

1 Introduction

2 Throughout the paper, every time we speak about matrices we mean binary
3 matrices (also called 01-matrices) and we omit the word binary. If we speak
4 about a *pattern*, we again mean a binary matrix and we use the word in order to
5 distinguish among more matrices as well as to indicate relationship.

6 When dealing with matrices, we always index rows and column starting with
7 one and when we speak about a row r , we simply mean a row with index r . A
8 *line* is a common word for both a row and a column. When we order a set of
9 lines, we first put all rows and then all columns. For $M \in \{0, 1\}^{m \times n}$, $[m]$ is a set
10 of all rows and $[m + n]$ is a set of all lines, where m -th element is the last row.
11 This goes with the usual notation.

12 **Notation 1.** For $n \in \mathbb{N}$ let $[n] := \{1, 2, \dots, n\}$ and for $m \in \mathbb{N}$, where $n \leq m$ let
13 $[n, m] := \{n, n + 1, \dots, m\}$.

14 **Notation 2.** For a matrix $M \in \{0, 1\}^{m \times n}$ and $L \subseteq [m + n]$ let $M[L]$ denote a
15 submatrix of M induced by lines in L .

16 **Notation 3.** For a matrix $M \in \{0, 1\}^{m \times n}$, $R \subseteq [m]$ and $C \subseteq [n]$ let $M[R, C]$
17 denote a submatrix of M induced by rows in R and columns in C . Furthermore,
18 for $r \in [m]$ and $c \in [n]$ let $M[r, c] := M[\{r\}, \{c\}] = M[\{r, c + m\}]$.

19 **Definition 1.** We say a matrix $M \in \{0, 1\}^{m \times n}$ *contains* a pattern $P \in \{0, 1\}^{k \times l}$
20 *as a submatrix* and denote it by $P \leq M$ if there are $R \in [m]$ and $C \in [n]$
21 such that $|R| = k$, $|C| = l$ and for every $r \in R$ and $c \in C$ if $P[r, c] = 1$, then
22 $M[R, C][r, c] = 1$.

23 This does not necessarily mean $P = M[R, C]$ as $M[R, C]$ can have more
24 one-entries than P does.

25 **Notation 4.** For a matrix $M \in \{0, 1\}^{m \times n}$ and $L \subseteq [m + n]$ let $M_{\preceq}[L]$ denote a
26 matrix acquired from M by applying following operation for each $l \in L$:

- 27 • If l is the first row in L then we replace the first l rows by one row that is
28 a bitwise OR of replaced rows.
- 29 • If l is the first column in L then we replace the first $l - m$ columns by one
30 column that is a bitwise OR of replaced columns.
- 31 • Otherwise, we take l 's predecessor $l' \in L$ in the standard ordering and
32 replace lines $[l' + 1, l]$ by one line that is a bitwise OR of replaced lines.

33 **Notation 5.** For a matrix $M \in \{0, 1\}^{m \times n}$, $R \subseteq [m]$ and $C \subseteq [n]$ let $M_{\preceq}[R, C] :=$
34 $M_{\preceq}[R \cup \{c + m | c \in C\}]$.

35 **Definition 2.** We say a matrix $M \in \{0, 1\}^{m \times n}$ *contains* a pattern $P \in \{0, 1\}^{k \times l}$
36 *as an interval minor* and denote it by $P \preceq M$ if there are $R \in [m]$ and $C \in [n]$
37 such that $|R| = k$, $|C| = l$ and for every $r \in R$ and $c \in C$ if $P[r, c] = 1$, then
38 $M_{\preceq}[R, C][r, c] = 1$.

39 **Observation 1.** For all matrices M and P , $P \leq M \Rightarrow P \preceq M$.

40 **Observation 2.** *For all matrices M and P , if P is a permutation matrix, then*
 41 *$P \leq M \Leftrightarrow P \preceq M$.*

42 *Proof.* If we have $P \preceq M$, then there is a partitioning of M into rectangles and for
 43 each one-entry of P there is at least one one-entry in the corresponding rectangle
 44 of M . Since P is a permutation matrix, it is sufficient to take rows and columns
 45 having at least one one-entry in the right rectangle and we can always do so.

46 Together with Observation 1 this gives us the statement. \square

0.1 Characterizations

Definition 3. A *walk* in a matrix M is a sequence of some of its entries, beginning in the top left corner and ending in the bottom right one. If an entry $M[i, j]$ is in the sequence, the next one is either $M[i + 1, j]$ or $M[i, j + 1]$.

Definition 4. We call a binary matrix M a *walking matrix* if there is a walk in M such that all one-entries of M are contained on the walk.

0.1.1 Patterns of size 2×2

Theorem 3. Let $P = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$, then for all M : $P \not\preceq M \Leftrightarrow M$ is a walking matrix.

Proof. Since P is a permutation matrix, $P \not\preceq M \Leftrightarrow P \not\preceq M$ and it is easy to see $P \not\preceq M \Leftrightarrow M$ is a walking matrix. \square

Theorem 4. Let $P = \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$, then for all $M \in \{0, 1\}^{m \times n}$: $P \not\preceq M \Leftrightarrow$ there exist a row r and a column c such that (see Figure 1)

- $M[[r - 1], [c - 1]]$ is empty,
- $M[[r - 1], [c + 1, n]]$ is empty,
- $M[[r + 1, m], [c - 1]]$ is empty and
- $M[[r, m], [c, n]]$ is a walking matrix.

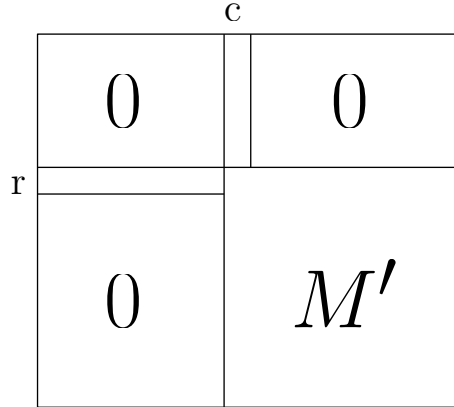


Figure 1: Characterization of a matrix avoiding $\begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$ as an interval minor.

Proof. \Rightarrow If $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \not\preceq M$, then M is a walking matrix and we set $r = c = 1$. Otherwise, there are one-entries $M[r, c']$ and $M[r', c]$ such that $r' < r$ and $c' < c$. If there is a one-entry in regions $M[[r - 1], [c - 1]]$, $M[[r - 1], [c + 1, n]]$ or $M[[r + 1, m], [c - 1]]$ then $P \preceq M$. If $M[[r, m], [c, n]]$ is not a walking matrix then it contains $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$ and we again get a contradiction.

\Leftarrow For contradiction, assume that M described in Figure 1 contains P as an interval minor. It means that there is a partition of the matrix into four quadrants such that there is at least one one-entry in each quadrant besides the bottom right one. If the matrix is partitioned above the r -th row, then

72 there is only one column containing one-entries and it is not possible for both
 73 top quadrants to have a one-entry. Similarly, if the matrix is partitioned to
 74 the left of the c -th column, there is only one row containing one-entries and
 75 there is no one-entry in either top-left or bottom-left quadrant. Therefore,
 76 the partitioning lies below the r -th row and to the right of the c -th column,
 77 but if the quadrants contain one-entries, there is a $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$ interval minor in
 78 M' , which is a contradiction with it being a walking matrix. \square
 79

80 To characterize matrices avoiding $\begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$ as an interval minor, we first need to
 81 define a few useful terms.

82 **Definition 5.** For $M \in \{0, 1\}^{m \times n}$ and $r \in [m], c \in [n]$ we say $M[r, c]$ is

- 83 • *top-left empty* if $M[[r - 1], [c - 1]]$ is an empty matrix,
- 84 • *top-right empty* if $M[[r - 1], [c + 1, n]]$ is empty,
- 85 • *bottom-left empty* if $M[[r + 1, m], [c - 1]]$ is empty,
- 86 • *bottom-right empty* if $M[[r + 1, m], [c + 1, n]]$ is empty.

87 **Lemma 5.** Let $P = \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$ and let $M \in \{0, 1\}^{m \times n}$ avoid P as an interval minor,
 88 then there exists a row r and a column c such that $M[r, c]$ is either

- 89 1. both top-left empty and bottom-right empty and $[r, c] \notin \{[1, n], [m, 1]\}$ or
- 90 2. both top-right empty and bottom-left empty and $[r, c] \notin \{[1, 1], [m, n]\}$.

91 *Proof.* \square

92 **Theorem 6.** Let $P = \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$, then for all M : $P \not\leq M \Leftrightarrow M$ looks like one of the
 93 matrices in Figure 2, where $\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \not\leq M_1$, $\begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix} \not\leq M_2$, $\begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix} \not\leq M_3$ and $\begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix} \not\leq M_4$.

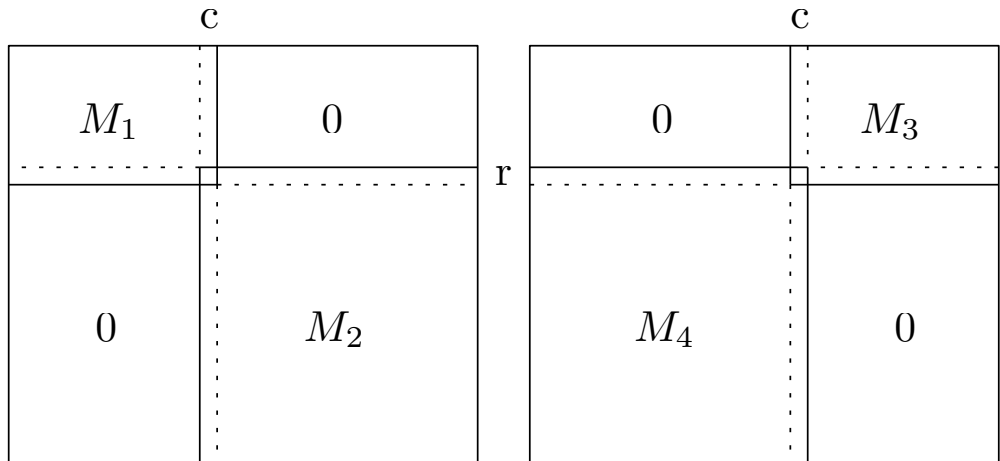


Figure 2: Characterization of a matrix avoiding $\begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$ as an interval minor.

94 *Proof.*

95 \Rightarrow We proceed by induction by the size of M .

96 If $M \in \{0, 1\}^{2 \times 2}$ then it either avoids $\begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix}$ or $\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}$ and we are done.

97 For bigger M there is, from Lemma 5, “the element”. Assume the first case
 98 (top-right and bottom-left empty (will change this when I have some notation)).
 99 If M_1 is non-empty, then $\begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix} \not\preceq M_2$; otherwise, $P \preceq M$. Similarly, $\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \not\preceq M_1$ if
 100 M_2 is non-empty. If one of them is empty, the other is a smaller matrix avoiding
 101 P as an interval minor and by induction hypothesis, it can be partitioned. Adding
 102 empty rows and columns does not break any condition and we get a partitioning
 103 of the whole M .

104 \Leftarrow Without loss of generality, let us assume M looks like the left matrix in Figure 2.
 105 For contradiction, assume $P \preceq M$. In that case, we can partition M into four
 106 quadrants such that there is at least one one-entry in each of them. It does not
 107 matter where we partition it, every time we either get $\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \preceq M_1$ or $\begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix} \preceq M_2$,
 108 which is a contradiction. \square

109 0.1.2 Matrices of size 2×3

110 **Theorem 7.** Let $P = \begin{pmatrix} 1 & 0 & 1 \\ 0 & 1 & 0 \end{pmatrix}$, then for all M : $P \not\preceq M \Leftrightarrow M = M_1 \oplus_h M_2$ where
 111 $\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \not\preceq M_1$ and $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \not\preceq M_2$.

112 *Proof.* \Rightarrow Let $e = [r, c]$ be the top-most one-entry of M . If $\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \preceq M[[m], [c -$
 113 $1]]$, together with e it forms P . If $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \not\preceq M[[m], [c, n]]$ then we are done.
 114 Let us assume it is not the case and let $e_{0,0}, e_{1,1}$ be any two one-entries
 115 forming the forbidden pattern. Symmetrically, let $\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \preceq M[[m], [c]]$ and
 116 let $e_{0,1}, e_{1,0}$ be any two one-entries forming the forbidden pattern. Now
 117 if we take $e_{0,0}, e_{0,1}$ and $e_{1,0}$ or $e_{1,1}$ with bigger row, we get the forbidden
 118 pattern P as an interval minor of M .

119 \Leftarrow For contradiction, let us assume $P \preceq M$ and $M = M_1 \oplus_h M_2$. If $P \preceq M$,
 120 look at the one-entry of M where the bottom one-entry of P is mapped. If
 121 it is in M_1 then $P \not\preceq M$ because $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \not\preceq M_1$. Otherwise, $P \not\preceq M$ because
 122 $\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \not\preceq M_2$. \square

124 **Lemma 8.** Let $P = \begin{pmatrix} 1 & 1 & 1 \\ 0 & 1 & 0 \end{pmatrix}$, then for all M : $P \not\preceq M \Rightarrow M = M_1 \oplus_h M_2$ where

125 1. $\begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix} \not\preceq M_1$ and $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \not\preceq M_2$ or

126 2. $\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \not\preceq M_1$ and $\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \not\preceq M_2$.

127 *Proof.* Let $e = [r, c]$ be the top-most one-entry of M . If $\begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix} \preceq M[[m], [c - 1]]$,
 128 together with e it would be the whole P . Similarly, $\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \not\preceq M[[m], [c + 1, n]]$.
 129 For contradiction with the statement, let $\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \preceq M[[m], [c]]$ and $e_{0,0}, e_{1,1}$ (none
 130 of them equal to e , since e lies in the top-right corner) be any two one-entries
 131 forming the pattern. Symmetrically, let $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \preceq M[[m], [c, n]]$ and $e_{0,1}, e_{1,0}$ be
 132 any two one-entries forming the pattern. In that case $e_{0,0}, e, e_{0,1}$ and $e_{1,0}$ or $e_{1,1}$
 133 with bigger row give us the forbidden pattern P as an interval minor of M . \square

134 **Theorem 9.** Let $P = \begin{pmatrix} 1 & 1 & 1 \\ 0 & 1 & 0 \end{pmatrix}$, then for all M : $P \not\preceq M \Leftrightarrow M$ looks like the matrix
 135 in Figure 3 and $\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \not\preceq M_1$ and $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \not\preceq M_2$.

0.2 Extremal function

Notation 6. Let M be a matrix. We denote $|M|$ the weight of M , the number of one-entries in M .

Usually $|M|$ stands for a determinant of matrix M . However, in this paper we do not work with determinants at all so the notation should not lead to misunderstanding.

Definition 6. For a matrix P we define $Ex(P, m, n) := \max\{|M| \mid M \in \{0, 1\}^{m \times n}, P \not\leq M\}$. We denote $Ex(P, n) := Ex(P, n, n)$.

Definition 7. For a matrix P we define $Ex_{\preceq}(P, m, n) := \max\{|M| \mid M \in \{0, 1\}^{m \times n}, P \not\preceq M\}$. We denote $Ex_{\preceq}(P, n) := Ex_{\preceq}(P, n, n)$.

Observation 10. For all P, m, n ; $Ex_{\preceq}(P, m, n) \leq Ex(P, m, n)$.

Observation 11. If $P \in \{0, 1\}^{k \times l}$ has a one-entry at position $[a, b]$, then

$$Ex(P, m, n) \geq \begin{cases} m \cdot n & k > m \vee l > m \\ (k-1)n + (l-1)m - (k-1)(l-1) & \text{otherwise.} \end{cases}$$

Observation 12. The same holds for $Ex_{\preceq}(P, m, n)$.

Definition 8. $P \in \{0, 1\}^{k \times l}$ is (strongly) minimalist if

$$Ex(P, m, n) = \begin{cases} m \cdot n & k > m \vee l > m \\ (k-1)n + (l-1)m - (k-1)(l-1) & \text{otherwise.} \end{cases}$$

Definition 9. $P \in \{0, 1\}^{k \times l}$ is weakly minimalist if

$$Ex_{\preceq}(P, m, n) = \begin{cases} m \cdot n & k > m \vee l > m \\ (k-1)n + (l-1)m - (k-1)(l-1) & \text{otherwise.} \end{cases}$$

Observation 13. If P is strongly minimalist, then P is weakly minimalist.

0.2.1 Known results

Fact 14. 1. $\begin{pmatrix} 1 \end{pmatrix}$ is strongly minimalist.

2. If $P \in \{0, 1\}^{k \times l}$ is strongly minimalist and there is a one-entry in the last row in the c -th column, then $P' \in \{0, 1\}^{k+1 \times l}$, which is created from P by adding a new row having a one-entry only in the c -th column, is strongly minimalist.

3. If P is strongly minimalist, then after changing a one-entry into a zero-entry it is still strongly minimalist.

Fact 15. Let $P = \begin{pmatrix} 1 & \dots & 1 \\ 1 & \dots & 1 \end{pmatrix}$ have l columns, then P is weakly minimalist.

Proof. Let $M \in \{0, 1\}^{m \times n}$ be a matrix avoiding $P = \{1\}^{2 \times l}$ as an interval minor and $A_i = \{j \in [n] \mid \text{weight of } M[[i, \{j\}] > 0 \wedge \text{weight of } M[[i + 1, m], \{j\}] > 0\}$. Clearly $|A_i| \leq l - 1$, otherwise $P \preceq M$. Let b_j denote the number of one-entries in the j -th column. Each column j of M appears in at least $b_j - 1$ of sets A_i , $0 \leq i \leq m - 2$. It follows that

$$\text{weight of } M = \sum_{j=0}^n b_j = \sum_{j=0}^n (b_j - 1) + n \leq \sum_{i=0}^{m-2} |A_i| + n \leq (l - 1)(m - 1) + n$$

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□

171 This result is indeed very important because it shows that there are matrices
172 like $\begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$, which are weakly minimalist, although it is known they are not strongly
173 minimalist.

174 **Fact 16.** Let $P = \begin{pmatrix} 1 & \dots & 1 \\ 1 & \dots & 1 \end{pmatrix}$ have l columns, then P is weakly minimalist.

Proof. Let $M \in \{0, 1\}^{m \times n}$ be a matrix avoiding $P = \{1\}^{3 \times l}$ as an interval minor and $A_i = \{j \in [n] \mid \text{weight of } M[[i - 1, \{j\}] > 0 \wedge \text{weight of } M[[i + 1, m], \{j\}] > 0 \wedge M[i, j] \text{ one-entry}\}$. Clearly $|A_i| \leq l - 1$, otherwise $P \preceq M$. Let b_j denote the number of one-entries in the j -th column. Each column j of M (for which $b_j \geq 2$) appears in exactly $b_j - 2$ of sets A_i , $1 \leq i \leq m - 1$. It follows that

$$\text{weight of } M = \sum_{j=0}^n b_j = \sum_{j=0}^n (b_j - 2) + 2n \leq \sum_{i=1}^{m-2} |A_i| + 2n \leq (l - 1)(m - 2) + 2n$$

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□