

Bellman Ford Algorithm Tracing.

main()

num_ver source A[][] Max_value = 999

Enter number of vertices

number

Enter the adjacency matrix

$A[s_n][d_n] \Rightarrow A[5][5]$

$A[1][1] = 0$

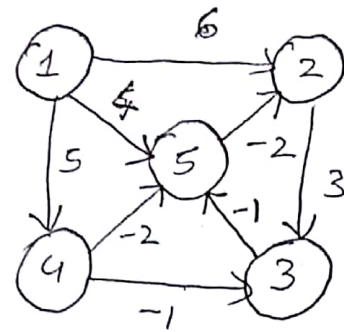
$A[1][2] = 6$

$A[1][3] = 0$

$A[1][4] = 5$

$A[1][5] = 4$

// indexing starts from 1



$A[2][1] = 0$

$A[2][2] = 0$

$A[2][3] = 3$

$A[2][4] = 0$

$A[2][5] = 0$

$A[3][1] = 0$

$A[3][2] = 0$

$A[3][3] = 0$

$A[3][4] = 0$

$A[3][5] = -1$

$A[4][1] = 0$

$A[4][2] = 0$

$A[4][3] = -1$

$A[4][4] = 0$

$A[4][5] = -2$

$A[5][1] = 0$

$A[5][2] = -2$

$A[5][3] = 0$

$A[5][4] = 0$

$A[5][5] = 0$

if $A[s_n][d_n] == 0$

then $A[s_n][d_n] = \text{Max_value} // 999$

All values of adjacency matrix are filled with 999 where.

$A[s_n][d_n] = 0$, other than condition where $s_n == d_n$,

if $s_n == d_n$

then $A[s_n][d_n] = 0$

$A[s_n][d_n]$	1	2	3	4	5
1	0	6	999	5	4
2	999	0	3	999	999
3	999	999	0	999	-1
4	999	999	-1	0	-2
5	999	-2	999	999	0

Enter source vertex:

source = 1

BellmanFord(S) // calls construction

{
 D[1] = 5 + 1 = 6 // D[6]
}

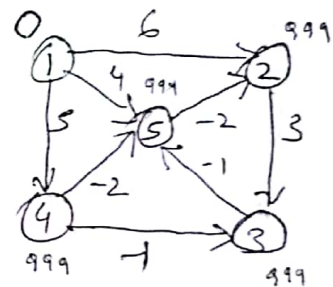
BellmanFordEvaluation(1, A[E][E]) // Function to evaluate

D[source] = 0; // D[1] = 0

Rest all nodes = ~~max~~ value \rightarrow 999

D[6]

0	1	2	3	4	5
0	999	999	999	999	999



If $D[dn] > D[sn] + \text{weight of edge } uv$

then $D[dn] = D[sn] + \text{weight of edge } uv$

This is executed when $A[sn][dn] \neq 999$ // Edge exist

Node 1
 $A[1][1] \Rightarrow 0 \neq 999 \checkmark$

$D[1] > D[1] + A[1][1]$

$0 > 0 + 0$ F

$D[1] = 0$

$A[1][2] \Rightarrow 6 \neq 999 \checkmark$

$D[2] > D[1] + A[1][2]$

$999 > 0 + 6$ T

$D[2] = 6$

$A[1][3] \Rightarrow 999 \neq 999 \times$

$D[3] = 999$

$A[1][4] \Rightarrow 5 \neq 999 \checkmark$

$D[4] > D[1] + A[1][4]$

$999 > 0 + 5$ T

$D[4] = 5$

$A[1][5] \Rightarrow 4 \neq 999 \checkmark$

$D[5] > D[1] + A[1][5]$

$999 > 4$ T

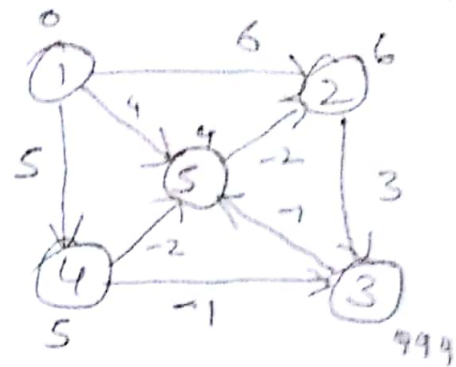
$D[5] = 4$

D[6]

0	1	2	3	4	5
0	6	999	5	4	

~~II $\Delta[2][1] = 999 = 999$ F~~
 ~~$\Delta[2][2] = 0 = 999$ T~~

Present Graph \rightarrow



II $\Delta[2][1] = 999 = 999$ F
 $\Delta[2][2] = 0 = 999$ T

$D[2] > D[2] + \Delta[2][2]$
 $6 > 6 + 0$

$D[1] = 0$

$D[2] = 6$

$\Delta[2][3] = 3 = 999$ T

$D[3] > D[2] + \Delta[2][3]$
 $999 > 6 + 3$ T

$D[3] = 9$

$\Delta[2][4] = 999 = 999$ F

$D[4] = 5$

$\Delta[2][5] = 999 = 999$ F

$D[5] = 4$

$D[6]$

0	1	2	3	4	5
0	6	9	5	4	

III $\Delta[3][1] = 999 = 999$ F

$D[1] = 0$

$\Delta[3][2] = 999 = 999$ F

$D[2] = 6$

$\Delta[3][3] = 0 = 999$ T

$D[3] > D[3] + \Delta[3][3]$
 $9 > 9 + 0$

$D[3] = 9$

$\Delta[3][4] = 999 = 999$ F

$D[4] = 5$

$\Delta[3][5] = -1 = 999$ T

$D[5] > D[3] + \Delta[3][5]$
 $4 > 9 - 1$ F

$D[5] = 4$

IV $\Delta[4][1] = 999 = 999$ F

$D[1] = 0$

$\Delta[4][2] = 999 = 999$ F

$D[2] = 6$

$\Delta[4][3] = -1 = 999$ T

$D[3] > D[4] + \Delta[4][3]$
 $9 > 5 - 1$

$D[3] = 4$

$\Delta[4][4] = 0 = 999$ T

$D[4] > D[4] + \Delta[4][4]$

$D[4] = 5$

$$DE[4][5] \quad -21 = 999 \quad T$$

$$DE[5] > DE[4] + A[4][5]$$

$$4 > 5 - 2$$

$$DE[5] = 2$$

$$DE[6]$$

0	1	2	3	4	5
0	6	4	5	3	

$$V \quad A[5][1] \quad 9991 = 999 \quad F$$

$$DE[1] = 0$$

$$A[5][2] \quad -21 = 999 \quad T$$

$$DE[2] > DE[5] + A[5][2]$$

$$6 > 3 - 2$$

$$DE[2] = 1$$

$$A[5][3] \quad 9991 = 999$$

$$DE[3] = 4$$

$$A[5][4] \quad 9991 = 999$$

$$DE[4] = 5$$

$$A[5][5] \quad 01 = 999$$

$$DE[5] > DE[5] + A[5][5]$$

$$3 > 3 + 0$$

$$DE[5] = 3$$

$$DE[6]$$

1	2	3	4	5
0	-1	4	5	3

$$mode = 2$$

$$I \quad A[1][1] \quad 01 = 999 \quad \checkmark$$

$$DE[1] > DE[1] + A[1][1]$$

$$DE[1] = 0$$

$$A[1][2] \quad 61 = 999$$

$$DE[2] > DE[1] + A[1][2]$$

$$1 > 0 + 6$$

$$DE[2] = 1$$

$$A[1][3] \quad 9991 = 999 \quad \times$$

$$DE[3] = 4$$

$$A[1][4] \quad 51 = 999 \quad \checkmark$$

$$DE[4] > DE[1] + A[1][4]$$

$$5 > 0 + 5$$

$$DE[4] = 5$$

$$A[1][5] \quad 41 = 999$$

$$DE[5] > DE[1] + A[1][5]$$

$$3 > 0 + 4 \quad F$$

$$DE[5] = 3$$

$$DE[6]$$

1	2	3	4	5
0	1	4	5	3

In similar way in $\Delta[2][1 \dots 5]$

$$D[6] = \begin{array}{|c|c|c|c|c|} \hline 1 & 2 & 3 & 4 & 5 \\ \hline 0 & 1 & 4 & 5 & 3 \\ \hline \end{array}$$

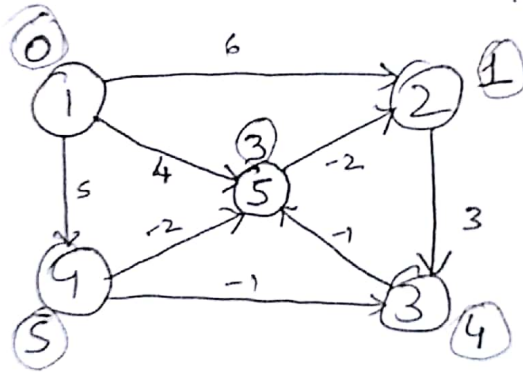
For $\Delta[3][1 \dots 5]$

$\Delta[4][1 \dots 5]$

$\Delta[5][1 \dots 5]$

$$D[6] = \begin{array}{|c|c|c|c|c|} \hline 1 & 2 & 3 & 4 & 5 \\ \hline 0 & 1 & 4 & 5 & 3 \\ \hline \end{array}$$

node = 3, 4

$$D[6] = \begin{array}{|c|c|c|c|c|} \hline 1 & 2 & 3 & 4 & 5 \\ \hline 0 & 1 & 4 & 5 & 3 \\ \hline \end{array}$$


distance of source 1 to 1 is $D[1] = 0$

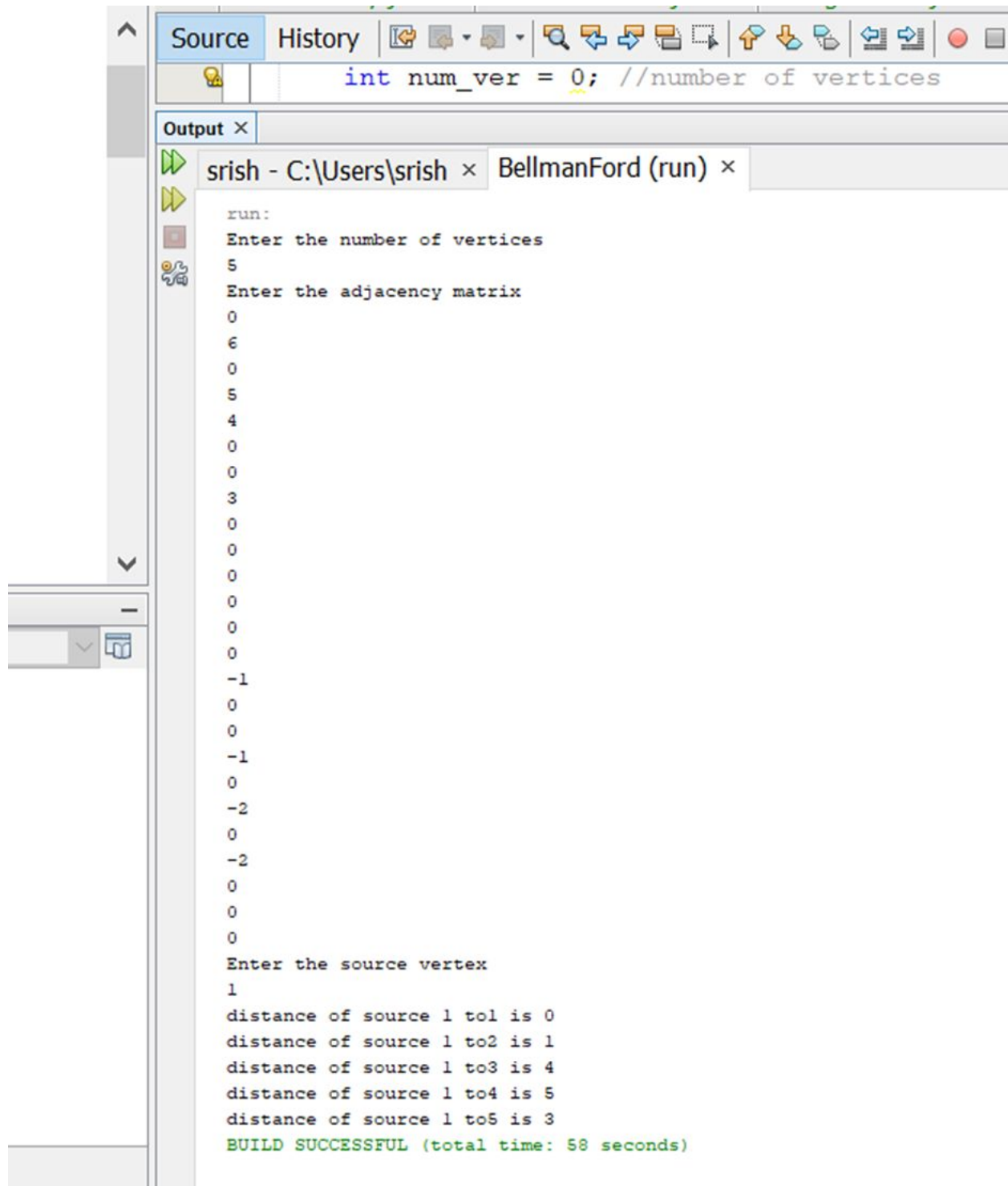
distance of source 1 to 2 is $D[2] = 1$

distance of source 1 to 3 is $D[3] = 4$

distance of source 1 to 4 is $D[4] = 5$

distance of source 1 to 5 is $D[5] = 3$

BellmanFord's Algorithm Output



The screenshot shows a C++ IDE with a 'Source' tab and an 'Output' tab. The 'Source' tab contains the following code:

```
int num_ver = 0; //number of vertices
```

The 'Output' tab shows the execution of the program. The output is as follows:

```
run:
Enter the number of vertices
5
Enter the adjacency matrix
0
6
0
5
4
0
0
3
0
0
0
0
0
0
-1
0
0
-1
0
-2
0
-2
0
0
0
Enter the source vertex
1
distance of source 1 to1 is 0
distance of source 1 to2 is 1
distance of source 1 to3 is 4
distance of source 1 to4 is 5
distance of source 1 to5 is 3
BUILD SUCCESSFUL (total time: 58 seconds)
```

CRC Program

Tracing

data bits \rightarrow ^{int} • divisor bits \rightarrow ^{int} total length \rightarrow ^{int}

data \rightarrow ^{array of int} div \rightarrow ^{array of int} divisor rem
 or

Enter number of data bits

\Rightarrow data bits = 6

Enter data bits

1
1
0
0
1
1

\Rightarrow data [6] =

1	1	0	0	1	1
---	---	---	---	---	---

Enter number of bits in divisor

\Rightarrow divisor bits = 4

Enter Divisor bits

1
0
1
1

\Rightarrow divisor [4] =

1	0	1	1
---	---	---	---

total length = 6 + 4 + 1 = 9

div[9] =

--	--	--	--	--	--	--	--	--

rem[9] =

--	--	--	--	--	--	--	--	--

or [9] =

--	--	--	--	--	--	--	--	--

for $i=0$ to data.length-1;

div[i] = data[i] = 110011

After appending 0's :- div[9] = 110011000

rem = divisor (divisor)

$$\text{rem}[] = \text{div}[] = 110011000$$

$\text{rem} = \text{divide}(\text{divisor}, \text{rem})$

$\text{divide}(1011, 110011000)$

Divide function

$\text{while}(\text{true})$

{

for $i = 0$ to $\text{divisor.length} // 4$:

$i = 0 \rightarrow \text{rem}[0] = \text{rem}[0] \wedge \text{divisor}[0]$

$\text{rem}[0] = 1 \wedge 1 \Rightarrow 0$

$i = 1 \rightarrow \text{rem}[1] = 1 \wedge 0 \Rightarrow 1$

$i = 2 \rightarrow \text{rem}[2] = 0 \wedge 1 \Rightarrow 1$

$i = 3 \rightarrow \text{rem}[3] = 0 \wedge 1 \Rightarrow 1$

$\Rightarrow \text{rem}[] = 0111$

$\text{while}(\text{rem}[\text{cur}] == 0 \text{ \& \& } \text{cur} != \text{rem.length} - 1)$

$\text{rem}[0] = 0 \text{ \& \& } 0 != 8 \quad \text{T}$

$\boxed{\text{cur} = 1}$

$\text{rem}[1] == 0 \text{ F}$

for $i = 0$ to 4 :

$\text{rem}[1] = \text{rem}[1] \wedge \text{divisor}[0]$

$1 \wedge 1 \Rightarrow 0$

$\text{rem}[2] = 1 \wedge 0 \Rightarrow 1$

$\text{rem}[3] = 1 \wedge 1 \Rightarrow 0$

$\text{rem}[4] = 1 \wedge 1 \Rightarrow 0$

$\Rightarrow \text{rem}[] = 00100$

$\text{while}(\text{rem}[\text{cur}] == 0 \text{ \& \& } \text{cur} != \text{rem.length} - 1) \quad \text{T}$

$\boxed{\text{cur} = 2}$

for $i = 0$ to 4 :

$\text{rem}[2] = 1 \wedge 1 \Rightarrow 0$

$\text{rem}[3] = 0 \wedge 0 \Rightarrow 0$

$\text{rem}[4] = 0 \wedge 1 \Rightarrow 1$

$\text{rem}[5] = 1 \wedge 1 \Rightarrow 0$

$\text{rem}[] = 000010$

$\text{while}(\text{rem}[\text{cur}] == 0 \text{ \& \& } \text{cur} != \text{rem.length} - 1)$

$0 \quad 2 != 8$

$\text{T} \quad \boxed{\text{cur} = 3}$

$\text{rem}[3] == 0 \quad \text{T}$

$\text{rem}[4] == 0 \quad \text{F} \quad \text{rem}$

$\boxed{\text{cur} = 4}$

for $i=0$ to 4:

$$\text{rem}[4] = \text{rem}[4] \wedge \text{divisor}[0]$$

$$\text{rem}[5] = 1 \wedge 1 \Rightarrow 0$$

$$\text{rem}[6] = 0 \wedge 0 \Rightarrow 0$$

$$\text{rem}[7] = 0 \wedge 1 \Rightarrow 1 \quad \text{rem} = 00000011$$

$$0 \wedge 1 \Rightarrow 1$$

while ($\text{rem}[4] = 0$ & $\text{rem.length} - 1$)
 $0 = 0$ & $4 = 4$ \Rightarrow 8 \Rightarrow T

cur = 5

$$\text{rem}[5] = 0 \quad T$$

cur = 6

if ($\text{rem.length} - \text{cur} \geq \text{divisor.length}$)
 $8 - 6 \geq 4 \quad T$

return 00000011 to rem. in CRC generator

CRC Generator

for $i=0$ to 9:

$$\text{crc}[i] = \text{div}[i] \wedge \text{rem}[i]$$

$\text{crc}[0]$	$=$	1	\wedge	0	\Rightarrow	1
$\text{crc}[1]$	$=$	1	\wedge	0	\Rightarrow	1
$\text{crc}[2]$	$=$	0	\wedge	0	\Rightarrow	0
$\text{crc}[3]$	$=$	0	\wedge	0	\Rightarrow	0
$\text{crc}[4]$	$=$	1	\wedge	0	\Rightarrow	1
$\text{crc}[5]$	$=$	1	\wedge	0	\Rightarrow	1
$\text{crc}[6]$	$=$	0	\wedge	1	\Rightarrow	1
$\text{crc}[7]$	$=$	0	\wedge	1	\Rightarrow	1
$\text{crc}[8]$	$=$	0	\wedge	0	\Rightarrow	0

CRC Code: 110011110

Error Detection

Enter CRC code of 9 bits:

1
1
0
0
1
1
1
1
1

$CRC[9] = 110011111$

$rem[9] = CNLC[9] = 110011111$

$rem = divide(divisor, rem)$ // 1011, 110011111
// Perform division calling divide func.

if returned $rem \neq 0$
then "Error detected"

if $i == 8$:
then "No error"

Here $rem = 1$ for 110011111
Hence "Error is detected"

CNC Generator.

$$\begin{array}{r}
 111010 \\
 \hline
 1011 \overline{) 110011000} \\
 \underline{1011} \downarrow \\
 \cancel{0}1111 \\
 \underline{1011} \downarrow \\
 \cancel{0}1001 \\
 \underline{1011} \downarrow \\
 \cancel{0}0100 \\
 \underline{0000} \downarrow \\
 \cancel{0}1000 \\
 \underline{1011} \downarrow \\
 \cancel{0}0110 \\
 \underline{0000} \\
 \hline
 \cancel{0}110
 \end{array}$$

Message to be transmitted 110011110

Error Detection

$$\begin{array}{r}
 1011 \overline{) 11001111} \\
 \underline{1011} \downarrow \\
 01111 \\
 \underline{1011} \downarrow \\
 01001 \\
 \underline{1011} \downarrow \\
 00101 \\
 \underline{0000} \downarrow \\
 01011 \\
 \underline{1011} \downarrow \\
 00001 \\
 \underline{0000} \downarrow \\
 1
 \end{array}$$

Here reminder 1 is returned

Cross Check of Inventory

CRC PROGRAM OUTPUT

```
Output - CRC (run) X
run:
Enter number of data bits :
6
Enter data bits :
1
1
0
0
1
1
Enter number of bits in divisor :
4
Enter Divisor bits :
1
0
1
1
Dividend (after appending 0's) are : 110011000

CRC code :
110011110
Enter CRC code of 9 bits :
1
1
0
0
1
1
1
1
1
Error
THANK YOU.... :)
BUILD SUCCESSFUL (total time: 34 seconds)
```

```
Output - CRC (run) X
run:
Enter number of data bits :
6
Enter data bits :
1
1
0
0
1
1
Enter number of bits in divisor :
4
Enter Divisor bits :
1
0
1
1
Dividend (after appending 0's) are : 110011000

CRC code :
110011110
Enter CRC code of 9 bits :
1
1
0
0
1
1
1
1
0
No Error
THANK YOU.... :)
BUILD SUCCESSFUL (total time: 35 seconds)
```