Εργαστήριο Τεχνητή Νοημοσύνη ΙΙ

Παύλος Πέππας

Τμήμα Ηλεκτρολόγων Μηχανικών και Τεχνολογίας Υπολογιστών

Ισχυρή Άρνηση σε ASP

Μπορούμε να υλοποιήσουμε την κλασική (ισχυρή) άρνηση σε ASP (μόνο) για ατομικούς τύπους.

П.χ.

a.

c :- a, -b.

Ισχυρή Άρνηση σε ASP

Μπορούμε να υλοποιήσουμε την κλασική (ισχυρή) άρνηση σε ASP (μόνο) για ατομικούς τύπους.

Π.χ.

a.

c :- a, -b.

είναι ουσιαστικά ισοδύναμο με το

a.

c :- a, -b.

- a :- a, -a. a :- b, -b. a :- c, -c.

- -a :- a, -a. -a :- b, -b. -a :- c, -c.
- b :- a, -a. b :- b, -b. b :- c, -c.

- -b :- a, -a. -b :- b, -b. -b :- c, -c.
- c :- a, -a. c :- b, -b. c :- c, -c.

- -c :- a, -a. -c :- b, -b. -c :- c, -c.

$$\Pi_1 = \begin{cases} bird(tweety). \\ flies(X) :- bird(X), -penguin(X). \end{cases} \qquad \Pi_2 = \begin{cases} bird(tweety). \\ flies(X) :- bird(X), not penguin(X). \end{cases}$$

$$\Pi_1 = \begin{cases} bird(tweety). \\ flies(X) :- bird(X), -penguin(X). \end{cases}$$

$$\Pi_2 = \begin{cases} bird(tweety). \\ flies(X) :- bird(X), not penguin(X). \end{cases}$$

M = { bird(tweety), flies(tweety) }

$$\Pi_1 = \begin{cases} bird(tweety). \\ flies(X) :- bird(X), -penguin(X). \end{cases} \qquad \Pi_2 = \begin{cases} bird(tweety). \\ flies(X) :- bird(X), not penguin(X). \end{cases}$$

$$M = \{ bird(tweety), flies(tweety) \}$$

$$M \not\models_{ASP} \Pi_1$$

$$\Pi_1 = \begin{cases} bird(tweety). \\ flies(X) :- bird(X), -penguin(X). \end{cases} \qquad \Pi_2 = \begin{cases} bird(tweety). \\ flies(X) :- bird(X), not penguin(X). \end{cases}$$

$$M = \{ bird(tweety), flies(tweety) \}$$

$$M \not\models_{ASP} \Pi_1$$

$$\Pi_2^{M} = \begin{cases}
bird(tweety). \\
flies(X) :- bird(X).
\end{cases}$$

$$\Pi_1 = \begin{cases}
bird(tweety). \\
flies(X) :- bird(X), -penguin(X).
\end{cases}$$

$$\Pi_2 = \begin{cases}
bird(tweety). \\
flies(X) :- bird(X), not penguin(X).
\end{cases}$$

$$M \nvDash_{ASP} \Pi_1$$

$$M \vDash_{ASP} \Pi_2$$

$$\Pi_2^M = \begin{cases} bird(tweety). \\ flies(X) :- bird(X). \end{cases}$$

```
bird(zippy).
bird(tweety).
-flies(zippy).
flies(X) :- bird(X), not -flies(X).
```

```
M = { bird(zippy), -flies(zippy), bird(tweety), flies(tweety) }
bird(zippy).
bird(tweety).
-flies(zippy).
flies(tweety) :- bird(tweety).
```

```
M = { bird(zippy), -flies(zippy), bird(tweety), flies(tweety) }
\Pi = \begin{cases} \text{bird(zippy).} \\ \text{bird(tweety).} \\ \text{-flies(zippy).} \\ \text{flies(X) :- bird(X), not -flies(X).} \end{cases} \qquad \Pi^{M} = \begin{cases} \text{bird(zippy).} \\ \text{-flies(zippy).} \\ \text{flies(tweety) :- bird(tweety).} \end{cases}
```

$$\Pi = \begin{cases} q(a). \\ r(a). \\ p(X) := q(X), \text{ not } -p(X). \\ -p(X) := r(X), \text{ not } p(X). \end{cases}$$

$$\Pi = \begin{cases} \text{bird(zippy).} \\ \text{bird(tweety).} \\ \text{-flies(zippy).} \\ \text{flies(X) :- bird(X), not -flies(X).} \end{cases} \qquad \Pi^{M} = \begin{cases} \text{bird(zippy).} \\ \text{bird(tweety).} \\ \text{-flies(zippy).} \\ \text{flies(tweety) :- bird(tweety).} \end{cases}$$

$$\Pi = \begin{cases} q(a). \\ r(a). \\ p(X) :- q(X), \text{ not } -p(X). \\ -p(X) :- r(X), \text{ not } p(X). \end{cases} \qquad \Pi^{M} = \begin{cases} q(a). \\ r(a). \\ p(a) :- q(a). \end{cases} \qquad \Pi^{M} = \begin{cases} q(a). \\ r(a). \\ -p(a) :- r(a). \end{cases}$$

 $M = \{ p(a), q(a), r(a) \}$

$$\Pi^{M} = \begin{cases}
q(a). \\
r(a). \\
-p(a) :- r(a).
\end{cases}$$

 $M = \{ -p(a), q(a), r(a) \}$

$$\Pi = \begin{cases} \text{bird(zippy).} \\ \text{bird(tweety).} \\ \text{-flies(zippy).} \\ \text{flies(X) :- bird(X), not -flies(X).} \end{cases} \qquad \Pi^{M} = \begin{cases} \text{bird(zippy).} \\ \text{bird(tweety).} \\ \text{-flies(zippy).} \\ \text{flies(tweety) :- bird(tweety).} \end{cases}$$

$$\Pi = \begin{cases} q(a). \\ r(a). \\ p(X) :- q(X), \text{ not } -p(X). \\ -p(X) :- r(X), \text{ not } p(X). \end{cases} \qquad \Pi^{M} = \begin{cases} q(a). \\ r(a). \\ p(a) :- q(a). \end{cases} \qquad \Pi^{M} = \begin{cases} q(a). \\ r(a). \\ -p(a) :- r(a). \end{cases}$$

 $M = \{ p(a), q(a), r(a) \}$

$$\Pi^{M} =
\begin{cases}
q(a). \\
r(a). \\
-p(a) :- r(a).
\end{cases}$$

 $M = \{ -p(a), q(a), r(a) \}$

- Εύπιστη (credulous) Προσέγγιση: Επιλέγουμε τυχαία ένα answer set.
- Δύσπιστη (skeptical) Προσέγγιση: Πιστεύουμε μόνο ότι ισχύει σε όλα τα answer sets.

```
penguin(zippy).
bird(X) :- penguin(X).
flies(X) :- bird(X), not -flies(X).
-flies(X) :- penguin(X), not flies(X).
```

```
penguin(zippy).
bird(X) :- penguin(X).
flies(X) :- bird(X), not -flies(X).
-flies(X) :- penguin(X), not flies(X).
```

```
M1 = { bird(zippy), penguin(zippy), flies(zippy) } X

M2 = { bird(zippy), penguin(zippy), -flies(zippy) }
```

$$\Pi = \begin{cases} \text{penguin(zippy).} \\ \text{bird(X) :- penguin(X).} \\ \text{flies(X) :- bird(X), not -flies(X).} \\ \text{-flies(X) :- penguin(X), not flies(X).} \end{cases}$$

Answer Sets:

```
M1 = { bird(zippy), penguin(zippy), flies(zippy) } 
M2 = { bird(zippy), penguin(zippy), -flies(zippy) }
```

```
\Pi = \begin{cases} \text{penguin(zippy).} \\ \text{bird(X) :- penguin(X).} \\ \text{flies(X) :- bird(X), -penguin(X), not -flies(X).} \\ \text{-flies(X) :- penguin(X), not flies(X).} \end{cases}
```

Answer Sets:

M2 = { bird(zippy), penguin(zippy), -flies(zippy) }

```
M1 = { bird(zippy), penguin(zippy), flies(zippy) }
M2 = { bird(zippy), penguin(zippy), -flies(zippy) }
```

```
\Pi = \begin{cases} \text{penguin(zippy).} \\ \text{bird(X) :- penguin(X).} \\ \text{flies(X) :- bird(X), -penguin(X), not -flies(X).} \\ \text{-flies(X) :- penguin(X), not flies(X).} \\ \text{hird(tweety).} \end{cases}
M = \{ \text{bird(zippy), bird(tweety), penguin(zippy), -flies(zippy) } \}
```

```
\Pi = \begin{cases} \text{penguin(zippy).} \\ \text{bird(X) :-penguin(X).} \\ \text{flies(X) :- bird(X), not -flies(X).} \\ \text{-flies(X) :- penguin(X), not flies(X).} \end{cases}
```

Answer Sets:

```
M1 = { bird(zippy), penguin(zippy), flies(zippy) }
M2 = { bird(zippy), penguin(zippy), -flies(zippy) }
```

```
penguin(zippy).
bird(X) :- penguin(X).
flies(X) :- bird(X), not abBird(X).
-flies(X) :- penguin(X), not abPeng(X).
bird(tweety).

abPeng(X) :- penguin(X), flies(X).
abBird(X) :- penguin(X), not abPeng(X).
```

```
M = { bird(zippy), penguin(zippy),
      bird(Tweety),
      abBird(zippy),
      -flies(zippy), flies(Tweety) }
```

```
\Pi = \begin{cases} \text{penguin(zippy).} \\ \text{bird(X) :-penguin(X).} \\ \text{flies(X) :- bird(X), not -flies(X).} \\ \text{-flies(X) :- penguin(X), not flies(X).} \end{cases}
```

Answer Sets:

```
M1 = { bird(zippy), penguin(zippy), flies(zippy) }
M2 = { bird(zippy), penguin(zippy), -flies(zippy) }
```

```
penguin(zippy).
bird(X) :- penguin(X).
flies(X) :- bird(X), not abBird(X).
-flies(X) :- penguin(X), not abPeng(X).
bird(tweety).
baby(tweety).

abBird(X) :- bird(X), baby(X).
abPeng(X) :- penguin(X), flies(X).
abBird(X) :- penguin(X), not abPeng(X).
```

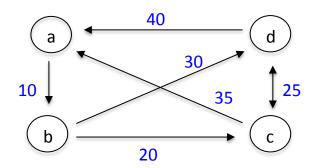
```
M = { bird(zippy), penguin(zippy),
      bird(Tweety), baby(tweety),
      abBird(zippy), abBird(tweety),
      -flies(zippy), flies(Tweety) }
```

```
penguin(zippy).
bird(X) :- penguin(X).
flies(X) :- bird(X), not abBird(X).
-flies(X) :- penguin(X), not abPeng(X).
bird(tweety).
baby(tweety).
veryFit(zippy).

abPeng(X) :- penguin(X), veryFit(X).
abBird(X) :- bird(X), baby(X).
abPeng(X) :- penguin(X), flies(X).
abBird(X) :- penguin(X), not abPeng(X).
```

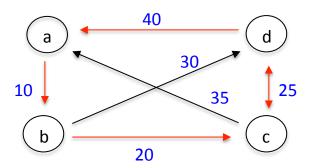
```
M = { bird(zippy), penguin(zippy), veryFit(zippy),
      bird(Tweety), baby(tweety),
      abBird(zippy), abBird(tweety), abPeng(zippy),
      -flies(zippy), flies(Tweety), flies(zippy) }
```

Travelling Salesman



Γράψτε πρόγραμμα ASP που θα επιλύει το πρόβλημα του πλανόδιου πωλητή.

Λύση

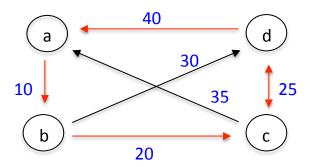


#const n=4.

city(a;b;c;d).

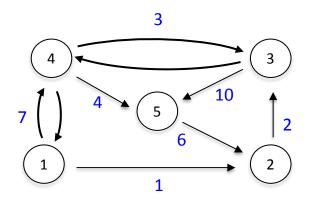
```
road(a,b,10). road(b,c,20). road(c,d,25). road(d,c,25). road(d,a,40). road(b,d,30). road(c,a,35).
% Επιλογή Κινήσεων
1 { at(C,T): city(C) } 1 :- T=1..n.
% Στοχος
visited(C) :- at(C,T), T=1..n.
:- city(C), not visited(C).
:- at(C1,T), at(C2,T+1), not road(C1,C2,_).
:- at(C1,n), at(C2,1), not road(C1,C2,_).
#minimize { D,C1,C2 : at(C1,T), at(C2,T+1), road(C1,C2,D); D,C1,C2 : at(C1,n), at(C2,1), road(C1,C2,D) }.
#show at/2.
```

Εκτέλεση



```
[<pavlos>$ clingo salesman.lp
clingo version 5.5.0
Reading from salesman.lp
Solving...
Answer: 1
at(a,4) at(b,1) at(c,3) at(d,2)
Optimization: 100
Answer: 2
at(a,4) at(b,1) at(c,2) at(d,3)
Optimization: 95
OPTIMUM FOUND
Models
             : 2
  Optimum
             : yes
Optimization: 95
Calls
             : 1
Time
             : 0.004s (Solving: 0.00s 1st Model: 0.00s Unsat: 0.00s)
CPU Time
             : 0.004s
```

Εύρεση Βέλτιστης Διαδρομής



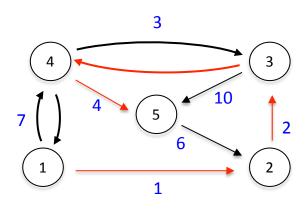
Γράψτε πρόγραμμα ASP που θα βρίσκει την βέλτιστη διαδρομή από την κορυφή 1 στην κορυφή 5.

Χρησιμοποιήσετε τα κατηγορήματα:

- node(X) για να δηλώσετε πως το X είναι κορυφή το γράφου,
- edge(X,Y,W) για να δηλώνει πως ο γράφος έχει ακμή που συνδέει την κορυφή X με την Y με κόστος W,

καθώς και όποια άλλα κατηγορήματα θεωρείται αναγκαία.

Λύση Α



#const n=5.

Το at(X,T) δηλώνει πως η βέλτιστη διαδρομή περνάει από την κορυφή X την στιγμή T.

```
#const start=1.

#const end=5.

node(1..n).

edge(1,2,1). edge(1,4,7).

edge(2,3,2).

edge(3,4,3). edge(3,5,10).

edge(4,1,7). edge(4,3,3). edge(4,5,4).

edge(5,2,6).
```

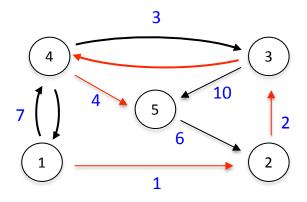
```
0 { at(X,T): node(X) } 1 :- T=2..n.

at(start,1).
:- not at(end,_).

:- at(X,T), at(Y,T+1), not edge(X,Y,_).
:- not at(_,T), at(_,T+1), T=1..n-1.

#minimize { D,X,Y: at(X,T), at(Y,T+1), edge(X,Y,D) }.
#show at/2.
```

Λύση Β



#const start=1.

```
• Το inPath(X,Y) δηλώνει πως η ακμή (X,Y) βρίσκεται στο βέλτιστο μονοπάτι.
```

Το reachable(X,Y) δηλώνει πως μέσω της βέλτιστης
 διαδρομής μπορούμε από την κορυφή X να φτάσουμε στην Y.

```
#const end=5.

edge(1,2,1). edge(1,4,7).

edge(2,3,2).

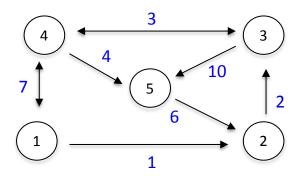
edge(3,4,3). edge(3,5,10).

edge(4,1,7). edge(4,3,3). edge(4,5,4).

edge(5,2,6).
```

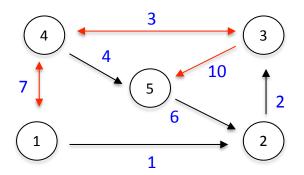
```
{ inPath(X,Y) : edge(X,Y,_) }.
reachable(X,Y) :- inPath(X,Y).
reachable(X,Y) :- reachable(X,Z), inPath(Z,Y).
:- not reachable(start,end).
#minimize{Z,X,Y: inPath(X,Y), edge(X,Y,Z)}.
#show inPath/2.
```

Εύρεση Χειρότερου Μονοπατιού



Γράψτε πρόγραμμα ASP που θα βρίσκει την χειρότερη ακυκλική διαδρομή από την 1 στην 5.

Λύση



```
#const n=5.
#const start=1.
#const end=5.

node(1..n).

edge(1,2,1). edge(1,4,7).
edge(2,3,2).
edge(3,4,3). edge(3,5,10).
edge(4,1,7). edge(4,3,3). edge(4,5,4).
edge(5,2,6).
```

```
0 { at(X,T): node(X) } 1 :- T=2..n.

at(start,1).
:- not at(end,_).

:- at(X,T), at(Y,T+1), not edge(X,Y,_).
:- at(X,T1), at(X,T2), T1 != T2.
:- at(end,T1), at(_,T2), T1<T2.

#maximize { D,X,Y: at(X,T), at(Y,T+1), edge(X,Y,D) }.
#show at/2.</pre>
```