

A SIMPLICIAL HOMOLOGY ALGORITHM FOR LIPSCHITZ OPTIMISATION

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A simplicial homology algorithm for Lipschitz optimisation

by

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Synopsis

The simplicial homology global optimisation (SHGO) algorithm is a general purpose global optimisation algorithm based on applications of simplicial integral homology and combinatorial topology. SHGO approximates the homology groups of a complex built on a hypersurface homeomorphic to a complex on the objective function. This provides both approximations of locally convex subdomains in the search space through Sperner's lemma ([1, 2]) and a useful visual tool for characterising and efficiently solving higher dimensional black and grey box optimisation problems. This complex is built up using sampling points within the feasible search space as vertices. The algorithm is specialised in finding all the local minima of an objective function with expensive function evaluations efficiently which is especially suitable to applications such as energy landscape exploration. SHGO was initially developed as an improvement on the topographical global optimisation (TGO) method first proposed by [3] ([4]). It is proven that the SHGO algorithm will always outperform TGO on function evaluations if the objective function is Lipschitz smooth. In this paper SHGO is applied to non-convex problems with linear and box constraints with bounds placed on the variables. Numerical experiments on linearly constrained test problems show that SHGO gives competitive results compared to TGO and the recently developed Lc-DISIMPL algorithm ([5, 6]) as well as the PSwarm and DIRECT-L1 algorithms. Furthermore SHGO is compared with the TGO, basinhopping (BH) and differential evolution (DE) global optimisation algorithms over a large selection of black-box problems with bounds placed on the variables from the SciPy ([7, 8]) benchmarking test suite. A Python implementation of the SHGO and TGO algorithms published under a MIT license can be found from <https://bitbucket.org/upiamcompthermo/shgo/>.

SLEUTELWOORDE: Global optimisation, SHGO, Computational homology

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CHAPTER 1

Introduction

1.1 Objective function statement and nomenclature

Consider a general optimisation problem of the form

$$\begin{aligned} \min_{\mathbf{x}} \quad & f(\mathbf{x}) \\ \text{s.t.} \quad & \mathbf{g}(\mathbf{x}) \geq 0 \end{aligned} \tag{1.1}$$

The continuous real objective function $f(\mathbf{x})$ of n dimensionality can be either smooth or non-smooth depending on the local minimisation method used. The variables \mathbf{x} are assumed to be bounded. In this publication we mainly consider real, smooth, but not necessarily convex functions with linear constraint functions. In addition we will assume that the objective function has a finite number of local minima

$$f : \mathbb{R}^n \rightarrow \mathbb{R} \tag{1.2}$$

\mathbf{g} maps the set of linear constraints

$$\mathbf{g} : [\mathbf{l}, \mathbf{u}]^n \rightarrow \mathbb{R}^m \tag{1.3}$$

for example if lower and upper bounds l_i and u_i are implemented for each variable then we have an initially defined hyperrectangle

$$\mathbf{x} \in \Omega \subseteq [\mathbf{l}, \mathbf{u}]^n = [l_1, u_1] \times [l_2, u_2] \times \dots \times [l_n, u_n] \subseteq \mathbb{R}^n \tag{1.4}$$

where Ω is the limited feasible subset excluding points outside the bounds and constraints.

$$\Omega = \{\mathbf{x} \in [\mathbf{l}, \mathbf{u}]^n \mid \mathbf{g}_i(\mathbf{x}) \geq 0, \forall i = 1, \dots, m\} \tag{1.5}$$

Since the constraints in \mathbf{g} are linear the set Ω is always a compact space.

In the development of SHGO several concepts from algebraic and combinatorial topology are required. The following definition was adapted from [?]: p. 9

Definition 1. A **k-simplex** is a set of $n+1$ vertices in a convex polyhedron of dimension n . Formally if the $n+1$ points are the $n+1$ standard basis vectors for $\mathbb{R}^{(n+1)}$. Then the n -dimensional k -simplex is the set

$$S^n = \left\{ (t_1, \dots, t_{n+1}) \in \mathbb{R}^{n+1} \mid \sum_{i=1}^{n+1} t_i = 1, t_i \geq 0 \right\}$$

For example, a 2-simplex is a triangle and a 3-simplex is a tetrahedron. We will use the following combinatorial definition of a simplicial complex ([?]: p. 107)

Definition 2. A **simplicial complex** \mathcal{H} is a set \mathcal{H}^0 of vertices together with sets \mathcal{H}^n of n -simplices, which are $(n+1)$ -element subsets of \mathcal{H}^0 . The only requirement is that each $(k+1)$ -element subset of the vertices of an n -simplex in \mathcal{H}^n is a k -simplex, in \mathcal{H}^k .

Thus each n -simplex has $n+1$ distinct vertices, and no other n -simplex has this same set of vertices.

In this publication the \mathcal{H} symbol will be used to represent a (finite) simplicial complex rather than the more standard Δ to avoid confusion with the difference and Laplacian operators common in optimisation. The superscript \mathcal{H}^k represents the subset of k -dimensional simplices where for an n dimensional problem the highest dimensional k -simplex contains $n+1$ vertices. Finally we define a k -chain?

Definition 3. A **k-chain** is a union of simplices.

For example a 0-chain is a set of vertices, a 1-chain is a set of edges and a 2-chain is a set of triangles. $C(\mathcal{H}^k)$ denotes a k -chain of k -simplices. A vertex in \mathcal{H}^0 is denoted by v_i . If v_i and v_j are two endpoints of a directed edge in \mathcal{H}^1 from v_i to v_j then the symbol $\overline{v_i v_j}$ represents the edge so that it is bounded by the 0-chain $\partial(\overline{v_i v_j}) = v_j - v_i$ and similarly for an edge directed from v_j to v_i , we have, $\partial(\overline{v_j v_i}) = \partial(-\overline{v_i v_j}) = v_i - v_j$. Higher dimensional simplices can be represented and directed in a similar manner, for example a triangle consisting of three vertices v_i, v_j and v_k directed as $\overline{v_i v_j v_k}$ has the boundary of directed edges $\partial(\overline{v_i v_j v_k}) = \overline{v_i v_j} + \overline{v_j v_k} + \overline{v_k v_i}$.

1.2 Multimodal objective functions and local minima mapping

Non-convex problems are commonly solved using global optimisation methods. One such example is the topographical global optimisation (TGO) method [????] which is a clustering algorithm that finds several local minima from which the (approximate) global

minimum is found. It is often desirable to find all the local minima of the objective function for example in applications such as energy landscape exploration of potential models wherein mapping the local minima of the potential functions can provide valuable insights into the system. Algorithms such as the basin-hopping global optimisation algorithm are typically used to find these points (?).

The graph extracted from the topographical global optimisation (TGO) to topograph (as described in chapter 2) is unsatisfactory in some ways. Primarily because several starting points in the same locally convex domain can be generated even when enough information from the objective function sampling is known to prevent this from occurring. This leads to superfluous function evaluations in the local minimisation step of the algorithm. Contrary to intuition, this problem is exacerbated by increasing the number of initial sampling points used in the algorithm as demonstrated in chapter 2. This can lead to a very large number of function evaluations required to solve the problem. In particular in multimodal energy surfaces where the local minima can often be located in short distances relative to the search space ? and thus requires a large number of initial sampling to locate all these domains. Some shortcomings in using the TGO method to map local minima are:

- Geometric information available from the sampling points is being disregarded by the graphs built up using only the Euclidean distance metric.
- Knowledge of the number and location of local minimisers in a given sampling set is not being used to the full extent.
- More than one minimiser might be produced in the same locally convex domain and there is no guarantee that a minimiser set produced by TGO will be in the locally convex domains of all local minima even if the number of local minima is known and a minimiser set of this cardinality is produced.

By constructing a directed simplicial complex we show that the simplicial homology global optimisation (SHGO) algorithm does not produce superfluous starting points for the class of all Lipschitz smooth functions resulting in more efficient performance for these problems compared to TGO. The directed complex is also used to approximate the homology group of the objective function hypersurface which, using integral homology version of the Invariance Theorem ?, allows for efficient mapping of optimisation problems where the number of local minima is known *a-priori*.

1.3 Derivative-free methods for Lipschitz optimisation problems

Both the SHGO and TGO algorithms only make use of function evaluations without requiring the derivatives of objective functions. This makes them applicable to black-box global optimisation problems. A recent review and experimental comparison of 22 derivative-free optimisation algorithms by [?] concluded that global optimisation solvers such as TOMLAB/MULTI-MIN, TOMLAB/GLCCLUSTER, MCS and TOMLAB/LGO perform better, on average, than other derivative-free solvers in terms of solution quality within 2500 function evaluations. Both the TOMLAB/GLCCLUSTER and MCS [?] implementations are based on the well-known DIRECT (DIviding RECTangle) algorithm [?].

The DISIMPL (DIviding SIMPLices) algorithm was recently proposed by [?]. The experimental investigation in [?] shows that the proposed simplicial algorithm gives very competitive results compared to the DIRECT algorithm. DISIMPL has been extended in [?]. The Gb-DISIMPL (Globally-biased DISIMPL) was compared in [?] to the DIRECT and DIRECTl methods in extensive numerical experiments on 800 multidimensional multiextremal test functions.

In a recent adaption of DISIMPL for linearly constrained optimisation problems, Lc-DISIMPL [?] showed extremely competitive results compared to the PSwarm [?] and DIRECT-L1 algorithms [?]. In particular the Lc-DISIMPL-v algorithm was shown to solve the problems in a fewer number of function evaluations on average and was the only algorithm to converge on all of the test problems. In this publication both the SHGO and TGO algorithms were tested on the same problem set and the results are compared to the data from [?] which also contains results on the PSwarm [?] and DIRECT-L1 algorithms [?].

The DISIMPL algorithm is the most similar to SHGO in the sense that both make use of a simplicial complex. DISIMPL uses a simplicial complex in a spatial partitioning of the initial search space. Since the geometric structure of the two algorithms are related, it is reasonable to expect some theoretical relation of its properties. In particular the graph structure in the DISIMPL-v algorithm [?] can be used to build the directed simplicial complex used by SHGO. In chapter 4 we also show how some of the same principles developed for SHGO can also be applied in the DISIMPL-v algorithm since the same information is readily available to the algorithm.

1.4 Overview of this publication

The TGO method is briefly reviewed in chapter 2 closely following the formalism developed by ?. In chapter 3 we provide numerical examples of TGO which is then used as an informal experimental motivation for extending the algorithm. These two sections are important for continuity and understanding of the improved features of SHGO, in particular 9 which will be used as a performance criterion. In section 3.1 we present the most immediately apparent extension of TGO and illustrate the shortcomings of that approach. The new SHGO method is then formally presented in chapter 4. In chapter 5 we provide experimental results of linearly constrained problems comparing the SHGO, TGO, Lc-DISIMPL ?, PSwarm ? and DIRECT-L1 ? algorithms. Furthermore SHGO is compared with the TGO, basinhopping (BH) and differential evolution (DE) global optimisation algorithms over a large selection of black-box problems from the SciPy (?, ?) global optimisation benchmarking test suite. We conclude with various recommendations for possible further improvements of SHGO.

CHAPTER 2

Topographical Global Optimisation (TGO)

The Topographical Global Optimisation (TGO) was originally conceived by ? and Henderson et al. ?? introduced new formalisms and empirical methods to determine hyperparameters described in this section. ? also presents the algorithm in an introductory fashion. It is in essence an iterative clustering algorithm that maps the hypersurface of the objective function into a topography matrix (called a t -matrix) and then finds a certain number of starting points referred to as local minimisers. A local search using the local minimisers as starting points is then used to find each minimum from which the global minimum is finally calculated. ? used the feasible direction interior-point method proposed by ? in this step. The feasible direction interior-point method allows for minimisation of problems with linear and/or nonlinear equality constraints; an extension by ? of the original applications of ?. The TGO method consists of three steps:

1. Uniform random sampling generation of N points in the search space.
2. Construction of the topograph, which is a directed graph with the sampled points as vertices on a k -nearest neighbours basis with the direction of the arc directed towards a point with a larger function value.
3. Local minimisation of topograph minimisers.

2.1 Step 1: Random Sampling Point generation

In order to generate the uniform sampling points within Ω the deterministic Sobol sequence is used in this publication ??. Other possible low discrepancy sequences such as the Halton and Van der Corput sequences ? can also be used in this step. An efficient Gray code implementation was proposed by ? wherein a single XOR operation for each dimension can be used to find the next sampling point in the sequence $x_{n,i} = x_{n-1,i} \oplus v_{k,i}$.

An adaptation of this method is available in the open source Python library UQToolbox?. The Sobol sequenced points are generated within the n dimensional hypercube $[0, 1]^n \in \mathbb{R}^n$, providing a uniform distribution on the hypersurface within this space. In the current implementation this set of points is stretched across the lower and upper bounds to form the hyperrectangle $[\mathbf{l}, \mathbf{u}]^n = [l_1, u_1] \times [l_2, u_2] \times \cdots \times [l_n, u_n] \subseteq \mathbb{R}^n$. The subset of feasible points contained in Ω is found by discarding any points lying outside the constraints $\mathbf{g}(\mathbf{x}) > 0$.

2.2 Step 2: Construction of the topograph

The topograph is constructed from the generated sampling points within Ω . From the topograph several global minimisers in f are found using the definitions developed in this section which are then used as starting points for local minimisation routines. First N points are selected from the uniformly generated sequence of points within the feasible domain of $\Omega \subset \mathbb{R}^n$. Points generated by the sequence that lie outside the constraints are excluded. The points are denoted by $\mathbf{p}_i, i = 1, 2, 3 \dots N$. Next for each point \mathbf{p}_i a reference list is constructed by ordering the other $N - 1$ points from their nearest to farthest Euclidean distances. These ordered lists make up the rows of the topography matrix (or topograph). Furthermore, for some point $\mathbf{p}_j \in \{1, 2, 3 \dots (N - 1)\}$ in the row with the first entry \mathbf{p}_i , a sign is assigned as follows:

$$\text{sign}(\mathbf{p}_j) = \begin{cases} f(\mathbf{p}_j) \geq f(\mathbf{p}_i) & \rightarrow + \\ f(\mathbf{p}_j) < f(\mathbf{p}_i) & \rightarrow - \end{cases}$$

In order to demonstrate this construction we will define this ordered list in such a way that the increasing indices represent an ordered list of the nearest points to \mathbf{p}_1 , that is $\|\mathbf{p}_i - \mathbf{p}_{i+1}\| \leq \|\mathbf{p}_{i+1} - \mathbf{p}_{i+2}\| \forall i$. Suppose for example that $f(\mathbf{p}_2) \geq f(\mathbf{p}_1)$, $f(\mathbf{p}_3) < f(\mathbf{p}_1)$ and $f(\mathbf{p}_N) \geq f(\mathbf{p}_1)$, the resulting topograph with the first row known is:

$$t\text{-matrix} = \begin{pmatrix} [c|cccc]\mathbf{p}_1 & +\mathbf{p}_2 & -\mathbf{p}_3 & \dots & +\mathbf{p}_N \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ \mathbf{p}_N & \mathbf{p}_j & \dots & \mathbf{p}_j & \mathbf{p}_j \end{pmatrix} \quad (2.1)$$

Note that the remaining rows (represented by unknown points and signs \mathbf{p}_j) are constructed similarly to the first row for every \mathbf{p}_i row. The topography matrix can be interpreted as a directed graph, where the signs represent the directed arcs on the graph. It should also be noted that if \mathbf{g} contains non-linear constraints then the graphs produced by the topograph may be connected across disconnected and/or non-convex subspaces of Ω . Example 1 in chapter 3 demonstrates the construction of the topograph numerically.

Given an integer $1 \leq k \leq (N - 1)$, the $N \times k$ submatrix obtained by considering

only the k -nearest neighbours is called the k - t -matrix. For example for $k = 1$:

$$\text{1-}t\text{-matrix} = \begin{pmatrix} [c|c]\mathbf{p}_1 & +\mathbf{p}_2 \\ \vdots & \vdots \\ \mathbf{p}_N & \mathbf{p}_j \end{pmatrix} \quad (2.2)$$

for $k = 2$:

$$\text{2-}t\text{-matrix} = \begin{pmatrix} [c|cc]\mathbf{p}_1 & +\mathbf{p}_2 & -\mathbf{p}_3 \\ \vdots & \vdots & \vdots \\ \mathbf{p}_N & \mathbf{p}_j & \mathbf{p}_j \end{pmatrix} \quad (2.3)$$

and so forth. The k - t -matrix is a representation of its k^+ -topograph where every row forms a directed subgraph.

The following definitions adapted from ? are used to find the global minimisers of the objective function

Definition 4. Given an integer $1 \leq k \leq (N - 1)$, the i th row of the k - t -matrix is said to be a positive row, if all its elements have a plus sign. That is iff $f(\mathbf{p}_j) \geq f(\mathbf{p}_i) \forall j$.

Definition 5. Given an integer $1 \leq k \leq (N - 1)$, a sampling point \mathbf{p}_i has a positive reference in the k - t -matrix, if there exists $j \neq i$ such that (a) the j^{th} row of the k - t -matrix is a positive row and (b) the number $+i$ is an element of this j^{th} row.

Definition 6. Given an integer $1 \leq k \leq (N - 1)$, the sample point \mathbf{p}_i is called a local minimiser of f in the k^+ -topograph if the i^{th} row of the k - t -matrix is a positive row.

Definition 7. Given an integer $1 \leq k \leq (N - 1)$, the sample point \mathbf{p}_i is a global minimiser of f in the k^+ -topograph if \mathbf{p}_i is a local minimiser of f in the k^+ -topograph and, in addition, \mathbf{p}_i has no positive references in the k - t -matrix.

The following propositions can be readily demonstrated to show the consistency of the aforementioned definitions ?.

Proposition 1. Given an integer $1 \leq k \leq (N - 1)$, the sample point \mathbf{p}_i is a global minimiser of f in the k^+ -topograph if and only if the sample point \mathbf{p}_i is the only minimiser of f in the k^+ -topograph which is global.

Proposition 2. Given an integer $1 \leq k \leq (N - 1)$, then the i th row of k - t -matrix is the only positive row of this matrix if and only if the sample point \mathbf{p}_i is the only minimiser of f in the k^+ -topograph which is global.

Corollary 1. Given an integer $1 \leq k \leq (N - 1)$, if the sample point \mathbf{p}_i is the only local minimiser of f in the k^+ -topograph, then \mathbf{p}_i is a global minimiser of f in this graph.

In this publication we will use the paradigm that all local minimisers of f in the k^+ -topograph will be used for the local search (Paradigm 2.2 in ?). As described in ? the number of local minimisers of f in the k^+ -topograph is greater than or equal to number of global minimisers in the topograph. We will therefore employ the following definition

Definition 8. *Given an integer $1 \leq k \leq (N - 1)$, the minimiser pool \mathcal{M}^k is the set containing all local minimisers \mathbf{p}_i in the in the k^+ -topograph. The total number of starting points used in the local search step is equal to the cardinality of the minimiser pool $|\mathcal{M}^k|$.*

The entire point of using k - t -matrices is because a t -matrix will always have at most one local (and thus global) minimiser. This is undesirable since this sampling point is not necessarily the starting point closest to the true global minimum of the objective function. ? developed a semi-empirical formula producing an integer value k_c which is used as an estimate for the optimal value for the integer k .

2.3 Step 3: Local minimisation

Each of the minimisers from the k_c - t -topograph is now used as a starting point in a local minimisation routine. The resulting minima are used to find the global minimum. Conceivably various local optimisation routines can be used to address a broad class of optimisations problems. For problems with non-linear inequality constraints ? used the feasible direction interior-point method proposed by ? minimising the objective function f subject to the set of inequality constraint functions \mathbf{g} using the minimiser set as the initial starting points for the algorithm. An algorithm used to solve the feasible direction interior-point method using the set of starting points calculated in step 2 is presented in detail by ?.

In this publication we will mainly be using the sequential least squares quadratic programming optimisation algorithm (SLSQP) contained in the SciPy library originally developed by Kraft ??. Our Python implementation of the TGO algorithm published under an open source licence uses this algorithm as implemented in the SciPy library ??.

CHAPTER 3

Motivation and a one-dimensional prelude

In this section we will demonstrate how the Euclidean distance criterion in the TGO method disregards useful information about the (approximate) geometry of the objective function and we show how known information can be used effectively both in global optimisation and in mapping the local minima of objective functions as efficiently as possible. We also show how two important hyperparameters used by TGO, namely the number of sampling points N and the choice of k can be iteratively selected by intelligently exploiting information known from the topograph. This draws parallels to other works on iterative versions of TGO (I-TGO) ? trying to extract information from black-box objective functions. The informal, but intuitive ideas developed here will later be extended more rigorously to higher dimensional surfaces. Note that from Equation (1.5) Ω is always a compact space, this fact is important in several proofs used in this Section.

Example 1 Consider the following objective function

$$\min_x f(x) = \frac{\sin(x)}{x}, \quad x \in \Omega = [1, 20] \quad (3.1)$$

In this instance of the bounded optimisation problem there are 3 local minima which we will try to map in as few function evaluations as possible.

Following the TGO procedure we start by generating low-discrepancy sampling points. The first $N = 10$ points in the 1-dimensional Sobol sequence is given by $\mathcal{P} = \{p_1 = 1.0, p_2 = 10.5, p_3 = 15.25, p_4 = 5.75, p_5 = 8.125, p_6 = 17.625, p_7 = 12.875, p_8 = 3.375, p_9 = 4.5625, p_{10} = 14.0625\} \subset \Omega$. After mapping the objective function at the set of sampling

points

$$f : \begin{bmatrix} p_1 = 1.0 \\ p_2 = 10.5 \\ p_3 = 15.25 \\ p_4 = 5.75 \\ p_5 = 8.125 \\ p_6 = 17.625 \\ p_7 = 12.875 \\ p_8 = 3.375 \\ p_9 = 4.5625 \\ p_{10} = 14.0625. \end{bmatrix} \rightarrow \begin{bmatrix} f_1 = 0.84147 \\ f_2 = -0.08378 \\ f_3 = 0.02899 \\ f_4 = -0.08840 \\ f_5 = 0.11858 \\ f_6 = -0.05337 \\ f_7 = 0.02359 \\ f_8 = -0.06853 \\ f_9 = -0.21672 \\ f_{10} = 0.07091 \end{bmatrix} \quad (3.2)$$

the corresponding topograph is constructed

$$\begin{bmatrix} [c|ccccccccc]p_1 & -p_8 & -p_9 & -p_4 & -p_5 & -p_2 & -p_7 & -p_{10} & -p_3 & -p_6 \\ p_2 & +p_5 & +p_7 & +p_{10} & +p_3 & -p_4 & -p_9 & +p_6 & +p_8 & +p_1 \\ p_3 & +p_{10} & -p_6 & -p_7 & -p_2 & +p_5 & -p_4 & -p_9 & -p_8 & +p_1 \\ p_4 & -p_9 & +p_5 & +p_8 & +p_1 & +p_2 & +p_7 & +p_{10} & +p_3 & +p_6 \\ p_5 & -p_2 & -p_4 & -p_9 & -p_7 & -p_8 & -p_{10} & +p_1 & -p_3 & -p_6 \\ p_6 & +p_3 & +p_{10} & +p_7 & -p_2 & +p_5 & -p_4 & -p_9 & -p_8 & +p_1 \\ p_7 & +p_{10} & -p_2 & +p_3 & +p_5 & -p_6 & -p_4 & -p_9 & -p_8 & +p_1 \\ p_8 & -p_9 & +p_1 & -p_4 & +p_5 & -p_2 & +p_7 & +p_{10} & +p_3 & +p_6 \\ p_9 & +p_4 & +p_8 & +p_1 & +p_5 & +p_2 & +p_7 & +p_{10} & +p_3 & +p_6 \\ p_{10} & -p_3 & -p_7 & -p_2 & -p_6 & +p_5 & -p_4 & -p_9 & -p_8 & +p_1 \end{bmatrix} \quad (3.3)$$

The sampling points together with the objective function evaluations are plotted in Figure 3.1. Using the empirical relation from ? the optimal k_c is calculated at $k_c = 8$. Using 6 we find that the resulting 8- t -matrix has only one minimiser; the global minimiser at $p_9 = 4.5625$. For the local minimisation we use the SLSQP method as implemented in the function `scipy.optimize.minimize` ? to find the approximate global minimum at $x = 4.4934$.

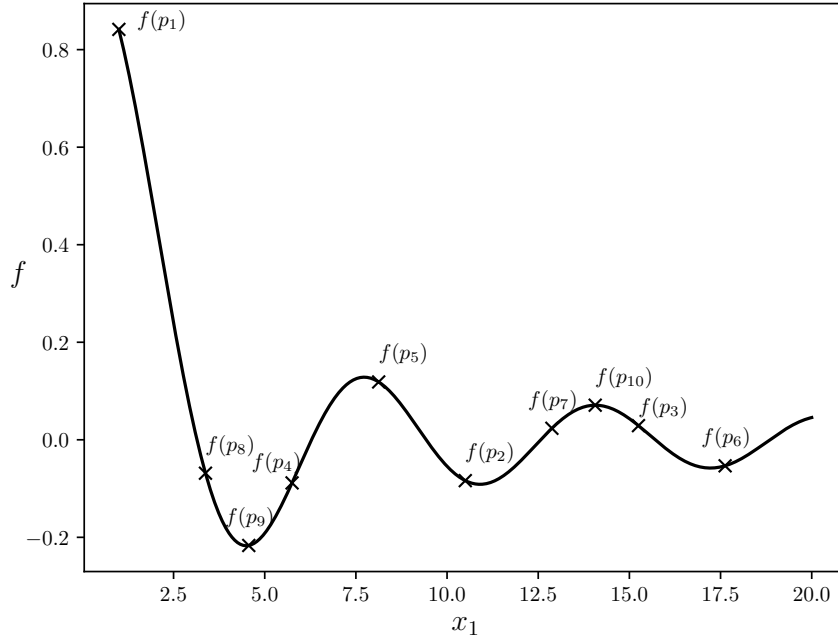


Figure 3.1: Test function give by Equation (3.1) with 10 Sobol sequenced sampling points

Observing Figure 3.1 it is immediately apparent that the set of 10 sampling points alone provides adequate information to deduce that there are at least 3 local minima. Observe that there are at least two other local minima since $f(p_5) < f(p_2) < f(p_7)$. So at least one local minimum exists in the domain $(p_5, p_7) \subset \mathbb{R}$ since between p_5 and p_2 we must have, by the mean value theorem (MVT), $\frac{df}{dx} < 0$ for some domain $x \in [p_5, p_2] \subset \mathbb{R}$. Similarly for $x \in [p_2, p_7] \subset \mathbb{R}$ we have by MVT $\frac{df}{dx} > 0$. Since f is a smooth, continuous function for $x \in (0, \infty)$ there must exist at least one stationary point $x \in (p_5, p_7) \subset \mathbb{R}$ where $\frac{df}{dx} = 0$. Furthermore we observe $f(p_6) < f(p_3)$ indicating another minimum in the domain $x \in (p_3, 20] \subset \mathbb{R}$ since the minimum must be either on the boundary or in $x \in (p_3, 20] \subset \mathbb{R}$ by the same argument as above.

The empirical relation by ? was mainly developed for the purpose of finding the global minimum. Therefore if only 10 sampling points are available, then to find more local minima using the TGO method is required to force a lower k value. Alternatively, since k_c is a function of N , simply sampling more points is sufficient to find all the local minima using Henderson's formula for this test problem. For example at $N = 16$ all 3 local minima are produced by TGO with Henderson's formula. Figure 3.2 shows the number of minimisers found at different k values for this example. The maximum minimiser set (other than using every sampling point as a starting point) can be trivially extracted by setting $k = 1$ and calculating $|\mathcal{M}^1|$. However, in this Example it leads to more starting points than optimal since at least two minimisers will be in the same convex basin domain and therefore converge to same minimum in the local minimisation

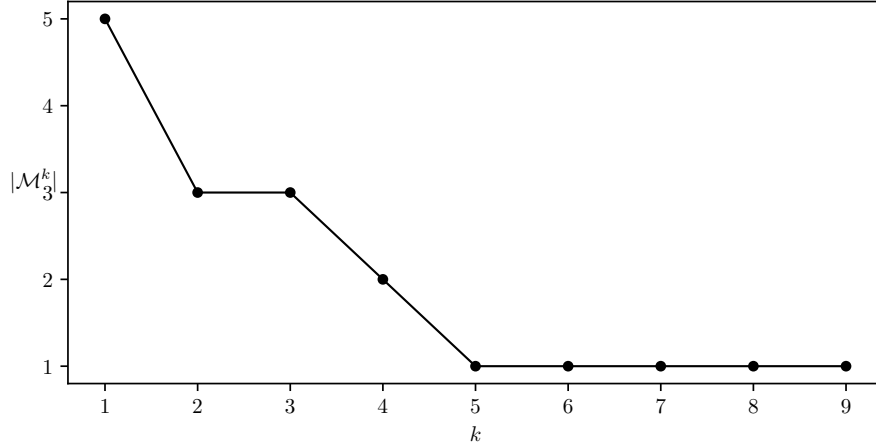


Figure 3.2: Number of minimisers $|\mathcal{M}^k|$ found using the TGO method for different k values at $N = 10$

step. This results in superfluous function evaluations without extracting more useful information from the objective function.

This idea drives the motivation behind the following definition.

Definition 9. For a given set \mathcal{P} of N sampling points, k_{opt} is any integer $1 \leq k \leq (N - 1)$ that will produce the optimal minimiser set $\mathcal{M}^{k_{opt}}$ containing the maximum set of minimisers such that no two starting points extracted from $\mathcal{M}^{k_{opt}}$ will lead to the same minimum in the local optimisation step for some tolerance ϵ . In other words every element contained in $\mathcal{M}^{k_{opt}}$ should lie in a unique locally convex sub-domain.

Note that for a given N , $\mathcal{M}^{k_{opt}}$ might not produce all the true local minima of an objective function. What's important is that, given the information known from the sampling, the maximum number of local minima are found. In addition, no function evaluations are wasted in the local minimisation step which lead to the same minimum.

In Example 1 for $N = 10$ the optimal k values are $k_{opt} = \{2, 3\}$ which will produce 3 minimisers $|\mathcal{M}^2| = |\mathcal{M}^3| = 3$. We will now show that these lower k values carry unexploited information on the best approximate geometry of the objective function. For example in Figure 3.3 we plot the $|\mathcal{M}^k|$ values corresponding to the set $k = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$ for every sampling point range $N \in [2, 50]$.

From Figure 3.3 we notice the special property of $k = 3$ for one dimensional objective functions sampled with the Sobol sequence.

Firstly, for a lower number of sampling points N it provides a higher number of starting minimisers than $k > 3$. Note that by inspection of Definition 6 it can be determined that any $k > 3$ value will always produce an equal or lower number of minimisers than $k = 3$ (this is also true for any $k > i$). When adding columns to a positive row there are only two possibilities: the next sampling point in the row can either have a positive or a negative sign. All other elements in the row have a positive sign by definition (see Definition 6).

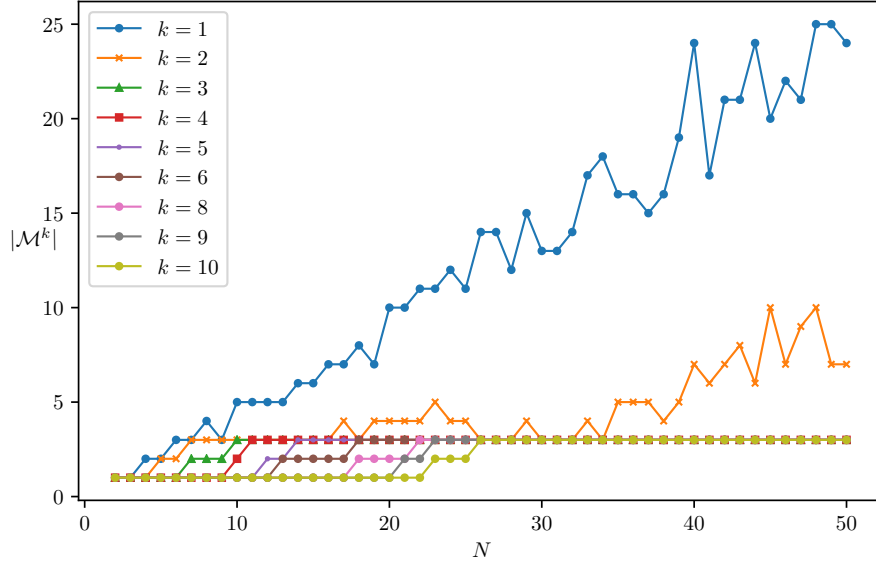


Figure 3.3: Number of minimisers $|\mathcal{M}^k|$ found using the TGO method for the given k values at various sampling points N

If the next sampling point in the row has a positive sign then the row will just remain a positive row and the number of minimisers remain the same. If the point is a negative reference point then the row will no longer be a positive row and thus the point is no longer a minimiser, lowering the total.

Secondly it can be observed that $k = 3$ never calculates a number of starting minimisers higher than optimal unlike $k < 3$. Therefore by using $k = 3$ in Example 1 TGO will always find as many minimisers in as few sampling¹ function evaluations as possible and furthermore all local minima will be found when $N \geq 10$. It should be noted that the total number of function evaluations depends on the particular local minimisation algorithm used. However, it is apparent that each minimiser starting point is in a unique locally convex domain. It is tempting for an optimisation practitioner to use the size of the set of minimisers $|\mathcal{M}^3|$ as a stopping criterion for iterative sampling N of one dimensional objective functions. The practical usefulness of this idea can be demonstrated with the following example:

Example 2 The following instance of the optimisation problem has 13 local minima in the given domain

$$\min_x f(x) = -x \sin(x), \quad x \in \Omega = [1, 80] \quad (3.4)$$

From Figure 3.4 we can deduce that the minimum number of sampling points required for $k = 3$ to find all local minima using the Sobol sequence is $N = 40$, this sampling is shown in Figure 3.5. If $N < 40$ then there aren't enough sampling points to deduce

¹not necessarily total function evaluations since starting points closer to the local minima may provide better performance for a given local minimisation routines

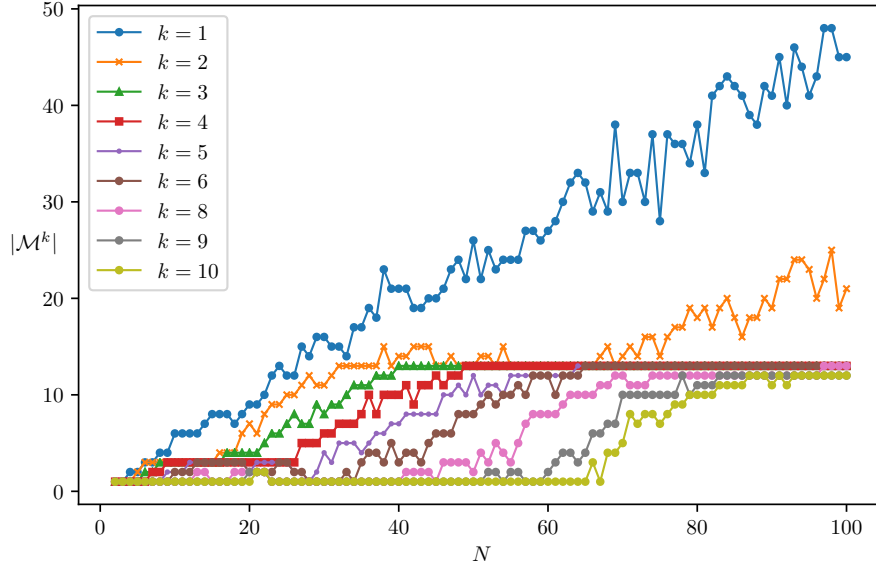


Figure 3.4: Number of minimisers $|\mathcal{M}^k|$ found using the TGO method for the given k values at various sampling points N

that there are at least 13 locally convex domains from using the same arguments as in Example 1. Note for example that if we used a sequence that skipped p_1 then $N = 39$ would be adequate since $l = 1 < p_{32} < p_{33}$. Using our Python implementation of TGO ? with $N = 40$ all 13 local minima of the objective function were found in a total of 285 function evaluations.

An example of a stopping criterion would be to stop sampling if $|\mathcal{M}^3|$ is unchanged after, say, 10 sampling point evaluations. The rate at which the number of elements in $|\mathcal{M}^3|$ grows with increasing N also provides a heuristic for characterising the multimodality and the geometry of the objective function. Objective functions that have a large number of local minima in a small domain (and relatively fewer minima in other larger domains) will have a much smaller growth in $|\mathcal{M}^3|$ for a given low-discrepancy sampling. This idea of continuously classifying and extracting approximate function characteristic information from the sampling points will be formalised and extended to higher dimensions in chapter 4.

There is a simple reason why the 3- t -matrix has this quality in the first dimension for the optimisation problem given in Equation (3.1). However, it is not guaranteed that this property holds for any sampling point distribution. In fact it holds true only under the following conditions:

1. Consider all points in the ordered sampling set from the smallest to greatest x value $\mathcal{P} = \{p_i \mid p_0 < p_1 < p_2 \dots < p_N - 1, p_i \in (x_l, x_u)\}$, excluding the supremum and infimum.
2. For any given point p_i the Euclidean distance between p_i and 2 of its nearest sam-

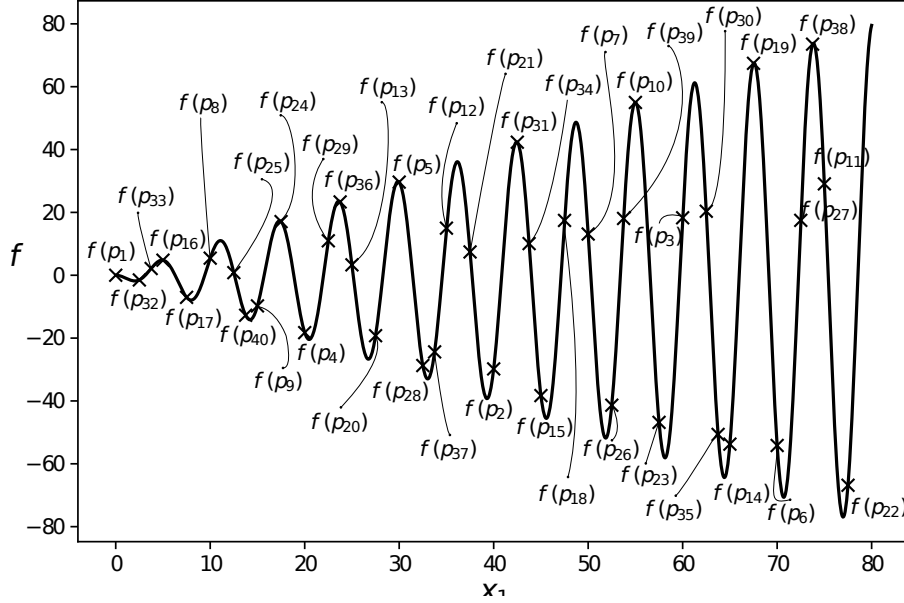


Figure 3.5: Plot of the objective function in Example 2 for $N = 40$ sampling points

pling points $p_{i-1} < p_i < p_{i+1}$ should be less than the relative difference between p_i and a fourth point in the sampling sequence $|p_i - p_j|$ where $j \neq i, i-1, i+1$.

In fact it is easy to prove both that for a locally, strictly convex domain of f the 3-topograph construction can produce a larger minimiser pool \mathcal{M}^3 than optimal. It can also be shown that a construction must exist where the optimal number of minimisers will *always* be extracted regardless of the sampling distribution. Furthermore it can be shown that at most 3 sampling points within a locally convex domain $x \in [x_l, x_u]$ is required to produce enough information so that only one minimiser in the domain is produced.

Theorem 1. *There exists a 1-dimensional sampling sequence such that $k = 3$ will produce a minimiser pool larger than optimal as defined by Definition 9.*

Proof. Consider a subdomain $x \in [x_l, x_u] \subset \mathbb{R}$ for which f is strictly convex. We define the set of N sampling points \mathcal{P} ordered in such a way that

$$\mathcal{P} = \{p_i \mid p_0 < p_1 < p_2 < \dots < p_{N-1}, p_i \in (x_l, x_u)\}$$

Let $\mathcal{F} = \{f_0, f_1, f_2, \dots, f_{N-1}\}$ be set of one-to-one function values corresponding to the points mapped by $f : \mathcal{P} \rightarrow \mathcal{F}$.

Suppose we have $f_1 < f_0$ and $f_1 < f_2 < f_3, \dots, f_{N-1}$. By construction we have $|p_1 - p_2| < |p_1 - p_3| < |p_1 - p_4| < |p_1 - p_5|$ then by the Definitions 4, 5 and 6 p_2 is a minimiser of the 3-t-topograph. Suppose we have a sampling distribution such that $|p_2 - p_3| < |p_1 - p_2|, |p_2 - p_4| < |p_1 - p_2|$ and $|p_2 - p_5| < |p_1 - p_2|$ then by the definitions 4, 5 and 6 p_3 is also a minimiser of the 3-t-topograph. Therefore more than two minimisers

are produced in the same locally convex sub-domain of $[x_l, x_u]$. We have shown that \mathcal{M}^3 can produce a minimiser pool larger than optimal which concludes the proof.

Lemma 1. *A construction exists that will always produce a minimiser pool larger than optimal as defined by Definition 9 for any given 1-dimensional sampling sequence.*

Now suppose that instead of using only the Euclidean distance metric we also invoke knowledge of the nearest point in every cartesian direction. We use the criterion that a minimiser point p_i is a minimiser iff with the ordering constructed in \mathcal{P} and \mathcal{F} we have $f_i < f_{i-1}$ and $f_i < f_{i+1}$. With this definition if the point p_i is a minimiser then no other point meets the criterion since by construction of the sampling in the locally convex domain $f_0 > f_1 > \dots > f_{i-1} > f_i$ and $f_{i+1} < f_{i+2} < f_{i+3} < \dots < f_{N-1}$. This proves Lemma 1.

Finally note that only information from the 3 points in the locally convex sub-domain of $[x_l, x_u]$ and their corresponding function values f_{i-1} , f_i and f_{i+1} are needed to produce a minimiser using this criterion. \square

An important consequence here is that for low discrepancy sequences in higher dimensions and for less well behaved objective functions the topographs connected with the Euclidean distance metrics will similarly discard available information about the local geometry. This produces larger than optimal minimiser pools leading to very high numbers of function evaluations needed to solve the problem.

In the following section we will develop a more efficient algorithm that will make use of this information. SHGO will always produce equivalent results to this algorithm in the one dimensional case.

3.1 Axially directed topograph

Based on the observations from chapter 3 we develop an algorithm that, for a given sampling set, always uses the optimal number of starting minimisers as defined for one dimensional objective functions without requiring *a-priori* specification of the k parameter. Here a new graph structure is proposed and attempts are made to directly extend the idea to higher dimensions by connecting every vertex to the nearest vertex in every cartesian axis direction. In Theorem 2 we show that the one dimensional properties of this algorithm does not extend to higher dimensions which finally leads us to the built up complexes in chapter 4. The main conclusion of this section is that simpler graph structures cannot be used to find locally convex sub-domains of a function in the same way that was accomplished in chapter 3.

The algorithm proceeds in the same way as TGO described in chapter 2 except for step 2 where a new structure described in Section 4.1 replaces the topograph.

3.1.1 Axially directed topograph

Let \mathcal{F} be the set of scalar outputs mapped by the objective function $f : \mathcal{P} \rightarrow \mathcal{F}$ for a given sampling set $\mathcal{P} \subseteq \Omega \subseteq \mathbb{R}^n$. The scalar elements $f_i \in \mathcal{F}$ have one-to-one correspondence with the vector elements $\mathbf{p}^i \in \mathcal{P}$ where the integer $i \in \{1, 2, 3, \dots, N\}$ indicates the sampling point index. The vector \mathbf{p}^i in turn has dimensional elements x_j^i where the integer $j \in \{1, 2, 3, \dots, n\}$ indicates the dimension of the scalar value $\forall i(x_1^i, x_2^i, x_3^i, \dots, x_n^i) \in \mathbf{p}^i$.

We wish to construct a graph that is ordered along the coordinate axes, this is done by formally defining the following related partially ordered sets.

Definition 10. *Given a finite structured set of N feasible ordered sampling points $\mathcal{P} = (\mathbf{p}^1, \mathbf{p}^2, \dots, \mathbf{p}^N)$ with its corresponding objective function outputs $\mathcal{F} = (f^1, f^2, \dots, f^N)$, the index set of \mathcal{P} is given as the ordered set $\mathcal{I} = (i = \{1, 2, 3, \dots, N\}, \leq)$*

Note that the initial ordering of the index set is arbitrary, what's important is that an ordered index set is defined. This ordering will allow us to keep track of any vertex in the graph to its corresponding sampling point in \mathcal{P} so that the corresponding objective function only needs to be evaluated once. Herein the order is taken as the order that is generated by the Sobol sequence.

Definition 11. *Given a set of feasible sampling points $\mathcal{P} \subseteq \Omega \subseteq \mathbb{R}^n$ define X_j for every dimension $j \in \{1, 2, 3, \dots, n\}$ as the partially ordered set $X_j = \{\mathbf{p}^i \mid \forall i(x_j^i < x_j^{i+1})\}$.*

The definition is demonstrated with the following numerical example:

Example 3 Given set of the first 5 points in the 2-dimensional Sobol sequence bounded by the 2-cube:

$$\mathcal{P} = ((0, 0), (0.5, 0.5), (0.75, 0.25), (0.25, 0.75), (0.375, 0.375)) \subseteq [0, 1] \times [0, 1] \subseteq \mathbb{R}^2$$

let $f(x) = x_1^2 + x_2^2$ so that

$$\mathcal{F} = (0, 0.5, 0.625, 0.625, 0.28125)$$

then

$$X_1 = ((0, 0), (0.25, 0.75), (0.375, 0.375), (0.5, 0.5), (0.75, 0.25))$$

and

$$X_2 = ((0, 0), (0.75, 0.25), (0.375, 0.375), (0.5, 0.5), (0.25, 0.75))$$

The corresponding index sets are $\mathcal{I}_1 = (1, 4, 5, 2, 3)$ and $\mathcal{I}_2 = (1, 3, 5, 2, 4)$.

Definition 12. For every dimension j , \mathcal{F}_j is the partially ordered set such that the position of the elements X_j correspond to the original index sampling of \mathcal{P} , $\mathcal{F}_j = \{f_j^{i,k} \mid \forall i(x_j^i < x_j^{i+1}), f_j^{i,k} = f_k \in \mathcal{F}, k \subseteq \mathcal{I}\}$

That is the first superscript i of the elements $f_j^{i,k}$ indicate the ordering in \mathcal{F}_j , while the second superscript k indicates the corresponding scalar value of $f_j^{i,k}$ in \mathcal{F} . Ordering the example we have $\mathcal{F}_1 = (0, 0.625, 0.28125, 0.5, 0.625)$ and $\mathcal{F}_2 = (0, 0.625, 0.28125, 0.5, 0.625)$.

Definition 13. For every dimension j , define the partially ordered sets of cardinality N such that $\mathcal{F}_j^+ = \{f_j^{i,k} - f_j^{i-1,k} \mid \forall i(x_j^i < x_j^{i+1}), f_j^{i,k} = f_k \in \mathcal{F}, i = \{1, 2, \dots, N, k \subseteq \mathcal{I}\}\}$ and $\mathcal{F}_j^- = \{f_j^{i,k} - f_j^{i+1,k} \mid \forall i(x_j^i < x_j^{i+1}), f_j^{i,k} = f_k \in \mathcal{F}, i = \{0, 1, \dots, N-1\}, k \subseteq \mathcal{I}\}$

These sets are essentially objective function differences between the sampling points along each dimensional Cartesian axis. Continuing from the numerical example we have

$$\mathcal{F}_1^+ = (0.625, -0.34375, 0.21875, 0.125)$$

$$\mathcal{F}_2^+ = (-0.625, 0.34375, -0.21875, -0.125)$$

$$\mathcal{F}_1^- = (0.625, -0.34375, 0.21875, 0.125)$$

$$\mathcal{F}_2^- = (-0.625, 0.34375, -0.21875, -0.125)$$

We denote the elements as $f_j^{+,i,k} \in \mathcal{F}_j^+$ and $f_j^{-,i,k} \in \mathcal{F}_j^-$ for every dimension $j \in \{1, 2, 3, \dots, n\}$, cartesian ordering $i \subseteq \mathcal{I}$ and corresponding sampling point $k \in \mathcal{I}$. The usefulness of these abstract constructions is apparent in the following definition.

Definition 14. For a given sampling set \mathcal{P} . The minimiser pool \mathcal{M} is defined as

$$\mathcal{M} = \mathcal{M}_c \cup \mathcal{M}_{lb} \cup \mathcal{M}_{ub}$$

where

$$\begin{aligned}\mathcal{M}_c &= \left\{ \mathbf{p}^i \mid \forall j \left((f_j^{+i} > 0) \wedge (f_j^{-(i+1)} > 0) \right), i = \{1, 2, 3, \dots, N-1\} \right\} \\ \mathcal{M}_{lb} &= \left\{ \mathbf{p}^i \mid \forall j \left(f_j^{-i} < 0 \right), i = \{0\} \right\} \\ \mathcal{M}_{ub} &= \left\{ \mathbf{p}^i \mid \forall j \left(f_j^{+i} < 0 \right), i = \{N\} \right\}\end{aligned}$$

That is, we simply check the finite difference between sampling points in every cartesian direction. In addition we check if the sampling points on the boundaries are minimisers.

Theorem 2. The minimiser pool \mathcal{M} from 14 always produces a set that is either smaller than or equal to the optimum minimiser pool as defined by 9 iff $j = 1$.

Proof. The proof for $j = 1$ follows the same argument from chapter 3. By Definition 10, 11 and 12 we have the ordering constructed as \mathcal{P} and \mathcal{F}_1 . If a given point \mathbf{p}^i is a minimiser with $f_1^{+i} > 0$ and $f_1^{-i} > 0$, then we have by Definition 13 $f^i < f^{i-1}$ and $f^i < f^{i+1}$, conversely if a given point \mathbf{p}^i is not a minimiser then either $f_1^{+i} < 0$ or $f_1^{-i} > 0$ so that regardless of the sampling method used and the Euclidean distance between points a minimiser will never be generated for any point that has $((f^i > f^{i-1}) \wedge (f^i > f^{i+1})) \vee ((f^i < f^{i-1}) \wedge (f^i < f^{i+1}))$.

If $j > 1$ we have no such guarantee for a higher dimensional locally convex domain. As a counter example consider the set of points

$$\mathcal{P} = ((0, 0), (0.25, 0.25), (0.75, 0.125), (0.125, 0.75))$$

on the same function as above, the minimiser set produced is $\mathcal{M} = \{(0, 0), (0.25, 0.25)\}$ which is clearly larger than optimal and will produce the same local minimum. \square

This unsatisfactory result for higher dimensions could still potentially show good performance for more regular spaced sampling such as grids, however, as we will see in the next section the SHGO algorithm can guarantee that the optimal minimiser set will be produced for any dimension.

3.1.2 Implementation

Algorithm 1 provides a high-level overview of the ATGO algorithm. A Python implementation of this algorithm can be found in ?.

Algorithm 1 ATGO algorithm

```

1: procedure INITIALISATION
2:   Input an objective function  $f$ , constraint functions  $\mathbf{g}$  and variable bounds and
    $[\mathbf{l}, \mathbf{u}]^n$ .
3:   Input  $N$  initial sampling points.
4:   Define a sampling sequence that generates a set  $\mathcal{X}$  of sampling points in the unit
   hypercube space  $[\mathbf{0}, \mathbf{1}]^n$ 
5: end procedure
6: procedure INITIAL SAMPLING
7:    $\mathcal{P} = \emptyset$ 
8:   while  $|\mathcal{P}| < N$  do
9:     Generate  $N - |\mathcal{P}|$  sequential sampling points  $\mathcal{X} \subset \mathbb{R}^n$ 
10:    Stretch  $\mathcal{X}$  over the lower and upper bounds  $[\mathbf{l}, \mathbf{u}]^n$ 
11:     $\mathcal{P} = \{\mathcal{X}_i \mid \mathbf{g}(\mathcal{X}_i) \geq 0, \forall \mathcal{X}_i \in \mathcal{X}\} \cup \mathcal{P}$   $\triangleright$  (Find  $\mathcal{P}$  in the feasible subset  $\Omega$ 
    by discarding any points mapped outside the linear constraints  $g$  and adding to the
    current set of  $\mathcal{P}$ .)
12:    Set  $\mathcal{X} = \emptyset$ 
13:   end while
14:   Find  $\mathcal{F}$  from the objective function  $f : \mathcal{P} \rightarrow \mathcal{F}$ 
15: end procedure
16: procedure CONSTRUCT  $\mathcal{M}$ 
17:   Calculate  $\mathcal{M}$  from the sets  $\mathcal{P}$  and  $\mathcal{F}$  using Definitions 11 through 14.
18: end procedure
19: procedure LOCAL MINIMISATION
20:   Calculate the approximate local minima of  $f$  using a local minimisation routine
   with the elements of  $\mathcal{M}$  as starting points.  $\triangleright$  These local minimisations can be
   performed in parallel.
21: end procedure
22: procedure PROCESS RETURN OBJECTS
23:   Order the final outputs of the minima of  $f$  found in the local minimisation step
   to find the approximate global minimum.
24: end procedure
25:
26: return the approximate global minimum and a list of all the minima found in the
   local minimisation step.

```

CHAPTER 4

Simplicial Homology Global Optimisation

4.1 Overview

The SHGO method strongly relies on constructing a simplicial complex using the sampled points of an objective function f as vertices. From this construction of the complex \mathcal{H} we use the resulting directed subgraph which contains the set of all 1-chains from the elements of $\mathcal{H}^1 \in \mathcal{H}$ to find minimiser pools using definitions similar to the methods demonstrated in the previous sections. This is accomplished by the application of Sperner's lemma ? allowing us to approximate the domains of stationary points for any objective function in the feasible search space Ω .

We prove that, if provided with an adequate sampling set, the construction of \mathcal{H} will produce the same homology groups. We use this result to show that for the given sampling set of vertices $\mathcal{H}^0 \in \mathcal{H}$ we always extract the optimal minimiser pool similar to the one-dimensional case described in chapter 3, but extended to higher dimensions.

The algorithm itself consists of four steps which will be described in detail:

1. Uniform sampling point generation of N vertices in the search space within the bounded and constrained subspace of Ω from which the 0-chains of \mathcal{H}^0 are constructed.
2. Construction of the directed simplicial complex \mathcal{H} by triangulation of the vertices.
3. Construction of the minimiser pool $\mathcal{M} \subset \mathcal{H}^0$ by repeated application of Sperner's lemma.
4. Local minimisation using the starting points defined in \mathcal{M} .

We will start by formally defining the construction of \mathcal{H} from a given set of feasible sampling points \mathcal{P} and proving its properties.

4.2 Directed simplicial complex approximation of the objective function

Consider again the general objective function mapping in the continuous domain $f : \mathbb{R}^n \rightarrow \mathbb{R}$. The purpose of this section is describe a discrete mapping $h : \mathcal{P} \rightarrow \mathcal{H}$ to provide a simplicial approximation for the surface of f . To guide the reader the methods will be demonstrated on the simple 2-dimensional optimisation problem defined in Example 4. The use of a 2-dimensional surface allows a demonstration of the techniques while the abstractions defined are readily extended to higher dimensions.

We start by formally defining the set of vertices from which 0-chains of the simplicial complex are built and the of edges from which the 1-chains of \mathcal{H} are built.

Definition 15. *Let \mathcal{X} be the set of sampling points generated by a sampling sequence in the bounded hyperrectangle $[\mathbf{l}, \mathbf{u}]^n$. The set $\mathcal{P} = \{\mathbf{x} \in \mathcal{X} \mid \mathbf{g}(\mathbf{x}) \geq 0\}$ is a set of points within the feasible set Ω .*

Definition 16. *For an objective function f , \mathcal{F} is the set of scalar outputs mapped by the objective function $f : \mathcal{P} \rightarrow \mathcal{F}$ for a given sampling set $\mathcal{P} \subseteq \Omega \subseteq \mathbb{R}^n$.*

Definition 17. *Let \mathcal{H} be a directed simplicial complex. Then $\mathcal{H}^0 := \mathcal{P}$ is the set of all vertices of \mathcal{H} .*

Definition 18. *For a given set of vertices \mathcal{H}^0 , the simplicial complex \mathcal{H} is constructed by a triangulation connecting every vertex in \mathcal{H}^0 . The triangulation supplies a set of undirected edges E .*

Definition 19. *The set \mathcal{H}^1 is constructed by directing every edge in E . A vertex $v_i \in \mathcal{H}^0$ is the connected to another vertex v_j by an edge contained in E . The edge is directed as $\overline{v_i v_j}$ from v_i to v_j iff $f(v_i) < f(v_j)$ so that $\partial(\overline{v_i v_j}) = v_j - v_i$. Similarly an edge is directed as $\overline{v_j v_i}$ from v_j to v_i iff $f(v_i) > f(v_j)$ so that $\partial(\overline{v_j v_i}) = v_i - v_j$.*

For practical computational reasons we must also consider the case where $f(v_i) = f(v_j)$. If neither v_i or v_j is already a minimiser we will make use of rule that the incidence direction of the connecting edge is always directed towards the vertex that was generated earliest by the sampling point sequence. If v_i is not connected to another vertex v_k then we leave the notation $\overline{v_i v_k}$ undefined and let $\partial(\overline{v_i v_k}) = 0$. We let the higher dimensional simplices of $\mathcal{H}^k, k = 2, 3, \dots, n+1$ be directed in any arbitrary direction which completes the construction of the complex $h : \mathcal{P} \rightarrow \mathcal{H}$. We can now use \mathcal{H} to find the minimiser pool for the local minimisation starting points used by the algorithm:

Definition 20. *A vertex v_i is a minimiser iff every edge connected to v_i is directed away from v_i , that is $\partial(\overline{v_i v_j}) = (v_{j \neq i} - v_i) \vee 0 \forall v_{j \neq i} \in \mathcal{H}^0$. The minimiser pool \mathcal{M} is the set of all minimisers.*

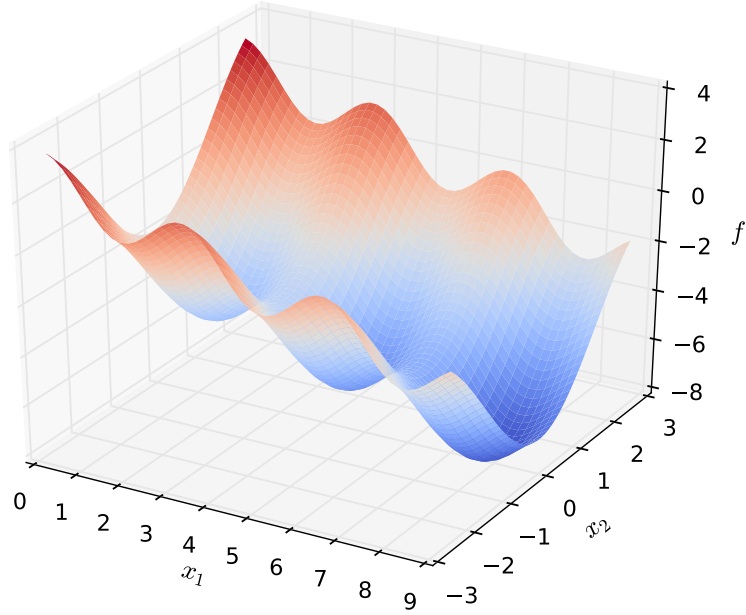


Figure 4.1: A 3-dimensional surface plot of the optimisation test function given in Example 4 $f(\mathbf{x}) = -\sin(2x_1 - 0.5\pi) - 3\cos(x_2) - 0.5x_1$ for the domain $\mathbf{x} \in \Omega = [0, 9] \times [-2.5, 2.5]$

We will also make extensive use of star notation ??:

Definition 21. *The star of a vertex v_i , written $st(v_i)$, is the set of points Q such that every simplex containing Q contains v_i .*

The k -chain $C(\mathcal{H}^k)$, $k = n + 1$ of simplices in $st(v_i)$ forms a boundary cycle $\partial(C(\mathcal{H}^{n+1}))$ with $\partial(\partial(C(\mathcal{H}^{n+1}))) = \emptyset$. The faces of $\partial(\mathcal{H}^{n+1})$ are the bounds of the domain defined by $st(v_i)$.

A visual demonstration of these constructions and notations is provided in the following numerical example:

Example 4 The Ursem01 function for two dimensions is defined as follows ?

$$\min f(\mathbf{x}) = -\sin(2x_1 - 0.5\pi) - 3\cos(x_2) - 0.5x_1, \quad x \in \Omega = [0, 9] \times [-2.5, 2.5]$$

Figure 4.1 provides a 3 dimensional plot of this function. The function has three local minima within the domain $\mathbf{x} \in [0, 9] \times [-2.5, 2.5]$.

We use a set \mathcal{P} of 15 sampling points from the 2-dimensional Sobol sequence. First

map out the objective function values:

$$f : \begin{bmatrix} [l]v_0 = (0.0, -2.5) \\ v_1 = (4.6, 0.0) \\ v_2 = (6.9, -1.25) \\ v_3 = (2.3, 1.25) \\ v_4 = (3.45, -0.625) \\ v_5 = (8.05, 1.875) \\ v_6 = (5.75, -1.875) \\ v_7 = (1.15, 0.625) \\ v_8 = (1.725, -0.9375) \\ v_9 = (6.325, 1.5625) \\ v_{10} = (8.625, -2.1875) \\ v_{11} = (4.025, 0.3125) \\ v_{12} = (2.875, -1.5625) \\ v_{13} = (7.475, 0.9375) \\ v_{14} = (5.175, -0.3125) \end{bmatrix} \rightarrow \begin{bmatrix} [l]f_0 = 3.403 \\ f_1 = -6.275 \\ f_2 = -4.0651 \\ f_3 = -2.208 \\ f_4 = -3.3429 \\ f_5 = -4.051 \\ f_6 = -1.493 \\ f_7 = -3.674 \\ f_8 = -3.591 \\ f_9 = -2.191 \\ f_{10} = -2.606 \\ f_{11} = -5.062 \\ f_{12} = -0.601 \\ f_{13} = -6.239 \\ f_{14} = -6.044 \end{bmatrix} \quad (4.1)$$

From 17 we find \mathcal{H}^0 from \mathcal{P} . Next we use Delaunay triangulation to find a set of connected edges according to 18. Any triangulation scheme resulting in a simplicial complex can be used. Next the edges are directed from the calculated values of \mathcal{F} using 19. Finally from 20 we find the minimiser set $\mathcal{M} = \{v_1, v_7, v_{13}\}$. The resulting structure is shown in Figure 4.2. Also shown in Figure 4.2 is the domain of $\text{st}(v_1)$ for a visual description of 21. Next we increase the sampling size to $N = 150$ points and repeat the procedure. The resulting complex is shown in Figure 4.3. Notice that while the minimiser vertices have changed (now closer to the true continuous local minima), the cardinality of the minimiser pool $|\mathcal{M}|$ remains unchanged. That is, given an adequate number sampling points $|\mathcal{M}|$ will cease to grow with increasing N , providing a heuristic for the number of sampling points needed to approximately map all minima of an objective function. This useful property of the SHGO algorithm is proven formally in section 4.4.

4.3 Guarantee of stationary points in sub-domains near minimiser points

This section is devoted to proving the following theorem:

Theorem 3. *Given a minimiser $v_i \in \mathcal{M} \subseteq \mathcal{H}^0$ on the surface of a continuous, Lipschitz smooth objective function f with a compact bounded domain in \mathbb{R}^n and range \mathbb{R} , there exists at least one stationary point of f within the domain defined by $\text{st}(v_i)$.*

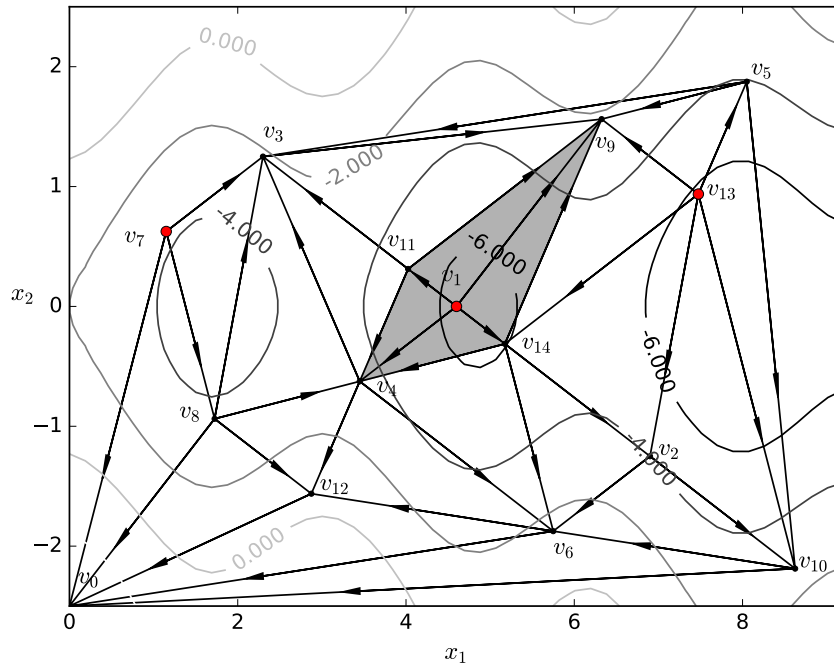


Figure 4.2: A directed complex \mathcal{H} forming a simplicial approximation for an objective function. There are three minimiser vertices v_1 , v_7 and v_{13} shown by the big red dots. The area shaded in grey represents the domain defined by $\text{st}(v_1)$

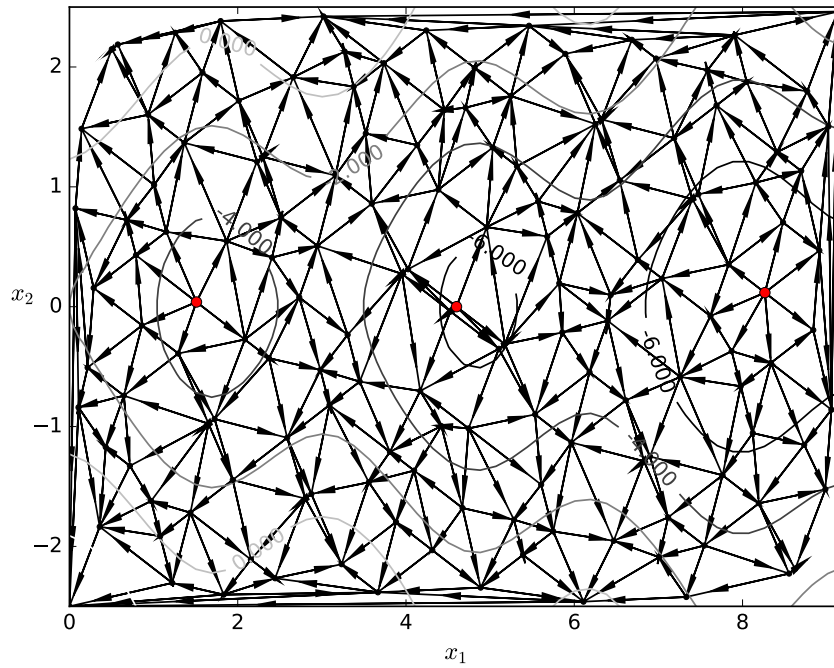


Figure 4.3: A directed complex \mathcal{H} forming a simplicial approximation for an objective function with 150 vertices. There are three minimiser vertices given by the big red dots

Proof. Our strategy relies on finding a simplex with a Sperner labelling where each label represents a different $n + 1$ label in every vector direction of the gradient vector field ∇f of f where of the $n + 1$ Cartesian directions we require only a vector pointing towards a section defined by $n + 1$ hyperplane cuts, the remainder of the proof then proceeds as usual for Brouwer's fixed point theorem ? found in for example ? : p. 40 utilising Sperner's lemma.

Theorem 4. (*Sperner's lemma (?)*) *Every Sperner labelling of a triangulation of a n -dimensional simplex contains a cell labelled with a complete set of labels: $1, 2, \dots, n+1$.*

Start with the observation that for any minimiser $v_i \in \mathcal{M} \subseteq \mathcal{H}^0$ we have by construction that for any vertex v_j with incidence on a connecting edge $\overline{v_i v_j}$ that $f(v_i) < f(v_j)$, so by the MVT there is at least one point on $\overline{v_i v_j}$ where ∇f points towards a Cartesian direction in a section that can receive a unique Sperner label. If we have $n + 1$ vertices with incidence on an edge $\overline{v_i v_j} \subseteq \mathcal{H}^1$ in every required Cartesian direction then we have a simplex within $\text{st}(v_i)$ with a Sperner labelling.

In the case where we do not have $n + 1$ vertices in every required section then by construction there is no vertex between v_i and the boundary of f defined by Ω in the required section. In the case where the constraint is not active and there exists at least one point v_k boundary where ∇f does not point towards the boundary and by the MVT v_k can receive a unique Sperner label from which we can construct a simplex within $\text{st}(v_i)$ with Sperner labelling.

Following the combinatorial version of Brouwer's fixed point theorem ? since ∇f is continuous and the domain $\text{st}(v_i)$ is compact we can produce a sequence of complete triangulations with arbitrarily small size in which the size of the simplices decreases toward zero. This sequence produces a sequence of vertices with gradients $\nabla f(V)$ pointing in every $n + 1$ direction. By continuity there is a vector $\nabla f(\mathbf{X})$ near the sequences, since the zero vector is the only vector pointing in all $n + 1$ directions we have a point \mathbf{X} bounded by the domain defined by $\text{st}(v_i)$ where $\nabla f(\mathbf{X}) = \bar{0}$. In the case where the constraint is active a local minimum lies on the constraint which is in the domain defined $\text{st}(v_i)$. This concludes the proof. \square

Figure 4.4 provides a visual demonstration of the proof using the complex from Example 4. Here we have divided the plane so that the 3 required directions are $[0, \frac{\pi}{2})$, $[\frac{\pi}{2}, \pi)$ and $[\pi, 2\pi)$. Note that this division is arbitrary and any $n + 1 = 3$ subdivisions can be chosen as long as all possible $n + 1 = 3$ directions can form a simplex in the space are covered. The three possible simplices are contained within the star domains of each minimiser $\text{st}(v_1)$, $\text{st}(v_7)$ and $\text{st}(v_{13})$.

First consider the minimiser v_{13} . There are three possible edges in $[\frac{\pi}{2}, \pi)$ on which a point exists that can be used as a vertex to receive a Sperner labelling for that direction namely $\overline{v_{13}v_{14}}$, $\overline{v_{13}v_2}$ and $\overline{v_{13}v_{10}}$. The only possible edges in the $[0, \frac{\pi}{2})$, $[\frac{\pi}{2}, \pi)$ directions are

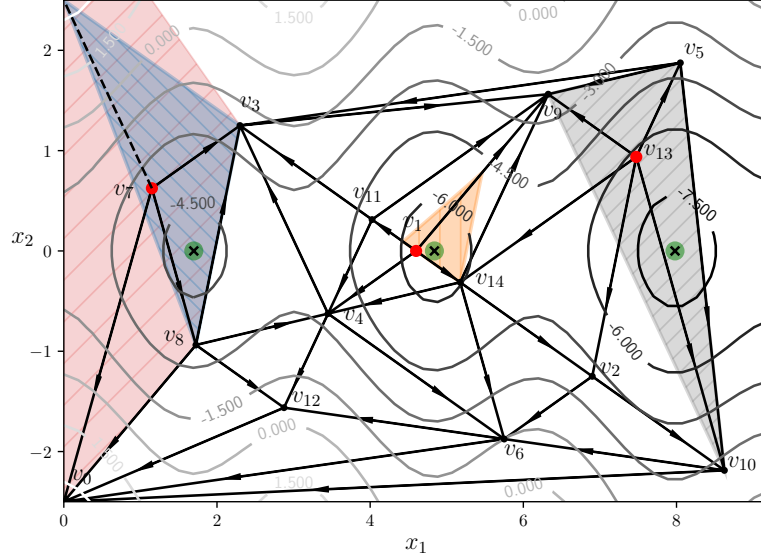


Figure 4.4: Visual demonstration of the proof by finding simplices with Sperner labellings. The three circled crosses are the (approximate) minimima of the objective function within the given bounds. The three possible Sperner simplices are contained within the star domains of each minimiser $\text{st}(v_1)$, $\text{st}(v_7)$ and $\text{st}(v_{13})$. v_7 is an example of simplices without complete Sperner labelings, the red shaded area around $\overline{v_7}$ is the bounded domain wherein at least one local minimum exist

$\overline{v_{13}v_5}$ and $\overline{v_{13}v_9}$ respectively. The simplex $\overline{v_5v_9v_{10}}$ drawn in Figure 4.4 is not necessarily the simplex with a Sperner labelling. The three vertices of the Sperner simplex which are proven to exist through the MVT exists on each of the edges $\overline{v_{13}v_{14}}$, $\overline{v_{13}v_2}$ and $\overline{v_{13}v_{10}}$ in a subdomain of this simplex $\overline{v_5v_9v_{10}}$. For example the simplex surrounding the minimiser v_1 is a possible Sperner simplex with vertices on the edges in every required direction.

Note that if the edge $\overline{v_{13}v_{14}}$ was chosen instead of $\overline{v_{13}v_{10}}$ then the local minimum of the function would be outside the domain of the simplex with the Sperner labelling. This is an important observation because it demonstrates that Theorem 3 cannot be used to further refine the location of the local minimum from the domain $\text{st}(v_{13})$ using mechanisms of the proof, it only states that at least one local minimum exists within $\text{st}(v_{13})$.

The boundaries of $\text{st}(v_{13})$ can be found using the 3-chain $C_{13}(\mathcal{H}^3)$ of simplices in $\text{st}(v_{13})$, recall that the directions of simplices higher than dimension 2 are undefined and so the directions can be arbitrarily chosen

$$C_{13}(\mathcal{H}^3) = \overline{v_{13}v_{10}v_5} + \overline{v_{13}v_5v_9} + \overline{v_{13}v_9v_{14}} + \overline{v_{13}v_{14}v_2} + \overline{v_{13}v_2v_{10}}$$

$C_{13}(\mathcal{H}^3)$ clearly forms a cycle, applying the boundary operator we find the faces defining the bounds of the domain of $\text{st}(v_i)$ which in this case is the chain of edges with

defined direction

$$\partial(C_{13}(\mathcal{H}^3)) = -\overline{v_{10}v_5} + \overline{v_5v_9} - \overline{v_9v_{14}} + \overline{v_{14}v_2} + \overline{v_2v_{10}}$$

thus $\partial(\partial(C(\mathcal{H}^3))) = \emptyset$.

$v_7 = (1.15, 0.625)$ is an example of a minimiser that does not have all three required directions for a Sperner labelling, the light red shaded area represents the area wherein a local minimum can exist. For example on the lines $x_1 = 0$ for $x_2 \in [0.625, 2.5]$ or $x_2 = 2.5$ for $x_1 \in [0, 1.15]$ there will either exist a point \mathbf{p} where the gradient $\nabla f(\mathbf{p})$ points in any direction pointing towards $[\frac{3}{2}\pi, 0)$ in which case and edge $\overline{v_{13}\mathbf{p}}$ exists that points in the $[\frac{\pi}{2}, \pi)$ direction and we have a simplex with a Sperner labelling. For example the dotted line on Figure 4.4 with the Sperner simplex represented by blue shaded around v_7 . If such a point does not exist then all points on those lines point in the $[0, \frac{3}{2}\pi)$ direction and so one or more local minimum lies somewhere on the boundary which is within the defined area.

There have been several developments in the extension of this lemma which could prove useful in applications extending the SHGO algorithm. One of the most interesting is by ? where they proved the Atanassov conjecture ? that for any polytope with N vertices there are $N - n$ simplices that receive a complete set of Sperner labels. ? further extended this theorem and more recently ? extended the theorems to a large class of manifolds with or without boundary. These theorems could prove useful for extending the algorithm to make use of this information. More explicitly, SHGO currently uses knowledge of the objective function evaluations, but only in a Boolean sense (in the form of directed edges). The theorems by Meunier and Musin allow us to extend Sperner's lemma to a simplicial complex built in a $(n + 1)$ -dimensional non-euclidean space. This would allow the application of ideas from discrete differential geometry. For example the Gauss-Bonnet theorem holds for discrete simplicial surfaces ?. The Gauss-Bonnet theorem provides a relation between the total Gaussian curvature and the Euler characteristic of a surface. By simple summation of the angle defect around every vertex we can determine the Euler characteristic of a continuous surface. As will be demonstrated in Section 5.4 the simplicial complex used by SHGO is homeomorphic to complexes built on other topological hypersurfaces. Therefore when using explicit coordinates of the expected homomorphism the summation can be used to compare the error with the Euler characteristic which provides a metric for how accurately the objective function surface has been sampled. In global optimisation theory a simplicial complex built in this space can be used for approximating local and global Lipschitz constants for an objective function while still retaining the ability to detect locally convex sub-domains in the search space.

4.4 Invariance of the directed complex within a bounded rectangle

We now have a guarantee of finding stationary points in sub-domains near stationary points. However, we would also like to ensure that SHGO does not generate more than one minimiser starting point per convex sub-domain. This can only be guaranteed when an objective function surface is "adequately sampled". For black box functions there is no way to know if the number and distribution of sampling points is adequate without more information (for example if the number of local minima are known in the problem). However, it is an important property of the algorithm that $|\mathcal{M}|$ will stop increasing with higher sampling after this point. First we define an adequately sampled surface.

Definition 22. *Consider a simplicial complex \mathcal{H} built on an objective function f with a compact feasible set Ω using Definitions 17 through 20. The surface is said to be **adequately sampled** if there is one and only one true stationary point within every domain defined by Theorem 3.*

The remainder of this section is devoted to proving the following theorem which holds in the case where $\Omega = [\mathbf{l}, \mathbf{u}]^n$.

Theorem 5. *(Invariance of an adequately sampled simplicial complex \mathcal{H}) For a given continuous objective function f that is adequately sampled by a sampling set of size N . If the cardinality of the minimiser pool extracted from the directed simplex \mathcal{H} is $|\mathcal{M}|$. Then any further increase of the sampling set N will not increase $|\mathcal{M}|$.*

Proof. The proof relies on a homomorphism between the simplicial complex \mathcal{H} constructed in the bounded hyperrectangle Ω and the homology (mod 2) groups of a constructed surface \mathcal{S} on which we can invoke the invariance theorem.

Define the n -torus \mathcal{S}_0 from the compact, bounded hyperrectangle Ω by identification of the opposite faces and all extreme vertices. Now for every strict local minimum point $\mathbf{p} \in \Omega$ puncture a hypersphere and after appropriate identification the resulting n -dimensional manifold \mathcal{S}_g is a connected g sum of g tori $\mathcal{S} := \mathcal{S}_0 \# \mathcal{S}_1 \# \cdots \# \mathcal{S}_{g-1}$ (g times)

Any triangulation \mathcal{K} of the topological space \mathcal{S} is homeomorphic to \mathcal{S} , $\mathbf{H}_k(\mathcal{K}) \cong \mathbf{H}_k(\mathcal{S}) \forall k \in \mathbb{Z}$. Note that this homomorphism is for a mod 2 homology between a triangulation \mathcal{K} and the surface \mathcal{S} and is thus undirected. A triangulation corresponding to all vertices and faces of \mathcal{K} can be directed according to 17, 18 and 19 providing the directed simplicial complex \mathcal{H} . By construction we have, for an adequately sampled simplicial complex \mathcal{H} , an equality which exists between the cardinality of \mathcal{M} and the Betti numbers of \mathcal{S} as $|\mathcal{M}| = h_1 = \text{rank}(\mathbf{H}_1(\mathcal{S})) = \text{rank}(\mathbf{H}_1(\mathcal{K}))$. Here we invoke the invariance theorem

Theorem 6. (*Invariance theorem?*) *The homology groups associated with a triangulation \mathcal{K} of the a compact, connected surface \mathcal{S} are independent of \mathcal{K} . In other words, the groups $\mathbf{H}_0(\mathcal{K})$, $\mathbf{H}_1(\mathcal{K})$ and $\mathbf{H}_2(\mathcal{K})$ do not depend on the simplices, incidence coefficients, or anything else arising from the choice of the particular triangulation \mathcal{K} ; they depend only on the surface \mathcal{S} itself.*

The invariance theorem can be extended to higher dimensional triangulable spaces using singular homology through the Eilenberg-Steenrod Axioms ???. As a direct consequence any triangulation of \mathcal{S} will produce the same homology groups for \mathcal{K} .

Adding any new sampling point within the corresponding subdomains of $\text{st}(v_i) \forall i (v_i \in \mathcal{M} \subseteq \mathcal{H}^0)$ as defined in Theorem 3 will by definitions 17 through 20 need to be connected directly to v_i by a new edge or the triangulation is no longer a simplicial complex and thus not increase $|\mathcal{M}|$ since only one vertex will be the new minimiser.

After adding any sampling point outside a domain $\text{st}(v_i)$ then, through the established homomorphism, any construction of \mathcal{H} will produce the same homology groups since $\text{rank}(\mathbf{H}_1(\mathcal{K}))$ remains unchanged and it is thus not possible for a new vertex to be wrongly identified as a minimiser in the triangulation \mathcal{H} .

This concludes the proof that any increase in N will not further increase $|\mathcal{M}|$. \square

It is important to note that Theorem 5 is only applicable to complexes with adequate sampling as defined, that is to say it is entirely possible that, in complexes with less than adequate sampling, two starting minimiser elements of \mathcal{M} will converge to the same local minimum. This flaw is inherent in the fact that there is insufficient information to completely identify the minima of a surface (and could be overcome if some extra information about f is known).

Theorem 3 and Theorem 5 also lead to the following corollary about an optimisation problem:

Corollary 2. *Consider any objective function $f : \Omega \subseteq \mathbb{R}^n \rightarrow \mathbb{R}$. Consider also a local minimisation routine that is guaranteed to converge to a local minimum in the same locally convex domain as the starting point inputted to the algorithm. Alternatively the local minimisation routine is guaranteed to converge to a point within a set of bounds (provided by the boundary of the k -chain around $\text{st}(v_i)$, $\partial(C(\mathcal{H}^k))$, $k = n + 1$). If such a local minimisation routine uses an element $v_i \in \mathcal{M}$ as a starting point and the routine leads to a minimum outside or on $\text{st}(v_i)$ and in addition the minimum is not contained in the set \mathcal{H}^0 . Then it can be concluded that either search space is not adequately sampled or f is not a Lipschitz smooth function.*

Therefore according to 2 if the number of local minima are known, as in for example phase equilibria problems, then we can extract valuable information about the objective function. In particular it can be determined whether or not the objective function is

Lipschitz smooth. Alternatively if the function is known to be Lipschitz smooth then 2 can be used to prove the sampling is insufficient when the condition is not met. When this happens it is also now known that there are more local minima to be found, one or more of which might possibly be the global minimum. 2 does not, however, provide any guarantee that the sampling is sufficient when the conditions are met.

4.5 Sampling generation

Using the Sobol sequence sampling point generation proceeds in a similar way as that described in section 2.1. However, rather than only generating an arbitrary number of pre-defined sampling points we will also consider heuristic methods starting with the minimum amount of sampling points required to triangulate an n dimensional space. For example start with the minimum amount of sampling points to construct an n -dimensional simplex and continue sampling while continuously calculating the $\mathbf{H}_1(\mathcal{H})$ homology groups of the complex. Using the definitions described in this section the sampling is continued until the growth rate of the approximated homology groups slows appreciably.

In this publication the Sobol sequenced sampling points are triangulated using De-launay triangulation as implemented in the SciPy library [\[1\]](#). A major disadvantage to this triangulation scheme is that it does not scale well to higher dimensions since it relies on solving convex hull using the quickhull method developed by [\[2\]](#). There are several possibilities for mitigating this problem. Since the Sobol sequence is deterministic the triangulations can be calculated and stored in a database. For SHGO another possibility whereby the convex hull does not need to be solved by using symmetry generated triangulation was developed. Building on the initial n -cube triangulation developed by [\[3\]](#) and using the symmetry groups S_n , $n = \{1, 2, 3, \dots, n\}$ to generate an initial triangulation. Subsequent uniform sampling that ensures a symmetrical triangulation is generated in the next generation of simplices. This is done by an ordering of edges and using the cycle $(123 \dots n-1)$ to ensure that we always split every simplex by a hyperplane that goes through a child vertex on the longest edge of simplex and every other vertex in the parent simplex that does not have incidence on the edge. Figure 4.5 demonstrates the symmetry of this sampling in $n = 2$ where the longest edge in the initial triangulation was sampled. Here an iteration is defined as any generation of sub-triangulations that provides a triangulation symmetrical to the initial triangulation. An implementation of this sampling sequence is available at [\[4\]](#).

In this publication we will use both the Sobol and the hypercube triangulation sampling sequences. Sobol provides a more direct comparison to the TGO algorithm while the second sequence is more similar to the DISIMPL-v algorithm. We will refer to the different uses of sampling sequences as SHGO-Sobol and SHGO-Simpl in the experimental results section in Section 5

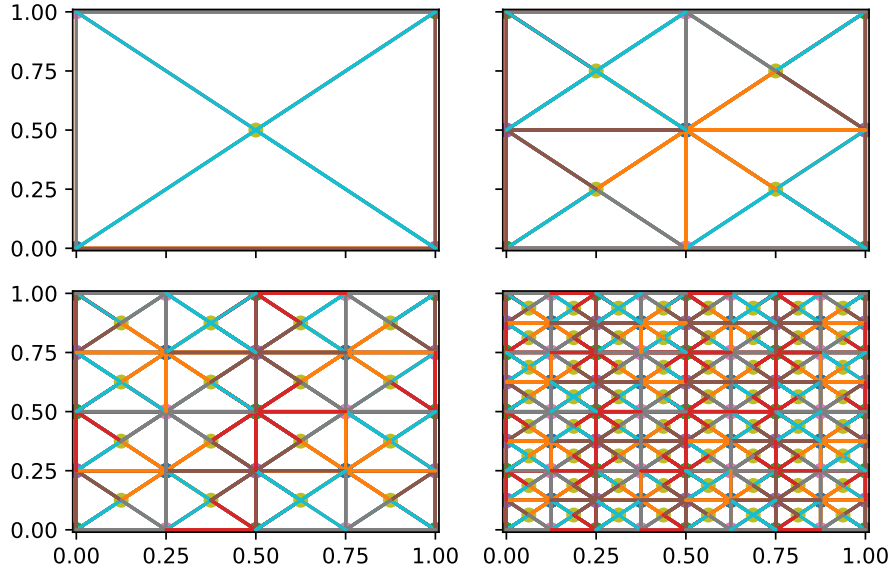


Figure 4.5: Triangulation of a unit hypercube shown in 2 dimensions for 4 iterations

4.6 Theoretical comparison to the DISIMPL algorithm

The DISIMPL algorithm developed by ? ??? is based on spatial partitioning of the search space. DISIMPL-v in particular should have a similar initial complex as SHGO-Simpl for box problems since this algorithm samples on the vertices of the simplicial complex (while DISIMPL-c samples at the geometric centre of the simplices which is more appropriate for higher dimensional problems). The graph structure of DISIMPL-v can thus be used to construct the directed complex \mathcal{H} and the homological properties can be calculated and applied. An example of one such application is given in the following paragraph.

At every iteration of the DISIMPL algorithm potentially optimal simplices are selected for refinement by considerations the Lipschitz properties of the optimisation problem. In general a combination of promising simplices with good function evaluations (related to local exploration of the search space) and simplices with larger hypervolumes (related to global exploration of the search space). Gb-DISIMPL ? is a very promising acceleration technique accomplished by switching between a "global phase" and a "usual phase". The global phase is focused on exploring simplices with larger hyper volumes and excludes smaller simplices which are potentially optimal in the usual phase. This technique prevents excessive evaluations near local minima as demonstrated in ?. Local minima can put a "drag" on the progress of refining the minimum because the algorithm selects many neighbouring simplices that are slightly worse on the function values, but also slightly larger in volume. A meta-parameter is used in Gb-DISIMPL to select the simplices to be excluded in the global phase and was shown in ? to be very efficient. However, using

knowledge from the directed complex of \mathcal{H} , the domain containing these simplices near the local minima could also be identified more explicitly through a Sperner labelling if the function is known to be Lipschitz smooth.

4.7 Algorithm implementation

We consider two modes for the SHGO algorithm. In the first a finite number of sampling points N are specified and sampling is continued until an Ω set of cardinality N is produced and no further sampling occurs. This method is demonstrated by Algorithm 2. The main reason for this algorithm is to present a more direct comparison to TGO that can be used in numerical experiments.

For the purposes of global optimisation and local minima exploration Algorithm 3 is more appropriate. By continuously calculating the $\mathbf{H}_1(\mathcal{H})$ homology group several termination criteria can be used to end the sampling. For example if the amount of local minima is known the sampling can be terminated once $|\mathcal{M}|$ is large enough. Another example with many possible heuristics is tracking the historical difference in $|\mathcal{M}|$ over $|\mathcal{P}|$ and terminating sampling if $|\mathcal{M}|$ is unchanged after a certain increase in $|\mathcal{P}|$. In optimisation problems where the global minimum is known we can also use the stopping criteria such as the one defined by ? .

$$pe = 100\% \times \begin{cases} \frac{\min\{\mathcal{F}\} - f^*}{|f^*|}, & f^* \neq 0 \\ \min\{\mathcal{F}\}, & f^* = 0 \end{cases}$$

Here $\min\{\mathcal{F}\}$ is the minimum function evaluation obtained including values obtained in the output of the local minimisation step as shown in the algorithm. Whatever termination criterion is used it requires an input $\mathbf{H}_1(\mathcal{H})$ or $\min\{\mathcal{F}\}$ and should output a Boolean, we will refer to this function as **TERM**($\mathbf{H}_1(\mathcal{H}), \min\{\mathcal{F}\}$) in Algorithm 3. In the practical implementation of the algorithm the user can also specify a finite number of iterations and/or sampling points. This functionality has been programmed into the **TERM**($\mathbf{H}_1(\mathcal{H}), \min\{\mathcal{F}\}$) function.

Open source python implementations of both of these algorithms are available and were published under a MIT compatible license ? .

Algorithm 2 SHGO finite sampling algorithm

```

1: procedure INITIALISATION
2:   Input an objective function  $f$ , constraint functions  $\mathbf{g}$  and variable bounds and
    $[\mathbf{l}, \mathbf{u}]^n$ .
3:   Input  $N$  initial sampling points.
4:   Define a sampling sequence that generates a set  $\mathcal{X}$  of sampling points in the unit
   hypercube space  $[\mathbf{0}, \mathbf{1}]^n$ 
5: end procedure
6: procedure INITIAL SAMPLING
7:    $\mathcal{P} = \emptyset$ 
8:   while  $|\mathcal{P}| < N$  do
9:     Generate  $N - |\mathcal{P}|$  sequential sampling points  $\mathcal{X} \subset \mathbb{R}^n$ 
10:    Stretch  $\mathcal{X}$  over the lower and upper bounds  $[\mathbf{l}, \mathbf{u}]^n$ 
11:     $\mathcal{P} = \{\mathcal{X}_i \mid \mathbf{g}(\mathcal{X}_i) \geq 0, \forall \mathcal{X}_i \in \mathcal{X}\} \cup \mathcal{P}$   $\triangleright$  (Find  $\mathcal{P}$  in the feasible subset  $\Omega$ 
    by discarding any points mapped outside the linear constraints  $g$  and adding to the
    current set of  $\mathcal{P}$ .)
12:    Set  $\mathcal{X} = \emptyset$ 
13:   end while
14:   Find  $\mathcal{F}$  from the objective function  $f : \mathcal{P} \rightarrow \mathcal{F}$ 
15: end procedure
16: procedure CONSTRUCT DIRECTED COMPLEX  $\mathcal{H}$ 
17:   Calculate  $\mathcal{H}$  from  $h : \mathcal{P} \rightarrow \mathcal{H}$ 
18: end procedure
19: procedure CONSTRUCT  $\mathcal{M}$ 
20:   Find  $\mathcal{M}$  from 20.
21: end procedure
22: procedure LOCAL MINIMISATION
23:   Calculate the approximate local minima of  $f$  using a local minimisation routine
   with the elements of  $\mathcal{M}$  as starting points.  $\triangleright$  These local minimisations can be
   performed in parallel.
24: end procedure
25: procedure PROCESS RETURN OBJECTS
26:   Order the final outputs of the minima of  $f$  found in the local minimisation step
   to find the approximate global minimum.
27: end procedure
28:
29: return the approximate global minimum and a list of all the minima found in the
   local minimisation step.

```

Algorithm 3 SHGO homology group growth algorithm

```

1: procedure INITIALISATION
2:   Input an objective function  $f$ , constraint functions  $\mathbf{g}$  and variable bounds and  $[\mathbf{l}, \mathbf{u}]^n$ .
3:   Input  $N$  initial sampling points.
4:   Define a sampling sequence that generates a set  $\mathcal{X}$  of sampling points in the unit hypercube space  $[\mathbf{0}, \mathbf{1}]^n$ 
5:   Define the empty set  $\mathcal{M}^E = \emptyset$  of vertices evaluated by a local minimisation.
6: end procedure
7: while  $\text{TERM}(\mathbf{H}_1(\mathcal{H}), \min\{\mathcal{F}\})$  is False do
8:   procedure SAMPLING
9:      $\mathcal{P} = \emptyset$ 
10:    while  $|\mathcal{P}| < N$  do
11:      Generate  $N - |\mathcal{P}|$  sequential sampling points  $\mathcal{X} \subset \mathbb{R}^n$ 
12:      Stretch  $\mathcal{X}$  over the lower and upper bounds  $[\mathbf{l}, \mathbf{u}]^n$ 
13:       $\mathcal{P} = \{\mathcal{X}_i \mid \mathbf{g}(\mathcal{X}_i) \geq 0, \forall \mathcal{X}_i \in \mathcal{X}\} \cup \mathcal{P}$   $\triangleright$  (Find  $\mathcal{P}$  in the feasible subset  $\Omega$  by discarding any points mapped outside the linear constraints  $\mathbf{g}$  and adding to the current set of  $\mathcal{P}$ .)
14:      Set  $\mathcal{X} = \emptyset$ 
15:    end while
16:    Find  $\mathcal{F}$  from the objective function  $f : \mathcal{P} \rightarrow \mathcal{F}$  for any new points in  $\mathcal{P}$ 
17:  end procedure
18:  procedure CONSTRUCT/APPEND DIRECTED COMPLEX  $\mathcal{H}$ 
19:    Calculate  $\mathcal{H}$  from  $h : \mathcal{P} \rightarrow \mathcal{H}$   $\triangleright$  (If  $\mathcal{H}$  was already constructed new points in  $\mathcal{P}$  are incorporated into the triangulation.)
20:    Calculate  $\mathbf{H}_1(\mathcal{H})$ 
21:  end procedure
22:  procedure CONSTRUCT  $\mathcal{M}$ 
23:    Find  $\mathcal{M}$  from 20.
24:  end procedure
25:  procedure LOCAL MINIMISATION
26:    Calculate the approximate local minima of  $f$  using a local minimisation routine with the elements of  $\mathcal{M} \setminus \mathcal{M}^E$  as starting points.  $\triangleright$  Process the most promising points first.
27:     $\mathcal{M}^E = \mathcal{M}^E \cup \mathcal{M}$   $\triangleright$  This excludes the evaluation any element  $v_i \in \mathcal{M}$  that is known to be the only point that in the domain  $\partial\text{st}(v_j)$  where  $v_j$  is known to any point already used as a starting point in Step 27. If any new  $v_i \in \mathcal{M}$  not in  $\mathcal{M}^E$  is known to be the only point  $\partial\text{st}(v_j)$  it can also be excluded.
28:    Add the function outputs of the local minimisation routine to  $\mathcal{F}$ 
29:  end procedure
30:  Find new value of  $\text{TERM}(\mathbf{H}_1)(\mathcal{H}, \min\{\mathcal{F}\})$ 
31: end while
32: procedure PROCESS RETURN OBJECTS
33:  Order the final outputs of the minima of  $f$  found in the local minimisation step to find the approximate global minimum.
34: end procedure
35:
36: return the approximate global minimum and a list of all the minima found in the local minimisation step.

```

CHAPTER 5

Experimental Results

5.1 Comparison to with linear constraints

In this section we provide experimental comparisons on 22 linearly constrained problems comparing the SHGO, TGO, Lc-DISIMPL ?, PSwarm ? and DIRECT-L1 ? algorithms. Note that the data for the Lc-DISIMPL, PSwarm and DIRECT-L1 algorithms was taken from ?. The same percentage error of $pe = 0.01\%$ used by ? was also used in this publication. To provide a fair comparison of TGO to SHGO and the other solvers the TGO algorithm was modified to stop sampling when it produced a minimiser that lead to the global minimum of the problem. Table 5.1 shows the results. Here f.e. is the total number of objective function evaluations required to solve the function and p.f.e. is the total number of penalty function evaluations. ? used DIRECT-L1 with the 3 different penalty parameters (p.p.) shown in the table. The PSwarm solver was run 10 times for each test problem.

The SHGO-Simpl, SHGO-Sobol and TGO (using Henderson’s formula for k_c) algorithms were able to solve all 22 problems. The lowest average number of function evaluations was achieved by SHGO-Simpl followed by SHGO-Sobol and TGO. It can be observed that Lc-DISIMPL-v achieved a better performance than any other algorithm for the horst-1 to horst-6, hs024, hs035, s232, s250 and bunnag2 problems. As noted in ? the initial triangulation of Lc-DISIMPL-v evaluates the function values at the vertices of the simplices and therefore for some of the tested problems the solutions were found after initial triangulation on one of the vertices of the feasible region. It is also possible to initiate SHGO with such an initial triangulation by definition the first few vertices in \mathcal{X} as the intercepts of the linear constraints in a similar way to ? and then continuing to add sampling points as normal.

Table 5.2 provides additional information for SHGO and TGO including the total number of function evaluations required by the algorithm to solve the problem (f.e.), the number of minimisers generated as starting points by the algorithm (nlmin), the number

of unique local minima mapped by the algorithm (nulmin) and the total processing time (runtime) in seconds.

It can be seen that neither of the SHGO algorithms produced more starting points leading to the same local minima as predicted by the theory for adequately sampled function surfaces. On the contrary TGO produced more than one starting point in the same locally convex domain on some test problems which lead to extra function evaluations, producing a poorer overall performance. While SHGO-Simpl had the lowest number of average function evaluations, a higher processing run time is observed compared to the other 2 algorithms. This can be explained by the fact the triangulation code for the sampling has not yet been optimised which consumed most of the run time. SHGO-Sobol and TGO use the same sampling generation code and it is observed that SHGO-Sobol has a lower processing run time as expected.

The source code used to produce these results including the scripts that run the test benchmarking suite is publically available at [?](#). The specifications of the system used to run the test problems can be found in Appendix .1.

Table 5.1: Function evaluation comparisons for test problems with linear constraints. The results for the Lc-DSIMPL, PSwarm and DIRECT-L1 algorithms were taken from ?

Problem	shgo-		tgo	Lc-DSIMPL- ^c		PSwarm ^c						DIRECT-L1 ^c		
	-simpl	-sobol		-v	-c	Minimum		Average		Maximum		p.p. = 10	p.p. = 10 ²	p.p. = 10 ⁶
	f.e.	f.e.	f.e.	f.e.	f.e.	f.e.	p.f.e	f.e.	p.f.e.	f.e.	p.f.e	f.e	f.e.	f.e.
horst-1	97	24	34	7	249	167	182	1329 ^{b(3)}	1343 ^{b(3)}	4100 ^{b(3)}	4101 ^{b(3)}	287 ^a	3689	100000
horst-2	10	11	11	5	171	160	176	424	492	768	867	265 ^a	10829	100000
horst-3	6	7	6	5	249	42	43	44	45	46	47	5 ^a	591	617
horst-4	10	25	24	8	260	90	179	114	194	129	211	58293 ^a	100000	100000
horst-5	20	15	15	8	259	106	150	134	192	214	302	7 ^a	100000	100000
horst-6	22	59	77	10	284	90	172	110	192	133	227	11 ^a	739 ^a	100000
horst-7	10	15	13	10	220	188	201	380	403	919	957	7 ^a	71 ^a	100000
hs021	24	23	23	189	133	110	110	189	192	392	405	97	97	97
hs024	24	15	36	3	141	101	153	118	172	138	195	19 ^a	57 ^a	100000
hs035	37	41	35	630	721	266	311	316	369	327	373	100000	100000	100000
hs036	105	20	103	8	314	179	179	396	401	561	574	25 ^a	49 ^a	100000
hs037	72	63	258	186	9129	127	131	160	167	201	574	7 ^a	7 ^a	100000
hs038	225	1029	389	3379	100000	53662	54445	58576	59821	65677	67660	7401	5885	6511
hs044	199	35	51	20	440	148 ^{b(9)}	218 ^{b(9)}	186 ^{b(9)}	281 ^{b(9)}	201 ^{b(9)}	299 ^{b(9)}	90283	100000	100000
hs076	56	37	44	548	4794	132	198	203	286	275	341	19135	100000	100000
s224	166	165	165	49	463	105	107	121	122	157	158	7 ^a	431	457
s231	99	99	383	2137	655	542	1011	2366	3020	4116	4800	1261	1209	43341
s232	24	15	22	3	141	105	144	119	171	162	236	19 ^a	57 ^a	100000
s250	105	20	103	8	314	296	296	367	375	495	498	25 ^a	49 ^a	100000
s251	72	63	258	186	9127	83	84	129	137	175	180	7 ^a	7 ^a	100000
bunnag1	34	47	39	630	721	132	142	214	228	411	438	1529	1495	1463
bunnag2	46	36	35	16	500	150	153	252	259	410	426	100000	100000	100000
Average	66	88	100	366	5877	2590	2672	3011	3130	3637	3812	17213	28421	75113

^a result is outside the feasible region

^{b(t)} t out of 10 times the global solution was not reached

^c results produced by ? (2)

Table 5.2: Total and average performance over all 22 test problems.

		f.e.	nlmin	nulmin	runtime (s)
problem	name				
All	shgo-simpl	1463	26	26	0.27294
	shgo-sobol	1864	23	23	0.11225
	tgo	2123	29	25	0.093607
Average	shgo-simplicial	65	1	1	0.012852
	shgo-sobol	88	1	1	0.004144
	tgo	100	1	1	0.004542

5.2 Function evaluations and comparison to other open source global optimisation algorithms

In this section we present numerical experiments comparing the SHGO and TGO algorithms with the SciPy implementations of basinhopping (BH) and differential evolution (DE). These algorithms were chosen both because the open source versions are readily available in the SciPy project and because BH is commonly used in energy surface optimisations from which the motivation for developing SHGO grew. DE has also been applied in optimising Gibbs energy surfaces for phase equilibria calculations. The optimisation problems in Appendix .1 were selected from the SciPy global optimisation benchmarking test suite (??). The test suite contains multi-modal problems with box constraints, they are described in detail in ?. We again used the stopping criteria $pe = 0.01\%$ for SHGO and TGO. For the stochastic algorithms (BH and DE) the starting points provided by the test suite were used. For every test the algorithm was terminated if the global minimum was not found after 10 minutes of processing time and the test was flagged as a fail. For comparisons we used normalised performance profiles using function evaluations and processing time as performance criteria. In total 180 test problems were used.

From Fig. 5.1 it can be observed that for this problem set SHGO-Sobol was the best performing algorithm, followed closely by TGO and SHGO-Simpl. Fig. 5.2 provides a clearer comparison between these three algorithms. While the performance of all 3 algorithms are comparable, SHGO-Sobol tends to outperform TGO, solving more problems for a given number of function evaluations. This is expected since, for the same sampling point sequence, TGO produced more than one starting point in the same locally convex domain on some test problems which leads to extra function evaluations. In total TGO produced 403 minima of which only 393 minima were unique while all of the 225 minima produced by SHGO-Sobol were unique. SHGO-Simpl produced 238 of which all 238

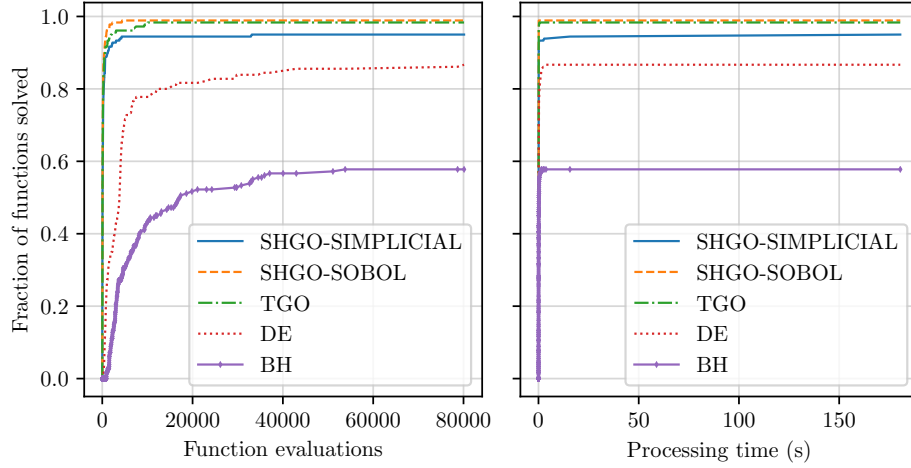


Figure 5.1: Performance profiles for SHGO, TGO, DE and BH on SciPy benchmarking test suite

were unique. It is apparent that SHGO-Simpl performed worse compared to the other sampling methods despite a better performance on the test problem set with linear constraints. There are two reasons for this result. First of all the uniformity properties of the Sobol sequence hold only for hypercubes, therefore it is lost for geometries defined by the search spaces inside linear constraints. Secondly the current code for the triangulation of the simplex cannot add only one sampling point per iteration, but must split all the simplices until the symmetry of the entire complex is restored. This leads to a much higher number of function evaluations during the sampling step of the algorithm.

The Table in Appendix .1 shows the raw numerical results. Note that, unlike the data in performance profiles, failed test runs did not get set to the worst case performance criteria by any solver (in order to preserve the raw data). Therefore the total and average function evaluations and processing times are misleading. The Table is mostly useful for comparisons on a particular test problem as well as comparing the total number of minima and unique minima found.

5.3 Invariance and optimum minimiser pool

The following 4 optimisation test problems were used to demonstrate the applications of Theorem 5 and to show the minimiser pool growth compared to TGO over a large number of sampling points. The results plotted in Figure 5.3 shows that SHGO performed as expected with the minimiser pool staying at the optimum cardinality to map all the local minima once the sampling is adequate as well as the shortcomings of the TGO especially in the higher dimensional test problems where the the minimiser pool tends to grow rapidly with the number sampling points N .

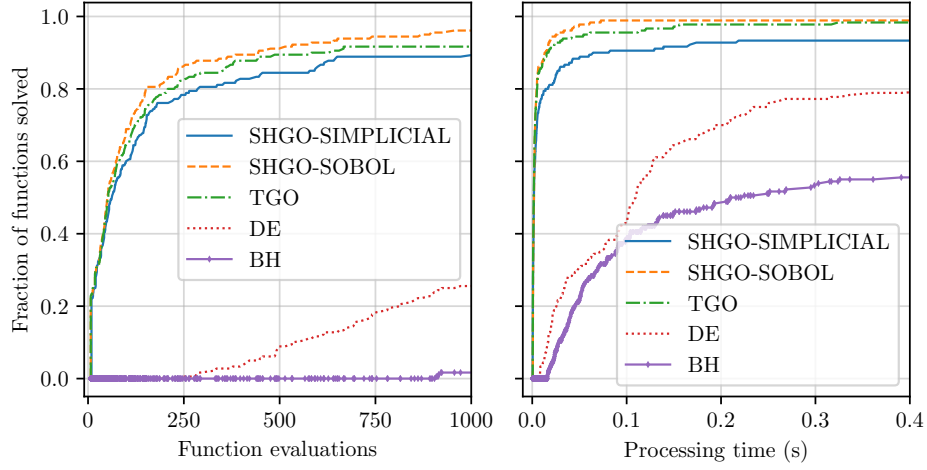


Figure 5.2: Performance profiles zoomed in to the range of $f.e. = [0, 1000]$ function evaluations and $[0, 0.4]$ seconds run time

The Ursem01 function for two dimensions is defined as follows ?

$$f(\mathbf{x}) = -\sin(2x_1 - 0.5\pi) - 3\cos(x_2) - 0.5x_1, \quad \mathbf{x} \in \Omega = [0, 9] \times [-2, 2] \quad (5.1)$$

The Paraboloid function for six dimensions is defined as follows

$$f(\mathbf{x}) = \sum_{i=1}^6 x_i^2, \quad \mathbf{x} \in \Omega = [-10, 10]^6 \quad (5.2)$$

The Bird function for two dimensions is defined as follows ?

$$f(\mathbf{x}) = (x_1 - x_2)^2 + e^{[1 - \sin(x_1)]^2} \cos(x_2) + e^{[1 - \cos(x_2)]^2} \sin(x_1), \\ \mathbf{x} \in \Omega = [-2\pi, 2\pi]^2 \quad (5.3)$$

The Schwefel01 function for six dimensions is defined as follows ?

$$f(\mathbf{x}) = \left(\sum_{i=1}^n x_i^2 \right)^{\sqrt{\pi}}, \quad \mathbf{x} \in \Omega = [-100, 100]^6 \quad (5.4)$$

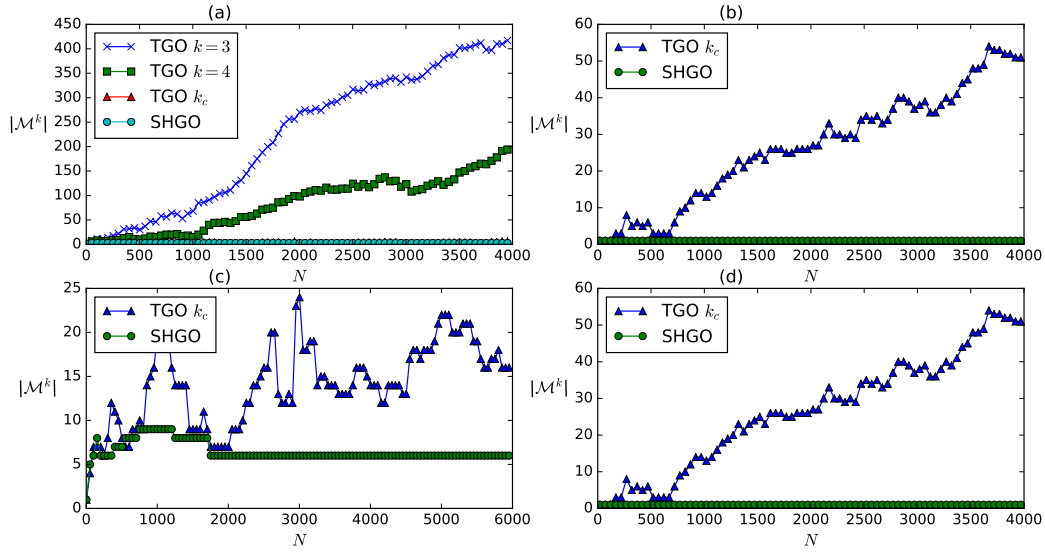


Figure 5.3: (a) The minimiser pool growth of the TGO and SHGO algorithms for the smooth objective function described in Example 3 and restated in Equation (5.1) for convenience, the SHGO never increases above the optimum of $|\mathcal{M}| = 3$, for TGO 3 different values of the k parameter are shown. (b) The minimiser pool growth for the six dimensional Paraboloid problem defined by Equation (5.2), note that even though the problem has only one minimum, the minimiser pool for TGO set at $k = k_c$ tends to increase for increasing sampling points N . In general this problem is exacerbated in higher dimensions while SHGO stays at the optimum $|\mathcal{M}| = 1$. The TGO minimiser pool for $k = 3$ and $k = 4$ are not shown here because the minimiser pool grows too rapidly. (c) The minimiser pool growth for the two dimensional Bird problem defined by Equation (5.3), an important observation here is that $|\mathcal{M}|$ is higher than optimum for SHGO before the sampling is adequate as defined by Equation (5) which happens at the after there are $N = 1722$ Sobol sequenced points after which $|\mathcal{M}|$ stays at the optimum value equal to the number of unique local minima with increasing N . (d) The minimiser pool growth for the six dimensional Schwefel01 problem defined by Equation (5.4), here again $|\mathcal{M}|$ for TGO set at k_c grows rapidly with N while $|\mathcal{M}|$ for SHGO stays constant at the optimum.

CHAPTER 6

Concluding remarks

The SHGO algorithm developed here shows promising properties and performance. On problems with linear constraints it was shown to provide competitive results to the TGO, Lc-DISIMPL, PSwarm and DIRECT-L1 algorithms. The use of a simplicial complex provides access to a wealth of tools from combinatorial topology and the growing field of computational homology. We are hopeful that these will drive further extensions and development of the algorithm. Many challenges remain such as finding the most appropriate sampling sequences for different classes of problems and finding computer resource efficient triangulation schemes. Due to the useful characterisations of objective function hypersurfaces provided by the homology groups of the simplicial complex SHGO allows an optimisation practitioner with a useful visual tool for understanding and efficiently solving higher dimensional black and grey box optimisation problems.

The main initial driving force behind the development of this algorithm grew out of a need for efficient, deterministic and reliable global optimisation methods for applications in phase equilibria modelling and calculations. However, the SHGO algorithm described here is appropriate for solving a wider class of global optimisation problems both those where mapping all the local minima is of interest and where only the global optimum is needed. It is especially appropriate for computationally expensive black and grey box functions common in science and engineering as described for example by ?.

Some key features of SHGO are that when the optimisation search space is adequately sampled and enough information is available to determine that all local minima have been mapped it is guaranteed that only one starting point for every locally convex domain will be produced by the algorithm. Note that in optimisation problems where the number of local minima is known, the sampling can stop and the local minimisation step started without superfluous function evaluations while for optimisation problems with an unknown number of local minima is unknown (and thus we can never truly know if all local minima has been found for any finite number of sampling) the guarantee still holds that that SHGO will not produce superfluous starting points that lead to the same stationary points. In addition because the homology groups can be calculated as sampling progresses

an optimisation practitioner can both visualise the extent of the optimisation problem's multi-modality and use intelligent stopping criteria for the sampling stage.

.1 Numerical results for selected optimisation problems

Table 1 show respectively: the name of the optimisation test problem (Problem), the name of the algorithm (Alg), number of dimensions (n) of the optimisation problem, the number of function evaluations required by the algorithm to solve the problem (nfev), the number of minimisers generated as starting points by the algorithm (nlmin), the number of unique local minima mapped by the algorithm (nulmin), whether successful convergence to the global minima was achieved (Success), the CPU run time measured in seconds (Runtime) and finally the number of function evaluations per unique local minima (nfev/nulmin). For all these test problems the algorithm was terminated if the algorithm ran for longer than 10 minutes.

The optimisation runs were done on a computer with the following specifications:

- CPU: Intel Core i7-6700K CPU @ 4.2GHz
- Kernel: x86_64 Linux 4.12.10-1-ARCH
- RAM: 15973MiB

Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
All	bh	0	1358408	0	0	NaN	NaN
	de	0	934804	0	0	NaN	NaN
	shgo-simplicial	0	72240	238	238	NaN	NaN
	shgo-sobol	0	29694	225	225	NaN	NaN
	tgo	0	63533	403	393	NaN	NaN
Average	bh	0	7546	0	0	NaN	0.10897
	de	0	5193	0	0	NaN	0.18817
	shgo-simplicial	0	401	1	1	NaN	1.11554
	shgo-sobol	0	164	1	1	NaN	0.00477
	tgo	0	352	2	2	NaN	0.00867

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
Ackley01	bh	2	16107	0	0	True	0.29883
	de	2	3423	0	0	True	0.19042
	shgo-simplicial	2	54	1	1	True	0.00175
	shgo-sobol	2	52	1	1	True	0.04189
	tgo	2	52	1	1	True	0.00199
Ackley02	bh	2	11844	0	0	True	0.09011
	de	2	456	0	0	True	0.01081
	shgo-simplicial	2	90	1	1	True	0.00190
	shgo-sobol	2	88	1	1	True	0.00173
	tgo	2	88	1	1	True	0.00161
Ackley03	bh	2	2370	0	0	False	0.04050
	de	2	421	0	0	True	0.01316
	shgo-simplicial	2	59	1	1	True	0.00144
	shgo-sobol	2	57	1	1	True	0.00152
	tgo	2	57	1	1	True	0.00143
Adjiman	bh	2	2070	0	0	False	0.04687
	de	2	532	0	0	True	0.03735
	shgo-simplicial	2	26	1	1	True	0.00362
	shgo-sobol	2	36	1	1	True	0.00490
	tgo	2	36	1	1	True	0.00441
Alpine01	bh	2	32928	0	0	True	0.30316
	de	2	4423	0	0	True	0.13895
	shgo-simplicial	2	55	1	1	True	0.00136
	shgo-sobol	2	53	1	1	True	0.00146
	tgo	2	53	1	1	True	0.00140
Alpine02	bh	2	1617	0	0	True	0.02474
	de	2	492	0	0	True	0.01472
	shgo-simplicial	2	153	5	5	True	0.00542

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
BartelsConn	shgo-sobol	2	62	1	1	True	0.00183
	tgo	2	108	3	3	True	0.00229
	bh	2	19857	0	0	True	0.20538
	de	2	1282	0	0	True	0.03618
	shgo-simplicial	2	55	1	1	True	0.00129
Beale	shgo-sobol	2	53	1	1	True	0.00149
	tgo	2	53	1	1	True	0.00130
	bh	2	6306	0	0	False	0.04513
	de	2	4803	0	0	True	0.12716
	shgo-simplicial	2	63	1	1	True	0.00122
BiggsExp02	shgo-sobol	2	61	1	1	True	0.00133
	tgo	2	61	1	1	True	0.00123
	bh	2	3009	0	0	True	0.07936
	de	2	4003	0	0	True	0.17757
	shgo-simplicial	2	147	2	2	True	0.00531
BiggsExp03	shgo-sobol	2	128	1	1	True	0.00432
	tgo	2	128	1	1	True	0.00413
	bh	3	5812	0	0	True	0.13472
	de	3	10564	0	0	True	0.49239
	shgo-simplicial	3	145	1	1	True	0.00708
BiggsExp04	shgo-sobol	3	151	1	1	True	0.00506
	tgo	3	151	1	1	True	0.00490
	bh	4	13095	0	0	True	0.29515
	de	4	29765	0	0	True	1.36838
	shgo-simplicial	4	1091	1	1	True	0.16711
BiggsExp05	shgo-sobol	4	384	1	1	True	0.01166
	tgo	4	384	1	1	True	0.01137
	bh	5	14346	0	0	False	0.46845
	de	5	7632	0	0	False	0.46143

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
Bird	shgo-simplicial	5	607	1	1	True	0.02376
	shgo-sobol	5	583	1	1	True	0.01982
	tgo	5	583	1	1	True	0.01971
	bh	2	2421	0	0	False	0.03866
	de	2	695	0	0	True	0.02148
	shgo-simplicial	2	42	1	1	True	0.00175
Bohachevsky1	shgo-sobol	2	43	1	1	True	0.00134
	tgo	2	43	1	1	True	0.00120
	bh	2	3510	0	0	True	0.03740
	de	2	2763	0	0	True	0.07746
	shgo-simplicial	2	9	1	1	True	0.00046
	shgo-sobol	2	7	1	1	True	0.00056
Bohachevsky2	tgo	2	7	1	1	True	0.00051
	bh	2	3471	0	0	True	0.03733
	de	2	2923	0	0	True	0.08088
	shgo-simplicial	2	9	1	1	True	0.00052
	shgo-sobol	2	7	1	1	True	0.00062
	tgo	2	7	1	1	True	0.00049
Bohachevsky3	bh	2	3438	0	0	True	0.03396
	de	2	3043	0	0	True	0.08109
	shgo-simplicial	2	9	1	1	True	0.00047
	shgo-sobol	2	7	1	1	True	0.00060
	tgo	2	7	1	1	True	0.00048
	bh	3	8096	0	0	True	0.18149
BoxBetts	de	3	11944	0	0	True	0.52510
	shgo-simplicial	3	89	1	1	True	0.00301
	shgo-sobol	3	76	1	1	True	0.00272
	tgo	3	76	1	1	True	0.00252
	bh	2	2229	0	0	True	0.02712

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
Branin02	de	2	615	0	0	True	0.01705
	shgo-simplicial	2	39	1	1	True	0.00098
	shgo-sobol	2	37	1	1	True	0.00106
	tgo	2	37	1	1	True	0.00093
	bh	2	2094	0	0	True	0.03192
	de	2	735	0	0	True	0.02204
	shgo-simplicial	2	111	2	2	True	0.00428
	shgo-sobol	2	44	1	1	True	0.00135
Brent	tgo	2	71	2	2	True	0.00166
	bh	2	915	0	0	True	0.01631
	de	2	6443	0	0	True	0.17647
	shgo-simplicial	2	9	1	1	True	0.00048
	shgo-sobol	2	7	1	1	True	0.00059
	tgo	2	7	1	1	True	0.00050
	bh	2	1857	0	0	True	0.03889
	de	2	4083	0	0	True	0.14645
Brown	shgo-simplicial	2	34	1	1	True	0.00116
	shgo-sobol	2	36	1	1	True	0.00133
	tgo	2	36	1	1	True	0.00124
	bh	2	663	0	0	False	0.01434
	de	2	815	0	0	True	0.02106
	shgo-simplicial	2	20	1	1	True	0.00066
	shgo-sobol	2	18	1	1	True	0.00076
	tgo	2	18	1	1	True	0.00064
Bukin02	bh	2	17166	0	0	True	0.09857
	de	2	4103	0	0	True	0.11085
	shgo-simplicial	2	26	1	1	True	0.00075
	shgo-sobol	2	24	1	1	True	0.00106
	tgo	2	24	1	1	True	0.00079
	bh	2	17166	0	0	True	0.09857
	de	2	4103	0	0	True	0.11085
	shgo-simplicial	2	26	1	1	True	0.00075
Bukin04	shgo-sobol	2	24	1	1	True	0.00106
	tgo	2	24	1	1	True	0.00079
	bh	2	17166	0	0	True	0.09857
	de	2	4103	0	0	True	0.11085
	shgo-simplicial	2	26	1	1	True	0.00075
	shgo-sobol	2	24	1	1	True	0.00106
	tgo	2	24	1	1	True	0.00079
	bh	2	17166	0	0	True	0.09857

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
Bukin06	bh	2	22014	0	0	False	0.17913
	de	2	2623	0	0	False	0.07575
	shgo-simplicial	2	1007	5	5	True	0.02179
	shgo-sobol	2	741	3	3	True	0.01237
	tgo	2	1169	5	5	True	0.01735
CarromTable	bh	2	1899	0	0	False	0.03518
	de	2	972	0	0	True	0.03484
	shgo-simplicial	2	36	1	1	True	0.00145
	shgo-sobol	2	31	1	1	True	0.00108
	tgo	2	31	1	1	True	0.00098
Cigar	bh	2	8193	0	0	True	0.08821
	de	2	3743	0	0	True	0.12489
	shgo-simplicial	2	20	1	1	True	0.00068
	shgo-sobol	2	18	1	1	True	0.00079
	tgo	2	18	1	1	True	0.00069
Colville	bh	4	12965	0	0	True	0.10312
	de	4	29685	0	0	True	0.83725
	shgo-simplicial	4	225	1	1	True	0.01405
	shgo-sobol	4	1029	2	2	True	0.06095
	tgo	4	3039	10	2	True	0.05422
Corana	bh	4	2555	0	0	False	0.07399
	de	4	4085	0	0	True	0.19297
	shgo-simplicial	4	23	1	1	True	0.00213
	shgo-sobol	4	12	1	1	True	0.00098
	tgo	4	12	1	1	True	0.00084
CosineMixture	bh	2	348	0	0	False	0.01460
	de	2	1166	0	0	True	0.03793
	shgo-simplicial	2	17	1	1	True	0.00112
	shgo-sobol	2	7	1	1	True	0.00064

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
CrossInTray	tgo	2	7	1	1	True	0.00055
	bh	2	1578	0	0	False	0.02898
	de	2	489	0	0	True	0.01929
	shgo-simplicial	2	69	1	1	True	0.00323
	shgo-sobol	2	33	1	1	True	0.00110
CrossLegTable	tgo	2	33	1	1	True	0.00102
	bh	2	17355	0	0	True	0.20905
	de	2	4783	0	0	False	0.15400
	shgo-simplicial	2	9	1	1	True	0.00054
	shgo-sobol	2	7	1	1	True	0.00065
CrownedCross	tgo	2	7	1	1	True	0.00064
	bh	2	17130	0	0	False	0.21095
	de	2	4263	0	0	False	0.13623
	shgo-simplicial	2	9	1	1	True	0.00055
	shgo-sobol	2	7	1	1	True	0.00064
Cube	tgo	2	7	1	1	True	0.00052
	bh	2	8529	0	0	True	0.05303
	de	2	5243	0	0	True	0.12590
	shgo-simplicial	2	146	1	1	True	0.00206
	shgo-sobol	2	144	1	1	True	0.00216
Damavandi	tgo	2	144	1	1	True	0.00207
	bh	2	1566	0	0	False	0.02617
	de	2	535	0	0	False	0.01555
	shgo-simplicial	2	578	2	2	True	0.03416
	shgo-sobol	2	60	2	2	True	0.00182
DeVilliersGlasser01	tgo	2	97	2	2	True	0.00202
	bh	4	16805	0	0	True	0.55462
	de	4	24265	0	0	True	1.20027
	shgo-simplicial	4	439	2	2	True	0.02469

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
Deb01	shgo-sobol	4	446	1	1	True	0.04439
	tgo	4	667	9	9	True	0.02304
	bh	2	5565	0	0	True	0.06559
	de	2	1532	0	0	True	0.04810
	shgo-simplicial	2	28	1	1	True	0.00133
Deb03	shgo-sobol	2	18	1	1	True	0.00082
	tgo	2	18	1	1	True	0.00074
	bh	2	5310	0	0	False	0.07614
	de	2	40103	0	0	False	1.45388
	shgo-simplicial	2	4234	4	4	True	0.08237
Decanomial	shgo-sobol	2	83	1	1	True	0.00257
	tgo	2	1476	2	2	True	0.02975
	bh	2	9465	0	0	True	0.11714
	de	2	3383	0	0	True	0.10716
	shgo-simplicial	2	256	1	1	True	0.00511
Deceptive	shgo-sobol	2	200	1	1	True	0.00430
	tgo	2	200	1	1	True	0.00408
	bh	2	441	0	0	False	0.01711
	de	2	913	0	0	True	0.02588
	shgo-simplicial	2	449	9	9	True	0.01786
DeckkersAarts	shgo-sobol	2	533	4	4	True	0.01272
	tgo	2	667	5	5	True	0.01531
	bh	2	2844	0	0	True	0.02659
	de	2	670	0	0	True	0.01610
	shgo-simplicial	2	86	1	1	True	0.00309
DeflectedCorrugatedSpring	shgo-sobol	2	99	2	2	True	0.00202
	tgo	2	167	2	2	True	0.00254
	bh	2	2055	0	0	True	0.04327
	de	2	892	0	0	True	0.03243

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
DixonPrice	shgo-simplicial	2	9	1	1	True	0.00069
	shgo-sobol	2	7	1	1	True	0.00067
	tgo	2	7	1	1	True	0.00056
	bh	2	3639	0	0	True	0.05977
	de	2	4563	0	0	True	0.16077
	shgo-simplicial	2	627	3	3	True	0.03906
Dolan	shgo-sobol	2	93	2	2	True	0.00286
	tgo	2	92	2	2	True	0.00241
	bh	5	51084	0	0	True	0.39090
	de	5	78692	0	0	True	2.47918
	shgo-simplicial	5	264	1	1	True	0.00755
	shgo-sobol	5	240	1	1	True	0.00405
DropWave	tgo	5	240	1	1	True	0.00383
	bh	2	2337	0	0	True	0.03663
	de	2	1012	0	0	True	0.03254
	shgo-simplicial	2	9	1	1	True	0.00052
	shgo-sobol	2	7	1	1	True	0.00063
	tgo	2	7	1	1	True	0.00052
Easom	bh	2	303	0	0	False	0.01313
	de	2	83	0	0	False	0.00192
	shgo-simplicial	2	2126	1	1	True	0.13901
	shgo-sobol	2	2210	1	1	True	0.04977
	tgo	2	2210	1	1	True	0.32387
	bh	2	1935	0	0	False	0.02526
EggCrate	de	2	3963	0	0	True	0.11058
	shgo-simplicial	2	9	1	1	True	0.00049
	shgo-sobol	2	7	1	1	True	0.00061
	tgo	2	7	1	1	True	0.00051
	bh	2	1983	0	0	False	0.04648
	de	2	3963	0	0	True	0.11058

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
ElAttarVidyasagarDutta	de	2	941	0	0	False	0.03674
	shgo-simplicial	2	616	2	2	True	0.04200
	shgo-sobol	2	233	4	4	True	0.00734
	tgo	2	293	5	5	True	0.00819
	bh	2	3219	0	0	True	0.03135
	de	2	1301	0	0	True	0.03457
Exp2	shgo-simplicial	2	33094	2	2	True	2.79985
	shgo-sobol	2	277	3	3	True	0.00557
	tgo	2	351	4	4	True	0.00526
	bh	2	2892	0	0	True	0.07888
	de	2	4003	0	0	True	0.18485
	shgo-simplicial	2	56	1	1	True	0.00259
Exponential	shgo-sobol	2	137	1	1	True	0.00475
	tgo	2	137	1	1	True	0.00476
	bh	2	1515	0	0	True	0.02607
	de	2	286	0	0	True	0.00840
	shgo-simplicial	2	9	1	1	True	0.00054
	shgo-sobol	2	7	1	1	True	0.00062
FreudensteinRoth	tgo	2	7	1	1	True	0.00053
	bh	2	5262	0	0	True	0.04146
	de	2	4103	0	0	True	0.10086
	shgo-simplicial	2	356	7	7	True	0.01209
	shgo-sobol	2	49	1	1	True	0.00119
	tgo	2	49	1	1	True	0.00100
Gear	bh	4	505	0	0	False	0.01387
	de	4	11445	0	0	True	0.33639
	shgo-simplicial	4	23	1	1	True	0.00146
	shgo-sobol	4	31	1	1	True	0.00192
	tgo	4	37	2	2	True	0.00113

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
Giunta	bh	2	2301	0	0	True	0.04630
	de	2	449	0	0	True	0.01573
	shgo-simplicial	2	34	2	2	True	0.00179
	shgo-sobol	2	31	1	1	True	0.00127
	tgo	2	31	1	1	True	0.00114
GoldsteinPrice	bh	2	4587	0	0	True	0.05059
	de	2	781	0	0	True	0.02110
	shgo-simplicial	2	85	1	1	True	0.00166
	shgo-sobol	2	83	1	1	True	0.00177
	tgo	2	83	1	1	True	0.00168
Griewank	bh	2	1872	0	0	False	0.03922
	de	2	3283	0	0	True	0.12409
	shgo-simplicial	2	9	1	1	True	0.00063
	shgo-sobol	2	7	1	1	True	0.00071
	tgo	2	7	1	1	True	0.00060
Gulf	bh	3	404	0	0	False	0.02584
	de	3	15244	0	0	True	0.92473
	shgo-simplicial	3	650	3	3	True	0.05025
	shgo-sobol	3	234	1	1	True	0.01089
	tgo	3	234	1	1	True	0.01077
Hansen	bh	2	3432	0	0	True	0.10406
	de	2	1341	0	0	True	0.07133
	shgo-simplicial	2	130	3	3	True	0.00494
	shgo-sobol	2	114	1	1	True	0.00423
	tgo	2	379	7	7	True	0.01404
Hartmann3	bh	3	7464	0	0	True	0.13259
	de	3	720	0	0	True	0.02682
	shgo-simplicial	3	70	1	1	True	0.00238
	shgo-sobol	3	54	1	1	True	0.00187

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
Hartmann6	tgo	3	53	1	1	True	0.00175
	bh	6	16625	0	0	True	0.26892
	de	6	2230	0	0	True	0.09149
	shgo-simplicial	6	181	1	1	True	0.02665
	shgo-sobol	6	153	1	1	True	0.00629
HelicalValley	tgo	6	443	3	2	True	0.01068
	bh	3	7688	0	0	True	0.07898
	de	3	12124	0	0	True	0.36106
	shgo-simplicial	3	1456	4	4	True	0.12962
	shgo-sobol	3	136	2	2	True	0.00324
HimmelBlau	tgo	3	137	2	2	True	0.00261
	bh	2	2529	0	0	True	0.02351
	de	2	4683	0	0	True	0.11484
	shgo-simplicial	2	66	1	1	True	0.00117
	shgo-sobol	2	45	1	1	True	0.00105
HolderTable	tgo	2	45	1	1	True	0.00095
	bh	2	1857	0	0	False	0.03147
	de	2	415	0	0	True	0.01261
	shgo-simplicial	2	179	2	2	True	0.01007
	shgo-sobol	2	97	2	2	True	0.00262
Hosaki	tgo	2	117	3	3	True	0.00264
	bh	2	2526	0	0	False	0.02915
	de	2	335	0	0	True	0.00849
	shgo-simplicial	2	47	1	1	True	0.00105
	shgo-sobol	2	29	1	1	True	0.00096
Infinity	tgo	2	29	1	1	True	0.00082
	bh	2	2583	0	0	True	0.04084
	de	2	3803	0	0	True	0.12671
	shgo-simplicial	2	13	1	1	True	0.00105

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
JennrichSampson	shgo-sobol	2	150	1	1	True	0.00414
	tgo	2	121	1	1	True	0.00282
	bh	2	10632	0	0	True	0.20981
	de	2	904	0	0	True	0.03388
	shgo-simplicial	2	52	1	1	True	0.00170
Judge	shgo-sobol	2	50	1	1	True	0.00176
	tgo	2	50	1	1	True	0.00164
	bh	2	3207	0	0	True	0.07254
	de	2	741	0	0	True	0.03109
	shgo-simplicial	2	53	1	1	True	0.00152
Katsuura	shgo-sobol	2	51	1	1	True	0.00192
	tgo	2	51	1	1	True	0.00181
	bh	2	474	0	0	False	0.02535
	de	2	2006	0	0	True	0.11021
	shgo-simplicial	2	9	1	1	True	0.00085
Keane	shgo-sobol	2	7	1	1	True	0.00078
	tgo	2	7	1	1	True	0.00069
	bh	2	1566	0	0	True	0.02399
	de	2	7523	0	0	True	0.21579
	shgo-simplicial	2	1502	2	2	True	0.02894
Kowalik	shgo-sobol	2	14	1	1	True	0.00078
	tgo	2	7	1	1	True	0.00050
	bh	4	14660	0	0	True	0.24895
	de	4	6755	0	0	True	0.26747
	shgo-simplicial	4	240	1	1	True	0.00643
Langermann	shgo-sobol	4	229	1	1	True	0.00565
	tgo	4	228	1	1	True	0.00553
	bh	2	2877	0	0	True	0.09868
	de	2	692	0	0	True	0.03566

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
LennardJones	shgo-simplicial	2	397	5	5	True	0.02323
	shgo-sobol	2	49	1	1	True	0.00248
	tgo	2	49	1	1	True	0.00227
	bh	6	10374	0	0	True	0.05865
	de	6	15748	0	0	True	0.45186
	shgo-simplicial	6	124	1	1	True	0.00193
Leon	shgo-sobol	6	81	1	1	True	0.00203
	tgo	6	173	1	1	True	0.00240
	bh	2	6207	0	0	True	0.04232
	de	2	5363	0	0	True	0.12909
	shgo-simplicial	2	99	1	1	True	0.00152
	shgo-sobol	2	97	1	1	True	0.00164
Levy03	tgo	2	97	1	1	True	0.00153
	bh	2	2670	0	0	False	0.06733
	de	2	3803	0	0	True	0.16395
	shgo-simplicial	2	45	1	1	True	0.00172
	shgo-sobol	2	43	1	1	True	0.00174
	tgo	2	43	1	1	True	0.00164
Levy13	bh	2	4491	0	0	True	0.05387
	de	2	3803	0	0	True	0.11468
	shgo-simplicial	2	20	1	1	True	0.00072
	shgo-sobol	2	18	1	1	True	0.00082
	tgo	2	18	1	1	True	0.00070
	bh	2	1803	0	0	True	0.01808
Matyas	de	2	4323	0	0	True	0.10309
	shgo-simplicial	2	9	1	1	True	0.00046
	shgo-sobol	2	7	1	1	True	0.00058
	tgo	2	7	1	1	True	0.00047
	bh	2	2073	0	0	False	0.02372
	de	2	4323	0	0	True	0.10309
McCormick	bh	2	2073	0	0	False	0.02372

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
Michalewicz	de	2	495	0	0	True	0.01285
	shgo-simplicial	2	42	1	1	True	0.00094
	shgo-sobol	2	40	1	1	True	0.00105
	tgo	2	40	1	1	True	0.00095
	bh	2	4320	0	0	True	0.07072
	de	2	498	0	0	True	0.01704
	shgo-simplicial	2	50	1	1	True	0.00151
	shgo-sobol	2	48	1	1	True	0.00159
MieleCantrell	tgo	2	48	1	1	True	0.00148
	bh	4	9270	0	0	True	0.08415
	de	4	42965	0	0	True	1.29797
	shgo-simplicial	4	455	1	1	True	0.00773
	shgo-sobol	4	444	1	1	True	0.00635
	tgo	4	443	1	1	True	0.00622
	bh	2	1830	0	0	False	0.02516
	de	2	406	0	0	True	0.01132
Mishra01	shgo-simplicial	2	0	0	0	False	0.00000
	shgo-sobol	2	11	1	1	True	0.00074
	tgo	2	11	1	1	True	0.00058
	bh	2	1752	0	0	False	0.02823
	de	2	566	0	0	True	0.01723
	shgo-simplicial	2	9	1	1	True	0.00051
	shgo-sobol	2	0	0	0	False	0.00000
	tgo	2	0	0	0	False	0.00000
Mishra02	bh	2	21105	0	0	False	0.19113
	de	2	2028	0	0	False	0.05798
	shgo-simplicial	2	70	1	1	True	0.00146
	shgo-sobol	2	68	1	1	True	0.00157
	tgo	2	68	1	1	True	0.00152
	bh	2	21105	0	0	False	0.19113
	de	2	2028	0	0	False	0.05798
	shgo-simplicial	2	70	1	1	True	0.00146
Mishra03	shgo-sobol	2	68	1	1	True	0.00157
	tgo	2	68	1	1	True	0.00152
	bh	2	21105	0	0	False	0.19113
	de	2	2028	0	0	False	0.05798
	shgo-simplicial	2	70	1	1	True	0.00146
	shgo-sobol	2	68	1	1	True	0.00157
	tgo	2	68	1	1	True	0.00152
	bh	2	21105	0	0	False	0.19113

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
Mishra04	bh	2	23679	0	0	False	0.21289
	de	2	1663	0	0	False	0.04837
	shgo-simplicial	2	599	3	3	True	0.01678
	shgo-sobol	2	4357	8	8	True	0.07262
	tgo	2	9363	18	18	True	0.14981
Mishra05	bh	2	7512	0	0	False	0.10062
	de	2	852	0	0	False	0.02696
	shgo-simplicial	2	50	1	1	True	0.00176
	shgo-sobol	2	142	2	2	True	0.00364
	tgo	2	263	3	3	True	0.00538
Mishra06	bh	2	2346	0	0	False	0.04210
	de	2	695	0	0	True	0.02279
	shgo-simplicial	2	62	1	1	True	0.00208
	shgo-sobol	2	121	2	2	True	0.00319
	tgo	2	170	2	2	True	0.00386
Mishra07	bh	2	1230	0	0	True	0.02805
	de	2	7043	0	0	True	0.25045
	shgo-simplicial	2	170	1	1	True	0.02487
	shgo-sobol	2	47	2	2	True	0.00179
	tgo	2	84	3	3	True	0.00226
Mishra08	bh	2	9831	0	0	True	0.12277
	de	2	2903	0	0	True	0.09230
	shgo-simplicial	2	243	1	1	True	0.00490
	shgo-sobol	2	235	1	1	True	0.00500
	tgo	2	235	1	1	True	0.00481
Mishra10	bh	2	303	0	0	False	0.01097
	de	2	923	0	0	True	0.02149
	shgo-simplicial	2	9	1	1	True	0.00047
	shgo-sobol	2	7	1	1	True	0.00058

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
Mishra11	tgo	2	7	1	1	True	0.00047
	bh	2	1140	0	0	True	0.02287
	de	2	6443	0	0	True	0.20703
	shgo-simplicial	2	9	1	1	True	0.00056
	shgo-sobol	2	7	1	1	True	0.00064
MultiModal	tgo	2	7	1	1	True	0.00052
	bh	2	21084	0	0	True	0.19264
	de	2	3643	0	0	True	0.11185
	shgo-simplicial	2	9	1	1	True	0.00054
	shgo-sobol	2	7	1	1	True	0.00063
NeedleEye	tgo	2	7	1	1	True	0.00051
	bh	2	10005	0	0	True	0.08801
	de	2	247	0	0	True	0.00480
	shgo-simplicial	2	9	1	1	True	0.00051
	shgo-sobol	2	7	1	1	True	0.00060
NewFunction01	tgo	2	7	1	1	True	0.00050
	bh	2	20874	0	0	False	0.17083
	de	2	1683	0	0	False	0.04564
	shgo-simplicial	2	3813	6	6	True	0.06441
	shgo-sobol	2	1569	6	6	True	0.02697
NewFunction02	tgo	2	10079	23	23	True	0.15573
	bh	2	22662	0	0	False	0.18609
	de	2	1721	0	0	False	0.04524
	shgo-simplicial	2	159	1	1	True	0.00448
	shgo-sobol	2	341	2	2	True	0.00578
OddSquare	tgo	2	361	2	2	True	0.00582
	bh	2	303	0	0	False	0.01561
	de	2	1238	0	0	True	0.04206
	shgo-simplicial	2	0	0	0	False	0.00000

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
Parsopoulos	shgo-sobol	2	204	1	1	True	0.00615
	tgo	2	487	8	8	True	0.01243
	bh	2	1608	0	0	True	0.02088
	de	2	4163	0	0	True	0.11071
	shgo-simplicial	2	30	1	1	True	0.00120
Pathological	shgo-sobol	2	39	1	1	True	0.00108
	tgo	2	39	1	1	True	0.00092
	bh	2	8106	0	0	True	0.19200
	de	2	2498	0	0	True	0.10988
	shgo-simplicial	2	20	1	1	True	0.00098
Paviani	shgo-sobol	2	18	1	1	True	0.00104
	tgo	2	18	1	1	True	0.00118
	bh	10	13970	0	0	False	0.23140
	de	10	6088	0	0	True	0.26143
	shgo-simplicial	10	1257	1	1	True	180.46720
PenHolder	shgo-sobol	10	364	1	1	True	0.01304
	tgo	10	358	1	1	True	0.00740
	bh	2	1512	0	0	False	0.02900
	de	2	532	0	0	True	0.01718
	shgo-simplicial	2	117	2	2	True	0.00450
Penalty01	shgo-sobol	2	86	2	2	True	0.00235
	tgo	2	74	2	2	True	0.00185
	bh	2	3300	0	0	True	0.10383
	de	2	3803	0	0	True	0.19280
	shgo-simplicial	2	45	1	1	True	0.00204
PermFunction01	shgo-sobol	2	43	1	1	True	0.00210
	tgo	2	43	1	1	True	0.00199
	bh	2	4599	0	0	True	0.13283
	de	2	4963	0	0	True	0.25058

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
PermFunction02	shgo-simplicial	2	83	1	1	True	0.00331
	shgo-sobol	2	70	1	1	True	0.00298
	tgo	2	70	1	1	True	0.00286
	bh	2	4665	0	0	True	0.13064
	de	2	4563	0	0	True	0.22856
	shgo-simplicial	2	73	1	1	True	0.00294
Pinter	shgo-sobol	2	71	1	1	True	0.00296
	tgo	2	71	1	1	True	0.00284
	bh	2	3075	0	0	False	0.12151
	de	2	4043	0	0	True	0.23087
	shgo-simplicial	2	9	1	1	True	0.00079
	shgo-sobol	2	7	1	1	True	0.00081
Plateau	tgo	2	7	1	1	True	0.00069
	bh	2	303	0	0	False	0.01253
	de	2	283	0	0	True	0.00751
	shgo-simplicial	2	9	1	1	True	0.00051
	shgo-sobol	2	7	1	1	True	0.00061
	tgo	2	7	1	1	True	0.00049
Powell	bh	4	14285	0	0	True	0.09099
	de	4	35285	0	0	True	0.93984
	shgo-simplicial	4	220	1	1	True	0.00355
	shgo-sobol	4	209	1	1	True	0.00296
	tgo	4	208	1	1	True	0.00280
	bh	4	53785	0	0	True	1.13694
PowerSum	de	4	80125	0	0	True	3.73696
	shgo-simplicial	4	386	1	1	True	0.01186
	shgo-sobol	4	641	1	1	True	0.01844
	tgo	4	640	1	1	True	0.01828
	bh	2	921	0	0	True	0.01621

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
Price02	de	2	4003	0	0	True	0.10630
	shgo-simplicial	2	14	1	1	True	0.00057
	shgo-sobol	2	12	1	1	True	0.00067
	tgo	2	12	1	1	True	0.00055
	bh	2	1614	0	0	False	0.02852
	de	2	732	0	0	False	0.02342
	shgo-simplicial	2	9	1	1	True	0.00051
	shgo-sobol	2	7	1	1	True	0.00065
Price03	tgo	2	7	1	1	True	0.00055
	bh	2	5136	0	0	True	0.04062
	de	2	4283	0	0	True	0.10703
	shgo-simplicial	2	58	1	1	True	0.00111
	shgo-sobol	2	56	1	1	True	0.00124
	tgo	2	56	1	1	True	0.00110
	bh	2	6984	0	0	True	0.05038
	de	2	40043	0	0	True	1.00106
Price04	shgo-simplicial	2	9	1	1	True	0.00048
	shgo-sobol	2	7	1	1	True	0.00059
	tgo	2	7	1	1	True	0.00048
	bh	2	1917	0	0	True	0.01907
	de	2	378	0	0	True	0.00840
	shgo-simplicial	2	35	1	1	True	0.00080
	shgo-sobol	2	33	1	1	True	0.00090
	tgo	2	33	1	1	True	0.00079
Quadratic	bh	2	35934	0	0	True	0.65335
	de	2	3823	0	0	True	0.15554
	shgo-simplicial	2	114	1	1	True	0.00343
	shgo-sobol	2	112	1	1	True	0.00347
	tgo	2	112	1	1	True	0.00337
	bh	2	35934	0	0	True	0.65335
	de	2	3823	0	0	True	0.15554
	shgo-simplicial	2	114	1	1	True	0.00343
Quintic	shgo-sobol	2	112	1	1	True	0.00347
	tgo	2	112	1	1	True	0.00337
	bh	2	35934	0	0	True	0.65335
	de	2	3823	0	0	True	0.15554
	shgo-simplicial	2	114	1	1	True	0.00343
	shgo-sobol	2	112	1	1	True	0.00347
	tgo	2	112	1	1	True	0.00337
	bh	2	35934	0	0	True	0.65335

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
Rana	bh	2	1851	0	0	False	0.04418
	de	2	1055	0	0	False	0.04140
	shgo-simplicial	2	292	5	5	True	0.01048
	shgo-sobol	2	651	10	10	True	0.01966
	tgo	2	1157	20	20	True	0.03224
Rastrigin	bh	2	3198	0	0	True	0.04510
	de	2	2323	0	0	True	0.07522
	shgo-simplicial	2	20	1	1	True	0.00075
	shgo-sobol	2	18	1	1	True	0.00083
	tgo	2	18	1	1	True	0.00072
Ratkowsky01	bh	4	5355	0	0	False	0.10861
	de	4	3595	0	0	False	0.14793
	shgo-simplicial	4	286	1	1	True	0.00856
	shgo-sobol	4	187	1	1	True	0.00514
	tgo	4	186	1	1	True	0.00501
Ratkowsky02	bh	3	18628	0	0	True	0.32585
	de	3	2308	0	0	True	0.08946
	shgo-simplicial	3	124	1	1	True	0.00534
	shgo-sobol	3	111	1	1	True	0.00297
	tgo	3	110	1	1	True	0.00285
Ripple01	bh	2	5961	0	0	False	0.12492
	de	2	824	0	0	False	0.03387
	shgo-simplicial	2	153	3	3	True	0.00682
	shgo-sobol	2	476	1	1	True	0.01604
	tgo	2	1879	29	29	True	0.06101
Ripple25	bh	2	3426	0	0	False	0.06860
	de	2	815	0	0	True	0.03085
	shgo-simplicial	2	106	3	3	True	0.00524
	shgo-sobol	2	99	1	1	True	0.00334

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
Rosenbrock	tgo	2	372	9	9	True	0.00987
	bh	2	6324	0	0	True	0.10065
	de	2	4923	0	0	True	0.17333
	shgo-simplicial	2	118	1	1	True	0.00299
	shgo-sobol	2	116	1	1	True	0.00302
RosenbrockModified	tgo	2	116	1	1	True	0.00296
	bh	2	5790	0	0	False	0.05810
	de	2	898	0	0	True	0.02411
	shgo-simplicial	2	128	2	2	True	0.00409
	shgo-sobol	2	66	1	1	True	0.00167
RotatedEllipse01	tgo	2	66	1	1	True	0.00143
	bh	2	1566	0	0	True	0.01985
	de	2	3923	0	0	True	0.10188
	shgo-simplicial	2	9	1	1	True	0.00052
	shgo-sobol	2	7	1	1	True	0.00060
RotatedEllipse02	tgo	2	7	1	1	True	0.00048
	bh	2	1542	0	0	True	0.01678
	de	2	3763	0	0	True	0.09109
	shgo-simplicial	2	9	1	1	True	0.00046
	shgo-sobol	2	7	1	1	True	0.00058
Salomon	tgo	2	7	1	1	True	0.00046
	bh	2	7743	0	0	True	0.09056
	de	2	1489	0	0	False	0.04805
	shgo-simplicial	2	40	1	1	True	0.00114
	shgo-sobol	2	38	1	1	True	0.00121
Sargan	tgo	2	38	1	1	True	0.00110
	bh	2	1536	0	0	True	0.06246
	de	2	4003	0	0	True	0.23429
	shgo-simplicial	2	9	1	1	True	0.00079

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
Schaffer01	shgo-sobol	2	7	1	1	True	0.00080
	tgo	2	7	1	1	True	0.00069
	bh	2	3273	0	0	False	0.03054
	de	2	2883	0	0	True	0.07698
	shgo-simplicial	2	9	1	1	True	0.00051
Schaffer02	shgo-sobol	2	7	1	1	True	0.00060
	tgo	2	7	1	1	True	0.00049
	bh	2	4806	0	0	False	0.04397
	de	2	2803	0	0	True	0.07734
	shgo-simplicial	2	9	1	1	True	0.00048
Schaffer03	shgo-sobol	2	7	1	1	True	0.00059
	tgo	2	7	1	1	True	0.00048
	bh	2	6183	0	0	False	0.06595
	de	2	3289	0	0	True	0.09439
	shgo-simplicial	2	0	0	0	False	0.00000
Schaffer04	shgo-sobol	2	870	2	2	True	0.01460
	tgo	2	6986	15	15	True	0.11268
	bh	2	4956	0	0	False	0.05469
	de	2	1769	0	0	True	0.05148
	shgo-simplicial	2	0	0	0	False	0.00000
Schwefel01	shgo-sobol	2	956	2	2	True	0.01607
	tgo	2	7424	16	16	True	0.12055
	bh	2	2592	0	0	True	0.03465
	de	2	4003	0	0	True	0.11927
	shgo-simplicial	2	9	1	1	True	0.00212
Schwefel02	shgo-sobol	2	7	1	1	True	0.00062
	tgo	2	7	1	1	True	0.00050
	bh	2	1530	0	0	True	0.05192
	de	2	4563	0	0	True	0.23140

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
Schwefel04	shgo-simplicial	2	9	1	1	True	0.00077
	shgo-sobol	2	7	1	1	True	0.00078
	tgo	2	7	1	1	True	0.00065
	bh	2	2856	0	0	True	0.04769
	de	2	4403	0	0	True	0.14890
Schwefel06	shgo-simplicial	2	47	1	1	True	0.00134
	shgo-sobol	2	45	1	1	True	0.00141
	tgo	2	45	1	1	True	0.00134
	bh	2	32622	0	0	True	0.22475
	de	2	4263	0	0	True	0.11486
Schwefel20	shgo-simplicial	2	150	1	1	True	0.00251
	shgo-sobol	2	148	1	1	True	0.00258
	tgo	2	148	1	1	True	0.00246
	bh	2	37062	0	0	True	0.25165
	de	2	3663	0	0	True	0.10227
Schwefel21	shgo-simplicial	2	72	1	1	True	0.00142
	shgo-sobol	2	70	1	1	True	0.00151
	tgo	2	70	1	1	True	0.00141
	bh	2	30786	0	0	True	0.15142
	de	2	4263	0	0	True	0.10354
Schwefel22	shgo-simplicial	2	23	1	1	True	0.00065
	shgo-sobol	2	21	1	1	True	0.00078
	tgo	2	21	1	1	True	0.00065
	bh	2	34536	0	0	True	0.31512
	de	2	3903	0	0	True	0.11968
Schwefel26	shgo-simplicial	2	75	1	1	True	0.00170
	shgo-sobol	2	73	1	1	True	0.00179
	tgo	2	73	1	1	True	0.00166
	bh	2	1377	0	0	False	0.02356

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
Schwefel36	de	2	1763	0	0	True	0.05906
	shgo-simplicial	2	85	2	2	True	0.06043
	shgo-sobol	2	46	1	1	True	0.00191
	tgo	2	46	1	1	True	0.00171
	bh	2	531	0	0	False	0.01282
	de	2	741	0	0	True	0.01720
	shgo-simplicial	2	593	1	1	True	0.02958
	shgo-sobol	2	514	1	1	True	0.01014
Shekel05	tgo	2	80	2	2	True	0.00142
	bh	4	4765	0	0	False	0.07772
	de	4	2505	0	0	True	0.09539
	shgo-simplicial	4	118	1	1	True	0.00375
	shgo-sobol	4	107	1	1	True	0.00292
	tgo	4	106	1	1	True	0.00276
	bh	4	5720	0	0	True	0.09156
	de	4	2590	0	0	True	0.09912
Shekel07	shgo-simplicial	4	134	1	1	True	0.00409
	shgo-sobol	4	123	1	1	True	0.00326
	tgo	4	122	1	1	True	0.00311
	bh	4	4365	0	0	False	0.07311
	de	4	2500	0	0	False	0.09704
	shgo-simplicial	4	142	1	1	True	0.00425
	shgo-sobol	4	131	1	1	True	0.00346
	tgo	4	130	1	1	True	0.00338
Shubert01	bh	2	3111	0	0	True	0.06710
	de	2	1467	0	0	True	0.06117
	shgo-simplicial	2	0	0	0	False	0.00000
	shgo-sobol	2	76	1	1	True	0.00353
	tgo	2	157	3	3	True	0.00549

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
Shubert03	bh	2	3000	0	0	True	0.06923
	de	2	1055	0	0	True	0.04516
	shgo-simplicial	2	0	0	0	False	0.00000
	shgo-sobol	2	51	1	1	True	0.00209
	tgo	2	51	1	1	True	0.00175
Shubert04	bh	2	3024	0	0	True	0.06949
	de	2	1175	0	0	True	0.04913
	shgo-simplicial	2	0	0	0	False	0.00000
	shgo-sobol	2	142	3	3	True	0.00491
	tgo	2	178	4	4	True	0.00544
SineEnvelope	bh	2	1416	0	0	False	0.03415
	de	2	1449	0	0	False	0.05351
	shgo-simplicial	2	9	1	1	True	0.00059
	shgo-sobol	2	7	1	1	True	0.00068
	tgo	2	7	1	1	True	0.00057
SixHumpCamel	bh	2	3030	0	0	True	0.02889
	de	2	615	0	0	True	0.01498
	shgo-simplicial	2	175	1	1	True	0.00934
	shgo-sobol	2	42	1	1	True	0.00118
	tgo	2	42	1	1	True	0.00096
Sodp	bh	2	3648	0	0	True	0.05151
	de	2	4043	0	0	True	0.13096
	shgo-simplicial	2	9	1	1	True	0.00050
	shgo-sobol	2	7	1	1	True	0.00060
	tgo	2	7	1	1	True	0.00051
Sphere	bh	2	909	0	0	True	0.01721
	de	2	3603	0	0	True	0.10187
	shgo-simplicial	2	9	1	1	True	0.00049
	shgo-sobol	2	7	1	1	True	0.00060

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
Step	tgo	2	7	1	1	True	0.00049
	bh	2	303	0	0	False	0.01252
	de	2	1083	0	0	True	0.03012
	shgo-simplicial	2	9	1	1	True	0.00048
	shgo-sobol	2	7	1	1	True	0.00060
Step2	tgo	2	7	1	1	True	0.00049
	bh	2	303	0	0	False	0.01331
	de	2	843	0	0	True	0.02558
	shgo-simplicial	2	9	1	1	True	0.00051
	shgo-sobol	2	7	1	1	True	0.00063
StretchedV	tgo	2	7	1	1	True	0.00050
	bh	2	1866	0	0	True	0.03963
	de	2	1529	0	0	True	0.05526
	shgo-simplicial	2	43	1	1	True	0.00155
	shgo-sobol	2	41	1	1	True	0.00148
StyblinskiTang	tgo	2	41	1	1	True	0.00137
	bh	2	2031	0	0	False	0.03692
	de	2	492	0	0	True	0.01594
	shgo-simplicial	2	41	1	1	True	0.00122
	shgo-sobol	2	48	1	1	True	0.00157
TestTubeHolder	tgo	2	48	1	1	True	0.00140
	bh	2	1563	0	0	False	0.02734
	de	2	852	0	0	True	0.02587
	shgo-simplicial	2	0	0	0	False	0.00000
	shgo-sobol	2	895	1	1	True	0.02486
ThreeHumpCamel	tgo	2	1101	31	30	True	0.02385
	bh	2	2247	0	0	True	0.02246
	de	2	4163	0	0	True	0.10308
	shgo-simplicial	2	9	1	1	True	0.00043

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
Treccani	shgo-sobol	2	7	1	1	True	0.00058
	tgo	2	7	1	1	True	0.00047
	bh	2	2658	0	0	True	0.02385
	de	2	2403	0	0	True	0.06263
	shgo-simplicial	2	9	1	1	True	0.00046
Trid	shgo-sobol	2	7	1	1	True	0.00059
	tgo	2	7	1	1	True	0.00047
	bh	6	9387	0	0	True	0.11500
	de	6	4178	0	0	True	0.14637
	shgo-simplicial	6	152	1	1	True	0.02563
Trigonometric01	shgo-sobol	6	98	1	1	True	0.00263
	tgo	6	94	1	1	True	0.00213
	bh	2	6288	0	0	True	0.14911
	de	2	6843	0	0	True	0.31674
	shgo-simplicial	2	9	1	1	True	0.00065
Tripod	shgo-sobol	2	7	1	1	True	0.00072
	tgo	2	7	1	1	True	0.00062
	bh	2	27252	0	0	False	0.20003
	de	2	3863	0	0	True	0.11160
	shgo-simplicial	2	163	2	2	True	0.00447
Ursem01	shgo-sobol	2	254	2	2	True	0.00440
	tgo	2	254	2	2	True	0.00404
	bh	2	1911	0	0	False	0.02379
	de	2	332	0	0	True	0.00828
	shgo-simplicial	2	22	1	1	True	0.00070
Ursem03	shgo-sobol	2	29	1	1	True	0.00092
	tgo	2	29	1	1	True	0.00081
	bh	2	5286	0	0	False	0.07254
	de	2	782	0	0	True	0.01990

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
Ursem04	shgo-simplicial	2	55	1	1	True	0.00141
	shgo-sobol	2	53	1	1	True	0.00148
	tgo	2	53	1	1	True	0.00137
	bh	2	16347	0	0	True	0.13874
	de	2	591	0	0	True	0.01374
UrsemWaves	shgo-simplicial	2	97	1	1	True	0.00174
	shgo-sobol	2	95	1	1	True	0.00185
	tgo	2	95	1	1	True	0.00174
	bh	2	420	0	0	False	0.01449
	de	2	498	0	0	False	0.01386
VenterSobieczczanskiSobieski	shgo-simplicial	2	13	2	2	True	0.00073
	shgo-sobol	2	19	1	1	True	0.00085
	tgo	2	19	1	1	True	0.00070
	bh	2	2448	0	0	False	0.03436
	de	2	655	0	0	True	0.01909
Vincent	shgo-simplicial	2	9	1	1	True	0.00049
	shgo-sobol	2	7	1	1	True	0.00061
	tgo	2	7	1	1	True	0.00050
	bh	2	2805	0	0	True	0.05658
	de	2	753	0	0	True	0.02150
Watson	shgo-simplicial	2	42	1	1	True	0.00105
	shgo-sobol	2	31	1	1	True	0.00101
	tgo	2	31	1	1	True	0.00092
	bh	6	33320	0	0	True	1.41551
	de	6	23095	0	0	True	1.64289
Wavy	shgo-simplicial	6	337	1	1	True	0.04292
	shgo-sobol	6	283	1	1	True	0.01689
	tgo	6	279	1	1	True	0.01508
	bh	2	3465	0	0	False	0.05412

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
WayburnSeader01	de	2	2603	0	0	True	0.08945
	shgo-simplicial	2	9	1	1	True	0.00055
	shgo-sobol	2	7	1	1	True	0.00065
	tgo	2	7	1	1	True	0.00054
	bh	2	6933	0	0	True	0.05294
	de	2	4823	0	0	True	0.12730
WayburnSeader02	shgo-simplicial	2	111	1	1	True	0.00171
	shgo-sobol	2	109	1	1	True	0.00185
	tgo	2	109	1	1	True	0.00174
	bh	2	6732	0	0	True	0.05500
	de	2	5043	0	0	True	0.12681
	shgo-simplicial	2	150	1	1	True	0.00218
Weierstrass	shgo-sobol	2	148	1	1	True	0.00232
	tgo	2	148	1	1	True	0.00221
	bh	2	30213	0	0	False	1.01388
	de	2	3623	0	0	True	0.21014
	shgo-simplicial	2	2225	1	1	True	0.21860
	shgo-sobol	2	0	0	0	False	0.00000
Whitley	tgo	2	0	0	0	False	0.00000
	bh	2	5244	0	0	False	0.12344
	de	2	1618	0	0	False	0.07122
	shgo-simplicial	2	34	1	1	True	0.00142
	shgo-sobol	2	32	1	1	True	0.00149
	tgo	2	32	1	1	True	0.00140
Wolfe	bh	3	1156	0	0	False	0.01560
	de	3	16444	0	0	True	0.42225
	shgo-simplicial	3	14	1	1	True	0.00116
	shgo-sobol	3	10	1	1	True	0.00068
	tgo	3	9	1	1	True	0.00055

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
XinSheYang01	bh	2	7632	0	0	False	0.09773
	de	2	6663	0	0	True	0.22097
	shgo-simplicial	2	153	1	1	True	0.00372
	shgo-sobol	2	77	1	1	True	0.00201
	tgo	2	149	1	1	True	0.00329
XinSheYang02	bh	2	2691	0	0	False	0.04671
	de	2	4223	0	0	True	0.14171
	shgo-simplicial	2	65	1	1	True	0.00168
	shgo-sobol	2	63	1	1	True	0.00177
	tgo	2	63	1	1	True	0.00167
XinSheYang03	bh	2	312	0	0	False	0.01709
	de	2	1409	0	0	True	0.05874
	shgo-simplicial	2	9	1	1	True	0.00064
	shgo-sobol	2	7	1	1	True	0.00071
	tgo	2	7	1	1	True	0.00059
XinSheYang04	bh	2	1791	0	0	False	0.04468
	de	2	1795	0	0	True	0.06512
	shgo-simplicial	2	77	1	1	True	0.00236
	shgo-sobol	2	75	1	1	True	0.00247
	tgo	2	75	1	1	True	0.00232
Xor	bh	9	7940	0	0	False	0.20712
	de	9	1520	0	0	False	0.06192
	shgo-simplicial	9	645	1	1	True	15.69089
	shgo-sobol	9	197	1	1	True	0.01734
	tgo	9	208	2	2	True	0.00647
YaoLiu04	bh	2	29316	0	0	True	0.17383
	de	2	3903	0	0	True	0.10074
	shgo-simplicial	2	23	1	1	True	0.00067
	shgo-sobol	2	21	1	1	True	0.00080

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Table 1: Comparison of the performance of SHGO, TGO, BH and DE over a wide selection of optimisation test problems. The columns are the name of the optimisation test problem, the algorithm (Alg) used, the number of dimensions (n) of the optimisation problem, the number of function evaluations (nfev), the number of minimisers generated (nlmin), the number of unique local minima (nulmin), successful convergence to the global minima and the total CPU run time (runtime).

Problem	Alg	ndim	nfev	nlmin	nulmin	success	runtime
YaoLiu09	tgo	2	21	1	1	True	0.00067
	bh	2	3300	0	0	True	0.04937
	de	2	2843	0	0	True	0.09378
	shgo-simplicial	2	20	1	1	True	0.00076
	shgo-sobol	2	18	1	1	True	0.00086
Zacharov	tgo	2	18	1	1	True	0.00075
	bh	2	2046	0	0	True	0.03903
	de	2	4043	0	0	True	0.14325
	shgo-simplicial	2	46	1	1	True	0.00184
	shgo-sobol	2	45	1	1	True	0.00150
ZeroSum	tgo	2	45	1	1	True	0.00139
	bh	2	20538	0	0	False	0.24523
	de	2	1743	0	0	False	0.05656
	shgo-simplicial	2	0	0	0	False	0.00000
	shgo-sobol	2	23	1	1	True	0.00142
Zettl	tgo	2	0	0	0	False	0.00000
	bh	2	4167	0	0	True	0.03328
	de	2	861	0	0	True	0.02009
	shgo-simplicial	2	116	1	1	True	0.00176
	shgo-sobol	2	114	1	1	True	0.00190
Zimmerman	tgo	2	114	1	1	True	0.00179
	bh	2	24543	0	0	False	0.27714
	de	2	6503	0	0	True	0.20711
	shgo-simplicial	2	3032	1	1	True	0.17307
	shgo-sobol	2	1585	1	1	True	0.03335
Zirilli	tgo	2	1585	1	1	True	0.04185
	bh	2	2562	0	0	False	0.02388
	de	2	575	0	0	True	0.01357
	shgo-simplicial	2	34	1	1	True	0.00077

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		ndim	nfev	nlmin	nulmin	success	runtime
Problem	Alg						
	shgo-sobol	2	32	1	1	True	0.00094
	tgo	2	32	1	1	True	0.00081

Bibliography