

# 1 Introduction

2 Throughout the paper, every time we speak about matrices we mean binary  
3 matrices (also called 01-matrices) and we omit the word binary. If we speak  
4 about a *pattern*, we again mean a binary matrix and we use the word in order to  
5 distinguish among more matrices as well as to indicate relationship.

6 When dealing with matrices, we always index rows and column starting with  
7 one and when we speak about a row  $r$ , we simply mean a row with index  $r$ . A  
8 *line* is a common word for both a row and a column. When we order a set of  
9 lines, we first put all rows and then all columns. For  $M \in \{0, 1\}^{m \times n}$ ,  $[m]$  is a set  
10 of all rows and  $[m + n]$  is a set of all lines, where  $m$ -th element is the last row.  
11 This goes with the usual notation.

12 **Notation 1.** For  $n \in \mathbb{N}$  let  $[n] := \{1, 2, \dots, n\}$  and for  $m \in \mathbb{N}$ , where  $n \leq m$  let  
13  $[n, m] := \{n, n + 1, \dots, m\}$ .

14 **Notation 2.** For a matrix  $M \in \{0, 1\}^{m \times n}$  and  $L \subseteq [m + n]$  let  $M[L]$  denote a  
15 submatrix of  $M$  induced by lines in  $L$ .

16 **Notation 3.** For a matrix  $M \in \{0, 1\}^{m \times n}$ ,  $R \subseteq [m]$  and  $C \subseteq [n]$  let  $M[R, C]$   
17 denote a submatrix of  $M$  induced by rows in  $R$  and columns in  $C$ . Furthermore,  
18 for  $r \in [m]$  and  $c \in [n]$  let  $M[r, c] := M[\{r\}, \{c\}] = M[\{r, c + m\}]$ .

19 **Definition 1.** We say a matrix  $M \in \{0, 1\}^{m \times n}$  *contains* a pattern  $P \in \{0, 1\}^{k \times l}$   
20 *as a submatrix* and denote it by  $P \leq M$  if there are  $R \in [m]$  and  $C \in [n]$   
21 such that  $|R| = k$ ,  $|C| = l$  and for every  $r \in R$  and  $c \in C$  if  $P[r, c] = 1$ , then  
22  $M[R, C][r, c] = 1$ .

23 This does not necessarily mean  $P = M[R, C]$  as  $M[R, C]$  can have more  
24 one-entries than  $P$  does.

25 **Notation 4.** For a matrix  $M \in \{0, 1\}^{m \times n}$  and  $L \subseteq [m + n]$  let  $M_{\preceq}[L]$  denote a  
26 matrix acquired from  $M$  by applying following operation for each  $l \in L$ :

- 27 • If  $l$  is the first row in  $L$  then we replace the first  $l$  rows by one row that is  
28 a bitwise OR of replaced rows.
- 29 • If  $l$  is the first column in  $L$  then we replace the first  $l - m$  columns by one  
30 column that is a bitwise OR of replaced columns.
- 31 • Otherwise, we take  $l$ 's predecessor  $l' \in L$  in the standard ordering and  
32 replace lines  $[l' + 1, l]$  by one line that is a bitwise OR of replaced lines.

33 **Notation 5.** For a matrix  $M \in \{0, 1\}^{m \times n}$ ,  $R \subseteq [m]$  and  $C \subseteq [n]$  let  $M_{\preceq}[R, C] :=$   
34  $M_{\preceq}[R \cup \{c + m | c \in C\}]$ .

35 **Definition 2.** We say a matrix  $M \in \{0, 1\}^{m \times n}$  *contains* a pattern  $P \in \{0, 1\}^{k \times l}$   
36 *as an interval minor* and denote it by  $P \preceq M$  if there are  $R \in [m]$  and  $C \in [n]$   
37 such that  $|R| = k$ ,  $|C| = l$  and for every  $r \in R$  and  $c \in C$  if  $P[r, c] = 1$ , then  
38  $M_{\preceq}[R, C][r, c] = 1$ .

39 **Observation 1.** For all matrices  $M$  and  $P$ ,  $P \leq M \Rightarrow P \preceq M$ .

40 **Observation 2.** *For all matrices  $M$  and  $P$ , if  $P$  is a permutation matrix, then*  
41  *$P \leq M \Leftrightarrow P \preceq M$ .*

42 *Proof.* If we have  $P \preceq M$ , then there is a partitioning of  $M$  into rectangles and for  
43 each one-entry of  $P$  there is at least one one-entry in the corresponding rectangle  
44 of  $M$ . Since  $P$  is a permutation matrix, it is sufficient to take rows and columns  
45 having at least one one-entry in the right rectangle and we can always do so.

46 Together with Observation 1 this gives us the statement.  $\square$

## 0.1 Characterizations

**Definition 3.** A *walk* in a matrix  $M$  is a sequence of some of its entries, beginning in the top left corner and ending in the bottom right one. If an entry  $M[i, j]$  is in the sequence, the next one is either  $M[i + 1, j]$  or  $M[i, j + 1]$ .

**Definition 4.** We call a binary matrix  $M$  a *walking matrix* if there is a walk in  $M$  such that all one-entries of  $M$  are contained on the walk.

### 0.1.1 Patterns of size $2 \times 2$

**Theorem 3.** Let  $P = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$ , then for all  $M$ :  $P \not\preceq M \Leftrightarrow M$  is a walking matrix.

*Proof.* Since  $P$  is a permutation matrix,  $P \not\preceq M \Leftrightarrow P \not\preceq M$  and it is easy to see  $P \not\preceq M \Leftrightarrow M$  is a walking matrix.  $\square$

**Theorem 4.** Let  $P = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}$ , then for all  $M \in \{0, 1\}^{m \times n}$ :  $P \not\preceq M \Leftrightarrow$  there exist a row  $r$  and a column  $c$ , such that  $M[[r - 1], [c - 1]]$ ,  $M[[r - 1], [c + 1, n]]$  and  $M[[r + 1, m], [c - 1]]$  are empty and  $M[[r, m], [c, n]]$  is a walking matrix (see Figure 1).

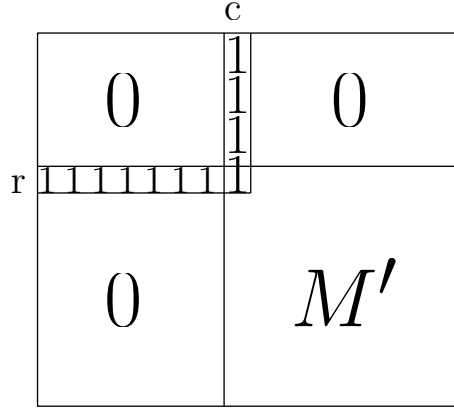


Figure 1: Characterization of a matrix avoiding  $\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}$  as an interval minor.

*Proof.*  $\Rightarrow$  If  $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \not\preceq M$ , then  $M$  is a walking matrix and we set  $r = c = 1$ . Otherwise, there are one-entries  $M[r, c']$  and  $M[r', c]$  such that  $r' < r$  and  $c' < c$ . If there is a one-entry in regions  $M[[r - 1], [c - 1]]$ ,  $M[[r - 1], [c + 1, n]]$  or  $M[[r + 1, m], [c - 1]]$  then  $P \preceq M$ . If  $M[[r, m], [c, n]]$  is not a walking matrix then it contains  $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$  and we again get a contradiction.

$\Leftarrow$  For contradiction, assume that  $M$  described in Figure 1 contains  $P$  as an interval minor. It means that there is a partition of the matrix into four quadrants such that there is at least one one-entry in each quadrant besides the bottom right one. If the matrix is partitioned above the  $r$ -th row, then there is only one column containing one-entries and it is not possible for both top quadrants to have a one-entry. Similarly, if the matrix is partitioned to the left of the  $c$ -th column, there is only one row containing one-entries and there is no one-entry in either top-left or bottom-left quadrant. Therefore,

74 the partitioning lies bellow the  $r$ -th row and to the right of the  $c$ -th column,  
 75 but if the quadrants contain one-entries, there is a  $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$  interval minor in  
 76  $M'$ , which is a contradiction with it being a walking matrix.  
 77 □

78 To characterize matrices avoiding  $\begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$  as an interval minor, we first need to  
 79 define a few useful terms.

80 **Definition 5.** For  $M \in \{0, 1\}^{m \times n}$  and  $r \in [m], c \in [n]$  we say  $M[r, c]$  is *top-left*  
 81 *empty* if  $M[[r - 1], [c - 1]]$  is an empty matrix. Similarly, it is *top-right empty* if  
 82  $M[[r - 1], [c + 1, n]]$  is empty and so on.

83 **Lemma 5.** Let  $P = \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$  and let  $M \in \{0, 1\}^{m \times n}$  avoid  $P$  as an interval minor,  
 84 then there exists a row  $r$  and a column  $c$  such that  $M[r, c]$  is either both *top-left*  
 85 *empty* and *bottom-right empty*, in which case  $[r, c] \notin \{[0, n - 1], [m - 1, 0]\}$  or both  
 86 *top-right empty* and *bottom-left empty*, in which case  $[r, c] \notin \{[0, 0], [m - 1, n - 1]\}$ .

87 *Proof.* □

88 **Theorem 6.** Let  $P = \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$ , then for all  $M$ :  $P \not\preceq M \Leftrightarrow M$  looks like one of the  
 89 matrices in Figure 2, where  $\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \not\preceq M_1$ ,  $\begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix} \not\preceq M_2$ ,  $\begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix} \not\preceq M_3$  and  $\begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix} \not\preceq M_4$ .

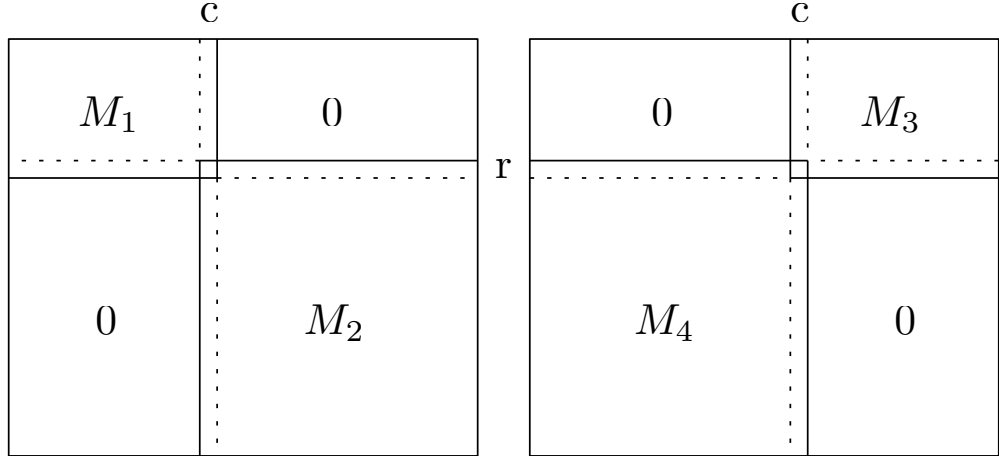


Figure 2: Characterization of a matrix avoiding  $\begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$  as an interval minor.

90 *Proof.*

91  $\Rightarrow$  We proceed by induction by the size of  $M$ .

92 If  $M \in \{0, 1\}^{2 \times 2}$  then it either avoids  $\begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix}$  or  $\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}$  and we are done.

93 For bigger  $M$  there is, from Lemma 5, “the element”. Assume the first case  
 94 (top-right and bottom-left empty (will change this when I have some notation)).  
 95 If  $M_1$  is non-empty, then  $\begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix} \not\preceq M_2$ ; otherwise,  $P \preceq M$ . Similarly,  $\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \not\preceq M_1$  if  
 96  $M_2$  is non-empty. If one of them is empty, the other is a smaller matrix avoiding  
 97  $P$  as an interval minor and by induction hypothesis, it can be partitioned. Adding  
 98 empty rows and columns does not break any condition and we get a partitioning  
 99 of the whole  $M$ .

100  $\Leftarrow$  Without loss of generality, let us assume  $M$  looks like the left matrix in Figure 2.  
 101 For contradiction, assume  $P \preceq M$ . In that case, we can partition  $M$  into four  
 102 quadrants such that there is at least one one-entry in each of them. It does not  
 103 matter where we partition it, every time we either get  $\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \preceq M_1$  or  $\begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix} \preceq M_2$ ,  
 104 which is a contradiction.  $\square$

## 105 0.1.2 Matrices of size $2 \times 3$

106 **Theorem 7.** Let  $P = \begin{pmatrix} 1 & 0 & 1 \\ 0 & 1 & 0 \end{pmatrix}$ , then for all  $M$ :  $P \not\preceq M \Leftrightarrow M = M_1 \oplus_h M_2$ , where  
 107  $\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \not\preceq M_1$  and  $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \not\preceq M_2$ .

108 *Proof.*  $\Rightarrow$  Let  $e = [r, c]$  be the top-most one-entry of  $M$ . If  $\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \preceq M[[m], [c -$   
 109  $1]]$ , together with  $e$  it would be the whole  $P$ . If  $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \not\preceq M[[m], [c, n]]$  then  
 110 we are done. Let us assume it is not the case and let  $e_{0,0}, e_{1,1}$  be any two  
 111 one-entries forming the forbidden pattern. Similarly, let  $\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \preceq M[[m], [c]]$   
 112 (else  $P \preceq M$ ) and let  $e_{0,1}, e_{1,0}$  be any two one-entries forming the forbidden  
 113 pattern. Now if we take  $e_{0,0}, e_{0,1}$  and  $e_{1,0}$  or  $e_{1,1}$  with bigger row, we get  
 114 the forbidden pattern  $P$  as an interval minor.

115  $\Leftarrow$  For contradiction, let us assume  $P \preceq M$  and  $M = M_1 \oplus_h M_2$ . If  $P \preceq M$ ,  
 116 look at the one-entry of  $M$  where the bottom one-entry of  $P$  is mapped. If  
 117 it is in  $M_1$  then  $P \not\preceq M$  because  $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \not\preceq M_1$ . Otherwise,  $P \not\preceq M$  because  
 118  $\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \not\preceq M_2$ .  $\square$

120 **Lemma 8.** Let  $P = \begin{pmatrix} 1 & 1 & 1 \\ 0 & 1 & 0 \end{pmatrix}$ , then for all  $M$ :  $P \not\preceq M \Rightarrow M = M_1 \oplus_h M_2$ , where  
 121  $\begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix} \not\preceq M_1$  and  $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \not\preceq M_2$  or  $\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \not\preceq M_1$  and  $\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \not\preceq M_2$ .

122 *Proof.* Let  $e = [r, c]$  be the top-most one-entry of  $M$ . If  $\begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix} \preceq M[[m], [c - 1]]$ ,  
 123 together with  $e$  it would be the whole  $P$ . Similarly,  $\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \not\preceq M[[m], [c + 1, n]]$ .  
 124 For contradiction with the statement, let  $\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \preceq M[[m], [c]]$  and  $e_{0,0}, e_{1,1}$  (non  
 125 of them equal to  $e$ , since  $e$  lies in the top-right corner) are any two one-entries  
 126 forming the pattern. Similarly, let  $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \preceq M[[m], [c, n]]$  and  $e_{0,1}, e_{1,0}$  are any  
 127 two one-entries forming the pattern. In that case  $e_{0,0}, e, e_{0,1}$  and  $e_{1,0}$  or  $e_{1,1}$   
 128 with bigger row give us the forbidden pattern  $P$  as an interval minor, which is a  
 129 contradiction.  $\square$

130 **Theorem 9.** Let  $P = \begin{pmatrix} 1 & 1 & 1 \\ 0 & 1 & 0 \end{pmatrix}$ , then for all  $M$ :  $P \not\preceq M \Leftrightarrow M$  looks like the matrix  
 131 in Figure 3 and  $\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \not\preceq M_1$  and  $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \not\preceq M_2$ .

132 *Proof.*  $\Rightarrow$  From Lemma 8 we know  $M = M'_1 \oplus_h M'_2$  where  $\begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix} \not\preceq M'_1$  and  
 133  $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \not\preceq M'_2$ . The second case would be dealt with symmetrically. From  
 134 Theorem 4 we have that  $M'_1$  can be characterized exactly like  $M[[m], [c_2 - 1]]$   
 135 and  $M[[m], [c_2, n]]$  forms a walking matrix. The only problem with our claim  
 136 would be if there were two different columns having a one-entry above the  
 137  $r$ -th row. In that case, those two one-entries together with a one-entry in  
 138 the  $r$ -th row between the columns  $c_1$  and  $c_2$  and a one-entry in the  $c_1$ -th  
 139 column above the  $r$ -th row form  $P$  as an interval minor.

		c <sub>1</sub>		c <sub>2</sub>	
	0	1	0	1	0
r	1	1	1	1	1
	$M_1$		0		$M_2$

Figure 3: Characterization of a matrix avoiding  $\begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$  as an interval minor.

140  $\Leftarrow$  The bottom-middle one-entry of  $P$  can not be mapped anywhere but to the  
141  $r$ -th row, but in that case there are at most two columns having one-entries  
142 above it.  
143 □

## 144 0.2 Extremal function

145 **Notation 6.** Let  $M$  be a matrix. We denote  $|M|$  the weight of  $M$ , the number  
146 of one-entries in  $M$ .

147 Usually  $|M|$  stands for a determinant of matrix  $M$ . However, in this paper  
148 we do not work with determinants at all so the notation should not lead to  
149 misunderstanding.

150 **Definition 6.** For a matrix  $P$  we define  $Ex(P, m, n) := \max\{|M| \mid M \in \{0, 1\}^{m \times n}, P \not\preceq$   
151  $M\}$ . We denote  $Ex(P, n) := Ex(P, n, n)$ .

152 **Definition 7.** For a matrix  $P$  we define  $Ex_{\preceq}(P, m, n) := \max\{|M| \mid M \in \{0, 1\}^{m \times n}, P \not\preceq$   
153  $M\}$ . We denote  $Ex_{\preceq}(P, n) := Ex_{\preceq}(P, n, n)$ .

154 **Observation 10.** For all  $P, m, n$ ;  $Ex_{\preceq}(P, m, n) \leq Ex(P, m, n)$ .

**Observation 11.** If  $P \in \{0, 1\}^{k \times l}$  has a one-entry at position  $[a, b]$ , then

$$Ex(P, m, n) \geq \begin{cases} m \cdot n & k > m \vee l > m \\ (k-1)n + (l-1)m - (k-1)(l-1) & \text{otherwise.} \end{cases}$$

155 **Observation 12.** The same holds for  $Ex_{\preceq}(P, m, n)$ .

**Definition 8.**  $P \in \{0, 1\}^{k \times l}$  is (strongly) minimalist if

$$Ex(P, m, n) = \begin{cases} m \cdot n & k > m \vee l > m \\ (k-1)n + (l-1)m - (k-1)(l-1) & \text{otherwise.} \end{cases}$$

**Definition 9.**  $P \in \{0, 1\}^{k \times l}$  is weakly minimalist if

$$Ex_{\preceq}(P, m, n) = \begin{cases} m \cdot n & k > m \vee l > m \\ (k-1)n + (l-1)m - (k-1)(l-1) & \text{otherwise.} \end{cases}$$

156 **Observation 13.** If  $P$  is strongly minimalist, then  $P$  is weakly minimalist.

### 157 0.2.1 Known results

158 **Fact 14.** 1.  $\begin{pmatrix} 1 \end{pmatrix}$  is strongly minimalist.

159 2. If  $P \in \{0, 1\}^{k \times l}$  is strongly minimalist and there is a one-entry in the last  
160 row in the  $c$ -th column, then  $P' \in \{0, 1\}^{k+1 \times l}$ , which is created from  $P$  by  
161 adding a new row having a one-entry only in the  $c$ -th column, is strongly  
162 minimalist.

163 3. If  $P$  is strongly minimalist, then after changing a one-entry into a zero-  
164 entry it is still strongly minimalist.

165 **Fact 15.** Let  $P = \begin{pmatrix} 1 & \dots & 1 \\ 1 & \dots & 1 \end{pmatrix}$  have  $l$  columns, then  $P$  is weakly minimalist.

*Proof.* Let  $M \in \{0, 1\}^{m \times n}$  be a matrix avoiding  $P = \{1\}^{2 \times l}$  as an interval minor and  $A_i = \{j \in [n] \mid \text{weight of } M[[i, \{j\}] > 0 \wedge \text{weight of } M[[i + 1, m], \{j\}] > 0\}$ . Clearly  $|A_i| \leq l - 1$ , otherwise  $P \preceq M$ . Let  $b_j$  denote the number of one-entries in the  $j$ -th column. Each column  $j$  of  $M$  appears in at least  $b_j - 1$  of sets  $A_i$ ,  $0 \leq i \leq m - 2$ . It follows that

$$\text{weight of } M = \sum_{j=0}^n b_j = \sum_{j=0}^n (b_j - 1) + n \leq \sum_{i=0}^{m-2} |A_i| + n \leq (l - 1)(m - 1) + n$$

166

□

167 This result is indeed very important because it shows that there are matrices  
168 like  $\begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$ , which are weakly minimalist, although it is known they are not strongly  
169 minimalist.

170 **Fact 16.** Let  $P = \begin{pmatrix} 1 & \dots & 1 \\ 1 & \dots & 1 \end{pmatrix}$  have  $l$  columns, then  $P$  is weakly minimalist.

*Proof.* Let  $M \in \{0, 1\}^{m \times n}$  be a matrix avoiding  $P = \{1\}^{3 \times l}$  as an interval minor and  $A_i = \{j \in [n] \mid \text{weight of } M[[i - 1], \{j\}] > 0 \wedge \text{weight of } M[[i + 1, m], \{j\}] > 0 \wedge M[i, j] \text{ one-entry}\}$ . Clearly  $|A_i| \leq l - 1$ , otherwise  $P \preceq M$ . Let  $b_j$  denote the number of one-entries in the  $j$ -th column. Each column  $j$  of  $M$  (for which  $b_j \geq 2$ ) appears in exactly  $b_j - 2$  of sets  $A_i$ ,  $1 \leq i \leq m - 1$ . It follows that

$$\text{weight of } M = \sum_{j=0}^n b_j = \sum_{j=0}^n (b_j - 2) + 2n \leq \sum_{i=1}^{m-2} |A_i| + 2n \leq (l - 1)(m - 2) + 2n$$

171

□