

Numerical Methods

Last time: def'n of normed vector space $(V, \|\cdot\|)$

and of convergence $x_n \rightarrow x \iff \|x_n - x\| \rightarrow 0$

Def'n: • A sequence x_n tending to x converges linearly if $\exists N \in \mathbb{N}$ and $0 < C < 1$ s.t. $\|x_{n+1} - x\| \leq C \|x_n - x\|$, $n \geq N$.

(In this case, $\exists \tilde{C} > 0$ s.t. $\|x_n - x\| \leq \tilde{C} C^n$, $n \in \mathbb{N}$.)

exercise: find \tilde{C}

"a priori error bound"

Note: $0 < C < 1 \Rightarrow C^n \rightarrow 0$ as $n \rightarrow \infty$

Exercise: explain why we don't need $C < 1$ here

• x_n converges quadratically to x if $\exists N \in \mathbb{N}$ and $C > 0$ s.t. $\|x_{n+1} - x\| \leq C \|x_n - x\|^2$, $n \geq N$.

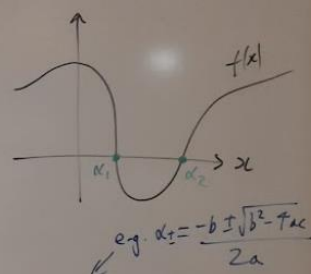
• x_n converges with order $q > 1$ to x if $\exists N \in \mathbb{N}$ and $C > 0$ s.t. $\|x_{n+1} - x\| \leq C \|x_n - x\|^q$, $n \geq N$.

not necessarily integer

§ 3: Nonlinear equations (root finding)

Given $f: \mathbb{R} \rightarrow \mathbb{R}$, find $\alpha \in \mathbb{R}$ s.t. $f(\alpha) = 0$.

(Can rephrase any equation $F(x) = G(x)$ in this way by setting $f(x) := F(x) - G(x)$.)



For simple functions (e.g. low order polynomials) we might have a closed form solution. But in general we need a numerical method to generate an approximate solution.

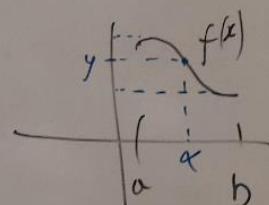
Existence of roots?

Recall Intermediate Value Theorem:

If $f: [a, b] \rightarrow \mathbb{R}$ is continuous then f is bounded and attains its bounds, and for every $y \in \mathbb{R}$ s.t.

$$\min_{x \in [a, b]} f(x) \leq y \leq \max_{x \in [a, b]} f(x)$$

$$\exists \alpha \in [a, b] \text{ s.t. } f(\alpha) = y$$



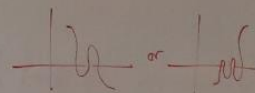
Corollary: If $f: [a, b] \rightarrow \mathbb{R}$ is continuous and $f(a)f(b) \leq 0$ then $\exists \alpha \in [a, b]$ s.t. $f(\alpha) = 0$.

Proof - see notes.

\leq here means $f(a)$ and $f(b)$ have opposite signs.

Note: this gives existence but not uniqueness.

There might be more than one root in $[a, b]$.



This leads naturally to a numerical method:

Bisection method: Assume $f: [a, b] \rightarrow \mathbb{R}$ is continuous and $f(a)f(b) \leq 0$.

- Set $k=0$
- If $f(a)f(b)=0$ then either $f(a)=0$ or $f(b)=0$. STOP.

If not:

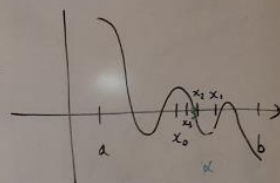
- Set $x_k = \frac{a+b}{2}$

- If $f(x_k)=0$, STOP

If not,

- if $f(a)f(x_k) < 0$, redefine $b = x_k$ and go to here.
- if $f(a)f(x_k) > 0$, redefine $a = x_k$ and go to here.

$\leftarrow \exists \text{ root in } [x_k, b]$

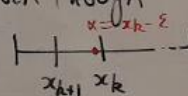


Theorem: If $f: [a, b] \rightarrow \mathbb{R}$ is continuous and $f(a)f(b) \leq 0$ then the bisection method converges to some $\alpha \in [a, b]$ with $f(\alpha) = 0$, and

$$(*) \quad |x_k - \alpha| \leq \frac{b-a}{2^{k+1}} = \underbrace{\frac{b-a}{2}}_{\varepsilon} \cdot \underbrace{\left(\frac{1}{2}\right)^k}_C$$

Proof: Straightforward from def'n of algorithm.

Note: Bisection method is not linearly convergent even though $(*)$ holds. $\left(\frac{\|x_{n+1} - \alpha\|}{\|x_n - \alpha\|} \text{ could be arbitrarily large!}\right)$

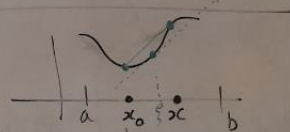


Bisection is nice as it only requires continuity of f .

But it's quite slow.

We can find faster methods, but in general they require more smoothness of f .

§ 2.4



Mean value theorem: Let $f: (a, b) \rightarrow \mathbb{R}$ be differentiable, then for $x, x_0 \in (a, b)$ with $x \neq x_0$, $\exists \xi$ between x and x_0 s.t.

$$f(x) = f(x_0) + f'(\xi)(x - x_0).$$

Taylor's Theorem: Let $f: (a, b) \rightarrow \mathbb{R}$ be k -times differentiable on (a, b) for some $k \in \mathbb{N}$. Then for $x, x_0 \in (a, b)$ with $x \neq x_0$,

$$f(x) = \sum_{j=0}^k \frac{f^{(j)}(x_0)}{j!} (x - x_0)^j + R_k(x; x_0), \quad \text{"asymptotic notation"}$$

where $R_k(x; x_0) = o(|x - x_0|^k)$ as $x \rightarrow x_0$. $\left(F(x) = o(g(x)) \text{ as } x \rightarrow x_0 \text{ means } \frac{F(x)}{g(x)} \rightarrow 0 \text{ as } x \rightarrow x_0, \text{ assuming } g(x) \neq 0 \right)$

If f is $(k+1)$ -times differentiable on (a, b) then

$$\exists \xi \text{ between } x_0 \text{ and } x \text{ s.t. } R_k(x; x_0) = \frac{f^{(k+1)}(\xi)}{(k+1)!} (x - x_0)^{k+1}.$$

Further, if $|f^{(k+1)}(x)| \leq M$ $\forall x \in (a, b)$ then

$$|R_k(x; x_0)| \leq \frac{M}{(k+1)!} |x - x_0|^{k+1} = O(|x - x_0|^{k+1}) \leftarrow \left(F(x) = O(g(x)) \text{ as } x \rightarrow x_0 \text{ means } \exists C > 0 \text{ s.t. } |F(x)| \leq C|g(x)| \text{ for } |x - x_0| \text{ suff. small} \right)$$

more generally, see notes.

9-3 Fixed point methods:

Suppose we can find $\phi: [a, b] \rightarrow \mathbb{R}$ s.t.

$$\phi(x) = x \iff f(x) = 0.$$

Then we say $\phi(x) = x$ is a fixed point reformulation of $f(x) = 0$.

Simple example that always holds: $\phi(x) = x - \beta f(x)$
for any $0 \neq \beta \in \mathbb{R}$.

But other choices are often possible.

Example: $f(x) = \sin 2x - 1 + x = 0$

Two choices are:

- $\phi_1(x) = x$ where $\phi_1(x) = 1 - \sin 2x$.
- $\phi_2(x) = x$ where $\phi_2(x) = \frac{1}{2} \arcsin(1-x)$ on $[0, 2]$.

we know roots of f lie in this interval since $f(x) = 0$

$$\Leftrightarrow 1-x = \sin 2x$$

$$\Leftrightarrow x = \frac{1 - \sin 2x}{2} \in [0, 2]$$

continuous

Existence of fixed points?

Brouwer's fixed point theorem: Suppose $\phi: [a, b] \rightarrow \mathbb{R}$ is continuous

and that $\phi(x) \in [a, b]$ for all $x \in [a, b]$. Then there exists $\alpha \in [a, b]$

$$\text{s.t. } \boxed{\phi(\alpha) = \alpha}$$

(Note: α may not be unique!)

Proof: Let $f(x) = x - \phi(x)$. Then since $\phi: [a, b] \rightarrow [a, b]$,
 $f(a) = a - \phi(a) \leq 0$ and $f(b) = b - \phi(b) \geq 0$, so $f(a)f(b) \leq 0$,
and by previous corollary $\exists \alpha \in [a, b]$ s.t. $f(\alpha) = 0$, i.e. $\phi(\alpha) = \alpha$. \square

To approximate fixed points we can try the fixed point iteration:

$$\boxed{x_{k+1} = \phi(x_k)}, \quad k \in \mathbb{N}_0, \quad \text{given some initial guess } x_0 \in [a, b].$$

$$\text{i.e. } x_1 = \phi(x_0), \quad x_2 = \phi(x_1) = \phi(\phi(x_0)), \text{ etc.}$$

This produces a sequence x_0, x_1, x_2, \dots which we hope converges to a fixed point of ϕ .