

# Elliptic Curve Generation

Peter Manohar, Xingyou Song

November 22, 2015

## Abstract

The goal of this report is to....

## Contents

<b>1</b>	<b>Introduction</b>	<b>1</b>
<b>2</b>	<b>Motivations and Applications</b>	<b>3</b>
<b>3</b>	<b>Generating Curves</b>	<b>3</b>

## 1 Introduction

In this section, we will introduce some of the properties of elliptic curves and their mathematical properties.

### 1.1 Elliptic Curves

The simplest explanation of an elliptic curve is using the Weierstrass Equation; i.e. assume we have a field  $K$ , and consider the equation (in Weierstrass Form)

$$y^2 = x^3 + Ax + B \tag{1}$$

Note that here, because (through Vieta formulas),

$$((r_1 - r_2)(r_1 - r_3)(r_2 - r_3))^2 = -(4A^3 + 27B^2) \tag{2}$$

We do not consider curves with multiple roots, and so we take the constraint  $4A^3 + 27B^2 \neq 0$ .

### 1.2 Group Law

A property of elliptic curves is an abelian group law which forms on its points, which we will describe below. (A more general proof over all curves is by Reimann-Roch)

Geometrically, consider two points  $P_1 = (x_1, y_1), P_2 = (x_2, y_2)$ . Define the operation of

$$P_3 = P_1 \oplus P_2 \tag{3}$$

where we draw a line in between  $P_1, P_2$  that hits the elliptic curve again at a new point  $P'$ , then reflect  $P'$  about the x-axis to obtain  $P_3$ .

More specifically, we can write this explicitly in terms of algebra:

$$x_3 = s^2 - 2x_1y_3 = \quad (4)$$

TODO

### 1.3 Pairings and Isogenies

Due to the abelian group law on curves, we introduce the concept of an isogeny, i.e. a homomorphism between the elliptic curve groups, that is also algebraic in nature. More formally, an isogeny  $\phi$  is a homomorphism with respect to the group laws

$$\phi : E_1 \rightarrow E_2, \quad \phi(O) = O$$

An example of an isogeny from  $E_1$  to itself is multiplication by  $m$ , i.e.

$$[m] : E_1 \rightarrow E_1, \quad [m]P = P + P + P + \dots + P \text{ (m times)} \quad (5)$$

### 1.4 j-invariant

A question arises when two elliptic curves over a field  $E_1(K), E_2(K)$  are have a bijective isogeny, i.e. are isomorphic with respect to the group law. An intuitive answer is that the points on elliptic curves be transformed algebraically, which will use different Weierstrass Equations

Note the transformation

$$\begin{aligned} x' &= \mu^2 x \\ y' &= \mu^3 y \end{aligned}$$

implies, after plugging into the Weierstrass equation,

$$(y')^2 = (x')^3 + \mu^4 A x' + \mu^6 B \implies (y')^2 = (x')^3 + A' x' + B' \quad (6)$$

where  $A' = \mu^4 A, B' = \mu^6 B$ .  $\mu$  may not exist in the field  $K$ , but only in its closure  $\bar{K}$ .

From here, we define the j-invariant  $j$  of a Weierstrass Form to be

$$j(E) = 1728 \frac{4A^3}{4A^3 + 27B^2} \quad (7)$$

Note that the  $j$  invariant is homogenous; scalings of the form from (5) will still leave  $j$  constant. Given a  $j$ , then the canonical elliptic curve associated with this j-invariant will be

$$y^2 = x^3 + \frac{3j}{1728 - j}x + \frac{2j}{1728 - j} \quad (8)$$

Going back to the concept of closed field, note that for non-closed fields, we may have non-isomorphic curves with the same j-invariant.

For instance,  $y^2 = x^3 - 25x$ ,  $y^2 = x^3 - 4x$  have  $j = 1728$ . The first curve has infinitely many points in  $Q$ , but the second has only finite. The transformation  $(x, y) \rightarrow (\mu^2 x, \mu^3 y)$  only exists when we consider  $Q\sqrt{10}$ , since  $\mu = \sqrt{10}/2$  is the scaling factor.

Specifically, however, from this example, we do not need the full closure  $\bar{K}$ ; we only need  $d = \mu^2$ , which means we only need  $K(\sqrt{d})$ .

## 1.5 Twist

From the above, we can give different "twists" of curves that are isomorphic in the closed field  $\bar{K}$ , but not in  $K$ . Then if  $D$  is a square free integer, then for a elliptic curve  $E : y^2 = x^3 + Ax + B$ , the given twist  $E(D)$ , from the above, will be

$$E(D) : y^2 = x^3 + AD^2x + BD^3 \tag{9}$$

# 2 Motivations and Applications

## 2.1 Discrete Log Problem

In the discrete log problem, we are given any group  $G$ , with a base generator element  $P$ , with a ciphertext  $Q$ , where the problem involves finding  $k$  such that  $P^k = Q$  in  $G$ .

For elliptic curves, this becomes using a base point in a finite field  $F_p$ . However, note that by the theorem of finite abelian groups, the

# 3 Generating Curves

## 3.1 CM Method