

155. Min Stack

Easy  8346  635  Add to List  Share

Design a stack that supports push, pop, top, and retrieving the minimum element in constant time.

Implement the `MinStack` class:

- `MinStack()` initializes the stack object.
- `void push(int val)` pushes the element `val` onto the stack.
- `void pop()` removes the element on the top of the stack.
- `int top()` gets the top element of the stack.
- `int getMin()` retrieves the minimum element in the stack.

Example 1:

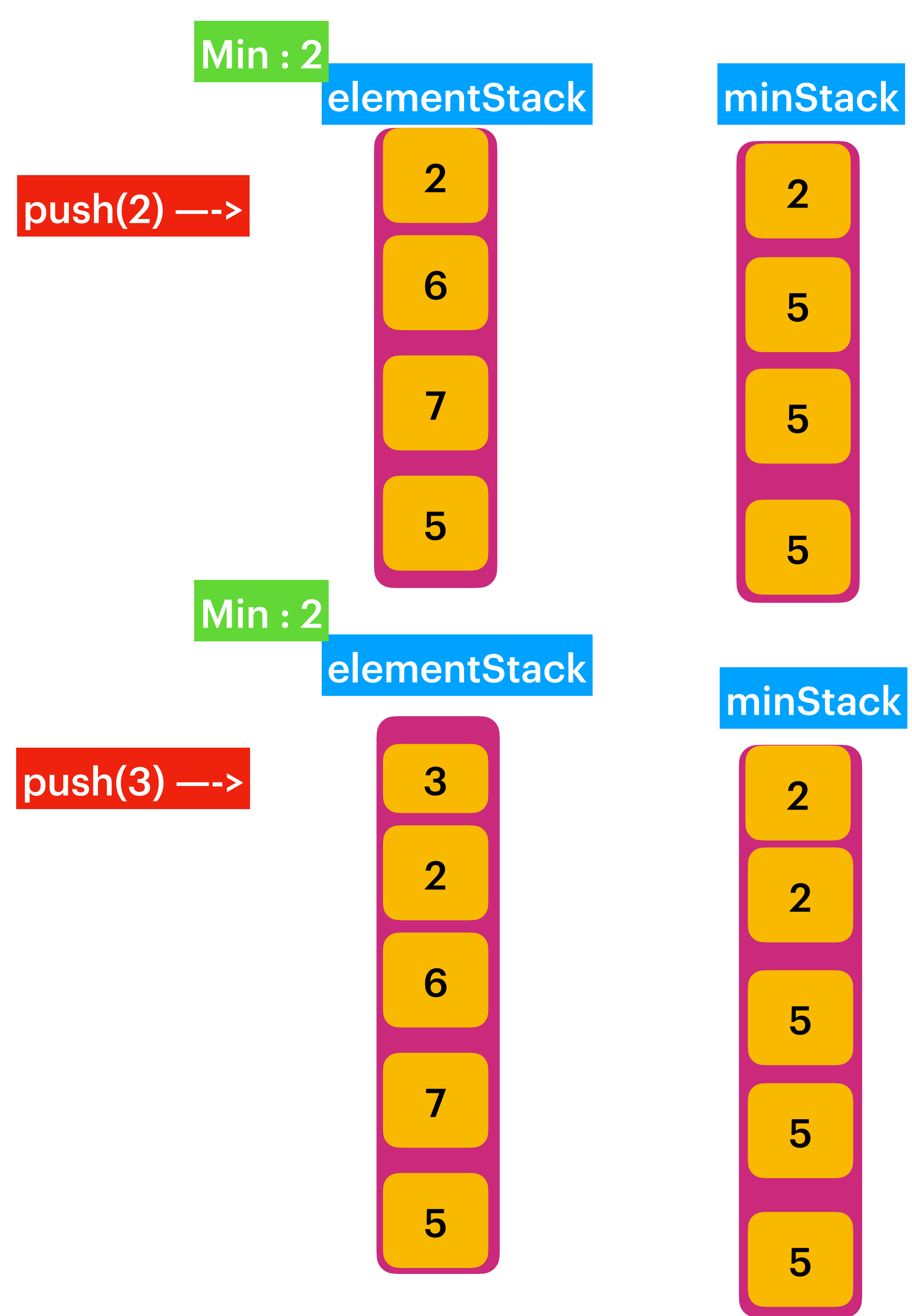
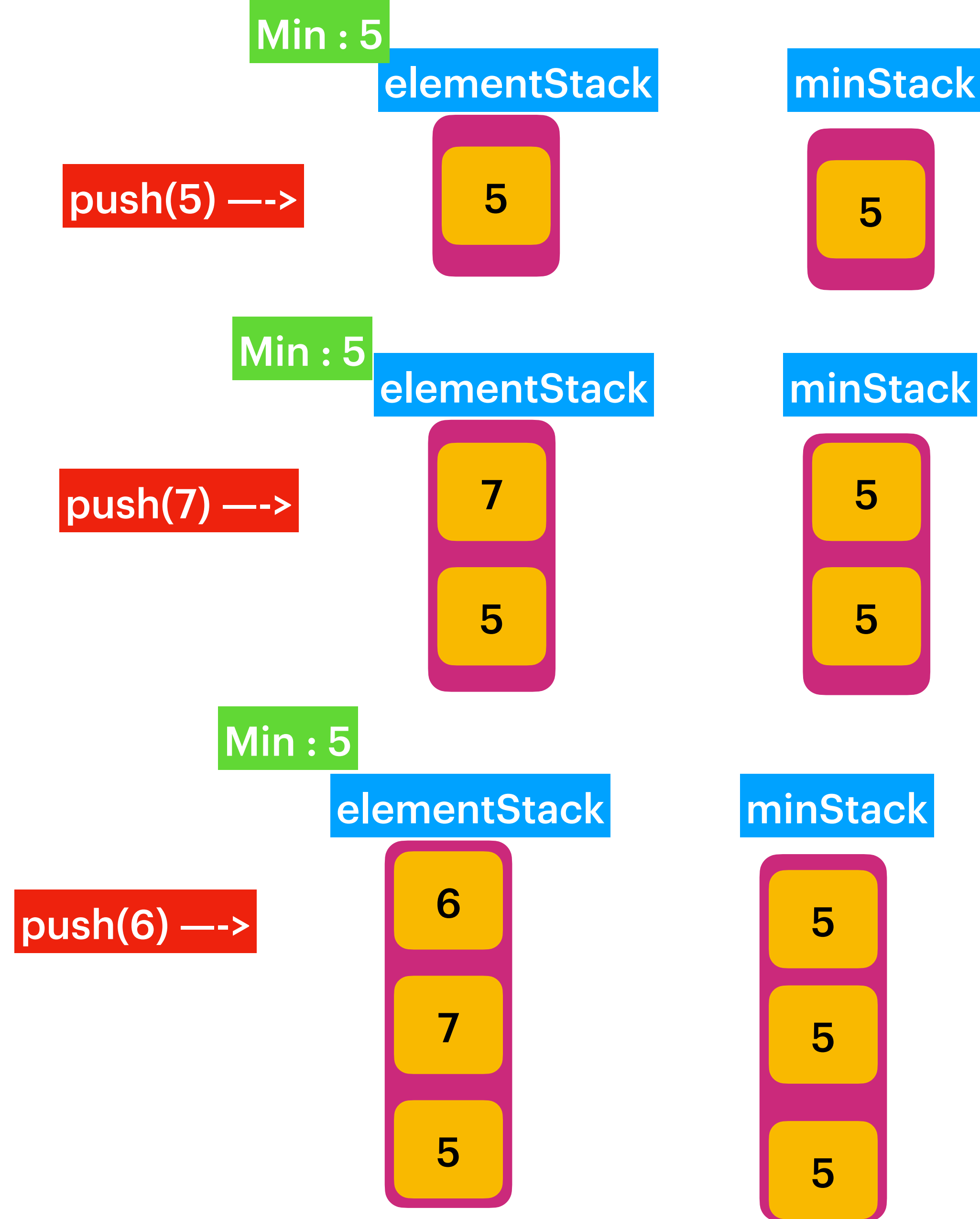
Input
["MinStack","push","push","push","getMin","pop","top","getMin"]
[[],[-2],[0],[-3],[],[],[],[]]

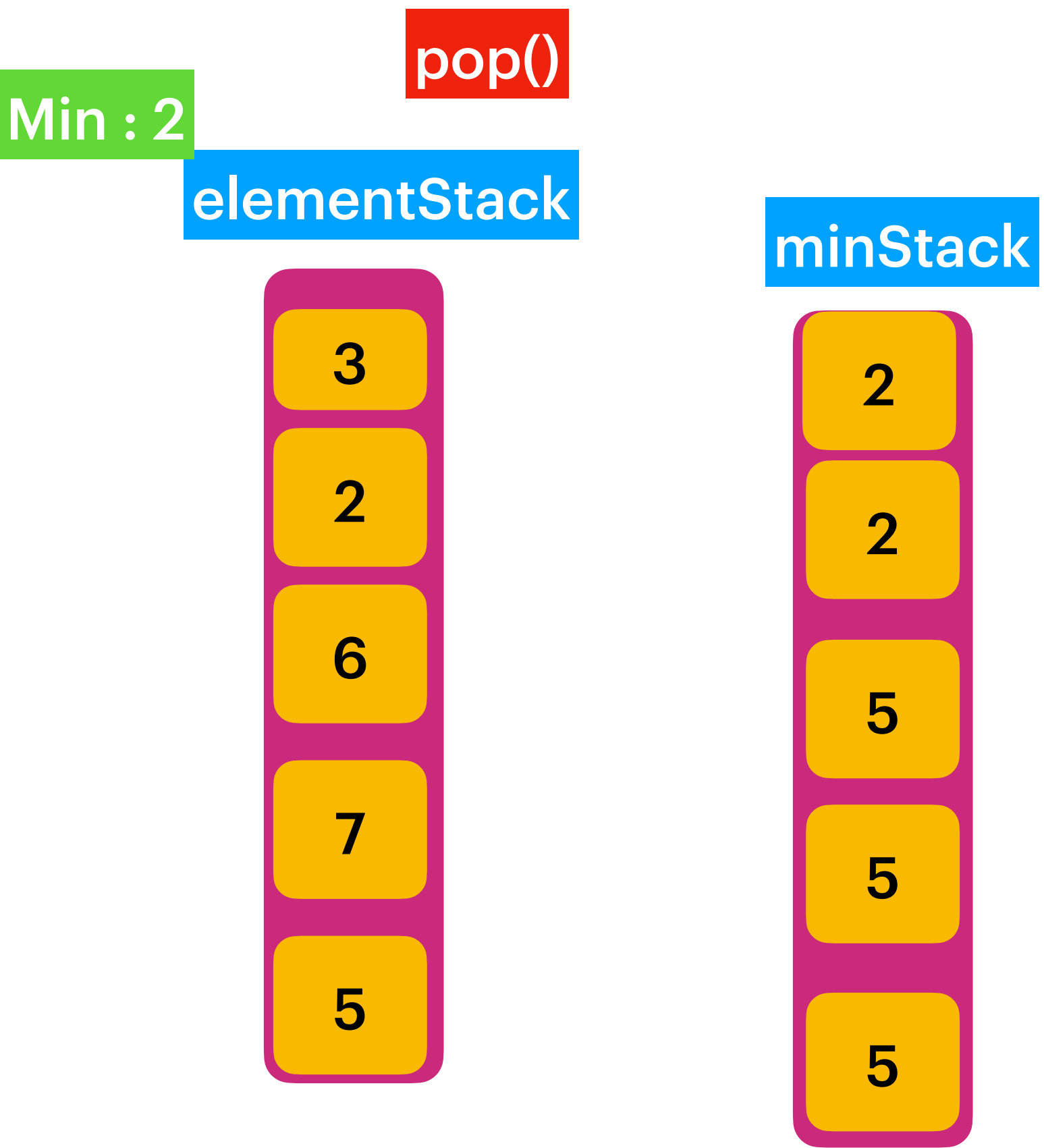
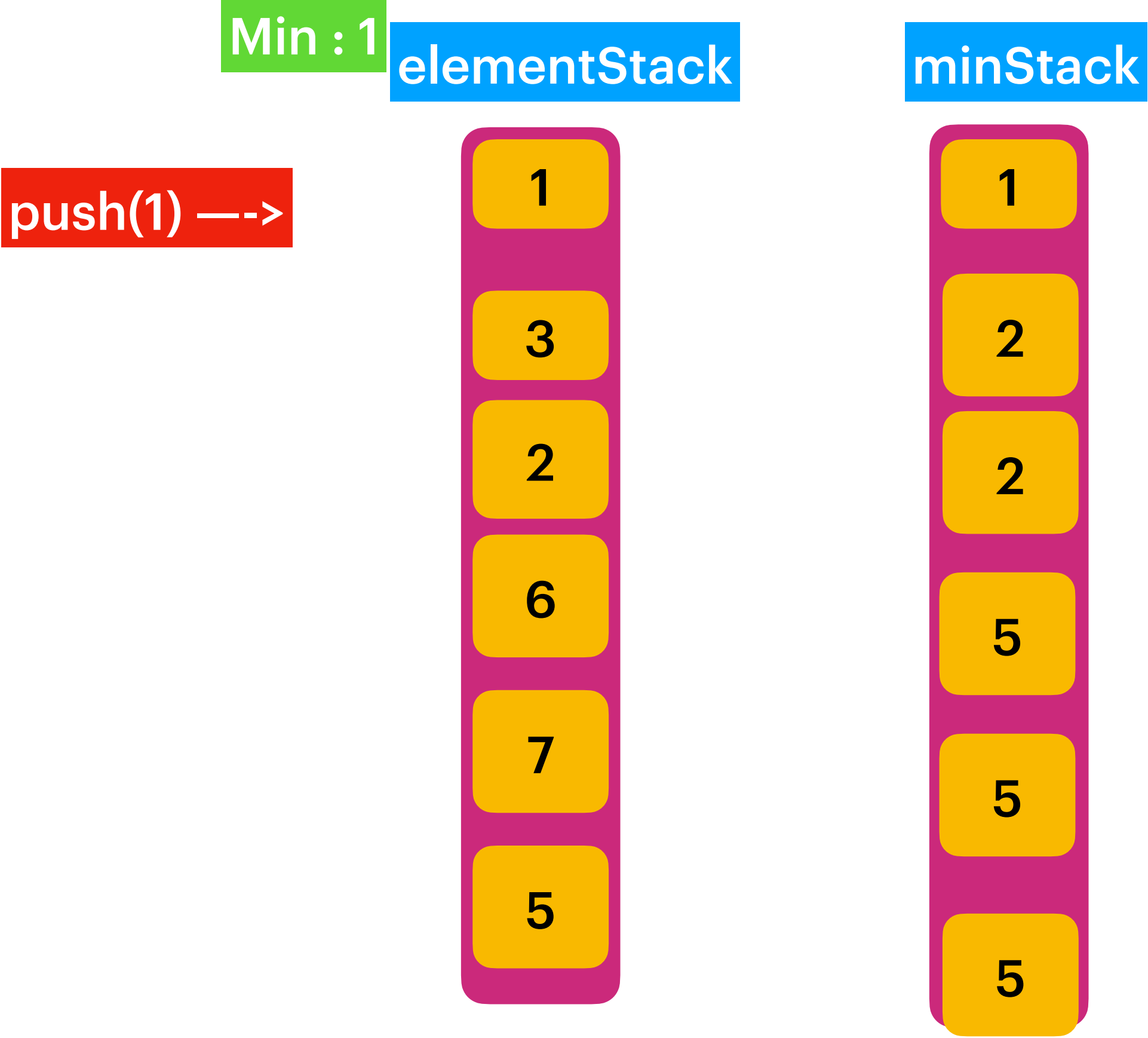
Output
[null,null,null,null,-3,null,0,-2]

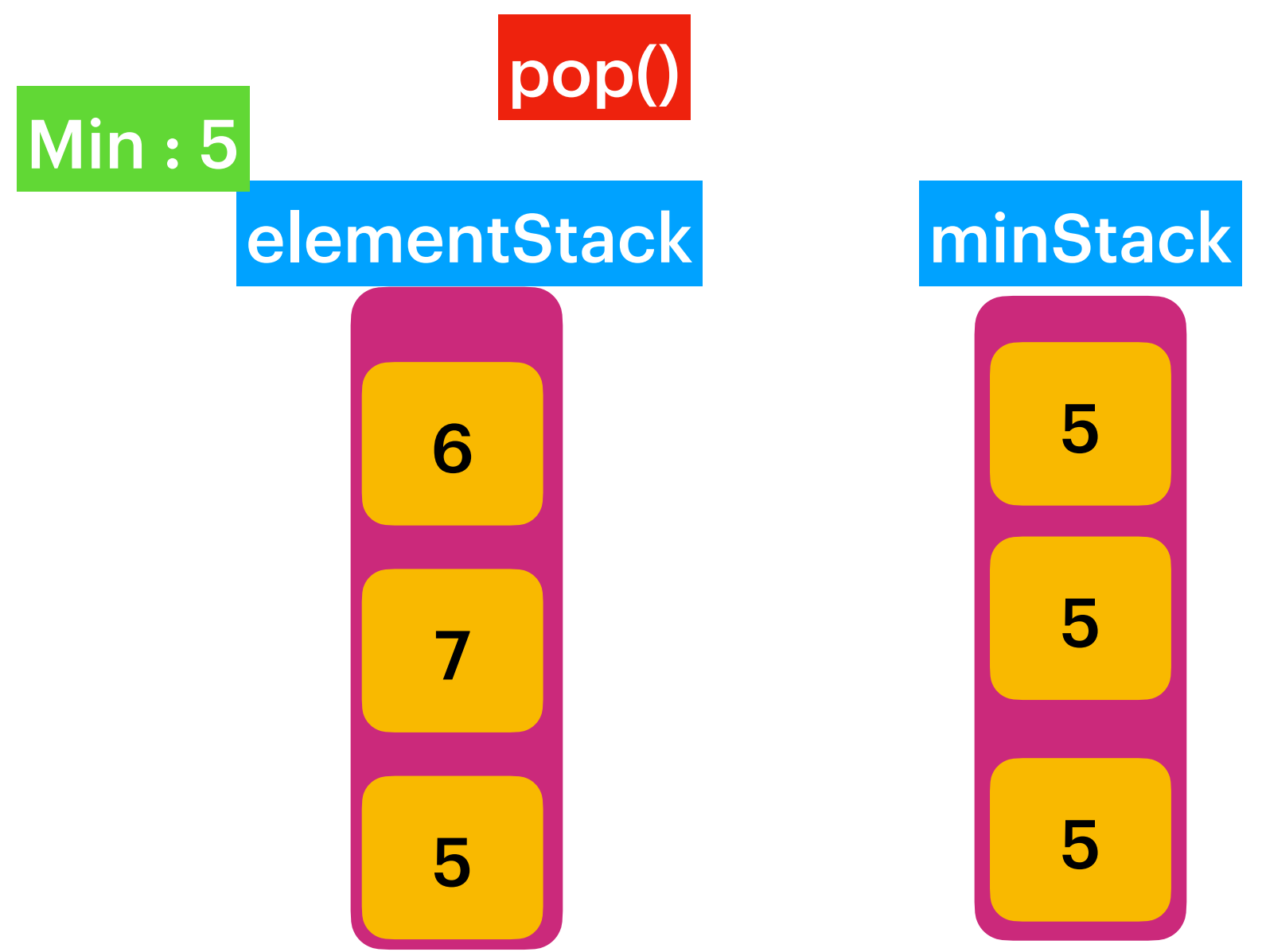
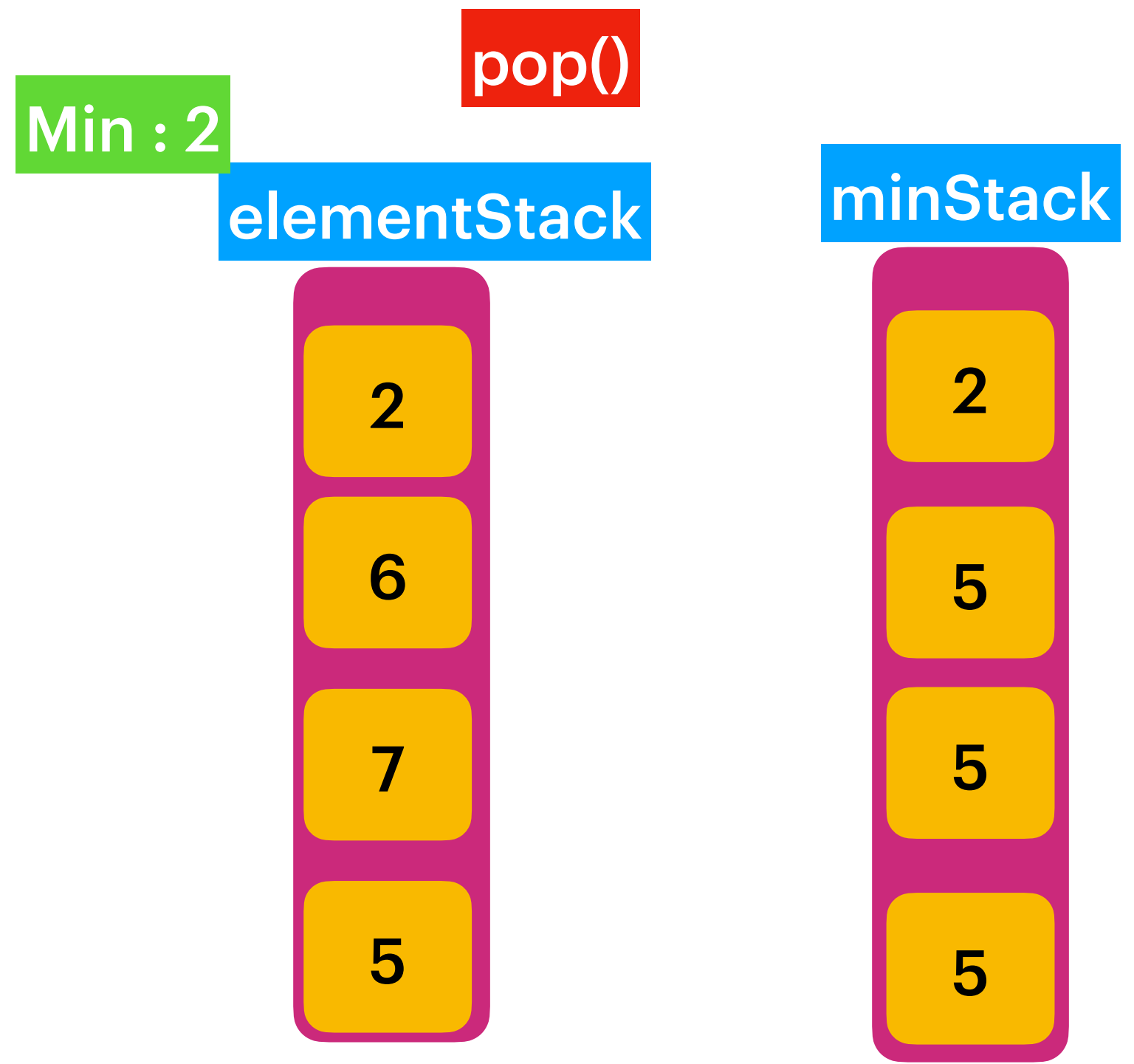
Explanation
`MinStack minStack = new MinStack();`
`minStack.push(-2);`
`minStack.push(0);`
`minStack.push(-3);`
`minStack.getMin(); // return -3`
`minStack.pop();`
`minStack.top(); // return 0`
`minStack.getMin(); // return -2`

Constraints:

- $-2^{31} \leq val \leq 2^{31} - 1$
- Methods `pop`, `top` and `getMin` operations will always be called on **non-empty** stacks.
- At most $3 * 10^4$ calls will be made to `push`, `pop`, `top`, and `getMin`.







`push(E)` → Time Complexity : $O(1)$
`pop()` → TimeComplexity : $O(1)$
`getMin()` → $O(1)$
`top()` → $O(1)$

20. Valid Parentheses

Easy  13517  611  Add to List  Share

Given a string `s` containing just the characters `'('`, `)'`, `'{'`, `'}'`, `'['` and `']'`, determine if the input string is valid.

An input string is valid if:

1. Open brackets must be closed by the same type of brackets.
2. Open brackets must be closed in the correct order.

Example 1:

Input: `s = "()"`
Output: `true`

Example 2:

Input: `s = "()[]{}"`
Output: `true`

Example 3:

Input: `s = "()["`
Output: `false`

Constraints:

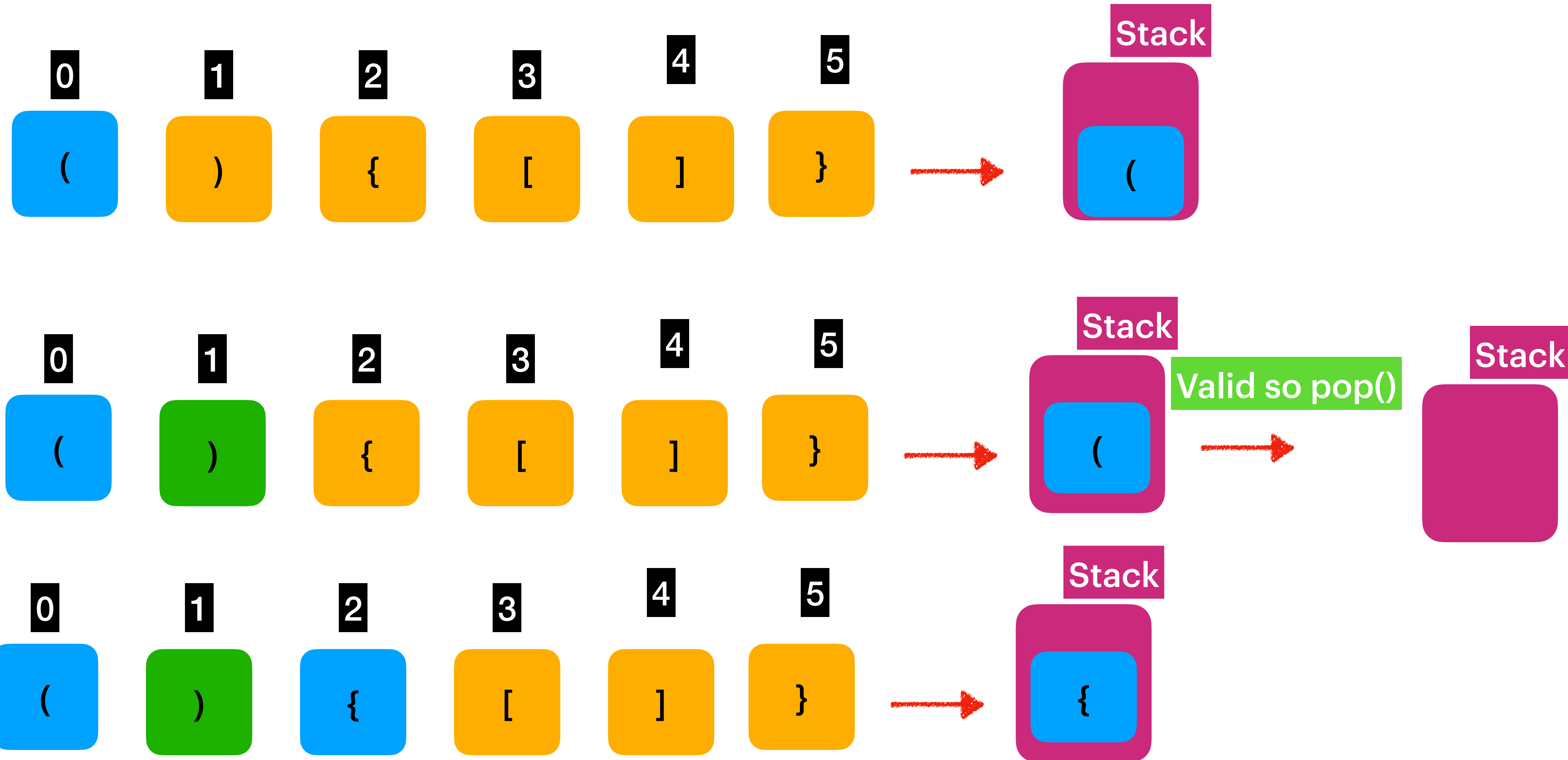
- `1 <= s.length <= 104`
- `s` consists of parentheses only `'()[]{}'`.

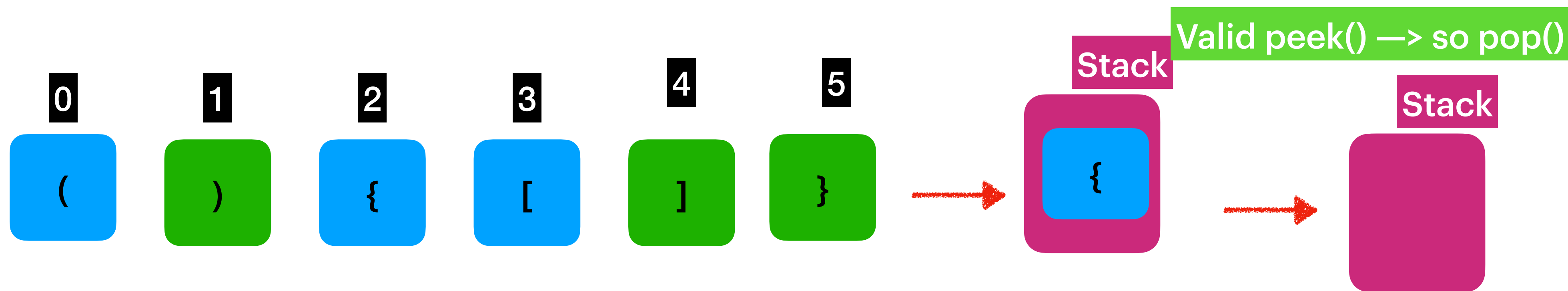
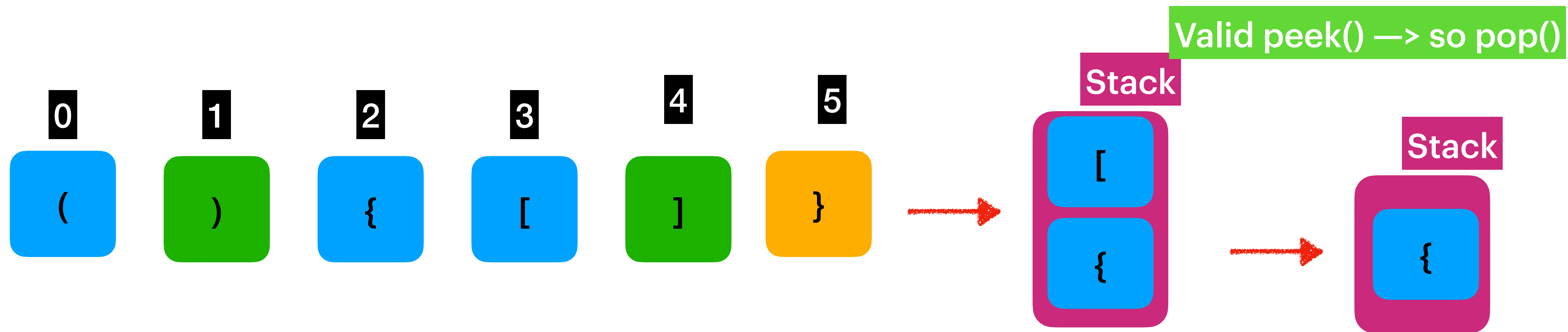
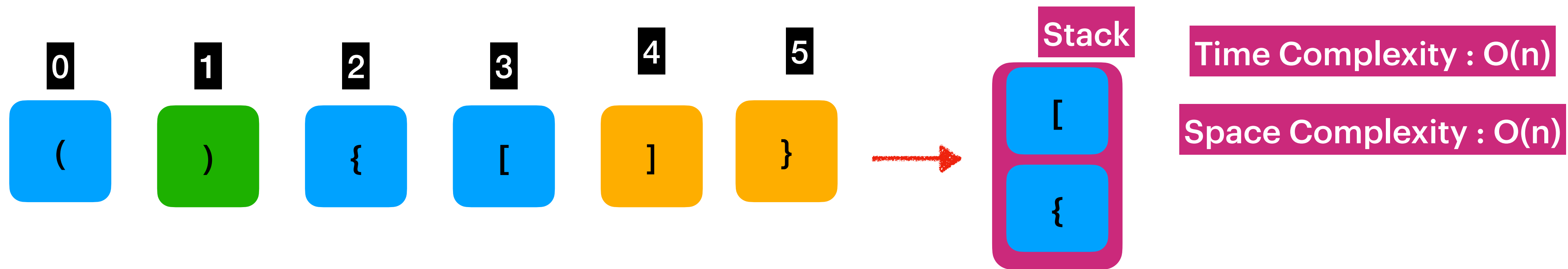
Algorithm :

When ever we find either "(" or "{" or "[" → Push into the stack.
Whenever we find either ")" or "}" or "]"
the element in the stack should be equivalent "(" or "{" or "["
if so pop() it otherwise return false.

Base Check : After processing all the elements return true if the stack is Empty.

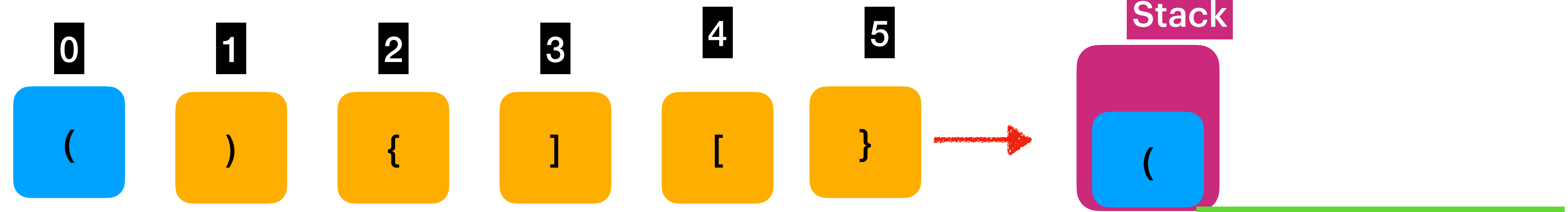
(){}[] → is valid



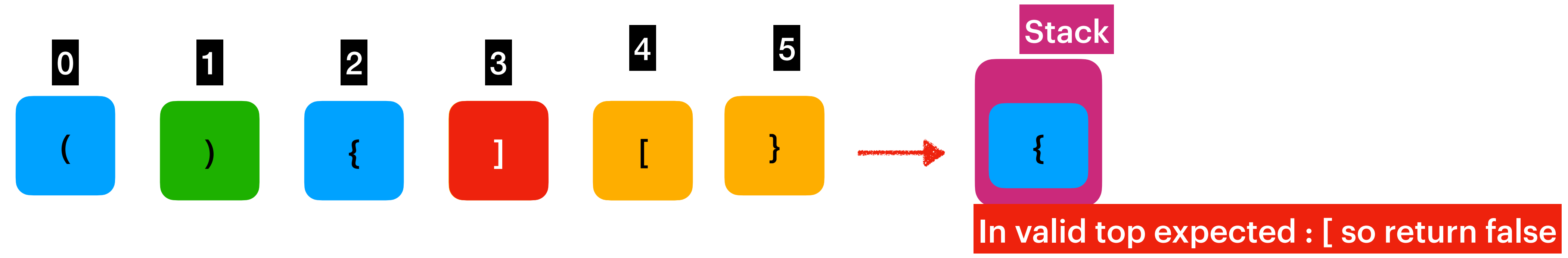
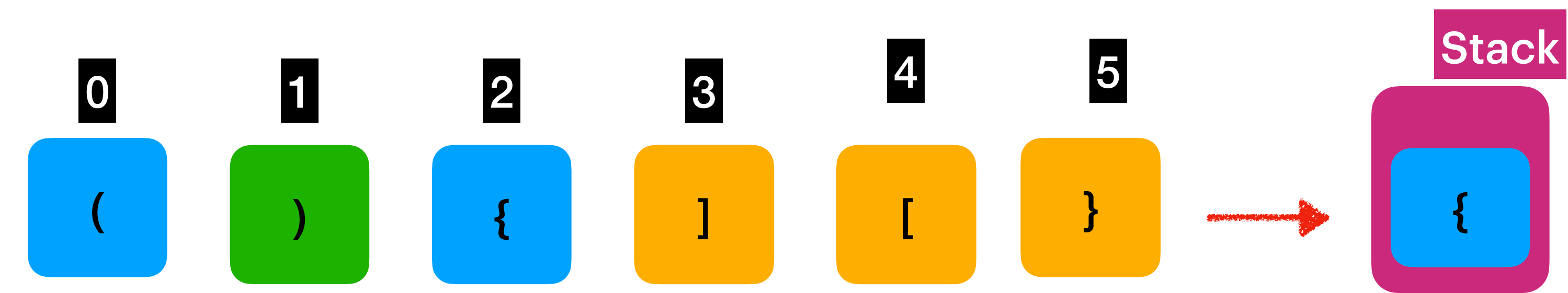
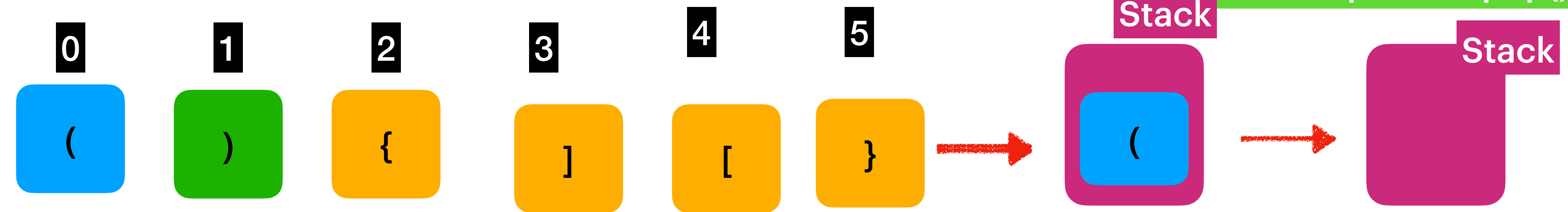


Stack is Empty() —> return true

() {] [} —> is not valid



Valid top —> so pop()



921. Minimum Add to Make Parentheses Valid

Medium 2612 148 Add to List Share

A parentheses string is valid if and only if:

- It is the empty string,
- It can be written as `AB` (`A` concatenated with `B`), where `A` and `B` are valid strings, or
- It can be written as `(A)` , where `A` is a valid string.

You are given a parentheses string `s` . In one move, you can insert a parenthesis at any position of the string.

- For example, if `s = "(())"` , you can insert an opening parenthesis to be `"((())"` or a closing parenthesis to be `"()))"` .

Return *the minimum number of moves required to make* `s` *valid*.

Example 1:

Input: `s = "()"`
Output: 1

Example 2:

Input: `s = "(((`
Output: 3

Constraints:

- `1 <= s.length <= 1000`
- `s[i]` is either `'('` or `')'` .

`"()))(("` —-> need 4 braces —> `() ' (') ' (') (') ' (')'`

Algorithm :

-> Take two stacks :: openBrachStack. & closedBraceStack.

-> Whenever we find openBrace add into openStack.

-> When ever we find closedBrace `')'`, check peek from OpenStack is it `' ('` if so pop() it otherwise add into closedBraceStack.

—> return openBraceStack.size() + closedBraceStack.size().

Time Complexity : $O(n)$

Space Complexity : $O(n)$

TimeComplexity : $O(n)$

Space Complexity : $O(n)$

For Explanation through Algonotes.

"()))((" —-> need 4 braces —> () '(') '(')(') (')'

Improvement on Space

Algorithm :

-> Take openCount, closedCount.

-> Whenever we find openBrace '(', add increment openCount.

-> When ever we find closedBrace ')', check openCount > 0 if so
Decrement openCount otherwise increment closedCount.

—> return openCount + closedCount.

Time Complexity : $O(n)$

Space Complexity : $O(1)$

TimeComplexity : $O(n)$

Space Complexity : $O(1)$