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1  |----- MODULE AJupiter -----|
   | Specification of the Jupiter protocol presented by Attiya and others.
6  |-----|
   | EXTENDS OT, Jupiter
7  |-----|
   | Messages between the Server and the Clients. There are two kinds of messages according to their
   | destinations.
12 Msg  $\triangleq$  [c : Client, ack : Int, op : Op  $\cup$  {Nop}]  $\cup$  messages sent to the Server from a client c  $\in$  Client
13               [ack : Int, op : Op  $\cup$  {Nop}] messages broadcast to Clients from the Server
14 |-----|
15 VARIABLES
   | For the client replicas:
19   cbuf,      cbuf[c]: buffer (of operations) at the client c  $\in$  Client
20   crec,      crec[c]: the number of new messages have been received by the client c  $\in$  Client
21               since the last time a message was sent
   | For the server replica:
25   sbuf,      sbuf[c]: buffer (of operations) at the Server, one per client c  $\in$  Client
26   srec,      srec[c]: the number of new messages have been ... , one per client c  $\in$  Client
   | For model checking:
30   chins     a set of chars to insert
31 |-----|
32 comm  $\triangleq$  INSTANCE CSComm WITH Msg  $\leftarrow$  Msg
33 |-----|
34 eVars  $\triangleq$  {chins}          variables for the environment
35 cVars  $\triangleq$  {cbuf, crec}    variables for the clients
36 ecVars  $\triangleq$  {eVars, cVars}  variables for the clients and the environment
37 sVars  $\triangleq$  {sbuf, srec}    variables for the server
38 commVars  $\triangleq$  {cincoming, sincoming} variables for communication
39 Vars  $\triangleq$  {eVars, cVars, sVars, commVars, state} all variables
40 |-----|
41 TypeOK  $\triangleq$ 
   | For the client replicas:
45    $\wedge$  cbuf  $\in$  [Client  $\rightarrow$  Seq(Op  $\cup$  {Nop})]
46    $\wedge$  crec  $\in$  [Client  $\rightarrow$  Int]
   | For the server replica:
50    $\wedge$  sbuf  $\in$  [Client  $\rightarrow$  Seq(Op  $\cup$  {Nop})]
51    $\wedge$  srec  $\in$  [Client  $\rightarrow$  Int]
   | For all replicas.
55    $\wedge$  state  $\in$  [Replica  $\rightarrow$  List]
   | For communication between the server and the clients:
59    $\wedge$  comm! TypeOK
   | For model checking:

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63 $\wedge chins \in \text{SUBSET } Char$

64 The *Init* predicate.

68 $Init \triangleq$

69 $\wedge chins = Char$

For the client replicas:

73 $\wedge cbuf = [c \in Client \mapsto \langle \rangle]$

74 $\wedge crec = [c \in Client \mapsto 0]$

For the server replica:

78 $\wedge sbuf = [c \in Client \mapsto \langle \rangle]$

79 $\wedge srec = [c \in Client \mapsto 0]$

For all replicas.

83 $\wedge state = [r \in Replica \mapsto InitState]$

For communication between the server and the clients:

87 $\wedge comm!Init$

88 Client $c \in Client$ issues an operation op .

92 $DoOp(c, op) \triangleq$

93 $\wedge state' = [state \text{ EXCEPT } ![c] = Apply(op, @)]$

94 $\wedge cbuf' = [cbuf \text{ EXCEPT } ![c] = Append(@, op)]$

95 $\wedge crec' = [crec \text{ EXCEPT } ![c] = 0]$

96 $\wedge comm!CSend([c \mapsto c, ack \mapsto crec[c], op \mapsto op])$

98 $DoIns(c) \triangleq$

99 $\exists ins \in Ins :$

100 $\wedge ins.pos \in 1 \dots (Len(state[c]) + 1)$

101 $\wedge ins.ch \in chins$

102 $\wedge ins.pr = Priority[c]$

103 $\wedge chins' = chins \setminus \{ins.ch\}$ We assume that all inserted elements are unique.

104 $\wedge DoOp(c, ins)$

105 $\wedge \text{UNCHANGED } sVars$

107 $DoDel(c) \triangleq$

108 $\exists del \in Del :$

109 $\wedge del.pos \in 1 \dots Len(state[c])$

110 $\wedge DoOp(c, del)$

111 $\wedge \text{UNCHANGED } \langle sVars, eVars \rangle$

113 $Do(c) \triangleq$

114 $\vee DoIns(c)$

115 $\vee DoDel(c)$

Client $c \in Client$ receives a message from the *Server*.

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120  $Rev(c) \triangleq$ 
121    $\wedge comm!CRev(c)$ 
122    $\wedge crec' = [crec \text{ EXCEPT } ![c] = @ + 1]$ 
123    $\wedge \text{LET } m \triangleq Head(cincoming[c])$ 
124      $cBuf \triangleq cbuf[c]$  the buffer at client  $c \in Client$ 
125      $cShiftedBuf \triangleq SubSeq(cBuf, m.ack + 1, Len(cBuf))$  buffer shifted
126      $xop \triangleq XformOpOps(m.op, cShiftedBuf)$  transform  $op$  vs. shifted buffer
127      $xcBuf \triangleq XformOpsOp(cShiftedBuf, m.op)$  transform shifted buffer vs.  $op$ 
128   IN    $\wedge cbuf' = [cbuf \text{ EXCEPT } ![c] = xcBuf]$ 
129        $\wedge state' = [state \text{ EXCEPT } ![c] = Apply(xop, @)]$  apply the transformed operation  $xop$ 
130    $\wedge \text{UNCHANGED } \langle sVars, eVars \rangle$ 
131 |-----|
132 | The Server receives a message. |
133 |-----|
134  $SRev \triangleq$ 
135    $\wedge comm!SRev$ 
136    $\wedge \text{LET } m \triangleq Head(sincoming)$  the message to handle with
137    $c \triangleq m.c$  the client  $c \in Client$  that sends this message
138    $cBuf \triangleq sbuf[c]$  the buffer at the Server for client  $c \in Client$ 
139    $cShiftedBuf \triangleq SubSeq(cBuf, m.ack + 1, Len(cBuf))$  buffer shifted
140    $xop \triangleq XformOpOps(m.op, cShiftedBuf)$  transform  $op$  vs. shifted buffer
141    $xcBuf \triangleq XformOpsOp(cShiftedBuf, m.op)$  transform shifted buffer vs.  $op$ 
142   IN    $\wedge srec' = [cl \in Client \mapsto$ 
143       IF  $cl = c$ 
144       THEN  $srec[cl] + 1$  receive one more operation from client  $c \in Client$ 
145       ELSE  $0]$  reset  $srec$  for other clients than  $c \in Client$ 
146    $\wedge sbuf' = [cl \in Client \mapsto$ 
147       IF  $cl = c$ 
148       THEN  $xcBuf$  transformed buffer for client  $c \in Client$ 
149       ELSE  $Append(sbuf[cl], xop)]$  store transformed  $xop$  into other clients' bufs
150    $\wedge state' = [state \text{ EXCEPT } ![Server] = Apply(xop, @)]$  apply the transformed operation
151    $\wedge comm!SSend(c, [cl \in Client \mapsto [ack \mapsto srec[cl], op \mapsto xop]])$ 
152    $\wedge \text{UNCHANGED } ecVars$ 
153 |-----|
154 | The safety properties to check: Eventual Convergence ( $EC$ ), Quiescent Consistency ( $QC$ ), Strong |
155 | Eventual Convergence ( $SEC$ ), Weak List Specification, ( $WLSpec$ ), and Strong List Specification, |
156 | ( $SLSpec$ ). |
157 |-----|
158 | Eventual Consistency ( $EC$ ) |
159 |-----|
160 | Quiescent Consistency ( $QC$ ) |
161 |-----|
162  $QC \triangleq comm!EmptyChannel \Rightarrow Cardinality(Range(state)) = 1$ 
163 |-----|
164 | Strong Eventual Consistency ( $SEC$ ) |
165 |-----|
166 | Instance  $JupiterH$  |
167 |-----|
168 | Strong Eventual Consistency ( $SEC$ ) |
169 |-----|
170 |

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* Modification History
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