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# ***Module 4 – Number System***

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# Module Overview

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- **Data in Computers**
- **Binary Numbering System**
- **Binary Arithmetic**
- **IEEE Floating Point Representation**



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# Data in Computers

# Data in Computers



- Computers understand only two things – 0 and 1
  - Data is stored only the form of combinations of 0s and 1s
- Computers use switches to store 0 or 1
  - **Stores 1 if switch is ON**
  - **Stores 0 if switch is OFF**
- Computers consists of Integrated circuits (IC) of billions of switches; allow storage of huge amounts of information.

# Data in Computers



- Kinds of data we store in computers:

- Integer numbers

- 1, 3, 94, -103, etc.

- Floating point numbers

- 1.00433, 54.9090354598, etc.

← We will study more deeply

- Characters

- 'A', 'a', '#', etc.

- Strings

- "Delhi", "Gopal", etc.

- etc.

All of them are stored as binary patterns, i.e. strings of 0s and 1s.



# Binary Numbering System

# The *human* numbering system



We use digits

0 to 9  $\rightarrow$  10 symbols (digits) (the decimal system)

What is so special about the number 10?

Nothing!



# The numbering system for computers



- Computers use binary numbering system
  - i.e. it has only two symbols to represent numbers – 0 and 1
- Just like humans have 10 symbols - (0 to 9)
- Other systems
  - Octal – 8 symbols
    - 0, 1, 2, 3, 4, 5, 6, 7
  - Hexadecimal – 16 symbols
    - 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, F



# Decimal vs Binary vs Octal vs Hexadecimal: An example



Decimal	Binary	Octal	Hexadecimal
0	0000	0	0
1	0001	1	1
2	0010	2	2
3	0011	3	3
4	0100	4	4
5	0101	5	5
6	0110	6	6
7	0111	7	7
8	1000	10	8
9	1001	11	9
10	1010	12	A
11	1011	13	B
12	1100	14	C
13	1101	15	D
14	1110	16	E
15	1111	17	F

# How a decimal number is interpreted?



Eg.: 357

Digits ->	3	5	7
Weights->	$10^2$	$10^1$	$10^0$
	MSD		LSD

Perform sum over Digits\*Weights:

$$3*10^2 + 5*10^1 + 7*10^0$$
$$= 357$$

# Range of binary numbers



One bit – up to decimal 1

Two bits – up to 3

Three bits – up to 7

In general:

$n$  bits -  $2^n - 1$  in decimal

# Examples (Worked out on Board)



- Convert 10101 from binary to decimal
- Convert 16 from decimal to binary
- Convert 23 from decimal to binary

# Negative Binary Numbers



- How do we signify negative numbers in arithmetic?
- The – symbol to the left of MSD
- So... 1111 is 15, then -1111 should -15
- But... computers cannot handle any symbol apart from 0 and 1



# Negative Binary Numbers

- To handle negative numbers: **left-most bit is used to indicate the sign.**
- Known by various names: Most Significant Bit (**MSB**), sign-bit or high-order bit
- Positive numbers: **MSB is 0**
- Negative numbers: **MSB is 1**
- For 8 bit unsigned numbers: **range of 0 to 255 ( $2^8 - 1$ )**
- What about signed numbers?
  - **-127 to 127 or -128 to 127**
- How does the computer know whether to treat a number as **signed or unsigned**?
  - **It cannot.** It's the programmer's job to tell.

# Negative Binary Numbers



Three schemes for representing negative binary numbers:

- Signed-magnitude representation
  - One's complement representation
  - Two's complement representation
- 
- We will discuss them with respect to 8-bit integers

# Signed-magnitude representation



MSB is the sign as usual

7 bits for the magnitude or the value

Range: -127 to 127

Examples:

109            01101101

-109           11101101

127            01111111

-127           11111111

**What about zero?**



# 1's complement



As before, MSB indicates the sign.

Negative no. = **One's complement** of the positive number

**One's complement** → Invert all the bits → 1s to 0s and 0s to 1s

Range: -127 to 127

Examples:

15      00001111

-15     11110000

85      01010101

-85     10101010

**What about zero?**

# 2's complement



**2's complement of a no. = it's 1's complement + 1**

Range: -128 to 127

Example: **calculating the two's complement representation of -15**

Decimal 15                      00001111

In one's complement        11110000

Adding 1                                +1

In two's complement        11110001

**What about zero?**

**Most widely used!!**



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# Binary Arithmetic

# Binary arithmetic rules

- Addition of binary bits

- $0 + 0 = 0$

- $0 + 1 = 1$

- $1 + 0 = 1$

- $1 + 1 = 0$      1 is carry. In binary addition carry is discarded

- Subtraction of binary bits:

- $0 - 0 = 0$

- $1 - 0 = 1$

- $1 - 1 = 0$

- $0 - 1 = 1$      Borrow 1 from next high order digit

**Note:** for a positive number, its binary representation in SMF, 1CT or 2CT – it is the same number itself

# Subtraction using 2 CT



- Subtraction using 2CT:

1. Write the 2CT of the subtrahend
2. Add it to the minuend
3. If there is a carry over discard it, the remaining bits gives the result
4. If there is no carry over, find the 2CT of the result and add –ve sign before it

# Subtraction using 2 CT: Example 1



Consider  $7 - 4$ :

1.  $2CT(0\ 1\ 0\ 0) = 1\ 0\ 1\ 1 + 0\ 0\ 0\ 1 \rightarrow 1\ 1\ 0\ 0$

2.  $0\ 1\ 1\ 1$

$$\begin{array}{r} +\ 1\ 1\ 0\ 0 \\ \hline \end{array}$$

$$1\ 0\ 0\ 1\ 1 \rightarrow +3$$

Discard

# Subtraction using 2CT: Example 2



Consider  $4 - 7$

1.  $2CT(7) = 1000 + 0001 = 1001$

2.  $0100$

$+ 1001$

---

$1101$

3.  $2CT(1101) = 0010 + 0001 \rightarrow 0011$

4. Final answer is -ve of the result obtained i.e., -3

# Addition in 2's Complement



- **Overflow**

- Only possible cases:

$$\begin{array}{r} 0111 (+7) \\ 0101 (+5) \\ \hline 01100 (-4) \times \end{array}$$

*Overflow* (indicated by an arrow pointing to the purple 0)

$$\begin{array}{r} 1111 (-1) \\ 0111 (+7) \\ \hline 10110 (+6) \end{array}$$

*Carry* (indicated by an arrow pointing to the purple 1)

$$\begin{array}{r} 1010 (-6) \\ 1001 (-7) \\ \hline 10011 (+3) \times \end{array}$$

*Overflow* (indicated by an arrow pointing to the purple 1)

$$\begin{array}{r} 1110 (-2) \\ 1101 (-3) \\ \hline 11011 (-5) \end{array}$$

*Carry* (indicated by an arrow pointing to the purple 1)



# Addition in 2's Complement



## Observations:

- When addition is performed on two numbers having same MSB
- When  $C_3 \neq C_4$

$$\begin{array}{r}
 C_4=1 \quad C_3=0 \quad C_2=0 \quad C_1=0 \quad C_0=0 \\
 \begin{array}{rcccccl}
 & 1 & 0 & 1 & 0 & (-6) \\
 & 1 & 0 & 0 & 1 & (-7) \\
 \hline
 1 & 0 & 0 & 1 & 1 & (+3)
 \end{array}
 \end{array}$$

$$\begin{array}{r}
 C_4=1 \quad C_3=1 \quad C_2=0 \quad C_1=0 \quad C_0=0 \\
 \begin{array}{rcccccl}
 & 1 & 1 & 1 & 0 & (-2) \\
 & 1 & 1 & 0 & 1 & (-3) \\
 \hline
 1 & 1 & 0 & 1 & 1 & (-5)
 \end{array}
 \end{array}$$

- Whenever an overflow occurs MSB gets modified
- Addition of two  $N$  bits integers  $B_1$  and  $B_2$  causes overflow only if
  - $B_1 + B_2 > +2^{N-1} - 1$  and  $B_1 + B_2 < -2^{N-1}$  addition is

➤ **Overflows occur when the sum does not fit in the given number of bits**



# IEEE Floating Point Representation

# Converting real numbers to binary



- Real numbers:  $X.Y$

- $X$  – Similar to the way integers are represented (Explained in previous slides)
- $Y$ 
  - $Y * 2 = X'.Y'$
  - Store  $X'$  and  $Y = Y'$
  - Repeat step 1 and 2, for the new fractional value obtained until  $X'.Y' = 1.0$
  - Print all the value of  $X'$ .

Example - Convert 10.6875 to binary

$$(10)_{10} = (1010)_2$$

$$(0.6875)_{10} =$$

$$0.6875 * 2 = 1.375 \text{ -----} \rightarrow 1$$

$$0.375 * 2 = 0.75 \text{ -----} \rightarrow 0$$

$$0.75 * 2 = 1.5 \text{ -----} \rightarrow 1$$

$$0.5 * 2 = 1.0 \text{ -----} \rightarrow 1$$

$$= (1010.1011)_2$$

# Converting binary real numbers to decimal



- Binary number: **X.Y**
  - **X** – Similar to the way integer values are obtained (Explained in previous slides)
  - **Y**
    - Multiply each bit with the weight of its position, where weight of positions are:  $2^{-1}$ ,  $2^{-2}$ ,  $2^{-3}$ , .....
    - Take the sum of the output

▪ E.g. 1 0 1 0 . 1 0 1 1 ←

↓                      ↓  
Integral part    Fractional part

**Exercise:** Convert this binary floating-point number to decimal system

# Storing Real number in Computers



Steps to convert a real number in decimal to binary:

1. First, we convert the real number in decimal to a binary real number.
  - E.g. **10.6875** is converted into **1010.1011**
2. Then we normalize the binary real number.
  - E.g. **1010.1011** can be written as  **$1.0101011 \times 2^3$**
3. Encode the normalized binary real number into **IEEE 754 32-bit or 64-bit floating point format**

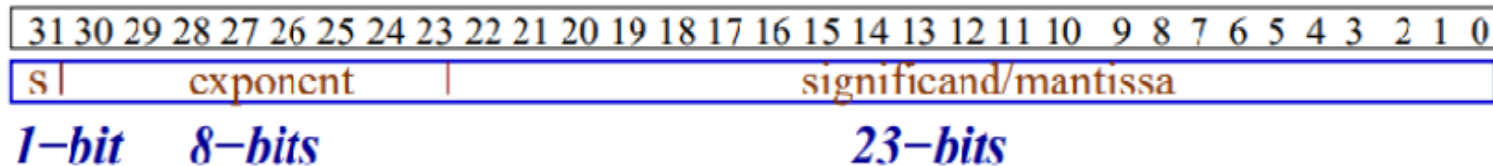
*Let us study the IEEE 754 Floating Point Representation*

# IEEE 754 floating point representation

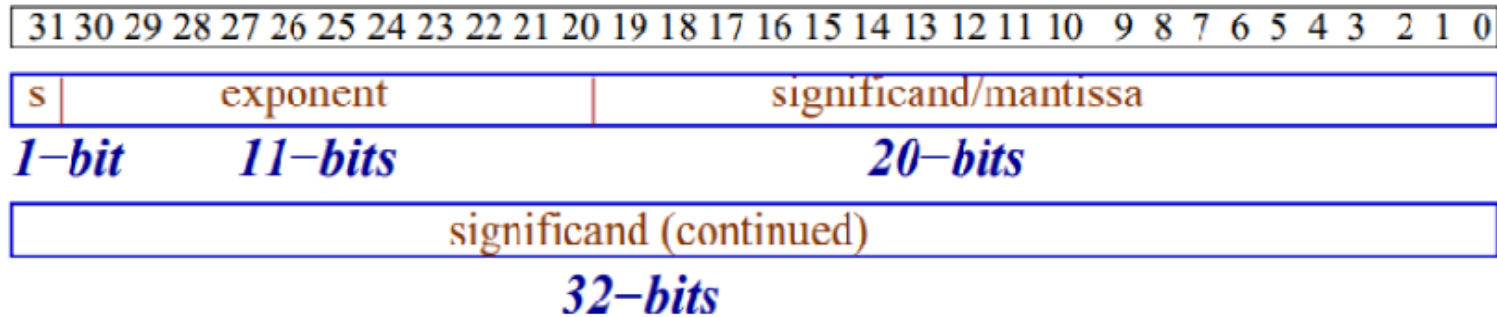


- Three pieces of information are required:
  - Sign of the number (**s**)
  - Significant value (**m**)
  - Signed exponent of 10 (**e**)
- The data type float uses IEEE 32-bit single precision format and the data type double uses IEEE 64-bit double precision format

# IEEE 754 floating point representation (contd.)



**Single precision (32 bit)**



**Double precision (64 bit)**

# IEEE floating point representation



- In single precision (32 bits):
  - The MSB indicates the sign bit (s): 1 is -ve and 0 is +ve.
  - The next 8 bits: Stores the value of the signed exponent (e)
  - Remaining 23: Stores the significand also known as *mantissa* (m).
  - Value of a floating number IN 32 bits is

$$(-1)^s \times 1.m \times 2^{e-127}$$

- **e – 127**: Excess 127 representation, meaning if the exponent value is 5, it will be stored as  $5 + 127 = 132$  in the exponent field



# Storing 10.6875 in IEEE 754 32-bit Floating Point Format



10.6875  $\rightarrow$  1010.1011  $\rightarrow$  1.0101011  $\times 2^3$

binary

Normalised binary



$s = 0$

$e = (3+127)_2 = (130)_2 = 10000010$

$m = 010101100000000000000000$

which is

0 10000010 010101100000000000000000

IEEE 754 32-bit Floating Point representation

# Conversion to decimal real number (Worked out on board)



Convert the following number to decimal:

1 00000111 110000000000000000000000000000

It is

$$-1.75 \times 2^{(7-127)} = -1.316554 \times 10^{-36}$$

# IEEE floating point representation



## Why excess 127?

- Helps in natural ordering in the values stored in the exponent field
- Consider  $2^{-126}$  and  $2^{+126}$ 
  - **Case 1:** In absence of Excess 127:
    - $-126 = 10000010$
    - $+126 = 01111110$ 
      - While  $-126$  is a smaller value the exponent field contains larger value ( $10000010 > 01111110$ )
  - **Case 2:** In presence of Excess 127:
    - $-126$  is written as  $1$  ( $1 - 127 = -126$ )  $\rightarrow 00000001$
    - $+126$  is written as  $253$  ( $253 - 127 = 126$ )  $\rightarrow 11111101$
    - $-126$  is a smaller value and so is the value stored in the exponent ( $00000001 < 11111101$ )

# Example 1



Convert 12.375 into 32-bit IEEE 754 Floating Point Format

1.  $(12)_{10} + (0.375)_{10} = (1100)_2 + (0.011)_2 = (1100.011)_2$
2. Shift  $(1100.011)_2$  by 3 digits to normalize the value

$$(12.375)_{10} = (1.100011)_2 * 2^3$$

$$s = 0$$

Exponent = 130 (represented in excess 127)

Mantissa = 100011

FP Representation  $\rightarrow$  0 | 10000010 | 1000110000.....

# Example 2



Convert the following binary number in 32-bit IEEE 754 floating point format into decimal:

1 1011 0110 011 0000 0000 0000 0000

$$(-1)^1 \times 2^{10110110 - 01111111} \times 1.011$$

$$= -1.375 \times 2^{55}$$

$$= -49539595901075456.0$$

$$= -4.9539595901075456 \times 10^{16}$$

# Example 3 (Work out on board)



Convert -10.7 into 32-bit IEEE 754 Floating Point Format

You will see a recurring pattern of bits that is never-ending.

What should you do?

Store only the bits that can fit your 23-bits mantissa. Rest are truncated.

This leads to approximation.



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***Thank you***  
**Q & A**