

```

BinarySearch(L, i, j, key) ←
  if j ≤ i < 0
  then
    return -1
  else
    k := ⌊(i+j)/2⌋
    if key = L[k]
    then
      return k
    if key < L[k]
    then
      return BinarySearch(L, i, k-1, key)
    else
      return BinarySearch(L, k+1, j, key)

```

$O(n^2)$

```

IndiceDelMinimo(A, i, n)
k := i
for j = i+1 to n-1
  if A[j] < A[k]
  then
    k := j
return k

```

$O(n \log n)$

```

InsertionSort(A, n)
for j = 1 to n-1 do
  i := j-1
  tmp := A[j]
  while (i > 0 AND A[i] > tmp) do
    A[i+1] := A[i]
    i := i-1
  A[i+1] := tmp

```

$teta(n^2)$

```

QUICKSORT(A, p, r)
if p < r
then
  q := PARTITION(A, p, r)
  QUICKSORT(A, p, q-1)
  QUICKSORT(A, q+1, r)

```

```

TopologicalSort(G)
S := new_stack()
for all v ∈ V
  visited[v] := FALSE
for all v ∈ V
  if visited[v] = FALSE
  then DFS-TS(G, v)
return S

```

$O(n \log n)$

$O(n^2)$

$O(n \log n)$

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$O(n^2)$

```

DFS_pre_order(t)
if NOT (t = NIL)
then
  // analisi di t
  print t.val
  DFS_pre_order(t.left)
  DFS_pre_order(t.right)

```

$O(n * \text{costoxnodo})$

```

BFS(T)
if (NOT T = NIL)
Q := new_queue()
enqueue(Q, T)
while NOT is_empty_queue(Q) do
  u := dequeue(Q)
  // analisi di u
  print u.val
  for all v ∈ children(t) do
    enqueue(Q, v)

```

```

DFS_pre_order(t)
if NOT (t = NIL)
then
  // analisi di t
  print t.val
  w := t.fistchild
  while (NOT w = NIL) do
    DFS_pre_order(w)
    w := w.nextsibling

```

$O(n * \text{costoxnodo})$

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```

DFS(G)
for all v ∈ V
  visited[v] := FALSE
for all v ∈ V
  if visited[v] = FALSE AND DFS-Visit(G, v)
  then return TRUE
return FALSE

```

$O(|V|+|E|)$

```

DFS-Visit(G, v)
visited[v] := TRUE
for all (v, u) ∈ E
  // archi incidenti in v
  if visited[u] = TRUE OR DFS-Visit(G, u)
  then return TRUE
return FALSE

```

```

COUNTINGSORT(A, k)
n := length(A) // cardinalità di A
crea C[0..k]
for i = 0 to k // inizializzazione C
  C[i] := 0
for j = 0 to n-1
  C[A[j]] := C[A[j]] + 1
// calcolo occorrenze: C[i] = numero di
// occorrenze di i
j = 0
for i = 0 to k // "riempimento" di A
  while C[i] > 0
    A[j] := i
    j := j + 1
    C[i] := C[i] - 1

```

non stabile $O(n+k)$

```

Dijkstra(G, c, s)
for all u ∈ V
  dist[u] := +∞
  prev[u] := 0
  dist[s] := 0
Q := Make_priority_queue((v, dist[v]) | v = 1..n)
// coda con priorità,
// nodi di V con pri
while NOT is_empty_queue(Q) do
  u := Dequeue(Q) // restituisce l'elemento con pri
  // massima (valore minimo) e lo
  // dalla coda
  for all (u, v) ∈ E do
    if dist[v] > dist[u] + c(u, v)
    then
      dist[v] := dist[u] + c(u, v)
      prev[v] := u
      Decrease_Priority(Q, v, dist[v])
  return prev[]

```

$O((|V|+|E|) \cdot \log |V|)$

```

Bellman-Ford(G, c, s)
for all u ∈ V
  dist[u] := +∞
  prev[u] := 0
  dist[s] := 0
for i = 1 to |V|-1 do
  for all (u, v) ∈ E do
    if dist[v] > dist[u] + c(u, v)
    then
      dist[v] := dist[u] + c(u, v)
      prev[v] := u
for all (u, v) ∈ E do
  if dist[v] > dist[u] + c(u, v)
  then return "esiste un ciclo negativo"
return prev[]

```

INIZIALIZZAZIONE

```

BFS(G, s)
for all u ∈ V
  dist[u] := +∞
  prev[u] := -1
  dist[s] := 0
Q := new_queue() // coda FIFO
enqueue(Q, s)
while NOT is_empty_queue(Q) do
  u := dequeue(Q)
  eventuale esame di u
  for all (u, v) ∈ E do
    if dist[v] = +∞
    then
      enqueue(Q, v)
      dist[v] := dist[u] + 1
      prev[v] := u

```

```

Kruscal(G=(V, E), c)
S := make_set(V) //disjoint set
T := new_list()
ordina gli archi di E per ordine non decrescente di costo c
count := 0
while count < n - 1 do
  scegli il prossimo arco (u, v) ∈ E nell'ordinamento
  if find-set(S, u) ≠ find-set(S, v)
  then p := new_list_node()
  p.val := (u, v)
  insert-head(T, p)
  union(S, u, v)
  count := count + 1
return T

```

$O(m \log n)$

```

Prim( $G=(V,E),c$ )
for all  $v \in V$  do
     $cost[v] := +\infty$ 
     $prev[v] := NIL$ 
     $S[v] := 0$ 
scegli un nodo  $s \in V$ 
 $cost[s] := 0$ 
 $S[s] := 1$ 
 $Q := make\_priority\_queue(V')$  // coppie nodi di  $V$  e  $cost[]$ 
while NOT is_empty_queue( $Q$ ) do
     $u := DeQueue(Q)$ 
     $S[u] := 1$ 
    for all  $(u,v) \in E$  do
        if  $S[v] = 0$  AND  $cost[v] > c(u,v)$ 
        then
             $cost[v] := c(u,v)$ 
             $prev[v] := u$ 
            Decrease_Priority( $Q,v,cost[v]$ )
return  $prev[]$ 

```

Prim $O(m \log n)$
