Given a string **s** of lowercase characters, return the **count of pairs (i,j)** such that **i < j** and **s[ i ] is 'a'** and **s[ j ] is 'g'**.

String s = "abegag" Ans ⇒ 3.

String s = "acgdgag" Ans ⇒ 4.

String s = "bcaggaag" Ans ⇒ 5.

1) Brute force :-

String s = "abegag"

n = 6.

s = a a a g g

T.C → O(n²)

S.C → O(1)

```
int ans = 0;
for (i = 0; i < n; i++) {
    if (s[i] == 'a') {
        for (j = i+1; j < n; j++) {
            if (s[j] == 'g') {
                ans++;
            }
        }
    }
}
return result;
```

## Optimized Solution :-

String s = "abegag"  (0 1 2 3 4 5)

## Observations :-

- For every **'g'**, we need to know the count of **'a'** on left side of **'g'**.
- We will store the count of **'a'** and whenever **'g'** is encountered, we will add the count of **'a'** to the result.

## Example :-

| | a | c | b | a | g | k | a | g | g |
|---|---|---|---|---|---|---|---|---|---|
| count a = 0 | 1 | 1 | 1 | 2 | 2 | 2 | 3 | 3 | 3 |
| Ans = 0 | 0 | 0 | 0 | 0 | 2 | 2 | 2 | 5 | 8 |

```
count a = 0, ans = 0;

for (i = 0; i < n; i++) {
    if (s[i] == 'a') {
        count a++;
    }
    else if (s[i] == 'g') {
        ans += count a;
    }
}

return ans;
```

T.C → O(n)
S.C → O(1).

## carry forward :-

## ← Subarrays →

A subarray is a contiguous part of an array. It is formed by selecting a range of elements from the array. A subarray can have one or more elements and must be a contiguous part of the original array.

arr[] → 4, 1, 2, 3, -1, 6, 9, 8, 12

(2, 3, -1, 6) ✓

(9) ✓

4, 1, 2, 3, -1, 6, 9, 8, 12 ✓

(4, 12) ✗

(1, 2, 6) ✗

(3, 2, 1, 4) ✗

### Quiz

**A[] = { 2, 4, 1, 6, -3, 7, 8, 4}**

## Representation of a subarray :-

→ Two ways

1) Start & end idx
2) Start Idx & len

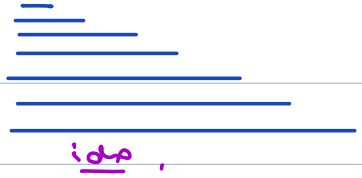
| | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
|---|---|---|---|---|---|---|---|---|---|
| arr[] → | 4, | 1, | 2, | 3, | -1, | 6, | 9, | 8, | 12 |

(2, 3, -1, 6)
- 1st way (2 to 5) — sidx, eidx
- 2nd ways (2 & 4) — sidx, len

(9)
- 1st way (6 to 6)
- 2nd way (6, 1).

**How many subarrays of the following array start from index 0**

**[4, 2, 10, 3, 12, -2, 15]**



idx. → # subarrays

| | |
|---|---|
| (4) | (0 - 0) |
| (4, 2) | (0 - 1) |
| (4, 2, 10) | (0 - 2) |
| (4, 2, 10, 3) | (0 - 3) |
| (4, 2, 10, 3, 12) | (0 - 4) |
| (4, 2, 10, 3, 12, -2) | (0 - 5) |
| (4, 2, 10, 3, 12, -2, 15) | (0 - 6) |

**How many subarrays of the following array start from index 1**

**[4, 2, 10, 3, 12, -2, 15]**

(0 1 2 3 4 5 6)

Subarray (s,e)

Ans => 6.

(1,1) → 2

(1,2) → 2, 10

(1,3) → 2, 10, 3

(1,4) → 2, 10, 3, 12

(1,5) → 2, 10, 3, 12, -2

(1,6) → 2, 10, 3, 12, -2, 15

arr[n]

0 1 2 3 4 ---- n-1

| Subarrays starting with 0 | Starting with 1 | Starting with 2 |
|---|---|---|
| (0-0) | (1-1) | (2-2) |
| (0-1) | (1-2) | (2-3) |
| (0-2) | (1-3) | (2-4) |
| . | . | . |
| . | . | . |
| . | . | . |
| (0-(n-1)) | (1-(n-1)) | (2-(n-1)) |
| n | n-1 | n-2 |

Subarrays starting with (n-1),

$$\frac{((n-1)-(n-1))}{1}$$

Total subarrays in an array

$$n + (n-1) + (n-2) + \dots 2 + 1 \Rightarrow \frac{n(n+1)}{2}$$

arr → (1, 2, 3) → n=3 $\quad \frac{3(3+1)}{2} \Rightarrow 6.$

Subarrays

(1)

(2)

(3)

(1,2)

(2,3)

(1,2,3)

Given an array of integers and two indices, a start index and an end index, we need to print the subarrays of the array that starts from the start index and ends at the end index (both inclusive).

0 1 2 3 4
(10, 20, 30, 40, 50)

s = 2, e = 4. ⟶ (30, 40, 50).

```
void printSubarray(int arr[], int start, int end) {
    for (int i = start; i <= end; i++) {
    print(arr[i] << " ";)
    }
    cout << endl;
}
```

T.C → O(n)

S.C → O(1).

Ques Print all Subarrays of a given array.

0 1 2
arr = [10, 20, 30]

$\frac{n(n+1)}{2} \Rightarrow \frac{3(3+1)}{2} \Rightarrow 6.$

(0-0) → 10

(0-1) → 10, 20

(0-2) → 10, 20, 30

(1-1) → 20

(1-2) → 20, 30

(2-2) = 30

arr = [10, 20, 30]
0 1 2

```
for ( s=0; s<n; s++) {
    for (e=s; e<n; e++) {
        print (s + " " + e);
    }
}
```

| s | e |
|---|---|
| 0 | 0 |
| 0 | 1 |
| 0 | 2 |
| 1 | 1 |
| 1 | 2 |
| 2 | 2 |

(0 0) → 10
(0 1) → 10, 20
(0 2) → 10, 20, 30
(1, 1) → 20
(1, 2) → 20, 30
(2, 2) → 30

T.C → O(n³)
S.C → O(1).

```
for ( s=0; s<n; s++) {
    for (e=s; e<n; e++) {
        for (i=s; i<=e; i++) {
            print (arr[i]);
        }
    }
}
```

## Ques

**Given an array of N integers, return the length of smallest subarray which contains both maximum and minimum element of the array.**

| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 |
|---|---|---|---|---|---|---|---|---|---|---|
| 2 | 2 | 6 | 4 | 5 | 1 | 5 | 2 | 6 | 4 | 1 |

Max = 6
Min = 1

Ans → 3.

Brute force :-

Check all subarrays & find ans.

T.C → $O(n^3)$

S.C → $O(1)$.

0 1 2 3 4 5 6 7 8 9 10

**2 2 6 4 5 1 5 2 6 4 1**

Maxm = 6

min = 1

Ans → 3.

Observations :-

1) we need only 1 max & 1 min in that.

2) Subarray ( no max & min will be there in b/w)

6 or 1

1 1 6 max max min

min max

1 6 5 ← min & max will be on the corners.

min max

max min

Dry Run :-

last_min_idx = ~~-1~~ 8 10
last_max_idx = ~~-1~~ 7 8
minValue = 1
MaxValue = 6
ans = ~~+∞~~ X 3

| idx → | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 |
|---|---|---|---|---|---|---|---|---|---|---|---|
| | 2 | 2 | 6 | 4 | 5 | 1 | 5 | 2 | 6 | 4 | 1 |

```
int min = Todo
int max = Todo
if (Max == Min) return 3
int last_min_idx = -1, last_max_idx = -1;
ans = Integer.Max_Value;
for (i=0; i<n; i++) {
    if (arr[i] == min) {
        if (last_Max_Idx != -1) {
            ans = Min(ans, i-last_Max_Idx+1);
        }
        last_min_idx = i
    }
    else if (arr[i] == max) {
        if (last_Min_Idx != -1) {
            ans = Min(ans, i-last_Min_Idx+1);
        }
        last_Max_idx = i;
    }
}
```

T.C → O(n)
S.C → O(1).

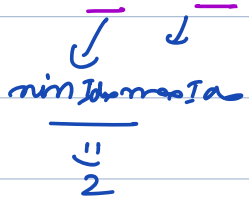
Return as ;

Edge Case :-

arr → 8, 8, 8 len = 1

max = 8
min = 8

| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 |
|---|---|---|---|---|---|---|---|---|---|---|
| 1 | 6 | 6 | 4 | 5 | 1 | 5 | 2 | 6 | 4 | 1 |



min Difference

= 2

| 10 | 20 | 30 | 40 | ... |
|---|---|---|---|---|