

HomeWork_3

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chapter 2

Exercise 2.19

$\$t0 = 0xAAAAAAAA$, $\$t1 = 0x12345678$

2.19.1

The first statement "sll $\$t2$, $\$t0$, 4", the value in $\$t2$ is $\$t0$ shifted left by 4 bits, 1 bytes.

$\$t2 = 0xAAAAAAAA \ll 4 = 0xAAAAAAAA0$

After first statement, $\$t2 = 0xAAAAAAAA0$

The second statement "or $\$t2$, $\$t2$, $\$t1$ ", compute $\$t2 = 0xAAAAAAAA0$ or $\$t1 = 0x12345678$, the answer is :

$0xAAAAAAAA0 \mid 0x12345678$

$= 10101010101010101010101010100000 \mid 00010010001101000101011001111000$

$= 101110101011111011111011111000$

$= 0xBABEFEF8$

After second statement, $\$t2 = 0xBABEFEF8$

2.19.2

The first statement "sll $\$t2$, $\$t0$, 4", the value in $\$t2$ is $\$t0$ shifted left by 4 bits, 1 bytes.

$\$t2 = 0xAAAAAAAA \ll 4 = 0xAAAAAAAA0$

After first statement, \$t2 = 0xAAAAAAAA0

The second statement "andi \$t2, \$t2, -1" ,compute \$t0=0xAAAAAAAA0 and 0xFFFFFFFF(-1), the answer is :

0xAAAAAAAA0 | 0xFFFFFFFF
= 10101010101010101010101010100000 & 11111111111111111111111111111111
= 10101010101010101010101010100000
= 0xAAAAAAAA0

After second statement, \$t2 = 0xAAAAAAAA0

2.19.3

The first statement "srl \$t2 \$t0 3" , the value in \$t2 is \$t0 shifted right by 3 bits,1bytes.

\$t2 = 0xAAAAAAAA >> 3 = 0x15555555

After first statement, \$t2 = 0x15555555

The second statement "andi \$t2, \$t2, 0xFFEF" ,compute \$t0=0x15555555 and 0xFFEF, the answer is :

0x15555555 | 0xFFEF
= 00010101010101010101010101010101 & 0000000000000000111111111101111
= 00000000000000000101010101000101
= 0x00005545

After second statement, \$t2 = 0x00005545

Exercise 2.26

```
1 题面 :  
2 LOOP: slt $t2, $0, $t1  
3 beq $t2, $0, DONE  
4 subi $t1, $t1, 1  
5 addi $s2, $s2, 2  
6 j LOOP  
7 DONE:
```

经分析，该代码的含义为若\$t1<0, 则结束循环，否则，每次循环执行两次操作，\$t1 = \$t1 -1, \$s2 = \$s2 + 2直到\$t1<0.

2.26.1

\$t1 = 10 \$s2 = 0

$$s2 = 0 + 2 \times (10 - 0) = 20$$

2.26.2

```

1  i = 10;
2  while(i > 0){
3      i = i - 1;
4      B += 2;
5  }
```

2.26.3

By analysis, because \$t1 is initialized to the value N, we have N+1 loops totally.

The first N loop has 5 instructions, and final one has 2 instructions

So totally 5N+2 instructions.

Exercise 2.31

To achieve the goal to computing the fibonacci number, we use recursion code by mips as follows:

```

1  fib: addi $sp, $sp, -12           # make room on stack
2  sw $ra, 8($sp)                  # push $ra
3  sw $s0, 4($sp)                  # push $s0
4  sw $a0, 0($sp)                  # push $a0 (N)
5  bgt $a0, $zero, equ1           # if n>0, turn to <equ1> to test (n==1?)
6  li $v0, 0                      # else fib(0) = 0
7  j return                        # turn to <return>
8
9  equ1: li $t0, 1
10 bne $a0, $t0, all              # if n>1, turn to <all>
11 li $v0, 1                      # else fib(1) = 1
12 j return                        # turn to <return>
13
14 all: subi $a0, $a0, 1           # n-1
15 jal fib                         # call fib(n-1)
16 addi $s0, $v0, 0               # add fib(n-1) to answer
17 subi $a0, $a0, 1              # n-2
18 jal fib                         # call fib(n-2)
19 add $v0, $v0, $s0              # add fib(n-2) to answer
20
21 return: lw $a0, 0($sp)          # pop $a0
22 lw $s0, 4($sp)                 # pop $s0
23 lw $ra, 8($sp)                 # pop $ra
24 addi $sp, $sp, 12              # restore sp
25 jr $ra                         # return to the position of instruction
ra
```

Instruction numbers :

fib(0) : 12 instructions

fib(1) : 14 instructions

fib(n) : $26 + 18n$ instructions ($n \geq 2$)