

## Discrete Mathematics

### Model Question Paper

#### Section - A

① (a) Show that  $P \rightarrow Q$  and  $\neg Q \rightarrow \neg P$  are logically Equivalent.

Sol Let us construct a truth table for  $P \rightarrow Q$  and  $\neg Q \rightarrow \neg P$ .

P	Q	$P \rightarrow Q$	$\neg Q$	$\neg P$	$\neg Q \rightarrow \neg P$
T	T	T	F	F	T
T	F	F	T	F	F
F	T	T	F	T	T
F	F	T	T	T	T

Since both the  $P \rightarrow Q$  &  $\neg Q \rightarrow \neg P$  columns are having identical truth values for all the combinations of the truth values of P and Q,  $P \rightarrow Q$  and  $\neg Q \rightarrow \neg P$  are logically equivalent.

(b) Show that the relation  $\leq$  (less than or equal) defined on the set of positive integers  $\mathbb{Z}^+$  is a partial order relation.

Sol Let  $a \in \mathbb{Z}^+$   
Since  $a \leq a$   $\forall a \in \mathbb{Z}^+$   
 $\therefore a \leq$  is a reflexive relation on  $\mathbb{Z}^+$

Let  $a, b \in \mathbb{Z}^+$

If  $a \leq b$  and  $b \leq a$  then  $a = b$  for every integer  $a, b$ .

$\therefore \leq$  is an anti-symmetric relation on  $\mathbb{Z}^+$

Let  $a, b, c \in \mathbb{Z}^+$

If  $a \leq b$  and  $b \leq c$  then

REDMI NOTE 6 PRO  
MI DUAL CAMERA for all  $a, b, c \in \mathbb{Z}^+$

$\leq$  is a transitive relation on  $\mathbb{Z}^+$

Since  $\leq$  is reflexive, anti-symmetric and transitive

$\leq$  is a partial order relation on  $\mathbb{Z}^+$ .

(c) Show that  $R \wedge (P \vee Q)$  is a valid conclusion from

The premises  $P \vee Q$ ,  $Q \rightarrow R$ ,  $P \rightarrow M$  and  $\neg M$ .

Sol

We can derive  $R \wedge (P \vee Q)$  from the given premises  
in the following steps.

$$\{1\} (1) P \rightarrow M$$

Rule P

$$\{2\} (2) \neg M$$

Rule P

$$\{1,2\} (3) \neg P$$

Rule T (1), (2)

~~$P \rightarrow M, \neg M \Rightarrow \neg P$~~

$$\{4\} (4) P \vee Q$$

Rule P.

$$\{1,2,4\} (5) Q$$

Rule T (3), (4)  $\neg P, P \vee Q \Rightarrow Q$

$$\{6\} (6) Q \rightarrow R$$

Rule P

$$\{1,2,4,6\} (7) R$$

Rule T (5), (6)  $Q, Q \rightarrow R \Rightarrow R$

$$\{1,2,4,6\} (8) R \wedge (P \vee Q)$$

Rule T (1), (4)  $R, P \vee Q \Rightarrow R \wedge (P \vee Q)$

i.e.  $R \wedge (P \vee Q)$  is a valid conclusion from the premises

$P \vee Q$ ,  $Q \rightarrow R$ ,  $P \rightarrow M$  and  $\neg M$ .

(d) If  $R$  be a relation on the set of integers  $\mathbb{Z}$ , defined by  $R = \{(x,y) / x \in \mathbb{Z}, y \in \mathbb{Z}, (x-y) \text{ is divisible by } 6\}$  is an equivalence relation.

Sol

$$R = \{(x-y) / x \in \mathbb{Z}, y \in \mathbb{Z}, (x-y) \text{ is divisible by } 6\}$$

Let  $x \in \mathbb{Z}$  then

$x-x=0$ , which is divisible by 6

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MI DUAL CAMERA  $\forall x, y \in \mathbb{Z}$ .

$\therefore R$  is reflexive

Let  $x, y \in \mathbb{Z}$

If  $(x-y)$  is divisible by 6 then  $(x,y) \in R$

Since  $x-y$  is divisible by 6

$\Rightarrow x-y = 6a$  ie multiple of 6

$\Rightarrow y-x = -6a$  which is also divisible by 6

$\therefore (y,x) \in R$  whenever  $(x,y) \in R$   $\forall x, y \in \mathbb{Z}$

$\therefore R$  is symmetric.

Let  $x, y, z \in \mathbb{Z}$

and  $(x,y) \in R$  and  $(y,z) \in R$

~~$x-y$  is of the form~~

$x-y = 6a$   $\dots$   $1-ND = ND$

$y-z$  is of the form  $y-z = 6b$   $\dots$

add ①, ②

$$x-y + y-z = 6a + 6b \quad 1-8$$

$$\Rightarrow x-z = 6(a+b) \quad 1-8$$

$\therefore (x-z)$  is divisible by 6.

$\therefore (x,z) \in R$  whenever  $(x,y) \in R$  &  $(y,z) \in R$   $\forall x, y, z \in \mathbb{Z}$

$\therefore R$  is transitive.

Since  $R$  is reflexive, symmetric and transitive

$\therefore R$  is an equivalence relation



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②(a) Solve the recurrence relation  $a_n = a_{n-1} + 2$ ,  $n \geq 1$   
 Subject to initial condition  $a_1 = 3$

Sol

$$a_n = a_{n-1} + 2$$

since the above recurrence relation is non-homogeneous  
 the required solution is

$$a_n = a_n^{[h]} + a_n^{[P]}$$

Here  $a_n^{[h]}$  is the solution for the homogeneous part  
 i.e. the associated homogeneous recurrence relation.

$$\text{i.e. } a_n = a_{n-1}$$

$$\Rightarrow a_n - a_{n-1} = 0 \quad \text{--- (2)}$$

The characteristic equation for eq(2) is

$$\lambda - 1 = 0$$

$$\Rightarrow \lambda = 1$$

Now, the solution for eq(2) is  $a_n^{[h]} = \lambda^n = 1^n = 1$

$$= \lambda^n = 1^n = 1 \quad \text{--- (3)}$$

Next, we have to find a particular polynomial solution

$$\text{i.e. } a_n^{[P]}$$

Since  $F(n) = 2$ , i.e. a polynomial in  $n$  of degree zero.

The reasonable solution is

$$a_n^{[P]} = n^m (A_0 + A_1 n + \dots + A_2 n^2 + \dots + A_m n^m)$$

Here  $m$  is the multiplicity of the root which is '1'

REDMI NOTE 6 PRO =  $A_n$  — (4)  
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Given  $a_n = a_{n-1} + 2$

Substitute  $a_n = A_0 n$

$$\Rightarrow A_0 n = A_0(n-1) + 2$$

$$\Rightarrow A_0 n = A_0 n - A_0 + 2$$

$$\Rightarrow A_0 = 2 \quad \text{--- (1)}$$

$$\therefore a_n^{[p]} = 2n \quad \text{--- (5)}$$

since from eq (1)  $a_n = a_n^{[h]} + a_n^{[p]}$

$$\Rightarrow a_n = d_1 + 2n \quad \text{--- (6)}$$

Since the initial condition  $a_1 = 3$

~~at 1st term~~ put  $n=1$  in eq (6)

$$\Rightarrow a_1 = d_1 + 2(1)$$

$$\Rightarrow a_1 = d_1 + 2 \quad \text{at 1st term}$$

Substitute  $a_1 = 3$  in eq (6)

$$\Rightarrow 3 = d_1 + 2$$

$$\Rightarrow d_1 = 3 - 2 = 1$$

$\therefore$  The required solution is  $a_n = 1 + 2n$

$$a_n = 1 + 2n$$

Iteration Method:

$$a_n = a_{n-1} + 2 \quad n \geq 2, a_1 = 3$$

$$a_n = a_{n-1} + 2 + 2 + \dots + 2 \quad (\text{n terms})$$

$$= (a_{n-2} + 2) + 2$$

$$= ((a_{n-3} + 2) + 2) + 2 \quad (\text{n terms})$$

$$= a_{n-(n-1)} + (n-1)2$$



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$$\begin{aligned}
 &= a_1 + (n-1)2 \\
 &= 3 + 2n - 2 \\
 &= 2n + 1 \\
 &\equiv 8 + (1-n)a_1 = nA
 \end{aligned}$$

(b) How many ways are there to assign five different jobs to four different employees if every employee assigned at least one job?

Sol We know that if  $m$  and  $n$  be positive integers with  $m > n$ . Then the expansion

$$n^m = nC_1(n-1)^{m-1} + nC_2(n-2)^{m-2} + \dots + (-1)^{n-1}C_{n-1}$$

gives the onto functions from a set with  $m$  elements to a set with  $n$  elements.

Now assume the assignment of jobs as a function from the set of five jobs to the set of four employees.

An assignment in which every employee gets atleast one job is the same as onto function from the set of jobs to the set of employees.

$$\therefore 4^5 = 4C_1(3)^5 + 4C_2(2)^5 - 4C_3(1)^5$$

$$= 1024 - 4 \cdot 243 + 6 \cdot 320 - 4 \cdot 1$$

$$= 1024 - 972 + 192 - 4 = 240$$

REDMI NOTE 6 PRO MI DUAL CAMERA assigned at least one job.

Q) Applying pigeon hole principle show that of any 14 integers are selected from the set  $S = \{1, 2, 3, \dots, 25\}$  there are at least two whose sum is 26. Also write a statement that generalizes this result.

Sol) Let the pigeons be the numbers selected.

Define 13 pigeon holes corresponding to the 13 sets:

$$\begin{aligned} & \{1, 25\}, \{2, 24\}, \{3, 23\}, \{4, 22\}, \{5, 21\}, \{6, 20\}, \{7, 19\}, \\ & \{8, 18\}, \{9, 17\}, \{10, 16\}, \{11, 15\}, \{12, 14\}, \{13\}. \end{aligned}$$

Notice that the numbers in each of the first 12 sets sum is 26 and there is no pair of distinct numbers containing 13 that sum to 26.

When a number is selected, it gets placed into the pigeonhole corresponding to the set that contains it. Atmost one number can go into the pigeonhole corresponding to  $\{13\}$ .

Since 14 numbers are selected and placed in 13 pigeonholes, some pigeonhole contains two numbers and those two numbers sum to 26.

$\therefore$  If  $n+1$  numbers are selected from  $\{1, 2, \dots, 2n\}$  then some two of these numbers sum to  $2n$



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② (d) In a class of 25 students, 12 have taken mathematics, 8 have taken mathematics but not biology. Find the number of students who have taken mathematics and biology and those who have taken biology but not mathematics.

Sol

Let  $n(M) = \text{no. of students taken Mathematics} = 12$ .

$n(B) = \text{no. of students taken Biology}$

Since there are 25 students in the class

Then  $n(M \cup B) = 25$ .

No. of students taken Mathematics but not Biology

$n(M \cap \bar{B}) = 8$ .

No. of students taken both Mathematics and Biology

$$\bullet n(M \cap B) = n(M) - n(M \cap \bar{B})$$

$$= 25 - 8$$

$$= 4.$$

Since N.K.T.  $n(M \cup B) = n(M) + n(B) - n(M \cap B)$

$$\Rightarrow 25 = 12 + n(B) - 4$$

$$\Rightarrow 25 = 8 + n(B)$$

$$\Rightarrow n(B) = 17$$

Now, No. of students taken Biology but not Mathematics

$$n(\bar{M} \cap B) = n(B) - n(M \cap B)$$

$$= 17 - 4$$

$$= \underline{\underline{13}}$$

∴ 4 students have taken both Mathematics & Biology.

O 13 REDMI NOTE 6 PRO MI DUAL CAMERA Ben Biology but not Mathematics.

③ (a) If  $G_1 = (V, E)$  be a directed graph with 'e' edges,  
then  $\sum_{v \in V} \deg_{G_1}^+(v) = \sum_{v \in V} \deg_{G_1}^-(v) = e$

Sol

Let  $G_1$  be a directed graph and it has 'e' edges.  
Since each edge incident on two vertices and  
each edge is a directed edge, it contributes '1' to  
the indegree of its terminal vertex and '1' to the  
outdegree of its initial vertex.

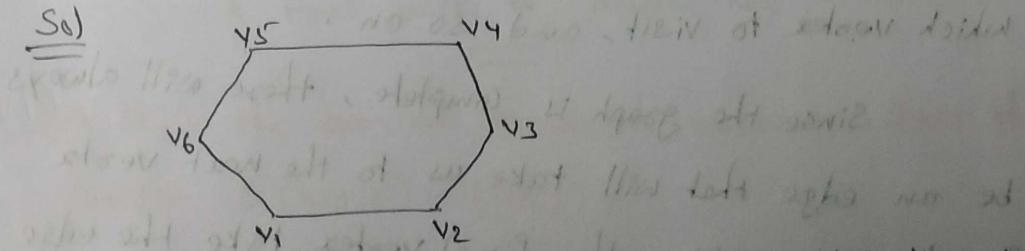
While counting the degree sum of the  
vertices of the graph, each edge is counted  
twice and each edge contributes '1' to the  
indegree sum and '1' to the outdegree sum of  
the vertices.

The sum of indegrees of the vertices of the  
graph  $G_1 = \sum_{v \in V} \deg_{G_1}^-(v) = e$

The sum of outdegrees of the vertices of the  
graph  $G_1 = \sum_{v \in V} \deg_{G_1}^+(v) = e$

(b) Show that  $C_6$  is a bipartite graph.

Sol



The above graph is a cycle on 6 vertices i.e.  $C_6$ .

We can partition its vertex set  $V = \{v_1, v_2, v_3, v_4, v_5, v_6\}$

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into two subsets  $V_1 = \{v_1, v_3, v_5\}$  and

$$V_2 = \{v_2, v_4, v_6\}$$

Observe here each and every edge of  $C_6$  connects a vertex in  $V_1$  and a vertex in  $V_2$ . There are no edges in between the vertices of  $V_1$ , as well as in  $V_2$ .  $\therefore C_6$  is a bipartite graph.

③

c) Show that the complete graph  $K_n$  has a Hamiltonian cycle.

Sol: Complete graph on  $n$  vertices  $K_n$  is a graph that has  $n$  vertices and each vertex is connected to every other vertex by an edge.

A hamiltonian cycle is a cycle which contains all the vertices in the given graph.

For  $K_n$ ,  $n \geq 3$ , there are  $n$  choices for where to begin the cycle, then  $(n-1)$  choices for which vertex to visit next, then  $(n-2)$  choices for which vertex to visit, and so on.

Since the graph is complete, there will always be an edge that will take us to the next vertex in the list. After the final vertex, take the edge that connects back to the starting vertex.

Hence we find a hamiltonian cycle.



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③ (d) Prove that a tree with 'n' vertices has  $n-1$  edges.

Sol We will prove this theorem using Mathematical Induction

Basis step:

$$\text{Let } n = 1$$

Then a tree with  $n=1$  vertex has no edges

~~base case~~  $\therefore$  It is true for  $n=1$ .

Induction Hypothesis:

Assume that every tree with  $K$  vertices has  $K-1$  edges, where  $K$  is a positive integer.

Induction Step:

Suppose  $T$  is a tree with  $K+1$  vertices.

Let  $v$  is a leaf of  $T$  and  $w$  be the parent of  $v$ .

Now, Remove vertex  $v$  and the edge connecting  $w$  and  $v$  from  $T$ . Now, the resultant tree  $T'$  is a tree with  $K$  vertices.

$K$  vertices.

By the induction hypothesis  $T'$  has  $K-1$  edges.

It follows that  $T$  has  $K$  edges, since it has one more than  $T'$ , the edge connecting  $v$  and  $w$ .

$\therefore$  A tree with 'n' vertices has  $n-1$  edges.

④ (a) Find the sum of products expansion for the

$$\text{function } F(x, y, z) = (x+y)\bar{z}$$

Sol REDMI NOTE 6 PRO  $(x+y)\bar{z}$   
MI DUAL CAMERA  $= x\bar{z} + y\bar{z}$   $\because$  Distributive Law

(6) (a) Define Hamilton circuit, Hamilton graph give examples

$$= x(y+z)\bar{z} + (x+\bar{x})y\bar{z} \quad [\because \text{Identity Law}]$$

$$= xy\bar{z} + x\bar{y}\bar{z} + \bar{x}y\bar{z} + \bar{x}\bar{y}\bar{z} \quad [\because \text{Distributive Law}]$$

$$= xy\bar{z} + x\bar{y}\bar{z} + \bar{x}y\bar{z} \quad [\because \text{Idempotent Law}]$$

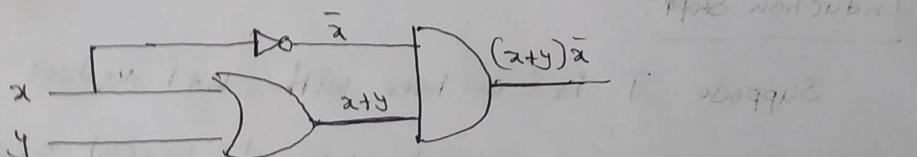
(4) (b) construct circuits that produce the following outputs.

(i)  $(x+y)\bar{z}$  (ii)  $\bar{x}(\bar{y}+\bar{z})$

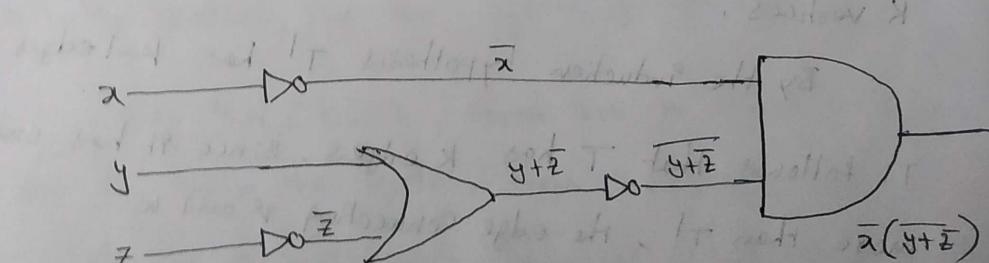
~~bit and follow the last four bits A~~

~~output will be a circuit. right~~

(i)  $(x+y)\bar{z}$



(ii)  $\bar{x}(\bar{y}+\bar{z})$



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④ (c) Show that distributive law  $x(y+z) = xy + xz$  is valid.

Sol Let us construct a truth table for both LHS & RHS.

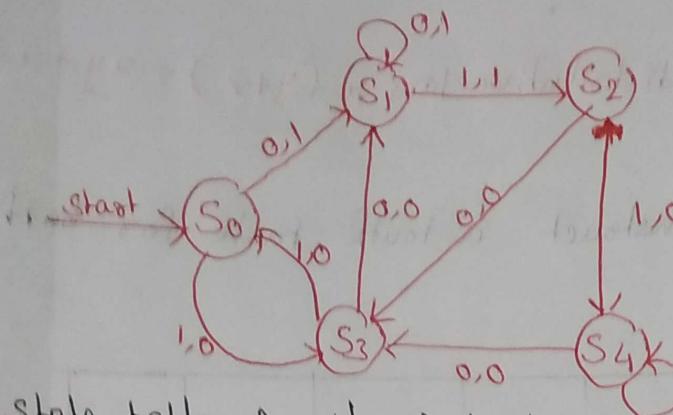
x	y	z	(y+z)	$x(y+z)$	$xy$	$xz$	$(xy+xz)$
0	0	0	0	0	0	0	0
0	0	1	1	0	0	0	0
0	1	0	1	0	0	0	0
0	1	1	1	0	0	0	0
1	0	0	0	0	0	0	0
1	0	1	1	0	0	0	0
1	1	0	1	1	0	1	1
1	1	1	1	1	1	1	1

Since both  $x(y+z)$  and  $(xy+xz)$  columns have identical truth values for all combinations of truth values of x, y and z  
 $x(y+z) = xy + xz$  is valid.

④ (d) construct the state table for the finite state machine with the state diagram shown in the following figure.



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State table for the finite state machine ↴

State	State Transition function		Output function g	
	Input=0	Input=1	Input=0	Input=1
S0	S1	S3	1 1	1 0
S1	S1	S2	1 0	0 1
S2	S3	S4	1 1	0 1
S3	S1	S0	1 0	0 1
S4	S3	S4	0	0

Section — B

Q) construct the truth table for  $P \wedge (\neg q \vee q)$

Sol

P	q	$\neg q$	$q \vee q$	$P \wedge (\neg q \vee q)$
T	T	F	T	T
T	F	T	T	T
F	T	F	T	F
F	F	T	T	F



⑤(b) write the following in symbolic form.

Every person is precious

Sol Let our univers of discourse is the people in  
the world.

Let  $P(x)$  be a predicate "x is Precious".

we can write the given statement in the following symbolic form

$\forall x \, P(x)$  will be true if and only if every  $x$  in the domain satisfies  $P(x)$ .

⑤ c) compute  $20!$  loops to answer

SOL 901 X

~~20 x 19 x 18~~

~~18~~ 18

20x19 without border

$$= \underline{\quad} \\ = \underline{\quad}$$

$$\text{Prove } A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$$

501 6 7 8 1 2 3 4 5 6 7 8 9 10 11 12

$$\Rightarrow AU(B \cap C) = \{x | x \in AU(B) \text{ and } x \in AU(C)\}$$

$$= \{x | x \in A \vee x \in B \wedge C\}$$

8-1-12 (1-7-10) 2

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$$= \{x | x \in (A \cup B) \cap x \in (A \cup C)\}$$

$$= \{x | x \in (A \cup B) \cap (A \cup C')\}$$

$$= (A \cup B) \cap (A \cup C')$$

⑤ (e) State and Prove Handshaking Theorem.

Sol: Let  $G_1(V, E)$  be an undirected graph with 'e' number of edges then  $\sum_{v \in V} \deg(v) = 2e$

i.e. The degree sum of the vertices of the given undirected graph is equal to twice the number of edges.

Proof: This is true because while counting the degree sum of the vertices, each edge is exactly counted twice. These pose the degree sum of vertices is equal to  $2 \times (\text{the no. of edges})$ .

⑤ (P) Define Hamilton circuit, Hamilton graph give examples to each.

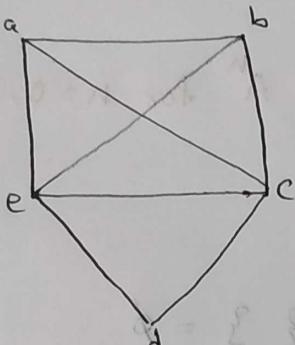
Sol

Hamilton Path: A simple path that covers all the vertices in the given graph is called Hamilton path.

Hamilton Circuit:

A simple circuit that covers all the vertices in the given graph exactly once, except starting vertex is called Hamilton circuit.

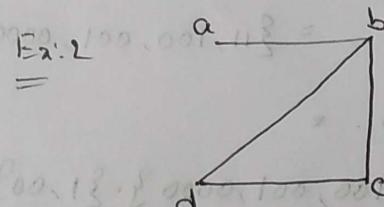
Ex: 1



In the above graph there exist a Hamilton circuit

for  $a-b-c-d-e-a$

Ex: 2



In the above graph no hamilton circuit, but a Hamilton path  $a-b-c-d$

A graph is a Hamilton graph if there exist a Hamilton circuit in the graph. Hence the graph in Ex-1 is a Hamilton graph.

⑤ ⑨ Find the duals of  $x(y+z)$  and  $\bar{x} \cdot 1 + (\bar{y}+z)$

Sol The dual of a Boolean expression is the expression that is obtained by interchanging '+' and '.', and interchanging 0's and 1's.

Now the dual of  $x(y+z) = x+y \cdot 1$

The dual of  $\bar{x} \cdot 1 + (\bar{y}+z) = (\bar{x}+0) \cdot (\bar{y} \cdot z)$

⑤ ⑩ Let  $A = \{1, 00\}$ , find  $A^n$  for  $n=0, 1, 2$  and 3

Sol

Since  $A = \{1, 00\}$

If  $n=0$  then  $A^0 = \{\}$  =  $\emptyset$

$n=1$  then  $A^1 = A = \{1, 00\}$

$n=2$  then  $A^2 = A \cdot A = \{1, 00\} \cdot \{1, 00\}$

=  $\{11, 100, 001, 0000\}$

$n=3$  then  $A^3 = A \cdot A^2$

=  $\{11, 100, 001, 0000\} \cdot \{1, 00\}$

=  $\{111, 1100, 1001, 10000, 0011, 00100, 00001, 000000\}$