

## 169. Majority Element

Easy

👍 9249

💬 331

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Given an array `nums` of size `n`, return *the majority element*.

The majority element is the element that appears more than  $\lfloor n / 2 \rfloor$  times.

You may assume that the majority element always exists in the array.

### Example 1:

**Input:** `nums = [3,2,3]`

**Output:** 3

### Example 2:

**Input:** `nums = [2,2,1,1,1,2,2]`

**Output:** 2

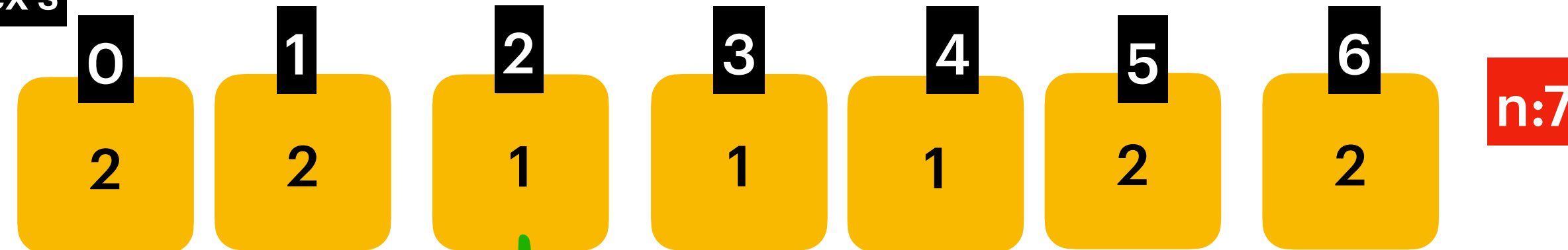
### Constraints:

- `n == nums.length`
- `1 <= n <= 5 * 104`
- `-109 <= nums[i] <= 109`

## 169 Majority Element

Key is Majority Element repeats  $> n/2$  times  
so that when we sort the array , we can always find Majority element in  $n/2$  index.

Index's

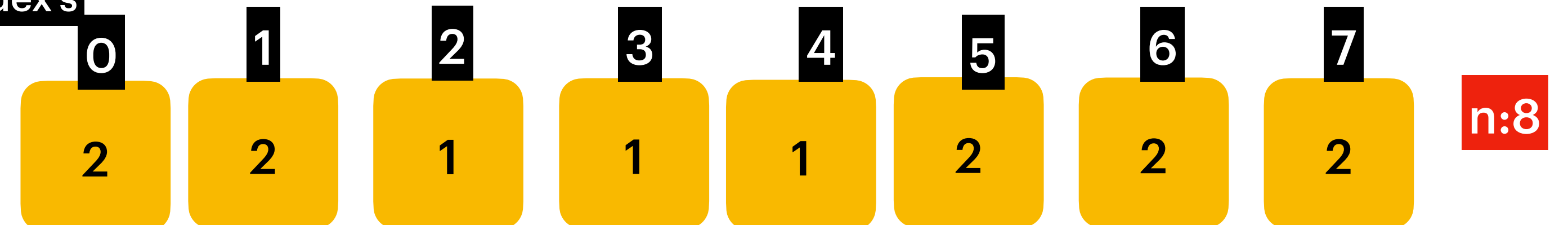


After Sort



Majority Element is on  $n/2$  index :  
`nums[3] = 2`

Index's



Time Complexity :  $O(n \log n)$

Merge Sort  $O(n)$

Sorting

Space Complexity :  $O(1)$

Quick Sort  $O(\log n)$



After Sort



Majority Element is on  $n/2$  index :  
`nums[4] = 2`

Can we improve the Time&Space Complexity ?

Yes By applying Math.

Lets apply voting:

As we know that within a given input the majority element presents  $> n/2$  times.

When we apply voting alway we left with the Majority Element.

When the element repeats increment the vote, other time decrement the vote.

When vote becomes zero, take the current integration element Takes the vote.

Time Complexity :  $O(n)$

Space Complexity :  $O(1)$

## Let's Apply Math

Initialise the current element and increment the vote

Element = nums[0] = 2  
Vote = 1

Current element is repeated increment the vote

Element = 2  
Vote = 2

Current element is not repeated  
Decrement the vote

Element = 2  
Vote = 1

Again Current element is not repeated  
Decrement the vote

Element = 2  
Vote = 0

As the Vote is zero in the previous step reinitialise the Element with current iterated value and increment the vote

Element = 1  
Vote = 1

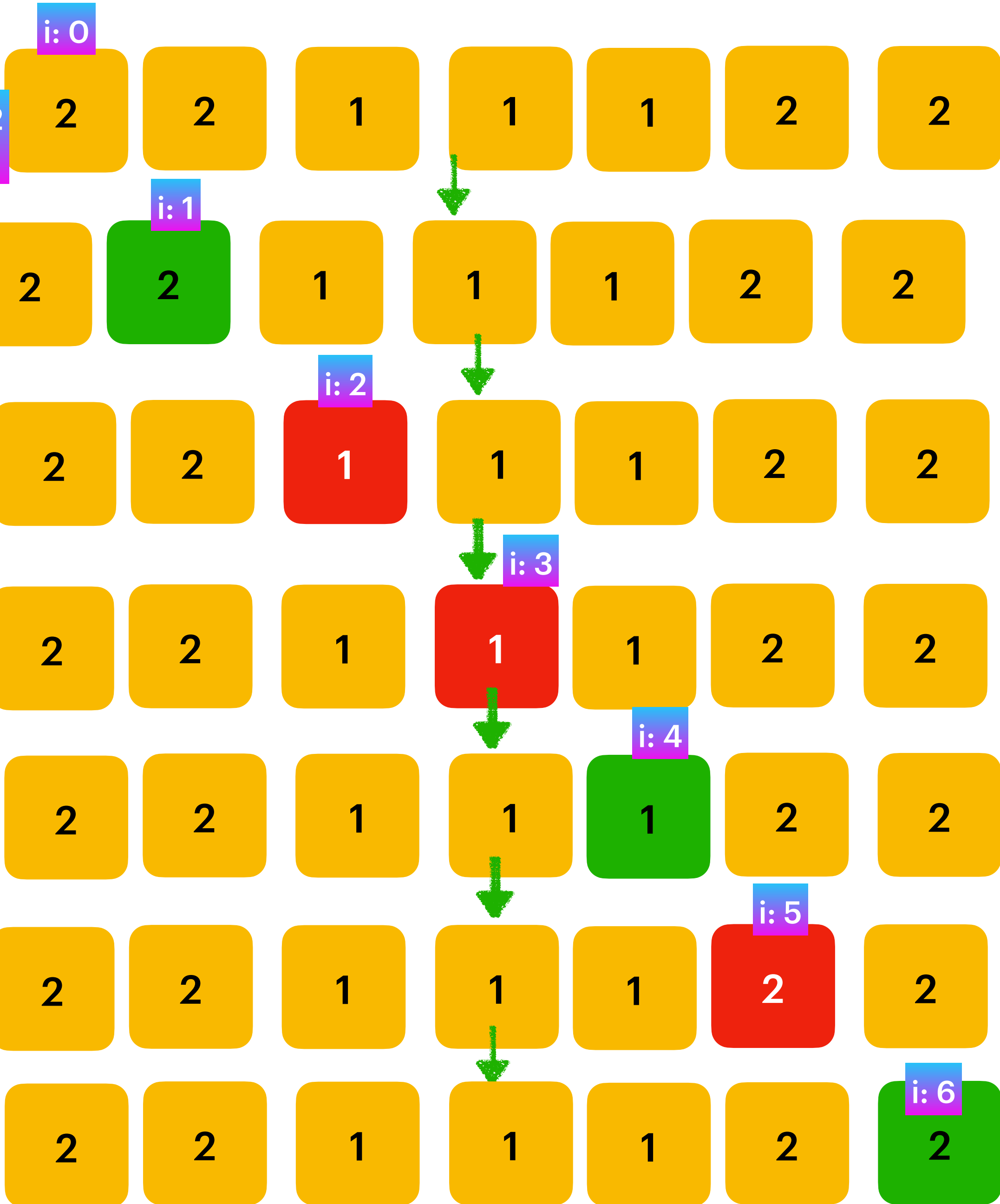
Current element is not repeated  
Decrement the vote

Element = 1  
Vote = 0

As the Vote is zero in the previous step reinitialise the Element with current iterated value and increment the vote

Element = 2  
Vote = 1

Return the Element i.e 2



## 4. Median of Two Sorted Arrays

Hard

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Given two sorted arrays `nums1` and `nums2` of size `m` and `n` respectively, return **the median** of the two sorted arrays.

The overall run time complexity should be  $O(\log(m+n))$ .

### Constraints:

- `nums1.length == m`
- `nums2.length == n`
- `0 <= m <= 1000`
- `0 <= n <= 1000`
- `1 <= m + n <= 2000`
- `-106 <= nums1[i], nums2[i] <= 106`

### Example 1:

**Input:** `nums1 = [1,3], nums2 = [2]`

**Output:** `2.00000`

**Explanation:** merged array = `[1,2,3]` and median is 2.

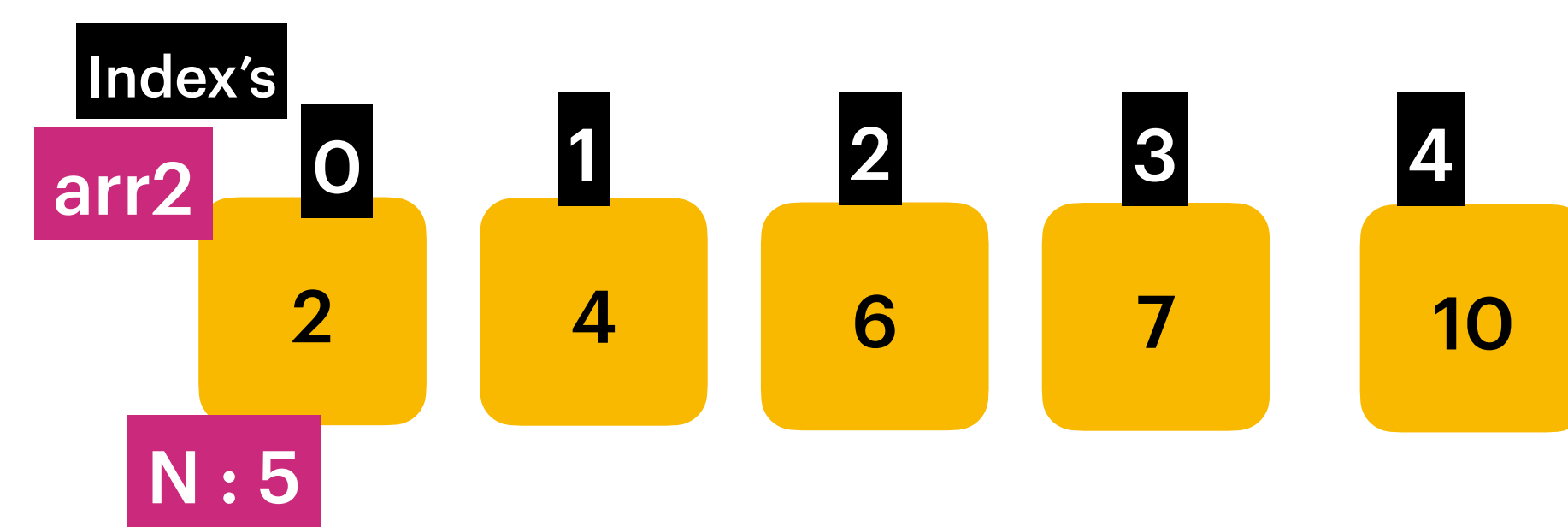
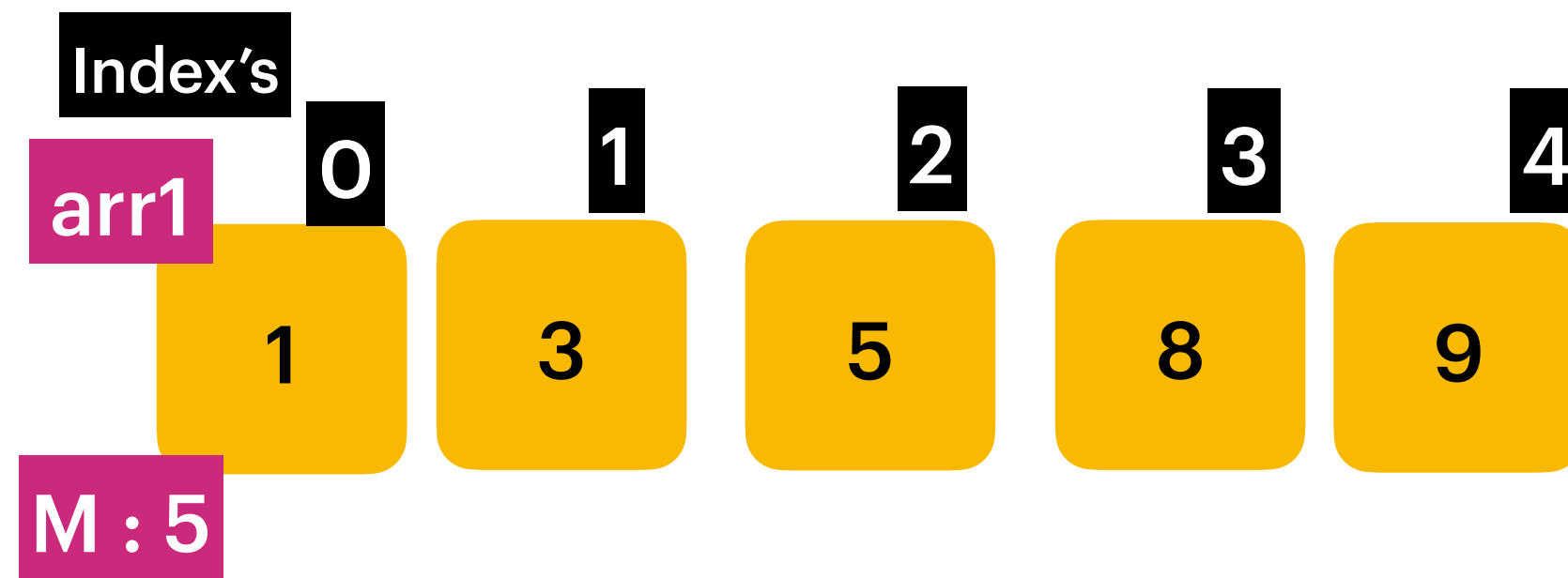
### Example 2:

**Input:** `nums1 = [1,2], nums2 = [3,4]`

**Output:** `2.50000`

**Explanation:** merged array = `[1,2,3,4]` and median is  $(2 + 3) / 2 = 2.5$ .

## 4 Median Of Two Sorted Arrays

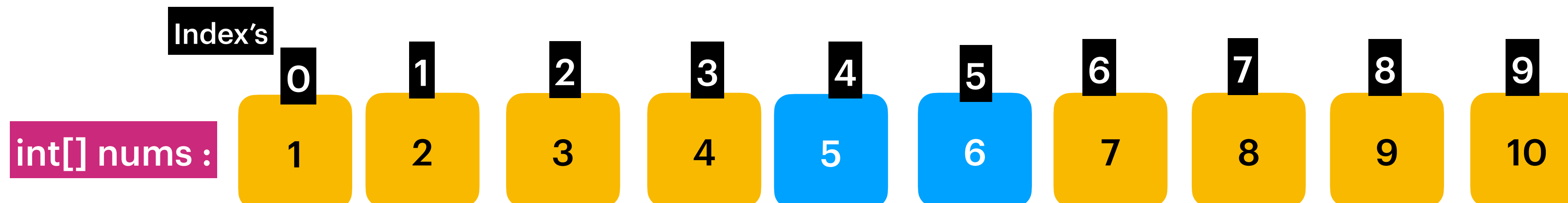


### Brute Force

Take an Array Of size M+N ,  
combine and merge both the arrays.

If the length is even then  $(\text{nums}[\text{mid}] + \text{nums}[\text{mid}-1] / 2)$   
is the Median.

If the length is odd then  $\text{nums}[\text{mid}]$  is the Median.



As the length is even then  $(\text{nums}[\text{mid}] + \text{nums}[\text{mid}-1] / 2)$   
is the Median.

$$= (5 + 6) / 2 = 11$$

Time Complexity :  $O(M+N)$

Space Complexity :  $O(M+N)$

Can we improve the Time&Space Complexity ?

Yes By applying Math.