

University of Technology, Jamaica

Theory of Computation: Worksheet

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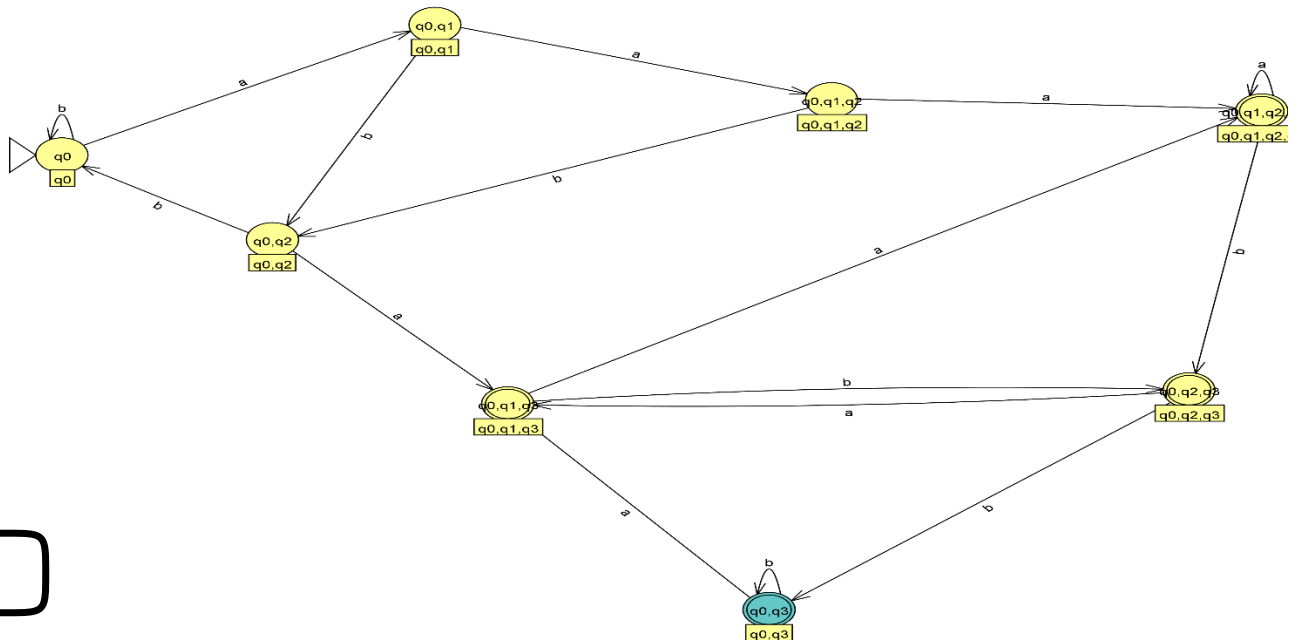
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1a.

	A	B
NULL	NULL	NULL
(Start){Q0}	{Q0,Q1}	{Q0} ✓
{Q1}	{Q2}	{Q2}
{Q2}	{Q3}	NULL
(Final){Q3}	{Q3}	{Q3}
{Q0,Q1}	{Q0,Q1,Q2}	{Q0,Q2} ✓
{Q0,Q2}	{Q0,Q1,Q3}	{Q0} ✓
{Q0,Q3}	{Q0,Q1,Q3}	{Q0,Q3} ✓
{Q1,Q2}	{Q2,Q3}	{Q2}
{Q1,Q3}	{Q2,Q3}	{Q2,Q3}
{Q2,Q3}	{Q3}	{Q3}
{Q0,Q1,Q2}	{Q0,Q1,Q2,Q3}	{Q0,Q2} ✓
{Q0,Q1,Q3}	{Q0,Q1,Q2,Q3}	{Q0,Q2,Q3} ✓
{Q0,Q2,Q3}	{Q0,Q1,Q3}	{Q0,Q3} ✓
{Q1,Q2,Q3}	{Q2,Q3}	{Q2,Q3}
{Q1,Q2,Q3,Q4}	{Q0,Q1,Q2,Q3}	{Q0,Q2,Q3} ✓



1b.

$$E(\{Q2\}) = (\{Q2, Q3, Q6, Q9, Q12\}) \quad \checkmark$$

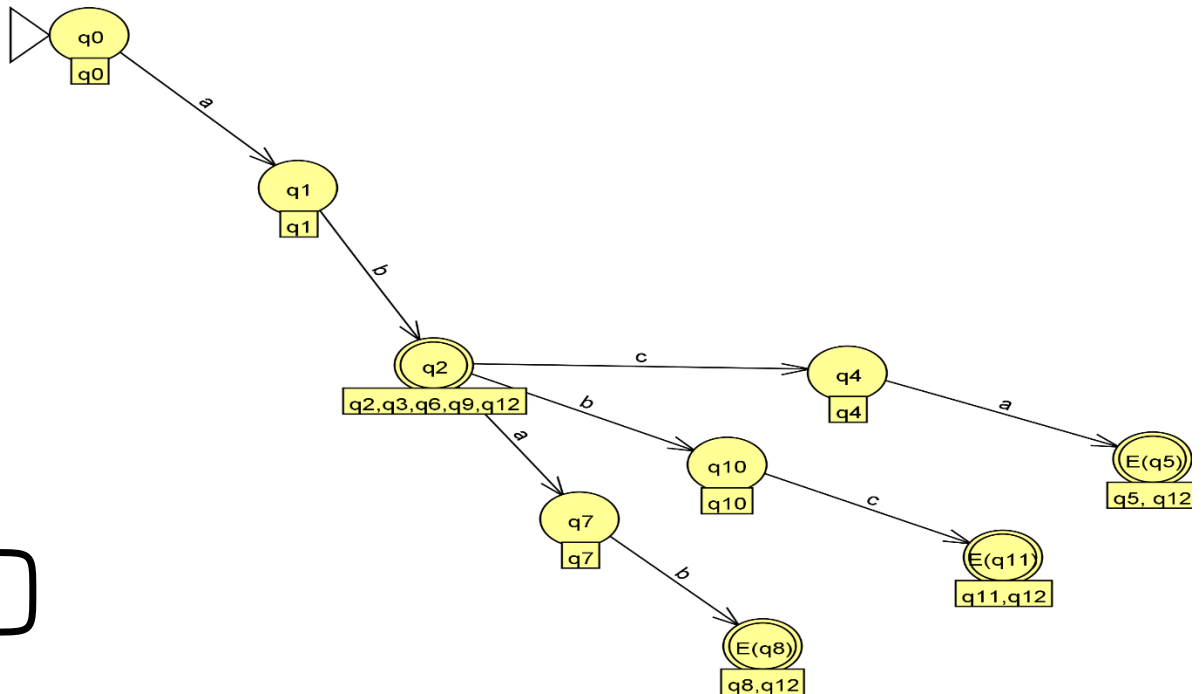
$$E(\{Q5\}) = (\{Q5, Q12\}) \quad \checkmark$$

$$E(\{Q8\}) = (\{Q8, Q12\}) \quad \checkmark$$

$$E(\{Q11\}) = (\{Q11, Q12\}) \quad \checkmark$$

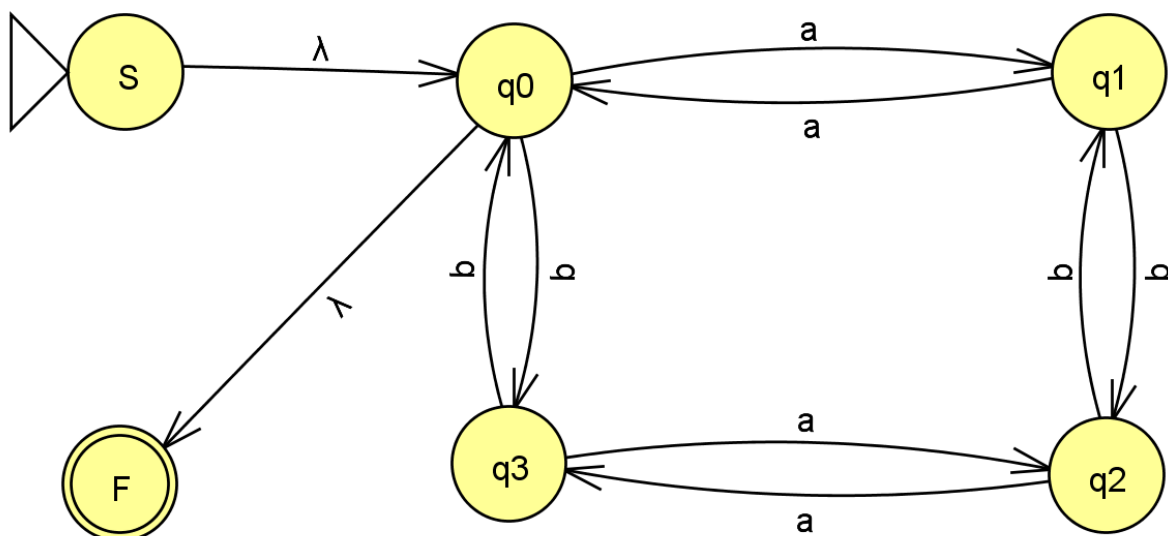
	A	B	C
(Start){Q0}	{Q1}	NULL	NULL
{Q1}	NULL	{Q2, Q3, Q6, Q9, Q12}	NULL
{Q2, Q3, Q6, Q9, Q12}	{Q7}	{Q10}	{Q4}
{Q4}	{Q5, Q12}	NULL	NULL
{Q7}	NULL	{Q8, Q12}	NULL
{Q10}	NULL	NULL	{Q11, Q12}
{Q5, Q12}	NULL	NULL	NULL
{Q8, Q12}	NULL	NULL	NULL
{Q11, Q12}	NULL	NULL	NULL

The null state is unable to reach a final state and is removed



2.

Add new start state and final state



Remove q0	
$S-q0-q1 \rightarrow a$	$S-q1 \rightarrow a$
$S-q0-q3 \rightarrow b$	$S-q3 \rightarrow b$
$S-q0-F \rightarrow \lambda$	$S-F \rightarrow \lambda$
$q1-q0-q3 \rightarrow ba$	$q1-q3 \rightarrow ba$
$q3-q0-q1 \rightarrow ab$	$q3-q1 \rightarrow ab$
$q1-q0-F \rightarrow a$	$q1-F \rightarrow a$
$q1-q0-F \rightarrow b$	$q3-F \rightarrow b$
$(q1-q0)^* \rightarrow aa$	$(q1)^* \rightarrow aa$
$(q3-q0)^* \rightarrow bb$	$(q3)^* \rightarrow bb$

✓

Remove q1	
$S-q1-q3 \rightarrow a(aa)^*ab$	$S-q3 \rightarrow b+a(aa)^*ab$
$S-q1-F \rightarrow a(aa)^*a$	$S-F \rightarrow \lambda+a(aa)^*b$
$q3-q1-F \rightarrow ba(aa)^*a$	$q3-F \rightarrow b+ba(aa)^*a$
$(q3-q1)^* \rightarrow ba(aa)^*ab$	$(q3)^* \rightarrow bb+ba(aa)^*ab$
$S-q1-q2 \rightarrow a(aa)^*b$	$S-q2 \rightarrow a(aa)^*b$
$q2-q1-F \rightarrow b(aa)^*a$	$q2-F \rightarrow b(aa)^*a$
$q2-q1-q3 \rightarrow b(aa)^*ab$	$q2-q3 \rightarrow b(aa)^*ab$
$(q2-q1)^* \rightarrow b(aa)^*b$	$(q2)^* \rightarrow b(aa)^*b$
$q3-q1-q2 \rightarrow ba(aa)^*b$	$q3-q2 \rightarrow ba(aa)^*b$

Remove q2	
$S-q2-F \rightarrow a(aa)^*b(b(aa)^*b)^*b(aa)^*a$	$S-F \rightarrow \lambda + a(aa)^*a + a(aa)^*b(b(aa)^*b)^*b(aa)^*a$
$S-q2-q3 \rightarrow$ $a(aa)^*b(b(aa)^*b)^*a + b(aa)^*ab$	$S-q3 \rightarrow$ $b + a(aa)^*ab + a(aa)^*b(b(aa)^*b)^*a + b(aa)^*ab$
$q3-q2-F \rightarrow$ $a + ba(aa)^*b(b(aa)^*b)^*b(aa)^*a$	$q3-F \rightarrow$ $b + ba(aa)^*a + a + ba(aa)^*b(b(aa)^*b)^*b(aa)^*a$
$(q3-q2)^* \rightarrow$ $a + ba(aa)^*b(b(aa)^*b)^*a + b(aa)^*ab$	$(q3)^* \rightarrow$ $bb + ba(aa)^*ab + a + ba(aa)^*b(b(aa)^*b)^*a + b(aa)^*ab$

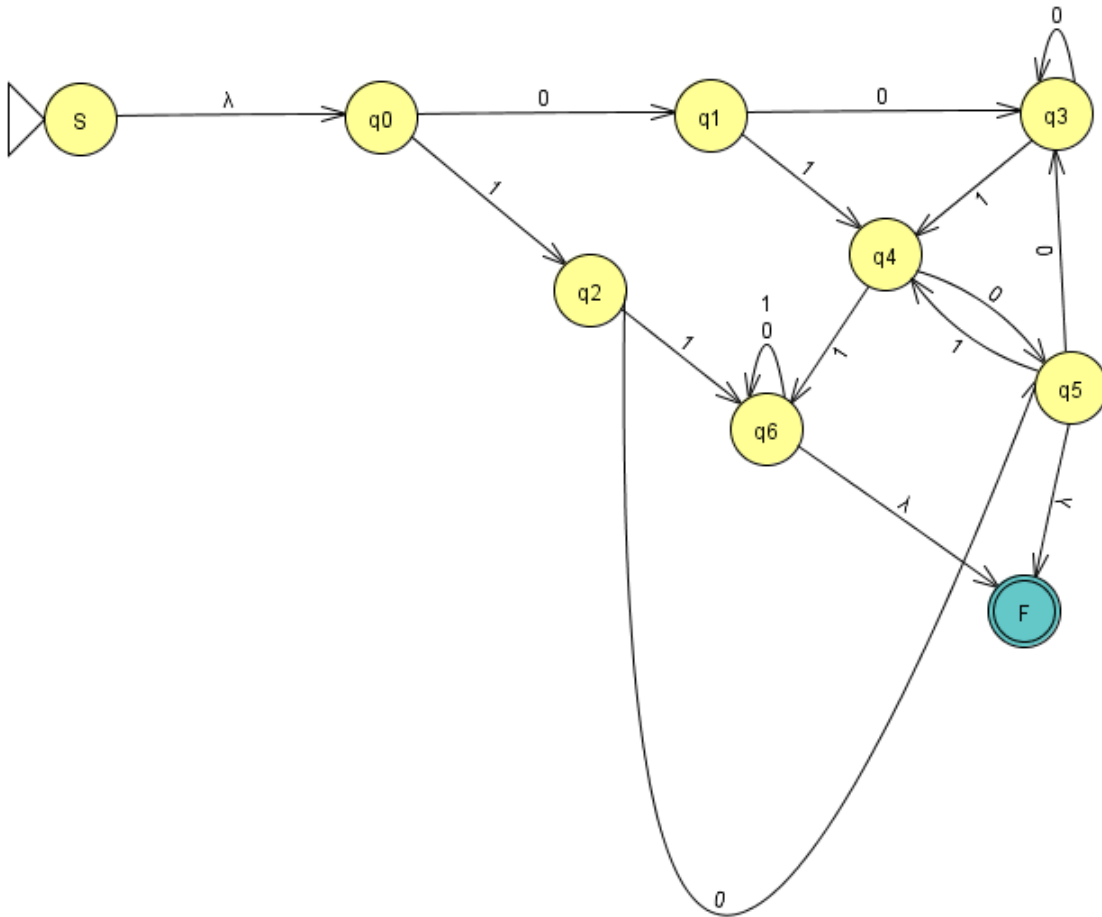
Remove q3	
$S-q3-F \rightarrow$ $b + a(aa)^*ab + a(aa)^*b(b(aa)^*b)^*a + b(aa)^*$ $*ab(bb + ba(aa)^*ab + a + ba(aa)^*b(b(aa)^*$ $b)^*a + b(aa)^*ab)^*$ $b + ba(aa)^*a + a + ba(aa)^*b(b(aa)^*b)^*b(a$ $a)^*a$	$S-F \rightarrow$ $\lambda + a(aa)^*a + a(aa)^*b(b(aa)^*b)^*b(aa)^*a + (b + a(aa)^*ab + a($ $aa)^*b(b(aa)^*b)^*a + b(aa)^*ab(bb + ba(aa)^*ab + a + ba(aa)^*$ $b(b(aa)^*b)^*a + b(aa)^*ab)^*$ $b + ba(aa)^*a + a + ba(aa)^*b(b(aa)^*b)^*b(aa)^*a)$

Answer:

$\lambda + a(aa)^*a + a(aa)^*b(b(aa)^*b)^*b(aa)^*a + (b + a(aa)^*ab + a(aa)^*b(b(aa)^*b)^*a + b(aa)^*ab(bb + ba(aa)^*ab + a + ba(aa)^*b(b(aa)^*b)^*a + b(aa)^*ab)^* b + ba(aa)^*a + a + ba(aa)^*b(b(aa)^*b)^*b(aa)^*a)$

Please show diagrams
with your steps next time. This is
not easy to follow

b. Add new start and final state



Remove q0	
S-q0-q1 → 0	S-q1 → 0
S-q0-q2 → 1	S-q2 → 1

✓

Remove q1	
S-q1-q3 → 00	S-q3 → 00
S-q1-q4 → 01	S-q4 → 01

✓

Remove q2	
S-q2-q5 → 10	S-q5 → 10
S-q2-q6 → 11	S-q6 → 11

✓

Remove q3	
S-q3-q4 → 000*1	S-q3 → 01+000*1
q5-q3-q4 → 00*1	S-q4 → 1+00*1

✓

Remove q4	
$S-q4-q5 \rightarrow 01+(000*1)0$	$S-q5 \rightarrow 10+(01+(000*1)0)$
$S-q4-q6 \rightarrow 01+(000*1)1$	$S-q6 \rightarrow 11+(01+(000*1)1)$
$q5-q4-q6 \rightarrow 1+(00*1)1$	$q5-q6 \rightarrow 1+(00*1)1$
$(q5-q4)^* \rightarrow 1+(00*1)0$	$(q5)^* \rightarrow 1+(00*1)0$

✓

✓

Remove q5	
$S-q5-F \rightarrow 10+(01+(000*1)0)(1+(00*1)0)^*$	$S-F \rightarrow 10+(01+(000*1)0)(1+(00*1)0)^*$
$S-q5-q6 \rightarrow 10+(01+(000*1)0)(1+(00*1)0)^*(1+(00*1)1)$	$S-q6 \rightarrow 11+(01+(000*1)1)+(10+(01+(000*1)0)(1+(00*1)0)^*(1+(00*1)1))$

✓

Remove q6	
$S-q6-F \rightarrow 11+(01+(000*1)1)+(10+(01+(000*1)0)(1+(00*1)0)^*(1+(00*1)1))(0+1)^*$	$S-F \rightarrow 10+(01+(000*1)0)(1+(00*1)0)^*+11+(01+(000*1)1)+(10+(01+(000*1)0)(1+(00*1)0)^*(1+(00*1)1))(0+1)^*$

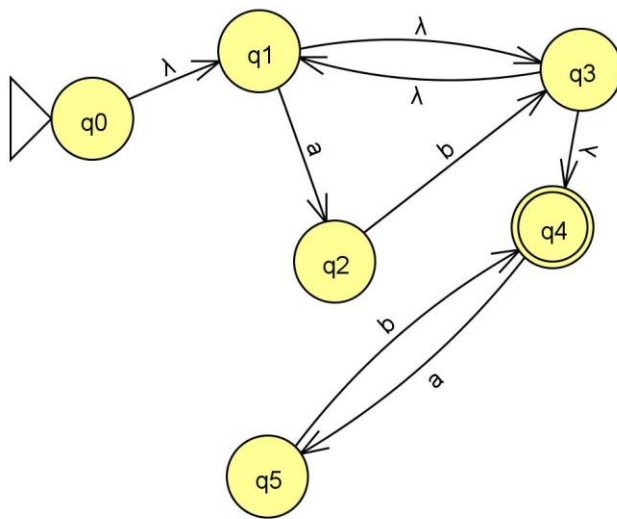
Answer:

$$10+(01+(000*1)0)(1+(00*1)0)^*+11+(01+(000*1)1)+(10+(01+(000*1)0)(1+(00*1)0)^*(1+(00*1)1))(0+1)^*$$

✓

3. RE to DFA

a)



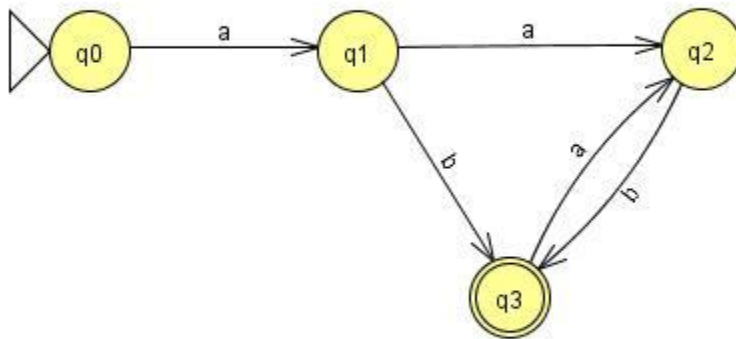
Should be able to start with b

C is mandatory, where is it?

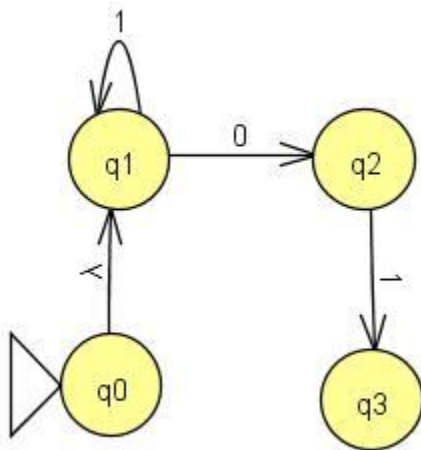
The question said less than 4 states, I count 6

And DFA does not recognize empty symbol

0/2



b)



You keep using the empty symbol but you were asked for a DFA

- 0 . 5

must start with 1 and where is your accept state?

- 1

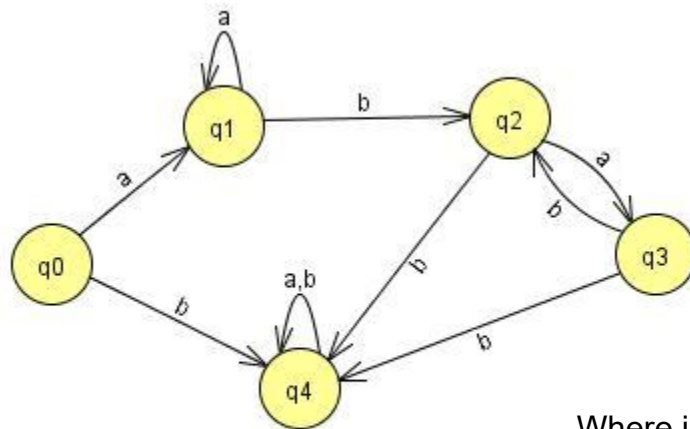
1.5/3

4. DFA Minimization

a)

q0	-	-	-	-	-	-
q1	-	-	-	-	-	X
q2	-	-	-	-	X	X
q3	-	-	-	X	X	X
q4	-	-	Yes	X	X	X
q5	-	X	X	X	X	X
	q5	q4	q3	q2	q1	q0

Minimized Diagram:



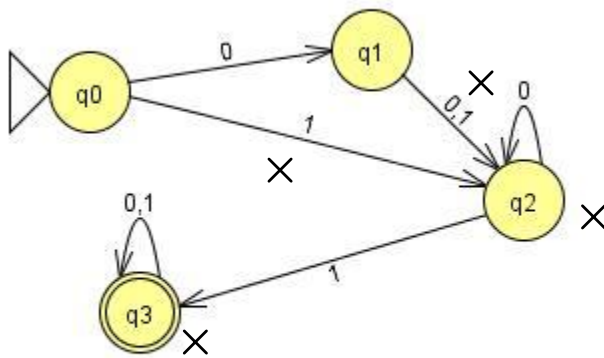
Where is the start and final states?

- 2

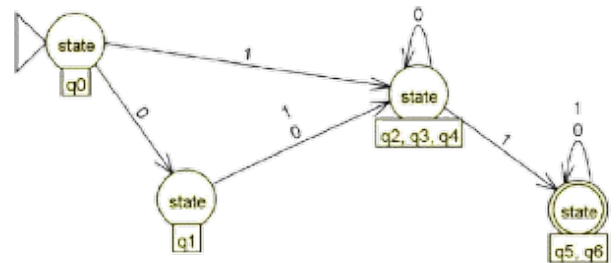
b)

q0	-	-	-	-	-	-	-
q1	-	-	-	-	-	-	X
q2	-	-	-	-	-	X	X
q3	-	-	-	-	YES	X	X
q4	-	-	-	YES	X	X	X
q5	-	-	X	X	X	X	X
q6	-	YES	X	X	X	X	X
	q6	q5	q4	q3	q2	q1	q0

Minimized Diagram:



Answer:



0.5/4

5a. There exists a language where $L = \{www \mid w \text{ is an element of } \{0,1\}^*\}$ is not regular. We will assume that L is a regular language and let $p \geq 1$ be the pumping length. A string s of $(0^p)(1^p)$ is established. $s \in A$ and $|s| = 2p \geq p$. Using the Pumping Lemma, $s = xyz$ where $y \neq \epsilon$, $|xy| \leq p$ and $xy^iz \in A$ for all $i \geq 0$. Since $|xy| \leq p$, the string y contains only 0s. Since $y \neq \epsilon$, y contains at least one 0. Therefore, the string $xy^3z = xyyyz$ contains more 0s than 1s which implies that the string is not contained in L and therefore is a contradiction, L is not a regular language.

What is this nonsense? where is your pumping length?

The string cannot be 01, it has to be a multiple of 3.

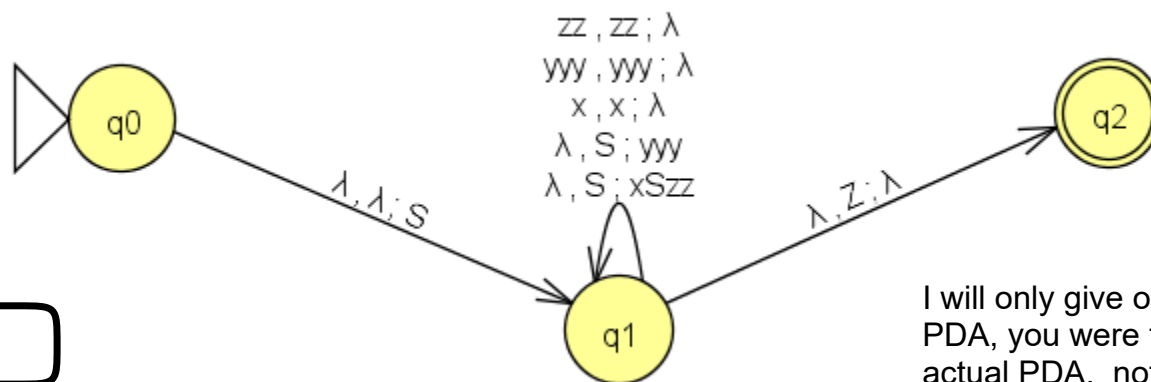
0/4

6a.

what value did you assign to k ? were you absent when I was teaching pumping lemma?

0/5

$S \rightarrow xSzz|yyy$ ✓



4/6

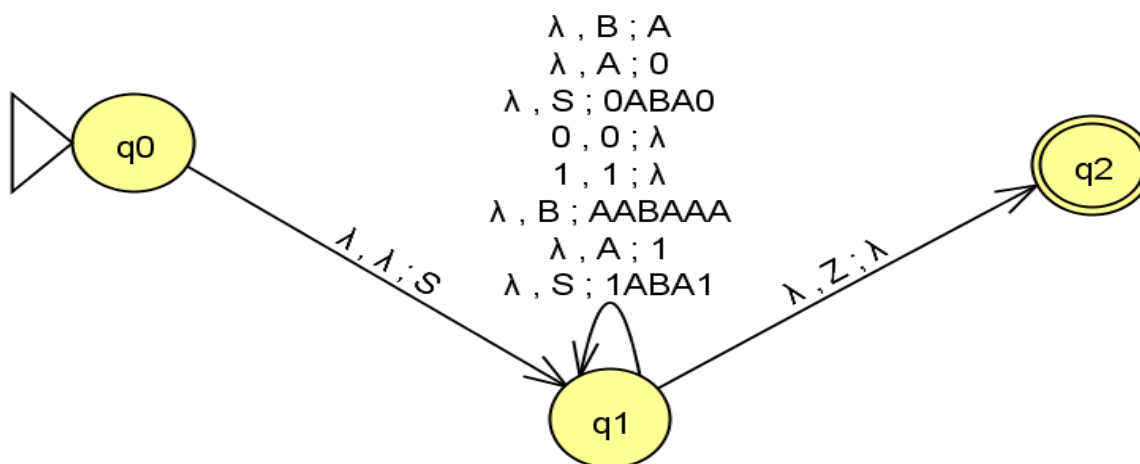
I will only give one mark for PDA, you were to create the actual PDA, not convert CFG to PDA

b.

$A \rightarrow 1|0$

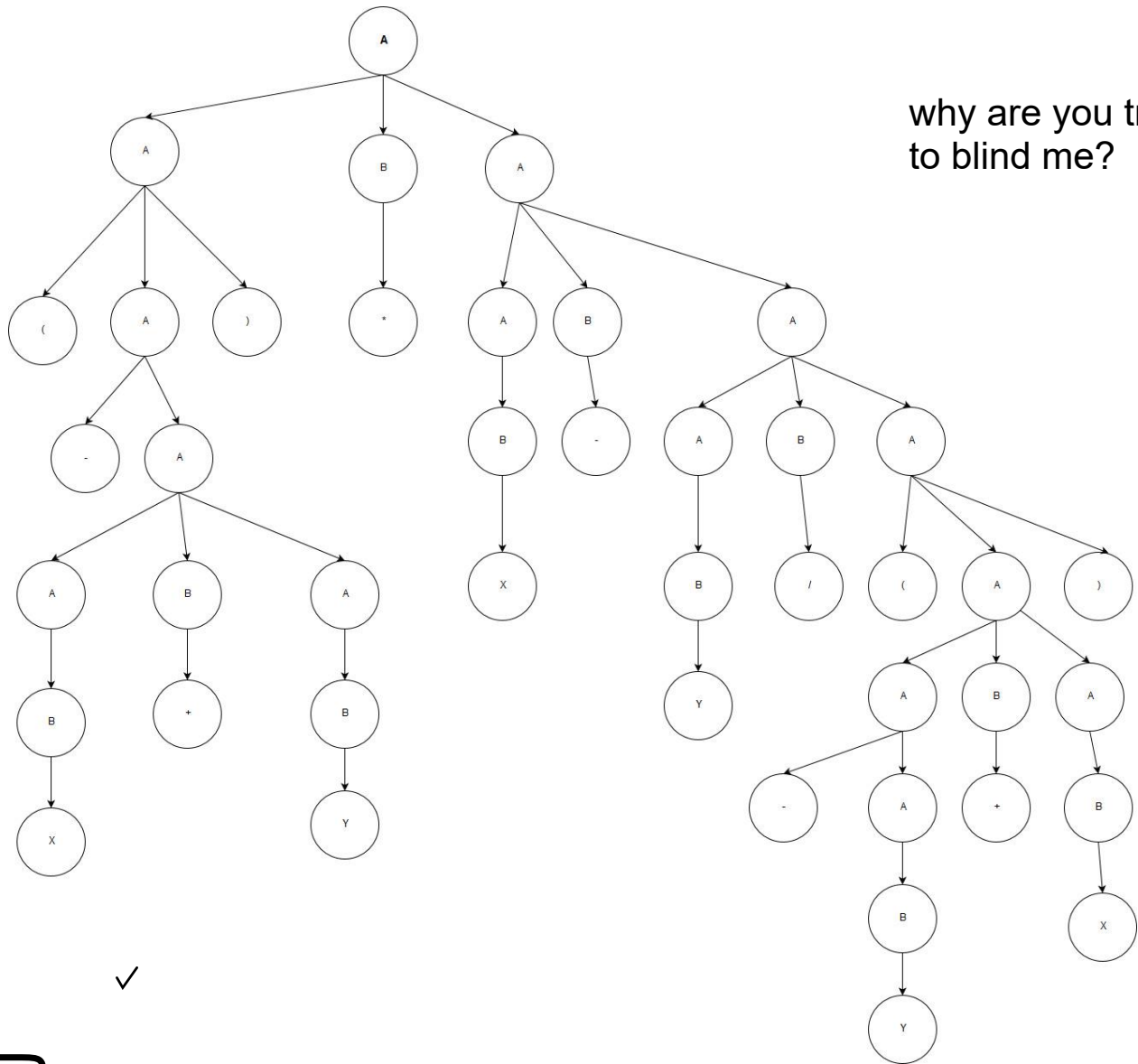
$B \rightarrow AABAAA|A$

$S \rightarrow 1ABA1|0ABA0$



5/8

7a.



6/6

b.)

$A \rightarrow ABA \rightarrow (A)BA \rightarrow (ABA)BA \rightarrow (BBA)BA \rightarrow (YBA)BA \rightarrow (Y/A)BA \rightarrow (Y/B)BA \rightarrow (Y/Y)BA \rightarrow (Y/Y)-ABA \rightarrow$
 $(Y/Y)-BBA \rightarrow (Y/Y)-XBA \rightarrow (Y/Y)-X^*A \rightarrow (Y/Y)-X^*(A) \rightarrow (Y/Y)-X^*(-A) \rightarrow (Y/Y)-X^*(-B) \rightarrow (Y/Y)-X^*(-Y)$

✓

4/4