	-1	-1
-1	-1	-1	-1

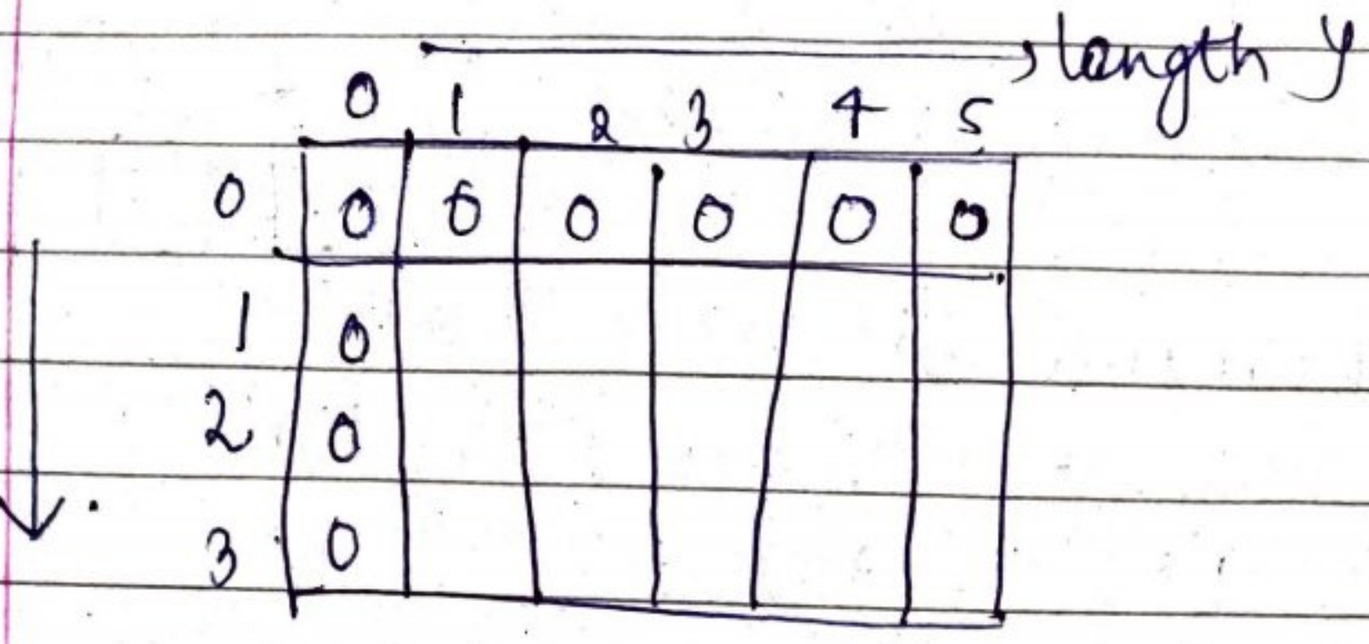
```
int main(){
    memset(t, -1, sizeof(t));
    LCS( )
}

int LCS( string X, string Y, int m, int n) {

    if (n==0 || m==0)
        return 0;
    if (t[m][n] != -1)
        return t[m][n];

    if (X[m-1] == Y[n-1])
return t[m][n] = 1+LCS(X, Y, m-1, n-1)
    else
return t[m][n] = max( LCS(X, Y, m, n-1),
                      LCS(X, Y, m-1, n))

}
```

exponential → $O(n^2)$

LCS Top-down Approach

x : a b c d a f
y : a c b c f

o/p → 4 (abcf)



Base cond^n → initialization
(RS) (Top down)

```
int LCS( string X, string Y, int n, int m) {
    int dp[m+1][n+1];

    for(int i=0; i<m+1; i++)
        for(int j=0; j<n+1; j++)
            if (i==0 || j==0)
                dp[i][j] = 0

    for (int i=1; i<m+1; i++)
        for (int j=1; j<n+1; j++)
            if( X[i-1] == Y[j-1])
                dp[i][j] = 1+ dp[i-1][j-1]
            else
                dp[i][j] = max(dp[i][j-1], dp[i-1][j]);
```

return dp[m][n];
}

Longest common Substring

I/p a :→ abcde
b :→ abfce

O/p → 2 (ab) ab, c, e

longest = ab

a b c d e
a b f c e

#② 0 1 1

length = 0

m →

0	0	0	0
0			
0			
0			

n ↓

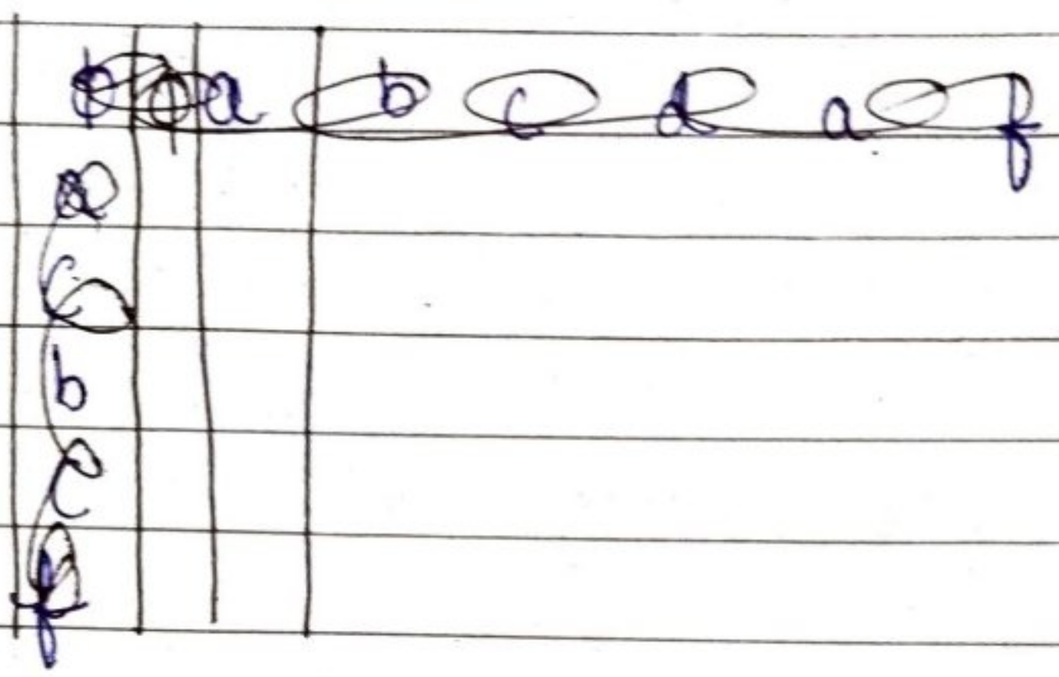
dis continuity = 0

```
if (a[i-1] == b[j-1])
    dp[i][j] = dp[i-1][j-1] + 1
else
    dp[i][j] = 0
// for max^m val.
int maxi = 0;
for (i = 0 → n+1)
    for (j = 0 → m+1)
        maxi = max(maxi, dp[i][j])
return maxi;
```

Print LCS btwn 2 string

a : @ c b c f
b : @ b c d a f

O/p → a b c f



	ϕ	a	b	c	d	a	f
ϕ	0	0	0	0	0	0	0
a	0	1	1	1	1	1	1
c	0	1	1	2	2	2	2
b	0	1	2	2	2	2	2
c	0	1	2	3	3	3	3
f	0	1	2	3	3	3	4

Now, we need to print LCS.

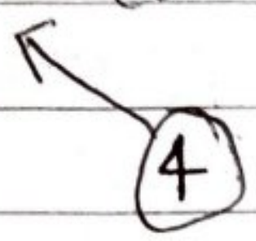
"f c b a"

abcf ← reverse.

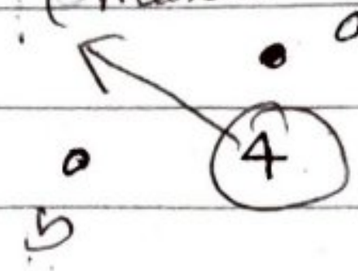
if equal i, j → (i--, j--)

if not equal i,j → max ((i-1, j), (i, j-1))

When equal
(i--, j--)



when not equal
(max betn a & b)



add in string : nothing to be added

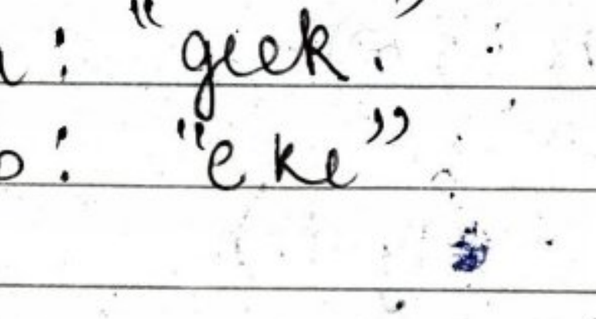
```
int i = m , j = n ;
String s = "" ;
$
while (i>0 && j>0)
{ a b
if (t[i-1] == t[j-1])
{
s.push_back ( t[i-1])
i--;
j--;
}
else {
if ( t [i][j-1] > t[i-1][j])
j--;
else
i--;
}
reverse ( s. begin( ), s.end ( ) ) ;
```


```
if (a[i-1] == b[j-1])
    S.push_back( a[i-1] )
        i--;
        j--;

else
    move in the direc^n of max'm
    reverse ( s.begin(), s.end() )
```

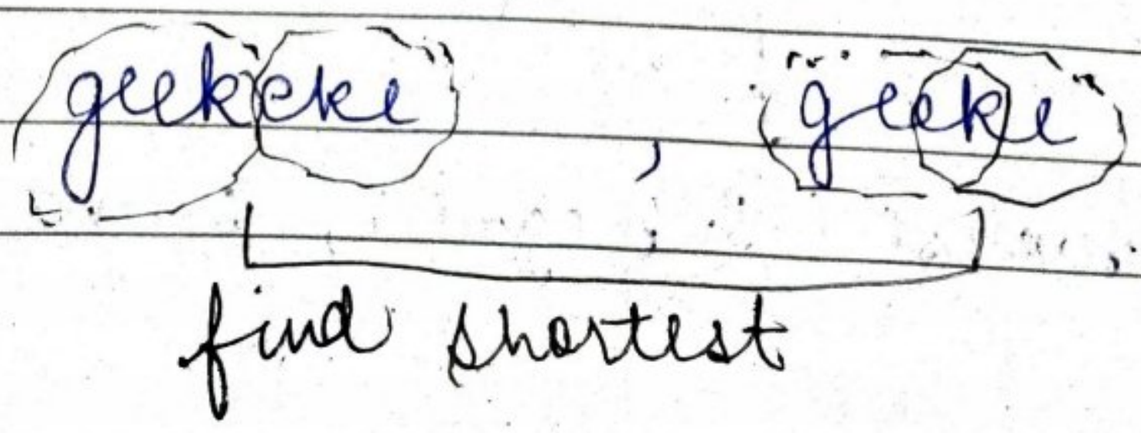
24. Shortest Common Supersequence

a : "geek"
b : "eke"

1) Problem statement
→ 2 string are given. We need to merge the string such that we get both the subsequence.
Mix 2 sequence to make it a super sequence.



o/p



Subsequence → Sequence of events

S_1 S_2 S_3

Order should be maintained but don't need to be continuous always

a : A G G T A B

b : G X T X A Y B

Super sequence : A G G T G X A B T X A Y B

A G G X T X A Y B → shortest supersequence

A G G T A B G X T X A Y B

Give the length in o/p.

The one letter which is common write once.

A G G T A B

G X T X A Y B

↓

G T A B is common in both i.e LCS.

Now thinking the brute force approach we can merge both the string & then subtract GTAB.

a: AGGTAB b: GXTXAYB

length of SS
a+b

Worst: a+b = AGGTABGXTXAYB

case

↓

AGGTABGXTXAYB – GTAB

↓

AGGXTXAYB

Shortest length = (m+n) – LCS

25. Minimum no of insertion & deletion to convert a string a to string b.

a: heap o/p = 1 (Insertion)
b: pea 2 (deletion)

a ⟶ b

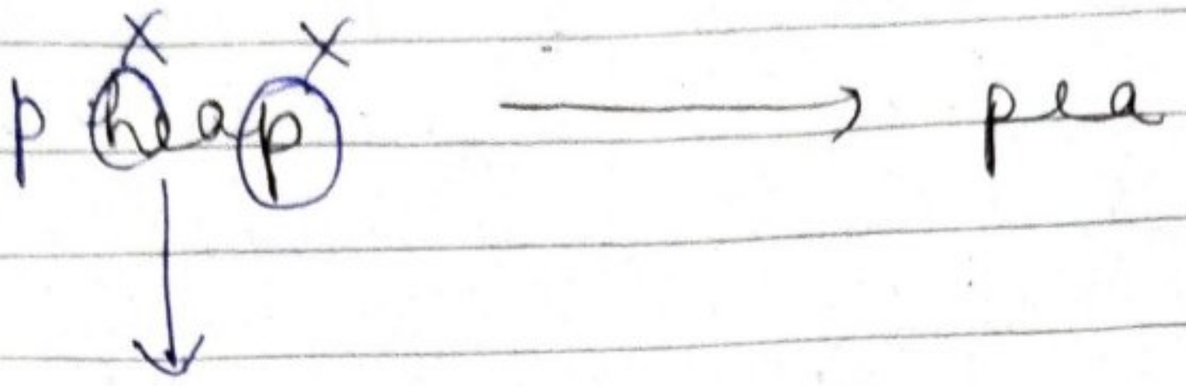
heap ⟶ pea

insertion deletion

When to use LCS.

	i/p	o/p	o/p	
LCS	a: ◯ b: ◯	LCS	int	→ Pattern matching algorithm
Given Ques	a: ◯ b: ◯	insertion/ deletion	int	

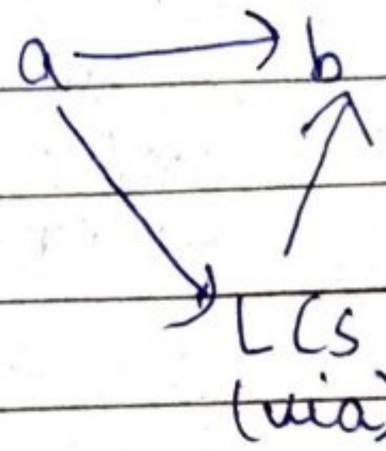
Convert heap to pea.



"ea" is the LCS of both string

heap → pea

Insertion/
deletion



heap —(2 deletion)→ ea (LCS) —(1 insertion (p))→ pea

No of deletion = a length - LCS
No of insertion = b length - LCS

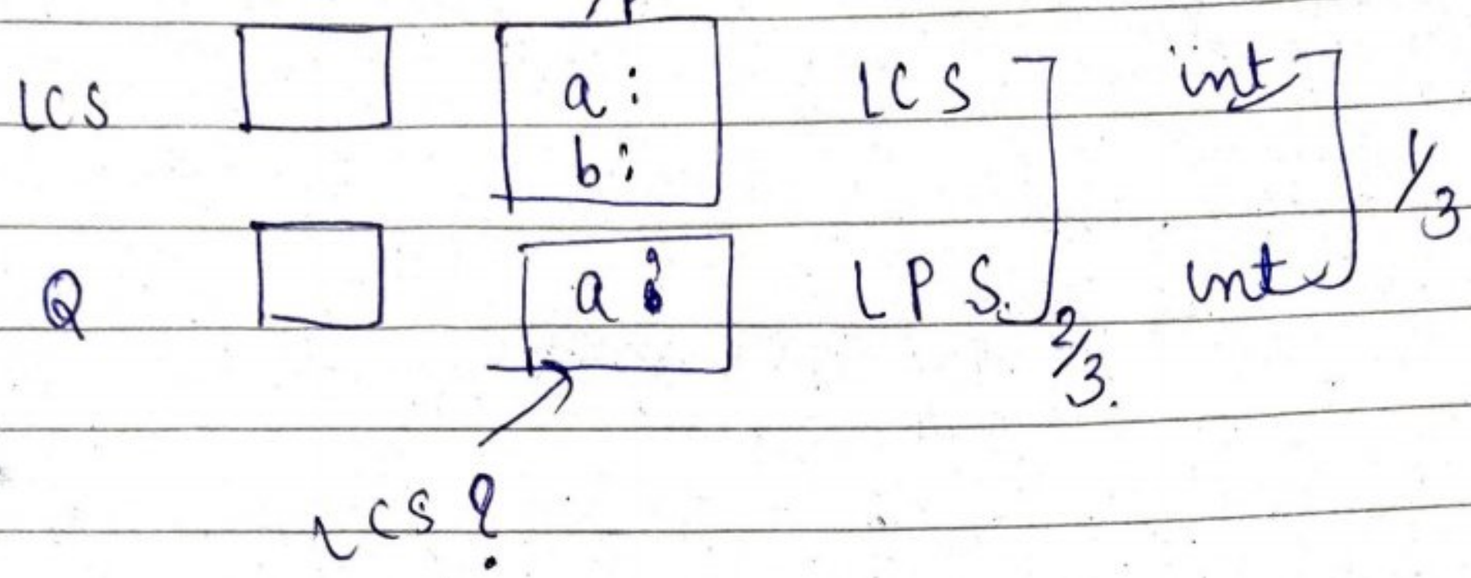
26. Longest Palindromic Subsequence

Problem Statement : S : agbcba
O/P = 5

5 ← longest [abcba ✓
bcb ✓
b ✓

O/p → 5 (abcba)

2) Is it LCS ?



LCS I/p a, b

LPS I/p a b=func(a)
(hidden string / redudant)

"agbcba"

↓

LPS

↙

a → agbcba

b → reverse(a) abcbga.

↓

LCS

abcbga
agbcba
> abcba

LPS(a) ≡ LCS (a, reverse(a))

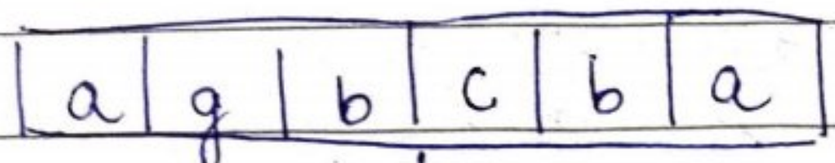
LPS(agbcba) → return (agbcba,

28. Minimum no. of deletion in a string to make it palindrome

a : "agbcba"

o/p = 1

S : agbcba



No of deletions (Minimize)

→ New string (palindrome)

~~a~~g~~b~~cb~~a~~

bcb (palindrome) 3

agcba (not palindrome)

agbcba

~~a~~g~~b~~cb~~a~~

bcb (3) (LPS)

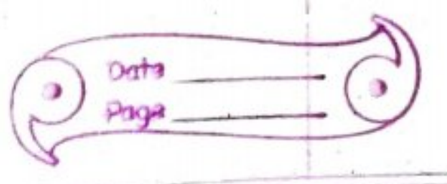
~~a~~g~~b~~c~~b~~a~~

C (5) (LPS)

agbcba

abcba (1) LPS

Minm.



↑ length of LPS ∝ $\frac{1}{\text{no of deletions}}$ ↓

LPS → ellipse → Min^m no of deletions.

ellipse → Palindromic subsequence.

↑ length of LPS ∝ $\frac{1}{\text{no of deletions}}$ ↓

(longest palindromic subsequence)

agbcba

↓

LCS(s, reverse(s))

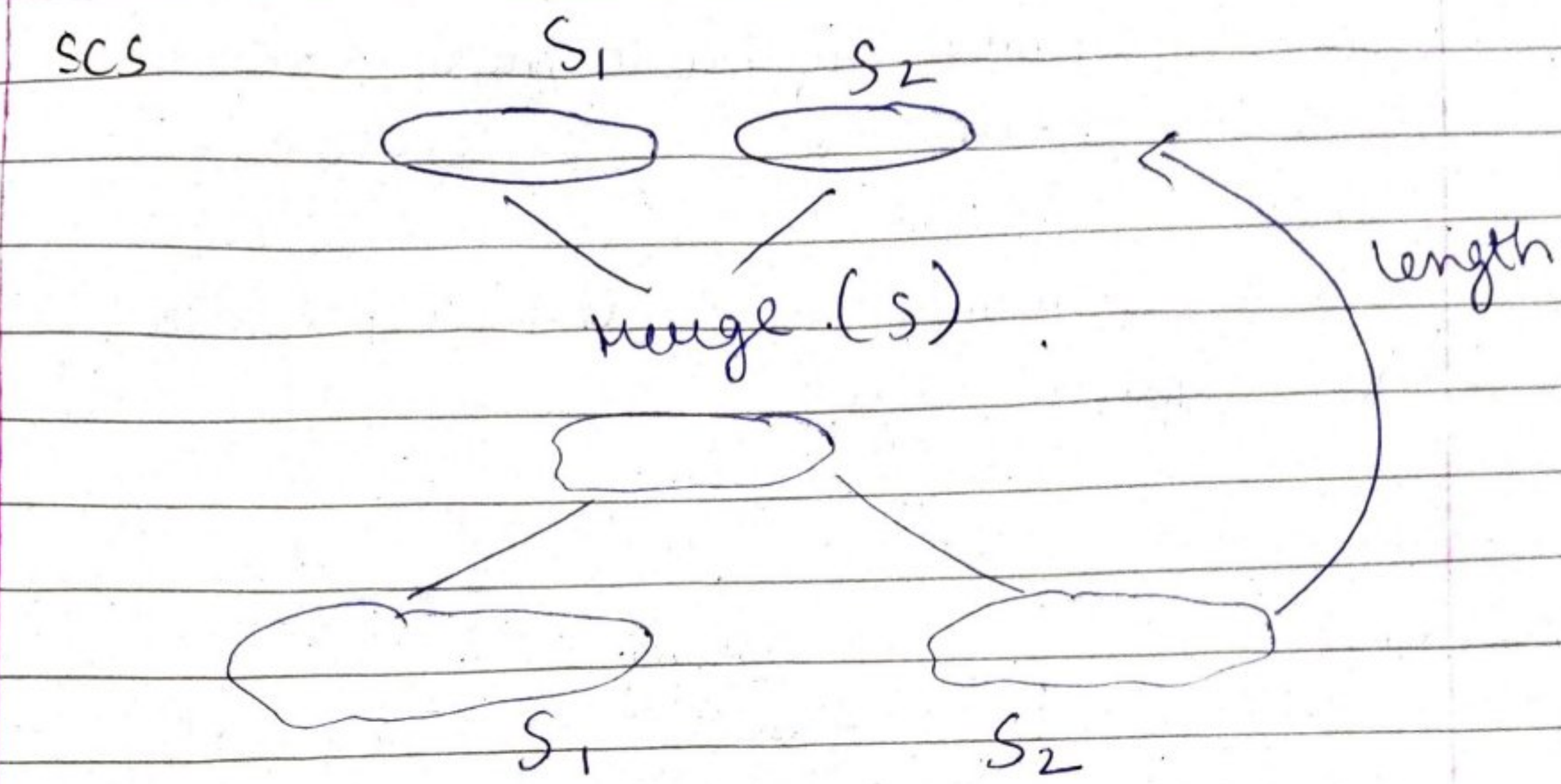
↓

abcba (LPS)

Min no of deletion = $\underset{(\text{length})}{S} - \underset{(\text{length})}{LPS}$

29. Print Shortest Common Subsequence

i/p a: acbcf o/p: acbcdaf
b: abcdaf



worst case

a c b c f a b c d a f

a c b c d a f → print the whole string

m+n

a c b c f a b c d a f

now

m+n - LCS.

LCS (similar) → SCS

Print LCS Print SCS

		a	b	c	d	a	f	
	0	0	0	0	0	0	0	→ LCS
a	0	1	1	1	1	1	1	
c	0	1	1	2	2	2	2	
b	0	1	2	2	2	2	2	
c	0	1	2	3	3	3	3	
f	0	1	2	3	3	3	4	

LCS → common ni hai to move karo jao
common hai to print karlo.

SCS → common hai to ek bar print karo
else print karo.

LCS.

```
string s = " ".
while (i>0 && j>0){
  if ( a[i-1] == b[j-1] ){
    s.push_back(a[i])
      i--;
      j--;
  }
  else{
  if(t[i][j-1] > t[i-1][j])
      j--;
  else if (t[i-1][j] > t[i][j-1])
      i--;
  }
```

SCS

```
string s = " ";
while (i>0 && j>0) {
  if(a[i-1] == b[j-1]) {
    s.push_back (a[i])
      i--, j--; }
  else {
  if (t[i][j-1]> t[i-1][j])
    s.pushback (b[j-1]);
  else  j--;

else if( t[i-1][j] > t[i][j-1])
    s.push_back (a[i-1]);
      i--, }
while(i>0)
  s.push-back (a[i-1]); i--;
while(j>0)
  s.push_back ( b[j-1])
    j--;
```

1) Now in SCS code we include the letter on moving at i-1 or j-1.
2) Now in (i>0 && j>0) we need to change because in case of LCS we don't need to stop at the topmost but in SCS we need to.

In case of LCS

	ϕ	a	b	f
ϕ	0	0	0	0
a	0			
c				
b	0			
d	0			

Some can't stop here.

	ϕ	a	b	b
ϕ	0	0	0	0
a	0			
c	0			
b	0			
d	0	.		

a c
" "] LCS (" ")

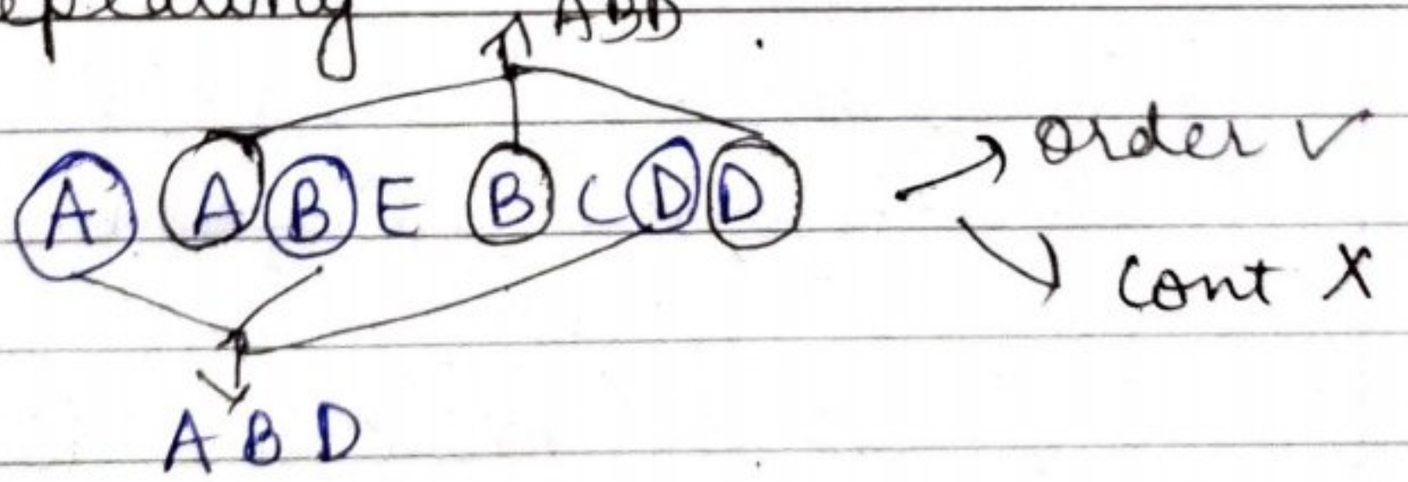
a c
" "] SCS ≠ (ab

Longest repeating Subsequence

Str = "A A B E B C D D"

O/P = 3

Problem Statement - Find a subsequence which is repeating



ABD (2X)
AB (2X)] longest → "ABD"

O/P = 3

S = A A B E B C D D

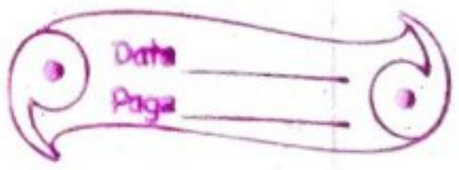
E → 3
C → 5

0 1 2 3 4 5 6 6
A A B E B C D D
A A B E B C D D

LCS = A A B E B C D D

A A B B D D
once ← → ans

E & C are occuring once.
letter at the same index don't take



E → 3 → i
C → 5 → j

i == j X

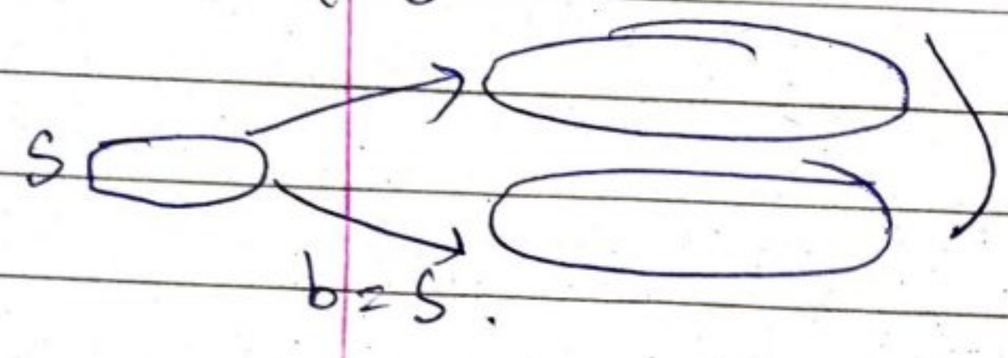
A → 0, 1
A → 0, 1

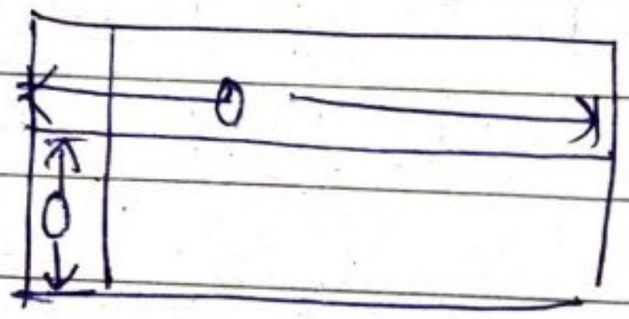
B → 2, 4
B → 2, 4

diff index
same index

So don't take A of same index take from other index.

Sum up:





We can't take same letter for both the string during LCS i.e. (i == j)

```
if (a[i-1] == b[j-1] && i != j)
    t[i][j] = t[i-1][j-1] + 1
else
    t[i][j] = max (t[i][j-1],
                   t[i-1][j])
```

31. Sequence Pattern Matching

a = "AXY"
b = "ADXCPY"

O/P → T/F

Problem statement: Is string "A" a subsequence of "B"

a : AXy b : ADXCPY

b : (A) D (X) C P (y)

a : AXy

O/p → True.

LCS : AXy

I/p

LCS	a b	sequence	int
Q	a b	sequence	T/F

I/p

b : ADXCPY b : ADXCPy

a : AXy a : AXyz

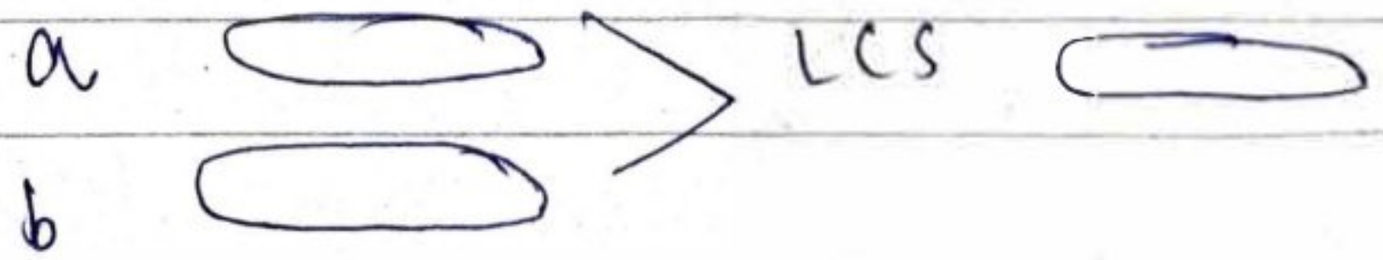
LCS : AXy (LCS == a) LCS : AXy (LCS != a)

if (LCS == A) return true

else

return false

Can it be done using length. ?



a: AXY (3)

b: ADXCPZ(6)

LCS = 0 to min(m,n)

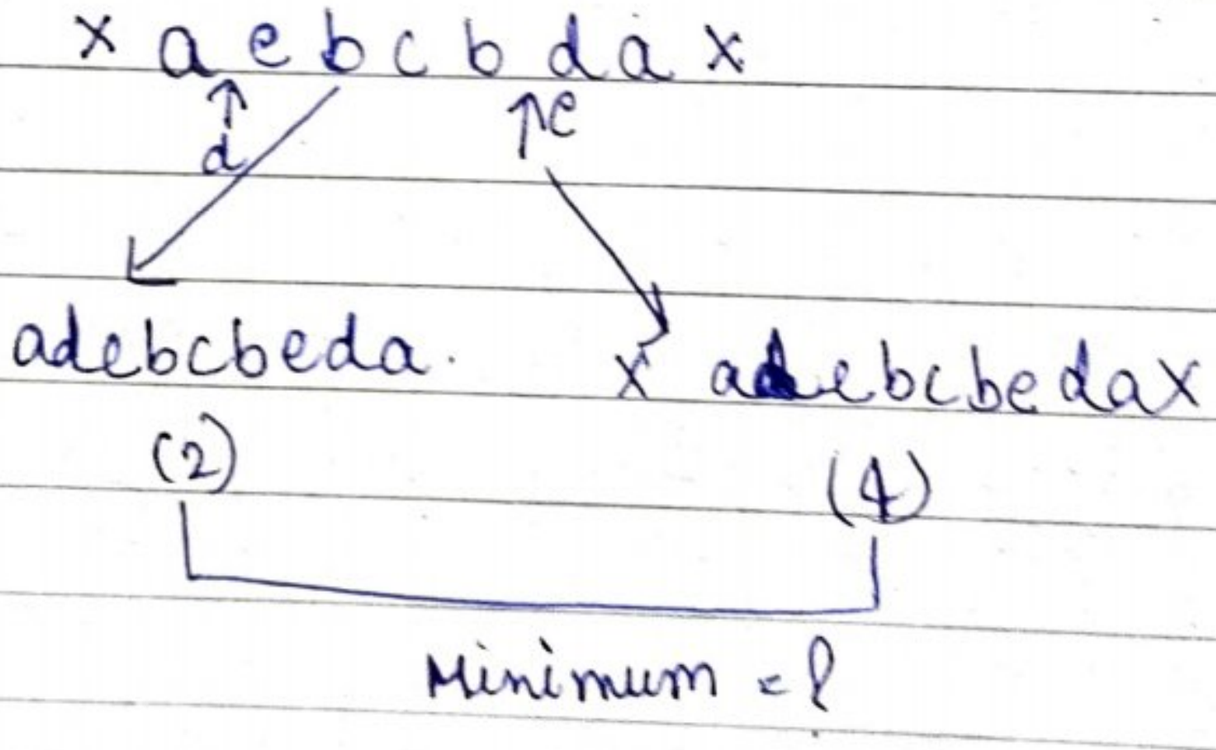
LCS = 2

```
if (LCS == a.length())
    return true
else
    return false.
```

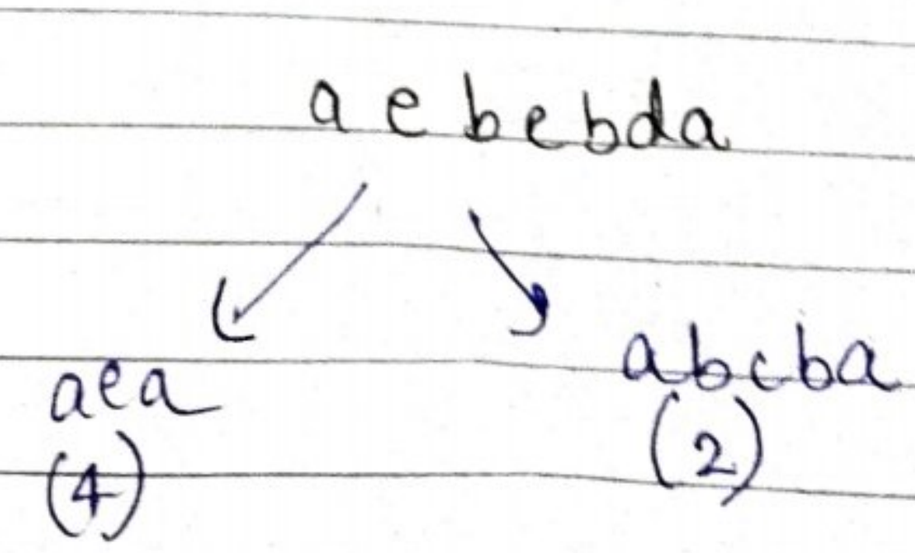
32. Min^m no of insertions in a string to make it palindrome

I/p: S: "aebcbda"

O/p → 2



Min^m no of deletion to make a string palindrom.



S.length() → LPS. (deletion)

S : "aebcbda"
(e, d)

Palindromic string

[a] e [b c b] d [a]

e ✗ { delete
d ✗ { e & d

a e ✓ { insert e & d & make their pair
d ✓

no of insertions = no of deletions
||
S.length - LPS.

Deletion → single (pair) → ✗
Insertion → single (pair) → ✓

33. MCM (Matrix Chain Multiplication)

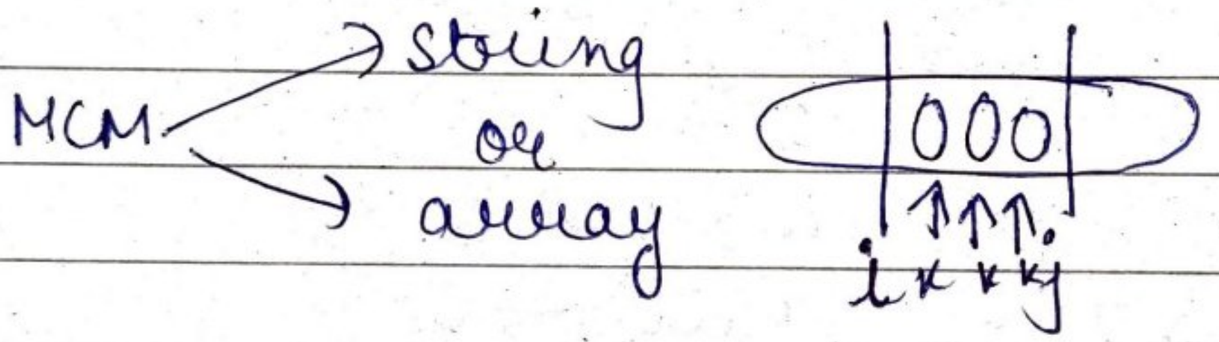
1) MCM
2) Printing MCM
3) Evaluate Exp. to True / Boolean Parenthesization
4) Min / Max value of an Expr.
5) palindrome partitioning
6) Scramble string
7) Egg Dropping problem

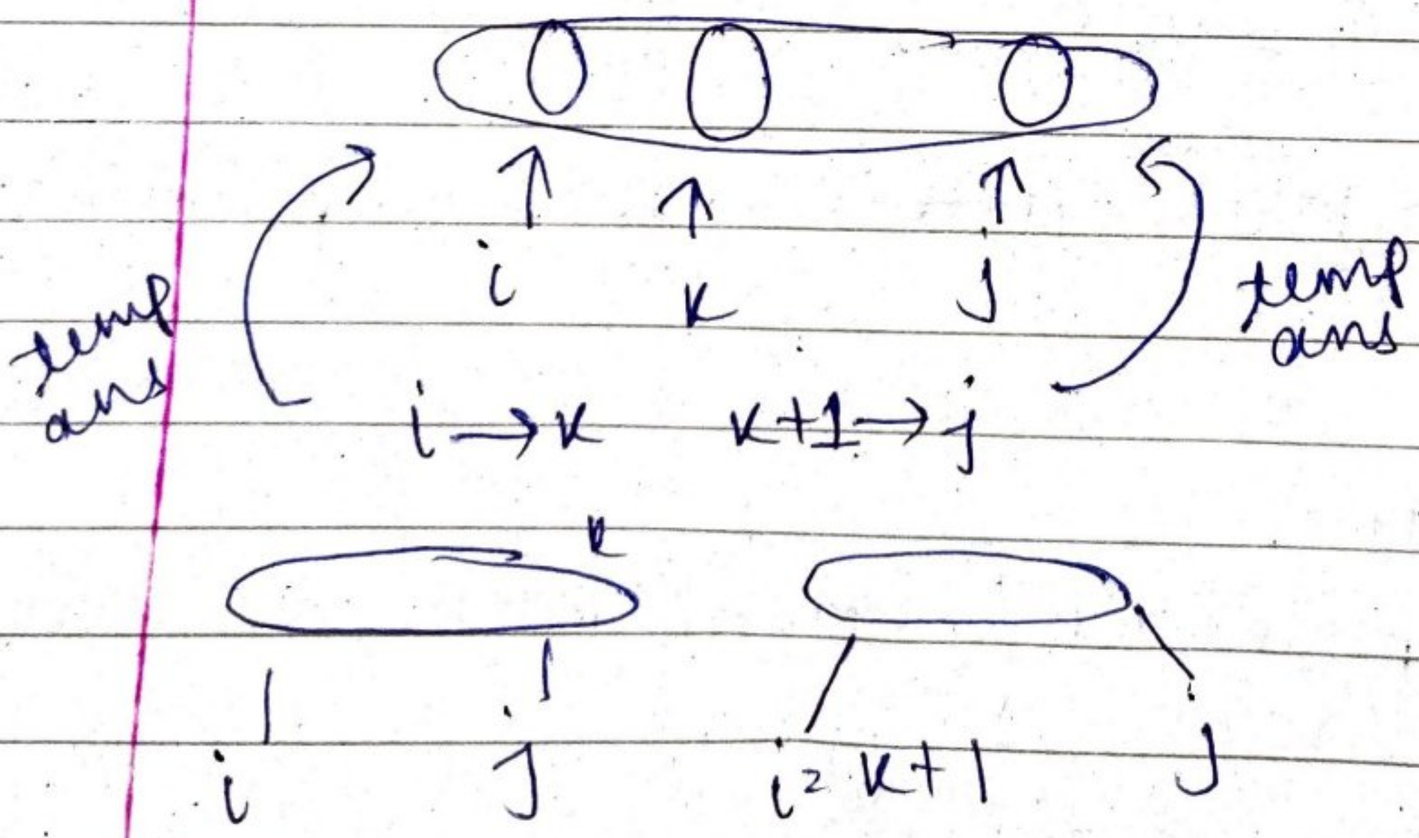
MCM format

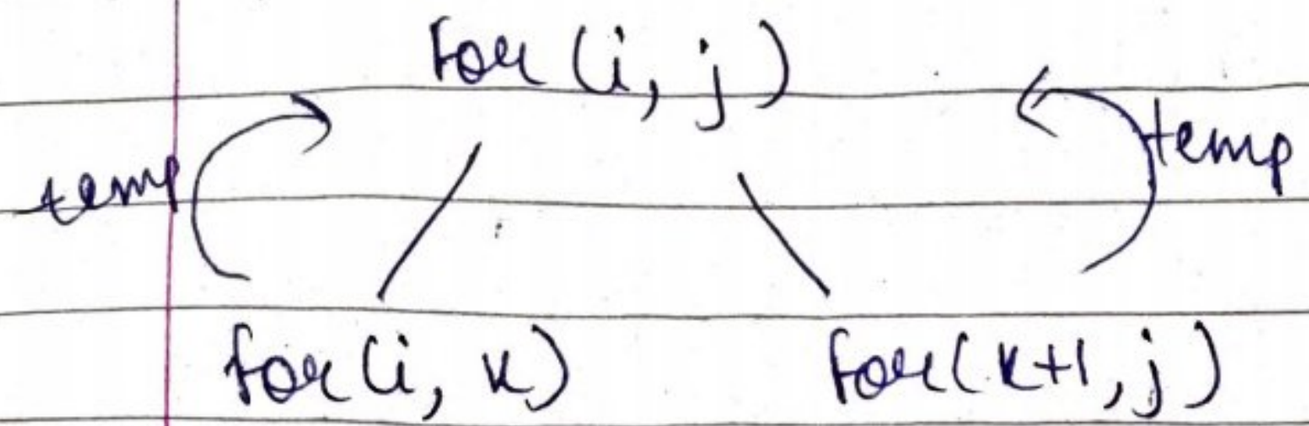
1) MCM
2) Printing MCM
3) Evaluate Exp^{r} to True / Boolean parenthesization
4) Min / Max value of an Expr
5) Palindrome partitioning
6) Scramble string
7) Egg dropping problem

Identify → MCM → basic format

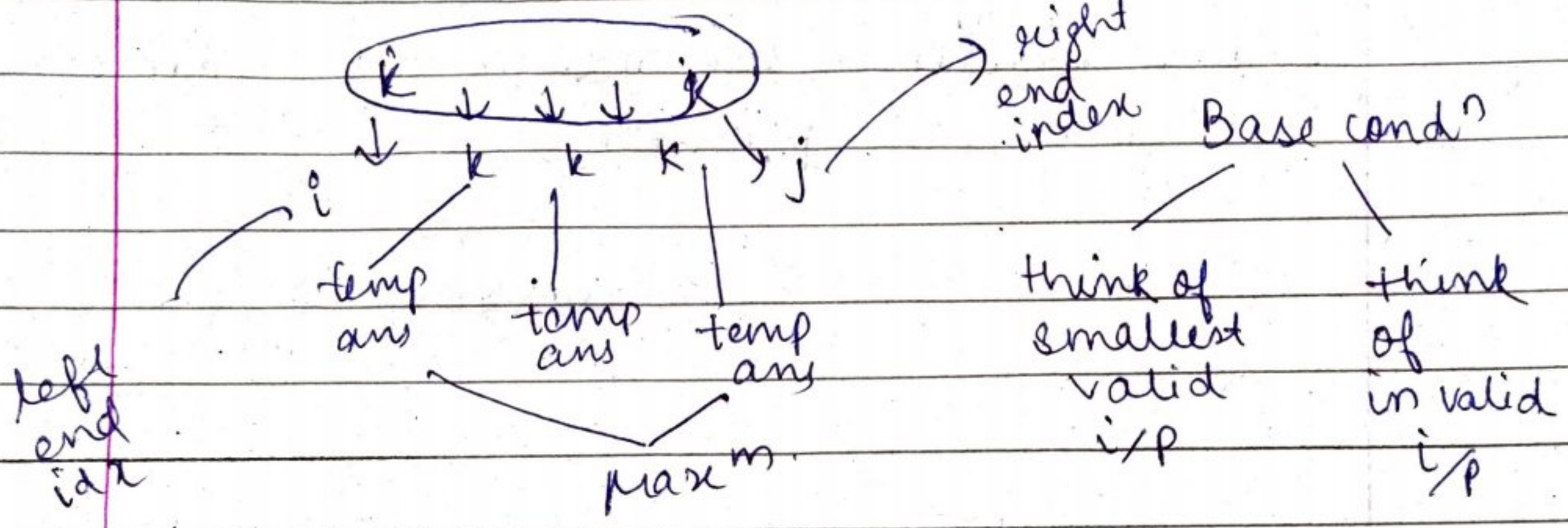
Identification + format







ans ← f (temp ans)




```
int solve ( int arr[], int i, int j)
  {
    if (i>j)  →  this may be diff acc^n. to ques
      return 0;

    for (int k=i ; k<j ; k++)
      {
        // calc. temp ans
        temp ans = solve(arr, i, k) +
                   solve(arr, k+1, j);

        ans =  func (temp ans)
      }
```

MCM (Recursive)

Problem

Statement : arr[] = {40, 20, 30, 10, 30}.

A_1 (m×n), A_2 (n×m), A_3 (m×n), A_4

Some matrix are given like A_1, A_2, A_3, A_4 we have to multiply the matrix to reduce the cost (no of multiplications)

b = c (then only we can multiply)

$\begin{bmatrix} & \\ & \end{bmatrix}_{2\times 3}$ $\begin{bmatrix} & \\ & \\ & \end{bmatrix}_{3\times 6}$

a×b b×d c

2 × 3 3 × 6

↓ ↓ ↓

2 3 6 = 36 cost (no of multiplications)

A_1 A_2 A_3 A_4

dimension [] = {x, y, z, w}.

$A_1\ A_2\ A_3\ A_4$

$(A_1\ (A_2 A_3))A_4 \quad ((A_1 A_2)(A_3 A_4)) \quad A_1(A_2(A_3\ A_4))$

$C_1 \qquad C_2 \qquad C_3$

$A \rightarrow 10 * 30$

$B \rightarrow 30 * 5$

$C \rightarrow 5 * 60$

$(A\ B)C = 10 * 30 * 30 * 5$

$(A\ B) \downarrow \quad C \downarrow$

$10 * 5 \qquad 5 * 60.$

$10 * \quad * 5 \qquad 5 * 0$

$(A\ B)C = (10 * 30 * 30 * 5)\ 5 * 60.$

$= 10 * 5 * 30 * 5 * 60$

$= 10 * 30 * 5 * 60 * 10 * 5$

$= 4500.$

$A_1 A_2 A_3\ A4$

↓ do parenthesiation in such a way cost is the least . min

$A_1(A_2 A_3)A_4 \quad C_1$

$A_1 A_2)(A_3 A_4) \quad C_2 \quad A(B\ C)$ → min cost

$A_1(A_2(A_3\ A_4)) \quad C_3.$

$30 * 5 \qquad 5 * 60$

$B\ C \rightarrow 30 * 5 * 60.$

$A * BC \rightarrow 10 * 30 \qquad 30 * 60$

$10 * 30 * 60$ ↓ t

Cost

Bracket lagane par alag alag cost aa rahi hai.

0	1	2	3	4
40	20	30	10	30

arr[] : {40, 20, 30, 10, 30}

arr size n n-1 matrix

$A_1 \rightarrow 40 * 20$
$A_2 \rightarrow 20 * 30$
$A_3 \rightarrow 30 * 10$
$A_4 \rightarrow 10 * 30$

} $A_1 * A_2 * A_3 * A_4$

Cost (minimum)

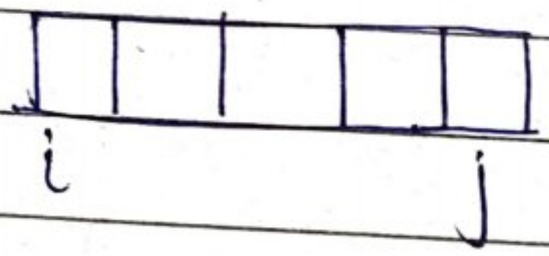
⇓

(no of multiplication should be less)

$A_i \rightarrow arr[i-1] * arr[i]$

$A[i] = arr[0] + arr[1]$

Next step → Format



(dimension)

arr[] : |40|20|30|10|30|

0 1 2 3 4

$A_1\ A_2\ A_3\ A_4$

↓

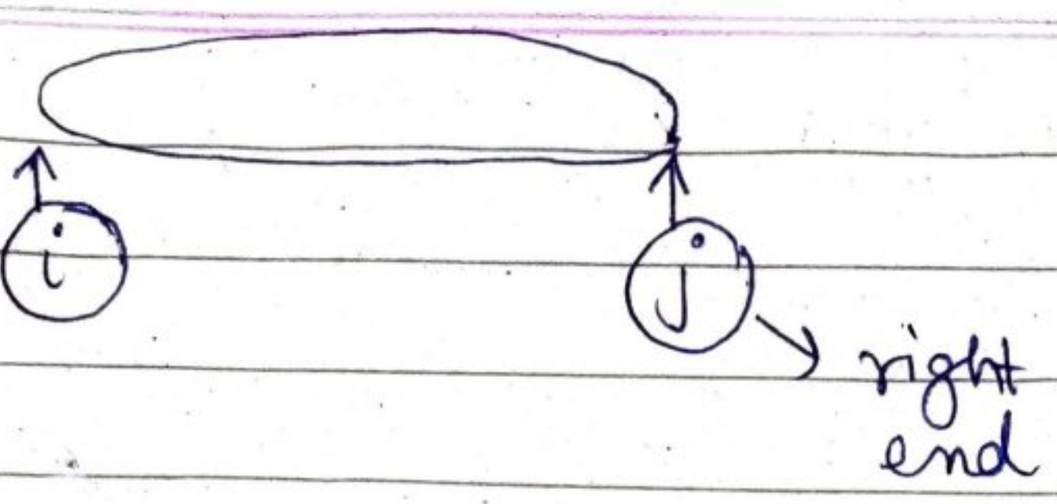
$(A_1)\ (A_2\ A_3\ A_4)$

↑ k ↓

Min cost + Min cost ← temp ans

$(A_1\ A_2\ A_3)\ (A_4)$

put brackets on different places to get the answer
Slide k at different positions.



40	20	30	10	30
i	i			j

Ai → arr[i-1] * arr[i]

when i=0 → Ai → arr[-1] * arr[0] ✗

when i=1 → Ai → arr[0] * arr[1]

Aj → arr[j-1] * arr[j]

arr[3] * arr[4]

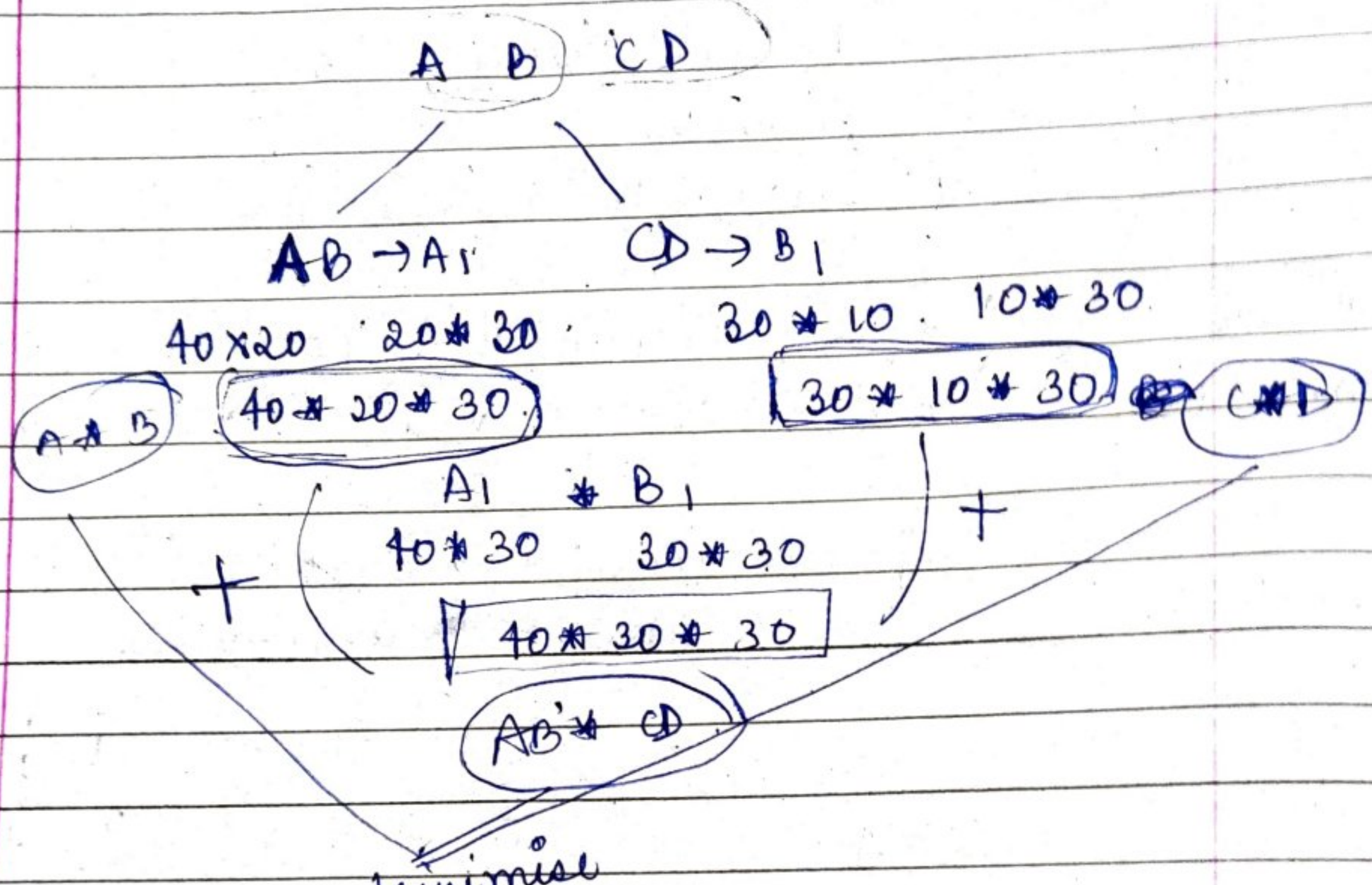
i>j → size=0

i==j → size=1 [(i-1) ni milega]

i = 1
j = n-1
} return cost

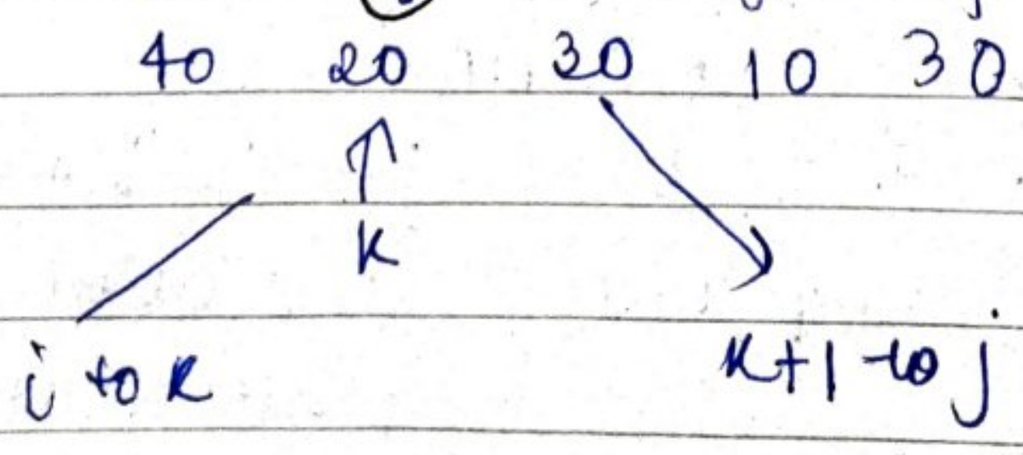
```
int solve ( int arr[], int i, int j)
{
if (i>=j)
    return 0;
    int mn = INT_MAX;
for(int k=i ; k <=j-1 ; k++){
    int tempans = solve(arr, i, k) +
                  solve(arr, k+1, j)
                + arr[i-1] * arr[k] * arr[j];

if(tempans)
    {
    mn = tempans;
}
```

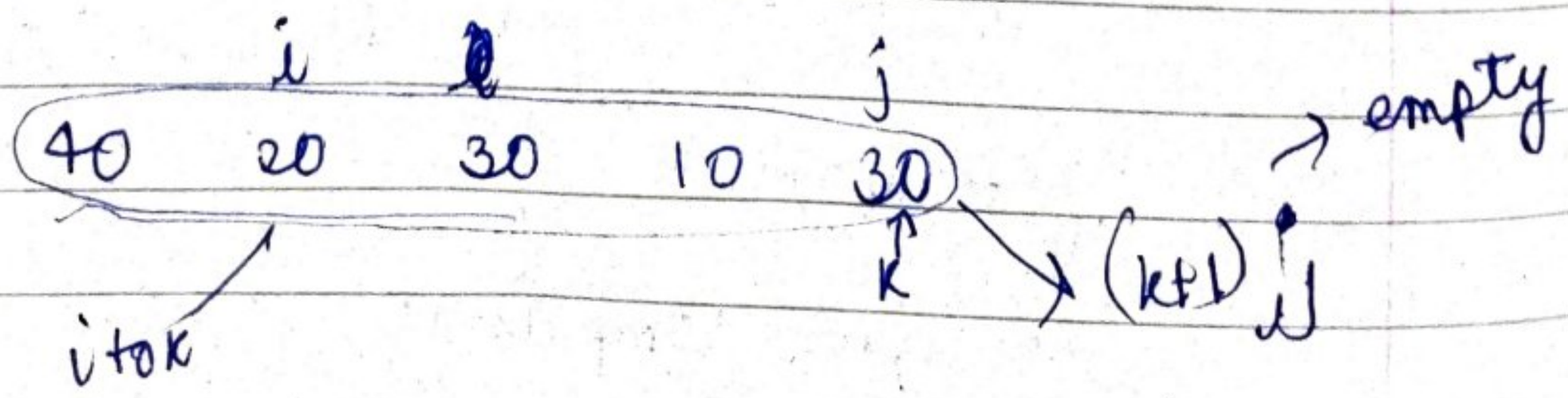
return mn.
}

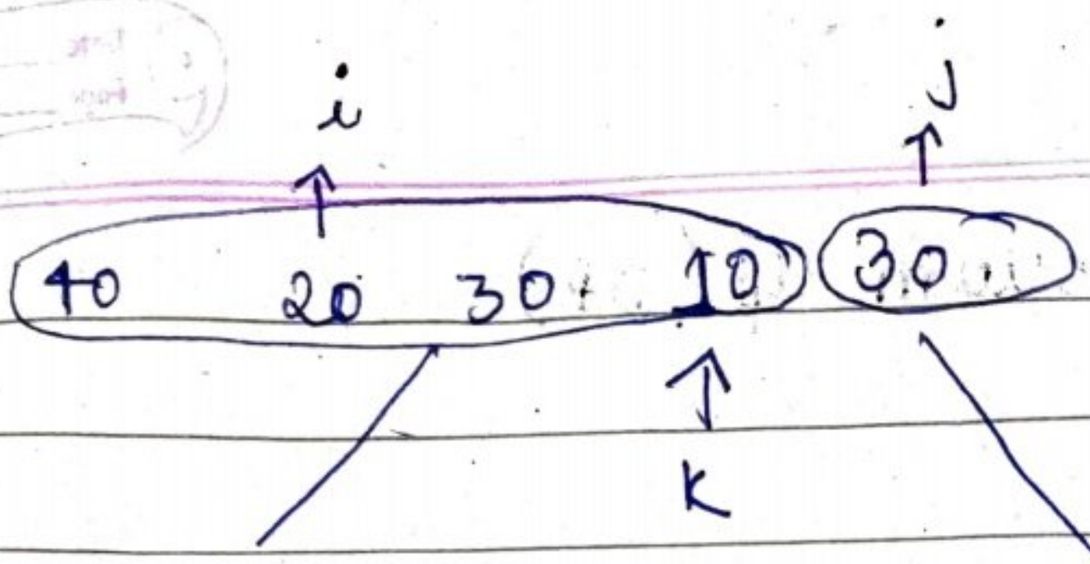


Steps for MCM

1) find i & j value
2) find right Base cond^n.
3) Move K → i to j. (find K-loop scheme)
4) Calculate (k) cost for (i) temp ans.





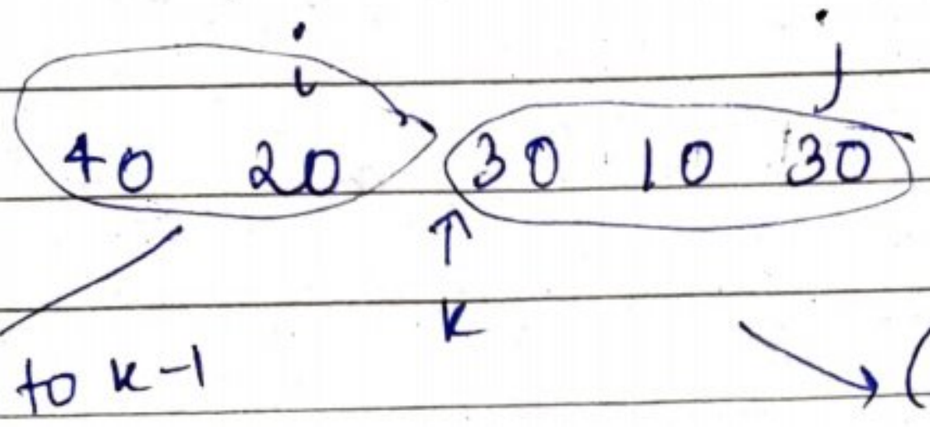


40 * 20

20 * 30

30 * 10.

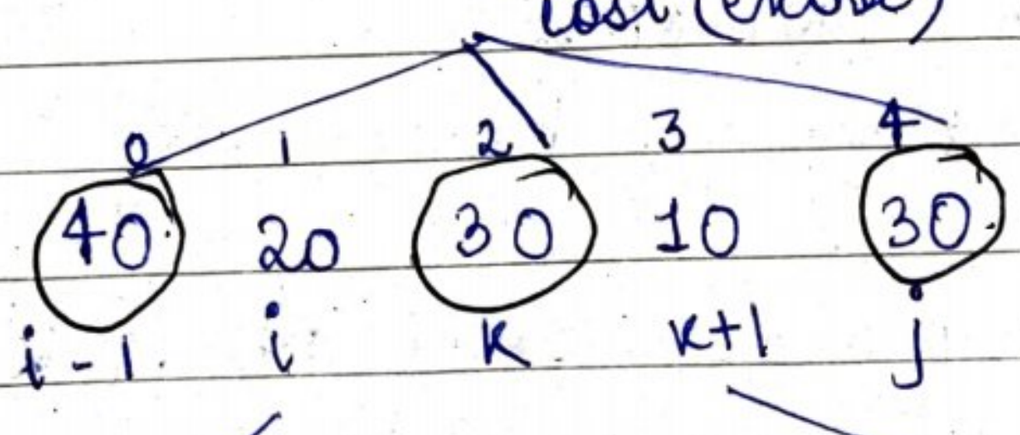
$k = i$ $k = j - 1$ ($i \to k$ $k+1 \to j$)



2 schemes.

$k = i$ to $k = j-1$ i to k k+1 to j

$k = i+1$ to $k = j$ i to k-1 k to j



for (i to k)

40 * 20 * 20 * 30

↓

40 * 20 * 30

solve (i to k)

for (k+1 to j)

30 * 10 10 * 30

↓

30 * 10 * 30

solve (k+1 to j)

40 * 20 * 30 * 30

↓

arr[i-1] ← 40 * 30 * 30 → arr[j]

arr[k]

dimension arr: [ ]

MCM Bottom Up (DP)

(dimension) arr[ ] : [ ]

```
int dp[100][100]

memset ( dp, -1 , sizeof(dp))

int solve ( int arr[], int i, int j)

if (i <= j)
    return 0;

if (dp[i][j] != -1)
    return dp[i][j];

int mn = INT_MAX;
for (int k = i; k<j; k++) {

int temp_ans = solve(arr, i, k) +
               solve (arr, k+1, j) +
               arr[i-1] * arr[k] * arr[j];

if (temp_ans < mn)
    mn = temp_ans;
}

dp[i][j] = mn;
return mn;
```

Palindrome Partitioning

```
int dp[1001][1001];

class Solution {
public:
    int Solve(int arr[], int i, int j) {
        if (i >= j)
            return 0;

        if (dp[i][j] != -1)
            return dp[i][j];

        int mn = INT_MAX;
        for(int k = i; k <= j-1; k++){
            int temp = Solve(arr, i, k) + solve(arr, k+1, j)
                       + arr[i-1] * arr[k] * arr[j];
            mn = min(temp, mn);
        }
        dp[i][j] = mn;
        return dp[i][j];
    }

int matrixmult(int N, int arr[]) {
    memset(dp, -1, sizeof(dp));
    int ans = solve(arr, 1, N-1);
    return ans;
}
};
```

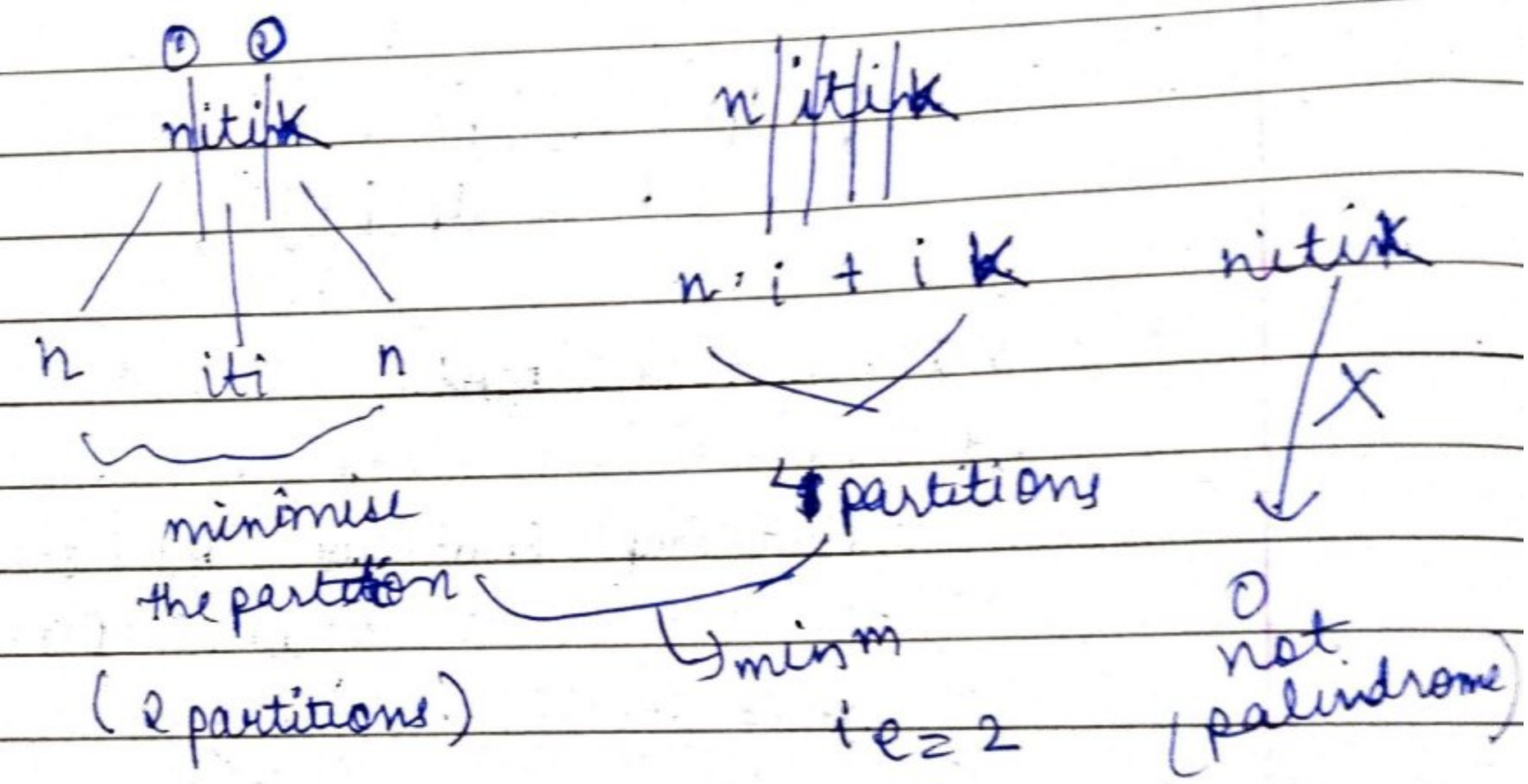
Palindrome partitioning

1) Problem statement

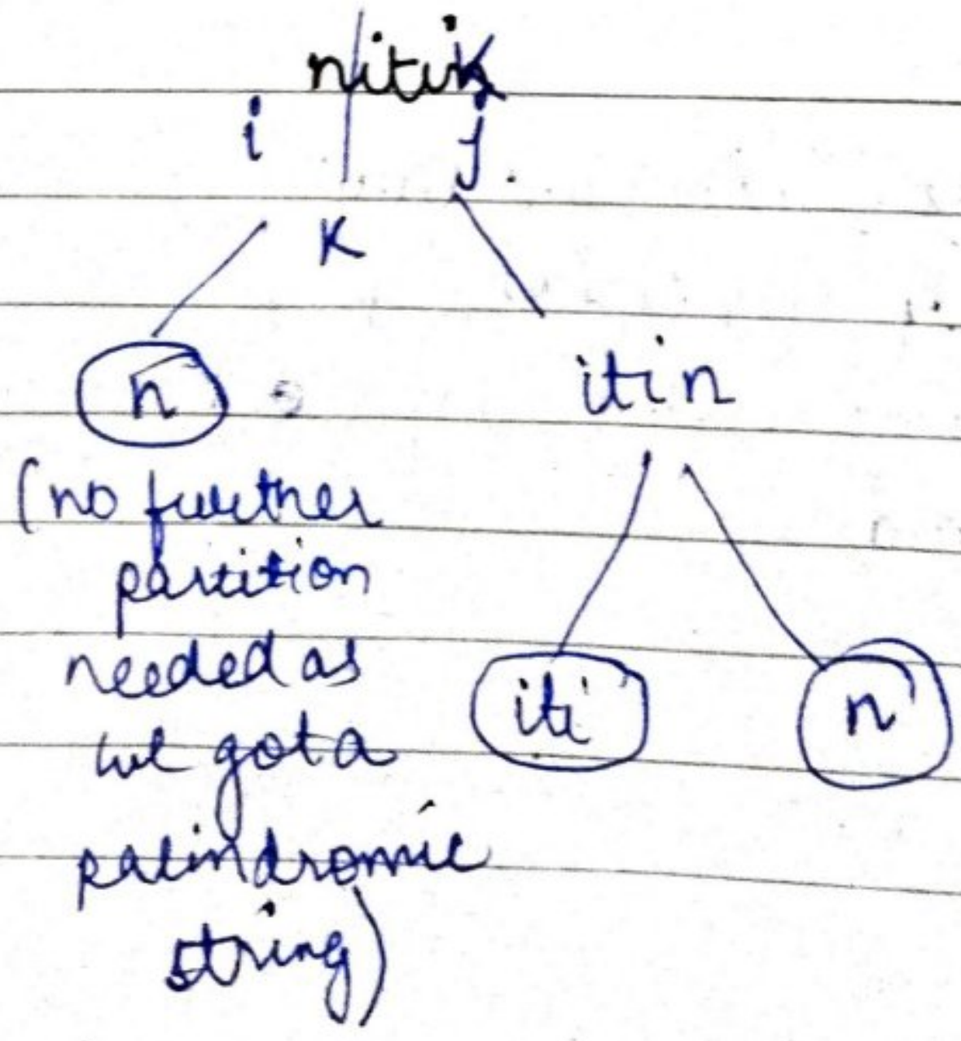
S : nitik → it is not a palindrome as a whole

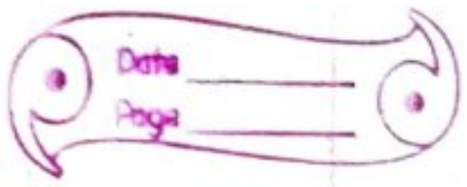
o/p : 2

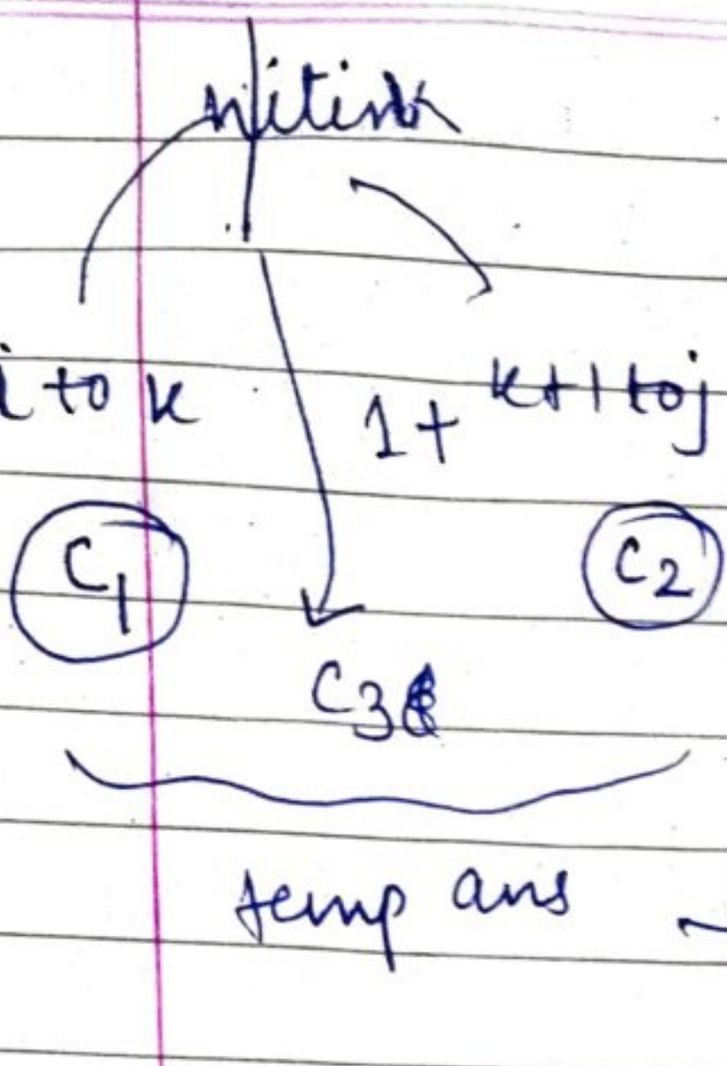
divide string in such a way all the strings are palindrome.

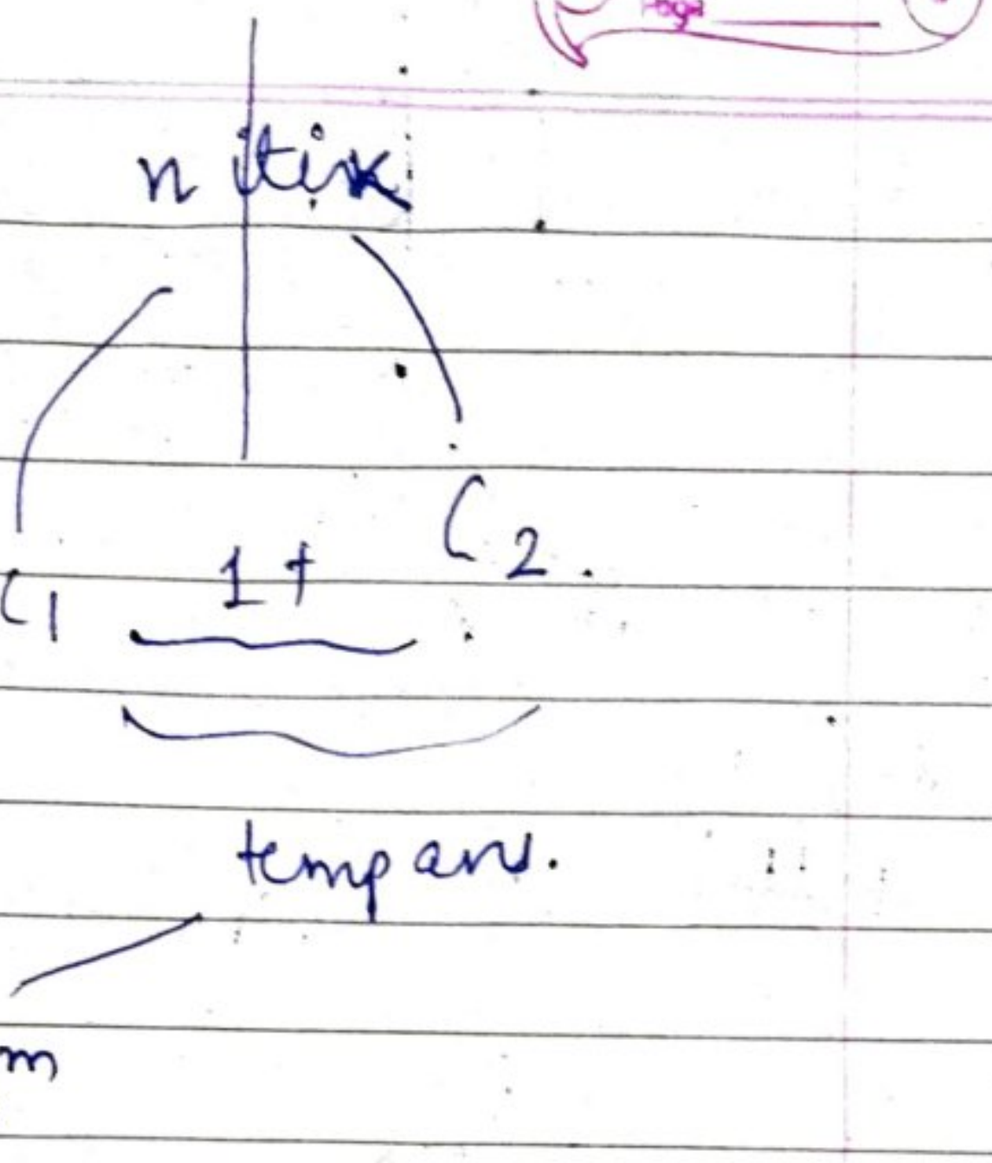


worst case partition : n-1









Format

1. Find i & j
2. Find BC.
3. Find K loop (scheme)

4 temp ans → min^m

Recursive code

n i t i k

i j (no need to start i from 1 since it doesn't require i-1)

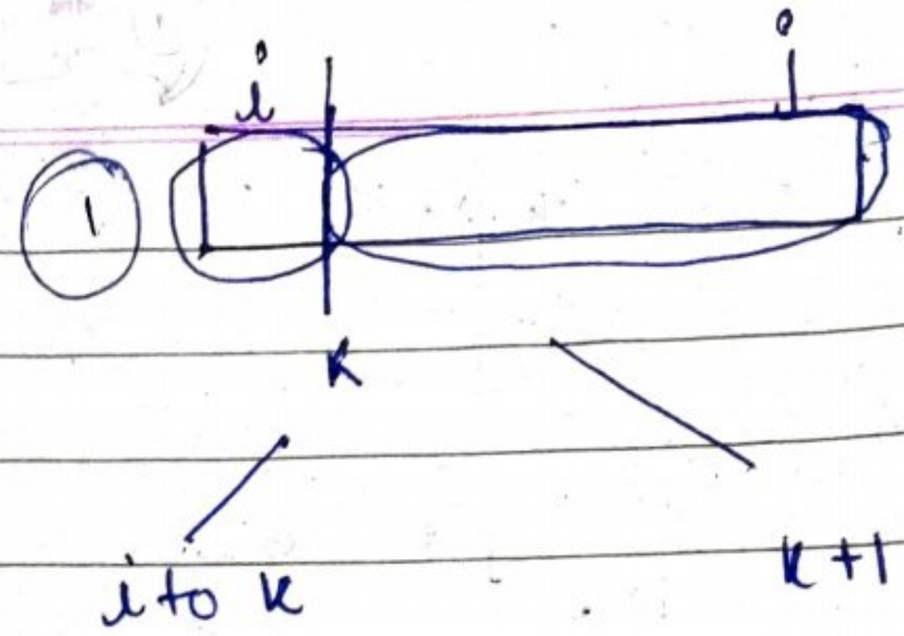
i=0 j=n-1

B.C i==j size=1

i>j size=0 } (if size is not given or 0 or equal no of partition would be 0 as string is empty)

& if palindrome partition should be 0.

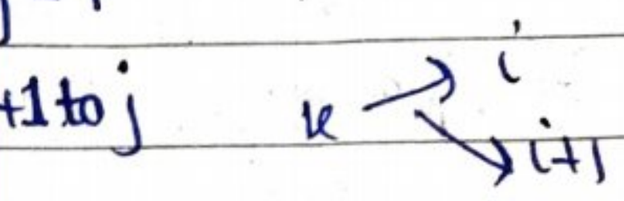
is palindrome (s, i, j)

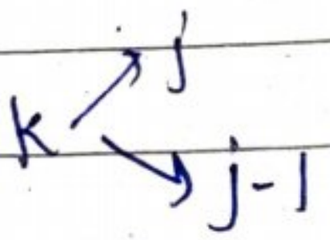


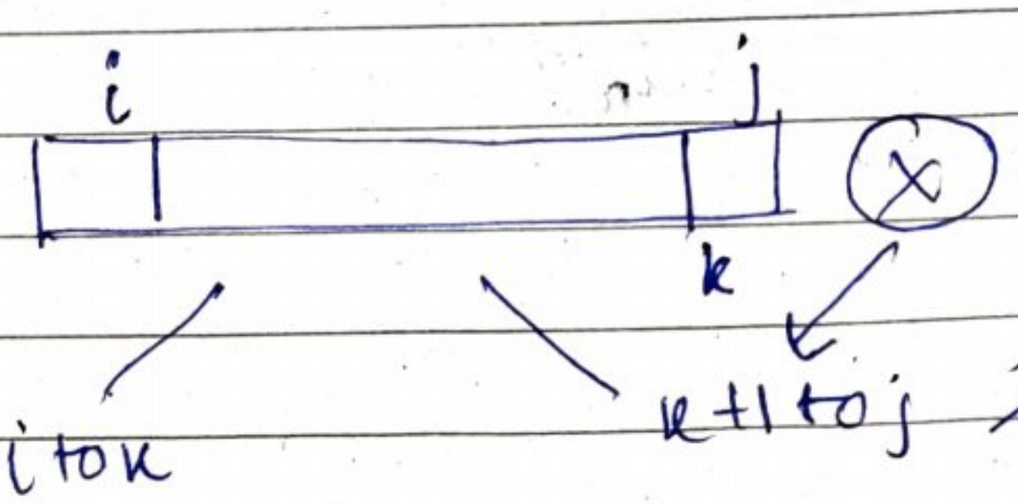
loop

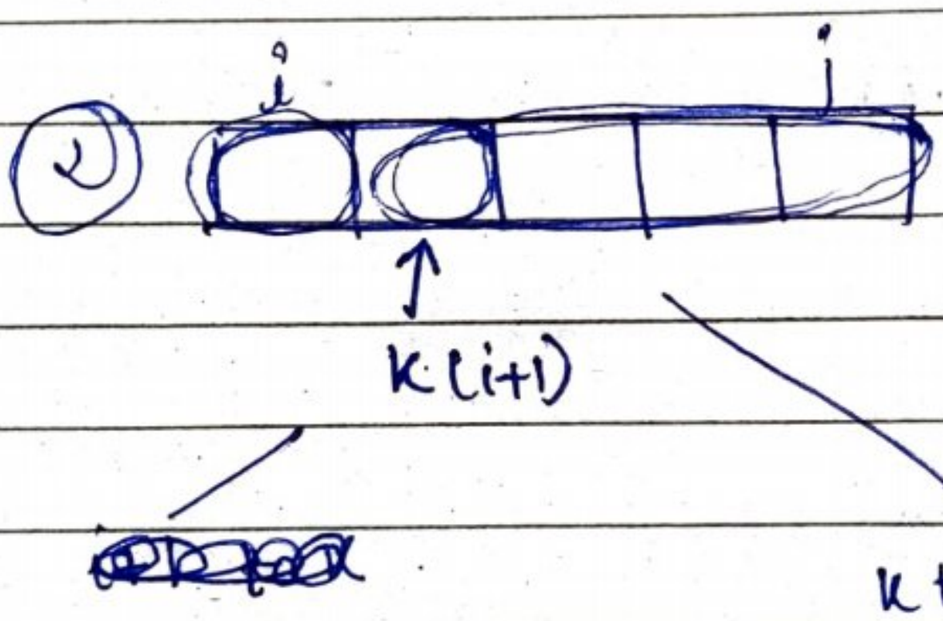
$k = i$ $k = j - 1$

i to k $k+1$ to j







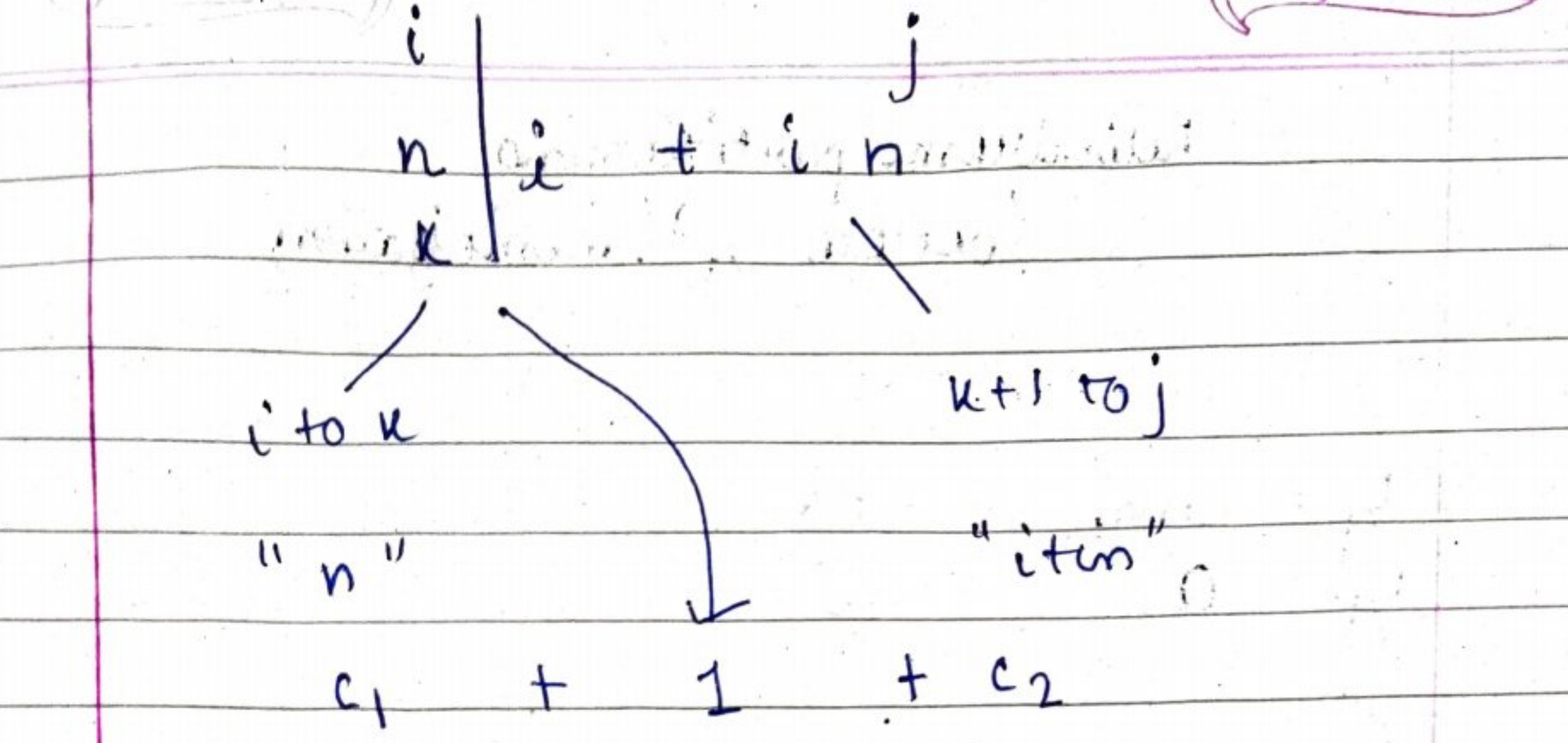


$k = i+1$

$k = j$

1) $k = i$	$k = j-1$	i to k	$k+1$ to j
2) $k = i+1$	$k = j$	i to $k-1$	k to j

```
for (int k = i ; k = j-1 ; k++)
    {
int temp = solve( s, i, k) + solve( s, k+1, j)
                                        +1
```



} ans = min (ans, temp)

return ans;

}

Final code

```
int Solve ( string s, int i, int j ) {
    if (i >= j)
        return 0;

    if( is palindrome (s, i, j) == = True)
        return 0;
    int mn = INT-MAX ;

    for(int k = i ; k <= j-1 ; k++) {
        int temp = 1+ solve (s, i, k) + solve(s, k+1, j );

        mn = min (mn, temp)
    }
    return mn;
}
```

memset < -1 / 0 > only allows 2 values

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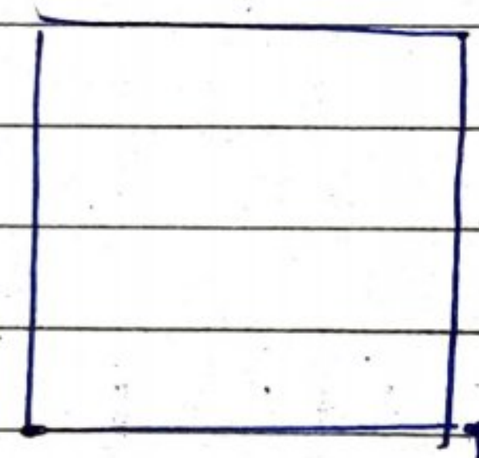
Palindrome partitioning (bottom up) (memoization)

i/p = nitin nitik
o/p = 0 2.

1) Initialise the matrix with -1.
2) Check if the value is
   - -1 (value isn't evaluated
   - ≠ -1 value is evaluated so we store & return the value.



nxm
↓
variables which changes basically

i & j changes.

```
int dp [1001][1001];
memset (dp, -1, sizeof(dp));
int solve ( string s, int i, int j) {
        if ( i >= j)
            return 0;

    if (is palindrome (s, i, j))
        return 0;
```

```
    if (dp[i][j] != -1)
        return dp[i][j];
    int mn = INT_MAX;
for(int k = i ; k <= j-1 ; k++) {
 int temp = solve(s, i, k) + solve(s, k+1, j) + 1
    mn = min (mn, temp);
    }.
    dp[i][j] = mn;
     return mn;
}


int palindromepart( string s) {
      int n = s.length() -1;
      int ans = solve( s, 0, n);
      return ans;
  }

bool ispalindrome (string s, int i, int j) {

  if (i == j)
      return True;

  if (i > j)
     return True;

  while (i<j)
  if (s[i] != s[j])
       return false
  else
       i++;
       j--;
```

Geeksforgeeks = ✓
Interviewbit = ✗

↳ Further Optimization

Why the above code is not most optimized?

Since we are calling

int temp = 1+ solve(s, i, k) + solve(s, k+1, j)
RC (left) RL (right)

There is a possibility, might be one of the RC is solved or called.

The code will remain same but the diff^n is

int temp = 1+ solve(s, i, k) + solve(s, k+1, j)

if (t[i][k] != -1)
    left = t[i][k]
else
    left = solve(s, i, k)
    t[i][k] = left

if (t[k+1][j] != -1)
    right = t[k+1][j]

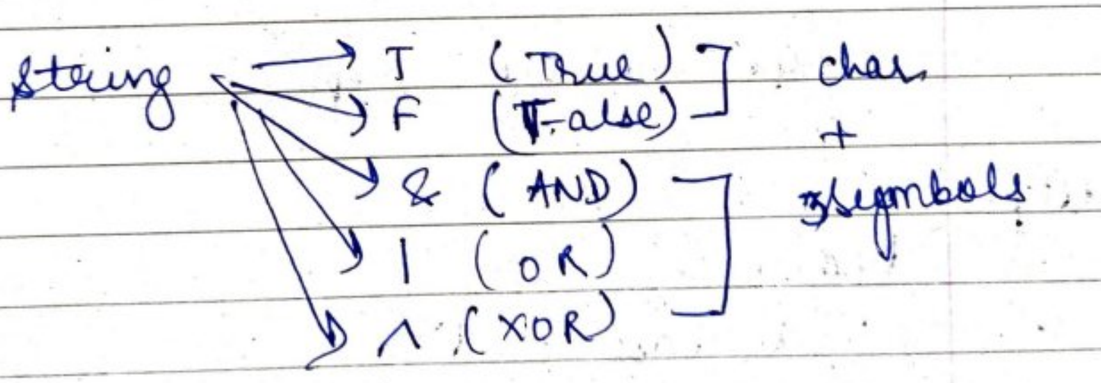
else
right = solve(s, k+1, j)
t[k+1][j] = right.

int temp = 1 + left + right;

Evaluate Expression to True
Boolean parenthesization

String : T or F and T
o/p : 2

Problem statement : A string is given. String might have some characters like T → true F → false



How to put bracket such that the expression evaluates to true.

Find the no of ways in which when bracket is put it evaluates to True.

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i/p : "T | F & T ∧ F"

no of ways

(( )) ( ) → T
( ) ( ) → F
( ) ( ) → T → 3 ways.
( ) ( ) → T

In MCM we put brackets for min^m cost & in this also we do the same.

(T | F & T ∧ F)

K-1 K K+1

We need to break bracket on operator

Expr^n operator Expr^n

4 steps:

1) find i & j
2) find base cond^n
3) Find K loop
4) temp ans ← func^n

↓

main ans.

T
F

i j

T | F & T ^ F

1) i = 0
   j = str.length() - 1

2) B C

i j

(T or F and T) xor (F)

i to k-1 | k | k+1 to j

(left) Expr^n1 XOR Expr^n2 (right)

i to k-1 | k | k+1 to j

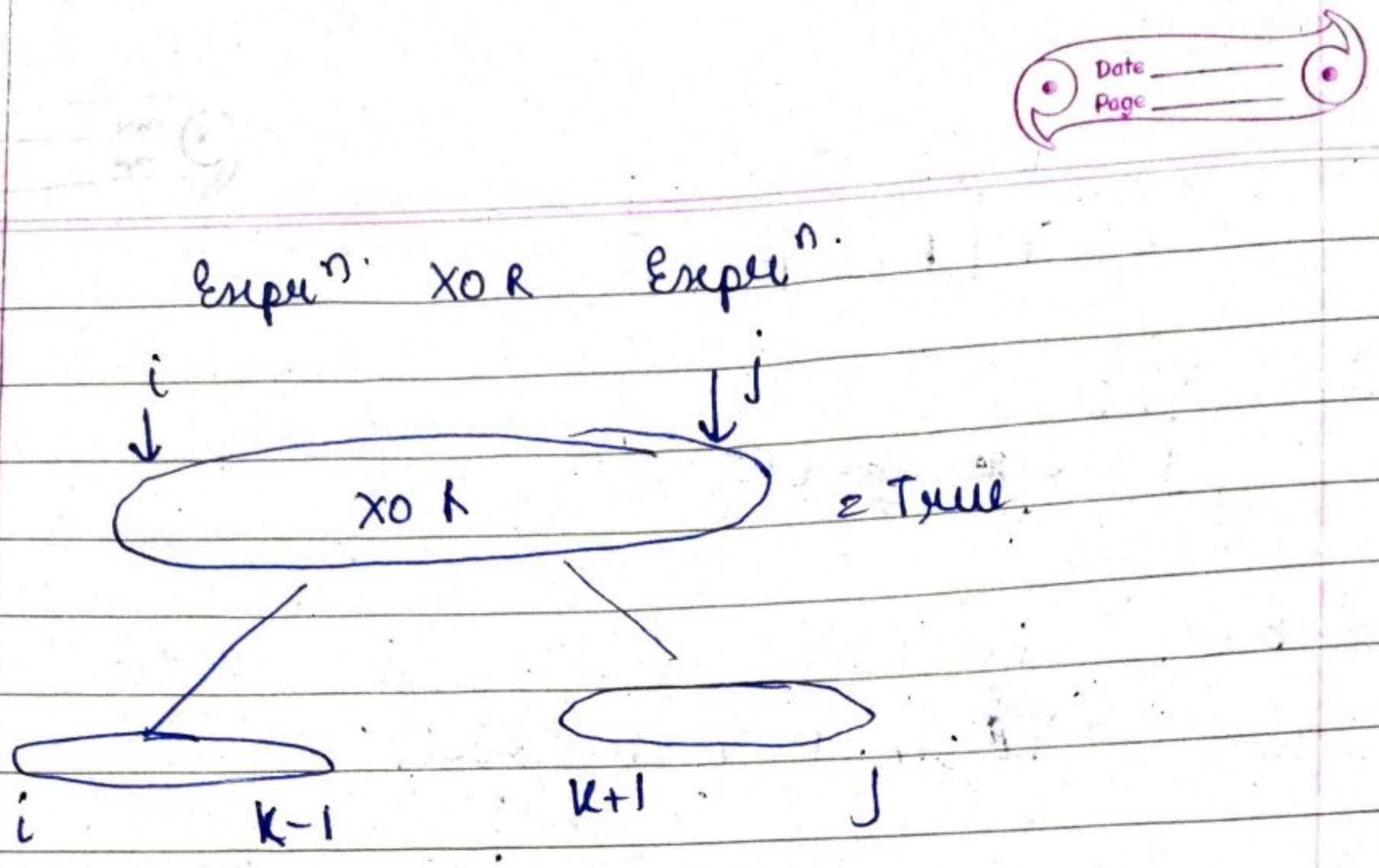
T ^ T = False
F ^ F = false
T ^ F = True
F ^ T = True

no of true ways = nT * nF + nF * nT

T ^ F = T
F ^ T = T

Expr^n. XOR Expr^n.

no of ways (True) → 2 * 4 ← no of ways false



int solve (string S, int i, int j, boolean to True or is False)

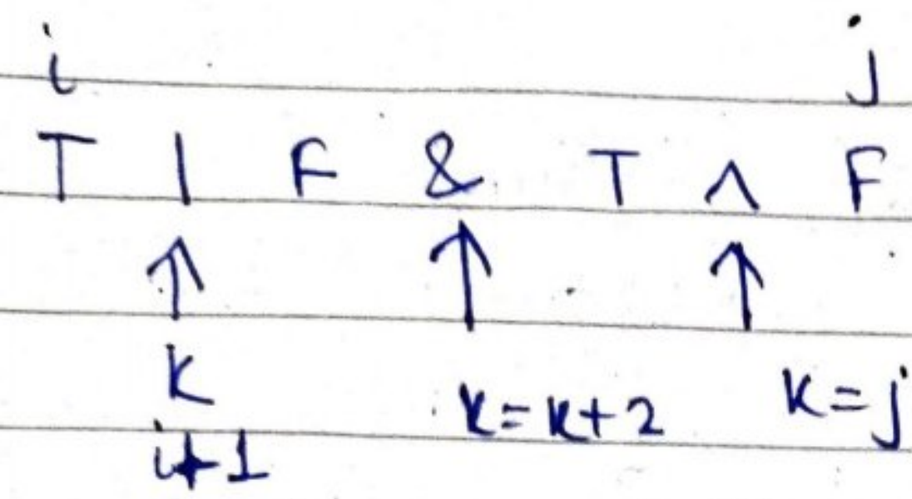
Base cond^n

i
T or F and T xor F
T | F & T ^ F

boolean is True :

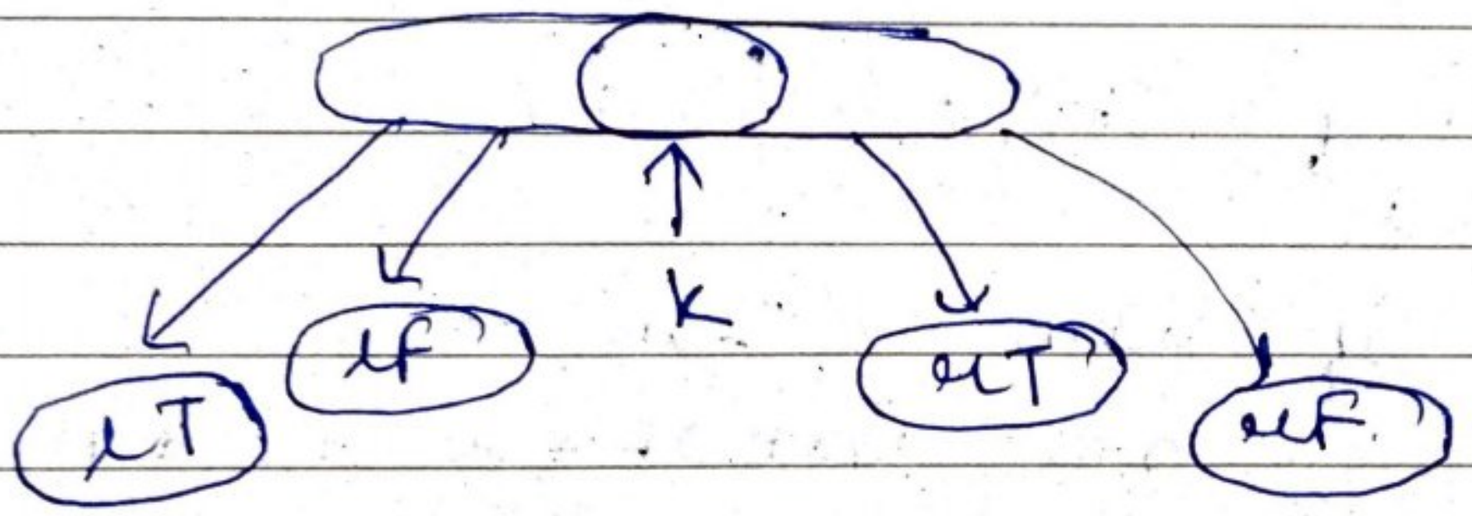
if (i > j) return false.
if (i == j)
    if (is True == True)
        return S[i] == 'T'
    else
        return S[i] == 'F'

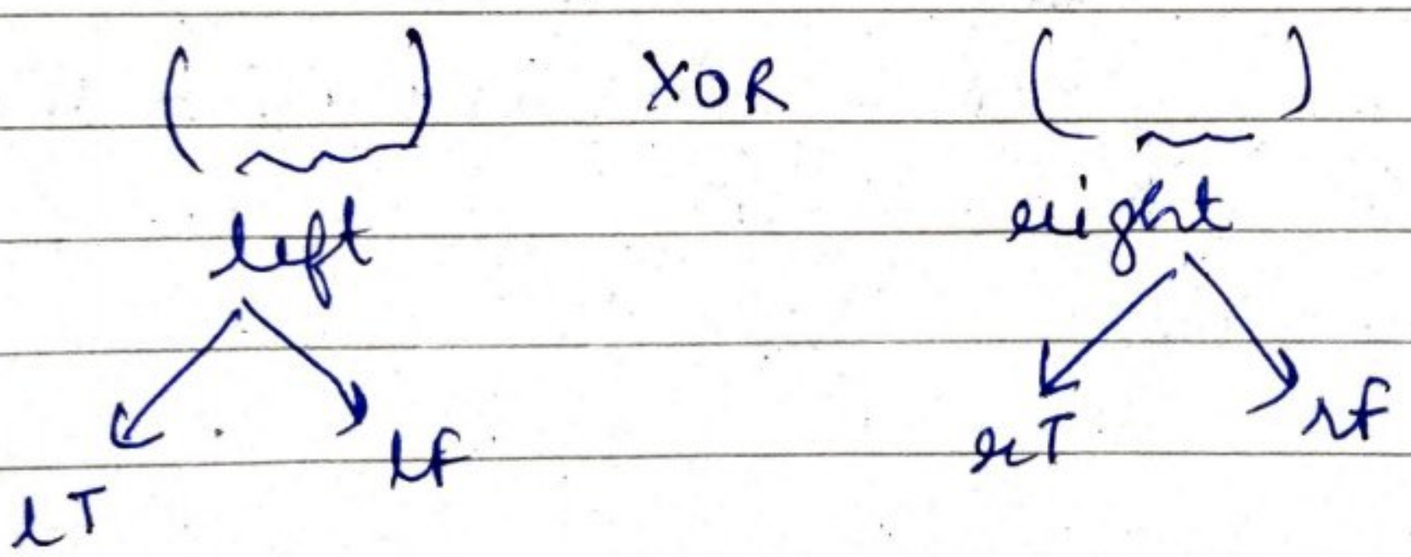
loop




```
for (int k = i+1 ; k <= j-1 ; k = k+2)
    int ans = 0;
    int lT = (s , i , k-1 , T)
    int lF = (s , i , k-1 , F)
    int rT = (s , k+1 , j , T)
    int rF = (s , k+1 , j , T)
```

} temp ans.





ans t = lT * rf + lf * rT

```
if (s[k] == '&')
{
if (isTrue == True)
  ans = ans + lT * rT
else {
ans = ans + lf * rT + lT * rf + lf * rf;
}
else if (s[k] == '|') {
  if (isTrue == True)
    ans = ans + lT * rT + lT * rF +
                lF * rT.
  else
    ans = ans + lF * rF.
}
else if (s[k] == '^')
  {
  if (isTrue == True)
    ans = ans + lF * rT + lT * rF
  else
    ans = ans + lT * rT + lF * rF
  }
return ans;
}
```

Whole code

```
int solve (string s ,int i, int j , bool isTrue ) {

    if (i > j) {
        return false ; }
    if (i == j) {
        if (isTrue == Trues)
            return s[i] == 'T';
        else
            return s[i] == 'F';
        }
for ( int k= i+1 ; k <= j-1 ; k+=2) {
        int ans = 0;
        int LT = solve (s, i, k-1, T);
       ,int LF = Solve (s, i, k-1, F);
       'int RT = Solve (s, k+1, j, T);
        int RF = solve (s, k+1, j, T);

    if (s[k] == '&') {
        if (isTrue == = Trues)
            ans = ans + LT * RT ;
        else
            ans = ans + LF * RT + LT * RF + LF * RF;
        }
    else if ( s[k] == '|') {
        if (isTrue == = Trues)
            ans += LT * RT + LT * RF + LF * RT;
        else
            ans += LF * RF;
        }
```

```
else if ( S[x] == 'A' ) {
        if (isTrue == true)
            ans += LF * RT + LT * RF;
        else
            ans += LT * RT + LF * RF;
        }
    }
    return ans;
}
```

Evaluate Expression To True Boolean Parenthesization - (memoization)

(Bottom Up - DP)

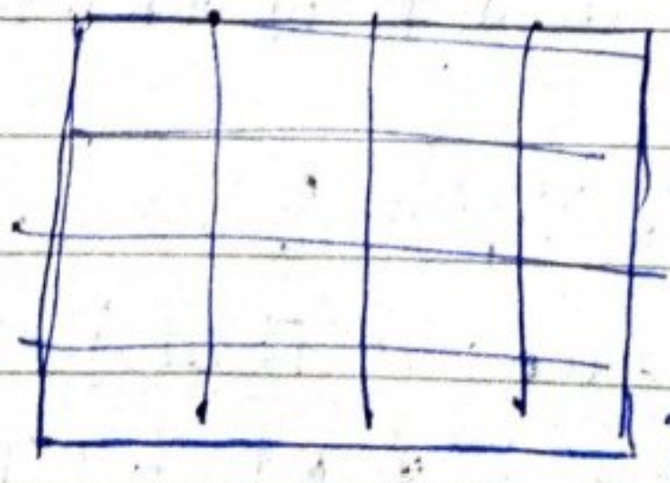
3D dp

Recursive → Recursive call (RC)

Top down → table

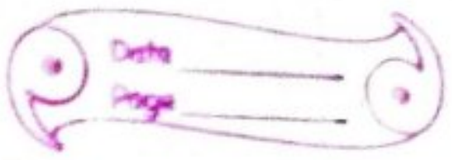
Bottom up → RC + table

PS → put brackets in such a way that the expression evaluates to true.



i * j * isTrue
(3d)

int t [100] [100] [2]



isTrue

True False

More better way: We can use map

map	value
i+" "+j+" "+isTrue	

unordered_map < string, int > mp;

```
int main(){
    mp.clear()
    Solve( )
}
```

After Base cond^n :-

```
String temp = to_string(i);
    temp.push_back(' ');
  temp.append ( to_string(j));
   temp.push_back(' ');
   temp.append ( to_string(isTrue));

if ( mp.find(temp) != mp.end()){
        return mp[temp];
    }
```

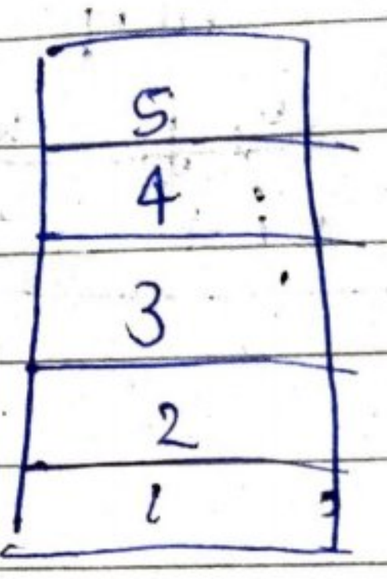
return mp[temp] = ans;

}

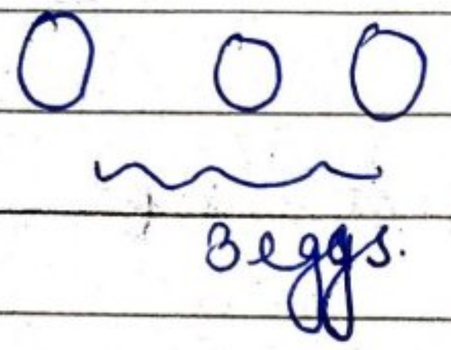
Egg Dropping Problem

I/p : e = 3
f = 5

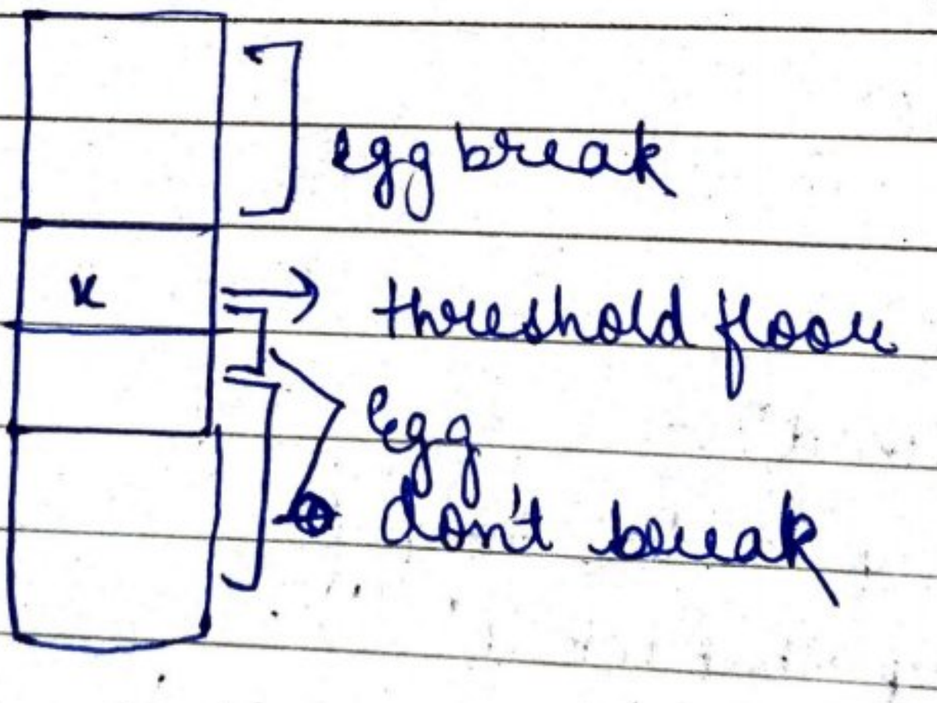
O/p : 3 → minmize no of attempts in worst case.



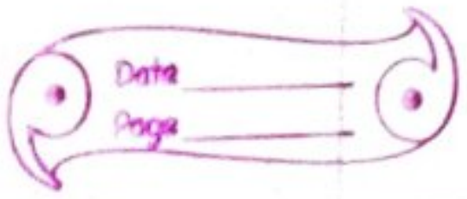
Problem statement :



3 eggs.

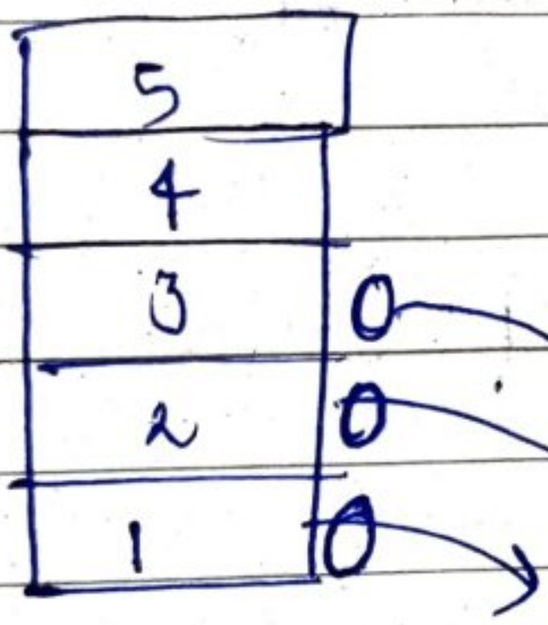


We are in a building, we need to find the minm no of attempts to find the critical floor.

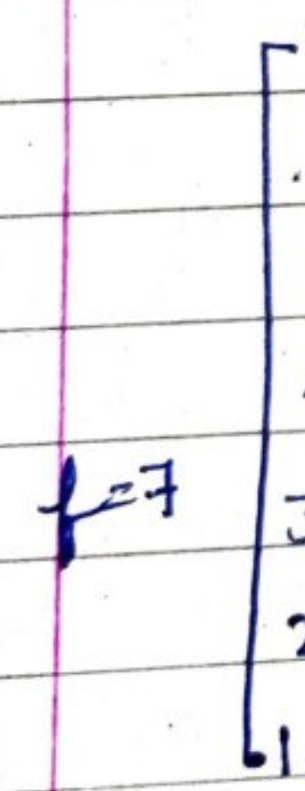


O O O 3 eggs

Safe strategy → Worst case (1 egg) · so drop from the last it wont break until threshold.



→ egg break (Threshold floor)



Worst-case → minimize no of attempts to find threshold floor.

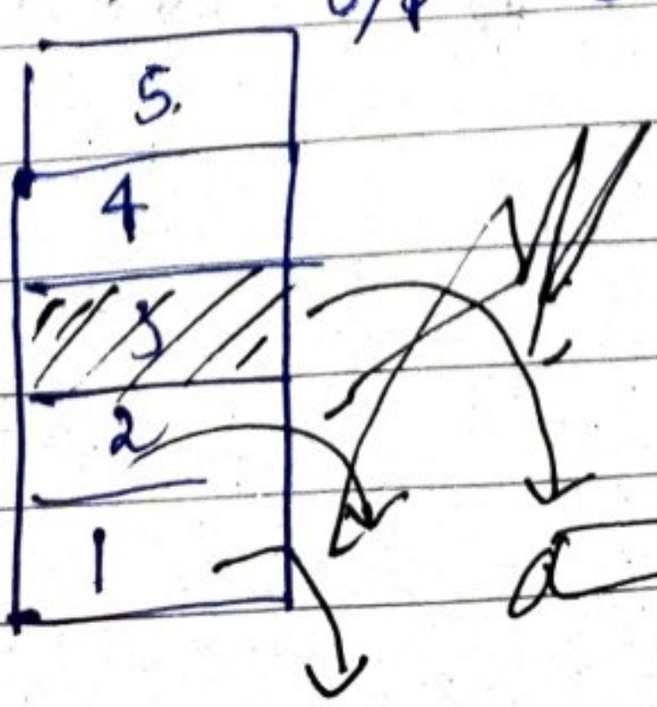
O O O

e = 3.

I/p e = 3

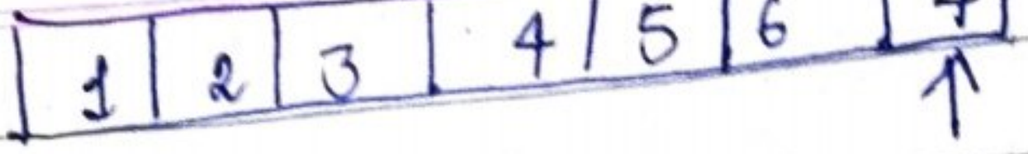
f = 5

O/p = 3. attempts



→ egg break.

(floor) f



→ Building looks like an array

from where to take k?

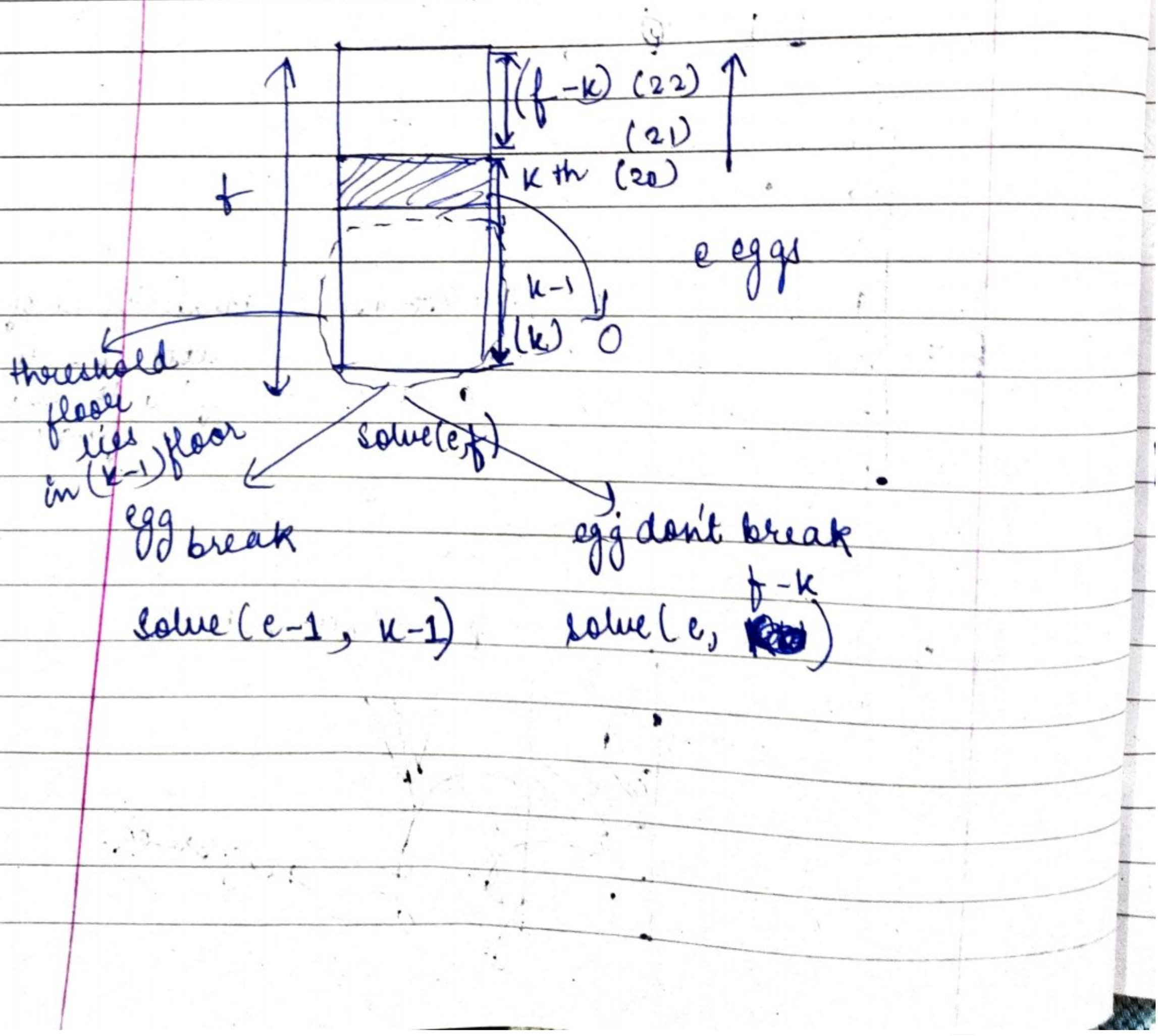
→ We will check for all k values.

for k = 1 k = f k++

Base condⁿ. (smallest valid i/p)

e = 1 return f

f = 0/1 return f

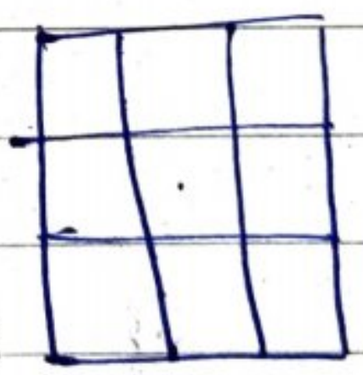


```
int solve (int e, int f)
  {
    if (f == 0 || f == 1)
        return f;

    if ( e == 1)
        return f;

  int  mn = INT_MAX;

  for(int k=1;  k<= f ; k++)
    {                          → for the worst case
  int temp = 1 +  max( solve (e-1, k-1),
                       solve (e, f-k));
      mn = min (mn, temp);
    }
  return mn;
}
```

Egg dropping Memoization

1 <= T <= 30
1 <= e <= 100
1 <= f <= 1000

```
// globally defined  int dp[1001][1001];
                     memset (dp, -1, sizeof(dp));
```



ex f

```
int solve ( int e, int f) {
    if (e == 1)
        return f;
    if (f == 0 || f == 1)
        return f;
```

```
if (t[e][f] != -1)
    return t[e][f];
int mn = INT_MAX;
for (int k=i ; k<=f ; k++) {

int temp = max(solve(e-1, k-1),
               solve(e, f-k)) + 1;

mn = min(mn, temp);
}

return t[e][f] = mn;
}
```

Further Optimization

Inside the loop.

```
if (t[e-1][k-1] != -1)
    int low = t[e-1][k-1]
else
    low = solve(e-1, k-1)
    t[e-1][k-1] = low;

if (t[e][f-k] != -1)
    int high = t[e][f-k];
else
    high = solve(e, f-k);
    t[e][f-k] = high;

int temp = 1+max(low, high);
```