

# LR(0) Items

- *LR(0) item* : An LR(0) item for a grammar  $G$  is a production rule of  $G$  with the symbol  $\bullet$  ( read as *dot* or *bullet*) inserted at some position in the rhs of the rule.
- Example of LR(0) items : For the rule given below

$decls \rightarrow decls\ decl$

the possible LR(0) items are :

$I_1 : decls \rightarrow \bullet decls\ decl$

$I_2 : decls \rightarrow decls\ \bullet decl$

$I_3 : decls \rightarrow decls\ decl\ \bullet$

The rule  $decls \rightarrow \epsilon$  has only one LR(0) item,

$I_4 : decls \rightarrow \bullet$

# LR(0) Items

- An LR(0) item is *complete* if the  $\bullet$  is the last symbol in the rhs.  
Example :  $I_3$  and  $I_4$  are complete items and  $I_1$  and  $I_2$  are incomplete items.
- An LR(0) item is called a *kernel item*, if the dot is not at the left end. However the item  $S' \rightarrow \bullet S$  is an exception and is defined to be a kernel item.  
Example :  $I_1, I_2, I_3$  are all kernel items and  $I_4$  is a non-kernel item.

# Canonical Collection of LR(0) Items

The construction of the SLR parsing table requires two functions *closure* and *goto*.

*closure*: Let  $U$  be the collection of all LR(0) items of a cfg  $G$ . Then  $\text{closure} : U \rightarrow 2^U$ .

- ①  $\text{closure}(I) = \{I\}$ , for  $I \in U$
- ② If  $A \rightarrow \alpha \bullet B\beta \in \text{closure}(I)$ , then the item  $B \rightarrow \bullet\eta$  is added to  $\text{closure}(I)$ .
- ③ Apply step (ii) above repeatedly till no more new items can be added to  $\text{closure}(I)$ .

Example: Consider the grammar  $A \rightarrow A a \mid b$

$$\text{closure}(A \rightarrow \bullet A a) = \{A \rightarrow \bullet A a, A \rightarrow \bullet b\}$$

# Canonical Collection of LR(0) Items

*goto*:  $goto : U \times X \rightarrow 2^U$ , where  $X$  is a grammar symbol.

$$goto(A \rightarrow \alpha \bullet X\beta, X) = closure(A \rightarrow \alpha X \bullet \beta).$$

Example:  $goto(A \rightarrow \bullet Aa, A) = closure(A \rightarrow A \bullet a) = \{A \rightarrow A \bullet a\}$   
*closure* and *goto* can be extended to a set  $S$  of  $LR(0)$  items by appropriate generalizations

- $closure(S) = \bigcup_{I \in S} \{closure(I)\}$
- $goto(S, X) = \bigcup_{I \in S} \{goto(I, X)\}$