

Space

In-Place

Out-Place

In-Place Algo:
Space Complexity : $O(1)$

```
int[] arr = {1,2,3,4};
```

square :

```
for(int i = 0 ; i < arr.length ; i++)  
    {  
        arr[i] = arr[i] * arr[i];  
    }
```

```
int[] arr = {2,4,8,16};
```

Out-Place Algo:

Space Complexity : $O(n)$

```
int[] arr = {1,2,3,4};  
int[] square = new int[arr.length];
```

```
for(int i = 0 ; i < arr.length ; i++)  
    {  
        square[i] = arr[i] * arr[i];  
    }
```

287. Find the Duplicate Number

Medium

👍 12229

💬 1366

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Given an array of integers `nums` containing `n + 1` integers where each integer is in the range `[1, n]` inclusive.

There is only **one repeated number** in `nums`, return *this repeated number*.

You must solve the problem **without** modifying the array `nums` and uses only constant extra space.

Example 1:

Input: `nums = [1,3,4,2,2]`

Output: `2`

Example 2:

Input: `nums = [3,1,3,4,2]`

Output: `3`

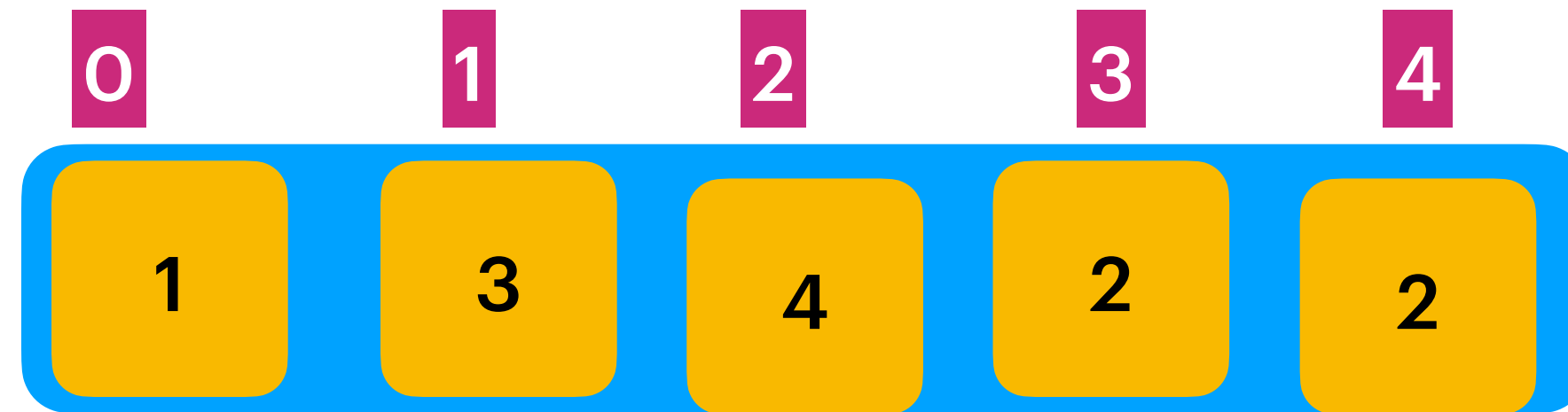
Constraints:

- `1 <= n <= 105`
- `nums.length == n + 1`
- `1 <= nums[i] <= n`
- All the integers in `nums` appear only **once** except for **precisely one integer** which appears **two or more** times.

Follow up:

- How can we prove that at least one duplicate number must exist in `nums` ?
- Can you solve the problem in linear runtime complexity?

OutPlace Algo



Time Complexity : $O(n)$

Excepted Output : 2

OutPlace Algo

Space Complexity : $O(n)$

-> Take a boolean[] array with size n. Has the default values are False.

-> For each iteration of input array, update the respective boolean[] array index value with True .

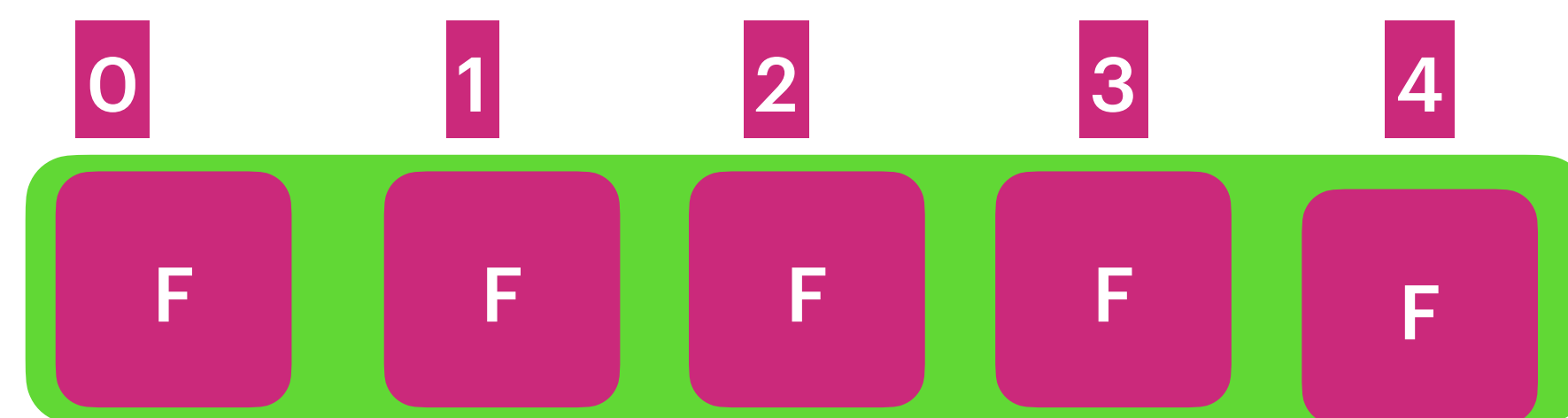
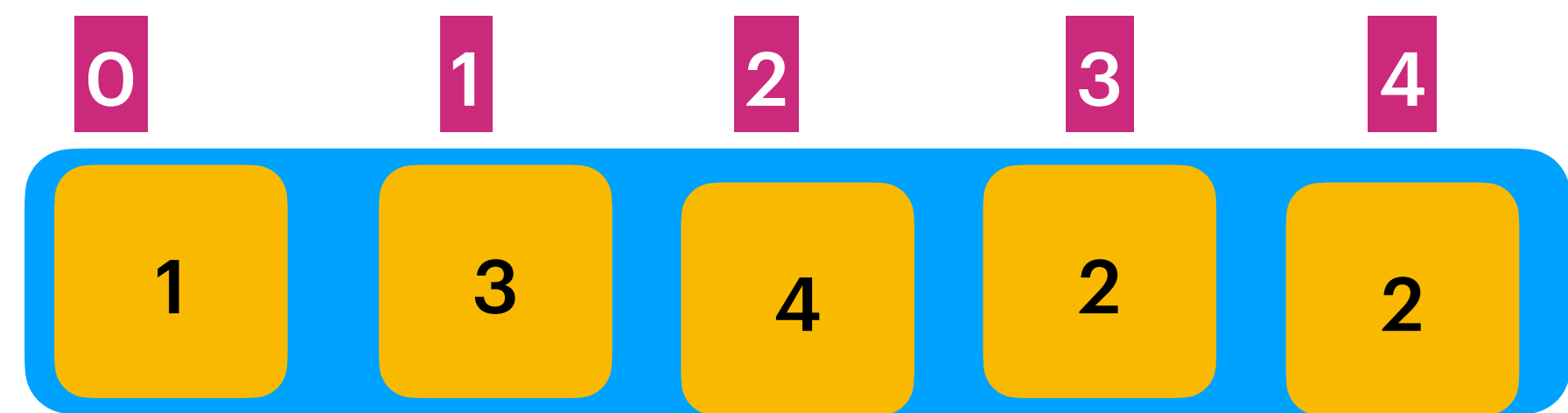
-> If the value is repeated for the second visit you already find true in boolean[] array so return the value.

Clue

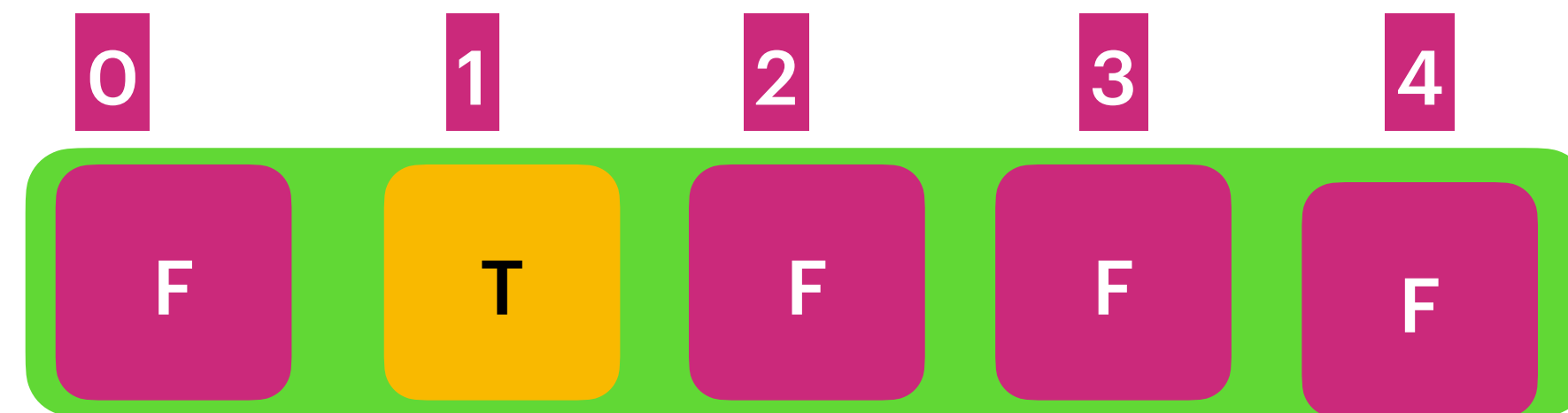
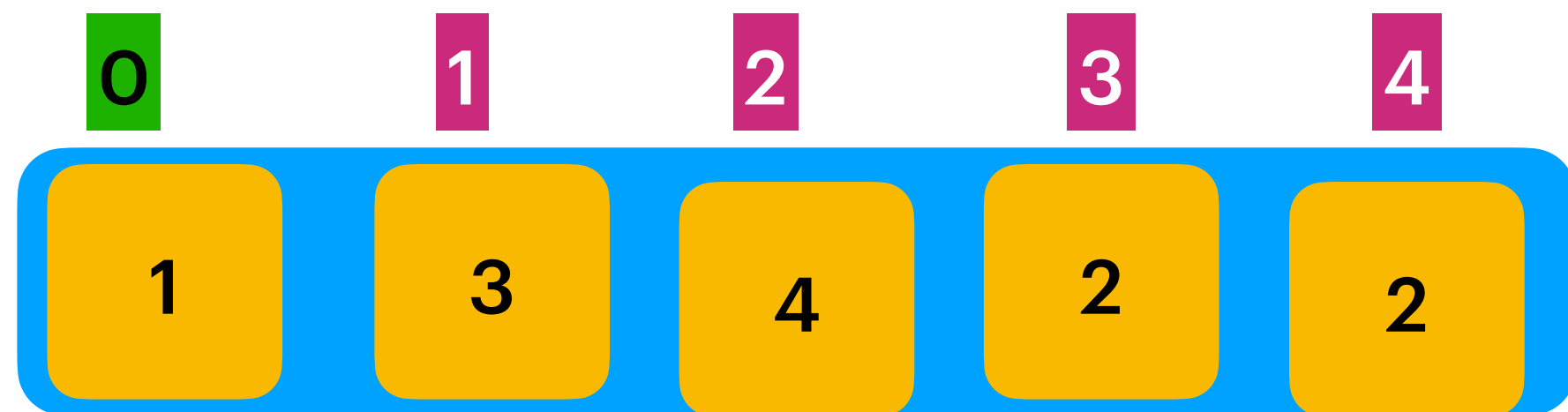
Hit is values are in the range of $[1, n]$ and the length of the array is : $n+1$

So that each value can be uniquely identified by an array index.

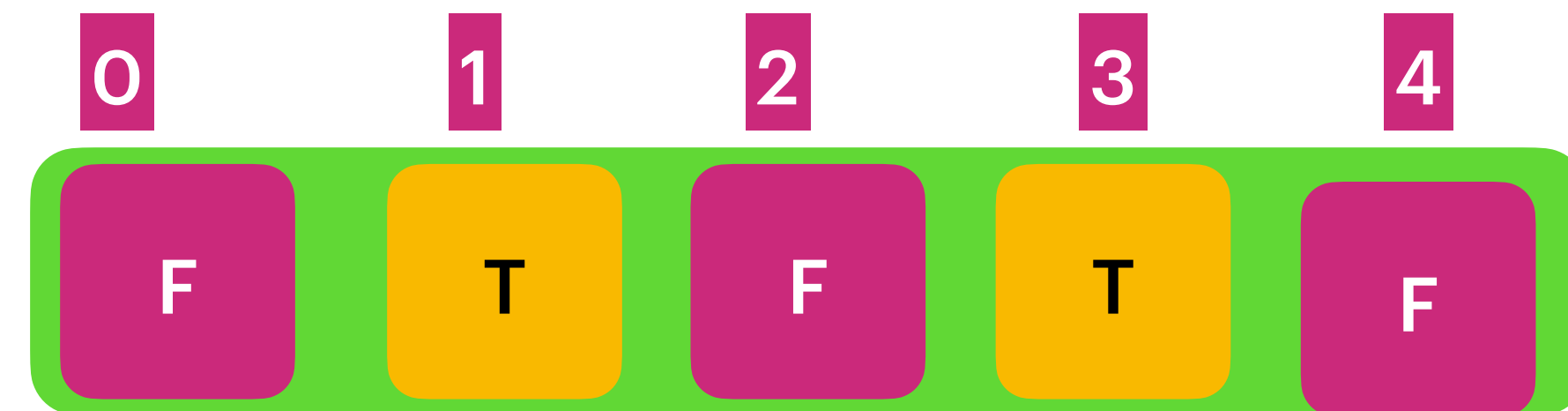
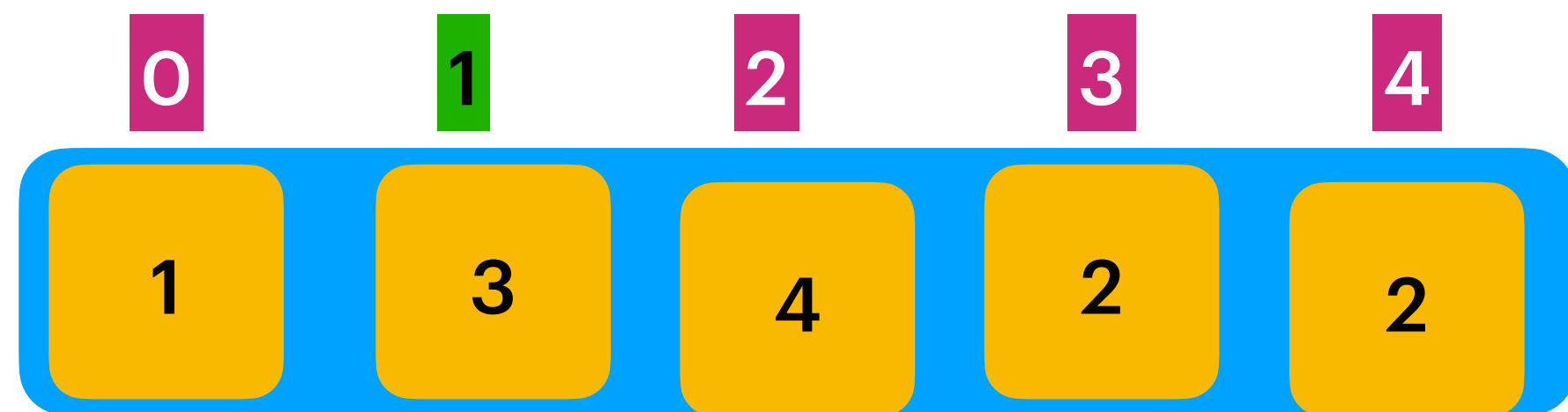
OutPlace Algo



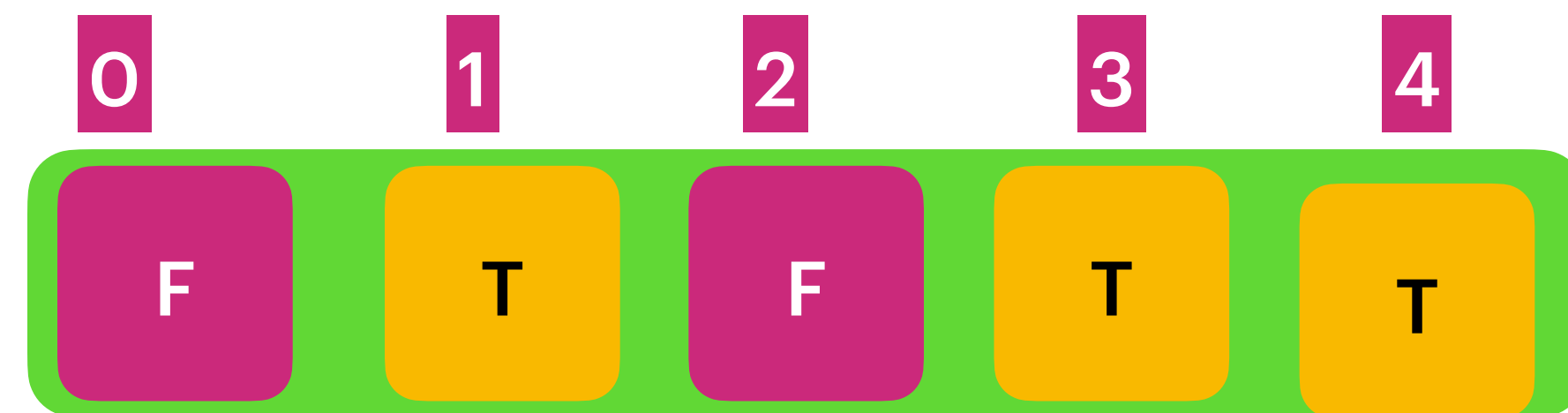
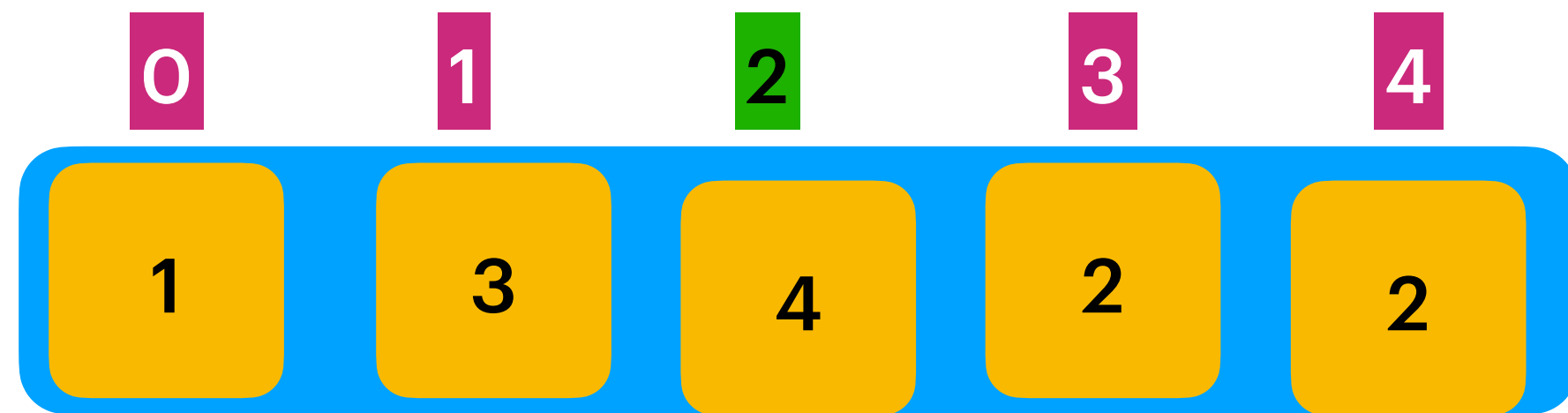
i=0

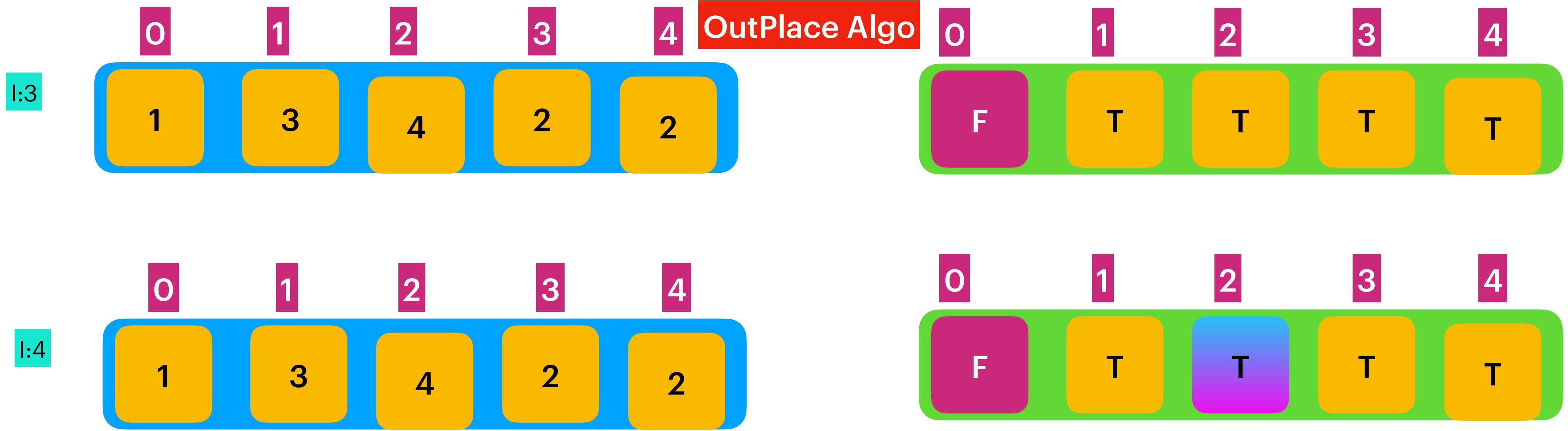


i=1

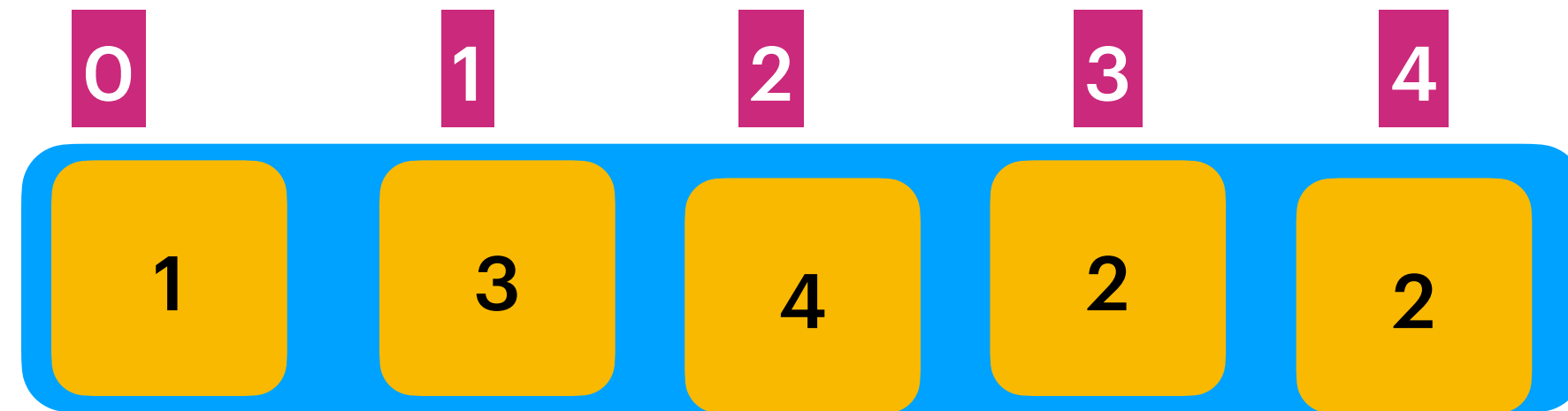


i=2





InPlace Algo



Time Complexity : $O(n)$

Excepted Output : 2

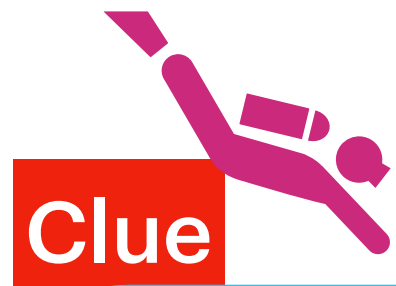
InPlace Algo

Space Complexity : $O(1)$

-> Here we replace the sign of index value mapping to current iterated value with in the same input array.

-> If the value is repeated then the sign would have already been updated so that it is the duplicate.

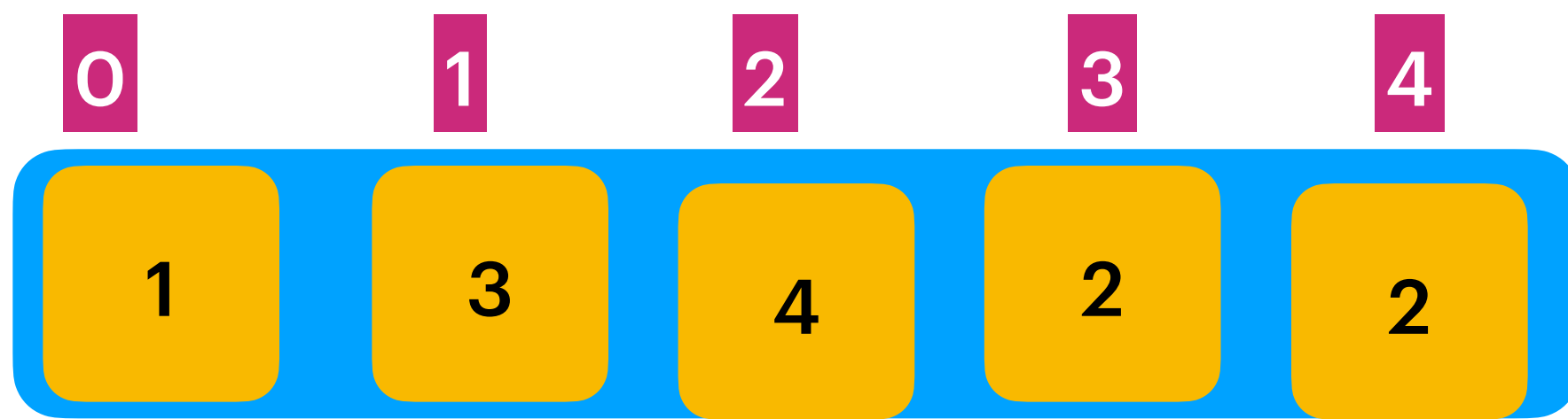
Note : As we are updating given input array, always take the absolute value for each iteration.



Clue

Hit is values are in the range of $[1,n]$ and the length of the array is : $n+1$
So that each value can be uniquely identified by an array index.

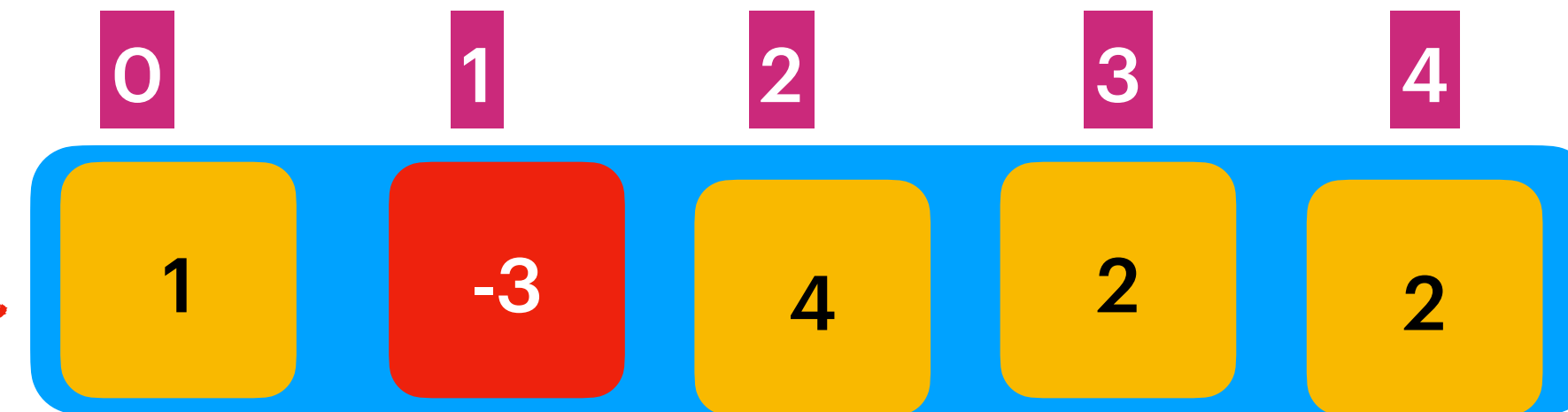
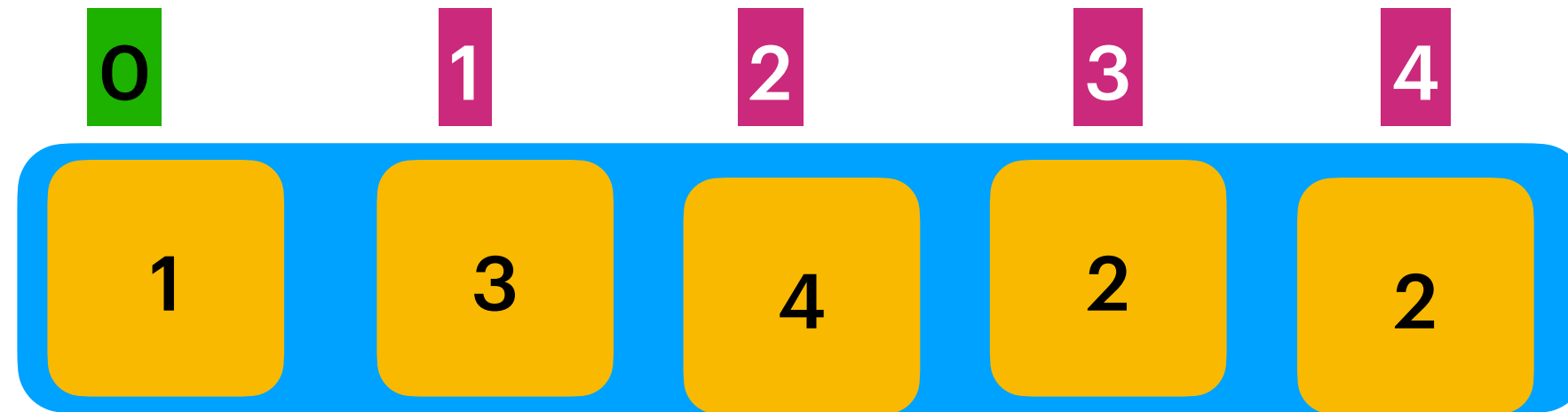
InPlace Algo



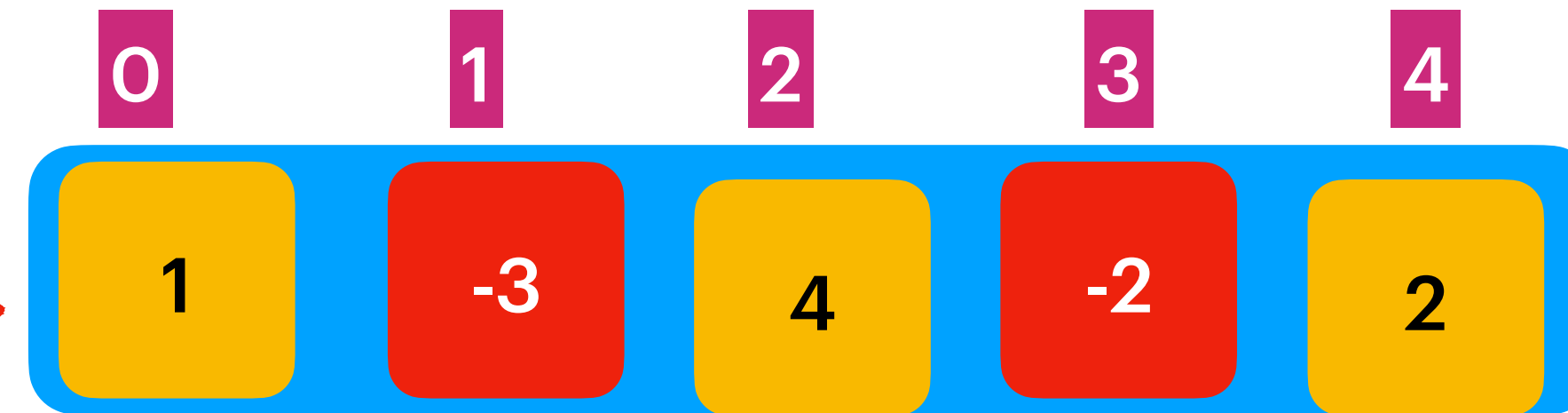
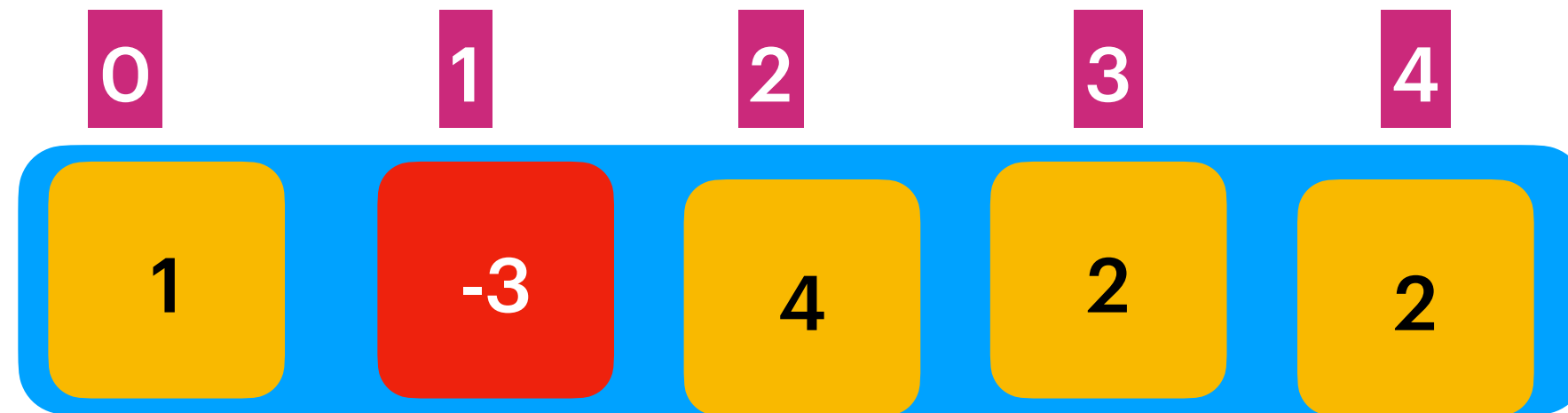
Input

Output

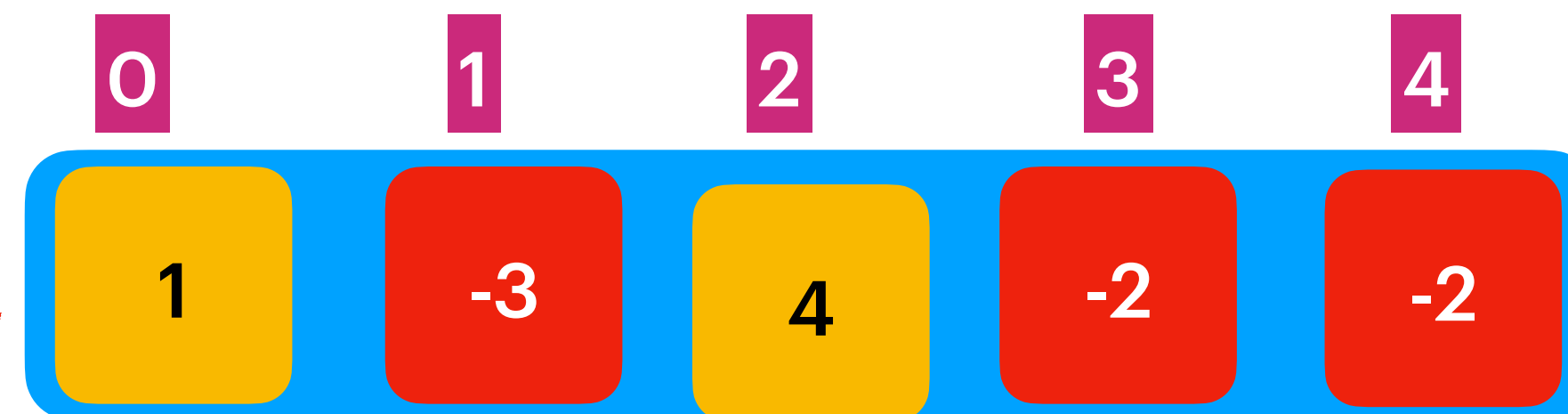
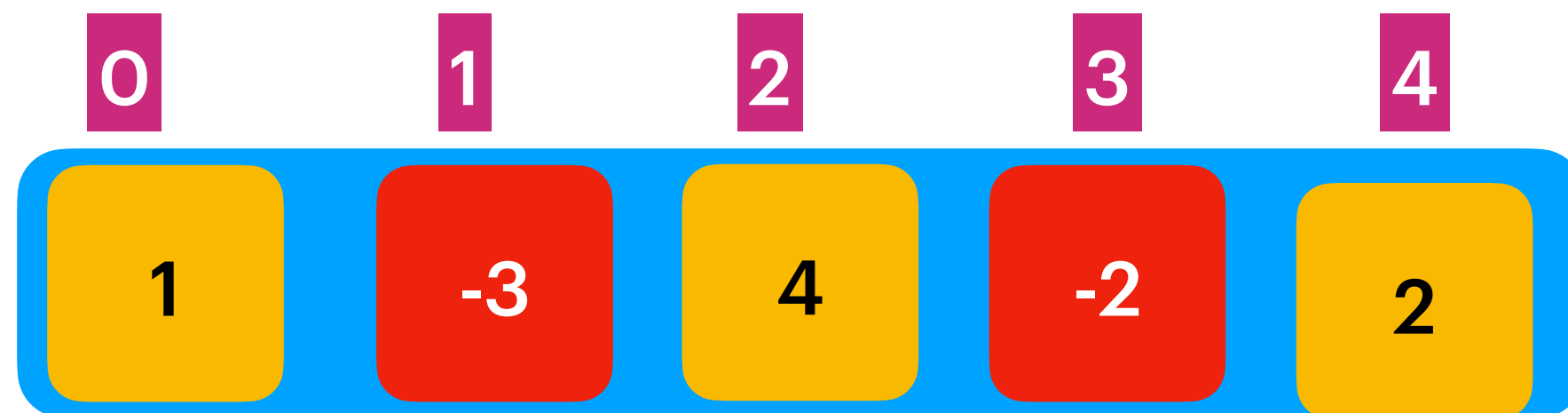
i=0



i=1

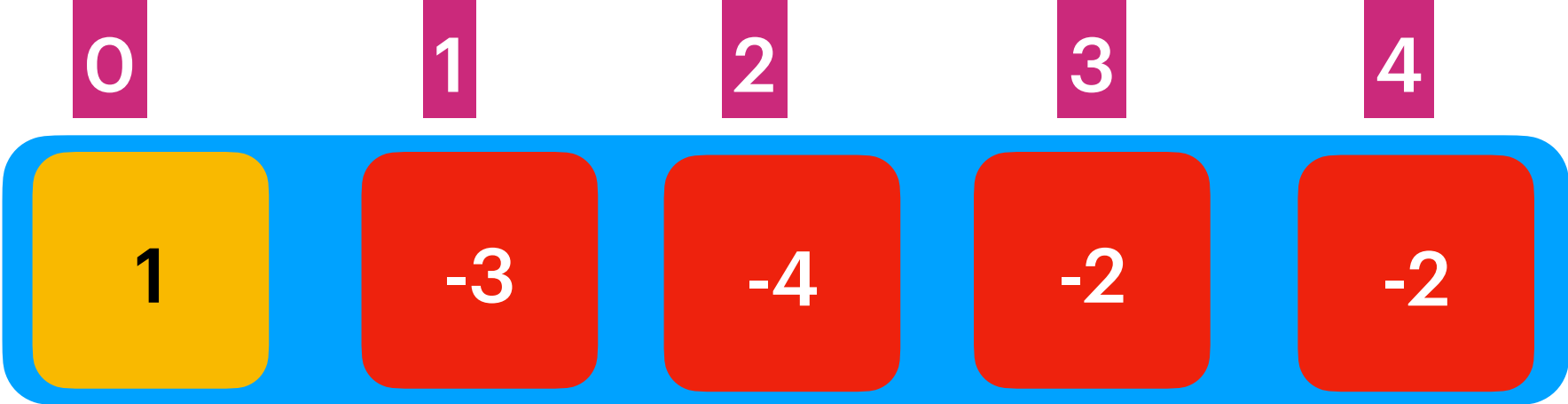
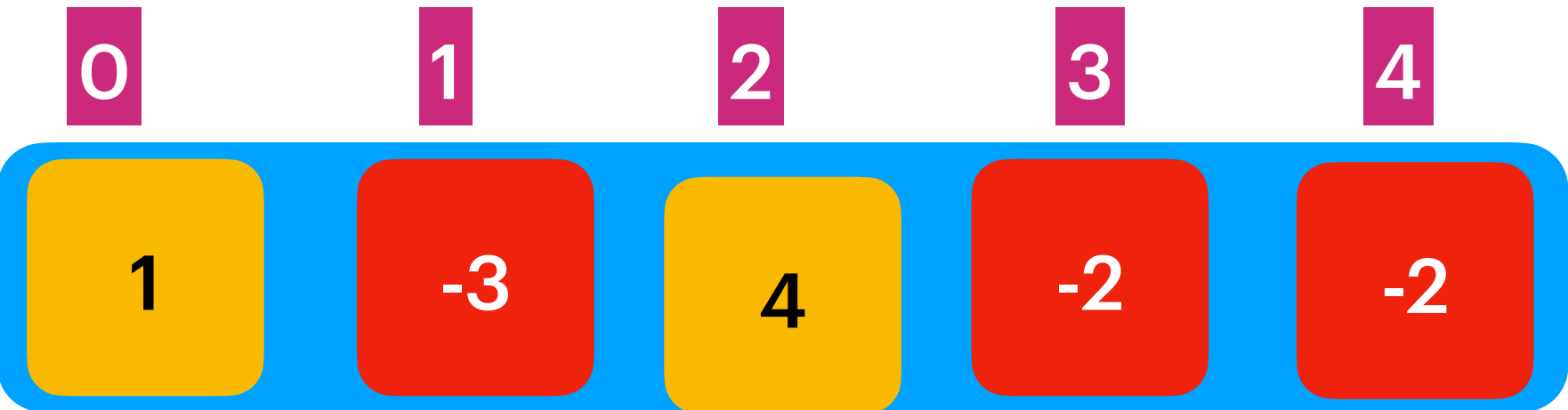


i=2

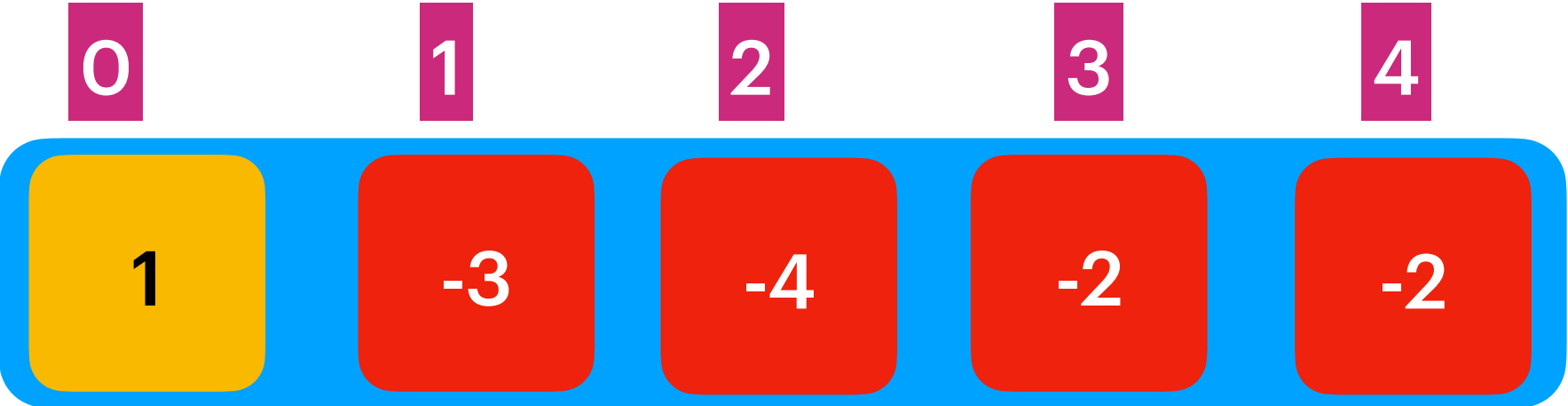
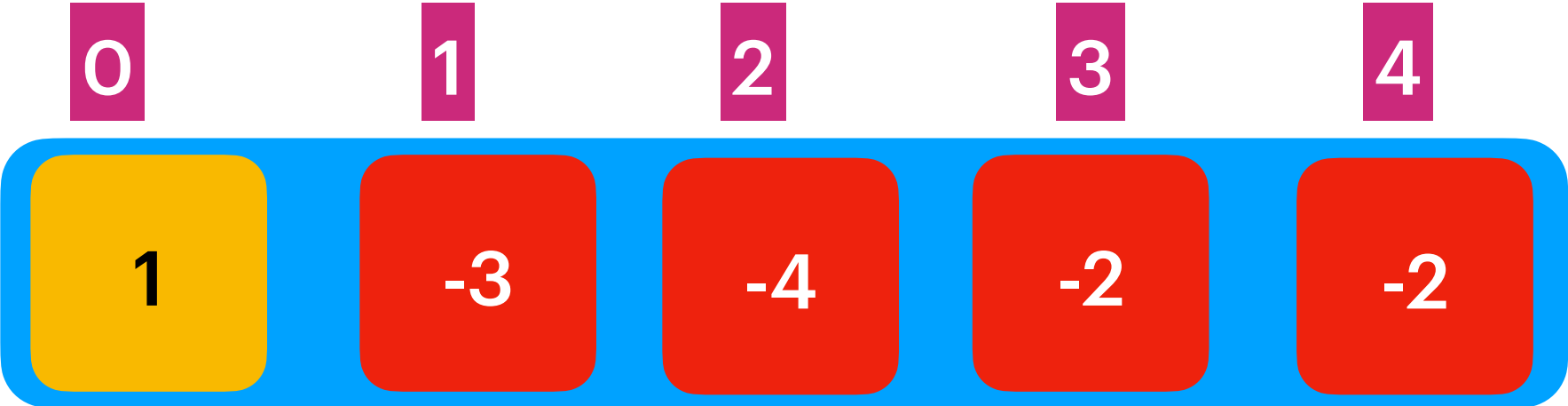


In-Place Algo

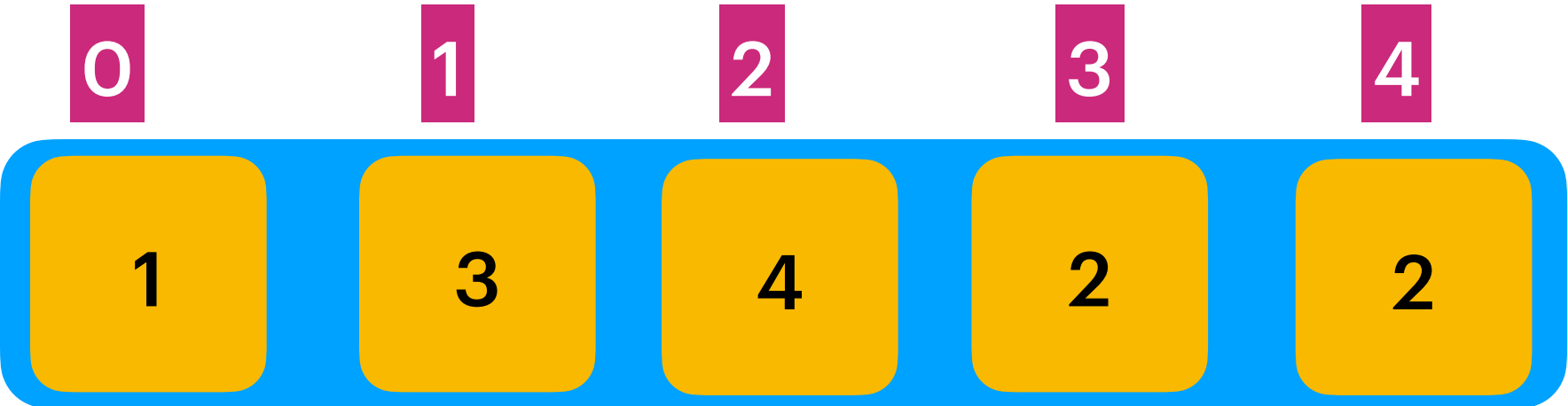
l:3



l:4



Time Complexity : $O(n)$
Excepted Output : 2
Space Complexity : $O(1)$



41. First Missing Positive

Hard

👍 8998

💬 1283

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Given an unsorted integer array `nums`, return the smallest missing positive integer.

You must implement an algorithm that runs in $O(n)$ time and uses constant extra space.

Example 1:

Input: `nums = [1,2,0]`

Output: 3

Example 2:

Input: `nums = [3,4,-1,1]`

Output: 2

Example 3:

Input: `nums = [7,8,9,11,12]`

Output: 1

Constraints:

- `1 <= nums.length <= 5 * 105`
- `-231 <= nums[i] <= 231 - 1`