
B Sets, Etc.

Many chapters of this book touch on the elements of discrete mathematics. This appendix reviews more completely the notations, definitions, and elementary properties of sets, relations, functions, graphs, and trees. If you are already well versed in this material, you can probably just skim this chapter.

B.1 Sets

A *set* is a collection of distinguishable objects, called its *members* or *elements*. If an object x is a member of a set S , we write $x \in S$ (read “ x is a member of S ” or, more briefly, “ x is in S ”). If x is not a member of S , we write $x \notin S$. We can describe a set by explicitly listing its members as a list inside braces. For example, we can define a set S to contain precisely the numbers 1, 2, and 3 by writing $S = \{1, 2, 3\}$. Since 2 is a member of the set S , we can write $2 \in S$, and since 4 is not a member, we have $4 \notin S$. A set cannot contain the same object more than once,¹ and its elements are not ordered. Two sets A and B are *equal*, written $A = B$, if they contain the same elements. For example, $\{1, 2, 3, 1\} = \{1, 2, 3\} = \{3, 2, 1\}$.

We adopt special notations for frequently encountered sets:

- \emptyset denotes the *empty set*, that is, the set containing no members.
- \mathbb{Z} denotes the set of *integers*, that is, the set $\{\dots, -2, -1, 0, 1, 2, \dots\}$.
- \mathbb{R} denotes the set of *real numbers*.
- \mathbb{N} denotes the set of *natural numbers*, that is, the set $\{0, 1, 2, \dots\}$.²

¹A variation of a set, which can contain the same object more than once, is called a *multiset*.

²Some authors start the natural numbers with 1 instead of 0. The modern trend seems to be to start with 0.

If all the elements of a set A are contained in a set B , that is, if $x \in A$ implies $x \in B$, then we write $A \subseteq B$ and say that A is a **subset** of B . A set A is a **proper subset** of B , written $A \subset B$, if $A \subseteq B$ but $A \neq B$. (Some authors use the symbol “ \subset ” to denote the ordinary subset relation, rather than the proper-subset relation.) For any set A , we have $A \subseteq A$. For two sets A and B , we have $A = B$ if and only if $A \subseteq B$ and $B \subseteq A$. For any three sets A , B , and C , if $A \subseteq B$ and $B \subseteq C$, then $A \subseteq C$. For any set A , we have $\emptyset \subseteq A$.

We sometimes define sets in terms of other sets. Given a set A , we can define a set $B \subseteq A$ by stating a property that distinguishes the elements of B . For example, we can define the set of even integers by $\{x : x \in \mathbb{Z} \text{ and } x/2 \text{ is an integer}\}$. The colon in this notation is read “such that.” (Some authors use a vertical bar in place of the colon.)

Given two sets A and B , we can also define new sets by applying **set operations**:

- The **intersection** of sets A and B is the set

$$A \cap B = \{x : x \in A \text{ and } x \in B\} .$$

- The **union** of sets A and B is the set

$$A \cup B = \{x : x \in A \text{ or } x \in B\} .$$

- The **difference** between two sets A and B is the set

$$A - B = \{x : x \in A \text{ and } x \notin B\} .$$

Set operations obey the following laws:

Empty set laws:

$$A \cap \emptyset = \emptyset ,$$

$$A \cup \emptyset = A .$$

Idempotency laws:

$$A \cap A = A ,$$

$$A \cup A = A .$$

Commutative laws:

$$A \cap B = B \cap A ,$$

$$A \cup B = B \cup A .$$

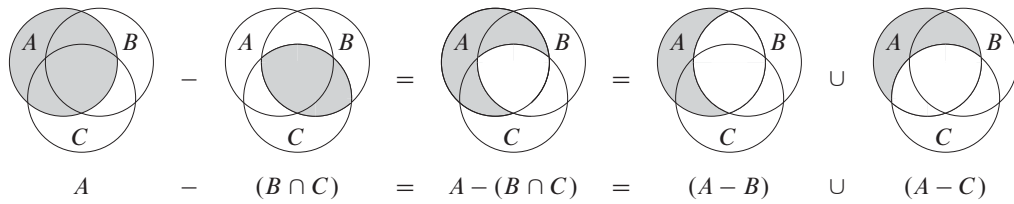


Figure B.1 A Venn diagram illustrating the first of DeMorgan's laws (B.2). Each of the sets A , B , and C is represented as a circle.

Associative laws:

$$\begin{aligned} A \cap (B \cap C) &= (A \cap B) \cap C, \\ A \cup (B \cup C) &= (A \cup B) \cup C. \end{aligned}$$

Distributive laws:

$$\begin{aligned} A \cap (B \cup C) &= (A \cap B) \cup (A \cap C), \\ A \cup (B \cap C) &= (A \cup B) \cap (A \cup C). \end{aligned} \tag{B.1}$$

Absorption laws:

$$\begin{aligned} A \cap (A \cup B) &= A, \\ A \cup (A \cap B) &= A. \end{aligned}$$

DeMorgan's laws:

$$\begin{aligned} A - (B \cap C) &= (A - B) \cup (A - C), \\ A - (B \cup C) &= (A - B) \cap (A - C). \end{aligned} \tag{B.2}$$

Figure B.1 illustrates the first of DeMorgan's laws, using a **Venn diagram**: a graphical picture in which sets are represented as regions of the plane.

Often, all the sets under consideration are subsets of some larger set U called the **universe**. For example, if we are considering various sets made up only of integers, the set \mathbb{Z} of integers is an appropriate universe. Given a universe U , we define the **complement** of a set A as $\bar{A} = U - A = \{x : x \in U \text{ and } x \notin A\}$. For any set $A \subseteq U$, we have the following laws:

$$\begin{aligned} \overline{\bar{A}} &= A, \\ A \cap \bar{A} &= \emptyset, \\ A \cup \bar{A} &= U. \end{aligned}$$

We can rewrite DeMorgan's laws (B.2) with set complements. For any two sets $B, C \subseteq U$, we have

$$\begin{aligned}\overline{B \cap C} &= \overline{B} \cup \overline{C}, \\ \overline{B \cup C} &= \overline{B} \cap \overline{C}.\end{aligned}$$

Two sets A and B are **disjoint** if they have no elements in common, that is, if $A \cap B = \emptyset$. A collection $\mathcal{S} = \{S_i\}$ of nonempty sets forms a **partition** of a set S if

- the sets are **pairwise disjoint**, that is, $S_i, S_j \in \mathcal{S}$ and $i \neq j$ imply $S_i \cap S_j = \emptyset$, and
- their union is S , that is,

$$S = \bigcup_{S_i \in \mathcal{S}} S_i.$$

In other words, \mathcal{S} forms a partition of S if each element of S appears in exactly one $S_i \in \mathcal{S}$.

The number of elements in a set is the **cardinality** (or **size**) of the set, denoted $|S|$. Two sets have the same cardinality if their elements can be put into a one-to-one correspondence. The cardinality of the empty set is $|\emptyset| = 0$. If the cardinality of a set is a natural number, we say the set is **finite**; otherwise, it is **infinite**. An infinite set that can be put into a one-to-one correspondence with the natural numbers \mathbb{N} is **countably infinite**; otherwise, it is **uncountable**. For example, the integers \mathbb{Z} are countable, but the reals \mathbb{R} are uncountable.

For any two finite sets A and B , we have the identity

$$|A \cup B| = |A| + |B| - |A \cap B|, \quad (\text{B.3})$$

from which we can conclude that

$$|A \cup B| \leq |A| + |B|.$$

If A and B are disjoint, then $|A \cap B| = 0$ and thus $|A \cup B| = |A| + |B|$. If $A \subseteq B$, then $|A| \leq |B|$.

A finite set of n elements is sometimes called an ***n*-set**. A 1-set is called a **singleton**. A subset of k elements of a set is sometimes called a ***k*-subset**.

We denote the set of all subsets of a set S , including the empty set and S itself, by 2^S ; we call 2^S the **power set** of S . For example, $2^{\{a,b\}} = \{\emptyset, \{a\}, \{b\}, \{a,b\}\}$. The power set of a finite set S has cardinality $2^{|S|}$ (see Exercise B.1-5).

We sometimes care about setlike structures in which the elements are ordered. An **ordered pair** of two elements a and b is denoted (a, b) and is defined formally as the set $(a, b) = \{a, \{a, b\}\}$. Thus, the ordered pair (a, b) is *not* the same as the ordered pair (b, a) .

The **Cartesian product** of two sets A and B , denoted $A \times B$, is the set of all ordered pairs such that the first element of the pair is an element of A and the second is an element of B . More formally,

$$A \times B = \{(a, b) : a \in A \text{ and } b \in B\} .$$

For example, $\{a, b\} \times \{a, b, c\} = \{(a, a), (a, b), (a, c), (b, a), (b, b), (b, c)\}$. When A and B are finite sets, the cardinality of their Cartesian product is

$$|A \times B| = |A| \cdot |B| . \tag{B.4}$$

The Cartesian product of n sets A_1, A_2, \dots, A_n is the set of **n -tuples**

$$A_1 \times A_2 \times \cdots \times A_n = \{(a_1, a_2, \dots, a_n) : a_i \in A_i \text{ for } i = 1, 2, \dots, n\} ,$$

whose cardinality is

$$|A_1 \times A_2 \times \cdots \times A_n| = |A_1| \cdot |A_2| \cdots |A_n|$$

if all sets are finite. We denote an n -fold Cartesian product over a single set A by the set

$$A^n = A \times A \times \cdots \times A ,$$

whose cardinality is $|A^n| = |A|^n$ if A is finite. We can also view an n -tuple as a finite sequence of length n (see page 1166).

Exercises

B.1-1

Draw Venn diagrams that illustrate the first of the distributive laws (B.1).

B.1-2

Prove the generalization of DeMorgan's laws to any finite collection of sets:

$$\begin{aligned} \overline{A_1 \cap A_2 \cap \cdots \cap A_n} &= \overline{A_1} \cup \overline{A_2} \cup \cdots \cup \overline{A_n} , \\ \overline{A_1 \cup A_2 \cup \cdots \cup A_n} &= \overline{A_1} \cap \overline{A_2} \cap \cdots \cap \overline{A_n} . \end{aligned}$$

B.1-3 ★

Prove the generalization of equation (B.3), which is called the *principle of inclusion and exclusion*:

$$\begin{aligned}
 |A_1 \cup A_2 \cup \cdots \cup A_n| = & \\
 & |A_1| + |A_2| + \cdots + |A_n| \\
 & - |A_1 \cap A_2| - |A_1 \cap A_3| - \cdots \quad (\text{all pairs}) \\
 & + |A_1 \cap A_2 \cap A_3| + \cdots \quad (\text{all triples}) \\
 & \vdots \\
 & + (-1)^{n-1} |A_1 \cap A_2 \cap \cdots \cap A_n| .
 \end{aligned}$$

B.1-4

Show that the set of odd natural numbers is countable.

B.1-5

Show that for any finite set S , the power set 2^S has $2^{|S|}$ elements (that is, there are $2^{|S|}$ distinct subsets of S).

B.1-6

Give an inductive definition for an n -tuple by extending the set-theoretic definition for an ordered pair.

B.2 Relations

A **binary relation** R on two sets A and B is a subset of the Cartesian product $A \times B$. If $(a, b) \in R$, we sometimes write $a R b$. When we say that R is a binary relation on a set A , we mean that R is a subset of $A \times A$. For example, the “less than” relation on the natural numbers is the set $\{(a, b) : a, b \in \mathbb{N} \text{ and } a < b\}$. An n -ary relation on sets A_1, A_2, \dots, A_n is a subset of $A_1 \times A_2 \times \cdots \times A_n$.

A binary relation $R \subseteq A \times A$ is **reflexive** if

$$a R a$$

for all $a \in A$. For example, “=” and “ \leq ” are reflexive relations on \mathbb{N} , but “<” is not. The relation R is **symmetric** if

$$a R b \text{ implies } b R a$$

for all $a, b \in A$. For example, “=” is symmetric, but “<” and “ \leq ” are not. The relation R is **transitive** if

$$a R b \text{ and } b R c \text{ imply } a R c$$

for all $a, b, c \in A$. For example, the relations “ $<$,” “ \leq ,” and “ $=$ ” are transitive, but the relation $R = \{(a, b) : a, b \in \mathbb{N} \text{ and } a = b - 1\}$ is not, since $3 R 4$ and $4 R 5$ do not imply $3 R 5$.

A relation that is reflexive, symmetric, and transitive is an **equivalence relation**. For example, “ $=$ ” is an equivalence relation on the natural numbers, but “ $<$ ” is not. If R is an equivalence relation on a set A , then for $a \in A$, the **equivalence class** of a is the set $[a] = \{b \in A : a R b\}$, that is, the set of all elements equivalent to a . For example, if we define $R = \{(a, b) : a, b \in \mathbb{N} \text{ and } a + b \text{ is an even number}\}$, then R is an equivalence relation, since $a + a$ is even (reflexive), $a + b$ is even implies $b + a$ is even (symmetric), and $a + b$ is even and $b + c$ is even imply $a + c$ is even (transitive). The equivalence class of 4 is $[4] = \{0, 2, 4, 6, \dots\}$, and the equivalence class of 3 is $[3] = \{1, 3, 5, 7, \dots\}$. A basic theorem of equivalence classes is the following.

Theorem B.1 (An equivalence relation is the same as a partition)

The equivalence classes of any equivalence relation R on a set A form a partition of A , and any partition of A determines an equivalence relation on A for which the sets in the partition are the equivalence classes.

Proof For the first part of the proof, we must show that the equivalence classes of R are nonempty, pairwise-disjoint sets whose union is A . Because R is reflexive, $a \in [a]$, and so the equivalence classes are nonempty; moreover, since every element $a \in A$ belongs to the equivalence class $[a]$, the union of the equivalence classes is A . It remains to show that the equivalence classes are pairwise disjoint, that is, if two equivalence classes $[a]$ and $[b]$ have an element c in common, then they are in fact the same set. Suppose that $a R c$ and $b R c$. By symmetry, $c R b$, and by transitivity, $a R b$. Thus, for any arbitrary element $x \in [a]$, we have $x R a$ and, by transitivity, $x R b$, and thus $[a] \subseteq [b]$. Similarly, $[b] \subseteq [a]$, and thus $[a] = [b]$.

For the second part of the proof, let $\mathcal{A} = \{A_i\}$ be a partition of A , and define $R = \{(a, b) : \text{there exists } i \text{ such that } a \in A_i \text{ and } b \in A_i\}$. We claim that R is an equivalence relation on A . Reflexivity holds, since $a \in A_i$ implies $a R a$. Symmetry holds, because if $a R b$, then a and b are in the same set A_i , and hence $b R a$. If $a R b$ and $b R c$, then all three elements are in the same set A_i , and thus $a R c$ and transitivity holds. To see that the sets in the partition are the equivalence classes of R , observe that if $a \in A_i$, then $x \in [a]$ implies $x \in A_i$, and $x \in A_i$ implies $x \in [a]$. ■

A binary relation R on a set A is **antisymmetric** if
 $a R b$ and $b R a$ imply $a = b$.

For example, the “ \leq ” relation on the natural numbers is antisymmetric, since $a \leq b$ and $b \leq a$ imply $a = b$. A relation that is reflexive, antisymmetric, and transitive is a **partial order**, and we call a set on which a partial order is defined a **partially ordered set**. For example, the relation “is a descendant of” is a partial order on the set of all people (if we view individuals as being their own descendants).

In a partially ordered set A , there may be no single “maximum” element a such that $b R a$ for all $b \in A$. Instead, the set may contain several **maximal** elements a such that for no $b \in A$, where $b \neq a$, is it the case that $a R b$. For example, a collection of different-sized boxes may contain several maximal boxes that don’t fit inside any other box, yet it has no single “maximum” box into which any other box will fit.³

A relation R on a set A is a **total relation** if for all $a, b \in A$, we have $a R b$ or $b R a$ (or both), that is, if every pairing of elements of A is related by R . A partial order that is also a total relation is a **total order** or **linear order**. For example, the relation “ \leq ” is a total order on the natural numbers, but the “is a descendant of” relation is not a total order on the set of all people, since there are individuals neither of whom is descended from the other. A total relation that is transitive, but not necessarily reflexive and antisymmetric, is a **total preorder**.

Exercises

B.2-1

Prove that the subset relation “ \subseteq ” on all subsets of \mathbb{Z} is a partial order but not a total order.

B.2-2

Show that for any positive integer n , the relation “equivalent modulo n ” is an equivalence relation on the integers. (We say that $a \equiv b \pmod{n}$ if there exists an integer q such that $a - b = qn$.) Into what equivalence classes does this relation partition the integers?

B.2-3

Give examples of relations that are

- a. reflexive and symmetric but not transitive,
- b. reflexive and transitive but not symmetric,
- c. symmetric and transitive but not reflexive.

³To be precise, in order for the “fit inside” relation to be a partial order, we need to view a box as fitting inside itself.

B.2-4

Let S be a finite set, and let R be an equivalence relation on $S \times S$. Show that if in addition R is antisymmetric, then the equivalence classes of S with respect to R are singletons.

B.2-5

Professor Narcissus claims that if a relation R is symmetric and transitive, then it is also reflexive. He offers the following proof. By symmetry, $a R b$ implies $b R a$. Transitivity, therefore, implies $a R a$. Is the professor correct?

B.3 Functions

Given two sets A and B , a **function** f is a binary relation on A and B such that for all $a \in A$, there exists precisely one $b \in B$ such that $(a, b) \in f$. The set A is called the **domain** of f , and the set B is called the **codomain** of f . We sometimes write $f : A \rightarrow B$; and if $(a, b) \in f$, we write $b = f(a)$, since b is uniquely determined by the choice of a .

Intuitively, the function f assigns an element of B to each element of A . No element of A is assigned two different elements of B , but the same element of B can be assigned to two different elements of A . For example, the binary relation

$$f = \{(a, b) : a, b \in \mathbb{N} \text{ and } b = a \bmod 2\}$$

is a function $f : \mathbb{N} \rightarrow \{0, 1\}$, since for each natural number a , there is exactly one value b in $\{0, 1\}$ such that $b = a \bmod 2$. For this example, $0 = f(0)$, $1 = f(1)$, $0 = f(2)$, etc. In contrast, the binary relation

$$g = \{(a, b) : a, b \in \mathbb{N} \text{ and } a + b \text{ is even}\}$$

is not a function, since $(1, 3)$ and $(1, 5)$ are both in g , and thus for the choice $a = 1$, there is not precisely one b such that $(a, b) \in g$.

Given a function $f : A \rightarrow B$, if $b = f(a)$, we say that a is the **argument** of f and that b is the **value** of f at a . We can define a function by stating its value for every element of its domain. For example, we might define $f(n) = 2n$ for $n \in \mathbb{N}$, which means $f = \{(n, 2n) : n \in \mathbb{N}\}$. Two functions f and g are **equal** if they have the same domain and codomain and if, for all a in the domain, $f(a) = g(a)$.

A **finite sequence** of length n is a function f whose domain is the set of n integers $\{0, 1, \dots, n-1\}$. We often denote a finite sequence by listing its values: $\langle f(0), f(1), \dots, f(n-1) \rangle$. An **infinite sequence** is a function whose domain is the set \mathbb{N} of natural numbers. For example, the Fibonacci sequence, defined by recurrence (3.22), is the infinite sequence $\langle 0, 1, 1, 2, 3, 5, 8, 13, 21, \dots \rangle$.

When the domain of a function f is a Cartesian product, we often omit the extra parentheses surrounding the argument of f . For example, if we had a function $f : A_1 \times A_2 \times \cdots \times A_n \rightarrow B$, we would write $b = f(a_1, a_2, \dots, a_n)$ instead of $b = f((a_1, a_2, \dots, a_n))$. We also call each a_i an **argument** to the function f , though technically the (single) argument to f is the n -tuple (a_1, a_2, \dots, a_n) .

If $f : A \rightarrow B$ is a function and $b = f(a)$, then we sometimes say that b is the **image** of a under f . The image of a set $A' \subseteq A$ under f is defined by

$$f(A') = \{b \in B : b = f(a) \text{ for some } a \in A'\}.$$

The **range** of f is the image of its domain, that is, $f(A)$. For example, the range of the function $f : \mathbb{N} \rightarrow \mathbb{N}$ defined by $f(n) = 2n$ is $f(\mathbb{N}) = \{m : m = 2n \text{ for some } n \in \mathbb{N}\}$, in other words, the set of nonnegative even integers.

A function is a **surjection** if its range is its codomain. For example, the function $f(n) = \lfloor n/2 \rfloor$ is a surjective function from \mathbb{N} to \mathbb{N} , since every element in \mathbb{N} appears as the value of f for some argument. In contrast, the function $f(n) = 2n$ is not a surjective function from \mathbb{N} to \mathbb{N} , since no argument to f can produce 3 as a value. The function $f(n) = 2n$ is, however, a surjective function from the natural numbers to the even numbers. A surjection $f : A \rightarrow B$ is sometimes described as mapping A **onto** B . When we say that f is onto, we mean that it is surjective.

A function $f : A \rightarrow B$ is an **injection** if distinct arguments to f produce distinct values, that is, if $a \neq a'$ implies $f(a) \neq f(a')$. For example, the function $f(n) = 2n$ is an injective function from \mathbb{N} to \mathbb{N} , since each even number b is the image under f of at most one element of the domain, namely $b/2$. The function $f(n) = \lfloor n/2 \rfloor$ is not injective, since the value 1 is produced by two arguments: 2 and 3. An injection is sometimes called a **one-to-one** function.

A function $f : A \rightarrow B$ is a **bijection** if it is injective and surjective. For example, the function $f(n) = (-1)^n \lfloor n/2 \rfloor$ is a bijection from \mathbb{N} to \mathbb{Z} :

$$\begin{aligned} 0 &\rightarrow 0, \\ 1 &\rightarrow -1, \\ 2 &\rightarrow 1, \\ 3 &\rightarrow -2, \\ 4 &\rightarrow 2, \\ &\vdots \end{aligned}$$

The function is injective, since no element of \mathbb{Z} is the image of more than one element of \mathbb{N} . It is surjective, since every element of \mathbb{Z} appears as the image of some element of \mathbb{N} . Hence, the function is bijective. A bijection is sometimes called a **one-to-one correspondence**, since it pairs elements in the domain and codomain. A bijection from a set A to itself is sometimes called a **permutation**.

When a function f is bijective, we define its **inverse** f^{-1} as $f^{-1}(b) = a$ if and only if $f(a) = b$.

For example, the inverse of the function $f(n) = (-1)^n \lceil n/2 \rceil$ is

$$f^{-1}(m) = \begin{cases} 2m & \text{if } m \geq 0, \\ -2m - 1 & \text{if } m < 0. \end{cases}$$

Exercises

B.3-1

Let A and B be finite sets, and let $f : A \rightarrow B$ be a function. Show that

- a.* if f is injective, then $|A| \leq |B|$;
- b.* if f is surjective, then $|A| \geq |B|$.

B.3-2

Is the function $f(x) = x + 1$ bijective when the domain and the codomain are \mathbb{N} ?
Is it bijective when the domain and the codomain are \mathbb{Z} ?

B.3-3

Give a natural definition for the inverse of a binary relation such that if a relation is in fact a bijective function, its relational inverse is its functional inverse.

B.3-4 ★

Give a bijection from \mathbb{Z} to $\mathbb{Z} \times \mathbb{Z}$.