${\bf IUT\ Serenity}, {\bf Islamic\ University\ of\ Technology}$

Contents		6	Math 1	6
			1	16
1		1	6.2 Combi	16
		1		16
	1.2 RandomGenerator	2	6.4 Pollard Rho	17
_				۱7
2	DP 2			17
		2	6.7 FFT	17
		2	6.8 NTT	18
	1 1	3	6.9 WalshHadamard 1	19
	1	3	6.10 Adaptive Simpsons	19
	2.5 Li Chao Tree	3	r	19
	2.6 SOS-DP	4	6.12 Fractional Binary Search 2	20
			6.13 Lagrange	20
3	Data Structure 4	4	6.14 Shanks' Baby Step, Giant Step 2	20
		4	6.15 Xor Basis	20
	3.2 Persistent Segment Tree 4	4		
	3.3 Implicit Segment Tree	4 7	String 2	
	3.4 HashTable	5		21
	3.5 DSU With Rollbacks	5		21
	3.6 Binary Trie	5		22
	3.7 BIT-2D	5		22
	3.8 Divide And Conquer Query Offline 5	5	ě	22
		6	v	22
		6		23
			7.8 Z Algo	23
4	Geometry 7		Eti	
	4.1 Point	7 8	Equations and Formulas 2	
	4.2 Linear	7		24
	4.3 Circular	8	9	24
	4.4 Convex	9	9	24
	4.5 Polygon	1		24
				24
5	•		8.6 GCD and LCM	24
	5.1 Graph Template	2		
	5.2 LCA	2		
	5.3 SCC	2		
	5.4 Euler Tour on Edge	3		
	5.5 LCA In O(1)	3		
	5.6 HLD	3		
	5.7 Centroid Decomposition	4		
	5.8 Dinic Max Flow			
	5.9 Min Cost Max Flow			
	5.10 Block Cut Tree			
	5.11 Bridge Tree			
	5.12 Tree Isomorphism			

```
Sublime Build
    "cmd" : ["g++ -std=c++14 -DSONIC $file_name -o
        $file_base_name && timeout 4s ./$file_base_name<</pre>
        inputf.in>outputf.in"],
    "selector" : "source.cpp",
    "file_regex": "^(..[^:]*):([0-9]+):?([0-9]+)?:? (.*)$
    "shell": true.
    "working_dir" : "$file_path"
}
vimrc
" Auto import & Compile
:autocmd BufNewFile *.cpp Or ~/template.cpp
autocmd filetype cpp nnoremap <F9> :w \| !g++ -std=gnu++2
    a -Wall -g % -o %:r && ./%:r<CR>
autocmd filetype cpp inoremap <F9> <ESC>:w \| !g++ -std=
    gnu++2a -Wall -g % -o %:r && ./%:r<CR>
autocmd filetype cpp nnoremap <F10> :w \| !g++ -std=c++17
     -Wall -g -DMUSHFIQ_DEBUG -DDBG_MACRO_NO_WARNING % -
    o %:r && ./%:r<CR>
autocmd filetype cpp inoremap <F10> <ESC>:w \| !g++ -std=
    c++17 -Wall -g -DMUSHFIQ_DEBUG -
    DDBG_MACRO_NO_WARNING % -o %:r && ./%:r<CR>
autocmd filetype python nnoremap <F9> :w \| !python3 %<CR
autocmd filetype python inoremap <F9> <ESC>:w \| !python3
     %<CR>
" Auto Completion
inoremap ( ()<left>
inoremap <expr> ) strpart(getline('.'), col('.')-1, 1) ==
     ")" ? "\<Right>" : ")"
inoremap { {}<left>
inoremap <expr> } strpart(getline('.'), col('.')-1, 1) ==
     "}" ? "\<Right>" : "}"
inoremap [ []<left>
inoremap <expr> ] strpart(getline('.'), col('.')-1, 1) ==
     "]" ? "\<Right>" : "]"
inoremap <expr> " strpart(getline('.'), col('.')-1, 1) ==
     "\"" ? "\<Right>" : "\"\"\<left>"
inoremap <expr> ' strpart(getline('.'), col('.')-1, 1) ==
     "\'" ? "\<Right>" : "\'\'\<left>"
inoremap <expr> <CR> <sid>insert_newline()
function s:insert_newline() abort
 let pair = strpart(getline('.'), col('.')-2, 2)
  return stridx('(){}[]', pair) % 2 == 0 && strlen(pair)
      == 2 ? "\<CR>\<ESC>\O" : "\<CR>"
```

```
endfunction
inoremap <expr> <space> <sid>insert_space()
function s:insert_space() abort
 let pair = strpart(getline('.'), col('.')-2, 2)
 return stridx('(){}[]', pair) % 2 == 0 && strlen(pair)
      == 2 ? "\<space>\<left>" : "\<space>"
endfunction
inoremap <expr> <bs> <sid>rm_pair()
function s:rm_pair() abort
let pair = strpart(getline('.'), col('.')-2, 2)
return stridx('(){}[]'''', pair) % 2 == 0 && strlen(
     pair) == 2 ? "\<del>\<bs>" : "\<bs>"
endfunction
set nocompatible
                           " be iMproved, required
                          " required
filetype on
filetype plugin on
filetype plugin indent on
syntax on
set mouse=a
set number
set relativenumber
set tabstop=2
set shiftwidth=2
set expandtab
set softtabstop=2
set smartindent
set smarttab
set autoindent
set cindent
set noerrorbells
set ruler
set guifont=*
set backspace=indent,eol,start
set ignorecase
set incsearch
set nowrap
set hlsearch
  set termguicolors
set foldmethod=indent
set nofoldenable
 set cursorline
set laststatus=2
set showcmd
set wildmenu
set clipboard=unnamedplus
set ttymouse=xterm2
```

```
nnoremap \langle S-k \rangle : m .-2 \langle CR \rangle ==
vnoremap <S-j> :m '>+1<CR>gv==gv
vnoremap <S-k> :m '<-2<CR>gv==gv
Stress-tester
#!/bin/bash
# Call as stresstester GENERATOR SOL1 SOL2 ITERATIONS [--
    count]
for i in $(seq 1 "$4"); do
    [[ $* == *--count* ]] && echo "Attempt $i/$4"
   1 > in.txt
   $2 < in.txt > out1.txt
   $3 < in.txt > out2.txt
   diff -y out1.txt out2.txt > diff.txt
   if [ $? -ne 0 ] ; then
       echo "Differing Testcase Found:"; cat in.txt
       echo -e "\nOutputs:"; cat diff.txt
       break
   fi
done
```

1 All Misc

1.1 All Macros

nnoremap $\langle S-j \rangle : m .+1 \langle CR \rangle ==$

```
//#pragma GCC optimize("Ofast")
//#pragma GCC optimization ("03")
//#pragma comment(linker, "/stack:200000000")
//#pragma GCC optimize("unroll-loops")
//#pragma GCC target("sse,sse2,sse3,sse3,sse4,popcnt,abm
    ,mmx,avx,tune=native")
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
   //find_by_order(k) --> returns iterator to the kth
        largest element counting from 0
   //order_of_key(val) --> returns the number of items
        in a set that are strictly smaller than our item
template <typename DT>
using ordered_set = tree <DT, null_type, less<DT>,
    rb_tree_tag, tree_order_statistics_node_update>;
/*--- DEBUG TEMPLATE STARTS HERE ---*/
#ifdef Serenity
void show(int x) {cerr << x;}</pre>
void show(long long x) {cerr << x;}</pre>
void show(double x) {cerr << x;}</pre>
void show(char x) {cerr << '\'' << x << '\'';}</pre>
void show(const string &x) {cerr << '\"' << x << '\"';}</pre>
void show(bool x) {cerr << (x ? "true" : "false");}</pre>
```

```
template<typename T, typename V>
void show(pair<T, V> x) { cerr << '\f'; show(x.first);</pre>
    cerr << ", "; show(x.second); cerr << '}'; }</pre>
template<typename T>
void show(T x) {int f = 0; cerr << "{"; for (auto &i: x)</pre>
    cerr << (f++ ? ", " : ""), show(i); cerr << "}";}</pre>
void debug_out(string s) {
   cerr << '\n';
template <typename T, typename... V>
void debug_out(string s, T t, V... v) {
   s.erase(remove(s.begin(), s.end(), ''), s.end());
                   "; // 8 spaces
   cerr << "
   cerr << s.substr(0, s.find(','));</pre>
   s = s.substr(s.find(',') + 1);
   cerr << " = ";
   show(t);
   cerr << endl;
   if(sizeof...(v)) debug_out(s, v...);
#define debug(x...) cerr << "LINE: " << __LINE__ << endl;</pre>
     debug_out(#x, x); cerr << endl;</pre>
#else
#define debug(x...)
#endif
mt19937 rnd(chrono::steady_clock::now().time_since_epoch
    ().count());
uniform_int_distribution<LL>gen(1, 100);
gen(rnd);
shuffle(a.begin(), a.end(), rnd);
```

1.2 RandomGenerator

```
#include <bits/stdc++.h>
using namespace std;

random_device rand_dev;
mt19937 generator(rand_dev());

int main() {
   ios::sync_with_stdio(0);
   cin.tie(0);

   uniform_int_distribution<long long> distr(1, 20);
   uniform_int_distribution<long long> dq(1, 10);
   set<int> st;
```

```
int n = distr(generator);
for(int i=1; i<=n; i++) st.insert(distr(generator));

cout<<st.size()<<endl;
uniform_int_distribution<long long> gg(1, st.size());

for(int x : st) cout<<x<' ';
cout << endl;

int q = dq(generator);
cout<<q<<endl;
while(q--){
   int l = gg(generator);
   int r = gg(generator);
   int r = gg(generator);
   int k = gg(generator);
   cout<< l << ' '<<r<' '<<k<<endl;
}</pre>
```

2 DP

2.1 1D-1D

```
/// Author: anachor
#include <bits/stdc++.h>
using namespace std;
/// Solves dp[i] = min(dp[j] + cost(j+1, i)) given that
    cost() is QF
long long solve1D(int n, long long cost(int, int)) {
 vector<long long> dp(n + 1), opt(n + 1);
  deque<pair<int, int>> dq;
  dq.push_back({0, 1});
 dp[0] = 0;
  for (int i = 1; i <= n; i++) {
   opt[i] = dq.front().first;
   dp[i] = dp[opt[i]] + cost(opt[i] + 1, i);
   if (i == n) break;
   dq[0].second++;
   if (dq.size() > 1 \&\& dq[0].second == dq[1].second) dq
        .pop_front();
   int en = n;
   while (dq.size()) {
     int o = dq.back().first, st = dq.back().second;
     if (dp[o] + cost(o + 1, st) >= dp[i] + cost(i + 1,
         st))
```

dq.pop_back();

```
else {
                              int lo = st, hi = en;
                              while (lo < hi) {</pre>
                                      int mid = (lo + hi + 1) / 2;
                                       if (dp[o] + cost(o + 1, mid) < dp[i] + cost(i +</pre>
                                                         1. mid))
                                              lo = mid;
                                       else
                                              hi = mid - 1;
                              if (lo < n) dq.push_back({i, lo + 1});</pre>
                              break;
                      en = st -1;
               if (dq.empty()) dq.push_back({i, i + 1});
      return dp[n];
/// Solves https://open.kattis.com/problems/
                  coveredwalkway
const int N = 1e6 + 7;
long long x[N];
long long cost(int 1, int r) { return (x[r] - x[1]) * (x[r] + x[r]) * (x[r] 
                 r] - x[1]) + c: }
int main() {
       ios::sync_with_stdio(false);
       cin.tie(0);
       int n;
       cin >> n >> c;
       for (int i = 1; i \le n; i++) cin >> x[i];
       cout << solve1D(n, cost) << endl;</pre>
```

2.2 Convex Hull Trick

```
struct line {
    ll m, c;
    line() {}
    line(ll m, ll c) : m(m), c(c) {}
};
struct convex_hull_trick {
    vector<line> lines;
    int ptr = 0;
    convex_hull_trick() {}
    bool bad(line a, line b, line c) {
```

```
return 1.0 * (c.c - a.c) * (a.m - b.m) < 1.0 * (b.c - const int K = 805, N = 4005;
         a.c) * (a.m - c.m);
  void add(line L) {
   int sz = lines.size();
   while (sz >= 2 && bad(lines[sz - 2], lines[sz - 1], L
    lines.pop_back();
     sz--;
   }
   lines.pb(L);
 ll get(int idx, int x) { return (1ll * lines[idx].m * x
       + lines[idx].c): }
 11 query(int x) {
   if (lines.empty()) return 0;
   if (ptr >= lines.size()) ptr = lines.size() - 1;
   while (ptr < lines.size() - 1 && get(ptr, x) > get(
        ptr + 1, x)) ptr++;
   return get(ptr, x);
};
11 sum[MAX];
11 dp[MAX];
int arr[MAX];
int main() {
 fastio:
 int t;
  cin >> t:
  while (t--) {
   int n, a, b, c;
   cin >> n >> a >> b >> c;
   for (int i = 1; i <= n; i++) cin >> sum[i];
   for (int i = 1; i <= n; i++) dp[i] = 0, sum[i] += sum
        [i - 1];
   convex_hull_trick cht;
   cht.add(line(0, 0));
   for (int pos = 1; pos <= n; pos++) {</pre>
     dp[pos] = cht.query(sum[pos]) - 111 * a * sqr(sum[
         pos]) - c;
     cht.add(line(211 * a * sum[pos], dp[pos] - a * sqr(
         sum[pos])));
   11 \text{ ans} = (-111 * dp[n]);
   ans += (111 * sum[n] * b);
   cout << ans << "\n";
```

2.3 Divide and Conquer dp

```
LL dp[2][N], _cost[N][N];
// 1-indexed for convenience
LL cost(int 1, int r) {
 return _cost[r][r] - _cost[1 - 1][r] - _cost[r][1 - 1]
      + cost[1 - 1][1 - 1] >> 1:
void compute(int cnt, int 1, int r, int optl, int optr) {
 if (1 > r) return;
 int mid = 1 + r >> 1:
 LL best = INT MAX:
  int opt = -1;
 for (int i = optl; i <= min(mid, optr); i++) {</pre>
   LL cur = dp[cnt ^1][i - 1] + cost(i, mid);
   if (cur < best) best = cur, opt = i;</pre>
  dp[cnt][mid] = best;
  compute(cnt, 1, mid - 1, optl, opt);
  compute(cnt, mid + 1, r, opt, optr);
LL dnc_dp(int k, int n) {
 fill(dp[0] + 1, dp[0] + n + 1, INT_MAX);
 for (int cnt = 1; cnt <= k; cnt++) {</pre>
   compute(cnt & 1, 1, n, 1, n);
 return dp[k & 1][n];
```

2.4 Knuth optimization

```
const int N = 1005;
LL dp[N][N], a[N];
int opt[N][N];
LL cost(int i, int j) { return a[j + 1] - a[i]; }
LL knuth_optimization(int n) {
 for (int i = 0; i < n; i++) {
   dp[i][i] = 0;
   opt[i][i] = i;
 for (int i = n - 2; i \ge 0; i--) {
   for (int j = i + 1; j < n; j++) {
     LL mn = LLONG_MAX;
     LL c = cost(i, j);
     for (int k = opt[i][j - 1]; k <= min(j - 1, opt[i +</pre>
          1][j]); k++) {
       if (mn > dp[i][k] + dp[k + 1][i] + c) {
         mn = dp[i][k] + dp[k + 1][j] + c;
         opt[i][j] = k;
     dp[i][j] = mn;
```

```
return dp[0][n - 1];
2.5 Li Chao Tree
struct line {
 LL m. c:
 line(LL m = 0, LL c = 0) : m(m), c(c) {}
LL calc(line L, LL x) { return 1LL * L.m * x + L.c; }
struct node {
 LL m, c;
 line L;
 node *lft, *rt;
 node(LL m = 0, LL c = 0, node *lft = NULL, node *rt =
     : L(line(m, c)), lft(lft), rt(rt) {}
struct LiChao {
 node *root:
 LiChao() { root = new node(); }
 void update(node *now, int L, int R, line newline) {
   int mid = L + (R - L) / 2;
   line lo = now->L, hi = newline;
   if (calc(lo, L) > calc(hi, L)) swap(lo, hi);
   if (calc(lo, R) <= calc(hi, R)) {</pre>
     now->L = hi:
     return;
```

if (calc(lo, mid) < calc(hi, mid)) {</pre> now->L = hi:if (now->rt == NULL) now->rt = new node(); update(now->rt, mid + 1, R, lo); } else { now->L = lo:if (now->lft == NULL) now->lft = new node(); update(now->lft, L, mid, hi); } LL query(node *now, int L, int R, LL x) { if (now == NULL) return -inf: int mid = L + (R - L) / 2; $if (x \le mid)$ return max(calc(now->L, x), query(now->lft, L, mid, x)); else

return max(calc(now->L, x), query(now->rt, mid + 1,

R, x);

}

2.6 SOS-DP for(int i = 0; i<(1<<N); ++i) F[i] = A[i]; for(int i = 0;i < N; ++i) for(int mask = 0; mask < (1<<N) ; ++mask){ if(mask & (1<<i)) F[mask] += F[mask^(1<<i)]; }</pre>

3 Data Structure

3.1 Segment Tree const int N = 1e5+7;

int a[N]:

};

```
LL tr[4*N]:
LL lz[4*N];
///1. Merge left and right
LL combine (LL left, LL right) {
 return left + right;
}
///2. Push lazy down and merge lazy
void propagate(int u, int st, int en) {
  if (!lz[u]) return;
  tr[u] += (en-st+1)*lz[u];
  if (st!=en) {
   1z[2*u] += 1z[u]:
   lz[2*u+1] += lz[u];
 }
  1z[u] = 0;
void build(int u, int st, int en) {
  if (st==en) {
    tr[u] = a[st];
                         ///3. Initialize
   lz[u] = 0:
  else {
    int mid = (st+en)/2;
   build(2*u, st, mid);
    build(2*u+1, mid+1, en);
    tr[u] = combine(tr[2*u], tr[2*u+1]);
   lz[u] = 0;
                         ///3. Initialize
 }
}
void update(int u, int st, int en, int l, int r, int x) {
```

```
propagate(u, st, en);
 if (r<st || en<1) return;</pre>
  else if (l<=st && en<=r) {</pre>
   lz[u] += x;
                          ///4. Merge lazy
   propagate(u, st, en);
 else {
   int mid = (st+en)/2:
   update(2*u, st, mid, 1, r, x);
   update(2*u+1, mid+1, en, l, r, x);
   tr[u] = combine(tr[2*u], tr[2*u+1]);
LL query(int u, int st, int en, int 1, int r) {
 propagate(u, st, en);
if (r<st || en<1) return 0;</pre>
                              /// 5. Proper null value
else if (1<=st && en<=r) return tr[u];
else {
 int mid = (st+en)/2;
 return combine(query(2*u, st, mid, 1, r), query(2*u+1,
      mid+1, en, l, r));
```

3.2 Persistent Segment Tree

```
struct Node{
 int val,left,right;
int a[500010],cnt=0,versions[500010];
Node tree[25*N+15]:
int build(int b,int e){
 cnt++:
 int curr=cnt:
 if(b==e){}
   tree[curr].val=1e9+7;
   return curr;
 int mid=(b+e)/2:
  tree[curr].left=build(b,mid);
  tree[curr].right=build(mid+1,e);
  tree[curr].val=min(tree[tree[curr].left].val,tree[tree[
      curr].right].val);
 return curr:
int update(int prev,int b,int e,int idx,int x){
 cnt++:
 int curr=cnt;
 if(b==e){
   tree[curr].val=x:
```

```
return curr:
 int mid=(b+e)/2;
 if(idx<=mid){</pre>
   tree[curr].left=update(tree[prev].left,b,mid,idx,x);
   tree[curr].right=tree[prev].right;
 else{
   tree[curr].left=tree[prev].left;
   tree[curr].right=update(tree[prev].right,mid+1,e,idx,
        x);
 tree[curr].val=min(tree[tree[curr].left].val,tree[tree[
      curr].right].val);
 return curr;
int query(int curr,int b,int e,int l,int r){
 if(e<1 || b>r) return 1e9+7;
 if(b>=1 && e<=r) return tree[curr].val;</pre>
 int mid=(b+e)/2;
 int q1=query(tree[curr].left,b,mid,l,r);
 int q2=query(tree[curr].right,mid+1,e,l,r);
 return min(q1,q2);
```

3.3 Implicit Segment Tree

```
struct node {
 int val:
node *lft, *rt;
 node() {}
node(int val = 0) : val(val), lft(NULL), rt(NULL) {}
struct implicit_segtree {
 node *root;
 implicit_segtree() {}
 implicit_segtree(int n) {
   root = new node(n);
 void update(node *now, int L, int R, int idx, int val)
   if (L == R) {
     now -> val += val;
     return;
   int mid = L + (R - L) / 2;
   if (now->lft == NULL) now->lft = new node(mid - L +
       1);
   if (now->rt == NULL) now->rt = new node(R - mid);
   if (idx <= mid) update(now->lft, L, mid, idx, val);
```

3.4 HashTable

```
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;

const int RANDOM = chrono::high_resolution_clock::now().
    time_since_epoch().count();
unsigned hash_f(unsigned x) {
    x = ((x >> 16) ^ x) * 0x45d9f3b;
    x = ((x >> 16) ^ x) * 0x45d9f3b;
    return x = (x >> 16) ^ x;
}

unsigned hash_combine(unsigned a, unsigned b) { return a
    * 31 + b; }

struct chash {
    int operator()(int x) const { return hash_f(x); }
};

typedef gp_hash_table<int, int, chash> gp;
gp table;
```

3.5 DSU With Rollbacks

```
struct Rollback_DSU {
  int n;
  vector<int> par, sz;
  vector<pair<int, int>> op;
  Rollback_DSU(int n) : par(n), sz(n, 1) {
    iota(par.begin(), par.end(), 0);
    op.reserve(n);
  }
  int Anc(int node) {
    for (; node != par[node]; node = par[node])
        ; // no path compression
    return node;
```

```
int Anc(int node) { // path compression, doesnt work
    will rollback
  if(par[node] == node) return node;
  return par[node] = Anc(par[node]);
}

void Unite(int x, int y) {
  if (sz[x = Anc(x)] < sz[y = Anc(y)]) swap(x, y);
  op.emplace_back(x, y);
  par[y] = x;
  sz[x] += sz[y];
}

void Undo(int t) {
  for (; op.size() > t; op.pop_back()) {
    par[op.back().second] = op.back().second;
    sz[op.back().first] -= sz[op.back().second];
}
}
;;
```

3.6 Binary Trie

int tc = 1:

struct node {

int vis = 0;

const int N = 1e7 + 5, b = 30;

```
int to [2] = \{0, 0\};
 int val[2] = {0, 0};
  void update() {
   to[0] = to[1] = 0;
   val[0] = val[1] = 0;
   vis = tc:
 }
T[N + 2];
node *root = T;
int ptr = 0;
node *nxt(node *cur, int x) {
 if (cur->to[x] == 0) cur->to[x] = ++ptr;
  assert(ptr < N);</pre>
  int idx = cur->to[x]:
  if (T[idx].vis < tc) T[idx].update();</pre>
  return T + idx;
int query(int j, int aj) {
 int ans = 0, jaj = j ^ aj;
 node *cur = root;
  for (int k = b - 1; k = -1) {
   maximize(ans, nxt(cur, (jaj >> k & 1) ^ 1)->val[1 ^ (
        aj >> k & 1)]);
   cur = nxt(cur, (jaj >> k & 1));
```

```
return ans;
}
void insert(int j, int aj, int val) {
  int jaj = j ^ aj;
  node *cur = root;
  for (int k = b - 1; ~k; k--) {
    cur = nxt(cur, (jaj >> k & 1));
    maximize(cur->val[j >> k & 1], val);
  }
}
void clear() {
  tc++;
  ptr = 0;
  root->update();
}
```

3.7 BIT-2D

3.8 Divide And Conquer Query Offline

```
namespace up {
int l[N], r[N], u[N], v[N], tm;
void push(int _l, int _r, int _u, int _v) {
    l[tm] = _l, r[tm] = _r, u[tm] = _u, v[tm] = _v;
    tm++;
}
} // namespace up
namespace que {
int node[N], tm;
LL ans[N];
void push(int _node) { node[++tm] = _node; }
} // namespace que
namespace edge_set {
void push(int i) { dsu ::merge(up ::u[i], up ::v[i]); }
```

```
void pop(int t) { dsu ::rollback(t); }
int time() { return dsu ::op.size(); }
LL query(int u) { return a[dsu ::root(u)]; }
} // namespace edge_set
namespace dncq {
vector<int> tree[4 * N];
void update(int idx, int 1 = 0, int r = que ::tm, int
    node = 1) {
 int ul = up ::1[idx], ur = up ::r[idx];
  if (r 
  if (ul <= l and r <= ur) {</pre>
   tree[node].push_back(idx);
   return;
 }
 int m = 1 + r >> 1;
  update(idx, 1, m, node << 1);
 update(idx, m + 1, r, node \langle\langle 1 | 1 \rangle\rangle;
void dfs(int 1 = 0, int r = que ::tm, int node = 1) {
 int cur = edge_set ::time();
 for (int e : tree[node]) edge_set ::push(e);
  if (1 == r) {
   que ::ans[1] = edge_set ::query(que ::node[1]);
 } else {
   int m = 1 + r >> 1;
   dfs(1, m, node << 1);
   dfs(m + 1, r, node << 1 | 1);
  edge_set ::pop(cur);
} // namespace dncg
void push_updates() {
 for (int i = 0; i < up ::tm; i++) dncq ::update(i);</pre>
```

3.9 MO with Update

```
int qi = 0. ui = 0:
int 1 = 1, r = 0, t = 0;
void add(int idx) {
 --cnt[freq[arr[idx]]];
 freq[arr[idx]]++;
  cnt[freq[arr[idx]]]++;
void remove(int idx) {
  --cnt[freq[arr[idx]]];
  freq[arr[idx]]--;
  cnt[freq[arr[idx]]]++;
void apply(int t) {
  const bool f = 1 <= Up[t].idx and Up[t].idx <= r;</pre>
  if (f) remove(Up[t].idx);
  arr[Up[t].idx] = Up[t].val;
  if (f) add(Up[t].idx);
void undo(int t) {
  const bool f = 1 <= Up[t].idx and Up[t].idx <= r;</pre>
  if (f) remove(Up[t].idx);
  arr[Up[t].idx] = Up[t].last;
  if (f) add(Up[t].idx);
int mex() {
 for (int i = 1; i <= N; i++)</pre>
   if (!cnt[i]) return i;
  assert(0);
int main() {
  int n, q;
  cin >> n >> q;
  int counter = 0;
  map<int, int> M;
  for (int i = 1; i <= n; i++) {</pre>
    cin >> arr[i]:
    if (!M[arr[i]]) M[arr[i]] = ++counter;
   arr[i] = M[arr[i]];
 }
  iota(id, id + N, 0);
  while (q--) {
   int tp, x, y;
    cin >> tp >> x >> y;
    if (tp == 1)
     Q[++qi] = query(x, y, ui);
    else {
     if (!M[y]) M[y] = ++counter;
     v = M[v];
     Up[++ui] = \{x, y, arr[x]\};
```

```
arr[x] = y;
}
t = ui;
cnt[0] = 3 * n;
sort(id + 1, id + qi + 1, [&](int x, int y) { return Q[
    x] < Q[v]; \});
for (int i = 1; i <= qi; i++) {</pre>
  int x = id[i];
  while (Q[x].t > t) apply(++t);
  while (Q[x].t < t) undo(t--);
  while (Q[x].1 < 1) add(--1);
  while (Q[x].r > r) add(++r);
  while (Q[x].1 > 1) remove(1++);
  while (Q[x].r < r) remove(r--);
  ans[x] = mex():
}
for (int i = 1; i <= qi; i++) cout << ans[i] << '\n';</pre>
```

3.10 SparseTable (Rectangle Query)

```
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 505;
const int LOGN = 9;
// O(n^2 (logn)^2
// Supports Rectangular Query
int A[MAXN][MAXN];
int M[MAXN][MAXN][LOGN][LOGN];
void Build2DSparse(int N) {
 for (int i = 1; i <= N; i++) {</pre>
   for (int j = 1; j <= N; j++) {</pre>
     M[i][j][0][0] = A[i][j];
   for (int q = 1; (1 << q) <= N; q++) {
     int add = 1 << (q - 1);
     for (int j = 1; j + add <= N; j++) {
       M[i][j][0][q] = max(M[i][j][0][q - 1], M[i][j +
           add][0][q - 1]);
   }
 }
 for (int p = 1; (1 << p) <= N; p++) {
   int add = 1 << (p - 1);
   for (int i = 1; i + add <= N; i++) {</pre>
```

for (int q = 0; (1 << q) <= \mathbb{N} ; q++) {

```
for (int j = 1; j <= N; j++) {
         M[i][j][p][q] = max(M[i][j][p - 1][q], M[i + add]
             ][j][p - 1][q]);
      }
     }
   }
    <=y2
int Query(int x1, int y1, int x2, int y2) {
 int kX = log2(x2 - x1 + 1);
 int kY = log2(y2 - y1 + 1);
 int addX = 1 << kX;</pre>
 int addY = 1 << kY;</pre>
 int ret1 = max(M[x1][y1][kX][kY], M[x1][y2 - addY + 1][ | Tf angleBetween(Point a, Point b) {
      kX][kY]);
 int ret2 = max(M[x2 - addX + 1][y1][kX][kY],
               M[x2 - addX + 1][y2 - addY + 1][kX][kY]);
 return max(ret1, ret2);
```

4 Geometry

4.1 Point

```
typedef double Tf;
typedef double Ti; /// use long long for exactness
const Tf PI = acos(-1), EPS = 1e-9;
int dcmp(Tf x) \{ return abs(x) < EPS ? 0 : (x < 0 ? -1 : 
    1): }
struct Point {
 Ti x, y;
 Point(Ti x = 0, Ti y = 0) : x(x), y(y) {}
 Point operator+(const Point& u) const { return Point(x
     + u.x, y + u.y); }
 Point operator-(const Point& u) const { return Point(x
      -u.x, y - u.y); }
 Point operator*(const LL u) const { return Point(x * u,
      y * u); }
 Point operator*(const Tf u) const { return Point(x * u,
      v * u): }
 Point operator/(const Tf u) const { return Point(x / u,
      y / u); }
 bool operator==(const Point& u) const {
   return dcmp(x - u.x) == 0 && dcmp(y - u.y) == 0;
```

```
bool operator!=(const Point& u) const { return !(*this
                                                                 == u): }
                                                            bool operator<(const Point& u) const {</pre>
                                                              return dcmp(x - u.x) < 0 \mid \mid (dcmp(x - u.x) == 0 \&\&
                                                                  dcmp(y - u.y) < 0);
                                                          };
                                                          Ti dot(Point a, Point b) { return a.x * b.x + a.y * b.y;
// returns max of all A[i][j], where x1<=i<=x2 and y1<=j | Ti cross(Point a, Point b) { return a.x * b.y - a.y * b.x
                                                          Tf length(Point a) { return sqrt(dot(a, a)); }
                                                          Ti sqLength(Point a) { return dot(a, a); }
                                                          Tf distance(Point a, Point b) { return length(a - b); }
                                                          Tf angle(Point u) { return atan2(u.y, u.x); }
                                                           // returns angle between oa, ob in (-PI, PI]
                                                            Tf ans = angle(b) - angle(a);
                                                            return ans <= -PI ? ans + 2 * PI : (ans > PI ? ans - 2
                                                                 * PI : ans):
                                                           // Rotate a ccw by rad radians, Tf Ti same
                                                           Point rotate(Point a, Tf rad) {
                                                            return Point(a.x * cos(rad) - a.y * sin(rad),
                                                                        a.x * sin(rad) + a.y * cos(rad));
                                                           // rotate a ccw by angle th with cos(th) = co && sin(th)
                                                               = si. tf ti same
                                                           Point rotatePrecise(Point a, Tf co, Tf si) {
                                                            return Point(a.x * co - a.y * si, a.y * co + a.x * si);
                                                          Point rotate90(Point a) { return Point(-a.y, a.x); }
                                                           // scales vector a by s such that length of a becomes s,
                                                               Tf Ti same
                                                          Point scale(Point a, Tf s) { return a / length(a) * s; }
                                                           // returns an unit vector perpendicular to vector a, Tf
                                                               Ti same
                                                           Point normal(Point a) {
                                                            Tf 1 = length(a);
                                                            return Point(-a.y / 1, a.x / 1);
                                                           // returns 1 if c is left of ab, 0 if on ab && -1 if
                                                               right of ab
                                                           int orient(Point a, Point b, Point c) { return dcmp(cross
                                                               (b - a, c - a)); }
                                                           /// Use as sort(v.begin(), v.end(), polarComp(0, dir))
                                                          /// Polar comparator around O starting at direction dir
                                                           struct polarComp {
                                                            Point O, dir;
```

```
polarComp(Point 0 = Point(0, 0), Point dir = Point(1,
     0)) : 0(0), dir(dir) {}
 bool half(Point p) {
   return dcmp(cross(dir, p)) < 0 ||</pre>
          (dcmp(cross(dir, p)) == 0 && dcmp(dot(dir, p))
              > 0):
 bool operator()(Point p, Point q) {
   return make_tuple(half(p), 0) < make_tuple(half(q),</pre>
       cross(p, q));
struct Segment {
 Point a. b:
 Segment(Point aa, Point bb) : a(aa), b(bb) {}
typedef Segment Line;
struct Circle {
 Point o;
 Tf r;
 Circle(Point o = Point(0, 0), Tf r = 0) : o(o), r(r) {}
 // returns true if point p is in || on the circle
 bool contains(Point p) { return dcmp(sqLength(p - o) -
      r * r) <= 0: 
 // returns a point on the circle rad radians away from
      +X CCW
 Point point(Tf rad) {
   static_assert(is_same<Tf, Ti>::value);
   return Point(o.x + cos(rad) * r, o.y + sin(rad) * r);
 // area of a circular sector with central angle rad
 Tf area(Tf rad = PI + PI) { return rad * r * r / 2; }
 // area of the circular sector cut by a chord with
      central angle alpha
 Tf sector(Tf alpha) { return r * r * 0.5 * (alpha - sin )
      (alpha)); }
```

4.2 Linear

```
// **** LINE LINE INTERSECTION START ****
// returns true if point p is on segment s
bool onSegment(Point p, Segment s) {
 return dcmp(cross(s.a - p, s.b - p)) == 0 && dcmp(dot(s
      .a - p, s.b - p)) <= 0;
// returns true if segment p && q touch or intersect
bool segmentsIntersect(Segment p, Segment q) {
 if (onSegment(p.a, q) || onSegment(p.b, q)) return true
```

```
if (onSegment(q.a, p) || onSegment(q.b, p)) return true
 Ti c1 = cross(p.b - p.a, q.a - p.a);
 Ti c2 = cross(p.b - p.a, q.b - p.a);
 Ti c3 = cross(q.b - q.a, p.a - q.a);
 Ti c4 = cross(q.b - q.a, p.b - q.a);
 return dcmp(c1) * dcmp(c2) < 0 && dcmp(c3) * dcmp(c4) <
bool linesParallel(Line p, Line q) {
 return dcmp(cross(p.b - p.a, q.b - q.a)) == 0;
// lines are represented as a ray from a point: (point,
// returns false if two lines (p, v) && (q, w) are
    parallel or collinear
// true otherwise, intersection point is stored at o via |}
    reference, Tf Ti Same
bool lineLineIntersection(Point p, Point v, Point q,
    Point w, Point& o) {
 if (dcmp(cross(v, w)) == 0) return false;
 Point u = p - q;
  o = p + v * (cross(w, u) / cross(v, w));
 return true;
// returns false if two lines p && q are parallel or
    collinear
// true otherwise, intersection point is stored at o via
bool lineLineIntersection(Line p, Line q, Point& o) {
 return lineLineIntersection(p.a, p.b - p.a, q.a, q.b -
      q.a, o);
// returns the distance from point a to line 1
// **** LINE LINE INTERSECTION FINISH ****
Tf distancePointLine(Point p, Line 1) {
 return abs(cross(1.b - 1.a, p - 1.a) / length(1.b - 1.a
      ));
}
// returns the shortest distance from point a to segment
Tf distancePointSegment(Point p, Segment s) {
  if (s.a == s.b) return length(p - s.a);
  Point v1 = s.b - s.a, v2 = p - s.a, v3 = p - s.b;
 if (dcmp(dot(v1, v2)) < 0)
   return length(v2);
  else if (dcmp(dot(v1, v3)) > 0)
   return length(v3);
  else
```

```
return abs(cross(v1, v2) / length(v1));
}
// returns the shortest distance from segment p to
    segment q
Tf distanceSegmentSegment(Segment p, Segment q) {
    if (segmentsIntersect(p, q)) return 0;
    Tf ans = distancePointSegment(p.a, q);
    ans = min(ans, distancePointSegment(p.b, q));
    ans = min(ans, distancePointSegment(q.a, p));
    ans = min(ans, distancePointSegment(q.b, p));
    return ans;
}
// returns the projection of point p on line l, Tf Ti
    Same
Point projectPointLine(Point p, Line l) {
    Point v = l.b - l.a;
    return l.a + v * ((Tf)dot(v, p - l.a) / dot(v, v));
}
```

4.3 Circular

```
// Extremely inaccurate for finding near touches
// compute intersection of line 1 with circle c
// The intersections are given in order of the ray (1.a,
    1.b), Tf Ti same
vector<Point> circleLineIntersection(Circle c, Line 1) {
 vector<Point> ret:
 Point b = 1.b - 1.a, a = 1.a - c.o;
 Tf A = dot(b, b), B = dot(a, b);
 Tf C = dot(a, a) - c.r * c.r, D = B * B - A * C;
 if (D < -EPS) return ret;</pre>
 ret.push_back(1.a + b * (-B - sqrt(D + EPS)) / A);
 if (D > EPS) ret.push_back(1.a + b * (-B + sqrt(D)) / A
      );
 return ret;
// signed area of intersection of circle(c.o, c.r) &&
// triangle(c.o, s.a, s.b) [cross(a-o, b-o)/2]
Tf circleTriangleIntersectionArea(Circle c, Segment s) {
 using Linear::distancePointSegment;
 Tf OA = length(c.o - s.a);
 Tf OB = length(c.o - s.b);
 // sector
 if (dcmp(distancePointSegment(c.o, s) - c.r) >= 0)
   return angleBetween(s.a - c.o, s.b - c.o) * (c.r * c.
       r) / 2.0;
 // triangle
 if (dcmp(OA - c.r) \le 0 \&\& dcmp(OB - c.r) \le 0)
   return cross(c.o - s.b, s.a - s.b) / 2.0;
 // three part: (A, a) (a, b) (b, B)
 vector<Point> Sect = circleLineIntersection(c, s);
```

```
return circleTriangleIntersectionArea(c, Segment(s.a,
      Sect[0])) +
        circleTriangleIntersectionArea(c, Segment(Sect
            [0], Sect[1])) +
        circleTriangleIntersectionArea(c, Segment(Sect
            [1], s.b));
// area of intersecion of circle(c.o, c.r) && simple
    polyson(p[])
Tf circlePolyIntersectionArea(Circle c, Polygon p) {
 Tf res = 0:
 int n = p.size();
 for (int i = 0; i < n; ++i)
   res += circleTriangleIntersectionArea(c, Segment(p[i
       ], p[(i + 1) % n]));
 return abs(res):
// locates circle c2 relative to c1
// interior
                     (d < R - r)
// interior tangents (d = R - r)
                                     ----> -1
// concentric
                   (d = 0)
// secants
                    (R - r < d < R + r) \longrightarrow 0
// exterior tangents (d = R + r)
                                     ----> 1
// exterior
                     (d > R + r)
                                        ----> 2
int circleCirclePosition(Circle c1, Circle c2) {
 Tf d = length(c1.o - c2.o);
 int in = dcmp(d - abs(c1.r - c2.r)), ex = dcmp(d - (c1.r))
     r + c2.r):
 return in < 0 ? -2 : in == 0 ? -1 : ex == 0 ? 1 : ex >
     0 ? 2 : 0:
// compute the intersection points between two circles c1
     && c2, Tf Ti same
vector<Point> circleCircleIntersection(Circle c1, Circle
    c2) {
 vector<Point> ret;
 Tf d = length(c1.o - c2.o):
 if (dcmp(d) == 0) return ret;
 if (dcmp(c1.r + c2.r - d) < 0) return ret;
 if (dcmp(abs(c1.r - c2.r) - d) > 0) return ret;
 Point v = c2.0 - c1.0:
 Tf co = (c1.r * c1.r + sqLength(v) - c2.r * c2.r) / (2)
      * c1.r * length(v));
 Tf si = sqrt(abs(1.0 - co * co));
 Point p1 = scale(rotatePrecise(v, co, -si), c1.r) + c1.
 Point p2 = scale(rotatePrecise(v, co, si), c1.r) + c1.o
```

```
ret.push_back(p1);
  if (p1 != p2) ret.push_back(p2);
  return ret;
}
// intersection area between two circles c1, c2
Tf circleCircleIntersectionArea(Circle c1, Circle c2) {
  Point AB = c2.o - c1.o;
 Tf d = length(AB);
  if (d >= c1.r + c2.r) return 0;
  if (d + c1.r <= c2.r) return PI * c1.r * c1.r;</pre>
  if (d + c2.r <= c1.r) return PI * c2.r * c2.r;</pre>
 Tf alpha1 = acos((c1.r * c1.r + d * d - c2.r * c2.r) /
      (2.0 * c1.r * d)):
 Tf alpha2 = acos((c2.r * c2.r + d * d - c1.r * c1.r) /
      (2.0 * c2.r * d)):
 return c1.sector(2 * alpha1) + c2.sector(2 * alpha2);
// returns tangents from a point p to circle c, Tf Ti
vector<Point> pointCircleTangents(Point p, Circle c) {
  vector<Point> ret;
  Point u = c.o - p;
  Tf d = length(u);
  if (d < c.r)
  else if (dcmp(d - c.r) == 0) {
   ret = {rotate(u, PI / 2)};
  } else {
   Tf ang = asin(c.r / d);
   ret = {rotate(u, -ang), rotate(u, ang)};
 return ret;
// returns the points on tangents that touches the circle
    , Tf Ti Same
vector<Point> pointCircleTangencyPoints(Point p, Circle c
    ) {
  Point u = p - c.o;
  Tf d = length(u);
  if (d < c.r)
   return {};
  else if (dcmp(d - c.r) == 0)
   return {c.o + u};
  else {
   Tf ang = acos(c.r / d);
   u = u / length(u) * c.r;
   return {c.o + rotate(u, -ang), c.o + rotate(u, ang)};
 }
}
```

```
// for two circles c1 && c2, returns two list of points a }
     && b
// such that a[i] is on c1 && b[i] is c2 && for every i
// Line(a[i], b[i]) is a tangent to both circles
// CAUTION: a[i] = b[i] in case they touch | -1 for c1 =
int circleCircleTangencyPoints(Circle c1, Circle c2,
    vector<Point> &a.
                            vector<Point> &b) {
  a.clear(), b.clear();
  int cnt = 0;
 if (dcmp(c1.r - c2.r) < 0) {
   swap(c1, c2);
   swap(a, b);
 Tf d2 = sqLength(c1.o - c2.o);
 Tf rdif = c1.r - c2.r, rsum = c1.r + c2.r;
  if (dcmp(d2 - rdif * rdif) < 0) return 0;</pre>
  if (dcmp(d2) == 0 \&\& dcmp(c1.r - c2.r) == 0) return -1;
  Tf base = angle(c2.o - c1.o);
  if (dcmp(d2 - rdif * rdif) == 0) {
   a.push_back(c1.point(base));
   b.push_back(c2.point(base));
   cnt++;
   return cnt;
 Tf ang = acos((c1.r - c2.r) / sqrt(d2));
  a.push_back(c1.point(base + ang));
  b.push_back(c2.point(base + ang));
  cnt++;
  a.push_back(c1.point(base - ang));
  b.push_back(c2.point(base - ang));
  cnt++;
  if (dcmp(d2 - rsum * rsum) == 0) {
   a.push_back(c1.point(base));
   b.push_back(c2.point(PI + base));
   cnt++;
 } else if (dcmp(d2 - rsum * rsum) > 0) {
   Tf ang = acos((c1.r + c2.r) / sqrt(d2));
   a.push_back(c1.point(base + ang));
   b.push_back(c2.point(PI + base + ang));
   a.push_back(c1.point(base - ang));
   b.push_back(c2.point(PI + base - ang));
   cnt++;
 return cnt;
```

```
4.4 Convex
/// minkowski sum of two polygons in O(n)
Polygon minkowskiSum(Polygon A, Polygon B) {
 int n = A.size(), m = B.size();
 rotate(A.begin(), min_element(A.begin(), A.end()), A.
      end());
 rotate(B.begin(), min_element(B.begin(), B.end()), B.
      end());
 A.push_back(A[0]);
 B.push_back(B[0]);
 for (int i = 0; i < n; i++) A[i] = A[i + 1] - A[i];
 for (int i = 0; i < m; i++) B[i] = B[i + 1] - B[i];</pre>
 Polygon C(n + m + 1);
 C[0] = A.back() + B.back();
 merge(A.begin(), A.end() - 1, B.begin(), B.end() - 1, C
      .begin() + 1,
       polarComp(Point(0, 0), Point(0, -1)));
 for (int i = 1; i < C.size(); i++) C[i] = C[i] + C[i -</pre>
 C.pop_back();
 return C;
// finds the rectangle with minimum area enclosing a
    convex polygon and
// the rectangle with minimum perimeter enclosing a
    convex polygon
// Tf Ti Same
pair<Tf, Tf> rotatingCalipersBoundingBox(const Polygon &p
 using Linear::distancePointLine;
 int n = p.size();
 int 1 = 1, r = 1, j = 1;
 Tf area = 1e100;
 Tf perimeter = 1e100;
 for (int i = 0; i < n; i++) {</pre>
   Point v = (p[(i + 1) \% n] - p[i]) / length(p[(i + 1)
       % n] - p[i]);
   while (dcmp(dot(v, p[r % n] - p[i]) - dot(v, p[(r +
        1) \% n] - p[i])) < 0)
     r++;
   while (j < r \mid | dcmp(cross(v, p[j % n] - p[i]) -
                       cross(v, p[(j + 1) % n] - p[i])) <
                            0)
     j++;
   while (1 < j ||
```

```
dcmp(dot(v, p[1 % n] - p[i]) - dot(v, p[(1 + 1)
               % n] - p[i])) > 0)
     1++;
   Tf w = dot(v, p[r \% n] - p[i]) - dot(v, p[1 \% n] - p[
       i]);
   Tf h = distancePointLine(p[j % n], Line(p[i], p[(i +
        1) % n]));
   area = min(area, w * h);
   perimeter = min(perimeter, 2 * w + 2 * h);
 return make_pair(area, perimeter);
// returns the left side of polygon u after cutting it by
     rav a->b
Polygon cutPolygon(Polygon u, Point a, Point b) {
  using Linear::lineLineIntersection;
  using Linear::onSegment;
  Polygon ret;
  int n = u.size();
  for (int i = 0; i < n; i++) {
   Point c = u[i], d = u[(i + 1) \% n];
   if (dcmp(cross(b - a, c - a)) >= 0) ret.push_back(c);
   if (dcmp(cross(b - a, d - c)) != 0) {
     Point t;
     lineLineIntersection(a, b - a, c, d - c, t);
     if (onSegment(t, Segment(c, d))) ret.push_back(t);
   }
 }
 return ret;
// returns true if point p is in or on triangle abc
bool pointInTriangle(Point a, Point b, Point c, Point p)
 return dcmp(cross(b - a, p - a)) >= 0 && dcmp(cross(c -
       b, p - b)) >= 0 &&
        dcmp(cross(a - c, p - c)) >= 0;
// pt must be in ccw order with no three collinear points
// returns inside = -1, on = 0, outside = 1
int pointInConvexPolygon(const Polygon &pt, Point p) {
 int n = pt.size();
  assert(n >= 3);
  int lo = 1, hi = n - 1;
  while (hi - lo > 1) {
   int mid = (lo + hi) / 2;
   if (dcmp(cross(pt[mid] - pt[0], p - pt[0])) > 0)
     lo = mid;
   else
```

```
hi = mid:
 bool in = pointInTriangle(pt[0], pt[lo], pt[hi], p);
 if (!in) return 1;
 if (dcmp(cross(pt[lo] - pt[lo - 1], p - pt[lo - 1])) == |}
       0) return 0:
 if (dcmp(cross(pt[hi] - pt[lo], p - pt[lo])) == 0)
      return 0:
 if (dcmp(cross(pt[hi] - pt[(hi + 1) % n], p - pt[(hi +
     1) % n (1) == 0)
   return 0;
 return -1;
// Extreme Point for a direction is the farthest point in
     that direction
// u is the direction for extremeness
int extremePoint(const Polygon &poly, Point u) {
 int n = (int)poly.size();
 int a = 0, b = n;
 while (b - a > 1) {
   int c = (a + b) / 2;
   if (dcmp(dot(poly[c] - poly[(c + 1) % n], u)) >= 0 &&
       dcmp(dot(poly[c] - poly[(c - 1 + n) % n], u)) >=
           0) {
     return c;
   bool a_up = dcmp(dot(poly[(a + 1) % n] - poly[a], u))
   bool c_{up} = dcmp(dot(poly[(c + 1) % n] - poly[c], u))
   bool a_above_c = dcmp(dot(poly[a] - poly[c], u)) > 0;
   if (a_up && !c_up)
     b = c:
   else if (!a_up && c_up)
   else if (a_up && c_up) {
     if (a_above_c)
       b = c;
     else
       a = c;
   } else {
     if (!a_above_c)
      b = c;
     else
       a = c;
```

```
}
 if (dcmp(dot(poly[a] - poly[(a + 1) % n], u)) > 0 &&
     dcmp(dot(poly[a] - poly[(a - 1 + n) % n], u)) > 0)
   return a:
 return b % n;
// For a convex polygon p and a line 1, returns a list of
// of p that touch or intersect line 1.
// the i'th segment is considered (p[i], p[(i + 1) modulo
// #1 If a segment is collinear with the line, only that
    is returned
// #2 Else if 1 goes through i'th point, the i'th segment
     is added
// Complexity: O(lg |p|)
vector<int> lineConvexPolyIntersection(const Polygon &p,
    Line 1) {
 assert((int)p.size() >= 3);
 assert(1.a != 1.b);
 int n = p.size();
 vector<int> ret;
 Point v = 1.b - 1.a;
 int lf = extremePoint(p, rotate90(v));
 int rt = extremePoint(p, rotate90(v) * Ti(-1));
 int olf = orient(l.a, l.b, p[lf]);
 int ort = orient(l.a, l.b, p[rt]);
 if (!olf || !ort) {
   int idx = (!olf ? lf : rt);
   if (orient(1.a, 1.b, p[(idx - 1 + n) \% n]) == 0)
     ret.push_back((idx - 1 + n) \% n);
     ret.push_back(idx);
   return ret;
 if (olf == ort) return ret;
 for (int i = 0; i < 2; ++i) {
   int lo = i ? rt : lf;
   int hi = i ? lf : rt;
   int olo = i ? ort : olf;
   while (true) {
     int gap = (hi - lo + n) % n;
     if (gap < 2) break;</pre>
```

```
int mid = (lo + gap / 2) % n;
     int omid = orient(l.a, l.b, p[mid]);
     if (!omid) {
       lo = mid;
       break;
     if (omid == olo)
       lo = mid:
     else
       hi = mid;
   ret.push_back(lo);
 return ret:
// Calculate [ACW, CW] tangent pair from an external
constexpr int CW = -1, ACW = 1;
bool isGood(Point u, Point v, Point Q, int dir) {
 return orient(Q, u, v) != -dir;
Point better(Point u, Point v, Point Q, int dir) {
 return orient(Q, u, v) == dir ? u : v;
}
Point pointPolyTangent(const Polygon &pt, Point Q, int
    dir, int lo, int hi) {
  while (hi - lo > 1) {
    int mid = (lo + hi) / 2;
   bool pvs = isGood(pt[mid], pt[mid - 1], Q, dir);
   bool nxt = isGood(pt[mid], pt[mid + 1], Q, dir);
    if (pvs && nxt) return pt[mid];
   if (!(pvs || nxt)) {
     Point p1 = pointPolyTangent(pt, Q, dir, mid + 1, hi
         );
     Point p2 = pointPolyTangent(pt, Q, dir, lo, mid -
         1):
     return better(p1, p2, Q, dir);
   }
   if (!pvs) {
     if (orient(Q, pt[mid], pt[lo]) == dir)
       hi = mid - 1;
     else if (better(pt[lo], pt[hi], Q, dir) == pt[lo])
       hi = mid - 1;
     else
       lo = mid + 1;
   }
   if (!nxt) {
     if (orient(Q, pt[mid], pt[lo]) == dir)
```

```
lo = mid + 1:
     else if (better(pt[lo], pt[hi], Q, dir) == pt[lo])
       hi = mid - 1;
     else
       lo = mid + 1;
   }
 }
 Point ret = pt[lo];
 for (int i = lo + 1; i <= hi; i++) ret = better(ret, pt</pre>
      [i], Q, dir);
 return ret;
// [ACW, CW] Tangent
pair<Point, Point> pointPolyTangents(const Polygon &pt,
    Point Q) {
 int n = pt.size();
 Point acw_tan = pointPolyTangent(pt, Q, ACW, 0, n - 1);
 Point cw_tan = pointPolyTangent(pt, Q, CW, 0, n - 1);
 return make_pair(acw_tan, cw_tan);
```

4.5 Polygon

```
typedef vector<Point> Polygon;
// removes redundant colinear points
// polygon can't be all colinear points
Polygon RemoveCollinear(const Polygon &poly) {
 Polygon ret;
 int n = poly.size();
 for (int i = 0; i < n; i++) {</pre>
   Point a = poly[i];
   Point b = poly[(i + 1) \% n];
   Point c = polv[(i + 2) \% n];
   if (dcmp(cross(b - a, c - b)) != 0 && (ret.empty() ||
         b != ret.back()))
     ret.push_back(b);
 }
 return ret;
// returns the signed area of polygon p of n vertices
Tf signedPolygonArea(const Polygon &p) {
 Tf ret = 0:
 for (int i = 0; i < (int)p.size() - 1; i++)</pre>
   ret += cross(p[i] - p[0], p[i + 1] - p[0]);
 return ret / 2;
// given a polygon p of n vertices, generates the convex
    hull in in CCW
// Tested on https://acm.timus.ru/problem.aspx?space=1&
    num=1185
```

```
// Caution: when all points are colinear AND
    removeRedundant == false
// output will be contain duplicate points (from upper
   hull) at back
Polygon convexHull(Polygon p, bool removeRedundant) {
 int check = removeRedundant ? 0 : -1;
 sort(p.begin(), p.end());
 p.erase(unique(p.begin(), p.end()), p.end());
 int n = p.size();
 Polygon ch(n + n);
 int m = 0; // preparing lower hull
 for (int i = 0; i < n; i++) {</pre>
   while (m > 1 &&
         dcmp(cross(ch[m-1]-ch[m-2], p[i]-ch[m-
               1])) <= check)
     m--;
   ch[m++] = p[i];
 int k = m; // preparing upper hull
 for (int i = n - 2; i \ge 0; i--) {
   while (m > k &&
         dcmp(cross(ch[m-1]-ch[m-2], p[i]-ch[m-
               21)) <= check)
     m--;
   ch[m++] = p[i];
 if (n > 1) m--;
 ch.resize(m):
 return ch;
// returns inside = -1, on = 0, outside = 1
int pointInPolygon(const Polygon &p, Point o) {
 using Linear::onSegment;
 int wn = 0, n = p.size();
 for (int i = 0; i < n; i++) {</pre>
   int j = (i + 1) \% n;
   if (onSegment(o, Segment(p[i], p[j])) || o == p[i])
       return 0;
   int k = dcmp(cross(p[i] - p[i], o - p[i]));
   int d1 = dcmp(p[i].y - o.y);
   int d2 = dcmp(p[j].y - o.y);
   if (k > 0 \&\& d1 \le 0 \&\& d2 > 0) wn++;
   if (k < 0 \&\& d2 \le 0 \&\& d1 > 0) wn--;
 return wn ? -1 : 1;
// Given a simple polygon p, and a line l, returns (x, y)
// x = longest segment of l in p, y = total length of l
    in p.
```

```
pair<Tf, Tf> linePolygonIntersection(Line 1, const
    Polygon &p) {
 using Linear::lineLineIntersection;
 int n = p.size();
 vector<pair<Tf, int>> ev;
 for (int i = 0; i < n; ++i) {</pre>
   Point a = p[i], b = p[(i + 1) \% n], z = p[(i - 1 + n)
        % n];
   int ora = orient(1.a, 1.b, a), orb = orient(1.a, 1.b,
       orz = orient(1.a, 1.b, z);
   if (!ora) {
     Tf d = dot(a - 1.a, 1.b - 1.a);
     if (orz && orb) {
       if (orz != orb) ev.emplace_back(d, 0);
       // else // Point Touch
     } else if (orz)
       ev.emplace_back(d, orz);
     else if (orb)
       ev.emplace_back(d, orb);
   } else if (ora == -orb) {
     Point ins;
     lineLineIntersection(l, Line(a, b), ins);
     ev.emplace_back(dot(ins - 1.a, 1.b - 1.a), 0);
   }
 sort(ev.begin(), ev.end());
 Tf ans = 0, len = 0, last = 0, tot = 0;
 bool active = false;
 int sign = 0;
 for (auto &qq : ev) {
   int tp = qq.second;
   Tf d = qq.first; /// current Segment is (last, d)
                 /// On Border
   if (sign) {
    len += d - last;
     tot += d - last:
     ans = max(ans, len);
     if (tp != sign) active = !active;
     sign = 0;
   } else {
     if (active) { /// Strictly Inside
       len += d - last;
       tot += d - last;
       ans = max(ans, len);
     if (tp == 0)
       active = !active;
     else
       sign = tp;
```

```
last = d;
   if (!active) len = 0;
  ans /= length(l.b - l.a);
  tot /= length(l.b - l.a);
 return {ans, tot};
    Graph
5.1 Graph Template
struct edge {
 int u, v;
 edge(int u = 0, int v = 0) : u(u), v(v) {}
 int to(int node) { return u ^ v ^ node; }
struct graph {
 int n;
  vector<vector<int>> adj;
  vector<edge> edges;
  graph(int n = 0) : n(n), adj(n) {}
  void addEdge(int u, int v, bool dir = 1) {
   adj[u].push_back(edges.size());
   if (dir) adj[v].push_back(edges.size());
   edges.emplace_back(u, v);
 int addNode() {
   adj.emplace_back();
   return n++;
  edge &operator()(int idx) { return edges[idx]; }
 vector<int> &operator[](int u) { return adj[u]; }
5.2 LCA
const int N = 3e5+7, K = 20;
vector<int> adj[N];
int anc[N][K]:
int level[N];
void setup(int u, int par) {
 level[u] = level[par]+1;
 anc[u][0] = par;
 for (int k=1; k<K; k++) anc[u][k] = anc[anc[u][k-1]][k]
      -1];
 for (int v: adj[u]) {
   if (v == par) continue;
   setup(v, u);
```

```
}
int lca(int u, int v) {
 if (level[u] > level[v]) swap(u, v);
 for (int k=K-1; k>=0; k--)
   if (level[u] + (1<<k) <= level[v]) v = anc[v][k];</pre>
 if (u == v) return u;
 for (int k=K-1; k>=0; k--)
   if (anc[u][k] != anc[v][k])
     u = anc[u][k], v = anc[v][k];
 return anc[u][0];
int getanc(int u, int d) {
 for (int k=0; k<K; k++)</pre>
   if (d & (1<<k))
     u = anc[u][k];
 return u;
int dist(int u, int v) {
 int g = lca(u, v);
 return level[u] + level[v] - 2*level[g];
5.3 SCC
typedef long long LL;
const LL N = 1e6 + 7;
bool vis[N];
vector<int> adj[N], adjr[N];
vector<int> order, component;
// tp = 0 ,finding topo order, tp = 1 , reverse edge
    traversal
void dfs(int u, int tp = 0) {
 vis[u] = true:
 if (tp) component.push_back(u);
 auto& ad = (tp ? adjr : adj);
 for (int v : ad[u])
   if (!vis[v]) dfs(v, tp);
 if (!tp) order.push_back(u);
int main() {
 for (int i = 1; i <= n; i++) {
   if (!vis[i]) dfs(i);
 memset(vis, 0, sizeof vis);
```

```
reverse(order.begin(), order.end());
for (int i : order) {
  if (!vis[i]) {
   // one component is found
   dfs(i, 1), component.clear();
 }
}
```

5.4 Euler Tour on Edge

```
// for simplicity, G[idx] contains the adjacency list of
    a node
// while G(e) is a reference to the e-th edge.
const int N = 2e5 + 5;
int in[N], out[N], fwd[N], bck[N];
int t = 0:
void dfs(graph &G, int node, int par) {
 out[node] = t;
 for (int e : G[node]) {
   int v = G(e).to(node);
   if (v == par) continue;
   fwd[e] = t++;
   dfs(G, v, node);
   bck[e] = t++;
 in[node] = t - 1;
void init(graph &G, int node) {
 t = 0:
 dfs(G, node, node);
```

5.5 LCA In O(1)

```
/* LCA in O(1)
* depth calculates weighted distance
* level calculates distance by number of edges
* Preprocessing in NlongN */
LL depth[N];
int level[N]:
int st[N], en[N], LOG[N], par[N];
int a[N], id[N], table[L][N];
vector<PII> adj[N];
int n, root, Time, cur;
void init(int nodes, int root_) {
 n = nodes, root = root_, LOG[0] = LOG[1] = 0;
 for (int i = 2; i <= n; i++) LOG[i] = LOG[i >> 1] + 1;
 for (int i = 0; i <= n; i++) adj[i].clear();</pre>
```

```
void addEdge(int u, int v, int w) {
 adj[u].push_back(PII(v, w));
 adj[v].push_back(PII(u, w));
int lca(int u, int v) {
 if (en[u] > en[v]) swap(u, v);
 if (st[v] <= st[u] && en[u] <= en[v]) return v;</pre>
 int 1 = LOG[id[v] - id[u] + 1];
 int p1 = id[u], p2 = id[v] - (1 << 1) + 1;
 int d1 = level[table[1][p1]], d2 = level[table[1][p2]];
 if (d1 < d2)
   return par[table[1][p1]];
   return par[table[1][p2]];
LL dist(int u, int v) {
 int 1 = lca(u, v);
 return (depth[u] + depth[v] - (depth[1] * 2));
/* Euler tour */
void dfs(int u, int p) {
 st[u] = ++Time, par[u] = p;
 for (auto [v, w] : adj[u]) {
   if (v == p) continue;
   depth[v] = depth[u] + w;
   level[v] = level[u] + 1;
   dfs(v, u);
 en[u] = ++Time:
 a[++cur] = u, id[u] = cur;
/* RMQ */
void pre() {
 cur = Time = 0, dfs(root, root);
 for (int i = 1; i <= n; i++) table[0][i] = a[i];</pre>
 for (int 1 = 0; 1 < L - 1; 1++) {
   for (int i = 1; i <= n; i++) {</pre>
     table[1 + 1][i] = table[1][i];
```

```
bool C1 = (1 << 1) + i <= n;
    bool C2 = level[table[1][i + (1 << 1)]] < level[</pre>
        table[1][i]];
    if (C1 && C2) table[1 + 1][i] = table[1][i + (1 <<</pre>
        1)];
}
```

```
5.6 HLD
const int N = 1e6 + 7;
template <typename DT>
struct Segtree {
 // write lazy segtree here
Segtree<int> tree(N);
vector<int> adj[N];
int depth[N], par[N], pos[N];
int head[N], heavy[N], cnt;
int dfs(int u, int p) {
 int SZ = 1, mxsz = 0, heavyc;
 depth[u] = depth[p] + 1;
 for (auto v : adj[u]) {
   if (v == p) continue;
   par[v] = u;
   int subsz = dfs(v, u);
   if (subsz > mxsz) heavy[u] = v, mxsz = subsz;
   SZ += subsz:
 }
 return SZ;
void decompose(int u, int h) {
 head[u] = h, pos[u] = ++cnt;
 if (heavy[u] != -1) decompose(heavy[u], h);
 for (int v : adj[u]) {
   if (v == par[u]) continue;
   if (v != heavy[u]) decompose(v, v);
 }
int query(int a, int b) {
 int ret = 0;
 for (; head[a] != head[b]; b = par[head[b]]) {
   if (depth[head[a]] > depth[head[b]]) swap(a, b);
   ret += tree.query(1, 0, cnt, pos[head[b]], pos[b]);
 }
```

```
if (depth[a] > depth[b]) swap(a, b);
ret += tree.query(1, 0, cnt, pos[a], pos[b]);
return ret;
}
```

5.7 Centroid Decomposition

```
class Centroid_Decomposition {
 vector<bool> blocked;
 vector<int> CompSize;
 int CompDFS(tree &T, int node, int par) {
   CompSize[node] = 1;
   for (int &e : T[node])
     if (e != par and !blocked[e]) CompSize[node] +=
         CompDFS(T, e, node);
   return CompSize[node];
 int FindCentroid(tree &T, int node, int par, int sz) {
   for (int &e : T[node])
     if (e != par and !blocked[e])
       if (CompSize[e] > sz / 2) return FindCentroid(T,
           e, node, sz);
   return node;
 pair<int, int> GetCentroid(tree &T, int entry) {
   int sz = CompDFS(T, entry, entry);
   return {FindCentroid(T, entry, entry, sz), sz};
 c_vector<LL> left[2], right[2];
 int GMin, GMax;
 void dfs(tree &T, int node, int par, int Min, int Max,
      int sum) {
   if (blocked[node]) return;
   right [Max < sum or Min > sum] [sum] ++;
   Max = max(Max, sum), Min = min(Min, sum);
   GMin = min(GMin, sum), GMax = max(GMax, sum);
   for (int i = 0; i < T[node].size(); i++)</pre>
     if (T[node][i] != par) {
       dfs(T, T[node][i], node, Min, Max, sum + T.col[
           nodel[i]):
     }
 }
 LL solve(tree &T, int c, int sz) {
   LL ans = 0;
   left[0].clear(-sz, sz), left[1].clear(-sz, sz);
   for (int i = 0; i < T[c].size(); i++) {</pre>
     GMin = 1, GMax = -1;
     dfs(T, T[c][i], c, GMin, GMax, T.col[c][i]);
     ans += right[0][0] + left[1][0] * right[1][0];
     for (int j : {0, 1})
      for (int k = GMin; k <= GMax; k++) {</pre>
```

```
ans += right[j][k] * (left[0][-k] + (j == 0) *
            left[1][-k]);
      }
    for (int j : {0, 1})
      for (int k = GMin; k <= GMax; k++) {</pre>
        left[j][k] += right[j][k];
        right[i][k] = 0;
      }
  }
  return ans;
public:
LL operator()(tree &T, int entry) {
  blocked.resize(T.n);
  CompSize.resize(T.n);
  for (int i : {0, 1})
    left[i].resize(2 * T.n + 5), right[i].resize(2 * T.
        n + 5):
  auto [c, sz] = GetCentroid(T, entry);
  LL ans = solve(T, c, sz);
  blocked[c] = true;
  for (int e : T[c])
    if (!blocked[e]) ans += (*this)(T, e);
  return ans;
}
```

5.8 Dinic Max Flow

```
/// flow with demand(lower bound) only for DAG
// create new src and sink
// add_edge(new src, u, sum(in_demand[u]))
// add_edge(u, new sink, sum(out_demand[u]))
// add_edge(old sink, old src, inf)
// if (sum of lower bound == flow) then demand satisfied
// flow in every edge i = demand[i] + e.flow
using Ti = long long;
const Ti INF = 1LL << 60:
struct edge {
 int v, u;
 Ti cap, flow = 0:
 edge(int v, int u, Ti cap) : v(v), u(u), cap(cap) {}
const int N = 1e5 + 50;
vector<edge> edges;
vector<int> adj[N];
int m = 0, n;
int level[N], ptr[N];
queue<int> q;
```

```
bool bfs(int s. int t) {
 for (q.push(s), level[s] = 0; !q.empty(); q.pop()) {
   for (int id : adj[q.front()]) {
     auto &ed = edges[id];
     if (ed.cap - ed.flow > 0 and level[ed.u] == -1)
       level[ed.u] = level[ed.v] + 1, q.push(ed.u);
  }
 }
 return level[t] != -1;
Ti dfs(int v, Ti pushed, int t) {
 if (pushed == 0) return 0;
 if (v == t) return pushed;
 for (int &cid = ptr[v]; cid < adj[v].size(); cid++) {</pre>
   int id = adj[v][cid];
   auto &ed = edges[id];
   if (level[v] + 1 != level[ed.u] || ed.cap - ed.flow <</pre>
        1) continue;
   Ti tr = dfs(ed.u, min(pushed, ed.cap - ed.flow), t);
   if (tr == 0) continue;
   ed.flow += tr:
   edges[id ^ 1].flow -= tr;
   return tr;
 }
 return 0;
void init(int nodes) {
 m = 0, n = nodes;
 for (int i = 0; i < n; i++) level[i] = -1, ptr[i] = 0,
      adj[i].clear();
void addEdge(int v, int u, Ti cap) {
 edges.emplace_back(v, u, cap), adj[v].push_back(m++);
 edges.emplace_back(u, v, 0), adj[u].push_back(m++);
Ti maxFlow(int s, int t) {
 Ti f = 0:
 for (auto &ed : edges) ed.flow = 0;
 for (; bfs(s, t); memset(level, -1, n * 4)) {
   for (memset(ptr, 0, n * 4); Ti pushed = dfs(s, INF, t
       ); f += pushed)
 }
 return f;
```

5.9 Min Cost Max Flow

```
struct edge {
  int v, rev;
  LL cap, cost, flow;
  edge() {}
  edge(int v, int rev, LL cap, LL cost)
     : v(v), rev(rev), cap(cap), cost(cost), flow(0) {}
};
struct mcmf {
  int src, sink, n;
  vector<int> par, idx, Q;
  vector<bool> inq;
  vector<LL> dis;
  vector<vector<edge>> g;
  mcmf() {}
  mcmf(int src, int sink, int n)
      : src(src), sink(sink), n(n), par(n), idx(n), inq(n
          ), dis(n), g(n), Q(10000005) {} // use Q(n) if
          not using random
  void add_edge(int u, int v, LL cap, LL cost, bool
      directed = true) {
    edge _u = edge(v, g[v].size(), cap, cost);
    edge _v = edge(u, g[u].size(), 0, -cost);
    g[u].pb(_u);
    g[v].pb(_v);
    if (!directed) add_edge(v, u, cap, cost, true);
  bool spfa() {
   for (int i = 0; i < n; i++) {
     dis[i] = inf, inq[i] = false;
    int f = 0, 1 = 0;
    dis[src] = 0, par[src] = -1, Q[1++] = src, inq[src] =
    while (f < 1) {
     int u = Q[f++];
     for (int i = 0; i < g[u].size(); i++) {</pre>
       edge &e = g[u][i];
       if (e.cap <= e.flow) continue;</pre>
       if (dis[e.v] > dis[u] + e.cost) {
         dis[e.v] = dis[u] + e.cost;
         par[e.v] = u, idx[e.v] = i;
         if (!inq[e.v]) inq[e.v] = true, Q[1++] = e.v;
         // if (!inq[e.v]) {
         // inq[e.v] = true;
         // if (f \&\& rnd() \& 7) Q[--f] = e.v;
         // else Q[1++] = e.v;
         // }
       }
     inq[u] = false;
```

```
}
   return (dis[sink] != inf);
 pair<LL, LL> solve() {
   LL mincost = 0, maxflow = 0;
   while (spfa()) {
     LL bottleneck = inf;
     for (int u = par[sink], v = idx[sink]; u != -1; v = |graph tree;
          idx[u], u = par[u]) {
       edge &e = g[u][v];
       bottleneck = min(bottleneck, e.cap - e.flow);
     for (int u = par[sink], v = idx[sink]; u != -1; v =
          idx[u], u = par[u]) {
       edge &e = g[u][v];
       e.flow += bottleneck:
       g[e.v][e.rev].flow -= bottleneck;
     mincost += bottleneck * dis[sink], maxflow +=
         bottleneck;
   }
   return make_pair(mincost, maxflow);
// want to minimize cost and don't care about flow
// add edge from sink to dummy sink (cap = inf, cost = 0)
// add edge from source to sink (cap = inf, cost = 0)
// run mcmf, cost returned is the minimum cost
```

5.10 Block Cut Tree

```
vector<vector<int> > components;
vector<int> cutpoints, start, low;
vector<bool> is_cutpoint;
stack<int> st;
void find_cutpoints(int node, graph &G, int par = -1, int
     d = 0) {
 low[node] = start[node] = d++;
 st.push(node);
 int cnt = 0;
 for (int e : G[node])
   if (int to = G(e).to(node); to != par) {
     if (start[to] == -1) {
       find_cutpoints(to, G, node, d + 1);
       cnt++;
       if (low[to] >= start[node]) {
         is_cutpoint[node] = par != -1 or cnt > 1;
         components.push_back({node}); // starting a new
             block with the point
         while (st.top() != node)
```

```
components.back().push_back(st.top()), st.pop
      }
     }
     low[node] = min(low[node], low[to]);
vector<int> id;
void init(graph &G) {
 int n = G.n;
 start.assign(n, -1), low.resize(n), is_cutpoint.resize(
     n), id.assign(n, -1);
 find_cutpoints(0, G);
 for (int u = 0; u < n; ++u)
   if (is_cutpoint[u]) id[u] = tree.addNode();
 for (auto &comp : components) {
   int node = tree.addNode();
   for (int u : comp)
     if (!is_cutpoint[u])
      id[u] = node;
      tree.addEdge(node, id[u]);
 }
 if (id[0] == -1) // corner - 1
   id[0] = tree.addNode();
```

5.11 Bridge Tree

```
vector<int>adj[100010];
bool vis[100010];
int start[100010],low[100010],dep[100010],par[100010],
    timer,up[100010][20];
vector<pair<int,int>>bridges;
int FindPar(int n){
 if(par[n]==n) return par[n];
 par[n]=FindPar(par[n]);
 return par[n];
void Merge(int u,int v){
 u=FindPar(u);
 v=FindPar(v);
 if(u!=v) par[v]=u;
void dfs(int s,int p){
 vis[s]=1;
 timer++;
 start[s]=timer;
 low[s]=timer;
 for(int i=0;i<adj[s].size();i++){</pre>
```

```
int v=adj[s][i];
    if(v==p) continue;
    if(vis[v]){
     low[s]=min(low[s],start[v]);
     continue;
   }
    else{
     dfs(v,s);
     low[s]=min(low[s],low[v]);
     if(low[v]>start[s]) bridges.push_back({s,v});
     else Merge(s,v);
   }
 }
}
void MakeBridgeTree(int n){
  for(int i=1;i<=n;i++) adj[i].clear();</pre>
  for(auto p: bridges){
    int u=FindPar(p.first);
    int v=FindPar(p.second);
    if(u!=v){
     adj[u].push_back(v);
     adj[v].push_back(u);
 }
void dfs2(int s,int p){
  if(p!=-1) up[s][0]=p;
  for(int i=1;i<=18;i++) up[s][i]=up[up[s][i-1]][i-1];</pre>
    for(int i=0;i<adj[s].size();i++){</pre>
     int v=adj[s][i];
     if(v==p) continue;
     dep[v]=dep[s]+1;
     dfs2(v,s);
int LCA(int u,int v){
  if(dep[u]>dep[v]) swap(u,v);
  int d=dep[v]-dep[u];
  for(int i=18;i>=0;i--){
    if(d&(1<<i)) v=up[v][i];</pre>
  }
  if(u==v) return u;
  for(int i=18;i>=0;i--){
    if(up[u][i]!=up[v][i]){
     u=up[u][i];
     v=up[v][i];
   }
  return up[u][0];
```

```
int GetDistance(int u,int v){
   return dep[u]+dep[v]-2*dep[LCA(u,v)];
}
int Intersect(int lca_ab,int lca_cd,int x,int y){ // x= a
      or b and y= c or d
   int lca_xy=LCA(x,y);
   if(dep[lca_xy]<dep[lca_ab] || dep[lca_xy]<dep[lca_cd])
      return 0;
   int ret=min(dep[lca_xy]-dep[lca_ab],dep[lca_xy]-dep[
      lca_cd]);
   return ret;
}</pre>
```

5.12 Tree Isomorphism

```
mp["01"] = 1;
ind = 1;
int dfs(int u, int p) {
 int cnt = 0;
 vector<int> vs;
 for (auto v : g1[u]) {
   if (v != p) {
     int got = dfs(v, u);
     vs.pb(got);
     cnt++;
   }
 }
 if (!cnt) return 1;
  sort(vs.begin(), vs.end());
  string s = "0";
 for (auto i : vs) s += to_string(i);
 vs.clear();
  s.pb('1');
 if (mp.find(s) == mp.end()) mp[s] = ++ind;
 int ret = mp[s];
 return ret;
```

6 Math

6.1 Matrix Exponentiation

```
int temp=((long long)(Mat[i][k]%mod)*((long long
             )other.Mat[k][j]%mod))%mod;
         product.Mat[i][j]=(product.Mat[i][j]%mod+temp%
             mod)%mod;
       }
     }
   }
   return product;
Matrix MatExpo(Matrix a, int p){
 Matrix product;
 for(int i=0;i<n;i++) product.Mat[i][i]=1;</pre>
 while(p>0){
   if(p%2) product=product*a;
   p/=2:
   a=a*a;
 return product;
```

6.2 Combi

6.3 Linear Sieve

```
const int N = 1e7;
vector<int> primes;
int spf[N + 5], phi[N + 5], NOD[N + 5], cnt[N + 5], POW[N + 5];
bool prime[N + 5];
int SOD[N + 5];
void init() {
  fill(prime + 2, prime + N + 1, 1);
  SOD[1] = NOD[1] = phi[1] = spf[1] = 1;
  for (LL i = 2; i <= N; i++) {
    if (prime[i]) {
      primes.push_back(i), spf[i] = i;
    }
}</pre>
```

```
phi[i] = i - 1;
   NOD[i] = 2, cnt[i] = 1;
   SOD[i] = i + 1, POW[i] = i;
 }
  for (auto p : primes) {
   if (p * i > N or p > spf[i]) break;
   prime[p * i] = false, spf[p * i] = p;
   if (i % p == 0) {
     phi[p * i] = p * phi[i];
     NOD[p * i] = NOD[i] / (cnt[i] + 1) * (cnt[i] + 2)
             cnt[p * i] = cnt[i] + 1;
     SOD[p * i] = SOD[i] / SOD[POW[i]] * (SOD[POW[i]]
         + p * POW[i]),
            POW[p * i] = p * POW[i];
     break:
   } else {
     phi[p * i] = phi[p] * phi[i];
     NOD[p * i] = NOD[p] * NOD[i], cnt[p * i] = 1;
     SOD[p * i] = SOD[p] * SOD[i], POW[p * i] = p;
   }
 }
}
```

6.4 Pollard Rho

```
LL mul(LL a, LL b, LL mod) {
 return (__int128)a * b % mod;
 // LL ans = a * b - mod * (LL) (1.L / mod * a * b);
 // return ans + mod * (ans < 0) - mod * (ans >= (LL)
      mod):
LL bigmod(LL num, LL pow, LL mod) {
 LL ans = 1;
 for (; pow > 0; pow >>= 1, num = mul(num, num, mod))
   if (pow & 1) ans = mul(ans, num, mod);
 return ans;
bool is_prime(LL n) {
 if (n < 2 or n % 6 % 4 != 1) return (n | 1) == 3;
 LL a[] = \{2, 325, 9375, 28178, 450775, 9780504,
      1795265022};
 LL s = \_builtin\_ctzll(n - 1), d = n >> s;
 for (LL x : a) {
   LL p = bigmod(x \% n, d, n), i = s;
   for (; p != 1 and p != n - 1 and x % n and i--; p =
       mul(p, p, n))
   if (p != n - 1 and i != s) return false;
```

```
return true:
LL get_factor(LL n) {
 auto f = [\&](LL x) \{ return mul(x, x, n) + 1; \};
 LL x = 0, y = 0, t = 0, prod = 2, i = 2, q;
 for (; t++ % 40 or gcd(prod, n) == 1; x = f(x), y = f(f)
      (v))) {
   (x == y) ? x = i++, y = f(x) : 0;
   prod = (q = mul(prod, max(x, y) - min(x, y), n)) ? q
 return gcd(prod, n);
map<LL, int> factorize(LL n) {
 map<LL, int> res;
 if (n < 2) return res;</pre>
 LL small_primes[] = {2, 3, 5, 7, 11, 13, 17, 19, 23,
      29, 31, 37, 41,
                     43, 47, 53, 59, 61, 67, 71, 73, 79,
                         83, 89, 97};
 for (LL p : small_primes)
   for (; n % p == 0; n /= p, res[p]++)
  auto _factor = [&](LL n, auto &_factor) {
   if (n == 1) return;
   if (is_prime(n))
     res[n]++;
   else {
     LL x = get_factor(n);
     _factor(x, _factor);
     _factor(n / x, _factor);
   }
 };
  _factor(n, _factor);
 return res;
```

6.5 Chinese Remainder Theorem

```
// given a, b will find solutions for
// ax + by = 1
tuple<LL, LL, LL> EGCD(LL a, LL b) {
  if (b == 0)
    return {1, 0, a};
  else {
    auto [x, y, g] = EGCD(b, a % b);
    return {y, x - a / b * y, g};
  }
}
// given modulo equations, will apply CRT
```

```
PLL CRT(vector<PLL> &v) {
   LL V = 0, M = 1;
   for (auto &[v, m] : v) { // value % mod
      auto [x, y, g] = EGCD(M, m);
      if ((v - V) % g != 0) return {-1, 0};
      V += x * (v - V) / g % (m / g) * M, M *= m / g;
      V = (V % M + M) % M;
   }
   return make_pair(V, M);
}
```

6.6 Mobius Function

```
const int N = 1e6 + 5;
int mob[N];
void mobius() {
  memset(mob, -1, sizeof mob);
  mob[1] = 1;
  for (int i = 2; i < N; i++)
    if (mob[i]) {
      for (int j = i + i; j < N; j += i) mob[j] -= mob[i
        ];
    }
}</pre>
```

6.7 FFT

```
using CD = complex<double>;
typedef long long LL;
const double PI = acos(-1.0L);
int N;
vector<int> perm;
vector<CD> wp[2];
void precalculate(int n) {
 assert((n \& (n - 1)) == 0), N = n;
 perm = vector<int>(N, 0);
 for (int k = 1; k < N; k <<= 1) {
   for (int i = 0; i < k; i++) {
     perm[i] <<= 1;
     perm[i + k] = 1 + perm[i];
 }
 wp[0] = wp[1] = vector < CD > (N);
 for (int i = 0; i < N; i++) {</pre>
   wp[0][i] = CD(cos(2 * PI * i / N), sin(2 * PI * i / N)
       ));
   wp[1][i] = CD(cos(2 * PI * i / N), -sin(2 * PI * i / N))
       N));
 }
void fft(vector<CD> &v, bool invert = false) {
```

```
if (v.size() != perm.size()) precalculate(v.size());
  for (int i = 0; i < N; i++)</pre>
   if (i < perm[i]) swap(v[i], v[perm[i]]);</pre>
  for (int len = 2; len <= N; len *= 2) {</pre>
   for (int i = 0, d = N / len; i < N; i += len) {</pre>
         ) {
       CD x = v[i + j];
       CD y = wp[invert][idx] * v[i + j + len / 2];
       v[i + j] = x + y;
       v[i + j + len / 2] = x - y;
     }
   }
  }
  if (invert) {
   for (int i = 0; i < N; i++) v[i] /= N;
 }
void pairfft(vector<CD> &a, vector<CD> &b, bool invert =
    false) {
  int N = a.size();
  vector<CD> p(N);
 for (int i = 0; i < N; i++) p[i] = a[i] + b[i] * CD(0,
      1);
 fft(p, invert);
  p.push_back(p[0]);
  for (int i = 0; i < N; i++) {</pre>
   if (invert) {
     a[i] = CD(p[i].real(), 0);
     b[i] = CD(p[i].imag(), 0);
   } else {
     a[i] = (p[i] + conj(p[N - i])) * CD(0.5, 0);
     b[i] = (p[i] - conj(p[N - i])) * CD(0, -0.5);
   }
 }
vector<LL> multiply(const vector<LL> &a, const vector<LL>
     &b) {
  int n = 1;
  while (n < a.size() + b.size()) n <<= 1;</pre>
  vector<CD> fa(a.begin(), a.end()), fb(b.begin(), b.end
      ());
 fa.resize(n);
  fb.resize(n);
           fft(fa); fft(fb);
  pairfft(fa, fb);
  for (int i = 0; i < n; i++) fa[i] = fa[i] * fb[i];</pre>
  fft(fa. true):
  vector<LL> ans(n);
```

```
for (int i = 0: i < n: i++) ans[i] = round(fa[i].real()</pre>
                                                             );
                                                        return ans:
                                                       const int M = 1e9 + 7, B = sqrt(M) + 1;
for (int j = 0, idx = 0; j < len / 2; j++, idx += d | vector<LL> anyMod(const vector<LL> &a, const vector<LL> &
                                                        int n = 1:
                                                         while (n < a.size() + b.size()) n <<= 1;</pre>
                                                         vector<CD> al(n), ar(n), bl(n), br(n);
                                                         for (int i = 0; i < a.size(); i++) al[i] = a[i] % M / B</pre>
                                                             , ar[i] = a[i] \% M \% B;
                                                        for (int i = 0; i < b.size(); i++) bl[i] = b[i] % M / B</pre>
                                                             , br[i] = b[i] % M % B;
                                                         pairfft(al, ar);
                                                         pairfft(bl. br):
                                                                  fft(al); fft(ar); fft(bl); fft(br);
                                                         for (int i = 0; i < n; i++) {</pre>
                                                          CD ll = (al[i] * bl[i]), lr = (al[i] * br[i]);
                                                          CD rl = (ar[i] * bl[i]), rr = (ar[i] * br[i]);
                                                          al[i] = 11;
                                                          ar[i] = lr;
                                                          bl[i] = rl;
                                                          br[i] = rr;
                                                         pairfft(al, ar, true);
                                                         pairfft(bl, br, true);
                                                                  fft(al, true); fft(ar, true); fft(bl, true);
                                                             fft(br. true):
                                                         vector<LL> ans(n);
                                                         for (int i = 0; i < n; i++) {</pre>
                                                          LL right = round(br[i].real()), left = round(al[i].
                                                               real());
                                                          LL mid = round(round(bl[i].real()) + round(ar[i].real
                                                           ans[i] = ((left \% M) * B * B + (mid \% M) * B + right)
                                                                % M;
                                                        }
                                                        return ans;
                                                       6.8 NTT
                                                       const LL N = 1 << 18;</pre>
                                                       const LL MOD = 786433;
                                                       vector<LL> P[N];
```

LL rev[N], w[N | 1], a[N], b[N], inv_n, g;

LL Pow(LL b, LL p) {

LL ret = 1;

```
while (p) {
   if (p & 1) ret = (ret * b) % MOD;
   b = (b * b) \% MOD;
   p >>= 1;
 return ret;
LL primitive_root(LL p) {
 vector<LL> factor;
 LL phi = p - 1, n = phi;
 for (LL i = 2; i * i <= n; i++) {
   if (n % i) continue;
   factor.emplace_back(i);
   while (n % i == 0) n /= i:
 if (n > 1) factor.emplace_back(n);
 for (LL res = 2; res <= p; res++) {</pre>
   bool ok = true;
   for (LL i = 0; i < factor.size() && ok; i++)</pre>
     ok &= Pow(res, phi / factor[i]) != 1;
   if (ok) return res:
 return -1;
void prepare(LL n) {
 LL sz = abs(31 - \_builtin\_clz(n));
 LL r = Pow(g, (MOD - 1) / n);
 inv_n = Pow(n, MOD - 2);
 w[0] = w[n] = 1:
 for (LL i = 1; i < n; i++) w[i] = (w[i-1] * r) % MOD;
 for (LL i = 1; i < n; i++)</pre>
   rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << (sz - 1));
void NTT(LL *a, LL n, LL dir = 0) {
for (LL i = 1; i < n - 1; i++)
   if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
 for (LL m = 2: m <= n: m <<= 1) {
   for (LL i = 0; i < n; i += m) {
     for (LL j = 0; j < (m >> 1); j++) {
       LL &u = a[i + j], &v = a[i + j + (m >> 1)];
       LL t = v * w[dir ? n - n / m * j : n / m * j] %
       v = u - t < 0 ? u - t + MOD : u - t;
       u = u + t >= MOD ? u + t - MOD : u + t;
   }
 }
   for (LL i = 0; i < n; i++) a[i] = (inv_n * a[i]) %
        MOD:
```

```
}
vector<LL> mul(vector<LL> p, vector<LL> q) {
  LL n = p.size(), m = q.size();
  LL t = n + m - 1, sz = 1;
  while (sz < t) sz <<= 1;
  prepare(sz);
  for (LL i = 0; i < n; i++) a[i] = p[i];</pre>
  for (LL i = 0; i < m; i++) b[i] = q[i];
  for (LL i = n; i < sz; i++) a[i] = 0;</pre>
  for (LL i = m; i < sz; i++) b[i] = 0;</pre>
  NTT(a, sz);
 NTT(b, sz);
  for (LL i = 0; i < sz; i++) a[i] = (a[i] * b[i]) % MOD;</pre>
  NTT(a, sz, 1);
 vector<LL> c(a, a + sz);
  while (c.size() && c.back() == 0) c.pop_back();
 return c;
```

6.9 WalshHadamard

```
#include <bits/stdc++.h>
using namespace std;
typedef long long LL;
#define bitwiseXOR 1
// #define bitwiseAND 2
// #define bitwiseOR 3
const LL MOD = 30011;
LL BigMod(LL b, LL p) {
 LL ret = 1;
  while (p > 0) {
   if (p % 2 == 1) {
     ret = (ret * b) % MOD;
   p = p / 2;
   b = (b * b) \% MOD;
 return ret % MOD;
void FWHT(vector<LL>& p, bool inverse) {
 LL n = p.size();
  assert((n & (n - 1)) == 0);
 for (LL len = 1; 2 * len <= n; len <<= 1) {</pre>
   for (LL i = 0; i < n; i += len + len) {</pre>
     for (LL j = 0; j < len; j++) {</pre>
```

```
LL u = p[i + j];
      LL v = p[i + len + j];
#ifdef bitwiseXOR
       p[i + j] = (u + v) \% MOD;
       p[i + len + j] = (u - v + MOD) \% MOD;
#endif // bitwiseXOR
#ifdef bitwiseAND
       if (!inverse) {
         p[i + j] = v \% MOD;
        p[i + len + j] = (u + v) \% MOD;
       } else {
         p[i + j] = (-u + v) \% MOD;
         p[i + len + j] = u \% MOD;
#endif // bitwiseAND
#ifdef bitwiseΩR
       if (!inverse) {
         p[i + j] = u + v;
        p[i + len + j] = u;
       } else {
         p[i + j] = v;
        p[i + len + j] = u - v;
#endif // bitwiseOR
     }
   }
 }
#ifdef bitwiseXOR
 if (inverse) {
   LL val = BigMod(n, MOD - 2); // Option 2: Exclude
   for (LL i = 0; i < n; i++) {</pre>
     // assert(p[i]%n==0); //Option 2: Include
     p[i] = (p[i] * val) % MOD; // Option 2: p[i]/=n;
   }
#endif // bitwiseXOR
```

6.10 Adaptive Simpsons

```
/*
   For finding the length of an arc in a range
   L = integrate(ds) from start to end of range
   where ds = sqrt(1+(d/dy(x))^2)dy
*/
const double SIMPSON_TERMINAL_EPS = 1e-12;
/// Function whose integration is to be calculated
```

```
double F(double x);
double simpson(double minx, double maxx) {
 return (\max - \min x) / 6 * (F(\min x) + 4 * F((\min x + \max x)))
      maxx) / 2.) + F(maxx));
double adaptive_simpson(double minx, double maxx, double
    c, double EPS) {
 // if(maxx - minx < SIMPSON_TERMINAL_EPS) return 0;</pre>
  double midx = (minx + maxx) / 2;
  double a = simpson(minx, midx);
  double b = simpson(midx, maxx);
  if (fabs(a + b - c) < 15 * EPS) return a + b + (a + b - c)
       c) / 15.0;
 return adaptive_simpson(minx, midx, a, EPS / 2.) +
        adaptive_simpson(midx, maxx, b, EPS / 2.);
double adaptive_simpson(double minx, double maxx, double
    EPS) {
 return adaptive_simpson(minx, maxx, simpson(minx, maxx,
       i), EPS);
```

6.11 Berlekamp Massey

```
struct berlekamp_massey { // for linear recursion
 typedef long long LL;
 static const int SZ = 2e5 + 5;
 static const int MOD = 1e9 + 7; /// mod must be a prime
 LL m , a[SZ] , h[SZ] , t_[SZ] , s[SZ] , t[SZ];
 // bigmod goes here
 inline vector <LL> BM( vector <LL> &x ) {
  LL lf , ld;
  vector <LL> ls , cur;
  for ( int i = 0; i < int(x.size()); ++i ) {</pre>
    LL t = 0:
    for ( int j = 0; j < int(cur.size()); ++j ) t = (t</pre>
         + x[i - j - 1] * cur[j]) % MOD;
    if ((t - x[i]) \% MOD == 0) continue;
    if (!cur.size()) {
      cur.resize( i + 1 );
      lf = i; ld = (t - x[i]) % MOD;
      continue;
    LL k = -(x[i] - t) * bigmod(ld, MOD - 2, MOD) %
          MOD;
    vector <LL> c(i - lf - 1);
    c.push_back( k );
```

```
for ( int j = 0; j < int(ls.size()); ++j ) c.</pre>
       push_back(-ls[j] * k % MOD);
   if ( c.size() < cur.size() ) c.resize( cur.size() )</pre>
   for ( int j = 0; j < int(cur.size()); ++j ) c[j] =</pre>
        (c[i] + cur[i]) % MOD;
   if (i - lf + (int)ls.size() >= (int)cur.size() ) ls
        = cur, lf = i, ld = (t - x[i]) % MOD;
   cur = c;
 }
 for ( int i = 0; i < int(cur.size()); ++i ) cur[i] =</pre>
      (cur[i] % MOD + MOD) % MOD;
 return cur;
inline void mull( LL *p , LL *q ) {
 for ( int i = 0; i < m + m; ++i ) t_[i] = 0;</pre>
 for ( int i = 0; i < m; ++i ) if ( p[i] )
     for ( int j = 0; j < m; ++j ) t_{i} = (t_{i} + j = (t_{i} +
                                                          Given a function f and n, finds the smallest fraction p /
          i] + p[i] * q[i]) % MOD;
 for ( int i = m + m - 1; i >= m; --i ) if ( t_{-}[i] )
     for ( int j = m - 1; \tilde{j}; --j ) t_[i - j - 1] = (
         t_{i} - i - 1 + t_{i} * h[j]) % MOD;
 for ( int i = 0; i < m; ++i ) p[i] = t_[i];</pre>
inline LL calc( LL K ) {
 for ( int i = m; ~i; --i ) s[i] = t[i] = 0;
 s[0] = 1; if ( m != 1 ) t[1] = 1; else t[0] = h[0];
 while ( K ) {
  if (K&1) mull(s,t);
   mull( t , t ); K >>= 1;
 }
 LL su = 0;
 for ( int i = 0; i < m; ++i ) su = (su + s[i] * a[i])
       % MOD;
 return (su % MOD + MOD) % MOD;
/// already calculated upto k , now calculate upto n.
inline vector <LL> process( vector <LL> &x , int n ,
    int k ) {
 auto re = BM( x );
 x.resize(n+1);
 for ( int i = k + 1; i <= n; i++ ) {
  for ( int j = 0; j < re.size(); j++ ) {</pre>
     x[i] += 1LL * x[i - j - 1] % MOD * re[j] % MOD; x
         [i] %= MOD;
   }
 }
 return x;
inline LL work( vector <LL> &x , LL n ) {
```

```
if ( n < int(x.size()) ) return x[n] % MOD:</pre>
    vector \langle LL \rangle v = BM( x ); m = v.size(); if ( !m )
    for ( int i = 0; i < m; ++i ) h[i] = v[i], a[i] = x[i</pre>
   return calc( n ) % MOD;
} rec:
vector <LL> v;
void solve() {
 int n;
  cin >> n;
  cout << rec.work(v, n - 1) << endl;</pre>
```

6.12 Fractional Binary Search

```
q in [0, 1] or [0,n]
such that f(p / q) is true, and p, q <= n.
Time: O(log(n))
**/
struct frac { long long p, q; };
bool f(frac x) {
return 6 + 8 * x.p >= 17 * x.q + 12;
frac fracBS(long long n) {
 bool dir = 1, A = 1, B = 1;
 frac lo{0, 1}, hi{1, 0}; // Set hi to 1/0 to search
      within [0, n] and {1, 1} to search within [0, 1]
 if (f(lo)) return lo;
  assert(f(hi)); //checking if any solution exists or not
  while (A | | B) {
   long long adv = 0, step = 1; // move hi if dir, else
       10
   for (int si = 0; step; (step *= 2) >>= si) {
     adv += step;
     frac mid{lo.p * adv + hi.p, lo.q * adv + hi.q};
     if (abs(mid.p) > n \mid\mid mid.q > n \mid\mid dir == !f(mid))
       adv -= step; si = 2;
   hi.p += lo.p * adv;
   hi.q += lo.q * adv;
   dir = !dir;
   swap(lo, hi);
   A = B; B = !!adv;
 return dir ? hi : lo;
```

6.13 Lagrange

```
// p is a polynomial with n points.
// p(0), p(1), p(2), \dots p(n-1) are given.
// Find p(x).
LL Lagrange(vector<LL> &p, LL x) {
 LL n = p.size(), L, i, ret;
 if (x < n) return p[x];</pre>
 L = 1:
 for (i = 1; i < n; i++) {
  L = (L * (x - i)) \% MOD;
  L = (L * bigmod(MOD - i, MOD - 2)) \% MOD;
 }
 ret = (L * p[0]) % MOD;
 for (i = 1; i < n; i++) {
   L = (L * (x - i + 1)) \% MOD;
   L = (L * bigmod(x - i, MOD - 2)) \% MOD;
   L = (L * bigmod(i, MOD - 2)) % MOD;
   L = (L * (MOD + i - n)) \% MOD;
   ret = (ret + L * p[i]) % MOD;
 return ret;
```

6.14 Shanks' Baby Step, Giant Step

```
// Finds a^x = b \pmod{p}
LL bigmod(LL b, LL p, LL m) {}
LL babyStepGiantStep(LL a, LL b, LL p) {
 LL i, j, c, sq = sqrt(p);
 map<LL, LL> babyTable;
 for (j = 0, c = 1; j \le sq; j++, c = (c * a) % p)
     babyTable[c] = j;
 LL giant = bigmod(a, sq * (p - 2), p);
 for (i = 0, c = 1; i <= sq; i++, c = (c * giant) % p) {
   if (babyTable.find((c * b) % p) != babyTable.end())
     return i * sq + babyTable[(c * b) % p];
 }
 return -1;
```

6.15 Xor Basis

```
struct XorBasis {
 static const int sz = 64;
```

```
array<ULL, sz> base = {0}, back;
  array<int, sz> pos;
  void insert(ULL x, int p) {
   ULL cur = 0;
   for (int i = sz - 1; ~i; i--)
     if (x >> i & 1) {
       if (!base[i]) {
         base[i] = x, back[i] = cur, pos[i] = p;
       } else x ^= base[i], cur |= 1ULL << i;</pre>
     }
  pair<ULL, vector<int>> construct(ULL mask) {
   ULL ok = 0, x = mask:
   for (int i = sz - 1; ~i; i--)
     if (mask >> i & 1 and base[i]) mask ^= base[i]. ok
          |= 1ULL << i;
   vector<int> ans;
   for (int i = 0; i < sz; i++)</pre>
     if (ok >> i & 1) {
       ans.push_back(pos[i]);
       ok ^= back[i];
   return {x ^ mask, ans};
 }
};
```

7 String

7.1 Aho Corasick

```
int trie[260010][26],fail[260010],freq[260010],node,pos
    [260010]:
void Clear(){
 for(int i=0;i<260000;i++){</pre>
   for(int j=0;j<26;j++) trie[i][j]=0;</pre>
     freq[i]=0;
   fail[i]=0;
   pos[i]=0;
 node=0:
int Insert(string s){
 int curr=0;
 for(int i=0;i<s.size();i++){</pre>
   int letter=s[i]-'a':
   if(!trie[curr][letter]){
     node++:
     trie[curr][letter]=node;
   curr=trie[curr][letter]:
```

```
return curr:
void bfs(){
 queue<int>q;
 for(int i=0;i<26;i++){</pre>
   if(trie[0][i]) q.push(trie[0][i]);
 while(!q.empty()){
   int u=q.front();
   q.pop();
   for(int i=0:i<26:i++){</pre>
     int v=trie[u][i];
     if(v==0) continue;
     q.push(v);
     int f=fail[u];
     while(f && trie[f][i]==0) f=fail[f]:
     fail[v]=trie[f][i];
   }
 }
void Aho_Corasick(string s){
 bfs();
 int curr=0:
 for(int i=0;i<s.size();i++){</pre>
   int letter=s[i]-'a';
   while(curr && trie[curr][letter]==0) curr=fail[curr];
   curr=trie[curr][letter]:
   int temp=curr;
   while(temp>0){
     freq[temp]++;
     temp=fail[temp];
 }
```

7.2 Double hash

```
7.2 Double hash
// define +, -, * for (PLL, LL) and (PLL, PLL), % for (
    PLL, PLL);
PLL base(1949313259, 1997293877);
PLL mod(2091573227, 2117566807);

PLL power(PLL a, LL p) {
    PLL ans = PLL(1, 1);
    for(; p; p >>= 1, a = a * a % mod) {
        if(p & 1) ans = ans * a % mod;
    }
    return ans;
}
```

```
PLL inverse(PLL a) { return power(a, (mod.ff - 1) * (mod.
    ss - 1) - 1); }
PLL inv_base = inverse(base);
PLL val:
vector<PLL> P;
void hash_init(int n) {
 P.resize(n + 1):
 P[0] = PLL(1, 1);
 for (int i = 1; i <= n; i++) P[i] = (P[i - 1] * base) %
       mod:
PLL append(PLL cur, char c) { return (cur * base + c) %
    mod: }
/// prepends c to string with size k
PLL prepend(PLL cur, int k, char c) { return (P[k] * c +
    cur) % mod; }
/// replaces the i-th (0-indexed) character from right
    from a to b:
PLL replace(PLL cur, int i, char a, char b) {
 cur = (cur + P[i] * (b - a)) \% mod:
 return (cur + mod) % mod;
/// Erases c from the back of the string
PLL pop_back(PLL hash, char c) {
 return (((hash - c) * inv_base) % mod + mod) % mod;
/// Erases c from front of the string with size len
PLL pop_front(PLL hash, int len, char c) {
 return ((hash - P[len - 1] * c) % mod + mod) % mod;
/// concatenates two strings where length of the right is
PLL concat(PLL left, PLL right, int k) { return (left * P
    [k] + right) % mod; }
/// Calculates hash of string with size len repeated cnt
    times
/// This is O(\log n). For O(1), pre-calculate inverses
PLL repeat(PLL hash, int len, LL cnt) {
 PLL mul = (P[len * cnt] - 1) * inverse(P[len] - 1);
 mul = (mul % mod + mod) % mod;
 PLL ret = (hash * mul) % mod;
  if (P[len].ff == 1) ret.ff = hash.ff * cnt;
 if (P[len].ss == 1) ret.ss = hash.ss * cnt;
 return ret:
LL get(PLL hash) { return ((hash.ff << 32) ^ hash.ss); }
struct hashlist {
 int len:
  vector<PLL> H. R:
```

```
hashlist() {}
  hashlist(string& s) {
   len = (int)s.size();
   hash_init(len);
   H.resize(len + 1, PLL(0, 0)), R.resize(len + 2, PLL
        (0, 0):
   for (int i = 1; i <= len; i++) H[i] = append(H[i -</pre>
       1], s[i - 1]);
   for (int i = len; i >= 1; i--) R[i] = append(R[i +
       1], s[i - 1]);
 /// 1-indexed
 PLL range_hash(int 1, int r) {
   return ((H[r] - H[1 - 1] * P[r - 1 + 1]) \% mod + mod)
        % mod;
 }
 PLL reverse_hash(int 1, int r) {
   return ((R[1] - R[r + 1] * P[r - 1 + 1]) \% mod + mod)
        % mod:
 }
 PLL concat_range_hash(int 11, int r1, int 12, int r2) {
   return concat(range_hash(11, r1), range_hash(12, r2),
        r2 - 12 + 1);
 }
 PLL concat_reverse_hash(int 11, int r1, int 12, int r2)
   return concat(reverse_hash(12, r2), reverse_hash(11,
       r1), r1 - 11 + 1);
 }
};
```

7.3 Manacher's

```
if (i + k > r) l = i - k - 1, r = i + k;
7.4 Palindromic Tree
int tree[100010][26],len[100010],link[100010],node,curr;
string s;
void Clear(int n){
 for(int i=0:i<=n:i++){</pre>
   len[i]=0, link[i]=0;
   for(int j=0;j<26;j++) tree[i][j]=0;</pre>
void Init(){
 node=2, curr=2, len[1]=-1, len[2]=0, link[1]=1, link
      [2]=1;
void Extend(int pos){
 while(s[pos-len[curr]-1]!=s[pos]) curr=link[curr];
 int c=s[pos]-'a',x=link[curr];
  while(s[pos-len[x]-1]!=s[pos]) x=link[x];
  if(!tree[curr][c]){
   node++;
   tree[curr][c]=node:
   len[node] = len[curr] + 2;
   if(len[node] == 1) link[node] = 2;
   else link[node] = tree[x][c];
 }
 curr=tree[curr][c];
```

7.5 String Match FFT

```
//find occurrences of t in s where '?'s are automatically
     matched with any character
//res[i + m - 1] = sum_j=0 to m - 1_{s[i + j]} * t[j] * (s]
    [i + i] - t[i]
vector<int> string_matching(string &s, string &t) {
 int n = s.size(), m = t.size();
 vector\langle int \rangle s1(n), s2(n), s3(n);
 for(int i = 0; i < n; i++) s1[i] = s[i] == '?' ? 0 : s[
      i] - 'a' + 1; //assign any non zero number for non
      171s
 for(int i = 0; i < n; i++) s2[i] = s1[i] * s1[i];</pre>
 for(int i = 0; i < n; i++) s3[i] = s1[i] * s2[i];
 vector<int> t1(m), t2(m), t3(m);
 for(int i = 0; i < m; i++) t1[i] = t[i] == '?' ? 0 : t[
      i] - 'a' + 1:
 for(int i = 0; i < m; i++) t2[i] = t1[i] * t1[i];</pre>
 for(int i = 0; i < m; i++) t3[i] = t1[i] * t2[i];</pre>
 reverse(t1.begin(), t1.end());
 reverse(t2.begin(), t2.end());
```

7.6 Suffix Array

```
vector<int> SortCyclicShifts(string s){
 int n=s.size():
 const int alphabet=256;
 vector<int>Start(n),Freq(max(alphabet,n),0),EClass(n);
 for(int i=0;i<n;i++) Freq[s[i]]++;</pre>
 for(int i=1;i<alphabet;i++) Freq[i]+=Freq[i-1];</pre>
 for(int i=0;i<n;i++) Start[--Freq[s[i]]]=i;</pre>
 int Class=1;
 EClass[Start[0]]=0:
 for(int i=1;i<n;i++){</pre>
   if(s[Start[i]]!=s[Start[i-1]]) Class++;
   EClass[Start[i]]=Class-1;
 for(int p=0;(1<<p)<n;p++){</pre>
   vector<int> STemp(n),ECTemp(n);
   for(int i=0;i<n;i++){</pre>
       STemp[i]=Start[i]-(1<<p);</pre>
       if(STemp[i]<0) STemp[i]+=n;</pre>
   fill(Freq.begin(),Freq.begin()+Class,0);
   for(int i=0;i<n;i++) Freq[EClass[STemp[i]]]++;</pre>
   for(int i=1;i<Class;i++) Freq[i]+=Freq[i-1];</pre>
   for(int i=n-1;i>=0;i--){
       Start[--Freq[EClass[STemp[i]]]]=STemp[i];
   ECTemp[Start[0]]=0;
   Class=1;
   for(int i=1;i<n;i++){</pre>
     pair<int,int>prev,curr;
     prev={EClass[Start[i-1]],EClass[(Start[i-1]+(1<<p))</pre>
         %n]};
     curr={EClass[Start[i]],EClass[(Start[i]+(1<<p))%n</pre>
     if(curr!=prev) Class++;
     ECTemp[Start[i]]=Class-1;
```

```
for(int i=0:i<n:i++){</pre>
     EClass[Start[i]] = ECTemp[Start[i]];
   }
 }
 return Start;
vector<int> BuildSuffixArray(string s){
   s+='$':
   vector<int>ret=SortCyclicShifts(s);
   ret.erase(ret.begin());
   return ret;
vector<int>BuildLCPArray(string s,vector<int>&v){
   int n=s.size():
   vector<int>rank(n),lcp(n-1,0);
   for(int i=0;i<n;i++) rank[v[i]]=i;</pre>
   int k=0;
   for(int i=0;i<n;i++){</pre>
     if(rank[i]==n-1){
         k=0;
         continue;
     int j=v[rank[i]+1];
     while (i+k< n \&\& j+k< n \&\& s[i+k]==s[j+k]) k++;
     lcp[rank[i]] = k;
     if (k) k--;
   return lcp;
```

7.7 Suffix Automata

```
const int MAXN = 1e5 + 7, ALPHA = 26;
int len[2 * MAXN], link[2 * MAXN], nxt[2 * MAXN][ALPHA];
int sz;
int last;

void sa_init() {
   memset(nxt, -1, sizeof nxt);
   len[0] = 0, link[0] = -1, sz = 1, last = 0;
}

void add(char ch) {
   int c = ch - 'a';
   int cur = sz++;
   len[cur] = len[last] + 1;

int u = last;
   while (u != -1 && nxt[u][c] == -1) {
      nxt[u][c] = cur;
      u = link[u];
}
```

```
}
 if (u == -1) link[cur] = 0;
 else {
   int v = nxt[u][c];
   if (len[v] == len[u] + 1) link[cur] = v;
     int clone = sz++;
     len[clone] = 1 + len[u], link[clone] = link[v];
     for (int i = 0; i < ALPHA; i++) nxt[clone][i] = nxt</pre>
         [v][i];
     while (u != -1 && nxt[u][c] == v) {
       nxt[u][c] = clone;
       u = link[u]:
     link[v] = link[cur] = clone;
 }
 last = cur:
vector<int> edge[2 * MAXN];
/// Optional, Call after adding all characters
void makeEdge() {
 for (int i = 0; i < sz; i++) {</pre>
   edge[i].clear();
   for (int j = 0; j < ALPHA; j++)
     if (nxt[i][j] != -1) edge[i].push_back(j);
 }
// Given a string S, you have to answer some queries:
// If all distinct substrings of string S were sorted
// lexicographically, which one will be the K-th smallest
long long dp[2 * MAXN];
bool vis[2 * MAXN];
void dfs(int u) {
 if (vis[u]) return;
 vis[u] = 1;
 dp[u] = 1;
 for (int i : edge[u]) {
   if (nxt[u][i] == -1) continue;
   dfs(nxt[u][i]);
   dp[u] += dp[nxt[u][i]];
```

```
void go(int u, long long rem, string &s) {
   if (rem == 1) return;
   long long sum = 1;
   for (int i : edge[u]) {
     if (nxt[u][i] == -1) continue;
     if (sum + dp[nxt[u][i]] < rem) {
        sum += dp[nxt[u][i]];
     } else {
        s += ('a' + i);
        go(nxt[u][i], rem - sum, s);
        return;
     }
   }
}</pre>
```

7.8 Z Algo

```
vector<int> calcz(string s) {
 int n = s.size();
 vector<int> z(n);
 int 1 = 0, r = 0;
for (int i = 1; i < n; i++) {</pre>
  if (i > r) {
    1 = r = i:
    while (r < n \&\& s[r] == s[r - 1]) r++;
    z[i] = r - 1, r--;
  } else {
     int k = i - 1;
    if (z[k] < r - i + 1) z[i] = z[k];
    else {
      1 = i:
      while (r < n \&\& s[r] == s[r - 1]) r++;
      z[i] = r - 1, r--;
  }
return z:
```

Equations and Formulas

Catalan Numbers

$$C_n = \frac{1}{n+1} {2n \choose n}$$
 $C_0 = 1, C_1 = 1$ and $C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}$

The number of ways to completely parenthesize n+1 factors. The number of triangulations of a convex polygon with n+2sides (i.e. the number of partitions of polygon into disjoint triangles by using the diagonals).

form n disjoint i.e. non-intersecting chords.

The number of rooted full binary trees with n+1 leaves (ver- $S^d(n,k) = S(n-d+1,k-d+1), n \ge k \ge d$ tices are not numbered). A rooted binary tree is full if every 8.4 Other Combinatorial Identities vertex has either two children or no children.

Number of permutations of $1, \ldots, n$ that avoid the pattern 123 (or any of the other patterns of length 3); that is, the number of permutations with no three-term increasing sub-sequence. For n = 3, these permutations are 132, 213, 231, 312 and 321

8.2 Stirling Numbers First Kind

The Stirling numbers of the first kind count permutations according to their number of cycles (counting fixed points as cycles of length one).

S(n,k) counts the number of permutations of n elements with k disjoint cycles.

$$S(n,k) = (n-1) \cdot S(n-1,k) + S(n-1,k-1), where, S(0,0) = \sum_{n=0}^{\infty} S(n,k) = \sum_{n=0}^{\infty} S(n,k) + \sum_{n=0}^{\infty} S($$

$$1, S(n,0) = S(0,n) = 0 \sum_{k=0}^{n} S(n,k) = n!$$

The unsigned Stirling numbers may also be defined algebraically, as the coefficient of the rising factorial:

$$x^{\bar{n}} = x(x+1)...(x+n-1) = \sum_{k=0}^{n} S(n,k)x^{k}$$

Lets [n, k] be the stirling number of the first kind, then

$$\begin{bmatrix} n & n \\ n & -k \end{bmatrix} = \sum_{0 \le i_1 \le i_2 \le i_k \le n} i_1 i_2 \dots i_k.$$

8.3 Stirling Numbers Second Kind

Stirling number of the second kind is the number of ways to partition a set of n objects into k non-empty subsets.

$$S(n,k) = k \cdot S(n-1,k) + S(n-1,k-1)$$
, where $S(0,0) = 1$, $S(n,0) = S(0,n) = 0$ $S(n,2) = 2^{n-1} - 1$ $S(n,k) \cdot k! = \text{number of ways to color } n \text{ nodes using colors from 1 to } k \text{ such that } f m \text{ is any integer, then } \gcd(a+m\cdot b,b) = \gcd(a,b)$ The gcd is a multiplicative function in the follow

ber of ways to partition a set of n objects into k subsets, with $\gcd(a_1,b) \cdot \gcd(a_2,b)$.

each subset containing at least r elements. It is denoted by $\gcd(a, \operatorname{lcm}(b, c)) = \operatorname{lcm}(\gcd(a, b), \gcd(a, c))$. $S_r(n,k)$ and obeys the recurrence relation. $S_r(n+1,k) = |\operatorname{lcm}(a,\operatorname{gcd}(b,c))| = \operatorname{gcd}(\operatorname{lcm}(a,b),\operatorname{lcm}(a,c)).$ $\left| kS_r(n,k) + \binom{n}{r-1} S_r(n-r+1,k-1) \right|$

Denote the n objects to partition by the integers $1, 2, \dots, n$. $\gcd(a, b) = \sum \phi(k)$ Define the reduced Stirling numbers of the second kind, denoted $S^d(n,k)$, to be the number of ways to partition the integers $1, 2, \ldots, n$ into k nonempty subsets such that all ele- $\sum [\gcd(i, n) = k] = \phi(\frac{n}{k})$ ments in each subset have pairwise distance at least d. That i=1The number of ways to connect the 2n points on a circle to is, for any integers i and j in a given subset, it is required that $|i-j| \geq d$. It has been shown that these numbers satisfy,

$$\binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}$$

$$\sum_{i=0}^{k} \binom{n+i}{i} = \sum_{i=0}^{k} \binom{n+i}{n} = \binom{n+k+1}{k}$$

$$n, r \in N, n > r, \sum_{i=r}^{n} \binom{i}{r} = \binom{n+1}{r+1}$$
If $P(n) = \sum_{k=0}^{n} \binom{n}{k} \cdot Q(k)$, then,

$$Q(n) = \sum_{k=0}^{n} (-1)^{n-k} \binom{n}{k} \cdot P(k)$$

If
$$P(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot Q(k)$$
, then,

$$Q(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot P(k)$$

8.5 Different Math Formulas

Picks Theorem : A = i + b/2 - 1

Deragements: $d(i) = (i-1) \times (d(i-1) + d(i-2))$

$$\frac{n}{ab}$$
 - $\left\{\frac{b'n}{a}\right\}$ - $\left\{\frac{a'n}{b}\right\}$ +

The gcd is a multiplicative function in the following sense: An r-associated Stirling number of the second kind is the num- if a_1 and a_2 are relatively prime, then $gcd(a_1 \cdot a_2, b) =$

For non-negative integers a and b, where a and b are not both zero, $gcd(n^a - 1, n^b - 1) = n^{gcd(a,b)} - 1$ k|a and k|b $\sum_{k=1}^{n} \gcd(k, n) = \sum_{d \mid n} d \cdot \phi\left(\frac{n}{d}\right)$ $\sum_{k=1}^{n} x^{\gcd(k,n)} = \sum_{d \mid n} x^d \cdot \phi\left(\frac{n}{d}\right)$ $\sum_{k=1}^{n} \frac{1}{\gcd(k,n)} = \sum_{n=1}^{\infty} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{1}{n} \sum_{n=1}^{\infty} d \cdot \phi(d)$ $\sum_{k=1}^{n} \frac{k}{\gcd(k,n)} = \frac{n}{2} \cdot \sum_{d|n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{n}{2} \cdot \frac{1}{n} \cdot \sum_{d|n} d \cdot \phi(d)$ $\left| \sum_{k=1}^{n} \frac{n}{\gcd(k,n)} \right| = 2 * \sum_{k=1}^{n} \frac{k}{\gcd(k,n)} - 1, \text{ for } n > 1$ $\sum_{i=1}^{n} \sum_{j=1}^{n} [\gcd(i,j) = 1] = \sum_{i=1}^{n} \mu(d) \left\lfloor \frac{n}{d} \right\rfloor^{2}$ $\sum_{i=1}^{n} \sum_{j=1}^{n} \gcd(i,j) = \sum_{d=1}^{n} \phi(d) \left\lfloor \frac{n}{d} \right\rfloor^{2}$

$$\sum_{i=1}^{n} \sum_{j=1}^{n} i \cdot j [\gcd(i,j) - \sum_{d=1}^{n} \phi(a) \lfloor d \rfloor$$

$$\sum_{i=1}^{n} \sum_{j=1}^{n} i \cdot j [\gcd(i,j) = 1] = \sum_{i=1}^{n} \phi(i) i^{2}$$

$$F(n) = \sum_{i=1}^{n} \sum_{j=1}^{n} \operatorname{lcm}(i,j) = \sum_{l=1}^{n} \left(\frac{\left(1 + \lfloor \frac{n}{l} \rfloor\right) \left(\lfloor \frac{n}{l} \rfloor\right)}{2} \right)^{2} \sum_{d \mid l} \mu(d) l d$$