

A Quadrotomy for Partial Orders

Alex Prieger
J. Stanley Warford
Pepperdine University
Malibu, CA 90263

Abstract

This paper uses a calculational system to formally prove a quadrotomy for partial orders \preceq , which states that any element x of the partial order must be related to any other element y by exactly one of (a) $x \prec y$, (b) $x = y$, (c) $y \prec x$, or (d) $\text{incomp}(x, y)$, where \prec is the reflexive reduction of \preceq and $\text{incomp}(x, y)$ is defined as $\text{incomp}(x, y) \equiv \neg(x \preceq y) \wedge \neg(y \preceq x)$. The calculational system, developed by Dijkstra and Scholten and extended by Gries and Schneider in their text *A Logical Approach to Discrete Math*, is based on only four inference rules – Substitution, Leibniz, Equanimity, and Transitivity. Inference rules in the older Hilbert-style systems, notably modus ponens, appear as theorems in this calculational deductive system, which is used to prove algorithm correctness in computer science. The theorem presented in this paper is a generalization of the trichotomy theorem for integers proved by Gries and Schneider, which states that any integer n must be related to any other integer m by exactly one of (a) $n < m$, (b) $n = m$, or (c) $m < n$.

1 Introduction

1.1 Background

A proof calculus is a system for proving propositions through the use of axioms and inference rules. Two of the most popular proof calculi are the Hilbert system and the natural deduction system. The former is sometimes used in the high-school environment to aid in the teaching of geometry; the latter is typically taught in undergraduate symbolic logic courses. There is a stark distinction between the two systems, due to their different approaches to axiomatization and their use of inference rules. Natural deduction attempts to mimic how people naturally reason, and contains few or no axioms and numerous inference rules. The Hilbert system makes no such attempt to mimic natural language – it contains only one inference rule, modus ponens, and a hefty amount of axiom schemes.

A third proof calculus, which will be used in this paper, is the equational system. This system is widely applicable in the realm of computer science, incorporating the usual propositional and predicate calculus of the other proof systems and seamlessly extending them

to reason about sets, sequences, functions, programs, and graphs. The equational system was developed in the late 1980's by Dijkstra and Scholten and was then largely expanded by Gries and Schneider. A distinctive feature of it is its use of only four inference rules - Substitution, Leibniz, Equanimity, and Transitivity - and its parsimonious use of axioms. Its axiomatization strikes a pleasant balance between the extensive use of axioms in the Hilbert style and the lack of them in natural deduction.

The equational system's cumulative structure and its resemblance to high-school algebra make it the optimal proof calculus to expound the case analysis proof technique for which this paper presents a new metatheorem. Case analysis is sometimes referred to as proof by exhaustion, for it attempts to prove that an exhaustive set of cases leads to a specific outcome. It is instrumental in natural deduction, where it is used as an inference rule called disjunction exploitation. The technique has two stages: the first involves proving that there is an exhaustive set of cases that altogether represents all possibilities; the second presents a proof for each individual case.

1.2 Case analysis

Describe Gries and Schneider's treatment of case analysis including Shannon as justification of the metatheorem.

Include the metatheorem and show how Gries justifies it using Shannon.

Present the new theorem and show how it is a more direct justification of the case analysis metatheorem.

2 Results

Lemma 1: $b \prec c \vee b = c \equiv b \preceq c$, where \prec is the reflexive reduction of \preceq .

Proof:

$$\begin{aligned}
 & b \prec c \vee b = c \\
 = & \langle (14.15.4) \text{ Notation, } \langle b, c \rangle \in \rho \text{ and } b \rho c \text{ are interchangeable notations.} \rangle \\
 & \langle b, c \rangle \in \prec \vee b = c \\
 = & \langle (14.15.3) \text{ Identity lemma, } \langle x, y \rangle \in i_B \equiv x = y \rangle \\
 & \langle b, c \rangle \in \prec \vee \langle b, c \rangle \in i_B \\
 = & \langle (11.20) \text{ Axiom, Union, } v \in S \cup T \equiv v \in S \vee v \in T \rangle \\
 & \langle b, c \rangle \in \prec \cup i_B \\
 = & \langle (14.49b) \text{ If } \rho \text{ is a quasi order over a set B, then } \rho \cup i_B \text{ is a partial order} \rangle \\
 & \langle b, c \rangle \in \preceq \\
 = & \langle (14.15.4) \text{ Notation} \rangle \\
 & b \preceq c \quad \blacksquare
 \end{aligned}$$

Lemma 2: $(p \vee q) \wedge \neg(q \vee r) \equiv p \wedge \neg q \wedge \neg r$

Proof:

$$\begin{aligned}
& (p \vee q) \wedge \neg(q \vee r) \\
= & \langle (3.47b) \text{ De Morgan, } \neg(p \vee q) \equiv \neg p \wedge \neg q \rangle \\
& (p \vee q) \wedge \neg q \wedge \neg r \\
= & \langle (3.44a) \text{ Absorption, } p \wedge (\neg p \vee q) \equiv (p \wedge q), \text{ with } p, q := \neg q, p \text{ and with} \\
& (3.12) \text{ Double negation} \rangle \\
& p \wedge \neg q \wedge \neg r \quad \blacksquare
\end{aligned}$$

Lemma 3: $(p \Rightarrow q) \Rightarrow (p \wedge q \equiv p)$

Proof:

$$\begin{aligned}
& p \wedge q \equiv p \\
= & \langle (3.60) \text{ Implication, } p \Rightarrow q \equiv p \wedge q \equiv p \rangle \\
& p \Rightarrow q \\
\Leftarrow & \langle (3.71) \text{ Reflexivity of } \Rightarrow, p \Rightarrow p \rangle \\
& p \Rightarrow q \quad \blacksquare
\end{aligned}$$

Lemma 4: ρ is irreflexive $\Rightarrow (b\rho c \Rightarrow \neg(b = c))$

Proof: The proof is by (4.4) Deduction.

$$\begin{aligned}
& b\rho c \Rightarrow \neg(b = c) \\
= & \langle (3.61) \text{ Contrapositive, } p \Rightarrow q \equiv \neg q \Rightarrow \neg p, \text{ with (3.12) Double negation} \rangle \\
& b = c \Rightarrow \neg(b\rho c) \\
= & \langle (3.84b) \text{ Substitution, } (e = f) \Rightarrow E_e^z \equiv (e = f) \Rightarrow E_f^z \rangle \\
& b = c \Rightarrow \neg(b\rho b) \\
= & \langle \text{Assume the antecedent, } \rho \text{ is irreflexive, or } (\forall b \mid : \neg(b\rho b)) \rangle \\
& b = c \Rightarrow \text{true} \\
= & \langle (3.72) \text{ Right Zero of } \Rightarrow, p \Rightarrow \text{true} \equiv \text{true} \rangle \\
& \text{true} \quad \blacksquare
\end{aligned}$$

Lemma 5: ρ is antisymmetric $\wedge \rho$ is reflexive $\Rightarrow (b\rho c \wedge c\rho b \equiv b = c)$

Proof: The proof is by (4.4) Deduction. We will prove the consequent, $(b\rho c \wedge c\rho b \equiv b = c)$, using the conjuncts of the antecedent as theorems. We will prove the consequent using (4.7) Mutual implication.

The proof in the first direction follows.

$$\begin{aligned}
& b\rho c \wedge c\rho b \\
= & \langle \text{Assume the conjunct of the antecedent } \rho \text{ is antisymmetric, or} \\
& (\forall b, c \mid : b\rho c \wedge c\rho b \Rightarrow b = c) \rangle \\
& b = c \\
\Rightarrow & \langle (3.71) \text{ Reflexivity of } \Rightarrow, p \Rightarrow p \rangle \\
& b = c \quad \blacksquare
\end{aligned}$$

The proof in the second direction, $b = c \Rightarrow b \rho c \wedge c \rho b$, is done by another application of (4.4) Deduction.

Proof:

$$\begin{aligned}
 & b \rho c \wedge c \rho b \\
 = & \langle \text{Assume the antecedent, } b = c \rangle \\
 & b \rho b \wedge b \rho b \\
 = & \langle (3.38) \text{ Idempotency of } \wedge, p \wedge p \equiv p \rangle \\
 & b \rho b \\
 = & \langle \text{Assume the conjunct of the antecedent } \rho \text{ is reflexive, or } (\forall b | : b \rho b) \rangle \\
 & \text{true} \quad \blacksquare
 \end{aligned}$$

Theorem (a): $b \prec c \equiv \neg(c \prec b \vee b = c \vee \text{incomp}(b, c))$

Proof:

$$\begin{aligned}
 & \neg(c \prec b \vee b = c \vee \text{incomp}(b, c)) \\
 = & \langle \text{Lemma 1, } b \prec c \vee b = c \equiv b \preceq c \rangle \\
 & \neg(c \preceq b \vee \text{incomp}(b, c)) \\
 = & \langle (14.47.1) \text{ Definition, Incomparable, } \text{incomp}(b, c) \equiv \neg(b \preceq c) \wedge \neg(c \preceq b) \rangle \\
 & \neg(c \preceq b \vee (\neg(b \preceq c) \wedge \neg(c \preceq b))) \\
 = & \langle (3.44b) \text{ Absorption, } p \vee (\neg p \wedge q) \equiv p \vee q \rangle \\
 & \neg(c \preceq b \vee \neg(b \preceq c)) \\
 = & \langle (3.47b) \text{ De Morgan, } \neg(p \vee q) \equiv \neg p \wedge \neg q, \text{ with (3.12) Double negation} \rangle \\
 & b \preceq c \wedge \neg(c \preceq b) \\
 = & \langle \text{Lemma 1, twice} \rangle \\
 & (b \prec c \vee b = c) \wedge \neg(c \prec b \vee b = c) \\
 = & \langle \text{Lemma 2, } (p \vee q) \wedge \neg(q \vee r) \equiv p \wedge \neg q \wedge \neg r \rangle \\
 & b \prec c \wedge \neg(b = c) \wedge \neg(c \prec b) \\
 = & \langle \text{Lemma 3, } (p \Rightarrow q) \Rightarrow (p \wedge q \equiv p), \text{ with } p, q := b \prec c, \neg(c \prec b). \text{ The antecedent, } \\
 & b \prec c \Rightarrow \neg(c \prec b), \text{ is true by the definition of asymmetry and the fact that strict} \\
 & \text{orders are asymmetric.} \rangle \\
 & b \prec c \wedge \neg(b = c) \\
 = & \langle \text{Lemma 3, with } p, q := b \prec c, \neg(b = c). \text{ The antecedent, } b \prec c \Rightarrow \neg(b = c), \text{ is} \\
 & \text{true by Lemma 4, } \rho \text{ is irreflexive } \Rightarrow (b \rho c \Rightarrow \neg(b = c)), \text{ and the fact that strict} \\
 & \text{orders are irreflexive.} \rangle \\
 & b \prec c \quad \blacksquare
 \end{aligned}$$

Theorem (b): $c \prec b \equiv \neg(b \prec c \vee b = c \vee \text{incomp}(b, c))$

Since $=$ and incomp are symmetric, the proof of (b) is identical to that of (a) with $b, c := c, b$.

Theorem (c): $b = c \equiv \neg(b \prec c \vee c \prec b \vee \text{incomp}(b, c))$

Proof:

$$\begin{aligned}
& \neg(b \prec c \vee c \prec b \vee \text{incomp}(b, c)) \\
= & \langle (3.26) \text{ Idempotency of } \vee, p \equiv p \vee p \rangle \\
& \neg(b \prec c \vee c \prec b \vee \text{incomp}(b, c) \vee \text{incomp}(b, c)) \\
= & \langle (14.48.2), \neg(b \preceq c) \equiv c \prec b \vee \text{incomp}(b, c), \text{ twice} \rangle \\
& \neg(\neg(b \preceq c) \vee \neg(c \preceq b)) \\
= & \langle (3.47a) \text{ De Morgan, } \neg(p \wedge q) \equiv \neg p \vee \neg q, \text{ with (3.12) Double negation} \rangle \\
& b \preceq c \wedge c \preceq b \\
= & \langle \text{Lemma 5, } \rho \text{ is antisymmetric} \wedge \rho \text{ is reflexive} \Rightarrow (b \rho c \wedge c \rho b \equiv b = c), \\
& \text{with the fact that partial orders are antisymmetric and reflexive} \rangle \\
& b = c \quad \blacksquare
\end{aligned}$$

Theorem (d): $\text{incomp}(b, c) \equiv \neg(b \prec c \vee c \prec b \vee b = c)$

Proof:

$$\begin{aligned}
& \neg(b \prec c \vee c \prec b \vee b = c) \\
= & \langle (3.26) \text{ Idempotency of } \vee, p \equiv p \vee p \rangle \\
& \neg(b \prec c \vee c \prec b \vee b = c \vee b = c) \\
= & \langle \text{Lemma 1, } b \prec c \vee b = c \equiv b \preceq c \text{ twice} \rangle \\
& \neg(b \preceq c \vee c \preceq b) \\
= & \langle (3.47b) \text{ De Morgan, } \neg(p \vee q) \equiv \neg p \wedge \neg q \rangle \\
& \neg(b \preceq c) \wedge \neg(c \preceq b) \\
= & \langle (14.47.1) \text{ Definition, Incomparable: } \text{incomp}(b, c) \equiv \neg(b \preceq c) \wedge \neg(c \preceq b) \rangle \\
& \text{incomp}(b, c) \quad \blacksquare
\end{aligned}$$

Four abbreviations for expressions will be introduced here to make the following proofs more concise. Note that replacing A, B, C , and D with $b \prec c, c \prec b, b = c$, and $\text{incomp}(b, c)$, in any order, will make the expressions identical to Theorems (a)-(d). However, the expressions in general describe any situation in which exactly one of A, B, C , and D is true.

$$\begin{aligned}
\alpha. A & \equiv \neg(B \vee C \vee D) \\
\beta. B & \equiv \neg(A \vee C \vee D) \\
\gamma. C & \equiv \neg(A \vee B \vee D) \\
\delta. D & \equiv \neg(A \vee B \vee C)
\end{aligned}$$

Theorem (e): $\alpha \wedge \beta \wedge \gamma \wedge \delta \Rightarrow A \vee B \vee C \vee D$

Proof: The proof is by (4.4) Deduction.

$$\begin{aligned}
& A \vee B \vee C \vee D \\
= & \langle \text{Assume conjunct } \alpha \text{ of the antecedent, } A \equiv \neg(B \vee C \vee D) \rangle \\
& \neg(B \vee C \vee D) \vee B \vee C \vee D \\
= & \langle (3.28) \text{ Excluded middle, } p \vee \neg p \rangle \\
& \text{true} \quad \blacksquare
\end{aligned}$$

Lemma 6: $\alpha \wedge \beta \wedge \gamma \wedge \delta \Rightarrow \neg(A \wedge B \wedge C)$

Proof: The proof is by (4.4) Deduction.

$$\begin{aligned}
& \neg(A \wedge B \wedge C) \\
= & \langle (3.47a) \text{ De Morgan, } \neg(p \wedge q) \equiv \neg p \vee \neg q, \text{ twice} \rangle \\
& \neg A \vee \neg B \vee \neg C \\
= & \langle \text{Assume conjunct } \alpha, A \equiv \neg(B \vee C \vee D), \text{ with (3.12) Double negation} \rangle \\
& B \vee C \vee D \vee \neg B \vee \neg C \\
= & \langle (3.28) \text{ Excluded middle, } p \vee \neg p, \text{ twice} \rangle \\
& \text{true} \vee \text{true} \vee D \\
= & \langle (3.29) \text{ Identity of } \vee, p \vee \text{true} \equiv \text{true} \rangle \\
& \text{true} \quad \blacksquare
\end{aligned}$$

Lemma 7: $\alpha \wedge \beta \wedge \gamma \wedge \delta \Rightarrow (A \wedge \neg B \wedge \neg C \equiv A)$

Proof: The proof is by (4.4) Deduction.

$$\begin{aligned}
& A \wedge \neg B \wedge \neg C \\
= & \langle \text{Assume conjunct } \alpha \text{ of the antecedent, } A \equiv \neg(B \vee C \vee D) \rangle \\
& \neg(B \vee C \vee D) \wedge \neg B \wedge \neg C \\
= & \langle (3.47b) \text{ De Morgan, } \neg(p \vee q) \equiv \neg p \wedge \neg q, \text{ twice} \rangle \\
& \neg B \wedge \neg C \wedge \neg D \wedge \neg B \wedge \neg C \\
= & \langle (3.40) \text{ Idempotency of } \wedge, p \wedge p \equiv p \rangle \\
& \neg B \wedge \neg C \wedge \neg D \\
= & \langle \text{Assume conjunct } \alpha \text{ of the antecedent, } A \equiv \neg(B \vee C \vee D) \rangle \\
& A \quad \blacksquare
\end{aligned}$$

Lemma 8: $(p \equiv q \equiv r) \equiv (p \wedge q \wedge r) \vee (p \wedge \neg q \wedge \neg r) \vee (\neg p \wedge q \wedge \neg r) \vee (\neg p \wedge \neg q \wedge r)$

Proof:

$$\begin{aligned}
& p \equiv q \equiv r \\
= & \langle (3.52) \text{ Equivalence, } p \equiv q \equiv (p \wedge q) \vee (\neg p \wedge \neg q) \rangle \\
& (p \wedge q) \vee (\neg p \wedge \neg q) \equiv r \\
= & \langle (3.52) \text{ Equivalence} \rangle \\
& (((p \wedge q) \vee (\neg p \wedge \neg q)) \wedge r) \vee (\neg((p \wedge q) \vee (\neg p \wedge \neg q)) \wedge \neg r) \\
= & \langle (3.46) \text{ Distributivity of } \wedge \text{ over } \vee, p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r) \rangle \\
& (p \wedge q \wedge r) \vee (\neg p \wedge \neg q \wedge r) \vee (\neg((p \wedge q) \vee (\neg p \wedge \neg q)) \wedge \neg r) \\
= & \langle (3.52) \text{ Equivalence} \rangle \\
& (p \wedge q \wedge r) \vee (\neg p \wedge \neg q \wedge r) \vee (\neg(p \equiv q) \wedge \neg r) \\
= & \langle (3.10) \text{ Definition of } \not\equiv, (p \not\equiv q) \equiv \neg(p \equiv q) \rangle \\
& (p \wedge q \wedge r) \vee (\neg p \wedge \neg q \wedge r) \vee ((p \not\equiv q) \wedge \neg r) \\
= & \langle (3.46) \text{ Exclusive or, } p \not\equiv q \equiv (\neg p \wedge q) \vee (p \wedge \neg q) \rangle \\
& (p \wedge q \wedge r) \vee (\neg p \wedge \neg q \wedge r) \vee (((\neg p \wedge q) \vee (p \wedge \neg q)) \wedge \neg r)
\end{aligned}$$

$$\begin{aligned}
&= \langle (3.46) \text{ Distributivity of } \wedge \text{ over } \vee \rangle \\
&\quad (p \wedge q \wedge r) \vee (\neg p \wedge \neg q \wedge r) \vee (p \wedge \neg q \wedge \neg r) \vee (\neg p \wedge q \wedge \neg r) \\
&= \langle (3.24) \text{ Symmetry of } \vee \rangle \\
&\quad (p \wedge q \wedge r) \vee (p \wedge \neg q \wedge \neg r) \vee (\neg p \wedge q \wedge \neg r) \vee (\neg p \wedge \neg q \wedge r) \quad \blacksquare
\end{aligned}$$

Lemma 9: $\text{incomp}(b, c) \equiv \text{false}$, where $\text{incomp}(b, c)$ refers to b and c being incomparable under \leq .

Proof:

$$\begin{aligned}
&\text{incomp}(b, c) \\
&= \langle (14.47.1) \text{ Definition, Incomparable, } \text{incomp}(b, c) \equiv \neg(b \preceq c) \wedge \neg(c \preceq b), \text{ where the partial order is } \leq \rangle \\
&\quad \neg(b \leq c) \wedge \neg(c \leq b) \\
&= \langle (3.47b) \text{ De Morgan, } \neg(p \vee q) \equiv \neg p \wedge \neg q \rangle \\
&\quad \neg(b \leq c \vee c \leq b) \\
&= \langle (14.50) \text{ Definition, Total Order: A partial order } \preceq \text{ over } B \text{ is called a total or linear order if } (\forall b, c \in B: b \preceq c \vee c \preceq b), \text{ with the fact that } \leq \text{ is a total order} \rangle \\
&\quad \neg \text{true} \\
&= \langle (3.8) \text{ Definition of false} \rangle \\
&\quad \text{false} \quad \blacksquare
\end{aligned}$$

(15.44) Trichotomy: $(a < b \equiv a = b \equiv a > b) \wedge \neg(a < b \wedge a = b \wedge a > b)$

Proof:

$$\begin{aligned}
&(a < b \equiv a = b \equiv a > b) \wedge \neg(a < b \wedge a = b \wedge a > b) \\
&= \langle \text{Lemma 6, } \alpha \wedge \beta \wedge \gamma \wedge \delta \Rightarrow \neg(A \wedge B \wedge C) \text{ with } A, B, C, D := a < b, a = b, a > b, \text{incomp}(a, b). \text{ The antecedent is true by Theorems (a)-(d)} \rangle \\
&\quad (a < b \equiv a = b \equiv a > b) \wedge \text{true} \\
&= \langle (3.39) \text{ Identity of } \wedge, p \wedge \text{true} \equiv p \rangle \\
&\quad (a < b \equiv a = b \equiv a > b) \\
&= \langle \text{Lemma 8, with } p, q, r := a < b, a = b, a > b \rangle \\
&\quad (a < b \wedge a = b \wedge a > b) \vee (a < b \wedge \neg(a = b) \wedge \neg(a > b)) \vee \\
&\quad (\neg(a < b) \wedge a = b \wedge \neg(a > b)) \vee (\neg(a < b) \wedge \neg(a = b) \wedge a > b) \\
&= \langle \text{Lemma 6 as in step 1, but with (3.12) Double negation} \rangle \\
&\quad \text{false} \vee (a < b \wedge \neg(a = b) \wedge \neg(a > b)) \vee (\neg(a < b) \wedge a = b \wedge \neg(a > b)) \vee \\
&\quad (\neg(a < b) \wedge \neg(a = b) \wedge a > b) \\
&= \langle (3.30) \text{ Identity of } \vee, p \vee \text{false} \equiv p \rangle \\
&\quad (a < b \wedge \neg(a = b) \wedge \neg(a > b)) \vee (\neg(a < b) \wedge a = b \wedge \neg(a > b)) \vee \\
&\quad (\neg(a < b) \wedge \neg(a = b) \wedge a > b) \\
&= \langle \text{Lemma 7, } \alpha \wedge \beta \wedge \gamma \wedge \delta \Rightarrow (A \wedge \neg B \wedge \neg C \equiv A), \text{ with } A, B, C, D := a < b, a = b, a > b, \text{incomp}(a, b). \text{ The conjuncts of the antecedent are satisfied by Theorems (a)-(d)} \rangle
\end{aligned}$$

$$\begin{aligned}
& a < b \vee (\neg(a < b) \wedge a = b \wedge \neg(a > b)) \vee (\neg(a < b) \wedge \neg(a = b) \wedge a > b) \\
= & \langle \text{Lemma 7, with } A, B, C, D := a = b, a < b, a > b, \text{incomp}(a, b). \text{ The conjuncts of} \\
& \text{the antecedent are satisfied as in the previous step. Note that the expressions} \\
& \text{replacing A and B are swapped compared to the previous step} \rangle \\
& a < b \vee a = b \vee (\neg(a < b) \wedge \neg(a = b) \wedge a > b) \\
= & \langle \text{Lemma 7, with } A, B, C, D := a < b, a = b, a > b, \text{incomp}(a, b). \text{ The conjuncts of} \\
& \text{the antecedent are satisfied as in the previous step. Note that the expressions are} \\
& \text{again swapped to make the lemma apply to the final disjunct} \rangle \\
& a < b \vee a = b \vee a > b \\
= & \langle \text{Theorem (d) with } \preceq := \leq, \text{ with (3.12) Double negation} \rangle \\
& \neg \text{incomp}(a, b) \\
= & \langle \text{Lemma 9} \rangle \\
& \text{true} \quad \blacksquare
\end{aligned}$$

3 Conclusion