#### Vitual Memory:

### 1. 作用:

Use main memory as a "cache" for secondary (disk) storage Programs share main memory

2. 概念: page page fault

由于page Fault Penalty太大,希望尽可能少miss,因此采取Fully associated

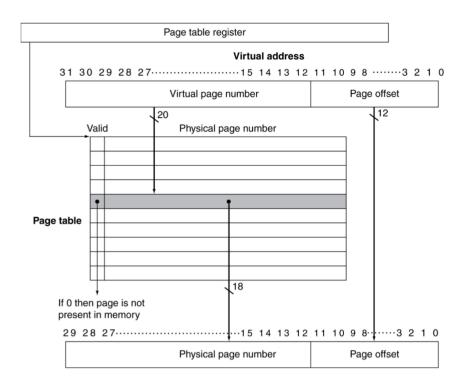
- Try to minimize page fault rate
  - Fully associative placement
  - Smart replacement algorithms
- 1. Vitual address如何与实际地址对应: Page Table

# Page Tables

- Where is the placement information? Page Table
  - Array of page table entries (PTE), indexed by virtual page number
  - Page table register in CPU points to page table in physical memory
- Each program has its page table. Page table is in memory
- If page is present in memory
  - PTE stores the physical page number
  - Plus other status bits (referenced, dirty, ...)
- If page is not present
  - PTE can refer to location in swap space on disk

Page table储存Vpn到ppn的映射

Page table储存在memory内,是一块memory



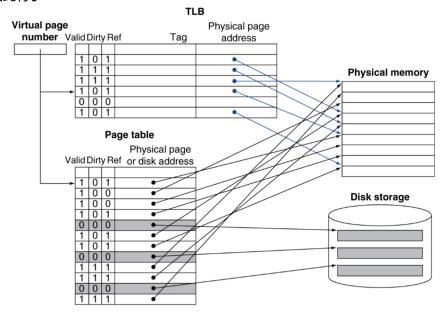
Valid 为0 代表此vitual address对应内容在disk中 还有dirty referrence等其他status bit Page table register 不同程序有不同的page table

## 2. Replacement and write

Use bit/referrence bit—replacement—LRU appromiate Dirty bit—write back

## 3. TLB

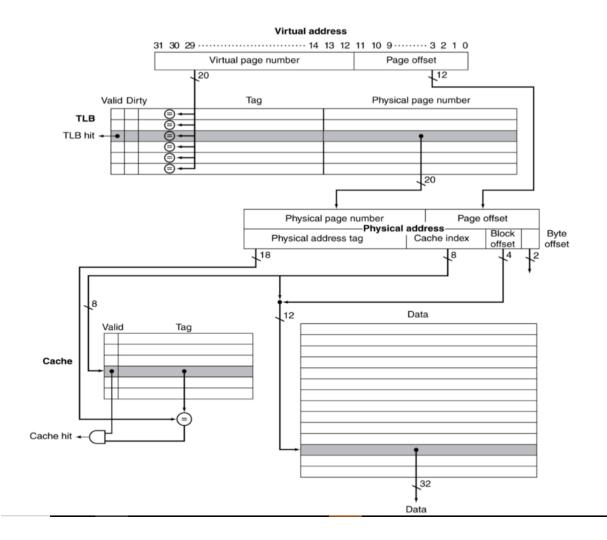
通过vitual address去Access memory需要两次memory access 因此有了TLB



#### TLB miss:

- If page is in memory
  - Load the PTE from memory and retry
  - Could be handled in hardware
    - Can get complex for more complicated page table structures
  - Or in software
    - Raise a special exception, with optimized handler
- If page is not in memory (page fault)
  - OS handles fetching the page and updating the page table
  - Then restart the faulting instruction

#### TLB interaction with cache:



# 4. Memory Protect

不同程序可以share相同的虚拟内存,但需要OS协助保证不发生错误访问 OS特权模式