



DFS(1) → pt pccore of
 $nv[i] = \text{niveau}$ in onb DFS
 $nv_min[i] = \dots$

i se poate prinde in ciclu
 • direct prin m. de inlocuire:
 (1) $A = \min\{nv[i], nv[j]\}$ | i, j m. de int.
 • printr-un desc (fiu)
 (2) $B = \min\{nv_min[j]\}$ | j este
 fiu d
 $nv_min[i] = \min\{A, B, nv[i]\}$

act. dupa rel (2)
 m.c.
 $nv_min[j] > nv[i]$
 $nv_min[j] \geq nv[i]$

PC:
 (3) fiu
 • radacina → $R^c (=)$ au ≥ 2 fi

Ex: Not $x_4 \rightarrow$ nr de fete de gr 4
 $x_6 \rightarrow$ nr de fete de gr 6
 $|F| = x_4 + x_6$

$$2m = \sum_{f \in F} d_m(f) = \underbrace{4 + \dots + 4}_{x_4 \text{ ori}} + \underbrace{6 + \dots + 6}_{x_6 \text{ ori}} = 4x_4 + 6x_6$$

$$\sum_v d_m(v) = 3n = 2m$$

$$\begin{cases} n - m + \overbrace{x_4 + x_6}^{|F|} = 2 \\ 4x_4 + 6x_6 = 2m \end{cases} \rightarrow 3n = 2m$$

