

13. $U_3 = (B \rightarrow A) \wedge (C \rightarrow A) \rightarrow (B \vee C \rightarrow A)$: Theorem \Rightarrow

U_3 is a theorem $\Leftrightarrow \text{CNF}(\neg U_3) \vdash_{\text{Res}} \square$

\downarrow
 U_3 is a theorem iff \square (the empty clause) can be derived from $\text{CNF}(\neg U_3)$ (the normal form of the negated U_3) using the RESOLUTION ALGORITHM.

$$U_3 = (B \xrightarrow{\text{I}} A) \wedge (C \xrightarrow{\text{II}} A) \xrightarrow{\text{III}} (B \vee C \xrightarrow{\text{IV}} A)$$

$$\text{I} : B \rightarrow A \equiv \neg B \vee A$$

$$\text{II} : C \rightarrow A \equiv \neg C \vee A$$

$$\text{III} : \neg((\neg B \vee A) \wedge (\neg C \vee A)) \vee (B \vee C \rightarrow A)$$

$$\text{IV} : \neg(B \vee C) \vee A$$

After the normalization of U_3 , we got:

$$U_3 = \neg((\neg B \vee A) \wedge (\neg C \vee A)) \vee (\neg(B \vee C) \vee A) \Rightarrow$$

$$\neg U_3 = \underbrace{(\neg B \vee A)}_{C_1} \wedge \underbrace{(\neg C \vee A)}_{C_2} \wedge \underbrace{(B \vee C)}_{C_3} \wedge \underbrace{\neg A}_{C_4}$$

The resolution algorithm is based on the fact that $\neg q$ and q are clashing clauses upon the literal q and resolve upon it to derive \square (empty cl.).

Therefore: The set S of clauses $\{C_1, C_2, C_3, C_4\}$ is inconsistent iff $S \vdash_{\text{Res}} \square \Leftrightarrow \vdash U_3$ iff $\text{CNF}(\neg U_3) \vdash_{\text{Res}} \square$

$$S = \{C_1 = \neg B \vee A, C_2 = \neg C \vee A, C_3 = B \vee C, C_4 = \neg A\}$$

$$S \vdash \frac{?}{\text{Res}} \square$$

$$C_1 = \neg B \vee A$$

$$C_2 = \neg C \vee A$$

$$C_3 = B \vee C$$

$$C_4 = \neg A$$

General α Resolution

$$C_5 = \text{Res}_A(C_1, C_4) = \neg B \quad (\text{Res}_A(\neg B \vee A, \neg A))$$

$$C_6 = \text{Res}_A(C_2, C_4) = \neg C \quad (\text{Res}_A(\neg C \vee A, \neg A))$$

$$C_7 = \text{Res}_B(C_3, C_5) = C \quad (\text{Res}_B(B \vee C, \neg B))$$

$$C_8 = \text{Res}_C(C_6, C_7) = \square \quad (\text{Res}_C(C, \neg C))$$

Conclusion: $\mathcal{Q}(\text{Res}_C(C_6, C_7)) = \square \Rightarrow$

$$\text{Res}_C(C, \neg C) = \square \Rightarrow$$

\mathcal{U}_3 - theorem

Upon the final clashing literal (C and $\neg C$), we denote that \mathcal{U}_3 is indeed a theorem, because we were able to prove the inconsistency of the negated clause $\neg \mathcal{U}_3$.