# ${\bf IUT\ Serenity}, {\bf Islamic\ University\ of\ Technology}$

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```
Sublime Build
   "cmd" : ["g++ -std=c++14 -DSONIC $file_name -o
        $file_base_name && timeout 4s ./$file_base_name<</pre>
        inputf.in>outputf.in"],
   "selector" : "source.cpp",
   "file_regex": "^(..[^:]*):([0-9]+):?([0-9]+)?:? (.*)$ |Stress-tester
   "shell": true,
   "working_dir" : "$file_path"
vimrc
set mouse=a
 set termguicolors
 filetype plugin indent on
 syntax on
" Some useful settings
 set smartindent expandtab ignorecase smartcase
      incsearch relativenumber nowrap autoread splitright
       splitbelow
 set tabstop=4
                      "the width of a tab
 set shiftwidth=4
                      "the width for indent
 colorscheme torte
"auto pair curlybraces
 inoremap {<ENTER> {}<LEFT><CR><ESC><S-o>
" mapping jj to esc
 inoremap jj <ESC>
 "compile and run using file input put
 autocmd filetype cpp map <F5> :wa<CR>:!clear && g++ % - #include <ext/pb_ds/tree_policy.hpp>
      D LOCAL -std=c++17 -Wall -Wextra -Wconversion -
      Wshadow -Wfloat-equal -o ~/Codes/prog && (timeout 5
      ~/Codes/prog < ~/Codes/in) > ~/Codes/out<CR>
 "copy to input file
 map <F4> :!xclip -o -sel clip > ~/Codes/in <CR><CR>
 map <F6> :vsplit ~/Codes/in<CR>:split ~/Codes/out<CR><C template <typename DT>
      -w >= 20 < C - w > < < C - w > < C - h >
 " Leader key
 let mapleader=',,'
 " Copy template
 noremap <Leader>t :!cp ~/Codes/temp.cpp %<CR><CR>
 :autocmd BufNewFile *.cpp Or ~/Codes/temp.cpp
```

```
"note if vim-features +clipboard is not found, it will
  "for fast check :echo has('clipboard) = 0 if clipboard
      features not present,
  "need vim-gtk / vim-gtk3 package for this
 set clipboard=unnamedplus
#!/bin/bash
# Call as stresstester GENERATOR SOL1 SOL2 ITERATIONS [--
    countl
for i in $(seq 1 "$4"); do
    [[ $* == *--count* ]] && echo "Attempt $i/$4"
   $1 > in.txt
   $2 < in.txt > out1.txt
   $3 < in.txt > out2.txt
   diff -y out1.txt out2.txt > diff.txt
   if [ $? -ne 0 ] ; then
       echo "Differing Testcase Found:"; cat in.txt
       echo -e "\nOutputs:"; cat diff.txt
       break
done
    All Misc
```

#### 1.1 All Macros

```
//#pragma GCC optimize("Ofast")
//#pragma GCC optimization ("03")
//#pragma comment(linker, "/stack:200000000")
//#pragma GCC optimize("unroll-loops")
//#pragma GCC target("sse,sse2,sse3,ssse3,sse4,popcnt,abm
    ,mmx,avx,tune=native")
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
   //find_by_order(k) --> returns iterator to the kth
        largest element counting from 0
   //order_of_key(val) --> returns the number of items
        in a set that are strictly smaller than our item
using ordered_set = tree <DT, null_type, less<DT>,
    rb_tree_tag,tree_order_statistics_node_update>;
/*--- DEBUG TEMPLATE STARTS HERE ---*/
#ifdef Serenity
void show(int x) {cerr << x;}</pre>
void show(long long x) {cerr << x;}</pre>
void show(double x) {cerr << x;}</pre>
void show(char x) {cerr << '\'' << x << '\'';}</pre>
```

```
void show(const string &x) {cerr << '\"' << x << '\"';}</pre>
void show(bool x) {cerr << (x ? "true" : "false");}</pre>
template<typename T, typename V>
void show(pair<T, V> x) { cerr << ','; show(x.first);</pre>
    cerr << ", "; show(x.second); cerr << '}'; }</pre>
template<typename T>
void show(T x) {int f = 0; cerr << "{"; for (auto &i: x)
    cerr << (f++ ? ", " : ""), show(i); cerr << "}";}</pre>
void debug_out(string s) {
   cerr << '\n';
template <typename T, typename... V>
void debug_out(string s, T t, V... v) {
   s.erase(remove(s.begin(), s.end(), ''), s.end());
   cerr << "
                    "; // 8 spaces
   cerr << s.substr(0, s.find(','));</pre>
   s = s.substr(s.find(',') + 1);
   cerr << " = ":
   show(t);
   cerr << endl:
   if(sizeof...(v)) debug_out(s, v...);
#define debug(x...) cerr << "LINE: " << __LINE__ << endl;</pre>
     debug_out(#x, x); cerr << endl;</pre>
#else
#define debug(x...)
#endif
mt19937 rnd(chrono::steady_clock::now().time_since_epoch
    ().count());
uniform_int_distribution<LL>gen(1, 100);
gen(rnd);
shuffle(a.begin(), a.end(), rnd);
```

```
1.2 RandomGenerator
#include <bits/stdc++.h>
using namespace std;
random_device rand_dev;
mt19937 generator(rand_dev());
int main() {
   ios::sync_with_stdio(0);
   cin.tie(0);
   uniform_int_distribution<long long> distr(1, 20);
   uniform_int_distribution<long long> dq(1, 10);
```

```
set<int> st;
int n = distr(generator);
for(int i=1; i<=n; i++) st.insert(distr(generator));</pre>
cout<<st.size()<<endl;</pre>
uniform_int_distribution<long long> gg(1, st.size());
for(int x : st) cout<<x<' ';</pre>
cout << endl;</pre>
int q = dq(generator);
cout<<q<<endl;</pre>
while(q--){
    int 1 = gg(generator);
    int r = gg(generator);
    if(1>r) swap(1,r);
    int k = gg(generator);
    cout<< 1 << ' '<<r<<' ' '<<k<<endl:
}
```

## 2 DP

## 2.1 1D-1D

```
/// Author: anachor
#include <bits/stdc++.h>
using namespace std;
/// Solves dp[i] = min(dp[j] + cost(j+1, i)) given that
    cost() is QF
long long solve1D(int n, long long cost(int, int)) {
  vector<long long> dp(n + 1), opt(n + 1);
  deque<pair<int, int>> dq;
 dq.push_back({0, 1});
 dp[0] = 0;
  for (int i = 1; i <= n; i++) {
   opt[i] = dq.front().first;
   dp[i] = dp[opt[i]] + cost(opt[i] + 1, i);
   if (i == n) break;
   dq[0].second++;
   if (dq.size() > 1 && dq[0].second == dq[1].second) dq
        .pop_front();
   int en = n;
   while (dq.size()) {
     int o = dq.back().first, st = dq.back().second;
```

```
if (dp[o] + cost(o + 1, st) >= dp[i] + cost(i + 1,
          st))
       dq.pop_back();
     else {
       int lo = st, hi = en;
       while (lo < hi) {</pre>
         int mid = (lo + hi + 1) / 2;
         if (dp[o] + cost(o + 1, mid) < dp[i] + cost(i +</pre>
              1, mid))
           lo = mid;
         else
           hi = mid - 1;
       if (lo < n) dq.push_back({i, lo + 1});</pre>
       break;
     en = st - 1;
    if (dq.empty()) dq.push_back({i, i + 1});
  return dp[n];
/// Solves https://open.kattis.com/problems/
    coveredwalkway
const int N = 1e6 + 7;
long long x[N];
int c;
long long cost(int 1, int r) { return (x[r] - x[1]) * (x[
    r] - x[1]) + c: }
int main() {
 ios::sync_with_stdio(false);
  cin.tie(0);
  int n;
  cin >> n >> c;
  for (int i = 1; i \le n; i++) cin >> x[i];
  cout << solve1D(n, cost) << endl;</pre>
```

## 2.2 Convex Hull Trick

```
struct line {
  ll m, c;
  line() {}
  line(ll m, ll c) : m(m), c(c) {}
};
struct convex_hull_trick {
  vector<line> lines;
  int ptr = 0;
```

```
convex hull trick() {}
 bool bad(line a, line b, line c) {
   return 1.0 * (c.c - a.c) * (a.m - b.m) < 1.0 * (b.c -
         a.c) * (a.m - c.m):
 void add(line L) {
   int sz = lines.size():
   while (sz >= 2 && bad(lines[sz - 2], lines[sz - 1], L
       )) {
     lines.pop_back();
     sz--;
   lines.pb(L);
 ll get(int idx, int x) { return (1ll * lines[idx].m * x
       + lines[idx].c); }
 11 query(int x) {
   if (lines.empty()) return 0;
   if (ptr >= lines.size()) ptr = lines.size() - 1;
   while (ptr < lines.size() - 1 && get(ptr, x) > get(
        ptr + 1, x)) ptr++;
   return get(ptr, x);
 }
11 sum[MAX];
11 dp[MAX];
int arr[MAX];
int main() {
 fastio;
 int t:
 cin >> t;
 while (t--) {
   int n, a, b, c;
   cin >> n >> a >> b >> c;
   for (int i = 1; i <= n; i++) cin >> sum[i];
   for (int i = 1; i <= n; i++) dp[i] = 0, sum[i] += sum</pre>
        [i - 1];
   convex hull trick cht:
   cht.add(line(0, 0));
   for (int pos = 1; pos <= n; pos++) {</pre>
     dp[pos] = cht.query(sum[pos]) - 111 * a * sqr(sum[
         pos]) - c;
     cht.add(line(211 * a * sum[pos], dp[pos] - a * sqr(
         sum[pos])));
   11 \text{ ans} = (-111 * dp[n]);
   ans += (111 * sum[n] * b);
   cout << ans << "\n";
```

## 2.3 Divide and Conquer dp const int K = 805, N = 4005; LL dp[2][N], \_cost[N][N]; // 1-indexed for convenience LL cost(int 1. int r) { return \_cost[r][r] - \_cost[1 - 1][r] - \_cost[r][1 - 1] + \_cost[1 - 1][1 - 1] >> 1; void compute(int cnt, int 1, int r, int opt1, int optr) { if (1 > r) return; int mid = 1 + r >> 1; LL best = INT\_MAX; int opt = -1; for (int i = optl; i <= min(mid, optr); i++) {</pre> LL cur = dp[cnt ^ 1][i - 1] + cost(i, mid); if (cur < best) best = cur, opt = i;</pre> dp[cnt][mid] = best; compute(cnt, 1, mid - 1, optl, opt); compute(cnt, mid + 1, r, opt, optr);

## 2.4 Knuth optimization

return dp[k & 1][n];

LL dnc\_dp(int k, int n) {

 $fill(dp[0] + 1, dp[0] + n + 1, INT_MAX);$ 

for (int cnt = 1; cnt <= k; cnt++) {</pre>

compute(cnt & 1, 1, n, 1, n);

```
const int N = 1005;
LL dp[N][N], a[N];
int opt[N][N];
LL cost(int i, int j) { return a[j + 1] - a[i]; }
LL knuth_optimization(int n) {
 for (int i = 0; i < n; i++) {</pre>
   dp[i][i] = 0;
   opt[i][i] = i;
 for (int i = n - 2; i \ge 0; i--) {
   for (int j = i + 1; j < n; j++) {
     LL mn = LLONG_MAX;
     LL c = cost(i, j);
     for (int k = opt[i][j - 1]; k <= min(j - 1, opt[i +</pre>
           1][j]); k++) {
       if (mn > dp[i][k] + dp[k + 1][i] + c) {
         mn = dp[i][k] + dp[k + 1][j] + c;
         opt[i][j] = k;
     }
```

```
dp[i][j] = mn;
   }
 }
 return dp[0][n - 1];
2.5 Li Chao Tree
struct line {
 LL m, c;
 line(LL m = 0, LL c = 0) : m(m), c(c) {}
LL calc(line L, LL x) { return 1LL * L.m * x + L.c; }
struct node {
 LL m, c;
 line L;
 node *lft, *rt;
 node(LL m = 0, LL c = 0, node *lft = NULL, node *rt =
     : L(line(m, c)), lft(lft), rt(rt) {}
struct LiChao {
 node *root:
 LiChao() { root = new node(); }
 void update(node *now, int L, int R, line newline) {
   int mid = L + (R - L) / 2;
   line lo = now->L, hi = newline;
   if (calc(lo, L) > calc(hi, L)) swap(lo, hi);
   if (calc(lo, R) <= calc(hi, R)) {</pre>
     now->L = hi;
     return;
   if (calc(lo, mid) < calc(hi, mid)) {</pre>
     now->L = hi;
     if (now->rt == NULL) now->rt = new node();
     update(now->rt, mid + 1, R, lo);
   } else {
     now->L = lo:
     if (now->lft == NULL) now->lft = new node();
     update(now->lft, L, mid, hi);
   }
 LL query(node *now, int L, int R, LL x) {
   if (now == NULL) return -inf:
   int mid = L + (R - L) / 2;
   if (x \le mid)
     return max(calc(now->L, x), query(now->lft, L, mid,
   else
     return max(calc(now->L, x), query(now->rt, mid + 1, |}
          R, x));
```

```
}
2.6 SOS-DP
for(int i = 0; i<(1<<N); ++i)</pre>
F[i] = A[i]:
for(int i = 0; i < N; ++i) for(int mask = 0; mask < (1<<N)
    : ++mask){
if(mask & (1<<i))</pre>
 F[mask] += F[mask^(1<< i)];
3 Data Structure
3.1 Segment Tree
const int N = 1e5+7;
int a[N];
LL tr[4*N];
LL lz[4*N];
///1. Merge left and right
LL combine (LL left, LL right) {
 return left + right;
///2. Push lazy down and merge lazy
void propagate(int u, int st, int en) {
 if (!lz[u]) return;
 tr[u] += (en-st+1)*lz[u];
 if (st!=en) {
  lz[2*u] += lz[u];
   lz[2*u+1] += lz[u];
 }
 lz[u] = 0;
void build(int u, int st, int en) {
 if (st==en) {
   tr[u] = a[st];
                         ///3. Initialize
   lz[u] = 0:
 }
 else {
   int mid = (st+en)/2;
   build(2*u, st, mid);
   build(2*u+1, mid+1, en);
   tr[u] = combine(tr[2*u], tr[2*u+1]);
   lz[u] = 0:
                         ///3. Initialize
```

```
void update(int u, int st, int en, int l, int r, int x) {
  propagate(u, st, en);
  if (r<st || en<1) return;</pre>
  else if (1<=st && en<=r) {
   lz[u] += x;
                          ///4. Merge lazy
    propagate(u, st, en);
  else {
    int mid = (st+en)/2:
    update(2*u, st, mid, 1, r, x);
    update(2*u+1, mid+1, en, l, r, x);
    tr[u] = combine(tr[2*u], tr[2*u+1]);
}
LL query(int u, int st, int en, int l, int r) {
  propagate(u, st, en);
if (r<st || en<1) return 0;</pre>
                                 /// 5. Proper null value int query(int curr, int b, int e, int l, int r){
else if (1<=st && en<=r) return tr[u]:
else {
  int mid = (st+en)/2;
  return combine(query(2*u, st, mid, 1, r), query(2*u+1,
      mid+1, en, l, r));
```

## 3.2 Persistent Segment Tree

```
struct Node{
 int val,left,right;
int a[500010],cnt=0,versions[500010];
Node tree[25*N+15]:
int build(int b,int e){
  cnt++:
 int curr=cnt;
  if(b==e){
   tree[curr].val=1e9+7;
   return curr;
  int mid=(b+e)/2:
  tree[curr].left=build(b,mid);
  tree[curr].right=build(mid+1,e);
  tree[curr].val=min(tree[tree[curr].left].val,tree[tree[
      curr].right].val);
 return curr;
int update(int prev,int b,int e,int idx,int x){
  cnt++;
 int curr=cnt;
 if(b==e){}
```

# 3.3 Implicit Segment Tree

curr].right].val);

if(e<1 || b>r) return 1e9+7;

if(b>=1 && e<=r) return tree[curr].val;</pre>

int q1=query(tree[curr].left,b,mid,l,r);

int q2=query(tree[curr].right,mid+1,e,l,r);

tree[curr].val=x:

return curr;

int mid=(b+e)/2:

if(idx<=mid){</pre>

return curr;

int mid=(b+e)/2;

return min(q1,q2);

}

else{

```
struct node {
 int val:
 node *lft, *rt;
 node() {}
 node(int val = 0) : val(val), lft(NULL), rt(NULL) {}
struct implicit_segtree {
 node *root;
 implicit_segtree() {}
 implicit_segtree(int n) {
   root = new node(n);
 void update(node *now, int L, int R, int idx, int val)
     {
   if (L == R) {
     now -> val += val;
     return;
   int mid = L + (R - L) / 2:
   if (now->lft == NULL) now->lft = new node(mid - L +
   if (now->rt == NULL) now->rt = new node(R - mid);
```

tree[curr].left=update(tree[prev].left,b,mid,idx,x);

tree[curr].right=update(tree[prev].right,mid+1,e,idx,

tree[curr].val=min(tree[tree[curr].left].val,tree[tree[

tree[curr].right=tree[prev].right;

tree[curr].left=tree[prev].left;

```
if (idx <= mid) update(now->lft, L, mid, idx, val);
  else update(now->rt, mid + 1, R, idx, val);
  now->val = (now->lft)->val + (now->rt)->val;
int query(node *now, int L, int R, int k) {
  if (L == R) return L:
  int mid = L + (R - L) / 2;
  if (now->lft == NULL) now->lft = new node(mid - L +
  if (now->rt == NULL) now->rt = new node(R - mid);
  if (k <= (now->lft)->val) return query(now->lft, L,
      mid, k);
  else return query(now->rt, mid + 1, R, k - (now->lft)
      ->val);
}
```

#### 3.4 HashTable

```
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
const int RANDOM = chrono::high_resolution_clock::now().
    time_since_epoch().count();
unsigned hash_f(unsigned x) {
 x = ((x >> 16) ^x) * 0x45d9f3b;
 x = ((x >> 16) ^x) * 0x45d9f3b;
 return x = (x >> 16) \hat{x};
unsigned hash_combine(unsigned a, unsigned b) { return a
    * 31 + b: }
struct chash {
 int operator()(int x) const { return hash_f(x); }
typedef gp_hash_table<int, int, chash> gp;
gp table;
```

#### 3.5 DSU With Rollbacks

```
struct Rollback DSU {
 int n;
 vector<int> par, sz;
 vector<pair<int, int>> op;
 Rollback_DSU(int n) : par(n), sz(n, 1) {
  iota(par.begin(), par.end(), 0);
  op.reserve(n);
 int Anc(int node) {
   for (; node != par[node]; node = par[node])
    ; // no path compression
```

```
return node;
  int Anc(int node) { // path compression, doesnt work
      will rollback
   if(par[node] == node) return node;
   return par[node] = Anc(par[node]);
  void Unite(int x, int y) {
   if (sz[x = Anc(x)] < sz[y = Anc(y)]) swap(x, y);
   op.emplace_back(x, y);
   par[v] = x;
   sz[x] += sz[y];
 void Undo(int t) {
   for (; op.size() > t; op.pop_back()) {
     par[op.back().second] = op.back().second;
     sz[op.back().first] -= sz[op.back().second];
   }
 }
};
```

## Binary Trie

const int N = 1e7 + 5, b = 30;

```
int tc = 1:
struct node {
  int vis = 0;
  int to [2] = \{0, 0\};
  int val[2] = \{0, 0\};
  void update() {
   to[0] = to[1] = 0;
   val[0] = val[1] = 0;
   vis = tc:
T[N + 2]:
node *root = T;
int ptr = 0;
node *nxt(node *cur, int x) {
  if (cur->to[x] == 0) cur->to[x] = ++ptr;
  assert(ptr < N);
  int idx = cur - to[x]:
  if (T[idx].vis < tc) T[idx].update();</pre>
 return T + idx;
int query(int j, int aj) {
 int ans = 0, jaj = j ^ aj;
  node *cur = root;
  for (int k = b - 1: "k: k--) {
   maximize(ans, nxt(cur, (jaj >> k & 1) ^ 1)->val[1 ^ (|void push(int _node) { node[++tm] = _node; }
        ai >> k & 1)]);
   cur = nxt(cur, (jaj >> k & 1));
```

```
return ans;
void insert(int j, int aj, int val) {
 int jaj = j ^ aj;
 node *cur = root;
 for (int k = b - 1; ~k; k--) {
   cur = nxt(cur, (jaj >> k & 1));
   maximize(cur->val[j >> k & 1], val);
void clear() {
 tc++:
 ptr = 0;
 root->update();
```

#### 3.7 BIT-2D

```
const int N = 1008:
int bit[N][N], n, m;
int a[N][N], q;
void update(int x, int y, int val) {
 for (; x < N; x += -x & x)
   for (int j = y; j < N; j \leftarrow -j & j) bit[x][j] += val;
int get(int x, int y) {
 int ans = 0:
 for (; x; x -= x & -x)
  for (int j = y; j; j -= j & -j) ans += bit[x][j];
 return ans:
int get(int x1, int y1, int x2, int y2) {
 return get(x2, y2) - get(x1 - 1, y2) - get(x2, y1 - 1)
      + get(x1 - 1, v1 - 1);
```

#### 3.8 Divide And Conquer Query Offline

```
namespace up {
int 1[N], r[N], u[N], v[N], tm;
void push(int _l, int _r, int _u, int _v) {
 1[tm] = _1, r[tm] = _r, u[tm] = _u, v[tm] = _v;
  tm++;
} // namespace up
namespace que {
int node[N], tm;
LL ans[N]:
} // namespace que
namespace edge_set {
```

```
void push(int i) { dsu ::merge(up ::u[i], up ::v[i]); }
void pop(int t) { dsu ::rollback(t); }
int time() { return dsu ::op.size(); }
LL query(int u) { return a[dsu ::root(u)]; }
} // namespace edge_set
namespace dncq {
vector<int> tree[4 * N];
void update(int idx, int 1 = 0, int r = que ::tm, int
    node = 1) {
 int ul = up ::1[idx], ur = up ::r[idx];
 if (r 
 if (ul <= l and r <= ur) {</pre>
   tree[node].push_back(idx);
   return;
 int m = 1 + r >> 1;
 update(idx, 1, m, node << 1);
 update(idx, m + 1, r, node << 1 | 1);
void dfs(int l = 0, int r = que ::tm, int node = 1) {
 int cur = edge_set ::time();
 for (int e : tree[node]) edge_set ::push(e);
 if (1 == r) {
   que ::ans[1] = edge_set ::query(que ::node[1]);
 } else {
   int m = 1 + r >> 1;
   dfs(1, m, node << 1);
   dfs(m + 1, r, node << 1 | 1);
 edge_set ::pop(cur);
} // namespace dncq
void push_updates() {
 for (int i = 0; i < up ::tm; i++) dncq ::update(i);</pre>
```

#### 3.9 MO with Update

```
const int N = 1e5 + 5, sz = 2700, bs = 25;
int arr[N], freq[2 * N], cnt[2 * N], id[N], ans[N];
struct query {
 int 1, r, t, L, R;
 query(int l = 1, int r = 0, int t = 1, int id = -1)
     : 1(1), r(r), t(t), L(1 / sz), R(r / sz) {}
 bool operator<(const query &rhs) const {</pre>
   return (L < rhs.L) or (L == rhs.L and R < rhs.R) or
          (L == rhs.L and R == rhs.R and t < rhs.t);
 }
} Q[N];
struct update {
 int idx, val, last;
```

```
} Up[N];
int qi = 0, ui = 0;
int 1 = 1, r = 0, t = 0;
void add(int idx) {
  --cnt[freq[arr[idx]]];
  freq[arr[idx]]++;
  cnt[freq[arr[idx]]]++;
void remove(int idx) {
  --cnt[freq[arr[idx]]];
  freq[arr[idx]]--;
  cnt[freq[arr[idx]]]++;
void apply(int t) {
  const bool f = 1 <= Up[t].idx and Up[t].idx <= r;</pre>
  if (f) remove(Up[t].idx);
  arr[Up[t].idx] = Up[t].val;
  if (f) add(Up[t].idx);
void undo(int t) {
  const bool f = 1 <= Up[t].idx and Up[t].idx <= r;</pre>
  if (f) remove(Up[t].idx);
  arr[Up[t].idx] = Up[t].last;
  if (f) add(Up[t].idx);
int mex() {
  for (int i = 1; i <= N; i++)</pre>
   if (!cnt[i]) return i;
  assert(0):
int main() {
  int n, q;
  cin >> n >> q;
  int counter = 0;
  map<int, int> M;
  for (int i = 1; i <= n; i++) {</pre>
    cin >> arr[i]:
   if (!M[arr[i]]) M[arr[i]] = ++counter;
    arr[i] = M[arr[i]];
  iota(id, id + N, 0);
  while (q--) {
   int tp, x, y;
    cin >> tp >> x >> y;
   if (tp == 1)
     Q[++qi] = query(x, y, ui);
    else {
     if (!M[y]) M[y] = ++counter;
     y = M[y];
```

```
Up[++ui] = \{x, y, arr[x]\};
    arr[x] = v;
 }
}
t = ui;
cnt[0] = 3 * n;
sort(id + 1, id + qi + 1, [&](int x, int y) { return Q[
    x] < Q[v]; \});
for (int i = 1; i <= qi; i++) {</pre>
  int x = id[i];
  while (Q[x].t > t) apply(++t);
  while (Q[x].t < t) undo(t--);
  while (Q[x].1 < 1) add(--1);
  while (Q[x].r > r) add(++r);
  while (Q[x].1 > 1) remove(1++);
  while (Q[x].r < r) remove(r--);
  ans[x] = mex():
for (int i = 1; i <= qi; i++) cout << ans[i] << '\n';</pre>
```

### 3.10 SparseTable (Rectangle Query)

```
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 505;
const int LOGN = 9;
// O(n^2 (logn)^2
// Supports Rectangular Query
int A[MAXN][MAXN];
int M[MAXN] [MAXN] [LOGN] [LOGN];
void Build2DSparse(int N) {
 for (int i = 1; i <= N; i++) {</pre>
   for (int j = 1; j <= N; j++) {</pre>
     M[i][j][0][0] = A[i][j];
   for (int q = 1; (1 << q) <= N; q++) {
     int add = 1 << (q - 1);
     for (int j = 1; j + add \le N; j++) {
       M[i][j][0][q] = max(M[i][j][0][q - 1], M[i][j +
            add][0][q - 1]);
     }
   }
 }
 for (int p = 1; (1 << p) <= N; p++) {
   int add = 1 << (p - 1);
   for (int i = 1; i + add <= N; i++) {</pre>
```

```
for (int q = 0; (1 << q) <= N; q++) {
      for (int j = 1; j <= N; j++) {</pre>
        M[i][j][p][q] = max(M[i][j][p - 1][q], M[i + add
            ][i][p - 1][q]);
      }
    }
// returns max of all A[i][j], where x1 \le i \le x2 and y1 \le j
int Query(int x1, int y1, int x2, int y2) {
 int kX = log2(x2 - x1 + 1);
 int kY = log2(y2 - y1 + 1);
 int addX = 1 << kX;
 int addY = 1 << kY:
 int ret1 = max(M[x1][y1][kX][kY], M[x1][y2 - addY + 1][
     kX][kY]):
 int ret2 = max(M[x2 - addX + 1][y1][kX][kY],
               M[x2 - addX + 1][y2 - addY + 1][kX][kY]);
 return max(ret1, ret2);
```

## 4 Geometry

## 4.1 Point

```
typedef double Tf;
typedef double Ti; /// use long long for exactness
const Tf PI = acos(-1), EPS = 1e-9;
int dcmp(Tf x) { return abs(x) < EPS ? 0 : (x < 0 ? -1 :
    1); }
struct Point {
 Ti x, y;
 Point(Ti x = 0, Ti y = 0) : x(x), y(y) {}
 Point operator+(const Point& u) const { return Point(x
     + u.x, y + u.y); }
 Point operator-(const Point& u) const { return Point(x
     -u.x, y - u.y); }
 Point operator*(const LL u) const { return Point(x * u,
 Point operator*(const Tf u) const { return Point(x * u,
      y * u); }
 Point operator/(const Tf u) const { return Point(x / u,
      y / u); }
 bool operator==(const Point& u) const {
   return dcmp(x - u.x) == 0 && dcmp(y - u.y) == 0;
```

```
bool operator!=(const Point& u) const { return !(*this
 bool operator<(const Point& u) const {</pre>
   return dcmp(x - u.x) < 0 \mid \mid (dcmp(x - u.x) == 0 \&\&
        dcmp(y - u.y) < 0);
 }
Ti dot(Point a, Point b) { return a.x * b.x + a.y * b.y;
Ti cross(Point a, Point b) { return a.x * b.y - a.y * b.x
Tf length(Point a) { return sqrt(dot(a, a)); }
Ti sqLength(Point a) { return dot(a, a); }
Tf distance(Point a, Point b) { return length(a - b); }
Tf angle(Point u) { return atan2(u.v, u.x); }
// returns angle between oa, ob in (-PI, PI]
Tf angleBetween(Point a, Point b) {
 Tf ans = angle(b) - angle(a);
 return ans <= -PI ? ans + 2 * PI : (ans > PI ? ans - 2
      * PI : ans):
// Rotate a ccw by rad radians, Tf Ti same
Point rotate(Point a. Tf rad) {
 return Point(a.x * cos(rad) - a.y * sin(rad),
             a.x * sin(rad) + a.y * cos(rad));
// rotate a ccw by angle th with cos(th) = co && sin(th)
    = si. tf ti same
Point rotatePrecise(Point a, Tf co, Tf si) {
 return Point(a.x * co - a.y * si, a.y * co + a.x * si);
Point rotate90(Point a) { return Point(-a.y, a.x); }
// scales vector a by s such that length of a becomes s,
    Tf Ti same
Point scale(Point a, Tf s) { return a / length(a) * s; }
// returns an unit vector perpendicular to vector a, Tf
    Ti same
Point normal(Point a) {
 Tf l = length(a);
 return Point(-a.y / 1, a.x / 1);
// returns 1 if c is left of ab, 0 if on ab && -1 if
int orient(Point a, Point b, Point c) { return dcmp(cross
    (b - a, c - a)); }
/// Use as sort(v.begin(), v.end(), polarComp(0, dir))
/// Polar comparator around O starting at direction dir
struct polarComp {
```

```
Point O, dir;
 polarComp(Point 0 = Point(0, 0), Point dir = Point(1,
      0)) : O(O), dir(dir) {}
 bool half(Point p) {
   return dcmp(cross(dir, p)) < 0 ||</pre>
          (dcmp(cross(dir, p)) == 0 && dcmp(dot(dir, p))
 bool operator()(Point p, Point q) {
   return make_tuple(half(p), 0) < make_tuple(half(q),</pre>
       cross(p, q));
 }
struct Segment {
 Point a, b;
 Segment(Point aa, Point bb) : a(aa), b(bb) {}
typedef Segment Line;
struct Circle {
 Point o:
 Tf r;
 Circle(Point o = Point(0, 0), Tf r = 0) : o(o), r(r) {}
 // returns true if point p is in || on the circle
 bool contains(Point p) { return dcmp(sqLength(p - o) -
      r * r) <= 0: 
 // returns a point on the circle rad radians away from
      +X CCW
 Point point(Tf rad) {
   static_assert(is_same<Tf, Ti>::value);
   return Point(o.x + cos(rad) * r, o.y + sin(rad) * r);
 // area of a circular sector with central angle rad
 Tf area(Tf rad = PI + PI) { return rad * r * r / 2; }
 // area of the circular sector cut by a chord with
      central angle alpha
 Tf sector(Tf alpha) { return r * r * 0.5 * (alpha - sin | // **** LINE LINE INTERSECTION FINISH ****
      (alpha)); }
4.2 Linear
```

```
// **** LINE LINE INTERSECTION START ****
// returns true if point p is on segment s
bool onSegment(Point p, Segment s) {
 return dcmp(cross(s.a - p, s.b - p)) == 0 && dcmp(dot(s
      .a - p, s.b - p)) <= 0;
// returns true if segment p && q touch or intersect
bool segmentsIntersect(Segment p, Segment q) {
 if (onSegment(p.a, q) || onSegment(p.b, q)) return true
```

```
if (onSegment(q.a, p) || onSegment(q.b, p)) return true
 Ti c1 = cross(p.b - p.a, q.a - p.a);
 Ti c2 = cross(p.b - p.a, q.b - p.a);
 Ti c3 = cross(q.b - q.a, p.a - q.a);
 Ti c4 = cross(q.b - q.a, p.b - q.a);
 return dcmp(c1) * dcmp(c2) < 0 && dcmp(c3) * dcmp(c4) <
bool linesParallel(Line p, Line q) {
 return dcmp(cross(p.b - p.a, q.b - q.a)) == 0;
// lines are represented as a ray from a point: (point,
// returns false if two lines (p, v) && (q, w) are
    parallel or collinear
// true otherwise, intersection point is stored at o via
    reference, Tf Ti Same
bool lineLineIntersection(Point p, Point v, Point q,
    Point w, Point& o) {
 if (dcmp(cross(v, w)) == 0) return false;
 Point u = p - q;
 o = p + v * (cross(w, u) / cross(v, w));
 return true:
// returns false if two lines p && q are parallel or
// true otherwise, intersection point is stored at o via
    reference
bool lineLineIntersection(Line p, Line q, Point& o) {
 return lineLineIntersection(p.a, p.b - p.a, q.a, q.b -
      q.a, o);
// returns the distance from point a to line 1
Tf distancePointLine(Point p, Line 1) {
 return abs(cross(l.b - l.a, p - l.a) / length(l.b - l.a
// returns the shortest distance from point a to segment
Tf distancePointSegment(Point p, Segment s) {
 if (s.a == s.b) return length(p - s.a);
 Point v1 = s.b - s.a, v2 = p - s.a, v3 = p - s.b;
 if (dcmp(dot(v1, v2)) < 0)
   return length(v2);
 else if (dcmp(dot(v1, v3)) > 0)
   return length(v3);
 else
```

```
return abs(cross(v1, v2) / length(v1));
}
// returns the shortest distance from segment p to
    segment q

Tf distanceSegmentSegment(Segment p, Segment q) {
    if (segmentsIntersect(p, q)) return 0;
    Tf ans = distancePointSegment(p.a, q);
    ans = min(ans, distancePointSegment(p.b, q));
    ans = min(ans, distancePointSegment(q.a, p));
    ans = min(ans, distancePointSegment(q.b, p));
    return ans;
}
// returns the projection of point p on line 1, Tf Ti
    Same
Point projectPointLine(Point p, Line 1) {
    Point v = 1.b - 1.a;
    return 1.a + v * ((Tf)dot(v, p - 1.a) / dot(v, v));
}
```

#### 4.3 Circular

```
// Extremely inaccurate for finding near touches
// compute intersection of line 1 with circle c
// The intersections are given in order of the ray (l.a,
    1.b), Tf Ti same
vector<Point> circleLineIntersection(Circle c, Line 1) {
  vector<Point> ret;
 Point b = 1.b - 1.a, a = 1.a - c.o;
 Tf A = dot(b, b), B = dot(a, b);
 Tf C = dot(a, a) - c.r * c.r, D = B * B - A * C;
 if (D < -EPS) return ret;</pre>
 ret.push_back(1.a + b * (-B - sqrt(D + EPS)) / A);
  if (D > EPS) ret.push_back(1.a + b * (-B + sqrt(D)) / A
      );
 return ret;
// signed area of intersection of circle(c.o, c.r) &&
// triangle(c.o, s.a, s.b) [cross(a-o, b-o)/2]
Tf circleTriangleIntersectionArea(Circle c, Segment s) {
  using Linear::distancePointSegment;
 Tf OA = length(c.o - s.a);
 Tf OB = length(c.o - s.b);
 // sector
  if (dcmp(distancePointSegment(c.o, s) - c.r) >= 0)
   return angleBetween(s.a - c.o, s.b - c.o) * (c.r * c.
       r) / 2.0;
  // triangle
  if (dcmp(OA - c.r) \le 0 \&\& dcmp(OB - c.r) \le 0)
   return cross(c.o - s.b, s.a - s.b) / 2.0;
  // three part: (A, a) (a, b) (b, B)
 vector<Point> Sect = circleLineIntersection(c, s);
```

```
return circleTriangleIntersectionArea(c, Segment(s.a,
      Sect[0])) +
        circleTriangleIntersectionArea(c, Segment(Sect
            [0]. Sect[1])) +
        circleTriangleIntersectionArea(c, Segment(Sect
            [1], s.b));
// area of intersecion of circle(c.o, c.r) && simple
    polyson(p[])
Tf circlePolyIntersectionArea(Circle c, Polygon p) {
 Tf res = 0:
 int n = p.size();
 for (int i = 0; i < n; ++i)</pre>
   res += circleTriangleIntersectionArea(c, Segment(p[i
       ], p[(i + 1) % n]));
 return abs(res);
// locates circle c2 relative to c1
// interior
                     (d < R - r)
// interior tangents (d = R - r)
// concentric
                   (d = 0)
// secants
                     (R - r < d < R + r) \longrightarrow 0
// exterior tangents (d = R + r)
// exterior
                     (d > R + r)
                                        ----> 2
int circleCirclePosition(Circle c1, Circle c2) {
 Tf d = length(c1.o - c2.o);
 int in = dcmp(d - abs(c1.r - c2.r)), ex = dcmp(d - (c1.r))
      r + c2.r):
 return in < 0 ? -2 : in == 0 ? -1 : ex == 0 ? 1 : ex >
      0 ? 2 : 0:
// compute the intersection points between two circles c1
     && c2, Tf Ti same
vector<Point> circleCircleIntersection(Circle c1, Circle
    c2) {
 vector<Point> ret;
 Tf d = length(c1.o - c2.o);
 if (dcmp(d) == 0) return ret;
 if (dcmp(c1.r + c2.r - d) < 0) return ret;
 if (dcmp(abs(c1.r - c2.r) - d) > 0) return ret;
 Point v = c2.0 - c1.0;
 Tf co = (c1.r * c1.r + sqLength(v) - c2.r * c2.r) / (2)
      * c1.r * length(v));
 Tf si = sqrt(abs(1.0 - co * co));
 Point p1 = scale(rotatePrecise(v, co, -si), c1.r) + c1.
 Point p2 = scale(rotatePrecise(v, co, si), c1.r) + c1.o
```

```
ret.push_back(p1);
 if (p1 != p2) ret.push_back(p2);
 return ret;
// intersection area between two circles c1, c2
Tf circleCircleIntersectionArea(Circle c1, Circle c2) {
 Point AB = c2.0 - c1.0:
 Tf d = length(AB);
 if (d \ge c1.r + c2.r) return 0;
 if (d + c1.r <= c2.r) return PI * c1.r * c1.r;</pre>
 if (d + c2.r <= c1.r) return PI * c2.r * c2.r;</pre>
 Tf alpha1 = acos((c1.r * c1.r + d * d - c2.r * c2.r) /
     (2.0 * c1.r * d));
 Tf alpha2 = acos((c2.r * c2.r + d * d - c1.r * c1.r) /
     (2.0 * c2.r * d));
 return c1.sector(2 * alpha1) + c2.sector(2 * alpha2);
// returns tangents from a point p to circle c, Tf Ti
vector<Point> pointCircleTangents(Point p, Circle c) {
 vector<Point> ret;
 Point u = c.o - p;
 Tf d = length(u);
 if (d < c.r)
 else if (dcmp(d - c.r) == 0) {
   ret = {rotate(u, PI / 2)};
 } else {
   Tf ang = asin(c.r / d);
   ret = {rotate(u, -ang), rotate(u, ang)};
 return ret;
// returns the points on tangents that touches the circle
    , Tf Ti Same
vector<Point> pointCircleTangencyPoints(Point p, Circle c
   ) {
 Point u = p - c.o;
 Tf d = length(u);
 if (d < c.r)
   return {};
 else if (dcmp(d - c.r) == 0)
   return {c.o + u};
 else {
   Tf ang = acos(c.r / d);
   u = u / length(u) * c.r;
   return {c.o + rotate(u, -ang), c.o + rotate(u, ang)};
```

```
// for two circles c1 && c2, returns two list of points a }
// such that a[i] is on c1 && b[i] is c2 && for every i
// Line(a[i], b[i]) is a tangent to both circles
// CAUTION: a[i] = b[i] in case they touch | -1 for c1 =
int circleCircleTangencyPoints(Circle c1, Circle c2,
    vector<Point> &a,
                            vector<Point> &b) {
  a.clear(), b.clear();
  int cnt = 0;
  if (dcmp(c1.r - c2.r) < 0) {
   swap(c1, c2);
   swap(a, b);
 Tf d2 = sqLength(c1.o - c2.o);
 Tf rdif = c1.r - c2.r, rsum = c1.r + c2.r;
  if (dcmp(d2 - rdif * rdif) < 0) return 0;</pre>
  if (dcmp(d2) == 0 \&\& dcmp(c1.r - c2.r) == 0) return -1;
 Tf base = angle(c2.o - c1.o);
  if (dcmp(d2 - rdif * rdif) == 0) {
   a.push_back(c1.point(base));
   b.push_back(c2.point(base));
   cnt++:
   return cnt;
 Tf ang = acos((c1.r - c2.r) / sqrt(d2));
  a.push_back(c1.point(base + ang));
  b.push_back(c2.point(base + ang));
  a.push_back(c1.point(base - ang));
  b.push_back(c2.point(base - ang));
  cnt++;
  if (dcmp(d2 - rsum * rsum) == 0) {
   a.push_back(c1.point(base));
   b.push_back(c2.point(PI + base));
   cnt++:
  } else if (dcmp(d2 - rsum * rsum) > 0) {
   Tf ang = acos((c1.r + c2.r) / sqrt(d2));
   a.push_back(c1.point(base + ang));
   b.push_back(c2.point(PI + base + ang));
   a.push_back(c1.point(base - ang));
   b.push_back(c2.point(PI + base - ang));
   cnt++;
 return cnt;
```

```
4.4 Convex
/// minkowski sum of two polygons in O(n)
Polygon minkowskiSum(Polygon A, Polygon B) {
 int n = A.size(), m = B.size();
 rotate(A.begin(), min_element(A.begin(), A.end()), A.
      end()):
 rotate(B.begin(), min_element(B.begin(), B.end()), B.
      end()):
  A.push_back(A[0]);
  B.push_back(B[0]);
  for (int i = 0; i < n; i++) A[i] = A[i + 1] - A[i];
 for (int i = 0; i < m; i++) B[i] = B[i + 1] - B[i];
 Polygon C(n + m + 1);
  C[0] = A.back() + B.back();
  merge(A.begin(), A.end() - 1, B.begin(), B.end() - 1, C
      .begin() + 1,
       polarComp(Point(0, 0), Point(0, -1)));
 for (int i = 1; i < C.size(); i++) C[i] = C[i] + C[i -</pre>
 C.pop_back();
 return C;
// finds the rectangle with minimum area enclosing a
    convex polygon and
// the rectangle with minimum perimeter enclosing a
    convex polygon
// Tf Ti Same
pair<Tf, Tf> rotatingCalipersBoundingBox(const Polygon &p
 using Linear::distancePointLine;
 int n = p.size();
 int 1 = 1, r = 1, j = 1;
 Tf area = 1e100;
 Tf perimeter = 1e100;
 for (int i = 0; i < n; i++) {</pre>
   Point v = (p[(i + 1) \% n] - p[i]) / length(p[(i + 1)
        % n] - p[i]);
   while (dcmp(dot(v, p[r \% n] - p[i]) - dot(v, p[(r +
        1) \% n] - p[i])) < 0)
     r++:
   while (j < r \mid | dcmp(cross(v, p[j % n] - p[i]) -
                       cross(v, p[(j + 1) % n] - p[i])) <
                            0)
     j++;
   while (1 < j ||
```

```
dcmp(dot(v, p[1 \% n] - p[i]) - dot(v, p[(1 + 1)
               n - p[i] > 0
     1++:
   Tf w = dot(v, p[r \% n] - p[i]) - dot(v, p[1 \% n] - p[
   Tf h = distancePointLine(p[j % n], Line(p[i], p[(i +
       1) % n]));
   area = min(area, w * h);
   perimeter = min(perimeter, 2 * w + 2 * h);
 }
 return make_pair(area, perimeter);
// returns the left side of polygon u after cutting it by
     ray a->b
Polygon cutPolygon(Polygon u, Point a, Point b) {
 using Linear::lineLineIntersection;
 using Linear::onSegment;
 Polygon ret;
 int n = u.size();
 for (int i = 0; i < n; i++) {</pre>
   Point c = u[i], d = u[(i + 1) \% n];
   if (dcmp(cross(b - a, c - a)) >= 0) ret.push_back(c);
   if (dcmp(cross(b - a, d - c)) != 0) {
     Point t:
     lineLineIntersection(a, b - a, c, d - c, t);
     if (onSegment(t, Segment(c, d))) ret.push_back(t);
 }
 return ret:
// returns true if point p is in or on triangle abc
bool pointInTriangle(Point a, Point b, Point c, Point p)
 return dcmp(cross(b - a, p - a)) >= 0 && dcmp(cross(c -
      b, p - b)) >= 0 &&
       dcmp(cross(a - c, p - c)) >= 0;
// pt must be in ccw order with no three collinear points
// returns inside = -1, on = 0, outside = 1
int pointInConvexPolygon(const Polygon &pt, Point p) {
 int n = pt.size();
 assert(n >= 3);
 int lo = 1, hi = n - 1;
 while (hi - lo > 1) {
   int mid = (lo + hi) / 2;
   if (dcmp(cross(pt[mid] - pt[0], p - pt[0])) > 0)
     lo = mid;
   else
```

```
hi = mid:
 bool in = pointInTriangle(pt[0], pt[lo], pt[hi], p);
 if (!in) return 1;
 if (dcmp(cross(pt[lo] - pt[lo - 1], p - pt[lo - 1])) == |}
       0) return 0;
 if (dcmp(cross(pt[hi] - pt[lo], p - pt[lo])) == 0)
      return 0;
 if (dcmp(cross(pt[hi] - pt[(hi + 1) % n], p - pt[(hi +
      1) % n)) == 0)
   return 0:
 return -1;
// Extreme Point for a direction is the farthest point in
     that direction
// u is the direction for extremeness
int extremePoint(const Polygon &poly, Point u) {
 int n = (int)poly.size();
 int a = 0, b = n;
 while (b - a > 1) {
   int c = (a + b) / 2;
   if (dcmp(dot(poly[c] - poly[(c + 1) % n], u)) >= 0 &&
       dcmp(dot(poly[c] - poly[(c - 1 + n) \% n], u)) >=
           0) {
     return c;
   }
   bool a_up = dcmp(dot(poly[(a + 1) % n] - poly[a], u))
   bool c_{up} = dcmp(dot(poly[(c + 1) % n] - poly[c], u))
   bool a_above_c = dcmp(dot(poly[a] - poly[c], u)) > 0;
   if (a_up && !c_up)
     b = c;
   else if (!a_up && c_up)
     a = c;
   else if (a_up && c_up) {
     if (a_above_c)
      b = c;
     else
       a = c;
   } else {
     if (!a_above_c)
      b = c;
     else
       a = c;
```

```
}
 if (dcmp(dot(poly[a] - poly[(a + 1) % n], u)) > 0 &&
     dcmp(dot(poly[a] - poly[(a - 1 + n) % n], u)) > 0)
   return a:
 return b % n;
// For a convex polygon p and a line 1, returns a list of
// of p that touch or intersect line 1.
// the i'th segment is considered (p[i], p[(i + 1) modulo
// #1 If a segment is collinear with the line, only that
    is returned
// #2 Else if 1 goes through i'th point, the i'th segment }
// Complexity: O(lg |p|)
vector<int> lineConvexPolyIntersection(const Polygon &p,
    Line 1) {
 assert((int)p.size() >= 3);
 assert(1.a != 1.b);
 int n = p.size();
 vector<int> ret;
 Point v = 1.b - 1.a;
 int lf = extremePoint(p, rotate90(v));
 int rt = extremePoint(p, rotate90(v) * Ti(-1));
 int olf = orient(l.a, l.b, p[lf]);
 int ort = orient(1.a, 1.b, p[rt]);
 if (!olf || !ort) {
   int idx = (!olf ? lf : rt);
   if (orient(1.a, 1.b, p[(idx - 1 + n) \% n]) == 0)
     ret.push_back((idx - 1 + n) \% n);
   else
     ret.push_back(idx);
   return ret;
 if (olf == ort) return ret;
 for (int i = 0; i < 2; ++i) {</pre>
   int lo = i ? rt : lf;
   int hi = i ? lf : rt;
   int olo = i ? ort : olf;
   while (true) {
     int gap = (hi - lo + n) \% n;
     if (gap < 2) break;</pre>
```

```
int mid = (lo + gap / 2) \% n;
     int omid = orient(l.a, l.b, p[mid]);
     if (!omid) {
       lo = mid:
       break;
     if (omid == olo)
       lo = mid;
     else
       hi = mid;
   ret.push_back(lo);
 return ret;
// Calculate [ACW, CW] tangent pair from an external
    point
constexpr int CW = -1, ACW = 1;
bool isGood(Point u, Point v, Point Q, int dir) {
 return orient(Q, u, v) != -dir;
Point better(Point u, Point v, Point Q, int dir) {
 return orient(Q, u, v) == dir ? u : v;
Point pointPolyTangent(const Polygon &pt, Point Q, int
    dir, int lo, int hi) {
 while (hi - lo > 1) {
   int mid = (lo + hi) / 2;
   bool pvs = isGood(pt[mid], pt[mid - 1], Q, dir);
   bool nxt = isGood(pt[mid], pt[mid + 1], Q, dir);
   if (pvs && nxt) return pt[mid];
   if (!(pvs || nxt)) {
     Point p1 = pointPolyTangent(pt, Q, dir, mid + 1, hi
     Point p2 = pointPolyTangent(pt, Q, dir, lo, mid -
     return better(p1, p2, Q, dir);
   }
   if (!pvs) {
     if (orient(Q, pt[mid], pt[lo]) == dir)
       hi = mid - 1:
     else if (better(pt[lo], pt[hi], Q, dir) == pt[lo])
       hi = mid - 1;
     else
       lo = mid + 1;
   }
   if (!nxt) {
     if (orient(Q, pt[mid], pt[lo]) == dir)
```

```
lo = mid + 1:
     else if (better(pt[lo], pt[hi], Q, dir) == pt[lo])
       hi = mid - 1;
     else
       lo = mid + 1;
   }
 }
 Point ret = pt[lo];
 for (int i = lo + 1; i <= hi; i++) ret = better(ret, pt</pre>
      [i], Q, dir);
 return ret;
// [ACW, CW] Tangent
pair<Point, Point> pointPolyTangents(const Polygon &pt,
    Point Q) {
 int n = pt.size();
  Point acw_tan = pointPolyTangent(pt, Q, ACW, 0, n - 1);
 Point cw_tan = pointPolyTangent(pt, Q, CW, 0, n - 1);
 return make_pair(acw_tan, cw_tan);
```

## 4.5 Polygon

```
typedef vector<Point> Polygon;
// removes redundant colinear points
// polygon can't be all colinear points
Polygon RemoveCollinear(const Polygon &poly) {
 Polygon ret;
 int n = poly.size();
 for (int i = 0; i < n; i++) {</pre>
   Point a = polv[i];
   Point b = poly[(i + 1) \% n];
   Point c = poly[(i + 2) \% n];
   if (dcmp(cross(b - a, c - b)) != 0 && (ret.empty() ||
         b != ret.back()))
     ret.push_back(b);
 }
 return ret;
// returns the signed area of polygon p of n vertices
Tf signedPolygonArea(const Polygon &p) {
 Tf ret = 0;
 for (int i = 0; i < (int)p.size() - 1; i++)</pre>
   ret += cross(p[i] - p[0], p[i + 1] - p[0]);
 return ret / 2:
// given a polygon p of n vertices, generates the convex
    hull in in CCW
// Tested on https://acm.timus.ru/problem.aspx?space=1&
    num=1185
```

```
// Caution: when all points are colinear AND
    removeRedundant == false
// output will be contain duplicate points (from upper
    hull) at back
Polygon convexHull(Polygon p, bool removeRedundant) {
 int check = removeRedundant ? 0 : -1;
 sort(p.begin(), p.end());
 p.erase(unique(p.begin(), p.end()), p.end());
 int n = p.size();
 Polygon ch(n + n);
 int m = 0; // preparing lower hull
 for (int i = 0; i < n; i++) {</pre>
   while (m > 1 &&
          dcmp(cross(ch[m-1]-ch[m-2], p[i]-ch[m
               1])) <= check)
     m--:
   ch[m++] = p[i];
 int k = m; // preparing upper hull
 for (int i = n - 2; i \ge 0; i--) {
   while (m > k &&
          dcmp(cross(ch[m-1]-ch[m-2], p[i]-ch[m
               2])) <= check)
     m--:
   ch[m++] = p[i];
 if (n > 1) m--:
 ch.resize(m);
 return ch:
// returns inside = -1, on = 0, outside = 1
int pointInPolygon(const Polygon &p, Point o) {
 using Linear::onSegment;
 int wn = 0, n = p.size();
 for (int i = 0; i < n; i++) {</pre>
   int j = (i + 1) \% n;
   if (onSegment(o, Segment(p[i], p[j])) || o == p[i])
   int k = dcmp(cross(p[j] - p[i], o - p[i]));
   int d1 = dcmp(p[i].y - o.y);
   int d2 = dcmp(p[j].y - o.y);
   if (k > 0 \&\& d1 \le 0 \&\& d2 > 0) wn++;
   if (k < 0 && d2 <= 0 && d1 > 0) wn--;
 return wn ? -1 : 1:
// Given a simple polygon p, and a line l, returns (x, y)
// x = longest segment of 1 in p, y = total length of 1
    in p.
```

```
pair<Tf, Tf> linePolygonIntersection(Line 1, const
    Polygon &p) {
 using Linear::lineLineIntersection;
 int n = p.size();
 vector<pair<Tf, int>> ev;
 for (int i = 0; i < n; ++i) {</pre>
   Point a = p[i], b = p[(i + 1) \% n], z = p[(i - 1 + n)
        % n];
   int ora = orient(l.a, l.b, a), orb = orient(l.a, l.b,
       orz = orient(l.a, l.b, z);
   if (!ora) {
     Tf d = dot(a - 1.a, 1.b - 1.a);
     if (orz && orb) {
       if (orz != orb) ev.emplace_back(d, 0);
       // else // Point Touch
     } else if (orz)
       ev.emplace_back(d, orz);
     else if (orb)
       ev.emplace_back(d, orb);
   } else if (ora == -orb) {
     Point ins:
     lineLineIntersection(1, Line(a, b), ins);
     ev.emplace_back(dot(ins - 1.a, 1.b - 1.a), 0);
 }
 sort(ev.begin(), ev.end());
 Tf ans = 0, len = 0, last = 0, tot = 0;
 bool active = false:
 int sign = 0;
 for (auto &qq : ev) {
   int tp = qq.second;
   Tf d = qq.first; /// current Segment is (last, d)
   if (sign) {
                   /// On Border
     len += d - last;
     tot += d - last;
     ans = max(ans, len);
     if (tp != sign) active = !active;
     sign = 0;
   } else {
     if (active) { /// Strictly Inside
       len += d - last;
       tot += d - last;
       ans = max(ans, len);
     if (tp == 0)
       active = !active;
       sign = tp;
```

```
}
  last = d;
  if (!active) len = 0;
}
  ans /= length(l.b - l.a);
  tot /= length(l.b - l.a);
  return {ans, tot};
}
```

## 5 Graph

## 5.1 Graph Template

```
struct edge {
 int u, v;
 edge(int u = 0, int v = 0) : u(u), v(v) {}
 int to(int node) { return u ^ v ^ node; }
struct graph {
 int n;
 vector<vector<int>> adj;
 vector<edge> edges;
 graph(int n = 0) : n(n), adj(n) {}
 void addEdge(int u, int v, bool dir = 1) {
   adj[u].push_back(edges.size());
   if (dir) adj[v].push_back(edges.size());
   edges.emplace_back(u, v);
 int addNode() {
   adj.emplace_back();
   return n++;
 edge &operator()(int idx) { return edges[idx]; }
 vector<int> &operator[](int u) { return adj[u]; }
```

## 5.2 LCA

```
}
int lca(int u. int v) {
 if (level[u] > level[v]) swap(u, v);
 for (int k=K-1; k>=0; k--)
   if (level[u] + (1<<k) <= level[v]) v = anc[v][k];</pre>
 if (u == v) return u:
 for (int k=K-1; k>=0; k--)
   if (anc[u][k] != anc[v][k])
     u = anc[u][k], v = anc[v][k];
 return anc[u][0];
int getanc(int u, int d) {
 for (int k=0; k<K; k++)</pre>
   if (d & (1<<k))
     u = anc[u][k];
 return u;
int dist(int u, int v) {
 int g = lca(u, v);
 return level[u] + level[v] - 2*level[g];
5.3 SCC
```

```
typedef long long LL;
const LL N = 1e6 + 7;
bool vis[N];
vector<int> adj[N], adjr[N];
vector<int> order, component;
// tp = 0 ,finding topo order, tp = 1 , reverse edge
    traversal
void dfs(int u, int tp = 0) {
 vis[u] = true;
 if (tp) component.push_back(u);
 auto& ad = (tp ? adjr : adj);
 for (int v : ad[u])
   if (!vis[v]) dfs(v, tp);
 if (!tp) order.push_back(u);
int main() {
 for (int i = 1; i <= n; i++) {</pre>
   if (!vis[i]) dfs(i);
 memset(vis, 0, sizeof vis);
```

```
reverse(order.begin(), order.end());
for (int i : order) {
   if (!vis[i]) {
      // one component is found
      dfs(i, 1), component.clear();
   }
}
```

## 5.4 Euler Tour on Edge

```
// for simplicity, G[idx] contains the adjacency list of
    a node
// while G(e) is a reference to the e-th edge.
const int N = 2e5 + 5;
int in[N], out[N], fwd[N], bck[N];
int t = 0:
void dfs(graph &G, int node, int par) {
 out[node] = t:
 for (int e : G[node]) {
   int v = G(e).to(node);
   if (v == par) continue;
   fwd[e] = t++;
   dfs(G, v, node);
   bck[e] = t++:
 in[node] = t - 1;
void init(graph &G, int node) {
 t = 0;
 dfs(G, node, node);
```

## 5.5 LCA In O(1)

```
/* LCA in O(1)

* depth calculates weighted distance

* level calculates distance by number of edges

* Preprocessing in NlongN */

LL depth[N];

int level[N];

int st[N], en[N], LOG[N], par[N];

int a[N], id[N], table[L][N];

vector<PII> adj[N];

int n, root, Time, cur;

void init(int nodes, int root_) {

n = nodes, root = root_, LOG[0] = LOG[1] = 0;

for (int i = 2; i <= n; i++) LOG[i] = LOG[i >> 1] + 1;

for (int i = 0; i <= n; i++) adj[i].clear();
```

```
}
void addEdge(int u, int v, int w) {
  adj[u].push_back(PII(v, w));
 adj[v].push_back(PII(u, w));
int lca(int u, int v) {
  if (en[u] > en[v]) swap(u, v);
  if (st[v] <= st[u] && en[u] <= en[v]) return v;</pre>
  int 1 = LOG[id[v] - id[u] + 1];
  int p1 = id[u], p2 = id[v] - (1 << 1) + 1;</pre>
  int d1 = level[table[1][p1]], d2 = level[table[1][p2]];
  if (d1 < d2)
   return par[table[1][p1]];
   return par[table[1][p2]];
LL dist(int u. int v) {
 int 1 = lca(u, v);
 return (depth[u] + depth[v] - (depth[1] * 2));
/* Euler tour */
void dfs(int u, int p) {
  st[u] = ++Time, par[u] = p;
  for (auto [v, w] : adi[u]) {
   if (v == p) continue;
   depth[v] = depth[u] + w;
   level[v] = level[u] + 1;
   dfs(v, u);
  en[u] = ++Time:
 a[++cur] = u, id[u] = cur;
/* RMQ */
void pre() {
  cur = Time = 0, dfs(root, root);
 for (int i = 1; i <= n; i++) table[0][i] = a[i];</pre>
 for (int 1 = 0; 1 < L - 1; 1++) {
   for (int i = 1; i <= n; i++) {
     table[1 + 1][i] = table[1][i];
```

```
bool C1 = (1 << 1) + i <= n;
     bool C2 = level[table[1][i + (1 << 1)]] < level[</pre>
         table[1][i]];
     if (C1 && C2) table[1 + 1][i] = table[1][i + (1 <<</pre>
         1)];
   }
 }
5.6 HLD
const int N = 1e6 + 7;
template <typename DT>
struct Segtree {
 // write lazy segtree here
Segtree<int> tree(N);
vector<int> adj[N];
int depth[N], par[N], pos[N];
int head[N], heavy[N], cnt;
int dfs(int u, int p) {
 int SZ = 1, mxsz = 0, heavyc;
 depth[u] = depth[p] + 1;
 for (auto v : adj[u]) {
   if (v == p) continue;
   par[v] = u;
   int subsz = dfs(v, u);
   if (subsz > mxsz) heavy[u] = v, mxsz = subsz;
   SZ += subsz;
 return SZ;
void decompose(int u, int h) {
 head[u] = h, pos[u] = ++cnt;
 if (heavy[u] != -1) decompose(heavy[u], h);
 for (int v : adj[u]) {
   if (v == par[u]) continue;
   if (v != heavy[u]) decompose(v, v);
int query(int a, int b) {
 int ret = 0:
 for (; head[a] != head[b]; b = par[head[b]]) {
   if (depth[head[a]] > depth[head[b]]) swap(a, b);
   ret += tree.query(1, 0, cnt, pos[head[b]], pos[b]);
 }
```

```
if (depth[a] > depth[b]) swap(a, b);
ret += tree.query(1, 0, cnt, pos[a], pos[b]);
return ret;
}
```

## 5.7 Centroid Decomposition

```
class Centroid_Decomposition {
 vector<bool> blocked;
 vector<int> CompSize;
 int CompDFS(tree &T, int node, int par) {
   CompSize[node] = 1;
  for (int &e : T[node])
    if (e != par and !blocked[e]) CompSize[node] +=
         CompDFS(T, e, node);
  return CompSize[node];
 int FindCentroid(tree &T, int node, int par, int sz) {
  for (int &e : T[node])
    if (e != par and !blocked[e])
      if (CompSize[e] > sz / 2) return FindCentroid(T,
           e, node, sz);
   return node;
 pair<int, int> GetCentroid(tree &T, int entry) {
  int sz = CompDFS(T, entry, entry);
   return {FindCentroid(T, entry, entry, sz), sz};
 c_vector<LL> left[2], right[2];
 int GMin, GMax;
 void dfs(tree &T, int node, int par, int Min, int Max,
     int sum) {
   if (blocked[node]) return:
   right[Max < sum or Min > sum][sum]++;
   Max = max(Max, sum), Min = min(Min, sum);
   GMin = min(GMin, sum), GMax = max(GMax, sum);
   for (int i = 0; i < T[node].size(); i++)</pre>
    if (T[node][i] != par) {
      dfs(T, T[node][i], node, Min, Max, sum + T.col[
           node][i]);
 }
LL solve(tree &T, int c, int sz) {
  LL ans = 0:
   left[0].clear(-sz, sz), left[1].clear(-sz, sz);
  for (int i = 0; i < T[c].size(); i++) {</pre>
    GMin = 1, GMax = -1;
     dfs(T, T[c][i], c, GMin, GMax, T.col[c][i]);
     ans += right[0][0] + left[1][0] * right[1][0];
     for (int j : {0, 1})
      for (int k = GMin; k <= GMax; k++) {</pre>
```

```
ans += right[j][k] * (left[0][-k] + (j == 0) *
             left[1][-k]);
     for (int j : {0, 1})
       for (int k = GMin; k <= GMax; k++) {</pre>
         left[j][k] += right[j][k];
        right[j][k] = 0;
       }
   }
   return ans;
public:
 LL operator()(tree &T, int entry) {
   blocked.resize(T.n);
   CompSize.resize(T.n);
   for (int i : {0, 1})
     left[i].resize(2 * T.n + 5), right[i].resize(2 * T.
         n + 5):
   auto [c, sz] = GetCentroid(T, entry);
   LL ans = solve(T, c, sz);
   blocked[c] = true:
   for (int e : T[c])
     if (!blocked[e]) ans += (*this)(T, e);
   return ans:
 }
};
```

#### 5.8 Dinic Max Flow

```
/// flow with demand(lower bound) only for DAG
// create new src and sink
// add_edge(new src, u, sum(in_demand[u]))
// add_edge(u, new sink, sum(out_demand[u]))
// add_edge(old sink, old src, inf)
// if (sum of lower bound == flow) then demand satisfied
// flow in every edge i = demand[i] + e.flow
using Ti = long long;
const Ti INF = 1LL << 60;</pre>
struct edge {
 int v, u;
 Ti cap, flow = 0;
  edge(int v, int u, Ti cap) : v(v), u(u), cap(cap) {}
const int N = 1e5 + 50:
vector<edge> edges;
vector<int> adj[N];
int m = 0, n;
int level[N], ptr[N];
queue<int> q;
```

```
bool bfs(int s, int t) {
 for (q.push(s), level[s] = 0; !q.empty(); q.pop()) {
   for (int id : adj[q.front()]) {
     auto &ed = edges[id];
     if (ed.cap - ed.flow > 0 and level[ed.u] == -1)
       level[ed.u] = level[ed.v] + 1, q.push(ed.u);
   }
 }
  return level[t] != -1;
Ti dfs(int v, Ti pushed, int t) {
 if (pushed == 0) return 0;
  if (v == t) return pushed;
 for (int &cid = ptr[v]; cid < adj[v].size(); cid++) {</pre>
   int id = adj[v][cid];
   auto &ed = edges[id];
    if (level[v] + 1 != level[ed.u] || ed.cap - ed.flow <</pre>
         1) continue;
    Ti tr = dfs(ed.u, min(pushed, ed.cap - ed.flow), t);
    if (tr == 0) continue;
    ed.flow += tr;
    edges[id ^ 1].flow -= tr;
   return tr;
  return 0;
void init(int nodes) {
 m = 0, n = nodes:
 for (int i = 0; i < n; i++) level[i] = -1, ptr[i] = 0,</pre>
      adj[i].clear();
void addEdge(int v, int u, Ti cap) {
  edges.emplace_back(v, u, cap), adj[v].push_back(m++);
  edges.emplace_back(u, v, 0), adj[u].push_back(m++);
Ti maxFlow(int s, int t) {
 Ti f = 0;
  for (auto &ed : edges) ed.flow = 0;
  for (; bfs(s, t); memset(level, -1, n * 4)) {
   for (memset(ptr, 0, n * 4); Ti pushed = dfs(s, INF, t
        ); f += pushed)
 }
  return f;
5.9 Min Cost Max Flow
```

```
struct edge {
 int v, rev;
 LL cap, cost, flow;
 edge() {}
 edge(int v, int rev, LL cap, LL cost)
     : v(v), rev(rev), cap(cap), cost(cost), flow(0) {}
struct mcmf {
 int src, sink, n;
 vector<int> par, idx, Q;
 vector<bool> inq;
 vector<LL> dis;
 vector<vector<edge>> g;
 mcmf() {}
 mcmf(int src, int sink, int n)
     : src(src), sink(sink), n(n), par(n), idx(n), inq(n
         ), dis(n), g(n), Q(10000005) {} // use Q(n) if
         not using random
 void add_edge(int u, int v, LL cap, LL cost, bool
      directed = true) {
   edge _u = edge(v, g[v].size(), cap, cost);
   edge _v = edge(u, g[u].size(), 0, -cost);
   g[u].pb(_u);
   g[v].pb(_v);
   if (!directed) add_edge(v, u, cap, cost, true);
 bool spfa() {
   for (int i = 0; i < n; i++) {
     dis[i] = inf, inq[i] = false;
   int f = 0, 1 = 0;
   dis[src] = 0, par[src] = -1, Q[1++] = src, inq[src] =
        true;
   while (f < 1) {
     int u = Q[f++];
     for (int i = 0; i < g[u].size(); i++) {</pre>
       edge &e = g[u][i];
       if (e.cap <= e.flow) continue;</pre>
       if (dis[e.v] > dis[u] + e.cost) {
        dis[e.v] = dis[u] + e.cost:
         par[e.v] = u, idx[e.v] = i;
        if (!inq[e.v]) inq[e.v] = true, Q[1++] = e.v;
        // if (!inq[e.v]) {
        // inq[e.v] = true;
        // if (f \&\& rnd() \& 7) Q[--f] = e.v;
        // else Q[1++] = e.v:
        // }
       }
     inq[u] = false;
```

```
}
   return (dis[sink] != inf);
  pair<LL, LL> solve() {
   LL mincost = 0, maxflow = 0;
   while (spfa()) {
     LL bottleneck = inf;
     for (int u = par[sink], v = idx[sink]; u != -1; v =
          idx[u], u = par[u]) {
       edge &e = g[u][v];
       bottleneck = min(bottleneck, e.cap - e.flow);
     for (int u = par[sink], v = idx[sink]; u != -1; v =
          idx[u], u = par[u]) {
       edge &e = g[u][v];
       e.flow += bottleneck;
       g[e.v][e.rev].flow -= bottleneck;
     mincost += bottleneck * dis[sink], maxflow +=
         bottleneck:
   return make_pair(mincost, maxflow);
};
// want to minimize cost and don't care about flow
// add edge from sink to dummy sink (cap = inf, cost = 0)
// add edge from source to sink (cap = inf, cost = 0)
// run mcmf, cost returned is the minimum cost
```

## 5.10 Block Cut Tree

```
vector<vector<int> > components;
vector<int> cutpoints, start, low;
vector<bool> is_cutpoint;
stack<int> st;
void find_cutpoints(int node, graph &G, int par = -1, int
     d = 0) {
 low[node] = start[node] = d++;
 st.push(node);
 int cnt = 0;
 for (int e : G[node])
   if (int to = G(e).to(node); to != par) {
     if (start[to] == -1) {
       find_cutpoints(to, G, node, d + 1);
       cnt++;
       if (low[to] >= start[node]) {
        is_cutpoint[node] = par != -1 or cnt > 1;
         components.push_back({node}); // starting a new
             block with the point
         while (st.top() != node)
```

```
components.back().push_back(st.top()), st.pop
               ();
       }
     low[node] = min(low[node], low[to]);
graph tree;
vector<int> id:
void init(graph &G) {
  int n = G.n;
  start.assign(n, -1), low.resize(n), is_cutpoint.resize(
      n), id.assign(n, -1);
  find_cutpoints(0, G);
  for (int u = 0; u < n; ++u)
    if (is_cutpoint[u]) id[u] = tree.addNode();
  for (auto &comp : components) {
    int node = tree.addNode();
    for (int u : comp)
      if (!is_cutpoint[u])
       id[u] = node;
      else
       tree.addEdge(node, id[u]);
  }
  if (id[0] == -1) // corner - 1
    id[0] = tree.addNode();
```

## 5.11 Bridge Tree

```
vector<int>adj[100010];
bool vis[100010];
int start[100010],low[100010],dep[100010],par[100010],
    timer, up [100010] [20];
vector<pair<int,int>>bridges;
int FindPar(int n){
 if(par[n]==n) return par[n];
  par[n]=FindPar(par[n]);
 return par[n];
void Merge(int u,int v){
 u=FindPar(u);
 v=FindPar(v);
 if(u!=v) par[v]=u;
void dfs(int s,int p){
 vis[s]=1;
  timer++:
  start[s]=timer;
 low[s]=timer;
 for(int i=0;i<adj[s].size();i++){</pre>
```

```
int v=adj[s][i];
   if(v==p) continue;
   if(vis[v]){
     low[s]=min(low[s],start[v]);
     continue;
   }
   else{
     dfs(v,s);
     low[s]=min(low[s],low[v]);
     if(low[v]>start[s]) bridges.push_back({s,v});
     else Merge(s,v);
void MakeBridgeTree(int n){
 for(int i=1;i<=n;i++) adj[i].clear();</pre>
 for(auto p: bridges){
   int u=FindPar(p.first);
   int v=FindPar(p.second);
   if(u!=v){
     adj[u].push_back(v);
     adj[v].push_back(u);
}
void dfs2(int s,int p){
 if(p!=-1) up[s][0]=p;
 for(int i=1;i<=18;i++) up[s][i]=up[up[s][i-1]][i-1];</pre>
   for(int i=0;i<adj[s].size();i++){</pre>
     int v=adj[s][i];
     if(v==p) continue;
     dep[v]=dep[s]+1;
     dfs2(v,s);
}
int LCA(int u,int v){
 if(dep[u]>dep[v]) swap(u,v);
 int d=dep[v]-dep[u];
 for(int i=18;i>=0;i--){
   if(d&(1<<i)) v=up[v][i];</pre>
 if(u==v) return u;
 for(int i=18;i>=0;i--){
   if(up[u][i]!=up[v][i]){
     u=up[u][i];
     v=up[v][i];
 return up[u][0];
```

```
int GetDistance(int u.int v){
 return dep[u]+dep[v]-2*dep[LCA(u,v)];
int Intersect(int lca_ab,int lca_cd,int x,int y){ // x= a }
     or b and v= c or d
 int lca_xy=LCA(x,y);
 if(dep[lca_xy] < dep[lca_ab] || dep[lca_xy] < dep[lca_cd])</pre>
      return 0;
  int ret=min(dep[lca_xy]-dep[lca_ab],dep[lca_xy]-dep[
      lca_cd]);
 return ret;
```

## 5.12 Tree Isomorphism

```
mp["01"] = 1;
ind = 1:
int dfs(int u, int p) {
 int cnt = 0;
 vector<int> vs:
 for (auto v : g1[u]) {
   if (v != p) {
    int got = dfs(v, u);
     vs.pb(got);
     cnt++;
   }
 if (!cnt) return 1;
  sort(vs.begin(), vs.end());
  string s = "0";
 for (auto i : vs) s += to_string(i);
 vs.clear();
  s.pb('1');
 if (mp.find(s) == mp.end()) mp[s] = ++ind;
 int ret = mp[s];
 return ret;
```

#### 6 Math

#### 6.1 Combi

```
array<int, N + 1> fact, inv, inv_fact;
void init() {
  fact[0] = inv_fact[0] = 1;
  for (int i = 1; i <= N; i++) {</pre>
   inv[i] = i == 1 ? 1 : (LL)inv[i - mod % i] * (mod / i | LL mul(LL a, LL b, LL mod) {
         + 1) % mod;
   fact[i] = (LL)fact[i - 1] * i % mod;
   inv_fact[i] = (LL)inv_fact[i - 1] * inv[i] % mod;
}
```

```
LL C(int n. int r) {
 return (r < 0 \text{ or } r > n) ? 0 : (LL)fact[n] * inv_fact[r]
       % mod * inv_fact[n - r] % mod;
```

```
6.2 Linear Sieve
const int N = 1e7:
vector<int> primes;
int spf[N + 5], phi[N + 5], NOD[N + 5], cnt[N + 5], POW[N]
bool prime[N + 5];
int SOD[N + 5];
void init() {
 fill(prime + 2, prime + N + 1, 1);
 SOD[1] = NOD[1] = phi[1] = spf[1] = 1;
 for (LL i = 2; i <= N; i++) {</pre>
   if (prime[i]) {
     primes.push_back(i), spf[i] = i;
     phi[i] = i - 1;
     NOD[i] = 2, cnt[i] = 1;
     SOD[i] = i + 1, POW[i] = i;
   for (auto p : primes) {
     if (p * i > N or p > spf[i]) break;
     prime[p * i] = false, spf[p * i] = p;
     if (i % p == 0) {
       phi[p * i] = p * phi[i];
       NOD[p * i] = NOD[i] / (cnt[i] + 1) * (cnt[i] + 2)
              cnt[p * i] = cnt[i] + 1;
       SOD[p * i] = SOD[i] / SOD[POW[i]] * (SOD[POW[i]]
           + p * POW[i]),
              POW[p * i] = p * POW[i];
       break;
     } else {
       phi[p * i] = phi[p] * phi[i];
       NOD[p * i] = NOD[p] * NOD[i], cnt[p * i] = 1;
       SOD[p * i] = SOD[p] * SOD[i], POW[p * i] = p;
   }
 }
```

#### 6.3 Pollard Rho

```
return (__int128)a * b % mod;
// LL ans = a * b - mod * (LL) (1.L / mod * a * b):
// return ans + mod * (ans < 0) - mod * (ans >= (LL)
    mod):
```

```
LL bigmod(LL num, LL pow, LL mod) {
 LL ans = 1;
 for (; pow > 0; pow >>= 1, num = mul(num, num, mod))
   if (pow & 1) ans = mul(ans, num, mod):
 return ans:
bool is_prime(LL n) {
 if (n < 2 or n % 6 % 4 != 1) return (n | 1) == 3;
 LL a[] = {2, 325, 9375, 28178, 450775, 9780504,
      1795265022};
 LL s = \_builtin\_ctzll(n - 1), d = n >> s;
 for (LL x : a) {
   LL p = bigmod(x \% n, d, n), i = s;
   for (; p != 1 and p != n - 1 and x % n and i--; p =
       mul(p, p, n))
   if (p != n - 1 and i != s) return false:
 return true;
LL get_factor(LL n) {
 auto f = [\&](LL x) \{ return mul(x, x, n) + 1; \};
 LL x = 0, y = 0, t = 0, prod = 2, i = 2, q;
 for (; t++ % 40 or gcd(prod, n) == 1; x = f(x), y = f(f(x))
      (v))) {
   (x == y) ? x = i++, y = f(x) : 0;
   prod = (q = mul(prod, max(x, y) - min(x, y), n)) ? q
        : prod;
 return gcd(prod, n);
map<LL, int> factorize(LL n) {
 map<LL, int> res;
 if (n < 2) return res;</pre>
 LL small_primes[] = {2, 3, 5, 7, 11, 13, 17, 19, 23,
      29, 31, 37, 41,
                     43, 47, 53, 59, 61, 67, 71, 73, 79,
                         83, 89, 971:
 for (LL p : small_primes)
   for (; n % p == 0; n /= p, res[p]++)
  auto _factor = [&](LL n, auto &_factor) {
   if (n == 1) return;
   if (is_prime(n))
     res[n]++:
   else {
     LL x = get_factor(n);
     _factor(x, _factor);
     _factor(n / x, _factor);
```

```
}
};
_factor(n, _factor);
return res;
}
```

### 6.4 Chinese Remainder Theorem

```
// given a, b will find solutions for
// ax + by = 1
tuple<LL, LL, LL> EGCD(LL a, LL b) {
  if (b == 0)
   return {1, 0, a};
  else {
    auto [x, y, g] = EGCD(b, a \% b);
   return \{y, x - a / b * y, g\};
}
// given modulo equations, will apply CRT
PLL CRT(vector<PLL> &v) {
  LL V = 0, M = 1;
  for (auto &[v, m] : v) { // value % mod
    auto [x, y, g] = EGCD(M, m);
   if ((v - V) % g != 0) return {-1, 0};
   V += x * (v - V) / g % (m / g) * M, M *= m / g;
   V = (V \% M + M) \% M;
  return make_pair(V, M);
```

#### 6.5 Mobius Function

```
const int N = 1e6 + 5;
int mob[N];
void mobius() {
  memset(mob, -1, sizeof mob);
  mob[1] = 1;
  for (int i = 2; i < N; i++)
    if (mob[i]) {
      for (int j = i + i; j < N; j += i) mob[j] -= mob[i
        ];
    }
}</pre>
```

## 6.6 FFT

```
using CD = complex<double>;
typedef long long LL;
const double PI = acos(-1.0L);
int N;
vector<int> perm;
vector<CD> wp[2];
```

```
void precalculate(int n) {
 assert((n & (n - 1)) == 0), N = n;
 perm = vector<int>(N, 0);
 for (int k = 1; k < N; k <<= 1) {</pre>
   for (int i = 0; i < k; i++) {</pre>
     perm[i] <<= 1;
     perm[i + k] = 1 + perm[i];
  wp[0] = wp[1] = vector < CD > (N);
 for (int i = 0; i < N; i++) {</pre>
   wp[0][i] = CD(cos(2 * PI * i / N), sin(2 * PI * i / N)
   wp[1][i] = CD(cos(2 * PI * i / N), -sin(2 * PI * i / N))
       N));
 }
void fft(vector<CD> &v, bool invert = false) {
 if (v.size() != perm.size()) precalculate(v.size());
 for (int i = 0; i < N; i++)</pre>
   if (i < perm[i]) swap(v[i], v[perm[i]]);</pre>
 for (int len = 2; len <= N; len *= 2) {</pre>
   for (int i = 0, d = N / len; i < N; i += len) {</pre>
     for (int j = 0, idx = 0; j < len / 2; j++, idx += d
         ) {
       CD x = v[i + i];
       CD y = wp[invert][idx] * v[i + j + len / 2];
       v[i + j] = x + y;
       v[i + j + len / 2] = x - y;
   }
 }
 if (invert) {
   for (int i = 0; i < N; i++) v[i] /= N;</pre>
void pairfft(vector<CD> &a, vector<CD> &b, bool invert =
    false) {
 int N = a.size();
 vector<CD> p(N);
 for (int i = 0; i < N; i++) p[i] = a[i] + b[i] * CD(0,
      1);
 fft(p, invert);
 p.push_back(p[0]);
 for (int i = 0; i < N; i++) {</pre>
   if (invert) {
     a[i] = CD(p[i].real(), 0);
     b[i] = CD(p[i].imag(), 0);
   } else {
     a[i] = (p[i] + conj(p[N - i])) * CD(0.5, 0);
```

```
b[i] = (p[i] - conj(p[N - i])) * CD(0, -0.5);
   }
}
vector<LL> multiply(const vector<LL> &a, const vector<LL>
     &b) {
 int n = 1;
 while (n < a.size() + b.size()) n <<= 1;</pre>
 vector<CD> fa(a.begin(), a.end()), fb(b.begin(), b.end
      ());
 fa.resize(n):
 fb.resize(n);
          fft(fa); fft(fb);
 pairfft(fa, fb);
 for (int i = 0; i < n; i++) fa[i] = fa[i] * fb[i];</pre>
 fft(fa, true);
 vector<LL> ans(n):
 for (int i = 0; i < n; i++) ans[i] = round(fa[i].real()</pre>
 return ans;
const int M = 1e9 + 7, B = sqrt(M) + 1;
vector<LL> anyMod(const vector<LL> &a, const vector<LL> &
    b) {
 int n = 1:
 while (n < a.size() + b.size()) n <<= 1;</pre>
 vector<CD> al(n), ar(n), bl(n), br(n);
 for (int i = 0; i < a.size(); i++) al[i] = a[i] % M / B</pre>
      , ar[i] = a[i] % M % B;
 for (int i = 0; i < b.size(); i++) bl[i] = b[i] % M / B</pre>
     , br[i] = b[i] % M % B;
 pairfft(al, ar);
 pairfft(bl, br);
          fft(al); fft(ar); fft(bl); fft(br);
 for (int i = 0; i < n; i++) {
   CD ll = (al[i] * bl[i]), lr = (al[i] * br[i]);
   CD rl = (ar[i] * bl[i]), rr = (ar[i] * br[i]);
   al[i] = 11:
   ar[i] = lr;
   bl[i] = rl;
   br[i] = rr:
 pairfft(al, ar, true);
 pairfft(bl, br, true);
          fft(al, true); fft(ar, true); fft(bl, true);
     fft(br. true):
 vector<LL> ans(n);
 for (int i = 0; i < n; i++) {</pre>
   LL right = round(br[i].real()), left = round(al[i].
       real());
```

```
LL mid = round(round(bl[i].real()) + round(ar[i].real
   ans[i] = ((left \% M) * B * B + (mid \% M) * B + right)
         % M;
 }
 return ans;
6.7 NTT
const LL N = 1 << 18:
const LL MOD = 786433;
vector<LL> P[N];
LL rev[N], w[N | 1], a[N], b[N], inv_n, g;
LL Pow(LL b, LL p) {
 LL ret = 1;
 while (p) {
   if (p & 1) ret = (ret * b) % MOD;
   b = (b * b) \% MOD;
   p >>= 1;
 return ret;
LL primitive_root(LL p) {
 vector<LL> factor;
 LL phi = p - 1, n = phi;
 for (LL i = 2; i * i <= n; i++) {</pre>
   if (n % i) continue;
   factor.emplace_back(i);
   while (n \% i == 0) n /= i;
 if (n > 1) factor.emplace_back(n);
 for (LL res = 2; res <= p; res++) {</pre>
   bool ok = true;
   for (LL i = 0; i < factor.size() && ok; i++)</pre>
     ok &= Pow(res, phi / factor[i]) != 1;
   if (ok) return res;
 }
 return -1;
void prepare(LL n) {
 LL sz = abs(31 - \_builtin\_clz(n));
 LL r = Pow(g, (MOD - 1) / n);
 inv_n = Pow(n, MOD - 2);
 w[0] = w[n] = 1;
 for (LL i = 1; i < n; i++) w[i] = (w[i - 1] * r) % MOD; const LL MOD = 30011;
 for (LL i = 1; i < n; i++)
   rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << (sz - 1));
```

```
void NTT(LL *a, LL n, LL dir = 0) {
for (LL i = 1; i < n - 1; i++)
   if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
 for (LL m = 2: m <= n: m <<= 1) {
   for (LL i = 0; i < n; i += m) {</pre>
     for (LL j = 0; j < (m >> 1); j++) {
       LL &u = a[i + j], &v = a[i + j + (m >> 1)];
       LL t = v * w[dir ? n - n / m * j : n / m * j] %
       v = u - t < 0 ? u - t + MOD : u - t;
       u = u + t >= MOD ? u + t - MOD : u + t;
   }
 }
 if (dir)
   for (LL i = 0; i < n; i++) a[i] = (inv_n * a[i]) %
       MOD:
vector<LL> mul(vector<LL> p, vector<LL> q) {
 LL n = p.size(), m = q.size();
 LL t = n + m - 1, sz = 1;
 while (sz < t) sz <<= 1;
 prepare(sz);
 for (LL i = 0; i < n; i++) a[i] = p[i];</pre>
 for (LL i = 0; i < m; i++) b[i] = q[i];
 for (LL i = n; i < sz; i++) a[i] = 0;</pre>
 for (LL i = m: i < sz: i++) b[i] = 0:
 NTT(a, sz);
 NTT(b, sz);
 for (LL i = 0; i < sz; i++) a[i] = (a[i] * b[i]) % MOD;</pre>
 NTT(a, sz, 1);
 vector<LL> c(a, a + sz);
 while (c.size() && c.back() == 0) c.pop_back();
 return c;
6.8 WalshHadamard
```

```
#include <bits/stdc++.h>
using namespace std;
typedef long long LL;
#define bitwiseXOR 1
// #define bitwiseAND 2
// #define bitwiseOR 3
const LL MOD = 30011;

LL BigMod(LL b, LL p) {
   LL ret = 1;
```

```
while (p > 0) {
   if (p % 2 == 1) {
     ret = (ret * b) \% MOD;
   p = p / 2;
   b = (b * b) \% MOD;
 return ret % MOD;
void FWHT(vector<LL>& p, bool inverse) {
 LL n = p.size();
 assert((n & (n - 1)) == 0);
 for (LL len = 1; 2 * len <= n; len <<= 1) {
   for (LL i = 0; i < n; i += len + len) {
     for (LL j = 0; j < len; j++) {</pre>
      LL u = p[i + j];
       LL v = p[i + len + j];
#ifdef bitwiseXOR
       p[i + j] = (u + v) \% MOD;
       p[i + len + j] = (u - v + MOD) \% MOD;
#endif // bitwiseXOR
#ifdef bitwiseAND
       if (!inverse) {
         p[i + j] = v \% MOD;
         p[i + len + j] = (u + v) \% MOD;
       } else {
         p[i + j] = (-u + v) \% MOD;
         p[i + len + j] = u \% MOD;
#endif // bitwiseAND
#ifdef bitwiseOR
       if (!inverse) {
         p[i + j] = u + v;
        p[i + len + j] = u;
       } else {
         p[i + j] = v;
         p[i + len + j] = u - v;
#endif // bitwiseOR
  }
 }
#ifdef bitwiseXOR
 if (inverse) {
```

```
LL val = BigMod(n, MOD - 2); // Option 2: Exclude
for (LL i = 0; i < n; i++) {
      // assert(p[i]%n==0); //Option 2: Include
      p[i] = (p[i] * val) % MOD; // Option 2: p[i]/=n;
    }
}
#endif // bitwiseXOR
}</pre>
```

## 6.9 Adaptive Simpsons

```
For finding the length of an arc in a range
   L = integrate(ds) from start to end of range
   where ds = sqrt(1+(d/dy(x))^2)dy
const double SIMPSON_TERMINAL_EPS = 1e-12;
/// Function whose integration is to be calculated
double F(double x);
double simpson(double minx, double maxx) {
 return (\max - \min x) / 6 * (F(\min x) + 4 * F((\min x + \max x)))
      maxx) / 2.) + F(maxx));
double adaptive_simpson(double minx, double maxx, double
    c, double EPS) {
  // if(maxx - minx < SIMPSON TERMINAL EPS) return 0:
  double midx = (minx + maxx) / 2;
  double a = simpson(minx, midx);
  double b = simpson(midx, maxx);
  if (fabs(a + b - c) < 15 * EPS) return a + b + (a + b - c)
       c) / 15.0:
 return adaptive_simpson(minx, midx, a, EPS / 2.) +
        adaptive_simpson(midx, maxx, b, EPS / 2.);
double adaptive_simpson(double minx, double maxx, double
    EPS) {
 return adaptive_simpson(minx, maxx, simpson(minx, maxx,
       i). EPS):
```

## 6.10 Berlekamp Massey

```
struct berlekamp_massey { // for linear recursion
  typedef long long LL;
  static const int SZ = 2e5 + 5;
  static const int MOD = 1e9 + 7; /// mod must be a prime
  LL m , a[SZ] , h[SZ] , t_[SZ] , s[SZ] , t[SZ];
  // bigmod goes here
  inline vector <LL> BM( vector <LL> &x ) {
```

```
LL lf , ld;
 vector <LL> ls , cur;
 for ( int i = 0; i < int(x.size()); ++i ) {</pre>
   LL t = 0:
   for ( int j = 0; j < int(cur.size()); ++j ) t = (t</pre>
       + x[i - j - 1] * cur[j]) % MOD;
   if ((t - x[i]) \% MOD == 0) continue;
   if (!cur.size()) {
     cur.resize( i + 1 );
     lf = i; ld = (t - x[i]) \% MOD;
     continue;
   LL k = -(x[i] - t) * bigmod(ld, MOD - 2, MOD) %
         MOD;
   vector <LL> c(i - lf - 1);
   c.push_back( k );
   for ( int j = 0; j < int(ls.size()); ++j ) c.</pre>
        push_back(-ls[j] * k % MOD);
   if ( c.size() < cur.size() ) c.resize( cur.size() )</pre>
   for ( int j = 0; j < int(cur.size()); ++j ) c[j] =</pre>
        (c[i] + cur[i]) % MOD;
   if (i - lf + (int)ls.size() >= (int)cur.size() ) ls
         = cur, lf = i, ld = (t - x[i]) \% MOD;
   cur = c;
 for ( int i = 0; i < int(cur.size()); ++i ) cur[i] =</pre>
      (cur[i] % MOD + MOD) % MOD;
 return cur;
inline void mull( LL *p , LL *q ) {
 for ( int i = 0; i < m + m; ++i ) t_{i} = 0;
 for ( int i = 0; i < m; ++i ) if ( p[i] )
     for ( int j = 0; j < m; ++j ) t_{i} = (t_{i} + t_{i})
          i] + p[i] * q[i]) % MOD;
 for ( int i = m + m - 1; i >= m; --i ) if ( t_[i] )
     for (int j = m - 1; i = m - 1) t_{i} = 1
         t_{i} = i - 1 + t_{i} * h[i] * h[i] % MOD;
 for ( int i = 0; i < m; ++i ) p[i] = t_[i];</pre>
inline LL calc( LL K ) {
 for ( int i = m; ~i; --i ) s[i] = t[i] = 0;
 s[0] = 1; if ( m != 1 ) t[1] = 1; else t[0] = h[0];
 while ( K ) {
   if (K&1) mull(s,t);
   mull( t , t ); K >>= 1;
 LL su = 0:
 for ( int i = 0; i < m; ++i ) su = (su + s[i] * a[i])
       % MOD;
```

```
return (su % MOD + MOD) % MOD;
 /// already calculated upto k , now calculate upto n.
 inline vector <LL> process( vector <LL> &x , int n ,
      int k ) {
   auto re = BM( x );
   x.resize(n+1):
   for ( int i = k + 1; i <= n; i++ ) {
     for ( int j = 0; j < re.size(); j++ ) {</pre>
       x[i] += 1LL * x[i - j - 1] % MOD * re[j] % MOD; x
            [i] %= MOD;
     }
   }
   return x;
 inline LL work( vector <LL> &x , LL n ) {
   if ( n < int(x.size()) ) return x[n] % MOD;</pre>
   vector \langle LL \rangle v = BM( x ); m = v.size(); if ( !m )
        return 0:
   for ( int i = 0; i < m; ++i ) h[i] = v[i], a[i] = x[i]
       ];
   return calc( n ) % MOD;
} rec;
vector <LL> v:
void solve() {
 int n;
 cin >> n:
 cout << rec.work(v, n - 1) << endl;</pre>
```

#### 6.11 Fractional Binary Search

```
/**
Given a function f and n, finds the smallest fraction p /
    q in [0, 1] or [0,n]
such that f(p / q) is true, and p, q <= n.
Time: O(log(n))
**/
struct frac { long long p, q; };
bool f(frac x) {
  return 6 + 8 * x.p >= 17 * x.q + 12;
}
frac fracBS(long long n) {
  bool dir = 1, A = 1, B = 1;
  frac lo{0, 1}, hi{1, 0}; // Set hi to 1/0 to search
    within [0, n] and {1, 1} to search within [0, 1]
if (f(lo)) return lo;
assert(f(hi)); //checking if any solution exists or not
while (A || B) {
```

```
long long adv = 0, step = 1; // move hi if dir, else
      10
  for (int si = 0; step; (step *= 2) >>= si) {
   adv += step:
   frac mid{lo.p * adv + hi.p, lo.q * adv + hi.q};
   if (abs(mid.p) > n \mid\mid mid.q > n \mid\mid dir == !f(mid))
        {
     adv -= step; si = 2;
   }
 }
  hi.p += lo.p * adv;
  hi.q += lo.q * adv;
  dir = !dir;
  swap(lo, hi);
  A = B; B = !!adv;
return dir ? hi : lo;
```

### 6.12 Lagrange

```
// p is a polynomial with n points.
// p(0), p(1), p(2), \dots p(n-1) are given.
// Find p(x).
LL Lagrange(vector<LL> &p, LL x) {
  LL n = p.size(), L, i, ret;
  if (x < n) return p[x];</pre>
 L = 1:
  for (i = 1; i < n; i++) {</pre>
   L = (L * (x - i)) \% MOD;
   L = (L * bigmod(MOD - i, MOD - 2)) \% MOD;
 ret = (L * p[0]) % MOD;
  for (i = 1; i < n; i++) {
   L = (L * (x - i + 1)) \% MOD;
   L = (L * bigmod(x - i, MOD - 2)) \% MOD;
   L = (L * bigmod(i, MOD - 2)) % MOD;
   L = (L * (MOD + i - n)) % MOD;
   ret = (ret + L * p[i]) % MOD;
 return ret;
```

## 6.13 Shanks' Baby Step, Giant Step

```
// Finds a^x = b (mod p)

LL bigmod(LL b, LL p, LL m) {}

LL babyStepGiantStep(LL a, LL b, LL p) {
    LL i, j, c, sq = sqrt(p);
    map<LL, LL> babyTable;
```

## 6.14 Xor Basis

```
struct XorBasis {
 static const int sz = 64;
 array<ULL, sz> base = {0}, back;
 array<int, sz> pos;
 void insert(ULL x, int p) {
   ULL cur = 0:
   for (int i = sz - 1; ~i; i--)
     if (x >> i & 1) {
       if (!base[i]) {
         base[i] = x, back[i] = cur, pos[i] = p;
       } else x ^= base[i], cur |= 1ULL << i;</pre>
     }
 }
 pair<ULL, vector<int>> construct(ULL mask) {
   ULL ok = 0, x = mask:
   for (int i = sz - 1; ~i; i--)
     if (mask >> i & 1 and base[i]) mask ^= base[i], ok
         |= 1ULL << i;
   vector<int> ans;
   for (int i = 0; i < sz; i++)</pre>
     if (ok >> i & 1) {
       ans.push_back(pos[i]);
       ok ^= back[i];
   return {x ^ mask, ans};
};
```

## 7 String

### 7.1 Aho Corasick

```
int trie[260010][26],fail[260010],freq[260010],node,pos
       [260010];
void Clear(){
  for(int i=0;i<260000;i++){</pre>
```

```
for(int j=0;j<26;j++) trie[i][j]=0;</pre>
    freq[i]=0;
   fail[i]=0;
  pos[i]=0;
 node=0;
int Insert(string s){
 int curr=0:
 for(int i=0;i<s.size();i++){</pre>
   int letter=s[i]-'a';
   if(!trie[curr][letter]){
     node++:
     trie[curr][letter]=node;
   curr=trie[curr][letter];
 return curr;
void bfs(){
 queue<int>q;
 for(int i=0;i<26;i++){
  if(trie[0][i]) q.push(trie[0][i]);
 while(!q.empty()){
   int u=q.front();
   q.pop();
   for(int i=0;i<26;i++){</pre>
     int v=trie[u][i];
     if(v==0) continue:
     q.push(v);
     int f=fail[u]:
     while(f && trie[f][i]==0) f=fail[f];
     fail[v]=trie[f][i];
   }
}
void Aho_Corasick(string s){
 bfs();
 int curr=0:
 for(int i=0;i<s.size();i++){</pre>
   int letter=s[i]-'a';
   while(curr && trie[curr][letter]==0) curr=fail[curr];
   curr=trie[curr][letter];
   int temp=curr;
   while(temp>0){
     freq[temp]++;
     temp=fail[temp];
}
```

## 7.2 Double hash

```
// define +, -, * for (PLL, LL) and (PLL, PLL), % for (
    PLL. PLL):
PLL base(1949313259, 1997293877);
PLL mod(2091573227, 2117566807);
PLL power(PLL a, LL p) {
 PLL ans = PLL(1, 1);
 for(; p; p >>= 1, a = a * a % mod) {
     if(p & 1) ans = ans * a % mod;
 return ans;
PLL inverse(PLL a) { return power(a, (mod.ff - 1) * (mod.
    ss - 1) - 1): }
PLL inv_base = inverse(base);
PLL val:
vector<PLL> P;
void hash init(int n) {
 P.resize(n + 1);
 P[0] = PLL(1, 1);
 for (int i = 1; i <= n; i++) P[i] = (P[i - 1] * base) %
      mod:
PLL append(PLL cur, char c) { return (cur * base + c) %
    mod; }
/// prepends c to string with size k
PLL prepend(PLL cur, int k, char c) { return (P[k] * c +
    cur) % mod: }
/// replaces the i-th (0-indexed) character from right
    from a to b;
PLL replace(PLL cur, int i, char a, char b) {
  cur = (cur + P[i] * (b - a)) \% mod;
 return (cur + mod) % mod;
/// Erases c from the back of the string
PLL pop_back(PLL hash, char c) {
 return (((hash - c) * inv_base) % mod + mod) % mod;
/// Erases c from front of the string with size len
PLL pop_front(PLL hash, int len, char c) {
 return ((hash - P[len - 1] * c) % mod + mod) % mod;
/// concatenates two strings where length of the right is
```

```
PLL concat(PLL left, PLL right, int k) { return (left * P // d[i] = number of palindromes taking s[i] as center
    [k] + right) % mod; }
/// Calculates hash of string with size len repeated cnt
    times
/// This is O(log n). For O(1), pre-calculate inverses
PLL repeat(PLL hash, int len, LL cnt) {
  PLL mul = (P[len * cnt] - 1) * inverse(P[len] - 1);
 mul = (mul % mod + mod) % mod;
  PLL ret = (hash * mul) % mod:
  if (P[len].ff == 1) ret.ff = hash.ff * cnt;
 if (P[len].ss == 1) ret.ss = hash.ss * cnt;
 return ret;
LL get(PLL hash) { return ((hash.ff << 32) ^ hash.ss); }
struct hashlist {
 int len;
  vector<PLL> H. R:
  hashlist() {}
  hashlist(string& s) {
   len = (int)s.size();
   hash_init(len);
   H.resize(len + 1, PLL(0, 0)), R.resize(len + 2, PLL
   for (int i = 1; i <= len; i++) H[i] = append(H[i -</pre>
        1]. s[i - 1]):
   for (int i = len; i >= 1; i--) R[i] = append(R[i +
        1], s[i - 1]);
 /// 1-indexed
  PLL range_hash(int 1, int r) {
   return ((H[r] - H[l - 1] * P[r - l + 1]) \% mod + mod)
         % mod:
  PLL reverse_hash(int 1, int r) {
   return ((R[1] - R[r + 1] * P[r - 1 + 1]) \% mod + mod)
         % mod;
  PLL concat_range_hash(int 11, int r1, int 12, int r2) {
   return concat(range_hash(l1, r1), range_hash(l2, r2),
         r2 - 12 + 1):
  PLL concat_reverse_hash(int 11, int r1, int 12, int r2)
   return concat(reverse_hash(12, r2), reverse_hash(11,
       r1), r1 - 11 + 1);
 }
```

7.3 Manacher's

vector<int> d1(n);

```
for (int i = 0, l = 0, r = -1; i < n; i++) {
 int k = (i > r) ? 1 : min(d1[1 + r - i], r - i + 1);
 while (0 \le i - k \&\& i + k \le n \&\& s[i - k] == s[i + k])
       k++:
 d1[i] = k--;
 if (i + k > r) 1 = i - k, r = i + k:
vector<int> d2(n):
// d[i] = number of palindromes taking s[i-1] and s[i] as
     center
for (int i = 0, l = 0, r = -1; i < n; i++) {
 int k = (i > r) ? 0 : min(d2[1 + r - i + 1], r - i + 1)
 while (0 \le i - k - 1 \&\& i + k \le n \&\& s[i - k - 1] == s
      [i + k]) k++;
 d2[i] = k--:
 if (i + k > r) l = i - k - 1, r = i + k;
7.4 Palindromic Tree
int tree[100010][26],len[100010],link[100010],node,curr;
void Clear(int n){
 for(int i=0:i<=n:i++){</pre>
   len[i]=0, link[i]=0;
   for(int j=0;j<26;j++) tree[i][j]=0;</pre>
 }
void Init(){
 node=2, curr=2, len[1]=-1, len[2]=0, link[1]=1, link
      [2]=1:
```

## 7.5 String Match FFT

curr=tree[curr][c];

void Extend(int pos){

if(!tree[curr][c]){

tree[curr][c]=node: len[node] = len[curr] + 2;

node++:

int c=s[pos]-'a',x=link[curr];

if(len[node] == 1) link[node] = 2;

else link[node] = tree[x][c];

//find occurrences of t in s where '?'s are automatically matched with any character

while(s[pos-len[curr]-1]!=s[pos]) curr=link[curr];

while(s[pos-len[x]-1]!=s[pos]) x=link[x];

```
//res[i + m - 1] = sum_j = 0 to m - 1_{s[i + j]} * t[j] * (s]
    [i + i] - t[i]
vector<int> string_matching(string &s, string &t) {
 int n = s.size(), m = t.size();
 vector\langle int \rangle s1(n), s2(n), s3(n);
 for(int i = 0; i < n; i++) s1[i] = s[i] == '?' ? 0 : s[
      i] - 'a' + 1; //assign any non zero number for non
      '?'s
  for(int i = 0; i < n; i++) s2[i] = s1[i] * s1[i];
 for(int i = 0; i < n; i++) s3[i] = s1[i] * s2[i];
  vector\langle int \rangle t1(m), t2(m), t3(m);
 for(int i = 0; i < m; i++) t1[i] = t[i] == '?' ? 0 : t[</pre>
      i] - 'a' + 1;
 for(int i = 0; i < m; i++) t2[i] = t1[i] * t1[i];</pre>
 for(int i = 0; i < m; i++) t3[i] = t1[i] * t2[i];</pre>
 reverse(t1.begin(), t1.end());
 reverse(t2.begin(), t2.end());
 reverse(t3.begin(), t3.end());
  vector<int> s1t3 = multiply(s1, t3);
  vector<int> s2t2 = multiply(s2, t2);
  vector<int> s3t1 = multiply(s3, t1);
 vector<int> res(n);
  for(int i = 0; i < n; i++) res[i] = s1t3[i] - s2t2[i] * }
       2 + s3t1[i]:
 vector<int> oc:
 for(int i = m - 1; i < n; i++) if(res[i] == 0) oc.
      push_back(i - m + 1);
 return oc;
```

#### 7.6 Suffix Array

```
vector<int> SortCyclicShifts(string s){
 int n=s.size();
 const int alphabet=256;
 vector<int>Start(n),Freq(max(alphabet,n),0),EClass(n);
 for(int i=0;i<n;i++) Freq[s[i]]++;</pre>
 for(int i=1;i<alphabet;i++) Freq[i]+=Freq[i-1];</pre>
  for(int i=0;i<n;i++) Start[--Freq[s[i]]]=i;</pre>
  int Class=1;
  EClass[Start[0]]=0:
  for(int i=1;i<n;i++){</pre>
   if(s[Start[i]]!=s[Start[i-1]]) Class++;
    EClass[Start[i]]=Class-1;
  for(int p=0;(1<<p)<n;p++){</pre>
   vector<int> STemp(n),ECTemp(n);
   for(int i=0;i<n;i++){</pre>
       STemp[i]=Start[i]-(1<<p);</pre>
       if(STemp[i]<0) STemp[i]+=n;</pre>
   }
```

```
fill(Freq.begin(),Freq.begin()+Class,0);
   for(int i=0;i<n;i++) Freq[EClass[STemp[i]]]++;</pre>
   for(int i=1;i<Class;i++) Freq[i]+=Freq[i-1];</pre>
   for(int i=n-1:i>=0:i--){
       Start[--Freq[EClass[STemp[i]]]]=STemp[i];
   ECTemp[Start[0]]=0;
   Class=1;
   for(int i=1;i<n;i++){</pre>
     pair<int,int>prev,curr;
     prev={EClass[Start[i-1]],EClass[(Start[i-1]+(1<<p))</pre>
     curr={EClass[Start[i]],EClass[(Start[i]+(1<<p))%n</pre>
          1}:
     if(curr!=prev) Class++;
     ECTemp[Start[i]]=Class-1;
   for(int i=0;i<n;i++){</pre>
     EClass[Start[i]] = ECTemp[Start[i]];
   }
 }
 return Start;
vector<int> BuildSuffixArray(string s){
   s+='$':
   vector<int>ret=SortCyclicShifts(s);
   ret.erase(ret.begin());
   return ret:
vector<int>BuildLCPArray(string s,vector<int>&v){
   int n=s.size();
   vector<int>rank(n),lcp(n-1,0);
   for(int i=0;i<n;i++) rank[v[i]]=i;</pre>
   int k=0:
   for(int i=0;i<n;i++){</pre>
     if(rank[i]==n-1){
         k=0;
         continue;
     int j=v[rank[i]+1];
     while (i+k< n \&\& j+k< n \&\& s[i+k]==s[j+k]) k++;
     lcp[rank[i]] = k;
     if (k) k--;
   return lcp;
7.7 Suffix Automata
```

```
const int MAXN = 1e5 + 7, ALPHA = 26;
int len[2 * MAXN], link[2 * MAXN], nxt[2 * MAXN][ALPHA];
```

```
int sz:
int last;
void sa init() {
 memset(nxt, -1, sizeof nxt);
 len[0] = 0, link[0] = -1, sz = 1, last = 0;
void add(char ch) {
 int c = ch - 'a';
 int cur = sz++;
 len[cur] = len[last] + 1;
 int u = last;
 while (u != -1 && nxt[u][c] == -1) {
   nxt[u][c] = cur;
   u = link[u]:
 if (u == -1) link[cur] = 0:
 else {
   int v = nxt[u][c]:
   if (len[v] == len[u] + 1) link[cur] = v;
   else {
     int clone = sz++:
     len[clone] = 1 + len[u], link[clone] = link[v];
     for (int i = 0; i < ALPHA; i++) nxt[clone][i] = nxt</pre>
          [v][i];
     while (u != -1 && nxt[u][c] == v) {
       nxt[u][c] = clone:
       u = link[u];
     link[v] = link[cur] = clone;
 }
 last = cur:
vector<int> edge[2 * MAXN];
/// Optional, Call after adding all characters
void makeEdge() {
 for (int i = 0; i < sz; i++) {</pre>
   edge[i].clear();
   for (int j = 0; j < ALPHA; j++)
     if (nxt[i][j] != -1) edge[i].push_back(j);
 }
// Given a string S, you have to answer some queries:
```

```
// If all distinct substrings of string S were sorted
// lexicographically, which one will be the K-th smallest
long long dp[2 * MAXN];
bool vis[2 * MAXN];
void dfs(int u) {
  if (vis[u]) return;
  vis[u] = 1;
  dp[u] = 1;
  for (int i : edge[u]) {
   if (nxt[u][i] == -1) continue;
    dfs(nxt[u][i]);
   dp[u] += dp[nxt[u][i]];
}
void go(int u, long long rem, string &s) {
  if (rem == 1) return;
  long long sum = 1;
 for (int i : edge[u]) {
   if (nxt[u][i] == -1) continue;
   if (sum + dp[nxt[u][i]] < rem) {</pre>
     sum += dp[nxt[u][i]];
   } else {
     s += ('a' + i);
     go(nxt[u][i], rem - sum, s);
     return;
   }
 }
```

## 7.8 Z Algo

```
vector<int> calcz(string s) {
 int n = s.size();
 vector<int> z(n);
 int 1 = 0, r = 0;
 for (int i = 1; i < n; i++) {</pre>
   if (i > r) {
    l = r = i;
     while (r < n \&\& s[r] == s[r - 1]) r++;
     z[i] = r - 1, r--;
   } else {
     int k = i - 1;
     if (z[k] < r - i + 1) z[i] = z[k];
     else {
       l = i;
       while (r < n \&\& s[r] == s[r - 1]) r++;
       z[i] = r - 1, r--;
```

```
}
}
return z;
```

## Equations and Formulas

## Catalan Numbers

$$C_n = \frac{1}{n+1} {2n \choose n} C_0 = 1, C_1 = 1 \text{ and } C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}$$

The number of ways to completely parenthesize n+1 factors. The number of triangulations of a convex polygon with n+2sides (i.e. the number of partitions of polygon into disjoint tegers  $1, 2, \ldots, n$  into k nonempty subsets such that all ele- $\sum [\gcd(i, n) = k] = \phi(\frac{n}{i})$ triangles by using the diagonals).

form n disjoint i.e. non-intersecting chords.

The number of rooted full binary trees with n+1 leaves (ver- $S^d(n,k) = S(n-d+1,k-d+1), n \ge k \ge d$ tices are not numbered). A rooted binary tree is full if every 8.4 Other Combinatorial Identities vertex has either two children or no children.

Number of permutations of  $1, \ldots, n$  that avoid the pattern 123 (or any of the other patterns of length 3); that is, the number of permutations with no three-term increasing sub-sequence. For n = 3, these permutations are 132, 213, 231, 312 and 321

## 8.2 Stirling Numbers First Kind

The Stirling numbers of the first kind count permutations according to their number of cycles (counting fixed points as cycles of length one).

S(n,k) counts the number of permutations of n elements with k disjoint cycles.

$$S(n,k) = (n-1) \cdot S(n-1,k) + S(n-1,k-1), where, S(0,0) =$$

$$1, S(n,0) = S(0,n) = 0 \sum_{k=0}^{n} S(n,k) = n!$$

The unsigned Stirling numbers may also be defined algebraically, as the coefficient of the rising factorial:

$$x^{\bar{n}} = x(x+1)...(x+n-1) = \sum_{k=0}^{n} S(n,k)x^{k}$$

Lets [n, k] be the stirling number of the first kind, then

$$[n - k] = \sum_{0 \le i_1 < i_2 < i_k < n} i_1 i_2 \dots i_k.$$

## 8.3 Stirling Numbers Second Kind

Stirling number of the second kind is the number of wavs to partition a set of n objects into k non-empty subsets.

$$S(n,k) = k \cdot S(n-1,k) + S(n-1,k-1)$$
, where  $S(0,0) = 1$ ,  $S(n,0) = S(0,n) = 0$   $S(n,2) = 2^{n-1} - 1$   $S(n,k) \cdot k! = \text{number of ways to color } n \text{ nodes using colors from } 1 \text{ to } k \text{ such that if } m \text{ is any integer, then } \gcd(a+m\cdot b,b) = \gcd(a,b)$  The gcd is a multiplicative function in the follow

ber of ways to partition a set of n objects into k subsets, with  $\gcd(a_1,b) \cdot \gcd(a_2,b)$ .

each subset containing at least r elements. It is denoted by  $\gcd(a, \operatorname{lcm}(b, c)) = \operatorname{lcm}(\gcd(a, b), \gcd(a, c))$ .  $S_r(n,k)$  and obeys the recurrence relation.  $S_r(n+1,k) = |\operatorname{lcm}(a,\operatorname{gcd}(b,c))| = \operatorname{gcd}(\operatorname{lcm}(a,b),\operatorname{lcm}(a,c)).$  $\left| kS_r(n,k) + \binom{n}{r-1} S_r(n-r+1,k-1) \right|$ 

Denote the n objects to partition by the integers  $1, 2, \dots, n$ .  $\gcd(a, b) = \sum_{a} \phi(k)$ Define the reduced Stirling numbers of the second kind, denoted  $S^d(n,k)$ , to be the number of ways to partition the inments in each subset have pairwise distance at least d. That  $\overline{i=1}$ The number of ways to connect the 2n points on a circle to is, for any integers i and j in a given subset, it is required that  $|i-j| \geq d$ . It has been shown that these numbers satisfy,

$$\binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}$$

$$\sum_{i=0}^{k} \binom{n+i}{i} = \sum_{i=0}^{k} \binom{n+i}{n} = \binom{n+k+1}{k}$$

$$n, r \in N, n > r, \sum_{i=r}^{n} \binom{i}{r} = \binom{n+1}{r+1}$$

If 
$$P(n) = \sum_{k=0}^{n} {n \choose k} \cdot Q(k)$$
, then,

$$Q(n) = \sum_{k=0}^{n} (-1)^{n-k} \binom{n}{k} \cdot P(k)$$

If 
$$P(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot Q(k)$$
, then,

$$Q(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot P(k)$$

## 8.5 Different Math Formulas

Picks Theorem: A = i + b/2 - 1

**Deragements:**  $d(i) = (i-1) \times (d(i-1) + d(i-2))$ 

$$\frac{n}{ab}$$
 -  $\left\{\frac{b\prime n}{a}\right\}$  -  $\left\{\frac{a\prime n}{b}\right\}$  +

The gcd is a multiplicative function in the following sense: An r-associated Stirling number of the second kind is the num- if  $a_1$  and  $a_2$  are relatively prime, then  $gcd(a_1 \cdot a_2, b) =$ 

For non-negative integers a and b, where a and b are not both zero,  $gcd(n^a - 1, n^b - 1) = n^{gcd(a,b)} - 1$  $\sum_{k=1}^{n} \gcd(k, n) = \sum_{d \mid n} d \cdot \phi\left(\frac{n}{d}\right)$  $\sum_{k=1}^{n} x^{\gcd(k,n)} = \sum_{l=1}^{n} x^{d} \cdot \phi\left(\frac{n}{d}\right)$  $\sum_{k=1}^{n} \frac{1}{\gcd(k,n)} = \sum_{\text{all}} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{1}{n} \sum_{\text{all}} d \cdot \phi(d)$  $\sum_{k=1}^{n} \frac{k}{\gcd(k,n)} = \frac{n}{2} \cdot \sum_{d|n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{n}{2} \cdot \frac{1}{n} \cdot \sum_{d|n} d \cdot \phi(d)$  $\sum_{k=1}^{n} \frac{n}{\gcd(k,n)} = 2 * \sum_{k=1}^{n} \frac{k}{\gcd(k,n)} - 1, \text{ for } n > 1$  $\left| \sum_{i=1}^{n} \sum_{j=1}^{n} [\gcd(i,j) = 1] = \sum_{d=1}^{n} \mu(d) \left\lfloor \frac{n}{d} \right\rfloor^{2} \right|$ 

$$\begin{split} &\sum_{i=1}^{n} \sum_{j=1}^{n} \gcd(i,j) = \sum_{d=1}^{n} \phi(d) \left\lfloor \frac{n}{d} \right\rfloor^{2} \\ &\sum_{i=1}^{n} \sum_{j=1}^{n} i \cdot j [\gcd(i,j) = 1] = \sum_{i=1}^{n} \phi(i) i^{2} \\ &F(n) = \sum_{i=1}^{n} \sum_{j=1}^{n} \operatorname{lcm}(i,j) = \sum_{i=1}^{n} \left( \frac{\left(1 + \left\lfloor \frac{n}{l} \right\rfloor\right) \left(\left\lfloor \frac{n}{l} \right\rfloor\right)}{2} \right)^{2} \sum_{i=1}^{n} \mu(d) l d \end{split}$$