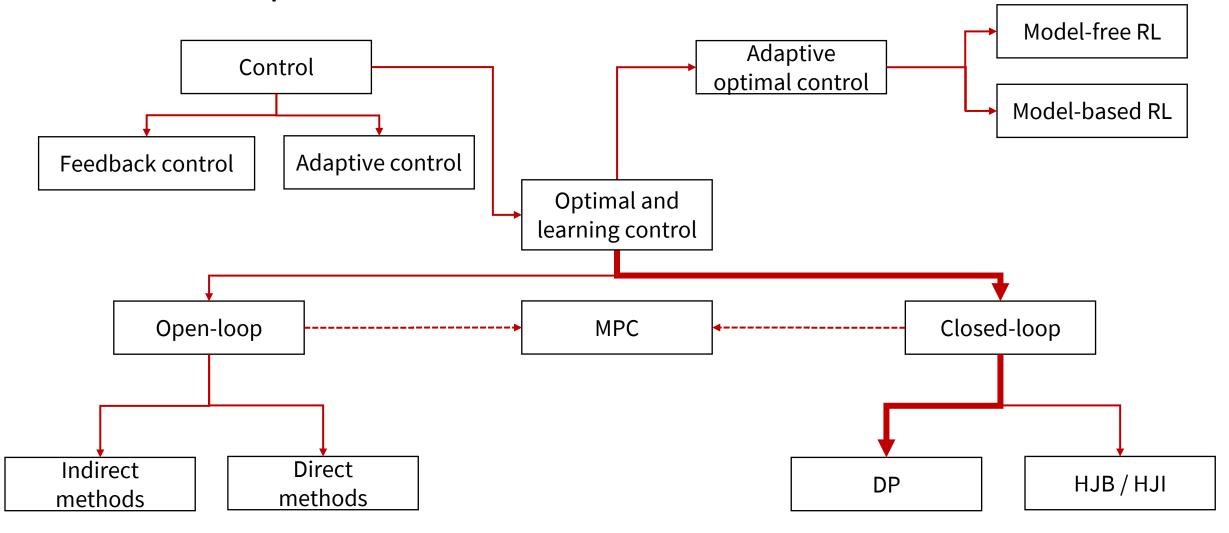
# AA203 Optimal and Learning-based Control

Intro to dynamic programming





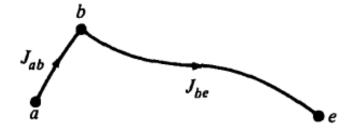
### Roadmap



## Principle of optimality

The key concept behind the dynamic programming approach is the principle of optimality

Suppose optimal path for a multi-stage decision-making problem is



- first decision yields segment a b with cost  $J_{ab}$
- remaining decisions yield segments b-e with cost  $J_{be}$
- optimal cost is then  $J_{ae}^* = J_{ab} + J_{be}$

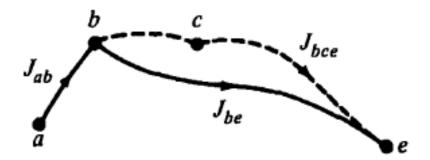
## Principle of optimality

- Claim: If a-b-e is optimal path from a to e, then b-e is optimal path from b to e
- Proof: Suppose b-c-e is the optimal path from b to e. Then

$$J_{bce} < J_{be}$$

and

$$J_{ab} + J_{bce} < J_{ab} + J_{be} = J_{ae}^*$$



**Contradiction!** 

## Principle of optimality

Principle of optimality (for discrete-time systems): Let  $\pi^*$ : =  $\{\pi_0^*, \pi_1^*, ..., \pi_{N-1}^*\}$  be an optimal policy. Assume state  $\mathbf{x}_k$  is reachable. Consider the subproblem whereby we are at  $\mathbf{x}_k$  at time k and we wish to minimize the cost-to-go from time k to time k. Then the truncated policy  $\{\pi_k^*, \pi_{k+1}^*, ..., \pi_{N-1}^*\}$  is optimal for the subproblem

- tail policies optimal for tail subproblems
- notation:  $\pi_k^*(\mathbf{x}_k) = \pi^*(\mathbf{x}_k, k)$

## Applying the principle of optimality

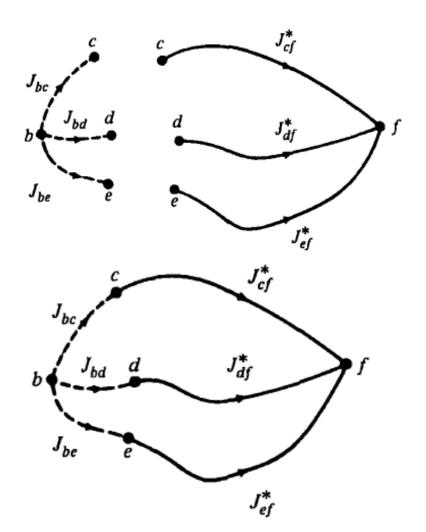
Principle of optimality: if b-c is the initial segment of the optimal path from b to f, then c-f is the terminal segment of this path

Hence, the optimal trajectory is found by comparing:

$$C_{bcf} = J_{bc} + J_{cf}^*$$

$$C_{bdf} = J_{bd} + J_{df}^*$$

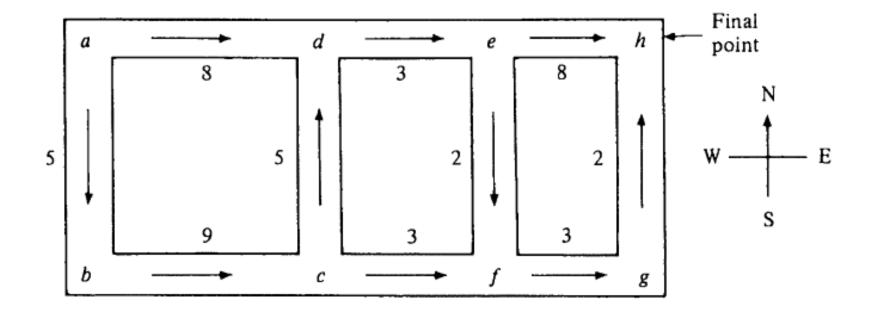
$$C_{bef} = J_{be} + J_{ef}^*$$



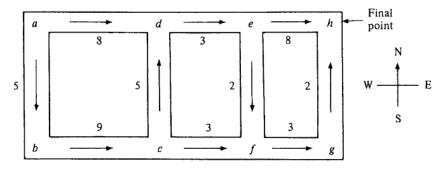
## Applying the principle of optimality

- need only to compare the concatenations of immediate decisions and optimal decisions
   → significant decrease in computation / possibilities
- in practice: carry out this procedure backward in time

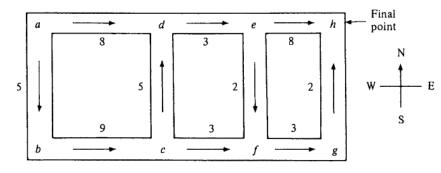
## Example



## Example



## Example



Optimal cost: 18; Optimal path:  $a \rightarrow d \rightarrow e \rightarrow f \rightarrow g \rightarrow h$ 

4/8/2021 AA 203 | Lecture 3

### DP Algorithm

- Model:  $\mathbf{x}_{k+1} = f(\mathbf{x}_k, \mathbf{u}_k, k), \quad \mathbf{u}_k \in U(\mathbf{x}_k)$
- Cost:  $J(\mathbf{x}_0) = h_N(\mathbf{x}_N) + \sum_{k=0}^{N-1} g(\mathbf{x}_k, \pi_k(\mathbf{x}_k), k)$

DP Algorithm: For every initial state  $\mathbf{x}_0$ , the optimal cost  $J^*(\mathbf{x}_0)$  is equal to  $J_0(\mathbf{x}_0)$ , given by the last step of the following algorithm, which proceeds backward in time from stage N-1 to stage 0:

$$J_N(\mathbf{x}_N) = h_N(\mathbf{x}_N)$$

$$J_k(\mathbf{x}_k) = \min_{\mathbf{u}_k \in U(\mathbf{x}_k)} g(\mathbf{x}_k, \mathbf{u}_k, k) + J_{k+1}(f(\mathbf{x}_k, \mathbf{u}_k, k)), \qquad k = 0, \dots, N-1$$

Furthermore, if  $\mathbf{u}_k^* = \pi_k^*(\mathbf{x}_k)$  minimizes the right hand side of the above equation for each  $\mathbf{x}_k$  and k, the policy  $\{\pi_0^*, \pi_1^*, \dots, \pi_{N-1}^*\}$  is optimal

#### Comments

- discretization (from differential equations to difference equations)
- quantization (from continuous to discrete state variables / controls)
- global minimum
- constraints, in general, simplify the numerical procedure
- optimal control in closed-loop form
- curse of dimensionality

#### Next time

- Canonical application: Discrete Linear Quadratic Regulator (LQR)
- Stochastic DP

$$V^*(x) = \max_{u} \left( R(x, u) + \gamma \sum_{x' \in \mathcal{X}} T(x'|x, u) V^*(x') \right)$$