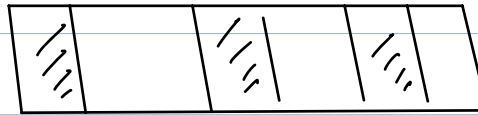


Issues with Array.



int \rightarrow 4 Bytes

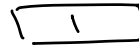
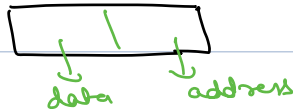
\downarrow
array \rightarrow 5

\downarrow
20 Bytes.

LL \rightarrow D.S., \rightarrow not necessarily stores the data contiguously.

LL

- A linear data structure that can utilize all the free memory
- We need not have continuous space to store nodes of a Linked List.



class Node

int data;

3

\rightarrow ok.

Node n = new Node();

Node temp = n;

data type \rightarrow variable.

ok

data = 0;

class Node

int data;

Node next;

Node (x) {

data = x;
next = null;

first node address is
called head of the LL.

10k

data = 100
next = 20k

20k

data = 200
next = 30k

30k

data = 300
next = null

Node t = new Node(100);

t.next = new Node(200);

t.next.next = new Node(300);

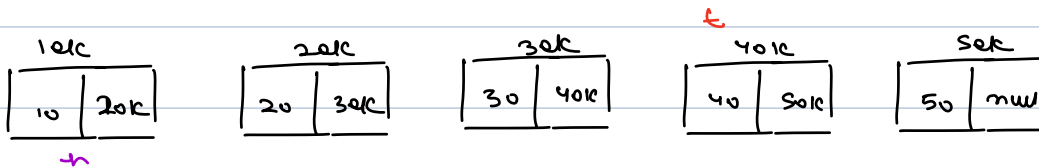
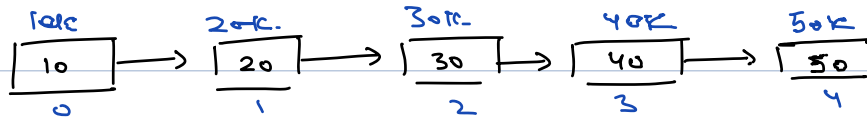
30k

func (t) {

}

Operations on Lk :-

$k=3 \rightarrow 40$
 $k=1 \rightarrow 20$



1) Access k^{th} idx:

T.C $\rightarrow O(n)$

int k^{th} (Node head, int k) {

Node temp = head;

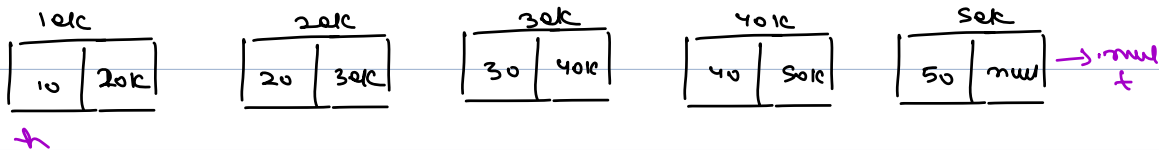
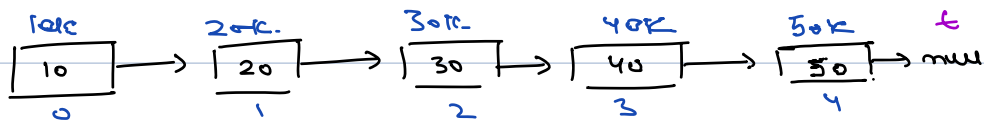
for (i = 0; i < k; i++) {

temp = temp->next;

return temp->data;

$i=0$
 $i=1$
 $i=2$

2) check for value x.



$x = 30 \rightarrow \text{True}$
 $x = 100 \rightarrow \text{false}$.

bool check (node head, int x) {

node temp = head;

while (temp != null) {

if (temp.data == x) {

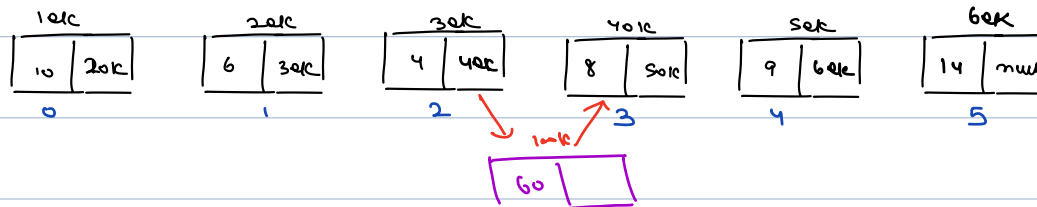
return True;

temp = temp.next;

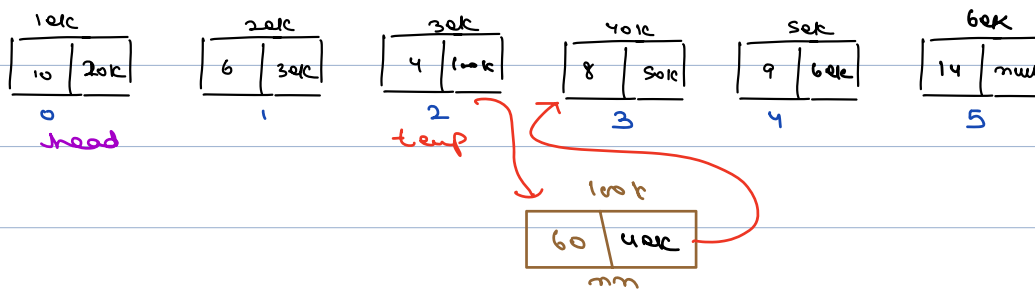
return false;

T.C $\rightarrow O(n)$

Ques Insert a new Node with Data.



Ex 1 $V = 60$, $P = 3$



Node nn = new Node(V);

Node temp = head;

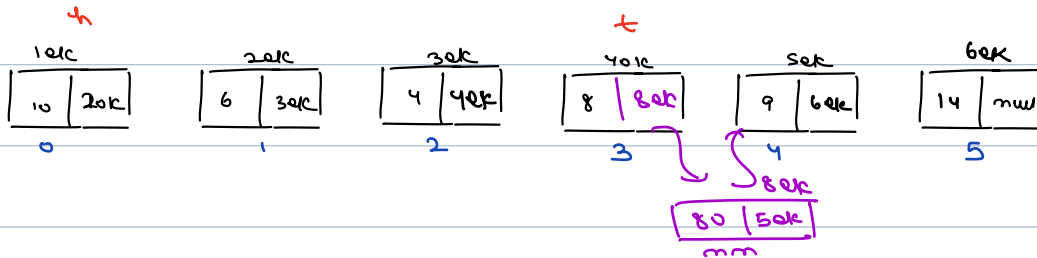
for (i = 0; i < P - 1; i++)

temp = temp.next

nn.next = temp.next

temp.next = nn;

ex1 $V = \underline{80}$, $P = \underline{\cancel{4}} \underline{6}$



Node nn = new Node(V);

Node temp = head;

for (i = 0; i < P - 1; i++) {

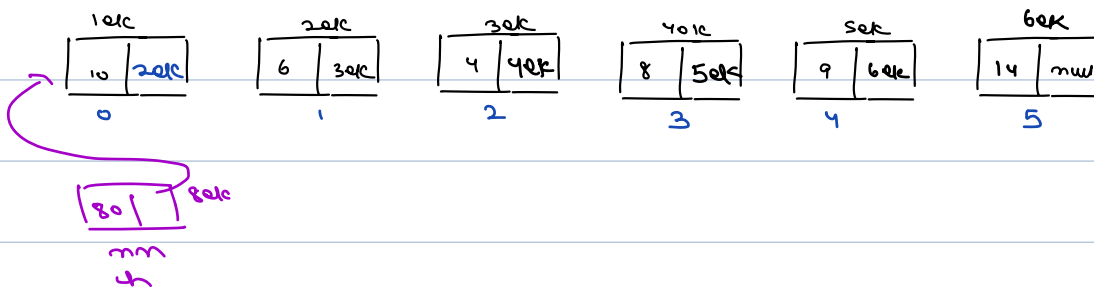
temp = temp->next;

nn->next = temp->next;

temp->next = nn;

Edge case :-

ex1 $V = \underline{80}$, $P = \underline{0}$



T.C $\rightarrow O(n)$.

Node insertAtK (Node head, int v, int p) {

Node nn = new Node(v)

if (p == 0) {

nn.next = head;
head = nn;
return head;

Node temp = head;

for (i = 0; i < p - 1; i++) {

temp = temp.next

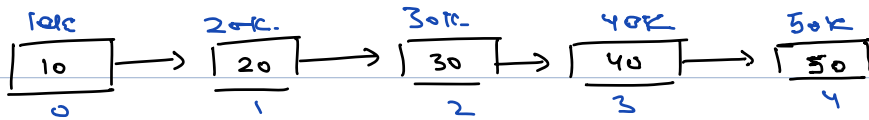
nn.next = temp.next ✓

temp.next = nn;

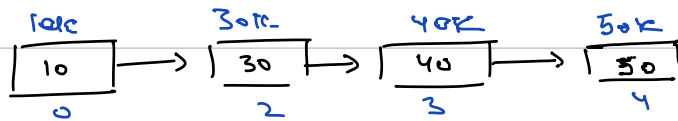
return head;

}

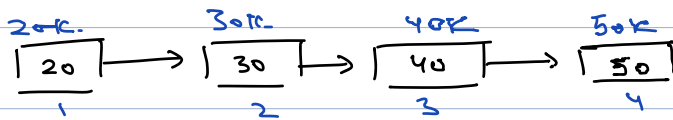
Ques Deletion in a LL.



e.g.1) $x = 20$

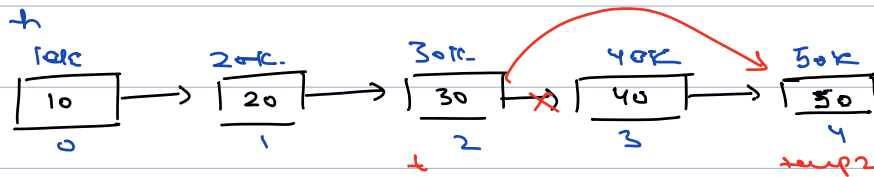


e.g.2) $x = 10$



Soln

$x = 40$.



Node temp = head;

while () {

if (temp->next->data == x) {

temp2 = temp->next->next;

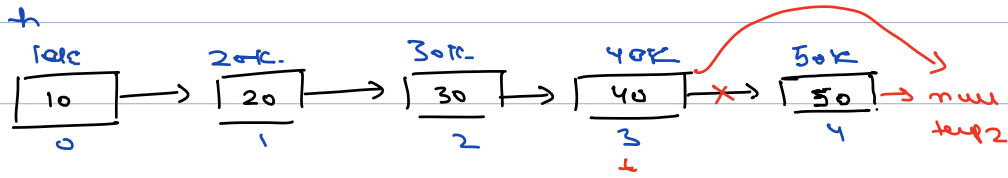
temp->next = temp2;

if (temp == head) head = temp2;

temp = temp->next;

}

$x = 50$.



Node temp = head;

while () {

if (temp->next->data == x) {

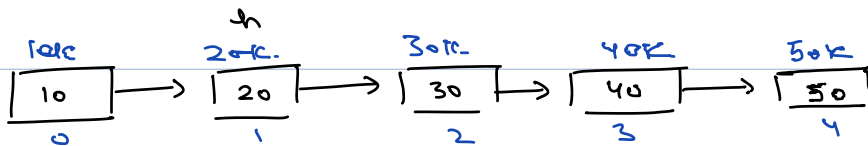
temp2 = temp->next->next;

temp->next = temp2;

return head;

temp = temp->next;

$x = 10$.

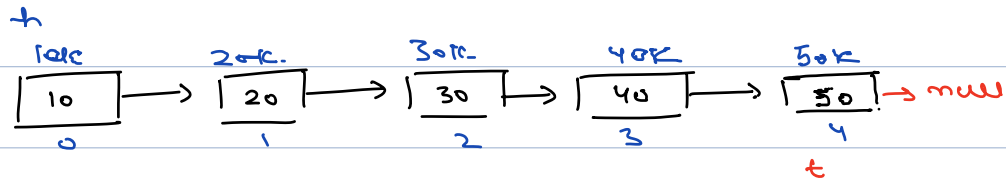


if (head->data == x) {

head = head->next;

return head;

$x = 50$.



Node temp = head;

while (temp.next != null) {

if (temp.next.data == x) {

temp2 = temp.next.next;

t.next = temp2;

delete head;

temp = temp.next;

node delete (node head, 'int x') {

if (head.data == x) {

T.C $\rightarrow O(n)$.

head = head.next;
return head;

Node temp = head;

while (temp.next != null) {

if (temp.next.data == x) {

temp2 = temp.next.next;

temp.next = temp2;

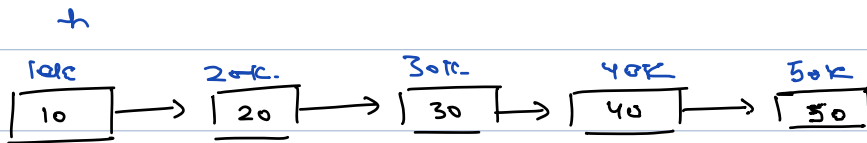
return head;

temp = temp.next;

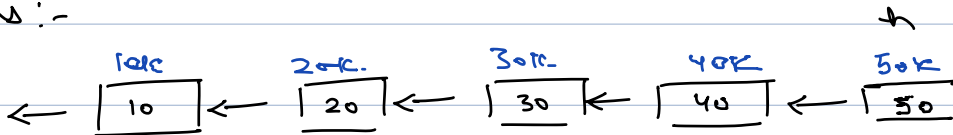
return head;

}

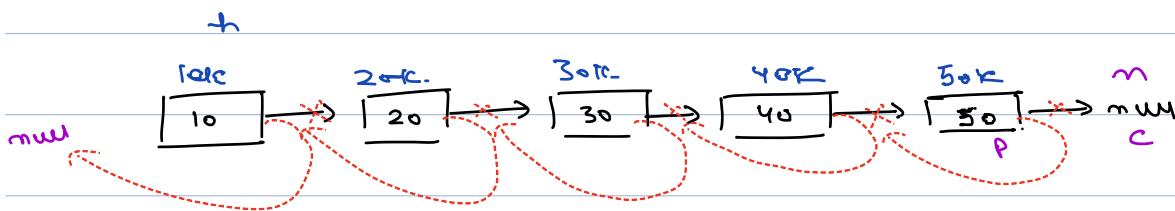
Reverse a LL



Ans:-



Soln



T.C $\Rightarrow O(n)$

node reverse (node head) {

node curr = head;

node prev = null;

while (curr != null) {

next = curr->next;

curr->next = prev;

prev = curr;

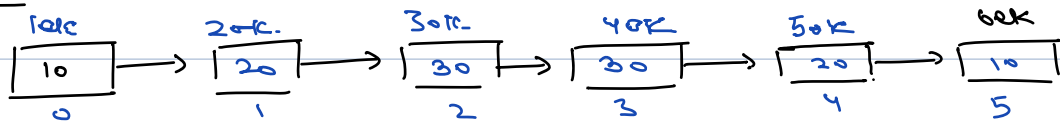
curr = next;

head = prev;

return head;

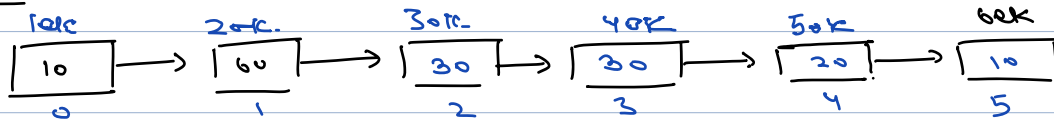
Ques Given a LL, check if it is palindrome.

e.g 1



return True

e.g 2



return false

Soln :-

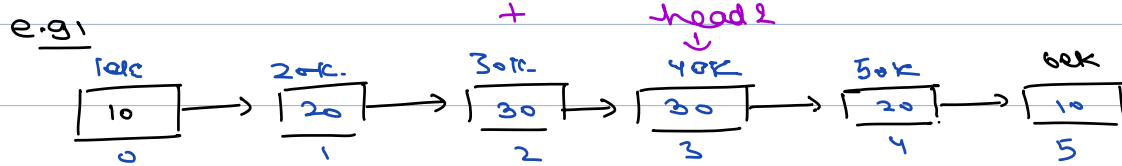
1) create a copy of LL,

2) reverse it

3) compare

T.C $\rightarrow O(n)$

S.C $\rightarrow O(n)$



step-1:-

find len of LL

n = 0;

temp = head;

while (temp != null) {

|
3

n++;
temp = temp.next;

// n = 6.

int halfLen = n / 2

node temp = head;

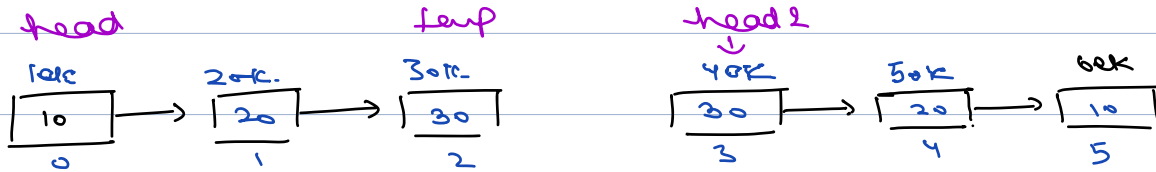
for (i = 0; i < halfLen - 1; i++) {

|
3

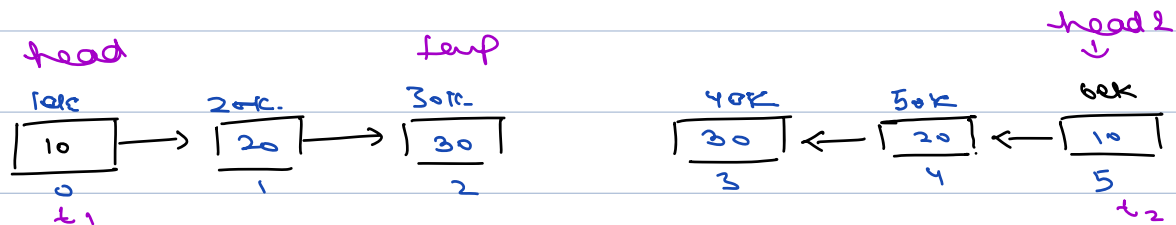
temp = temp.next;

node head2 = temp.next;

temp.next = null;



head 2 = reverse (head 2);



t₁ = head;
t₂ = head 2;

while (t₁ != null & t₂ != null) {

if (t₁.data != t₂.data) {

return false;

t₁ = t₁.next;

t₂ = t₂.next;

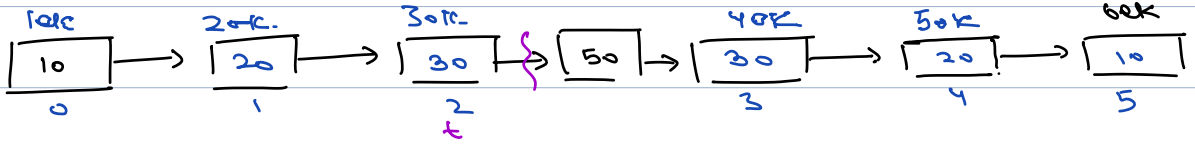
return true;

T.C → O(n)

S.C → O(1)

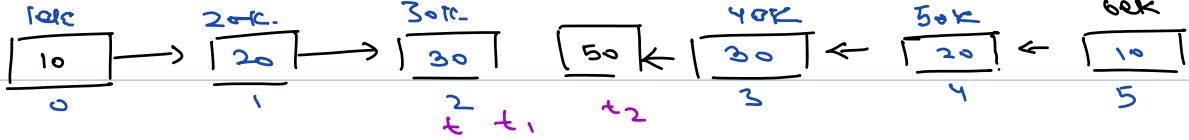
$n = 7$
half len = 3

head



head 2

head



head 2

人 心 二