



Johnson Algorithm

Parallel Implementation in Rust

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Overview

Johnson's algorithm is a way to find the shortest paths between all pairs of vertices in a sparse, edge weighted, directed graph. It allows some of the edge weights to be negative numbers, but no negative-weight cycles may exist.

Goals

1. Formulate a parallel approach for Johnson algorithm and implement in Rust
2. Analyze and improve the performance.

Algorithm

```
1.  procedure JOHNSON_SINGLE_SOURCE_SP( $V, E, s$ )
2.  begin
3.     $Q := V$ ;
4.    for all  $v \in Q$  do
5.       $l[v] := \infty$ ;
6.     $l[s] := 0$ ;
7.    while  $Q \neq \emptyset$  do
8.      begin
9.         $u := \text{extract\_min}(Q)$ ;
10.       for each  $v \in \text{Adj}[u]$  do
11.         if  $v \in Q$  and  $l[u] + w(u, v) < l[v]$  then
12.            $l[v] := l[u] + w(u, v)$ ;
13.       endwhile
14.    end JOHNSON_SINGLE_SOURCE_SP
```

Implementation in Rust

```

while let Some(que_node { cost , vertex }) = heap.pop() {
    if cost > distance[vertex as usize] {
        continue;
    }

    for i in 0..nodes {
        let mut weigh = adj_matrix[vertex as usize][i as usize];
        if(i != vertex && weigh != 0 ){
            let next = que_node { cost: cost + weigh, vertex: i };
            if next.cost < distance[next.vertex as usize] {
                heap.push(next);
                distance[next.vertex as usize] = next.cost;
            }
        }
    }
}

```

The algorithm

Parallelizing Step

```

for mut j in 0..limit {
    for i in j..NTHREADS + j {
        let adj_list = adj_list.clone();
        children.push(thread::spawn(move || {
            Johnson(i as u32, nodes as u32, edges as u32 , &adj_list);
        }));
    }
}

```

How to run the code

git clone <https://github.com/praneelrathore/etig>

cd etig/src

uncomment `parallel::johnson(&mut gra, na);`

cargo build

cargo run input_file

Make sure the main.rs contains `extern crate algorithms; use algorithms::parallel;`

Sample Input and Output

```
~/rust-algo master*
> more input.in
5
3 1 0 0 3
1 4 4 2 4
2 1 0 4 3
2 2 4 0 2
0 0 2 0 0

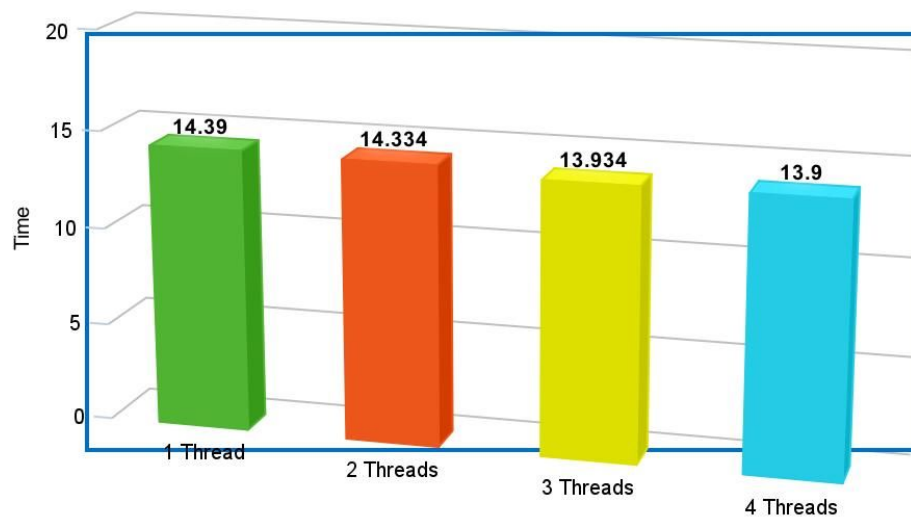
~/rust-algo master*
> cargo run input.in
    Running `target/debug/etig input.in`
min distance of 0 = 0
min distance of 1 = 1
min distance of 2 = 5
min distance of 3 = 3
min distance of 4 = 3
```

Analysis of Performance

Nodes = 500

Threads	Time Taken
1	14.39
2	14.33
3	13.93
4	13.90

Time Vs Parallel Threads 500 Nodes



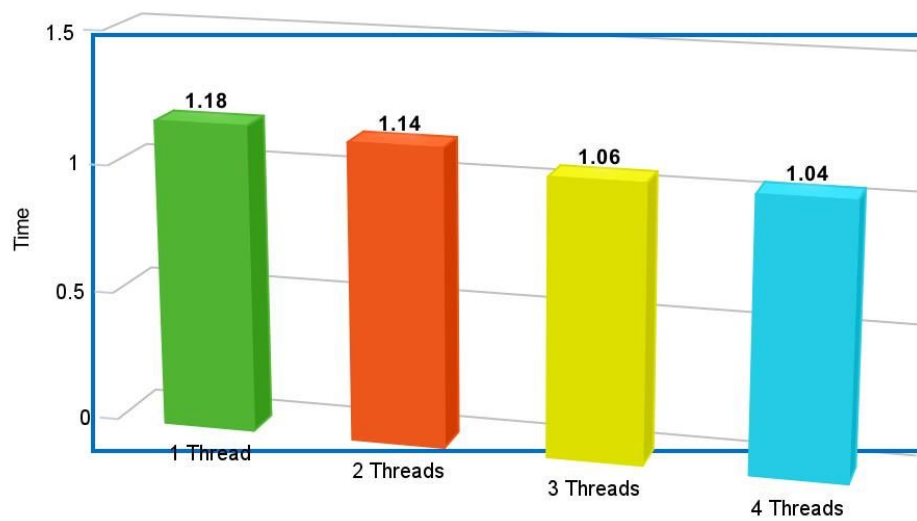
Series 1

Analysis of Performance

Nodes = 100

Threads	Time Taken
1	1.18
2	1.14
3	1.06
4	1.04

Time Vs Parallel Threads 100 Nodes



Series 1

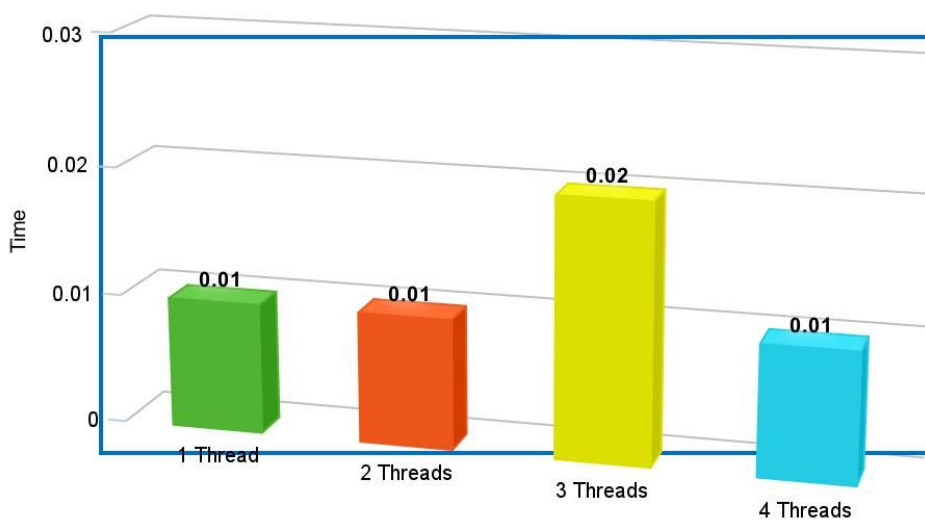
Analysis of Performance

Nodes = 10

Thread creation overhead comes to play as time for 3 threads became greater than for 2

Threads	Time Taken
1	0.01
2	0.01
3	0.02
4	0.01

Time Vs Parallel Threads 10 Nodes



Series 1

Time Complexity

Without parallelization

The time complexity of this algorithm, using Fibonacci heaps in the implementation of Dijkstra's algorithm, is $O(V^2 \log V + VE)$ and $O(V \log V + E)$ for each of V instantiations of Dijkstra's algorithm. Thus, when the graph is sparse, the total time can be faster than the Floyd–Warshall algorithm, which solves the same problem in time $O(V^3)$.

After Parallelization

The time complexity after parallelization is

$$O(K^2 * \text{Log} V + KE) + O(K * \text{Log} V + V)$$

$$K = V / \text{NO_OF_THREADS}$$

$$V = \text{No of Vertices}$$