



CMPT 225

Lecture 5 – Comparing **List** **implementations** - Tables

List ADT class invariant?

- Element duplication allowed?
- For sake of simplicity, we shall allow duplicated elements to be inserted in our List

Comparing various implementations of the position-oriented List ADT class using Big O notation

► Time efficiency of their operations (worst case scenario):

Operations	array-based (heap)	link-based (SHSL list)
getElementCount	$O(1)$	$O(1)$
insert	$O(n)$ needs to shift right to make space for newElement	$O(n)$ insert at the end (can be $O(1)$ if use DHSL list)
remove	$O(n)$ needs to shift left to overwrite (remove) the element to be removed	$O(n)$ remove last element
removeAll/clear	$O(1)$ if elements have not been stored on heap (e.g. <code>int</code>) – otherwise each element needs to be deleted -> $O(n)$	$O(n)$ delete each Node object on heap and perhaps each element object on the heap
retrieve/get	$O(1)$	$O(n)$ must traverse to reach last element

List ADT class invariant?

- Sorted – yes!
- Element duplication allowed?
- For sake of simplicity, we shall allow duplicated elements to be inserted in our List

Comparing various implementations of the value-oriented List ADT class using Big O notation

➤ Time efficiency of their operations (worst case scenario):

Operations	array-based (heap)	link-based (SHSL list)
getElementCount	$O(1)$	$O(1)$
insert	$O(n)$ needs to shift right to make space for newElement and keep elements sorted	$O(n)^*$
remove	$O(n)$ needs to shift left to overwrite (remove) the element to be removed and keep elements sorted	$O(n)^*$
removeAll/clear	$O(1)$ if elements have not been stored on heap (e.g. <code>int</code>) – otherwise each element needs to be deleted -> $O(n)$	$O(n)$
retrieve/get	$O(\log_2 n)$	$O(n)^*$



- There are several **binary search algorithms** (and programs) that successfully search a sorted linked list (SHSL list) for a target element (search key).
- They have different time efficiencies:
 - The **binary search algorithm** that performs in **$O(\log_2 n)$** has a space efficiency of **$O(n)$** , which is not very efficient. To achieve a time efficiency of **$O(\log_2 n)$** , it uses an array in which the memory address of each of the nodes are stored (hence **$O(n)$**).
 - **Bottom Line:** Since this implementation of Binary Search requires an array, perhaps, using an array, as opposed to a linked list, would be better: same time efficiency, but better space efficiency and easier implementation.
 - The **binary search algorithm** that performs in **$O(n \log_2 n)$** uses the slow/fast pointer algorithm to find the middle node/element. The first middle element is found in $n/2$ steps (**$O(n)$**), then $n/4$, then $n/8$, etc... There are **$O(\log_2 n)$** steps, hence **$O(n \log_2 n)$**
 - **Bottom Line:** Using linear search would be better: better time efficiency and easier implementation.