Table of Contents

[1. The Compiler Phases 9](#_Toc481936105)

[1.1.1. The Preprocessor 10](#_Toc481936106)

[1.1.2. Scanning and Parsing 10](#_Toc481936107)

[1.1.1. Syntax Tree Optimization 10](#_Toc481936108)

[1.1.2. Middle Code Generation 10](#_Toc481936109)

[1.1.3. Middle Code Optimization 11](#_Toc481936110)

[1.1.4. Register Preparation and Allocation 12](#_Toc481936111)

[1.1.5. Register Allocation 12](#_Toc481936112)

[1.1.6. Object Code Generation 12](#_Toc481936113)

[1.1.1. Object Code Optimization 12](#_Toc481936114)

[1.1.2. Linking 12](#_Toc481936115)

[1.1.3. The Standard Library 12](#_Toc481936116)

[2. An Introduction to Bison and JFlex 13](#_Toc481936117)

[2.1.1. The Language 13](#_Toc481936118)

[2.1.2. The Grammar 13](#_Toc481936119)

[2.1.3. Bison 15](#_Toc481936120)

[2.1.4. JFlex 19](#_Toc481936121)

[2.1.5. Main 22](#_Toc481936122)

[3. Main 24](#_Toc481936123)

[3.1. Reading the Source File 25](#_Toc481936124)

[3.2. Checking the Object File 26](#_Toc481936125)

[3.3. Reading the Object File 26](#_Toc481936126)

[3.4. The main method 26](#_Toc481936127)

[4. The Preprocessor 29](#_Toc481936128)

[4.1. The Expression Parser 29](#_Toc481936129)

[4.1.1. The Grammar 29](#_Toc481936130)

[4.1.2. The Scanner 33](#_Toc481936131)

[4.2. The Preprocessor 36](#_Toc481936132)

[4.2.1. Tri Graphs 37](#_Toc481936133)

[4.2.1. Comments, Strings, and Characters 38](#_Toc481936134)

[4.2.2. Slash Codes 40](#_Toc481936135)

[4.2.3. The Line List 43](#_Toc481936136)

[4.2.4. Lines 46](#_Toc481936137)

[4.2.5. Includes 46](#_Toc481936138)

[4.2.6. Macros 47](#_Toc481936139)

[4.2.7. Conditional Programming 51](#_Toc481936140)

[4.2.8. Macro Expansion 53](#_Toc481936141)

[4.2.9. String Merging 56](#_Toc481936142)

[5. Scanning 58](#_Toc481936143)

[5.1.1. The typedef-name Problem 58](#_Toc481936144)

[5.2. The Scanner 58](#_Toc481936145)

[6. Parsing 64](#_Toc481936146)

[6.1. Declarations 65](#_Toc481936147)

[6.1.1. Function Definitions 65](#_Toc481936148)

[6.1.2. Declarations 69](#_Toc481936149)

[6.1.3. Struct and Union 75](#_Toc481936150)

[6.1.4. Enum 77](#_Toc481936151)

[6.1.5. Declarator 78](#_Toc481936152)

[6.1.6. Pointers 82](#_Toc481936153)

[6.1.7. Parameter List 84](#_Toc481936154)

[6.1.8. Initializer 85](#_Toc481936155)

[6.1.9. Type Name 86](#_Toc481936156)

[6.1.10. Direct and Abstract Declarator 86](#_Toc481936157)

[6.2. Statements 90](#_Toc481936158)

[6.2.1. The Forward-Jump Problem (Backpatching) 90](#_Toc481936159)

[6.2.2. The if-else Problem (Open and Closed Statements) 92](#_Toc481936160)

[6.2.3. Header Rules 93](#_Toc481936161)

[6.2.4. If Statements 94](#_Toc481936162)

[6.2.5. Switch Statements 94](#_Toc481936163)

[6.2.6. Case Statements 95](#_Toc481936164)

[6.2.7. While Statements 96](#_Toc481936165)

[6.2.8. For Statements 97](#_Toc481936166)

[6.2.9. Label Statements 97](#_Toc481936167)

[6.2.10. Closed statement 98](#_Toc481936168)

[6.2.11. Do Statements 98](#_Toc481936169)

[6.2.12. Block Statements 99](#_Toc481936170)

[6.2.13. Default Statements 99](#_Toc481936171)

[6.2.14. Break Statements 100](#_Toc481936172)

[6.2.1. Continue Statements 100](#_Toc481936173)

[6.2.2. Goto Statements 100](#_Toc481936174)

[6.2.3. Return Statements 101](#_Toc481936175)

[6.2.4. Expression Statements 102](#_Toc481936176)

[6.2.1. Declaration Statements 102](#_Toc481936177)

[6.2.2. Load Register Statements 102](#_Toc481936178)

[6.2.3. Store Register Statements 103](#_Toc481936179)

[6.2.4. Clear Registers Statements 104](#_Toc481936180)

[6.2.1. Store Flagbyte Statements 104](#_Toc481936181)

[6.2.1. Jump Register Statements 104](#_Toc481936182)

[6.2.2. Interrupt Statements 105](#_Toc481936183)

[6.3. Expressions 105](#_Toc481936184)

[6.3.1. Precedence and Associativity 105](#_Toc481936185)

[6.3.2. Syntax Tree Generation Methods 106](#_Toc481936186)

[6.3.3. Comma Expression 106](#_Toc481936187)

[6.3.4. Logical Expression 107](#_Toc481936188)

[6.3.5. Integer Expressions 108](#_Toc481936189)

[6.3.6. Integral and Floating Promotion 109](#_Toc481936190)

[6.3.7. Arithmetic Conversion 109](#_Toc481936191)

[6.3.8. Assignment Expression 110](#_Toc481936192)

[6.3.9. Conditional Expression 110](#_Toc481936193)

[6.3.10. Constant Expression 111](#_Toc481936194)

[6.3.11. Logical Expression 111](#_Toc481936195)

[6.3.12. Bitwise Expression 112](#_Toc481936196)

[6.3.13. Equality Expressions 112](#_Toc481936197)

[6.3.14. Relation Expressions 113](#_Toc481936198)

[6.3.15. Shift Expression 114](#_Toc481936199)

[6.3.16. Arithmetic Expression 115](#_Toc481936200)

[6.3.17. Cast Expression 116](#_Toc481936201)

[6.3.18. Prefix Expression 116](#_Toc481936202)

[6.3.19. Unary Arithmetic Expression 117](#_Toc481936203)

[6.3.1. Prefix Increment and Decrement 117](#_Toc481936204)

[6.3.2. Logical Not 117](#_Toc481936205)

[6.3.3. Bitwise Not 118](#_Toc481936206)

[6.3.4. sizeof 118](#_Toc481936207)

[6.3.5. Address Expression 119](#_Toc481936208)

[6.3.6. Dereference Expression 119](#_Toc481936209)

[6.3.7. Postfix Expression 119](#_Toc481936210)

[6.3.8. Postfix Increment and Decrement 120](#_Toc481936211)

[6.3.9. Dot 120](#_Toc481936212)

[6.3.10. Arrow 121](#_Toc481936213)

[6.3.11. Index 121](#_Toc481936214)

[6.3.12. Function Call 122](#_Toc481936215)

[6.3.1. Parameter Argument List 123](#_Toc481936216)

[6.3.2. Primary Expression 124](#_Toc481936217)

[7. The Type System 125](#_Toc481936218)

[7.1. The Type Class 125](#_Toc481936219)

[7.1.1. Arithmetic Types 125](#_Toc481936220)

[7.1.2. Enumeration Types 126](#_Toc481936221)

[1.1.1. Bitfield Types 126](#_Toc481936222)

[7.1.3. Logical Types 126](#_Toc481936223)

[7.1.4. Pointer Types 127](#_Toc481936224)

[7.1.5. Array Types 127](#_Toc481936225)

[7.1.6. Function Types 128](#_Toc481936226)

[7.1.7. Struct and Union Types 129](#_Toc481936315)

[Type 130](#_Toc481936316)

[7.1.8. Size 130](#_Toc481936354)

[7.1.9. Completeness 131](#_Toc481936355)

[7.1.1. Constant and Volatile Types 132](#_Toc481936356)

[7.1.2. Equality 132](#_Toc481936357)

[7.1.3. Test Methods 133](#_Toc481936398)

[7.2. Type Casting 136](#_Toc481936399)

[7.3. Type Promotion 140](#_Toc481936400)

[7.4. The Value Conversion Class 142](#_Toc481936401)

[8. The Symbol Table 147](#_Toc481936402)

[8.1. The Symbol Class 147](#_Toc481936403)

[8.1.1. Symbol Storage 147](#_Toc481936404)

[8.2. The Symbol Table Class 158](#_Toc481936405)

[8.2.1. Add Symbols 162](#_Toc481936406)

[8.2.2. Looking Up Symbols 162](#_Toc481936407)

[1.1.1. Generate Offset 163](#_Toc481936408)

[1.1.1. Table Text 163](#_Toc481936409)

[9. The Syntax Tree 164](#_Toc481936410)

[10. Syntax Tree Optimization 167](#_Toc481936411)

[11. The Middle Code 183](#_Toc481936412)

[12. Middle Code Generation 194](#_Toc481936413)

[13. Initializer 216](#_Toc481936414)

[13.1.1. Static Initialization 216](#_Toc481936415)

[13.1.2. Stack Initialization 228](#_Toc481936416)

[14. Middle Code Optimization 233](#_Toc481936417)

[14.1. Optimization 233](#_Toc481936418)

[14.1.1. Next Goto Statements 234](#_Toc481936419)

[14.1.2. Next-Double Goto Statements 234](#_Toc481936420)

[14.1.3. Goto Chains 234](#_Toc481936421)

[14.1.4. Clear Unreachable Code 235](#_Toc481936422)

[14.1.5. Remove Cleared Code 236](#_Toc481936423)

[14.2. Generation 237](#_Toc481936424)

[14.2.1. Base Blocks Generation 237](#_Toc481936425)

[14.2.2. Live Sets Generation 239](#_Toc481936426)

[14.2.3. Path Sets 241](#_Toc481936427)

[14.2.4. Removal of Unused Code 244](#_Toc481936428)

[15. Object Code Preparation 245](#_Toc481936429)

[15.1. Register Preparation 245](#_Toc481936430)

[15.2. Register Allocation 265](#_Toc481936431)

[16. Object Code Generation 269](#_Toc481936432)

[16.1. The Object Operator 269](#_Toc481936433)

[16.2. The Object Code 269](#_Toc481936434)

[16.2.1. Object Code Generation 283](#_Toc481936435)

[16.3. Object Code Generation 283](#_Toc481936436)

[17. The Linker 322](#_Toc481936437)

[17.1. The Linker Class 322](#_Toc481936438)

[17.1.1. Generation 323](#_Toc481936439)

[17.1.2. Main Trace Generation 325](#_Toc481936440)

[18. Auxiliary Classes 329](#_Toc481936441)

[18.1. Error Handling 329](#_Toc481936442)

[18.2. Container Classes 330](#_Toc481936443)

[18.2.1. Ordered Pair 330](#_Toc481936444)

[18.2.2. Unordered Pair 331](#_Toc481936445)

[18.2.3. Triple 332](#_Toc481936446)

[18.2.4. Tuple 333](#_Toc481936447)

[18.3. Undirected Graph 333](#_Toc481936448)

[18.3.1. Addition and Removal of Vertices and Edges 334](#_Toc481936449)

[18.3.2. Graph Partition 335](#_Toc481936450)

[18.4. Directed Graph 336](#_Toc481936451)

[19. The Standard Library 339](#_Toc481936452)

[19.1. Integral and Floating Limits 339](#_Toc481936453)

[19.2. The Assert Macro 340](#_Toc481936454)

[19.3. Locale Data 340](#_Toc481936455)

[19.4. Character Types 342](#_Toc481936456)

[19.5. Strings 344](#_Toc481936457)

[19.6. Long Jumps 350](#_Toc481936458)

[19.7. Mathematical Functions 351](#_Toc481936459)

[19.8. Standard Input/Output 359](#_Toc481936460)

[19.8.1. Printing 359](#_Toc481936461)

[19.8.2. Scanning 371](#_Toc481936462)

[19.8.3. File Handling 381](#_Toc481936463)

[19.9. The Standard Library 391](#_Toc481936464)

[19.10. Time 396](#_Toc481936465)

[20. Main.java 404](#_Toc481936466)

[21. The C Grammar 405](#_Toc481936467)

[21.1. The Preprocessor Grammar 405](#_Toc481936468)

[21.2. The Language Grammar 405](#_Toc481936469)

[22. A Crash Course in C 412](#_Toc481936470)

[22.1. The Compiler and the Linker 412](#_Toc481936471)

[22.2. The Hello-World Program 413](#_Toc481936472)

[22.3. Comments 414](#_Toc481936473)

[1.1.1 Types and Variables 414](#_Toc481936474)

[2.1.1 Simple Types 415](#_Toc481936475)

[3.1.1 Variables 415](#_Toc481936476)

[4.1.1 Constants 416](#_Toc481936477)

[5.1.1 Output 416](#_Toc481936478)

[6.1.1 Input 417](#_Toc481936479)

[7.1.1 Enumerations 418](#_Toc481936480)

[8.1.1 Arrays 419](#_Toc481936481)

[9.1.1 Characters and Strings 419](#_Toc481936482)

[10.1.1 The ASCII Table 420](#_Toc481936483)

[11.1.1 Pointers 420](#_Toc481936484)

[12.1.1 Pointers and Dynamic Memory 421](#_Toc481936485)

[13.1.1 Defining our own types 423](#_Toc481936486)

[14.1.1 The Type Size 424](#_Toc481936487)

[22.4. Expressions and Operators 424](#_Toc481936488)

[15.1.1 Arithmetic Operators 424](#_Toc481936489)

[16.1.1 Pointer Arithmetic 425](#_Toc481936490)

[17.1.1 Increment and decrement 426](#_Toc481936491)

[18.1.1 Relational Operators 426](#_Toc481936492)

[19.1.1 Logical Operators 426](#_Toc481936493)

[20.1.1 Bitwise Operators 427](#_Toc481936494)

[21.1.1 Assignment 428](#_Toc481936495)

[22.1.1 The Condition Operator 428](#_Toc481936496)

[23.1.1 Precedence and Associatively 428](#_Toc481936497)

[22.5. Statements 429](#_Toc481936498)

[24.1.1 Selection statements 429](#_Toc481936499)

[25.1.1 Iteration statements 431](#_Toc481936500)

[26.1.1 Jump statements 433](#_Toc481936501)

[27.1.1 Expression statements 433](#_Toc481936502)

[28.1.1 Volatile 433](#_Toc481936503)

[22.6. Functions 433](#_Toc481936504)

[22.6.1. Void Functions 435](#_Toc481936505)

[22.6.2. Local and global variables 435](#_Toc481936506)

[22.6.3. Inner Blocks with Local Variables 436](#_Toc481936507)

[22.6.4. Call-by-Value and Call-by-Reference 437](#_Toc481936508)

[22.6.5. Static, Extern, and Register Variables 439](#_Toc481936509)

[22.6.6. Recursion 440](#_Toc481936510)

[22.6.7. Definition and Declaration 440](#_Toc481936511)

[22.6.8. Higher-Order Functions 441](#_Toc481936512)

[22.7. Structures and Linked Lists 443](#_Toc481936513)

[22.7.1. Stacks and Linked Lists 443](#_Toc481936514)

[22.7.2. Unions 446](#_Toc481936515)

[22.7.3. Bitfields 446](#_Toc481936516)

[22.8. Structures with function pointers 447](#_Toc481936517)

[22.9. The Preprocessor 447](#_Toc481936518)

[22.10. The Standard Library 448](#_Toc481936519)

[22.10.1. File Management 448](#_Toc481936520)

[22.10.2. Program Parameters 450](#_Toc481936521)

[22.10.3. Environment Functions 450](#_Toc481936522)

[22.10.4. Searching and Sorting 451](#_Toc481936523)

[22.10.5. Functions with Variable Number of Parameters 452](#_Toc481936524)

[22.10.6. String Management 452](#_Toc481936525)

[22.10.7. Memory Management 454](#_Toc481936526)

[22.10.8. Mathematical Functions 454](#_Toc481936527)

[22.10.9. Long Jumps 454](#_Toc481936528)

[22.10.10. Time 455](#_Toc481936529)

[22.10.11. Random Numbers 456](#_Toc481936530)

[22.10.12. Limits of Integral and Floating Types 456](#_Toc481936531)

[22.10.13. Character Functions 457](#_Toc481936532)

[22.11. Further Reading 458](#_Toc481936533)

[23. A Crash Course in Java 459](#_Toc481936534)

[Foreword 460](#_Toc481936535)

# The Compiler Phases

Register Preparator

Register Allocator

Parser

Middle Code with Register Information

Object Code Generation

Object Code Optimizer

Object Code

Optimized Object Code

Middle Code with Allocated Registers

### The Preprocessor

First, we need to perform preprocessing: comments are removed, files are included, macros (with or without parameters) are expanded, the conditional programming is evaluated.

### Scanning and Parsing

The scanner is responsible for putting together sequences of characters into tokens, which is the least significant parts of the source code. For instance, the characters ‘i’ and ‘f’ are interpreted as the key word if and the characters ‘1’, ‘2’, and ‘3’ are interpreted as the integer 123.

Every programming language has a syntax, which in some cases is defined by a grammar. The C grammar of this books is described in Chapter 21. The parser checks whether the source code applies to the grammar by requesting tokes from the scanner when needed. The C grammar can be divided into three parts: declarations and definitions, statements and expressions. When the declarations are parsed the symbol table is generated, which holds information of the declared variables, types, and functions that is used when parsing the statements and expressions. During the parsing, comprehensive type checking and type conversions are also performed to determine whether the source code is semantically correct.

The output of the parser is a syntax holding the structure of the program, with references to the symbols of the symbol table.

### Syntax Tree Optimization

The purpose of the syntax tree optimization phase is to make it as small as possible. To some extent, it reduces the tree caused by the code of the programmer. In the following example, the side effects (function call and increment) are kept, but the addition is removed.

|  |  |
| --- | --- |
| f(x) + i++; | f(x); i++; |
| (a) Before | (b) After |

However, in many cases the it is subtrees added by the parser that are removed. For instance, the C standard function printf returns an integer, which in many cases are ignored. The parser stores the value in a temporary variable, which is removed in case it is not used.

|  |  |
| --- | --- |
| temp = printf(“Hello”); | printf(“Hello”); |
| (a) Before | (b) After |

### Middle Code Generation

When the syntax tree has been appropriate optimized, is translated into middle code list. Each middle code instruction is made of three-address-code, holding one operator and three addresses. Each address may refer to a name, a symbol, a value, or a jump address.

ymbols and values. As the name implies, the instruction it is The output of sequence of middle code. That is, a simple code format suitable for further optimization and object code generation. In this book, we use three-address-code. As the name implies, the code is made up of an operator and three addresses, which are symbols or jump targets.

### Middle Code Optimization

The purpose of the middle code optimization phase is to make the middle code more effective. One example is the next-jump instruction. Naturally, an instruction that jumps to the next instruction shall be removed:

|  |  |
| --- | --- |
| 1. goto 2  2. ... | 1. empty  2. ... |
| (a) Before | (b) After |

In the same way, if a conditional instruction jump can be reduced:

|  |  |
| --- | --- |
| 1. if a < b goto 3  2. goto 10  3. ... | 1. if a >= b goto 10  2. empty  3. ... |
| (a) Before | (b) After |

Goto-chains are also reduced. In the following example, all goto instructions jumps to the final target:

|  |  |
| --- | --- |
| 1. goto 3  2. ...  3. goto 5  4. ...  5. goto 10 | 1. goto 10  2. ...  3. goto 10  4. ...  5. goto 10 |
| (a) Before | (b) After |

Finally, unreachable code is reduced. That is, every instruction that cannot be reached from the first instruction of the function is removed:

if (0) printf(“False”); else printf(“True”);

|  |  |
| --- | --- |
| 1. goto 2  2. // *call printf(“False”);*  3. goto 5  4. // *call printf(“True”);*  5. ... | 1. empty  2. // *call printf(“False”);*  3. empty  4. empty  5. ... |
| (a) Before | (b) After |

simplify the middle code. When the middle code has been generated, the functions holds code that is not reachable and variables not used, introduced both by the programmer and the scanning and parsing phase above, that need to be removed. Moreover, there are also complicated jump instruction that can be simplified.

This phase is also responsible for the generation of base blocks and live sets. A base block is the longest sequence of middle code instructions that is not the target of a jump instruction. A live set is a set connected to each middle code instruction that holds the variable accessed later on in the same base block, that information is used when generating object code in the next phase.

### Register Preparation and Allocation

Before we do the actual register allocation, we traverse the middle code list and track the use of (so far anonyms) registers. The result is a set of register usages that forms the vertex set of a graph. An edge is added for each pair of usages that overlaps. Then the registers are allocated by a graph-coloring algorithm.

### Register Allocation

A graph is a mathematical structure made up by a non-empty set of vertices, where some pairs are connected by edges. Two vertices connected by an edge are neighborhoods. A graph can be colored in such way that that no neighbors have the same color. The color number of a graph is the minimum number of colors needed to color the graph. When allocating the registers we use a graph, where the tracks are the vertices and two vertices are connected by an edge if the tracks overlaps. Then we use graph coloring to assign each track a register.

### Object Code Generation

When the registers have been allocated, we are ready to generate the object code. In this book, there are actually two versions of object code. This phase generates both 16-bits executable com-code and assembler code that can be further assembled into com-code. The reason for the generation of both assembler code and com-code is that it is virtually impossible to debug the generated com-code, it is much easier to study the assembler code instead. When generating the target code we use the registers allocated, we also use registers holding address to values.

### Object Code Optimization

When the object code has been generated, it needs to be optimized. Foremost, we need to consider the conditional jump instructions. An unconditional jump instruction may be eight of sixteen bits, while a condition jump is always eight bits. We also translates addition or subtraction with one into incremement or decrement, and we reduse pop-push-chains.

### Linking

When the source files have compiled into object files they need to be linked together into an executable file. The task of the linker is to modify all references to static and extern variables and functions as well as make sure that only the parts reachable from the main functions are included in the final executable file.

### The Standard Library

The C language comes with a set of functions, macros, and instructions for many services. It is mostly implemented in C. However, there are some additional non-standard functions for accessing registers.

# An Introduction to Bison and JFlex

In this chapter, we use the JFlex scanner generator and the Bison parser generator, and construct a simple programming language.

### The Language

A program is made up by a non-empty sequence of statements, each terminated by a semicolon. There are five kinds of statements:

1. Read. A non-empty sequence of variables are assigned values read from the input stream, with an optional prompt.
2. Write. The non-empty sequence of values are is written to the output stream, with an optional text.
3. Assign. A variable is assigned the value of an expression.

The mathematical constants pi and e are stored in the variable map when the execution starts. An expression can be made up of the four rules of arithmetic, parentheses as well as the functions sin, cos, tan, log, exp, log10, exp10, and sqrt (square root). Division by zero, invalid function arguments, or invalid input generate error messages.

Below is an example:

// The area and circumference of a circle

read "Input the radius of a circle: ", radius;

assign area = pi \* radius \* radius;

assign circumference = 2 \* pi \* radius;

write "The area and circumference are: ", area, circumference;

newline;

// Fahrenheit to Celsius

read "Input a temperature in degrees Fahrenheit: ", farenheit;

assign celsius = (farenheit - 32) / 1.8;

write "In Celsius degrees: ", celsius;

newline;

// Pythagorean Theorem

read "Input the two cathetuses of a right-angled triangle: ", a, b;

assign c = sqrt(a \* a + b \* b);

write "The hypothenuse is: ", c;

When executed (the underlined text represent input values):

Input the radius of a circle: 10

The area and circumference are: 62.83185307179586, 314.1592653589793

Input a temperature in degrees Fahrenheit: 212

In Celsius degrees: 100.0

Input the two cathetuses of a right-angled triangle: 3,4

The hypothenuse is: 5.0

### The Grammar

Every programming language has a syntax, which may be defined by a grammar. The grammar of our language is given below. A grammar is made up by a set of *rules*, among which one is the start rule (in this case statement\_list). The grammar rules for our language follows below, the vertical bar at the left of the rules can be read as “or”. The first rule statement\_list can be read as “a statement list is a statement or a statement followed by a statement list.”

statement\_list ::=

statement

| statement statement\_list

statement ::=

**assign** **identifier** EQUAL expression **;**

| **read** optional\_output identifier\_list **;**

| **write** optional\_output expression\_list **;**

| **newline** **;**

optional\_output ::=

/\* Empty. \*/

| **string** **,**

identifier\_list ::=

**identifier**

| **identifier** **,** identifier\_list

expression\_list ::=

expression

| expression **,** expression\_list

expression ::=

binary\_expression

| expression **+** binary\_expression

| expression **-** binary\_expression

binary\_expression ::=

unary\_expression

| binary\_expression **\*** unary\_expression

| binary\_expression **/** unary\_expression

unary\_expression ::=

primary\_expression

| **+** unary\_expression

| **-** unary\_expression

| **sin** unary\_expression

| **cos** unary\_expression

| **tan** unary\_expression

| **log** unary\_expression

| **exp** unary\_expression

| **log10** unary\_expression

| **exp10** unary\_expression

| **sqrt** unary\_expression

primary\_expression ::=

**(** expression **)**

| **identifier**

| **value**

The words and character in boldface is the end products of the grammar. The rules shall be applied in such a way that finally only tokens remain. The first line of the example in the previous section are made up of the following tokens:

**read string , identifier ;**

### Bison

Bison is a tool based on Yacc[[1]](#footnote-1), which is a classic parser generator. Even though it is possible to write a parser in pure Java, it is easier to use a tool like Bison. Simply put, Bison takes a grammar as input and generators Java code as output. However, it does not only check whether the input string comply with the grammar, each rule is also allowed to perform actions, defined by Java code.

The parser is divided into two parts: one part stating the rules and tokens of the grammar, and one part defining the grammar rules. The tokens written in boldface in the previous section are written in capital letters in CUP.

MainParser.gppg

%{

package calculator;

import java.io.\*;

import java.util.\*;

class ErrorMessage {

public static void print(String message) {

Main.ErrorStream.println(message);

System.exit(-1);

}

}

%}

%name-prefix = ""

The following tokens do not holds any attributes.

%token ASSIGN READ WRITE NEWLINE EQUAL PLUS MINUS TIMES DIVIDE

SIN COS TAN LOG EXP LOG10 EXP10 SQRT

LEFT\_PAREN RIGHT\_PAREN COMMA SEMICOLON

STRING, IDENTIFIER, and VALUE are tokens with attributes.

%token <String> STRING IDENTIFIER

%token <Double> VALUE

%type <List> identifier\_list

%type <List> expression\_list

%type <Double> expression binary\_expression unary\_expression primary\_expression

There are also the rules statement\_list, statement, and optional\_output in the rules section below. However, they do not return values and do not need to be defined in this section.

%start statement\_list;

%%

statement\_list:

statement

| statement statement\_list;

The right-hand side are numbered from one, and the attribute values are accessable with the dollar-notation. In the rule below, the value of the identifier is stored in $2 and the value of the expression is stored in $4.

statement:

ASSIGN IDENTIFIER EQUAL expression SEMICOLON {

Main.VariableMap.put($2, $4);

}

| READ optional\_output identifier\_list SEMICOLON {

try {

String buffer = Main.BufferedReader.readLine();

String textArray[] = buffer.split(",");

if ($3.size() != textArray.length) {

ErrorMessage.print("Invalid number of values.");

}

for (int index = 0; index < $3.size(); ++index) {

String name = (String) $3.get(index), text = textArray[index];

try {

double value = Double.valueOf(text);

Main.VariableMap.put(name, value);

}

catch (Exception exception) {

ErrorMessage.print("Invalid input.");

}

}

}

catch (IOException exception) {

ErrorMessage.print(exception.getMessage());

}

}

| WRITE optional\_output expression\_list SEMICOLON {

boolean first = true;

for (Double value : (List<Double>) $3) {

Main.OutputStream.print((first ? "" : ", ") + value);

first = false;

}

Main.OutputStream.println();

}

| NEWLINE SEMICOLON {

Main.OutputStream.println();

};

optional\_output:

/\* Empty. \*/

| STRING COMMA {

Main.OutputStream.print($1);

};

The value of the rule is stored in $$. However, the $$ value is always a reference to an Object. Therefore, we need to transform it into a list before we can call the add method.

identifier\_list:

IDENTIFIER {

$$ = new LinkedList<String>();

((List<String>) $$).add($1);

}

| IDENTIFIER COMMA identifier\_list {

$$ = $3;

((List<String>) $$).add($1);

};

expression\_list:

expression {

$$ = new LinkedList<Double>();

((List<Double>) $$).add($1);

}

| expression COMMA expression\_list {

$$ = $3;

((List<Double>) $$).add($1);

};

expression:

binary\_expression {

$$ = $1;

}

| expression PLUS binary\_expression {

$$ = $1 + $3;

}

| expression MINUS binary\_expression {

$$ = $1 - $3;

};

binary\_expression:

unary\_expression {

$$ = $1;

}

| binary\_expression TIMES unary\_expression {

$$ = $1 \* $3;

}

| binary\_expression DIVIDE unary\_expression {

if ($3 == 0) {

ErrorMessage.print("Division by Zero.");

}

$$ = $1 / $3;

};

unary\_expression:

primary\_expression {

$$ = $1;

}

| PLUS unary\_expression {

$$ = $2;

}

| MINUS unary\_expression {

$$ = -$2;

}

| SIN unary\_expression {

$$ = Math.sin($2);

}

| COS unary\_expression {

$$ = Math.cos($2);

}

| TAN unary\_expression {

$$ = Math.tan($2);

}

| LOG unary\_expression {

if ($2 <= 0) {

ErrorMessage.print("Logarithm of Non-Positive Value.");

}

$$ = Math.log($2);

}

| EXP unary\_expression {

$$ = Math.exp($2);

}

| LOG10 unary\_expression {

if ($2 <= 0) {

ErrorMessage.print("Logarithm of Non-Positive Value.");

}

$$ = Math.log10($2);

}

| EXP10 unary\_expression {

$$ = Math.pow(10, $2);

}

| SQRT unary\_expression {

if ($2 < 0) {

ErrorMessage.print("Square Root of Negativ Value.");

}

$$ = Math.sqrt($2);

};

primary\_expression:

LEFT\_PAREN expression RIGHT\_PAREN {

$$ = $2;

}

| IDENTIFIER {

if (Main.VariableMap.containsKey($1)) {

$$ = Main.VariableMap.get($1);

}

else {

ErrorMessage.print("Unknown Name: \"" + $1 + "\".");

}

}

| VALUE {

$$ = $1;

};

### JFlex

For the parser of the previous section to work properly it also needs a scanner that identifies the tokens. JFlex is a scanner-generator, which task is to identify and transform groups of characters into tokens.

First we need to Token class that handles a token generated by the scanner, it holds its identity and a potential attribute value.

Token.java

package calculator;

public class Token {

public int m\_id;

public Object m\_value;

public Token(int id) {

m\_id = id;

m\_value = null;

}

public Token(int id, Object value) {

m\_id = id;

m\_value = value;

}

public int getId() {

return m\_id;

}

public Object getValue() {

return m\_value;

}

}

When it comes to complicated tokens, JFlex applies regular expressions. Some characters have special meanings in accordance with the following table. However, their special status can be canceled by preceeding it with a backslash.

|  |  |
| --- | --- |
| **Character** | **Meaning** |
| ( and ) | Grouping subexpressions |
| [ and ] | Any character within the brackets |
| [^ and ] | Any character except the characters within the brackets |
| \* | Zero or more characters |
| + | One or moree characters |
| . | Any character except new line |

Scanner.jflex

package calculator;

%%

%public

%class Scanner

%type Token

%notunix

//%unicode

%line

%char

%eofval{

return null;

%eofval}

%yylexthrow{

java.lang.Exception

%yylexthrow}

%eofval{

return null;

%eofval}

%yylexthrow{

java.lang.Exception

%yylexthrow}

%{

public Parser.Lexer toLexer() {

return (new ScannerToLexer(this));

}

%}

In the scanner code below, a line comment is two slashes followed by zero or more occurrences of any character except newline. Since both the slash has special meaning we must insert a backslash before it for it to mean just a slash.

LINE\_COMMENT = \/\/.\*

A string is a double quotation mark, followed by zero or more occurrences of any character (except double quotation mark) followed by a double quotation mark.

STRING = \"[^\"]\*\"

An identifier is a capital or small letter, or a underline, followed by zero or more capital or small letters, or digits, or underlines.

IDENTIFIER = [a-zA-Z\_][a-zA-Z0-9\_]\*

A value is as least one digit, potentially followed by more digits, a dot, and even more digits.

VALUE = [0-9]+|([0-9]+\.[0-9]+)

A white space is either a space, a tabulator, a return, a newline, or a form-feed.

WHITE\_SPACE = [ \t\r\n\f]

%%

"assign" { return (new Token(Parser.ASSIGN)); }

"read" { return (new Token(Parser.READ)); }

"write" { return (new Token(Parser.WRITE)); }

"newline" { return (new Token(Parser.NEWLINE)); }

"sin" { return (new Token(Parser.SIN)); }

"cos" { return (new Token(Parser.COS)); }

"tan" { return (new Token(Parser.TAN)); }

"log" { return (new Token(Parser.LOG)); }

"exp" { return (new Token(Parser.EXP)); }

"log10" { return (new Token(Parser.LOG10)); }

"exp10" { return (new Token(Parser.EXP10)); }

"sqrt" { return (new Token(Parser.SQRT)); }

"=" { return (new Token(Parser.EQUAL)); }

"+" { return (new Token(Parser.PLUS)); }

"-" { return (new Token(Parser.MINUS)); }

"\*" { return (new Token(Parser.TIMES)); }

"/" { return (new Token(Parser.DIVIDE)); }

"(" { return (new Token(Parser.LEFT\_PAREN)); }

")" { return (new Token(Parser.RIGHT\_PAREN)); }

"," { return (new Token(Parser.COMMA)); }

";" { return (new Token(Parser.SEMICOLON)); }

The yytext method return the scanned text, which is useful with identifiers and strings.

{IDENTIFIER} {

return (new Token(Parser.IDENTIFIER, yytext()));

}

{STRING} {

String text = yytext();

return (new Token(Parser.STRING, text.substring(1, text.length() - 1)));

}

{VALUE} {

return (new Token(Parser.VALUE, Double.valueOf(yytext())));

}

{LINE\_COMMENT} {

// Empty.

}

{WHITE\_SPACE} {

// Empty.

}

Finally, the last part of the scanner is the dot rule, which catch the case not caught by any of the previous tokens (equivalent to default in a switch statement). In that case, we just exit the execution with an error message.

. {

Main.ErrorStream.println("Unknown character: \'" + yytext() + "\'.");

System.exit(-1);

}

The ScannerToLexer class is necessary to make the parser work with the scanner. It works as link between the scanner and the parser and implements the Lexer interface in the Parser class. The yylex method reads the next token from the scanner and returns it in the form preferred by the parser.

ScannerToLexer.java

package calculator;

public class ScannerToLexer implements Parser.Lexer {

private Scanner m\_scanner;

private Object m\_nextValue;

public ScannerToLexer(Scanner scanner) {

m\_scanner = scanner;

}

public int yylex() throws java.io.IOException {

try {

Token nextToken = m\_scanner.yylex();

if (nextToken != null) {

m\_nextValue = nextToken.getValue();

return nextToken.getId();

}

}

catch (Exception exception) {

exception.printStackTrace();

yyerror(exception.getMessage());

}

return 0;

}

public Object getLVal () {

return m\_nextValue;

}

public void yyerror (String message) {

Main.ErrorStream.println(message);

System.exit(-1);

}

}

### Main

Bison and JFlex generates the Parser and Scanner classes. The main class creates a scanner object which is used by the parser object. The parse method is called with invoke the parsing. The input is read from BufferedReader, the output is written to OutputStream, error messages are written to ErrorStream, and VariableMap is used to keep track of the variables read and assigned by the program.

Main.java

package calculator;

import java.io.\*;

import java.util.\*;

public class Main {

private static InputStream InputStream = System.in;

public static PrintStream OutputStream = System.out;

public static PrintStream ErrorStream = System.err;

public static BufferedReader BufferedReader = null;

public static final Map<String,Double> VariableMap = new HashMap<>();

public static void main(String args[]) {

if (args.length != 1) {

ErrorStream.printf("Usage: calculator inputfile");

}

try {

InputStreamReader reader = new InputStreamReader(Main.InputStream);

BufferedReader = new BufferedReader(reader);

InputStreamReader scannerReader =

new InputStreamReader(new FileInputStream(args[0]));

Scanner scanner = new Scanner(scannerReader);

Parser parser = new Parser(scanner.toLexer()); // bison

parser.parse();

}

catch (Exception exception) {

exception.printStackTrace();

}

}

static {

Main.VariableMap.put("pi", Math.PI);

Main.VariableMap.put("e", Math.E);

}

}

Finally, below follows a small script file that generates the Java source code for the parser and scanner.

Parser.bat

@c:

@cd "C:\Users\Stefan\Documents\Calculator\_Bison\src\calculator"

@"C:\Program Files (x86)\GnuWin32\bin\bison" --language=Java -o Parser.java MainParser.gppg

@"C:\Program Files\Java\jre1.8.0\_73\bin\java" -jar C:\Users\Stefan\Documents\YaccForJava\JFlex\jflex-1.6.1\lib\jflex-1.6.1.jar Scanner.jflex

# Main

The Main class handles the overall compiling and linking.

# The Preprocessor

Basically, the preprocessor is made up of two parts, where the first part is a scanner and a parser and the second part takes care of comments, string, and characters as well as handling macros and conditional programming.

## The Expression Parser

### The Grammar

Before we consider the preprocessor itself, we need to find a way to decide the whether the expressions given in the #if or #elif directives equals zero. To do so, we need a parser and a scanner, and we use JFlex and Bison from Chapter 2. The grammar of an expression follows below.

expression ::=

logical\_or\_expression

| logical\_or\_expression **?** expression **:** expression

logical\_or\_expression ::=

logical\_and\_expression

| logical\_or\_expression **||** logical\_and\_expression

logical\_and\_expression ::=

BITWISE\_OR\_expression

| bitwise\_and\_expression **&&** BITWISE\_OR\_expression

BITWISE\_OR\_expression ::=

bitwise\_xor\_expression

| BITWISE\_OR\_expression **|** bitwise\_xor\_expression

bitwise\_xor\_expression ::=

bitwise\_and\_expression

| bitwise\_xor\_expression **^** bitwise\_and\_expression

bitwise\_and\_expression ::=

equality\_expression

| bitwise\_and\_expression **&** equality\_expression

equality\_expression ::=

relation\_expression

| equality\_expression **==** relation\_expression

| equality\_expression **!=** relation\_expression

relation\_expression ::=

shift\_expression

| relation\_expression **<** shift\_expression

| relation\_expression **<=** shift\_expression

| relation\_expression **>** shift\_expression

| relation\_expression **>=** shift\_expression

shift\_expression ::=

add\_expression

| shift\_expression **<<** add\_expression

| shift\_expression **>>** add\_expression

add\_expression ::=

multiply\_expression

| add\_expression **+** multiply\_expression

| add\_expression **-** multiply\_expression

multiply\_expression ::=

prefix\_expression

| multiply\_expression **\*** prefix\_expression

| multiply\_expression **/** prefix\_expression

| multiply\_expression **%** prefix\_expression

prefix\_expression ::=

primary\_expression

| **+** prefix\_expression

| **-** prefix\_expression

| **~** prefix\_expression

| **!** prefix\_expression

primary\_expression ::=

**identifier**

| **value**

| **defined** **identifier**

| **defined** **( identifier )**

| **(** expression **)**

The parser is divided into two parts: one part that defines the rules and tokens, and one part that defines the grammar rules. As mentioned in Section **Fel! Hittar inte referenskälla.**, the parser defined in Bison decides whether the given code agrees with the grammar. Moreover, the parser can also evaluate a resulting value. In our case, we want to calculate the integer value of the expression to decide whether it is zero. In this parser, an identifier is a string, a value is an integer and the value of each non-terminal is also an integer (real values are not used by the preprocessor).

PreMainParser.gppg

%{

package c\_compiler;

import java.math.\*;

import java.util.\*;

%}

%name-prefix = "Pre"

%token QUESTION\_MARK COLON LEFT\_PAREN RIGHT\_PAREN

PLUS MINUS MULTIPLY DIVIDE MODULO DEFINED

EQUAL NOT\_EQUAL LESS\_THAN LESS\_THAN\_EQUAL

GREATER\_THAN GREATER\_THAN\_EQUAL LEFT\_SHIFT

RIGHT\_SHIFT LOGICAL\_OR LOGICAL\_AND LOGICAL\_NOT

BITWISE\_XOR BITWISE\_OR BITWISE\_AND BITWISE\_NOT

%token <String> IDENTIFIER

%token <Integer> VALUE

%type <Integer> expression logical\_or\_expression

logical\_and\_expression BITWISE\_OR\_expression

bitwise\_xor\_expression bitwise\_and\_expression

equality\_expression relation\_expression

shift\_expression add\_expression

multiply\_expression prefix\_expression

primary\_expression

%start expression

// --------------------------------------------------------------------------

%%

expression:

logical\_or\_expression {

$$ = $1;

PreProcessor.PreProcessorResult = $$;

}

| logical\_or\_expression QUESTION\_MARK expression

COLON expression {

$$ = ($1 != 0) ? $3 : $5;

PreProcessor.PreProcessorResult = $$;

};

logical\_or\_expression:

logical\_and\_expression {

$$ = $1;

}

| logical\_or\_expression LOGICAL\_OR logical\_and\_expression {

$$ = (($1 != 0) || ($3 != 0)) ? 1 : 0;

};

logical\_and\_expression:

BITWISE\_OR\_expression {

$$ = $1;

}

| bitwise\_and\_expression LOGICAL\_AND BITWISE\_OR\_expression {

$$ = (($1 != 0) && ($3 != 0)) ? 1 : 0;

};

BITWISE\_OR\_expression:

bitwise\_xor\_expression {

$$ = $1;

}

| BITWISE\_OR\_expression BITWISE\_OR bitwise\_xor\_expression {

$$ = $1 | $3;

};

bitwise\_xor\_expression:

bitwise\_and\_expression {

$$ = $1;

}

| bitwise\_xor\_expression BITWISE\_XOR bitwise\_and\_expression {

$$ = $1 ^ $3;

};

bitwise\_and\_expression:

equality\_expression {

$$ = $1;

}

| bitwise\_and\_expression BITWISE\_AND equality\_expression {

$$ = $1 & $3;

};

equality\_expression:

relation\_expression {

$$ = $1;

}

| equality\_expression EQUAL relation\_expression {

$$ = ($1 == $3) ? 1 : 0;

}

| equality\_expression NOT\_EQUAL relation\_expression {

$$ = ($1 != $3) ? 1 : 0;

};

relation\_expression:

shift\_expression { $$ = $1; }

| relation\_expression LESS\_THAN shift\_expression {

$$ = ($1 < $3) ? 1 : 0;

}

| relation\_expression LESS\_THAN\_EQUAL shift\_expression {

$$ = ($1 <= $3) ? 1 : 0;

}

| relation\_expression GREATER\_THAN shift\_expression {

$$ = ($1 > $3) ? 1 : 0;

}

| relation\_expression GREATER\_THAN\_EQUAL shift\_expression {

$$ = ($1 >=$3) ? 1 : 0;

};

shift\_expression:

add\_expression {

$$ = $1;

}

| shift\_expression LEFT\_SHIFT add\_expression {

$$ = $1 << $3;

}

| shift\_expression RIGHT\_SHIFT add\_expression {

$$ = $1 >> $3;

};

add\_expression:

multiply\_expression {

$$ = $1;

}

| add\_expression PLUS multiply\_expression {

$$ = $1 + $3;

}

| add\_expression MINUS multiply\_expression {

$$ = $1 - $3;

};

multiply\_expression:

prefix\_expression {

$$ = $1;

}

| multiply\_expression MULTIPLY prefix\_expression {

$$ = $1 \* $3;

}

| multiply\_expression DIVIDE prefix\_expression {

$$ = $1 / $3;

}

| multiply\_expression MODULO prefix\_expression {

$$ = $1 % $3;

};

prefix\_expression:

primary\_expression {

$$ = $1;

}

| PLUS prefix\_expression {

$$ = $2;

}

| MINUS prefix\_expression {

$$ = -$2;

}

| BITWISE\_NOT prefix\_expression {

$$ = ~$2;

}

An expression is cosidered to be true if it does not equal zero.

| LOGICAL\_NOT prefix\_expression {

$$ = ($2 == 0) ? 1 : 0;

};

In the preprocessor directive, every identifier is treated as the value zero.

primary\_expression:

IDENTIFIER {

$$ = 0;

}

| VALUE {

$$ = $1;

}

The defined directive works both with or without parenthesis.

|  |  |
| --- | --- |
| #if defined (x) | #if defined x |
| With parenthesis | Without parenthesis |

| DEFINED IDENTIFIER {

$$ = Main.MacroMap.containsKey($2) ? 1 : 0;

}

| DEFINED LEFT\_PAREN IDENTIFIER RIGHT\_PAREN {

$$ = Main.MacroMap.containsKey($3) ? 1 : 0;

}

| LEFT\_PAREN expression RIGHT\_PAREN {

$$ = $2;

};

### The Scanner

As mention in Section 2.1.4, tokens are the smallest parts of the language. However, the parser cannot know that, for instance, a sequence of letters is an identifier, or that a sequence of digits is an integer value. Instead, that is the task of a scanner.

PreScanner.jflex

package c\_compiler;

import java.math.\*;

%%

%public

%class PreScanner

%type Token

%notunix

//%unicode

%line

%char

%eofval{

return null;

%eofval}

%yylexthrow{

java.lang.Exception

%yylexthrow}

%eofval{

return null;

%eofval}

%yylexthrow{

java.lang.Exception

%yylexthrow}

%{

public PreParser.Lexer toLexer() {

return (new PreScannerToPreLexer(this));

}

%}

OCTAL\_VALUE = [0][0-7]\*([uU]|[sSlL]|[uU][sSlL]|[sSlL][uU])?

DECIMAL\_VALUE = [1-9][0-9]\*([uU]|[sSlL]|[uU][sSlL]|[sSlL][uU])?

HEXADECIMAL\_VALUE = [0][xX][0-9a-fA-F]+([uU]|[sSlL]|[uU][sSlL]|[sSlL][uU])?

IDENTIFIER = [a-zA-Z\_][a-zA-Z0-9\_]\*

WHITE\_SPACE = [ \t\r\n\f\0]

%%

"defined" { return (new Token(PreParser.DEFINED)); }

"?" { return (new Token(PreParser.QUESTION\_MARK)); }

":" { return (new Token(PreParser.COLON)); }

"||" { return (new Token(PreParser.LOGICAL\_OR)); }

"&&" { return (new Token(PreParser.LOGICAL\_AND)); }

"!" { return (new Token(PreParser.LOGICAL\_NOT)); }

"&" { return (new Token(PreParser.BITWISE\_AND)); }

"^" { return (new Token(PreParser.BITWISE\_XOR)); }

"|" { return (new Token(PreParser.BITWISE\_OR)); }

"~" { return (new Token(PreParser.BITWISE\_NOT)); }

"==" { return (new Token(PreParser.EQUAL)); }

"!=" { return (new Token(PreParser.NOT\_EQUAL)); }

"<" { return (new Token(PreParser.LESS\_THAN)); }

"<=" { return (new Token(PreParser.LESS\_THAN\_EQUAL)); }

">" { return (new Token(PreParser.GREATER\_THAN)); }

">=" { return (new Token(PreParser.GREATER\_THAN\_EQUAL)); }

"<<" { return (new Token(PreParser.LEFT\_SHIFT)); }

">>" { return (new Token(PreParser.RIGHT\_SHIFT)); }

"+" { return (new Token(PreParser.PLUS)); }

"-" { return (new Token(PreParser.MINUS)); }

"\*" { return (new Token(PreParser.MULTIPLY)); }

"/" { return (new Token(PreParser.DIVIDE)); }

"%" { return (new Token(PreParser.MODULO)); }

"(" { return (new Token(PreParser.LEFT\_PAREN)); }

")" { return (new Token(PreParser.RIGHT\_PAREN)); }

In case of an octal, decimal, or hexadecimal value, we just remove the letters u (unsigned), s (short), l (long), and x (hexadecimal value marker) from the text and let the Integer class interpret the value.

{IDENTIFIER} {

return (new Token(PreParser.IDENTIFIER, yytext()));

}

{OCTAL\_VALUE} {

String text = yytext().toLowerCase().replace("x", "").replace("u", "").

replace("s", "").replace("l", "");

int value = Integer.valueOf(text, 8);

return (new Token(PreParser.VALUE, value));

}

{DECIMAL\_VALUE} {

String text = yytext().toLowerCase().replace("x", "").replace("u", "").

replace("s", "").replace("l", "");

int value = Integer.valueOf(text, 10);

return (new Token(PreParser.VALUE, value));

}

{HEXADECIMAL\_VALUE} {

String text = yytext().toLowerCase().replace("x", "").replace("u", "").

replace("s", "").replace("l", "");

int value = Integer.valueOf(text, 16);

return (new Token(PreParser.VALUE, value));

}

It is important to detect white spaces, otherwise the user would not be allowed to add any spaces to the code.

{WHITE\_SPACE} {

// Empty.

}

If we detect a character we do not recognize, we just stop the execution with an error message.

. { Assert.error(yytext(), "unknown character"); }

## The Preprocessor

The preprocessor has several tasks:

* **Tri Graphs**. When C was originally introduced, some keyboards had a limited set of keys. Therefore, a special set of double question mark character sequence was introduced, which are replaced by their modern equivalents.
* **Comments**. Each block comment or line comment is replaced by blank characters.
* **Slashes in strings and characters**. Each backslash is processed and transformed into a regular character. Thereafter, to ease the macro expansion later we transform each character into the format of a backslash followed by three octal digits.
* **Include files**. The system include files (encapsulated by ‘<’ and ‘>’) and internal include files (encapsulated by quotes) are read and included in the final code.
* **Conditional programming**. The #if, #ifdef, #ifndef, #elif, and #endif directives are processed.
* **Macro expansion**. The source code is traversed and each macro is expanded. Since the contents of each character and string constant has been translated into slash codes we do not need to check whether the macro to be expanded is part of a character of string.
* **String sequences**. Finally, every sequence of string is concatenated into one string. For instance, the sequence "ab" "cd" "ef" is concatenated into "abcded".

The first phase of the preprocessor is to read the source file and place in a text buffer (a StringBuilder object), and then replace the tri graph sequences, remove the comments and replace every character in string and characters with its equivalent backslash code. Then the buffer is divided into lines and each line starting with a preprocessor directive is evaluated and every line is search for macros to expand.

Preprocessor.java

package c\_compiler;

import java.io.\*;

import java.util.\*;

public class PreProcessor {

public static enum IfStatus {If, Else};

public static void doProcess(File file, StringBuilder buffer)

throws Exception {

readFile(file, buffer);

triGraphs(buffer);

traverse(buffer);

Main.Path = file.getPath();

Main.Line = 0;

buffer.insert(0, Main.SeparatorId + Main.Path + ",0$\n");

int stackSize = Main.IfStack.size();

List<String> lineList = generateLineList(buffer);

lineList = modifyLineList(lineList);

traverseLineList(lineList);

Assert.error(Main.IfStack.size() == stackSize,

"unbalanced #if and #endif strucure");

buffer.replace(0, buffer.length(), "");

for (String line : lineList) {

buffer.append(line + "\n");

}

mergeStrings(buffer);

replaceAll(buffer, "\0", "");

}

private static void readFile(File file, StringBuilder buffer)

throws IOException {

FileInputStream fileStream = new FileInputStream(file);

InputStreamReader streamReader = new InputStreamReader(fileStream);

BufferedReader reader = new BufferedReader(streamReader);

String line;

while ((line = reader.readLine()) != null) {

buffer.append(line.trim() + "\n");

}

reader.close();

}

### Tri Graphs

The following table shows the tri graphs character sequences and their modern equivalences:

|  |  |
| --- | --- |
| **Tri-Graph Character Sequence** | **Modern Equivalence** |
| ??= | # |
| ??/ | \ |
| ??´ | ^ |
| ??( | [ |
| ??) | ] |
| ??! | | |
| ??< | { |
| ??> | } |
| ??- | ~ |

The tri graphs sequences are easy to replace, since we do not need to take any consideration in whether the tri graphs are placed inside strings, characters, or comments.

private static void triGraphs(StringBuilder buffer) {

replaceAll(buffer, "?=", "#");

replaceAll(buffer, "?/", "\\");

replaceAll(buffer, "?'", "^");

replaceAll(buffer, "?(", "[");

replaceAll(buffer, "?)", "]");

replaceAll(buffer, "?!", "|");

replaceAll(buffer, "?<", "{");

replaceAll(buffer, "?>", "}");

replaceAll(buffer, "?-", "~");

}

private static void replaceAll(StringBuilder buffer,

String oldText, String newText) {

int oldSize = oldText.length(), index;

while ((index = buffer.indexOf(oldText)) != -1) {

buffer.replace(index, index + oldSize, newText);

}

}

### Comments, Strings, and Characters

The next step is to take care of block comments (/\* to \*/) and line comments (// to the end of the line), strings, and characters. Since they may be nested, they need to be evaluated in the same phase.

The idea is that we go through the buffer, and when we find the beginning of a line comment, a block comment, a string, or a character we call the corresponding method, which modify the buffer and returns the index of the next character to be inspected.

To make the inspection of the source code buffer easier we start by adding three zero characters, which are removed at the end of the method.

public static void traverse(StringBuilder buffer) {

buffer.append("\0\0\0");

for (int index = 0; buffer.charAt(index) != '\0'; ++index) {

if (buffer.substring(index, index + 2).equals("//")) {

index = traverseLineComment(index, buffer);

}

else if (buffer.substring(index, index + 2).equals("/\*")) {

index = traverseBlockComment(index, buffer);

}

else if (buffer.charAt(index) == '\"') {

index = traverseString(index + 1, buffer, true);

}

else if (buffer.charAt(index) == '\'') {

index = traverseCharacter(index + 1, buffer, true);

}

else {

Assert.error(buffer.charAt(index) != '$',

Main.SeparatorId, "invalid character");

}

}

replaceAll(buffer, "\0", "");

}

The traverseLineComment method is quite easy, we just continue until we find the end of the line (‘\n’) or the end of the buffer (‘\0’), and replace every preceding character with a blank.

private static int traverseLineComment(int index, StringBuilder buffer) {

while ((buffer.charAt(index) != '\n') &&

(buffer.charAt(index) != '\0')) {

buffer.setCharAt(index++, ' ');

}

return index;

}

The traverseBlockComment method is a little bit more complicated. We continue until we find the end of the comment (\*/). If we instead find the end of the buffer (‘\0’) the comment is unterminated, which result in an error message. Each character is replaced by a space character. However, newlines (‘\n’) are not replaced, in order for the line count to work properly.

private static int traverseBlockComment(int index, StringBuilder buffer) {

while (true) {

if (buffer.charAt(index) == '\0') {

Assert.error("Unfinished commant block.");

}

else if (buffer.substring(index, index + 2).equals("\*/")) {

buffer.setCharAt(index, ' ');

buffer.setCharAt(index + 1, ' ');

return (index + 1);

}

else if (buffer.charAt(index) == '\n') {

buffer.setCharAt(index++, '\n');

}

else {

buffer.setCharAt(index++, ' ');

}

}

}

The traverseString method needs to keep track of the end of the string (‘\”’), the end of the buffer (‘\0’), or newline in the string (‘\n’). The end of the buffer or newline result in error messages. Moreover, when we find a backslash (‘\\’) we call doSlash, that transforms the slash code into a regular character. Finally, when the string has been traversed each regular character is modified into a backslash followed by three octal digits, in order to ease the macro replacement process in a later phase.

public static int traverseString(int index, StringBuilder buffer, boolean octal) {

while (true) {

switch (buffer.charAt(index)) {

case '\"':

return index;

case '\n':

Assert.error("new line in string.");

break;

case '\0':

Assert.error("unfinished string.");

break;

case '\\':

slashToChar(index, buffer);

break;

}

if (octal) {

charToOctal(index, buffer);

index += 4;

}

else {

++index;

}

}

}

The traverseCharacter method is similar to traverseString above, it generates error messages for end-of-line or newline, calls doSlash for each backslash, and finally changes each regular character into a backslash followed by three octal digits.

public static int traverseCharacter(int index, StringBuilder buffer,

boolean octal) {

int charCount = 0;

while (true) {

switch (buffer.charAt(index)) {

case '\'':

Assert.error(charCount == 1, "invalid character sequence");

return index;

case '\n':

Assert.error("new line in character.");

break;

case '\0':

Assert.error("unfinished character.");

break;

case '\\':

slashToChar(index, buffer);

break;

}

if (octal) {

charToOctal(index, buffer);

index += 4;

}

else {

++index;

}

++charCount;

}

}

### Slash Codes

The slashToChar method inspect the character succeeding the slash.

If the characters following the slash are a small or capital x and one or two hexadecimal characters, or three octal the following two characters are inspected, and their value calculated and the two hexadecimal digits are replaced three octal digits. However, if the two following characters are not hexadecimal digits, an error message occurs.

public static void slashToChar(int index, StringBuilder buffer) {

char char1 = buffer.charAt(index + 1),

char2 = buffer.charAt(index + 2),

char3 = buffer.charAt(index + 3);

If the three, two, or one characters following the slash are octal digits, the value of the digits is calculated and the slash sequence is replaced by the character with the ASCII value of the octal digits. If the value of the digits exceeds 255, an error is reported.

if (isOctal(char1) && isOctal(char2) && isOctal(char3)) {

int octValue = 64 \* charToOctal(char1) +

8 \* charToOctal(char2) +

charToOctal(char3);

Assert.error(octValue < 256, "invalid octal sequence");

buffer.replace(index, index + 4, String.valueOf((char) octValue));

}

else if (isOctal(char1) && isOctal(char2)) {

int octValue = 8 \* charToOctal(char1) + charToOctal(char2);

buffer.replace(index, index + 3, String.valueOf((char) octValue));

}

else if (isOctal(char1)) {

int octValue = charToOctal(char1);

buffer.replace(index, index + 2, String.valueOf((char) octValue));

}

If the character following the slash is a small or capital x and two or one hexadecimal digits, the value of the digits is calculated, and the slash sequence is replaced by the character with the ASCII value of the hexadecimal digits. A small or capital x not followed by at least one hexadecimal digit results in an error being reported.

else if ((char1 == 'x') || (char1 == 'X')) {

if (isHex(char1) && isHex(char2)) {

int hexValue = 16 \* charToHex(char1) + charToHex(char2);

buffer.replace(index, index + 3, String.valueOf((char) hexValue));

}

else if (isHex(char1)) {

int hexValue = charToHex(char1);

buffer.replace(index, index + 2, String.valueOf((char) hexValue));

}

else {

Assert.error("invalid hexadecimal sequence");

}

}

If none of the cases above applies, an error is reported.

else {

Assert.error(buffer.charAt(index + 1), "invalid slash sequence");

}

}

private static final Map<Character,Character> m\_escapeMap = new ListMap<>();

static {

m\_escapeMap.put('0','\0');

m\_escapeMap.put('n','\n');

m\_escapeMap.put('t','\t');

m\_escapeMap.put('v', (char) 11);

m\_escapeMap.put('b','\b');

m\_escapeMap.put('r','\r');

m\_escapeMap.put('f','\f');

m\_escapeMap.put('a',(char) 7);

m\_escapeMap.put('\\','\\');

m\_escapeMap.put('?','?');

m\_escapeMap.put('\'','\'');

m\_escapeMap.put('"','\"');

}

private static boolean isOctal(char c) {

return "01234567".contains(String.valueOf(c));

}

private static boolean isHex(char c) {

return "0123456789abcdefABCDEF".contains(String.valueOf(c));

}

private static int charToOctal(char c) {

return "01234567".indexOf(c);

}

private static int charToHex(char c) {

return "0123456789abcdef".indexOf(Character.toLowerCase(c));

}

When the character or string has been translated from slash codes to regular characters, charToOctal is called to translate them to octal slash codes.

public static void charToOctal(int index, StringBuilder buffer) {

int asciiValue = buffer.charAt(index);

char firstChar = "01234567".charAt(asciiValue / 64),

secondChar = "01234567".charAt((asciiValue % 64) / 8),

thirdChar = "01234567".charAt(asciiValue % 8);

String text = "\\" + String.valueOf(firstChar) +

String.valueOf(secondChar) +

String.valueOf(thirdChar) ;

buffer.replace(index, index + 1, text);

}

The octalToChar method is called by the scanner to tralform the octal shalch codes into regular characters.

public static String octalToChar(String text) {

StringBuilder buffer = new StringBuilder(text);

for (int index = (text.length() - 1); index >= 0; --index) {

if (buffer.charAt(index) == '\\') {

PreProcessor.slashToChar(index, buffer);

}

}

return buffer.toString();

}

### The Line List

When the buffer has been modified in accordance with the methods above, it is time to transform it into a list of lines to process the preprocessors directives. The transformation is simple, we just use the split method. We add an extra blank line in the last line is a macro definition that ends with two slashes (“\\”).

private static List<String> generateLineList(StringBuilder buffer) {

String lineArray[] = buffer.toString().split("\n");

List<String> lineList = new LinkedList<>();

for (String line : lineArray) {

lineList.add(line.trim());

}

lineList.add("");

return lineList;

}

However, we also need to merge lines that ends with two backslashes together:

|  |  |
| --- | --- |
| #define min(x,y) \\  (((x) < (y)) ? \\  (x) : (y)) | #define min(x,y) (((x) < (y)) ? (x) : (y)) |
| (a) Before | (b) After |

If we encounter lines ending with two backslashes, we continue to append the lines to lineBuffer. Since an empty line was added at the end of the line list in generateLineList above, there is no risk that that the last line ends with a backslash.

private static List<String> modifyLineList(List<String> lineList) {

List<String> resultList = new LinkedList<>();

for (int index = 0; index < lineList.size(); ++index) {

String line = lineList.get(index);

StringBuilder lineBuffer = new StringBuilder();

while (line.endsWith("\\\\")) {

lineBuffer.append(line.substring(0, line.length() - 2) + " ");

line = lineList.get(++index);

}

resultList.add(lineBuffer.toString() + line + "\0");

}

return resultList;

}

When traversing the lines, we check whether the line starts with a sharp (‘#’). If it does, we consider the word following the sharp sign (#). If it is a preprocessor directive, we call the corresponding method. If the line does not start with a sharp and we are in a visible part of source code, we call searchForMacros, which expands macros. If we are not in a visible part of the source code, the line is replaced by an empty line.

public static void traverseLineList(List<String> lineList)throws Exception{

for (int index = 0; index < lineList.size(); ++index) {

String line = lineList.get(index);

if (line.startsWith("$")) {

// Empty.

}

else if (line.startsWith("#")) {

line = trimLeft(line.substring(1));

if (line.startsWith("line")) {

String rest = line.substring("line".length());

lineList.set(index, doLine(trimLeft(rest)));

}

else if (isVisible() && line.startsWith("include")) {

String rest = line.substring("include".length());

String result = doInclude(rest.trim());

lineList.set(index, result);

}

else if (isVisible() && line.startsWith("error")) {

String message = line.substring("error".length()).trim();

System.err.println(message);

System.exit(-1);

}

else if (isVisible() && line.startsWith("define")) {

String rest = line.substring("define".length());

doDefine(trimLeft(rest));

lineList.set(index, "");

}

else if (isVisible() && line.startsWith("undef")) {

String rest = line.substring("undef".length());

doUndef(trimLeft(rest));

lineList.set(index, "");

}

else if (line.startsWith("ifdef")) {

String rest = line.substring("ifdef".length());

doIfDefined(trimLeft(rest));

lineList.set(index, "");

}

else if (line.startsWith("ifndef")) {

String rest = line.substring("ifndef".length());

doIfNotDefined(trimLeft(rest));

lineList.set(index, "");

}

else if (line.startsWith("if")) {

String rest = line.substring("if".length());

doIf(trimLeft(rest));

lineList.set(index, "");

}

else if (line.startsWith("elif")) {

String rest = line.substring("elif".length());

doElseIf(trimLeft(rest));

lineList.set(index, "");

}

else if (line.startsWith("else")) {

String rest = line.substring("else".length());

doElse(trimLeft(rest));

lineList.set(index, "");

}

else if (line.startsWith("endif")) {

String rest = line.substring("endif".length());

doEndIf(trimLeft(rest));

lineList.set(index, "");

}

else {

lineList.set(index, "");

}

}

else if (isVisible()) {

String result = searchForMacros(line, new Stack<String>());

lineList.set(index, result);

}

else {

lineList.set(index, "");

}

Main.Line += countChar(line, '\n') + 1;

}

}

The isVisible method checks whether we are in a visible part of the source code; that is, a part that has not been earlier omitted by a #if, #ifdef, #ifndef, #elif, or #else preprocessor directive. We use the Main.IfStack stack since conditional programming can be nested. If the source code is not visible at any nesting level (that is, anywhere in the stack), the area is not visible and we return false.

private static boolean isVisible() {

for (Triple<Boolean,Boolean,IfStatus> ifTriple : Main.IfStack) {

boolean currTrue = ifTriple.getFirst();

if (!currTrue) {

return false;

}

}

return true;

}

The trimLeft and trimRight methods simple removes all white-spaces to the left or right of the string.

private static String trimLeft(String text) {

int index = 0;

while ((index < text.length()) &&

Character.isWhitespace(text.charAt(index))) {

++index;

}

return text.substring(index);

}

private static String trimRight(String text) {

int index = text.length() - 1;

while ((index >= 0) && Character.isWhitespace(text.charAt(index))) {

--index;

}

return text.substring(0, index + 1);

}

The countChar method simple counts the number of occurrences of the character.

private static int countChar(String text, char c) {

int count = 0;

for (char c2 : text.toCharArray()) {

if (c == c2) {

++count;

}

}

return count;

}

### Lines

The doLine method handles the #line directive by setting the Main.Path and Main.Line fields and returns a text with the path and line that starts and ends with dollar signs (‘$’).

|  |  |  |
| --- | --- | --- |
| #line 100 C:\Temp\Test.c |  |  |
| $C:\Temp\Test.c,100$ |  |  |

private static String doLine(String line) {

String lineText, pathText;

int spaceIndex = line.indexOf(" ");

if (spaceIndex != -1) {

lineText = line.substring(0, spaceIndex);

pathText = line.substring(spaceIndex + 1).trim();

}

else {

lineText = line;

pathText = null;

}

try {

Main.Line = Integer.parseInt(lineText);

}

catch (NumberFormatException exception) {

Assert.error(lineText, "invalid line number");

}

if (pathText != null) {

Main.Path = pathText;

}

return Main.SeparatorId + Main.Path + "," + Main.Line + Main.SeparatorId;

}

### Includes

There are two kinds of include files: system files (‘<’ and ‘>’) and internal files (‘\”’). The systems files are included from the path given by Main.IncludePath, and the internal files are included locally.

The Main.IncludeSet set is used to track the included file so that the source file include the files shall be recompiled if it or any of its included files has been changed after the object file. The Main.IncludeStack is used to make sure that there is no circular inclusion.

The Main.Path and Main.Line fields are temporary replaced by the path and line of the included file. After the include line has been read, the original path and line is added to the included source code, in the dollar sign format (for example: “$C:\Temp\Test.c,100$”). The method returns the source code of the included file, which will replace the original include line.

private static String doInclude(String line) throws Exception {

File includeFile = null;

if (line.startsWith("<") && line.endsWith(">")) {

includeFile = new File(Main.IncludePath +

line.substring(1, line.length() - 1));

}

else if (line.startsWith("\"") && line.endsWith("\"")) {

String name = line.substring(1, line.length() - 1);

includeFile = new File(octalToChar(name));

}

else {

Assert.error(line, "invalid include format");

}

Assert.error(!Main.IncludeStack.contains(includeFile),

"nestled include directives: " + includeFile.getPath());

Main.IncludeStack.push(includeFile);

Main.IncludeSet.add(includeFile);

String oldPath = Main.Path;

int oldLine = Main.Line;

StringBuilder buffer = new StringBuilder();

doProcess(includeFile, buffer);

Main.Line = oldLine;

Main.Path = oldPath;

buffer.append(Main.SeparatorId + Main.Path + "," +

Main.Line + Main.SeparatorId);

Main.IncludeStack.pop();

return buffer.toString();

}

### Macros

The Macro class keep track of a macro, a macro has a possible empty list of parameters and a text.

Macro.java

package c\_compiler;

public class Macro {

private int m\_parameters;

private String m\_body;

public Macro(int parameters, String body) {

m\_parameters = parameters;

m\_body = body;

}

public Macro(String body) {

m\_parameters = 0;

m\_body = body;

}

public int parameters() {

return m\_parameters;

}

public String body() {

return m\_body;

}

public boolean equals(Object object) {

if (object instanceof Macro) {

Macro macro = (Macro) object;

return m\_body.equals(macro.m\_body) &&

(m\_parameters == macro.m\_parameters);

}

return false;

}

}

The doDefine method splits the line into the name of the macro, a potential parameter list, and a body. When the first character that is not a letter, digit, or an underline is reached, we check whether the line so far is a proper identifier. Then we have three cases:

1. We have reached the end of the line, in which case we have a macro with a name, but without parameters or body.

2. We have reached a space, in which case we have a macro with a name and a body, but no parameters.

3. We have reached a left parenthesis, in which case we have a list of parameters and we call doParameterDefine.

4. If none of the above applies, we generate an error message.

Preprocessor.java

public static void doDefine(String line) {

int index;

for (index = 0; Character.isLetterOrDigit(line.charAt(index)) ||

(line.charAt(index) == '\_'); ++index) {

// Empty.

}

String name = line.substring(0, index);

checkIdentifier(name);

switch (line.charAt(index)) {

case '\0': {

Macro oldMacro = Main.MacroMap.get(name),

newMacro = new Macro("");

Assert.error((oldMacro == null) || oldMacro.equals(newMacro),

name, "macro already defined");

Main.MacroMap.put(name, newMacro);

}

break;

case ' ': {

String text = line.substring(index + 1).trim();

Macro oldMacro = Main.MacroMap.get(name),

newMacro = new Macro(text);

Assert.error((oldMacro == null) || oldMacro.equals(newMacro),

name, "macro already defined");

Main.MacroMap.put(name, newMacro);

}

break;

case '(': {

String text = line.substring(index);

doParameterDefine(name, text);

}

break;

default: {

Assert.error(line, "invalid macro definition");

}

break;

}

}

The doParameterDefine method look into macros with parameters in two steps. First we extract the parameters by calling scanParameters. Then we go through the parameter list and check that the parameters are identifiers and that no parameters is repeated.

private static void doParameterDefine(String name, String text) {

List<String> paramList = new LinkedList<>();

StringBuilder buffer = new StringBuilder(text);

int lastIndex = scanParameters(buffer, 0, paramList);

buffer.delete(0, lastIndex + 1);

Set<String> paramSet = new ListSet<>();

for (int index = 0; index < paramList.size(); ++index) {

String param = paramList.get(index).trim();

checkIdentifier(param);

Assert.error(paramSet.add(param), param, "repeated macro parameter");

paramList.set(index, param);

}

Then we extract the identifiers from the body and replace each occurrence of a parameters with the text “$parameter\_index$”.

List<Pair<String,Integer>> identList =

getIdentifierList(buffer.toString());

for (int index = (identList.size() - 1); index >= 0; --index) {

Pair<String,Integer> identPair = identList.get(index);

String ident = identPair.getFirst();

int identIndex = identPair.getSecond();

int paramIndex = paramList.indexOf(ident);

if (paramIndex != -1) {

buffer.replace(identIndex, identIndex + ident.length(),

Main.SeparatorId + paramIndex + Main.SeparatorId);

}

}

Finally, we look for the merge operator ##. When we find one, we remove the operator and trim the text to its left and right.

int index;

String body = buffer.toString();

while ((index = body.indexOf("##")) != -1) {

String left = body.substring(0, index),

right = body.substring(index + 2);

body = left.trim() + right.trim();

}

It is allowed to redefine a macro if it is has the same parameters list and macro.

Macro oldMacro = Main.MacroMap.get(name),

newMacro = new Macro(paramList.size(), body);

Assert.error((oldMacro == null) || oldMacro.equals(newMacro),

name, "macro already defined");

Main.MacroMap.put(name, newMacro);

}

An identifier is a text starting with a letter or an underline and continuing with letters, digits, or underlines. If the given name is not an identifier, an error occurs.

private static void checkIdentifier(String name) {

Assert.error(!name.isEmpty() && (Character.isLetter(name.charAt(0)) ||

(name.charAt(0) == '\_')),

name, "invalid identifier");

for (int index = 1; index < name.length(); ++index) {

Assert.error(Character.isLetterOrDigit(name.charAt(index)) ||

(name.charAt(index) == '\_'),

name, "invalid identifier");

}

}

The getIdentifierList method extracts identifiers from a text and returns a list holding the identifiers and their indexes in the text.

private static List<Pair<String,Integer>> getIdentifierList(String text) {

StringBuilder buffer = new StringBuilder(text);

List<Pair<String,Integer>> identList = new LinkedList<>();

if (!text.isEmpty()) {

int index = 0;

while (buffer.charAt(index) != '\0') {

char c = buffer.charAt(index);

if (Character.isLetter(buffer.charAt(index)) ||

(buffer.charAt(index) == '\_')) {

int nextIndex = index;

while (Character.isLetterOrDigit(buffer.charAt(nextIndex)) ||

(buffer.charAt(nextIndex) == '\_')) {

++nextIndex;

}

String ident = buffer.substring(index, nextIndex);

identList.add(new Pair<>(ident, index));

index = nextIndex;

}

else {

++index;

}

}

}

return identList;

}

The doUndef method removes a macro from the MacroMap map. If the macro does not exists a warning is given.

public static void doUndef(String line) {

String name = line;

checkIdentifier(name);

Assert.warning(Main.MacroMap.remove(name) != null,

name, "macro not defined");

}

### Conditional Programming

With conditional programming, it is possible to exclude part of the source code. The doIf method uses the parser of Section 4.1. It calls parseExpression, which creates a scanner and parser, parse the line, and returns a Boolean value that is true in case of a non-zero value.

private static void doIf(String line) {

boolean result = parseExpression(line);

Main.IfStack.push(new Triple<>(result, result, IfStatus.If));

}

public static Object PreProcessorResult;

private static boolean parseExpression(String line) {

StringReader reader = new StringReader(line);

int result = 0;

try {

PreScanner preScanner = new PreScanner(reader);

PreParser preParser = new PreParser(preScanner.toLexer());

Assert.error(preParser.parse(), "pre parser");

result = (int) PreProcessorResult;

}

catch (Exception exception) {

Assert.error(line, "invalid expression");

}

reader.close();

return (result != 0);

}

The doIfDefined and doIfNotDefined checks if the macro is defined or not, respectively, and add the result to the Main.IfStack stack.

private static void doIfDefined(String line) {

String name = line.trim();

checkIdentifier(name);

boolean result = Main.MacroMap.containsKey(name);

Main.IfStack.push(new Triple<>(result, result, IfStatus.If));

}

private static void doIfNotDefined(String line) {

String name = line.trim();

checkIdentifier(name);

boolean result = !Main.MacroMap.containsKey(name);

Main.IfStack.push(new Triple<>(result, result, IfStatus.If));

}

The doIfElseIf method is a little bit more complicated. It checks whether there is a preceding #if directive (Main.IfStack is not empty). Thereafter, it looks up the status of the if-stack. The stack value is a triple, the first value indicates whether the current if-status is true, and the second value gives whether there has been an earlier true if-status. The third value is of type IfStatus, that can hold the value If and Else. If the value is Else, and #elseif has already occurred, and another on is not allowed. Finally, if this if-chain has ever been true (lastTrue is true) is does not matter whether the condition of directive is true, the new status is false. If not, we consider if this value is true. In both cases, we push the result on the stack.

private static void doElseIf(String line) {

boolean result = parseExpression(line);

Assert.error(!Main.IfStack.isEmpty(),

"#elif without preceeding #if, #ifdef, or #ifndef");

Triple<Boolean,Boolean,IfStatus> ifTriple = Main.IfStack.pop();

boolean currTrue = ifTriple.getFirst(),

lastTrue = ifTriple.getSecond();

IfStatus lastIf = ifTriple.getThird();

Assert.error(lastIf == IfStatus.If,

"#elif after #else if not allowed");

if (!lastTrue) {

Main.IfStack.push(new Triple<>(result, result, IfStatus.If));

}

else {

Main.IfStack.push(new Triple<>(false, true, IfStatus.If));

}

}

The doElse method is a little bit easier. It checks whether there has ever been a true part of the if-chain. If not, it pushes true on the stack.

private static void doElse(String line) {

Assert.error(!Main.IfStack.isEmpty(),

"#else without preceeding #if, #ifdef, or #ifndef");

Triple<Boolean,Boolean,IfStatus> ifTriple = Main.IfStack.pop();

boolean lastTrue = ifTriple.getSecond();

IfStatus lastIf = ifTriple.getThird();

Assert.error(lastIf == IfStatus.If,

"#else after #else if not allowed");

Main.IfStack.push(new Triple<>(!lastTrue, true /\* doesn't matter \*/,

IfStatus.Else));

}

Finally, the doEndIf just pops the if-stack, subject to it is not already empty.

private static void doEndIf(String line) {

Assert.error(!Main.IfStack.isEmpty(),

"#endif without preceeding #if, #ifdef, or #ifndef");

Main.IfStack.pop();

}

### Macro Expansion

In the regular source code (line not starting with ‘#’), we need to replace macro calls with their resulting bodies. We also need to recursively replace macro calls in the bodies. However, unlike function, recursive macro calls are not allowed.

We go through the text, and look for identifiers, we do not have to worry about identifier in strings or characters, since every letter or underline (an identifier must begin with a letter or underline) in strings or characters has been changed to a three-octal-digit backslash code by slashToChar in Section 4.2.2.

When we find and identifier that is a stored macro name (either in m\_specialMacroSet or MacroMap), we scan the parameters if the name is followed by a left parenthesis, each parameter is recursively searched for macros. We then look up the macro body and replace its marked parameter locations with the parameters. Finally, we replace the macro call with its body. The nameStack stack is used to prevent recursive macro calls.

private static final Set<String> m\_specialMacroSet = new ListSet<>();

static {

m\_specialMacroSet.add("\_\_STDC\_\_");

m\_specialMacroSet.add("\_\_FILE\_\_");

m\_specialMacroSet.add("\_\_LINE\_\_");

m\_specialMacroSet.add("\_\_DATE\_\_");

m\_specialMacroSet.add("\_\_TIME\_\_");

}

private static String searchForMacros(String text,Stack<String> nameStack){

StringBuilder buffer = new StringBuilder(text + "\0");

for (int index = 0; buffer.charAt(index) != '\0'; ++index) {

if (Character.isLetter(buffer.charAt(index)) ||

(buffer.charAt(index) == '\_')) {

int nextIndex = index;

while (Character.isLetterOrDigit(buffer.charAt(nextIndex)) ||

(buffer.charAt(nextIndex) == '\_')) {

++nextIndex;

}

String name = buffer.substring(index, nextIndex);

if (!nameStack.contains(name) && (Main.MacroMap.containsKey(name) ||

m\_specialMacroSet.contains(name))){

nameStack.push(name);

List<String> paramList = new LinkedList<>();

if (buffer.charAt(nextIndex) == '(') {

nextIndex = scanParameters(buffer, nextIndex, paramList) + 1;

}

for (int paramIndex = 0; paramIndex < paramList.size();

++paramIndex) {

String oldParam = paramList.get(paramIndex) + "\0";

String newParam = searchForMacros(oldParam, nameStack);

paramList.set(paramIndex, newParam);

}

String macroText = lookupMacro(name, paramList);

macroText = searchForMacros(macroText, nameStack);

buffer.replace(index, nextIndex, macroText);

index += macroText.length();

nameStack.pop();

}

else {

index = nextIndex - 1;

}

}

}

return buffer.toString().replace("\0", "");

}

The scanParameters extracts the actual parameters of a macro call. The parameters are separated by commas and the list is terminated by a right parenthesis. It becomes a bit complicated since the parameters themselves can hold parenthesis and commas. The field paranCount counts the level of parentheses nesting and only takes parameter into consideration when the nesting level is zero.

private static int scanParameters(StringBuilder buffer, int index,

List<String> paramList) {

int parenCount = 0;

StringBuilder paramBuffer = new StringBuilder();

for (; true; ++index) {

char c = buffer.charAt(index);

switch (c) {

case '(':

if (++parenCount > 1) {

paramBuffer.append(c);

}

break;

case ')':

if (--parenCount == 0) {

paramList.add(paramBuffer.toString());

return index;

}

else {

paramBuffer.append(c);

}

break;

case ',':

if (parenCount == 1) {

paramList.add(paramBuffer.toString());

paramBuffer.delete(0, paramBuffer.length());

}

else {

paramBuffer.append(c);

}

break;

case '\0':

Assert.error("unfinished macro parameter sequence");

break;

default:

paramBuffer.append(c);

break;

}

}

}

The lookupMacro method looks up a given macro. If it stored in MacroMap, the number of parameters is checked and then the parameter locations in the macro body are replaced by the actual parameters.

private static String lookupMacro(String name, List<String> paramList) {

if (Main.MacroMap.containsKey(name)) {

Macro macro = Main.MacroMap.get(name);

Assert.error(paramList.size() == macro.parameters(),

"invalid number of parameters");

StringBuilder buffer = new StringBuilder(macro.body());

for (int index = 0; index < macro.parameters(); ++index) {

replaceAll(buffer, "#$" + index + Main.SeparatorId,

"\"" + paramList.get(index) + "\"");

replaceAll(buffer, Main.SeparatorId + index + Main.SeparatorId,

paramList.get(index));

}

return buffer.toString();

}

If the macro name is not stored in MacroMap, it may be one of the predefined special macros. The \_\_STDC\_\_ macro is replaced by the one integer value since the compiler of this book supports standard C.

else if (name.equals("\_\_STDC\_\_")) {

Assert.error(paramList.isEmpty(), "\_\_STDC\_\_",

"invalid macro parameter list");

return "1";

}

The \_\_FILE\_\_ macro is replaced by the current path (Main.Path), the \_\_LINE\_\_ macro is replaced by the current line number (Main.Line).

else if (name.equals("\_\_FILE\_\_")) {

Assert.error(paramList.isEmpty(), "\_\_FILE\_\_",

"invalid macro parameter list");

StringBuilder fileBuffer =

new StringBuilder("\"" + Main.Path.replace("\\", "\\\\") + "\"\0");

traverseString(1, fileBuffer, true);

return fileBuffer.toString().replace("\0", "");

}

else if (name.equals("\_\_LINE\_\_")) {

Assert.error(paramList.isEmpty(), "\_\_LINE\_\_",

"invalid macro parameter list");

return Integer.toString(Main.Line);

}

The \_\_DATE\_\_ macro is replaced by the current date on the format “Jan 1, 1970” and the \_\_TIME\_\_ macro is replaced by the current time on the format “01:02:03”.

else if (name.equals("\_\_DATE\_\_")) {

Assert.error(paramList.isEmpty(), "\_\_DATE\_\_",

"invalid macro parameter list");

Calendar calendar = new GregorianCalendar();

String monthArray[] = {"Jan", "Feb", "Mar", "Apr", "May", "Jun",

"Jul", "Aug", "Sep", "Oct", "Nov", "Dec"};

return (String.format("\"%s %02d %04d\"",

monthArray[calendar.get(Calendar.MONTH)],

calendar.get(Calendar.DAY\_OF\_MONTH),

calendar.get(Calendar.YEAR)));

}

else if (name.equals("\_\_TIME\_\_")) {

Assert.error(paramList.isEmpty(), "\_\_TIME\_\_",

"invalid macro parameter list");

Calendar calendar = new GregorianCalendar();

return (String.format("\"%02d:%02d:%02d\"",

calendar.get(Calendar.HOUR\_OF\_DAY),

calendar.get(Calendar.MINUTE),

calendar.get(Calendar.SECOND)));

}

If the macro is not stored in the macro map and is not a special macro, null is returned.

else {

return null;

}

}

### String Merging

When the all the macros have been expanded, it is finally time to check whether there are pairs of string constants that need to be merged:

|  |  |  |
| --- | --- | --- |
| "Hello" "World" | "HelloWorld" |  |
| (a) Before | (b) After |  |

private static void mergeStrings(StringBuilder buffer) {

int lastIndex = -1, index = buffer.length() - 1;

while (index >= 0) {

char c = buffer.charAt(index);

if (c == '\"') {

if (lastIndex != -1) {

buffer.replace(index, lastIndex + 1, "");

}

--index;

while ((index >= 0) && (buffer.charAt(index) != '\"')) {

--index;

}

lastIndex = index--;

}

else {

if (!Character.isWhitespace(c)) {

lastIndex = -1;

}

--index;

}

}

}

}

# Scanning

In this project we use Garden Point Lex, which is based on the classic scanner generator tool Lex. See appendix A for a crash course.

### The typedef-name Problem

Let us take look at the tow source code lines below. Intuitively, the first line looks like an expression statement where x and y are variables of integral or floating types. Admittedly, the expression lacks side effects and is therefore meaningless, but it is still an expression statement. On the other hand, the second line looks like a pointer declaration where T is a type defined by typedef and p is a pointer to that type.

x \* y;

T\* p;

However, syntactically it is the same thing: an identifier followed by an asterisk, another identifier and a semicolon. The parser cannot distinguish between the two cases, which is a problem. The only reasonably solution is to give the scanner access to the symbol table, so it can look up the whether the identifier is a typedef. If it is, the typedef token is returned; otherwise, the name token is returned. With these changes, the parser can distinguish between the two cases.

## The Scanner

The scanner of this section is similar to the scanners of Sections 2.1.4 and 4.1.2. However, there are more keyword and operators as well as more complicated regular expressions.

Scanner.gplex

%namespace CCompiler\_Main

%using CCompiler;

%using System.Numerics;

%{

public static FileInfo Path = null;

public static int Line = 1;

The first problem for the scanner to solve is to transform the escape sequences of characters and string into their proper characters. For instance, the character sequence ’\n’ shall be transformed into the character newline, with ASCII value 10.

The escape map (m\_escapeMap) holds the escape sequences of ANSI C:

|  |  |  |
| --- | --- | --- |
| **Escape Sequence** | **Code** | **ASCII value** |
| Zero Character | \0 | 0 |
| Alert (Beep, Bell) | \a | 7 |
| Backspace | \b | 8 |
| Form Feed Page Break | \f | 12 |
| Newline (Line Feed) | \n | 10 |
| Carriage Return | \r | 13 |
| Horizontal Tabulator | \t | 9 |
| Vertical Tabulator | \v | 11 |
| Single Quotation Mark | \' | 39 |
| Double Quotation Mark | \" | 34 |
| Question Mark | \? | 63 |
| Backslash | \\ | 92 |

private static IDictionary<char,char> m\_escapeMap =

new Dictionary<char,char>();

static Scanner() { // ASCII value

m\_escapeMap.Add('0', '\0'); // 0

m\_escapeMap.Add('a', '\a'); // 7

m\_escapeMap.Add('b', '\b'); // 8

m\_escapeMap.Add('f', '\f'); // 12

m\_escapeMap.Add('n', '\n'); // 10

m\_escapeMap.Add('r', '\r'); // 13

m\_escapeMap.Add('t', '\t'); // 9

m\_escapeMap.Add('v', '\v'); // 11

m\_escapeMap.Add('\'', '\''); // 39

m\_escapeMap.Add('\"', '\"'); // 34

m\_escapeMap.Add('?', '?'); // 63

m\_escapeMap.Add('\\', '\\'); // 92

}

The method SlashToChar takes a string with slash codes and returns the same string with the slash codes replaced by the corresponding characters.

public static string SlashToChar(string text) {

StringBuilder buffer = new StringBuilder(text);

We add three zero-characters at the end of the string to make sure the char1, char2, and char3 values are valid.

buffer.Append("\0\0\0");

for (int index = 0; buffer[index] != '\0'; ++index) {

if (buffer[index] == '\\') {

char char1 = buffer[index + 1],

char2 = buffer[index + 2],

char3 = buffer[index + 3];

If the slash is followed by an character stored in the slash map, we replace the slash and the character with the character corresponding to the slash code.

if (m\_escapeMap.ContainsKey(char1)) {

buffer.Remove(index, 2);

buffer.Insert(index, m\_escapeMap[char1]);

}

If the three characters following the slash are octal digits, we calculate the ASCII value for the character.

else if (IsOctal(char1) && IsOctal(char2) && IsOctal(char3)) {

int octValue = 64 \* CharToOctal(char1) +

8 \* CharToOctal(char2) +

CharToOctal(char3);

An ASCII value cannot exceed 255.

Assert.Error(octValue < 256, Message.Invalid\_octal\_sequence);

We remove four characters; that is, the slash and the following three octal digits, and then we insert the resulting character.

buffer.Remove(index, 4);

buffer.Insert(index, (char) octValue);

}

If the slash is followed by two octal digits only, we use them to calculate the ASCII value. In this case, we do not need to check whether the ASCII value is less than 256, since the highest possible value with two octal digits is 63 (778 = 8 \* 7 + 7 = 6310).

else if (IsOctal(char1) && IsOctal(char2)) {

int octValue = 8 \* CharToOctal(char1) +

CharToOctal(char2);

buffer.Remove(index, 3);

buffer.Insert(index, (char)octValue);

}

If the slash is followed by one octal digit only, we use it to calculate the ASCII value.

else if (IsOctal(char1)) {

int octValue = CharToOctal(char1);

buffer.Remove(index, 2);

buffer.Insert(index, (char) octValue);

}

If the slash is followed by a lowercase or uppercase ‘x’ and two hexadecimal digits, we use them to calculate the ASCII value. Neither in this case we need to check that the ASCII value does not exceed 255, since the highest possible value with two hexadecimal digits is 255 (FF16 = 15 \* 16 + 15 = 25510).

else if (char.ToLower(char1) == 'x') {

if (IsHex(char1) && IsHex(char2)) {

int hexValue = 16 \* CharToHex(char1) + CharToHex(char2);

buffer.Remove(index, 3);

buffer.Insert(index, (char) hexValue);

}

If the slash is followed by a lowercase or uppercase ‘x’ and hexadecimal digit only, we use it to calculate the ASCII value.

else if (IsHex(char1)) {

int hexValue = CharToHex(char1);

buffer.Remove(index, 2);

buffer.Insert(index, (char) hexValue);

}

If the slash is followed by capital or small ‘x’ without at least one hexadecimal digit, we report an error.

else {

Assert.Error(char1.ToString(),

Message.Invalid\_hexadecimal\_code);

}

}

Finally, if the slash is not followed by a character in the slash map, a capital or small ‘x’, or a octal digit, we report an error.

else {

Assert.Error(buffer[index + 1].ToString(),

Message.Invalid\_slash\_sequence);

}

}

}

When we have traversed through the string, we remove the three zero-character we added at the beginning and return the string.

buffer.Remove(buffer.Length - 3, 3);

return buffer.ToString();

}

The IsOctal and IsHex methods return true if the given character is and octal or hexadecimal digit, respectively.

private static bool IsOctal(char c) {

return "01234567".Contains(c.ToString());

}

private static bool IsHex(char c) {

return "0123456789abcdef".Contains(c.ToString().ToLower());

}

The CharToOctal and CharToHex methods return the numerical value corresponding to the given character.

private static int CharToOctal(char c) {

return "01234567".IndexOf(c);

}

private static int CharToHex(char c) {

return "0123456789abcdef".IndexOf(c.ToString().ToLower());

}

%}

Now we have reached the part of the scanner where we define the scanner rules. We need rules for names as well as octal, decimal, and hexadecimal, floating, string and character values. However, the regular expressions for values are a bit more complicated. There are several kinds of values:

* Octal. An octal value starts with a zero that is followed by zero or more digits between zero and eight, inclusive.
* Decimal. A decimal value is made up by one or more digits.
* Hexadecimal. A hexadecimal value starts with zero, followed by a small or capital x and one of more hexadecimal digits.
* Floating. A floating value starts with one or more digits followed by a dot and zero or more digits or starts with a dot followed by one or more digits. It is possible to add an exponent part that starts with a small or capital e, followed by a potential plus or minus sign and one or more digits. Finally, a floating value can also be made up by one or more digits without a dot, followed by a small or capital e, a possible plus or minus sign, and one or more digits.

The decimal values can be appended by the small or capital letter ‘u’ to indicate an unsigned value and the small or capital letter ‘s’ or ‘l’ to indicate a short or long value. In the same way, a floating value can be appended by a small or capital f or l to indicate a float or long double.

DECIMAL\_VALUE [\+\-]?[1-9][0-9]\*

OCTAL\_VALUE [\+\-]?0[0-7]\*

HEXADECIMAL\_VALUE [\+\-]?0[xX][0-9a-fA-F]+

POSTFIX ([uU]?[sSlL]?)|([sSlL]?[uU]?)

INTEGRAL\_VALUE ({DECIMAL\_VALUE}|{OCTAL\_VALUE}|{HEXADECIMAL\_VALUE}){POSTFIX}

DECIMAL\_PART [\+\-]?([0-9]+|[0-9]+\.[0-9]\*|\.[0-9]+)

EXPONENT\_PART (([eE][\+\-]?[0-9]+)?|[0-9]+[eE][\+\-]?[0-9]+)([fF]|[lL])?

FLOATING\_VALUE {DECIMAL\_PART}{EXPONENT\_PART}

A character is either two single quotation enclosing another single quotation preceded by a backslash (‘\’’) or two single quotation mark enclosing everything except single quotation marks.

CHAR\_VALUE (\'\\\'\')|(\'[^\']\*\')

A character starts with single quotation mark followed by one or more occurrences of everything except a single quotation mark and is terminated by another single quotation mark. However, there may be a single quotation mark if it is preceded by a backslash. Since the quotation mark and backslash are character with special meaning, we need to precede them with backslashes. However, note that we do not need the backslash when the quotation mark is enclosed by brackets.

CHAR\_VALUE \'(\\\'|[^'])\*\'

In the same way, a string starts with a double quotation mark followed by anything except double quotation mark and is terminated by another double quotation mark. However, the string may hold double quotations marks if they are proceeded by backslashes.

STRING\_VALUE \"(\\\"|[^"])\*\"

Register names are used internally only, when performing system calls. They start with the text “register\_” followed by the name of the register.

REGISTER\_NAME "register\_"[a-z]+

Names are used to identify variables, constants, struct and unions, enumerations, functions, and macros. They start with a letter or an underscore (’\_’), that is potentially followed by letters, digits, or underscores.

NAME [a-zA-Z\_][a-zA-Z0-9\_]\*

The path line is used to keeping track of the current line number. It is used by the preprocessor macro \_\_LINE\_\_, and when reporting errors. The path line starts with a dollar sign followed by anything except newline (the dot represents every character except newline) and ends with another dollar sign.

PATH\_LINE \$.\*\$

A white space is a space or any character that can substitute as a space; that is horizontal tabulator, return, newline, and form feed.

WHITE\_SPACE [ \t\r\n\f]

The next section holds the actions of keywords, operators, and the rules defined above.

%%

"auto" { return ((int) Tokens.AUTO); }

"break" { return ((int) Tokens.BREAK); }

"case" { return ((int) Tokens.CASE); }

"carry\_flag" { return ((int) Tokens.CARRY\_FLAG); }

"char" { return ((int) Tokens.CHAR); }

"const" { return ((int) Tokens.CONSTANT); }

"continue" { return ((int) Tokens.CONTINUE); }

"default" { return ((int) Tokens.DEFAULT); }

"do" { return ((int) Tokens.DO); }

"double" { return ((int) Tokens.DOUBLE); }

"else" { return ((int) Tokens.ELSE); }

"enum" { return ((int) Tokens.ENUM); }

"extern" { return ((int) Tokens.EXTERN); }

"float" { return ((int) Tokens.FLOAT); }

"for" { return ((int) Tokens.FOR); }

"goto" { return ((int) Tokens.GOTO); }

"int" { return ((int) Tokens.INT); }

"interrupt" { return ((int) Tokens.INTERUPT); }

"if" { return ((int) Tokens.IF); }

"jump\_register" { return ((int) Tokens.JUMP\_REGISTER); }

"long" { return ((int) Tokens.LONG); }

"register" { return ((int) Tokens.REGISTER); }

"return" { return ((int) Tokens.RETURN); }

"short" { return ((int) Tokens.SHORT); }

"signed" { return ((int) Tokens.SIGNED); }

"sizeof" { return ((int) Tokens.SIZEOF); }

"static" { return ((int) Tokens.STATIC); }

"struct" { return ((int) Tokens.STRUCT); }

"switch" { return ((int) Tokens.SWITCH); }

"syscall" { return ((int) Tokens.SYSCALL); }

"typedef" { return ((int) Tokens.TYPEDEF); }

"union" { return ((int) Tokens.UNION); }

"unsigned" { return ((int) Tokens.UNSIGNED); }

"while" { return ((int) Tokens.WHILE); }

"void" { return ((int) Tokens.VOID); }

"volatile" { return ((int) Tokens.VOLATILE); }

";" { return ((int) Tokens.SEMICOLON); }

":" { return ((int) Tokens.COLON); }

"," { return ((int) Tokens.COMMA); }

"." { return ((int) Tokens.DOT); }

"->" { return ((int) Tokens.ARROW); }

"..." { return ((int) Tokens.ELLIPSE); }

"(" { return ((int) Tokens.LEFT\_PAREN); }

")" { return ((int) Tokens.RIGHT\_PAREN); }

"{" { return ((int) Tokens.LEFT\_BLOCK); }

"}" { return ((int) Tokens.RIGHT\_BLOCK); }

"[" { return ((int) Tokens.LEFT\_SQUARE); }

"]" { return ((int) Tokens.RIGHT\_SQUARE); }

"\*" { return ((int) Tokens.ASTERRISK); }

"?" { return ((int) Tokens.QUESTION\_MARK); }

"||" { return ((int) Tokens.LOGICAL\_OR); }

"&&" { return ((int) Tokens.LOGICAL\_AND); }

"!" { return ((int) Tokens.LOGICAL\_NOT); }

"&" { return ((int) Tokens.AMPERSAND); }

"^" { return ((int) Tokens.BITWISE\_XOR); }

"|" { return ((int) Tokens.BITWISE\_OR); }

"~" { return ((int) Tokens.BITWISE\_NOT); }

"==" { return ((int) Tokens.EQUAL); }

"!=" { return ((int) Tokens.NOT\_EQUAL); }

"<" { return ((int) Tokens.LESS\_THAN); }

"<=" { return ((int) Tokens.LESS\_THAN\_EQUAL); }

">" { return ((int) Tokens.GREATER\_THAN); }

">=" { return ((int) Tokens.GREATER\_THAN\_EQUAL); }

"<<" { return ((int) Tokens.LEFT\_SHIFT); }

">>" { return ((int) Tokens.RIGHT\_SHIFT); }

"+" { return ((int) Tokens.PLUS); }

"-" { return ((int) Tokens.MINUS); }

"/" { return ((int) Tokens.DIVIDE); }

"%" { return ((int) Tokens.MODULO); }

"++" { return ((int) Tokens.INCREMENT); }

"--" { return ((int) Tokens.DECREMENT); }

"=" { return ((int) Tokens.ASSIGN); }

"+=" { return ((int) Tokens.ADD\_ASSIGN); }

"-=" { return ((int) Tokens.SUBTRACT\_ASSIGN); }

"\*=" { return ((int) Tokens.MULTIPLY\_ASSIGN); }

"/=" { return ((int) Tokens.DIVIDE\_ASSIGN); }

"%=" { return ((int) Tokens.MODULO\_ASSIGN); }

"<<=" { return ((int) Tokens.LEFT\_SHIFT\_ASSIGN); }

">>=" { return ((int) Tokens.RIGHT\_SHIFT\_ASSIGN); }

"&=" { return ((int) Tokens.AND\_ASSIGN); }

"^=" { return ((int) Tokens.XOR\_ASSIGN); }

"|=" { return ((int) Tokens.OR\_ASSIGN); }

When we encounter a register name, we look up the register, and report an error if the name is not the name of a register.

{REGISTER\_NAME} {

{ Register register;

string text = yytext.Substring(9);

if (Enum.TryParse<Register>(text, out register)) {

yylval.register = register;

return ((int) Tokens.REGISTER\_NAME);

}

Assert.Error(text, Message.Unknown\_register);

}

}

When we encounter a name, we need to check whether it represent a typedef name or a regular name; that is, the name of a variable, constant, struct or union, enumeration, or a function. We look up the name in the current symbol table and return the typedef name token in case of a typedef name. In case of a regular name, we return the name token.

{NAME} {

{ Symbol symbol = SymbolTable.CurrentTable.LookupSymbol(yytext);

if ((symbol != null) && symbol.IsTypedef()) {

yylval.type = symbol.Type;

return ((int) Tokens.TYPEDEF\_NAME);

}

else {

yylval.name = yytext;

return ((int) Tokens.NAME);

}

}

}

When it comes to integral values, we need to decide the sign, base, and type of the value.

{INTEGRAL\_VALUE} {

{ string text = yytext.Trim().ToLower();

The ToUInt64 method below does not accept plus or minus signs. Therefore, we need to remove the sign and assign true to the minus variable in case of a negative value.

bool minus = false;

if (text.StartsWith("+")) {

text = text.Substring(1);

}

else if (text.StartsWith("-")) {

minus = true;

text = text.Substring(1);

}

Then we decide the base of the value. If the text start with “0x” or “0X”m the value is a hexadecimal with the base 16. If the text start with “0”, the value is an octal value with the base 8. Otherwise, it is a decimal value with the base 10.

int fromBase;

if (text.StartsWith("0x")) {

fromBase = 16;

text = text.Substring(2);

}

else if (text.StartsWith("0")) {

fromBase = 8;

}

else {

fromBase = 10;

}

Then we decide the type of the value. If the text ends with any combination of small or capitel “s” and “u”, the type is unsigned small integer.

“s” or “S”, the type is small integer. If the text ends with “l” or “L”, the type is large integer. If the text ends with “u” or “U”, the type is unsigned. Otherwise, the type is signed integer.

CCompiler.Type type;

if (text.EndsWith("us") || text.EndsWith("su")) {

type = CCompiler.Type.UnsignedShortIntegerType;

text = text.Substring(0, text.Length - 2);

}

If the text ends with any combination of small or capital “l” and “u”, the type is large unsigned integer.

else if (text.EndsWith("ul") || text.EndsWith("lu")) {

type = CCompiler.Type.UnsignedLongIntegerType;

text = text.Substring(0, text.Length - 2);

}

If the text ends with a single small or capital “s”, the type is signed short integer.

else if (text.EndsWith("s")) {

type = CCompiler.Type.SignedShortIntegerType;

text = text.Substring(0, text.Length - 1);

}

If the text ends with a single small or capital “l”, the type is signed long integer.

else if (text.EndsWith("l")) {

type = CCompiler.Type.SignedLongIntegerType;

text = text.Substring(0, text.Length - 1);

}

If the text ends with a single small or capital “u”, the type is unsigned integer.

else if (text.EndsWith("u")) {

type = CCompiler.Type.UnsignedIntegerType;

text = text.Substring(0, text.Length - 1);

}

Otherwise, the type is signed integer.

else {

type = CCompiler.Type.SignedIntegerType;

}

Finally, we decide the value by calling the ToUInt64 method in the Convert standard class.

try {

ulong unsignedValue = Convert.ToUInt64(text, fromBase);

We convert the value to an object of the standard BigInteger class, and if minus is true, we make the value negative.

BigInteger bigValue = new BigInteger(unsignedValue);

if (minus) {

bigValue = -bigValue;

}

We then create a symbol of the value that we return, We also store a static symbol in the global static set.

yylval.symbol = new Symbol(type, bigValue);

yylval.symbol.StaticSymbol =

ConstantExpression.Value(yylval.symbol.UniqueName, type, bigValue);

SymbolTable.StaticSet.Add(yylval.symbol.StaticSymbol);

}

catch (OverflowException) {

Assert.Error("X " + type + ": " + text, Message.Value\_overflow);

}

return ((int) Tokens.VALUE);

}

}

When it comes to floating values, we do not need to look for its plus or minus sign. However, we do notice if the text ends with small or capital ‘f’ or ‘l’, in which case the type becomes float or long double, respectively. If the text does not end with ‘f’ or ‘l’, the type becomes double.

{FLOATING\_VALUE} {

{ string text = yytext.ToLower();

CCompiler.Type type = CCompiler.Type.DoubleType;

if (text.EndsWith("f")) {

type = CCompiler.Type.FloatType;

text = text.Substring(0, text.Length - 1);

}

else if (text.EndsWith("l")) {

type = CCompiler.Type.LongDoubleType;

text = text.Substring(0, text.Length - 1);

}

We call the Parse method of the decimal standard class to obtain the value. Then we create and return a symbol the value. We also add a static value to the global static set.

try {

decimal value = decimal.Parse(text, NumberStyles.Float);

yylval.symbol = new Symbol(type, value);

yylval.symbol.StaticSymbol =

ConstantExpression.Value(yylval.symbol.UniqueName, type, value);

SymbolTable.StaticSet.Add(yylval.symbol.StaticSymbol);

}

catch (OverflowException) {

Assert.Error("Y " + type + ": " + text, Message.Value\_overflow);

}

return ((int) Tokens.VALUE);

}

}

In the case of a character value, we call the SlashToChar method above to clear the text from escape characters. We then check that it holds three characters (including the single quotation marks), and return a symbol holding the value.

{CHAR\_VALUE} {

{ CCompiler.Type type = new CCompiler.Type(Sort.Signed\_Char);

string text = SlashToChar(yytext);

Assert.Error(text.Length == 3, yytext, Message.Invalid\_char\_sequence);

yylval.symbol = new Symbol(type, (BigInteger) ((int) text[1]));

return ((int) Tokens.VALUE);

}

}

In the case of a string value, we also call the SlashToChar method above to clear the text from escape characters. We then return a symbol holding the value. In this case we do not need to check the length of the string.

{STRING\_VALUE} {

{ CCompiler.Type type = new CCompiler.Type(Sort.String);

string text = SlashToChar(yytext);

object value = text.Substring(1, text.Length - 2);

yylval.symbol = new Symbol(type, value);

yylval.symbol.StaticSymbol =

ConstantExpression.Value(yylval.symbol.UniqueName, type, value);

SymbolTable.StaticSet.Add(yylval.symbol.StaticSymbol);

return ((int) Tokens.VALUE);

}

}

Path lines are generated by the preprocessor for the \_\_FILE\_\_ and \_\_LINE\_\_ macros to hold the correct file name and line number. It is also used by the compiler when reporting error. Its text is made up by two dollar signs (‘$’) at the beginning and end, and the file name and line number separated by a comma (‘,’). The text also holds plus signs (‘+’) in place of spaces.

{PATH\_LINE} {

{ string text = yytext.Substring(1, yyleng - 2);

int index = text.IndexOf(',');

Path = new FileInfo(text.Substring(0, index).Replace("+", " "));

Line = int.Parse(text.Substring(index + 1));

}

}

In case of a white space, we only update the Line variable if the character is a newline. Note that we do not return anything in this case, which causes the scanner to proceed with the next character in the input stream.

{WHITE\_SPACE} {

if (yytext.Equals("\n")) {

++Line;

}

}

Finally, when the scanner finds a character that has not been handled by the rules above, a error is reported.

. { Assert.Error(yytext, Message.Unknown\_character); }

# Parsing

The parser can be considered the heart of the compiler. It requests token from the scanner, checks that the source code complies with the grammar, builds the symbol table, performs type checking, and generate the middle code.

This chapter describes the parser of the compiler, how it parses the source code to test whether is complies to the grammar. However, the parser does also performs type checking, store symbols in the symbol table and generates middle code.

MainParser.gppg

%namespace CCompiler\_Main

%partial

%using CCompiler;

%using System.Numerics;

%{

public static Stack<Specifier> SpecifierStack = new Stack<Specifier>();

public static Stack<BigInteger> EnumValueStack = new Stack<BigInteger>();

public static int CallDepth = 0;

%}

The parser is made up by tokens and rules, where tokens corrensponds to operators, keyword and characters. The following tokens are used by the parser, and they are also returned by the scanner in the previous chapter.

A variable or function has auto, register, static, extern, or typedef storage.

%token AUTO REGISTER STATIC EXTERN TYPEDEF

A variable can also be qualified as constant or volatile.

CONSTANT VOLATILE

An integral type can by signed or unsigned, and char, short int, int, or long int.

SIGNED UNSIGNED CHAR SHORT INT LONG

An floating type can be float, double, or long double.

FLOAT DOUBLE

Moreover, there is

ENUM STRUCT UNION

Finally, there is also the void type, that marks the absence of a type. It is used to to mark the absence of a parameter list or return type of a function, and as pointer to void.

VOID

The asterrisk (’\*’) is used both for multiplication and dereferenceereing pointers.

PLUS MINUS ASTERRISK DIVIDE MODULO INCREMENT DECREMENT

EQUAL NOT\_EQUAL LESS\_THAN LESS\_THAN\_EQUAL GREATER\_THAN

GREATER\_THAN\_EQUAL

The assigmnent is both simple and compund.

ASSIGN ADD\_ASSIGN SUBTRACT\_ASSIGN MULTIPLY\_ASSIGN

DIVIDE\_ASSIGN MODULO\_ASSIGN AND\_ASSIGN OR\_ASSIGN XOR\_ASSIGN

LEFT\_SHIFT\_ASSIGN RIGHT\_SHIFT\_ASSIGN

LOGICAL\_OR LOGICAL\_AND LOGICAL\_NOT

The ampersand (’&’) is used both bitwise and as well as address operator.

AMPERSAND BITWISE\_XOR BITWISE\_OR BITWISE\_NOT

The ellipse is made up by threee dots ’...’ and is used when defining functions with a varriable number of parameters, sush as printf or scanf.

QUESTION\_MARK COLON COMMA SEMICOLON ELLIPSE DOT ARROW SIZEOF

In C, we have regular parentheses ’(’ and ’)’, blocks ’{’ and ’}’, and squares ’[’ and ’]’

LEFT\_PARENTHESIS RIGHT\_PARENTHESIS LEFT\_BLOCK RIGHT\_BLOCK

LEFT\_SQUARE RIGHT\_SQUARE

We need a set of keyword for the statements of C.

IF ELSE SWITCH CASE DEFAULT FOR WHILE DO BREAK CONTINUE RETURN GOTO

The interrupt, jump\_register, syscall, and carry flag tokens are used for internal system calls only.

INTERRUPT JUMP\_REGISTER SYSCALL CARRY\_FLAG

Note that we do not need any tokens for comments, since they have been taken care of by the preprocessor.

The next part of the scanner is the union section, where we define the attributes of the rules below.

%union {

public string name;

The Register enumeration holds the registers of the target machine.

public Register register;

Each symbol has a type, which is described by the Type class.

public CCompiler.Type type;

public List<CCompiler.Type> type\_list;

public Sort sort;

public Symbol symbol;

public IDictionary<string,Symbol> symbol\_map;

public ISet<Pair<Symbol,bool>> symbol\_bool\_pair\_set;

public Pair<Symbol,bool> symbol\_bool\_pair;

public Pair<string,Symbol> string\_symbol\_pair;

public List<Pair<string,Symbol>> string\_symbol\_pair\_list;

public List<string> string\_list;

public Declarator pointer\_declarator;

public List<Declarator> declarator\_list;

public MiddleOperator middleOperator;

public Expression expression;

public List<Expression> expression\_list;

public Statement statement;

public Pair<List<Pair<string,Symbol>>,Boolean> parameter\_pair;

public List<MiddleCode> middle\_code\_list;

public object obj;

public List<object> object\_list;

}

%token <name> NAME

%token <register> REGISTER\_NAME

%token <type> TYPEDEF\_NAME

%token <symbol> VALUE

%type <obj> declaration\_specifier declaration\_specifier\_list\_x

%type <object\_list> declaration\_specifier\_list

%type <name> optional\_name

%type <type> struct\_or\_union\_specifier

%type <sort> struct\_or\_union

%type <type> enum\_specifier

%type <symbol\_bool\_pair\_set> enum\_list

%type <symbol\_bool\_pair> enum

%type <middle\_code\_list> declarator\_list

initialization\_bitfield\_pointer\_declarator declaration

%type <pointer\_declarator> pointer\_declarator direct\_pointer\_declarator

%type <type\_list> optional\_pointer\_list pointer\_list

%type <type> pointer

%type <parameter\_pair> optional\_parameter\_ellipse\_list

parameter\_ellipse\_list

%type <string\_symbol\_pair\_list> parameter\_list

%type <string\_symbol\_pair> parameter\_declaration

%type <string\_list> optional\_name\_list identifier\_list

%type <object\_list> initializer\_list

%type <obj> initializer

%type <type> type\_name

%type <pointer\_declarator> abstract\_pointer\_declarator direct\_abstract\_pointer\_declarator

%type <middleOperator> assignment\_operator equality\_operator

relation\_operator add\_operator shift\_operator

multiply\_operator prefix\_add\_operator

increment\_operator

%type <expression> optional\_constant\_integral\_expression

constant\_integral\_expression optional\_expression

expression assignment\_expression

condition\_expression logical\_or\_expression

logical\_and\_expression BITWISE\_OR\_expression

bitwise\_xor\_expression bitwise\_and\_expression

equality\_expression relation\_expression

shift\_expression add\_expression

multiply\_expression type\_cast\_expression

prefix\_expression postfix\_expression

primary\_expression

%type <statement> optional\_statement\_list statement

closed\_statement opened\_statement

%type <expression\_list> optional\_argument\_expression\_list

argument\_expression\_list

%start translation\_unit

## Declarations

%%

Finally, we have reached the last (and largest) part of the parser, where the rules are defined. The translation\_unit rule is the start of the grammar, and represent the whole of a source code file.

translation\_unit:

external\_declaration

| translation\_unit external\_declaration;

external\_declaration:

function\_definition

| declaration;

The rules of the parser calls corrensponding methods of the MiddleCodeGenerator class to perform type checking, build the symbol table, and generate the middle code.

MiddleCodeGenerator.cs

using System;

using System.IO;

using System.Linq;

using System.Numerics;

using System.Collections.Generic;

namespace CCompiler {

public class MiddleCodeGenerator {

public static MiddleCode AddMiddleCode(List<MiddleCode> codeList,

MiddleOperator op, object operand0 = null,

object operand1 = null, object operand2 = null) {

MiddleCode middleCode = new MiddleCode(op, operand0, operand1, operand2);

codeList.Add(middleCode);

return middleCode;

}

### Backpatching

When generating middle code, there are numerous occasions that we generate jump instructions without knowing where to jump. In those cases, we store the instructions in sets and later go back and fill in the jump address. This process is called *backpatching*.

The first Backpatch methods take set of jump instructions and a list of instructions, where its first instruction is the target of the jump instructions. If the list is empty, we add an empty instruction in order to make sure there is always a jump target. Then we call the second Backpatch method with the set of jump instruction and the target instruction.

MiddleCodeGenerator.cs

public static void Backpatch(ISet<MiddleCode> sourceSet,

List<MiddleCode> list) {

if (list.Count == 0) {

AddMiddleCode(list, MiddleOperator.Empty);

}

Backpatch(sourceSet, list[0]);

}

The second Backpatch method iterates through the set of jump instructions and set its first operand (index 0) to the target. The first operand of each jump instruction is its target address.

public static void Backpatch(ISet<MiddleCode> sourceSet,

MiddleCode target) {

foreach (MiddleCode source in sourceSet) {

source[0] = target;

}

}

### Function Definition

A function definition is made up by declarator, possible preceded by a declaration specifier list, followed by an optional declaration list, and a block with an optional statement list.

A function definition is made up by pointer\_declarator, possible preceded by a declaration specifier list, followed by an optional declaration list, and a block with an optional statement list.

For instance, in the function below unsigned long int is the declaration specifier list, square(int value) is the pointer\_declarator, and return value \* value is the statement list.

unsigned long int square(int value) { return value \* value; }

In C, there are two ways to define a function, the old way where the parameter list hold holds the name of the parameters and where there types are defined afterwards, and new way where the parameter list holds the names and types of the parameters. In the following examples f is defined by the old way and g is defined by the new way.

int f(a, b)

int a;

double b; {

return a + b;

}

int f(int a, double b) {

return a + b;

}

The first part of the function definition rule handles the case where there is specifier list of at least specifier present. There are three methods to handle the function definition: FunctionHeader that XXX, CheckFunctionDefinition that makes sure the function is defined in either the old or new way, and FunctionEnd that generates the assembly code of the function and saves it in the static set.

MainParser.gppg

function\_definition:

declaration\_specifier\_list\_x pointer\_declarator {

MiddleCodeGenerator.FunctionHeader

(SpecifierStack.Pop(), $2);

}

optional\_declaration\_list {

MiddleCodeGenerator.CheckFunctionDefinition();

}

LEFT\_BLOCK optional\_statement\_list RIGHT\_BLOCK {

MiddleCodeGenerator.BackpatchGoto();

MiddleCodeGenerator.FunctionEnd($7);

};

The second rule handles the case where there is no spcifier list. In that case, the function return type is assumed to be signed int.

| pointer\_declarator {

MiddleCodeGenerator.FunctionHeader(null, $1);

}

optional\_declaration\_list {

MiddleCodeGenerator.CheckFunctionDefinition();

}

LEFT\_BLOCK optional\_statement\_list RIGHT\_BLOCK {

MiddleCodeGenerator.BackpatchGoto();

MiddleCodeGenerator.FunctionEnd($6);

};

optional\_declaration\_list:

/\* Empty \*/

| declaration\_list;

The generateFunctionHeader is called before the middle code of the function has been generated.

The generateAssignmentDeclarator method is called for a declarator with an assignment. We check that the declarator is not located inside a struct or union (the scope of the symbol table is struct or union), since it is not allowed to initialize struct or union members. We also check that the declarator is not a function, since functions cannot be initialized. If the declarator is defined in the global scope and does not have the static storage it has external linkage and is visible from other source files, the linker will make it visible in a later phase.

An extern or typedef declarator cannot be initialized. However, if it is auto or register, we call GenerateInitializer.generateStack that generates a suitable series of assignment instructions. If it static, we instead call ValueConversion.storeStaticSymbol that will store the initialized value (that is always constant) in a block of bytes, that will be paced in the final execution file by the linker.

The rule handles a definition of a function. A function *definition* is made up by the specifier list and declaration as well as the body holding the code, while a function *declaration* is made up by the specifier list and declaration only. For instance, in the following examples, f is a function definition while g is a function declaration:

int f(int i) { return i + 1; }

int g(int i);

MiddleCodeGenerator.cs

A general function pointer\_declarator may be unnamed; however, a function definition must be named and of course be a function declaration.

public static void FunctionHeader(Specifier specifier,

Declarator declarator) {

Storage? storage;

Type returnType;

If specifier is not null, we extract the storage and return type of the function from the specifier.

if (specifier != null) {

storage = specifier.Storage;

returnType = specifier.Type;

}

If specifier is null, there was no return type defined. In that case, we set the storage to extern and the return type to signed int.

else {

storage = Storage.Extern;

returnType = Type.SignedIntegerType;

}

We add the return type to the pointer declarator. In this way the type become complete: a function with a return type.

declarator.Add(returnType);

Then we perform some error checking. Technically, in a function declaration, the function may lack a name. However, in definition, the function must have a name.

Assert.Error(declarator.Name != null,

Message.Unnamed\_function\_definition);

Technically, the declaration may have any type, and we must check that it really is a function declaration.

Assert.Error(declarator.Type.IsFunction(), declarator.Name,

Message.Not\_a\_function);

The public and static member variable CurrentFunction is a reference to the symbol holding the current function. This reference is referred on several occasion during the parsing process.

SymbolTable.CurrentFunction =

new Symbol(declarator.Name, storage, declarator.Type);

A function definition must be static or extern, it cannot be auto, register, or typedef.

Assert.Error(SymbolTable.CurrentFunction.IsStaticOrExtern(),

declarator.Name,

Message.A\_function\_must\_be\_static\_or\_extern);

Every function definition, as well as static variables, are added to the static set. In will eventually be translated into assembly code.

SymbolTable.CurrentTable.AddSymbol(SymbolTable.CurrentFunction);

If this function is the main function, the return type must be void or (signed or unsigned) integer.

if (SymbolTable.CurrentFunction.UniqueName.Equals("main")) {

Assert.Error(returnType.IsVoid() || returnType.IsInteger(), "main",

Message.Function\_main\_must\_return\_void\_or\_integer);

}

Finally, we create a new symbol for the function. Every symbol of declared in the function will be stored in the new symbol table, or in its sub tables.

SymbolTable.CurrentTable =

new SymbolTable(SymbolTable.CurrentTable, Scope.Function);

}

The CheckFunctionDefinition method checks that the definition is correct, especially in case of old-style definition. In the old-style function definition, the parameter list must match the declaration list. In the new-style function definition, the declaration list must be empty since it is not allowed to mix the old and new style.

public static void CheckFunctionDefinition() {

Type funcType = SymbolTable.CurrentFunction.Type;

if (funcType.Style == Type.FunctionStyle.Old) {

List<string> nameList = funcType.NameList;

IDictionary<string,Symbol> entryMap =

SymbolTable.CurrentTable.EntryMap;

In the old-style case, we must make sure the number of parameters equals the number of declarations.

Assert.Error(nameList.Count == entryMap.Count,

SymbolTable.CurrentFunction.Name, Message.

Unmatched\_number\_of\_parameters\_in\_old\_\_style\_function\_definition);

Then we iterate through the parameter list to make sure they are all declared, and to assigned them the correct offset. When the parameters were declared they were given offsets. However, since the parameters and declaration may come in different order, we need to reassign the declarations’ offsets.

int offset = SymbolTable.FunctionHeaderSize;

foreach (string name in nameList) {

Symbol symbol;

If the parameter has not been declared, we report an error.

if (!entryMap.TryGetValue(name, out symbol)) {

Assert.Error(name, Message.

Undefined\_parameter\_in\_old\_\_style\_function\_definition);

}

symbol.Offset = offset;

offset += symbol.Type.Size();

}

}

In case of a new-style function definition, we need to make sure the entry map of the current function is empty. If it is not empty, we have a mix of old-style and new-style function definition, which is not allowed.

else {

Assert.Error(SymbolTable.CurrentTable.EntryMap.Count == 0,

Message.New\_and\_old\_style\_mixed\_function\_definition);

Then we iterate through the parameters list and add each parameter to the current symbol table of the function. In this way, each parameter is located at a memory offset so they can be assigned in function calls and be treated like regular variables inside the function.

foreach (Pair<string,Symbol> pair in funcType.ParameterList){

SymbolTable.CurrentTable.AddSymbol(pair.Second);

}

}

When the parameter list has been settled, we check if the function is the main function. In that case, some special rules apply.

if (SymbolTable.CurrentFunction.UniqueName.Equals("main")) {

First, we call the InitializationCodeList method that adds a few lines of code that initializes the system and will be placed before the code of the main function.

AssemblyCodeGenerator.InitializationCodeList();

We extract the parameter type list from the current function. In case of old-style definition, this list will be null. In case of new-style definition, there is only two allowed parameter list:

* No parameters, marked with void.

int main(void) { /\* ... \*/ }

* Command line arguments.

int main(int argc, char \*argv[]) { /\* ... \*/ } or int main(int argc, char \*\*argv) { /\* ... \*/ }

The two cases are actually the same case, since arrays are changed to pointers in parameter types. Also note that do not check the names of the parameters, the parameters can have any names. However, the names cannot be omitted.

List<Type> typeList =

SymbolTable.CurrentFunction.Type.TypeList;

If typelist is not null, and there are two parameters, we check if the command line argument case applies. The first parameter shall be a signed or unsigned integer while the second parameter shall be a pointer to pointer to signed or unsigned character.

if ((typeList != null) && (typeList.Count == 2)) {

Assert.Error(typeList[0].IsInteger() &&

typeList[1].IsPointer() &&

typeList[1].PointerType.IsPointer() &&

typeList[1].PointerType.PointerType.IsChar(),

"main", Message.Invalid\_parameter\_list);

AssemblyCodeGenerator.ArgumentCodeList();

}

If typelist is null, we have the old-style function definition, in which case we do not perform any type checking at all. If typelist is empty, we have the void-marked empty parameter list.

else {

Assert.Error((typeList == null) || (typeList.Count == 0),

"main", Message.Invalid\_parameter\_list);

}

}

}

The FunctionEnd method is called when the statements of the function have been parsed. Its task is to check the return statement at the end of the function and generate a static symbol holding the assembly code of the function.

The first step is to backpatch the next set of the statement. Each statement has a next set holding jump instruction out of the statement. For instance, a for or while statements has a next sets holding jumps out of the statement. We add an new empty statement and backpatch the next set to it.

public static void FunctionEnd(Statement statement) {

MiddleCode nextCode =

AddMiddleCode(statement.CodeList, MiddleOperator.Empty);

Backpatch(statement.NextSet, nextCode);

If the return type of the function is void, we need to add a return instruction at the end of the function. However, If the current function is the main function, we instead add an exit statement that will return the control to the enclosing system with exit code zero.

if (SymbolTable.CurrentFunction.Type.ReturnType.IsVoid()) {

if (SymbolTable.CurrentFunction.UniqueName.Equals("main")) {

Type signedShortType = new Type(Sort.Signed\_Short\_Int);

Symbol zeroSymbol = new Symbol(signedShortType, (BigInteger.Zero));

AddMiddleCode(statement.CodeList, MiddleOperator.Exit,

null, zeroSymbol);

}

If the function is not the main function, we add a return middle code instruction that will return control to the caller function.

else {

AddMiddleCode(statement.CodeList, MiddleOperator.Return);

}

}

We then add a FunctionEnd middle code instruction. Its only purpose is that the middle code optimizer will report an error it it is possible to reach the EndFunction instruction in a function that does not return void.

AddMiddleCode(statement.CodeList, MiddleOperator.FunctionEnd,

SymbolTable.CurrentFunction);

When we finally have added all middle code instructions, we optimize the middle code of the function. The optimizer transforms the middle code in several ways. See chapter XXX.

MiddleCodeOptimizer middleCodeOptimizer =

new MiddleCodeOptimizer(statement.CodeList);

middleCodeOptimizer.Optimize();

When the middle code has been optimized, we generate assembly code of the function. See chapter XXX.

List<AssemblyCode> assemblyCodeList = new List<AssemblyCode>();

AssemblyCodeGenerator.GenerateAssembly(statement.CodeList,

assemblyCodeList);

As mention in the introduction chapter of this book, we generate two kinds of target code. In the first part of the book we generate assembly code text for 64-bit Linux on Intel processors, which is than assembled and linked in traditional manner. In the second part of the book we generate an executable target file for 16-bit Windows. In the case where we need to distinguish between the two cases, we test whether the Start.Linux or Start.Windows static variable is true. In this chapter, we perform the actions for the Linux target machine. See chapter XXX for the Windows target machine.

In the Linux case, we need to generate the assembly text from the assmebly code list. We also need to generate the extern set; that is, the set of accesses of static variables with external linkage and the calls to all functions with external linkage.

if (Start.Linux) {

List<string> textList = new List<string>();

ISet<string> externSet = new HashSet<string>();

GenerateStaticInitializerLinux.LinuxTextList(assemblyCodeList,

textList, externSet);

When the text list and access set has been generated, we create a static symbol that we add to the static set. In this way, the assembly code text of the function will be added to final assembly file.

StaticSymbol staticSymbol =

new StaticSymbolLinux(StaticSymbolLinux.TextOrData.Text,

SymbolTable.CurrentFunction.UniqueName,

textList, externSet);

SymbolTable.StaticSet.Add(staticSymbol);

}

Finally, we set the former symbol table, which is the parent table of this function’s table, to be the current table. We also set the current function to null, which strictly speaking in not necessary. However, it is logical to set it to null when the parser is not parser the code of a function.

SymbolTable.CurrentTable = SymbolTable.CurrentTable.ParentTable;

SymbolTable.CurrentFunction = null;

}

### Specifier List

A declaration is made up by a declaration specifier list, an optional pointer\_declarator list, and a semicolon. For instance, in the following declaration, static long int is the specifier list while \*p, and a[3] are pointer declarators.

static long int \*p, a[3];

There is possible to define a declaration without a pointer\_declarator list. This feature is useful when declare a name struct or union:

struct \_s {int i};

However, this feature makes it also possible to declare unnamed specifier lists. The following declaration are syntactically correct but without meaning, they will be ignored by the compiler.

unsigned short int;

union {short s;}

However, the following declaration has meaning since the enumeration values can be used.

enum {ZERO = 0, ONE = 1, TEN = 10};

MainParser.gppg

declaration:

declaration\_specifier\_list SEMICOLON {

SpecifierStack.Push(Specifier.SpecifierList($1));

$$ = new List<MiddleCode>();

}

| declaration\_specifier\_list\_x declarator\_list SEMICOLON {

SpecifierStack.Pop();

$$ = $2;

};

optional\_declaration\_specifier\_list\_x:

/\* empty \*/

| declaration\_specifier\_list\_x;

declaration\_specifier\_list\_x:

declaration\_specifier\_list {

SpecifierStack.Push(Specifier.SpecifierList($1));

};

A declaration specifier list is made up by one or more declaration specifiers, which are stored in a list.

declaration\_specifier\_list:

declaration\_specifier {

$$ = new List<object>();

$$.Add($1);

}

| declaration\_specifier declaration\_specifier\_list {

$2.Add($1);

$$ = $2;

};

There a set of different kind of specifiers. The following specifiers return a value of the Mask enumeration.

* Storage: auto, register, static, extern, or typedef
* Type quailifiers: constant or volatile
* Type specifiers: void, char, short, int, long, float, double, signed, or unsigned

The following specifiers return a reference to the Type class.

* Struct or union specifier
* Enum specifier
* Typedef name

declaration\_specifier:

AUTO { $$ = Mask.Auto; }

| REGISTER { $$ = Mask.Register; }

| STATIC { $$ = Mask.Static; }

| EXTERN { $$ = Mask.Extern; }

| TYPEDEF { $$ = Mask.Typedef; }

| CONSTANT { $$ = Mask.Constant; }

| VOLATILE { $$ = Mask.Volatile; }

| VOID { $$ = Mask.Void; }

| CHAR { $$ = Mask.Char; }

| SHORT { $$ = Mask.Short; }

| INT { $$ = Mask.Int; }

| LONG { $$ = Mask.Long; }

| FLOAT { $$ = Mask.Float; }

| DOUBLE { $$ = Mask.Double; }

| SIGNED { $$ = Mask.Signed; }

| UNSIGNED { $$ = Mask.Unsigned; }

| struct\_or\_union\_specifier { $$ = $1; }

| enum\_specifier { $$ = $1; }

| TYPEDEF\_NAME { $$ = $1; };

#### Struct and Union

A struct or union specifier can be stated with an optional name and a declaration list within brackets, or just a name. When the declaration list is parsed, each member of the struct or union becomes added to the entry map of the current symbol table. Before the parsing of the declaration list we create a new symbol table for the members of the declaration list.

MainParser.gppg

struct\_or\_union\_specifier:

struct\_or\_union optional\_name {

The call to StructOrUnionHeader adds the struct or union to the symbol table, if the optional name is not null.

MiddleCodeGenerator.StructUnionHeader($2, $1);

SymbolTable.CurrentTable =

new SymbolTable(SymbolTable.CurrentTable, (Scope) $1);

}

The declaration list add the members of the struct or union to the current symbol table.

LEFT\_BLOCK declaration\_list RIGHT\_BLOCK {

The call to StructOrUnionSpecifier adds the member map of the struct or union to the symbol table, if the optional name is not null.

$$ = MiddleCodeGenerator.StructUnionSpecifier($2, $1);

SymbolTable.CurrentTable =

SymbolTable.CurrentTable.ParentTable;

}

In case of a struct of union without a declaration list, we look up the name.

| struct\_or\_union NAME {

$$ = MiddleCodeGenerator.LookupStructUnionSpecifier($2, $1);

};

struct\_or\_union:

STRUCT { $$ = Sort.Struct; }

| UNION { $$ = Sort.Union; };

optional\_name:

/\* Empty \*/ { $$ = null; }

| NAME { $$ = $1; };

declaration\_list:

declaration

| declaration\_list declaration;

The StructOrUnionHeader adds the struct or union to the tag map if the optional name is not null. We need to do this in order for recursive pointers to work. For instance:

struct Cell {

int value;

struct Cell\* next;

}

If we do not perform this action, there is a risk that next in the code above will point to another struct with the same name, defined in a surronding block.

MiddleCodeGenerator.cs

public static void StructUnionHeader(string optionalName, Sort sort) {

if (optionalName != null) {

Type type = new Type(sort);

SymbolTable.CurrentTable.AddTag(optionalName, type);

}

}

The StructUnionSpecifier method adds a struct or union. The sort parameter is either Sort.Struct or Sort.Union. If the optional name is not null, we look up the type (it has been added by StructOrUnionHeader above) and sets its member map, which is given by the entry map of the current symbol table.

public static Type StructUnionSpecifier(string optionalName, Sort sort) {

if (optionalName != null) {

Type type = SymbolTable.CurrentTable.LookupTag(optionalName, sort);

type.MemberMap = SymbolTable.CurrentTable.EntryMap;

return type;

}

If the optional name is null, we create and return a type with the entry map of the current symbol table.

else {

return (new Type(sort, SymbolTable.CurrentTable.EntryMap));

}

}

The LookupStructUnion method looks a struct or union. The sort parameter is either Sort.Struct or Sort.Union.

public static Type LookupStructUnion(string name, Sort sort) {

Type type = SymbolTable.CurrentTable.LookupTag(name, sort);

If the struct or union exists, we simply return its type.

if (type != null) {

return type;

}

If the struct or union does not exist, we create a new type and add it to the tag map. However, the type lacks a member map, which means that it is not yet possible to define variables of the type. It is only when the struct or union becomes properly define, with a member map, this it will be possible to define variables.

else {

type = new Type(sort);

SymbolTable.CurrentTable.AddTag(name, type);

return type;

}

}

#### Enumeration

A enumeration declaration may holds a list of enumeration items, which are assigned values. Each item is declared as a constant signed integer, with a value that is given implicitly or explicitly assigned.

MainParser.gppg

enum\_specifier:

ENUM optional\_name {

EnumValueStack.Push(BigInteger.Zero);

}

LEFT\_BLOCK enum\_list RIGHT\_BLOCK {

EnumValueStack.Pop();

$$ = MiddleCodeGenerator.EnumSpecifier($2, $5);

}

A enumeration declaration may also hold a name without a enumeration item list. In that case, we look up the name and returns its type. The type will actually be constant signed integer. However, the name of the enumeraion must exist,; otherwise, a error is reported.

| ENUM NAME {

$$ = MiddleCodeGenerator.LookupEnum($2);

};

The result of enum\_list is a set holding the symbols of the enumeration items and a boolean value indicating whether the item was explicitly assignd.

enum\_list:

enum {

ISet<Pair<Symbol,bool>> memberSet =

new HashSet<Pair<Symbol,bool>>();

memberSet.Add($1);

$$ = memberSet;

}

| enum\_list COMMA enum {

ISet<Pair<Symbol,bool>> memberSet = $1;

memberSet.Add($3);

$$ = memberSet;

};

A enumeration item might be assigned a value. Otherwise, a values is implicitly assigned.

When it comes to enumeration we have a potential problem. Each enumeration item is stored in the symbol table as a signed integer with a value. The value may be assigned or given implicity. If the storage of the enumeration is extern, we are not allowed to assigne values to the enumeration items. However, it is in C allowed to state the declaration specifiers in arbitary order. For instance, the following declaration is valid

enum {a, b} extern;

The following declaration is invalid:

enum {a = 1, b} extern; // Invalid declaration

The problem is that when we parse enumeration declaration we might not yet know if its storage is extern. Therefore we must also store, for each item, whether it has been explicity assigned. The handling of the storage is maganed by the Specifier class after all declaration specifiers have been parsed. The result is that for each item a pair is returned, holding the symbol of the item and a boolean value indicating whether the item has been explicitly assigned.

enum:

NAME {

Symbol symbol = MiddleCodeGenerator.EnumItem($1, null);

$$ = new Pair<Symbol,bool>(symbol, false);

}

| NAME ASSIGN constant\_integral\_expression {

Symbol symbol = MiddleCodeGenerator.EnumItem($1, $3.Symbol);

$$ = new Pair<Symbol,bool>(symbol, true);

};

The EnumItem method stores the item in the symbol table and checks the optional initialization symbol. The type of the item is constant signed integer.

MiddleCodeGenerator.cs

public static Symbol EnumItem(string itemName, Symbol optInitializerSymbol) {

Type itemType = new Type(Sort.Signed\_Int);

itemType.IsConstant = true;

The value of the item is either set by the initialization symbol value, if present, or by the enumeration value stack. The stack holds the value of the current enumeration list. It is popped and pushed with the item value plus one. In that way will each item hold the value of the previous item plus one, the first item holds the value zero.

BigInteger value;

if (optInitializerSymbol != null) {

Assert.Error(optInitializerSymbol.Type.IsIntegral(), itemName,

Message.Non\_\_integral\_enum\_value);

Assert.Error(optInitializerSymbol.IsValue(), itemName,

Message.Non\_\_constant\_enum\_value);

CCompiler\_Main.Parser.EnumValueStack.Pop();

value = (BigInteger) optInitializerSymbol.Value;

}

else {

value = CCompiler\_Main.Parser.EnumValueStack.Pop();

}

Symbol itemSymbol = new Symbol(itemName, null, itemType, value);

SymbolTable.CurrentTable.AddSymbol(itemSymbol);

CCompiler\_Main.Parser.EnumValueStack.Push(value + 1);

return itemSymbol;

}

The EnumSpecifier method add the enumerator to the current symbol table if optional name is not null.

public static Type EnumSpecifier(string optionalName,

ISet<Pair<Symbol,bool>> enumSet) {

Type enumType = new Type(Sort.Enumeration, enumSet);

if (optionalName != null) {

SymbolTable.CurrentTable.AddTag(optionalName, enumType);

}

return enumType;

}

The LookupEnum method

public static Type LookupEnum(string name) {

Type type = SymbolTable.CurrentTable.LookupTag(name, Sort.Enumeration);

Assert.Error(type != null, name, Message.Tag\_not\_found);

return type;

}

#### Declaration

A declaration is made up by a declaration specifier list, an optional pointer\_declarator list, and a semicolon. For instance, in the following declaration, static long int is the specifier list while \*p, and a[3] are pointer declarators.

static long int \*p, a[3];

There is possible to define a declaration without a pointer\_declarator list. This feature is useful when declare a name struct or union:

struct \_s {int i};

However, this feature makes it also possible to declare unnamed specifier lists. The following declaration are syntactically correct but without meaning, they will be ignored by the compiler.

unsigned short int;

union {short s;}

However, the following declaration has meaning since the enumeration values can be used.

enum {ZERO = 0, ONE = 1, TEN = 10};

declaration:

declaration\_specifier\_list SEMICOLON {

SpecifierStack.Push(Specifier.SpecifierList($1));

$$ = new List<MiddleCode>();

}

| declaration\_specifier\_list\_x declarator\_list SEMICOLON {

SpecifierStack.Pop();

$$ = $2;

};

declaration\_specifier\_list\_x:

declaration\_specifier\_list {

SpecifierStack.Push(Specifier.SpecifierList($1));

};

declaration\_specifier\_list:

declaration\_specifier {

$$ = new List<object>();

$$.Add($1);

}

| declaration\_specifier declaration\_specifier\_list {

$2.Add($1);

$$ = $2;

};

There a set of different kind of specifiers. The following specifiers return a value of the Mask enumeration.

* Storage: auto, register, static, extern, or typedef
* Type quailifiers: constant or volatile
* Type specifiers: void, char, short, int, long, float, double, signed, or unsigned

The following specifiers return a reference to the Type class.

* Struct or union specifier
* Enum specifier
* Typedef name

declaration\_specifier:

CONSTANT { $$ = Mask.Constant; }

| VOLATILE { $$ = Mask.Volatile; }

| AUTO { $$ = Mask.Auto; }

| REGISTER { $$ = Mask.Register; }

| STATIC { $$ = Mask.Static; }

| EXTERN { $$ = Mask.Extern; }

| TYPEDEF { $$ = Mask.Typedef; }

| VOID { $$ = Mask.Void; }

| CHAR { $$ = Mask.Char; }

| SHORT { $$ = Mask.Short; }

| INT { $$ = Mask.Int; }

| LONG { $$ = Mask.Long; }

| FLOAT { $$ = Mask.Float; }

| DOUBLE { $$ = Mask.Double; }

| SIGNED { $$ = Mask.Signed; }

| UNSIGNED { $$ = Mask.Unsigned; }

| struct\_or\_union\_specifier { $$ = $1; }

| enum\_specifier { $$ = $1; }

| TYPEDEF\_NAME { $$ = $1; };

### Declarator

A declarator list is made up by initialization-bitfield declarators separated by commas.

MainParser.gppg

declarator\_list:

initialization\_bitfield\_pointer\_declarator {

$$ = $1;

}

| declarator\_list COMMA initialization\_bitfield\_pointer\_declarator {

$1.AddRange($3);

$$ = $1;

};

A intialization-bitfield declarator is be a pointer declarator, a declarator initialized with a value, or marked as a bitfield. However, a declarator cannot be both initialized and markd as bitfield. In the bitfield case the declarator can be omitted. For instance:

int a;

int b = 1;

int c : 8;

int : 8;

In each case we call the matchning method.

MainParser.gppg

initialization\_bitfield\_pointer\_declarator:

pointer\_declarator {

MiddleCodeGenerator.Declarator(SpecifierStack.Peek(), $1);

$$ = new List<MiddleCode>();

}

| pointer\_declarator ASSIGN initializer {

$$ = MiddleCodeGenerator.AssignmentDeclarator

(SpecifierStack.Peek(), $1, $3);

}

| optional\_pointer\_declarator COLON constant\_integral\_expression {

MiddleCodeGenerator.BitfieldDeclarator

(SpecifierStack.Peek(), $1, $3.Symbol);

$$ = new List<MiddleCode>();

};

#### Pointer Declarator

optional\_pointer\_declarator:

/\* Empty \*/ { $$ = null; }

| pointer\_declarator { $$ = $1; };

A pointer declarator is made up by an optional list of pointer declarators and a direct pointer declarator.

pointer\_declarator:

optional\_pointer\_list direct\_pointer\_declarator {

$$ = MiddleCodeGenerator.PointerDeclarator($1, $2);

};

optional\_pointer\_list:

/\* Empty \*/ { $$ = new List<CCompiler.Type>(); }

| pointer\_list { $$ = $1; };

pointer\_list:

pointer {

$$ = new List<CCompiler.Type>();

$$.Add($1);

}

| pointer\_list pointer {

$1.Add($2);

$$ = $1;

};

pointer:

ASTERRISK optional\_qualifier\_list {

$$ = Specifier.QualifierList($2);

};

optional\_qualifier\_list:

/\* Empty \*/ {

$$ = new List<Mask>();

}

| optional\_qualifier\_list qualifier {

$$ = $1;

$$.Add($2);

};

qualifier:

CONSTANT { $$ = Mask.Constant; }

VOLATILE { $$ = Mask.Volatile; };

#### Direct Declarator

A direct declarator is a name, a declarator inside a parenthesis pair, another direct declarator followed by a parameter-ellipse list inside a parenthesis pair.

A direct declarator in its simplest form is just name. For instance,

direct\_declarator:

NAME {

$$ = new Declarator($1);

}

Another form of direct declarator is another declarator. In that case we simply return the declarator. The parentheses is only present to change the precidence of the declarator. For instance:

| LEFT\_PARENTHESIS declarator RIGHT\_PARENTHESIS {

$$ = $2;

}

A direct declarator can be an array declarator; that is,

| direct\_declarator LEFT\_SQUARE

optional\_constant\_integral\_expression RIGHT\_SQUARE {

$$ = MiddleCodeGenerator.ArrayType($1, $3);

}

| direct\_declarator

LEFT\_PARENTHESIS parameter\_ellipse\_list RIGHT\_PARENTHESIS {

$$ = MiddleCodeGenerator.

NewFunctionDeclaration($1, $3.First, $3.Second);

}

| direct\_declarator LEFT\_PARENTHESIS

optional\_name\_list RIGHT\_PARENTHESIS {

$$ = MiddleCodeGenerator.OldFunctionDeclaration($1, $3);

};

optional\_parameter\_ellipse\_list:

/\* Empty \*/ { $$ = null; }

| parameter\_ellipse\_list { $$ = $1; };

parameter\_ellipse\_list:

parameter\_list {

$$ = new Pair<List<Pair<string,Symbol>>,Boolean>($1, false);

}

| parameter\_list COMMA ELLIPSE {

$$ = new Pair<List<Pair<string,Symbol>>,Boolean>($1, true);

};

parameter\_list:

{ ++CallDepth; }

parameter\_declaration {

--CallDepth;

$$ = new List<Pair<string,Symbol>>();

$$.Add($2);

}

| parameter\_list COMMA {

++CallDepth;

}

parameter\_declaration {

--CallDepth;

$1.Add($4);

$$ = $1;

};

parameter\_declaration:

declaration\_specifier\_list {

$$ = MiddleCodeGenerator.Parameter(Specifier.SpecifierList($1), null);

}

| declaration\_specifier\_list\_x declarator {

$$ = MiddleCodeGenerator.Parameter(SpecifierStack.Pop(), $2);

}

| declaration\_specifier\_list\_x abstract\_declarator {

$$ = MiddleCodeGenerator.Parameter(SpecifierStack.Pop(), $2);

};

optional\_name\_list:

/\* Empty \*/ { $$ = new List<string>(); }

| name\_list { $$ = $1; };

name\_list:

NAME {

$$ = new List<string>();

$$.Add($1);

}

| name\_list COMMA NAME {

$1.Add($3);

$$ = $1;

};

#### Initialization

The initializer can be either an expression or a block. The assignment expression makes sure that the expression cannot hold assignment or comma.

initializer:

assignment\_expression {

$$ = $1;

}

The block holds a list of initializer, which means that initializeralizer lists can be nested. But in the end there is always expressions in the initialization lists.

| LEFT\_BLOCK initializer\_list optional\_comma RIGHT\_BLOCK {

$$ = $2;

};

optional\_comma:

/\* Empty \*/

| COMMA;

initializer\_list:

initializer {

$$ = new List<object>();

$$.Add($1);

}

| initializer\_list COMMA initializer {

$1.Add($3);

$$ = $1;

};

#### Abstract Declarator

abstract\_declarator:

pointer\_list {

$$ = MiddleCodeGenerator.PointerDeclarator($1, null);

}

| optional\_pointer\_list direct\_abstract\_declarator {

$$ = MiddleCodeGenerator.PointerDeclarator($1, $2);

};

direct\_abstract\_declarator:

LEFT\_PARENTHESIS abstract\_declarator RIGHT\_PARENTHESIS {

$$ = $2;

}

| LEFT\_SQUARE optional\_constant\_integral\_expression RIGHT\_SQUARE {

$$ = MiddleCodeGenerator.ArrayType(null, $2);

}

| direct\_abstract\_declarator

LEFT\_SQUARE optional\_constant\_integral\_expression RIGHT\_SQUARE {

$$ = MiddleCodeGenerator.ArrayType($1, $3);

}

| LEFT\_PARENTHESIS optional\_parameter\_ellipse\_list RIGHT\_PARENTHESIS {

$$ = MiddleCodeGenerator.

NewFunctionDeclaration(null, $2.First, $2.Second);

}

| direct\_abstract\_declarator

LEFT\_PARENTHESIS optional\_parameter\_ellipse\_list RIGHT\_PARENTHESIS {

$$ = MiddleCodeGenerator.

NewFunctionDeclaration($1, $3.First, $3.Second);

};

## Statements

The generateFunctionCode is called after the middle code of the function has been generated. Its task is to genereate the assembler and object code of the function. If the function returns void and is the main function, the Exit middle code is added to middle code list. If the it returns void and is not the main, the Return middle code instruction is added.

Then the middle code is being optimized and a list of base blocks is generated.

The object code is actually generated twice, since we want to know which symbols are not used and can be removed from the symbol table.

If the function is the main function we have to modify the entry point of the function.

Finally, the function is added to the linker map and the symbol table is popped.

The generateAssemblerText method generates the assembler code for the function.

When a function containing goto statements has been parsed, the addresses need to be set with the help of the label map.

### The Nested-If Problem

The if-else problem[[2]](#footnote-2) is the problem of syntactically interpret the leftmost source code below. Semantically, the middle interpretation of the left statement is the correct one, each else shall be connected to the latest preceding if.

|  |  |  |
| --- | --- | --- |
| if (a < b)  if (c < d)  e = 1;  else  f = 2; | if (a < b) {  if (c < d)  e = 1;  else  f = 2;  } | if (a < b)  if (c < d) {  e = 1;  }  else  f = 2; |

Below is a simple set of statement rules. Unfortunately, they are ambiguous in that way that the an else does not have to be connected to the latest preceding if, resulting in both the middle and rightmost semantically interpretation above, depending in which order the rules are applied.

statement ::=

IF LEFT\_PAREN logical\_expression RIGHT\_PAREN statement

| IF LEFT\_PAREN logical\_expression RIGHT\_PAREN statement ELSE statement

| ...

To solve the problem, we need a more complicated set of rules that works with open and closed statements. The following set is unambiguous in that way that it always connects each else with the latest preceding if, with compiles to the semantics of C. A theoretical explanation of the open-closed statement solution is beyond the scope of this book, but I recommend the Dragon Book by Aho et al. for a closer look.

When it comes to the if-else statement we have a problem. Given the following rules:

statement:

**if** ( expression ) statement

| **if** ( expression ) statement **else** statement

| *all other statements*

Let us look at the following two examples:

if (a < b)

if (c < d)

a = 1;

else

b = 1;

if (a < b)

if (c < d)

a = 1;

else

b = 1;

The question is whether the else part matches the first of second if part. The first example suggests that it matches the first part, while the second example suggests that it matches the second part. The first suggestion is the right one, each else shall be matched to the closest if. However, the parser above supports both suggestions. We must rewrite the parser so that it only accepts the first case, which we do by introducing the opened and closed statements:

statement:

opened\_statement

| closed\_statement

opened\_statement:

if ( expression ) statement

| if ( expression ) closed\_statement else opened\_statement

| switch ( expression ) opened\_statement

| case constant\_integral\_expression : opened\_statement

| while ( expression ) opened\_statement

| for ( optional\_expression ; optional\_expression ; optional\_expression )

opened\_statement {

| name : opened\_statement

closed\_statement:

if ( expression ) closed\_statement else closed\_statement

| switch ( expression ) closed\_statement

| while ( expression ) closed\_statement

| for ( optional\_expression ; optional\_expression ; optional\_expression )

closed\_statement

| do statement while ( expression ) ;

| case constant\_integral\_expression : closed\_statement

| default : closed\_statement

| continue ;

| break ;

| { optional\_statement\_list }

| goto name ;

| return optional\_expression ;

| optional\_expression ;

| declaration ;

| jump\_register ( register\_name ) ;

| interrupt ( constant\_integral\_expression ) ;

| syscall ( ) ;

A theoretical explanation of the open-closed statement solution is beyond the scope of this book, but I recommend the Dragon Book by Aho et al. for a closer look.

So, let us return to the if-else statement.

MainParser.cs

statement:

opened\_statement { $$ = $1; }

| closed\_statement { $$ = $1; };

opened\_statement:

IF LEFT\_PARENTHESIS expression RIGHT\_PARENTHESIS statement {

$$ = MiddleCodeGenerator.IfStatement($3, $5);

}

| IF LEFT\_PARENTHESIS expression RIGHT\_PARENTHESIS closed\_statement

ELSE opened\_statement {

$$ = MiddleCodeGenerator.IfElseStatement($3, $5, $7);

}

| // ...

closed\_statement:

IF LEFT\_PARENTHESIS expression RIGHT\_PARENTHESIS closed\_statement

ELSE closed\_statement {

$$ = MiddleCodeGenerator.IfElseStatement($3, $5, $7);

}

| // ...

First, let us look at the Statement class. A statement holds a list of middle code instructions, and the next set, which is a set om jump instructions that shall be backpatched to jump to the next instruction after the statement.

Statement.cs

using System.Collections.Generic;

namespace CCompiler {

public class Statement {

private List<MiddleCode> m\_list;

private ISet<MiddleCode> m\_nextSet;

public Statement(List<MiddleCode> list, ISet<MiddleCode> nextSet) {

Assert.ErrorXXX(list != null);

m\_list = list;

m\_nextSet = (nextSet != null) ? nextSet : (new HashSet<MiddleCode>());

}

public List<MiddleCode> CodeList {

get { return m\_list; }

}

public ISet<MiddleCode> NextSet {

get { return m\_nextSet; }

}

}

}

Then we look at the Expression class. It handles a symbol, a register, and a short and a long list. The register parameter is used in system calls when we need to access or assign specific register. The difference between the short and long list is that the short list holds only the side effects of the expression. For instance, in the following expression the long list holds the whole expression: the function call, decrement, and addition. The short list holds only the side effects: the function call and decrement, but not the addition.

f(x) + a--;

The following statement is correct, but in holds no side-effects, its short list will be empty, and it will eventually be ignored by the compiler.

a + (b \* c);

Expression.cs

using System.Numerics;

using System.Collections.Generic;

namespace CCompiler {

public class Expression {

private Symbol m\_symbol;

private List<MiddleCode> m\_shortList;

private List<MiddleCode> m\_longList;

private Register? m\_register;

public Expression(Symbol symbol, List<MiddleCode> shortList,

List<MiddleCode> longList, Register? register = null) {

m\_symbol = symbol;

m\_shortList = (shortList != null) ? shortList : (new List<MiddleCode>());

m\_longList = (longList != null) ? longList : (new List<MiddleCode>());

m\_register = register;

}

public Symbol Symbol {

get { return m\_symbol; }

}

public List<MiddleCode> ShortList {

get { return m\_shortList; }

}

public List<MiddleCode> LongList {

get { return m\_longList; }

}

public Register? Register {

get { return m\_register; }

}

public override string ToString() {

return m\_symbol.ToString();

}

}

}

The if statement backpatches the true set to the beginning of the statement and the false set to the middle code instruction following the statement. The if-else statement backpatches the true set to the beginning of the true statement and the false set to the beginning of the false statement.

The IfStatement method handles the if statement of the parser. Since the if-statement accepts logical expressions only, we need to start with making sure the expression is or becomes a logical expression.

MiddleCodeGenerator.cs

public static Statement IfStatement(Expression expression,

Statement innerStatement){

expression = TypeCast.ToLogical(expression);

Since we use the value of the expression in the if-statement, we only use the long list.

List<MiddleCode> codeList = expression.LongList;

AddMiddleCode(codeList, MiddleOperator.CheckTrackMapFloatStack);

We backpatch the true value of the expression to the first instruction in the middle code list of the inner statement. This means that when the expression is true, the program will jump to the beginning of the code list of the inner statement.

Backpatch(expression.Symbol.TrueSet(), innerStatement.CodeList);

codeList.AddRange(innerStatement.CodeList);

We add a jump instruction at the end of the middle code that shall jump to the instruction following this if-statement.

MiddleCode nextCode = AddMiddleCode(codeList, MiddleOperator.Goto);

We define the next set of the if-statement, it holds the instructions that jump to the instruction after the if- statement. It is the union of the next set of the inner statement, the false set of the expression, and the jump instruction just added to the list.

ISet<MiddleCode> nextSet = new HashSet<MiddleCode>();

nextSet.UnionWith(innerStatement.NextSet);

nextSet.UnionWith(expression.Symbol.FalseSet());

nextSet.Add(nextCode);

Finally, we return on object of the Statement class holding the code list and next set of the if-statement.

return (new Statement(codeList, nextSet));

}

The IfElseStatement methods is like the IfStatement method above. The difference is that we have two inter statements. On statement that the program jumps to in case the expression is evaluated to a true value, and one statement where it jumps to in case of a false value.

public static Statement IfElseStatement(Expression expression,

Statement trueStatement,

Statement falseStatement) {

expression = TypeCast.ToLogical(expression);

List<MiddleCode> codeList = expression.LongList;

AddMiddleCode(codeList, MiddleOperator.CheckTrackMapFloatStack);

We backpatch the true and false set to the beginning of the code list of the true and false statement, respectively.

Backpatch(expression.Symbol.TrueSet(), trueStatement.CodeList);

Backpatch(expression.Symbol.FalseSet(), falseStatement.CodeList);

codeList.AddRange(trueStatement.CodeList);

We add a jump instruction after the end of the true statement code list for it to jump out of the if-else-statement.

MiddleCode gotoCode = AddMiddleCode(codeList, MiddleOperator.Goto);

codeList.AddRange(falseStatement.CodeList);

The next set of the if-else-statement is the union of the next sets of the true and false statement, and the jump instruction just added to the code.

ISet<MiddleCode> nextSet = new HashSet<MiddleCode>();

nextSet.UnionWith(trueStatement.NextSet);

nextSet.UnionWith(falseStatement.NextSet);

nextSet.Add(gotoCode);

Finally, we return on object of the Statement class holding the code list and next set of the if-statement.

return (new Statement(codeList, nextSet));

}

### The Forward-Jump Problem (Backpatching)

When generating middle code instructions for expressions or statements, a common situation is that we generate forward jump instructions without yet knowing the target address. One way to solve the problem is to work with backpatching, which means that we use sets to keep track of jump instructions with yet unknown target addresses that eventually becomes filled in when the target addresses have become known.

For instance, let us look at the while statement to the left. The parsing of expression a < b will generate the middle code in the middle and the sets to the right. The middle code line will a have undefined targets since we at the moment does not know where to jump if the expression is true or false.

|  |  |  |
| --- | --- | --- |
| while (a < b)  ++a; | 1. if a < b goto ?  2. goto ? | a < b: true set = {1}  false set = {2} |

When parsing the while statement, the sets of the expression becomes backpatched: If the a < b expression is true, line 1 shall jump to line 3, and if it is false, line 2 shall jump to line 5. Whatever comes after the while statement.

|  |  |  |
| --- | --- | --- |
| 1. if a < b goto 3  2. goto 5  3. ++a  4. goto 1  5. ... |  |  |

If we instead let the while statement be surrounded by an if-else statement, we have both the true and false sets of the if expression x < y and the sets of the while expression a < b:

|  |  |  |
| --- | --- | --- |
| if (x < y)  while (a < b)  ++a;  else  ++b; | 1. if x < y goto ?  2. goto ?  3. if a < b goto ?  4. goto ?  5. ++a  6. goto 3  7. goto 9  8. ++b  9. ... | x < y: true set: {1}  false set: {2}  a < b: true set: {3}  false set: {4} |

When the if and while statements become parsed the following call will occur:

backpatch({1}, 3);

backpatch({2}, 8);

backpatch({3}, 5);

backpatch({4}, 7);

|  |  |  |
| --- | --- | --- |
| 1. if x < y goto 3  2. goto 8  3. if a < b goto 5  4. goto 7  5. ++a  6. goto 3  7. goto 9  8. ++b  9. ... |  |  |

The generated middle code above may appear ineffective, but the middle code optimizer of Chapter 14 will change it into:

|  |  |  |
| --- | --- | --- |
| 1. if x >= y goto 5  2. if a >= b goto 6  3. ++a  4. goto 2  5. ++b  6. ... |  |  |

Another example is the evaluation of a logical-and expression. Since C takes advantage of lazy evaluation, we have to insert jump instructions in the middle code.

|  |  |  |
| --- | --- | --- |
| if ((a < b) && (c < d))  ++a;  else  ++b; | 1. if a < b goto ?  2. goto ?  3. if c < d goto ?  4. goto ?  5. ++a  6. goto 7  7. ++b  8. ... |  |

When the first expression (a < b) is evaluated, its true set is {1} and its false set is {2}. When the second expression is evaluated the true set of the first expression is backpatched to the beginning of the second expression. This means that if the first expression is true we continue to evaluate the second expression. However, if the first expression is false the second expression will never be evaluated since it not necessary, the total expression will be false regardless of the value of the second expression. The result will be that the true set of the total expression is the one of the second expression and the false set is the union of the falses sets of both expressions.

|  |  |  |
| --- | --- | --- |
| 1. if a < b goto 3  2. goto ?  3. if c < d goto ?  4. goto ?  5. ++a  6. goto 7  7. ++b | true set: {3}  false set: {2, 4} |  |

In a similar way, the true set of an expression with logical-or will be the union of both true sets while the false set will the false set of the second expression.

|  |  |  |
| --- | --- | --- |
| if ((a < b) || (c < d))  ++a;  else  ++b; | 1. if a < b goto ?  2. goto 3  3. if c < d goto ?  4. goto ?  5. ++a  6. goto 7  7. ++b  8. ... | true set: {1, 3}  false set: {4} |

The Backpatch class holds a set of static methods that backpatch a single address or a set of addresses. The first two methods call the last two methods with the current size of ther middle code list as target address. In this way, the target address becomes the next middle code instruction to become generated. More specifically, the backpatching is performed by setting the third address of the middle code, since the third address is always the jump address in conditional and unconditional jump instructions.

### Header Rules

The open statements include the compound statements if, if-else, switch, case, while, for, and label statements. The switch, while, and for statements need to do some preparation before the parsing, which means that they need one header rule each. Each rule calls its corresponding method in the Generate class that does the actual work of type checking and middle code generation.

### The Switch Statement

The generateSwitchStatement method generates middle code jump instructions for each of the case statements, which are stored in the Main.CaseMapStack stack. Each entry in the stack holds an integer value (each case expression value must be possible to evaluate in compile-time) and a line number. If the original switch expression equals the constant case expression, we jump to the start line of the statement following the case expression.

If there is a default statements, we finally jump to its start line if the switch statement does not match any of the case statements. Note that the default statement does not have to be placed after the case statements, even though I can see no good reason not to do so.

MainParser.gppg

switch\_header:

/\* Empty. \*/ { MiddleCodeGenerator.SwitchHeader(); };

opened\_statement:

// ...

| SWITCH switch\_header LEFT\_PARENTHESIS expression RIGHT\_PARENTHESIS

opened\_statement {

$$ = MiddleCodeGenerator.SwitchStatement($4, $6);

}

| CASE constant\_integral\_expression COLON opened\_statement {

$$ = MiddleCodeGenerator.CaseStatement($2, $4);

}

// ...

closed\_statement:

// ...

| SWITCH switch\_header LEFT\_PARENTHESIS expression RIGHT\_PARENTHESIS

closed\_statement {

$$ = MiddleCodeGenerator.SwitchStatement($4, $6);

}

| CASE constant\_integral\_expression COLON closed\_statement {

$$ = MiddleCodeGenerator.CaseStatement($2, $4);

}

| DEFAULT COLON closed\_statement {

$$ = MiddleCodeGenerator.DefaultStatement($3);

}

// ...

We need a map to keep track of the case statements. Since the switch statements can be nested, we need a stack to keep track of the case maps (m\_caseMapStack). In the same way, we need a stack keep track of the default statements (m\_caseMapStack), and a stack of sets to keep track of the break statements (m\_breakStack).

MiddleCodeGenerator.cs

private static Stack<IDictionary<BigInteger, MiddleCode>> m\_caseMapStack =

new Stack<IDictionary<BigInteger, MiddleCode>>();

private static Stack<MiddleCode> m\_defaultStack = new Stack<MiddleCode>();

private static Stack<ISet<MiddleCode>> m\_breakSetStack =

new Stack<ISet<MiddleCode>>();

The SwitchHeader method is called before the parsing of the switch statement. Its task is to push a new map to the case map stack, a null reference to the default stack, and a new set to the break set stack.

public static void SwitchHeader() {

m\_caseMapStack.Push(new Dictionary<BigInteger,MiddleCode>());

DefaultStack.Push(null);

BreakSetStack.Push(new HashSet<MiddleCode>());

}

The SwitchStatement method is called after the parsing of the switch statement. Its task is to generate middle code for the case and default statements, and to backpatch the break statements.

public static Statement SwitchStatement(Expression switchExpression,

Statement innerStatement) {

Since the switch statement need an integral value rather than a logical value we need to, contrary to the if- statement above, type cast a potential logical value to an integral value.

switchExpression = TypeCast.LogicalToIntegral(switchExpression);

Since we need the value of the switch statement, we are only interested of the long list of the switch expression.

List<MiddleCode> codeList = switchExpression.LongList;

Each case value is type casted to the type of the switch expression.

Type switchType = switchExpression.Symbol.Type;

foreach (KeyValuePair<BigInteger,MiddleCode> entry

in m\_caseMapStack.Pop()) {

BigInteger caseValue = entry.Key;

MiddleCode caseTarget = entry.Value;

Symbol caseSymbol = new Symbol(switchType, caseValue);

We add code where we compare the case value with the switch value and jump to the matching middle code instruction if the values are equal.

AddMiddleCode(codeList, MiddleOperator.Case, caseTarget,

switchExpression.Symbol, caseSymbol);

}

After the case values, we add a case end instruction. The reason for this is that the register used for the case comparison in the generated assembly code shall be unallocated.

AddMiddleCode(codeList, MiddleOperator.CaseEnd,

switchExpression.Symbol);

Similar to the if- statement, the switch- statement has a next set; that is, a set of middle code instruction that shall jump to the instruction following the switch statement.

ISet<MiddleCode> nextSet = new HashSet<MiddleCode>();

If there is a default statement, we jump to it if the switch statement does not match any of the case statement. In that case, the next set becomes empty.

MiddleCode defaultCode = DefaultStack.Pop();

if (defaultCode != null) {

AddMiddleCode(codeList, MiddleOperator.Goto, defaultCode);

}

If there is no default statement, we instead add a jump instruction to the next set.

else {

nextSet.Add(AddMiddleCode(codeList, MiddleOperator.Goto));

}

We finally add the code list of the inner statement to the resulting code list.

codeList.AddRange(innerStatement.CodeList);

The next set is the union of the next set of the inner statement and the break statements of the switch statement.

nextSet.UnionWith(innerStatement.NextSet);

nextSet.UnionWith(BreakSetStack.Pop());

return (new Statement(codeList, nextSet));

}

### The Case Statement

MainParser.gppg

| CASE constant\_integral\_expression COLON closed\_statement {

$$ = MiddleCodeGenerator.CaseStatement($2, $4);

}

The case statement must comply with the following demands:

* Main.CaseMapStack must not be empty. If it is empty, the case statements misses a surrounding switch statements.
* The case expression must be integral and constant and thereby possible to evaluate by the compiler.
* Since Main.CaseMapStack is not empty, there is at least one surrounding switch statement. We call peek on the stack to obtain the map holding the case expression value of the closest surrounding switch statement, add the case value with the line number of the case statement, and check that there is no earlier cases expression with the same value (put returns the old value if the key is already present in the map, null if is not[[3]](#footnote-3)).

If the case map stack is empty, we have a case statement without a enclosing switch statement and an error is reported.

public static Statement CaseStatement(Expression expression,

Statement statement) {

Assert.Error(m\_caseMapStack.Count > 0, Message.Case\_without\_switch);

expression = TypeCast.LogicalToIntegral(expression);

If the value of the symbol is null, we have a non-constant case value and an error is reported.

Assert.Error(expression.Symbol.Value != null, expression.Symbol.Name,

Message.Non\_\_constant\_case\_value);

BigInteger caseValue = (BigInteger) expression.Symbol.Value;

We extract the current case map from the top of the case map stack and adds the current case value to the map.

IDictionary<BigInteger, MiddleCode> caseMap = m\_caseMapStack.Peek();

I the map already contains the value we have two case values with the same value and an error is reported.

Assert.Error(!caseMap.ContainsKey(caseValue), caseValue,

Message.Repeated\_case\_value);

Finally, we add the case value and the first instruction of the statement’s code

caseMap.Add(caseValue, GetFirst(statement.CodeList));

return statement;

}

The GetFirst method returns the first instruction in the middle code instruction list. It also adds an empty instruction in case the list is empty. This is done for the backpatching to always have an instruction to set as the target.

private static MiddleCode GetFirst(List<MiddleCode> list) {

if (list.Count == 0) {

AddMiddleCode(list, MiddleOperator.Empty);

}

return list[0];

}

### The Default Statement

The default statement have to comply with the following demands:

* Main.m\_defaultStack must not be empty. If it is empty, the default statements misses a surrounding switch statements.
* Since Main.m\_defaultStack is not empty, there is at least one surrounding switch statement. We call pop to see if the top value is minus one. If it is not, there has already been a default statement in the closest surrounding switch statement. If it is minus one, there has not been a earlier default statement and we push the line number at the beginning of the statement following the default statement. Since the line numbers are numbered from zero, it cannot be minus one.

MainParser.gppg

| DEFAULT COLON closed\_statement {

$$ = MiddleCodeGenerator.DefaultStatement($3);

}

Technically, the default statement does not have to be placed after the last case statements. However, we cannot have more than one default statement in a switch statement.

public static Statement DefaultStatement(Statement statement) {

If the default stack is empty, we have a default statement without an enclosing switch statement and an error is reported.

Assert.Error(DefaultStack.Count > 0, Message.Default\_without\_switch);

If the value on the top of the stack is not null, we have a switch statement with two default statement and an error is reported.

Assert.Error(DefaultStack.Pop() == null, Message.Repeted\_default);

Finally, we add the default middle code instruction at the top of the stack.

DefaultStack.Push(GetFirst(statement.CodeList));

return statement;

}

### The While Statement

The while header prepares the while statement by pushing break and continue stacks with integers sets. The continue set will eventually be backpatched to the beginning of the while statement and break set will be backpatched to the statement following the while statement.

The true set of the while expression is backpatched to the beginning of the statement (or the first statement of a block statement) inside the while loop while the false set is backpatched to the statement following the while statement, similar to the break set. In order to make sure that the while statement is followed by jump statement, the jump\_marker rule adds a jump statement at the end of the statement surrounded by the while statement, which is backpatched to the beginning of the while expression, similar to the continue set.

MainParser.gppg

loop\_header:

/\* Empty. \*/ { MiddleCodeGenerator.LoopHeader(); };

opened\_statement:

// ...

| WHILE loop\_header LEFT\_PARENTHESIS expression RIGHT\_PARENTHESIS

opened\_statement {

$$ = MiddleCodeGenerator.WhileStatement($4, $6);

}

// ...

closed\_statement:

// ...

| WHILE loop\_header LEFT\_PARENTHESIS expression RIGHT\_PARENTHESIS

closed\_statement {

$$ = MiddleCodeGenerator.WhileStatement($4, $6);

}

// ...

Similar to the break set stack above, we also need a continue set satck.

MiddleCodeGenerator.cs

private static Stack<ISet<MiddleCode>> m\_continueSetStack =

new Stack<ISet<MiddleCode>>();

The LoopHeader method is called before the parsing of a while, do, or for statement. Its task is to add empty sets at the top of the break set stack and continue set stack.

public static void LoopHeader() {

m\_breakSetStack.Push(new HashSet<MiddleCode>());

m\_continueSetStack.Push(new HashSet<MiddleCode>());

}

The WhileStatement method is called after the while statement has been parsed. Similar to the if statement above, it starts by type casting the expression into logical type.

public static Statement WhileStatement(Expression expression,

Statement innerStatement) {

expression = TypeCast.ToLogical(expression);

List<MiddleCode> codeList = expression.LongList;

AddMiddleCode(codeList, MiddleOperator.CheckTrackMapFloatStack);

We backpatchar the true set of the expression to the code list of the inner statement; that is, if the expression is true the program shall jump to the beginning of the inner statement.

Backpatch(expression.Symbol.TrueSet, innerStatement.CodeList);

codeList.AddRange(innerStatement.CodeList);

We define the next set of the while statement, and a jump instruction.

ISet<MiddleCode> nextSet = new HashSet<MiddleCode>();

nextSet.Add(AddMiddleCode(codeList, MiddleOperator.Goto,

GetFirst(codeList)));

We also add the false set of the expression and the break set of the inner statement to the next set.

nextSet.UnionWith(expression.Symbol.FalseSet);

nextSet.UnionWith(m\_breakSetStack.Pop());

We backpatch the next set of the inner statement and its continue set to the code list; that is, the jump instruction of the sets shall jump back to the beginning of the while statement.

Backpatch(innerStatement.NextSet, codeList);

Backpatch(m\_continueSetStack.Pop(), codeList);

return (new Statement(codeList, nextSet));

}

### The Do Statement

The do statement is a weaker version of the while statement. The difference is that the while expression may be false from the beginning resulting in zero iterations while the do expression is located at the end, resulting in at least one iteration. However, the do statement may in some cases present a simpler and more intuitive solution.

Like the while and for cases, the do statement is prepared by pushing the continue and break stacks. The true set of the do expression is backpatched to the beginning of the statement surrounded by the do statement, the false set is backpatched to the statement followed the do statement and finally the continue and break stacks are popped.

MainParser.cs

closed\_statement:

// ...

| DO loop\_header statement WHILE

LEFT\_PARENTHESIS expression RIGHT\_PARENTHESIS SEMICOLON {

$$ = MiddleCodeGenerator.DoStatement($3, $6);

}

// ...

MiddleCodeGenerator.cs

public static Statement DoStatement(Statement innerStatement,

Expression expression) {

List<MiddleCode> codeList = innerStatement.CodeList;

Backpatch(innerStatement.NextSet, codeList);

AddMiddleCode(codeList, MiddleOperator.CheckTrackMapFloatStack);

codeList.AddRange(expression.LongList);

Backpatch(expression.Symbol.TrueSet, codeList);

Backpatch(m\_continueSetStack.Pop(), codeList);

ISet<MiddleCode> nextSet = new HashSet<MiddleCode>();

nextSet.UnionWith(expression.Symbol.FalseSet);

nextSet.UnionWith(m\_breakSetStack.Pop());

AddMiddleCode(codeList, MiddleOperator.Goto, GetFirst(codeList));

return (new Statement(codeList, nextSet));

}

### The For Statement

The for statement holds three optional expression: the initialization expression, the test expression, and the increment expression. The true and false sets of the initialization and increment expressions are all backpatched to beginning of the test expression. Like the while case, in order to make sure that the for statement is followed by jump statement the jump\_marker rule add a jump statement at the end of the statement surrounded by the for statement, which is also backpatched to the beginning of the test expression. There is also a jump line inserted after the test optional expression, which is backpatched to the beginning of the for statements to make sure that the for loop works properly even if the test expression has been omitted. An omitted test expression is equivalent to an infinitive loop.

Finally, like the while case, the continue set is backpatched to the beginning of the test expression and the break set is backpatched to the beginning of the statement following the while statement.

The for statement is more complicated than the do and while statements. It holds three optional expression: the initialization expression, the test expression, and the increment expression. The true and false sets of the initialization and increment expressions are all backpatched to beginning of the test expression. An omitted test expression is equivalent to an infinitializere loop.

opened\_statement:

// ...

| FOR loop\_header LEFT\_PARENTHESIS optional\_expression SEMICOLON

optional\_expression SEMICOLON optional\_expression RIGHT\_PARENTHESIS

opened\_statement {

$$ = MiddleCodeGenerator.ForStatement($4, $6, $8, $10);

}

// ...

closed\_statement:

// ...

| FOR loop\_header LEFT\_PARENTHESIS optional\_expression SEMICOLON

optional\_expression SEMICOLON optional\_expression RIGHT\_PARENTHESIS

closed\_statement {

$$ = MiddleCodeGenerator.ForStatement($4, $6, $8, $10);

}

// ...

In the ForStatement method, we make difference of the short and long list of the expressions. If the test expression is present, we want its value, and therefore we re interested in its long list. On the other hand, if the initializer or next expressions are present we are only interested in their potential side effekts.

MiddleCodeGenerator.cs

public static Statement ForStatement(Expression initializerExpression,

Expression testExpression, Expression nextExpression,

Statement innerStatement) {

List<MiddleCode> codeList = new List<MiddleCode>();

ISet<MiddleCode> nextSet = new HashSet<MiddleCode>();

We add an empty instruction as the target of the test expression, since we do not know if the test expression is not null.

MiddleCode testTarget = AddMiddleCode(codeList, MiddleOperator.Empty);

If the initializer expression is not null, we add its short list. Note that we the short list instead of the long list since we are only interested of the side effects of the expression, not its value. We also backpatch both its true set and false set to the test target; that is, the beginning of the test expression code.

if (initializerExpression != null) {

codeList.AddRange(initializerExpression.ShortList);

Backpatch(initializerExpression.Symbol.TrueSet, testTarget);

Backpatch(initializerExpression.Symbol.FalseSet, testTarget);

}

If the test expression is not null, we start by type casting it into a logical value and add its long list to the code. Not that we add the long list rather than the short list in this case, since we need the value of the expression, rather than just its side effects.

if (testExpression != null) {

testExpression = TypeCast.ToLogical(testExpression);

codeList.AddRange(testExpression.LongList);

AddMiddleCode(codeList, MiddleOperator.CheckTrackMapFloatStack);

We backpatch the true set of the test expression to the beginning or the code list of the inner statement.

Backpatch(testExpression.Symbol.TrueSet, innerStatement.CodeList);

nextSet.UnionWith(testExpression.Symbol.FalseSet);

}

We add the code list of the inner statement. We then add the next target instruction and backpatch the next set of the code of the inner statement since we do not know if the next expression is not null.

codeList.AddRange(innerStatement.CodeList);

MiddleCode nextTarget = AddMiddleCode(codeList, MiddleOperator.Empty);

Backpatch(innerStatement.NextSet, nextTarget);

If the next expression is not null, we add its short list and backpatch its true set and false set to the

if (nextExpression != null) {

codeList.AddRange(nextExpression.ShortList);

Backpatch(nextExpression.Symbol.TrueSet, testTarget);

Backpatch(nextExpression.Symbol.FalseSet, testTarget);

}

Finally, like the while case, the continue set is backpatched to the beginning of the test expression and the break set is backpatched to the beginning of the statement following the while statement.

AddMiddleCode(codeList, MiddleOperator.Goto, testTarget);

Backpatch(m\_continueSetStack.Pop(), nextTarget);

nextSet.UnionWith(m\_breakSetStack.Pop());

return (new Statement(codeList, nextSet));

}

### Label and Goto Statement

The label statement is quite simple, we just add the name of the label together with the line number of the beginning of the statement following the label to Main.LabelMap and check that the label has not been added already. The labels are used as targets of the goto statements.[[4]](#footnote-4)

MainParser.cs

opened\_statement:

NAME COLON opened\_statement {

$$ = MiddleCodeGenerator.LabelStatement($1, $3);

};

closed\_statement:

GOTO NAME SEMICOLON {

$$ = MiddleCodeGenerator.GotoStatement($2);

}

To keep track of the label statements we have the label map m\_labelMap with the label name as key and the first middle code instruction of the following statement as value.

MiddleCodeGenerator.cs

public static IDictionary<string, MiddleCode> m\_labelMap =

new Dictionary<string, MiddleCode>();

public static Statement LabelStatement(string labelName,

Statement statement) {

If the label map already hold the label name as a key, we have two labels with the same name and an error is reported.

Assert.Error(!m\_labelMap.ContainsKey(labelName),

labelName, Message.Defined\_twice);

m\_labelMap.Add(labelName, GetFirst(statement.CodeList));

return statement;

}

We also have the goto set map m\_gotoSetMap with the label name as key and the goto middle code instruction as value.

public static IDictionary<string, ISet<MiddleCode>> m\_gotoSetMap =

new Dictionary<string, ISet<MiddleCode>>();

In the GotoStatement method we start by adding middle code holding a goto instruction.

public static Statement GotoStatement(string labelName) {

List<MiddleCode> gotoList = new List<MiddleCode>();

MiddleCode gotoCode = AddMiddleCode(gotoList, MiddleOperator.Goto);

If the label name is already a key in the goto set map, we look up the goto set and a the goto middle code instruction. This situation occurs in when we encounter the first jump to the label.

if (m\_gotoSetMap.ContainsKey(labelName)) {

ISet<MiddleCode> gotoSet = m\_gotoSetMap[labelName];

gotoSet.Add(gotoCode);

}

If the label name is not already a key in the goto set map, we create a new goto set to which we add the goto instruction, and then we add the label name and goto set to the the goto set map. This situation occurs when we encounter the the jumps following the first jump to the label. In this way, we have a map where each label name is associated with the goto instruction jump to that label.

else {

ISet<MiddleCode> gotoSet = new HashSet<MiddleCode>();

gotoSet.Add(gotoCode);

m\_gotoSetMap.Add(labelName, gotoSet);

}

return (new Statement(gotoList));

}

The BackpatchGoto method is called at the end of each function definition. Its iterates through the goto set map and, for each label name, we look up the label middle code instruction and backpatch the goto set to that instruction.

public static void BackpatchGoto() {

foreach (KeyValuePair<string,ISet<MiddleCode>> entry in m\_gotoSetMap) {

string labelName = entry.Key;

ISet<MiddleCode> gotoSet = entry.Value;

If the label name is not a key in the label map, we have a goto statement to an unknown label and we report an error.

MiddleCode labelCode;

Assert.Error(m\_labelMap.TryGetValue(labelName, out labelCode),

labelName, Message.Missing\_goto\_address);

Backpatch(gotoSet, labelCode);

}

}

### Return Statement

The return statement may have an optional expression.

The generateReturnStatement method is surprisely complicated, which due to the fact the return statement shall be interpreted as an exit statement in case of the main function.

If the surrounding function returns a pointer to a function and the return expression is a or a pointer to a function, we check the they are equal; that is, that they return equal types and have equal parameter lists.

If the sorrounding function returns a pointer and the expression type is array, we compare the pointer type with array type. If the return expression is a string we check that the surrounding function’s return type is a pointer to a (signed or unsigned) character. Otherwise, we just cast the return expression to the surrounding function’s return type.

If the surrounding function is the main function, we cast the return expression to a signed short integer, since we already checked that main does not return void, it has to return an integral type. If the surrounding function is not main, we generate middle code instructions for the return value and function return. Note that there are two different instructions, where the first one sets the return value and the second one performs the return jump.

On the other hand, if there is no return expression we check that the surrounding function returns void. If it is the main function, we add the exit middle code instruction with argument zero. If it is not, we just add the return middle code instruction without precede it with the setting of a return value.

MainParser.cs

closed\_statement:

RETURN optional\_expression SEMICOLON {

$$ = MiddleCodeGenerator.ReturnStatement($2);

}

We need to regard whether there is an expression and whether the return statement is located inside the main function.

MiddleCodeGenerator.cs

public static Statement ReturnStatement(Expression expression) {

List<MiddleCode> codeList;

If the expression is not null, we need the check that the function does not return void. If it does we report an error.

if (expression != null) {

Assert.Error(!SymbolTable.CurrentFunction.Type.ReturnType.IsVoid(),

Message.Non\_\_void\_return\_from\_void\_function);

We cast the return expression to the return type of the function.

expression = TypeCast.ImplicitCast(expression,

SymbolTable.CurrentFunction.Type.ReturnType);

codeList = expression.LongList;

If the function is the main function, we shall not return a value. Instead, we shall exit the program execution and return an integer value to the enclosing system.

if (SymbolTable.CurrentFunction.UniqueName.Equals("main")) {

AddMiddleCode(codeList, MiddleOperator.Exit,

null, expression.Symbol);

}

If the function is not the main function, we return the value of the expression.

else {

AddMiddleCode(codeList, MiddleOperator.SetReturnValue,

null, expression.Symbol);

AddMiddleCode(codeList, MiddleOperator.Return);

}

}

If the expression is null, we check that the function returns void. If it does not we report and error.

else {

Assert.Error(SymbolTable.CurrentFunction.Type.ReturnType.IsVoid(),

Message.Void\_returned\_from\_non\_\_void\_function);

codeList = new List<MiddleCode>();

If the function is the main function, we exit the execution of the program without a return value.

if (SymbolTable.CurrentFunction.UniqueName.Equals("main")) {

AddMiddleCode(codeList, MiddleOperator.Exit);

}

If the function is not the main function, return the function without a return value.

else {

AddMiddleCode(codeList, MiddleOperator.Return);

}

}

return (new Statement(codeList));

}

### Optional Expression Statement

A statement can also be made up by an optional expression; that is, the statement is an expression followed by a semicolon, or simply a semicolon.

An expression statement is an optional expression followed by a semicolon. The expression is evaluated and its true and false sets are backpatched to the beginning of the next statement. Note that we do not use the result of the expression, the only interested part is its potential side effects. If it has no side effects, the statements will be removed by the middle code optimizer in Chapter 14.

MainParser.gppg

closed\_statement:

optional\_expression SEMICOLON {

$$ = MiddleCodeGenerator.ExpressionStatement($1);

}

If there is an expression, we add its short list to the code list. We choose the short list rather than the long list since we are only interested in the side effects of the expression.

MiddleCodeGenerator.cs

public static Statement ExpressionStatement(Expression expression) {

List<MiddleCode> codeList = new List<MiddleCode>();

if (expression != null) {

codeList.AddRange(expression.ShortList);

}

return (new Statement(codeList));

}

### Block Statement

A statement can be an optional sequence of statements enclosed in brackets. The sequence is parsed with a new symbol table.

The block statement simple, a block is just a list of statements surrounded by brackets. Since a two variables with the same name can be defined in different blocks, we push the symbol table before the statement list.

MainParser.cs

closed\_statement:

LEFT\_BLOCK {

SymbolTable.CurrentTable =

new SymbolTable(SymbolTable.CurrentTable, Scope.Block);

}

optional\_statement\_list RIGHT\_BLOCK {

SymbolTable.CurrentTable =

SymbolTable.CurrentTable.ParentTable;

$$ = $3;

}

In case of an empty statement list, we return a statement with an empty code list and an empty next set.

optional\_statement\_list:

/\* Empty \*/ {

$$ = new Statement(new List<MiddleCode>(),

new HashSet<MiddleCode>());

}

In case of an non-empty statement list, we add the code list of the statements of the statement list. For each statement in the list, except the last statement, we backpatch the next set to the beginning of the code list of the next set. The result is a statement with the total middle code list and the next set of the last statement.

| optional\_statement\_list statement {

MiddleCodeGenerator.Backpatch($1.NextSet, $2.CodeList);

List<MiddleCode> codeList = new List<MiddleCode>();

codeList.AddRange($1.CodeList);

codeList.AddRange($2.CodeList);

$$ = new Statement(codeList, $2.NextSet);

};

### Declaration Statement

A statement can be a declaration. We return a statement with the middle code list of the declaration. The middle code list may hold code for initialization of variables or constants.

MainParser.cs

closed\_statement:

declaration {

$$ = new Statement($1);

}

### Jump Register Statements

When performing a call to a function which address is stored in a pointer variables we need to store the address in a register and jump to that register.

MainParser.gppg

| JUMP\_REGISTER LEFT\_PARENTHESIS REGISTER\_NAME RIGHT\_PARENTHESIS SEMICOLON {

$$ = MiddleCodeGenerator.JumpRegisterStatement($3);

}

MiddleCodeGenerator.cs

public static Statement JumpRegisterStatement(Register register) {

List<MiddleCode> codeList = new List<MiddleCode>();

AddMiddleCode(codeList, MiddleOperator.JumpRegister, register);

return (new Statement(codeList));

}

### Interrupt Statements

When making system calls, an interrupt occurs. The operand is an integral value of short size (1 byte).

MainParser.gppg

| INTERRUPT LEFT\_PARENTHESIS constant\_integral\_expression RIGHT\_PARENTHESIS

SEMICOLON {

$$ = MiddleCodeGenerator.InterruptStatement($3);

}

MiddleCodeGenerator.cs

public static Statement InterruptStatement(Expression expression) {

List<MiddleCode> codeList = new List<MiddleCode>();

AddMiddleCode(codeList, MiddleOperator.Interrupt,

expression.Symbol.Value);

return (new Statement(codeList));

}

### System Call Statements

System calls for Linux.

MainParser.gppg

| SYSCALL LEFT\_PARENTHESIS RIGHT\_PARENTHESIS SEMICOLON {

$$ = MiddleCodeGenerator.SyscallStatement();

};

MiddleCodeGenerator.cs

public static Statement SyscallStatement() {

List<MiddleCode> codeList = new List<MiddleCode>();

AddMiddleCode(codeList, MiddleOperator.SysCall);

return (new Statement(codeList));

}

## Expressions

The third part of the parser is the expressions. The idea is that we start with the expression of lowest precedence and add a new rule for each precedence.

MainParser.cs

optional\_expression:

/\* Empty \*/ { $$ = null; }

| expression { $$ = $1; };

An expression may be an assignment expression or two expressions separated by a comma.

expression:

assignment\_expression {

$$ = $1;

}

| expression COMMA assignment\_expression {

$$ = MiddleCodeGenerator.CommaExpression($1, $3);

};

The value of a comma expression is the value of the right expression, the value of the left expression is discarded. However, we keep the side effects of both the left and right expressions.

MiddleCodeGenerator.cs

public static Expression CommaExpression(Expression leftExpression,

Expression rightExpression) {

List<MiddleCode> shortList = new List<MiddleCode>();

shortList.AddRange(leftExpression.ShortList);

shortList.AddRange(rightExpression.ShortList);

Note that we add the short list of the left expression to the middle code list. The reason for this is that we are not interested in the value of the left expression, only its side effects.

List<MiddleCode> longList = new List<MiddleCode>();

longList.AddRange(leftExpression.ShortList);

longList.AddRange(rightExpression.LongList);

return (new Expression(rightExpression.Symbol, shortList, longList));

}

### The Assignment Expression

There are two kinds of assignment: simple and compound. In the case of compound assignment, we first evaluate the corresponding binary expression (addition in the example below), assign its result to a temporary variable that we then assign to original left symbol.

|  |  |  |
| --- | --- | --- |
| x += y; | $1 = x + y  x = $1 |  |

An assignment expression can be condition expression, a simple assignment, or a compound assignment.

MainParser.cs

assignment\_expression:

condition\_expression {

$$ = $1;

}

| prefix\_expression assignment\_operator assignment\_expression {

$$ = MiddleCodeGenerator.AssignmentExpression($2, $1, $3);

};

assignment\_operator:

ASSIGN { $$ = MiddleOperator.Assign; }

| ADD\_ASSIGN { $$ = MiddleOperator.BinaryAdd; }

| SUBTRACT\_ASSIGN { $$ = MiddleOperator.BinarySubtract; }

| MULTIPLY\_ASSIGN { $$ = MiddleOperator.SignedMultiply; }

| DIVIDE\_ASSIGN { $$ = MiddleOperator.SignedDivide; }

| MODULO\_ASSIGN { $$ = MiddleOperator.SignedModulo; }

| AND\_ASSIGN { $$ = MiddleOperator.BitwiseAnd; }

| OR\_ASSIGN { $$ = MiddleOperator.BitwiseOr; }

| XOR\_ASSIGN { $$ = MiddleOperator.BitwiseXOr; }

| LEFT\_SHIFT\_ASSIGN { $$ = MiddleOperator.ShiftLeft; }

| RIGHT\_SHIFT\_ASSIGN { $$ = MiddleOperator.ShiftRight; };

MiddleCodeGenerator.cs

public static Expression AssignmentExpression(MiddleOperator middleOp,

Expression leftExpression,

Expression rightExpression){

In case of compound assignment, we call the matching method to perform the operation, and

switch (middleOp) {

case MiddleOperator.Assign:

return Assignment(leftExpression, rightExpression, true);

In cases of compound assignment, we call the matching method to perform the operation, and Assignment for the final assignment.

case MiddleOperator.BinaryAdd:

case MiddleOperator.BinarySubtract:

return Assignment(leftExpression,

AdditionExpression(middleOp, leftExpression,

rightExpression));

case MiddleOperator.SignedMultiply:

case MiddleOperator.SignedDivide:

case MiddleOperator.SignedModulo:

return Assignment(leftExpression,

MultiplyExpression(middleOp, leftExpression,

rightExpression));

case MiddleOperator.BitwiseAnd:

case MiddleOperator.BitwiseOr:

case MiddleOperator.BitwiseXOr:

return Assignment(leftExpression,

BitwiseExpression(middleOp, leftExpression,

rightExpression));

default: // shift left, shift right

return Assignment(leftExpression,

ShiftExpression(middleOp, leftExpression,

rightExpression));

}

}

The Assignment method performs simple or compound assignment.

public static Expression Assignment(Expression leftExpression,

Expression rightExpression,

bool simpleAssignment = false) {

Register? register = leftExpression.Register;

In cases of system calls, a specific register is assigned a value.

if (register != null) {

Symbol rightSymbol = rightExpression.Symbol;

We check that the register has the same size (in bytes) as the expression it is assigned to. If it is not we report an error.

Assert.Error(AssemblyCode.SizeOfRegister(register.Value) ==

rightExpression.Symbol.Type.Size(),

Message.Unmatched\_register\_size);

We add the long list of the right expression, since we need its value, to the final middle code list.

List<MiddleCode> codeList = new List<MiddleCode>();

codeList.AddRange(rightExpression.LongList);

We add the assignment of the value of right expression to the register.

AddMiddleCode(codeList, MiddleOperator.AssignRegister,

register, rightExpression.Symbol);

Since this an assignment, the code list holds both the assignment of the register and the side effect of the assignment. Therefore, we add the long list as both the short list and the long list of the final expression.

return (new Expression(rightExpression.Symbol, codeList, codeList));

}

If the left expression is not a register, we first check that the left expression is assignable. If it is not, we report an error.

else {

Assert.Error(leftExpression.Symbol.Assignable,

leftExpression, Message.Not\_assignable);

List<MiddleCode> codeList = new List<MiddleCode>();

In case of a simple assignment we add the long list of the left expression to the code list. In case of a compound assignment, the long list of the left expression has already been added to long list of the right expression. If we added the long list in case of a compound assignment, we would add the same code twice.

if (simpleAssignment) {

codeList.AddRange(leftExpression.LongList);

If the left expression is of floating type, we need to pop the floating value stack in case of a simple assignment. In case of a compound assignment, we let the value stay on the stack.

if (leftExpression.Symbol.Type.IsFloating()) {

AddMiddleCode(codeList, MiddleOperator.PopFloat);

}

}

We type cast the right expression to the type of the left expression and add its long list to the final code list.

rightExpression = TypeCast.ImplicitCast(rightExpression,

leftExpression.Symbol.Type);

codeList.AddRange(rightExpression.LongList);

If the expressions hold floating type, we really do not perform an assignment, we simply top the current value on the floating value stack to the left expression.

if (leftExpression.Symbol.Type.IsFloating()) {

AddMiddleCode(codeList, MiddleOperator.TopFloat,

leftExpression.Symbol);

List<MiddleCode> shortList = new List<MiddleCode>();

shortList.AddRange(codeList);

In case of floating type, there is a difference between the short list and long list. In the long list, the resulting value of the assignment (which equals the new value of the left expression) shall be used. Therefore, it is preserved on the stack. In the short list, the value is popped from the stack, since the value shall not be used, only the side effect of the assignment.

AddMiddleCode(shortList, MiddleOperator.PopFloat);

return (new Expression(leftExpression.Symbol, shortList, codeList));

}

In case of integral type, we add an assign instruction to the final code.

else {

AddMiddleCode(codeList, MiddleOperator.Assign,

leftExpression.Symbol, rightExpression.Symbol);

BigInteger? bitFieldMask =

leftExpression.Symbol.Type.BitfieldMask();

In case of bitfield, we add an add operation to set the non-relevant bits to zero.

if (bitFieldMask != null) {

Symbol maskSymbol = new Symbol(leftExpression.Symbol.Type,

bitFieldMask);

AddMiddleCode(codeList, MiddleOperator.BitwiseAnd,

leftExpression.Symbol, leftExpression.Symbol,

maskSymbol);

}

In the integral case, we set the code list as both the short list and long list in the resulting expression, since the assignment is a side effect.

return (new Expression(leftExpression.Symbol, codeList, codeList));

}

}

}

### The Condition Expression

The conditional operator applies lazy evaluation, which means that only one of the true or false expression will be evaluated.

MainParser.gppg

condition\_expression:

logical\_or\_expression {

$$ = $1;

}

| logical\_or\_expression QUESTION\_MARK expression COLON condition\_expression{

$$ = MiddleCodeGenerator.ConditionalExpression($1, $3, $5);

};

MiddleCodeGenerator.cs

public static Expression ConditionalExpression(Expression testExpression,

Expression trueExpression,

Expression falseExpression) {

We type cast the test expression to logical type.

testExpression = TypeCast.ToLogical(testExpression);

If the test expression is constant, we simple return the true or false expression depending on whether the test expression is true.

if (ConstantExpression.IsConstant(testExpression)) {

return ConstantExpression.IsTrue(testExpression)

? falseExpression : trueExpression;

}

If both the true and false expressions hold logical types, we keep their types.

if (trueExpression.Symbol.Type.IsLogical() &&

falseExpression.Symbol.Type.IsLogical()) {

We start by backpathing the true and false set of the test expression to the beginning of the true and false expression’s code list.

Backpatch(testExpression.Symbol.TrueSet, trueExpression.LongList);

Backpatch(testExpression.Symbol.FalseSet, falseExpression.LongList);

The resulting true set is the union of the true sets of the true and false expression, and the resulting false set is the union of the false sets of the true and false expression.

ISet<MiddleCode> trueSet = new HashSet<MiddleCode>(),

falseSet = new HashSet<MiddleCode>();

trueSet.UnionWith(trueExpression.Symbol.TrueSet);

trueSet.UnionWith(falseExpression.Symbol.TrueSet);

falseSet.UnionWith(trueExpression.Symbol.FalseSet);

falseSet.UnionWith(falseExpression.Symbol.FalseSet);

List<MiddleCode> shortList = new List<MiddleCode>();

If the short list of both the true and false expression is empty, it does not matter if the test expression is true or false, and we let the resulting short list be the short list of the true expression (which may be empty).

if (IsEmpty(trueExpression.ShortList) &&

IsEmpty(falseExpression.ShortList)) {

shortList.AddRange(testExpression.ShortList);

}

If the short list of the test expression is not empty, the situation becomes a bit more complicated. We add the long list, rather than the short list, of the test expression to the final short list since we need the value of the test expression in order to jump to the short list of the true of false expression.

else {

shortList.AddRange(testExpression.LongList);

shortList.AddRange(trueExpression.ShortList);

shortList.AddRange(falseExpression.ShortList);

}

We add the long list of the test, true, and false expression to the final long list.

List<MiddleCode> longList = new List<MiddleCode>();

longList.AddRange(testExpression.LongList);

longList.AddRange(trueExpression.LongList);

longList.AddRange(falseExpression.LongList);

Finally, we create a new symbol with logical type and the resulting true and false sets.

Symbol symbol = new Symbol(trueSet, falseSet);

return (new Expression(symbol, shortList, longList));

}

If at least one of the true or false expression does not hold logical type, we define maxType as their largest type, and type cast both expressions to that type.

else {

Type maxType = TypeCast.MaxType(trueExpression.Symbol.Type,

falseExpression.Symbol.Type);

trueExpression = TypeCast.ImplicitCast(trueExpression, maxType);

Backpatch(testExpression.Symbol.TrueSet, trueExpression.LongList);

We create a new temporary symbol to hold the result of the condition expression.

Symbol symbol = new Symbol(maxType);

In case of non-floating type, we add the assignment instruction. In case of a floating type, the value is already placed at the floating value stack and we do not need to do anything.

if (!maxType.IsFloating()) {

AddMiddleCode(trueExpression.LongList, MiddleOperator.Assign,

symbol, trueExpression.Symbol);

}

We add the jump to a target code to both the short and long list of the true expression. The target code will be added after the code of the false expression. Its purpose is to jump over the false expression code.

MiddleCode targetCode = new MiddleCode(MiddleOperator.Empty);

AddMiddleCode(trueExpression.ShortList,

MiddleOperator.Goto, targetCode);

AddMiddleCode(trueExpression.LongList,

MiddleOperator.Goto, targetCode);

Similar to the true expression, we type cast and false expression, and backpatch the true set of the test expression to the beginning of the false expression code.

falseExpression = TypeCast.ImplicitCast(falseExpression, maxType);

Backpatch(testExpression.Symbol.FalseSet, falseExpression.LongList);

We also assign the value of the false expression to the temporary symbol if it does not hold floating type. If it holds floating type, the value is already placed at the top of the floating value stack, and we do nothing.

if (!maxType.IsFloating()) {

AddMiddleCode(falseExpression.LongList, MiddleOperator.Assign,

symbol, falseExpression.Symbol);

}

If both the short lists of the true and false expressions are empty, we just add the short list of the test expression (which may be empty).

List<MiddleCode> shortList = new List<MiddleCode>();

if (IsEmpty(trueExpression.ShortList) &&

IsEmpty(falseExpression.ShortList)) {

shortList.AddRange(testExpression.ShortList);

}

If not both the short lists of the true and false expressions are empty, we add the long list of the test expression, rather than the short list, since we need to value of the test expression, as well as the short lists of the true and false expressions.

else {

shortList.AddRange(testExpression.LongList);

shortList.AddRange(trueExpression.ShortList);

shortList.AddRange(falseExpression.ShortList);

shortList.Add(targetCode);

}

Finally, add the long lists and return the expression.

List<MiddleCode> longList = new List<MiddleCode>();

longList.AddRange(testExpression.LongList);

longList.AddRange(trueExpression.LongList);

longList.AddRange(falseExpression.LongList);

longList.Add(targetCode);

return (new Expression(symbol, shortList, longList));

}

}

### Constant Expression

On several occasions, such array limits and enumeration values, we need to parse constant integral expressions. The expression is a conditional expression, since it cannot hold commas or assignments.

MainParser.cs

optional\_constant\_integral\_expression:

/\* Empty \*/ { $$ = null; }

| constant\_integral\_expression { $$ = $1; };

constant\_integral\_expression:

condition\_expression {

$$ = MiddleCodeGenerator.ConstantIntegralExpression($1);

};

The ConstantIntegralExpression methods makes sure the expression is indeed constant and either integral or pointer.

MiddleCodeGenerator.cs

public static Expression ConstantIntegralExpression(Expression expression)

{ expression = ConstantExpression.Cast(expression,

Type.SignedLongIntegerType);

Assert.Error(expression != null, expression,

Message.Non\_\_constant\_expression);

Assert.Error(expression.Symbol.Type.IsIntegralOrPointer(),

expression.Symbol, Message.Non\_\_integral\_expression);

return expression;

}

### Logical Expressions

The logical-and expression backpatch the true set of the first operand to the beginning of the right operand. This can be interpreted as if the left expression is true, we try the right expression instead. If both the expressions are true, the result is true. However, if the left expression is false, the result is determined to be false without evaluating the right expression.

The logical-and expression backpatch the true set of the first operand to the beginning of the right operand. This can be interpreted as if the left expression is true, we try the right expression instead. If both the expressions are true, the result is true. However, if the left expression is false, the result is determined to be false without evaluating the right expression.

MainParser.cs

logical\_or\_expression:

logical\_and\_expression {

$$ = $1;

}

| logical\_or\_expression LOGICAL\_OR logical\_and\_expression {

$$ = MiddleCodeGenerator.LogicalOrExpression($1, $3);

};

The LogicalOrExpression method returns a true expression is at least one of the left or right expression is true, and a false expression if both the left and right expression is false.

MiddleCodeGenerator.cs

public static Expression LogicalOrExpression(Expression leftExpression,

Expression rightExpression) {

Since we have passed the constant integral expression parsing, we need to check if the expression is constant. If it is constant, we return the constant expression.

Expression constantExpression =

ConstantExpression.Logical(MiddleOperator.LogicalOr,

leftExpression, rightExpression);

if (constantExpression != null) {

return constantExpression;

}

We type cast both the expression to logical type.

leftExpression = TypeCast.ToLogical(leftExpression);

rightExpression = TypeCast.ToLogical(rightExpression);

For the resulting expression to be true, it is enough that one of the left or right expression is true. Therefore, the true set of the resulting expression is the union of the true set of the left and right expression. This results in laze evaluation, if the left expression is evaluated to true, the right expression (including its side effects) shall not evaluated. The false sets, on the other hand, are different. If the left expression is evaluated to false, we need to evaluate the right expression. Therefore, we backpatch the false set of the left expression to the beginning of the right expression code. The false set of the resulting expression is the false set of the right expression.

ISet<MiddleCode> trueSet = new HashSet<MiddleCode>();

trueSet.UnionWith(leftExpression.Symbol.TrueSet);

trueSet.UnionWith(rightExpression.Symbol.TrueSet);

Backpatch(leftExpression.Symbol.FalseSet, rightExpression.LongList);

Symbol symbol = new Symbol(trueSet, rightExpression.Symbol.FalseSet);

List<MiddleCode> longList = new List<MiddleCode>();

longList.AddRange(leftExpression.LongList);

longList.AddRange(rightExpression.LongList);

List<MiddleCode> shortList = new List<MiddleCode>();

shortList.AddRange(leftExpression.ShortList);

shortList.AddRange(rightExpression.ShortList);

return (new Expression(symbol, shortList, longList));

}

MainParser.cs

logical\_and\_expression:

bitwise\_or\_expression {

$$ = $1;

}

| logical\_and\_expression LOGICAL\_AND bitwise\_or\_expression {

$$ = MiddleCodeGenerator.LogicalAndExpression($1, $3);

};

The LogicalAndExpression method returns a true expression both the left and right expression is true, and a false expression if at least one of the left or right expression is false.

MiddleCodeGenerator.cs

public static Expression LogicalAndExpression(Expression leftExpression,

Expression rightExpression){

Similar to the logical or expression above, we need to check if the expression is constant.

Expression constantExpression =

ConstantExpression.Logical(MiddleOperator.LogicalAnd,

leftExpression, rightExpression);

if (constantExpression != null) {

return constantExpression;

}

We type cast both the expression to logical type.

leftExpression = TypeCast.ToLogical(leftExpression);

rightExpression = TypeCast.ToLogical(rightExpression);

The true and false sets of the works in an opposite way compared to the logical or expression above. The final expression is false if at least one of the left or right expression is false. If the left expression is evaluated to false, the right expression shall not be evaluated. Therefore, the false set of the resulting expression is the union of the false sets of the left and right expression. For the resulting expression to be true, both the left and right expression must be true. We therefore backpatch the true set of the left expression to the beginning of the right expression code. The true set of the resulting expression is the true set of the right expression.

ISet<MiddleCode> falseSet = new HashSet<MiddleCode>();

falseSet.UnionWith(leftExpression.Symbol.FalseSet);

falseSet.UnionWith(rightExpression.Symbol.FalseSet);

Backpatch(leftExpression.Symbol.TrueSet, rightExpression.LongList);

Symbol symbol = new Symbol(rightExpression.Symbol.TrueSet, falseSet);

List<MiddleCode> longList = new List<MiddleCode>();

longList.AddRange(leftExpression.LongList);

longList.AddRange(rightExpression.LongList);

List<MiddleCode> shortList = new List<MiddleCode>();

shortList.AddRange(leftExpression.ShortList);

shortList.AddRange(rightExpression.ShortList);

return (new Expression(symbol, shortList, longList));

}

### Bitwise Expressions

Unlike the logical expressions in the previous section, the inclucise-or, exclusive-or and and bitwise expression does not support lazy evaluation, which means that the right expressions is always evaluated regardless of the value of the left expression. Both the operands have to be integral and the result is integral.

The are three bitwise operators: or, xor (exclusive or), and and. They take two integer values and perform operations on each bit in the values:

* or: one if at least one value is one, zero otherwise
* xor: one if at exact one value is one, zero otherwise
* and: one if at both values are one, zero otherwise

The expressions have one rule each since they have different precedence, but they call the same method: BitwiseExpression in MiddleCodeGenerator.

MainParser.gppg

bitwise\_or\_expression:

bitwise\_xor\_expression {

$$ = $1;

}

| bitwise\_or\_expression BITWISE\_OR bitwise\_xor\_expression {

$$ = MiddleCodeGenerator.BitwiseExpression

(MiddleOperator.BitwiseOr, $1, $3);

};

bitwise\_xor\_expression:

bitwise\_and\_expression {

$$ = $1;

}

| bitwise\_xor\_expression BITWISE\_XOR bitwise\_and\_expression {

$$ = MiddleCodeGenerator.BitwiseExpression

(MiddleOperator.BitwiseXOr, $1, $3);

};

bitwise\_and\_expression:

equality\_expression {

$$ = $1;

}

| bitwise\_and\_expression AMPERSAND equality\_expression {

$$ = MiddleCodeGenerator.BitwiseExpression

(MiddleOperator.BitwiseAnd, $1, $3);

};

MiddleCodeGenerator.cs

public static Expression BitwiseExpression(MiddleOperator middleOp,

Expression leftExpression,

Expression rightExpression) {

If the expresion can be evaluated to a constant value, we return the constant expression.

Expression constantExpression = ConstantExpression.

Arithmetic(middleOp, leftExpression, rightExpression);

if (constantExpression != null) {

return constantExpression;

}

We find the maximal type of the left and right expression.

Type maxType = TypeCast.MaxType(leftExpression.Symbol.Type,

rightExpression.Symbol.Type);

Symbol resultSymbol = new Symbol(maxType);

The type of the left and right expression must be integral or pointer. Array, string, and functions are converted to pointers.

Assert.Error(maxType.IsIntegralPointerArrayStringOrFunction(),

maxType, Message.Invalid\_type\_in\_bitwise\_expression);

We type cast both expressions into the maximal type.

leftExpression = TypeCast.ImplicitCast(leftExpression, maxType);

rightExpression = TypeCast.ImplicitCast(rightExpression, maxType);

The short list of the resulting expression is simply the short list of the expressions.

List<MiddleCode> shortList = new List<MiddleCode>();

shortList.AddRange(leftExpression.ShortList);

shortList.AddRange(rightExpression.ShortList);

The long list of the resulting expression is the long list of the expressions, and the bitwise operation.

List<MiddleCode> longList = new List<MiddleCode>();

longList.AddRange(leftExpression.LongList);

longList.AddRange(rightExpression.LongList);

AddMiddleCode(longList, middleOp, resultSymbol,

leftExpression.Symbol, rightExpression.Symbol);

return (new Expression(resultSymbol, shortList, longList));

}

### Shift Expression

The left and right expression of a shift expression must hold integral types. The right expression is type cast to a one-byte integer value.

MainParser.gppg

shift\_expression:

add\_expression {

$$ = $1;

}

| shift\_expression shift\_operator add\_expression {

$$ = MiddleCodeGenerator.ShiftExpression($2, $1, $3);

};

shift\_operator:

LEFT\_SHIFT { $$ = MiddleOperator.ShiftLeft; }

| RIGHT\_SHIFT { $$ = MiddleOperator.ShiftRight; };

**MiddleCodeGenerator.cs**

public static Expression ShiftExpression(MiddleOperator middleOp,

Expression leftExpression,

Expression rightExpression) {

First we check that the left expression hols integral or pointer type. An array, a string, or a function is cast to a pointer.

Assert.Error(leftExpression.Symbol.Type.

IsIntegralPointerArrayStringOrFunction(),

leftExpression, Message.Invalid\_type\_in\_shift\_expression);

If the expression can be evaluated to a constant expression, we return the constant expression.

Expression constantExpression =

ConstantExpression.Arithmetic(middleOp, leftExpression,

rightExpression);

if (constantExpression != null) {

return constantExpression;

}

The right expression is type cast to an unsigned character, which always have a size of one byte.

rightExpression =

TypeCast.ImplicitCast(rightExpression, Type.UnsignedCharType);

The final short list is simple the short lists of the left and right expression.

List<MiddleCode> shortList = new List<MiddleCode>();

shortList.AddRange(leftExpression.ShortList);

shortList.AddRange(rightExpression.ShortList);

The final short list is simple the short lists of the left and right expression.

List<MiddleCode> longList = new List<MiddleCode>();

longList.AddRange(leftExpression.LongList);

longList.AddRange(rightExpression.LongList);

The final long list is the short lists of the left and right expression, and the shift operation.

Symbol resultSymbol = new Symbol(leftExpression.Symbol.Type);

AddMiddleCode(longList, middleOp, resultSymbol,

leftExpression.Symbol, rightExpression.Symbol);

return (new Expression(resultSymbol, shortList, longList));

}

### Equality and Relation Expressions

Values of all types except structs or unions can be compared with equality and inequality operator. All arithmetic and pointer values can be compared with relation operators. The equality operators have higher precedence than the relation operators, but they call the same method: RelationalExpression in MiddleCodeGenerator.

MainParser.gppg

equality\_expression:

relation\_expression {

$$ = $1;

}

| equality\_expression equality\_operator relation\_expression {

$$ = MiddleCodeGenerator.RelationalExpression($2, $1, $3);

};

equality\_operator:

EQUAL { $$ = MiddleOperator.Equal; }

| NOT\_EQUAL { $$ = MiddleOperator.NotEqual; };

relation\_expression:

shift\_expression {

$$ = $1;

}

| relation\_expression relation\_operator shift\_expression {

$$ = MiddleCodeGenerator.RelationalExpression ($2, $1, $3);

};

relation\_operator:

LESS\_THAN { $$ = MiddleOperator.SignedLessThan; }

| LESS\_THAN\_EQUAL { $$ = MiddleOperator.SignedLessThanEqual; }

| GREATER\_THAN { $$ = MiddleOperator.SignedGreaterThan; }

| GREATER\_THAN\_EQUAL { $$ = MiddleOperator.SignedGreaterThanEqual; };

MiddleCodeGenerator.cs

public static Expression RelationalExpression(MiddleOperator middleOp,

Expression leftExpression,

Expression rightExpression){

First, we check the types of the expression. Everything except struct or union can be compared.

Assert.Error(!leftExpression.Symbol.Type.IsStructOrUnion(),

leftExpression,

Message.Invalid\_type\_in\_expression);

Assert.Error(!rightExpression.Symbol.Type.IsStructOrUnion(),

rightExpression,

Message.Invalid\_type\_in\_expression);

The find the maximal type of the left and right expression types, and type cast both the expressions to that type.

Type maxType = TypeCast.MaxType(leftExpression.Symbol.Type,

rightExpression.Symbol.Type);

leftExpression = TypeCast.ImplicitCast(leftExpression, maxType);

rightExpression = TypeCast.ImplicitCast(rightExpression, maxType);

The final short list is simply the short lists of the left end right expression.

List<MiddleCode> shortList = new List<MiddleCode>();

shortList.AddRange(leftExpression.ShortList);

shortList.AddRange(rightExpression.ShortList);

If the maximal type is unsigned, we change to operator from signed to unsigned.

if (maxType.IsUnsigned()) {

string name = Enum.GetName(typeof(MiddleOperator), middleOp);

middleOp = (MiddleOperator) Enum.Parse(typeof(MiddleOperator),

name.Replace("Signed", "Unsigned"));

}

The final long list is the long lists of the left end right expression, to begin with.

List<MiddleCode> longList = new List<MiddleCode>();

longList.AddRange(leftExpression.LongList);

longList.AddRange(rightExpression.LongList);

We add two instruction to the long list. The first is an if-goto instruction that we add to the true set. If the expression is true, this instruction will jump to the target that later will be backpatched into the instruction. In the same way, we add a goto instruction to the false set. If the expression is false, the instruction will jump to the target later backpatched to the instruction.

ISet<MiddleCode> trueSet = new HashSet<MiddleCode>(),

falseSet = new HashSet<MiddleCode>();

trueSet.Add(AddMiddleCode(longList, middleOp, null,

leftExpression.Symbol,

rightExpression.Symbol));

falseSet.Add(AddMiddleCode(longList, MiddleOperator.Goto));

The final symbol holds logical type with true and false set.

Symbol symbol = new Symbol(trueSet, falseSet);

return (new Expression(symbol, shortList, longList));

}

### Addition and Subtraction Expression

The addition and subtraction expression are a bit more complicated, since we need to take pointer arithmetic into consideration.

MainParser.gppg

add\_expression:

multiply\_expression {

$$ = $1;

}

| add\_expression addition\_operator multiply\_expression {

$$ = MiddleCodeGenerator.AdditionExpression($2, $1, $3);

};

add\_operator:

PLUS { $$ = MiddleOperator.BinaryAdd; }

| MINUS { $$ = MiddleOperator.BinarySubtract; };

MiddleCodeGenerator.cs

public static Expression AdditionExpression(MiddleOperator middleOp,

Expression leftExpression,

Expression rightExpression) {

Type leftType = leftExpression.Symbol.Type,

rightType = rightExpression.Symbol.Type;

Similar to the other cases above, the return the possible constant expression.

Expression constantExpression =

ConstantExpression.Arithmetic(MiddleOperator.BinaryAdd,

leftExpression, rightExpression);

if (constantExpression != null) {

return constantExpression;

}

In the addition case, we also check if the expression is a static expression, in which case we return the expression.

Expression staticExpression =

StaticExpression.Binary(MiddleOperator.BinarySubtract,

leftExpression, rightExpression);

if (staticExpression != null) {

return staticExpression;

}

The final short list is simple the short lists of the left and right expression.

List<MiddleCode> shortList = new List<MiddleCode>();

shortList.AddRange(leftExpression.ShortList);

shortList.AddRange(rightExpression.ShortList);

If at least on the expressions has pointer or array type, we call the PointerArithmetic methods that handles pointer arithmetic.

if (leftType.IsPointerOrArray()) {

return PointerArithmetics(middleOp, leftExpression, rightExpression);

}

else if (rightType.IsPointerOrArray()) {

In the right expression has pointer type, the operator must be addition, not subtraction. Note that we swap the expressions in the call to PointerArithmetic, so that the pointer is the first expression

Assert.Error(middleOp == MiddleOperator.BinaryAdd, middleOp,

Message.Invalid\_types\_in\_subtraction\_expression);

return PointerArithmetics(middleOp, rightExpression, leftExpression);

}

If none of the left or right expressions are pointers, we perform an arithmetic operation. First, we make sure that both expressions have arithmetic types (integral or floating).

else {

Assert.Error(leftExpression.Symbol.Type.IsArithmetic(),

leftExpression, Message.Non\_\_arithmetic\_expression);

Assert.Error(rightExpression.Symbol.Type.IsArithmetic(),

rightExpression, Message.Non\_\_arithmetic\_expression);

The we find the maximal type and cast both expressions to the type. The resulting expression does also have the maximal type.

Type maxType = TypeCast.MaxType(leftType, rightType);

leftExpression = TypeCast.ImplicitCast(leftExpression, maxType);

rightExpression = TypeCast.ImplicitCast(rightExpression, maxType);

Symbol resultSymbol = new Symbol(maxType);

The final long list is the long lists of the two expressions, and the operation (addition or subtraction).

List<MiddleCode> longList = new List<MiddleCode>();

longList.AddRange(leftExpression.LongList);

longList.AddRange(rightExpression.LongList);

AddMiddleCode(longList, middleOp, resultSymbol,

leftExpression.Symbol, rightExpression.Symbol);

return (new Expression(resultSymbol, shortList, longList));

}

}

The PointerArithmetic methods performs pointer arithmetic. We have two cases: addition of a pointer and an integral value, and subtraction of two pointer values.

private static Expression PointerArithmetic(MiddleOperator middleOp,

Expression leftExpression,

Expression rightExpression) {

List<MiddleCode> shortList = new List<MiddleCode>();

shortList.AddRange(leftExpression.LongList);

shortList.AddRange(rightExpression.LongList);

Type leftType = leftExpression.Symbol.Type,

rightType = rightExpression.Symbol.Type;

If the left type is pointer and the right type is integral, we add an integer values to and pointer. What makes this different compared to regular addition is that we need take the size of the pointer type into consideration.

if (leftType.IsPointerOrArray() && rightType.IsIntegral()) {

We make sure that the left pointer type is not void, since pointer arithmetic is not allowed on void pointers.

Assert.Error(!leftType.PointerOrArrayType.IsVoid(),

leftExpression, Message.Pointer\_to\_void);

We type cast the integral right expression to the left pointer type for the arithmetic to be performed on values of the same size.

rightExpression = TypeCast.ImplicitCast(rightExpression, leftType);

The final long list is the long lists of the left and right expression, and the pointer arithmetic.

List<MiddleCode> longList = new List<MiddleCode>();

longList.AddRange(leftExpression.LongList);

longList.AddRange(rightExpression.LongList);

The result symbol has the same type as the pointer expression.

Symbol resultSymbol =

new Symbol(new Type(leftType.PointerOrArrayType));

The pointer arithmetic is performed in two steps. First, we multiply the integral value with the pointer size, and then we add the product to the

int pointerSize = leftType.PointerOrArrayType.Size();

Symbol multSymbol = new Symbol(Type.PointerTypeX),

sizeSymbol = new Symbol(Type.PointerTypeX,

(BigInteger) pointerSize);

We multiply the integral value with the size of the pointer type and add the product to the pointer value. The multSymbol value be the result of two constant values, in case of a constant integral value, or the integral value times one, in case of a pointer size of one, which is inefficient. However, the Middle Code Optimizer in Chapter XXX will take care of those cases.

AddMiddleCode(longList, MiddleOperator.UnsignedMultiply, multSymbol,

rightExpression.Symbol, sizeSymbol);

AddMiddleCode(longList, middleOp, resultSymbol,

leftExpression.Symbol, multSymbol);

return (new Expression(resultSymbol, shortList, longList));

}

If it is not the case of a pointer and an integral value, then there must be two pointer values. We make sure that none of them pointes at void. Moreover, we also make sure that their pointer types have the same size.

else {

Assert.Error(!leftType.PointerOrArrayType.IsVoid(),

leftExpression, Message.Pointer\_to\_void);

Assert.Error(!rightType.PointerOrArrayType.IsVoid(),

rightExpression, Message.Pointer\_to\_void);

Assert.Error(leftType.PointerOrArrayType.Size() ==

rightType.PointerOrArrayType.Size(),

leftType + " and " + rightType,

Message.Invalid\_expression);

List<MiddleCode> longList = new List<MiddleCode>();

longList.AddRange(leftExpression.LongList);

longList.AddRange(rightExpression.LongList);

Symbol resultSymbol = new Symbol(Type.VoidPointerType);

int pointerSize = rightType.PointerOrArrayType.Size();

Symbol subtractSymbol = new Symbol(Type.PointerTypeX),

sizeSymbol = new Symbol(Type.PointerTypeX,

(BigInteger) pointerSize);

We first subtract the pointer values, and then divide the difference by the size of the pointer types. There may be a division by one, in case the pointer types have size one, which is inefficient. However, the Middle Code Optimizer in Chapter XXX will take care of the case.

AddMiddleCode(longList, MiddleOperator.BinarySubtract, subtractSymbol,

leftExpression.Symbol, rightExpression.Symbol);

AddMiddleCode(longList, MiddleOperator.UnsignedDivide, resultSymbol,

subtractSymbol, sizeSymbol);

Expression resultExpression = new Expression(resultSymbol,

shortList, longList);

Finally, we type cast the result of the division into signed integer.

return TypeCast.ImplicitCast(resultExpression,

Type.SignedIntegerType);

}

}

### Multiplication Expressions

MainParser.gppg

multiply\_expression:

type\_cast\_expression {

$$ = $1;

}

| multiply\_expression multiply\_operator type\_cast\_expression {

$$ = MiddleCodeGenerator.MultiplyExpression($2, $1, $3);

};

The multiplcation operators are multiply, division, and modulo.

multiply\_operator:

ASTERRISK { $$ = MiddleOperator.SignedMultiply; }

| DIVIDE { $$ = MiddleOperator.SignedDivide; }

| MODULO { $$ = MiddleOperator.SignedModulo; };

MiddleCodeGenerator.cs

public static Expression MultiplyExpression(MiddleOperator middleOp,

Expression leftExpression,

Expression rightExpression) {

List<MiddleCode> constLongList = new List<MiddleCode>();

constLongList.AddRange(leftExpression.LongList);

constLongList.AddRange(rightExpression.LongList);

Expression constantExpression =

ConstantExpression.Arithmetic(middleOp, leftExpression,

rightExpression);

if (constantExpression != null) {

return constantExpression;

}

Type maxType = TypeCast.MaxType(leftExpression.Symbol.Type,

rightExpression.Symbol.Type);

if (MiddleCode.IsModulo(middleOp)) {

Assert.Error(maxType.IsIntegralPointerArrayStringOrFunction(),

maxType, Message.Invalid\_type\_in\_expression);

}

else {

Assert.Error(maxType.IsArithmeticPointerArrayStringOrFunction(),

maxType, Message.Invalid\_type\_in\_expression);

}

leftExpression = TypeCast.ImplicitCast(leftExpression, maxType);

rightExpression = TypeCast.ImplicitCast(rightExpression, maxType);

Symbol resultSymbol = new Symbol(maxType);

if (maxType.IsUnsigned()) {

string name = Enum.GetName(typeof(MiddleOperator), middleOp);

middleOp = (MiddleOperator) Enum.Parse(typeof(MiddleOperator),

name.Replace("Signed", "Unsigned"));

}

List<MiddleCode> shortList = new List<MiddleCode>(),

longList = new List<MiddleCode>();

shortList.AddRange(leftExpression.ShortList);

shortList.AddRange(rightExpression.ShortList);

longList.AddRange(leftExpression.LongList);

longList.AddRange(rightExpression.LongList);

AddMiddleCode(longList, middleOp, resultSymbol,

leftExpression.Symbol, rightExpression.Symbol);

return (new Expression(resultSymbol, shortList, longList));

}

### Cast Expressions

A type cast expression is a type name within parentheses followed by a type cast expression.

MainParser.gppg

type\_cast\_expression:

prefix\_expression {

$$ = $1;

}

| LEFT\_PARENTHESIS type\_name RIGHT\_PARENTHESIS type\_cast\_expression {

$$ = MiddleCodeGenerator.CastExpression($2, $4);

};

The type name is a declaration specifier list with or without an abstract declarator.

type\_name:

declaration\_specifier\_list {

$$ = MiddleCodeGenerator.TypeName(Specifier.SpecifierList($1), null);

}

| declaration\_specifier\_list {

SpecifierStack.Push(Specifier.SpecifierList($1));

}

abstract\_declarator {

$$ = MiddleCodeGenerator.TypeName(SpecifierStack.Pop(), $3);

};

The specifier parameter holds the resulting type of the declaration specifier list.

MiddleCodeGenerator.cs

public static Type TypeName(Specifier specifier, Declarator declarator) {

Type specifierType = specifier.Type;

If the declarator is not null, we add the specifier’s type to the declarator, and returns its type.

if (declarator != null) {

declarator.Add(specifierType);

return declarator.Type;

}

If the declarator is null, we just return the specifer’s type.

else {

return specifierType;

}

}

### Prefix Expression

In C, there are several prefix expressions:

* Unary add and minus
* Logical not
* Bitwise not
* The sizeof operator
* The address operator
* The dereferenceence operator

### Unary Addition Expressions

MainParser.gppg

prefix\_expression:

postfix\_expression {

$$ = $1;

}

| prefix\_add\_operator type\_cast\_expression {

$$ = MiddleCodeGenerator.UnaryExpression($1, $2);

};

prefix\_add\_operator:

PLUS { $$ = MiddleOperator.UnaryAdd; }

| MINUS { $$ = MiddleOperator.UnarySubtract; };

MiddleCodeGenerator.cs

public static Expression UnaryExpression(MiddleOperator middleOp,

Expression expression) {

Type type = expression.Symbol.Type;

Assert.Error(type.IsLogical() ||

type.IsArithmeticPointerArrayStringOrFunction(),

expression, Message.Non\_\_arithmetic\_expression);

Expression constantExpression =

ConstantExpression.Arithmetic(middleOp, expression);

if (constantExpression != null) {

return constantExpression;

}

Symbol resultSymbol = new Symbol(expression.Symbol.Type);

AddMiddleCode(expression.LongList, middleOp,

resultSymbol, expression.Symbol);

return (new Expression(resultSymbol, expression.ShortList,

expression.LongList));

}

### Logical Not Expression

The logical-not operator just swaps the true and false sets of its operand.

MainParser.gppg

prefix\_expression:

LOGICAL\_NOT type\_cast\_expression {

$$ = MiddleCodeGenerator.LogicalNotExpression($2);

};

MiddleCodeGenerator.cs

public static Expression LogicalNotExpression(Expression expression) {

Expression constantExpression =

ConstantExpression.LogicalNot(expression);

if (constantExpression != null) {

return constantExpression;

}

The resulting expression is the original expression with swapped true and false sets.

expression = TypeCast.ToLogical(expression);

Symbol notSymbol =

new Symbol(expression.Symbol.FalseSet, expression.Symbol.TrueSet);

return (new Expression(notSymbol, expression.ShortList,

expression.LongList));

}

### Bitwise Not Expression

MainParser.gppg

prefix\_expression:

BITWISE\_NOT type\_cast\_expression {

$$ = MiddleCodeGenerator.BitwiseNotExpression($2);

};

MiddleCodeGenerator.cs

public static Expression BitwiseNotExpression(Expression expression) {

Expression constantExpression =

ConstantExpression.Arithmetic(MiddleOperator.BitwiseNot, expression);

if (constantExpression != null) {

return constantExpression;

}

expression = TypeCast.LogicalToIntegral(expression);

Symbol resultSymbol = new Symbol(expression.Symbol.Type);

AddMiddleCode(expression.LongList, MiddleOperator.BitwiseNot,

resultSymbol, expression.Symbol);

return (new Expression(resultSymbol, expression.ShortList,

expression.LongList));

}

### The Sizeof Expression

The sizeof-operator takes an expression or a type inside parenthesis as operand. The operator is not allowed on function or bitfields.

MainParser.gppg

prefix\_expression:

SIZEOF prefix\_expression {

$$ = MiddleCodeGenerator.SizeOfExpression($2);

};

MiddleCodeGenerator.cs

public static Expression SizeOfExpression(Expression expression) {

We check that the storage of the expression is not register, since it is not allowed.

Assert.Error(!expression.Symbol.IsRegister(), expression,

Message.Register\_storage\_not\_allowed\_in\_sizof\_expression);

Moreover, we check that it is not a function or a bitfield, since they do not have sizes.

Type type = expression.Symbol.Type;

Assert.Error(!type.IsFunction(),

Message.Sizeof\_applied\_to\_function\_not\_allowed);

Assert.Error(!type.IsBitfield(),

Message.Sizeof\_applied\_to\_bitfield\_not\_allowed);

We create a symbol of signed integer types with the size as its value.

Symbol symbol = new Symbol(Type.SignedIntegerType,

(BigInteger) (expression.Symbol.Type.Size()));

symbol.StaticSymbol =

ConstantExpression.Value(symbol.UniqueName, Type.SignedIntegerType,

(BigInteger) (expression.Symbol.Type.Size()));

SymbolTable.StaticSet.Add(symbol.StaticSymbol);

return (new Expression(symbol, new List<MiddleCode>(),

new List<MiddleCode>()));

}

The sizeof operator can also be applied to a type name within parentheses.

MainParser.gppg

prefix\_expression:

SIZEOF LEFT\_PARENTHESIS type\_name RIGHT\_PARENTHESIS {

$$ = MiddleCodeGenerator.SizeOfType($3);

};

In the case of the type name, we alsoe check that the expression is not a function or a bitfield.

MiddleCodeGenerator.cs

public static Expression SizeOfType(Type type) {

Assert.Error(!type.IsFunction(),

Message.Sizeof\_applied\_to\_function\_not\_allowed);

Assert.Error(!type.IsBitfield(),

Message.Sizeof\_applied\_to\_bitfield\_not\_allowed);

Symbol symbol =

new Symbol(Type.SignedIntegerType, (BigInteger) type.Size());

return (new Expression(symbol, new List<MiddleCode>(),

new List<MiddleCode>()));

}

### Address Expression

The address operator takes the address of its operand and the return type is a pointer to the operand type, unless it is an array, in which case the result is a pointer to the array type. The operand cannot have register storage.

MainParser.gppg

prefix\_expression:

AMPERSAND type\_cast\_expression {

$$ = MiddleCodeGenerator.AddressExpression($2);

};

The address operator does not apply to bitfields symbols with register storage.

MiddleCodeGenerator.cs

public static Expression AddressExpression(Expression expression) {

Assert.Error(expression.Symbol.Addressable, expression,

Message.Not\_addressable);

Assert.Error(!expression.Symbol.IsRegister(), expression,

Message.Invalid\_address\_of\_register\_storage);

The address operator may result in a static address.

Expression staticExpression =

StaticExpression.Unary(MiddleOperator.Address, expression);

if (staticExpression!= null) {

return staticExpression ;

}

If the expression has floating type, we need to pop the value from the floating value stack.

if (expression.Symbol.Type.IsFloating()) {

AddMiddleCode(expression.LongList, MiddleOperator.PopFloat);

}

The type of the resulting expression is a pointer to the type of the original expression

Type pointerType = new Type(expression.Symbol.Type);

Symbol resultSymbol = new Symbol(pointerType);

AddMiddleCode(expression.LongList, MiddleOperator.Address,

resultSymbol, expression.Symbol);

return (new Expression(resultSymbol, expression.ShortList,

expression.LongList));

}

### Dereference Expression

The dereference operator takes a pointer or array as operand.

MainParser.gppg

prefix\_expression:

ASTERRISK type\_cast\_expression {

$$ = MiddleCodeGenerator.DereferenceExpression($2);

};

The dereference operator applies to pointers, arrays, strings, and functions. In the function case, the expression is regarded as a pointer to a function.

MiddleCodeGenerator.cs

public static Expression DereferenceExpression(Expression expression) {

Assert.Error(expression.Symbol.Type.IsPointerArrayStringOrFunction(),

Message.Invalid\_dereference\_of\_non\_\_pointer);

bool assignable =

!expression.Symbol.Type.PointerOrArrayType.IsConstantRecursive() &&

!expression.Symbol.Type.PointerOrArrayType.IsArrayOrFunction(),

addressable =

!expression.Symbol.Type.PointerOrArrayType.IsBitfield();

Symbol resultSymbol =

new Symbol(expression.Symbol.Type.PointerOrArrayType,

assignable, addressable);

return Dereference(expression, resultSymbol, 0);

}

### Prefix Increment and Decrement Expression

The postfix increment and decrement change the operand and return the original value. The operand has to be a left-value of arithmetic or pointer type.

The increment and decrement operators come in two forms: prefix and postfix. The difference is that in the prefix case the resulting value is value after the operator has been applied, while in the postfix case the resulting value is the original value before the operator has been applied. However, the side effects of the operators are the same: in the increment case the value is added by one, and in the decrement case the value is subtracted by one.

MainParser.gppg

prefix\_expression:

increment\_operator prefix\_expression {

$$ = MiddleCodeGenerator.PrefixIncrementExpression($1, $2);

};

increment\_operator:

INCREMENT { $$ = MiddleOperator.Increment; }

| DECREMENT { $$ = MiddleOperator.Decrement; };

MiddleCodeGenerator.cs

private static IDictionary<MiddleOperator,MiddleOperator> m\_incrementMap =

new Dictionary<MiddleOperator, MiddleOperator>();

private static IDictionary<MiddleOperator,MiddleOperator>

m\_incrementInverseMap = new Dictionary<MiddleOperator,MiddleOperator>();

static MiddleCodeGenerator() {

m\_incrementMap.Add(MiddleOperator.Increment,

MiddleOperator.BinaryAdd);

m\_incrementMap.Add(MiddleOperator.Decrement,

MiddleOperator.BinarySubtract);

m\_incrementInverseMap.Add(MiddleOperator.Increment,

MiddleOperator.BinarySubtract);

m\_incrementInverseMap.Add(MiddleOperator.Decrement,

MiddleOperator.BinaryAdd);

}

public static Expression PrefixIncrementExpression

(MiddleOperator middleOp, Expression expression){

Symbol symbol = expression.Symbol;

Assert.Error(symbol.Assignable, Message.Not\_assignable);

Assert.Error(symbol.Type.IsArithmeticOrPointer(),

expression, Message.Invalid\_type\_in\_increment\_expression);

if (symbol.Type.IsIntegralOrPointer()) {

AddMiddleCode(expression.ShortList, middleOp, null, symbol);

AddMiddleCode(expression.LongList, middleOp, null, symbol);

BigInteger? bitFieldMask = symbol.Type.BitfieldMask();

if (bitFieldMask != null) {

Symbol maskSymbol = new Symbol(symbol.Type, bitFieldMask.Value);

MiddleCode maskCode = new MiddleCode(MiddleOperator.BitwiseAnd,

symbol, symbol, maskSymbol);

expression.ShortList.Add(maskCode);

expression.LongList.Add(maskCode);

}

Symbol resultSymbol = new Symbol(symbol.Type);

AddMiddleCode(expression.LongList, MiddleOperator.Assign,

resultSymbol, symbol);

return (new Expression(resultSymbol, expression.ShortList,

expression.LongList));

}

The increment and decrement operator apply not only to integral values, but also to floating values. The operation. We start by pushing the value one at the floating value stack, the value to be incremented or decremented has already been pushed at the stack by earlier operations. We perform the operation, which is addition or subtraction. We preform the operation on both the short list and list of the expression. The difference is that in the short list case we pop the value off the stack since we do not need it anymore. In the long list case we do nothing, we just let the value stay on the stack to be used by later operations.

else {

AddMiddleCode(expression.ShortList, MiddleOperator.PushOne);

Symbol oneSymbol = new Symbol(symbol.Type, (decimal) 1);

AddMiddleCode(expression.ShortList, m\_incrementMap[middleOp],

symbol, symbol, oneSymbol);

AddMiddleCode(expression.ShortList, MiddleOperator.PopFloat, symbol);

AddMiddleCode(expression.LongList, MiddleOperator.PushOne);

AddMiddleCode(expression.LongList, m\_incrementMap[middleOp],

symbol, symbol, oneSymbol);

AddMiddleCode(expression.LongList, MiddleOperator.TopFloat, symbol);

Symbol resultSymbol = new Symbol(symbol.Type);

return (new Expression(resultSymbol, expression.ShortList,

expression.LongList));

}

}

### Postfix Increment and Decrement Expression

MainParser.gppg

postfix\_expression:

primary\_expression {

$$ = $1;

}

| postfix\_expression increment\_operator {

$$ = MiddleCodeGenerator.PostfixIncrementExpression($2, $1);

};

MiddleCodeGenerator.cs

public static Expression PostfixIncrementExpression

(MiddleOperator middleOp, Expression expression) {

Symbol symbol = expression.Symbol;

Assert.Error(symbol.Assignable, Message.Not\_assignable);

Assert.Error(symbol.Type.IsArithmeticOrPointer(),

expression, Message.Invalid\_type\_in\_increment\_expression);

if (symbol.Type.IsIntegralOrPointer()) {

Symbol resultSymbol = new Symbol(symbol.Type);

AddMiddleCode(expression.LongList, MiddleOperator.Assign,

resultSymbol, symbol);

AddMiddleCode(expression.ShortList, middleOp, null, symbol);

AddMiddleCode(expression.LongList, middleOp, null, symbol);

BigInteger? bitFieldMask = symbol.Type.BitfieldMask();

if (bitFieldMask != null) {

Symbol maskSymbol = new Symbol(symbol.Type, bitFieldMask.Value);

AddMiddleCode(expression.ShortList, MiddleOperator.BitwiseAnd,

symbol, symbol, maskSymbol);

AddMiddleCode(expression.LongList, MiddleOperator.BitwiseAnd,

symbol, symbol, maskSymbol);

}

return (new Expression(resultSymbol, expression.ShortList,

expression.LongList));

}

else {

AddMiddleCode(expression.ShortList, MiddleOperator.PushOne);

Symbol oneSymbol = new Symbol(symbol.Type, (decimal) 1);

AddMiddleCode(expression.ShortList, m\_incrementMap[middleOp],

symbol, symbol, oneSymbol);

AddMiddleCode(expression.ShortList, MiddleOperator.PopFloat, symbol);

In the long list case, the situation becomes a little more complicated, since the result of the operations shall be the original value, not the resulting value. Therefore, we must perform the inverse operation on the value on the stack in order to return to the original value.

AddMiddleCode(expression.LongList, MiddleOperator.PushOne);

AddMiddleCode(expression.LongList, m\_incrementMap[middleOp],

symbol, symbol, oneSymbol);

AddMiddleCode(expression.ShortList, MiddleOperator.TopFloat, symbol);

AddMiddleCode(expression.LongList, MiddleOperator.PushOne);

AddMiddleCode(expression.LongList, m\_incrementInverseMap[middleOp],

symbol, symbol, oneSymbol);

Symbol resultSymbol = new Symbol(symbol.Type);

return (new Expression(resultSymbol, expression.ShortList,

expression.LongList));

}

}

### Arrow Expression

The arrow operator takes a pointer to a struct or union as operand. The result symbol has the operand as address symbol and the member offset as offset.

MainParser.gppg

postfix\_expression:

postfix\_expression ARROW NAME {

$$ = MiddleCodeGenerator.ArrowExpression($1, $3);

};

MiddleCodeGenerator.cs

public static Expression ArrowExpression(Expression expression,

string memberName) {

Assert.Error(expression.Symbol.Type.IsPointer() &&

expression.Symbol.Type.PointerType.IsStructOrUnion(),

expression,

Message.Not\_a\_pointer\_to\_a\_struct\_or\_union\_in\_arrow\_expression);

Symbol memberSymbol =

expression.Symbol.Type.PointerType.MemberMap[memberName];

Assert.Error(memberSymbol != null, memberName,

Message.Unknown\_member\_in\_arrow\_expression);

bool assignable = !expression.Symbol.Type.PointerOrArrayType.IsConstant

&& !memberSymbol.Type.IsConstantRecursive() &&

!memberSymbol.Type.IsArrayOrFunction(),

addressable = !memberSymbol.Type.IsBitfield();

Symbol resultSymbol =

new Symbol(memberSymbol.Type, assignable, addressable);

return Dereference(expression, resultSymbol, memberSymbol.Offset);

}

### Index Expression

The index operator takes a pointer or array and an integral. We have three cases:

1. The index is constant, in which case we treat the operator as if it was the arrow operator.
2. The index is not constant and the array or pointer type size is one, in which case we generate that add the index to the pointer or array.
3. The index is not constant and the array or pointer type size is greater than, in which case we also need to generate code that multiply the index with the type size.

MainParser.gppg

postfix\_expression:

postfix\_expression LEFT\_SQUARE expression RIGHT\_SQUARE {

$$ = MiddleCodeGenerator.IndexExpression($1, $3);

};

MiddleCodeGenerator.cs

public static Expression IndexExpression(Expression leftExpression,

Expression rightExpression) {

Expression staticExpression =

StaticExpression.Binary(MiddleOperator.Index, leftExpression,

rightExpression);

if (staticExpression != null) {

return staticExpression;

}

Expression arrayExpression, indexExpression;

if (leftExpression.Symbol.Type.IsPointerOrArray()) {

arrayExpression = leftExpression;

indexExpression = rightExpression;

}

else {

indexExpression = leftExpression;

arrayExpression = rightExpression;

}

Type arrayType = arrayExpression.Symbol.Type,

indexType = indexExpression.Symbol.Type;

Assert.Error(arrayType.IsPointerOrArray() &&

!arrayType.PointerOrArrayType.IsVoid(),

arrayType, Message.Invalid\_type\_in\_index\_expression);

Assert.Error(indexType.IsIntegral(), indexExpression,

Message.Invalid\_type\_in\_index\_expression);

List<MiddleCode> shortList = new List<MiddleCode>();

shortList.AddRange(arrayExpression.ShortList);

shortList.AddRange(indexExpression.ShortList);

bool assignable =

!arrayExpression.Symbol.Type.PointerOrArrayType.IsConstantRecursive()

&& !arrayExpression.Symbol.Type.PointerOrArrayType.IsArrayOrFunction();

bool addressable =

!arrayExpression.Symbol.Type.PointerOrArrayType.IsBitfield();

Symbol resultSymbol =

new Symbol(arrayExpression.Symbol.Type.PointerOrArrayType,

assignable, addressable);

if (indexExpression.Symbol.Value is BigInteger) {

int indexValue = (int) ((BigInteger)indexExpression.Symbol.Value),

indexSize = arrayExpression.Symbol.Type.PointerOrArrayType.Size();

return Dereference(arrayExpression, resultSymbol,

indexValue \* indexSize);

}

else {

indexExpression =

TypeCast.ImplicitCast(indexExpression, arrayExpression.Symbol.Type);

List<MiddleCode> longList = new List<MiddleCode>();

longList.AddRange(arrayExpression.LongList);

longList.AddRange(indexExpression.LongList);

Symbol arraySymbol = arrayExpression.Symbol,

indexSymbol = indexExpression.Symbol;

Symbol sizeSymbol = new Symbol(arraySymbol.Type, (BigInteger)

arraySymbol.Type.PointerOrArrayType.Size()),

multSymbol = new Symbol(arraySymbol.Type);

AddMiddleCode(longList, MiddleOperator.UnsignedMultiply,

multSymbol, indexSymbol, sizeSymbol);

Symbol addSymbol = new Symbol(arraySymbol.Type);

AddMiddleCode(longList, MiddleOperator.BinaryAdd,

addSymbol, arraySymbol, multSymbol);

Expression addExpression =

new Expression(addSymbol, shortList, longList);

return Dereference(addExpression, resultSymbol, 0);

}

}

private static Expression Dereference(Expression expression,

Symbol resultSymbol, int offset) {

resultSymbol.AddressSymbol = expression.Symbol;

resultSymbol.AddressOffset = offset;

AddMiddleCode(expression.LongList, MiddleOperator.Dereference,

resultSymbol, expression.Symbol, 0);

If the resulting expression holds floating type, we need to pop the value from the floating value stack.

if (resultSymbol.Type.IsFloating()) {

AddMiddleCode(expression.LongList, MiddleOperator.PushFloat,

resultSymbol);

}

return (new Expression(resultSymbol, expression.ShortList,

expression.LongList));

}

### Dot Expression

The dot operator generates different result is the operand has and address symbol. If it has, it is the result of the arrow, dereference, or index operator. In that case, it inherits the address symbol and increase the address offset with the member offset.

If the operand does not have an address symbol, it is a regular variable. Then we just create a temporary symbol with the same storage and increase the offset with the member offset.

MainParser.gppg

postfix\_expression:

postfix\_expression DOT NAME {

$$ = MiddleCodeGenerator.DotExpression($1, $3);

};

MiddleCodeGenerator.cs

public static Expression DotExpression(Expression expression,

string memberName) {

Symbol parentSymbol = expression.Symbol;

Assert.Error(parentSymbol.Type.IsStructOrUnion(), expression,

Message.Not\_a\_struct\_or\_union\_in\_dot\_expression);

Symbol memberSymbol = parentSymbol.Type.MemberMap[memberName];

Assert.Error(memberSymbol != null, memberName,

Message.Unknown\_member\_in\_dot\_expression);

Symbol resultSymbol;

if (parentSymbol.AddressSymbol != null) {

string name = parentSymbol.Name + "." + memberSymbol.Name +

Symbol.SeparatorId + memberSymbol.Offset;

resultSymbol = new Symbol(name, parentSymbol.ExternalLinkage,

parentSymbol.Storage, memberSymbol.Type,

parentSymbol.IsParameter());

resultSymbol.UniqueName = parentSymbol.UniqueName;

resultSymbol.AddressSymbol = parentSymbol.AddressSymbol;

resultSymbol.AddressOffset = parentSymbol.AddressOffset;

resultSymbol.Offset = parentSymbol.Offset + memberSymbol.Offset;

resultSymbol.Assignable = !parentSymbol.Type.IsConstant &&

!memberSymbol.Type.IsConstantRecursive() &&

!memberSymbol.Type.IsArrayOrFunction();

resultSymbol.Addressable = !parentSymbol.IsRegister() &&

!memberSymbol.Type.IsBitfield();

}

else {

bool assignable = !parentSymbol.Type.IsConstant &&

!memberSymbol.Type.IsConstantRecursive() &&

!memberSymbol.Type.IsArrayOrFunction(),

addressable = !parentSymbol.IsRegister() &&

!memberSymbol.Type.IsBitfield();

resultSymbol = new Symbol(memberSymbol.Type, assignable, addressable);

resultSymbol.Name = parentSymbol.Name + Symbol.SeparatorId +

memberName;

resultSymbol.UniqueName = parentSymbol.UniqueName;

resultSymbol.Storage = parentSymbol.Storage;

resultSymbol.Offset = parentSymbol.Offset + memberSymbol.Offset;

}

return (new Expression(resultSymbol, expression.ShortList,

expression.LongList));

}

### Function Call Expression

When calling a function, first we need to check the number of actual and formal parameters, which shall be equal. Unless the function is elliptic, in which case the actual parameter list can be longer.

Each regular actual parameter is transformed to the type of the formal parameter, and each extra actual parameter is subjected to the type promotion described in Section 6.3.6. Unless the return type is void, we place the result value in a temporary symbol.

Each actual parameter is checked. For a regular parameter (paramType is not null), a function is converted to a pointer to a function, otherwise implicit type conversion occurs. For an extra parameter, a function, array, or string is converted to a pointer, otherwise argument promotion occurs.

For an extra parameter: a char is converted to an integer and a float is converted to a double.

MainParser.gppg

postfix\_expression:

postfix\_expression {

MiddleCodeGenerator.FunctionPreCall($1);

}

LEFT\_PARENTHESIS optional\_argument\_expression\_list RIGHT\_PARENTHESIS {

$$ = MiddleCodeGenerator.CallExpression($1, $4);

};

The FunctionPreCall method is called before the argument list is parsed. It checks that the expression is either a function or a pointer to a function, adds the parameter type list of the function to the type list stack, adds zero to the current offset stack, and adds the call header instruction to the short and long list of the expression.

MiddleCodeGenerator.cs

public static void FunctionPreCall(Expression expression) {

Type type = expression.Symbol.Type;

Assert.Error((type.IsPointer() && type.PointerType.IsFunction()) ||

type.IsFunction(), expression.Symbol,

Message.Not\_a\_function);

Type functionType = type.IsFunction() ? type : type.PointerType;

TypeListStack.Push(functionType.TypeList);

ParameterOffsetStack.Push(0);

AddMiddleCode(expression.ShortList, MiddleOperator.FunctionPreCall,

SymbolTable.CurrentTable.CurrentOffset);

AddMiddleCode(expression.LongList, MiddleOperator.FunctionPreCall,

SymbolTable.CurrentTable.CurrentOffset);

}

The CallExpression method is called after the argument list is parsed. It checks that the expression is either a function or a pointer to a function, adds the parameter type list of the function to the type list stack, adds zero to the current offset stack, and adds the call header instruction to the short and long list of the expression.

public static Expression CallExpression(Expression functionExpression,

List<Expression> argumentList){

List<Type> typeList = TypeListStack.Pop();

ParameterOffsetStack.Pop();

The reason we keep track of the current offset is that we in case of nested function calls need to allocate space for the previous arguments of the functions call outside the nested function. If the current offset is more than zero, we add it together with the size of the standard function top the total offset. If the current offset is zero, no argument has yet been passed, and we do not need to increase the total offset.

int totalOffset = 0;

foreach (int currentOffset in ParameterOffsetStack) {

if (currentOffset > 0) {

totalOffset += (SymbolTable.FunctionHeaderSize + currentOffset);

}

}

Type functionType = functionExpression.Symbol.Type.IsPointer() ?

functionExpression.Symbol.Type.PointerType :

functionExpression.Symbol.Type;

We check that there are not too few parameters in the call. If the type list is null, the number of arguments does not matter. However, if the type list is not null, the number of arguments must be at least the number of parameters in the type list.

Assert.Error((typeList == null) ||

(argumentList.Count >= typeList.Count),

functionExpression,

Message.Too\_few\_actual\_parameters\_in\_function\_call);

We also check that there are not too many parameters in the call. If the type list is null or if the callee function has an elliptic parameter list, the number of arguments does not matter. There may be more arguments than parameters, we already know that the arguments are not fewer than the parameters. However, if the type list is not null and the function does not have an elliptic parameter list, the number of arguments must equal the number of parameters.

Assert.Error(functionType.IsEllipse() || (typeList == null) ||

(argumentList.Count == typeList.Count),

functionExpression,

Message.Too\_many\_parameters\_in\_function\_call);

List<MiddleCode> longList = new List<MiddleCode>();

longList.AddRange(functionExpression.LongList);

When iterate through the, we keep track of the extra size. In elliptic function calls we need to allocate memory for the extra parameters that is not included in the regular parameter list.

int count = 0, offset = SymbolTable.FunctionHeaderSize, extra = 0;

foreach (Expression argumentExpression in argumentList) {

longList.AddRange(argumentExpression.LongList);

Type type;

if ((typeList != null) && (count++ < typeList.Count)) {

type = typeList[count];

}

else {

type = ParameterType(argumentExpression.Symbol);

extra += type.Size();

}

We add the argument to the parameter list and update the current offset.

AddMiddleCode(longList, MiddleOperator.Parameter, type,

argumentExpression.Symbol, SymbolTable.CurrentTable.

CurrentOffset + totalOffset + offset);

offset += type.Size();

}

We add both the function call and post call instructions to the long list. The post call instruction is for

Symbol functionSymbol = functionExpression.Symbol;

AddMiddleCode(longList, MiddleOperator.Call, functionSymbol,

SymbolTable.CurrentTable.CurrentOffset + totalOffset,

extra);

AddMiddleCode(longList, MiddleOperator.PostCall, null, null,

SymbolTable.CurrentTable.CurrentOffset + totalOffset);

Type returnType = functionType.ReturnType;

Symbol returnSymbol = new Symbol(returnType);

We make the short list be a copy of the long list. However, there is one difference between the lists if the function returns a floating value. In that case, we pop the value off the floating value stack in the short list, since the value is not interesting in the short list, only the side effect of the function call. However, we keep the value on the stack in the long list, since the value shall be used in the long list.

List<MiddleCode> shortList = new List<MiddleCode>();

shortList.AddRange(longList);

if (!returnType.IsVoid()) {

if (returnType.IsStructOrUnion()) {

Type pointerType = new Type(returnType);

Symbol addressSymbol = new Symbol(pointerType);

returnSymbol.AddressSymbol = addressSymbol;

}

AddMiddleCode(longList, MiddleOperator.GetReturnValue,

returnSymbol);

if (returnType.IsFloating()) {

AddMiddleCode(shortList, MiddleOperator.PopEmpty);

}

}

return (new Expression(returnSymbol, shortList, longList));

}

### Argument Expression List

Each parameter is converted from a possible logical type to an integer by calling Generate.generateIntegerExpression.

MainParser.gppg

argument\_expression\_list:

assignment\_expression {

$$ = new List<Expression>();

$$.Add(MiddleCodeGenerator.ArgumentExpression(0, $1));

}

| argument\_expression\_list COMMA assignment\_expression {

$1.Add(MiddleCodeGenerator.ArgumentExpression($1.Count, $3));

$$ = $1;

};

The ArgumentExpression method type casts the argument into an appropriate type.

MiddleCodeGenerator.cs

public static Expression ArgumentExpression(int index,

Expression expression) {

We need to type list on top of the type list stack to compare the argument expression with the parameter type.

List<Type> typeList = TypeListStack.Peek();

If the argument is within the parameter list, we type cast the argument into the parameter type.

if ((typeList != null) && (index < typeList.Count)) {

expression = TypeCast.ImplicitCast(expression, typeList[index]);

}

else {

If the argument is not within the parameter list (elliptic call), we type cast the argument into an appropriate type: character and short integer are cast into (signed or unsigned) integer, and float is cast into double. These type casts apply to the elliptic function call rules of ANSI C.

Type type = expression.Symbol.Type;

if (type.IsChar() || type.IsShort()) {

if (type.IsSigned()) {

expression =

TypeCast.ImplicitCast(expression, Type.SignedIntegerType);

}

else {

expression =

TypeCast.ImplicitCast(expression, Type.UnsignedIntegerType);

}

}

else if (type.IsFloat()) {

expression = TypeCast.ImplicitCast(expression, Type.DoubleType);

}

}

We need to update current offset to allocate space in case of nested function calls.

int offset = ParameterOffsetStack.Pop();

ParameterOffsetStack.Push(offset +

ParameterType(expression.Symbol).Size());

return (new Expression(expression.Symbol, expression.LongList,

expression.LongList));

}

The ParameterType method returns the type of the parameter. In case of array, string, or function, it returns pointer types. Otherwise, it returns the regular type.

private static Type ParameterType(Symbol symbol) {

switch (symbol.Type.Sort) {

case Sort.Array:

return (new Type(symbol.Type.ArrayType));

case Sort.Function:

return (new Type(symbol.Type));

case Sort.String:

return (new Type(new Type(Sort.Signed\_Char)));

default:

return symbol.Type;

}

}

### Primary Expressions

We have finally reached the last kind of expression. A primary expression may be another expression within parentheses, a value, a symbol, a register, or the carry flag.

A primary expression is an identifier, which symbol we look up in the symbol table, a value, which we store in the global space, or an expression surrounded by parenthesis.

MainParser.gppg

primary\_expression:

LEFT\_PARENTHESIS expression RIGHT\_PARENTHESIS {

$$ = $2;

}

VALUE {

$$ = MiddleCodeGenerator.ValueExpression($1);

};

The value holds floating type, we push it at the floating value stack.

MiddleCodeGenerator.cs

public static Expression ValueExpression(Symbol symbol) {

List<MiddleCode> longList = new List<MiddleCode>();

if (symbol.Type.IsFloating()) {

AddMiddleCode(longList, MiddleOperator.PushFloat, symbol);

}

return (new Expression(symbol, new List<MiddleCode>(), longList));

}

MainParser.gppg

primary\_expression:

NAME {

$$ = MiddleCodeGenerator.SymbolExpression($1);

};

Given a name, we look it up in the current symbol table. If there is no symbol, we add a function without a parameter list, that returns a signed integer, to the symbol table.

MiddleCodeGenerator.cs

public static Expression SymbolExpression(string name) {

Symbol symbol = SymbolTable.CurrentTable.LookupSymbol(name);

if (symbol == null) {

Type type = new Type(Type.SignedIntegerType, null, false);

symbol = new Symbol(name, Storage.Extern, type);

}

List<MiddleCode> shortList = new List<MiddleCode>(),

longList = new List<MiddleCode>();

if (symbol.Type.IsFloating()) {

AddMiddleCode(shortList, MiddleOperator.PushFloat, symbol);

AddMiddleCode(longList, MiddleOperator.PushFloat, symbol);

}

return (new Expression(symbol, shortList, longList));

}

The registers are only used internally, in conjunction with system calls. On some occasion the interrupt call returns information stored in a register.

The closed statements do also include a set of internal statements; that is, statements not included in standard C, but necessary for the standard library functionality. The first statement is the load register statement, which stores a value in a register. The operands are the name of a register and an expression of integral or pointer type with same size as the register.

MainParser.gppg

primary\_expression:

REGISTER\_NAME {

$$ = MiddleCodeGenerator.RegisterExpression($1);

};

MiddleCodeGenerator.cs

public static Expression RegisterExpression(Register register) {

List<MiddleCode> longList = new List<MiddleCode>();

int size = AssemblyCode.SizeOfRegister(register);

Type type = TypeSize.SizeToUnsignedType(size);

Symbol symbol = new Symbol(type);

AddMiddleCode(longList, MiddleOperator.InspectRegister,

symbol, register);

return (new Expression(symbol, new List<MiddleCode>(),

longList, register));

}

The carry flag is also only used internally in conjunction with system calls. On some occasion the interrupt call returns information stored in the carry flag.

MainParser.gppg

primary\_expression:

CARRY\_FLAG {

$$ = MiddleCodeGenerator.CarryFlagExpression();

};

MiddleCodeGenerator.cs

public static Expression CarryFlagExpression() {

ISet<MiddleCode> trueSet = new HashSet<MiddleCode>(),

falseSet = new HashSet<MiddleCode>();

List<MiddleCode> longList = new List<MiddleCode>();

trueSet.Add(AddMiddleCode(longList, MiddleOperator.Carry));

falseSet.Add(AddMiddleCode(longList, MiddleOperator.Goto));

Symbol symbol = new Symbol(trueSet, falseSet);

return (new Expression(symbol, new List<MiddleCode>(), longList));

}

MainParser.gppg

# The Middle Code

The middle of the compiler in this book is three-address code. As the name implies, each instruction is made up by one operator and three addresses of the Object class, which can be register, symbols or jump addresses.

There is also the fields m\_visited and m\_leader, which is used by MiddleCodeOptimizer to search for non-reachable instruction and when creating base block (base blocks are described in Section 14.2.1). The m\_liveSet is generated by MiddleCode­Optimizer and used by ObjectCodeGenerator to select suitable object code instructions, it holds the live symbols of the base block; that is, the symbols that will be accessed later in the base block.

InverseMap holds the inverses of the equality and relation operators and is used when optimizing conditional jump instructions in MiddleCodeOptimizer.

The toString method is not strictly necessary; nevertheless, it called when generating the assembler code and added as comments.

MiddleCode.cs

# Initializer

In C, it is possible to initialize simple and compound variables. Therefore, we need to check that the initialized value has the correct type. Basically, there are two kinds of initialization: static and stack. Static initialization results in a block of bytes, and stack initialization results in a sequence of assignment middle code instructions.

### Static Initialization

Static initialization occurs when a static variable becomes initialized. The initialization value has to be constant and known at compile time. No code is generated, instead a memory block holding the value is created. There are several cases to consider:

|  |  |  |
| --- | --- | --- |
| **Variable** | **Initializer** | **Example** |
| Pointer | Address | static int array[3];  static int\* p = &a[2]; |
| Pointer to signed or unsigned char | String | static char \*p = “Hello”; |
| Pointer | List | int \*p = {1,2,3}; |
| Array | String | char s[] = “World”; |
| Array | List | Char st[] = {‘a’, ’b’, ‘c’}; |
| Struct | List |  |
| Union | List |  |
| Integral or Pointer |  |  |
| Floating |  |  |

GenerateInitializer.java

If the variable is a pointer and the initializer is an address we add its offset to the block and store the name of the address in the access map, the address value will later be looked up and added by the linker.

If the variable type is a pointer (signed or unsigned) char and the initializator is a string, the address is stored in the access map and a zero address is stored in the block, it will later be properly looked up and set by the linker.

### Stack Initialization

Stack initialization occurs when an auto or register variable becomes initialized. Since the initialization value can be non-constant, no memory block is created. Instead, a sequence of assignment instructions is added to the middle code.

# Declaration Specifiers and Declarators

The Mask class is used to distinguish between the different type specifiers and is used by the Specifier class below.

Mask.cs

namespace CCompiler {

public enum Mask {StorageMask = 0x0000FFF,

Auto = 0x0000010,

Register = 0x0000020,

Static = 0x0000040,

Extern = 0x0000080,

Typedef = 0x0000100,

QualifierMask = 0x000F000,

Constant = 0x0001000,

Volatile = 0x0002000,

SortMask = 0xFFF0000,

Signed = 0x0010000,

Unsigned = 0x0020000,

Char = 0x0040000,

Short = 0x0100000,

Int = 0x0200000,

Long = 0x0400000,

Float = 0x0800000,

Double = 0x1000000,

Void = 0x2000000,

SignedChar = Signed | Char,

UnsignedChar = Unsigned | Char,

ShortInt = Short | Int,

SignedShort = Signed | Short,

SignedShortInt = Signed | Short | Int,

UnsignedShort = Unsigned | Short,

UnsignedShortInt = Unsigned | Short | Int,

SignedInt = Signed | Int,

UnsignedInt = Unsigned | Int,

LongInt = Long | Int,

SignedLong = Signed | Long,

SignedLongInt = Signed | Long | Int,

UnsignedLong = Unsigned | Long,

UnsignedLongInt = Unsigned | Long | Int,

LongDouble = Long | Double};

}

The Specifier class is used to generate the type specified by a list of specifiers. The declaration specifiers are made up by keywords or struct, union, or enumeration declarations. First, we need the maskToStorageMap and maskToSortMap to map between specifiers and storage and sorts.

Specifier.cs

using System;

using System.Text;

using System.Collections.Generic;

namespace CCompiler {

public class Specifier {

The specifier decides the storage and type of the symbol, it also decide whether the symbol has external linkage.

private bool m\_externalLinkage;

private Storage? m\_storage;

private Type m\_type;

The mask-to-sort map maps the masks of the Mask enumeration to the types of the Sort enumeration. However, the masks are represented by unsigned integer values since we use unsigned integer values to mask the keywords of the declaration specifiers in the SpecifierList method below.

private static IDictionary<int, Sort> m\_maskToSortMap =

new Dictionary<int, Sort>() {

{(int) Mask.Void, Sort.Void},

{(int) Mask.Char, Sort.Signed\_Char},

{(int) Mask.SignedChar, Sort.Signed\_Char},

{(int) Mask.UnsignedChar, Sort.Unsigned\_Char},

{(int) Mask.Short, Sort.Signed\_Short\_Int},

{(int) Mask.ShortInt, Sort.Signed\_Short\_Int},

{(int) Mask.SignedShort, Sort.Signed\_Short\_Int},

{(int) Mask.SignedShortInt, Sort.Signed\_Short\_Int},

{(int) Mask.UnsignedShort, Sort.Unsigned\_Short\_Int},

{(int) Mask.UnsignedShortInt, Sort.Unsigned\_Short\_Int},

{(int) Mask.Int, Sort.Signed\_Int},

{(int) Mask.Signed, Sort.Signed\_Int},

{(int) Mask.SignedInt, Sort.Signed\_Int},

{(int) Mask.Unsigned, Sort.Unsigned\_Int},

{(int) Mask.UnsignedInt, Sort.Unsigned\_Int},

{(int) Mask.Long, Sort.Signed\_Long\_Int},

{(int) Mask.LongInt, Sort.Signed\_Long\_Int},

{(int) Mask.SignedLong, Sort.Signed\_Long\_Int},

{(int) Mask.SignedLongInt, Sort.Signed\_Long\_Int},

{(int) Mask.UnsignedLong, Sort.Unsigned\_Long\_Int},

{(int) Mask.UnsignedLongInt, Sort.Unsigned\_Long\_Int},

{(int) Mask.Float, Sort.Float},

{(int) Mask.Double, Sort.Double},

{(int) Mask.LongDouble, Sort.Long\_Double}

};

A specifier is made up by the external linkage, storage, and type.

public Specifier(bool externalLinkage, Storage? storage, Type type) {

m\_externalLinkage = externalLinkage;

m\_storage = storage;

m\_type = type;

}

public bool ExternalLinkage {

get { return m\_externalLinkage; }

}

public CCompiler.Storage? StorageX {

get { return m\_storage; }

}

public Type Type {

get { return m\_type; }

}

The static SpecifierList method takes a list of declaration specifiers and returns a Specifier object, holding the storage, type, and external linkage of the symbol.

public static Specifier SpecifierList(List<object> specifierList) {

int totalMaskValue = 0;

Type compoundType = null;

We iterate through the declaration specifier list and gather the keywords of the Mask enumeration. If a keyword occurs twice, we report an error.

foreach (object obj in specifierList) {

if (obj is Mask) {

int maskValue = (int) obj;

if ((maskValue & totalMaskValue) != 0) {

Assert.Error(MaskToString(maskValue),

Message.Keyword\_defined\_twice);

}

totalMaskValue |= maskValue;

}

If the element of the list is not a Mask enumeration, it must be a type. If more than one type occurs in the list, we report an error.

else {

if (compoundType != null) {

Assert.Error(MaskToString(totalMaskValue),

Message.Invalid\_specifier\_sequence);

}

compoundType = (Type) obj;

}

}

When we have the final mask, we begin by extracting the storage. The storage is initially set to null, and we mask out the storage part of the total mask. It is possible that there was no storage in the specifier list. However, if there is a storage it must match one of the storage masks. If it does not match, more than one storage specifiers were present in the declaration specifier list, and we report an error.

Storage? storage = null;

{ int totalStorageValue = totalMaskValue & ((int) Mask.StorageMask);

if (totalStorageValue != 0) {

if (!Enum.IsDefined(typeof(Mask), totalStorageValue)) {

Assert.Error(MaskToString(totalStorageValue),

Message.Invalid\_specifier\_sequence);

}

storage = (Storage) totalStorageValue;

}

}

The next step is to determine the external linkage. The symbol has externa linkage if the scope of the current symbol table is global (the symbol is defined in global space) and if the symbol is not a function parameter. The CallDepth is increased for each function parameter definition, so it is zero if there is no function definition. Moreover, the storage shall be null (is has not been stated in the declaration specifier list) or extern. If the symbol is static in global scope it has not external linkage.

bool externalLinkage = (SymbolTable.CurrentTable.Scope == Scope.Global)

&& (CCompiler\_Main.Parser.CallDepth == 0) &&

((storage == null) ||

(storage == CCompiler.Storage.Extern));

When the storage and external linkage has been taken care of, we perform a series of error checks. If the symbol is a function parameter (the CallDepth variable is more than zero) the storage must be null, auto or register.

if (CCompiler\_Main.Parser.CallDepth > 0) {

Assert.Error((storage == null) || (storage == Storage.Auto) ||

(storage == Storage.Register), storage, Message.

Only\_auto\_or\_register\_storage\_allowed\_in\_parameter\_declaration);

}

If the scope of the current symbol table is a struct or union, the storage must also be null, auto, or register.

else if ((SymbolTable.CurrentTable.Scope == Scope.Struct) ||

(SymbolTable.CurrentTable.Scope == Scope.Union)) {

Assert.Error((storage == null) || (storage == Storage.Auto) ||

(storage == Storage.Register), storage, Message.

Only\_auto\_or\_register\_storage\_allowed\_for\_struct\_or\_union\_scope);

}

If the scope of the current symbol table is global, the storage must also be null, extern, static, or typedef.

else if (SymbolTable.CurrentTable.Scope == Scope.Global) {

Assert.Error((storage == null) || (storage == Storage.Extern) ||

(storage == Storage.Static) ||

(storage == Storage.Typedef), storage, Message.

Only\_extern\_\_\_\_static\_\_\_\_or\_typedef\_storage\_allowed\_in\_global\_scope);

}

If there is a compound type and if it is an enumeration, we need to add its members as signed integers to the symbol table.

if ((compoundType != null) && (compoundType.EnumerationItemSet != null)){

if (storage != null) {

foreach (Pair<Symbol,bool> pair in compoundType.EnumerationItemSet){

Symbol enumSymbol = pair.First;

switch (enumSymbol.Storage = storage.Value) {

case CCompiler.Storage.Static:

SymbolTable.StaticSet.Add(ConstantExpression.

Value(enumSymbol));

break;

case CCompiler.Storage.Extern: {

bool enumInitializer = pair.Second;

Assert.Error(!enumInitializer,

enumSymbol + " = " + enumSymbol.Value,

Message.Extern\_enumeration\_item\_cannot\_be\_initialized);

}

break;

case CCompiler.Storage.Auto:

case CCompiler.Storage.Register:

SymbolTable.CurrentTable.SetOffset(enumSymbol);

break;

}

}

}

else {

foreach (Pair<Symbol,bool> pair in compoundType.EnumerationItemSet) {

Symbol enumSymbol = pair.First;

switch (SymbolTable.CurrentTable.Scope) {

case Scope.Global:

enumSymbol.Storage = CCompiler.Storage.Static;

SymbolTable.StaticSet.Add(ConstantExpression.

Value(enumSymbol));

break;

default:

enumSymbol.Storage = CCompiler.Storage.Auto;

SymbolTable.CurrentTable.SetOffset(enumSymbol);

break;

}

}

}

}

We check if the total mask holds the masks representing the constant or volatile keywords.

{ bool isConstant = (totalMaskValue & ((int) Mask.Constant)) != 0;

bool isVolatile = (totalMaskValue & ((int) Mask.Volatile)) != 0;

If the type is constant or volatile we face a rather complicated situation. XXX

if ((isConstant || isVolatile) && (compoundType != null) &&

compoundType.IsStructOrUnion() /\*&& compoundType.HasTag()\*/) {

compoundType = new Type(compoundType.Sort, compoundType.MemberMap);

}

Finally, we extract the sort mask and test whether it represent a valid type. If the declaration specifier list does not representing a valid combination, we report an error. For instance, signed short double is an invalid combination of declaration specifiers.

Sort? sort = null;

int sortMaskValue = totalMaskValue & ((int) Mask.SortMask);

if (sortMaskValue != 0) {

if (!m\_maskToSortMap.ContainsKey(sortMaskValue)) {

Assert.Error(MaskToString(sortMaskValue),

Message.Invalid\_specifier\_sequence);

}

sort = m\_maskToSortMap[sortMaskValue];

}

When defining the final type, there are four cases. If the compound type is not null and the sort is null, the result type is the compound type, with the possible addition of constant or volatile qualifiers.

Type type = null;

if ((compoundType != null) && (sort == null)) {

compoundType.IsConstant = (compoundType.IsConstant || isConstant);

compoundType.Volatile = (compoundType.Volatile || isVolatile);

type = compoundType;

}

If the compound type is null and the sort is not null, the result type is based on the sort.

else if ((compoundType == null) && (sort != null)) {

type = new Type(sort.Value);

type.IsConstant = isConstant;

type.Volatile = isVolatile;

}

If the compound type and sort are both null, there is no type. This is valid in function definitions, in which case the type shall be a signed integer.

else if ((compoundType == null) && (sort == null)) {

type = new Type(Sort.Signed\_Int);

type.IsConstant = isConstant;

type.Volatile = isVolatile;

}

If both the compound type and the sort are not null, the type specification is invalid. For instance, unsigned short struct {int i;} is an invalid combination.

else {

Assert.Error(MaskToString((int) sortMaskValue), Message.

Invalid\_specifier\_sequence\_together\_with\_type);

}

Finally, the specifier with the external linkage, storage, and type is returned.

return (new Specifier(externalLinkage, storage, type));

}

}

The QualifierList method is called when a pointer is declared. It exams whether the pointer is constant or volatile.

public static Type QualifierList(List<Mask> qualifierList) {

int totalMaskValue = 0;

foreach (Mask mask in qualifierList) {

int maskValue = (int) mask;

if ((maskValue & totalMaskValue) != 0) {

Assert.Error(MaskToString(maskValue),

Message.Keyword\_defined\_twice);

}

totalMaskValue |= maskValue;

}

Type type = Type.VoidPointerType;

type.IsConstant = (totalMaskValue & ((int) Mask.Constant)) != 0;

type.IsVolatile = (totalMaskValue & ((int) Mask.Volatile)) != 0;

return type;

}

The MaskToString method is called when reporting an error. It returns a string holding the declaration specifiers as text.

private static string MaskToString(int totalMaskValue) {

StringBuilder maskBuffer = new StringBuilder();

We iterate throught the bits of the mask and add the text of the Mask enumeration matching each true bit (bit with value one).

for (int maskValue = 1; maskValue != 0; maskValue <<= 1) {

if ((maskValue & totalMaskValue) != 0) {

string maskName = Enum.GetName(typeof(Mask), maskValue).ToLower();

maskBuffer.Append(((maskBuffer.Length > 0) ? " " : "") + maskName);

}

}

return maskBuffer.ToString();

}

}

}

## Declarators

A declarator follows the declaration specifiers. A declarator can be initialized with a value or marked with its size in bits. The bold part of the following code are examples of declarators. Only members of a struct can be marked with bits. Struct or union member as well as function parameters cannot be initialized.

int i, j = 2;

double x = 3.14;

int f(a, b, c);

char s[] = "Hello";

int \*p, \*q = &i;

long int value:7;

A declarator can be a simple variable, an array, or a function with old-style or new-style parameter list.

Declarator.cs

namespace CCompiler {

public class Declarator {

private string m\_name;

private Type m\_firstType, m\_lastType;

public Declarator(string name) {

m\_name = name;

}

public string Name {

get { return m\_name; }

set { m\_name = value; }

}

public Type Type {

get { return m\_firstType; }

}

If we add a type to the declarator where there is no previous type, we set both the first and last type to that type.

public void Add(Type newType) {

if (m\_firstType == null) {

m\_firstType = m\_lastType = newType;

}

else {

switch (m\_lastType.Sort) {

If the last type of the declarator is a pointer, we set it to point at the new type.

case Sort.Pointer:

m\_lastType.PointerType = newType;

m\_lastType = newType;

break;

If the last type of the declarator is an array, we set the type of the array to be the new type. However, the array must be complete; that is, it must have a stated size. Moreover, the new type cannot be a function, since arrays of functions are not allowed.

case Sort.Array:

Assert.Error(newType.IsComplete(),

Message.Array\_of\_incomplete\_type\_not\_allowed);

Assert.Error(!newType.IsFunction(),

Message.Array\_of\_function\_not\_allowed);

m\_lastType.ArrayType = newType;

m\_lastType = newType;

break;

If the last type of the declarator is a function, set the return type of the function to be the new type. However, the new type cannot be an array or a function, since functions are not allowed to return an arrays or functions.

case Sort.Function:

Assert.Error(!newType.IsArray(),

Message.Function\_cannot\_return\_array);

Assert.Error(!newType.IsFunction(),

Message.Function\_cannot\_return\_function);

m\_lastType.ReturnType = newType;

m\_lastType = newType;

break;

}

}

}

}

}

# The Symbol Table

The symbol table holds symbols as well as struct and union tags, but not enumeration tags. Moreover, there is one symbol table for the global space, for each function, for each block in a function, and for each struct or union. In this way, the symbol tables form a hierarchy where each table holds its parent table.

As the name implies, the symbol table is a table holding symbols. It also holds struct and union tags (but not enumeration tags). Moreover, there is one symbol table for the global space, one table for each function, and one table for each block in a function. In this way, the symbol tables form a hierarchy where each table has a reference to its parent table (except the table for the global space, which table parent reference is null), and a set of references to its child tables.

The symbol table keeps track of the values, types, function, and variables of the code, variables defined by the programmer and well as temporary variables introduced by the compiler. The term symbol table is misleading since there actually is a stack of tables, where is table has a reference to the parent table, the parent reference table for the global space table is null. The current table is stored in Main.CurrentTable (see Chapter 20). The symbol table has set of features:

* Table Identity. Each table is given a unique consecutive integer value.
* Scope. A table can the scopes Global, Function, Block, Struct, Union, or Enumeration.
* Parent Table. Each table has apparent table which Main.CurrentTable is set to when the table reaches its end-of-scope.
* Entry List. The entries of the symbol table is stored in the entry list. An entry can be a symbol or a block, represented by another entry list.
* Tag List. Parallel to the entry list, there is also a tag list, holding tags of structs or unions, but not enums.
* Size. The size of the symbol table is the sum of the sizes of all its entries and is of interest when dimensioning the size of a function’s activation record.

Scope.cs

namespace CCompiler {

public enum Scope { Struct = (int) Sort.Struct,

Union = (int) Sort.Union,

Function, Global, Block};

}

SymbolTable.cs

using System;

using System.Collections.Generic;

namespace CCompiler {

public class SymbolTable {

For each function call, an activation record is allocated at the call stack. The first three entries are the return address (the address of the instruction following the function call), the regular frame pointer (the address of the current activation record), and the ellipse frame pointer (the regular frame pointer plus the size of the extra parameters in a call of an elliptic function). The function header size is the sum of the size of those three entries.

public static int ReturnAddressOffset;

public static int RegularFrameOffset;

public static int EllipseFrameOffset;

public static int FunctionHeaderSize;

Each table holds a reference to its parent table, except the table of global space, which table parent reference is null.

private SymbolTable m\_parentTable;

The scope of a table is global, function, block, struct, or union.

private Scope m\_scope;

The activation record is organized in that way that after the three first entries at holding the return address, regular frame pointer, and ellipse frame pointer comes the function parameters, the extra parameters in case of an elliptic case, and variables of auto or register storage, as well as temporary variables. Each one of those symbols holds an offset on the activation record.

private int m\_currentOffset;

All parameters and variables of a function or a block, as well as the members of a struct or a union, are stored in the entry map. Note that it is an object of the ListMap rather than Dictionary, because we need to keep track of the of the struct or union members in case of initialization.

private IDictionary<string,Symbol>m\_entryMap=new ListMap<string,Symbol>();

The tag map holds the types of the structs or unions that are marked with a tag.

private IDictionary<string,Type> m\_tagMap = new Dictionary<string,Type>();

The current table and current function hold references to the table in currently processed by the parser. We need the current table when adding or looking up symbols, and the current function when calling a function or returning a function value. In global scope, the current function is null.

public static SymbolTable CurrentTable = null;

public static Symbol CurrentFunction = null;

The static set holds all static symbols. Note that symbol of static, extern, or typedef storage are not given an offset, since they ar not stored at the activation record.

public static ISet<StaticSymbol> StaticSet = new HashSet<StaticSymbol>();

The static area of the SymbolTable class initializes the offsets of the return address, regular frame pointer, and ellipse frame pointer. As the compiler in this book can be set to generate code for different target machines, the offset can also be given different values as the pointer size may vary.

static SymbolTable() {

ReturnAddressOffset = 0;

RegularFrameOffset = TypeSize.PointerSize;

EllipseFrameOffset = 2 \* TypeSize.PointerSize;

FunctionHeaderSize = 3 \* TypeSize.PointerSize;

}

The constructor of the SymbolTable class takes the parent table and the scope of the table. The parent table is null in global scope.

public SymbolTable(SymbolTable parentTable, Scope scope) {

m\_parentTable = parentTable;

In case of global scope, we create an object of the static set. There is only one table of global scope and one static set of each source code file. In this case we do not need to set the offset field since there are no variables of auto or register storage in global scope.

switch (m\_scope = scope) {

case Scope.Global:

SymbolTable.StaticSet = new HashSet<StaticSymbol>();

break;

In case of struct of union scope, the current offset is initialized to zero.

case Scope.Struct:

case Scope.Union:

m\_currentOffset = 0;

break;

In case of function scope, the current offset is initialized to the size of the header size. The return address, regular frame pointer, and ellipse frame pointer holds the first entries, and the parameters are given the following offsets

case Scope.Function:

m\_currentOffset = FunctionHeaderSize;

break;

In case of block scope, the current offset is initialized to the current offset of its parent table. When the parsing leaves a block the parent table again becomes the current table, and a following block is again initialized to the current offset of the parent table. In this way. Different block can share the same memory space on the activation record.

case Scope.Block:

m\_currentOffset = m\_parentTable.m\_currentOffset;

break;

}

}

The Scope, ParentTable, EntryMap, and CurrentOffset properties returns the scope, parent table, entry map, and current offset of the symbol table.

public Scope Scope {

get { return m\_scope; }

}

public SymbolTable ParentTable {

get { return m\_parentTable; }

}

public IDictionary<string,Symbol> EntryMap {

get { return m\_entryMap; }

}

public int CurrentOffset {

get { return m\_currentOffset; }

}

The AddSymbol method adds a symbol to the symbol table. However, the process is a bit complicated since there is possible to add two symbols with the same name, given that at least one them holds extern storage.

When adding a symbol with a name we have to check a few things:

* If the symbol is extern it cannot be initialized (its value is not-null).
* We look up the entry list, and if we find another symbol with the same name, we check whether at least one of the symbols is extern. If it is, the symbols shall have equals types. If the new symbols is extern, weare finished and we just return. If the old symbol is extern, we remove it and stored the new symbol in the entry list. If not at least one of the symbols is extern, an error message is reported.

It is easier to add or erase a tag. If it is already added an error message is reported. When removing the tag, the tag associated to the name is returned, or null is it not stored in the map.

public void AddSymbol(Symbol newSymbol) {

string name = newSymbol.Name;

The name of the symbol may be null in case of unnamed function parameter.

if (name != null) {

Symbol oldSymbol;

If there is a sembol in the table with the same name, at least one of the4 old and new symbols must hold extern storage. Moreover, they must also have the same type.

if (m\_entryMap.TryGetValue(name, out oldSymbol)) {

Assert.Error(oldSymbol.IsExtern() || newSymbol.IsExtern(),

name, Message.Name\_already\_defined);

Assert.Error(oldSymbol.Type.Equals(newSymbol.Type),

name, Message.Different\_types\_in\_redeclaration);

}

m\_entryMap[name] = newSymbol;

}

If the variable is auto or register, it shall be given an offset in the activation record. Unless in union scope, in which each member always have offset zero.

if (newSymbol.IsAutoOrRegister()) {

if (m\_scope == Scope.Union) {

newSymbol.Offset = 0;

}

If the new symbol is an enumeration constant item, is shall not be given an offset. This is because we do not know at this point if the items really holds auto or register scope. An item isa given auto storage from the beginning, but that may change when the whole enumeration is given its storage, which is done after the item is stored in the symbol table.

else if (!newSymbol.Type.EnumeratorItem) {

newSymbol.Offset = m\_currentOffset;

m\_currentOffset += newSymbol.Type.Size();

}

}

}

The SetOffset method is called just when an enumeration constant item has fianelly been given auto or register storage.

public void SetOffset(Symbol symbol) {

symbol.Offset = m\_currentOffset;

m\_currentOffset += symbol.Type.Size();

}

The LookupSymbol looks up the symbol with the given in the current symbol table, or any of its parent tables, recursively. If the tag is not found, null is returned. Note that this method looks in the parent tables while AddSymbol only checks for symbols with the same name in the current table.

The lookupList method searches the current table for a symbol with the given name, while lookupSymbol searches the table hierarchy recursively upwards by calling lookupList until it finds a symbol with the given name or the parent table is null, which will happen when the table of the global space has been searched.

A tag is recursively looked up in a similar manner.

public Symbol LookupSymbol(string name) {

Symbol symbol;

if (m\_entryMap.TryGetValue(name, out symbol)) {

return symbol;

}

else if (m\_parentTable != null) {

return m\_parentTable.LookupSymbol(name);

}

return null;

}

The AddTag method adds a struct, union, or enumerator tag to the current symbol table. It is possible to add a new struct or union tag (but not an enumeration tag) with an already occupied name. However, in that case, the old and the new tag must have the same sort (they must both be struct or they must both be union) and the member map of at least one of them must be null.

public void AddTag(string name, Type newType) {

if (m\_tagMap.ContainsKey(name)) {

Type oldType = m\_tagMap[name];

Assert.Error(!oldType.IsEnumerator() &&

(oldType.Sort == newType.Sort), name,

Message.Name\_already\_defined);

If the member map of the old tag is null, we assign it the member map of then

if (oldType.MemberMap == null) {

oldType.MemberMap = newType.MemberMap;

}

If the neither the old nor the new member is null, we report an error since it is not allowed top have two structs or unions with the same name tags with non-null member maps.

else {

Assert.Error(newType.MemberMap == null, name,

Message.Name\_already\_defined);

}

}

If there is no struct or union tag previous added to the tag map, we simply add the new tag to the map.

else {

m\_tagMap.Add(name, newType);

}

}

The LookUp method looks up a tag in the current table or its parent tables. If the tag is not found, null is returned.

public Type LookupTag(string name, Sort sort) {

Type type;

if (m\_tagMap.TryGetValue(name, out type) && (type.Sort == sort)) {

return type;

}

else if (m\_parentTable != null) {

return m\_parentTable.LookupTag(name, sort);

}

return null;

}

}

}

## The Symbol

A symbol has extern, static, typedef, auto, or register storage. If the storage is not specifically stated, a symbol becomes static in global scope and auto in a function. If the storage specifier is omitted, a function becomes extern while a global variable becomes static and a local variable becomes auto. In this book, the only difference between an auto or register symbol is that the address operator cannot be applied to a register symbol.

A symbol can hold the storage Typedef, Static, Extern, Auto, and Register. If the storage specifier is omitted, a global variable is regarded as static, and a local variable is regarded as auto. The Register storage has been included for the sake of completeness, but in this book is has no effect. On every occasion, a symbol with the Auto or Register storage are treated equally. Moreover, a symbol has a name, type, and potential value.

Storage.cs

namespace CCompiler {

public enum Storage {Auto = (int) Mask.Auto,

Register = (int) Mask.Register,

Static = (int) Mask.Static,

Extern = (int) Mask.Extern,

Typedef = (int) Mask.Typedef};

}

The Symbol class describes a variable, a value, a function, or a type specified with typedef.

* Name. A symbol actually has two names: the regular name it has been defined with and a unique name, which is the same for symbols with external linkage and prefixed with the number of its block is has been defined in for symbols without external linkage.
* Storage. A symbol may be auto, register, static, extern or typedef. If the storage specifier is omitted a global variable is regarded as static and a local variable is regarded as auto. In this book, the only difference between an auto or register symbol is that the address operator cannot be applied to a register symbol.
* Type. Each symbol has a type as described in Section 7.1.
* Value. A symbol may have a value of integral,
* External. A symbol may have external linkage; that is, it is visible from other source code files.
* Address Symbol. In case of the address, arrow, or index operator the address symbol is the dereferred symbol. In case of the arrow operator or constant index we also store the offset of the address symbol.
* Address Offset.
* Offset. Each auto or register symbol has an offset on the activation record. Each struct member has also an offset relative the beginning of the struct (the first member has offset zero), all union members have offset zero.
* When the final object code is generated, we mark each auto or register symbol as used. In the end, all non-used symbols are being removed.

The name is somewhat misleading since we also store functions and types, defined by typedef, in the table. In some compiler technique textbooks, the table is called the symbol table. However, I think that term is too wide for our purposes.

Symbol.cs

using System;

using System.IO;

using System.Text;

using System.Numerics;

using System.Globalization;

using System.Collections.Generic;

namespace CCompiler {

public class Symbol {

public const string NumberId = "#";

public const string TemporaryId = "£";

public const string SeparatorId = "$";

public const string SeparatorDot = ".";

public const string FileMarker = "@";

A symbol has actually two names: the regular name it has been defined with and is looked up by the symbol table, and a unique name. The names are the same for symbols with external linkage. For symbols without external linkage, the unique name is really unique. The idea is that the no two names without external linkage in any shall have the same name. For instance, two static symbols in two different functions, or two global symbols in two different files.

private bool m\_externalLinkage;

private string m\_name, m\_uniqueName;

A symbol may be a parameter in a function definition or a temporary variable, added to the symbol table by the middle code generator.

private bool m\_temporary, m\_parameter;

A symbol has a storage and a type, it may also have a value.

private Storage m\_storage;

private Type m\_type;

private object m\_value;

Symbols with auto and register storage are given an offset on the current activation record.

private int m\_offset;

When dereferencing a symbol with the derefernce, arrow, and index (with a constant index value) expression, the dereferenced symbol is stored in the address symbol, with its offset. The offset is zero in the dereference case, the offset of the field in the arrow case, and the index value times the size of the array type in the index case. In case of the index operator with a non-constant index value, the value of the address offset is calculated in run-time.

private Symbol m\_addressSymbol;

private int m\_addressOffset;

A symbol can be assignable and addressable. It is assignable if it is not constant, and none of its members are (recursively) constant in case of struct or union and is not an array of a function. It is addressable if it does not have register storage and is not a bitfield in a struct.

private bool m\_assignable, m\_addressable;

A symbol may hold logical type, in which case the true set and false set is stored in the symbol.

private ISet<MiddleCode> m\_trueSet, m\_falseSet;

Finally, we have static counters for unique names and temporary names.

private static int UniqueNameCount = 0, TemporaryNameCount = 0;

The first constructor is used for defined symbols, variables, constants, and functions. In case of enumeration constants, they also hold a value.

public Symbol(string name, bool externalLinkage, Storage storage,

Type type, bool parameter = false, object value = null) {

m\_name = name;

m\_externalLinkage = externalLinkage;

m\_storage = storage;

If the symbol has external linkage its unique name is its regular name.

if (m\_externalLinkage) {

m\_uniqueName = m\_name;

}

If the symbol does not have external linkage, its unique name is a file marker, a unique number, and the regular name. The file marker will by the linker be replaced by a number unique for the object file. In this way, every symbol of the same file will have a unique number and each file will also have a unique number.

else {

m\_uniqueName = Symbol.FileMarker + (UniqueNameCount++) +s

Symbol.SeparatorId + m\_name;

}

m\_type = type;

m\_parameter = parameter;

m\_temporary = false;

The symbol is assignable if its type not recursively constant and is not an array or function. A type is recursively constant if it is constant or, in case of a struct or union, any of its members is constant.

m\_assignable = !m\_type.IsConstantRecursive() &&

!m\_type.IsArrayOrFunction();

A symbol is addressable if it does not have register storage and is not a bitfield in a struct or union.

m\_addressable = !IsRegister() && !m\_type.IsBitfield();

m\_value = value;

CheckValue(m\_type, m\_value);

}

If the value of the symbol is a BigInteger value, we check that it does not exceed its limits, which is defined for each integral type.

private static void CheckValue(Type type, object value) {

if (value is BigInteger) {

BigInteger bigValue = (BigInteger) value;

Assert.Error((bigValue >= TypeSize.GetMinValue(type.Sort)) &&

(bigValue <= TypeSize.GetMaxValue(type.Sort)),

type + ": " + value, Message.Value\_overflow);

}

}

In case of a dereference, arrow, dot, or index expression, the resulting symbol may be assignable and addressable.

public Symbol(Type type, bool assignable, bool addressable = false) {

m\_name = Symbol.TemporaryId + "field" + (TemporaryNameCount++);

m\_externalLinkage = false;

m\_storage = Storage.Auto;

m\_type = type;

m\_temporary = false;

m\_parameter = false;

m\_assignable = assignable;

m\_addressable = addressable;

}

In many expressions, the resulting value is a temporary symbol.

public Symbol(Type type) {

m\_name = Symbol.TemporaryId + "temporary" + (TemporaryNameCount++);

m\_externalLinkage = false;

m\_storage = Storage.Auto;

m\_type = type;

m\_temporary = true;

m\_parameter = false;

m\_assignable = false;

m\_addressable = false;

}

The resulting symbol of a logical expression.

public Symbol(ISet<MiddleCode> trueSet, ISet<MiddleCode> falseSet) {

m\_name = Symbol.TemporaryId + "logical" + (TemporaryNameCount++);

m\_storage = Storage.Auto;

m\_type = new Type(Sort.Logical);

m\_trueSet = (trueSet != null) ? trueSet : (new HashSet<MiddleCode>());

m\_falseSet = (falseSet != null) ? falseSet : (new HashSet<MiddleCode>());

m\_temporary = false;

m\_parameter = false;

m\_assignable = false;

m\_addressable = false;

}

We need to store a constant value as the denominator in an integral multiplication and division, and when pushing an integral or floating value to the floating value stack.

public Symbol(Type type, object value) {

Assert.ErrorXXX(!(value is bool));

m\_uniqueName = ValueName(type, value);

m\_storage = Storage.Static;

m\_type = type;

m\_value = value;

m\_temporary = false;

m\_parameter = false;

m\_assignable = false;

m\_addressable = false;

CheckValue(m\_type, m\_value);

}

The value is given a name the reflects its type and its actual value.

public static string ValueName(CCompiler.Type type, object value) {

Assert.ErrorXXX(value != null);

if (value is string) {

string text = (string) value;

StringBuilder buffer = new StringBuilder();

for (int index = 0; index < text.Length; ++index) {

if (char.IsLetterOrDigit(text[index]) ||

(text[index] == '\_')) {

buffer.Append(text[index]);

}

else if (text[index] != '\0') {

int asciiValue = (int) text[index];

char hex1 = "0123456789ABCDEF"[asciiValue / 16],

hex2 = "0123456789ABCDEF"[asciiValue % 16];

buffer.Append(hex1.ToString() + hex2.ToString());

}

}

return "string\_" + buffer.ToString() + Symbol.NumberId;

}

else if (value is StaticAddress) {

StaticAddress staticAddress = (StaticAddress) value;

return "staticaddress" + Symbol.SeparatorId + staticAddress.UniqueName

+ Symbol.SeparatorId + staticAddress.Offset + Symbol.NumberId;

}

else if (type.IsArray()) {

return "Array\_" + value.ToString() + Symbol.NumberId;

}

else if (type.IsFloating()) {

return "float" + type.Size().ToString() + Symbol.SeparatorId +

value.ToString().Replace("-", "minus") + Symbol.NumberId;

}

else if (type.IsLogical()) {

return "int" + type.Size().ToString() + Symbol.SeparatorId +

value.ToString().Replace("-", "minus") + Symbol.NumberId;

}

else {

return "int" + type.Size().ToString() + Symbol.SeparatorId +

value.ToString().Replace("-", "minus") + Symbol.NumberId;

}

}

public string Name {

get { return m\_name; }

set { m\_name = value; }

}

public string UniqueName {

get { return m\_uniqueName; }

set { m\_uniqueName = value; }

}

public bool ExternalLinkage {

get { return m\_externalLinkage; }

}

public Storage Storage {

get { return m\_storage; }

set { m\_storage = value; }

}

public Type Type {

get { return m\_type; }

set { m\_type = value; }

}

public int Offset {

get { return m\_offset; }

set { m\_offset = value; }

}

public bool IsExtern() {

return (m\_storage == Storage.Extern);

}

public bool IsStatic() {

return (m\_storage == Storage.Static);

}

public bool IsExternOrStatic() {

return IsExtern() || IsStatic();

}

public bool IsTypedef() {

return (m\_storage == Storage.Typedef);

}

public bool IsAuto() {

return (m\_storage == Storage.Auto);

}

public bool IsRegister() {

return (m\_storage == Storage.Register);

}

public bool IsAutoOrRegister() {

return IsAuto() || IsRegister();

}

public ISet<MiddleCode> TrueSet {

get { return m\_trueSet; }

}

public ISet<MiddleCode> FalseSet {

get { return m\_falseSet; }

}

public bool IsParameter() {

return m\_parameter;

}

public bool IsTemporary() {

return m\_temporary;

}

public object Value {

get { return m\_value; }

set { m\_value = value; }

}

public Symbol AddressSymbol {

get { return m\_addressSymbol; }

set { m\_addressSymbol = value; }

}

public int AddressOffset {

get { return m\_addressOffset; }

set { m\_addressOffset = value; }

}

public bool Assignable {

get { return m\_assignable; }

set { m\_assignable = value; }

}

public bool Addressable {

get { return m\_addressable; }

set { m\_addressable = value; }

}

public override string ToString() {

if (m\_name != null) {

if (m\_addressSymbol != null) {

return ((m\_name != null) ? m\_name : "") + " -> " +

m\_addressSymbol.ToString();

}

else {

return ((m\_name != null) ? m\_name : "");

}

}

else {

if (m\_addressSymbol != null) {

return ((m\_uniqueName != null) ? m\_uniqueName : "") + " -> " +

m\_addressSymbol.ToString();

}

else {

return ((m\_uniqueName != null) ? m\_uniqueName : "");

}

}

}

public static string SimpleName(string name) {

int index = name.LastIndexOf(Symbol.SeparatorId);

return (index != -1) ? name.Substring(index + 1).Replace("#", "") : name;

}

}

}

# The Type System

C has a rather large set of types. The simple types are made up of the integral types signed and unsigned char, short, int, and long as well as the floating types float, double, and long double. The compound type are pointers, arrays, structs, unions, and functions. Values of enumeration types (enum) are stored as signed integer. Moreover, a type can also be constant or volatile.

The Type class hold information about a type. The default constructor is private, which makes it impossible to explicitly create type objects. Instead, there is a series of static create-methods that create and return a Type object.

The type system of C is rather large. There is the integer types and floating types, structs and union, pointers, array, and function as well as void.

There is no logical type in C, but we use it to store temporary values. Moreover, there is not string type in C, but

Sort.cs

namespace CCompiler {

public enum Sort {Void, Boolean, Signed\_Char, Unsigned\_Char,

Signed\_Short\_Int, Unsigned\_Short\_Int, Signed\_Int,

Unsigned\_Int, Signed\_Long\_Int, Unsigned\_Long\_Int,

Float, Double, Long\_Double, String,

Pointer, Array, Struct, Union, // Enumeration,

Function, Logical};

}

Type.cs

using System;

using System.Linq;

using System.Numerics;

using System.Collections.Generic;

namespace CCompiler {

public class Type {

private Sort m\_sort;

Each type holds a sort, it may be an integral sort (signed or unsigned character, short integer, integer, or long integer) a floating sort (float, double, or long double), a struct or union, a pointer or array, a function, or void.

public Sort Sort {

get { return m\_sort; }

}

The following constructor takes an integral or arithmetic type, or void.

public Type(Sort sort) { // arithmetic or logical

m\_sort = sort;

}

Contrary to some other languages, there is no logical type in C. However, as C applies lazy evaluation, we need a logical type. Lazy evaluation means that an expression shall be not be evaluated more than necessary to determine its value. More specifically, if the left operand of the logical or-operator is true the right operand shall not be evaluated. In the same way, if the left operand of the logical and-operator is false the right operand shall not be evaluated. The same goes for the conditional operator, depending on whether the first expression is true or false, only the second or the third expression is evaluated. The logical type holds the two sets m\_trueSet and m\_falseSet, which hold middle code addresses to jumps instructions that will later be filled with the addresses to jump to if the expression is true or false.

// -----------------------------------------------------------------------

In case of a bitfield member in a struct, we store its bitfield mask. The mask is technically two to the power of the bits minus one or, informally, the bits number of ones, counting from the least significant bit.

The bitfield type is basically an integral type, with the addition of the bitfield mask that is used when a bitfield variable is assigned a value to set the unused bits to zero.

private BigInteger? m\_bitfieldMask = null;

public void SetBitfieldMask(int bits) {

m\_bitfieldMask = ((BigInteger) Math.Pow(2, bits) - (BigInteger.One));

}

public bool IsBitfield() {

return (m\_bitfieldMask != null);

}

public BigInteger? BitfieldMask() {

return m\_bitfieldMask;

}

// -----------------------------------------------------------------------

The type pointed at is null when the type is created, it will later be set by the declarator type.

In case of a pointer type the type its point at is given in the constructor. The sort is Sort.Pointer.

private Type m\_pointerType;

public Type(Type pointerType) {

m\_sort = Sort.Pointer;

m\_pointerType = pointerType;

}

public Type PointerType {

get { return m\_pointerType; }

set { m\_pointerType = value; }

}

// -----------------------------------------------------------------------

In case of an array, the constructor takes its size and the array type.

The array size can be zero, when the type is created. In that case it will later be set by the length of its initialization list or become uncomplete.

private int m\_arraySize;

private Type m\_arrayType;

public Type(int arraySize, Type arrayType) {

m\_sort = Sort.Array;

m\_arraySize = arraySize;

m\_arrayType = arrayType;

}

public int ArraySize {

get { return m\_arraySize; }

set { m\_arraySize = value; }

}

public Type ArrayType {

get { return m\_arrayType; }

set { m\_arrayType = value; }

}

The PointerOrArrayType method returns the pointer or array type, which case it may be.

public Type PointerOrArrayType {

get {

return (m\_sort == Sort.Pointer) ? m\_pointerType : m\_arrayType;

}

}

// -----------------------------------------------------------------------

C support both old-style and new-style function declarations. The old-style declaration takes a parameter list of identifiers, that is matched to a set of declarations, while the new-style declaration takes a parameter list of declarations.

The function constructor is a bit more complicated because we must take into consideration the fact that there are both the old-style and new-style function definition. However, both styles have in common that the function has a return type.

public enum FunctionStyle {Old, New};

private FunctionStyle m\_functionStyle;

private Type m\_returnType;

The old-style function has name list, while the new style function has a parameter list, made up by name-symbol pairs.

private List<string> m\_nameList;

private List<Pair<string,Symbol>> m\_parameterList;

The parameter list is transformed into a type list. Holding the types of each parameter.

private List<Type> m\_typeList;

Som functions, like printf and scanf, are elliptic, which means that they can take a various number of parameters.

private bool m\_ellipse;

The following constructor takes an old-style function. The name in the list must be unique.

public Type(Type returnType, List<string> nameList) {

m\_sort = Sort.Function;

m\_functionStyle = FunctionStyle.Old;

m\_returnType = returnType;

m\_nameList = nameList;

m\_parameterList = null;

m\_ellipse = false;

We add the names of the list to a set, and check that they holds the same size. If the do not, there is at least two names with the sam name, which is not allowed. If the same name is added to set twice, only the first is stored in the set, and the set will have fewer names than the list.

Assert.Error(nameList.Count == new HashSet<string>(nameList).Count,

null, Message.Duplicate\_name\_in\_parameter\_list);

}

The following construct takes a new-style function, with return type, parameter list, and ellipse status.

public Type(Type returnType, List<Pair<string,Symbol>> parameterList,

bool ellipse) {

m\_sort = Sort.Function;

m\_functionStyle = FunctionStyle.New;

m\_returnType = returnType;

m\_nameList = null;

m\_parameterList = parameterList;

m\_ellipse = ellipse;

m\_typeList = null;

if (parameterList != null) {

m\_typeList = new List<Type>();

foreach (Pair<string,Symbol> pair in parameterList) {

m\_typeList.Add(pair.Second.Type);

}

}

}

public FunctionStyle Style {

get { return m\_functionStyle; }

}

public List<string> NameList {

get { return m\_nameList; }

}

public List<Pair<string,Symbol>> ParameterList {

get { return m\_parameterList; }

}

public List<Type> TypeList {

get { return m\_typeList; }

}

public Type ReturnType {

get { return m\_returnType; }

set { m\_returnType = value; }

}

public bool IsEllipse() {

return m\_ellipse;

}

// -----------------------------------------------------------------------

A struct and union takes a list of symbols and an indication whether the struct or union is tagged. A struct or union is tagged it is given a name:

|  |  |
| --- | --- |
| struct {  int a, b;  }; | struct s {  int a, b;  }; |
| (a) An untagged struct | (b) A struct tagged with the name “s” |

The only reason for including the m\_hasTag field is that an untagged struct or union with no defined variables shall raise a warning, since it is a meaningless declaration:

struct s {int i;}; // Meaningful, s can later be used to define variables.

struct {int i;} t; // Meaningful, t is a defined variable.

struct {int i;}; // Meaningless, but allowed.

private IDictionary<string,Symbol> m\_memberMap;

public Type(Sort sort, IDictionary<string,Symbol> symbolMap) {

m\_sort = sort;

m\_memberMap = symbolMap;

}

public IDictionary<string,Symbol> MemberMap {

get { return m\_memberMap; }

set { m\_memberMap = value; }

}

// -----------------------------------------------------------------------

The enumeration type (enum) is stored as an assigned integer. However, the Specifier class of Section 6.1.2 needs to know if the type is enum to initialize its value. Therefore, we add the m\_enum field and the isEnumeration method.

private ISet<Pair<Symbol,bool>> m\_enumeratorItemSet;

private bool m\_enumeratorItem;

public Type(Sort sort, bool enumeratorItem) {

m\_sort = sort;

m\_enumeratorItem = enumeratorItem;

}

public bool EnumeratorItem {

get { return m\_enumeratorItem; }

}

public Type(Sort sort, ISet<Pair<Symbol,bool>> enumSet) {

m\_sort = sort;

m\_enumeratorItemSet = enumSet;

}

public ISet<Pair<Symbol,bool>> EnumerationItemSet {

get { return m\_enumeratorItemSet; }

}

// -----------------------------------------------------------------------

public static bool IsSigned(Sort sort) {

return (sort == Sort.Signed\_Char) || (sort == Sort.Signed\_Short\_Int) ||

(sort == Sort.Signed\_Int) || (sort == Sort.Signed\_Long\_Int);

}

public int SizeArray() {

switch (m\_sort) {

case Sort.Array:

return TypeSize.PointerSize;

default:

return Size();

}

}

Each type has a size, even though void and function have size zero. The size of an array is its size times the size of its type, the size of a struct is the sum of the sizes of its members, and the size of a union is the size of its largest member. Note that a pointer always has the same size, regardless of what it points at.

public int Size() {

switch (m\_sort) {

case Sort.Array:

return m\_arraySize \* m\_arrayType.Size();

case Sort.Struct: {

int size = 0;

foreach (Symbol symbol in m\_memberMap.Values) {

size += symbol.Type.Size();

}

return size;

}

case Sort.Union: {

int size = 0;

foreach (Symbol symbol in m\_memberMap.Values) {

size = Math.Max(size, symbol.Type.Size());

}

return size;

}

case Sort.Logical:

return TypeSize.SignedIntegerSize;

default:

return TypeSize.Size(m\_sort);

}

}

public int ConvertedSize() {

switch (m\_sort) {

case Sort.Array:

case Sort.Function:

return TypeSize.PointerSize;

default:

return Size();

}

}

It is possible to define an array without stating its size, in which case the array is given the size zero. In that case, the array size must be determined by the size of its initialization list. However, if the array definition lacks an initialization list, the array keeps the size zero and is considered incomplete. In the same way, it is possible to define only the name tag of a struct or union, with its member map to be defined later. In that case, the member map is given the value null, which is keep if the member map is not defined, and the struct or union is consider incomplete. Variables can only have complete types, and the pointer type, array type or function return type must be also complete.

XXX It is possible to define an array without stating its size, in which case the array is given the size zero. In that case, the array size must be determined by the size of its initialization list. However, if the array definition lacks an initialization list, the array keeps the size zero and is considered incomplete. In the same way, it is possible to define only the name tag of a struct or union, with its member map to be defined later. In that case, the member map is given the value null, which is keep if the member map is not defined, and the struct or union is consider incomplete. Variables can only have complete types, and the pointer type, array type or function return type must be also complete.

public bool IsComplete() {

switch (m\_sort) {

case Sort.Array:

return (m\_arraySize > 0);

case Sort.Struct:

case Sort.Union:

return (m\_memberMap != null);

default:

return true;

}

}

// -----------------------------------------------------------------------

Simply put, a type is constant if its field m\_constant is true. However, a struct or union is regarded as constant if it is constant (m\_constant is true) or is any of its members if (recursively) constant.

The idea of the volatile qualifier is to prevent optimization, and since this book is focused on optimization techniques we have no real use for the volatile qualifier. However, for the sake of completeness we include the m\_volatile field in the Type class.

private bool m\_constant;

private bool m\_volatile;

public bool IsConstant {

get { return m\_constant; }

set { m\_constant = value; }

}

public bool IsVolatile {

get { return m\_volatile; }

set { m\_volatile = value; }

}

public bool IsConstantRecursive() {

if (m\_constant) {

return true;

}

else if (IsStructOrUnion() && (m\_memberMap != null)) {

foreach (Symbol symbol in m\_memberMap.Values) {

if (symbol.Type.IsConstantRecursive()) {

return true;

}

}

}

return false;

}

public bool Volatile {

get { return m\_volatile; }

set { m\_volatile = value; }

}

// -----------------------------------------------------------------------

public override int GetHashCode() {

return base.GetHashCode();

}

Two pointers are considered to be equal if the types they point at are equal, two arrays are equal if (1) their types are equal and (2) they have the same size or both are incomplete (their sizes are zero). Two structs or unions are equals if they both are incomplete (their member maps are null) or if their member maps are equal. Note that they must not only have the same members, the members must also appear in the same order. Two functions are equal if their return types are equal, they have the same style (old or new) and their name lists or type lists, respectively, are equal.

public override bool Equals(object obj) {

if (obj is Type) {

Type type = (Type) obj;

if ((m\_constant == type.m\_constant) &&

(m\_volatile == type.m\_volatile) && (m\_sort == type.m\_sort)) {

switch (m\_sort) {

case Sort.Pointer:

return m\_pointerType.Equals(type.m\_pointerType);

case Sort.Array:

return ((m\_arraySize == 0) || (type.m\_arraySize == 0) ||

(m\_arraySize == type.m\_arraySize)) &&

m\_arrayType.Equals(type.m\_arrayType);

case Sort.Struct:

case Sort.Union:

return ((m\_memberMap == null) && (type.m\_memberMap == null)) ||

((m\_memberMap != null) && (type.m\_memberMap != null) &&

m\_memberMap.Equals(type.m\_memberMap));

case Sort.Function:

return m\_returnType.Equals(type.m\_returnType) &&

(((m\_typeList == null) && (type.m\_typeList == null)) ||

((m\_typeList != null) && (type.m\_typeList != null) &&

m\_typeList.SequenceEqual(type.m\_typeList)));

default:

return true;

}

}

}

return false;

}

// -----------------------------------------------------------------------

public bool IsVoid() {

return (m\_sort == Sort.Void);

}

public bool IsChar() {

return (m\_sort == Sort.Signed\_Char) || (m\_sort == Sort.Unsigned\_Char);

}

public bool IsShort() {

return (m\_sort == Sort.Signed\_Short\_Int) ||

(m\_sort == Sort.Unsigned\_Short\_Int);

}

public bool IsInteger() {

return (m\_sort == Sort.Signed\_Int) || (m\_sort == Sort.Unsigned\_Int);

}

public bool IsIntegral() {

return IsSigned() || IsUnsigned();

}

public bool IsSigned() {

switch (m\_sort) {

case Sort.Signed\_Char:

case Sort.Signed\_Short\_Int:

case Sort.Signed\_Int:

case Sort.Signed\_Long\_Int:

return true;

default:

return false;

}

}

public bool IsUnsigned() {

switch (m\_sort) {

case Sort.Unsigned\_Char:

case Sort.Unsigned\_Short\_Int:

case Sort.Unsigned\_Int:

case Sort.Unsigned\_Long\_Int:

return true;

default:

return false;

}

}

public bool IsFloat() {

return (m\_sort == Sort.Float);

}

public bool IsFloating() {

switch (m\_sort) {

case Sort.Float:

case Sort.Double:

case Sort.Long\_Double:

return true;

default:

return false;

}

}

public bool IsLogical() {

return (m\_sort == Sort.Logical);

}

public bool IsLogicalOrIntegral() {

return IsLogical() || IsIntegral();

}

public bool IsPointer() {

return (m\_sort == Sort.Pointer);

}

public bool IsArray() {

return (m\_sort == Sort.Array);

}

public bool IsPointerOrArray() {

return IsPointer() || IsArray();

}

public bool IsFunction() {

return (m\_sort == Sort.Function);

}

public bool IsArrayOrFunction() {

return IsArray() || IsFunction();

}

public bool IsPointerArrayStringOrFunction() {

return IsPointer() || IsArrayOrFunction() || IsString();

}

public bool IsString() {

return (m\_sort == Sort.String);

}

public bool IsArrayFunctionOrString() {

return IsArrayOrFunction() || IsString();

}

public bool IsFunctionPointer() {

return (m\_sort == Sort.Pointer) && m\_pointerType.IsFunction();

}

public bool IsFunctionOrArray() {

return IsFunction() || IsArray();

}

public bool IsArrayStringOrFunction() {

return IsFunctionOrArray() || IsString();

}

public bool IsArrayPointerStringOrFunction() {

return IsPointer() || IsFunctionOrArray();

}

public bool IsArrayFunctionStringStructOrUnion() {

return IsFunctionOrArray() || IsString() || IsStructOrUnion();

}

public bool IsArithmetic() {

return IsIntegral() || IsFloating();

}

public bool IsIntegralOrPointer() {

return IsIntegral() || IsPointer();

}

public bool IsIntegralLogicalOrPointer() {

return IsIntegral() || IsLogical() || IsPointer();

}

public bool IsIntegralOrArray() {

return IsIntegral() || IsArray();

}

public bool IsIntegralArrayOrPointer() {

return IsIntegral() || IsPointer() || IsArray();

}

public bool IsIntegralPointerArrayOrString() {

return IsIntegralArrayOrPointer() || IsString();

}

public bool IsIntegralPointerArrayStringOrFunction() {

return IsIntegralPointerArrayOrString() || IsFunction();

}

public bool IsIntegralPointerOrFunction() {

return IsIntegralOrPointer() || IsFunction();

}

public bool IsIntegralPointerArrayOrFunction() {

return IsIntegralArrayOrPointer() || IsFunction();

}

public bool IsArithmeticOrPointer() {

return IsArithmetic() || IsPointer();

}

public bool IsArithmeticPointerOrArray() {

return IsArithmeticOrPointer() || IsArray();

}

public bool IsArithmeticPointerArrayOrFunction() {

return IsArithmeticPointerOrArray() || IsFunction();

}

public bool IsArithmeticPointerArrayStringOrFunction() {

return IsArithmeticPointerArrayOrFunction() || IsString();

}

public bool IsStruct() {

return (m\_sort == Sort.Struct);

}

public bool IsUnion() {

return (m\_sort == Sort.Union);

}

public bool IsStructOrUnion() {

return IsStruct() || IsUnion();

}

public bool IsArithmeticPointerStructOrUnion() {

return IsArithmeticOrPointer() || IsStructOrUnion();

}

public bool IsEnumerator() {

return (m\_enumeratorItemSet != null);

}

public override string ToString() {

return Enum.GetName(typeof(Sort), m\_sort).

Replace("\_\_", "-").Replace("\_", " ").ToLower();

}

public static Type SignedShortIntegerType =

new Type(Sort.Signed\_Short\_Int);

public static Type UnsignedShortIntegerType =

new Type(Sort.Unsigned\_Short\_Int);

public static Type SignedIntegerType = new Type(Sort.Signed\_Int);

public static Type UnsignedIntegerType = new Type(Sort.Unsigned\_Int);

public static Type SignedLongIntegerType = new Type(Sort.Signed\_Long\_Int);

public static Type UnsignedLongIntegerType =

new Type(Sort.Unsigned\_Long\_Int);

public static Type FloatType = new Type(Sort.Float);

public static Type DoubleType = new Type(Sort.Double);

public static Type LongDoubleType = new Type(Sort.Long\_Double);

public static Type SignedCharType = new Type(Sort.Signed\_Char);

public static Type UnsignedCharType = new Type(Sort.Unsigned\_Char);

public static Type StringType = new Type(Sort.String);

public static Type VoidType = new Type(Sort.Void);

public static Type PointerTypeX = new Type(SignedIntegerType);

public static Type VoidPointerType = new Type(new Type(Sort.Void));

public static Type LogicalType = new Type(Sort.Logical);

}

}

The TypeSize class holds the sizes of the types and the minimum and maximum values of each type.

Each type has a size, even though void and function have size zero. The size of an array is its size times the size of its type, the size of a struct is the sum of the sizes of its members, and the size of a union is the size of its largest member. Note that a pointer always has the same size, regardless of what it points at.

TypeSize.cs

using System.Numerics;

using System.Collections.Generic;

namespace CCompiler {

class TypeSize {

public static int PointerSize;

public static int SignedIntegerSize;

public static IDictionary<Sort,int>

m\_sizeMap = new Dictionary<Sort,int>();

public static IDictionary<int,Type>

m\_signedMap = new Dictionary<int,Type>(),

m\_unsignedMap = new Dictionary<int,Type>();

private static IDictionary<int, BigInteger>

m\_maskMap = new Dictionary<int, BigInteger>();

private static IDictionary<Sort,BigInteger>

m\_minValueMap = new Dictionary<Sort,BigInteger>(),

m\_maxValueMap = new Dictionary<Sort,BigInteger>();

private static IDictionary<Sort,decimal>

m\_minValueFloatMap = new Dictionary<Sort,decimal>(),

m\_maxValueFloatMap = new Dictionary<Sort,decimal>();

static TypeSize() {

m\_maskMap.Add(1, BigInteger.Zerox000000FF);

m\_maskMap.Add(2, BigInteger.Zerox0000FFFF);

m\_maskMap.Add(4, BigInteger.ZeroxFFFFFFFF);

The sizes and values of the types depends on whether the compiler generates code for the Linux or Windows environment.

if (Start.Windows) {

PointerSize = 2;

SignedIntegerSize = 2;

m\_sizeMap.Add(Sort.Void, 0);

m\_sizeMap.Add(Sort.Function, 0);

m\_sizeMap.Add(Sort.Logical, 1);

m\_sizeMap.Add(Sort.Array, 2);

m\_sizeMap.Add(Sort.Pointer, 2);

m\_sizeMap.Add(Sort.String, 2);

m\_sizeMap.Add(Sort.Signed\_Char, 1);

m\_sizeMap.Add(Sort.Unsigned\_Char, 1);

m\_sizeMap.Add(Sort.Signed\_Short\_Int, 1);

m\_sizeMap.Add(Sort.Unsigned\_Short\_Int, 1);

m\_sizeMap.Add(Sort.Signed\_Int, 2);

m\_sizeMap.Add(Sort.Unsigned\_Int, 2);

m\_sizeMap.Add(Sort.Signed\_Long\_Int, 4);

m\_sizeMap.Add(Sort.Unsigned\_Long\_Int, 4);

m\_sizeMap.Add(Sort.Float, 4);

m\_sizeMap.Add(Sort.Double, 8);

m\_sizeMap.Add(Sort.Long\_Double, 8);

m\_signedMap.Add(1, Type.SignedCharType);

m\_signedMap.Add(2, Type.SignedIntegerType);

m\_signedMap.Add(4, Type.SignedLongIntegerType);

m\_unsignedMap.Add(1, Type.UnsignedCharType);

m\_unsignedMap.Add(2, Type.UnsignedIntegerType);

m\_unsignedMap.Add(4, Type.UnsignedLongIntegerType);

m\_minValueMap.Add(Sort.Logical, 0);

m\_minValueMap.Add(Sort.Signed\_Char, -128);

m\_minValueMap.Add(Sort.Unsigned\_Char, 0);

m\_minValueMap.Add(Sort.Signed\_Short\_Int, -128);

m\_minValueMap.Add(Sort.Unsigned\_Short\_Int, 0);

m\_minValueMap.Add(Sort.Signed\_Int, -32768);

m\_minValueMap.Add(Sort.Unsigned\_Int, 0);

m\_minValueMap.Add(Sort.Array, 0);

m\_minValueMap.Add(Sort.Pointer, 0);

m\_minValueMap.Add(Sort.Signed\_Long\_Int, -2147483648);

m\_minValueMap.Add(Sort.Unsigned\_Long\_Int, 0);

m\_maxValueMap.Add(Sort.Logical, 1);

m\_maxValueMap.Add(Sort.Signed\_Char, 127);

m\_maxValueMap.Add(Sort.Unsigned\_Char, 255);

m\_maxValueMap.Add(Sort.Signed\_Short\_Int, 127);

m\_maxValueMap.Add(Sort.Unsigned\_Short\_Int, 255);

m\_maxValueMap.Add(Sort.Signed\_Int, 32767);

m\_maxValueMap.Add(Sort.Unsigned\_Int, 65535);

m\_maxValueMap.Add(Sort.Array, 65535);

m\_maxValueMap.Add(Sort.Pointer, 65535);

m\_maxValueMap.Add(Sort.Signed\_Long\_Int, 2147483647);

m\_maxValueMap.Add(Sort.Unsigned\_Long\_Int, 4294967295);

/\*m\_minValueFloatMap.Add(Sort.Float, decimal.

Parse("1.2E-38", NumberStyles.Float));

m\_minValueFloatMap.Add(Sort.Double, decimal.

Parse("2.3E-308", NumberStyles.Float));

m\_minValueFloatMap.Add(Sort.Long\_Double, decimal.

Parse("2.3E-308", NumberStyles.Float));

m\_maxValueFloatMap.Add(Sort.Float, decimal.

Parse("3.4E+38", NumberStyles.Float));

m\_maxValueFloatMap.Add(Sort.Double, decimal.

Parse("1.7E+308", NumberStyles.Float));

m\_maxValueFloatMap.Add(Sort.Long\_Double, decimal.

Parse("1.7E+308", NumberStyles.Float));\*/

/\*m\_maskMap.Add(Sort.Unsigned\_Char, 0x00000000000000FF);

m\_maskMap.Add(Sort.Unsigned\_Short\_Int, 0x00000000000000FF);

m\_maskMap.Add(Sort.Unsigned\_Int, 0x000000000000FFFF);

m\_maskMap.Add(Sort.Unsigned\_Long\_Int, 0x00000000FFFFFFFF);\*/

}

if (Start.Linux) {

PointerSize = 8;

SignedIntegerSize = 4;

m\_sizeMap.Add(Sort.Void, 0);

m\_sizeMap.Add(Sort.Function, 0);

m\_sizeMap.Add(Sort.Logical, 1);

m\_sizeMap.Add(Sort.Pointer, 8);

m\_sizeMap.Add(Sort.Array, 8);

m\_sizeMap.Add(Sort.String, 4);

m\_sizeMap.Add(Sort.Signed\_Char, 1);

m\_sizeMap.Add(Sort.Unsigned\_Char, 1);

m\_sizeMap.Add(Sort.Signed\_Short\_Int, 2);

m\_sizeMap.Add(Sort.Unsigned\_Short\_Int, 2);

m\_sizeMap.Add(Sort.Signed\_Int, 4);

m\_sizeMap.Add(Sort.Unsigned\_Int, 4);

m\_sizeMap.Add(Sort.Signed\_Long\_Int, 8);

m\_sizeMap.Add(Sort.Unsigned\_Long\_Int, 8);

m\_sizeMap.Add(Sort.Float, 4);

m\_sizeMap.Add(Sort.Double, 8);

m\_sizeMap.Add(Sort.Long\_Double, 8);

m\_signedMap.Add(1, Type.SignedCharType);

m\_signedMap.Add(2, Type.SignedShortIntegerType);

m\_signedMap.Add(4, Type.SignedIntegerType);

m\_signedMap.Add(8, Type.SignedLongIntegerType);

m\_unsignedMap.Add(1, Type.UnsignedCharType);

m\_unsignedMap.Add(2, Type.UnsignedShortIntegerType);

m\_unsignedMap.Add(4, Type.UnsignedIntegerType);

m\_unsignedMap.Add(8, Type.UnsignedLongIntegerType);

m\_minValueMap.Add(Sort.Logical, 0);

m\_minValueMap.Add(Sort.Signed\_Char, -128);

m\_minValueMap.Add(Sort.Unsigned\_Char, 0);

m\_minValueMap.Add(Sort.Signed\_Short\_Int, -32768);

m\_minValueMap.Add(Sort.Unsigned\_Short\_Int, 0);

m\_minValueMap.Add(Sort.Signed\_Int, -2147483648);

m\_minValueMap.Add(Sort.Unsigned\_Int, 0);

m\_minValueMap.Add(Sort.Array, 0);

m\_minValueMap.Add(Sort.Pointer, 0);

m\_minValueMap.Add(Sort.Signed\_Long\_Int, -9223372036854775808);

m\_minValueMap.Add(Sort.Unsigned\_Long\_Int, 0);

m\_maxValueMap.Add(Sort.Logical, 1);

m\_maxValueMap.Add(Sort.Signed\_Char, 127);

m\_maxValueMap.Add(Sort.Unsigned\_Char, 255);

m\_maxValueMap.Add(Sort.Signed\_Short\_Int, 32767);

m\_maxValueMap.Add(Sort.Unsigned\_Short\_Int, 65535);

m\_maxValueMap.Add(Sort.Signed\_Int, 2147483647);

m\_maxValueMap.Add(Sort.Unsigned\_Int, 4294967295);

m\_maxValueMap.Add(Sort.Array, 4294967295);

m\_maxValueMap.Add(Sort.Pointer, 4294967295);

m\_maxValueMap.Add(Sort.Signed\_Long\_Int, 9223372036854775807);

m\_maxValueMap.Add(Sort.Unsigned\_Long\_Int, 18446744073709551615);

/\*m\_minValueFloatMap.Add(Sort.Float, decimal.

Parse("1.2E-38", NumberStyles.Float));

m\_minValueFloatMap.Add(Sort.Double, decimal.

Parse("2.3E-308", NumberStyles.Float));

m\_minValueFloatMap.Add(Sort.Long\_Double, decimal.

Parse("2.3E-308", NumberStyles.Float));

m\_maxValueFloatMap.Add(Sort.Float, decimal.

Parse("3.4E+38", NumberStyles.Float));

m\_maxValueFloatMap.Add(Sort.Double, decimal.

Parse("1.7E+308", NumberStyles.Float));

m\_maxValueFloatMap.Add(Sort.Long\_Double, decimal.

Parse("1.7E+308", NumberStyles.Float));\*/

/\* m\_maskMap.Add(Sort.Unsigned\_Char, 0x00000000000000FF);

m\_maskMap.Add(Sort.Unsigned\_Short\_Int,

0x00000000000000FF);

m\_maskMap.Add(Sort.Unsigned\_Int, 0x000000000000FFFF);

m\_maskMap.Add(Sort.Unsigned\_Long\_Int,

0x0FFFFFFFFFFFFFFF);\*/

}

}

public static BigInteger GetMinValue(Sort sort) {

return m\_minValueMap[sort];

}

public static BigInteger GetMaxValue(Sort sort) {

return m\_maxValueMap[sort];

}

public static BigInteger GetMask(Sort sort) {

return m\_maskMap[m\_sizeMap[sort]];

}

public static Type SizeToSignedType(int size) {

return m\_signedMap[size];

}

public static Type SizeToUnsignedType(int size) {

return m\_unsignedMap[size];

}

public static int Size(Sort sort) {

return m\_sizeMap[sort];

}

}

}

## Type Casting

Type casting occurs on several occasions in C. There are two types of cast: explicit (manual) conversation where type cast is stated and implicit (automatic) where it the type cast is omitted:

double x = 3.1;

int y = x, // implicit cast

z = (int) x; // explicit cast

Even though implicit casts are allowed in C, a warning is often raised when the cast causes the value to be truncated in some form, like in the code above where the decimal part of the double value is lost when casted into an integer value.

Naturally, if the types of the source and target trees are the same, we just return the source tree.

In case of two pointers, we create and return new pointer in the explicit case. The reason for this is the resulting type shall not be a left-value. We want to prevent that an explicit cast is assigned a value. An expression such as “((char\*) p) = &q;” is not allowed.

In case of a cast from an integral type of one byte (pointers are always two bytes) to a floating type, it becomes a little bit more complicated since there is no assembly converting operator. Instead, we need to first convert the small integral value (one byte) to a larger integral value (two bytes), which we in turn convert to a floating value.

The same goes for the other way around, when converting a floating value to a small (one byte) integral value, we first need to convert the floating value to an integral value of two bytes, which we in turn convert to an integral value of one byte.

The remaining cases are trivial, we just generate the appropriate converting operation.

When converting a string to a character pointer we just return the source tree.

When converting between pointers and arrays we check that they have the same type and (in the arrays case) the same size.

It is possible to explicitly convert a value to void. In that case, we just return null.

If none of the cases above applies, we raise an error.

The implicitCast method does basically the same thing as explicitCast. The difference is that is raises warnings in case of converting from a larger type to a smaller type. Another difference is that we do not need to create a new type if it has the same size as the source type.

In case of two pointers, we create and return new pointer in the explicit case. The reason for this is the resulting type shall not be a left-value. We want to prevent that an explicit cast is assigned a value. An expression such as “((char\*) p) = &q;” is not allowed.

TypeCast.cs

using System.Numerics;

using System.Collections.Generic;

namespace CCompiler {

public class TypeCast {

public static Expression LogicalToIntegral(Expression expression) {

if (expression.Symbol.Type.IsLogical()) {

return ImplicitCast(expression, Type.SignedIntegerType);

}

return expression;

}

public static Expression LogicalToFloating(Expression expression) {

if (expression.Symbol.Type.IsLogical()) {

return ImplicitCast(expression, Type.DoubleType);

}

return expression;

}

public static Expression ToLogical(Expression expression) {

if (!expression.Symbol.Type.IsLogical()) {

return ImplicitCast(expression, Type.LogicalType);

}

return expression;

}

public static Expression ImplicitCast(Expression fromExpression,

Type toType) {

Expression constantExpression =

ConstantExpression.Cast(fromExpression, toType);

if (constantExpression != null) {

return constantExpression;

}

Type fromType = fromExpression.Symbol.Type;

if (fromType.Equals(toType) ||

(fromType.IsLogical() && toType.IsLogical()) ||

(fromType.IsPointerArrayStringOrFunction() &&

toType.IsPointerOrArray()) ||

(((fromType.IsFloating() && toType.IsFloating()) ||

(fromType.IsIntegralPointerOrFunction() &&

toType.IsIntegralPointerArrayOrFunction())) &&

(fromType.ConvertedSize() == toType.ConvertedSize()))) {

return fromExpression;

}

else {

return ExplicitCast(fromExpression, toType);

}

}

public static Expression ExplicitCast(Expression fromExpression,

Type toType) {

Expression constantExpression =

ConstantExpression.Cast(fromExpression, toType);

if (constantExpression != null) {

return constantExpression;

}

Symbol fromSymbol = fromExpression.Symbol;

Type fromType = fromSymbol.Type;

if (toType.IsVoid()) {

return (new Expression(new Symbol(toType), null, null));

}

else if (fromType.IsStructOrUnion() && toType.IsStructOrUnion()) {

Assert.Error(fromType.Equals(toType), fromType + " to " + toType,

Message.Invalid\_type\_cast);

return (new Expression(new Symbol(toType), fromExpression.ShortList,

fromExpression.LongList));

}

else if (fromType.IsLogical()) {

if (toType.IsFloating()) {

List<MiddleCode> codeList = fromExpression.LongList;

Symbol resultSymbol = new Symbol(toType);

MiddleCode oneCode = MiddleCodeGenerator.AddMiddleCode(codeList,

MiddleOperator.PushOne);

MiddleCodeGenerator.Backpatch(fromSymbol.TrueSet, oneCode);

MiddleCode toCode = new MiddleCode(MiddleOperator.Empty);

MiddleCodeGenerator.AddMiddleCode(codeList, MiddleOperator.Goto,

toCode);

MiddleCode zeroCode = MiddleCodeGenerator.AddMiddleCode(codeList,

MiddleOperator.PushZero);

MiddleCodeGenerator.Backpatch(fromSymbol.FalseSet, zeroCode);

codeList.Add(toCode);

return (new Expression(resultSymbol, fromExpression.ShortList,

codeList));

}

else if (toType.IsIntegralArrayOrPointer()) {

List<MiddleCode> codeList = fromExpression.LongList;

Symbol resultSymbol = new Symbol(toType);

Symbol oneSymbol = new Symbol(toType, (BigInteger.One));

MiddleCode assignTrue =

MiddleCodeGenerator.AddMiddleCode(codeList, MiddleOperator.Assign,

resultSymbol, oneSymbol);

MiddleCodeGenerator.Backpatch(fromSymbol.TrueSet, assignTrue);

MiddleCode toCode = new MiddleCode(MiddleOperator.Empty);

MiddleCodeGenerator.AddMiddleCode(codeList, MiddleOperator.Goto,

toCode);

Symbol zeroSymbol = new Symbol(toType, (BigInteger.Zero));

MiddleCode assignFalse =

MiddleCodeGenerator.AddMiddleCode(codeList, MiddleOperator.Assign,

resultSymbol, zeroSymbol);

MiddleCodeGenerator.Backpatch(fromSymbol.FalseSet, assignFalse);

codeList.Add(toCode);

return (new Expression(resultSymbol, fromExpression.ShortList,

codeList));

}

Assert.Error(fromType + " to " + toType, Message.Invalid\_type\_cast);

}

The same goes for the other way around, when converting a floating value to a small (one byte) integral value, we first need to convert the floating value to an integral value of two bytes, which we in turn convert to an integral value of one byte.

else if (fromType.IsFloating()) {

if (toType.IsFloating()) {

return (new Expression(new Symbol(toType), fromExpression.ShortList,

fromExpression.LongList));

}

else if (toType.IsIntegralArrayOrPointer()) {

List<MiddleCode> codeList = fromExpression.LongList;

Symbol resultSymbol = new Symbol(toType);

if (toType.Size() == 1) {

Type tempType = fromType.IsSigned() ? Type.SignedIntegerType

: Type.UnsignedIntegerType;

Symbol tempSymbol = new Symbol(tempType);

MiddleCode tempCode =

new MiddleCode(MiddleOperator.FloatingToIntegral, tempSymbol,

fromSymbol);

codeList.Add(tempCode);

MiddleCode resultCode =

new MiddleCode(MiddleOperator.IntegralToIntegral, resultSymbol,

tempSymbol);

codeList.Add(resultCode);

}

else {

MiddleCode resultCode =

new MiddleCode(MiddleOperator.FloatingToIntegral, resultSymbol,

fromSymbol);

codeList.Add(resultCode);

}

return (new Expression(resultSymbol, fromExpression.ShortList,

codeList));

}

else if (toType.IsLogical()) {

List<MiddleCode> codeList = fromExpression.LongList;

ISet<MiddleCode> trueSet = new HashSet<MiddleCode>();

Symbol zeroSymbol = new Symbol(toType, (decimal) 0);

MiddleCode testCode =

new MiddleCode(MiddleOperator.NotEqual, null,

fromExpression.Symbol, zeroSymbol);

codeList.Add(testCode);

trueSet.Add(testCode);

ISet<MiddleCode> falseSet = new HashSet<MiddleCode>();

MiddleCode gotoCode = new MiddleCode(MiddleOperator.Goto);

codeList.Add(gotoCode);

falseSet.Add(testCode);

Symbol symbol = new Symbol(trueSet, falseSet);

return (new Expression(symbol, null, codeList));

}

Assert.Error(fromType + " to " + toType, Message.Invalid\_type\_cast);

}

In case of a cast from an integral type of one byte (pointers are always two bytes) to a floating type, it becomes a little bit more complicated since there is no assembly converting operator. Instead, we need to first convert the small integral value (one byte) to a larger integral value (two bytes), which we in turn convert to a floating value.

else if (fromType.IsIntegralArrayOrPointer()) {

if (toType.IsFloating()) {

List<MiddleCode> codeList = fromExpression.LongList;

Symbol resultSymbol = new Symbol(toType);

if (fromType.Size() == 1) {

Type tempType = fromType.IsSigned() ? Type.SignedIntegerType

: Type.UnsignedIntegerType;

Symbol tempSymbol = new Symbol(tempType);

MiddleCodeGenerator.

AddMiddleCode(codeList, MiddleOperator.IntegralToIntegral,

tempSymbol, fromSymbol);

MiddleCodeGenerator.

AddMiddleCode(codeList, MiddleOperator.IntegralToFloating,

resultSymbol, tempSymbol);

}

else {

MiddleCodeGenerator.

AddMiddleCode(codeList, MiddleOperator.IntegralToFloating,

resultSymbol, fromSymbol);

}

return (new Expression(resultSymbol, fromExpression.ShortList,

codeList));

}

else if (toType.IsIntegralArrayOrPointer()) {

List<MiddleCode> codeList = fromExpression.LongList;

Symbol resultSymbol = new Symbol(toType);

MiddleCodeGenerator.

AddMiddleCode(codeList, MiddleOperator.IntegralToIntegral,

resultSymbol, fromSymbol);

return (new Expression(resultSymbol, fromExpression.ShortList,

codeList));

}

else if (toType.IsLogical()) {

List<MiddleCode> codeList = fromExpression.LongList;

ISet<MiddleCode> trueSet = new HashSet<MiddleCode>();

Symbol zeroSymbol = new Symbol(toType, BigInteger.Zero);

MiddleCode testCode =

new MiddleCode(MiddleOperator.NotEqual, null,

fromExpression.Symbol, zeroSymbol);

codeList.Add(testCode);

trueSet.Add(testCode);

ISet<MiddleCode> falseSet = new HashSet<MiddleCode>();

MiddleCode gotoCode = new MiddleCode(MiddleOperator.Goto);

codeList.Add(gotoCode);

falseSet.Add(gotoCode);

Symbol symbol = new Symbol(trueSet, falseSet);

return (new Expression(symbol, null, codeList));

}

Assert.Error(fromType + " to " + toType, Message.Invalid\_type\_cast);

}

else {

Assert.Error(fromType + " to " + toType, Message.Invalid\_type\_cast);

}

return null;

}

Type promotion is the process of converting the “smaller” type to the “larger” type in a binary expression. For instance, in the expression below, the integer i shall be converted to a double value:

int i = 3;

double x = 3.14, y;

y = x + i;

Below is a table holding the relations of the types. Note that float is considered to be larger than long, event thought they have the same size. The same goes for short and char.

|  |  |  |
| --- | --- | --- |
|  |  | Size |
| Largest type | long double | 8 |
|  | double | 8 |
|  | float | 4 |
|  | signed long int  unsigned long int | 4 |
|  | signed int  unsigned int | 2 |
|  | signed short int  unsigned short int | 1 |
| Smallest type | signed char  unsigned char | 1 |

In accordance with the Ansi C standard the signed and unsigned values have the same size, and a character is always one byte[[5]](#footnote-5). However, the standard does not state whether a signed value shall be converted to an unsigned value or the other way around in the expression below. In book we always convert the unsigned value to the matching signed type (if the hold the same size), with a warning in case of implicit conversion.

signed int i = -3;

unsigned int u = 6;

i + u;

public static Type MaxType(Type leftType, Type rightType) {

if ((leftType.IsFloating() && !rightType.IsFloating()) ||

((leftType.Size() == rightType.Size()) &&

leftType.IsSigned() && !rightType.IsSigned())) {

return leftType;

}

else if ((!leftType.IsFloating() && rightType.IsFloating()) ||

((leftType.Size() == rightType.Size()) &&

!leftType.IsSigned() && rightType.IsSigned())) {

return rightType;

}

else {

return (leftType.Size() > rightType.Size()) ? leftType : rightType;

}

}

}

}

# Constant Expression

There are advantages by calculating the values of constant expression during the compilation rather than during the execution. Besides the optimization benefits, there are some occasions when we need to calculate the value. In array definitions, we need to evaluate the size of the array, in initialized enumeration values we need to evaluate the value to be able to assign the next uninitialized enumeration its correct value, and the preprocessor need to evaluate the value of if-directives.

ConstantExpression.cs

using System.Numerics;

using System.Collections.Generic;

namespace CCompiler {

public class ConstantExpression {

The IsConstant method returns true if the expression can be evaluated to a value, which it can if it holds logical type and one of the true set or the false set is empty. If one set is empty, the expression holds no if-statement, only a goto-statement. It can also be evaluated to a constant value if it holds integral, pointer type, or floating type with a not-null value.

public static bool IsConstant(Expression expression) {

Symbol resultSymbol = expression.Symbol;

Type type = resultSymbol.Type;

return (type.IsLogical() && ((resultSymbol.TrueSet.Count == 0) ||

(resultSymbol.FalseSet.Count == 0))) ||

(type.IsIntegralOrPointer() &&

(resultSymbol.Value is BigInteger)) ||

(type.IsFloating() && (resultSymbol.Value is decimal));

}

The IsTrue method returns true if the (assumed constant) expression can be evaluated to a true value., which it can if the value is a reference to a BigInteger object that is not zero, or a decimal object that is not zero, or if the true set is not empty.

public static bool IsTrue(Expression expression) {

return ((expression.Symbol.Value is BigInteger) &&

!((BigInteger) expression.Symbol.Value).IsZero) ||

((expression.Symbol.Value is decimal) &&

!((decimal) expression.Symbol.Value).Equals((decimal) 0)) ||

(expression.Symbol.TrueSet.Count > 0);

}

The LogicalToIntegral method makes sure the expression holds integral type. If it holds logical type, it is cast from logical type to signed integer type.

private static Expression LogicalToIntegral(Expression expression) {

if (expression.Symbol.Type.IsLogical()) {

return Cast(expression, Type.SignedIntegerType);

}

return expression;

}

The LogicalToFloating method makes sure the expression holds floating type. If it holds logical type, it is cast from logical type to double type.

private static Expression LogicalToFloating(Expression expression) {

if (expression.Symbol.Type.IsLogical()) {

return Cast(expression, Type.DoubleType);

}

return expression;

}

The ToLogical method makes sure the expression has a logical type. The expression is cast to logical type if it does not already hold logical type.

private static Expression ToLogical(Expression expression) {

if (!expression.Symbol.Type.IsLogical()) {

return Cast(expression, Type.LogicalType);

}

return expression;

}

The Relation method returns the calculated value of the expression, if possible. Otherwise, it returns null.

public static Expression Relation(MiddleOperator middleOp,

Expression leftExpression,

Expression rightExpression) {

If at least one of the expressions is not constant, we just return null.

if (!IsConstant(leftExpression) || !IsConstant(rightExpression)) {

return null;

}

If at least one of the expressions have floating type, we call the RelationFloating method. If not at least one of the expressions have floating type, both expressions must have integral type and we call RelationIntegral instead.

else if (leftExpression.Symbol.Type.IsFloating() ||

rightExpression.Symbol.Type.IsFloating()) {

return RelationFloating(middleOp, leftExpression, rightExpression);

}

else {

return RelationIntegral(middleOp, leftExpression, rightExpression);

}

}

The RelationIntegral method evaluates the constant value of a relation expression. We assume that both operands are constant.

private static Expression RelationIntegral(MiddleOperator middleOp,

Expression leftExpression,

Expression rightExpression) {

First, we need to cast the expression from a possible logical type into an integral type (signed integer).

leftExpression = LogicalToIntegral(leftExpression);

rightExpression = LogicalToIntegral(leftExpression);

This method is only called if both expressions have values. Therefore, we do not have to check that the values are not null.

BigInteger leftValue = (BigInteger) leftExpression.Symbol.Value,

rightValue = (BigInteger) rightExpression.Symbol.Value;

Since the expression has a constant value, we only add one unconditional jump instruction.

List<MiddleCode> longList = new List<MiddleCode>();

MiddleCode jumpCode = new MiddleCode(MiddleOperator.Goto);

longList.Add(jumpCode);

ISet<MiddleCode> jumpSet = new HashSet<MiddleCode>();

jumpSet.Add(jumpCode);

Then we need to evaluate the value of the expression. Since the values are object of the BigInteger class we can use the regular relation operators.

bool resultValue = false;

switch (middleOp) {

case MiddleOperator.Equal:

resultValue = (leftValue == rightValue);

break;

case MiddleOperator.NotEqual:

resultValue = (leftValue != rightValue);

break;

case MiddleOperator.SignedLessThan:

case MiddleOperator.UnsignedLessThan:

resultValue = (leftValue < rightValue);

break;

case MiddleOperator.SignedLessThanEqual:

case MiddleOperator.UnsignedLessThanEqual:

resultValue = (leftValue <= rightValue);

break;

case MiddleOperator.SignedGreaterThan:

case MiddleOperator.UnsignedGreaterThan:

resultValue = (leftValue > rightValue);

break;

case MiddleOperator.SignedGreaterThanEqual:

case MiddleOperator.UnsignedGreaterThanEqual:

resultValue = (leftValue >= rightValue);

break;

}

Finally, we define the logical symbol with the jump set, depending on the value of the expression. If the value is true the jump set becomes the true set of the symbol, and if the value is false the jump set becomes the false set of the symbol.

Symbol resultSymbol = resultValue ? (new Symbol(jumpSet, null))

: (new Symbol(null, jumpSet));

The resulting expression has no short list, and its long list is made up be one instruction only: the unconditional jump.

return (new Expression(resultSymbol, null, longList));

}

The RelationFloating method makes sure the expression holds integral type. If it holds logical type, it is cast from logical type to signed integer type.

private static Expression RelationFloating(MiddleOperator middleOp,

Expression leftExpression,

Expression rightExpression) {

First, we need to cast the potential logical or integral expressions to floating type (double).

leftExpression = LogicalToFloating(leftExpression);

rightExpression = LogicalToFloating(leftExpression);

Then we cast left value and right from BigInteger to decimal, if necessary.

decimal leftValue;

if (leftExpression.Symbol.Value is BigInteger) {

leftValue = (decimal) ((BigInteger) leftExpression.Symbol.Value);

}

else {

leftValue = (decimal) leftExpression.Symbol.Value;

}

decimal rightValue;

if (rightExpression.Symbol.Value is BigInteger) {

rightValue = (decimal) ((BigInteger) rightExpression.Symbol.Value);

}

else {

rightValue = (decimal) rightExpression.Symbol.Value;

}

We add an unconditional jump instruction to the long list and jump set. The jump set will take the place of true set of false set, depending on whether the value is true.

List<MiddleCode> longList = new List<MiddleCode>();

MiddleCode jumpCode = new MiddleCode(MiddleOperator.Goto);

longList.Add(jumpCode);

ISet<MiddleCode> jumpSet = new HashSet<MiddleCode>();

jumpSet.Add(jumpCode);

Since the values are object of the decimal class, we can use the regular relation operators to evaluate the value.

bool resultValue = false;

switch (middleOp) {

case MiddleOperator.Equal:

resultValue = (leftValue == rightValue);

break;

case MiddleOperator.NotEqual:

resultValue = (leftValue != rightValue);

break;

case MiddleOperator.SignedLessThan:

resultValue = (leftValue < rightValue);

break;

case MiddleOperator.SignedLessThanEqual:

resultValue = (leftValue <= rightValue);

break;

case MiddleOperator.SignedGreaterThan:

resultValue = (leftValue > rightValue);

break;

case MiddleOperator.SignedGreaterThanEqual:

resultValue = (leftValue >= rightValue);

break;

}

The resulting symbol has a logical value. The jump set becomes its true set or false set, depending on whether the result value is true.

Symbol resultSymbol = resultValue ? (new Symbol(jumpSet, null))

: (new Symbol(null, jumpSet));

return (new Expression(resultSymbol, null, longList));

}

The Logical method evaluate the value of two constant logical expressions, with the or or and operator.

public static Expression Logical(MiddleOperator middleOp,

Expression leftExpression,

Expression rightExpression) {

if (!IsConstant(leftExpression) || !IsConstant(rightExpression)) {

return null;

}

We make sure that the expressions have logical type by calling ToLogical.

Expression leftLogicalExpression = ToLogical(leftExpression),

rightLogicalExpression = ToLogical(rightExpression);

We decide the value of the expression by looking at their tre sets. If the value is true, the unconditional jump instruction is stored in its true set. Otherwise, the value is stored in the false set. In this way, we can conclude that the value is true if the true set is not empty, and false if the true set is empty.

bool leftValue = leftLogicalExpression.Symbol.TrueSet.Count > 0,

rightValue = rightLogicalExpression.Symbol.TrueSet.Count > 0;

The long list of the final expression holds only one unconditional jump instruction that is added to the jump set, which will be the true set or false of the final expression.

List<MiddleCode> longList = new List<MiddleCode>();

MiddleCode jumpCode = new MiddleCode(MiddleOperator.Goto);

longList.Add(jumpCode);

ISet<MiddleCode> jumpSet = new HashSet<MiddleCode>();

jumpSet.Add(jumpCode);

We compare the expression with the given operator, and receive the final value.

bool resultValue = false;

switch (middleOp) {

case MiddleOperator.LogicalAnd:

resultValue = leftValue && rightValue;

break;

case MiddleOperator.LogicalOr:

resultValue = leftValue || rightValue;

break;

}

The resulting symbol has a logical value. The jump set becomes its true set or false set, depending on whether the result value is true.

Symbol resultSymbol = resultValue ? (new Symbol(jumpSet, null))

: (new Symbol(null, jumpSet));

The final expression does not have a short list, and its long list is made up by the unconditional jump instruction only.

return (new Expression(resultSymbol, null, longList));

}

The Arithmetic method evaluates the value of a constant arithmetic expression.

public static Expression Arithmetic(MiddleOperator middleOp,

Expression leftExpression,

Expression rightExpression) {

If at least one of the operands is not constant, we just return null.

if (!IsConstant(leftExpression) || !IsConstant(rightExpression)) {

return null;

}

If at least one of the expressions has floating type, we call ArithmeticFloating.

else if (leftExpression.Symbol.Type.IsFloating() ||

rightExpression.Symbol.Type.IsFloating()) {

return ArithmeticFloating(middleOp, leftExpression, rightExpression);

}

If neither of the expressions has floating type, we call ArithmeticIntegral.

else {

return ArithmeticIntegral(middleOp, leftExpression, rightExpression);

}

}

The ArithmeticIntegral method evaluate the value of a constant integral arithmetic expression. It calls in turn the ArithmeticIntegral method that takes two symbols as parameters. The reason for using two methods is that the latter method is called by the middle code optimizer.

private static Expression ArithmeticIntegral(MiddleOperator middleOp,

Expression leftExpression,

Expression rightExpression){

We cast potential logical expression to integral type (signed integer).

leftExpression = LogicalToIntegral(leftExpression);

rightExpression = LogicalToIntegral(rightExpression);

Symbol symbol = ArithmeticIntegral(middleOp, leftExpression.Symbol,

rightExpression.Symbol);

The integral arithmetic evaluation does not generate any midddle code, which is why the expression has neither a short nor a long list.

return (new Expression(symbol, null, null));

}

The integral arithmetic evaluation is a bit more complicated since we must take pointer arithmetic in consideration.

public static Symbol ArithmeticIntegral(MiddleOperator middleOp,

Symbol leftSymbol,

Symbol rightSymbol) {

Type leftType = leftSymbol.Type,

rightType = rightSymbol.Type;

Since this method has been called, we know that none of the expression have floating type, and the potiential logical types has been cast to integral types. Therefore, we can assume that their values are BigInteger objects.

BigInteger leftValue = (BigInteger) leftSymbol.Value,

rightValue = (BigInteger) rightSymbol.Value,

resultValue = 0;

In case of addition, we must check whether the type of the left expression is a pointer or array type. In that case we need to multiply the value of the right expression with the size of the pointer or array type.

switch (middleOp) {

case MiddleOperator.BinaryAdd:

if (leftType.IsPointerOrArray()) {

resultValue = leftValue +

(rightValue \* leftType.PointerOrArrayType.Size());

}

We must also check whether the right type of the right expression is a pointer or array, in which case we multiply the left value with the size of the pointer or array type.

else if (leftType.IsPointerOrArray()) {

resultValue = (leftValue \* rightType.PointerOrArrayType.Size()) +

rightValue;

}

If neither the left nor the right expression holds pointer or array types, we simply add the values.

else {

resultValue = leftValue + rightValue;

}

break;

In case of subtraction, to begin with we must check if both the expressions hold pointer or array types. If they do, we subtract their values and divide the difference with size of the pointer or array type.

case MiddleOperator.BinarySubtract:

if (leftType.IsPointerOrArray() && rightType.IsPointerOrArray()) {

resultValue = (leftValue - rightValue) /

leftType.PointerOrArrayType.Size();

}

If the left expression (but not the right expression) holds pointer or array type, we must multiply the right value with the size of the pointer or array type before we subtract t it from the left value.

else if (leftType.IsPointerOrArray()) {

resultValue = leftValue -

(rightValue \* leftType.PointerOrArrayType.Size());

}

If neither the left nor the right expression holds pointer or array type, we simply subtract the values.

else {

resultValue = leftValue - rightValue;

}

break;

In case of multiplication, division, or modulo, bitwise, or shift operators, we simply perform the operation.

case MiddleOperator.SignedMultiply:

case MiddleOperator.UnsignedMultiply:

resultValue = leftValue \* rightValue;

break;

case MiddleOperator.SignedDivide:

case MiddleOperator.UnsignedDivide:

resultValue = leftValue / rightValue;

break;

case MiddleOperator.SignedModulo:

case MiddleOperator.UnsignedModulo:

resultValue = leftValue % rightValue;

break;

case MiddleOperator.BitwiseOr:

resultValue = leftValue | rightValue;

break;

case MiddleOperator.BitwiseXOr:

resultValue = leftValue ^ rightValue;

break;

case MiddleOperator.BitwiseAnd:

resultValue = leftValue & rightValue;

break;

In case of shift operators, we need to type cast the right value from BigInteger to int, simply because the BigInteger operations want an integer in those cases.

case MiddleOperator.ShiftLeft:

resultValue = leftValue << ((int) rightValue);

break;

case MiddleOperator.ShiftRight:

resultValue = leftValue >> ((int) rightValue);

break;

}

When we have evaluated the value, we need to find its type, which we do by calling the MaxType method in the TypeCast class.

Type maxType = TypeCast.MaxType(leftSymbol.Type, rightSymbol.Type);

Finally, we return a symbol holding the maximum type and the resulting value.

return (new Symbol(maxType, resultValue));

}

The evaluation of a floating constant expression is simplier than the integral expression, since we do not need to take pointer arithmetic in consideration.

private static Expression ArithmeticFloating(MiddleOperator middleOp,

Expression leftExpression,

Expression rightExpression) {

leftExpression = LogicalToFloating(leftExpression);

rightExpression = LogicalToFloating(rightExpression);

decimal leftValue, rightValue, resultValue = 0;

The value of each expression can be BigInteger or decimal. In the case of BigInteger, we first cast the value to BigInteger and then to decimal.

if (leftExpression.Symbol.Value is BigInteger) {

leftValue = (decimal) ((BigInteger) leftExpression.Symbol.Value);

}

else {

leftValue = (decimal) leftExpression.Symbol.Value;

}

if (rightExpression.Symbol.Value is BigInteger) {

rightValue = (decimal) ((BigInteger) rightExpression.Symbol.Value);

}

else {

rightValue = (decimal) rightExpression.Symbol.Value;

}

Since the values are object of the decimal class, we can user the regular operators.

switch (middleOp) {

case MiddleOperator.BinaryAdd:

resultValue = leftValue + rightValue;

break;

case MiddleOperator.BinarySubtract:

resultValue = leftValue - rightValue;

break;

case MiddleOperator.SignedMultiply:

resultValue = leftValue \* rightValue;

break;

case MiddleOperator.SignedDivide:

resultValue = leftValue / rightValue;

break;

}

When we have evaluated the value, we must decide its type, which we do by calling the MaxType method in the TypeCast class.

Type maxType = TypeCast.MaxType(leftExpression.Symbol.Type,

rightExpression.Symbol.Type);

The final symbol is a symbol with the maximum type and resulting value.

Symbol resultSymbol = new Symbol(maxType, resultValue);

In the floating evaluation case, we need to push the resulting value to the floating value stack.

List<MiddleCode> longList = new List<MiddleCode>();

MiddleCodeGenerator.AddMiddleCode(longList, MiddleOperator.PushFloat,

resultSymbol);

The final expression has no short list, and the long list hold the stack pushing instruction only.

return (new Expression(resultSymbol, null, longList));

}

The LogicalNot method evaluates the logical invers of the expression if it is constant.

public static Expression LogicalNot(Expression expression) {

if (IsConstant(expression)) {

expression = ToLogical(expression);

The resulting symbol is simply the original symbol with its true set and false set swapped.

Symbol resultSymbol = new Symbol(expression.Symbol.FalseSet,

expression.Symbol.TrueSet);

return (new Expression(resultSymbol, null, expression.LongList));

}

return null;

}

The Arithmetic method performs a unary arithmetic operation on a constant expression. If the expression is constant and floating ArithmeticFloating is called, and if it is constant and integral ArithmeticIntegral is called.

public static Expression Arithmetic(MiddleOperator middleOp,

Expression expression) {

if (!IsConstant(expression)) {

return null;

}

else if (expression.Symbol.Type.IsFloating()) {

return ArithmeticFloating(middleOp, expression);

}

else {

return ArithmeticIntegral(middleOp, expression);

}

}

The ArithmeticIntegral method performs XXX

private static Expression ArithmeticIntegral(MiddleOperator middleOp,

Expression expression) {

expression = LogicalToIntegral(expression);

BigInteger value = (BigInteger) expression.Symbol.Value,

resultValue = 0;

The unary addition operator does in fact nothing. The unary subtraction operator and bitwise not-operator perform the corresponding operations.

switch (middleOp) {

case MiddleOperator.UnaryAdd:

resultValue = value;

break;

case MiddleOperator.UnarySubtract:

resultValue = -value;

break;

case MiddleOperator.BitwiseNot:

resultValue = ~value;

break;

}

The resulting symbol hold the type of the origional expression and the resulting value.

Symbol resultSymbol = new Symbol(expression.Symbol.Type, resultValue);

The final expression does not have short list and long list, since the integral operations do not produce any middle code.

return (new Expression(resultSymbol, null, null));

}

The ArithmeticFloating method performs unary addition and subtraction. In case of addition, nothing happens, a new symbol with the same value is created.

private static Expression ArithmeticFloating(MiddleOperator middleOp,

Expression expression) {

expression = LogicalToFloating(expression);

decimal value = (decimal) expression.Symbol.Value;

Symbol resultSymbol = null;

switch (middleOp) {

case MiddleOperator.UnaryAdd:

resultSymbol = new Symbol(expression.Symbol.Type, value);

break;

case MiddleOperator.UnarySubtract:

resultSymbol = new Symbol(expression.Symbol.Type, -value);

break;

}

In the case of floating operators, we push the resulting value on the floating value stack.

List<MiddleCode> longList = new List<MiddleCode>();

MiddleCodeGenerator.AddMiddleCode(longList, MiddleOperator.PushFloat,

resultSymbol);

The resulting expression has no short list, and the long list holds only the floating stack push instruction.

return (new Expression(resultSymbol, null, longList));

}

The Cast method cast a constant expression into a new type. If the expression if not constant, we just return null.

public static Expression Cast(Expression sourceExpression,

Type targetType) {

if (!IsConstant(sourceExpression)) {

return null;

}

Symbol sourceSymbol = sourceExpression.Symbol;

Type sourceType = sourceSymbol.Type;

object sourceValue = sourceSymbol.Value;

If the source type equals the target type, we just return the origional expression.

if (sourceType.Equals(targetType)) {

return sourceExpression;

}

If the source and target types have the same size and they both are either integral, array, or pointers, or they both are floating, we just return a new expression with the target symbol and the short list and long list of the source expression.

else if ((sourceType.Size() == targetType.Size()) &&

((sourceType.IsIntegralArrayOrPointer() &&

targetType.IsIntegralArrayOrPointer()) ||

(sourceType.IsFloating() && targetType.IsFloating()))) {

Symbol targetSymbol = new Symbol(targetType, sourceValue);

return (new Expression(targetSymbol, sourceExpression.ShortList,

sourceExpression.LongList));

}

If the source type is logical and the target type is in integral, array, or pointer. We look into the true and false sets of the expression.

else if (sourceType.IsLogical() &&

targetType.IsIntegralArrayOrPointer()) {

If the true set is not empty, the expression is true and we return a new expression with the value one. We do not need to exam the false set, since we know that the expression is constant, in which case one of the sets is always empty.

if (sourceSymbol.TrueSet.Count > 0) {

Symbol targetSymbol = new Symbol(targetType, BigInteger.One);

return (new Expression(targetSymbol, null, null));

}

If the true set is empty, we return a new expression with the value zero, since the false set must be non-empty.

else {

Symbol targetSymbol = new Symbol(targetType, BigInteger.Zero);

return (new Expression(targetSymbol, null, null));

}

}

If the source type is integral, array, pointer, or floating, and the target type is logical, we start by constructing a jump set holding one unconditional jump instruction.

else if ((sourceType.IsIntegralArrayOrPointer() ||

sourceType.IsFloating()) &&

targetType.IsLogical()) {

List<MiddleCode> longList = new List<MiddleCode>();

MiddleCode jumpCode = new MiddleCode(MiddleOperator.Goto);

longList.Add(jumpCode);

ISet<MiddleCode> jumpSet = new HashSet<MiddleCode>();

jumpSet.Add(jumpCode);

If the value of the source symbol is zero, we create a target symbol with jump set as its false set. If it not zero, we create the target symbol with the jump set as its true set.

Symbol targetSymbol;

if (sourceValue.Equals(BigInteger.Zero) ||

sourceValue.Equals((decimal) 0)) {

targetSymbol = new Symbol(null, jumpSet);

}

else {

targetSymbol = new Symbol(jumpSet, null);

}

The long list holds the unconditional jump instruction of the jump set.

return (new Expression(targetSymbol, null, longList));

}

else {

Assert.ErrorXXX((sourceValue is BigInteger) ||

(sourceValue is decimal));

Finally, we have reached the point where both the source type and target type are integral, array, pointer, or floating. The only thing that remains to be done is to make sure the final value is an object of the correct class: BigInteger or decimal.

Symbol targetSymbol = null;

If the source type is integral, array, or pointer, and the target type is floating, we need to type cast the value from BigInteger to decimal.

if (sourceType.IsIntegralArrayOrPointer() &&targetType.IsFloating()) {

targetSymbol =

new Symbol(targetType, (decimal) ((BigInteger) sourceValue));

}

In the same way, if the source type is floating and the target type is integral, array, or pointer, we need to type cast the value from decimal to BigInteger.

else if (sourceType.IsFloating() &&

targetType.IsIntegralArrayOrPointer()) {

targetSymbol =

new Symbol(targetType, (BigInteger) ((decimal) sourceValue));

}

Finally, if the both the source and target types are integral, array, or pointer, or if they both are floating, we just create a new symbol with target type and source value.

else {

targetSymbol = new Symbol(targetType, sourceValue);

}

List<MiddleCode> longList = new List<MiddleCode>();

if (targetType.IsFloating()) {

MiddleCodeGenerator.

AddMiddleCode(longList, MiddleOperator.PushFloat, targetSymbol);

}

return (new Expression(targetSymbol, null, longList));

}

}

The Value method takes a symbol and returns a static symbol.

public static StaticSymbol Value(Symbol symbol) {

return Value(symbol.UniqueName, symbol.Type, symbol.Value);

}

public static StaticSymbol Value(string uniqueName, Type type,

object value) {

List<MiddleCode> middleCodeList = new List<MiddleCode>();

If the value is not null, we add the initializer instruction to the code list.

if (value != null) {

middleCodeList.Add(new MiddleCode(MiddleOperator.Initializer,

type.Sort, value));

}

If the value is null, we add the initializer ero instruction, with the type size.

else {

middleCodeList.Add(new MiddleCode(MiddleOperator.InitializerZero,

type.Size()));

}

We translate the middle code instructions to assembly code.

List<AssemblyCode> assemblyCodeList = new List<AssemblyCode>();

AssemblyCodeGenerator.GenerateAssembly(middleCodeList,

assemblyCodeList);

For the Linux target machine, we generate a list of text holding the final assembly code. For Windows target code generation, se Chapter XXX.

if (Start.Linux) {

List<string> textList = new List<string>();

textList.Add("\n" + uniqueName + ":");

ISet<string> externSet = new HashSet<string>();

AssemblyCodeGenerator.LinuxTextList(assemblyCodeList, textList, externSet);

return (new StaticSymbolLinux(StaticSymbolLinux.TextOrData.Data,

uniqueName, textList, externSet));

}

if (Start.Windows) {

List<byte> byteList = new List<byte>();

IDictionary<int,string> accessMap = new Dictionary<int,string>();

AssemblyCodeGenerator.

GenerateTargetWindows(assemblyCodeList, byteList,

accessMap, null, null);

return (new StaticSymbolWindows(uniqueName, byteList, accessMap));

}

return null;

}

}

}

# Static Address Expression

If an expression is not constant, the next step is to decide whether it is static.

The StaticValue and StaticAddress classes are identical, they hold a name and an offset. In the end, only static addresses are allowed, but a static value can hold a middle value. For instance, in the static address &a[3], a[3] is temporary stored as a static value.

StaticValue.cs

namespace CCompiler {

public class StaticValue {

private string m\_uniqueName;

private int m\_offset;

public StaticValue(string name, int offset) {

m\_uniqueName = name;

m\_offset = offset;

}

public string UniqueName {

get { return m\_uniqueName; }

}

public int Offset {

get { return m\_offset; }

}

}

}

StaticAddress.cs

namespace CCompiler {

public class StaticAddress {

private string m\_uniqueName;

private int m\_offset;

public StaticAddress(string name, int offset) {

m\_uniqueName = name;

m\_offset = offset;

}

public string UniqueName {

get { return m\_uniqueName; }

}

public int Offset {

get { return m\_offset; }

}

}

}

The StaticExpression class hold the two methods Binary and Unary, which takes a binary or unary expression and returns a static expression.

StaticExpression.cs

using System.Numerics;

namespace CCompiler {

public class StaticExpression {

The Binary method exams the expressions if the operator is binary addition or subtraction, index, or dot.

public static Expression Binary(MiddleOperator middleOp,

Expression leftExpression,

Expression rightExpression) {

Type leftType = leftExpression.Symbol.Type,

rightType = rightExpression.Symbol.Type;

object leftValue = leftExpression.Symbol.Value,

rightValue = rightExpression.Symbol.Value;

In the addition case, the operand values must be a static address and a constant integer value, or a extern or static array and a constant value. For instance, &i + 2, 2 + &i, a + 2 or 2 + a.

switch (middleOp) {

case MiddleOperator.BinaryAdd: // &i + 2, a + 2

if (((leftValue is StaticAddress) ||

(leftExpression.Symbol.IsExternOrStatic() &&

leftType.IsArray())) && (rightValue is BigInteger)) {

return GenerateAddition(leftExpression.Symbol,

(BigInteger) rightValue);

}

else if ((leftValue is BigInteger) && // 2 + &i, 2 + a

((rightValue is StaticAddress) || (rightType.IsArray() &&

rightExpression.Symbol.IsExternOrStatic()))) {

return GenerateAddition(rightExpression.Symbol,

(BigInteger) leftValue);

}

break;

In the subtraction case, the left operand must a static address as value or be an extern of static array, and the right operand must be a constant integer value. For instance, &i - 2 or a - 2.

case MiddleOperator.BinarySubtract: // &i - 2, a - 2

if (((leftValue is StaticAddress) ||

(leftExpression.Symbol.IsExternOrStatic() &&

leftType.IsArray())) &&

(rightValue is BigInteger)) {

return GenerateAddition(leftExpression.Symbol,

-((BigInteger) rightValue));

}

break;

In the index case, the operands must be an extern or static array or a static address, and a constant integer value.

case MiddleOperator.Index:

if (((leftValue is StaticAddress) || (leftType.IsArray() &&

leftExpression.Symbol.IsExternOrStatic())) &&

(rightValue is BigInteger)) { // a[2]

return GenerateIndex(leftExpression.Symbol,

(BigInteger) rightValue);

}

else if ((leftValue is BigInteger) && ((rightType.IsArray() &&

rightExpression.Symbol.IsExternOrStatic()) ||

(rightValue is StaticAddress))) {

return GenerateIndex(rightExpression.Symbol,

(BigInteger) leftValue);

}

break;

In the dot case, the XXX

case MiddleOperator.Dot:

if (leftExpression.Symbol.IsExternOrStatic()) {

object resultValue =

new StaticValue(leftExpression.Symbol.UniqueName,

rightExpression.Symbol.Offset); // s.i

Symbol resultSymbol = new Symbol(leftType, resultValue);

return (new Expression(resultSymbol, null, null));

}

break;

}

return null;

}

private static Expression GenerateAddition(Symbol symbol,

BigInteger value) {

int offset = ((int) value) \* symbol.Type.PointerOrArrayType.Size();

StaticAddress resultValue;

if (symbol.Value is StaticAddress) {

StaticAddress staticAddress = (StaticAddress) symbol.Value;

resultValue = new StaticAddress(staticAddress.UniqueName,

staticAddress.Offset + offset);

}

else {

resultValue = new StaticAddress(symbol.UniqueName, offset);

}

Symbol resultSymbol = new Symbol(symbol.Type, resultValue);

return (new Expression(resultSymbol, null, null));

}

private static Expression GenerateIndex(Symbol symbol,

BigInteger value) {

int offset = ((int) value) \* symbol.Type.ArrayType.Size();

StaticValue resultValue;

if (symbol.Value is StaticAddress) {

StaticAddress staticAddress = (StaticAddress) symbol.Value;

resultValue = new StaticValue(staticAddress.UniqueName,

staticAddress.Offset + offset);

}

else {

resultValue = new StaticValue(symbol.UniqueName, offset);

}

Symbol resultSymbol = new Symbol(symbol.Type, resultValue);

resultSymbol.Addressable = !symbol.IsRegister() &&

!symbol.Type.ArrayType.IsBitfield();

resultSymbol.Assignable =

!symbol.Type.ArrayType.IsConstantRecursive() &&

!symbol.Type.ArrayType.IsArrayFunctionOrString();

return (new Expression(resultSymbol, null, null));

}

Finally, we have to unary case. There is only one relevant operator: the address operator.

public static Expression Unary(MiddleOperator middleOp,

Expression expression) {

Symbol symbol = expression.Symbol;

If the symbol in the address case is a static value, we create a static address with the same name and offset.

switch (middleOp) {

case MiddleOperator.Address:

if (symbol.Value is StaticValue) { // &a[i], &s.i

StaticValue staticValue = (StaticValue) symbol.Value;

StaticAddress staticAddress =

new StaticAddress(staticValue.UniqueName, staticValue.Offset);

Symbol resultSymbol =

new Symbol(new Type(symbol.Type), staticAddress);

return (new Expression(resultSymbol, null, null));

}

If the symbol is extern or static, we create a static address, with the symbol name and offset zero.

else if (symbol.IsExternOrStatic()) { // &i

StaticAddress staticAddress =

new StaticAddress(symbol.UniqueName, 0);

Symbol resultSymbol =

new Symbol(new Type(symbol.Type), staticAddress);

return (new Expression(resultSymbol, null, null));

}

break;

}

}

return null;

}

}

}

# Initialization

There are two kinds of initialization: auto and static. Basically, they perform the same task: matches a type against an initializer. The initializer may be an expression or a list of expressions or lists.

GenerateAutoInitializer.cs

using System.Numerics;

using System.Collections.Generic;

namespace CCompiler {

class GenerateAutoInitializer {

public static List<MiddleCode> GenerateAuto(Symbol toSymbol,

object fromInitializer) {

Assert.ErrorXXX((fromInitializer is Expression) ||

(fromInitializer is List<object>));

Type toType = toSymbol.Type;

List<MiddleCode> codeList = new List<MiddleCode>();

If the initializer is an expression, we have two cases. If the type is an array of characters and the initializer is a string, we change the initializer from a string to a list of characters, and call GenerateAutoInitializer recursively with the list of characters instream of the string. In C, when initializing an array of characters, a string or a list of characters are equivalent. For instance, the initializations char s[] = "Hi" and char s[] = {'H', 'i', '\0'} are equivalent. Note the terminating zero-character is added implicitly and the string case.

if (fromInitializer is Expression) {

Expression fromExpression = (Expression) fromInitializer;

if (toType.IsArray() && toType.ArrayType.IsChar() &&

fromExpression.Symbol.Type.IsString()) {

string text = ((string) fromExpression.Symbol.Value) + "0";

List<object> list = new List<object>();

foreach (char c in text) {

Symbol charSymbol =

new Symbol(toType.ArrayType, (BigInteger) ((int) c));

Expression charExpression = new Expression(charSymbol, null, null);

list.Add(charExpression);

}

return GenerateAuto(toSymbol, list);

}

In all other cases, we type cast the expression into the defined type and generate middle code instructions. In case of a floating value, we pop the value from the floating value stack to the symbol to be initialized. In case of any other type we generate an assignment instruction.

else {

fromExpression = TypeCast.ImplicitCast(fromExpression, toType);

codeList.AddRange(fromExpression.LongList);

if (toSymbol.Type.IsFloating()) {

codeList.Add(new MiddleCode(MiddleOperator.PopFloat, toSymbol));

}

else {

codeList.Add(new MiddleCode(MiddleOperator.Assign, toSymbol,

fromExpression.Symbol));

}

}

}

If the initializer is not an expression is must be a list of expressions or other lists, recursively. In that case the type must be an array, a struct, or a union, since only arrays, structs, and unions can be initialized with a list.

else {

Assert.Error(toType.IsArray() ||toType.IsStructOrUnion(),

toType, Message.

Only\_array\_struct\_or\_union\_can\_be\_initialized\_by\_a\_list);

List<object> fromList = (List<object>) fromInitializer;

If the type is an array, we set its size if it has not yet been defined or check that the list size does not exceed the array size if the size has been defined. Note that in the case of an array we do not add a terminating zero character.

switch (toType.Sort) {

case Sort.Array: {

fromList = ModifyInitializer.ModifyArray(toType, fromList);

if (toType.ArraySize == 0) {

toType.ArraySize = fromList.Count;

}

else {

Assert.Error(fromList.Count <= toType.ArraySize,

toType, Message.Too\_many\_initializers);

}

We iterate through the list and, for each element in the list, define a symbol for the array index and call GenerateAuto recursively, so that potential sub list are properly processed. If the list does not cover the whole array, the remaining part of the array remains uninitialized. The offset of each index symbol is sum of the offset of the symbol and the offset of the index symbol. Note that we must multiply the index of the array with the size of the array type to determine the correct index. However, it is not strictly necessary give the index symbol a name. That is for readability reasons only.

for (int index = 0; index < fromList.Count; ++index) {

Symbol indexSymbol = new Symbol(toType.ArrayType, true);

indexSymbol.Offset = toSymbol.Offset +

(index \* toType.ArrayType.Size());

indexSymbol.Name = toSymbol.Name + "[" + index + "]";

codeList.AddRange(GenerateAuto(indexSymbol, fromList[index]));

}

}

break;

In case of a struct or a union, we check that the list size does not exceed its allowed size. In case of a struct, the list size must not exceed the number of struct members. In case of a union, the list must hold exactly on element.

case Sort.Struct:

case Sort.Union: {

IDictionary<string,Symbol> memberMap = toType.MemberMap;

Assert.Error((toType.IsStruct() &&

(fromList.Count <= memberMap.Count)) ||

(toType.IsUnion() && (fromList.Count == 1)),

toType, Message.Too\_many\_initializers);

Like the array case, we iterate through the list and, for each element in the list, create a symbol for the member and call GenerateAuto recursively, so that potential sub lists are properly processed. The offset of each member symbol is the sum of the offset of the struct or union and the member.

IEnumerator<Symbol> enumerator =

memberMap.Values.GetEnumerator();

foreach (object fromInitializor in fromList) {

enumerator.MoveNext();

Symbol memberSymbol = enumerator.Current;

Symbol subSymbol = new Symbol(memberSymbol.Type, true);

subSymbol.Name = toSymbol.Name + "." + memberSymbol.Name;

subSymbol.Offset = toSymbol.Offset + memberSymbol.Offset;

codeList.AddRange(GenerateAuto(subSymbol, fromInitializor));

}

}

break;

}

}

return codeList;

}

}

}

The GenerateStaticInitializer class basically performs the same task as the GenerateAutoInitializer class: it interprets the initialization and compare it to the defined type. However, it does not generate assignment instruction, but rather initialization instructions. One difference compared to the auto initialization case is that if the array, struct, or union, is not completely initialized we need to generate code for initialization of the remaining part.

GenerateStaticInitializer.cs

using System.Collections.Generic;

namespace CCompiler {

public class GenerateStaticInitializer {

public static List<MiddleCode> GenerateStatic(Type toType,

object fromInitializer) {

Assert.ErrorXXX((fromInitializer is Expression) ||

(fromInitializer is List<object>));

List<MiddleCode> codeList = new List<MiddleCode>();

If the initializer is an expression, we have two special cases: that the defined type is a character array and the initializer is a string, the defined type is a pointer and the initializer is an array, a function, or a string.

if (fromInitializer is Expression) {

Expression fromExpression = (Expression) fromInitializer;

Symbol fromSymbol = fromExpression.Symbol;

Assert.Error(fromSymbol.IsExternOrStatic(), fromSymbol,

Message.Non\_\_static\_initializer);

Type fromType = fromSymbol.Type;

If the defined type is a character array and the initializer is a string, we first check the size of the string. If the array size is undefined, we set its size to the size of the string, plus one for the terminating zero-character. If the array size is defined, we check that the string (including the terminating zero-character) fits withing the array.

if (toType.IsArray() && toType.ArrayType.IsChar() &&

fromType.IsString()) {

string text = (string) fromSymbol.Value;

Note that we add space for an extra character, the terminating zero-character.

if (toType.ArraySize == 0) {

toType.ArraySize = text.Length + 1;

}

Note that the string length must be strictly less than the array size, for the terminating zero-character to fit in the string.

else {

Assert.Error(text.Length < toType.ArraySize,

toType, Message.Too\_many\_initializers);

}

codeList.Add(new MiddleCode(MiddleOperator.Initializer,

fromSymbol.Type.Sort, text));

}

If the defined type is a pointer and the initializer is an array, a function, or a string, we create a static address and add an initialization middle code instruction.

else if (toType.IsPointer() && fromType.IsArrayFunctionOrString()) {

Assert.ErrorXXX((fromType.IsString() && toType.PointerType.IsChar())

||(fromType.IsArray() &&

fromType.ArrayType.Equals(toType.PointerType)) ||

(fromType.IsFunction() &&

fromType.Equals(toType.PointerType)));

StaticAddress staticAddress =

new StaticAddress(fromSymbol.UniqueName, 0);

codeList.Add(new MiddleCode(MiddleOperator.Initializer,

toType.Sort, staticAddress));

}

In all other cases, we type cast the expression into the defined type, check that the expression is constant, and add an initialization middle code instruction.

else {

Expression toExpression =

TypeCast.ImplicitCast(fromExpression, toType);

Symbol toSymbol = toExpression.Symbol;

If the value is not null, the expression is constant. However, the value may be a static address, which we consider to be constant in this context.

Assert.Error(toSymbol.Value != null, toSymbol,

Message.Non\_\_constant\_expression);

codeList.Add(new MiddleCode(MiddleOperator.Initializer,

toSymbol.Type.Sort, toSymbol.Value));

}

}

If the initializer is a list, the defined type must be an array, a struct, or a union.

else {

Assert.Error(toType.IsArray() || toType.IsStructOrUnion(),

toType, Message.

Only\_array\_struct\_or\_union\_can\_be\_initialized\_by\_a\_list);

List<object> fromList = (List<object>) fromInitializer;

In case of an arry, we set the array size to the list size if undefined. If the array is defined, we check thet the list size does not exceed the array size.

switch (toType.Sort) {

case Sort.Array: {

fromList = ModifyInitializer.ModifyArray(toType, fromList);

if (toType.ArraySize == 0) {

toType.ArraySize = fromList.Count;

}

else {

Assert.Error(fromList.Count <= toType.ArraySize,

toType, Message.Too\_many\_initializers);

}

We iterate the array values and call GenerateStatic for each value. In the static case we do not need to create index symbol since the code to be generated are initializations rather than assignments.

foreach (object value in fromList) {

codeList.AddRange(GenerateStatic(toType.ArrayType, value));

}

However, in the static case we must add an instruction for the initialization of the potential remaining part of the array. The InitializerZero instruction adds a sequence of zero-bytes.

int restSize = toType.Size() -

(fromList.Count \* toType.ArrayType.Size());

if (restSize > 0) {

codeList.Add(new MiddleCode(MiddleOperator.InitializerZero,

restSize));

}

}

break;

Like the auto case, the initializer list size must not exceed the number of members in a struct. The initializer list must hold exactly one element in case of a union.

case Sort.Struct:

case Sort.Union: {

IDictionary<string,Symbol> memberMap = toType.MemberMap;

Assert.Error((toType.IsStruct() &&

(fromList.Count <= memberMap.Count)) ||

(toType.IsUnion() && (fromList.Count == 1)),

toType, Message.Too\_many\_initializers);

Like the auto case, we iterate throught the inielization list and call GenerateStatic recursivly for each element in the list. However, in the static case we must sum the size of the initializer list in order to add zero-bytes if the list does not cover the struct or union.

int size = 0;

IEnumerator<Symbol> enumerator =

memberMap.Values.GetEnumerator();

foreach (object fromInitializor in fromList) {

enumerator.MoveNext();

Symbol memberSymbol = enumerator.Current;

codeList.AddRange(GenerateStatic(memberSymbol.Type,

fromInitializor));

size += memberSymbol.Type.Size();

}

int restSize = toType.Size() - size;

if (restSize > 0) {

codeList.Add(new MiddleCode(MiddleOperator.InitializerZero,

restSize));

}

}

break;

}

}

return codeList;

}

}

}

The ModifyInitializer class changes an initialzation list into the form of the defined type. For instance: int [3][2] = {1, 2, 3, 4, 5, 6} is changed to int [3][2] = {{1, 2}, {3, 4}, {5, 6}}.

ModifyInitializer.cs

using System;

using System.IO;

using System.Collections.Generic;

namespace CCompiler {

class ModifyInitializer {

public static List<object> ModifyArray(Type type, List<object> list) {

First, we define the dimension-to-size map; that is, each dimension in the defined array is assigned an array size. For instance, in int [2][3][4] dimension 1 has size 4, dimension 2 has size 3, and dimension 3 has size 2. Dimension zero refers to the final array type, and its size is always zero. In int [][3][4], dimension 3 is undefined and set to zero, but that does not matter since we do not exam the highest dimension.

IDictionary<int,int> dimensionToSizeMap = new Dictionary<int,int>();

int maxDimension = DimensionToSizeMap(type, dimensionToSizeMap);

Then we define the initializer-to-dimension map; that is, each list and sub list in the initializer is assigned a dimension.

IDictionary<object,int> initializerToDimensionMap =

new Dictionary<object,int>();

InitializerToDimensionMap(list, initializerToDimensionMap);

Then we iterate throught the dimensions of the array, from dimension one to the second-highest dimension, inclusive. For each dimension we define the total list, where each element holds the current dimension.

for (int dimension = 1; dimension < maxDimension; ++dimension) {

List<object> totalList =

new List<object>(), currentList = new List<object>();

int arraySize = dimensionToSizeMap[dimension];

Assert.ErrorXXX(arraySize > 0);

foreach (object member in list) {

if (initializerToDimensionMap[member] < dimension) {

currentList.Add(member);

}

else {

if (currentList.Count > 0) {

initializerToDimensionMap[currentList] = dimension;

totalList.Add(currentList);

}

totalList.Add(member);

currentList = new List<object>();

}

if (currentList.Count == arraySize) {

initializerToDimensionMap[currentList] = dimension;

totalList.Add(currentList);

currentList = new List<object>();

}

}

if (currentList.Count > 0) {

initializerToDimensionMap[currentList] = dimension;

totalList.Add(currentList);

}

list = totalList;

}

return list;

}

The DimensionToSizeMap method assigns the array size of each dimension in the nested array.

private static int DimensionToSizeMap(Type type,

IDictionary<int,int> dimensionToSizeMap) {

if (type.IsArray()) {

int dimension =

DimensionToSizeMap(type.ArrayType, dimensionToSizeMap) + 1;

dimensionToSizeMap[dimension] = type.ArraySize;

}

return 0;

}

The InitializerToDimensionMap method assigns the dimension to each list in the nested list.

private static int InitializerToDimensionMap(object initializer,

IDictionary<object,int> initializerToDimensionMap) {

if (initializer is List<object>) {

List<object> list = (List<object>) initializer;

int maxDimension = 0;

If the initializer is a list, we iterate throught the list, and assign the list the dimesion of the sub list with dhe higest dimension, plus one.

foreach (object member in list) {

int dimension =

InitializerToDimensionMap(member, initializerToDimensionMap);

maxDimension = Math.Max(maxDimension, dimension);

}

initializerToDimensionMap[list] = maxDimension + 1;

return maxDimension + 1;

}

return 0;

}

}

}

# Middle Code Optimization

When the middle code has been generated, we need to perform several optimizations since the code may be ineffective. Some ineffective parts may be introduced by the programmer, but most parts are likely to be introduced by the parser. The MiddleCodeOptimization class takes care of the optimization.

There is a set of optimizations available:

* We clear goto-next-statements: jumps to the next line.
* We also modify goto-next-double-statements: an conditional jump instruction that jumps two steps ahead and is followed by an unconditional jump instruction.
* We trace goto-chains: jump instructions that jump to other jump instructions.
* We clear unreachable code: code that is not reachable from the first function instruction.
* We remove empty code: code that has been cleared by the optimization above.

Since the optimization methods may reveal new optimization opportunities we need to repeat them until we do not detect any more opportunity. The field m\_update is set to true if we find an optimization opportunity.

MiddleCodeOptimizor.cs

using System.Numerics;

using System.Collections.Generic;

namespace CCompiler {

public class MiddleCodeOptimizer {

private bool m\_update;

private List<MiddleCode> m\_middleCodeList;

public MiddleCodeOptimizer(List<MiddleCode> middleCodeList) {

m\_middleCodeList = middleCodeList;

}

public void Optimize() {

ObjectToIntegerAddresses();

do {

m\_update = false;

ClearGotoNextStatements();

ClearDoubleRelationStatements();

TraceGotoChains();

ClearUnreachableCode();

RemovePushPop();

MergePopPushToTop();

MergeTopPopToPop();

//MergeBinary(); // XXX

//MergeDoubleAssign(); // XXX

SematicOptimization();

OptimizeRelation();

OptimizeCommutative();

OptimizeBinary();

CheckIntegral(); // XXX

CheckFloating(); // XXX

RemoveClearedCode();

} while (m\_update);

}

### Object to Integer Addresses

When we generated the middle code, we use middle code objects as target for jump instructions. The first we need to do is to change the middle code targets to integer targets. The reason for this is that the optimization process of this chapter will result in the removal of middle code instruction, among which some may be jump targets.

MiddleCodeOptimizor.cs

public void ObjectToIntegerAddresses() {

We start by defining the address map, which we load with the index of each instruction.

IDictionary<MiddleCode,int> addressMap =

new Dictionary<MiddleCode,int>();

for (int index = 0; index < m\_middleCodeList.Count; ++index) {

addressMap.Add(m\_middleCodeList[index], index);

}

We then iterate through the middle code list and, with the help of the address map, change the targets from middle code instructions to integer values.

for (int index = 0; index < m\_middleCodeList.Count; ++index) {

MiddleCode sourceCode = m\_middleCodeList[index];

if (sourceCode.IsGoto() || sourceCode.IsCarry() ||

sourceCode.IsRelation()) {

Assert.ErrorXXX(sourceCode[0] is MiddleCode);

MiddleCode targetCode = (MiddleCode) sourceCode[0];

Assert.ErrorXXX(addressMap.ContainsKey(targetCode));

sourceCode[0] = addressMap[targetCode];

}

}

}

### Goto Next Instructions

In some cases, there may be goto instructions that just jump to the next instructions. Those instructions are meaningless and shall be removed. For instance:

1. goto 2

2. ...

MiddleCodeOptimizor.cs

private void ClearGotoNextStatements() {

for (int index = 1; index < (m\_middleCodeList.Count - 1); ++index) {

MiddleCode middleCode = m\_middleCodeList[index];

if (middleCode.IsRelationCarryOrGoto()) {

int target = (int) middleCode[0];

When we encounter a jump instruction that jumps to the next instruction, we remove it

if (target == (index + 1)) {

middleCode.Clear();

m\_update = true;

}

}

}

}

### Next-Double Goto Statements

A conditional jump instruction that jumps two steps ahead and is followed by an unconditional jump instruction can be modified in that we reverse the condition and the target.

|  |  |  |
| --- | --- | --- |
| 1. if a < b goto 3  2. goto 10  3. ... | 1. if a >= b goto 10  2. *removed*  3. ... |  |

To begin with, we need the inverse map to change the condition. The swap map is used in OptimizeRelation below. The reason we define it at this point in the code is that we cannot have more than one static area in the same class.

MiddleCodeOptimizor.cs

public static IDictionary<MiddleOperator, MiddleOperator> m\_inverseMap =

new Dictionary<MiddleOperator, MiddleOperator>() {

{MiddleOperator.Equal, MiddleOperator.NotEqual},

{MiddleOperator.NotEqual, MiddleOperator.Equal},

{MiddleOperator.Carry, MiddleOperator.NotCarry},

{MiddleOperator.NotCarry, MiddleOperator.Carry},

{MiddleOperator.SignedLessThan,

MiddleOperator.SignedGreaterThanEqual},

{MiddleOperator.SignedLessThanEqual,

MiddleOperator.SignedGreaterThan},

{MiddleOperator.SignedGreaterThan,

MiddleOperator.SignedLessThanEqual},

{MiddleOperator.SignedGreaterThanEqual,

MiddleOperator.SignedLessThan},

{MiddleOperator.UnsignedLessThan,

MiddleOperator.UnsignedGreaterThanEqual},

{MiddleOperator.UnsignedLessThanEqual,

MiddleOperator.UnsignedGreaterThan},

{MiddleOperator.UnsignedGreaterThan,

MiddleOperator.UnsignedLessThanEqual},

{MiddleOperator.UnsignedGreaterThanEqual,

MiddleOperator.UnsignedLessThan}

};

We iterate through the middle code list, from the first to the next to last instruction since we inspect the current instruction and the instruction following it.

private void ClearDoubleRelationStatements() {

for (int index = 0; index < (m\_middleCodeList.Count - 1); ++index) {

MiddleCode thisCode = m\_middleCodeList[index],

nextCode = m\_middleCodeList[index + 1];

if ((thisCode.IsRelation() || thisCode.IsCarry()) &&

nextCode.IsGoto()) {

int target1 = (int) thisCode[0],

target2 = (int) nextCode[0];

If the instruction jumps over the next instruction, we use the inverse map to change the current instruction, and we clear the next instruction.

if (target1 == (index + 2)) {

MiddleOperator operator1 = thisCode.Operator;

thisCode.Operator = m\_inverseMap[operator1];

thisCode[0] = target2;

nextCode.Clear();

m\_update = true;

}

}

}

}

### Goto Chains

A goto chain is a sequence of unconditional jump instructions where each instruction in the chain except the last one jumps to another unconditional jump instruction. It would be more effective if they all jump to the same target as the last instruction.

|  |  |  |
| --- | --- | --- |
| 1. goto 3  2. ..  3. goto 5  4. ...  5. goto 7  6. ...  7. a = 1 | 1. goto 7  2. ..  3. goto 7  4. ...  5. goto 7  6. ...  7. a = 1 |  |

First, we trace all the unconditional jump instructions by calling TraceGotoChains, that in turn calls TraceGoto. then we change all the jump targets.

MiddleCodeOptimizor.cs

private void TraceGotoChains() {

for (int index = 1; index < m\_middleCodeList.Count; ++index) {

MiddleCode middleCode = m\_middleCodeList[index];

if (middleCode.IsRelationCarryOrGoto()) {

ISet<int> sourceSet = new HashSet<int>();

sourceSet.Add(index);

For each jump instruction we trace the goto chain and if the first and last jump instructions differs, we replace all the jumps in the chain with the last jump.

int firstTarget = (int) middleCode[0];

int finalTarget = TraceGoto(firstTarget, sourceSet);

if (firstTarget != finalTarget) {

foreach (int source in sourceSet) {

MiddleCode sourceCode = m\_middleCodeList[source];

sourceCode[0] = finalTarget;

}

m\_update = true;

}

}

}

}

The TraceGoto method follows the jump instructions as long as they continue to jump to new targets.

private int TraceGoto(int target, ISet<int> sourceSet) {

MiddleCode objectCode = m\_middleCodeList[target];

if (!sourceSet.Contains(target) && objectCode.IsGoto()) {

sourceSet.Add(target);

int nextTarget = (int) objectCode[0];

return TraceGoto(nextTarget, sourceSet);

}

else {

return target;

}

}

### Remove Unreachable Code

Code that is not reachable from the first instruction of the function shall be removed. We call SearchReachableCode that follows all conditional and unconditional jumps instructions and save their line numbers in the visited set. We also warn if we reach the last instruction of a function that do not returns void.

MiddleCodeOptimizor.cs

private void ClearUnreachableCode() {

ISet<MiddleCode> visitedSet = new HashSet<MiddleCode>();

SearchReachableCode(0, visitedSet);

for (int index = 0; index < (m\_middleCodeList.Count - 1); ++index) {

MiddleCode middleCode = m\_middleCodeList[index];

if (!visitedSet.Contains(middleCode)) {

m\_middleCodeList[index].Clear();

m\_update = true;

}

}

}

private void SearchReachableCode(int index, ISet<MiddleCode> visitedSet) {

for (; index < m\_middleCodeList.Count; ++index) {

MiddleCode middleCode = m\_middleCodeList[index];

if (visitedSet.Contains(middleCode)) {

return;

}

visitedSet.Add(middleCode);

if (middleCode.IsRelation() || middleCode.IsCarry()) {

int target = (int) middleCode[0];

SearchReachableCode(target, visitedSet);

}

else if (middleCode.IsGoto()) {

int target = (int) middleCode[0];

SearchReachableCode(target, visitedSet);

return;

}

else if ((middleCode.Operator == MiddleOperator.Exit) ||

(middleCode.Operator == MiddleOperator.Return)) {

return;

}

else if (middleCode.Operator == MiddleOperator.FunctionEnd) {

Symbol funcSymbol = (Symbol) middleCode[0];

Assert.Error(funcSymbol.Type.ReturnType.IsVoid(),

funcSymbol.Name,

Message.Reached\_the\_end\_of\_a\_non\_\_void\_function);

return;

}

}

}

### Remove Push-Pop Chains

Sometimes there may be a push followed by a pop of the same symbol, or no symbol at all, which in meaningless and shall be removed.

|  |  |  |
| --- | --- | --- |
| 1. push x  2. pop x  1. push x  2. pop | 1. *empty*  2. *empty*  1. *empty*  2. *empty* |  |

However, we do not remove code where different symbols are pushed and popped.

|  |  |  |
| --- | --- | --- |
| 1. push x  2. pop y |  |  |

MiddleCodeOptimizor.cs

public void MergePopPushToTop() {

for (int index = 0; index < (m\_middleCodeList.Count - 1); ++index) {

MiddleCode thisCode = m\_middleCodeList[index],

nextCode = m\_middleCodeList[index + 1];

if ((thisCode.Operator == MiddleOperator.PopFloat) &&

(nextCode.Operator == MiddleOperator.PushFloat) &&

(thisCode[0] == nextCode[0])) {

thisCode.Operator = MiddleOperator.TopFloat;

nextCode.Clear();

m\_update = true;

}

}

}

### Change Top-Pop to Pop

Sometimes there may be a push followed by a pop of the same symbol, or no symbol at all, which in meaningless and shall be removed.

|  |  |  |
| --- | --- | --- |
| 1. top x  2. pop | 1. pop x  2. *empty* |  |

MiddleCodeOptimizor.cs

public void MergeTopPopToPop() {

for (int index = 0; index < (m\_middleCodeList.Count - 1); ++index) {

MiddleCode thisCode = m\_middleCodeList[index],

nextCode = m\_middleCodeList[index + 1];

if ((thisCode.Operator == MiddleOperator.TopFloat) &&

(nextCode.Operator == MiddleOperator.PopFloat)) {

Assert.ErrorXXX(nextCode[0] == null);

thisCode.Operator = MiddleOperator.PopFloat;

nextCode.Clear();

m\_update = true;

}

}

}

### Sematic Optimization

The semantic opimization handles

MiddleCodeOptimizor.cs

private void SematicOptimization() {

for (int index = 0; index < m\_middleCodeList.Count; ++index) {

MiddleCode thisCode = m\_middleCodeList[index];

if (thisCode.IsBinary()) {

Symbol resultSymbol = (Symbol) thisCode[0],

leftSymbol = (Symbol) thisCode[1],

rightSymbol = (Symbol) thisCode[2],

newSymbol = null;

The pointer arithmentic during the parsing may have caused constant expression, expression where both operands are constant values. In that case call ArithmeticIntegral in ConstantExpression to evaluate the value of the expression.

if ((leftSymbol.Value is BigInteger) && // t0 = 2 \* 3

(rightSymbol.Value is BigInteger)) {

newSymbol =

ConstantExpression.ArithmeticIntegral(thisCode.Operator,

leftSymbol, rightSymbol);

}

In case of additions where the left operand is zero, 0 + i, we keep the right operand.

// t0 = 0 + i

else if ((thisCode.Operator == MiddleOperator.BinaryAdd) &&

(leftSymbol.Value is BigInteger) &&

(leftSymbol.Value.Equals(BigInteger.Zero))) {

newSymbol = rightSymbol;

}

In case of additions or subtractions where the right operand is zero, i + 0 or i - 0, we keep the left operand.

// t0 = i + 0

// t0 = i - 0

else if (((thisCode.Operator == MiddleOperator.BinaryAdd) ||

(thisCode.Operator == MiddleOperator.BinarySubtract)) &&

(rightSymbol.Value is BigInteger) &&

(rightSymbol.Value.Equals(BigInteger.Zero))) {

newSymbol = leftSymbol;

}

In case of multiplications where the left operand is zero, the result is zero.

// t0 = 0 \* i

else if (((thisCode.Operator == MiddleOperator.SignedMultiply) ||

(thisCode.Operator == MiddleOperator.UnsignedMultiply)) &&

(leftSymbol.Value is BigInteger) &&

(leftSymbol.Value.Equals(BigInteger.Zero))) {

newSymbol = new Symbol(resultSymbol.Type, BigInteger.Zero);

}

In case of multiplications where the left operand is one, we keep the right operand.

// t0 = 1 \* i

else if (((thisCode.Operator == MiddleOperator.SignedMultiply) ||

(thisCode.Operator == MiddleOperator.UnsignedMultiply)) &&

(leftSymbol.Value is BigInteger) &&

(leftSymbol.Value.Equals(BigInteger.One))) {

newSymbol = rightSymbol;

}

In case of multiplications where the right operand is zero, the result is zero.

// t0 = i \* 0

else if (((thisCode.Operator == MiddleOperator.SignedMultiply) ||

(thisCode.Operator == MiddleOperator.UnsignedMultiply)) &&

(rightSymbol.Value is BigInteger) &&

(rightSymbol.Value.Equals(BigInteger.Zero))) {

newSymbol = new Symbol(resultSymbol.Type, BigInteger.Zero);

}

In case of multiplications or division where the right operand is one, we keep the left operand.

// t0 = i \* 1

// t0 = i / 1

else if (((thisCode.Operator == MiddleOperator.SignedMultiply) ||

(thisCode.Operator == MiddleOperator.UnsignedMultiply) ||

(thisCode.Operator == MiddleOperator.SignedDivide) ||

(thisCode.Operator == MiddleOperator.UnsignedDivide)) &&

(rightSymbol.Value is BigInteger) &&

(rightSymbol.Value.Equals(BigInteger.One))) {

newSymbol = leftSymbol;

}

If the new symbol is not null, we replace the expression with the new symbol. If the result symbol is a temporary symbol. We iterate until we find the place where the result symbol is accessed, and replace the result symbol with the new symbol at that place.

if (newSymbol != null) {

if (resultSymbol.Temporary) {

thisCode.Operator = MiddleOperator.Empty;

int index2;

for (index2 = index + 1; index2 < m\_middleCodeList.Count;

++index2) {

MiddleCode nextCode = m\_middleCodeList[index2];

if (nextCode[1] == resultSymbol) {

nextCode[1] = newSymbol;

break;

}

if (nextCode[2] == resultSymbol) {

nextCode[2] = newSymbol;

break;

}

}

Assert.ErrorXXX(index2 < m\_middleCodeList.Count);

}

If the result symbol is not temporary symbol, we change the instruction into an assignment, where the result symbol is assigned to the new symbol.

else {

thisCode.Operator = MiddleOperator.Assign; // i = 0 + j;

thisCode[1] = newSymbol; // i = j;

thisCode[2] = null;

}

m\_update = true;

}

}

}

}

### Optimize Relation Expression

In order to optimize the final assembly code generation, we shall swap the operator in expression where the left operand is a value.

|  |  |  |
| --- | --- | --- |
| 1. if 10 < i goto 20 | 1. if i > 10 goto 20 |  |

When generating the final assembly code, we cannot have an integer value as the left expression in a relational expression. Therefore, we swap the operands if the left operand holds an integer value. The expression cannot hold two integer values, in that case the ConstantExpression class of Section XXX would have reduced the expression to its resulting value.

Moreover, if the left expression is an array, function, or string, we want to use its address rather than its value. Similar to the integer value case, we cannot use the address directly in the assembly code. Therefore, if the left expression is an array, function or string, and the right expression is not an array, function, or string, or an integer value, we also swap the expressions.

We use the m\_swapMap map to swap the operator. Note that this is not the same case as the m\_inverseMap we used in the Next-Goto-Double case above. When inversing an expression, we change its meaning, when we swap the expression it still has the same meaning, we just arrange the operands so that we cane generate more efficient assembly code.

MiddleCodeOptimizor.cs

public static IDictionary<MiddleOperator, MiddleOperator> m\_swapMap =

new Dictionary<MiddleOperator, MiddleOperator>() {

{MiddleOperator.Equal, MiddleOperator.Equal},

{MiddleOperator.NotEqual, MiddleOperator.NotEqual},

{MiddleOperator.SignedLessThan, MiddleOperator.SignedGreaterThan},

{MiddleOperator.SignedGreaterThan, MiddleOperator.SignedLessThan},

{MiddleOperator.SignedLessThanEqual,

MiddleOperator.SignedGreaterThanEqual},

{MiddleOperator.SignedGreaterThanEqual,

MiddleOperator.SignedLessThanEqual},

{MiddleOperator.UnsignedLessThan, MiddleOperator.UnsignedGreaterThan},

{MiddleOperator.UnsignedGreaterThan, MiddleOperator.UnsignedLessThan},

{MiddleOperator.UnsignedLessThanEqual,

MiddleOperator.UnsignedGreaterThanEqual},

{MiddleOperator.UnsignedGreaterThanEqual,

MiddleOperator.UnsignedLessThanEqual}

};

private void SwapRelation() {

foreach (MiddleCode middleCode in m\_middleCodeList) {

if (middleCode.IsRelation()) {

Symbol leftSymbol = (Symbol) middleCode[1],

rightSymbol = (Symbol) middleCode[2];

if ((leftSymbol.Value is BigInteger) ||

(leftSymbol.Type.IsArrayFunctionOrString()) &&

!rightSymbol.Type.IsArrayFunctionOrString() &&

!(rightSymbol.Value is BigInteger)) {

middleCode.Operator = m\_swapMap[middleCode.Operator];

middleCode[1] = rightSymbol;

middleCode[2] = leftSymbol;

}

}

}

}

### Optimize Communicative Expression

MiddleCodeOptimizor.cs

private void OptimizeCommutative() {

foreach (MiddleCode middleCode in m\_middleCodeList) {

if (middleCode.IsCommutative()) {

Symbol leftSymbol = (Symbol) middleCode[1],

rightSymbol = (Symbol) middleCode[2];

// not 1 - i

if (leftSymbol.Type.IsIntegralPointerArrayOrFunction() &&

(leftSymbol.Value is BigInteger)) {

middleCode[1] = rightSymbol;

middleCode[2] = leftSymbol;

}

}

}

}

### Remove Trivial Assignment

Trivial assignment, where a symbol is assigned the value of the same symbol, shall be removed.

MiddleCodeOptimizor.cs

private void RemoveTrivialAssign() {

foreach (MiddleCode middleCode in m\_middleCodeList) {

MiddleOperator middleOperator = middleCode.Operator;

if (middleOperator == MiddleOperator.Assign) {

Symbol resultSymbol = (Symbol) middleCode[0],

assignSymbol = (Symbol) middleCode[1];

if (resultSymbol == assignSymbol) {

middleCode.Operator = MiddleOperator.Empty;

m\_update = true;

}

}

}

}

### Remove Cleared Code

Technically, it is not strictly necessary to remove cleared middle code instructions since it does not generate target code. However, we can simplify the optimization methods if we do so.

MiddleCodeOptimizor.cs

public void RemoveClearedCode() {

We iterate backwards through the middle code list for performance reasons.

for (int index = (m\_middleCodeList.Count - 1); index > 0;--index){

if (m\_middleCodeList[index].Operator == MiddleOperator.Empty) {

For each empty instruction we must go through all other instruction and decrease all target by one that is greater than the index of the instruction to be removed.

foreach (MiddleCode middleCode in m\_middleCodeList) {

if (middleCode.IsRelationCarryOrGoto()) {

int target = (int) middleCode[0];

if (target > index) {

middleCode[0] = target - 1;

}

}

}

m\_middleCodeList.RemoveAt(index);

}

}

}

# Assembly Code Generation

When we have generated and optimized the middle code, it is time to generate the final assembly code. The first step is to generate the assembly code with tracks. A track is a place hold for a register that follows the (yet unknown) register through the code. The track is then replaced by a proper register by the register allocation process. Finally, the assembly code instruction are written in plain text.

## Assembly Operator

Like the middle code operators, we also have the assembly code operators. Several operators comes in several varieties. For instance, the add instruction comes in the base form, which is used when a register is assigned a value, like add ax, 123. The other varieties are used when there is no register to assign, but we still have to specify the size of the value to be assigned, like add word [bp + 2], 123.

AssemblyOperator.cs

namespace CCompiler {

public enum AssemblyOperator {

add, add\_byte, add\_dword, add\_qword, add\_word, return\_address,

and, and\_byte, and\_dword, and\_qword, and\_word, call,

cmp, cmp\_byte, cmp\_dword, cmp\_qword, cmp\_word, comment,

dec, dec\_byte, dec\_dword, dec\_qword, dec\_word,

define\_address, define\_value, define\_zero\_sequence,

div, div\_byte, div\_dword, div\_qword, div\_word, empty,

fabs, fadd, faddp, fchs, fcompp, fdiv, fdivp, fdivr,

fdivrp, fild\_dword, fild\_qword, fild\_word,

fist\_dword, fist\_qword, fist\_word,

fistp\_dword, fistp\_qword, fistp\_word,

fld1, fld\_dword, fld\_qword, fldcw, fldz,

fmul, fmulp, fst\_dword, fst\_qword, fstcw,

fstp\_dword, fstp\_qword, fstsw,

fsub, fsubp, fsubr, fsubrp, ftst,

idiv, idiv\_byte, idiv\_dword, idiv\_qword, idiv\_word,

imul, imul\_byte, imul\_dword, imul\_qword, imul\_word,

inc, inc\_byte, inc\_dword, inc\_qword, inc\_word, interrupt,

ja, jae, jb, jbe, jc, je, jg, jge, jl, jle, jmp, jnc, jne, jnz, jz,

label, lahf, mov, mov\_byte, mov\_dword, mov\_qword, mov\_word,

mul, mul\_byte, mul\_dword, mul\_qword, mul\_word,

neg, neg\_byte, neg\_dword, neg\_qword, neg\_word, new\_middle\_code,

nop, not, not\_byte, not\_dword, not\_qword, not\_word,

or, or\_byte, or\_dword, or\_qword, or\_word, pop, ret, sahf, set\_track\_size,

shl, shl\_byte, shl\_dword, shl\_qword, shl\_word,

shr, shr\_byte, shr\_dword, shr\_qword, shr\_word,

sub, sub\_byte, sub\_dword, sub\_qword, sub\_word, syscall,

xor, xor\_byte, xor\_dword, xor\_qword, xor\_word

};

};

## Assembly Code

The AssemblyCode class handles one assembly code instruction. It holds methods for initialization and optimization of a single instruction, a set of methods for testing and conversion of values and register, and the ToString method that writes the instruction in plain text.

AssemblyCode.cs

using System;

using System.Text;

using System.Numerics;

using System.Collections.Generic;

First, there is a set of special registers:

* FrameRegister. Hold the pointer to the current activation record.
* EllipseRegister. Does also hold the pointer to the current activation record in an elliptic function. However, if the function has been called with elliptic arguments, it value has been increased by the size (in bytes) of elliptic arguments, in order to correctly access the non-parameter variables of the function.
* ReturnValueRegister. Holds the function return value in case of integral or register type.
* ReturnPointerRegister. Holds the address of the function return value in case of struct or union type.
* ShiftRegister. Holds the value of the right operand in a short operation, is always cl.

namespace CCompiler {

public class AssemblyCode {

public static Register FrameRegister;

public static Register EllipseRegister;

public static Register ReturnValueRegister;

public static Register ReturnPointerRegister;

public const Register ShiftRegister = Register.cl;

static AssemblyCode() {

FrameRegister = RegisterToSize(Register.bp, TypeSize.PointerSize);

EllipseRegister = RegisterToSize(Register.di, TypeSize.PointerSize);

ReturnValueRegister = RegisterToSize(Register.bx, TypeSize.PointerSize);

ReturnPointerRegister =RegisterToSize(Register.bx,TypeSize.PointerSize);

}

private AssemblyOperator m\_operator;

private object[] m\_operandArray = new object[3];

The constructor initializes the instruction, and calls FromAdditionToIncrement, which changes addition and subtraction to increment and decrement, and CheckSize, which changes the operators in accordance with the size of the values of the instruction.

public AssemblyCode(AssemblyOperator objectOp, object operand0,

object operand1, object operand2 = null,

int size = 0) {

m\_operator = objectOp;

m\_operandArray[0] = operand0;

m\_operandArray[1] = operand1;

m\_operandArray[2] = operand2;

FromAdditionToIncrement();

CheckSize(size);

}

public AssemblyOperator Operator {

get { return m\_operator; }

set { m\_operator = value; }

}

public object this[int index] {

get { return m\_operandArray[index]; }

set { m\_operandArray[index] = value; }

}

The FromAdditionToIncrement method changes the addition of one or subtraction of minus one to increment, and subtraction of one and addition of minus on to decrement. The first condition is that operator indeed is addition or subtraction and that the first operand is a track, register, or string.

private void FromAdditionToIncrement() {

if (((Operator == AssemblyOperator.add) ||

(Operator == AssemblyOperator.sub)) &&

((m\_operandArray[0] is Track) || (m\_operandArray[0] is Register) ||

(m\_operandArray[0] is string))) {

Then we check whether the second operand is a BigInteger and the third operand is null. In that case the value to be inspected has index one and we call CheckIncrement for further processing.

if ((m\_operandArray[1] is BigInteger) && (m\_operandArray[2] == null)){

CheckIncrement(1);

}

If the second operand is an integer and operand is a BigInteger, the value to be inspected has index two.

else if ((m\_operandArray[1] is int) &&

(m\_operandArray[2] is BigInteger)) {

CheckIncrement(2);

}

}

}

The CheckIncrement method inspects the value to be added or subtracted. If the operation is addition and the value is one or if the operation is subtraction and the value is minus one, we replace the operator with increment and set the value to null.

private void CheckIncrement(int valueIndex) {

int value = (int) ((BigInteger) m\_operandArray[valueIndex]);

if (((Operator == AssemblyOperator.add) && (value == 1)) ||

((Operator == AssemblyOperator.sub) && (value == -1))) {

m\_operator = AssemblyOperator.inc;

m\_operandArray[valueIndex] = null;

}

In the same way, if the operation is addition and the value is minus one or if the operation is subtraction and the value is one, we replace the operator with decrement and set the value to null.

else if (((Operator == AssemblyOperator.add) && (value == -1)) ||

((Operator == AssemblyOperator.sub) && (value == 1))) {

m\_operator = AssemblyOperator.dec;

m\_operandArray[valueIndex] = null;

}

}

The CheckSize method changes the operator to a size operator.

private void CheckSize(int size) {

if ((size != 0) && ((m\_operandArray[0] is Register) ||

(m\_operandArray[0] is Track)|| (m\_operandArray[0] is String)) &&

(m\_operandArray[1] is int) &&

(IsUnary() || ((m\_operandArray[2] is BigInteger) ||

(m\_operandArray[2] is String)))) {

m\_operator = OperatorToSize(m\_operator, size);

}

}

public bool IsUnary() {

string operatorName = Enum.GetName(typeof(AssemblyOperator), Operator);

return operatorName.StartsWith("neg") ||

operatorName.StartsWith("not") ||

operatorName.StartsWith("inc") ||

operatorName.StartsWith("dec") ||

operatorName.StartsWith("mul") ||

operatorName.StartsWith("imul") ||

operatorName.StartsWith("div") ||

operatorName.StartsWith("idiv") ||

operatorName.StartsWith("fst") ||

operatorName.StartsWith("fld") ||

operatorName.StartsWith("fist") ||

operatorName.StartsWith("fild");

}

public bool IsJumpRegister() {

return (Operator == AssemblyOperator.jmp) &&

(m\_operandArray[0] is Register);

}

public bool IsJumpNotRegister() {

return (Operator == AssemblyOperator.jmp) &&

!(m\_operandArray[0] is Register);

}

public bool IsCallRegister() {

return (Operator == AssemblyOperator.call) &&

(m\_operandArray[0] is Register);

}

public bool IsCallNotRegister() {

return (Operator == AssemblyOperator.call) &&

(m\_operandArray[0] is string);

}

public bool IsRelationNotRegister() {

switch (Operator) {

case AssemblyOperator.je:

case AssemblyOperator.jne:

case AssemblyOperator.jl:

case AssemblyOperator.jle:

case AssemblyOperator.jg:

case AssemblyOperator.jge:

case AssemblyOperator.jb:

case AssemblyOperator.jbe:

case AssemblyOperator.ja:

case AssemblyOperator.jae:

case AssemblyOperator.jc:

case AssemblyOperator.jnc:

return true;

default:

return false;

}

}

Two register overlaps if they are non-null and have the same name after we have removed the letter ‘l’, ‘h’, ‘x’, ‘e’, and ‘r’ characters. However, if one of the registers contains ‘l’ and the other register contains ‘h’, they do not overlap.

public static bool RegisterOverlap(Register? register1,

Register? register2) {

if ((register1 == null) || (register2 == null)) {

return false;

}

if (register1.Equals(register2)) {

return true;

}

string name1 = Enum.GetName(typeof(Register), register1),

name2 = Enum.GetName(typeof(Register), register2);

if ((name1.Contains("h") && name2.Contains("l")) ||

(name1.Contains("l") && name2.Contains("h"))) {

return false;

}

name1 = name1.Replace("h", "").Replace("l", "").Replace("x", "")

.Replace("e", "").Replace("r", "");

name2 = name2.Replace("h", "").Replace("l", "").Replace("x", "")

.Replace("e", "").Replace("r", "");

return name1.Equals(name2);

}

If the name of the register contains the letter ‘r’, its size is eight bytes. If it contains ‘e’, its size is four bytes. If it contains ‘l’ or ‘h’, its size is one byte. Otherwise, its size is two bytes.

The SizeOfRegister method return the size of a register.

public static int SizeOfRegister(Register register) {

string name = Enum.GetName(typeof(Register), register);

if (name.Contains("r")) {

return 8;

}

else if (name.Contains("e")) {

return 4;

}

else if (name.Contains("l") || name.Contains("h")) {

return 1;

}

else {

return 2;

}

}

The RegisterToSize method return the register with is stated size.

public static Register RegisterToSize(Register register, int size) {

string name = Enum.GetName(typeof(Register), register);

If the size is one byte, we replace ‘x’ with ‘l’ and remove ‘e’ and ‘r’.

switch (size) {

case 1:

name = name.Replace("x", "l").Replace("e", "").Replace("r", "");

break;

If the size is two bytes, we replace ‘l’ with ‘x’ and remove ‘e’ and ‘r’.

case 2:

Assert.ErrorXXX(!name.Contains("h"));

name = name.Replace("l", "x").Replace("e", "").Replace("r", "");

break;

If the size is four bytes, we replace ‘l’ with ‘x’ and remove ‘e’ and ‘r’, and add ‘e’ at the beginning.

case 4:

Assert.ErrorXXX(!name.Contains("h"));

name = "e" + name.Replace("l", "x").Replace("e", "")

.Replace("r", "");

break;

If the size is eight bytes, we replace ‘l’ with ‘x’ and remove ‘e’ and ‘r’, and add ‘r’ at the beginning.

case 8:

Assert.ErrorXXX(!name.Contains("h"));

name = "r" + name.Replace("l", "x").Replace("e", "")

.Replace("r", "");

break;

}

return ((Register) Enum.Parse(typeof(Register), name));

}

The SizeOfOperator method returns the size of the operator.

public static int SizeOfOperator(AssemblyOperator objectOp) {

string name = Enum.GetName(typeof(AssemblyOperator), objectOp);

if (name.Contains("\_byte")) {

return 1;

}

else if (name.Contains("\_word")) {

return 2;

}

else if (name.Contains("\_dword")) {

return 4;

}

else {

Assert.ErrorXXX(name.Contains("\_qword"));

return 8;

}

}

The OperatorToSize method changes a plain operator to a size operator.

public static AssemblyOperator OperatorToSize

(AssemblyOperator objectOp, int size) {

string name = Enum.GetName(typeof(AssemblyOperator), objectOp);

Assert.ErrorXXX(objectOp != AssemblyOperator.interrupt);

switch (size) {

case 1:

name = name + "\_byte";

break;

case 2:

name = name + "\_word";

break;

case 4:

name = name + "\_dword";

break;

case 8:

name = name + "\_qword";

break;

}

Assert.ErrorXXX(name.Contains("\_"));

return ((AssemblyOperator) Enum.Parse(typeof(AssemblyOperator), name));

}

The SizeOfValue method return a size of a value, with some exceptions. In case of a mov-operator, the size is given by the operator, and in case of a compare-operator with value zero, the size is one. Otherwise, we just call the second SizeOfValue to get the value size.

public static int SizeOfValue(BigInteger value, AssemblyOperator op) {

string name = Enum.GetName(typeof(AssemblyOperator), op);

if (name.StartsWith("mov")) {

return SizeOfOperator(op);

}

else if (name.StartsWith("cmp") && (value == 0)) {

return 1;

}

else {

return SizeOfValue(value);

}

}

public static int SizeOfValue(BigInteger value) {

if (value == 0) {

return 0;

}

else if ((-128 <= value) && (value <= 255)) {

return 1;

}

else if ((-32768 <= value) && (value <= 65535)) {

return 2;

}

else if ((-2147483648 <= value) && (value <= 4294967295)) {

return 4;

}

else {

return 8;

}

}

### ToString

The ToString method returns the instruction in plain text. We have six categories of operations:

public override string ToString() {

object operand0 = m\_operandArray[0],

operand1 = m\_operandArray[1],

operand2 = m\_operandArray[2];

string operatorName = Enum.GetName(typeof(AssemblyOperator),

Operator).Replace("\_", " ");

Instructions that perform computations are nullary (without operands), unary (with one operand), or binary (with two operands). In binary operations, the left operand must be a register or an address; a static name or an integer value is not allowed. The operations can be further divided into six categories:

|  |  |  |  |
| --- | --- | --- | --- |
| **Type** | **Category** | **Description** | **Examples** |
| Nullary | 1 | operator | fadd |
| Unary | 2 | operator register | neg ax |
| 3 | operator [register + offset]  operator [static name + offset] | inc [bp + 2]  dec [stdin + 4] |
| Binary | 5 | operator register, register  operator register, static name  operator register, integer value | mov ax, bx  add ax, stdin  sub ax, 123 |
| 6 | operator [register + offset], register  operator [register + offset], static name  operator [register + offset], integer value  operator [static name + offset], register  operator [static name + offset], static name  operator [static name + offset], integer value | mov [bp + 2], bx  add [bp + 2], stdin  sub [bp + 2], 123  mov [stdin + 4], bx  add [stdin + 4], stdin  sub [stdin + 4], 123 |
| 6 | operator register, [register + offset]  operator register, [static name + offset] | mov ax, [bp + 2]  add ax, [stdin +4] |

In case of a nullary operator, we simply return its name.

if (IsNullary()) {

return "\t" + operatorName;

}

Unary operations of category two and three, where the operand is an address or a register. In case of an address, the first operand is a register or a string (static name), the second operand is an integer (offset), and the third operand is null. In case of a register, the second and third operands are null.

else if (IsUnary()) {

if (((operand0 is Register) || (operand0 is string)) &&

(operand1 is int) && (operand2 == null)) {

return "\t" + operatorName + " [" + operand0 +

WithSign(operand1) + "]";

}

else if ((operand0 is Register) && (operand1 == null) &&

(operand2 == null)) {

return "\t" + operatorName + " " + operand0;

}

Next, we handle binary operations. In category four, the first operand is a register and the second operand is a register, a string (static name), or a BigInteger (integer value).

else if (IsBinary()) {

if ((operand0 is Register) && ((operand1 is Register) ||

(operand1 is string) || (operand1 is BigInteger)) &&

(operand2 == null)) {

Assert.ErrorXXX(!(operand0 is string));

return "\t" + operatorName + " " + operand0 + ", " + operand1;

}

In category five, the first operand is a register or a string (static name), the second operand is an integer (offset), and the third operand is register, a string (static name), or a BigInteger (integer value).

else if (((operand0 is Register) || (operand0 is string)) &&

(operand1 is int) && ((operand2 is Register) ||

(operand2 is string) || (operand2 is BigInteger))) {

return "\t" + operatorName +

" [" + operand0 + WithSign(operand1) + "], " + operand2;

}

In category six, the first operand is a register, the second operand is register or string (static name), and the third operand is an integer (offset).

else if ((operand0 is Register) && ((operand1 is Register) ||

(operand1 is string)) && (operand2 is int)) {

return "\t" + operatorName + " " + operand0 +

", [" + operand1 + WithSign(operand2) + "]";

}

}

In case of a label or comment, we simple return them. Note that we add a colon (‘:’) to the label and insert a semicolon (‘;’).

else if (Operator == AssemblyOperator.label) {

return "\n " + operand0 + ":";

}

else if (Operator == AssemblyOperator.comment) {

return "\t; " + operand0;

}

When defining a value, we call the ToVisibleString method in case of a string.

else if (Operator == AssemblyOperator.define\_value) {

Sort sort = (Sort) operand0;

object value = operand1;

if (sort == Sort.String) {

return "\tdb " + ToVisibleString((string) operand1);

}

If the value is not a string, it is an integral or floating value. In case of floating value, we add a dot (‘.’), unless it already has a dot.

else {

string text = operand1.ToString();

if (((sort == Sort.Float) || (sort == Sort.Double) ||

(sort == Sort.Long\_Double)) && !text.Contains(".")) {

text += ".0";

}

When we define the value, we use different directives depending on the size of the value: db (define byte) for one byte, dw (define word) for two bytes, dd (define double-word) for four bytes, and dq (define quarto) for eight bytes.

switch (TypeSize.Size(sort)) {

case 1:

return "\tdb " + text;

case 2:

return "\tdw " + text;

case 4:

return "\tdd " + text;

case 8:

return "\tdq " + text;

}

}

}

When defining an address, we use the dq directive, since the size of addresses is eight bytes. The WithSign method write a positive or negative offset, unless the offset is zero.

else if (Operator == AssemblyOperator.define\_address) {

string name = (string) operand0;

int offset = (int) operand1;

return "\tdq " + name + WithSign(offset);

}

When defining a sequence of zeros, we use the times and db directive, which repeats the following definition, which is the zero value of one-byte size.

else if (Operator == AssemblyOperator.define\_zero\_sequence) {

int size = (int) operand0;

return "\ttimes " + size + " db 0";

}

For a jump or function call, we use the jmp instruction.

else if (IsJumpRegister() || IsCallRegister() ||

IsCallNotRegister()) {

return "\tjmp " + operand0;

}

For address return, we load the target into the address.

else if (Operator == AssemblyOperator.return\_address) {

string target = SymbolTable.CurrentFunction.UniqueName +

Symbol.SeparatorId + operand2;

return "\tmov qword [" + operand0 + WithSign(operand1) + "], " +

target;

}

In case of conditional and unconditional jump instructions, we have three cases. If the third operand (operand2) is an integer, we jump to the address of a middle code instruction. More specific, we jump to a label, which is made up by the name of the current function, and the middle code index given by the third operand.

else if (IsRelationNotRegister() || IsJumpNotRegister()) {

if (operand2 is int) {

string label = SymbolTable.CurrentFunction.UniqueName +

Symbol.SeparatorId + operand2;

return "\t" + operatorName + " " + label;

}

If the third operand is string, we have a special case: the jump instruction is part of the handling of the command line arguments. We simply jump to the address given by the string.

else if (operand2 is string) {

return "\t" + operatorName + " " + operand2;

}

Otherwise, we have jump in a memory copy of a struct or union.

else {

int labelIndex = (int) operand1;

string labelText = MakeMemoryLabel(labelIndex);

return "\t" + operatorName + " " + labelText;

}

}

}

The ToVisibleString method returns the text with non-graphical character written with their ascii numbers. For instance, the text “Hello\nWorld\r”, will be translated to “Hello”, 10, “World”, 13.

private static string ToVisibleString(string text) {

StringBuilder buffer = new StringBuilder();

The insideString variable is true as long as the current character is locaed inside a string.

bool insideString = false;

We iterate throught the text and, for each character, check if it is not a graphical character. In that case we end the string, if necessary, and add the ascii value of the character.

foreach (char c in text) {

if (Char.IsControl(c) || (c == '\"') || (c == '\'')) {

if (insideString) {

buffer.Append("\", " + ((int) c).ToString() + ", ");

insideString = false;

}

else {

buffer.Append(((int) c).ToString() + ", ");

}

}

If the character is a graphical character, we begin the string, if necessary, and add the character to the string.

else {

if (insideString) {

buffer.Append(c);

}

else {

buffer.Append("\"" + c.ToString());

insideString = true;

}

}

}

We end the string by adding the terminating zero character.

if (insideString) {

buffer.Append("\", 0");

}

else {

buffer.Append("0");

}

return buffer.ToString();

}

The MakeMemory method return a label for the memory copy method.

public static string MakeMemoryLabel(int labelIndex) {

return "memorycopy" + labelIndex;

}

The WithSign method returns the offset of an address preceded by a plus or minus sign. In case of a zero value, an empty string is returned.

private string WithSign(object value) {

int offset = (int) value;

if (offset > 0) {

return " + " + offset;

}

else if (offset < 0) {

return " - " + (-offset);

}

else {

return "";

}

}

public List<byte> ByteList() {

object operand0 = m\_operandArray[0],

operand1 = m\_operandArray[1],

operand2 = m\_operandArray[2];

if ((Operator == AssemblyOperator.empty) ||

(Operator == AssemblyOperator.label) ||

(Operator == AssemblyOperator.comment)) {

return (new List<byte>());

}

else if (Operator == AssemblyOperator.define\_address) {

int offset = (int) operand1;

List<byte> byteList = new List<byte>(new byte[TypeSize.PointerSize]);

AssemblyCode.LoadByteList(byteList, 0, TypeSize.PointerSize,

(BigInteger) offset);

return byteList;

}

else if (Operator == AssemblyOperator.define\_zero\_sequence) {

int size = (int) operand0;

return (new List<byte>(new byte[size]));

}

else if (Operator == AssemblyOperator.define\_value) {

Sort sort = (Sort) operand0;

object value = operand1;

if (sort == Sort.Pointer) {

List<byte> byteList = new List<byte>(new byte[TypeSize.PointerSize]);

if (value is string) {

AssemblyCode.LoadByteList(byteList, 0, TypeSize.PointerSize,

BigInteger.Zero);

}

else if (value is StaticAddress) {

StaticAddress staticAddress = (StaticAddress) value;

int offset = staticAddress.Offset;

AssemblyCode.LoadByteList(byteList, 0, TypeSize.PointerSize,

(BigInteger) offset);

}

else {

AssemblyCode.LoadByteList(byteList, 0, TypeSize.PointerSize,

(BigInteger) value);

}

return byteList;

}

else if (sort == Sort.Float) {

float floatValue = (float) ((decimal) operand0);

return (new List<byte>(BitConverter.GetBytes(floatValue)));

}

else if ((sort == Sort.Double) || (sort == Sort.Long\_Double)) {

double doubleValue = (double) ((decimal) value);

return (new List<byte>(BitConverter.GetBytes(doubleValue)));

}

else if (sort == Sort.String) {

string text = (string) value;

List<byte> byteList = new List<byte>();

foreach (char c in text) {

byteList.Add((byte) c);

}

byteList.Add((byte) 0);

return byteList;

}

else {

int size = TypeSize.Size(sort);

List<byte> byteList = new List<byte>(new byte[size]);

AssemblyCode.LoadByteList(byteList, 0, size, (BigInteger) value);

return byteList;

}

}

else if (IsJumpRegister() || IsCallRegister()) {

Register register = (Register) operand0;

return LookupByteArray(AssemblyOperator.jmp, register);

}

else if (IsCallNotRegister()) {

List<byte> byteList =

LookupByteArray(AssemblyOperator.jmp, TypeSize.PointerSize);

LoadByteList(byteList, byteList.Count - TypeSize.PointerSize,

TypeSize.PointerSize, 0);

return byteList;

}

else if (Operator == AssemblyOperator.return\_address) {

Register register = (Register) operand0;

int offset = (int) operand1;

int size = SizeOfValue(offset);

int address = (int)((BigInteger)operand2);

List<byte> byteList =

LookupByteArray(AssemblyOperator.mov\_word, register,

size, TypeSize.PointerSize);

LoadByteList(byteList, byteList.Count - (size + TypeSize.PointerSize),

size, offset);

LoadByteList(byteList, byteList.Count - TypeSize.PointerSize,

TypeSize.PointerSize, address);

return byteList;

}

else if (IsRelationNotRegister() || IsJumpNotRegister()) {

int address = (int) operand0;

int size = ((address >= -128) && (address <= 127)) ? 1 : 2;

List<byte> byteList = LookupByteArray(Operator, size);

LoadByteList(byteList, byteList.Count - size, size, address);

return byteList;

}

|  |  |  |
| --- | --- | --- |
| **Category** | **Description** | **Examples** |
| 1 | binary\_operator register, register | mov ax, bx |
| 2 | binary\_operator register, static\_name | add ax, stdin |
| 3 | binary\_operator register, integer\_value | sub ax, 123 |
| 2 | binary\_operator [register + offset], register  binary\_operator [register + offset], static\_name  binary\_operator [register + offset], integer\_value  binary\_operator [static\_name + offset], register  binary\_operator [static\_name + offset], static\_name  binary\_operator [static\_name + offset], integer\_value | mov [bp + 2], bx  add [bp + 2], stdin  sub [bp + 2], 123  mov [stdin + 4], bx  add [stdin + 4], stdin  sub [stdin + 4], 123 |
| 3 | binary\_operator register, [register + offset]  binary\_operator register, [static\_name + offset]  binary\_operator static\_name, [register + offset]  binary\_operator static\_name, [static\_name + offset] | mov ax, [bp + 2]  add ax, [stdin +4]  cmp stdin, [bp + 2]  cmp stdin, [stdin + 4] |
| 4 | unary\_operator register  unary\_operator integer\_value | inc ax  int 33 |
| 5 | unary\_operator [register + offset]  unary\_operator [static\_name + offset]  unary\_operator [offset] |  |
| 6 | operator\_without\_operands |  |

// mov ax, [bp + 2]

else if ((operand0 is Register) && (operand1 is Register) &&

(operand2 is int)) {

Register toRegister = (Register) operand0,

baseRegister = (Register) operand1;

int offset = (int) operand2;

int size = SizeOfValue(offset);

List<byte> byteList =

LookupByteArray(Operator, toRegister, baseRegister, size);

LoadByteList(byteList, byteList.Count - size, size, offset);

return byteList;

}

// mov ax, [global + 4]; mov ax, [null + 4]

else if (((operand0 is Register) && (operand1 is string) &&

(operand2 is int)) || ((operand0 is Register) &&

(operand1 == null) && (operand2 is int))) {

Register toRegister = (Register) operand0;

int offset = (int) operand2;

List<byte> byteList =

LookupByteArray(Operator, toRegister, null, TypeSize.PointerSize);

LoadByteList(byteList, byteList.Count - TypeSize.PointerSize,

TypeSize.PointerSize, offset);

return byteList;

}

// mov [bp + 2], ax

else if ((operand0 is Register) && (operand1 is int) &&

(operand2 is Register)) {

Register baseRegister = (Register) operand0,

fromRegister = (Register) operand2;

int offset = (int) operand1;

int size = SizeOfValue(offset);

List<byte> byteList =

LookupByteArray(Operator, baseRegister, size, fromRegister);

LoadByteList(byteList, byteList.Count - size, size, offset);

return byteList;

}

// mov [global + 4], ax; mov [null + 4], ax

else if ((operand1 is int) && (operand2 is Register)) {

Assert.ErrorXXX((operand0 is string) || (operand0 == null));

int offset = (int)operand1;

Register fromRegister = (Register)operand2;

List<byte> byteList =

LookupByteArray(Operator, null, TypeSize.PointerSize, fromRegister);

LoadByteList(byteList, byteList.Count - TypeSize.PointerSize,

TypeSize.PointerSize, offset);

return byteList;

}

// mov [bp + 2], 123

else if ((operand0 is Register) && (operand1 is int) &&

(operand2 is BigInteger)) {

Register baseRegister = (Register) operand0;

int offset = (int) operand1;

BigInteger value = (BigInteger) operand2;

int offsetSize = SizeOfValue(offset),

valueSize = SizeOfValue(value, Operator);

List<byte> byteList =

LookupByteArray(Operator, baseRegister, offsetSize, valueSize);

LoadByteList(byteList, byteList.Count - (offsetSize + valueSize),

offsetSize, offset);

LoadByteList(byteList, byteList.Count - valueSize,

valueSize, value);

return byteList;

}

// mov [bp + 2], global

else if ((operand0 is Register) && (operand1 is int) &&

(operand2 is string)) {

Register baseRegister = (Register) operand0;

int offset = (int) operand1;

int offsetSize = SizeOfValue(offset);

List<byte> byteList =

LookupByteArray(Operator, baseRegister, offsetSize,

TypeSize.PointerSize);

LoadByteList(byteList, byteList.Count -

(offsetSize + TypeSize.PointerSize), offsetSize, offset);

LoadByteList(byteList, byteList.Count - TypeSize.PointerSize,

TypeSize.PointerSize, 0);

return byteList;

}

// mov [global + 4], 123; mov [null + 4], 123

else if ((operand1 is int) && (operand2 is BigInteger)) {

Assert.ErrorXXX((operand0 is string) || (operand0 == null));

int offset = (int) operand1;

BigInteger value = (BigInteger) operand2;

int valueSize = SizeOfValue(value, Operator);

List<byte> byteList =

LookupByteArray(Operator, null, TypeSize.PointerSize, valueSize);

LoadByteList(byteList, byteList.Count - (TypeSize.PointerSize +

valueSize), TypeSize.PointerSize, offset);

LoadByteList(byteList, byteList.Count - valueSize,

valueSize, value);

return byteList;

}

// mov ax, bx

else if ((operand0 is Register) && (operand1 is Register)) {

Assert.ErrorXXX(operand2 == null);

Register toRegister = (Register) operand0,

fromRegister = (Register) operand1;

return LookupByteArray(Operator, toRegister, fromRegister);

}

// mov ax, 123

else if ((operand0 is Register) && (operand1 is BigInteger)) {

Assert.ErrorXXX(operand2 == null);

Register register = (Register) operand0;

BigInteger value = (BigInteger) operand1;

int size = ((Operator == AssemblyOperator.mov) ||

(Operator == AssemblyOperator.and))

? SizeOfRegister(register) : SizeOfValue(value);

List<byte> byteList = LookupByteArray(Operator, register, size);

LoadByteList(byteList, byteList.Count - size, size, value);

return byteList;

}

// mov ax, global

else if ((operand0 is Register) && (operand1 is string)) {

Assert.ErrorXXX(operand2 == null);

Register register = (Register) operand0;

int size = SizeOfRegister(register);

List<byte> byteList = LookupByteArray(Operator, register, size);

LoadByteList(byteList, byteList.Count - size, size, 0);

return byteList;

}

// inc [bp + 2]

else if ((operand0 is Register) && (operand1 is int)) {

Assert.ErrorXXX(operand2 == null);

Register baseRegister = (Register) operand0;

int offset = (int) operand1;

int size = SizeOfValue(offset);

List<byte> byteList = LookupByteArray(Operator, baseRegister, size);

LoadByteList(byteList, byteList.Count - size, size, offset);

return byteList;

}

// inc [global + 4]; inc [null + 4]

else if (operand1 is int) {

Assert.ErrorXXX((operand0 is string) || (operand0 == null));

Assert.ErrorXXX(operand2 == null);

int offset = (int) operand1;

List<byte> byteList =

LookupByteArray(Operator, null, TypeSize.PointerSize);

LoadByteList(byteList, byteList.Count - TypeSize.PointerSize,

TypeSize.PointerSize, offset);

return byteList;

}

// inc ax

else if (operand0 is Register) {

Assert.ErrorXXX((operand1 == null) && (operand2 == null));

Register register = (Register) operand0;

return LookupByteArray(Operator, register);

}

// int 33

else if (operand0 is BigInteger) {

Assert.ErrorXXX((operand1 == null) && (operand2 == null));

BigInteger value = (BigInteger) operand0;

int size = SizeOfValue(value);

List<byte> byteList = LookupByteArray(Operator, size);

LoadByteList(byteList, byteList.Count - size, size, value);

return byteList;

}

// lahf

else {

Assert.ErrorXXX((operand0 == null) && (operand1 == null) &&

(operand2 == null));

return LookupByteArray(Operator);

}

}

public static void LoadByteList(IList<byte> byteList, int index,

int size, BigInteger value) {

while (byteList.Count < (index + size)) {

byteList.Add((byte) 0);

}

switch (size) {

case 1: {

if (value < 0) {

byteList[index] = (byte) ((sbyte) value);

}

else {

byteList[index] = (byte) value;

}

}

break;

case 2: {

if (value < 0) {

short shortValue = (short) value;

byteList[index] = (byte) ((sbyte) shortValue);

byteList[index + 1] = (byte) ((sbyte) (shortValue >> 8));

}

else {

ushort ushortValue = (ushort) value;

byteList[index] = (byte) ushortValue;

byteList[index + 1] = (byte) (ushortValue >> 8);

}

}

break;

case 4: {

if (value < 0) {

int intValue = (int) value;

byteList[index] = (byte) ((sbyte) intValue);

byteList[index + 1] = (byte) ((sbyte) (intValue >> 8));

byteList[index + 2] = (byte) ((sbyte) (intValue >> 16));

byteList[index + 3] = (byte) ((sbyte) (intValue >> 24));

}

else {

uint uintValue = (uint) value;

byteList[index] = (byte) uintValue;

byteList[index + 1] = (byte) (uintValue >> 8);

byteList[index + 2] = (byte) (uintValue >> 16);

byteList[index + 3] = (byte) (uintValue >> 24);

}

}

break;

}

}

public static List<byte> LookupByteArray(AssemblyOperator objectOp,

object operand1 = null, object operand2 = null,

object operand3 = null) {

if ((objectOp == AssemblyOperator.shl) ||

(objectOp == AssemblyOperator.shr)) {

operand1 = (operand1 is BigInteger) ? 0L : operand1;

operand2 = (operand2 is BigInteger) ? 0L : operand2;

operand3 = (operand3 is BigInteger) ? 0L : operand3;

}

ObjectCodeInfo info =

new ObjectCodeInfo(objectOp, operand1, operand2, operand3);

byte[] byteArray = ObjectCodeTable.MainArrayMap[info];

Assert.ErrorXXX(byteArray != null);

List<byte> byteList = new List<byte>();

foreach (byte b in byteArray) {

byteList.Add(b);

}

return byteList;

}

}

}

## Tracks

A track is a place holder for a register through the assembly code.

Track.cs

using System;

using System.Collections.Generic;

namespace CCompiler {

public class Track {

private Register? m\_register = null;

private List<TrackEntry> m\_entryList = new List<TrackEntry>();

private bool m\_pointer;

private int m\_currentSize, m\_maxSize;

public Track(Symbol symbol, Register? register = null) {

m\_register = register;

Assert.ErrorXXX(symbol != null);

Assert.ErrorXXX(!symbol.Type.IsStructOrUnion());

m\_currentSize = m\_maxSize = symbol.Type.SizeArray();

}

public Track(Type type) {

Assert.ErrorXXX(type != null);

Assert.ErrorXXX(!type.IsStructOrUnion());

Assert.ErrorXXX(!type.IsArrayFunctionOrString());

m\_currentSize = m\_maxSize = type.Size();

}

public void Replace(List<AssemblyCode> assemblyCodeList, Track newTrack) {

foreach (TrackEntry entry in m\_entryList) {

assemblyCodeList[entry.Line][entry.Position] = newTrack;

}

}

public int CurrentSize {

get { return m\_currentSize; }

set {

m\_currentSize = value;

m\_maxSize = Math.Max(m\_maxSize, m\_currentSize);

}

}

public void AddEntry(int position, int line) {

m\_entryList.Add(new TrackEntry(position, line, m\_currentSize));

}

public int MaxSize {

get {return m\_maxSize; }

}

public Register? Register {

get { return m\_register; }

set {

Assert.ErrorXXX((value == null) || (m\_register == null) ||

AssemblyCode.RegisterOverlap(value, m\_register));

m\_register = value;

}

}

public bool Pointer {

get { return m\_pointer; }

set { m\_pointer = value; }

}

public static bool Overlaps(Track track1, Track track2) {

if ((track1.m\_entryList.Count == 0) || (track2.m\_entryList.Count == 0)){

return false;

}

TrackEntry minEntry1 = track1.m\_entryList[0],

minEntry2 = track2.m\_entryList[0],

maxEntry1 = track1.m\_entryList[track1.m\_entryList.Count - 1],

maxEntry2 = track2.m\_entryList[track2.m\_entryList.Count - 1];

return !(((maxEntry1.Line < minEntry2.Line) ||

(maxEntry2.Line < minEntry1.Line)));

}

public void Generate(List<AssemblyCode> objectCodeList) {

foreach (TrackEntry entry in m\_entryList) {

Register sizeRegister =

AssemblyCode.RegisterToSize(m\_register.Value, entry.Size);

AssemblyCode objectCode = objectCodeList[entry.Line];

objectCode[entry.Position] = sizeRegister;

}

}

}

}

A track entry is one entry in a track, it represents one occurrence of a register in the assembly code. In hold the line number of the assembly code instruction, the position (zero, one, or two, representing the left, middle, and right operand), and the size of the register.

TrackEntry.cs

namespace CCompiler {

public class TrackEntry {

private int m\_line, m\_position, m\_size;

public TrackEntry(int position, int line, int size) {

m\_line = line;

m\_position = position;

m\_size = size;

}

public int Line {

get { return m\_line; }

}

public int Position {

get { return m\_position; }

}

public int Size {

get { return m\_size; }

}

}

}Register Allocation

The RegisterAllocator method take the total track set and total assembly list; that is, the track set for the whole function, and the assembly list for the whole function. The first task is the find out which tracks overlaps each other. We construct the graph in this way: the tracks are the vertices of the graph and two vertices have an edge if the two tracks overlaps. The register allocation is then performed as a graph coloring, two tracks that overlaps cannot be assigned the same register.

RegisterAllocator.cs

using System.Collections.Generic;

namespace CCompiler {

public class RegisterAllocator {

public RegisterAllocator(ISet<Track> totalTrackSet,

List<AssemblyCode> assemblyCodeList) {

Graph<Track> totalTrackGraph = new Graph<Track>(totalTrackSet);

foreach (Track track1 in totalTrackSet) {

foreach (Track track2 in totalTrackSet) {

if (!track1.Equals(track2) && Track.Overlaps(track1, track2)) {

totalTrackGraph.AddEdge(track1, track2);

}

}

}

ISet<Graph<Track>> split = totalTrackGraph.Split();

int index = 0;

foreach (Graph<Track> trackGraph in split) {

List<Track> trackList = new List<Track>(trackGraph.VertexSet);

Assert.Error(DeepFirstSearch(trackList, 0, trackGraph),

Message.Out\_of\_registers);

++index;

}

foreach (Track track in totalTrackSet) {

track.Generate(assemblyCodeList);

}

}

The DeepFirstSearch method searches the graph in a deep-first manner. It takes the track list and the current index in that list as well as the track graph.

private bool DeepFirstSearch(List<Track> trackList, int listIndex,

Graph<Track> trackGraph) {

If the index equals the size of the track list, we return true because we have iterated through the list and found a match; that is, each track has been assigned a register and no overlapping tracks have the same register.

if (listIndex == trackList.Count) {

return true;

}

If the current track has already been assigned a register, we just call DeepFirstSearch with the next index.

Track track = trackList[listIndex];

if (track.Register != null) {

return DeepFirstSearch(trackList, listIndex + 1, trackGraph);

}

If the current track has not been assigned a register, we look up the set of possible register and the set of neighbor vertices; that is, the set of overlapping tracks.

ISet<Register> possibleSet = GetPossibleSet(track);

ISet<Track> neighbourSet = trackGraph.GetNeighbourSet(track);

We iterate through the set of possible register and, for each register that does not cause an overlapping, we assign the track the register and call DeepFirstSearch recursively with the next index. If the call returns true, we have found a total mapping of registers to the track, and we just return true. In this way, every call to DeepFirstSearch, including the first call in RegisterAllocator. However, if the call does not return false, we just try with another of the possible register. If the none of the registers causes a match, we clear the register of the track and return false.

foreach (Register possibleRegister in possibleSet) {

if (!OverlapNeighbourSet(possibleRegister, neighbourSet)) {

track.Register = possibleRegister;

if (DeepFirstSearch(trackList, listIndex + 1, trackGraph)) {

return true;

}

track.Register = null;

}

}

track.Register = null;

return false;

}

The OverlapNeighbourSet method return true if the register overlaps any of its neighbors. The RegisterOverlap method in the AssemblyCode class test whether two register overlaps.

private bool OverlapNeighbourSet(Register register,

ISet<Track> neighbourSet) {

foreach (Track neighbourTrack in neighbourSet) {

if (AssemblyCode.RegisterOverlap(register, neighbourTrack.Register)) {

return true;

}

}

return false;

}

The PointerRegisterSetWithEllipse set holds the possible pointer registers of an elliptic function while PointerRegisterSetWithoutEllipse holds the possible pointer registers of an non-elliptic function. The Byte1RegisterSet ste holds all registers of one byte while Byte2RegisterSet holds all registers of two bytes.

public static ISet<Register>

PointerRegisterSetWithEllipse = new HashSet<Register>(),

PointerRegisterSetWithoutEllipse = new HashSet<Register>(),

Byte1RegisterSet = new HashSet<Register>(),

Byte2RegisterSet = new HashSet<Register>();

static RegisterAllocator() {

PointerRegisterSetWithEllipse.

Add(AssemblyCode.RegisterToSize(Register.si, TypeSize.PointerSize));

PointerRegisterSetWithEllipse.

Add(AssemblyCode.RegisterToSize(Register.di, TypeSize.PointerSize));

PointerRegisterSetWithEllipse.

Add(AssemblyCode.RegisterToSize(Register.bx, TypeSize.PointerSize));

The PointerRegisterSetWithoutEllipse holds the registers of PointerRegisterSetWithEllipse minus the EllipseRegister register since we need it to keep track of the ellipse frame pointer in elliptic functions.

PointerRegisterSetWithoutEllipse.

UnionWith(PointerRegisterSetWithEllipse);

PointerRegisterSetWithoutEllipse.Remove(AssemblyCode.EllipseRegister);

Byte1RegisterSet.Add(Register.al);

Byte1RegisterSet.Add(Register.ah);

Byte1RegisterSet.Add(Register.bl);

Byte1RegisterSet.Add(Register.bh);

Byte1RegisterSet.Add(Register.cl);

Byte1RegisterSet.Add(Register.ch);

Byte1RegisterSet.Add(Register.dl);

Byte1RegisterSet.Add(Register.dh);

Byte2RegisterSet.Add(Register.ax);

Byte2RegisterSet.Add(Register.bx);

Byte2RegisterSet.Add(Register.cx);

Byte2RegisterSet.Add(Register.dx);

}

The GetPossibleSet method returns the possible set a track, depending on whether the track holds a pointer, or the size of the track.

private static ISet<Register> GetPossibleSet(Track track) {

if (track.Pointer) {

if (SymbolTable.CurrentFunction.Type.IsEllipse()) {

return PointerRegisterSetWithoutEllipse;

}

else {

return PointerRegisterSetWithEllipse;

}

}

If the track does not hold a pointer wee look into its size. If the size is one, we have a larger set to choose from. There are eight non-pointer registers of size one while there is four registers of the other sizes.

else if (track.MaxSize == 1) {

return Byte1RegisterSet;

}

We return the set of registers of size two, even if the size is actually larger than two. In that case, the RegisterToSize method in the AssemblyCode class will find the register of correct size.

else {

return Byte2RegisterSet;

}

}

}

}

## Assembly Code Generation

AssemblyCodeGenerator.cs

using System.Linq;

using System.Numerics;

using System.Collections.Generic;

namespace CCompiler {

public class AssemblyCodeGenerator {

public IDictionary<Symbol,Track> m\_trackMap =

new ListMap<Symbol,Track>();

public List<AssemblyCode> m\_assemblyCodeList;

private int m\_floatStackSize = 0;

public const int FloatingStackMaxSize = 7;

public static string MainName = "main";

public static string InitializerName = Symbol.SeparatorId + "initializer";

public static string ArgsName = Symbol.SeparatorId + "args";

public static string PathName = Symbol.SeparatorId + "PathName";

public static string PathText = "";

public AssemblyCodeGenerator(List<AssemblyCode> assemblyCodeList) {

m\_assemblyCodeList = assemblyCodeList;

}

private void RegisterAllocation(ISet<Track> trackSet) {

new RegisterAllocator(trackSet, m\_assemblyCodeList);

}

public static void GenerateAssembly(List<MiddleCode> middleCodeList,

List<AssemblyCode> assemblyCodeList) {

AssemblyCodeGenerator objectCodeGenerator =

new AssemblyCodeGenerator(assemblyCodeList);

objectCodeGenerator.AssemblyCodeList(middleCodeList);

ISet<Track> trackSet = objectCodeGenerator.TrackSet();

objectCodeGenerator.RegisterAllocation(trackSet);

}

public static void GenerateTargetWindows

(List<AssemblyCode> assemblyCodeList, List<byte> byteList,

IDictionary<int,string> accessMap, IDictionary<int,string> callMap,

ISet<int> returnSet) {

AssemblyCodeGenerator objectCodeGenerator =

new AssemblyCodeGenerator(assemblyCodeList);

objectCodeGenerator.WindowsJumpInfo();

objectCodeGenerator.WindowsByteList

(byteList, accessMap, callMap, returnSet);

}

The first AddAssemblyCode method adds a new assembly instruction to the assembly code list m\_assemblyCodeList.

public AssemblyCode AddAssemblyCode(AssemblyOperator objectOp,

object operand0 = null, object operand1 = null,

object operand2 = null, int size = 0) {

AssemblyCode assemblyCode =

new AssemblyCode(objectOp, operand0, operand1, operand2, size);

m\_assemblyCodeList.Add(assemblyCode);

return assemblyCode;

}

The second AddAssemblyCode method adds a new assembly instruction to the given assembly code list. It is static and is called when generating special assembly code, such as initialization code and code for command line arguments.

public static AssemblyCode AddAssemblyCode(List<AssemblyCode> list,

AssemblyOperator objectOp, object operand0 = null,

object operand1 = null, object operand2 = null,

int size = 0) {

AssemblyCode assemblyCode =

new AssemblyCode(objectOp, operand0, operand1, operand2, size);

list.Add(assemblyCode);

return assemblyCode;

}

### The Long Switch Statement

The constructor iterates through the middle code list and calls an appropriate method for each kind of middle code instruction.

public void AssemblyCodeList(List<MiddleCode> middleCodeList){

for (int middleIndex = 0; middleIndex < middleCodeList.Count;

++middleIndex) {

MiddleCode middleCode = middleCodeList[middleIndex];

AddAssemblyCode(AssemblyOperator.new\_middle\_code, middleIndex);

If the current function is not null (if the code is in function scope rather global scope), we add a label. If the middle code instruction is an initializer, we just add a label with the text of the middle code instruction.

if (SymbolTable.CurrentFunction != null) {

if ((middleCode.Operator == MiddleOperator.Initializer) ||

(middleCode.Operator == MiddleOperator.InitializerZero)) {

AddAssemblyCode(AssemblyOperator.label, null,

middleCode.ToString());

}

If the middle code instruction is not an initializer, we add a label with the unique name of the function. If the middle code index is greater the zero, we also add it to the label.

else {

string label = SymbolTable.CurrentFunction.UniqueName;

if (middleIndex > 0) {

label += Symbol.SeparatorId + middleIndex;

}

AddAssemblyCode(AssemblyOperator.label, label,

middleCode.ToString());

}

}

Now follows the long switch statement, where each middle code instruction is matched to a function call that generates the corresponding assembly code.

switch (middleCode.Operator) {

case MiddleOperator.CallHeader:

FunctionPreCall(middleCode);

break;

case MiddleOperator.Call:

FunctionCall(middleCode, middleIndex);

break;

case MiddleOperator.PostCall:

FunctionPostCall(middleCode);

break;

case MiddleOperator.Return:

Return(middleCode);

break;

case MiddleOperator.Exit:

Exit(middleCode);

break;

case MiddleOperator.Goto:

Goto(middleCode);

break;

case MiddleOperator.AssignRegister:

LoadToRegister(middleCode);

break;

case MiddleOperator.InspectRegister:

InspectRegister(middleCode);

break;

case MiddleOperator.JumpRegister:

JumpToRegister(middleCode);

break;

case MiddleOperator.Interrupt:

Interrupt(middleCode);

break;

case MiddleOperator.SysCall:

SystemCall(middleCode);

break;

case MiddleOperator.Initializer:

Initializer(middleCode);

break;

case MiddleOperator.InitializerZero:

InitializerZero(middleCode);

break;

In case of assignment, we inspect the type of the symbol to be assigned. If it is a struct or union, we call StructUnionAssign. Otherwise, the type must be integral, and we call IntegralAssign. Note that the assignment middle code instruction is not used for assignment of floating values. In those case, we instead top or pop the floating value stack.

case MiddleOperator.Assign: {

Symbol symbol = (Symbol)middleCode[0];

if (symbol.Type.IsStructOrUnion()) {

StructUnionAssign(middleCode, middleIndex);

}

else {

IntegralAssign(middleCode);

}

}

break;

case MiddleOperator.BitwiseAnd:

case MiddleOperator.BitwiseOr:

case MiddleOperator.BitwiseXOr:

case MiddleOperator.ShiftLeft:

case MiddleOperator.ShiftRight:

IntegralBinary(middleCode);

break;

In case of addition or subtract, we check the type of the left operand. If it is a floating value, we call FloatingBinary. Otherwise, the type is integral, and we call IntegralBinary.

case MiddleOperator.BinaryAdd:

case MiddleOperator.BinarySubtract: {

Symbol resultSymbol = (Symbol) middleCode[1];

if (resultSymbol.Type.IsFloating()) {

FloatingBinary(middleCode);

}

else {

IntegralBinary(middleCode);

}

}

break;

We have a similar case when performing multiplication or division. If case of floating values, we call FloatingBinary (the same method as in the floating addition and subtracting case). Otherwise, the type is integral, and we call IntegralMultiply.

case MiddleOperator.SignedMultiply:

case MiddleOperator.SignedDivide:

case MiddleOperator.SignedModulo:

case MiddleOperator.UnsignedMultiply:

case MiddleOperator.UnsignedDivide:

case MiddleOperator.UnsignedModulo: {

Symbol resultSymbol = (Symbol) middleCode[0];

if (resultSymbol.Type.IsFloating()) {

FloatingBinary(middleCode);

}

else {

IntegralMultiply(middleCode);

}

}

break;

case MiddleOperator.Carry:

case MiddleOperator.NotCarry:

CarryExpression(middleCode);

break;

In case of equality and relational operators, we call FloatingRelation in case of floating values and IntegralBinary otherwise.

case MiddleOperator.Equal:

case MiddleOperator.NotEqual:

case MiddleOperator.SignedLessThan:

case MiddleOperator.SignedLessThanEqual:

case MiddleOperator.SignedGreaterThan:

case MiddleOperator.SignedGreaterThanEqual:

case MiddleOperator.UnsignedLessThan:

case MiddleOperator.UnsignedLessThanEqual:

case MiddleOperator.UnsignedGreaterThan:

case MiddleOperator.UnsignedGreaterThanEqual: {

Symbol leftSymbol = (Symbol) middleCode[1];

if (leftSymbol.Type.IsFloating()) {

FloatingRelation(middleCode);

}

else {

IntegralBinary(middleCode);

}

}

break;

case MiddleOperator.Case:

Case(middleCode);

break;

case MiddleOperator.CaseEnd:

CaseEnd(middleCode);

break;

The same goes for the unary operations. In case of floating type we call FloatingUnary, and in case of integral type we call InategralUnary.

case MiddleOperator.UnaryAdd:

case MiddleOperator.UnarySubtract:

case MiddleOperator.BitwiseNot: {

Symbol resultSymbol = (Symbol) middleCode[0];

if (resultSymbol.Type.IsFloating()) {

FloatingUnary(middleCode);

}

else {

IntegralUnary(middleCode);

}

}

break;

case MiddleOperator.Address:

Address(middleCode);

break;

case MiddleOperator.Dereference: {

Symbol symbol = (Symbol) middleCode[1];

if (symbol.Type.IsFloating()) {

FloatingDereference(middleCode);

}

else {

IntegralDereference(middleCode);

}

}

break;

case MiddleOperator.DecreaseStack:

Assert.ErrorXXX((--m\_floatStackSize) >= 0);

break;

case MiddleOperator.CheckTrackMapFloatStack:

Assert.ErrorXXX((m\_trackMap.Count == 0) &&

(m\_floatStackSize == 0));

break;

case MiddleOperator.PushZero:

PushSymbol(new Symbol(Type.DoubleType, (decimal) 0));

break;

case MiddleOperator.PushOne:

PushSymbol(new Symbol(Type.DoubleType, (decimal) 1));

break;

case MiddleOperator.PushFloat:

PushSymbol((Symbol) middleCode[0]);

break;

case MiddleOperator.TopFloat:

TopPopSymbol((Symbol) middleCode[0], TopOrPop.Top);

break;

case MiddleOperator.PopFloat:

TopPopSymbol((Symbol) middleCode[0], TopOrPop.Pop);

break;

case MiddleOperator.PopEmpty:

PopEmpty();

break;

case MiddleOperator.IntegralToIntegral:

IntegralToIntegral(middleCode, middleIndex);

break;

case MiddleOperator.IntegralToFloating:

IntegralToFloating(middleCode);

break;

case MiddleOperator.FloatingToIntegral:

FloatingToIntegral(middleCode);

break;

case MiddleOperator.Parameter: {

Symbol paramSymbol = (Symbol) middleCode[1];

if (paramSymbol.Type.IsFloating()) {

FloatingParameter(middleCode);

}

else if (paramSymbol.Type.IsStructOrUnion()) {

StructUnionParameter(middleCode, middleIndex);

}

else {

IntegralParameter(middleCode);

}

}

break;

case MiddleOperator.GetReturnValue: {

Symbol returnSymbol = (Symbol) middleCode[0];

if (returnSymbol.Type.IsStructOrUnion()) {

StructUnionGetReturnValue(middleCode);

}

else if (returnSymbol.Type.IsFloating()) {

Assert.Error((++m\_floatStackSize) <= FloatingStackMaxSize,

null, Message.Floating\_stack\_overflow);

}

else {

IntegralGetReturnValue(middleCode);

}

}

break;

case MiddleOperator.SetReturnValue: {

Symbol returnSymbol = (Symbol) middleCode[1];

if (returnSymbol.Type.IsStructOrUnion()) {

StructUnionSetReturnValue(middleCode);

}

else if (returnSymbol.Type.IsFloating()) {

Assert.ErrorXXX((--m\_floatStackSize) == 0);

}

else {

IntegralSetReturnValue(middleCode);

}

}

break;

}

}

}

### Track Set Generation

When the assembly code has been generated, a lot of tracks has been added to the assembly code. The task of the TrackSet method is to pick up the tracks from the and generate the track set; that is, a set holding the tracks of the code. Each track holds a list of entries, specifying the positions of the track in the code.

private ISet<Track> TrackSet() {

ISet<Track> trackSet = new HashSet<Track>();

for (int index = 0; index < m\_assemblyCodeList.Count; ++index) {

AssemblyCode assemblyCode = m\_assemblyCodeList[index];

object operand0 = assemblyCode[0],

operand1 = assemblyCode[1],

operand2 = assemblyCode[2];

In case of the set-track-size instruction, we set the current size of its track and replace the instruction with an empty instruction. The current size is then used when adding entries to the track so that the registers in the end are given the correct size.

if (assemblyCode.Operator == AssemblyOperator.set\_track\_size) {

Track track = (Track) operand0;

if (operand1 is int) {

track.CurrentSize = (int) operand1;

}

else {

track.CurrentSize = ((Track) operand1).CurrentSize;

}

assemblyCode.Operator = AssemblyOperator.empty;

}

In all other cases, we exam the operators of the instruction by calling CheckTrack.

else {

CheckTrack(trackSet, operand0, 0, index);

CheckTrack(trackSet, operand1, 1, index);

CheckTrack(trackSet, operand2, 2, index);

}

}

return trackSet;

}

The CheckTrack method takes the track set, one of the operands in an assembly code instruction, the position of the operand (zero, one, or two), and the index of the assembly code instruction in the assembly code list.

private void CheckTrack(ISet<Track> trackSet, object operand,

int position, int index) {

if (operand is Track) {

Track track = (Track) operand;

trackSet.Add(track);

track.AddEntry(position, index);

}

}

### Function Calls

The BaseRegister method returns the base register of a symbol that represents an auto or register variable or parameter. However, the symbol can also be null, in case of an unnamed parameter. The method returns the regular frame pointer or the ellipse frame pointer. Note that we always use the regular frame pointer in non-elliptic functions.

public Register BaseRegister(Symbol symbol) {

Assert.ErrorXXX((symbol == null) || symbol.IsAutoOrRegister());

We return the elliptic frame pointer is the function is elliptic and the symbol is not a parameter. If the symbol is null, it represents an unnamed parameter in which case the elliptic frame pointer is returned.

if (SymbolTable.CurrentFunction.Type.IsEllipse() &&

(symbol != null) && !symbol.IsParameter()) {

return AssemblyCode.EllipseRegister;

}

If the function is not elliptic, or if the symbol is a parameter, we return the regular frame pointer. The activation record is organized so that the parameters are located before the variables, with potential extra parameter between the regular parameters and the variables. Therefore, we use the regular frame pointer for parameters in both elliptic and non-elliptic functions, and the elliptic frame pointer for variables in elliptic functions since they are located after the extra parameters.

else {

return AssemblyCode.FrameRegister;

}

}

The FunctionPreCall methods is called before the actual function call, and before the parameter list of the function call. The idea is that each expression in the parameter shall have access to all the registers and the whole floating-point value stack. Therefore, we need to store the values currently stored in registers and on the floating-point value stack at the activation record. The FunctionPostCall method below restores the values after the function call.

The m\_totalExtraSize field keeps track of the current record size, which increases with nested calls. The m\_recordSizeStack field holds the record sizes of nested calls. Tech, we could manage with only the stack, and sum the record sizes to find the total record size. However, we include m\_totalExtraSize for effency reasons.

private int m\_totalExtraSize = 0;

private Stack<int> m\_recordSizeStack = new Stack<int>();

As mentioned above, each parameter expression shall have access to all registers. Therefore, we need to save the track map of each nested call in the m\_trackMapStack field. We need also keep track of the position of each saved register on the activation record in m\_registerMapStack. Both stacks are popped in the PostFunctionCall method below.

private Stack<IDictionary<Symbol,Track>> m\_trackMapStack =

new Stack<IDictionary<Symbol,Track>>();

private Stack<IDictionary<Track,int>> m\_registerMapStack =

new Stack<IDictionary<Track,int>>();

In the FunctionPreCall method, we start by obtaining the base register, which is the regular or elliptic frame pointer, depending on whether the function is elliptic. The recordSize variable is the original size of the activation record while the extraSize field holds the size of the extra values int the registers and on the floating-point value stack.

public void FunctionPreCall(MiddleCode middleCode) {

Register baseRegister = BaseRegister(null);

int recordSize = (int) middleCode[0], extraSize = 0;

We load the post map with the tracks of the register with their locations on the activation record. For each register to be stored on the activation record, we add a mov-instruction. The extraSize field is increased for each register.

IDictionary<Track,int> registerMap = new Dictionary<Track,int>();

foreach (KeyValuePair<Symbol, Track> pair in m\_trackMap) {

Track track = pair.Value;

AddAssemblyCode(AssemblyOperator.mov, baseRegister,

recordSize + extraSize, track);

registerMap.Add(track, recordSize + extraSize);

Symbol symbol = pair.Key;

extraSize += symbol.Type.Size();

}

After the registers have been stored on the activation record, we store the values of the floating-point value stack on the activation record. Again, the extraSize field is being increased with the size of the values.

for (int count = 0; count < m\_floatStackSize; ++count) {

AddAssemblyCode(AssemblyOperator.fstp\_qword, baseRegister,

recordSize + extraSize);

extraSize += 8;

}

When the values have been stored on the activation record, we need to store the register map for the registers to be restored after the function call in the PostFunctionCall method below.

m\_registerMapStack.Push(registerMap);

m\_recordSizeStack.Push(extraSize);

We also use the m\_totalExtraSize field to store the total extra size for the in case of nested calls. Technically, we do not need this field since we could obtain the same value by summarizing the values pushed on the m\_recordSizeStack stack. However, it has been included for efficiency reasons.

m\_totalExtraSize += extraSize;

The track map is pushed on the stack, and a new object is instanciated as the new track map. The track map will be restored after the function call.

m\_trackMapStack.Push(m\_trackMap);

m\_trackMap = new Dictionary<Symbol, Track>();

}

The m\_returnFloating is set to true if the function returns a floating value. In that case the value is stored on the floating-point value stack and need to be considered by the PostCallFunction method below.

private bool m\_returnFloating = false;

When calling a function there are several cases to consider: the function symbol holding the function to be called may be a proper function or a pointer to a function, and the caller or callee function (or both) may be elliptic.

public void FunctionCall(MiddleCode middleCode, int index) {

int recordSize = ((int) middleCode[1]) +

m\_totalExtraSize;

Symbol calleeSymbol = (Symbol) middleCode[0];

int extraSize = (int) middleCode[2];

The type is the function or, in case of a pointer the pointer type, the function the pointer points at.

Type calleeType = calleeSymbol.Type.IsFunction()

? calleeSymbol.Type : calleeSymbol.Type.PointerType;

Both the caller and callee function may be elliptic.

bool callerEllipse = SymbolTable.CurrentFunction.Type.IsEllipse(),

calleeEllipse = calleeType.IsEllipse();

The frame register is the elliptic frame pointer register if the caller function is elliptic. If it not, it is the regular frame pointer register.

Register frameRegister = callerEllipse ? AssemblyCode.EllipseRegister

: AssemblyCode.FrameRegister;

We start by adding the assembly code instruction for the return address. To begin with, we set it to the index of the next middle code instruction. However, the address will be changed by the WindowsJumpInfo method to a proper assembly code address.

AddAssemblyCode(AssemblyOperator.return\_address, frameRegister,

recordSize + SymbolTable.ReturnAddressOffset,

(BigInteger) (index + 1));

We set the current regular frame pointer and, in case of an elliptic caller function, the elliptic pointer to their offset in the activations record of the callee function. In this way, we can reset their values in accordance with the caller function returning from the function call.

AddAssemblyCode(AssemblyOperator.mov, frameRegister,

recordSize + SymbolTable.RegularFrameOffset,

AssemblyCode.FrameRegister);

if (callerEllipse) {

AddAssemblyCode(AssemblyOperator.mov, frameRegister,

recordSize + SymbolTable.EllipseFrameOffset,

AssemblyCode.EllipseRegister);

}

We add the size of the caller function’s activation record to the frame pointer, so that it pointes at the activation record of the callee function.

AddAssemblyCode(AssemblyOperator.add, frameRegister, // add di, 10

(BigInteger) recordSize);

If the caller function is elliptic, the frame pointer that we just added is the elliptic frame pointer, and we also need to set the regular frame pointer to the same value.

if (callerEllipse) { // mov bp, di

AddAssemblyCode(AssemblyOperator.mov, AssemblyCode.FrameRegister,

AssemblyCode.EllipseRegister);

}

If the callee function, but not the caller function, is elliptic, the frame pointer that we just added is the elliptic frame pointer, and we also need to set the regular frame pointer to the same value.

else if (calleeEllipse) {

AddAssemblyCode(AssemblyOperator.mov, AssemblyCode.EllipseRegister,

AssemblyCode.FrameRegister);

}

If the callee function is elliptic, and the extra size is more than zero (there are extra arguments in the function calls that are not matched to the declared parameters), we add the size to the elliptic frame pointer. The callee function will have both a regular frame pointer and an elliptic frame pointer. However, the elliptic pointer will point to a higher address to give space to the extra arguments.

if (calleeEllipse && (extraSize > 0)) {

AddAssemblyCode(AssemblyOperator.add, AssemblyCode.EllipseRegister,

(BigInteger) extraSize);

}

Finally. if the callee function is a function (not a pointer to a function) we a middle code call instruction. The m\_returnFloating field is set to true if the callee function returns a floating value. It will later be inspected by the FunctionPostCall method.

if (calleeSymbol.Type.IsFunction()) {

AddAssemblyCode(AssemblyOperator.call, calleeSymbol.UniqueName);

m\_returnFloating = calleeSymbol.Type.ReturnType.IsFloating();

}

If the callee function is a pointer to a function, we load its value into the jumpTrack track. The call becomes in this case a jump to the address stored in jumpTrack.

else {

Track jumpTrack = LoadValueToRegister(calleeSymbol);

AddAssemblyCode(AssemblyOperator.jmp, jumpTrack);

m\_returnFloating =

calleeSymbol.Type.PointerType.ReturnType.IsFloating();

}

}

The FunctionPostCall method is called after a function call. Its task is to restore the actions performed by FunctionPreCall above.

public void FunctionPostCall(MiddleCode middleCode) {

Register baseRegister = BaseRegister(null);

m\_trackMap = m\_trackMapStack.Pop();

IDictionary<Track,int> registerMap = m\_registerMapStack.Pop();

First, we iterate through the post map defined in FunctionPreCall before the function call. We load the values stored in the map into the registers.

foreach (KeyValuePair<Track,int> pair in registerMap) {

Track track = pair.Key;

int offset = pair.Value;

AddAssemblyCode(AssemblyOperator.mov, track, baseRegister,offset);

}

Then we investigate the floating value stack. If floating stack value before the function call was non-empty, we need to restore it.

if (m\_floatStackSize > 0) {

int recordOffset = (int) middleCode[2];

int recordSize = m\_recordSizeStack.Pop();

If the return value of the previous function call is a floating value, we need to temporary restore it when we restore the floating value stack, since the return value shall be placed at the top of the floating stack.

if (m\_returnFloating) {

AddAssemblyCode(AssemblyOperator.fstp\_qword, baseRegister,

recordOffset + recordSize);

}

We iterate through the floating values, as they are located above the activation record. We begin by pushing the top-most value and iterate to the bottom-most value.

int currentOffset = recordOffset + recordSize;

for (int count = 0; count < m\_floatStackSize; ++count) {

currentOffset -= 8;

AddAssemblyCode(AssemblyOperator.fld\_qword, baseRegister,

currentOffset);

}

If the function returned a floating value, we push it from its temporary location to the top of the floating value stack.

if (m\_returnFloating) {

AddAssemblyCode(AssemblyOperator.fld\_qword, baseRegister,

recordOffset + recordSize);

}

m\_totalExtraSize -= recordSize;

}

else {

m\_totalExtraSize -= m\_recordSizeStack.Pop();

}

}

### Loading Values into Registers

The LoadValueToRegister method loads the value of a symbol into a register. A specific register may be given. However, the register is represented by a track object that will be replaced by a proper register by the register allocator.

public Track LoadValueToRegister(Symbol symbol,

Register? register = null) {

If the register is specified, we must check that no other value is stored in that register. If it is, we must move it to another register.

if (register != null) {

CheckRegister(symbol, register.Value);

}

Then we check if the value is already stored in a register. If it is, we look up it track and compare the specified register with the register of the track. If the

Track track;

if (m\_trackMap.TryGetValue(symbol, out track)) {

m\_trackMap.Remove(symbol);

If the registers are not null do not overlap, we need to remove the previous register from the track. If we find that the register is associated to a track, we add a mov-instruction that moves the value to a new (yet unknown) register. The track holding that register is then returned.

if ((register != null) && (track.Register != null) &&

!AssemblyCode.RegisterOverlap(register, track.Register)) {

Track newTrack = new Track(symbol, register.Value);

AddAssemblyCode(AssemblyOperator.set\_track\_size,

newTrack, track);

AddAssemblyCode(AssemblyOperator.mov, newTrack, track);

return newTrack;

}

If there is no non-overlapping register we set the register of the track, if it is not null, and return the track.

else {

if (register != null) {

track.Register = register;

}

return track;

}

}

If the track register is null, we assign it the specified register (which may or may not be null).

However, if both the specified register and the track register is non-null, and they do not overlap each other, we need to move the value of the track register to specified register. We create and return a new track for the specified register.

If the symbol does not already have a track, we create a new track for it.

Basically, we have three different cases: the symbol holds an integer value, it holds and address, or a value stored at an address. In case of an integer value, we simply load it int to the register.

if (symbol.Value is BigInteger) {

AddAssemblyCode(AssemblyOperator.mov, track, symbol.Value);

}

If the symbol is an array, function, or string, or if it holds a static address, we load the address of the value into the register, not the value itself. We load the base of the symbol into the register. The base is the symbol’s unique name if it is extern or static, the regular or ellipse frame register in case of an auto or register array (a function or string is always static), and the name of a static address. If the offset of the address is non-zero, we add it to the register.

else if (symbol.Type.IsArrayFunctionOrString() ||

(symbol.Value is StaticAddress)) {

AddAssemblyCode(AssemblyOperator.mov, track,

Base(symbol));

int offset = Offset(symbol);

if (offset != 0) {

AddAssemblyCode(AssemblyOperator.add, track,

(BigInteger) offset);

}

}

In all other cases, we load the value of the symbol into the register. The base is the symbol’s unique name if it is extern or static, and the regular or ellipse frame register is it is auto or register.

else {

AddAssemblyCode(AssemblyOperator.mov, track,

Base(symbol), Offset(symbol));

}

return track;

}

}

If the symbol is a static address, we load its address into the register. At the moment, we load its name into the register and add its offset, if more than zero. The name will later be replaced by the address by the linker.

If the symbol is an array, we load its base into the register and add its offset, if more than zero. The base may be the unique name of a static symbol, or the regular or elliptic frame pointer in case of a parameter or local variable.

If the symbol is not a constant, an array, or a static function, struct, or union, we load the value of the symbol to the register with the help of the symbol’s base and offset. The base may be the unique name of a static symbol, or the regular or elliptic frame pointer.

The CheckRegister method checks whether a symbol value is stored in the register

public void CheckRegister(Symbol symbol, Register register) {

foreach (KeyValuePair<Symbol,Track> entry in m\_trackMap) {

Symbol oldSymbol = entry.Key;

Track oldTrack = entry.Value;

We iterate through the track map and check if any track holds a register overlapping the specified register.

if (!oldSymbol.Equals(symbol) &&

AssemblyCode.RegisterOverlap(register, oldTrack.Register)) {

If the find a match, we create a new track and insert instructions for setting the size of the track and moving the value from the previous track to the new track.

Track newTrack = new Track(oldSymbol);

m\_trackMap[oldSymbol] = newTrack;

int lastLine;

for (lastLine = m\_assemblyCodeList.Count - 1; lastLine >= 0;

--lastLine) {

AssemblyCode assemblyCode = m\_assemblyCodeList[lastLine];

if (oldTrack.Equals(assemblyCode[0])) {

break;

}

}

Assert.ErrorXXX(lastLine >= 0);

AssemblyCode setCode =

new AssemblyCode(AssemblyOperator.set\_track\_size,

newTrack, oldTrack);

AssemblyCode movCode =

new AssemblyCode(AssemblyOperator.mov, newTrack, oldTrack);

m\_assemblyCodeList.Insert(lastLine + 1, setCode);

m\_assemblyCodeList.Insert(lastLine + 2, movCode);

break;

}

}

}

The LoadAddressToRegister loads the address of value, rather the value itself, into a register.

public Track LoadAddressToRegister(Symbol symbol,

Register? register = null) {

If the address symbol of the symbol is not null, we simply call LoadValueToRegister with the address symbol, and returns the track. However, we must mark the track as a pointer, which means that the set of possible register re will be restricted.

if (symbol.AddressSymbol != null) {

Track addressTrack = LoadValueToRegister(symbol.AddressSymbol);

addressTrack.Pointer = true;

return addressTrack;

}

If the address symbol is null, we need to obtain the address manually. We start by mov the base of the symbol to the register. The base is the regular or ellipse frame pointer in case of an auto or register storage, since the value is located on the activation record, or the name of the symbol in case of extern or static storage. However, the symbol may have a static address its value. In that case, the base is the name of the static address.

else {

Symbol addressSymbol = new Symbol(new Type(symbol.Type));

Track addressTrack = new Track(addressSymbol, register);

Assert.ErrorXXX((addressTrack.Register == null) ||

RegisterAllocator.PointerRegisterSetWithEllipse.

Contains(addressTrack.Register.Value));

addressTrack.Pointer = true;

Assert.ErrorXXX(!(symbol.Value is BigInteger));

AddAssemblyCode(AssemblyOperator.mov, addressTrack, Base(symbol));

Then we add the offset of the symbol, which is always non-zero in case of auto or register storage and zero in case of extern or static storage. In case a static address, the offset may be zero or non-zero.

int offset = Offset(symbol);

if (offset != 0) {

AddAssemblyCode(AssemblyOperator.add, addressTrack,

(BigInteger) offset);

}

return addressTrack;

}

}

### Return, Exit, and Goto

The Return method generates the code for returning from a function call. We catch the return address (which we must do before we reset the pointers), reset the regular frame pointer and elliptic frame pointer, and jump back to the return address.

public void Return(MiddleCode middleCode) {

Assert.ErrorXXX(m\_floatStackSize == 0);

Track track = new Track(Type.VoidPointerType);

AddAssemblyCode(AssemblyOperator.mov, track,

AssemblyCode.FrameRegister,

SymbolTable.ReturnAddressOffset);

AddAssemblyCode(AssemblyOperator.mov, AssemblyCode.EllipseRegister,

AssemblyCode.FrameRegister,

SymbolTable.EllipseFrameOffset);

AddAssemblyCode(AssemblyOperator.mov, AssemblyCode.FrameRegister,

AssemblyCode.FrameRegister,

SymbolTable.RegularFrameOffset);

AddAssemblyCode(AssemblyOperator.jmp, track);

}

The IntegralGetReturnValue method stores the return register in a track.

public void IntegralGetReturnValue(MiddleCode middleCode) {

Symbol returnSymbol = (Symbol) middleCode[0];

Register returnRegister =

AssemblyCode.RegisterToSize(AssemblyCode.ReturnValueRegister,

returnSymbol.Type.Size());

The call to CheckRegister makes sure that if the register already has a value, it is moved to another register.

CheckRegister(returnSymbol, returnRegister);

Track returnTrack = new Track(returnSymbol, returnRegister);

m\_trackMap.Add(returnSymbol, returnTrack);

AddAssemblyCode(AssemblyOperator.empty, returnTrack);

}

The IntegralSetReturnValue loads the symbol of the expression into the return register.

public void IntegralSetReturnValue(MiddleCode middleCode) {

Symbol returnSymbol = (Symbol) middleCode[1];

Register returnRegister =

AssemblyCode.RegisterToSize(AssemblyCode.ReturnValueRegister,

returnSymbol.Type.Size());

LoadValueToRegister(returnSymbol, returnRegister);

m\_trackMap.Remove(returnSymbol);

}

The Exit method generates code for exit the execution of the program. However, the code is different depending on whether we generate Windows or Linux code, and whether there is an exit symbol given.

public void Exit(MiddleCode middleCode) {

Symbol exitSymbol = (Symbol) middleCode[0];

If the exit symbol is not null, we load its value into the rdi register If it is null, we just load zero into the register. This value will be returned to the surrounding operation system.

if (Start.Linux) {

if (exitSymbol != null) {

LoadValueToRegister(exitSymbol, Register.rdi);

}

else {

AddAssemblyCode(AssemblyOperator.mov, Register.rdi,

BigInteger.Zero);

}

The Linux exit code is to load value 60 to register rax and do a system call.

AddAssemblyCode(AssemblyOperator.mov, Register.rax,

(BigInteger) 60); // 0x3C

AddAssemblyCode(AssemblyOperator.syscall);

}

In case of Windows code, we load the exit symbol value into register al.

if (Start.Windows) {

if (exitSymbol != null) {

LoadValueToRegister(exitSymbol, Register.al);

}

else {

AddAssemblyCode(AssemblyOperator.mov, Register.al,

BigInteger.Zero);

}

The Windows exit code is to load value 76 to register ah and do interrupt number 33.

AddAssemblyCode(AssemblyOperator.mov, Register.ah,

(BigInteger) 76); // 0x4C

AddAssemblyCode(AssemblyOperator.interrupt, (BigInteger) 33); // 0x21

}

}

The Goto method simply add a jump instruction with the index a middle code instruction as target. The target will be changed by the WindowsJumpInfo method to a proper assembly code target.

public void Goto(MiddleCode middleCode) {

int jumpTarget = (int) middleCode[0];

AddAssemblyCode(AssemblyOperator.jmp, null, null, jumpTarget);

}

### Load and Inspect Registers

The LoadToRegister method adds assembly code that loads the value of a symbol into a specified register. This method, as well as InspectRegister, CarryExpression, JumpRegister, Interrupt, and SystemCall, is only used by the standard library in internal system calls.

private ISet<Track> m\_syscallSet = new HashSet<Track>();

public void LoadToRegister(MiddleCode middleCode) {

Register register = (Register) middleCode[0];

Symbol symbol = (Symbol) middleCode[1];

Track track = LoadValueToRegister(symbol, register);

m\_syscallSet.Add(track);

}

The InspectRegister method loads the value of a register to a symbol. More specifically, it adds a track holding the symbol and register to the track map.

public void InspectRegister(MiddleCode middleCode) {

Symbol symbol = (Symbol) middleCode[0];

Register register = (Register) middleCode[1];

Track track = new Track(symbol, register);

m\_trackMap.Add(symbol, track);

}

The CarryExpression method adds assembly code that jumps to a middle code target if the carry flag is set. The target will later be changed by the WindowsJumpInfo method to a proper assembly code target.

public void CarryExpression(MiddleCode middleCode) {

AssemblyOperator objectOperator =

m\_middleToIntegralMap[middleCode.Operator];

int jumpTarget = (int) middleCode[0];

AddAssemblyCode(objectOperator, null, null, jumpTarget);

}

The JumpRegister method adds assembly code that jumps to the address stored in a register.

public void JumpToRegister(MiddleCode middleCode) {

Register jumpRegister = (Register) middleCode[0];

AddAssemblyCode(AssemblyOperator.jmp, jumpRegister);

}

The Interrupt method adds code that performs an interrupt call. Before the call, we clear the system call set.

public void Interrupt(MiddleCode middleCode) {

foreach (Track track in m\_syscallSet) {

AddAssemblyCode(AssemblyOperator.empty, track);

}

AddAssemblyCode(AssemblyOperator.interrupt,

(BigInteger) middleCode[0]);

m\_trackMap.Clear();

}

The SystemCall method adds code that performs an interrupt call. Like the interrupt case, we clear the system call set before the call.

public void SystemCall(MiddleCode middleCode) {

foreach (Track track in m\_syscallSet) {

AddAssemblyCode(AssemblyOperator.empty, track);

}

AddAssemblyCode(AssemblyOperator.syscall);

m\_trackMap.Clear();

}

### Initialization

The Initializer method adds code initializing a value. The value becomes static block, or a part of a static block, which

private void Initializer(MiddleCode middleCode) {

Sort sort = (Sort) middleCode[0];

object value = middleCode[1];

If the value is a static address, we add code for defining the address with its name and offset.

if (value is StaticAddress) {

StaticAddress staticAddress = (StaticAddress) value;

string name = staticAddress.UniqueName;

int offset = staticAddress.Offset;

// dw name + offset

AddAssemblyCode(AssemblyOperator.define\_address, name, offset);

}

Otherwise, we add code for defining the value. The initialization instruction will later be transformed to a sequence of instructions in the Linux case and memory block in the Windows case.

else {

AddAssemblyCode(AssemblyOperator.define\_value, sort, value);

}

}

The InitializerZero adds an instruction for a sequence of zero values, each holding one byte. Like the Initializer method above, the initialization instruction will later be transformed to a sequence of instructions in the Linux case and memory block in the Windows case.

private void InitializerZero(MiddleCode middleCode) {

int size = (int) middleCode[0];

Assert.ErrorXXX(size > 0);

AddAssemblyCode(AssemblyOperator.define\_zero\_sequence, size);

}

### Integral Multiplication, Division, and Modulo

The IntegralMultiply method adds assembly code for the integral operators multiplication, division, and module. These operations differ from the previous integral operations in that way that the value of left symbol is always assumed to be stored in a specific register. However, the right symbol is given in the operation. The value of the right symbol is given in the operations. It can be stored in a register or on a memory address. However, unlike the previous integral operations, we cannot give an integer value directly. Instead, we have to store the value on an address, and give the address to the operation.

Moreover, there are in fact four registers involved in the operation. The specific registers depend on the type size.

* The register where we store the value of the left symbol.
* The register holding the upper part of the left symbol value. We ignore the value in this register, but we must set it to zero and mark in the track map we have in fact altered its value.
* The result of multiplication or division, which is ignored by modulo.
* The result of module, which is ignored by multiplication or division.

Even if one of the last two registers is ignored, we still need to mark in the track map that is has in fact been assigned a new value.

The m\_leftRegisterMap map hold the registers where we store the value of the left symbol. The m\_zeroRegisterMap map holds the registers that era to be set to zero before the operation. The m\_productQuintentRegisterMap map holds the registers where the product or quintet is stored after a multiplication or division operations, while m\_remainderRegisterMap holds the registers of the remainder of a modulo operations.

public static IDictionary<int,Register> m\_leftRegisterMap =

new Dictionary<int,Register>() {{1, Register.al}, {2, Register.ax},

{4, Register.eax}, {8, Register.rax}};

public static IDictionary<int,Register> m\_zeroRegisterMap =

new Dictionary<int,Register>() {{1, Register.ah}, {2, Register.dx},

{4, Register.edx}, {8, Register.rdx}};

public static IDictionary<int,Register> m\_productQuintentRegisterMap =

new Dictionary<int,Register>() {{1, Register.al}, {2, Register.ax},

{4, Register.eax}, {8, Register.rax}};

public static IDictionary<int,Register> m\_remainderRegisterMap =

new Dictionary<int,Register>() {{1, Register.ah}, {2, Register.dx},

{4, Register.edx}, {8, Register.rdx}};

The IntegralMultiply method adds code for multiplication operations. To begin with, we load the value of the left symbol into the register given by m\_leftRegisterMap.

public void IntegralMultiply(MiddleCode middleCode) {

Symbol leftSymbol = (Symbol) middleCode[1];

int typeSize = leftSymbol.Type.SizeArray();

Register leftRegister = m\_leftRegisterMap[typeSize];

Track leftTrack = LoadValueToRegister(leftSymbol, leftRegister);

Then we clear the zero register by perfrom the exlusive-xor operation on itself, which is an effetive way to clear a register, in order to provide current input to the operation.

Register zeroRegister = m\_zeroRegisterMap[typeSize];

Track zeroTrack = new Track(leftSymbol, zeroRegister);

AddAssemblyCode(AssemblyOperator.xor, zeroTrack, zeroTrack);

The we call IntegralUnary to load the right symbol into a suitable storage and perform the operation.

Symbol rightSymbol = (Symbol) middleCode[2];

IntegralUnary(middleCode.Operator, rightSymbol, rightSymbol);

When the operation is done, we need to find out whiech register to keep and which one to discard. We only use one of the registers, depending on the operation. But we also need to mark that the value of the discarded register has been modified, to prevent that another value is stored in the register during the operation.

Register resultRegister, discardRegister;

In case of the modulo operation, the resulting register is picked from m\_remainderRegisterMap, and the register to be discarded m\_productQuintentRegisterMap. In case of multiplication or division, it is the other way around.

if ((middleCode.Operator == MiddleOperator.SignedModulo) ||

(middleCode.Operator == MiddleOperator.UnsignedModulo)) {

resultRegister = m\_remainderRegisterMap[typeSize];

discardRegister = m\_productQuintentRegisterMap[typeSize];

}

else {

resultRegister = m\_productQuintentRegisterMap[typeSize];

discardRegister = m\_remainderRegisterMap[typeSize];

}

We create a new track holding the result symbol and register. If the result symbol is a temporary variable, we store the track in the track map. We also add an empty instruction, to mark that the register has been altered by the operation.

Symbol resultSymbol = (Symbol) middleCode[0];

Track resultTrack = new Track(resultSymbol, resultRegister);

if (resultSymbol.IsTemporary()) {

Assert.ErrorXXX(resultSymbol.AddressSymbol == null);

m\_trackMap.Add(resultSymbol, resultTrack);

AddAssemblyCode(AssemblyOperator.empty, resultTrack);

}

If the result symbol is a regular variable, we store the value in the register at the address of the variable. In this case we do not need to add an empty instruction since the loading of the register marks that is has been in use.

else {

AddAssemblyCode(AssemblyOperator.mov, Base(resultSymbol),

Offset(resultSymbol), resultTrack);

}

However, we need to add an empty instruction for the discarded register, to mark that it has been altered.

Track discaredTrack = new Track(resultSymbol, discardRegister);

AddAssemblyCode(AssemblyOperator.empty, discaredTrack);

}

### Integral Assignment and Parameters

The IntegralAssign adds code assignment of an integral value. What makes it a bit complicated is that we must take into consideration of the assignment of a conditional expression. Both the values of the true and false expression shall be assigned to the same register.

public void IntegralAssign(MiddleCode middleCode) {

Symbol resultSymbol = (Symbol) middleCode[0],

assignSymbol = (Symbol) middleCode[1];

IntegralAssign(resultSymbol, assignSymbol);

}

The IntegralParameter method is quite simple. We only need to create the symbol representing the parameter, and then call IntegralAssign.

public void IntegralParameter(MiddleCode middleCode) {

Type toType = (Type) middleCode[0];

Symbol toSymbol = new Symbol(null, false, Storage.Auto, toType);

Symbol fromSymbol = (Symbol) middleCode[1];

int parameterOffset = (int) middleCode[2];

toSymbol.Offset = m\_totalExtraSize + parameterOffset;

IntegralAssign(toSymbol, fromSymbol);

}

The IntegralAssign method is central in the assembly code generator. It is called by IntegralAssign, IntegralParameter. We have main two cases: the value of left symbol may be stored in a register,

public void IntegralAssign(Symbol resultSymbol, Symbol assignSymbol) {

Track resultTrack = null, assignTrack = null;

m\_trackMap.TryGetValue(resultSymbol, out resultTrack);

m\_trackMap.TryGetValue(assignSymbol, out assignTrack);

int typeSize = assignSymbol.Type.SizeArray();

If the result symbol is temporary and its address symbol is null, we are facing the assignment of a conditional expression. The only other cases when a temporary symbol is assigned is in case of a dereferred, index, or arrow expression. However, in those cases, the address symbol is not null.

if (resultSymbol.IsTemporary()) {

Assert.ErrorXXX(resultSymbol.AddressSymbol == null);

If the assignment symbol is stored in the track map, which means that we assign the true expression of the conditional expression, we start by copying its track to the assignment track. Then we add the assignment track to the track map.

if (assignTrack != null) {

if (resultTrack != null) {

resultTrack.Replace(m\_assemblyCodeList, assignTrack);

}

m\_trackMap[resultSymbol] = assignTrack;

m\_trackMap.Remove(assignSymbol);

}

If the assignment symbol is not stored in the track map, which means that we assign the false expression of the conditional expression, we begin by creating a track for the result symbol.

else {

if (resultTrack == null) {

resultTrack = new Track(resultSymbol);

m\_trackMap.Add(resultSymbol, resultTrack);

}

If the assign symbol is an integer value, we move it into the register.

if (assignSymbol.Value is BigInteger) {

AddAssemblyCode(AssemblyOperator.mov, resultTrack,

assignSymbol.Value);

}

If the assign symbol is an array, function, string, or static address, we move its address rather than ist value into the register.

else if (assignSymbol.Type.IsArrayFunctionOrString() ||

(assignSymbol.Value is StaticAddress)) {

AddAssemblyCode(AssemblyOperator.mov, resultTrack,

Base(assignSymbol));

int offset = Offset(assignSymbol);

if (offset != 0) {

AddAssemblyCode(AssemblyOperator.add, resultTrack,

(BigInteger) offset);

}

}

Otherwise, we move the value of the symbol into the register.

else {

AddAssemblyCode(AssemblyOperator.mov, resultTrack,

Base(assignSymbol), Offset(assignSymbol));

}

}

}

else {

IntegralBinary(MiddleOperator.Assign, resultSymbol,

resultSymbol, assignSymbol);

}

}

If the result symbol is not temporary, we move the value directly into the address of the result symbol.

In case of an integer value, there is a special case due to technical limitations if the size is eight bytes. In that case we have to load the value into a register and load the register into the address, since we cannot load an eight-byte value directly into an address.

If the value size is not eight-byte, we load the value directly into the address.

If the symbol is an array, function, string, or static address, we load the address rather than the value onto the address.

Otherwise, we need to load the value into a register, which we than load onto the address.

### Unary Integral Operations

The Unary method is called to add code for unary subtraction and bitwise not. There is also unary addition, but it generates no code.

First, we need a map from middle code operators to assembly code operators for integral types. We will use the map in the IntegralUnary and IntegralBinary methods.

public static IDictionary<MiddleOperator,AssemblyOperator>

m\_middleToIntegralMap =

new Dictionary<MiddleOperator,AssemblyOperator>() {

{MiddleOperator.BitwiseNot, AssemblyOperator.not},

{MiddleOperator.UnarySubtract, AssemblyOperator.neg},

{MiddleOperator.SignedMultiply, AssemblyOperator.imul},

{MiddleOperator.SignedDivide, AssemblyOperator.idiv},

{MiddleOperator.SignedModulo, AssemblyOperator.idiv},

{MiddleOperator.UnsignedMultiply, AssemblyOperator.mul},

{MiddleOperator.UnsignedDivide, AssemblyOperator.div},

{MiddleOperator.UnsignedModulo, AssemblyOperator.div},

{MiddleOperator.Assign, AssemblyOperator.mov},

{MiddleOperator.BinaryAdd, AssemblyOperator.add},

{MiddleOperator.BinarySubtract, AssemblyOperator.sub},

{MiddleOperator.BitwiseAnd, AssemblyOperator.and},

{MiddleOperator.BitwiseOr, AssemblyOperator.or},

{MiddleOperator.BitwiseXOr, AssemblyOperator.xor},

{MiddleOperator.ShiftLeft, AssemblyOperator.shl},

{MiddleOperator.ShiftRight, AssemblyOperator.shr},

{MiddleOperator.Equal, AssemblyOperator.je},

{MiddleOperator.NotEqual, AssemblyOperator.jne},

{MiddleOperator.Carry, AssemblyOperator.jc},

{MiddleOperator.NotCarry, AssemblyOperator.jnc},

{MiddleOperator.Compare, AssemblyOperator.cmp},

{MiddleOperator.SignedLessThan, AssemblyOperator.jl},

{MiddleOperator.SignedLessThanEqual,AssemblyOperator.jle},

{MiddleOperator.SignedGreaterThan, AssemblyOperator.jg},

{MiddleOperator.SignedGreaterThanEqual, AssemblyOperator.jge},

{MiddleOperator.UnsignedLessThan, AssemblyOperator.jb},

{MiddleOperator.UnsignedLessThanEqual, AssemblyOperator.jbe},

{MiddleOperator.UnsignedGreaterThan, AssemblyOperator.ja},

{MiddleOperator.UnsignedGreaterThanEqual, AssemblyOperator.jae}};

public void IntegralUnary(MiddleCode middleCode) {

Symbol resultSymbol = (Symbol)middleCode[0],

unarySymbol = (Symbol)middleCode[1];

IntegralUnary(middleCode.Operator, resultSymbol, unarySymbol);

}

public void IntegralUnary(MiddleOperator middleOperator,

Symbol resultSymbol, Symbol unarySymbol) {

AssemblyOperator objectOperator =

m\_middleToIntegralMap[middleOperator];

int typeSize = unarySymbol.Type.SizeArray();

Track unaryTrack = null;

m\_trackMap.TryGetValue(unarySymbol, out unaryTrack);

if (unarySymbol.Value is BigInteger) {

SymbolTable.StaticSet.Add(ConstantExpression.Value(unarySymbol));

AddAssemblyCode(objectOperator, unarySymbol.UniqueName,

0, null, typeSize);

}

We have to cases. If the result symbol equals the unary symbol, we do not need to use a register. We can just perform the operation directly on the address of the symbol. Unless the operator is unary addition, in which case we do nothing.

else if (resultSymbol == unarySymbol) {

Assert.ErrorXXX(unaryTrack == null);

if (middleOperator != MiddleOperator.UnaryAdd) {

AddAssemblyCode(objectOperator, Base(unarySymbol),

Offset(unarySymbol), null, typeSize);

}

}

If the result symbol does not equal the unary symbol, we cannot perform the operation directly. We need to store the value in a register, perform the operation on the register, and store its value in the result symbol. Again, unless the operator is unary addition, in which case we do not perform the operations. However, the value is stored in the same way as the other operators.

else {

unaryTrack = LoadValueToRegister(unarySymbol);

if (middleOperator != MiddleOperator.UnaryAdd) {

AddAssemblyCode(objectOperator, unaryTrack);

}

If the result symbol is a temporary variable, we add the unary track to the result symbol in the track map.

if (resultSymbol.IsTemporary()) {

Assert.ErrorXXX(resultSymbol.AddressSymbol == null);

m\_trackMap.Add(resultSymbol, unaryTrack);

}

If the result symbol is a regular variable, we store the resulting value of the register on its address.

else {

AddAssemblyCode(AssemblyOperator.mov, Base(resultSymbol),

Offset(resultSymbol), unaryTrack);

}

m\_trackMap.Remove(unarySymbol);

}

}

### Integral Binary

public void IntegralBinary(MiddleCode middleCode) {

Symbol resultSymbol = (Symbol)middleCode[0],

leftSymbol = (Symbol)middleCode[1],

rightSymbol = (Symbol)middleCode[2];

IntegralBinary(middleCode.Operator, resultSymbol,

leftSymbol, rightSymbol);

}

public void IntegralRelation(MiddleCode middleCode) {

Symbol leftSymbol = (Symbol) middleCode[1],

rightSymbol = (Symbol) middleCode[2];

IntegralBinary(MiddleOperator.Compare, null, leftSymbol, rightSymbol);

AssemblyOperator objectOperator =

m\_middleToIntegralMap[middleCode.Operator];

int target = (int) middleCode[0];

AddAssemblyCode(objectOperator, null, null, target);

}

The IntegralBinary method generates assembly code for binary integral, addition, bitwise, and shift expressions. Basically, we have two cases: we to load the left operand into a register, or we perform the operation on the address of the left operand. In both cases, the right operand may be stored in a register or on an address, it may also be an integer value or and address.

public void IntegralBinary(MiddleOperator middleOperator,

Symbol resultSymbol, Symbol leftSymbol,

Symbol rightSymbol) {

Track leftTrack = null, rig htTrack = null;

m\_trackMap.TryGetValue(leftSymbol, out leftTrack);

m\_trackMap.TryGetValue(rightSymbol, out rightTrack);

If the left or right symbol is a static address with a non-zero offset or auto or register array (which offset is always non-zero), we have to load to their symbol into a register. XXX

if ((leftTrack == null) &&

(middleOperator != MiddleOperator.Assign) &&

(((resultSymbol != null) && (resultSymbol != leftSymbol)) ||

(leftSymbol.Type.IsArrayFunctionOrString() ||

(leftSymbol.Value is StaticAddress)))) {

leftTrack = LoadValueToRegister(leftSymbol);

}

if ((rightTrack == null) &&

(middleOperator != MiddleOperator.BinaryAdd) &&

(middleOperator != MiddleOperator.BinarySubtract) &&

(middleOperator != MiddleOperator.Assign) &&

((rightSymbol.Type.IsArrayFunctionOrString() &&

(rightSymbol.Offset != 0)) ||

((rightSymbol.Value is StaticAddress) &&

(((StaticAddress) rightSymbol.Value).Offset != 0)))) {

rightTrack = LoadValueToRegister(rightSymbol);

}

if ((rightTrack == null) && (rightSymbol.Value is BigInteger) &&

((middleOperator != MiddleOperator.Assign) ||

((leftTrack == null) &&

!(leftSymbol.Type.IsArrayFunctionOrString() ||

(leftSymbol.Value is StaticAddress))))) {

BigInteger bigValue = (BigInteger) rightSymbol.Value;

if (!((-2147483648 <= bigValue) && (bigValue <= 2147483647))) {

rightTrack = LoadValueToRegister(rightSymbol);

}

}

In case of the left or right shift operator, we need to load the right symbol into a specific register due to technical limitations.

if (MiddleCode.IsShift(middleOperator)) {

rightTrack =

LoadValueToRegister(rightSymbol, AssemblyCode.ShiftRegister);

}

int typeSize = leftSymbol.Type.Size();

AssemblyOperator objectOperator = m\_middleToIntegralMap[middleOperator];

Assert.ErrorXXX(!(leftSymbol.Value is BigInteger));

If the left symbol if not stored in a register and equals the result value, we do not need to load it into a register. We can just use the address of the left symbol. However, the right symbol needs in some cases to be loaded into a register.

If the right symbol is already stored in a register, we simply use that register.

if (leftTrack != null) {

if (rightTrack != null) {

AddAssemblyCode(objectOperator, leftTrack, rightTrack);

}

If the right symbol is an integer value, we use that value directly.

else if (rightSymbol.Value is BigInteger) {

AddAssemblyCode(objectOperator, leftTrack, rightSymbol.Value);

}

If the right symbol is an array, function, string, or static address, we use its name. Note that we do not need to look into the offset. We can be sure that it is zero, since the symbol would have been loaded into a register at the beginning of this method if the offset had been non-zero.

else if (rightSymbol.Type.IsArrayFunctionOrString() ||

(rightSymbol.Value is StaticAddress)) {

AddAssemblyCode(objectOperator, leftTrack, Base(rightSymbol));

int rightOffset = Offset(rightSymbol);

if (middleOperator == MiddleOperator.Assign) {

AddAssemblyCode(AssemblyOperator.add, leftTrack,

(BigInteger) rightOffset);

}

else {

AddAssemblyCode(objectOperator, leftTrack,

(BigInteger) rightOffset);

}

}

If none of the above cases apply, we need to load the right symbol into a register.

else {

if (middleOperator == MiddleOperator.Assign) {

leftTrack = new Track(resultSymbol);

}

AddAssemblyCode(objectOperator, leftTrack,

Base(rightSymbol), Offset(rightSymbol));

}

The only remaining question is what to do with the resulting value. If the result is a temporary variable, we add it to the track map to be used by succeeding instructions.

if (resultSymbol != null) {

if (resultSymbol.IsTemporary()) {

Assert.ErrorXXX(resultSymbol.AddressSymbol == null);

m\_trackMap.Add(resultSymbol, leftTrack);

}

If the result is not a temporary variable, we store the value on its address.

else {

AddAssemblyCode(AssemblyOperator.mov, Base(resultSymbol),

Offset(resultSymbol), leftTrack);

}

}

}

If the right symbol is stored in a register, we simply use it.

else {

if (rightTrack != null) {

AddAssemblyCode(objectOperator, Base(leftSymbol),

Offset(leftSymbol), rightTrack);

}

If the right symbol is an integer value, we simply use it.

else if (rightSymbol.Value is BigInteger) {

AddAssemblyCode(objectOperator, Base(leftSymbol),

Offset(leftSymbol), rightSymbol.Value, typeSize);

}

If the right symbol is an array, function, string, or static value, we use its address.

else if (rightSymbol.Type.IsArrayFunctionOrString() ||

(rightSymbol.Value is StaticAddress)) {

AddAssemblyCode(objectOperator, Base(leftSymbol),

Offset(leftSymbol), Base(rightSymbol), typeSize);

int rightOffset = Offset(rightSymbol);

if (rightOffset != 0) {

if (middleOperator == MiddleOperator.Assign) {

AddAssemblyCode(AssemblyOperator.add, Base(leftSymbol),

Offset(leftSymbol), (BigInteger) rightOffset,

typeSize);

}

else {

AddAssemblyCode(objectOperator, Base(leftSymbol),

Offset(leftSymbol), (BigInteger) rightOffset,

typeSize);

}

}

}

If none of the above cases apply, we use the base and offset of the right symbol.

else {

rightTrack = LoadValueToRegister(rightSymbol);

AddAssemblyCode(objectOperator, Base(leftSymbol),

Offset(leftSymbol), rightTrack);

}

}

Finally, we need to remove the left and right symbol from the track map, since they shall not be used any more.

m\_trackMap.Remove(leftSymbol);

m\_trackMap.Remove(rightSymbol);

}

### Base and Offset

The Base method returns the base of a symbol, which may be a register, a track, or a name. If the symbol’s value if a static address we return its name, which will later be replaced by a proper address by the linker.

private object Base(Symbol symbol) {

Assert.ErrorXXX(!(symbol.Value is BigInteger));

if (symbol.Value is StaticAddress) {

StaticAddress staticAddress = (StaticAddress) symbol.Value;

return staticAddress.UniqueName;

}

If the symbol’s address symbol is not null, the symbol represents a dereferred symbol in a dereference, index, or arrow expression, and we return the register holding the address. We look up the address track (the track of the address symbol) by calling LoadValueToRegister and mark the track as pointer, which means that the register allocator will choose the register from a smaller set. Since the track is no longer needed, we delete it from the track map before we return it.

else if (symbol.AddressSymbol != null) {

Track addressTrack = LoadValueToRegister(symbol.AddressSymbol);

Assert.ErrorXXX((addressTrack.Register == null) ||

RegisterAllocator.PointerRegisterSetWithEllipse.

Contains(addressTrack.Register.Value));

addressTrack.Pointer = true;

m\_trackMap.Remove(symbol.AddressSymbol);

return addressTrack;

}

If the symbol is extern or static, we return its unique name, which will later be replaced by a proper address by the linker.

else if (symbol.IsExternOrStatic()) {

return symbol.UniqueName;

}

Finally, if the symbol is auto or register, we call BaseRegister to obtain the regular or ellipse frame pointer.

else { //resultSymbol.IsAutoOrRegister()

return BaseRegister(symbol);

}

}

The Offset method returns the offset of the symbol. If its value is a static address, we return the offset of the static address.

private int Offset(Symbol symbol) {

Assert.ErrorXXX(!(symbol.Value is BigInteger));

if (symbol.Value is StaticAddress) {

StaticAddress staticAddress = (StaticAddress) symbol.Value;

return staticAddress.Offset;

}

If the symbol’s address symbol is not null, we return the offset of the address symbol.

else if (symbol.AddressSymbol != null) {

return symbol.AddressOffset;

}

Finally, in all other case we return the offset of the symbol. The offset is significant in case of auto or register parameters and variables. In case of an extern or static symbol, the offset is always zero and we will not use the value.

else {

return symbol.Offset;

}

}

### Case

The Case is method basically checks whether the switch symbol equals a value. Normally, we load the value of a symbol into a register, which we add to the track map. After the operation has been performed, we remove the register from the track map. However, in case of a sequence of case statements, we load the value of the switch symbol into a register, and then we keep it in the register during all the case test instructions.

public void Case(MiddleCode middleCode) {

Symbol switchSymbol = (Symbol) middleCode[1];

Track switchTrack = LoadValueToRegister(switchSymbol);

Symbol caseSymbol = (Symbol) middleCode[2];

BigInteger caseValue = (BigInteger) caseSymbol.Value; // cmp ax, 123

AddAssemblyCode(AssemblyOperator.cmp, switchTrack, caseValue);

int target = (int) middleCode[0];

AddAssemblyCode(AssemblyOperator.je, null, null, target);

Note that we do add the symbol to the track map, in order to keep the value in the register thought out all the succeeding case instruction of the switch statement.

m\_trackMap.Add(switchSymbol, switchTrack);

}

The CaseEnd method is called after the last case instruction of the switch statement. We only remove the switch symbol from the track map, since we do not need to keep the value in the register after the last case instruction.

public void CaseEnd(MiddleCode middleCode) {

Symbol symbol = (Symbol) middleCode[0];

m\_trackMap.Remove(symbol);

}

### Address

The Address method adds assembly code instructions for the address operator. It is actually quite simple, since we only gave to call LoadAddressToRegister to load the address of the symbol into a register.

public void Address(MiddleCode middleCode) {

Symbol resultSymbol = (Symbol) middleCode[0],

addressSymbol = (Symbol) middleCode[1];

Track track = LoadAddressToRegister(addressSymbol);

m\_trackMap.Add(resultSymbol, track);

m\_trackMap.Remove(addressSymbol);

}

The IntegralDereference method adds assembly code instructions for derefereeing a pointer to an integral value. We only have to call LoadValueToRegister for the address symbol.

public void IntegralDereference(MiddleCode middleCode) {

Symbol resultSymbol = (Symbol) middleCode[0];

Assert.ErrorXXX(resultSymbol.AddressSymbol != null);

Track addressTrack = LoadValueToRegister(resultSymbol.AddressSymbol);

m\_trackMap.Add(resultSymbol.AddressSymbol, addressTrack);

}

The FloatingDereference method adds assembly code instructions for derefereing a pointer to a floating value.

public void FloatingDereference(MiddleCode middleCode) {

Symbol resultSymbol = (Symbol) middleCode[0];

Assert.ErrorXXX(resultSymbol.AddressSymbol != null);

Track addressTrack = LoadValueToRegister(resultSymbol.AddressSymbol);

m\_trackMap.Add(resultSymbol.AddressSymbol, addressTrack);

}

### Floating Binary

Similar to the integral operations, we need the m\_middleToFloatingMap map to translate the middle code operators to their equivalent assembly code operators.

public static IDictionary<MiddleOperator, AssemblyOperator>

m\_middleToFloatingMap =

new Dictionary<MiddleOperator, AssemblyOperator>() {

{MiddleOperator.UnarySubtract, AssemblyOperator.fchs},

{MiddleOperator.BinaryAdd, AssemblyOperator.fadd},

{MiddleOperator.BinarySubtract, AssemblyOperator.fsub},

{MiddleOperator.SignedMultiply, AssemblyOperator.fmul},

{MiddleOperator.SignedDivide, AssemblyOperator.fdiv},

{MiddleOperator.Equal, AssemblyOperator.je},

{MiddleOperator.NotEqual, AssemblyOperator.jne},

{MiddleOperator.SignedLessThan, AssemblyOperator.ja},

{MiddleOperator.SignedLessThanEqual, AssemblyOperator.jae},

{MiddleOperator.SignedGreaterThan, AssemblyOperator.jb},

{MiddleOperator.SignedGreaterThanEqual, AssemblyOperator.jbe}

};

The methods for the floating-point operations are in fact mush simplier than the corresponding integral methods. The FloatingUnary method just adds the instruction to the assembly code. The operand of the operator has already been pushed to the floating-point stack, the operator pops that value, and pushes the result to the stack.

public void FloatingUnary(MiddleCode middleCode) {

AddAssemblyCode(m\_middleToFloatingMap[middleCode.Operator]);

}

The FloatingBinary works in the same way, the operator pops the two operands from the floating-point stack and pushes the result on the stack.

public void FloatingBinary(MiddleCode middleCode) {

Assert.ErrorXXX((--m\_floatStackSize) >= 0);

AddAssemblyCode(m\_middleToFloatingMap[middleCode.Operator]);

}

In the FloatingParameter we do not need to actual parameter symbol, since its value has already been pushed on the floating-point stack. We only need to type and offset of the parameter when we pop it on the floating-point stack by calling TopPopSymbol.

public void FloatingParameter(MiddleCode middleCode) {

Type paramType = (Type) middleCode[0];

Symbol paramSymbol = new Symbol(paramType);

int paramOffset = (int) middleCode[2];

paramSymbol.Offset = m\_totalExtraSize + paramOffset;

TopPopSymbol(paramSymbol, TopOrPop.Pop);

}

### Floating Relation

The FloatingRelation method generate code for relation operators for floating types. The fcompp (float compare pop) assembly code instruction performs a comparation on the two values on top of the floating value stack, stores the result in the internal float status word, and pops the values from the stack. The fstsw (float store status word) assembly code instruction stores the floating status word in the ax register. The sahf (store ah into flags) stores the value of register ah (which hold the higher byte of register ax) into the integral flags. Finally, we add a jump instruction that matches the original middle code instruction and is looked up in the middle to floating map.

public void FloatingRelation(MiddleCode middleCode) {

Assert.ErrorXXX((m\_floatStackSize -= 2) >= 0);

int target = (int) middleCode[0];

AddAssemblyCode(AssemblyOperator.fcompp);

AddAssemblyCode(AssemblyOperator.fstsw, Register.ax);

AddAssemblyCode(AssemblyOperator.sahf);

AssemblyOperator objectOperator =

m\_middleToFloatingMap[middleCode.Operator];

AddAssemblyCode(objectOperator, null, null, target);

}

### Floating Push and Pop

When pushing floating values to the floating-point stack, we need the m\_floatPushMap map that maps the type sizes of the symbols to assembly code operators.

public static IDictionary<Sort, AssemblyOperator> m\_floatPushMap =

new Dictionary<Sort, AssemblyOperator>() {

{Sort.Signed\_Int, AssemblyOperator.fild\_word},

{Sort.Unsigned\_Int, AssemblyOperator.fild\_word},

{Sort.Signed\_Long\_Int, AssemblyOperator.fild\_dword},

{Sort.Unsigned\_Long\_Int, AssemblyOperator.fild\_dword},

{Sort.Float, AssemblyOperator.fld\_dword},

{Sort.Double, AssemblyOperator.fld\_qword},

{Sort.Long\_Double, AssemblyOperator.fld\_qword}

};

The PushSymbol method generates code that pushes the value of a symbol to the floating-point stack.

public void PushSymbol(Symbol symbol) {

Assert.ErrorXXX((++m\_floatStackSize) <= FloatingStackMaxSize);

AssemblyOperator objectOperator = m\_floatPushMap[symbol.Type.Sort];

Track track;

There are two special assembly code instructions for pushing the value zero or one to the floating-point stack: fldz (float load zero) fld1 (float load one). If the value is (integral or floating) zero, we use fldz to load the value zero, and if the value is (integral or floating) one we use fld1 to load the value zero.

if (((symbol.Value is BigInteger) &&

(((BigInteger) symbol.Value).IsZero)) ||

((symbol.Value is decimal) && (((decimal) symbol.Value) == 0))) {

AddAssemblyCode(AssemblyOperator.fldz);

}

else if (((symbol.Value is BigInteger) &&

(((BigInteger) symbol.Value).IsOne)) ||

((symbol.Value is decimal) &&

(((decimal) symbol.Value) == 1))) {

AddAssemblyCode(AssemblyOperator.fld1);

}

For any other integral or floating value, we cannot load it directly. Instead, we need store the value on a memory address and load the value from the address. We add the symbol to the static set for the linker to be able to find the value, and we add the push operation with the name of symbol, which will be changed to a proper address by the linker.

else if ((symbol.Value is BigInteger) || (symbol.Value is decimal)) {

SymbolTable.StaticSet.Add(ConstantExpression.Value(symbol));

AddAssemblyCode(objectOperator, symbol.UniqueName, 0);

}

If the symbol does not hold an integral or a floating value, we check if the value is stored in a register. In that case we cannot store the register directly, but like the value case above we need to store the value at the memory address. We use the special integral storage to load the value.

else if (m\_trackMap.TryGetValue(symbol, out track)) {

m\_trackMap.Remove(symbol);

string containerName = AddStaticContainer(symbol.Type);

AddAssemblyCode(AssemblyOperator.mov, containerName, 0, track);

AddAssemblyCode(objectOperator, containerName, 0);

}

If the symbol is an array, function, or string, we also use the integral storage name to store the value. However, in this case will shall load the address of the symbol, rather than its value. We start by loading the base address of the symbol to the integral storage. Then we add its offset, unless it is zero.

else if (symbol.Type.IsArrayFunctionOrString()) {

string containerName = AddStaticContainer(symbol.Type);

AddAssemblyCode(AssemblyOperator.mov, containerName, 0,

Base(symbol), TypeSize.PointerSize);

int offset = Offset(symbol);

if (offset != 0) {

AddAssemblyCode(AssemblyOperator.add, containerName, 0,

(BigInteger) offset, TypeSize.PointerSize);

}

AddAssemblyCode(objectOperator, containerName, 0);

}

If none of the above apply, we perform the operation on the address of the symbol.

else {

AddAssemblyCode(objectOperator, Base(symbol), Offset(symbol));

}

}

private static string AddStaticContainer(Type type) {

string containerName = "container" + type.Size() +

"bytes" + Symbol.NumberId;

SymbolTable.StaticSet.Add(ConstantExpression.Value(containerName,

type, null));

return containerName;

}

When popping or topping a value from the floating-point stack, we also need the m\_floatPopMap and m\_floatTopMap maps that maps from the type sizes of the symbols to assembly code instructions.

public static IDictionary<Sort, AssemblyOperator> m\_floatPopMap =

new Dictionary<Sort, AssemblyOperator>() {

{Sort.Signed\_Int, AssemblyOperator.fistp\_word},

{Sort.Unsigned\_Int, AssemblyOperator.fistp\_word},

{Sort.Pointer, AssemblyOperator.fistp\_word},

{Sort.Signed\_Long\_Int, AssemblyOperator.fistp\_dword},

{Sort.Unsigned\_Long\_Int, AssemblyOperator.fistp\_dword},

{Sort.Float, AssemblyOperator.fstp\_dword},

{Sort.Double, AssemblyOperator.fstp\_qword},

{Sort.Long\_Double, AssemblyOperator.fstp\_qword}

};

public static IDictionary<Sort, AssemblyOperator> m\_floatTopMap =

new Dictionary<Sort, AssemblyOperator>() {

{Sort.Signed\_Int, AssemblyOperator.fist\_word},

{Sort.Unsigned\_Int, AssemblyOperator.fist\_word},

{Sort.Pointer, AssemblyOperator.fist\_word},

{Sort.Signed\_Long\_Int, AssemblyOperator.fist\_dword},

{Sort.Unsigned\_Long\_Int, AssemblyOperator.fist\_dword},

{Sort.Float, AssemblyOperator.fst\_dword},

{Sort.Double, AssemblyOperator.fst\_qword},

{Sort.Long\_Double, AssemblyOperator.fst\_qword}

};

public enum TopOrPop {Top, Pop};

The PopEmpty method pops the floating-point stack without storing the value anywhere. However, there is no matching assembly code instruction. The instructions always store the value somewhere. Therefore, we use the integral storage, simply because we have to use some storage.

public void PopEmpty() {

string containerName = AddStaticContainer(Type.LongDoubleType);

AddAssemblyCode(AssemblyOperator.fistp\_word, containerName, 0);

}

The TopPopSymbol method tops or pops the top-most value of the floating-point stack.

public void TopPopSymbol(Symbol symbol, TopOrPop topOrPop) {

Assert.ErrorXXX(symbol != null);

AssemblyOperator objectOperator;

Depending on the value opf the topOrPop parameter, we extract the assembly code instruction from the m\_floatPopMap or m\_floatTopMap map.

if (topOrPop == TopOrPop.Pop) {

objectOperator = m\_floatPopMap[symbol.Type.Sort];

Assert.ErrorXXX((--m\_floatStackSize) >= 0);

}

else {

objectOperator = m\_floatTopMap[symbol.Type.Sort];

}

If the symbol is a temporary variable, we store its value in a register. However, we cannot store the value directly in a register. Therefore, we pop or top the value to the integral storage, and loads the value from there into the register.

if (symbol.IsTemporary() && (symbol.AddressSymbol == null) &&

(symbol.Offset == 0)) {

string containerName = AddStaticContainer(symbol.Type);

AddAssemblyCode(objectOperator, containerName, 0);

Track track = new Track(symbol);

AddAssemblyCode(AssemblyOperator.mov, track, containerName, 0);

m\_trackMap.Add(symbol, track);

}

If none of the above cases apply, we pop or top the value to the address of the symbol.

else {

AddAssemblyCode(objectOperator, Base(symbol), Offset(symbol));

}

}

### Type Conversion

The IntegralToIntegral method adds assembly code instructions that converts an integral value from one size to another size.

public void IntegralToIntegral(MiddleCode middleCode, int index) {

Symbol targetSymbol = (Symbol) middleCode[0],

sourceSymbol = (Symbol) middleCode[1];

The sizeArray method returns pointer size for arrays, functions, and strings.

Type targetType = targetSymbol.Type, sourceType = sourceSymbol.Type;

int targetSize = targetType.SizeArray(),

sourceSize = sourceType.SizeArray();

We need target add a set-track-size instruction for

Track sourceTrack = LoadValueToRegister(sourceSymbol);

AddAssemblyCode(AssemblyOperatargetr.set\_track\_size,

sourceTrack, targetSize);

If the source size is smaller than the target size, we add code to mask the upper bits; that is, we set to zero the bits of the target value that do not fit in the source value.

if (sourceSize != targetSize) {

if (sourceSize < targetSize) {

If the target size is eight bytes, we have a special case. We need to load the mask to a register, and use the and instruction to do the actual masking.

if (targetSize == 8) {

Track targetTrack = new Track(targetSymbol);

AddAssemblyCode(AssemblyOperatargetr.mov, targetTrack,

TypeSize.GetMask(sourceType.Sort));

AddAssemblyCode(AssemblyOperatargetr.and,

sourceTrack, targetTrack);

}

If the target size is not eight bytes, we can use the and instruction directly.

else {

AddAssemblyCode(AssemblyOperatargetr.and, sourceTrack,

TypeSize.GetMask(sourceType.Sort));

}

}

If both the source type and target type is signed, we need to preserve the signed status throught the conversion process. If the value is negative, the top-most bit is set to one. Since the source size does not equals the target size, the top-most bits differ between the values. Therefore, if the value is negative, we need to change it to a positive value with the source size and change it back to a negative value with the target size. In this way, the signed status will be preserved. However, if the value is non-negative (zero or positive), we do nothing. We just add to the track map that the values holds a new size.

if (sourceType.IsSigned() && targetType.IsSigned()) {

AddAssemblyCode(AssemblyOperatargetr.set\_track\_size,

sourceTrack, sourceSize);

AddAssemblyCode(AssemblyOperatargetr.cmp, sourceTrack,

BigInteger.Zero);

AddAssemblyCode(AssemblyOperatargetr.jge, null, null, index + 1);

AddAssemblyCode(AssemblyOperatargetr.neg, sourceTrack);

AddAssemblyCode(AssemblyOperatargetr.set\_track\_size,

sourceTrack, targetSize);

AddAssemblyCode(AssemblyOperatargetr.neg, sourceTrack);

}

}

If the target type and source type have the same size, we only add the target symbol with the source track in the track map.

m\_trackMap.Add(targetSymbol, sourceTrack);

}

The IntegralToFloating and FloatingToIntegral are quite simple. In the integral-to-floating-case we push the integral value at the floating-point stack, and in the floating-to-integral case, we pop the value from the stack. Note that there is no floating-to-floating case. In that case, we just keep the value on the stack, the type conversion occurs when the value is pushed, popped, or topped.

public void IntegralToFloating(MiddleCode middleCode) {

Symbol fromSymbol = (Symbol) middleCode[1];

PushSymbol(fromSymbol);

}

public void FloatingToIntegral(MiddleCode middleCode) {

Symbol toSymbol = (Symbol) middleCode[0];

TopPopSymbol(toSymbol, TopOrPop.Pop);

}

### Struct and Union

This section describes the handling of struct and unions, assignment, parameter, and function return value. Note that these methods make no difference between struct and unions, in all cases the task is to move the data of a memory block between two addresses.

The StructUnionAssign method copies the data from the address of the target symbol to address of the source symbol by calling MemoryCopy.

public void StructUnionAssign(MiddleCode middleCode, int index) {

Symbol targetSymbol = (Symbol) middleCode[0],

sourceSymbol = (Symbol) middleCode[1];

Track targetAddressTrack = LoadAddressToRegister(targetSymbol),

sourceAddressTrack = LoadAddressToRegister(sourceSymbol);

MemoryCopy(targetAddressTrack, sourceAddressTrack,

targetSymbol.Type.Size(), index);

}

The StructUnionParameter copies the data from the address of the parameter symbol to the address on the activation record by calling MemoryCopy.

public void StructUnionParameter(MiddleCode middleCode, int index) {

Symbol sourceSymbol = (Symbol) middleCode[1];

Symbol targetSymbol = new Symbol(Type.PointerTypeX);

int paramOffset = (int) middleCode[2];

targetSymbol.Offset = m\_totalExtraSize + paramOffset;

Track sourceAddressTrack = LoadAddressToRegister(sourceSymbol);

Track targetAddressTrack = LoadAddressToRegister(targetSymbol);

MemoryCopy(targetAddressTrack, sourceAddressTrack,

sourceSymbol.Type.Size(), index);

}

The StructUnionGetReturnValue loads the address symbol of the symbol into the return pointer register, which is then added to the track map. In this way, is the value of the address symbol of the struct or union stored in the track, and will be used in later operations.

public void StructUnionGetReturnValue(MiddleCode middleCode) {

Symbol targetSymbol = (Symbol) middleCode[0];

CheckRegister(targetSymbol, AssemblyCode.ReturnPointerRegister);

Track targetAddressTrack =

new Track(targetSymbol.AddressSymbol,

AssemblyCode.ReturnPointerRegister);

m\_trackMap.Add(targetSymbol.AddressSymbol, targetAddressTrack);

}

The StructUnionSetReturnValue method

public void StructUnionSetReturnValue(MiddleCode middleCode) {

Symbol returnSymbol = (Symbol) middleCode[1];

LoadAddressToRegister(returnSymbol, AssemblyCode.ReturnPointerRegister);

}

The m\_labelCount field is used by MemoryCopy to generate a new label for each copy process.

private static int m\_labelCount = 0;

The MemoryCopy method adds code to copy the data of a memory block between two memory addresses.

private void MemoryCopy(Track targetAddressTrack,

Track sourceAddressTrack, int size, int index) {

We create tracks for the register to count of the data block and the value being copied. If the size of the data block is less than 256, we use the character types (1 byte) for the count register, otherwise we use the integer type (more than one byte) XXX.

Type countType = (size < 256) ? Type.UnsignedCharType

: Type.UnsignedIntegerType;

Track countTrack = new Track(countType),

valueTrack = new Track(Type.UnsignedCharType);

We start by loading the size of the data block into the count register, the loop continues until it reaches zero.

AddAssemblyCode(AssemblyOperator.mov, countTrack, (BigInteger) size);

Then the loop begins, we use a label that is unique in the assembly file with the m\_labelCount field.

int labelIndex = m\_labelCount++;

string labelText = AssemblyCode.MakeMemoryLabel(labelIndex);

AddAssemblyCode(AssemblyOperator.label, labelText);

We move a value from the source address to the target address.

AddAssemblyCode(AssemblyOperator.mov, valueTrack,

sourceAddressTrack, 0);

AddAssemblyCode(AssemblyOperator.mov, targetAddressTrack,

0, valueTrack);

We increment the source and target address.

AddAssemblyCode(AssemblyOperator.inc, sourceAddressTrack);

AddAssemblyCode(AssemblyOperator.inc, targetAddressTrack);

We decrement the count register.

AddAssemblyCode(AssemblyOperator.dec, countTrack);

If the count register does not euqla zero, we jump to the label. If it does equal zero, the copy process is done.

AddAssemblyCode(AssemblyOperator.cmp, countTrack, BigInteger.Zero);

AddAssemblyCode(AssemblyOperator.jne, null, labelIndex);

}

### Initialization Code

The initialization of data blocks differs between the Linux and Windows environment.

public static void InitializerializationCodeList() {

List<AssemblyCode> assemblyCodeList = new List<AssemblyCode>();

AddAssemblyCode(assemblyCodeList, AssemblyOperator.comment,

"Initialize Stack Pointer");

AddAssemblyCode(assemblyCodeList, AssemblyOperator.mov,

AssemblyCode.FrameRegister, LinkerWindows.StackTopName);

if (Start.Windows) {

AddAssemblyCode(assemblyCodeList, AssemblyOperator.comment,

"Initialize Heap Pointer");

AddAssemblyCode(assemblyCodeList, AssemblyOperator.mov\_word,

null, 65534, (BigInteger) 65534);

}

if (Start.Linux) {

AddAssemblyCode(assemblyCodeList, AssemblyOperator.comment,

"Initialize Heap Pointer");

AddAssemblyCode(assemblyCodeList, AssemblyOperator.mov\_dword,

LinkerWindows.StackTopName, 65534,

LinkerWindows.StackTopName);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.add\_dword,

LinkerWindows.StackTopName, 65534,

(BigInteger) 65534);

}

AddAssemblyCode(assemblyCodeList, AssemblyOperator.comment,

"Initialize FPU Control Word, truncate mode " +

"=> set bit 10 and 11.");

AddAssemblyCode(assemblyCodeList, AssemblyOperator.fstcw,

AssemblyCode.FrameRegister, 0);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.or\_word,

AssemblyCode.FrameRegister, 0, (BigInteger) 0x0C00);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.fldcw,

AssemblyCode.FrameRegister, 0);

if (Start.Linux) {

List<string> textList = new List<string>();

ISet<string> externSet = new HashSet<string>();

AssemblyCodeGenerator.LinuxTextList(assemblyCodeList, textList, externSet);

SymbolTable.StaticSet.Add(new StaticSymbolLinux

(StaticSymbolLinux.TextOrData.Text,

AssemblyCodeGenerator.InitializerName, textList, externSet));

}

if (Start.Windows) {

List<byte> byteList = new List<byte>();

IDictionary<int, string> accessMap = new Dictionary<int, string>();

IDictionary<int, string> callMap = new Dictionary<int, string>();

ISet<int> returnSet = new HashSet<int>();

AssemblyCodeGenerator.GenerateTargetWindows(assemblyCodeList,

byteList, accessMap, callMap, returnSet);

StaticSymbol staticSymbol =

new StaticSymbolWindows(AssemblyCodeGenerator.InitializerName,

byteList, accessMap, callMap, returnSet);

SymbolTable.StaticSet.Add(staticSymbol);

}

}

### Command Line Arguments

The code for obtaining the command line arguments differs between the Linux and Windows environment.

public static void ArgumentCodeList() {

List<AssemblyCode> assemblyCodeList = new List<AssemblyCode>();

When to execution starts, the stack is loaded with values corresponding to the command line arguments.

if (Start.Linux) {

AddAssemblyCode(assemblyCodeList, AssemblyOperator.comment,

"Initialize Command Line Arguments");

AddAssemblyCode(assemblyCodeList, AssemblyOperator.pop, Register.rbx);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.mov,

Register.rax, Register.rbx);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.mov,

Register.rdx, Register.rbp);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.label, "$args$loop");

AddAssemblyCode(assemblyCodeList, AssemblyOperator.cmp,

Register.rbx, BigInteger.Zero);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.je,

null, null, "$args$exit");

AddAssemblyCode(assemblyCodeList, AssemblyOperator.pop, Register.rsi);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.mov,

Register.rbp, 0, Register.rsi);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.add, Register.rbp,

(BigInteger) TypeSize.PointerSize);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.dec, Register.rbx);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.jmp,

null, null, "$args$loop");

AddAssemblyCode(assemblyCodeList, AssemblyOperator.label,

"$args$exit");

AddAssemblyCode(assemblyCodeList, AssemblyOperator.mov\_qword,

Register.rbp, 0, BigInteger.Zero);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.add, Register.rbp,

(BigInteger) TypeSize.PointerSize);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.mov, Register.rbp,

SymbolTable.FunctionHeaderSize, Register.eax);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.mov, Register.rbp,

SymbolTable.FunctionHeaderSize +

TypeSize.SignedIntegerSize, Register.rdx);

List<string> textList = new List<string>();

ISet<string> externSet = new HashSet<string>();

AssemblyCodeGenerator.LinuxTextList(assemblyCodeList, textList,

externSet);

SymbolTable.StaticSet.

Add(new StaticSymbolLinux(StaticSymbolLinux.TextOrData.Text,

AssemblyCodeGenerator.ArgsName, textList, externSet));

}

else if (Start.Windows) {

AddAssemblyCode(assemblyCodeList, AssemblyOperator.mov,

Register.si, Register.bp);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.mov\_word,

Register.bp, 0, AssemblyCodeGenerator.PathName);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.add,

Register.bp, (BigInteger) 2);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.mov,

Register.ax, BigInteger.One);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.mov,

Register.bx, (BigInteger) 129);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.cmp\_byte,

Register.bx, 0, (BigInteger) 32);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.je,

null, assemblyCodeList.Count + 5);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.cmp\_byte,

Register.bx, 0, (BigInteger) 13);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.je,

null, assemblyCodeList.Count + 13);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.inc, Register.bx);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.jmp,

null, assemblyCodeList.Count - 5);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.cmp,

Register.ax, BigInteger.One);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.je,

null, assemblyCodeList.Count + 2);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.mov\_byte,

Register.bx, 0, BigInteger.Zero);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.inc, Register.bx);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.cmp\_byte,

Register.bx, 0, (BigInteger) 32);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.je,

null, assemblyCodeList.Count - 2);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.mov,

Register.bp, 0, Register.bx);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.add,

Register.bp, (BigInteger) 2);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.inc, Register.ax);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.jmp,

null, assemblyCodeList.Count - 15);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.mov\_byte,

Register.bx, 0, BigInteger.Zero);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.mov\_word,

Register.bp, 0, BigInteger.Zero);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.add,

Register.bp, (BigInteger) 2);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.mov,

Register.bp, 6, Register.ax);

AddAssemblyCode(assemblyCodeList, AssemblyOperator.mov,

Register.bp, 8, Register.si);

List<byte> byteList = new List<byte>();

IDictionary<int, string> accessMap = new Dictionary<int, string>();

IDictionary<int, string> callMap = new Dictionary<int, string>();

ISet<int> returnSet = new HashSet<int>();

AssemblyCodeGenerator.

GenerateTargetWindows(assemblyCodeList, byteList,

accessMap, callMap, returnSet);

StaticSymbol staticSymbol =

new StaticSymbolWindows(AssemblyCodeGenerator.ArgsName, byteList,

accessMap, callMap, returnSet);

SymbolTable.StaticSet.Add(staticSymbol);

}

}

### Linux Text List

When the assembly code has been generated, it needs to be written as a text file, which is the task of LinuxTextList. It iterates through the assembly code list and writes the assembly code instruction at one line each by calling ToString in AssemblyCode. It also adds names add the externSet parameter for each name in the instructions.

public static void LinuxTextList(IList<AssemblyCode> assemblyCodeList,

IList<string> textList,

ISet<string> externSet) {

foreach (AssemblyCode assemblyCode in assemblyCodeList) {

AssemblyOperator assemblyOperator = assemblyCode.Operator;

object operand0 = assemblyCode[0],

operand1 = assemblyCode[1],

operand2 = assemblyCode[2];

if (assemblyOperator == AssemblyOperator.define\_value) {

Sort sort = (Sort) operand0;

if ((sort != Sort.String) && (operand1 is string)) {

string name1 = (string) operand1;

if (!name1.Contains(Symbol.SeparatorId)) {

externSet.Add(name1);

}

}

}

else if ((assemblyOperator != AssemblyOperator.label) &&

(assemblyOperator != AssemblyOperator.comment)) {

if (operand0 is string) {

string name0 = (string) operand0;

if (!name0.Contains(Symbol.SeparatorId)) {

externSet.Add(name0);

}

}

if (operand1 is string) {

string name1 = (string) operand1;

if (!name1.Contains(Symbol.SeparatorId)) {

externSet.Add(name1);

}

}

if (operand2 is string) {

string name2 = (string) operand2;

if (!name2.Contains(Symbol.SeparatorId)) {

externSet.Add(name2);

}

}

}

string text = assemblyCode.ToString();

if (text != null) {

textList.Add(text);

}

}

}

### Windows Jump Info

The WindowsJumpInfo method changes the jump addresses from middle code line numbers to assembly code line numbers, and then to relative byte size addresses. To begin with we need the middleToAssemblyMap map to keep track of the assembly code addresses.

private void WindowsJumpInfo() {

IDictionary<int,int> middleToAssemblyMap = new Dictionary<int,int>();

We iterate through the assembly code list and add the middle code line number the assembly code line number each time we encounter a new-middle-code instruction. We really need just the line numbers of first of the assembly line that corresponds to a middle code instruction, since all jumps are to those addresses.

for (int assemblyIndex = 0;

assemblyIndex < m\_assemblyCodeList.Count; ++assemblyIndex) {

AssemblyCode assemblyCode = m\_assemblyCodeList[assemblyIndex];

if (assemblyCode.Operator == AssemblyOperator.new\_middle\_code) {

int middleIndex = (int) assemblyCode[0];

middleToAssemblyMap.Add(middleIndex, assemblyIndex);

assemblyCode.Operator = AssemblyOperator.empty;

}

}

Then we iterate through the assembly code list and change all jump targets from middle code line numbers to assembly code line numbers. Note that the middle code line number was stored in the third position (index 2) of the instruction, and that we set the assembly line number to the second position (index 1).

for (int line = 0; line < m\_assemblyCodeList.Count; ++line) {

AssemblyCode assemblyCode = m\_assemblyCodeList[line];

if (assemblyCode.IsRelationNotRegister() ||

assemblyCode.IsJumpNotRegister()) {

if (!(assemblyCode[1] is int)) {

int middleTarget = (int) assemblyCode[2];

assemblyCode[1] = middleToAssemblyMap[middleTarget];

}

}

}

When we have changed the jump targets from middle code line numbers to assembly code line numbers, the second step is to change the targets into the number of bytes between the next assembly instruction (the instruction following the current instruction) and the target instruction, which is a negative value in case of backward jumps.

First, we need the assemblyToByteMap map to keep track of the address of each assembly code instruction, counted in bytes from the beginning of the code. Then we iterate through the assembly code list and add the byte addresses each instruction. At the moment, all jump instruction are equally large (in bytes), but that may change later on.

IDictionary<int,int> assemblyToByteMap = new Dictionary<int,int>();

{ int byteSize = 0, line = 0;

foreach (AssemblyCode assemblyCode in m\_assemblyCodeList) {

assemblyToByteMap.Add(line++, byteSize);

if (!(assemblyCode.IsRelationNotRegister() ||

assemblyCode.IsJumpNotRegister())) {

byteSize += assemblyCode.ByteList().Count;

}

}

assemblyToByteMap.Add(m\_assemblyCodeList.Count, byteSize);

}

When we change the assembly code targets to relative byte targets, we may have to iterate several times because short jumps (between -128 and 127, inclusive) instructions are smaller than long jumps instructions. At the beginning, all jumps are long jumps, but some of them may be changed to short jump during the iteration.

while (true) {

for (int line = 0; line < (m\_assemblyCodeList.Count - 1); ++line) {

AssemblyCode thisCode = m\_assemblyCodeList[line],

nextCode = m\_assemblyCodeList[line + 1];

if (thisCode.IsRelationNotRegister() ||

thisCode.IsJumpNotRegister()) {

int assemblyTarget = (int) thisCode[1];

The byteSource variable is the byte address of the target assembly instruction following the current instruction, byteTarget is the byte address of the target instruction, and byteDistance is difference between byteTarget and byteSource, which becomes the final target address.

int byteSource = assemblyToByteMap[line + 1],

byteTarget = assemblyToByteMap[assemblyTarget];

int byteDistance = byteTarget - byteSource;

Assert.ErrorXXX(byteDistance != 0);

thisCode[0] = byteDistance;

}

}

When the jump instruction has been given proper target address, we iterate through the assembly list and update the assemblyToByteMap. Some byte address may have been changed because some long jump instructions have been changed to short jump instructions. If any address has been changed, update is set to true and the overall loop will perform another iteration.

bool update = false;

{ int byteSize = 0, line = 0;

foreach (AssemblyCode objectCode in m\_assemblyCodeList) {

if (assemblyToByteMap[line] != byteSize) {

assemblyToByteMap[line] = byteSize;

update = true;

}

byteSize += objectCode.ByteList().Count;

++line;

}

}

If nothing has been changed from the previous iteration, every assembly jump instruction holds the correct target address and we break the loop.

if (!update) {

break;

}

}

Finally, we need to iterate through the assembly code list, and change the return addresses in the same way as the byte target addresses above.

for (int line = 0; line < m\_assemblyCodeList.Count; ++line) {

AssemblyCode assemblyCode = m\_assemblyCodeList[line];

if (assemblyCode.Operator == AssemblyOperator.return\_address) {

int middleAddress = (int) ((BigInteger) assemblyCode[2]);

int assemblyAddress = middleToAssemblyMap[middleAddress];

int byteAddress = assemblyToByteMap[assemblyAddress];

assemblyCode[2] = (BigInteger) byteAddress;

}

}

}

### Windows Byte List

The WindowsByteList method iterates through the assembly code list and generates the byte list for each instruction by calling ByteList in AssemblyCode. It also adds names to the accessMap, callMap, and returnSet parameters.

private void WindowsByteList(List<byte> byteList,

IDictionary<int,string> accessMap,

IDictionary<int,string> callMap,

ISet<int> returnSet) {

foreach (AssemblyCode assemblyCode in m\_assemblyCodeList) {

We add the byte list of the instruction to the byte list parameter.

byteList.AddRange(assemblyCode.ByteList());

Unless the instruction is a label, comment, or zero sequence, we check whether there are any names in the code that shall be added to the access map or the call map, or the return set.

if ((assemblyCode.Operator != AssemblyOperator.label) &&

(assemblyCode.Operator != AssemblyOperator.comment) &&

(assemblyCode.Operator != AssemblyOperator.define\_zero\_sequence)){

If the instruction is an address definition, we add its name to the access map.

if (assemblyCode.Operator == AssemblyOperator.define\_address) {

string name = (string) assemblyCode[0];

accessMap.Add(byteList.Count - TypeSize.PointerSize, name);

}

If the instruction is a value definition, we need to look into its sort. If it is a pointer or a static address, we add its name.

else if (assemblyCode.Operator == AssemblyOperator.define\_value) {

Sort sort = (Sort) assemblyCode[0];

object value = assemblyCode[1];

if (sort == Sort.Pointer) {

if (value is string) {

accessMap.Add(byteList.Count - TypeSize.PointerSize,

(string) value);

}

else if (value is StaticAddress) {

StaticAddress staticAddress = (StaticAddress) value;

accessMap.Add(byteList.Count - TypeSize.PointerSize,

staticAddress.UniqueName);

}

}

}

In case of a function call, add the name of the function to be called to the call map.

else if ((assemblyCode.Operator == AssemblyOperator.call) &&

(assemblyCode[0] is string)) {

string calleeName = (string) assemblyCode[0];

int address = byteList.Count - TypeSize.PointerSize;

callMap.Add(address, calleeName);

}

In case of a function return, we add the address the return set.

else if (assemblyCode.Operator == AssemblyOperator.return\_address) {

int address = byteList.Count - TypeSize.PointerSize;

returnSet.Add(address);

}

If the first operand is a string, we have two cases: if the third operand is a value of BigInteger, we need decrease the address with its size and with the size of a pointer. Otherwise, we just decrease the address with the pointer size.

else if (assemblyCode[0] is string) { // Add [g + 1], 2

if (assemblyCode[2] is BigInteger) {

int size = AssemblyCode.SizeOfValue((BigInteger)assemblyCode[2],

assemblyCode.Operator);

int address = byteList.Count - TypeSize.PointerSize - size;

accessMap.Add(address, (string) assemblyCode[0]);

}

else {

int address = byteList.Count - TypeSize.PointerSize;

accessMap.Add(address, (string) assemblyCode[0]);

}

}

If the second or third operand is a string, we add its name the the address, which is the size of the instruction minus the pointer size.

else if (assemblyCode[1] is string) { // mov ax, [g + 1]; mov ax, g

int address = byteList.Count - TypeSize.PointerSize;

accessMap.Add(address, (string) assemblyCode[1]);

}

else if (assemblyCode[2] is string) { // Add [bp + 2], g

int address = byteList.Count - TypeSize.PointerSize;

accessMap.Add(address, (string) assemblyCode[2]);

}

}

}

}

}

}

## Generation

When the middle code has been optimized, we need to generate base blocks and live sets.

public void generate() {

{ generateDefAndUseLineSets();

Map<Integer,BaseBlock> leaderMap = new ListMap<>();

Set<BaseBlock> baseBlockSet = generateBaseBlockSet(leaderMap);

redundantExpressions(baseBlockSet);

removeClearedCode();

}

BaseBlock.clearCount();

Map<Integer,BaseBlock> leaderMap = new ListMap<>();

Set<BaseBlock> baseBlockSet = generateBaseBlockSet(leaderMap);

Set<Pair<BaseBlock,BaseBlock>> edgeSet =

generateEdgeSet(baseBlockSet, leaderMap);

m\_controlFlowGraph = new DirectedGraph<>(baseBlockSet, edgeSet);

generateLiveMaps();

generatePathSets();

if (Main.SmallWarning) {

generateDefAndUseLineSets();

generateDefAndUseBlockSets();

generatePathSets();

generateOutPathSets();

}

for (BaseBlock baseBlock : m\_controlFlowGraph.getVertexSet()) {

int firstIndex = baseBlock.getFirst(), lastIndex = baseBlock.getLast();

new RegisterPreparator(m\_middleCodeList, firstIndex, lastIndex);

for (int index = firstIndex; index <= lastIndex; ++index) {

Main.DebugStream.println(index + ": " + m\_middleCodeList.get(index));

}

}

}

### Base Blocks Generation

A base block is a maximal sequence of middle code where only the first instruction is a jump target. The idea is that we shall be able to trust that the registers hold the expected values thorugh the base block. In the first step, we look into all conditional and unconditional jump instructions and mark their targets as leaders, which will form the beginning of a base block. The first instruction of a function is always a leader and all instructions following a function all an interrupts system call is also marked as leaders, since a function or system call may affect the registers of the object code.

When all leaders have been established in the first step, we generate a list of pairs holding the first and last line number of each base block.

private Set<BaseBlock> generateBaseBlockSet

(Map<Integer,BaseBlock> leaderMap) {

m\_middleCodeList.get(0).setLeader(); // Function Start

for (int index = 0; index < (m\_middleCodeList.size() - 1); ++index) {

MiddleCode middleCode = m\_middleCodeList.get(index);

if (Main.CheckStackHeap &&

(middleCode.getOperator() == MiddleOperator.FunctionStart)) {

m\_middleCodeList.get(index + 1).setLeader();

}

if (middleCode.isRelationOrGoto()) {

int target = (Integer) middleCode.getOperand(1);

m\_middleCodeList.get(target).setLeader();

m\_middleCodeList.get(index + 1).setLeader();

}

else if ((middleCode.getOperator() == MiddleOperator.Call) ||

(middleCode.getOperator() == MiddleOperator.Interrupt)) {

m\_middleCodeList.get(index + 1).setLeader();

}

}

int startIndex = 0;

Set<BaseBlock> baseBlockSet = new ListSet<>();

for (int index = 1; index < m\_middleCodeList.size(); ++index) {

if (m\_middleCodeList.get(index).isLeader()) {

BaseBlock baseBlock = new BaseBlock(startIndex, index - 1);

leaderMap.put(startIndex, baseBlock);

baseBlockSet.add(baseBlock);

startIndex = index;

}

}

{ BaseBlock baseBlock =

new BaseBlock(startIndex, m\_middleCodeList.size() - 1);

leaderMap.put(startIndex, baseBlock);

baseBlockSet.add(baseBlock);

}

return baseBlockSet;

}

private Set<Pair<BaseBlock,BaseBlock>> generateEdgeSet(Set<BaseBlock> baseBlockSet,

Map<Integer,BaseBlock> leaderMap) {

Set<Pair<BaseBlock,BaseBlock>> edgeSet = new ListSet<>();

for (BaseBlock baseBlock : baseBlockSet) {

int firstIndex = baseBlock.getFirst(), lastIndex = baseBlock.getLast();

for (int index = firstIndex; index <= lastIndex; ++index) {

MiddleCode middleCode = m\_middleCodeList.get(index);

switch (middleCode.getOperator()) {

case Goto: {

int targetAddress = (int) middleCode.getOperand(1);

BaseBlock targetBlock = leaderMap.get(targetAddress);

edgeSet.add(new Pair<>(baseBlock, targetBlock));

}

break;

case Equal:

case NotEqual:

case LessThan:

case LessThanEqual:

case GreaterThan:

case GreaterThanEqual: {

BaseBlock nextBlock = leaderMap.get(index + 1);

edgeSet.add(new Pair<>(baseBlock, nextBlock));

int targetAddress = (int) middleCode.getOperand(1);

BaseBlock targetBlock = leaderMap.get(targetAddress);

edgeSet.add(new Pair<>(baseBlock, targetBlock));

}

break;

case FunctionEnd:

case Return:

case Exit:

break;

default:

if (index == lastIndex) {

BaseBlock nextBlock = leaderMap.get(index + 1);

edgeSet.add(new Pair<>(baseBlock, nextBlock));

}

break;

}

}

}

return edgeSet;

}

### Live Sets Generation

We equip each middle code instruction with a live set; that is, the set of all symbols which values are in demand later in the base block. A symbol is live if it is an operand in a unary or binary expression, or as a parameter or return value. However, it is not live if it is the result of an expression.

private void generateLiveMaps() {

for (BaseBlock baseBlock : m\_controlFlowGraph.getVertexSet()) {

int firstIndex = baseBlock.getFirst(), lastIndex = baseBlock.getLast();

Map<Symbol,Integer> liveMap = new ListMap<>();

for (int index = lastIndex; index >= firstIndex; --index) {

MiddleCode middleCode = m\_middleCodeList.get(index);

middleCode.setLiveMap(((ListMap) liveMap).clone());

MiddleOperator middleOp = middleCode.getOperator();

switch (middleOp) {

case Address:

break;

default: {

Object resultObject = middleCode.getOperand(1),

leftObject = middleCode.getOperand(2),

rightObject = middleCode.getOperand(3);

Symbol resultSymbol, leftSymbol, rightSymbol;

if (middleCode.getOperator() == MiddleOperator.Assign) {

resultSymbol = (leftObject instanceof Symbol) ? ((Symbol) leftObject) : null;

leftSymbol = (rightObject instanceof Symbol) ? ((Symbol) rightObject) : null;

rightSymbol = null;

}

else {

resultSymbol = (resultObject instanceof Symbol) ? ((Symbol) resultObject) : null;

leftSymbol = (leftObject instanceof Symbol) ? ((Symbol) leftObject) : null;

rightSymbol = (rightObject instanceof Symbol) ? ((Symbol) rightObject) : null;

}

Symbol resultAddressSymbol = (resultSymbol != null) ? resultSymbol.getAddressSymbol() : null,

leftAddressSymbol = (leftSymbol != null) ? leftSymbol.getAddressSymbol() : null,

rightAddressSymbol = (rightSymbol != null) ? rightSymbol.getAddressSymbol() : null;

removeIfNotNull(resultSymbol, liveMap);

addIfNotNull(resultAddressSymbol, liveMap);

addIfNotNull(leftSymbol, liveMap);

addIfNotNull(leftAddressSymbol, liveMap);

addIfNotNull(rightSymbol, liveMap);

addIfNotNull(rightAddressSymbol, liveMap);

/\* if (middleCode.isMultiply()) {

Set<Symbol> removeSet = new ListSet<>();

for (Map.Entry<Symbol,Pair<Integer,Boolean>> entry : liveMap.entrySet()) {

Symbol symbol = entry.getKey();

Pair<> pair = entry.getValue();

boolean isMul = pair.getSecond();

if (isMul) {

removeSet.add(symbol);

}

}

for (Symbol symbol : removeSet) {

liveMap.removeParentsAndChildSymbols(symbol);

}

}\*/

}

}

}

}

}

private void removeIfNotNull(Symbol symbol, Set<Symbol> liveSet) {

if (symbol != null) {

liveSet.remove(symbol);

}

}

private void addIfNotNull(Symbol symbol, Set<Symbol> liveSet) {

if (symbol != null) {

liveSet.add(symbol);

}

}

### Path Sets

private void generatePathSets() {

BaseBlock firstBlock = m\_controlFlowGraph.getVertexSet().iterator().next();

Assert.error(m\_controlFlowGraph.getInNeighbourSet(firstBlock).isEmpty());

int size = m\_controlFlowGraph.getVertexSet().size();

for (BaseBlock startBlock : m\_controlFlowGraph.getVertexSet()) {

for (BaseBlock inBlock:m\_controlFlowGraph.getInNeighbourSet(startBlock)) {

generateInPathSet(firstBlock, startBlock, inBlock,

new LinkedList<Object>(), startBlock.getInPathSet());

}

for (BaseBlock outBlock:m\_controlFlowGraph.getOutNeighbourSet(startBlock)){

generateOutPathSet(startBlock, outBlock, new LinkedList<Object>(),

startBlock.getOutPathSet());

}

}

}

private void generateInPathSet(BaseBlock firstBlock, BaseBlock startBlock,

BaseBlock currBlock, List<Object> inPath,

Set<List<Object>> inPathSet){

if (currBlock == firstBlock) {

inPath.add(0, currBlock);

inPathSet.add(new LinkedList<>(inPath));

inPath.remove(0);

}

else if ((currBlock != startBlock) && !inPath.contains(currBlock)) {

inPath.add(0, currBlock);

for (BaseBlock inBlock:m\_controlFlowGraph.getInNeighbourSet(currBlock)) {

generateInPathSet(firstBlock, startBlock, inBlock, inPath, inPathSet);

}

inPath.remove(0);

}

}

private void generateOutPathSet(BaseBlock startBlock, BaseBlock currBlock,

List<Object> outPath, Set<List<Object>> outPathSet) {

Set<BaseBlock> outNeighbourSet =

m\_controlFlowGraph.getOutNeighbourSet(currBlock);

if ((currBlock == startBlock) || outNeighbourSet.isEmpty()) {

outPath.add(currBlock);

outPathSet.add(new LinkedList<>(outPath));

outPath.remove(outPath.size() - 1);

}

else if (outPath.contains(currBlock)) {

outPathSet.add(new LinkedList<>(outPath));

}

else {

outPath.add(currBlock);

Set<List<Object>> longPathSet = new ListSet<>();

for (BaseBlock outBlock : outNeighbourSet) {

generateOutPathSet(startBlock, outBlock, outPath, longPathSet);

}

if (!longPathSet.isEmpty()) {

Set<List<Object>> shortPathSet = new ListSet<>();

for (List<Object> longPath : longPathSet) {

shortPathSet.add(longPath.subList(outPath.size(), longPath.size()));

}

List<Object> newPath = new LinkedList<>(outPath);

if (shortPathSet.size() == 1) {

newPath.addAll(shortPathSet.iterator().next());

outPathSet.add(newPath);

}

else {

newPath.add(shortPathSet);

outPathSet.add(newPath);

}

}

outPath.remove(outPath.size() - 1);

}

}

private void generateDefAndUseBlockSets() {

for (BaseBlock baseBlock : m\_controlFlowGraph.getVertexSet()) {

int firstIndex = baseBlock.getFirst(), lastIndex = baseBlock.getLast();

for (int index = lastIndex; index >= firstIndex; --index) {

MiddleCode middleCode = m\_middleCodeList.get(index);

addParentsAndChildSymbols(baseBlock.getHardDefSymbolSet(),

middleCode.getDefSymbolSet());

addParentsAndChildSymbols(baseBlock.getHardUseSymbolSet(),

middleCode.getUseSymbolSet());

addParentsAndChildSymbols(baseBlock.getSoftDefSymbolSet(),

middleCode.getDefSymbolSet());

removeParentsAndChildSymbols(baseBlock.getSoftDefSymbolSet(),

baseBlock.getHardUseSymbolSet());

removeParentsAndChildSymbols(baseBlock.getSoftUseSymbolSet(),

middleCode.getDefSymbolSet());

addParentsAndChildSymbols(baseBlock.getSoftUseSymbolSet(),

middleCode.getUseSymbolSet());

}

}

}

private void addParentsAndChildSymbols(Set<Symbol> mainSet,

Set<Symbol> addSet) {

for (Symbol symbol : addSet) {

Symbol parentSymbol = symbol;

while (parentSymbol != null) {

mainSet.add(parentSymbol);

parentSymbol = parentSymbol.getParentSymbol();

}

addChildSet(mainSet, symbol.getChildSymbolSet());

}

}

private void addChildSet(Set<Symbol> mainSet, Set<Symbol> childSet) {

for (Symbol childSymbol : childSet) {

mainSet.add(childSymbol);

addChildSet(mainSet, childSymbol.getChildSymbolSet());

}

}

private void removeParentsAndChildSymbols(Set<Symbol> mainSet,

Set<Symbol> removeSet) {

for (Symbol removeSymbol : removeSet) {

Symbol parentSymbol = removeSymbol;

while (parentSymbol != null) {

mainSet.remove(parentSymbol);

parentSymbol = parentSymbol.getParentSymbol();

}

removeChildSet(mainSet, removeSymbol.getChildSymbolSet());

}

}

private void removeChildSet(Set<Symbol> mainSet, Set<Symbol> childSet) {

for (Symbol childSymbol : childSet) {

mainSet.remove(childSymbol);

removeChildSet(mainSet, childSymbol.getChildSymbolSet());

}

}

private void removeIfNotNull(Symbol symbol, Map<Symbol,Integer> liveMap) {

if (symbol != null) {

liveMap.remove(symbol);

}

}

private void addIfNotNull(Symbol symbol, Map<Symbol,Integer> liveMap) {

if ((symbol != null) && !symbol.hasValue() && !symbol.isSolid()) {

int count = liveMap.containsKey(symbol) ? liveMap.get(symbol) : 0;

liveMap.put(symbol, count + 1);

}

}

}

### Removal of Unused Code

Unused code is code that does not perform any meaningful task, such as:

|  |  |  |
| --- | --- | --- |
| x \* y; | 1. $0 = x \* y |  |

We use the live sets and remove all instructions that temporary symbols that are not a member of the live set.

# Object Code Preparation

## Register Preparation

## Register Allocation

Graph coloring is a NP-complete problem, which means that there is no known algorithm that works on polynomial time. To be sure that we have found the best solution, we have to exam every possible solution, which would require an unrealistic amount of time. In this book we perform a deep search, where we sort the vertices into a list and for each vertices try to found a register not already allocated by any of its neighbors. If we cannot find such register we backtrack and try another combination. If we finally find a solution where each vertex is mapped to a register not mapped by any of its neighbors, we have found an optimal solution.

If the deep search does not found such a solution, we have to do something about the graph. In this book we remove one edge and try the deep search again. In this way we continue to remove edges until we have found a solution. The benefit of removing edges is that the likeliness of founding a solution increases for each removed edge. The drawback is that each removed edge means that two tracks have to share we one register and that we have to add store and load instruction for the code to work. When we need to remove an edge we sort the edges by counting the number of load and store instructions necessary to add as compensation for the removal of the edge. There is no guaranty that this method generates the optimal solution but since shortage of registers only occur on rare occasions it shall be good enough.

# Object Code Generation

The object code generation[[6]](#footnote-6) is the …

## The Object Operator

ObjectOperator.java

package c\_compiler;

public enum ObjectOperator {add, add\_byte, add\_dword, add\_word,

and, and\_byte, and\_dword, and\_word,

cmp, cmp\_byte, cmp\_dword, cmp\_word,

dec, dec\_byte, dec\_dword, dec\_word,

div, div\_byte, div\_dword, div\_word,

fabs, faddp, fadd, fchs, fcompp,

fdivp, fdivrp, fdivr, fdiv,

fild\_dword, fild\_word, finitializer,

fistp\_dword, fistp\_word, fist\_dword,

fist\_word, fld1, fldcw, fldz,

fld\_dword, fld\_qword, fmulp, fmul,

fstcw, fstp\_dword, fstp\_qword,

fstsw, fst\_dword, fst\_qword, fsubp,

fsubrp, fsubr, fsub, ftst, fwait,

idiv, idiv\_byte, idiv\_dword,

idiv\_word, imul, imul\_byte,

imul\_dword, imul\_word, inc,

inc\_byte, inc\_dword, inc\_word,

interrupt, jae, ja, jbe, jb, jc, je,

jge, jg, jle, jl, jmp, short\_jmp,

long\_jmp, jnc, jne, jnz, jz, lahf,

mov, mov\_byte, mov\_dword, mov\_word,

mul, mul\_byte, mul\_dword, mul\_word,

neg, neg\_byte, neg\_dword, neg\_word,

nop, not, not\_byte, not\_dword,

not\_word, or, or\_byte, or\_dword,

or\_word, sahf, shl, shr, sub,

sub\_byte, sub\_dword, sub\_word, xor,

xor\_byte, xor\_dword, xor\_word, call,

empty, label, label2, comment};

## The Object Code

ObjectCode.java

### Object Code Generation

|  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- |
| Case | Operator | Operand 1 | Operand 2 | Operand 3 | Full Name | Example |
| 1 | mov | null | null | Integer | null |  |
| 2 | mov\_wordX | Register | Integer | null |  |  |
| 3 |  | Integer | Null | Null | Null | int 21 |
| 4 |  |  |  |  |  |  |
| 5 |  |  |  |  |  |  |
|  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |

## Object Code Generation

# Executable Code Generation

In this book, the target code so far has been assembly code that need to be assembled and linked before execution. However, in this chapter we look into an alternative way, we generate target code in the format of executable code, that we link into an executable file. I have chosen the 16-bit .com format, which is a simple file format. Unfortunately, the format is no longer supported by Windows. Therefore, we need use a simulator the supports the format. I have chosen DosBox, which is a simple simulator capable of executing files in the .com format.

In the book so far, we have encountered the Start.Linux condition on several occasions, indicating code specific for the Linux target machine. However, there are also Start.Windows conditions, indicating code specific for the Windows target machine, which we will look into in this chapter.

### Command Line Arguments

## Linking

The linker if the final part of the compilation process. Its merges together the compiled files and generates an executable file. To be more exact, it has the following tasks:

* To resolve access to static symbols
* To resolve function calls.
* To resolve return addresses.

### The Linker Class

The LinkerInfo class holds the information of a function definition or a static variable:

* m\_fullName: the full name of the function or static variable. The full name of a function or global variable with external linkage is just its regular name, while the full name of
* m\_entryPoint: the function call entry point, which may differ from the start address of the function[[7]](#footnote-7).
* m\_accessMap: the map of all accesses of function or static variables.
* m\_callMap: the map of all function calls. The reason we distinguish between m\_accessMap and m\_callMap is that accesses are absolute while function calls are relative.
* m\_returnSet: the set of all return address assignments.
* m\_byteList: the actual code of the function or initialization data of the static variable. If the static variable is uninitialized, all data is zero.

The Linker class merge code of the object files and modifies the global accesses, function calls, and return assignments and saves the final code in a com-file, which is executable in 16-bits Windows. More specific, the linker has the following tasks:

* First we load all function and static variables from all the object files and to make sure that two elements do not share the same name.
* The we identify the main function and trace all functions and global variables reachable from main. All elements not reachable from main are discarded.
* For each function or global variable we modifies all accesses, and for each function we modify all calls and return assignments.
* Finally, we save the code and data for all elements reachable from main in the com-file.

There are several maps:

* m\_globalMap: holds all functions and static variables of all object files.
* m\_globalList: holds all functions and static variables reachable from main, where main is always placed first.
* m\_addressMap holds the beginning of each function or static variables in the final code.
* m\_entryMap holds the entry point of each function.[[8]](#footnote-8)

Linker.cs

using System;

using System.IO;

using System.Collections.Generic;

namespace CCompiler {

public class WindowsLinker {

The m\_globalMap map holds all the static symbols of the source code, while m\_globalList holds only the symbols that are reachable (directly or indirectly) from the main function. The final code is generated from the symbols of m\_globalList, the other symbols are omitted from the final code. The m\_addressMap holds the address of each symbol in m\_globalList. Finally, m\_totalSize holds the current total size of the symbols of m\_globalList and is used to add the positions of the symbols to m\_addressMap.

private IDictionary<string,StaticSymbolWindows> m\_globalMap =

new Dictionary<string,StaticSymbolWindows>();

private List<StaticSymbolWindows> m\_globalList =

new List<StaticSymbolWindows>();

private IDictionary<string,int> m\_addressMap =

new Dictionary<string,int>();

private int m\_totalSize = 256;

The Add method is called for each static symbol of the source code. If two symbols have the same unique name, that does not end with the number identification character (‘#’), we have two symbols with external linkage, which is not allowed. However, if the name ends with the identification character, it only means that two numerical constants with the same value is present in two different files, with is allowed. We simply refrain from adding the second symbol to the map.

public void Add(StaticSymbolWindows staticSymbol) {

string uniqueName = staticSymbol.UniqueName;

if (!m\_globalMap.ContainsKey(uniqueName)) {

m\_globalMap.Add(uniqueName, staticSymbol);

}

else {

Assert.Error(uniqueName.EndsWith(Symbol.NumberId),

SimpleName(uniqueName), Message.Duplicate\_global\_name);

}

}

The Generate method writes the final executable code to the target file. Its task is to identify the symbols reachable from the main function, and to replace the names of accessed symbols with called functions with proper addresses as well as to replace the function returns addresses with proper addresses. However, we need to start by including the code for the initialization execution code and, optionally, the code for handling the command line arguments.

public void Generate(FileInfo targetFile) {

The code for initialization shall be added to the code at the beginning, before the main function

{ StaticSymbolWindows initializerInfo =

m\_globalMap[AssemblyCodeGenerator.InitializerName];

m\_globalList.Add(initializerInfo);

m\_totalSize += initializerInfo.ByteList.Count;

m\_addressMap.Add(AssemblyCodeGenerator.InitializerName, 0);

}

In case of command line arguments, we add its symbol after the initialization code symbol and before the main function code.

StaticSymbolWindows pathNameSymbol = null;

if (m\_globalMap.ContainsKey(AssemblyCodeGenerator.ArgsName)) {

{ StaticSymbolWindows argsInfo =

m\_globalMap[AssemblyCodeGenerator.ArgsName];

m\_globalList.Add(argsInfo);

m\_totalSize += argsInfo.ByteList.Count;

m\_addressMap.Add(AssemblyCodeGenerator.ArgsName, 0);

}

We also need to add the path name of final executable file.

{ List<byte> byteList = new List<byte>();

IDictionary<int, string> accessMap = new Dictionary<int, string>();

pathNameSymbol = (StaticSymbolWindows)

ConstantExpression.Value(AssemblyCodeGenerator.PathName,

Type.StringType, targetFile.FullName);

m\_globalMap.Add(AssemblyCodeGenerator.PathName, pathNameSymbol);

}

}

We look up the main function symbol and report an error if it is not present. If it is present, we add it after the initiation symbol and potential command line symbol. Then we call GenerateTrace, which traces all symbols reachable from the main function, and adds them to m\_globalList and m\_addressMap.

StaticSymbolWindows mainInfo;

Assert.Error(m\_globalMap.TryGetValue(AssemblyCodeGenerator.MainName,

out mainInfo), "non-static main",

Message.Function\_missing);

GenerateTrace(mainInfo);

In case of command line arguments, we add the symbol for the executable file path name to m\_globalList and m\_addressMap.

if (pathNameSymbol != null) {

Assert.ErrorXXX(!m\_globalList.Contains(pathNameSymbol));

m\_globalList.Add(pathNameSymbol);

m\_addressMap.Add(pathNameSymbol.UniqueName, m\_totalSize);

m\_totalSize += (int) pathNameSymbol.ByteList.Count;

}

Finally, we add the stack top name, which will be the address of the activation record of the main function.

m\_addressMap.Add(AssemblyCodeGenerator.StackTopName, m\_totalSize);

We iterate through the global list and, for each of its symbols, we call GenerateAccess, GenerateCall, GenerateReturn, which replace the names of accessed static symbols and called functions as well as return addresses with proper addresses.

foreach (StaticSymbolWindows staticSymbol in m\_globalList) {

List<byte> byteList = staticSymbol.ByteList;

int startAddress = m\_addressMap[staticSymbol.UniqueName];

GenerateAccess(staticSymbol.AccessMap, byteList);

GenerateCall(startAddress, staticSymbol.CallMap, byteList);

GenerateReturn(startAddress, staticSymbol.ReturnSet, byteList);

}

Finally, we write the byte list of each symbol in the global list to the target file.

{ Console.Out.WriteLine("Generating \"" + targetFile.FullName + "\".");

targetFile.Delete();

BinaryWriter targetStream =

new BinaryWriter(File.OpenWrite(targetFile.FullName));

foreach (StaticSymbolWindows staticSymbol in m\_globalList) {

foreach (sbyte b in staticSymbol.ByteList) {

targetStream.Write(b);

}

}

targetStream.Close();

}

}

The GenerateTrace method adds the symbol to the m\_globalList iterates through the names of access name set and function call name set and calls GenerateTrace recursively for each symbol.

private void GenerateTrace(StaticSymbolWindows staticSymbol) {

if (!m\_globalList.Contains(staticSymbol)) {

m\_globalList.Add(staticSymbol);

m\_addressMap.Add(staticSymbol.UniqueName, m\_totalSize);

m\_totalSize += (int) staticSymbol.ByteList.Count;

ISet<string> accessNameSet =

new HashSet<string>(staticSymbol.AccessMap.Values);

foreach (string accessName in accessNameSet) {

StaticSymbolWindows accessSymbol;

Assert.Error(m\_globalMap.TryGetValue(accessName, out accessSymbol),

accessName, Message.Object\_missing\_in\_linking);

Assert.ErrorXXX(accessSymbol != null);

GenerateTrace(accessSymbol);

}

ISet<string> callNameSet =

new HashSet<string>(staticSymbol.CallMap.Values);

foreach (string callName in callNameSet) {

StaticSymbolWindows funcSymbol;

Assert.Error(m\_globalMap.TryGetValue(callName, out funcSymbol),

callName, Message.Function\_missing\_in\_linking);

Assert.Error(funcSymbol != null, SimpleName(callName),

Message.Missing\_external\_function);

GenerateTrace(funcSymbol);

}

}

}

The GenerateAccess methods iterators through the access map and modifies the addresses.

private void GenerateAccess(IDictionary<int,string> accessMap,

List<byte> byteList) {

foreach (KeyValuePair<int,string> entry in accessMap) {

int sourceAddress = entry.Key;

string targetName = entry.Value;

We obtain the target address from the low and high byte of the source address. Note that the target address does not need to be zero, it may hold an offset. For instance, the target address may be a pointer to a value in an array of a non-zero index, in which case its offset is non-zero.

byte lowByte = byteList[sourceAddress],

highByte = byteList[sourceAddress + 1];

int targetAddress = ((int) highByte << 8) + lowByte;

We modify the target address by adding the address of the name to target address, which we restore to the source address.

targetAddress += m\_addressMap[targetName];

byteList[sourceAddress] = (byte) targetAddress;

byteList[sourceAddress + 1] = (byte) (targetAddress >> 8);

}

}

The GenerateCall method iterates through the names of the functions to called. There are two differences between this method and GenerateAccess above. First, we do not need to take into consideration that the original value on the address of the function call may be non-zero, since it is not possible to add an offset to a function call. Second, the address of the call is relative the next assembly code instruction.

private void GenerateCall(int callerStartAddress,

IDictionary<int,string> callMap,

List<byte> byteList) {

foreach (KeyValuePair<int,string> entry in callMap) {

int sourceAddress = entry.Key;

string sourceName = entry.Value;

The source address is the last two bytes of the function assembly code instruction, why the address of the next instruction is the source address plus two. However, the source address is relative the beginning of the function. Therefore, we need to add the global start address of the caller function to the caller address, in order to make the address relative the whole executable file.

int nextAddress = (sourceAddress + 2) + callerStartAddress;

int calleeAddress = m\_addressMap[sourceName];

The relative address of the call is the difference of the address of the callee function and the next instruction (the address of the instruction following the call).

int relativeAddress = calleeAddress - nextAddress;

byteList[sourceAddress] = (byte) ((sbyte) relativeAddress);

byteList[sourceAddress + 1] = (byte) ((sbyte) (relativeAddress >> 8));

}

}

The GenerateReturn method changes the return address from the address relative the beginning of the function

private void GenerateReturn(int functionStartAddress, ISet<int> returnSet,

List<byte> byteList) {

foreach (int sourceAddress in returnSet) {

int lowByte = byteList[sourceAddress],

highByte = byteList[sourceAddress + 1];

int targetAddress = (highByte << 8) + lowByte;

targetAddress += functionStartAddress;

byteList[sourceAddress] = (byte) targetAddress;

byteList[sourceAddress + 1] = (byte) (targetAddress >> 8);

}

}

The SimpleName method returns the left part of the name if it includes the separator identity character (‘$’).

public static string SimpleName(string name) {

int index = name.LastIndexOf(Symbol.SeparatorId);

return (index != -1) ? name.Substring(0, index) : name;

}

}

}

To begin with, addAll is called by Main class of Section 20. If the same symbol is loaded from two object files, unless the name starts with “x#” (in that case it is a numerical constant, which is allowed in more than one object file), an error message is reported. All loaded elements are stored in m\_globalMap.

public void add(Symbol symbol) {

String name = symbol.getUniqueName();

if (name.equals(Main.StackTopName)) {

m\_stackTop = symbol;

}

else {

Symbol result = m\_globalMap.put(name, symbol);

Assert.error((result == null) || name.endsWith(Main.NumberId),

simpleName(name), "duplicate global identifier");

}

}

The generate method traces the elements reachable from main by calling generateTrace, and sets the total size of the code at the beginning of the main code to initialize the stack.

public void generate() throws Exception {

m\_globalMap.put(Main.PathName, generatePathSymbol());

{ Symbol mainInfo = m\_globalMap.get("main");

Assert.error(mainInfo != null, "main", "function missing");

generateTrace(mainInfo);

m\_addressMap.put(m\_stackTop.getUniqueName(), m\_totalSize);

m\_globalList.add(m\_stackTop);

m\_totalSize += m\_stackTop.getByteList().size();

}

for (Symbol symbol : m\_globalList) {

List<Byte> byteList = symbol.getByteList();

int startAddress = m\_addressMap.get(symbol.getUniqueName());

generateAccess(symbol.getAccessMap(), byteList);

generateCall(startAddress, symbol.getCallMap(), byteList,

symbol.getTextList(), symbol.getByteToTextMap());

generateReturn(startAddress, symbol.getReturnSet(), byteList);

}

for (Symbol symbol : m\_globalList) {

int startAddress = m\_addressMap.get(symbol.getUniqueName());

splitTextList(symbol.getTextList());

generateTextAddress(startAddress, symbol.getByteList(),

symbol.getTextList());

}

{ System.out.println("Generating \"" + m\_comFile.getPath() + "\".");

OutputStream comStream = new FileOutputStream(m\_comFile);

for (Symbol symbol : m\_globalList) {

for (Byte b : symbol.getByteList()) {

comStream.write(b);

}

}

comStream.close();

}

{ System.out.println("Generating \"" + m\_asmFile.getPath() + "\".");

PrintStream asmStream = new PrintStream(m\_asmFile);

asmStream.println("\torg 100h");

for (Symbol symbol : m\_globalList) {

for (String text : symbol.getTextList()) {

if (text != null) {

asmStream.println(text);

}

}

}

asmStream.println("x" + m\_totalSize + ":");

asmStream.close();

}

}

private Symbol generatePathSymbol() {

List<Byte> byteList = new LinkedList<>();

Map<Integer,String> accessMap = new ListMap<>(),

callMap = new ListMap<>();

Set<Integer> returnSet = new ListSet<>();

Map<Integer,Integer> byteToTextMap = new ListMap<>();

List<String> textList = new LinkedList<>();

GenerateInitializerStatic.generateByteTextList

(Type.StringType, m\_comFile.getPath(),byteList, accessMap, textList);

Symbol symbol =

Symbol.createStaticSymbol(Main.PathName, byteList, accessMap,

callMap, returnSet, byteToTextMap, textList);

symbol.setByteList(byteList);

symbol.setTextList(textList);

return symbol;

}

### Main Trace Generation

The generateTrace method looks into the globalInfo parameter, and adds it to m\_globalList. However, we must check that the element is not already a member since recursive function calls may occur. The start address and entry point of the element is added m\_addressMap and m\_entryMap, respectively. The m\_totalSize field keeps track of the current size of the code.

Then we continue to call generateTrace recursively for each function call in the call map and global variables access in the access map. If any function or variable is missing, an error message is reported.

private void generateTrace(Symbol symbol) {

if (!m\_globalList.contains(symbol)) {

m\_globalList.add(symbol);

m\_addressMap.put(symbol.getUniqueName(), m\_totalSize);

m\_entryMap.put(symbol.getUniqueName(), m\_totalSize +

symbol.getEntryPoint());

m\_totalSize += symbol.getByteList().size();

Set<String> accessNameSet =

new ListSet<>(symbol.getAccessMap().values());

for (String accessName : accessNameSet) {

Symbol accessSymbol = m\_globalMap.get(accessName);

if (!accessName.equals(Main.StackTopName)) {

Assert.error(accessSymbol != null, "Linkage",

"missing external variable \"" + simpleName(accessName)+

"\" in \"" + simpleName(symbol.getUniqueName()) +"\"");

generateTrace(accessSymbol);

}

}

Set<String> callNameSet = new ListSet<>(symbol.getCallMap().values());

for (String callName : callNameSet) {

Symbol funcSymbol = m\_globalMap.get(callName);

Assert.error(funcSymbol != null, callName,

"missing extern function \"" +

simpleName(callName) + "\" in \"" +

simpleName(symbol.getUniqueName()) + "\"");

generateTrace(funcSymbol);

}

}

}

For each element, reachable from main, we need to modify its accesses, function calls and return assignment (for a global variable, the call map and return set are empty) by calling generateAccess, generateCall, and generateReturn, respectivly.

The generateAccess method looks up the name and address of each global access. The address of the name is looked up in accessMap, and original address is calculated by adding the low anad high byte of the address. Note that the access may be not zero in some cases, such as pointer arithmentic with a static pointer and constant integer. In the code below, the initializerial value of p will be two, since the address is two bytes higher than the address of a; generateAccess will then look up and add the address of a in order to properly initialize p.

static char arr[10];

static int p = &a[2];

private void generateAccess(Map<Integer,String> accessMap,

List<Byte> byteList) {

for (Map.Entry<Integer,String> entry : accessMap.entrySet()) {

int address = entry.getKey();

String name = entry.getValue();

int oldLowByte = checkValue(byteList.get(address)),

oldHighByte = checkValue(byteList.get(address + 1));

int oldTarget = (oldHighByte << 8) + oldLowByte;

int newTarget = oldTarget + m\_addressMap.get(name);

byte newLowByte = (byte) newTarget,

newHighByte = (byte) (newTarget >> 8);

byteList.set(address, newLowByte);

byteList.set(address + 1, newHighByte);

}

}

private static int checkValue(int i) {

return (i < 0) ? (i + 256) : i;

}

The generateCall method sets the address of each function call by looking up the function. Unlike the access case above, we do not have to look up the original address, it is always zero. However, since jumps are relative we need to calculate the relative address between the source and target. Note that the caller address is the address of the instruction following the caller instruction, which is the address stored in callMap plus the start address of the calling function plus another two bytes.

Moreover, we have two special cases and two general cases. ObjectCodeGenerator of Section 16.3 generated long jumps for each function call; that is, a long jump instruction followed two bytes holding the relative address. However, if the relative address only one byte, we change to a short jump instruction followed by one byte holding the relative address and insert a nop instruction in order to keep the size of the code[[9]](#footnote-9).

The first special case occurs when the relative address is 127, in which case we set 127 as the low byte jump address and 0 as the high byte. The second special case is when the relative address is -129, in which case we change the jump instruction from long jump to short jump with the relative address -128.

The first general case occurs when the relative address is between -128 and 126, inclusive. In that case, we change the jump instruction from long the short jump and insert a nop instruction. Finally, in all other cases we set the low and high byte of the relative jump address.

private void generateCall(int startAddress, Map<Integer,String> callMap,

List<Byte> byteList, List<String> textList,

Map<Integer,Integer> byteToTextMap) {

for (Map.Entry<Integer,String> entry : callMap.entrySet()) {

int address = entry.getKey();

int callerAddress = startAddress + address + 2;

int calleeAddress = m\_entryMap.get(entry.getValue());

int relativeAddress = calleeAddress - callerAddress;

if (relativeAddress == -129) {

byteList.set(address - 1,(byte) ObjectCode.ShortJumpOperator);

byteList.set(address, (byte) -128);

byteList.set(address + 1, (byte) ObjectCode.NopOperator);

}

else if ((relativeAddress >= -128) && (relativeAddress <= 127)) {

byteList.set(address - 1, (byte) ObjectCode.NopOperator);

byteList.set(address, (byte) ObjectCode.ShortJumpOperator);

byteList.set(address + 1, (byte) relativeAddress);

int textLine = byteToTextMap.get(address);

Assert.error(textLine < textList.size(), "text line");

String text = textList.get(textLine);

textList.set(textLine,"\tnop\t; 1\n" + text.replace("\t; 3", "\t; 2"));

}

else {

byteList.set(address, (byte) relativeAddress);

byteList.set(address + 1, (byte) (relativeAddress >> 8));

}

}

}

The generateReturn method modifies all return address assignments. When a function call is generated in Section **Fel! Hittar inte referenskälla.**, the return assignment address is relative to the assignment itself, we need to add the function start address and the address of the return assignment to the relative address.

private void generateReturn(int startAddress, Set<Integer> returnSet,

List<Byte> byteList) {

for (int address : returnSet) {

int relativeLowByte = checkValue(byteList.get(address)),

relativeHighByte = checkValue(byteList.get(address + 1));

int relativeAddress = (relativeHighByte << 8) + relativeLowByte;

int globalAddress = startAddress + address + relativeAddress;

byte globalLowByte = (byte) globalAddress;

byte globaHighByte = (byte) (globalAddress >> 8);

byteList.set(address, globalLowByte);

byteList.set(address + 1, globaHighByte);

}

}

private void generateTextAddress(int address, List<Byte> byteList,

List<String> textList) {

int byteIndex = 0;

for (int line = 0; line < textList.size(); ++line) {

String text = textList.get(line);

if ((text != null) && !text.isEmpty() && !text.contains(":\t;") &&

!text.startsWith(";") && !text.endsWith(":")) {

StringBuilder buffer = new StringBuilder(text);

buffer.insert(0, "x" + address + ":");

int semiIndex = buffer.lastIndexOf("; "), size = 0;

try {

size = Integer.valueOf(buffer.substring(semiIndex + 2));

}

catch (NumberFormatException exception) {

Assert.error(exception.getMessage());

}

buffer.append(":");

for (int index = 0; index < size; ++index) {

buffer.append(" " + byteList.get(byteIndex++));

}

textList.set(line, buffer.toString());

address += size;

}

}

Assert.error(byteIndex == byteList.size(), "linker byte list");

}

private void splitTextList(List<String> textList) {

for (int index = (textList.size() - 1); index >= 0; --index) {

String text = textList.get(index);

int newLineIndex = text.indexOf("\n");

if (newLineIndex != -1) {

textList.set(index, text.substring(0, newLineIndex));

textList.add(index + 1, text.substring(newLineIndex + 1));

}

}

}

String simpleName(String name) {

int index = name.lastIndexOf(Main.SeparatorId);

return ((index != -1)) ? name.substring(0, index) : name;

}

}

# Auxiliary Classes

Java comes with a large class library, with functionality for almost all requirements. However, there are some classes that are missing, which we need to write ourselves: a class for error handling, container classes, and a graph class.

## Error Handling

To begin with, we need error handling, a way to notify the programmer of errors and warnings. In the C Standard Library there is a macro called assert, and in this application we imitate that macro by defining a class with the corresponding behavor. It has the two method sets error and warning, that writes a message. The difference between them is that error stops the execution.

Assert.java

## Container Classes

Java has a large class library holding many container classes. However, there are no classes for ordered or unordered pairs, there is also no classes for triples or tuples.

### Ordered Pair

In the preprocessor (and in several other classes in the following chapters) we also need the auxiliary class Pair.

Pair.java

package c\_compiler;

public class Pair<FirstType,SecondType> {

protected FirstType m\_first;

protected SecondType m\_second;

public Pair(FirstType first, SecondType second) {

m\_first = first;

m\_second = second;

}

public FirstType getFirst() {

return m\_first;

}

public void setFirst(FirstType first) {

m\_first = first;

}

public SecondType getSecond() {

return m\_second;

}

public void setSecond(SecondType second) {

m\_second = second;

}

Two pairs are considered equal if (for both the first and second value) the values ar both null, or they are both not null and the values are equal (the equals method returns true).

public boolean equals(Object object) {

if (object instanceof Pair) {

Pair pair = (Pair) object;

return equalsFirst(pair) && equalsSecond(pair);

}

return false;

}

private boolean equalsFirst(Pair pair) {

return (((m\_first == null) && (pair.m\_first == null)) ||

((m\_first != null) && (pair.m\_first != null) &&

m\_first.equals(pair.m\_first)));

}

private boolean equalsSecond(Pair pair) {

return (((m\_second == null) && (pair.m\_second == null)) ||

((m\_second != null) && (pair.m\_second != null) &&

m\_second.equals(pair.m\_second)));

}

public String toString() {

String firstText = (m\_first != null) ? m\_first.toString() : "<null>",

secondText = (m\_second != null) ? m\_second.toString() : "<null>";

return "(" + firstText + "," + secondText + ")";

}

}

### Unordered Pair

UnPair is a subclass of Pair, the difference is that the pair is unordered and that two pair are considered equals even if their values are switched.

UnorderdPair.java

### Triple

Moreover, we also need the class Triple, to hold three values. The Triple class is a direct sub class of Pair, it only adds the third value.

Triple.java

## Graph

Graph implements a general unordered graph. It holds of set of vertices and a set of edges.

Graph.java

package c\_compiler;

import java.util.\*;

public class Graph<VertexType> {

private Set<VertexType> m\_vertexSet = new ListSet<>();

private Set<UnorderedPair<VertexType,VertexType>> m\_edgeSet = new ListSet<>();

public Graph() {

m\_vertexSet = new ListSet<>();

m\_edgeSet = new ListSet<>();

}

public Graph(Set<VertexType> vertexSet) {

m\_vertexSet = vertexSet;

m\_edgeSet = new ListSet<>();

}

public Graph(Set<VertexType> vertexSet,

Set<UnorderedPair<VertexType,VertexType>> edgeSet) {

m\_vertexSet = vertexSet;

m\_edgeSet = edgeSet;

}

public Set<VertexType> getVertexSet() {

return m\_vertexSet;

}

public Set<UnorderedPair<VertexType,VertexType>> getEdgeSet() {

return m\_edgeSet;

}

The neighbourSet method goes through all edges and add each found neighbor to the vertex.

public Set<VertexType> getNeighbourSet(VertexType vertex) {

Set<VertexType> neighbourSet = new ListSet<>();

for (UnorderedPair<VertexType,VertexType> edge : m\_edgeSet) {

if (edge.getFirst() == vertex) {

neighbourSet.add(edge.getSecond());

}

if (edge.getSecond() == vertex) {

neighbourSet.add(edge.getFirst());

}

}

return neighbourSet;

}

### Addition and Removal of Vertices and Edges

public void addVertex(VertexType vertex) {

m\_vertexSet.add(vertex);

}

public void eraseVertex(VertexType vertex) {

m\_vertexSet.remove(vertex);

}

public void addEdge(VertexType vertex1, VertexType vertex2) {

m\_edgeSet.add(new UnorderedPair<>(vertex1, vertex2));

}

public void addEdge(UnorderedPair<VertexType,VertexType> edge) {

m\_edgeSet.add(edge);

}

public void eraseEdge(VertexType vertex1, VertexType vertex2) {

m\_edgeSet.remove(new UnorderedPair<>(vertex1, vertex2));

}

public void eraseEdge(UnorderedPair<VertexType,VertexType> edge) {

m\_edgeSet.remove(edge);

}

### Graph Partition

The method partitionate divides the graph into free subgraphs; that is, subgraphs which vertices have no neighbors in any of the other free subgraphs. First we go through the vertices and perform a deep search to find all vertices reachable from the vertex. Then we generate a subgraph for each such vertex set.

public Set<Graph<VertexType>> partitionate() {

Set<Set<VertexType>> qliqueSet = new ListSet<>();

for (VertexType vertex : m\_vertexSet) {

Set<VertexType> vertexSet = new ListSet<>();

DeepFirstSearch(vertex, vertexSet);

qliqueSet.add(vertexSet);

}

Set<Graph<VertexType>> graphSet = new ListSet<>();

for (Set<VertexType> vertexSet : qliqueSet) {

graphSet.add(generateSubGraph(vertexSet));

}

return graphSet;

}

The DeepFirstSearch method search recursively through the graph in order to find all vertices reachable from the given vertex. To avoid cyclic search, the search is terminated if the vertex is already a member of the result set.

private void DeepFirstSearch(VertexType vertex, Set<VertexType> resultSet) {

if (!resultSet.contains(vertex)) {

resultSet.add(vertex);

Set<VertexType> neighbourSet = getNeighbourSet(vertex);

for (VertexType neighbour : neighbourSet) {

DeepFirstSearch(neighbour, resultSet);

}

}

}

The generateSubGraph method goes through the edge set and add all edges which both end vertices are members of the vertex set.

public Graph<VertexType> generateSubGraph(Set<VertexType> vertexSet) {

Set<UnorderedPair<VertexType,VertexType>> resultEdgeSet = new ListSet<>();

for (UnorderedPair<VertexType,VertexType> edge : m\_edgeSet) {

if (vertexSet.contains(edge.getFirst()) &&

vertexSet.contains(edge.getSecond())) {

resultEdgeSet.add(edge);

}

}

return (new Graph(vertexSet, resultEdgeSet));

}

public String toString() {

return "<" + m\_vertexSet + "," + m\_edgeSet + ">";

}

}

## Directed Graph

Besides the undirected graph, we also need the directed graph.

DirectedGraph.java

package c\_compiler;

import java.util.\*;

public class DirectedGraph<VertexType> {

private Set<VertexType> m\_vertexSet = new ListSet<>();

private Set<Pair<VertexType,VertexType>> m\_edgeSet = new ListSet<>();

public DirectedGraph() {

m\_vertexSet = new ListSet<>();

m\_edgeSet = new ListSet<>();

}

public DirectedGraph(Set<VertexType> vertexSet) {

m\_vertexSet = vertexSet;

m\_edgeSet = new ListSet<>();

}

public DirectedGraph(Set<VertexType> vertexSet,

Set<Pair<VertexType,VertexType>> edgeSet) {

m\_vertexSet = vertexSet;

m\_edgeSet = edgeSet;

}

public Set<VertexType> getVertexSet() {

return m\_vertexSet;

}

public Set<VertexType> getInNeighbourSet(VertexType vertex) {

Set<VertexType> inNeighbourSet = new ListSet<>();

for (Pair<VertexType,VertexType> edge : m\_edgeSet) {

if (edge.getSecond() == vertex) {

inNeighbourSet.add(edge.getFirst());

}

}

return inNeighbourSet;

}

public Set<VertexType> getOutNeighbourSet(VertexType vertex) {

Set<VertexType> outNeighbourSet = new ListSet<>();

for (Pair<VertexType,VertexType> edge : m\_edgeSet) {

if (edge.getFirst() == vertex) {

outNeighbourSet.add(edge.getSecond());

}

}

return outNeighbourSet;

}

public Set<Pair<VertexType,VertexType>> getEdgeSet() {

return m\_edgeSet;

}

// ------------------------------------------------------------------------

public void addVertex(VertexType vertex) {

m\_vertexSet.add(vertex);

}

public void eraseVertex(VertexType vertex) {

m\_vertexSet.remove(vertex);

}

public void addEdge(VertexType vertex1, VertexType vertex2) {

m\_edgeSet.add(new Pair<>(vertex1, vertex2));

}

public void addEdge(Pair<VertexType,VertexType> edge) {

m\_edgeSet.add(edge);

}

public void eraseEdge(VertexType vertex1, VertexType vertex2) {

m\_edgeSet.remove(new Pair<>(vertex1, vertex2));

}

public void eraseEdge(Pair<VertexType,VertexType> edge) {

m\_edgeSet.remove(edge);

}

public Set<DirectedGraph<VertexType>> partitionate() {

Set<Set<VertexType>> qliqueSet = new ListSet<>();

for (VertexType vertex : m\_vertexSet) {

Set<VertexType> vertexSet = new ListSet<>();

DeepFirstSearch(vertex, vertexSet);

qliqueSet.add(vertexSet);

}

Set<DirectedGraph<VertexType>> graphSet = new ListSet<>();

for (Set<VertexType> vertexSet : qliqueSet) {

graphSet.add(generateSubGraph(vertexSet));

}

return graphSet;

}

private void DeepFirstSearch(VertexType vertex, Set<VertexType> resultSet) {

if (!resultSet.contains(vertex)) {

resultSet.add(vertex);

Set<VertexType> outNeighbourSet = getOutNeighbourSet(vertex);

for (VertexType outNeighbour : outNeighbourSet) {

DeepFirstSearch(outNeighbour, resultSet);

}

}

}

public DirectedGraph<VertexType> generateSubGraph(Set<VertexType>vertexSet){

Set<Pair<VertexType,VertexType>> resultEdgeSet = new ListSet<>();

for (Pair<VertexType,VertexType> edge : m\_edgeSet) {

if (vertexSet.contains(edge.getFirst()) &&

vertexSet.contains(edge.getSecond())) {

resultEdgeSet.add(edge);

}

}

return (new DirectedGraph(vertexSet, resultEdgeSet));

}

public String toString() {

return "<" + m\_vertexSet + "," + m\_edgeSet + ">";

}

}

# The Standard Library

## Integral and Floating Limits

Limits.h

#ifndef \_\_LIMITS\_H\_\_

#define \_\_LIMITS\_H\_\_

#define CHAR\_BIT 8

#define CHAR\_MAX 127s

#define CHAR\_MIN -128s

#define SCHAR\_MAX 127s

#define SCHAR\_MIN -128s

#define UCHAR\_MAX 255us

#define SHRT\_MAX 127s

#define SHRT\_MIN -128s

#define USHRT\_MAX 255us

#define INT\_MAX 32767

#define INT\_MIN -32768

#define UINT\_MAX 65535u

#define LONG\_MAX 2147483647l

#define LONG\_MIN -2147483647l

#define ULONG\_MAX 4294967295lu

#endif

Float.h

#ifndef \_\_FLOAT\_H\_\_

#define \_\_FLOAT\_H\_\_

#define FLT\_RADIX 2

#define FLT\_ROUNDS 1

#define FLT\_DIG 6

#define FLT\_EPSILON 1e-6

#define FLT\_MANT\_DIG 2

#define FLT\_MAX 1e6

#define FLT\_MAX\_EXP 6

#define FLT\_MIN 1e-6

#define FLT\_MIN\_EXP -6

#define DBL\_DIG 6

#define DBL\_EPSILON 1e-6

#define DBL\_MANT\_DIG 2

#define DBL\_MAX 1e+6

#define DBL\_MAX\_EXP 6

#define DBL\_MIN 1e-6

#define DBL\_MIN\_EXP -6

#endif

## The Assert Macro

Assert.h

#ifndef \_\_ASSERT\_H\_\_

#define \_\_ASEERT\_H\_\_

#ifndef NDEBUG

#include <StdIO.h>

#include <StdLib.h>

#define assert(expression) if (!(expression)) { \\

fprintf(stderr, "Assertion failed: \"%s\", file \"%s\", line %i.\n", \\

#expression, \_\_FILE\_\_, \_\_LINE\_\_); abort(); }

#else

#define assert(expression)

#endif

#endif

## Locale Data

Locale.h

#ifndef \_\_LOCALE\_H\_\_

#define \_\_LOCALE\_H\_\_

#define LC\_COLLATE 0x01

#define LC\_CTYPE 0x02

#define LC\_MONETARY 0x04

#define LC\_NUMERIC 0x08

#define LC\_TIME 0x10

#define LC\_ALL 0x1F

struct lconv {

int summerTimeZone, // time

winterTimeZone;

char \*\*shortDayList;

char \*\*longDayList;

char \*\*shortMonthList;

char \*\*longMonthList;

char \*lowerCase; // ctype

char \*upperCase;

char \*\*messageList; // errno

};

extern char\* enMessageList[];

extern char\* setlocale(int flag, char\* name);

extern struct lconv \*localeconv(void);

#endif

Locale.c

#include <StdIO.h>

#include <StdLib.h>

#include <String.h>

#include <Locale.h>

static char\* enShortDayList[] = { "Sun", "Mon", "Tue", "Wed", "Thu", "Fri", "Sat" };

static char\* enLongDayList[] = { "Sunday", "Monday", "Tuesday", "Wednesday", "Thursday", "Friday", "Saturday" };

static char\* enShortMonthList[] = {"Jan", "Feb", "Mar", "Apr", "May", "Jun", "Jul", "Aug", "Sep", "Oct", "Nov", "Dec"};

static char\* enLongMonthList[] = {"January", "February", "March", "April", "May", "June",

"July", "August", "September", "October", "November", "December"};

char\* enMessageList[] = {"no error", "function number invalid",

"file not found", "path not found", "no handle available", "access denied",

"out of domain", "out of range", "invalid multibyte sequence",

"error while opening", "error while flushing", "error while closing", "open mode invalid",

"error while writing", "error while reading", "error while seeking",

"error while telling", "error while removing file", "error while renaming file"};

static struct lconv en\_US\_utf8 = {-5, -4, enShortDayList, enLongDayList, enShortMonthList, enLongMonthList,

"abcdefghijklmnopqrstuvwxyz", "ABCDEFGHIJKLMNOPQRSTUVWXYZ", enMessageList};

static char\* swShortDayList[] = {"Son", "Man", "Tis", "Ons", "Tor", "Fre", "Lor"};

static char\* swLongDayList[] = {"Sondag", "Mandag", "Tisdag", "Onsdag", "Torsdag", "Fredag", "Lordag"};

static char\* swShortMonthList[] = {"Jan", "Feb", "Mar", "Apr", "Maj", "Jun", "Jul", "Aug", "Sep", "Okt", "Nov", "Dec"};

static char\* swLongMonthList[] = {"Januari", "Februari", "Mars", "April", "Maj", "Juni", "Juli", "Augusit", "September", "Oktober", "November", "December"};

static char\* swMessageList[] = {"inga fel", "felaktigt functionsnummer",

"hittar ej filen", "hittar ej sokvagen", "inget handtag tillgangligt", "atkomst nekad",

"utanfor doman", "utanfor range", "felaktig multibyte-sekvens",

"fel vid oppning", "fel vid flushing", "fel vid stangning", "fel oppningslage",

"fel vid skrivning", "fel vid lasning", "fel vid sokning",

"fel vid telling", "fel vid borttagning av fil", "fel vid namnbyte av fil"};

static struct lconv sw\_EN\_utf8 = {1, 2, swShortDayList, swLongDayList, enShortMonthList, swLongMonthList,

"abcdefghijklmnopqrstuvwxyz", "ABCDEFGHIJKLMNOPQRSTUVWXYZ", swMessageList};

static struct \_s {

char\* name;

struct lconv\* localePtr;

} sArray[] = {{"", &en\_US\_utf8}, {"C", &en\_US\_utf8}, {"US", &en\_US\_utf8}, {"SE", &sw\_EN\_utf8}};

static int sSize = (sizeof sArray) / (sizeof sArray[0]);

static struct \_s\* g\_currStructPtr = &sArray[0];

#define PRINT(x,y) { printf(#x " = <%" #y ">\n", (x)); }

char\* setlocale(int /\*flag\*/, char\* newName) {

int index;

char \*oldName = (g\_currStructPtr != NULL) ? g\_currStructPtr->name : NULL;

g\_currStructPtr = NULL;

if (newName != NULL) {

for (index = 0; index < sSize; ++index) {

if (strcmp(newName, sArray[index].name) == 0) {

g\_currStructPtr = &sArray[index];

break;

}

}

}

return oldName;

}

struct lconv\* localeconv(void) {

return (g\_currStructPtr != NULL) ? g\_currStructPtr->localePtr : NULL;

}

## Character Types

CType.h

#ifndef \_\_CTYPE\_H\_\_

#define \_\_CTYPE\_H\_\_

extern int islower(int c);

extern int isupper(int c);

extern int isalpha(int c);

extern int isdigit(int c);

extern int isalnum(int c);

extern int isxdigit(int c);

extern int isgraph(int c);

extern int isprint(int c);

extern int ispunct(int c);

extern int iscntrl(int c);

extern int isspace(int c);

extern int tolower(int c);

extern int toupper(int c);

#endif

CType.c

#include <CType.h>

#include <StdIO.h>

#include <Locale.h>

#include <String.h>

#include <StdDef.h>

int islower(int c) {

struct lconv\* localeConvPtr = localeconv();

if (localeConvPtr != NULL) {

return (strchr(localeConvPtr->lowerCase, c) != NULL);

}

else {

return ((c >= 'a') && (c <= 'z'));

}

}

int isupper(int c) {

struct lconv\* localeConvPtr = localeconv();

if (localeConvPtr != NULL) {

return (strchr(localeConvPtr->upperCase, c) != NULL);

}

else {

return ((c >= 'A') && (c <= 'Z'));

}

}

int isalpha(int c) {

return islower(c) || isupper(c);

}

int isdigit(int c) {

return (c >= '0') && (c <= '9');

}

int isalnum(int c) {

return isalpha(c) || isdigit(c);

}

int isxdigit(int c) {

return isdigit(c) || ((c >= 'a') && (c <= 'f'))

|| ((c >= 'A') && (c <= 'F'));

}

int isgraph(int c) {

return (c >= 32) && (c <= 126);

}

int isprint(int c) {

return isgraph(c) && (c != ' ');

}

int ispunct(int c) {

return isgraph(c) && !isalnum(c);

}

int iscntrl(int c) {

return !isprint(c);

}

int isspace(int c) {

return (c == ' ') || (c == '\f') || (c == '\n') ||

(c == '\r') || (c == '\t') || (c == '\v');

}

int tolowerX(int c) {

return (isupper(c)) ? (c + 32) : c;

}

int tolower(int c) {

if (isupper(c)) {

struct lconv\* localeConvPtr = localeconv();

if (localeConvPtr != NULL) {

char \*lowerCase = localeConvPtr->lowerCase, \*upperCase = localeConvPtr->upperCase;

int index = (strchr(upperCase, c) - upperCase);

return ((int) lowerCase[index]);

}

else {

return (c + 32);

}

}

else {

return c;

}

}

int toupperX(int c) {

return (islower(c)) ? (c - 32) : c;

}

int toupper(int c) {

if (islower(c)) {

struct lconv\* localeConvPtr = localeconv();

if (localeConvPtr != NULL) {

char \*lowerCase = localeConvPtr->lowerCase,

\*upperCase = localeConvPtr->upperCase;

int index = (strchr(lowerCase, c) - lowerCase);

return ((int) upperCase[index]);

}

else {

return (c - 32);

}

}

else {

return c;

}

}

## Strings

String.h

#ifndef \_\_STRING\_H\_\_

#define \_\_STRING\_H\_\_

#define size\_t int

extern char\* strcpy(char\* target, const char\* source);

extern char\* strncpy(char\* target, const char\* source, size\_t size);

extern char\* strcat(char\* target, const char\* source);

extern char\* strncat(char\* target, const char\* source, size\_t size);

extern int strcmp(const char\* left, const char\* right);

extern int strncmp(const char\* left, const char\* right, size\_t size);

extern char\* strchr(const char\* text, int i);

extern char\* strrchr(const char\* text, int i);

extern size\_t strspn(const char\* mainString, const char\* charSet);

extern size\_t strcspn(const char\* mainString, const char\* charSet);

extern char\* strpbrk(const char\* mainString, const char\* charSet);

extern char\* strstr(const char\* mainString, const char\* subString);

extern size\_t strlen(const char\* string);

extern char\* strerror(int error);

extern char\* strtok(char\* string, const char\* charSet);

extern void\* memcpy(void\* target, const void\* source, size\_t size);

extern void\* memmove(void\* target, const void\* source, size\_t size);

extern int memcmp(const void\* left, const void\* right, size\_t size);

extern void\* memchr(const void\* block, int i, size\_t size);

extern void\* memset(void\* block, int i, size\_t size);

#endif

String.c

#include <ErrNo.h>

#include <StdIO.h>

#include <StdDef.h>

#include <String.h>

#include <Locale.h>

char\* strcpy(char\* target, const char\* source) {

size\_t index;

for (index = 0; source[index] != '\0'; ++index) {

target[index] = source[index];

}

target[index] = '\0';

return target;

}

char\* strncpy(char\* target, const char\* source, size\_t size) {

size\_t index;

for (index = 0; (index < size) && (source[index] != '\0'); ++index) {

target[index] = source[index];

}

for (; index < size; ++index) {

target[index] = '\0';

}

return target;

}

char\* strcat(char\* target, const char\* source) {

size\_t index;

size\_t targetLength = strlen(target);

for (index = 0; source[index] != '\0'; ++index) {

target[targetLength + index] = source[index];

}

target[targetLength + index] = '\0';

return target;

}

char\* strncat(char\* target, const char\* source, size\_t size) {

size\_t index;

size\_t targetLength = strlen(target);

for (index = 0; (index < (size - 1)) && (source[index] != '\0'); ++index) {

target[targetLength + index] = source[index];

}

target[targetLength + size - 1] = '\0';

return target;

}

int strcmp(const char\* left, const char\* right) {

size\_t index;

for (index = 0; TRUE; ++index) {

if ((left[index] == '\0') && (right[index] == '\0')) {

return 0;

}

else if (left[index] == '\0') {

return -1;

}

else if (right[index] == '\0') {

return 1;

}

else if (left[index] < right[index]) {

return -1;

}

else if (left[index] > right[index]) {

return 1;

}

}

}

int strncmp(const char\* left, const char\* right, size\_t size) {

size\_t index;

for (index = 0; index < size; ++index) {

if ((left[index] == '\0') && (right[index] == '\0')) {

return 0;

}

else if (left[index] == '\0') {

return -1;

}

else if (right[index] == '\0') {

return 1;

}

else if (left[index] < right[index]) {

return -1;

}

else if (left[index] > right[index]) {

return 1;

}

}

return 0;

}

char\* strchr(const char\* text, int i) {

size\_t index;

char c = (char) i;

for (index = 0; text[index] != '\0'; ++index) {

if (text[index] == c) {

return &text[index];

}

}

return NULL;

}

char\* strrchr(const char\* text, int i) {

size\_t index;

char\* result = NULL;

char c = (char) i;

for (index = 0; text[index] != '\0'; ++index) {

if (text[index] == c) {

result = &text[index];

}

}

return result;

}

size\_t strspn(const char\* mainString, const char\* charSet) {

size\_t index;

for (index = 0; mainString[index] != '\0'; ++index) {

if (strchr(charSet, mainString[index]) == NULL) {

return index;

}

}

return -1;

}

size\_t strcspn(const char\* mainString, const char\* charSet) {

size\_t index;

for (index = 0; mainString[index] != '\0'; ++index) {

if (strchr(charSet, mainString[index]) != NULL) {

return index;

}

}

return -1;

}

char\* strpbrk(const char\* mainString, const char\* charSet) {

size\_t index;

for (index = 0; mainString[index] != '\0'; ++index) {

if (strchr(charSet, mainString[index]) != NULL) {

return &mainString[index];

}

}

return NULL;

}

char\* strstr(const char\* mainString, const char\* subString) {

size\_t index;

for (index = 0; mainString[index] != '\0'; ++index) {

if (strcmp(mainString + index, subString) == 0) {

return &mainString[index];

}

}

return NULL;

}

size\_t strlen(const char\* string) {

int index;

for (index = 0; string[index] != '\0'; ++index) {

// Empty.

}

return index;

}

extern char\* enMessageList[];

char\* strerror(int errno) {

struct lconv\* localeConvPtr = localeconv();

char\*\* messageList = (localeConvPtr != NULL) ?

localeConvPtr->messageList : NULL;

messageList = (messageList != NULL) ? messageList : enMessageList;

return messageList[errno];

}

char\* token = NULL;

char\* strtok(char\* string, const char\* charSet) {

int index;

char\* tokenStart;

if (string != NULL) {

if (string[0] == '\0') {

return NULL;

}

for (index = 0; string[index] != '\0'; ++index) {

if (strchr(charSet, string[index]) != NULL) {

string[index] = '\0';

token = &string[index + 1];

return string;

}

}

token = &string[index];

return string;

}

else if (token == NULL) {

return NULL;

}

else {

if (token[0] == '\0') {

return NULL;

}

for (index = 0; token[index] != '\0'; ++index) {

if (strchr(charSet, token[index]) != NULL) {

char\* tokenStart2 = token;

token[index] = '\0';

token = &token[index + 1];

return tokenStart2;

}

}

tokenStart = token;

token = &token[index];

return tokenStart;

}

}

void\* memcpy(void\* target, const void\* source, size\_t size) {

char\* charTarget = (char\*) target;

const char\* charSource = (const char\*) source;

size\_t index;

for (index = 0; index < size; ++index) {

charTarget[index] = charSource[index];

}

return (void\*) target;

}

void\* memmove(void\* target, const void\* source, size\_t size) {

char\* charTarget = (char\*) target;

const char\* charSource = (const char\*) source;

size\_t index;

if (source < target) {

for (index = (size - 1); index >= 0; --index) {

charTarget[index] = charSource[index];

}

}

else {

for (index = 0; index < size; ++index) {

charTarget[index] = charSource[index];

}

}

return (void\*) target;

}

int memcmp(const void\* left, const void\* right, size\_t size) {

const char\* charLeft = (const char\*) left;

const char\* charRight = (const char\*) right;

size\_t index;

for (index = 0; index < size; ++index) {

if (charLeft[index] < charRight[index]) {

return -1;

}

else if (charLeft[index] > charRight[index]) {

return 1;

}

}

return 0;

}

void\* memchr(const void\* block, int i, size\_t size) {

size\_t index;

char\* charBlock = (char\*) block;

char c = (char) i;

for (index = 0; index < size; ++index) {

if (charBlock[index] == c) {

return (void\*) &charBlock[index];

}

}

return NULL;

}

void\* memset(void\* block, int i, size\_t size) {

char\* charBlock = (char\*) block;

char c = (char) i;

size\_t index;

for (index = 0; index < size; ++index) {

charBlock[index] = c;

}

return block;

}

## Long Jumps

SetJmp.h

typedef int jmp\_buf[3];

int setjmp(jmp\_buf env);

void longjmp(jmp\_buf env, int value);

SetJmp.c

#include <SetJmp.h>

#include <StdIO.h>

int setjmp(jmp\_buf buf) {

int\* bp\_pointer;

store\_register(bp, bp\_pointer);

buf[0] = bp\_pointer[0]; // return address

buf[1] = bp\_pointer[1]; // normal stack

buf[2] = bp\_pointer[2]; // ellipse stack

return 0;

}

void longjmp(jmp\_buf buf, int return\_value) {

load\_register(bx, return\_value);

load\_register(cx, buf[0]); // return address

load\_register(bp, buf[1]); // normal stack

load\_register(di, buf[2]); // ellipse stack

jump\_register(cx);

}

## Mathematical Functions

Math.h

#ifndef \_\_MATH\_H\_\_

#define \_\_MATH\_H\_\_

#define PI 3.1415926535897932384626433

#define E 2.7182818284590452353602874

#define EPSILON 1e-9

extern double exp(double value);

extern double log(double value);

extern double log10(double value);

extern int log10\_int(double value);

extern double pow(double base, double exponent);

extern double pow\_int(double base, int exponent);

extern double ldexp(double value, int exponent);

extern double frexp(double value, int\* exponent);

extern double sin(double value);

extern double cos(double value);

extern double tan(double value);

extern double sinh(double value);

extern double cosh(double value);

extern double tanh(double value);

extern double sqrt(double value);

extern double asin(double value);

extern double acos(double value);

extern double atan(double value);

extern double atan2(double num, double denum);

extern double floor(double value);

extern double ceil(double value);

extern double round(double value);

extern double fabs(double value);

extern double modf(double value, double\* integralPart);

extern double fmod(double num, double denum);

#endif

Math.c

#include <Math.h>

#include <ErrNo.h>

#include <StdIO.h>

#include <StdDef.h>

#include <StdlIB.h>

double exp(double x) {

double n = 0, faculty = 1, power = 1, term, sum = 0;

do {

term = power / faculty;

sum += term;

power \*= x;

faculty \*= ++n;

} while (fabs(term) >= EPSILON);

return sum;

}

static double logX(double x) { // [1/E,1]

double n = 1, plusMinusOne = 1, x\_minus\_1 = x - 1,

term, sum = 0, power = x\_minus\_1;

do {

term = plusMinusOne \* (power / n++);

sum += term;

power \*= x\_minus\_1;

plusMinusOne \*= -1.0;

} while (fabs(term) > EPSILON);

return sum;

}

double log(double x) {

if (x <= 0) {

errno = EDOM;

return 0;

}

else if (x == 1) {

return 0;

}

else if (x > 1) {

int expo = 0;

while (x > 1) {

x /= E;

++expo;

}

return (logX(x) + expo);

}

else {

int expo = 0;

while (x < 1) {

x \*= E;

++expo;

}

return (logX(x / E) - (expo - 1));

}

}

double log10(double x) {

return 0.4342944820 \* log(x);

}

int log10\_int(double x) {

if (x > 0) {

if (x == 1) {

return 0;

}

else if (x > 1) {

int count = 0;

while (x > 1) {

x /= 10;

++count;

}

return (count - 1);

}

else {

int count = 0;

while (x < 1) {

x \*= 10;

++count;

}

return -count;

}

}

else {

errno = EDOM;

return 0;

}

}

double pow\_int(double x, int y) {

BOOL minus = FALSE;

if (y < 0) {

y = -y;

minus = TRUE;

}

double product = 1;

int index;

for (index = 0; index < y; ++index) {

product \*= x;

}

return minus ? (1 / product) : product;

}

double pow(double x, double y) {

if (x > 0) {

return exp(y \* log(x));

}

else {

errno = EDOM;

exit(-1);

return 0;

}

}

double ldexp(double x, int n) {

return x \* pow(2, (double)n);

}

double frexp(double x, int\* p) {

if (p != NULL) {

if (x == 0) {

\*p = 0;

return 0;

}

else {

\*p = (int) (log(fabs(x)) / log(2)) + 1;

double quotient = fabs(x) / pow(2, \*p);

return (x < 0) ? -quotient : quotient;

}

}

else {

if (x == 0) {

return 0;

}

else {

int n = (int)(log(fabs(x)) / log(2)) + 1;

double a = fabs(x) / pow(2, n);

return (x < 0) ? -a : a;

}

}

}

double sin(double x) {

double n = 0, plusMinusOne = 1, faculty = 1, power = x, term, sum = 0;

do {

term = plusMinusOne \* (power / faculty);

sum += term;

plusMinusOne \*= -1;

power \*= x \* x;

faculty \*= (n + 2) \* (n + 3);

n += 2;

} while (fabs(term) >= EPSILON);

return sum;

}

double cos(double x) {

double n = 0, plusMinusOne = 1, faculty = 1, power = 1, term, sum = 0;

do {

term = plusMinusOne \* (power / faculty);

sum += term;

plusMinusOne \*= -1;

power \*= x \* x;

faculty \*= (n + 1) \* (n + 2);

n += 2;

} while (fabs(term) >= EPSILON);

return sum;

}

double tan(double x) {

double cos\_value = cos(x);

if (cos\_value != 0) {

return sin(x) / cos\_value;

}

else {

errno = EDOM;

return 0;

}

}

double sinh(double x) {

return (exp(x) - exp(-x)) / 2;

}

double cosh(double x) {

return (exp(x) + exp(-x)) / 2;

}

double tanh(double x) {

return sinh(x) / cosh(x);

}

double sqrt(double v) {

if (v >= 0) {

double x\_nplus1 = 1, x;

do {

x = x\_nplus1;

x\_nplus1 = (x + (v / x)) / 2;

} while (fabs(x\_nplus1 - x) >= EPSILON);

return x\_nplus1;

}

else {

errno = EDOM;

return 0;

}

}

double asin(double v) {

if (v == 1) {

return PI / 2;

}

else if (v == -1) {

return -PI / 2;

}

else if (fabs(v) <= 1) {

double x\_nplus1 = 1, x;

do {

x = x\_nplus1;

x\_nplus1 = x - tan(x) + (v / cos(x));

} while (fabs(x\_nplus1 - x) >= EPSILON);

return x\_nplus1;

}

else {

errno = EDOM;

return 0;

}

}

double acos(double v) {

if (v == 1) {

return 0;

}

else if (fabs(v) <= 1) {

double x\_nplus1 = 1, x;

do {

x = x\_nplus1;

x\_nplus1 = x + (1 / tan(x)) - (v / sin(x));

} while (fabs(x\_nplus1 - x) >= EPSILON);

return x\_nplus1;

}

else {

errno = EDOM;

return 0;

}

}

double square(double x) {

return x \* x;

}

double atan(double v) {

if (v == 0) {

return 0;

}

else if (v == 1) {

return PI / 4;

}

else if (v == -1) {

return -PI / 4;

}

else if (fabs(v) < 0.5) {

int sign = 1, denominator = 1; // count = 0;

double product = v, term, sum = 0;

do {

term = sign \* product / denominator;

sum += term;

sign = -sign;

product \*= v \* v;

denominator += 2;

} while (fabs(term) >= EPSILON);

return sum;

}

else if (fabs(v) < 1) {

double x\_nplus1 = 1, x;

do {

x = x\_nplus1;

x\_nplus1 = x - ((tan(x) - v) \* square(cos(2 \* x) + 1)) / 2;

} while (fabs(x\_nplus1 - x) >= EPSILON);

return x\_nplus1;

}

else {

errno = EDOM;

return 0;

}

}

double atanX(double v) {

if (v == 0) {

return 0;

}

else if (fabs(v) <= 1) {

double x\_nplus1 = 1, x;

do {

x = x\_nplus1;

x\_nplus1 = x - ((tan(x) - v) \* square(cos(2 \* x) + 1)) / 2;

} while (fabs(x\_nplus1 - x) >= EPSILON);

return x\_nplus1;

}

else {

errno = EDOM;

return 0;

}

}

double atan2(double num, double denum) {

if (denum > 0) {

return atan(num / denum);

}

else if ((num >= 0) && (denum < 0)) {

return PI + atan(num / denum);

}

else if ((num < 0) && (denum < 0)) {

return (-PI) + atan(num / denum);

}

else if ((num > 0) && (denum == 0)) {

return PI / 2;

}

else if ((num < 0) && (denum == 0)) {

return (-PI) / 2;

}

else {

errno = EDOM;

return 0;

}

}

double floor(double x) {

if (x < 0) {

return -ceil(-x);

}

return (double) (long) x;

}

double ceil(double x) {

if (x < 0) {

return -floor(-x);

}

return (double) (long) (x + 0.999999999999);

}

double round(double x) {

return (double) (long) ((x < 0) ? (x - 0.5) : (x + 0.5));

}

double fabs(double x) {

return (x < 0) ? -x : x;

}

double modf(double x, double\* p) {

double integral = (double) (long) fabs(x);

if (p != NULL) {

\*p = (x > 0) ? (fabs(x) - integral) : -(fabs(x) - integral);

}

return (x > 0) ? integral : -integral;

}

double fmod(double x, double y) {

if (y != 0) {

double quotient = x / y;

double remainder = fabs(quotient - (double) (long) quotient);

return (x > 0) ? remainder : -remainder;

}

else {

errno = EDOM;

return 0;

}

}

## Standard Input/Output

StdIO.h

#ifndef \_\_STDIO\_H\_\_

#define \_\_STDIO\_H\_\_

#include <Math.h>

#include <CType.h>

#include <StdArg.h>

#include <StdDef.h>

#include <File.h>

#include <Temp.h>

#include <Scanf.h>

#include <Printf.h>

#endif

### Printing

Printf.h

#ifndef \_\_PRINTF\_H\_\_

#define \_\_PRINTF\_H\_\_

#define DEVICE 0

#define STRING 1

extern int g\_outStatus, g\_charCount;

extern void\* g\_outDevice;

int putc(int c, FILE\* stream);

int fputc(int c, FILE\* stream);

int putchar(int c);

void printChar(char c);

void printString(char\* s, int precision);

void printIntRec(int intValue);

void printInt(int intValue, BOOL plus, BOOL space);

void printDoublePlain(double doubleValue, BOOL plus, BOOL space,

BOOL grid, int precision);

void printLongDoublePlain(long double doubleValue, BOOL plus,

BOOL space, BOOL grid, int precision);

void printInt(int intValue, BOOL plus, BOOL space);

void printLongInt(long longIntValue, BOOL plus, BOOL space);

void printLongDoubleFraction(long double longDoubleValue,

BOOL grid, int precision);

void printLongDoublePlain(long double longDoubleValue, BOOL plus,

BOOL space, BOOL grid, int precision);

int printFormat(char\* format, va\_list arg\_list);

int printf(char\* format, ...);

int vprintf(char\* format, va\_list arg\_list);

int fprintf(FILE\* outStream, char\* format, ...);

int vfprintf(FILE\* outStream, char\* format, va\_list arg\_list);

int sprintf(char\* outString, char\* format, ...);

int vsprintf(char\* outString, char\* format, va\_list arg\_list);

#endif

Printf.c

#include <Math.h>

#include <CType.h>

#include <StdIO.h>

#include <StdDef.h>

#include <StdArg.h>

#include <StdLib.h>

#include <Scanf.h>

#include <Printf.h>

#define DEVICE 0

#define STRING 1

#define BLANK 2

#define PRINT(x,y) { printf(#x " = <%" #y ">\n", (x)); }

int g\_outStatus, g\_outChars;

void\* g\_outDevice;

#define DEFAULT\_PRECISION 10

int putc(int i, FILE\* stream) {

g\_outStatus = DEVICE;

g\_outDevice = (void\*) stream;

printChar((char) i);

return 1;

}

int fputc(int i, FILE\* stream) {

g\_outStatus = DEVICE;

g\_outDevice = (void\*) stream;

printChar((char) i);

return 1;

}

int putchar(int i) {

g\_outStatus = DEVICE;

g\_outDevice = (void\*) stdout;

printChar((char) i);

return 1;

}

void printChar(char c) {

int handle;

char\* outString;

switch (g\_outStatus) {

case DEVICE: {

FILE\* stream = (FILE\*) g\_outDevice;

handle = stream->handle;

load\_register(ah, 0x40s);

load\_register(bx, handle);

load\_register(cx, 1);

load\_register(dx, &c);

interrupt(0x21s);

++g\_outChars;

break;

}

case STRING: {

outString = (char\*) g\_outDevice;

outString[g\_outChars++] = c;

}

break;

case BLANK:

g\_outChars++;

break;

}

}

void printString(char\* s, int precision) {

if (s != NULL) {

int index;

if (precision == 0) {

for (index = 0; s[index] != '\0'; ++index) {

printChar(s[index]);

}

}

else {

for (index = 0; (precision-- > 0) && (s[index] != '\0'); ++index) {

printChar(s[index]);

}

}

}

else {

printChar('<');

printChar('N');

printChar('U');

printChar('L');

printChar('L');

printChar('>');

}

}

void printLongIntRec(long longValue) {

if (longValue != 0) {

int digit = (int) (longValue % 10L);

printLongIntRec(longValue / 10L);

printChar((char) (digit + '0'));

}

}

void printLongInt(long longValue, BOOL plus, BOOL space) {

if (longValue < 0L) {

longValue = -longValue;

printChar('-');

}

else if (space) {

printChar(' ');

}

else if (plus) {

printChar('+');

}

if (longValue == 0L) {

printChar('0');

}

else {

printLongIntRec(longValue);

}

}

char digitToChar(int digit, BOOL capital) {

if (digit < 10) {

return ((char) ('0' + digit));

}

else if (capital) {

return ((char) ('A' + (digit - 10)));

}

else {

return ((char) ('a' + (digit - 10)));

}

}

void printUnsignedLongRec(unsigned long unsignedValue, unsigned long base, BOOL capital) {

if (unsignedValue > 0ul) {

int digit = (int) (unsignedValue % base);

printUnsignedLongRec(unsignedValue / base, base, capital);

char c = digitToChar(digit, capital);

printChar(c);

}

}

void printUnsignedLong(unsigned long unsignedValue, BOOL plus, BOOL space,

BOOL grid, unsigned long base, BOOL capital) {

if (plus) {

printChar('+');

}

if (space) {

printChar(' ');

}

if (grid) {

if (base == 8ul) {

printChar('0');

}

if (base == 16ul) {

printChar('0');

printChar(capital ? 'X' : 'x');

}

}

if (unsignedValue == 0ul) {

printChar('0');

}

else {

printUnsignedLongRec(unsignedValue, base, capital);

}

}

void printLongDoubleFraction(long double longDoubleValue, BOOL grid, int precision) {

longDoubleValue -= (long) longDoubleValue;

if (precision == 0) {

precision = DEFAULT\_PRECISION;

}

if (grid || (precision > 0)) {

printChar('.');

}

while (precision-- > 0) {

long double longDoubleValue10 = 10.0L \* longDoubleValue;

int digitValue = (int) longDoubleValue10;

char c = (char) (digitValue + '0');

printChar(c);

longDoubleValue = longDoubleValue10 - ((long double) digitValue);

}

}

void printLongDoublePlain(long double longDoubleValue, BOOL plus, BOOL space, BOOL grid, int precision) {

if (longDoubleValue < 0.0L) {

printChar('-');

longDoubleValue = -longDoubleValue;

plus = FALSE;

space = FALSE;

}

long longValue = (long) longDoubleValue;

printLongInt(longValue, plus, space);

longDoubleValue -= (long double) longValue;

printLongDoubleFraction(longDoubleValue, grid, precision);

}

void printLongDoubleExpo(long double value, BOOL plus, BOOL space,

BOOL grid, int precision, BOOL capital) {

if (value == 0.0L) {

printChar('0');

printLongDoubleFraction(0.0L, precision, grid);

printChar(capital ? 'E' : 'e');

printChar('0');

}

else {

if (value < 0.0L) {

printChar('-');

value = -value;

}

int expo = (int) log10(value);

value /= pow(10, expo);

printLongDoublePlain(value, plus, space, grid, precision);

printChar(capital ? 'E' : 'e');

printLongInt(expo, TRUE, FALSE);

}

}

va\_list checkWidthAndPrecision(va\_list arg\_list, int\* widthPtr, int\* precisionPtr) {

if ((widthPtr != NULL) && (\*widthPtr == -1)) {

\*widthPtr = va\_arg(arg\_list, int);

}

if ((precisionPtr != NULL) && (\*precisionPtr == -1)) {

\*precisionPtr = va\_arg(arg\_list, int);

}

return arg\_list;

}

va\_list printArgument(char\* format, va\_list arg\_list, BOOL plus, BOOL space,

BOOL grid, int\* widthPtr, int precision, BOOL shortInt,

BOOL longInt, BOOL longDouble, BOOL sign, BOOL\* negativePtr) {

char c = format[0], charValue;

int /\*logValue, \*/ \*intPtr;

long double longDoubleValue;

void\* ptrValue;

switch (c) {

case 'i':

case 'd': {

long longValue;

if (shortInt) {

longValue = (long) (short) va\_arg(arg\_list, int);

}

else if (longInt) {

longValue = va\_arg(arg\_list, long);

}

else {

longValue = (long) va\_arg(arg\_list, int);

}

if (negativePtr != NULL) {

\*negativePtr = (longValue < 0);

}

if (!sign) {

longValue = labs(longValue);

}

arg\_list = checkWidthAndPrecision(arg\_list, widthPtr, &precision);

printLongInt(longValue, plus, space);

}

break;

case 'c':

charValue = (char) va\_arg(arg\_list, int);

arg\_list = checkWidthAndPrecision(arg\_list, widthPtr, &precision);

printChar(charValue);

break;

case 's': {

char\* stringValue = va\_arg(arg\_list, char\*);

arg\_list = checkWidthAndPrecision(arg\_list, widthPtr, &precision);

printString(stringValue, precision);

}

break;

case 'u':

case 'o':

case 'b':

case 'x':

case 'X': {

unsigned long base = ((c == 'u') ? 10ul : ((c == 'o') ? 8ul : ((c == 'b') ? 2ul : 16ul)));

unsigned long value;

if (shortInt) {

value = (unsigned long) (unsigned short) va\_arg(arg\_list, unsigned int);

}

else if (longInt) {

value = va\_arg(arg\_list, unsigned long);

}

else {

value = (unsigned long) va\_arg(arg\_list, unsigned int);

}

arg\_list = checkWidthAndPrecision(arg\_list, widthPtr, &precision);

printUnsignedLong(value, plus, space, grid, base, isupper(c));

}

break;

case 'f':

case 'e':

case 'E':

case 'g':

case 'G':

if (longDouble) {

longDoubleValue = va\_arg(arg\_list, long double);

}

else {

longDoubleValue = (long double) va\_arg(arg\_list, double);

}

if (negativePtr != NULL) {

\*negativePtr = (longDoubleValue < 0);

}

if (!sign) {

longDoubleValue = fabs(longDoubleValue);

}

arg\_list = checkWidthAndPrecision(arg\_list, widthPtr, &precision);

if (c == 'f') {

printLongDoublePlain(longDoubleValue, plus, space, grid, precision);

}

else if (tolower(c) == 'e') {

printLongDoubleExpo(longDoubleValue, plus, space, grid, precision, isupper(c));

}

else {

int expo = (int) log10(fabs(longDoubleValue));

if ((expo >= -3) && (expo <= 2)) {

printLongDoublePlain(longDoubleValue, plus, space, grid, precision);

}

else {

printLongDoubleExpo(longDoubleValue, plus, space, grid, precision, isupper(c));

}

}

break;

case 'p':

ptrValue = va\_arg(arg\_list, void\*);

arg\_list = checkWidthAndPrecision(arg\_list, widthPtr, &precision);

printUnsignedLong((unsigned long) ptrValue, FALSE, FALSE, FALSE, 10u, FALSE);

break;

case 'n':

ptrValue = va\_arg(arg\_list, void\*);

intPtr = va\_arg(arg\_list, int\*);

arg\_list = checkWidthAndPrecision(arg\_list, widthPtr, &precision);

\*intPtr = g\_outChars;

break;

case '%':

//printf("<X>");

arg\_list = checkWidthAndPrecision(arg\_list, widthPtr, &precision);

printChar('%');

break;

}

return arg\_list;

}

int printFormat(char\* format, va\_list arg\_list) {

int index, count = 0, width = 0, precision = 0;

BOOL percent = FALSE, plus = FALSE, minus = FALSE, space = FALSE,

zero = FALSE, grid = FALSE, widthStar = FALSE,

period = FALSE, precisionStar = FALSE,

shortInt = FALSE, longInt = FALSE, longDouble = FALSE;

g\_outChars = 0;

for (index = 0; format[index] != '\0'; ++index) {

char c = format[index];

if (percent) {

switch (c) {

case '+':

plus = TRUE;

break;

case '-':

minus = TRUE;

break;

case ' ':

space = TRUE;

break;

case '0':

zero = TRUE;

break;

case '#':

grid = TRUE;

break;

case '.':

period = TRUE;

break;

case '\*':

if (!period) {

width = -1;

}

else {

precision = -1;

}

break;

case 'h':

shortInt = TRUE;

break;

case 'l':

longInt = TRUE;

break;

case 'L':

longDouble = TRUE;

break;

case 'i':

case 'd':

case 'u':

case 'b':

case 'o':

case 'x':

case 'X':

case 'c':

case 's':

case 'f':

case 'e':

case 'E':

case 'g':

case 'G':

case 'p':

case 'n':

case '%': {

if (minus) {

int startChars = g\_outChars;

arg\_list = printArgument(&format[index], arg\_list, plus, space, grid, &width,

precision, shortInt, longInt, longDouble, TRUE, NULL);

int field = g\_outChars - startChars; // fill after

while (field++ < width) {

printChar(' ');

}

}

else if (zero) {

int startChars = g\_outChars, oldOutStatus = g\_outStatus;

BOOL negative = FALSE;

g\_outStatus = BLANK;

printArgument(&format[index], arg\_list, FALSE, FALSE, grid, &width,

precision, shortInt, longInt, longDouble, FALSE, &negative);

g\_outStatus = oldOutStatus;

int field = g\_outChars - startChars;

g\_outChars = startChars;

if (negative) {

printChar('-');

++field;

}

else if (plus) {

printChar('+');

++field;

}

else if (space) {

printChar(' ');

++field;

}

while (field++ < width) {

printChar('0');

}

arg\_list = printArgument(&format[index], arg\_list, FALSE, FALSE, grid, NULL,

precision, shortInt, longInt, longDouble, FALSE, NULL);

}

else {

int startChars = g\_outChars, oldOutStatus = g\_outStatus;

g\_outStatus = BLANK;

printArgument(&format[index], arg\_list, plus, space, grid, &width,

precision, shortInt, longInt, longDouble, TRUE, NULL);

g\_outStatus = oldOutStatus;

int field = g\_outChars - startChars;

g\_outChars = startChars;

while (field++ < width) {

printChar(' ');

}

arg\_list = printArgument(&format[index], arg\_list, plus, space, grid, NULL,

precision, shortInt, longInt, longDouble, TRUE, NULL);

}

percent = FALSE;

}

break;

default: {

int value = 0;

while (isdigit(c)) {

value = (10 \* value) + (c - '0');

c = format[++index];

}

--index;

if (!period) {

width = value;

}

else {

precision = value;

}

}

break;

}

}

else {

if (c == '%') {

percent = TRUE;

plus = FALSE;

minus = FALSE;

space = FALSE;

zero = FALSE;

grid = FALSE;

widthStar = FALSE;

period = FALSE;

precisionStar = FALSE;

shortInt = FALSE;

longInt = FALSE;

longDouble = FALSE;

width = 0;

precision = 0;

}

else {

printChar(c);

}

}

}

if (g\_outStatus == STRING) {

char\* outString = (char\*) g\_outDevice;

outString[g\_outChars] = '\0';

}

return g\_outChars;

}

int printf(char\* format, ...) {

va\_list arg\_list;

va\_start(arg\_list, format);

return vprintf(format, arg\_list);

}

int vprintf(char\* format, va\_list arg\_list) {

return vfprintf(stdout, format, arg\_list);

}

int fprintf(FILE\* outStream, char\* format, ...) {

va\_list arg\_list;

va\_start(arg\_list, format);

return vfprintf(outStream, format, arg\_list);

}

int vfprintf(FILE\* outStream, char\* format, va\_list arg\_list) {

g\_outStatus = DEVICE;

g\_outDevice = (void\*) outStream;

return printFormat(format, arg\_list);

}

int sprintf(char\* outString, char\* format, ...) {

va\_list arg\_list;

va\_start(arg\_list, format);

return vsprintf(outString, format, arg\_list);

}

int vsprintf(char\* outString, char\* format, va\_list arg\_list) {

g\_outStatus = STRING;

g\_outDevice = (void\*) outString;

return printFormat(format, arg\_list);

}

### Scanning

Scanf.h

#ifndef \_\_SCANF\_H\_\_

#define \_\_SCANF\_H\_\_

#define DEVICE 0

#define STRING 1

#define EOF -1

extern int g\_inStatus, g\_inChars;

extern void\* g\_inDevice;

char scanChar(void);

void unscanChar(char c);

void scanString(char\* string, int precision);

long scanLongInt(void);

unsigned long scanUnsignedLongInt(unsigned long base);

long double scanLongDouble(void);

int scanf(char\* format, ...);

int vscanf(char\* format, va\_list arg\_list);

int fscanf(FILE\* inStream, char\* format, ...);

int vfscanf(FILE\* inStream, char\* format, va\_list arg\_list);

int sscanf(char\* inString, char\* format, ...);

int vsscanf(char\* inString, char\* format, va\_list arg\_list);

#endif

Scanf.c

#include <Math.h>

#include <CType.h>

#include <StdIO.h>

#include <StdDef.h>

#include <StdArg.h>

#include <String.h>

#include <Scanf.h>

#include <Printf.h>

#define PRINT(x,y) { printf(#x " = <%" #y ">\n", (x)); }

int g\_inStatus, g\_inChars;

void\* g\_inDevice;

int g\_inCount;

char scanChar(void) {

char c = '\0';

FILE\* stream;

int handle;

char\* inString;

switch (g\_inStatus) {

case DEVICE:

stream = (FILE\*) g\_inDevice;

if ((g\_inDevice != stdin) && feof(stream)) {

return ((char) EOF);

}

else {

handle = stream->handle;

load\_register(ah, 0x3Fs);

load\_register(bx, handle);

load\_register(cx, 1);

load\_register(dx, &c);

interrupt(0x21s);

++g\_inChars;

return c;

}

case STRING:

inString = (char\*) g\_inDevice;

return inString[g\_inChars++];

default:

return '\0';

}

}

void unscanChar(char /\* c \*/) {

switch (g\_inStatus) {

case DEVICE:

/\* stream = (FILE\*) g\_inDevice;

handle = stream->handle;

load\_register(ah, 0x3F);

load\_register(bx, handle);

load\_register(cx, 1);

load\_register(dx, &c);

interrupt(0x21s);

\*/

--g\_inChars;

break;

case STRING:

--g\_inChars;

break;

}

}

void scanPattern(char\* string, char\* pattern, BOOL not) {

int index = 0;

char input = scanChar();

while (isspace(input)) {

input = scanChar();

}

if (string != NULL) {

while ((!not && strchr(pattern, input)) ||

(not && !strchr(pattern, input))) {

string[index++] = input;

input = scanChar();

}

string[index] = '\0';

}

else {

while ((!not && strchr(pattern, input)) ||

(not && !strchr(pattern, input))) {

input = scanChar();

}

}

}

void scanString(char\* string, int precision) {

int index = 0;

char input = scanChar();

BOOL found = FALSE;

while (isspace(input)) {

input = scanChar();

}

if (string != NULL) {

if (precision == 0) {

while (!isspace(input) && (input != EOF)) {

string[index++] = input;

input = scanChar();

found = TRUE;

++g\_inChars;

}

string[index] = '\0';

++g\_inChars;

}

else {

while ((precision-- > 0) && (!isspace(input) && (input != EOF))) {

string[index++] = input;

input = scanChar();

found = TRUE;

++g\_inChars;

}

if (precision > 0) {

string[index] = '\0';

++g\_inChars;

}

}

}

else {

if (precision == 0) {

while (!isspace(input) && (input != EOF)) {

input = scanChar();

found = TRUE;

++g\_inChars;

}

++g\_inChars;

}

else {

while ((precision-- > 0) && (!isspace(input) && (input != EOF))) {

input = scanChar();

found = TRUE;

++g\_inChars;

}

if (precision > 0) {

++g\_inChars;

}

}

}

if (found) {

++g\_inCount;

}

}

unsigned long digitToValue(char input) {

if (isdigit(input)) {

return (input - '0');

}

else if (islower(input)) {

return ((input - 'a') + 10ul);

}

else {

return ((input - 'A') + 10ul);

}

}

long scanLongInt(void) {

long longValue = 0l;

BOOL minus = FALSE, found = FALSE;

char input = scanChar();

while (isspace(input)) {

input = scanChar();

}

if (input == '+') {

input = scanChar();

}

else if (input == '-') {

minus = TRUE;

input = scanChar();

}

while (isdigit(input)) {

longValue = (10l \* longValue) + (input - '0');

input = scanChar();

found = TRUE;

}

if (minus) {

longValue = -longValue;

}

if (found) {

++g\_inCount;

}

unscanChar(input);

return longValue;

}

unsigned long scanUnsignedLongInt(unsigned long base) {

unsigned long unsignedLongValue = 0ul, digit;

char input = scanChar();

BOOL found = TRUE;

while (isspace(input)) {

input = scanChar();

}

if (input == '0') {

input = scanChar();

if (tolower(input) == 'x') {

base = (base == 0ul) ? 16ul : base;

input = scanChar();

}

else {

base = (base == 0ul) ? 8ul : base;

}

}

if (base == 0ul) {

base = 10ul;

}

while (isxdigit(input)) {

digit = digitToValue(input);

if (digit >= base) {

break;

}

unsignedLongValue = (unsignedLongValue \* base) + digit;

found = TRUE;

input = scanChar();

}

if (found) {

++g\_inCount;

}

unscanChar(input);

return unsignedLongValue;

}

long double scanLongDouble(void) {

BOOL minus = FALSE, found = FALSE;

long double value = 0.0L, factor = 1.0L;

char input = scanChar();

while (isspace(input)) {

input = scanChar();

}

if (input == '+') {

input = scanChar();

}

else if (input == '-') {

minus = TRUE;

input = scanChar();

}

while (isdigit(input)) {

value = (10.0L \* value) + ((long double) (input - '0'));

input = scanChar();

found = TRUE;

}

if (input == '.') {

input = scanChar();

while (isdigit(input)) {

factor /= 10.0L;

value += factor \* ((long double) (input - '0'));

input = scanChar();

found = TRUE;

}

}

if (tolower(input) == 'e') {

double exponent = (double) scanLongInt();

value \*= pow(10.0, exponent);

}

else {

unscanChar(input);

}

if (minus) {

value = -value;

}

if (found) {

++g\_inCount;

}

return value;

}

int scanFormat(char\* format, va\_list arg\_list) {

char c, \*charPtr;

BOOL percent = FALSE, shortInt = FALSE, longIntOrDouble = FALSE, longDouble = FALSE, star = FALSE;

long longValue, \*longPtr;

short\* shortPtr;

int index, \*intPtr, \*charsPtr;

unsigned long unsignedLongValue, \*unsignedLongPtr;

unsigned short\* unsignedShortPtr;

unsigned int\* unsignedIntPtr;

long double longDoubleValue;

g\_inCount = 0;

g\_inChars = 0;

for (index = 0; format[index] != '\0'; ++index) {

c = format[index];

if (percent) {

switch (c) {

case 'h':

shortInt = TRUE;

break;

case 'l':

longIntOrDouble = TRUE;

break;

case 'L':

longDouble = TRUE;

break;

case '\*':

star = TRUE;

break;

case 'c': {

char charValue = scanChar();

if (!star) {

charPtr = va\_arg(arg\_list, char\*);

\*charPtr = charValue;

}

percent = FALSE;

if (charValue != EOF) {

++g\_inCount;

}

}

break;

case 's':

if (!star) {

charPtr = va\_arg(arg\_list, char\*);

scanString(charPtr, 0);

}

else {

scanString(NULL, 0);

}

percent = FALSE;

break;

case 'i':

case 'd':

longValue = scanLongInt();

if (!star) {

if (shortInt) {

shortPtr = va\_arg(arg\_list, short\*);

\*shortPtr = (short) longValue;

}

else if (!longIntOrDouble) {

intPtr = va\_arg(arg\_list, int\*);

\*intPtr = (int) longValue;

}

else {

longPtr = va\_arg(arg\_list, long\*);

\*longPtr = longValue;

}

}

percent = FALSE;

break;

case 'o':

unsignedLongValue = scanUnsignedLongInt(8ul);

if (!star) {

if (shortInt) {

unsignedShortPtr = va\_arg(arg\_list, unsigned short\*);

\*unsignedShortPtr = (short) unsignedLongValue;

}

else if (!longIntOrDouble) {

unsignedIntPtr = va\_arg(arg\_list, unsigned int\*);

\*unsignedIntPtr = (int) unsignedLongValue;

}

else {

unsignedLongPtr = va\_arg(arg\_list, unsigned long\*);

\*unsignedLongPtr = unsignedLongValue;

}

}

percent = FALSE;

break;

case 'x':

unsignedLongValue = scanUnsignedLongInt(16ul);

if (!star) {

if (shortInt) {

unsignedShortPtr = va\_arg(arg\_list, unsigned short\*);

\*unsignedShortPtr = (short) unsignedLongValue;

}

else if (!longIntOrDouble) {

unsignedIntPtr = va\_arg(arg\_list, unsigned int\*);

\*unsignedIntPtr = (int) unsignedLongValue;

}

else {

unsignedLongPtr = va\_arg(arg\_list, unsigned long\*);

\*unsignedLongPtr = unsignedLongValue;

}

}

percent = FALSE;

break;

case 'u':

unsignedLongValue = scanUnsignedLongInt(0ul);

if (!star) {

if (shortInt) {

unsignedShortPtr = va\_arg(arg\_list, unsigned short\*);

\*unsignedShortPtr = (short) unsignedLongValue;

}

else if (!longIntOrDouble) {

unsignedIntPtr = va\_arg(arg\_list, unsigned int\*);

\*unsignedIntPtr = (int) unsignedLongValue;

}

else {

unsignedLongPtr = va\_arg(arg\_list, unsigned long\*);

\*unsignedLongPtr = unsignedLongValue;

}

}

percent = FALSE;

break;

case 'e':

case 'f':

case 'g':

longDoubleValue = scanLongDouble();

if (!star) {

if (longIntOrDouble) {

double\* doublePtr = va\_arg(arg\_list, double\*);

\*doublePtr = (double) longDoubleValue;

}

else if (longDouble) {

long double\* longDoublePtr = va\_arg(arg\_list, long double\*);

\*longDoublePtr = longDoubleValue;

}

else {

float\* floatPtr = va\_arg(arg\_list, float\*);

\*floatPtr = (float) longDoubleValue;

}

}

percent = FALSE;

break;

case '[': {

BOOL not = FALSE;

++index;

if (format[index] == '^') {

not = TRUE;

++index;

}

int startIndex = index;

while (format[index] != ']') {

++index;

}

format[index] = '\0';

if (!star) {

char\* string = va\_arg(arg\_list, char\*);

scanPattern(string, &format[startIndex], not);

}

else {

scanPattern(NULL, &format[startIndex], not);

}

}

break;

case 'n':

charsPtr = va\_arg(arg\_list, int\*);

\*charsPtr = g\_inChars;

percent = FALSE;

break;

}

}

else {

if (c == '%') {

percent = TRUE;

shortInt = FALSE;

longIntOrDouble = FALSE;

longDouble = FALSE;

star = FALSE;

}

}

}

return g\_inCount;

}

int scanf(char\* format, ...) {

va\_list arg\_list;

va\_start(arg\_list, format);

return vscanf(format, arg\_list);

}

int vscanf(char\* format, va\_list arg\_list) {

return vfscanf(stdin, format, arg\_list);

}

int fscanf(FILE\* inStream, char\* format, ...) {

va\_list arg\_list;

va\_start(arg\_list, format);

return vfscanf(inStream, format, arg\_list);

}

int vfscanf(FILE\* inStream, char\* format, va\_list arg\_list) {

g\_inStatus = DEVICE;

g\_inDevice = (void\*) inStream;

return scanFormat(format, arg\_list);

}

int sscanf(char\* inString, char\* format, ...) {

va\_list arg\_list;

va\_start(arg\_list, format);

return vsscanf(inString, format, arg\_list);

}

int vsscanf(char\* inString, char\* format, va\_list arg\_list) {

g\_inStatus = STRING;

g\_inDevice = (void\*) inString;

return scanFormat(format, arg\_list);

}

### File Handling

File.h

#ifndef \_\_FILE\_H\_\_

#define \_\_FILE\_H\_\_

#define FOPEN\_MAX 16

#define FILENAME\_MAX 16

#define fpos\_t int

#define CARRY\_FLAG 0x01u

#define EOF -1

typedef unsigned int UINTX;

typedef struct {

BOOL open;

unsigned int handle;

char name[FILENAME\_MAX], ungetc;

int errno;

unsigned long position, size;

BOOL temporary;

} FILE;

extern FILE\* stdin;

extern FILE\* stdout;

extern FILE\* stderr;

extern enum {EEXIST, ENOENT, EACCES};

extern enum {SEEK\_SET, SEEK\_CUR, SEEK\_END};

#define getc(stream) fgetc(stream)

BOOL fileexists(const char\* name);

FILE\* fopen(const char\* filename, const char\* mode);

FILE\* freopen(const char\* filename, const char\* mode, FILE\* stream);

int fflush(FILE\* stream);

int fclose(FILE\* stream);

int remove(const char\* name);

int rename(const char\* oldName, const char\* newName);

int setvbuf(FILE\* stream, char\* buffer, int mode, size\_t size);

void setbuf(FILE\* stream, char\* buffer);

int fgetc(FILE\* stream);

char\* fgets(char\* s, int n, FILE\* stream);

int fputc(int i, FILE\* stream);

int fputs(const char\* s, FILE\* stream);

int getchar(void);

char\* gets(char\* s);

int putchar(int c);

int puts(const char\* s);

int ungetc(int c, FILE\* stream);

size\_t fread(void\* ptr, size\_t size, size\_t nobj, FILE\* stream);

size\_t fwrite(const void\* ptr, size\_t size, size\_t nobj, FILE\* stream);

long fseek(FILE\* stream, long offset, int origin);

long ftell(FILE\* stream);

void rewind(FILE\* stream);

int fgetpos(FILE\* stream, fpos\_t\* ptr);

int fsetpos(FILE\* stream, const fpos\_t\* ptr);

void clearerr(FILE\* stream);

BOOL feof(FILE\* stream);

int ferror(FILE\* stream);

void perror(const char\* s);

#endif

File.c

#include <StdIO.h>

#include <ErrNo.h>

#include <Locale.h>

#include <String.h>

FILE g\_fileArray[FOPEN\_MAX] = {{TRUE, 0}, {TRUE, 1}, {TRUE, 2}};

FILE\* stdin = &g\_fileArray[0];

FILE\* stdout = &g\_fileArray[1];

FILE\* stderr = &g\_fileArray[2];

static enum { EEXIST, ENOENT, EACCES };

static enum { SEEK\_SET = 0, SEEK\_CUR = 1, SEEK\_END = 2 };

static enum {READ = 0x40, WRITE = 0x41, READ\_WRITE = 0x42};

#define MAX(a,b) (((a) > (b)) ? (a) : (b))

#define FILE\_NOT\_FOUND 0x02

#define PRINT(x,y) { printf(#x " = <%" #y ">\n", (x)); }

long filesize(int handle) {

load\_register(ah, 0x42s);

load\_register(bx, handle);

load\_register(cx, 0);

load\_register(dx, 0);

interrupt(0x21s);

int high, low;

store\_register(dx, high);

store\_register(ax, low);

clear\_registers();

return (((long) high) << 4) + ((long) low);

}

static int filecreate(const char\* name) {

load\_register(ah, 0x3Cs);

load\_register(cx, 0x00);

load\_register(dx, name);

interrupt(0x21s);

int handle;

store\_register(ax, handle);

return handle;

}

BOOL fileexists(const char\* name) {

load\_register(ah, 0x43s);

load\_register(al, 0x00s);

load\_register(dx, name);

interrupt(0x21s);

{ int code;

store\_register(ax, code);

clear\_registers();

return (code != FILE\_NOT\_FOUND);

}

}

static int fileopen(const char\* name, unsigned short mode) {

load\_register(ah, 0x3Ds);

load\_register(al, mode);

load\_register(dx, name);

interrupt(0x21s);

int handle;

store\_register(ax, handle);

if (handle > 2) {

return handle;

}

else {

errno = handle;

return -1;

}

}

FILE\* fopen(const char\* filename, const char\* mode) {

int index;

for (index = 0; index < FOPEN\_MAX; ++index) {

if (!g\_fileArray[index].open) {

return freopen(filename, mode, &g\_fileArray[index]);

}

}

return NULL;

}

FILE\* freopen(const char\* filename, const char\* mode, FILE\* stream) {

int handle = -1;

if (strcmp(mode, "r") == 0) {

handle = fileopen(filename, (unsigned short) READ);

}

else if (strcmp(mode, "w") == 0) {

handle = filecreate(filename);

}

else if (strcmp(mode, "a") == 0) {

handle = fileopen(filename, (unsigned short) WRITE);

if (handle != -1) {

fseek(stream, 0L, (int) SEEK\_END);

}

else {

handle = filecreate(filename);

}

}

else if (strcmp(mode, "r+") == 0) {

handle = fileopen(filename, (unsigned short) READ\_WRITE);

}

else if (strcmp(mode, "w+") == 0) {

handle = filecreate(filename);

}

else if (strcmp(mode, "a+") == 0) {

handle = fileopen(filename, (unsigned short) READ\_WRITE);

if (handle != -1) {

fseek(stream, 0L, (int) SEEK\_END);

}

else {

handle = filecreate(filename);

}

}

if (handle != -1) {

stream->open = TRUE;

stream->handle = handle;

stream->size = 0l; // filesize(handle);

strcpy(stream->name, filename);

stream->temporary = FALSE;

return stream;

}

else {

stream->open = FALSE;

return NULL;

}

}

int fflush(FILE\* stream) {

if (stream == NULL) {

int index;

for (index = 0; index < FOPEN\_MAX; ++index) {

if (g\_fileArray[index].open) {

if (fflush(&g\_fileArray[index]) == EOF) {

return EOF;

}

}

}

}

return 0;

}

int fclose(FILE\* stream) {

if (stream != NULL) {

int handle = stream->handle;

load\_register(ah, 0x3Es);

load\_register(bx, handle);

interrupt(0x21s);

{ int handle;

short flagbyte;

store\_register(ax, handle);

store\_flagbyte(flagbyte);

if (!(flagbyte & CARRY\_FLAG)) {

return 0;

}

else {

errno = handle;

return -1;

}

}

}

else {

int index;

for (index = 0; index < FOPEN\_MAX; ++index) {

if (g\_fileArray[index].open) {

if (fclose(&g\_fileArray[index]) == EOF) {

return EOF;

}

}

}

return 0;

}

}

int remove(const char\* name) {

load\_register(ah, 0x41s);

load\_register(dx, name);

load\_register(cl, 0s);

interrupt(0x21s);

{/\* short flagbyte;

store\_flagbyte(flagbyte);

if (flagbyte & CARRY\_FLAG) {

errno = FREMOVE;

return EOF;

}

\*/

return 0;

}

}

int rename(const char\* oldName, const char\* newName) {

load\_register(ah, 0x56s);

load\_register(cl, 0s);

load\_register(dx, oldName);

load\_register(di, newName);

interrupt(0x21s);

{ /\*short flagbyte;

store\_flagbyte(flagbyte);

if (flagbyte & CARRY\_FLAG) {

errno = errno = FRENAME;

return EOF;

}\*/

return 0;

}

}

int setvbuf(FILE\* /\* stream \*/, char\* /\* buffer \*/, int /\* mode \*/, size\_t /\* size \*/) {

return 0;

/\*

if (buffer != NULL) {

buffer = malloc(size);

if (buffer == NULL) {

return -1;

}

}

stream->buffer = buffer;

stream->bufferMode = mode;

stream->bufferMaxSize = size;

return 0;\*/

}

void setbuf(FILE\* /\* stream \*/, char\* /\* buffer \*/) {

/\*

if (buffer == NULL) {

free(stream->buffer);

stream->buffer = NULL;

stream->bufferMode = \_IONBF;

}

else {

setvbuf(stream, buffer, \_IOFBF, BUFSIZ);

}\*/

}

int fgetc(FILE\* stream) {

char c = '\0';

if (fread(&c, sizeof (char), 1, stream) > 0) {

return (int) c;

}

return EOF;

}

char\* fgets(char\* text, int size, FILE\* stream) {

int count = 0;

char prevChar = '\0';

while ((count < (size - 1))) {

char currChar = '\0';

fscanf(stream, "%c", &currChar);

if ((prevChar == '\r') && (currChar == '\n')) {

text[count] = '\0';

break;

}

if ((prevChar != '\r') && (currChar == '\n')) {

text[count] = '\0';

break;

}

if (currChar == EOF) {

text[count] = '\0';

break;

}

if ((currChar != '\r') && (currChar != '\n')) {

text[count++] = currChar;

}

prevChar = currChar;

}

return text;

}

int fputs(const char\* s, FILE\* stream) {

int size = strlen(s) \* sizeof (char);

return (fwrite(s, size, 1, stream) == size) ? 0 : EOF;

}

int getchar(void) {

return fgetc(stdin);

}

char\* gets(char\* s) {

if (fgets(s, -1, stdin) != NULL) {

int size = strlen(s);

if (size > 0) {

s[size - 1] = '\0';

}

return s;

}

else {

return NULL;

}

}

int puts(const char\* s) {

if (fputs(s, stdout) != 0) {

return fputc('\n', stdout);

}

return EOF;

}

int ungetc(int c, FILE\* stream) {

if (stream->ungetc != EOF) {

stream->ungetc = (char) c;

}

return c;

}

size\_t fread(void\* ptr, size\_t size, size\_t nobj, FILE\* stream) {

/\* if (((size \* nobj) > 0) && (stream->ungetc != EOF)) {

char c = stream->ungetc;

memcpy(ptr, &c, 1);

stream->ungetc = EOF;

return fread(((char\*) ptr) + 1, (size \* nobj) - 1, 1, stream);

}\*/

int handle = stream->handle, total = size \* nobj;

load\_register(ah, 0x3Fs);

load\_register(bx, handle);

load\_register(cx, total);

load\_register(dx, ptr);

interrupt(0x21s);

{ short flagbyte;

store\_flagbyte(flagbyte);

if (flagbyte & CARRY\_FLAG) {

stream->errno = errno = FREAD;

return 0;

}

else {

int count;

store\_register(ax, count);

// stream->position += count;

return count;

}

}

}

size\_t fwrite(const void\* ptr, size\_t size, size\_t nobj, FILE\* stream) {

int handle = stream->handle, total = size \* nobj;

load\_register(ah, 0x40s);

load\_register(bx, handle);

load\_register(cx, total);

load\_register(dx, ptr);

interrupt(0x21s);

{ short flagbyte;

store\_flagbyte(flagbyte);

if (flagbyte & CARRY\_FLAG) {

stream->errno = errno = FWRITE;

return 0;

}

else {

int count;

store\_register(ax, count);

// stream->position += count;

// stream->size = MAX(stream->size, stream->position);

return count;

}

}

}

long fseek(FILE\* stream, long offset, int origin) {

load\_register(al, (short) origin);

load\_register(ah, 0x42s);

load\_register(bx, stream->handle);

load\_register(cx, (int) (offset >> 16));

load\_register(dx, (int) offset);

interrupt(0x21s);

{ int highPosition, lowPositionHandle;

store\_register(dx, highPosition);

store\_register(ax, lowPositionHandle);

short flagbyte;

clear\_registers();

store\_flagbyte(flagbyte);

if (!(flagbyte & CARRY\_FLAG)) {

stream->position = (((unsigned long) highPosition) << 16) +

((unsigned long) lowPositionHandle);

return stream->position;

}

else {

stream->errno = (errno = lowPositionHandle);

return -1l;

}

}

}

long ftell(FILE\* stream) {

return fseek(stream, 0L, (int) SEEK\_CUR);

}

void rewind(FILE\* stream) {

(void) fseek(stream, 0L, (int) SEEK\_SET);

}

int fgetpos(FILE\* stream, fpos\_t\* ptr) {

\*ptr = (fpos\_t) ftell(stream);

return 0;

}

int fsetpos(FILE\* stream, const fpos\_t\* ptr) {

return ((int) fseek(stream, \*ptr, (int) SEEK\_SET));

}

void clearerr(FILE\* stream) {

stream->errno = errno = 0;

}

BOOL feof(FILE\* stream) {

long unsigned currPosition = fseek(stream, 0L, (int) SEEK\_CUR);

long unsigned lastPosition = fseek(stream, 0L, (int)SEEK\_END);

fseek(stream, currPosition, (int) SEEK\_SET);

{ BOOL endOfFile = (currPosition == lastPosition);

return endOfFile;

}

}

int ferror(FILE\* stream) {

return stream->errno;

}

void perror(const char\* s) {

printf("%s: %s.\n", s, strerror(errno));

}

## The Standard Library

StdLib.h

#ifndef \_\_STDLIB\_H\_\_

#define \_\_STDLIB\_H\_\_

#define NULL ((void\*) 0)

double atof(char\* s);

int atoi(char\* s);

long atol(char\* s);

double strtod(char\* s, char\*\* endp);

long strtol(char\* s, char\*\* endp, int base);

unsigned long strtoul(char\* s, char\*\* endp, int base);

int rand(void);

void srand(unsigned int seed);

char\* getenv(const char\* name);

int system(const char\* command);

void abort(void);

void exit(int status);

#define FUNC\_MAX 256

#define OPEN\_MAX 16

typedef void(\*FCN)(void);

int atexit(FCN fcn);

int abs(int value);

long labs(long value);

void\* malloc(size\_t size);

void\* realloc(void\* ptr, size\_t newSize);

void\* calloc(size\_t num, size\_t size);

void free(void\* ptr);

void qsort(const void\* valueList, size\_t listSize, size\_t valueSize,

int (\*compare)(const void\*, const void\*), ...);

void\* bsearch(const void\* key, const void\* valueList, size\_t listSize,

size\_t valueSize, int (\*compare)(const void\*, const void\*));

int abs(int value);

long labs(long value);

typedef struct {

int quot, rem;

} div\_t;

div\_t div(int num, int denum);

typedef struct {

long quot, rem;

} ldiv\_t;

ldiv\_t ldiv(long num, long denum);

#endif

StdLib.c

#include <Math.h>

#include <CType.h>

#include <ErrNo.h>

#include <StdArg.h>

#include <StdDef.h>

#include <String.h>

#include <StdLib.h>

#include <StdIO.h>

extern FILE g\_fileArray[];

int atoi(char\* s) {

return (int) strtol(s, (char\*\*) NULL, 10);

}

long atol(char\* s) {

return strtol(s, (char\*\*) NULL, 10);

}

long strtol(char\* s, char\*\* endp, int /\* base \*/) {

int chars = 0;

long value = 0;

sscanf(s, "%li%n", &value, &chars);

if (endp != NULL) {

\*endp = s + chars;

}

return value;

}

unsigned long strtoul(char\* s, char\*\* endp, int /\* base \*/) {

int chars = 0;

unsigned long value = 0;

sscanf(s, "%lu%n", &value, &chars);

if (endp != NULL) {

\*endp = s + chars;

}

return value;

}

double atof(char\* s) {

return strtod(s, (char\*\*) NULL);

}

double strtod(char\* s, char\*\* endp) {

int chars = '\0';

double value = 0;

sscanf(s, "%lf%n", &value, &chars);

if (endp != NULL) {

\*endp = s + chars;

}

return value;

}

void abort(void) {

load\_register(ah, 0x4Cs);

load\_register(al, -1s);

interrupt(0x21s);

}

char\* getenv(char\* /\* name \*/) {

return NULL;

}

int system(const char\* /\* command \*/) {

return -1;

}

void memswp(void\* value1, void\* value2, int valueSize) {

char\* charValue1 = (char\*) value1;

char\* charValue2 = (char\*) value2;

int index;

for (index = 0; index < valueSize; ++index) {

char tempValue = charValue1[index];

charValue1[index] = charValue2[index];

charValue2[index] = tempValue;

}

}

void\* bsearch(const void\* keyPtr, const void\* valueList, size\_t listSize,

size\_t valueSize, int (\*compare)(const void\*, const void\*)) {

int firstIndex = 0, lastIndex = listSize - 1;

if (listSize == 0) {

return NULL;

}

while (TRUE) {

{ char\* firstValuePtr = ((char\*) valueList) + (firstIndex \* valueSize);

int firstCompare = compare(keyPtr, firstValuePtr);

if (firstCompare < 0) {

return NULL;

}

else if (firstCompare == 0) {

return firstValuePtr;

}

}

{ char\* lastValuePtr = ((char\*) valueList) + (lastIndex \* valueSize);

int lastCompare = compare(keyPtr, lastValuePtr);

if (lastCompare > 0) {

return NULL;

}

else if (lastCompare == 0) {

return lastValuePtr;

}

}

{ int middleIndex = (firstIndex + lastIndex) / 2;

char\* middleValuePtr = ((char\*) valueList) + (middleIndex \* valueSize);

int middleCompare = compare(keyPtr, middleValuePtr);

if (middleCompare < 0) {

lastIndex = middleIndex;

}

else if (middleCompare > 0) {

firstIndex = middleIndex;

}

else {

return middleValuePtr;

}

}

}

}

static long g\_randValue;

#define A 1664525l

#define C 1013904223l

#define RAND\_MAX 127 // 32767

int rand(void) {

g\_randValue = ((A \* g\_randValue) + C) % RAND\_MAX;

return (int) g\_randValue;

}

void srand(unsigned int seed) {

g\_randValue = (long) seed;

}

FCN g\_funcArray[FUNC\_MAX] = {NULL};

int atexit(FCN fcn) {

int index;

for (index = 0; index < FUNC\_MAX; ++index) {

if (g\_funcArray[index] == NULL) {

g\_funcArray[index] = fcn;

return 0;

}

}

return -1;

}

void exit(int status) {

int index;

for (index = 0; index < OPEN\_MAX; ++index) {

if (g\_fileArray[index].open) {

fclose(&g\_fileArray[index]);

}

}

for (index = (FUNC\_MAX - 1); index >= 0; --index) {

if (g\_funcArray[index] != NULL) {

g\_funcArray[index]();

}

}

load\_register(ah, 0x4Cs);

load\_register(al, (short) status);

interrupt(0x21s);

}

void swap(char\* leftValuePtr, char\* rightValuePtr, int valueSize) {

int index;

for (index = 0; index < valueSize; ++index) {

char tempValue = leftValuePtr[index];

leftValuePtr[index] = rightValuePtr[index];

rightValuePtr[index] = tempValue;

}

}

void qsort(const void\* valueList, size\_t listSize, size\_t valueSize,

int (\*compare)(const void\*, const void\*)) {

BOOL update;

char\* charList = (char\*) valueList;

int index1;

for (index1 = (listSize - 1); index1 > 0 ; --index1) {

int index2;

update = FALSE;

for (index2 = 0; index2 < index1; ++index2) {

char\* valuePtr1 = charList + (index2 \* valueSize);

char\* valuePtr2 = charList + ((index2 + 1) \* valueSize);

if (compare(valuePtr1, valuePtr2) > 0) {

swap(valuePtr1, valuePtr2, valueSize);

update = TRUE;

}

}

if (!update) {

break;

}

}

}

int abs(int value) {

return (value < 0) ? -value : value;

}

long labs(long value) {

return (value < 0l) ? -value : value;

}

div\_t div(int num, int denum) {

div\_t result = {0, 0};

if (denum == 0) {

errno = EDOM;

return result;

}

result.quot = num / denum;

result.rem = num % denum;

return result;

}

ldiv\_t ldiv(long num, long denum) {

ldiv\_t result = {0, 0};

if (denum == 0) {

errno = EDOM;

return result;

}

result.quot = num / denum;

result.rem = num % denum;

return result;

}

## Time

Time.h

#ifndef \_\_TIME\_H\_\_

#define \_\_TIME\_H\_\_

#define size\_t int

#define time\_t unsigned long

#define clock\_t long

struct tm {

int tm\_sec;

int tm\_min;

int tm\_hour;

int tm\_mday;

int tm\_mon;

int tm\_year;

int tm\_wday;

int tm\_yday;

int tm\_isdst;

};

extern clock\_t clock(void);

extern time\_t time(time\_t\* time);

extern double difftime(time\_t time2, time\_t time1);

extern time\_t mktime(struct tm\* timeStruct);

extern char\* asctime(const struct tm\* timeStruct);

extern char\* ctime(const time\_t\* time);

extern struct tm\* gmtime(const time\_t\* time);

extern struct tm\* localtime(const time\_t\* time);

extern struct tm\* localtimeX(const time\_t\* time);

extern size\_t strftime(char\* buffer, size\_t size,

const char\* format, const struct tm\* timeStruct);

#endif

Time.c

#include <Time.h>

#include <StdIO.h>

#include <StdLib.h>

#include <String.h>

#include <Locale.h>

#include <Assert.h>

clock\_t clock(void) {

return -1;

}

time\_t time(time\_t\* timePtr) {

int year;

short month, monthDay;

load\_register(ah, 0x2As);

interrupt(0x21s);

store\_register(cx, year);

store\_register(dh, month);

store\_register(dl, monthDay);

clear\_registers();

year -= 1900;

--month;

clear\_registers();

short hour, min, sec;

load\_register(ah, 0x2Cs);

interrupt(0x21s);

store\_register(ch, hour);

store\_register(cl, min);

store\_register(dh, sec);

clear\_registers();

struct tm s = {sec, min, hour, monthDay, month, year, 0, 0, 1};

time\_t time = mktime(&s);

if (timePtr != NULL) {

\*timePtr = time;

}

return time;

}

time\_t mktime(struct tm\* tp) {

int leapDays = ((tp->tm\_year - 69) / 4);

int totalDays = 365 \* (tp->tm\_year - 69) + leapDays + (tp->tm\_mday - 1);

BOOL leapYear = ((tp->tm\_year % 4) == 0);

if (leapYear) {

int leapMonths[] = {0, 31, 60, 91, 121, 152, 182, 213, 243, 274, 304, 335};

totalDays += leapMonths[tp->tm\_mon];

}

else {

int regularMonths[] = {0, 31, 59, 90, 120, 151, 181, 212, 242, 273, 303,334};

totalDays += regularMonths[tp->tm\_mon];

}

unsigned long totalSeconds = (86400lu \* totalDays) + (3600lu \* tp->tm\_hour)+

(60lu \* tp->tm\_min) + tp->tm\_sec;

unsigned long result = (((unsigned long)tp->tm\_isdst) << 31) |

(0x7FFFFFFF & totalSeconds);

return result;

}

static struct tm g\_timeStruct;

#define JAN\_1\_1969 3 // Wednesday January 1, 1969

#define DAYS\_OF\_FOUR\_YEARS 1461

#define DAYS\_OF\_ONE\_YEAR 365

struct tm\* localtime(const time\_t\* timePtr) {

if (timePtr != NULL) {

time\_t time = \*timePtr;

g\_timeStruct.tm\_isdst = (time >> 31);

unsigned long totalSeconds = (time & 0x7FFFFFFF);

{ unsigned long totalDays = totalSeconds / 86400lu;

g\_timeStruct.tm\_wday = (int) ((totalDays + JAN\_1\_1969) % 7);

int fourYears = totalDays / DAYS\_OF\_FOUR\_YEARS;

totalDays %= DAYS\_OF\_FOUR\_YEARS;

int years = totalDays / DAYS\_OF\_ONE\_YEAR;

totalDays %= DAYS\_OF\_ONE\_YEAR;

g\_timeStruct.tm\_year = 69 + (4 \* fourYears) + years;

g\_timeStruct.tm\_yday = totalDays;

BOOL leapYear = ((g\_timeStruct.tm\_year % 4) == 0);

if (leapYear) {

int leapMonths[] = {0, 31, 60, 91, 121, 152,

182, 213, 243, 274, 304, 335, 366},

arraySize = (sizeof leapMonths) / (sizeof leapMonths[0]), index;

for (index = 1; index < arraySize; ++index) {

if (totalDays < leapMonths[index]) {

g\_timeStruct.tm\_mon = index - 1;

g\_timeStruct.tm\_mday = (totalDays % leapMonths[index - 1]) + 1;

break;

}

}

}

else {

int regularMonths[] = {0, 31, 59, 90, 120, 151,

181, 212, 242, 273, 303, 334, 365},

arraySize = (sizeof regularMonths) / (sizeof regularMonths[0]),

index;

for (index = 1; index < arraySize; ++index) {

if (totalDays < regularMonths[index]) {

g\_timeStruct.tm\_mon = index - 1;

g\_timeStruct.tm\_mday = (totalDays % regularMonths[index - 1]) + 1;

break;

}

}

}

}

{ unsigned long daySeconds = totalSeconds % 86400lu;

g\_timeStruct.tm\_hour = (int) (daySeconds / 3600);

g\_timeStruct.tm\_min = (int) ((daySeconds % 3600) / 60);

g\_timeStruct.tm\_sec = (int) (daySeconds % 60);

}

return &g\_timeStruct;

}

return NULL;

}

struct tm\* localtimeX(const time\_t\* timePtr) {

if (timePtr != NULL) {

time\_t time = \*timePtr;

g\_timeStruct.tm\_isdst = (time >> 31);

unsigned long totalSeconds = (time & 0x7FFFFFFF);

{ unsigned long totalDays = totalSeconds / 86400lu;

g\_timeStruct.tm\_wday = (int) ((totalDays + JAN\_1\_1969) % 7);

int year;

BOOL leapYear;

for (year = 69;; ++year) {

leapYear = ((year % 4) == 0);

int yearDays = leapYear ? 366 : 365;

if (totalDays >= yearDays) {

totalDays -= yearDays;

}

else {

break;

}

}

g\_timeStruct.tm\_year = year;

g\_timeStruct.tm\_yday = totalDays;

int month, months[12] = {31, leapYear ? 29 : 28, 31, 30, 31, 30,

31, 31, 30, 31, 30, 31};

for (month = 0; month < 12; ++month) {

if (totalDays >= months[month]) {

totalDays -= months[month];

}

else {

break;

}

}

g\_timeStruct.tm\_mon = month;

g\_timeStruct.tm\_mday = totalDays + 1;

}

{ unsigned long daySeconds = totalSeconds % 86400lu;

g\_timeStruct.tm\_hour = (int) (daySeconds / 3600lu);

g\_timeStruct.tm\_min = (int) ((daySeconds % 3600lu) / 60lu);

g\_timeStruct.tm\_sec = (int) (daySeconds % 60);

}

return &g\_timeStruct;

}

return NULL;

}

double difftime(time\_t time1, time\_t time2) {

long totalSeconds1 = (time1 &0x7FFFFFFF),

totalSeconds2 = (time2 &0x7FFFFFFF);

return (double) (totalSeconds2 - totalSeconds1);

}

static char g\_timeString[256];

static char\* g\_defaultShortDayList[] =

{"Sun", "Mon", "Tue", "Wed", "Thu", "Fri", "Sat"};

static char\* g\_defaultLongDayList[] = {"Sunday", "Monday", "Tuesday",

"Wednesday", "Thursday", "Friday", "Saturday"};

static char\* g\_defaultShortMonthList[] = {"Jan", "Feb", "Mar", "Apr", "May",

"Jun", "Jul", "Aug", "Sep", "Oct", "Nov", "Dec"};

static char\* g\_defaultLongMonthList[] = {"January", "February", "March",

"April", "May", "June", "July", "August",

"September", "October", "November", "December"};

void default\_test() {

printf("<%p>", g\_defaultShortDayList);

printf("<%s>", g\_defaultShortDayList[1]);

}

char\* asctime(const struct tm\* tp) {

struct lconv\* localeConvPtr = NULL;//localeconv();

char\*\* shortDayList = (localeConvPtr != NULL) ?

localeConvPtr->shortDayList : NULL;

char\*\* shortMonthList = (localeConvPtr != NULL) ?

localeConvPtr->shortMonthList : NULL;

shortDayList = (shortDayList != NULL) ? shortDayList

: g\_defaultShortDayList;

shortMonthList = (shortMonthList != NULL) ? shortMonthList

: g\_defaultShortMonthList;

sprintf(g\_timeString, "%s %s %2i %02i:%02i:%02i %04i\n",

shortDayList[tp->tm\_wday], shortMonthList[tp->tm\_mon],

tp->tm\_mday, tp->tm\_hour, tp->tm\_min,

tp->tm\_sec, tp->tm\_year + 1900);

return g\_timeString;

}

char\* ctime(const time\_t\* time) {

return asctime(localtime(time));

}

struct tm\* gmtime(const time\_t\* timePtr) {

time\_t t = \*timePtr;

int timeZone = 0;

struct tm\* tmPtr = localtime(timePtr);

struct lconv\* localeConvPtr = localeconv();

if (localeConvPtr != NULL) {

timeZone = tmPtr->tm\_isdst ? localeConvPtr->summerTimeZone

: localeConvPtr->winterTimeZone;

}

t -= 3600l \* timeZone;

return localtime(&t);

}

size\_t strftime(char\* s, size\_t smax, const char\* fmt, const struct tm\* tp) {

struct lconv\* localeConvPtr = localeconv();

char\*\* shortDayList = (localeConvPtr != NULL) ?

(localeConvPtr->shortDayList) : NULL;

char\*\* shortMonthList = (localeConvPtr != NULL) ?

(localeConvPtr->shortMonthList) : NULL;

char\*\* longDayList = (localeConvPtr != NULL) ?

(localeConvPtr->longDayList) : NULL;

char\*\* longMonthList = (localeConvPtr != NULL) ?

(localeConvPtr->longMonthList) : NULL;

strcpy(s, "");

shortDayList = (shortDayList != NULL) ? shortDayList : g\_defaultShortDayList;

longDayList = (longDayList != NULL) ? longDayList : g\_defaultLongDayList;

shortMonthList = (shortMonthList != NULL) ? shortMonthList

: g\_defaultShortMonthList;

longMonthList = (longMonthList != NULL) ? longMonthList

: g\_defaultLongMonthList;

int index;

for (index = 0; fmt[index] != '\0'; ++index) {

char add[20];

if (fmt[index] == '%') {

switch (fmt[++index]) {

case 'a':

strcpy(add, shortDayList[tp->tm\_wday]);

break;

case 'A':

strcpy(add, longDayList[tp->tm\_wday]);

break;

case 'b':

strcpy(add, shortMonthList[tp->tm\_mon]);

break;

case 'B':

strcpy(add, longMonthList[tp->tm\_mon]);

break;

case 'c':

sprintf(add, "%04d-%02d-%02d %02d:%02d:%02d", 1900 + tp->tm\_year,

tp->tm\_mon + 1, tp->tm\_mday, tp->tm\_hour, tp->tm\_min,

tp->tm\_sec);

break;

case 'd':

sprintf(add, "%02d", tp->tm\_mday);

break;

case 'H':

sprintf(add, "%02d", tp->tm\_hour);

break;

case 'I':

sprintf(add, "%02d", (tp->tm\_hour % 12) + 1);

break;

case 'j':

sprintf(add, "%03d", tp->tm\_yday);

break;

case 'm':

sprintf(add, "%02d", tp->tm\_mon);

break;

case 'M':

sprintf(add, "%02d", tp->tm\_min);

break;

case 'p':

sprintf(add, "%s", (tp->tm\_hour < 12) ? "AM" : "PM");

break;

case 'S':

sprintf(add, "%02d", tp->tm\_sec);

break;

case 'U':

sprintf(add, "%02d", 0 /\*WeekNumberSunday(tp) \*/);

break;

case 'w':

sprintf(add, "%02d", tp->tm\_wday);

break;

case 'W':

sprintf(add, "%02d", 0 /\* WeekNumberMonday(tp) \*/);

break;

case 'x':

sprintf(add, "%04d-%02d-%02d", 1900 + tp->tm\_year, tp->tm\_mon,

tp->tm\_mday);

break;

case 'X':

sprintf(add, "%02d:%02d:%02d", tp->tm\_hour, tp->tm\_min, tp->tm\_sec);

break;

case 'y':

sprintf(add, "%02d", tp->tm\_year % 100);

break;

case 'Y':

sprintf(add, "%04d", 1900 + tp->tm\_year);

break;

case 'Z':

strcpy(add, "");

break;

case '%':

strcpy(add, "%");

}

}

else {

add[0] = fmt[index];

add[1] = '\0';

}

if ((strlen(s) + strlen(add)) < smax) {

strcat(s, add);

}

else {

break;

}

}

return strlen(s);

}

# The C Grammar

This grammar of the C compiler of this book is based on (and to a large extent identical with) the grammar defined in the *American National Standard for Information systems – Programming Language C, X3.159-1989* standard, which is described in the second edition of *The C Programming Language* by Brian W. Kernighan and Dennis M. Ritchie.

## The Preprocessor Grammar

The following grammar is implemented in CUP in Chapter 3.

control\_line ::=

# define identifier token-sequence

| # define identifier( identifier, … , identifier ) token-sequence

| # undef identifier

| # include <filename>

| # include “filename“

| # line constant

| # line constant “filename”

| # error optional\_token\_sequence

| # pragma optional\_token\_sequence

| #

| preprocessor\_conditional

optional\_token\_sequence ::=

/\* empty \*/

| token\_sequence

preprocessor\_conditional ::=

if-line text optional\_elif-parts optional\_else\_line # endif

if-line ::=

# if constant\_expression

| # ifdef identifier

| # ifndef identifier

optional\_elif-parts ::=

/\* empty \*/

| elif-part optional\_elif-parts

elif-part ::=

# elif constant\_expression text

optional\_else\_part ::=

/\* empty \*/

| #else text

## The Language Grammar

The following grammar is implemented in Bison in Chapter 6.

*translation*\_*unit* ::=

*external*\_*declaration*

| *translation*\_*unit* *external*\_*declaration*

*external*\_*declaration* ::=

*function*\_*definition*

| *declaration*

*function*\_*definition* ::=

*pointer\_declarator* *optional*\_*declaration*\_*list* **(** *optional*\_*statement*\_*list* **)**

| *declaration*\_*specifier*\_*list* *pointer\_declarator* *optional*\_*declaration*\_*list*

**(** *optional*\_*statement*\_*list* **)**

*optional*\_*declaration*\_*list* ::=

/\* *empty* \*/

| *optional*\_*declaration*\_*list* *declaration*

*declaration* ::=

*declaration*\_*specifier*\_*list* **;**

| *declaration*\_*specifier*\_*list* *declarator\_list* **;**

*declaration*\_*specifier*\_*list* ::=

*declaration*\_*specifier*

| *declaration*\_*specifier*\_*list* *declaration*\_*specifier*

*declaration*\_*specifier* ::=

**const** | **volatile** | **auto** | **register** | **static** | **extern** | **typedef** | **void**

| **char** | **wchar\_t** | **short** | **int** | **long** | **float** | **double** | **signed** | **unsigned**

| *struct\_union*\_*specifier |* *enum*\_*specifier* | *typedef*\_*name*

*struct\_union*\_*specifier* ::=

struct\_union *optional*\_*identifier* **(** *declaration*\_*list* **)**

| *struct\_union* **identifier**

*struct\_union ::=*

**struct** | **union**

*optional*\_*identifier* ::=

/\* *empty* \*/

| **identifier**

*declaration*\_*list* ::=

*declaration*

| *declaration*\_*list* *declaration*

*enum*\_*specifier* ::=

**enum** *optional*\_*identifier* **(** *enum*\_*list* **)**

| **enum** **identifier**

*enum*\_*list* ::=

*enum*

| *enum*\_*list* **,** *enum*

*enum* ::=

**identifier**

| **identifier** **=** *const*\_*expression*

*declarator\_list* ::=

*initialization\_bitfield\_pointer\_declarator*

| *declarator\_list* **,** *initialization\_bitfield\_pointer\_declarator*

*initialization\_bitfield\_pointer\_declarator* ::=

*pointer\_declarator*

| *pointer\_declarator* **=** *initializer*

| *pointer\_declarator* **:** *const*\_*expression*

| **:** *const*\_*expression*

*pointer\_declarator* ::=

*optional*\_*pointer*\_*list* *direct*\_*pointer\_declarator*

*direct*\_*pointer\_declarator* ::=

**identifier**

| **(** *pointer\_declarator* **)**

| *direct*\_*pointer\_declarator* **[** *optional*\_*const*\_*expression* **]**

| *direct*\_*pointer\_declarator* **(** *parameter*\_*ellipse*\_*list* **)**

| *direct*\_*pointer\_declarator* **(** *optional*\_*identifier*\_*list* **)**

*optional*\_*pointer*\_*list* ::=

/\* *empty* \*/

| *pointer*\_*list*

*pointer*\_*list* ::=

*pointer*

| *pointer*\_*list* *pointer*

*pointer* ::=

**\*** | **\*** *declaration*\_*specifier*\_*list*

*parameter*\_*ellipse*\_*list* ::=

*parameter*\_*list*

| *parameter*\_*list* **,** **...**

*parameter*\_*list* ::=

*parameter*\_*declaration*

| *parameter*\_*list* **,** *parameter*\_*declaration*

*parameter*\_*declaration* ::=

*declaration*\_*specifier*\_*list*

| *declaration*\_*specifier*\_*list* *pointer\_declarator*

| *declaration*\_*specifier*\_*list* *abstract*\_*pointer\_declarator*

*optional*\_*identifier*\_*list* ::=

/\* *empty* \*/

| *identifier*\_*list*

*identifier*\_*list* ::=

**identifier**

| *identifier*\_*list* **,** **identifier**

*initializer*\_*list* ::=

*initializer*

| *initializer*\_*list* **,** *initializer*

*initializer* ::=

*assignment*\_*expression*

| **(** *initializer*\_*list* **)**

| **(** *initializer*\_*list* **,** **)**

*type*\_*name* ::=

*declaration*\_*specifier*\_*list*

| *declaration*\_*specifier*\_*list* *abstract*\_*pointer\_declarator*

*abstract*\_*pointer\_declarator* ::=

*pointer*\_*list*

| *optional*\_*pointer*\_*list* *direct*\_*abstract*\_*pointer\_declarator*

*direct*\_*abstract*\_*pointer\_declarator* ::=

**(** *abstract*\_*pointer\_declarator* **)**

| **[** *optional*\_*const*\_*expression* **]**

| *direct*\_*abstract*\_*pointer\_declarator* **[** *optional*\_*const*\_*expression* **]**

| **(** *optional*\_*parameter*\_*ellipse*\_*list* **)**

| *direct*\_*abstract*\_*pointer\_declarator* **(** *optional*\_*parameter*\_*ellipse*\_*list* **)**

*optional*\_*statement*\_*list* ::=

/\* *empty* \*/

| *statement*\_*list*

*statement*\_*list* ::=

*statement*

| *statement*\_*list* *statement*

*statement* ::=

*open*\_*statement*

| *closed*\_*statement*

*open*\_*statement* ::=

**if** **(** *expression* **)** *statement*

| **if** **(** *expression* **)** *closed*\_*statement* **else** *open*\_*statement*

| **switch** **(** *expression* **)** *open*\_*statement*

| **case** *const*\_*expression* **:** *open*\_*statement*

| **while** **(** *expression* **)** *open*\_*statement*

| **for** **(** *optional*\_*expression* **;** *optional*\_*expression* **;**

*optional*\_*expression* **)** *open*\_*statement*

| **identifier** **:** *open*\_*statement*

*closed*\_*statement* ::=

**if** **(** *expression* **)** *closed*\_*statement* **else** *closed*\_*statement*

| **switch** **(** *expression* **)** *closed*\_*statement*

| **while** **(** *expression* **)** *closed*\_*statement*

| **for** **(** *optional*\_*expression* **;** *optional*\_*expression* **;**

*optional*\_*expression* **)** *closed*\_*statement*

| *declaration*

| **case** *const*\_*expression* **:** *closed*\_*statement*

| **default** **:** *closed*\_*statement*

| *optional*\_*expression* **;**

| **(** *optional*\_*statement*\_*list* **)**

| **do** *statement* **while** **(** *expression* **)** **;**

| **goto** **identifier ;**

| **continue ;**

| **break ;**

| **return** *optional*\_*expression* **;**

| **load\_register** **(** **identifier ,** *expression* **) ;**

| **store\_register** **(** **identifier** **,** *expression* **) ;**

| **store\_flagbyte** **(** *expression* **) ;**

| **jump\_register** **(** **identifier** **) ;**

| **interupt** **(** *const*\_*expression* **) ;**

*optional*\_*expression* ::=

/\* *empty* \*/

| *expression*

*expression* ::=

*assignment*\_*expression*

| *expression* **,** *assignment*\_*expression*

*assignment*\_*expression* ::=

*conditional*\_*expression*

| *prefix*\_*expression* *assignment*\_*operator* *assignment*\_*expression*

*assignment*\_*operator* ::=

**=** | \***=** | /**=** | %**=** | +**=** | -**=** | <<**=** | <<**=** | &**=** | |**=** | ^**=**

*conditional*\_*expression* ::=

*logical*\_*or*\_*expression*

| *logical*\_*or*\_*expression*

**?** *expression* **:** *conditional*\_*expression*

*optional*\_*const*\_*expression* ::=

/\* *empty* \*/

| *const*\_*expression*

*const*\_*expression* ::=

*conditional*\_*expression*

*logical*\_*or*\_*expression* ::=

*logical*\_*and*\_*expression*

| *logical*\_*or*\_*expression* *||* *logical*\_*and*\_*expression*

*logical*\_*and*\_*expression* ::=

*BITWISE\_OR*\_*expression*

| *bitwise*\_*and*\_*expression* **&&** *BITWISE\_OR*\_*expression*

*BITWISE\_OR*\_*expression* ::=

*bitwise*\_*xor*\_*expression*

| *BITWISE\_OR*\_*expression* **|** *bitwise*\_*xor*\_*expression*

*bitwise*\_*xor*\_*expression* ::=

*bitwise*\_*and*\_*expression*

| *bitwise*\_*xor*\_*expression* **^** *bitwise*\_*and*\_*expression*

*bitwise*\_*and*\_*expression* ::=

*equality*\_*expression*

| *bitwise*\_*and*\_*expression* **&** *equality*\_*expression*

*equality*\_*expression* ::=

*relation*\_*expression*

| *equality*\_*expression* *equality*\_*operator* *relation*\_*expression*

*equality*\_*operator* ::=

**==** | **!=**

*relation*\_*expression* ::=

*shift*\_*expression*

| *relation*\_*expression* *relation*\_*operator* *shift*\_*expression*

*relation*\_*operator* ::=

**<** | **<=** | **>** | **>=**

*shift*\_*expression* ::=

*add*\_*expression*

| *shift*\_*expression* *shift*\_*operator* *add*\_*expression*

*shift*\_*operator* ::=

**<<** | **>>**

*add*\_*expression* ::=

*multiply*\_*expression*

| *add*\_*expression* *binary*\_*add*\_*operator* *multiply*\_*expression*

*binary*\_*add*\_*operator* ::=

**+** | **-**

*multiply*\_*expression* ::=

*type\_cast\_expression*

| *multiply*\_*expression* *multiply*\_*operator* *type\_cast\_expression*

*multiply*\_*operator* ::=

**\*** | **/** | **%**

*type\_cast\_expression* ::=

*prefix*\_*expression*

| **(** *type*\_*name* **)** *type\_cast\_expression*

*prefix*\_*expression* ::=

*postfix*\_*expression*

| *increment*\_*operator* *prefix*\_*expression*

| *prefix*\_*add*\_*operator* *type\_cast\_expression*

| **!** *type\_cast\_expression*

| **~** *type\_cast\_expression*

| **&** *type\_cast\_expression*

| **\*** *type\_cast\_expression*

| **sizeof** *prefix*\_*expression*

| **sizeof** **(** *type*\_*name* **)**

*prefix*\_*add*\_*operator* ::=

**+** | **-**

*increment*\_*operator* ::=

**++** | **--**

*postfix*\_*expression* ::=

*primary*\_*expression*

| *postfix*\_*expression* **[** *expression* **]**

| *postfix*\_*expression* **(** *optional*\_*argument*\_*expression*\_*list* **)**

| *postfix*\_*expression* **.** **identifier**

| *postfix*\_*expression* **->** **identifier**

| *postfix*\_*expression* *increment*\_*operator*

*primary*\_*expression* ::=

**identifier**

| **value**

| **(** *expression* **)**

*optional*\_*argument*\_*expression*\_*list* ::=

/\* *empty* \*/

| *argument*\_*expression*\_*list*

*argument*\_*expression*\_*list* ::=

*assignment*\_*expression*

| *argument*\_*expression*\_*list* , *assignment*\_*expression*

# A Crash Course in C

This appendix gives a crash course in the C programming language. The story of C began in the late sixties with the Multics project, which was a predecessor to the the UNIX operating system. The first two versions of Multics were written in the assembler language. However, since programming in assembly languages was rather cumbersome, an appropriate high-level language was called for.

The researchers Dennis Ritchie and Ken Thompon created the language B by simplifying the research language BCPL[[10]](#footnote-10) in 1969. The new versions, called New B, were introduced in 1972. The name did soon evolve into C. In 1978, Brian Kernighan and Dennis Ritchie wrote the book The C Programming Language, which introduced the informal K&R C standard, named after the authors. In 1983, a working group was formed by the American National Standard Institute (ANSI) in order to define a C standard, which was finally adopted in 1989. The ANSI C standard is described in an appendix of the second edition of The C Programming Languages.

## The Compiler and the Linker

The source code is the the actual text the programmer writes. The compiler is the program that translates the source code into object code, which the computer can read. The linker puts several compiled files into an executable file

Let us say we have a C program in the source code file Program.c and a routine used by the program in Routine.c. Furthermore, the program calls a function in the standard library. In this case, the compiler translates the source code into object code and the linker joins the code into the executable file Program.exe.

Compiler

Program.c

Linker

Program.obj

Routine.c

Routine.obj

Compiler

Standard Library

Program.exe

If the compiler reports an error we refer to it as compile-time error, and if an error occurs during the execution of the program we call it a run-time error.

## The Hello-World Program

In the first edition of The C Programming Language the hello-world example was introduced; that is, a small program that prints the text “Hello, World!”[[11]](#footnote-11) on the screen. To write a hello-world programoften means to write a small program to test that the compiler and linker work properly. The execution of the program does always start with the function main.

#include <stdio.h>

void main(void) {

puts("Hello, World!");

}

When executing, the program above generates the following printout.



## Comments

In C, it is possible to insert comments to describe and clarify the meaning of the program. The comments are ignored by the compiler[[12]](#footnote-12). There are two types of comments: line comments and block comments. Line comments start with two slashes and ends at the end of the line.

printf("Hello, World!"); // Prints "Hello, World!".

Block comments begin with a slash and an asterisk and ends with an asterisk and a slash. A block comment may range over several lines.

/\* This is an example of a C program. It writes the text

"Hello, World!" at the screen. \*/

#include <stdio.h>

void main(void) {

printf("Hello, World!");

}

Block comments cannot be nested. The following example will result in a compile-time error.

/\* A block comment cannot be /\* nested inside \*/ another

block comment. \*/

A piece of advice is that you use the line comments for regular comments, and save the block comments for situations when you need to comment a whole block of code for debugging purposes. Actually, there is a third level of comments if you use conditional programming, see macros (Section 22.9) for details.

#if 0

This is an example of a C program. It writes the text

"Hello, World!" at the screen.

#endif

#include <stdio.h>

void main(void) {

printf("Hello, World!");

}

### Types and Variables

There are several types in C. They can be divided into two groups: simple and compunded. The simple types can be further classified into integral, floating, and logical types. The compunded types are arrays, structures, unions, and pointers. They all (directly or indirectly) composed by simple types. We can also define a type with our own integer values, called the enumeration type.

### Simple Types

There are four simple types intended for storing integers: char, short int, int, and long int. They are called the integral types. The types short int and long int may be abbreviated to short and long, respectively. As the names implies, they are designed for storing characters, small integers, normal integers, and large integers, respectively. The exact limits of the values possible to store varies between different compilers. Such property that is not defined in the standard, but is rather defined by the compiler, linker, or operating system, is said to implementation dependent.

Furthermore, the integral types may be signed or unsigned. An unsigned type must not have negative values. If the word signed or unsigned is left out, a short int, int, and long int will be signed. Weather a char will be signed or unsigned is implementation dependent. However, a char is always one byte long, which means that it always holds a single character, regardless whether it is unsigned or not.

The next category of simple types is the floating types, which are used to store decimal numbers. The types are float, double, and long double; where float stores the smallest value and long double the largest one. How large values each type can store is implementation dependent. A floating type cannot be unsigned.

In many languages, there is a logical type (named bool, boolean, or logical) that stores logical value true or false. In C, however, there is no such type[[13]](#footnote-13). Instead, int is used as a logical type, where zero represents false and all other values represent true.

### Variables

A variable can be regarded as a box in the memory. In almost every case, we do not need to know its exact memory address. A variable always has a name, a type, and a value. We define a variable by simple writing its type and name. If we want to, we can initialize the variable; that is, assign it a value. If we do not, the variable’s value is undefined[[14]](#footnote-14). In this chapter, that value is denoted by a question mark (?).

int i = 123, j;

double d = 3.14;

char c = 'a';

123

i

?

j

3.14

d

'a'

c

We can transform values between the types by stating the new type within parentheses. The use of stating the type is not strictly necessary. However, if omitted the compiler may give warnings[[15]](#footnote-15).

int i = 123;

double x = 1.23;

int j = (int) x;

double y = (double) i;

### Constants

As the name implies, a constant is a variable whose value cannot be altered since it has been initialized. Unlike variables, constants must always be initialized and are often written in capital letters.

const double PI = 3.14;

### Output

As noted in the first section of this chapter, the standard function for output is printf (formatted print) writes to the standard output device (normally a text window).

The first argument is the format strings. With it, we can write every kind of value. However, there are many format codes to keep track of. A format code begins with the per cent character (%) followed with optional digits and signs ending with a character. We can write all the types we have went through so far. The character '\n' represents a new line. Below follows a list of the most common format codes.

|  |  |
| --- | --- |
| Code | Type |
| d | decimal integer |
| o | octal integer |
| x | hexadecimal integer |
| f | double |
| lf | long double |
| c | char |
| s | char \* |
| p | Prints the address of the variable in hexadecimal form. |

In the program below, *i* represents an integer, *f[[16]](#footnote-16)* a double, and *c* a character. Note that it is the responsible of the programmer to match the values following the format string with values of the correct types. There is no built-in error handling and mistakes would likely result in a run-time error.

#include <stdio.h>

void main() {

char c = 'a';

int i = 123;

double d = 1.23;

printf("char %c, int %d, double %f\n", c, i, d);

}

There is also the puts function that prints a string followed by a new-line. The following two lines are equivalent.

printf("Hello, World!\n");

puts("Hello, world!");

### Input

The standard functions for input is and scanf (formatted scan), it reads from the standard input device (normally the keyboard).

Similar to printf, its first argument is the format string. The format codes mostly agree with those of printf, but not completely. See the table at the end of the chapter for a list of format codes. For instance, scanf expects lf to read a double value, in contrast to printf that expects f*[[17]](#footnote-17)*.

Another difference is that while printf wants the values corresponding to the format codes, scanf wants the addresses of the variables that are going to store the values. This is accomplished by using the ampersand (&), see the Section 11.1.1 on pointers. Below follows a list of the most common format codes.

|  |  |
| --- | --- |
| Code | Type |
| d | decimal integer |
| o | octal integer |
| x | hexadecimal integer |
| f | float |
| lf | double |
| ld | long double |
| c | Char |
| S | char \* |

In the program below, note that scanf expects lf to represent a double float, in contrast to printf that expects f.

#include <stdio.h>

void main() {

char c;

int i;

double d;

printf("char: ");

scanf("%c", &c);

printf("int: ");

scanf("%d", &i);

printf("double: ");

scanf("%lf", &d);

printf("char %c, int %d, double %f\n", c, i, d);

}

There is also the gets function that reads a string until it reaches a new-line, which differs from scanf that reads a string until it reaches a white-space character[[18]](#footnote-18). It the user prints “Hello World” followed by a new-line twice when executing the code below, gets will place “Hello world” is s, while scanf will place “Hello” in t. Note that we do not give the addresses of s or t, since they are pointer to characters in themselves.

char s[20], t[20];

gets(s);

scanf("%s", t);

### Enumerations

An enumeration is a way to create our own integral type. We can define which values a variable of the type can store. In practice, however, enumerations are essentially an easy way to define integer constants[[19]](#footnote-19).

enum Cars {FORD, VOLVO, TOYOTA, VOLKSWAGEN};

Unless we state otherwise, the constants are assigned to zero, one, two, three, and so on. In the example above, FORD is an integer constant with the value zero; VOLVO has the value one, TOYOTA three, and VOLKSWAGEN four.

We do not have to name the enumeration type. In the example above, Cars can be omitted. We can also assign an integer value to some (or all) of the constants. In the example below, TOYOTA is assigned the value 10. The constants without assigned values will be given the value of the preceding constant before, adding one. This implies that VOLKSWAGEN will be assigned the value 11.

enum {FORD, VOLVO, TOYOTA = 10, VOLKSWAGEN};

### Arrays

An array is a variable compiled by several values of the same type. The values are stored on consecutive locations in memory. The array size may be stated explicitly of implicitly by a list of values. In the following example, a is given the size three, b is given size two and c is given size four, even though only its first two values are defined, which may cause the compiler to emit a warning. Naturally, d is given the size three.

int a[3] = {11, 12, 13};

double b[] = {1.2, 3.4};

char c[4] = {'a', 'b'}, d[3];

d

11

a

b

c

12

13

1.2

3.4

'a'

'b'

?

?

?

?

?

0

1

2

0

1

0

1

2

3

0

1

2

A value of an array can be accessed by index (brackets) notation.

int i = a[2];

double x = b[0];

char t = c[1];

13

i

1.2

x

'b'

t

### Characters and Strings

As stated in Section 2.1.1, a char is a small integer type intended to store exactly one character. A string stores a (possible empty) sequence of characters. There is no built in type for describing a string. Instead, a string is defined by an array of characters. Note that characters are enclosed by single quotations while strings are enclosed by double quotations.

char c = 'a';

char s[] = "Hello, World!";

'a'

c

"Hello, World!"

s

When defining a string, the compiler adds a special zero character '\0' at the end. Its task is to mark the end of the string. If the code above, the character array s has a size of fourteen (thirteen regular characters, including space, comma, and exlamation mark as well as the finishing zero character). The two following definitions are equivalent, they both define character arrays of size four (three regular characters and the finishing zero character).

char s1[] = "abc";

char s2[] = {'a', 'b', 'c', '\0'};

### The ASCII Table

So far we have assigned characters to char variables. Technically, the compiler translates a character literal into an integer in accordance to the ASCII table (see Appendix **Fel! Hittar inte referenskälla.**), where character has a corresponding integer value. For instance, ‘a’ corresponds to 97, ‘b’ to 98, and so on. The line above could have been written as the following code instead. The possibility to write character and string literal is for convenience only.

char s2[] = {97, 98, 99, 0};

### Pointers

A pointer is a variable containing the address of value. For example, let us say that the integer i has the value 123 that is stored on the memory address 10,000, and that the size of each value is four bytes. As p is a pointer to i, it holds the value 10,000.

10000

p

123

10000

9992

9996

10004

In almost all cases, we do not really need to know the address of the value. Therefore, a simpler (and clearer) way to illustrate the same thing is to draw an arrow from the pointer to the value.

p

123

i

The following code gives rise to the diagram above[[20]](#footnote-20), where the ampersand (&) denotes the address of the variable.

int i = 123;

int \*p = &i;

If we want to access the value pointed at, we use the asterisk (\*), which dereferenceers the pointer (“follows the arrow”). As the address operator (&) gives the address of a variable (“against the arrow”), they address and dereferenceerring operators can be regarded as each other reverses. Note that the asterisk is used on two occasions, when we define a pointer variable and when we dereferenceer a pointer. The asterisk is in fact used on a third occasion, when multiplying two values. The ampersand is used on two occasions: address and bitwise and[[21]](#footnote-21).

int i = 999;

int \*p = &i; // Pointer definition.

int j = \*p; // Dereferenceereeing of a pointer.

int a = 1, b = 2;

int c = a \* b; // Multiplication.

### Pointers and Dynamic Memory

Pointers can also be used to allocate dynamic memory. There is a section of the memory called the heap that is used for dynamically allocated memory blocks. The standard functions malloc, calloc, realloc and free used to allocate, reallocate, and deallocate the memory, respectively. Memory not dynamically allocated is referred to static memory.

p

123

int \*p = malloc(sizeof(int));

\*p = 123;

free(p);

We can also allocate memory for a whole array. Even though p is a pointer in the example below, we can use the array index notation to access a value of the array in the allocated memory block.

p

123

124

125

0

1

2

int \*p = calloc(3, sizeof(int));

p[0] = 123;

p[1] = 124;

p[2] = 125;

free(p);

Once we have allocated a memory, we can allocate a block of different size by calling realloc, which copies to contents of the memory block to new address, is necessary.

int \*p = malloc(sizeof(int));

p[0] = 123;

p = realloc(2 \* sizeof(int));

p[1] = 456;

free(p);

p

123

456

0

1

The difference between malloc and calloc is that calloc resets the bytes of the memory blocks to zero, which malloc do not. Otherwise, the following two calls are equivalent.

int \*p = calloc(3, sizeof(int));

int \*q = malloc(3 \* sizeof(int));

The predefined constant NULL holds the pointer equivalence of the zero value. We say that the pointer is set to null. In a diagram, we can simple write null.

p

NULL

int \*p = NULL;

Sometimes, the electric ground symbol is used to symbolize a null pointer. For this reason, a null pointer is sometimes said to be a grounded pointer.

p

There is a special type void, which actually marks the absence of a type. We can define a pointer to void; we cannot, however, dereferenceer the pointer. It is useful in low-level application where we want to exam a specific location in memory.

void\* voidPtr = (void\*) 10000;

The void type is useful on one further occasion: to mark that a function does not return a value or misses parameters, see the function section later in this chapter.

In the example below, the memory block has been deallocated, but p has not been set to null. It has become a dangling pointer; it is not null and does not really points at anything. In spite of that, we try to access the value p points at. This is a dangerous operation and would most likely result in a run-time error.

p

dangling pointer

int \*p = malloc(sizeof(int));

\*p = 1;

free(p);

\*p = 2; // Undefined behavor

In the example below, we allocate memory for two pointers, p and q. Then we assign p to q, by doing so we have created a memory leak. There is no way we can access or deallocate the memory block that was pointed at by p. In fact, we deallocate the same memory block twice as both pointers by now points at the same memory block. This dangerous operation will most likely also result in a run-time error.

p

(a)

q

1

2

p

(b)

q

1

2

int \*p = malloc(sizoef(int)); // (a)

int \*q = malloc(sizeof(int));

\*p = 1;

\*q = 2;

p = q; // (b)

free(p); // Undefined behavor. Deallocates the same memory block twice,

free(q); // as p and q points at the same memory block. Also memory leak,

// since the memory block holding 1 cannot be accessed any more.

### Defining our own types

It is possible to define our own type with typedef, which is a great tool for increasing the readability of the code. However, too many defined types tend to make the code less readable. Therefore, a piece of advice would be to use typedef with care.

int i = 1;

typedef unsigned int unsigned\_int;

unsigned\_int u = 2;

typedef int\* int\_ptr;

int\_ptr ip = &i;

typedef unsigned\_int\* uint\_ptr;

uint\_ptr up = &u;

### The Type Size

The operator sizeof gives us the size of a type[[22]](#footnote-22), either by taking the type surrounded by parentheses or a value of the type. The size of a character is always one byte and the signed and unsigned forms of each integral type always have the same size. Otherwise, the sizes are decided by the compiler. In these cases, when a property is not specified by the standard by rather the compiler, we say the the property is implementation dependant. The operator returns a value of the type size\_t, which exact definition is also implementation dependant. However, it is often an unsigned integer.

#include <stdio.h>

void main() {

int intSize1 = sizeof (int);

int intSize2 = sizeof intSize1;

int\* intPtr = &intSize1;

int ptrSize = sizeof intPtr;

int array[3] = {1, 2, 3};

int arraySize = sizeof array;

printf("Integer sizes: %d and %d.\n", intSize1, intSize2);

printf("Pointer size: %d\n", ptrSize);

printf("Array size: %d\n\n", arraySize);

}

## Expressions and Operators

The operations of C are divided into the arithmetic, relational, logical, and bitwise operators as well as simple and compound assignment. Moreover, there is the conditional operator.

In the figure below, + is an operator, 1 and 2 are operands and the whole term is an expression.



Operator

Operand

Operand

### Arithmetic Operators

The arithmetic operators are addition (+), subtraction (-), multiplication (\*), division (/), and modulo (%). The first four operators are equivalent to the four fundamental rules of arithmetic. The operators can take operands of integral and floating types. The last operator, modulo, gives the remainder of integer division. If we mix integral and floating types in the expression, the result will have floating type. The modulo operator, however, only works with integral operands. The last assignment in the following code may give rise to a compiler warning as the result of the division is a double and is converted into an integer.

int a = 10, b = 3, c;

c = a + b; // 13

c = a - b; // 7

c = a \* b; // 30

c = a / b; // 3, integer division

c = a % 3; // 1, remainder

double d = 3.0, e;

e = a / d; // 3.333, floating type

### Pointer Arithmetic

The addition and subtraction operators are also applicable on pointers. It is called pointer arithmetic. An integral value can be added to or subtracted from a pointer. The value of the pointer is then changed by the integral value times the size of the type the pointer points at. As the void type is not really a type, but rather the absence of a type, it has no size. Therefore, we cannot perform pointer arithmetic on void pointers.

In the code below, let us assume that number is stored at memory location 10,000 and that the integer type has the size of four bytes. Then the pointer numberPtr will assume the values 10,000, 10,004, 10,008, and 10,012, not the values 10,000, 10,001, 10,002, and 10,003, as pointer arithmetic always takes the size of the type into consideration.

numberPtr

102

100

101

103

10008

10000

10004

10012

int number = 100;

int\* numberPtr = &number;

numberPtr = numberPtr + 1;

\*numberPtr = number + 1;

numberPtr = numberPtr + 1;

\*numberPtr = number + 2;

numberPtr = numberPtr + 1;

\*numberPtr = number + 3;

There is also possible to subtract two pointers pointing at the same type. The result will be the difference in bytes between their two memory locations divided with the size of the type. Again, pointer arithmetic always considers the type size.

int array[] = {1, 2, 3};

int\* p1 = &array[0];

int\* p2 = &array[2];

int diff = p2 - p1; // 2

The index notation for arrays is equivalent to the dereferenceereeing of pointers together with pointer arithmetic. The second and third lines of the following code are by definition interchangeable.

int array[] = {1, 2, 3};

array[1] = array[2] + 1;

\*(array + 1) = \*(array + 2) + 1;

### Increment and decrement

There are two rather special operators: increment (++) and decrement (--)[[23]](#footnote-23). They add or subtract one from its operand. The operator can be placed before (prefix) or after (postfix) its operand.

int a = 1, b = 1;

++a; // 2, prefix increment

b++; // 2, postfix increment

However, there is a difference between prefix and postfix increment or decrement. In the prefix case below, the subtraction occurs first and the new value is returned; in the postfix case, after the subtraction the original value is returned.

int a = 1, b = 1, c, d;

c = --a; // c = 0, prefix decrement

d = b--; // d = 1, postfix decrement

### Relational Operators

There are six relational operators: equal to (==), not equal to (!=), less than (<), less than or equal to (<=), greater than (>), greater than or equal to (>=)[[24]](#footnote-24). Observe that the equal to operator is constituted by two equal signs rather than one. The operators give a logical value, true (an integer value not equal zero) or false (zero). The operands shall be of integral of floating type.

int i = 3;

double x = 1.2;

int b = i > 0; // true (an integer value not equal to zero, usually one)

int c = (x == 2); // false (always zero)

### Logical Operators

There are three logical operators: not (!), or (||), and and (&&). They take and return numerical values interpreted as logical values.

int i = 3;

int b, c, d, e;

b = (i == 3); // true (not zero)

c = !b; // false (zero)

d = b || c; // true (not zero)

e = b && c; // false (zero)

C applies lazy evaluation (also called short-circuit evaluation), which means that an expression will not be evaluated to a further extent than necessary to find its value. In the following example, the evaluation of the expression is completed when the left expression (i != 0) is evaluated to false. If the left expression is false, the whole expression must also be false because it needs both the left and right sub expressions to be true for the whole expression to be true. This gives that the right expression (1 / i == 1) will never be evaluated and the division with zero will never occur.

int i = 0;

int b = (i != 0) && (1 / i == 1); // false (zero);

### Bitwise Operators

An integer value can be viewed as a bit pattern. Our familiar decimal system has base 10; it can be marked with an index 10.



An integer value can also be viewed with the binary system, which has base 2. A single digit is called a bit, and the integer value is called a bit pattern. A bit may have the values one and zero.



There are four bitwise operations in C: inverse (~), and (&), or (|), and exclusive or (^). Exclusive or means that the result is one if one of its operand bits, but not both, is one. They all operate on integral values on bit level; that is, they examine each individual bit of an integer value.

101010102 101010102 101010102

& 100101102 | 100101102 ^ 100101102 ~ 100101102

----------- ------------ ------------ ------------

= 100000102 = 101111102 = 001111002 = 011010012

int a = 170; // 101010102

int b = 150; // 100101102

int c = a & b; // 100000102 = 13010

int d = a | b; // 101111102 = 19010

int e = a ^ b; // 001111002 = 6010

int f = ~b; // 011010012 = 10510

An integer value can also be shifted to the left (<<) or to the right (>>). Each left shift is equivalent with doubling the value, and each right shift is equivalent with (integer) dividing the value with 2. Overflowing bits are dropped for unsigned values; the behaviour of signed values is implementation dependent. Therefore, I recommend you only perform bitwise operations on unsigned values.

unsigned char a = 172; // 101011002

unsigned char b = a << 2; // 101100002 = 17610

unsigned char c = 166; // 101001102

unsigned char d = c >> 2; // 001010012 = 4110

### Assignment

There are two kinds of assignment operators: simple and compound. The simple variant is quite trivial, one or several variables are assigned the value of an expression. In the example below, a, b, and c are all assigned the value 123.

int a, b, c, d = 123;

a = d;

b = c = d;

The compound variant is more complicated. Let us start with the additional assignment operator. In the example below, the value of a is increased by the value of b; that is, a is given the value 4.

int a = 2, b = 4, c = 2;

a += c; // 4, equivalent to a = a + c.

b -= c; // 2, equivalent to a = a - c.

In a similar manner, there are operations -=, \*=, /=, %=, |=, &=, ^=, <<=, and >>=. Note the difference between a -= 2 and a =- 2. In the former case, a is decreased by 2. In the latter case, a is assigned -2.

### The Condition Operator

The condition operator resembles the if-else statement of the next section. It is the only C operator taking three operands. The first expression is evaluated. If it is true, the second expression is evaluated and its value is returned. If the first expression instead is false, the third expression is evaluated and its value is returned.

int a = 1, b = 2, max;

max = (a > b) ? a : b; // The maximal value of a and b.

Too frequent use of this operator tends to make the code compact and hard to read. A piece of advice is that you restrict your use of the operator to the trivial cases.

### Precedence and Associatively

Given the expression 1 + 2 \* 5, what is its value? It is 11 because we first multiply two with five and then add one. We say that multiplication has a higher precedence than addition.

What if we limit ourselves to one operator? Let us pick subtraction. What value has the expression 8 – 4 – 2? As we first subtract 4 from 8 and then subtract 2, the result is 2. As we evaluate the value from left to right, we say that subtraction is left associative.

Below follows a table over the prorities and associative of the C operators. The first operator in the table has the highest precedence.

|  |  |  |
| --- | --- | --- |
| Group | Operators | Associatively |
| Brackets and fields | () [] -> . | Left to Right |
| Unary operator | ! ~ ++ -- + - (type) sizeof | Right to Left |
| Arithmetic operators | \* / %  + - | Left to Right  Left to Right |
| Shift operators | << >> | Left to Right |
| Relation operators | < <= > >=  == != | Left to Right |
| Bitwise operators | &  ^  | | Left to Right |
| Logical operators | &&  || | Left to Right |
| Conditional operator | ?: | Right to Left |
| Assignment operators | = += -= \*/ /= %= &= ^= |= <<= >>= | Right to Left |
| Comma operator | , | Left to Right |

Note that unary +, -, &, and \* have higher precedence than their binary forms. Note that we always can change the evaluation order of an expression by inserting parantheses at appropriate posititions. The expression (1 + 2) \* 5 has the value 15.

## Statements

There are four kinds of statements in C: the selection, iteration, jump, and expression statements.

|  |  |
| --- | --- |
| Kind | Statements |
| Selection | if, if-else, switch |
| Iteration | for, while, do-while |
| Jump | break, continue, goto, return |
| Expression | expression ; |

### Selection statements

The if statement needs in its simplest form a logical expression to decide whether to execute the following statement. The example below means that the text will be output if i not equal zero.

if (i != 0) {

puts("i does not equal zero");

}

As all values not equal to zero are interpreted as true, the following statement is equivalent with the previous one.fexacltly

if (i) {

puts("i does not equal zero");

}

We can also attach an else part, which is executed if the expression is false.

if (i > 0) {

puts("i is greater than zero");

}

else { puts("i is not greater than zero");

}

Between the if and else part we can insert one or more else if parts.

if (i > 0) {

puts("i is greater than zero");

}

else if (i < 0) {

puts("i is less than zero");

}

else {

puts("i is equal to zero");

}

In the examples above, there is not strictly necessary to surround the output statements with brackets. However, in this book brackets are always used for sake of clarity. Moreover, it is always necessary in case of several statements. The brackets and the code in between are called a block.

if (i > 0) { int j = i + 1;

printf("j is %d\n", j);

}

A warning may be in order. In an if statement, it is perfectly legal to use one equal sign instead of two signs when comparing two values. As one equals sign is used for assignment, not comparison, the variable i in the following code will be assign the value one, and the expression will always be true.

if (i = 1) { // Always true.

// ...

}

One way to avoid the mistake above is to swap the variable and the value. As a value can be compared but not assigned, the compiler will issue an error message and the error will be recognized.

if (1 = i) { // Compile-time error.

// ...

}

The switch statement is simpler than the if statement, but not as powerful. It evaluates the switch expression and jumps to a case statement with the matching (constant) value. If no value matches, it jumps to the default statement, if present. The break statement is used to jump out of a switch or iteration statements. The switch expression must have an integral or pointer type and two case statements cannot hold the same value. We can only have one default statement, and it can be omitted. However, it needs not to be placed at the end of the switch statement, even though I recommend you to do so.

switch (i) { case 1:

puts("i is equal to 1");

break;

case 2:

puts("i is equal to 2");

break;

case 3:

puts("i is equal to 3");

break;

default:

puts("i is not equal to 1, 2, or 3.");

break;

}

It is important to remember the break statement. Otherwise, the execution would simple continue with the code attached to the next case statement[[25]](#footnote-25). However, in C we can use the fact that an omitted break statement makes the execution continue with the next statement to group several case statements together.

switch (i) { case 1:

case 2:

case 3:

printf("i is equal to 1, 2, or 3");

break;

// ...

}

### Iteration statements

Iteration statements, also called loops, iterate one or several statements as long as a certain condition is true; while is the simplest loop. It repeats one statement or a block of statements as long as the given expression is true. The example below uses the while statement to write the numbers one to ten.

int i = 1;

while (i <= 10) { printf("%d\n", i);

++i;

}

The same thing can be done with a do-while statement.

int i = 1;

do { printf("%d\n", i);

++i;

}

while (i <= 10);

However, the do-while statement is less powerful. If the expression is false to begin with, the while statement just skips the repetitions altogether, but the do-while statement must always execute the iteration statement at least once in order to reach the continuation expression at the end.

We can also use the for statement, which is a more compact variant of while. It takes three expressions, separated by semicolons. The first expression initializes the variable, the repetition continues as long as the second expression is true and the third expression is executed at the end of each repetition.

int i;

for (i = 1; i <= 10; ++i) { printf("%d\n", i);

}

Similar to the switch statement, an iteration statement can be interrupted by the break statement. Below is an infinitializere loop (remember that every non-zero value is regarded as true) that is stopped by the break statement.

int i = 1;

while (1) { printf("%d\n", i);

++i;

if (i > 10) { break;

}

}

The same effect can be archived by omitting the second expression of the for statement.

int i;

for (i = 1; ; ++i) { printf("%d\n", i);

if (i == 10) { break;

}

}

An iteration statement can also include a continue statement. It skips the rest of the current repetition. The following example prints the numbers from 1 to 10 with the exception of 5.

int i;

for (i = 1; i <= 10; ++i) { if (i == 5) { continue;

}

printf("%d\n", i);

}

The following example, however, will not work as intended. Because the continue statement will skip the rest of the while block, i will never be updated and we will be stuck in an infinitializere loop. Therefore, I recommend that you use continue with care, or avoid it altogether.

int i = 1;

while (i <= 10) { if (i == 5) { continue;

}

printf("%d\n", i);

++i;

}

### Jump statements

We can jump from one location to another inside the same function by marking the latter location with a label inside the block with the goto statement. The following code is yet another example of how to print the numbers from one to ten, inclusive. As you can, it is not as clear as the previous examples.

int i = 1;

label: printf("%d\n", i);

++i;

if (i <= 10) { goto label;

}

The goto statement is considered to give rise to unstructured code, so called “spagetti code.” I strongly recommended that you omit goto altogether. To put it bluntly: now when you seen it, forget it.

### Expression statements

An expression can form a statement.

a = b + 1; // Assignment operator.

puts("Hello, World!"); // Function call.

In the above examples, we are only interested in the side effects; that a is assigned a new value and that a text is written. We are allowed to write expression statements without side effects; even thought it has no meaning and is likely to be erased by the compiler.

a + b \* c;

### Volatile

A variable can be volatile, which has no functional effect but is rather a notification to the compiler that the variable shall not be subjected to optimization. Let us say we that the memory location 10000 is connected to an input port in a low-level application, and that we want to wait until we receive a signal through the port. If p is not marked as volatile, there is a risk that the compiler would replace the for-loop with another statement since the value of \*p does not seem to change inside the loop.

volatile int\* p = (void\*) 10000;

while (\*p != 0) {

// Wait.

}

## Functions

A function can be compared to a black box. We send in information (input) and we receive information (output). In C, the input values are called parameters and the output value is called the return value.

Input (Parameters)

Function

Output (Return Value)

To start with, let us try the function Square; it takes an integer and returns its square.

int Square(int n) {

return n \* n;

}

void main() {

int i = Square(2); // Square returns 4.

}

2

Square

4

In the example above, the parameter n in the Square definition is called a formal parameter, and the value 3 in the Square call in main is called an actual parameter or an argument.

Now, let us try a more complicated function: SquareRoot takes a value of double type and returns its square root. The idea is that the function iterates and calculates increasingly better root values by taking the mean value of the original value divided with the current root value and the previous root value. The process continues until the difference between two consecutive root values has reached an acceptable tolerance. Just like main[[26]](#footnote-26), a function can have local variables; root and oldRoot hold the current and previous value of the root, respectively.

#include <stdio.h>

double SquareRoot(double value) {

const double EPSILON = 1e-12;

double root = value, oldRoot = value;

while (1) { root = ((value / root) + root) / 2;

if ((oldRoot - root) <= EPSILON) { return root;

}

oldRoot = root;

}

}

void main() {

double input = 16;

printf("SquareRoot of %f: %f\n", input,

SquareRoot(input));

}

### Void Functions

A function does not have to return a value; in that case, we set void as the return type. As mentioned above, void is used to state the absence of a type rather than a type in itself. We can return from a void function by just stating return without a value.

void PrintSign(int value) {

if (value < 0) { puts("Negative.");

return;

}

if (value > 0) { puts("Positive.");

return;

}

puts("Zero");

}

There is not a problem if the execution of a void function reaches the end of the code, it just jump back to the calling function. However, a function not returning void shall always return a value before reaching the end of the function. Otherwise, the compiler often gives a warning.

### Local and global variables

There are two kinds of variables: local and global. A global variable is defined outside a function and a local variable is defined inside a function.

#include <stdio.h>

int global = 1;

void main() {

int local = 2;

printf("Global variable: %d, Local variable: %d.\n",

global, local);

}

A global and a local variable can have the same name. In that case, the name inside the function refers to the local variable. We cannot access the global variable in the inner scope[[27]](#footnote-27) when a local variable with same name is present.

#include <stdio.h>

int number = 1;

void main() {

int number = 2;

printf("Variable: %d.\n", number); // 2

}

One way to distinguish them is to precede the global variable name with ’g\_’.

#include <stdio.h>

int g\_number = 1;

void main() {

int number = 2;

printf("Global variable: %d, Local variable: %d.\n",

g\_number, number); // 1, 2

}

Global variables can only be initialized with constant values.

int g\_number1 = 1; // Right.

int g\_number2 = g\_number1 + 1; // Wrong.

### Inner Blocks with Local Variables

In C, you are allowed to use the bracket notation when initializerializing structures, unions, and arrays, but not when assigning[[28]](#footnote-28). Therefore, the following code will not work.

INT\_PAIR Block(int\* p) {

INT\_PAIR result;

if (p != NULL) {

result = {\*p - 1, \*p + 1}; // Wrong.

return result;

}

else {

result = {0, 0}; // Wrong.

return result;

}

}

One solution would be to a more clumsy code.

INT\_PAIR Block(int\* p) {

INT\_PAIR result;

if (p != NULL) {

result.a = \*p - 1;

result.b = \*p + 1;

return result;

}

else {

result.a = 0;

result.b = 0;

return result;

}

}

However, a more elegant solution would be to use inner blocks with local variables. Each block can define its own variable, which is especially valuable when initializerializing structures.

struct IntPair {

int a, b;

};

struct IntPair Block(int\* p) {

if (p != NULL) {

struct IntPair result = {\*p - 1, \*p + 1};

return result;

}

else {

struct IntPair empty = {0, 0};

return empty;

}

}

An alternative way is to use typedef.

typedef struct {

int a, b;

} INT\_PAIR;

INT\_PAIR Block(int\* p) {

if (p != NULL) {

INT\_PAIR result = {\*p - 1, \*p + 1};

return result;

}

else {

INT\_PAIR empty = {0, 0};

return empty;

}

}

### Call-by-Value and Call-by-Reference

Say we want to write a function for switching the values of two variables.

#include <stdio.h>

void Swap(int number1, int number2) {

int temp = number1; // (a)

number1 = number2; // (b)

number2 = temp; // (c)

}

void main() {

int num1 = 1, num2 = 2;

printf("Before: %d, %d\n", num1, num2); // 1, 2

Swap(num1, num2);

printf("After: %d, %d\n", num1, num2); // 1, 2

}

Unfortunately, it will not work; the variables keep their values. The explanation is that the values of firstNum and secondNum in main are copied into num1 and num2 in Swap. Then num1 and num2 exchange values with the help of temp. However, their values are not copied back into firstNum and secondNum in main.

1

2

main:

1

firstNum

2

Swap:

num2

1

temp

secondNum

num1

2

2

num2

1

temp

num1

2

1

num2

1

temp

num1

(a)

(b)

(c)

The problem can be solved with reference calls. Instead of sending the values of the actual parameters, we send theirs addresses by define the formal parameters as pointers to the type.

#include <stdio.h>

void Swap(int\* numberPtr1, int\* numberPtr2) {

int temp = \*numberPtr1; // (a)

\*numberPtr1 = \*numberPtr2; // (b)

\*numberPtr2 = temp; // (c)

}

void main() {

int num1 = 1, num2 = 2;

printf("Before: %d, %d\n", num1, num2); // 1, 2

Swap(&num1, &num2);

printf("After: %d, %d\n", num1, num2); // 2, 1

}

In this case, we do not send the values of firstNum and secondNum, but rather their addresses. Therefore, num1 and num2 in Swap does in fact contain the addresses of firstNum and secondNum on main. As in the reference section above, we illustrate this with dashed arrows. Therefore, when num1 and num2 exchanges values, in fact the values of firstNum and secondNum are exchanged.

1

2

firstNum

num2

1

temp

secondNum

num1

(a)

2

2

firstNum

num2

1

temp

num1

(b)

2

1

firstNum

num2

1

temp

num1

(c)

### Static, Extern, and Register Variables

In the function below, s\_count (the prefix ‘\_s’ is used for static variables) is a static local variable, which means that it is initialized when the execution of the program starts rather when the function is called. If s\_count was a regular local variable (without the keyword static), the function would, at every call, write that the function has been called once, as s\_count would be initialized to zero at every call.

void KeepCount() {

static int s\_count = 0;

s\_count++;

printf("This function has been called %d times.", s\_count);

}

If we define a global variable in one file, it will be accessable in another file if we declare[[29]](#footnote-29) it as extern. If we omit the extern keyword in the second file, the linker would complain about us defining two global variables with the same name.

File1.c

int g;

File2.c

extern int g;

Finally, a variable can also be marked with the keyword register, which is a notification to the compiler that the variable is suitable to place in a process register rather them the memory. The only function difference is that we cannot get the address of a register variable, since it might be located in a register.

register int i;

int\* p = &i; // Wrong

### Recursion

A function may call itself; it is called recursion. In the following example, the mathematical function factorial (n!) is implemented. It can be defined in two ways. The first definition is rather straightforward. The result of the function applied to a positive integer n is the product of all positive integers up to and including n.



int Factorial(int n) {

int result = 1, i;

for (i = 1; i <= n; ++ i) { result \*= i;

}

return result;

}

An equivalent definition involves a recursive call that is easier to implement.



int Factorial(int n) {

if (n == 1) {

return 1;

}

else { return n \* Factorial(n - 1);

}

}

### Definition and Declaration

It important to distinguish between definition and declaration. For a function, its definition generates code while the declaration is merely an item of information to the compiler. A function declaration is also called a prototype.

When it comes to mutual recursion (two functions calling each other), at least the second of them must have a prototype to avoid compiler warnings. I recommend that you put prototypes for all functions at the beginning of the file (or in a header file, see Section 22.7). In the following example, we use two functions to decide whether a given non-negative integer is even or odd according to the formulas below.





Note that it is not necessary to name the parameters in a function prototype. If names after all are given, they will be ignored by the compiler, which means that the parameters do not need to have the same names in the function declaration and definition. The only restriction is that two parameters cannot have the same name.

int IsEven(int);

int IsOdd(int n);

As C does not include a logical type, we use int to represent true (1) or false (0).

int IsEven(int num) {

if (num == 0) {

return 1;

}

else { return IsOdd(num - 1);

}

}

int IsOdd(int num) {

if (num == 0) {

return 0;

}

else {

return IsEven(num - 1);

}

}

One peculiar thing about prototypes in C is that they can have an unspecified parameter list. If the parameter list is completely omitted in a function prototype, every parameter list is allowed in the call. The following code is will not result in any compile-time errors. However, there is likely to be run-time errors.

void Print();

void main(void) {

Print();

Print(1);

Print(1, 2);

Print(1, 2, 3);

}

The obvious way to avoid the run-time errors is to always state the parameter list in function prototypes. If the function does not have any parameters, it can be stated by the void type. In that case, the function can only be called without parameters.

void Print(void);

### Higher-Order Functions

A function that takes another function as a parameter is called a higher-order function. Technically, it does not take the function itself as a parameter, but rather a pointer to the function. However, the pointer marker (\*) may be omitted. The parameter function is also called a callback function. The following example takes an array of the given size and applies the given function to each integer in the array. ApplyArray a higher order function and Apply is a callback function.

#include <stdio.h>

void ApplyArray(int intArray[], int size, int Apply(int)) {

int index;

for (index = 0; index < size; ++index) { intArray[index] = Apply(intArray[index]);

}

}

int Double(int number) {

return 2 \* number;

}

int Square(int number) {

return number \* number;

}

void PrintArray(int intArray[], int size) {

int index;

for (index = 0; index < size; ++index) { printf("%d ", intArray[index]);

}

printf("\n");

}

void main() {

int numberArray[] = {1, 2, 3, 4, 5, 6, 7, 8, 9, 10};

int arraySize = sizeof numberArray / sizeof numberArray[0];

PrintArray(numberArray, arraySize);

// Doubles every value in the array.

ApplyArray(numberArray, arraySize, Double);

PrintArray(numberArray, arraySize);

// Squares every value in the array.

ApplyArray(numberArray, arraySize, Square);

PrintArray(numberArray, arraySize);

}

An alternative way is to define the type of the callback function. In the following code segment, the APPLY\_FUNC type is a pointer to a function that takes an integer parameter and returns an integer value.

typedef int (\*APPLY\_FUNC)(int);

void ApplyArray(int intArray[], int size,

APPLY\_FUNC pApplyFunc) {

int index;

for (index = 0; index < size; ++index) { intArray[index] = pApplyFunc(intArray[index]);

}

}

One additional point of the example above is how to of find the size of an array: we divide the size of the array with the size of its first value. This function does only work on static arrays, not on dynamically allocated arrays or array given as parameters to functions. A parameter array is in fact converted to a pointer to the array type. The following two function definitions are (by definition) equivalent. This gives that the function must have an extra integer parameter stating the array size.

void PrintArray(int intArray[], int arraySize) {

// ...

}

void PrintArray(int\* intArray, int arraySize) {

// ...

}

## Structures and Linked Lists

A struct is a compound type. Like arrays, it holds several values, called fields. However, in contrast to arrays, the values can hold different types.

A pointer may point at a struct value as well as a simple value, and a struct may have a pointer to another struct value as a field, which in turn points at another struct value, and so one. In this way, a linked list can be constructed. The list must end eventually, so the last pointer points at null. A pointer to the next cell in the list is called a link.

struct Cell {

int value;

struct Cell\* nextPtr;

};

struct Cell cell3 = {3, NULL};

struct Cell cell2 = {2, &cell3};

struct Cell cell1 = {1, &cell2};

1

value

nextPtr

2

value

nextPtr

3

value

nextPtr

NULL

### Stacks and Linked Lists

A stack is very valuable in a number of applications and it can be implemented with a linked list. We can add a value on top of the stack, we can inspect or remove the topmost value, and we can check whether the stack is empty. However, we cannot do anything to the values that are not on top. The function adding a new value on top of the stack is called push and the function removing it is called pop. Let us say that we push our stack three times with the values 1, 2, and 3. Then we can only access the topmost value, 3, and not the two below, 1 or 2.

1

2

3

push 1

push 2

push 3

Cell.h

struct Cell { int m\_value;

struct Cell\* m\_nextPtr;

};

typedef struct Cell CELL;

void CellInitializer(CELL\* cellPtr, int value, CELL\* nextPtr);

void CellSetValue(CELL\* cellPtr, int value);

int CellGetValue(CELL\* cellPtr);

void CellSetNext(CELL\* cellPtr, CELL\* nextPtr);

CELL\* CellGetNext(CELL\* cellPtr);

Cell.c

#include "Cell.h"

void CellInitializer(CELL\* cellPtr, int value, CELL\* nextPtr) {

cellPtr->m\_value = value;

cellPtr->m\_nextPtr = nextPtr;

}

void CellSetValue(CELL\* cellPtr, int value) {

cellPtr->m\_value = value;

}

int CellGetValue(CELL\* cellPtr) {

return cellPtr->m\_value;

}

void CellSetNext(CELL\* cellPtr, CELL\* nextPtr) {

cellPtr->m\_nextPtr = nextPtr;

}

CELL\* CellGetNext(CELL\* cellPtr) {

return cellPtr->m\_nextPtr;

}

Main.c

#include <stdlib.h>

#include "Cell.h"

void main() {

CELL\* cellPtr1 = malloc(sizeof(CELL));

CELL\* cellPtr2 = malloc(sizeof(CELL));

CELL\* cellPtr3 = malloc(sizeof(CELL));

CellInitializer(cellPtr3, 3, NULL);

CellInitializer(cellPtr2, 2, cellPtr3);

CellInitializer(cellPtr1, 1, cellPtr2);

}

The following struct is created by main above.

1

m\_value

m\_nextPtr

2

m\_value

m\_nextPtr

3

m\_value

m\_nextPtr

NULL

cellPtr1

cellPtr2

cellPtr3

However, there is one more thing to think about. What happens if we run out of dynamic memory or try to access the topmost value of an empty stack? We can deal with the problem in some different ways, everything from ignoring it to abort the execution. In this book, I have limited the error handling of memory shortage to the use of assert, which is a macro (see Section 22.9) that takes a logical value and aborts the execution if the value is false, adding an error message with information about the file name and the code line number. To keep things simple, let us use that function in the stack structure too.

Stack.h

struct Stack { CELL\* m\_firstCellPtr;

};

typedef struct Stack STACK;

void StackInitializer(STACK\* stackPtr);

void StackClear(STACK\* stackPtr);

void StackPush(STACK\* stackPtr, int value);

void StackPop(STACK\* stackPtr);

int StackTop(STACK\* stackPtr);

int StackIsEmpty(STACK\* stackPtr);

Stack.c

#include <stdlib.h>

#include <assert.h>

#include "Cell.h"

#include "Stack.h"

void StackInitializer(STACK\* stackPtr) {

stackPtr->m\_firstCellPtr = NULL;

}

void StackClear(STACK\* stackPtr) {

CELL\* cellPtr = stackPtr->m\_firstCellPtr;

while (cellPtr != NULL) { CELL\* tempCellPtr = cellPtr;

cellPtr = CellGetNext(cellPtr);

free(tempCellPtr);

}

}

void StackPush(STACK\* stackPtr, int value) {

CELL\* newCellPtr= malloc(sizeof(CELL));

assert(newCellPtr!= NULL);

CellInitializer(newCellPtr, value, stackPtr->m\_firstCellPtr);

stackPtr->m\_firstCellPtr = newCellPtr;

}

void StackPop(STACK\* stackPtr) {

CELL\* tempCellPtr;

assert(stackPtr->m\_firstCellPtr != NULL);

tempCellPtr = stackPtr->m\_firstCellPtr;

stackPtr->m\_firstCellPtr = CellGetNext(stackPtr->m\_firstCellPtr);

free(tempCellPtr);

}

int StackTop(STACK\* stackPtr) {

assert(stackPtr->m\_firstCellPtr != NULL);

return CellGetValue(stackPtr->m\_firstCellPtr);

}

int StackIsEmpty(STACK\* stackPtr) {

return (stackPtr->m\_firstCellPtr == NULL);

}

Main.c

#include <stdlib.h>

#include "Cell.h"

#include "Stack.h"

void main() {

STACK stack;

StackInitializer(&stack);

StackPush(&stack, 1);

StackPush(&stack, 2);

StackPush(&stack, 3);

StackClear(&stack);

}

The stack in main above gives rise to the following struct. However, when StackClear is called, it deallocates the allocated memory. All memory allocated with malloc or calloc must (or at least should) be deallocated with free.

m\_nextPtr

3

2

m\_value

m\_nextPtr

1

m\_value

m\_nextPtr

NULL

m\_value

m\_firstCellPtr

### Unions

A union is a compound type like struct. The difference is that while a struct places each of its members on different memory addresses in order for them to not intervene with each other, a union locates each of its members on the same address. This implies that when one member is assigned a value, it overwrites the other values. Unions are mostly used in low-level applications. In the code below, the two structures overlap each other.

struct \_Bits8Registers { // A character always has the size of 1 byte (8 bits).

unsigned char ah, al, bh, bl, ch, cl, dh, dl;

};

struct \_Bits16Registers { // Assuming short has a size of 2 bytes (16 bits).

unsigned short ax, bx, cx, dx;

};

union \_Registers {

struct \_Bits8Registers bits8Registers;

struct \_Bits16Registers bits16Registers;

};

### Bitfields

A bitfield is a struct with the possibility to define the size of the members in bits, which is useful in low-level applications. Only integral types can have bits, and the total number of bits must not exceed the number of bits an integer (eight times sizeof (int)). In the code below, the whole struct shares one byte (eight bits). The second field does not have a value; it only marks that the second bit shall be ignored.

struct Communicate {

int alert : 1;

int :1;

int read : 3;

int write : 3;

};

## Structures with function pointers

A struct can hold pointers to functions as well as values. However, in order to obtain the specific struct value, we have to add its address a pointer parameter[[30]](#footnote-30) in the function call.

#include <stdio.h>

struct Person { char name[100];

int age;

void (\*Print)(struct Person\*);

};

void PrintNameAndAge(struct Person\* personPtr) {

printf("Name: %s\n", personPtr->name);

printf("Age: %d\n\n", personPtr->age);

}

void PrintOnlyName(struct Person\* personPtr) {

printf("Name: %s\n\n", personPtr->name);

}

void PrintOnlyAge(struct Person\* personPtr) {

printf("Age: %d\n\n", personPtr->age);

}

void main() {

struct Person adam = {"Adam", 10, &PrintNameAndAge};

struct Person bertil = {"Bertil", 20, &PrintOnlyName};

struct Person ceasar = {"Ceasar", 30, &PrintOnlyAge};

adam.Print(&adam); // Prints the name and age.

bertil.Print(&bertil); // Prints only the name.

ceasar.Print(&ceasar); // Prints only the age.

}

## The Preprocessor

The preprocessor is a tool that precedes the compiler in interpreting the code. The #include directive is one of its parts. It opens the file and includes its text. So far, we have only included system header files, whose names are surrounded by arrow brackets (for instance, <StdIO.h>). Later on, we will include our own header files. Then we will use parentheses instead of arrow brackets. The difference is that the pre-processor looks for the system header files in a special system directory while it looks for our header files in the local directory (usually the source code directory).

Another part of the pre-processor is the macros. They acts like functions with the difference that they do not perform any type checking, they just replace text. We have already used the assert macro. A macro is introduced with the #define directive and is often written with capitals.

#define ADD(a, b) ((a) + (b))

printf("%d\n", ADD(1 + 2, 3 \* 4)); // 15

Unlike regular C code, if we add a page-break we have to mark it by a backslash.

#define ADD(a, b) ((a) + \

(b))

It is also possible to perform conditional programming by checking the value of macros. In the following example, we define a system integer according to the underlying operational system.

#ifdef WINDOWS

#define SYSINT int

#endif

#ifdef LINUX

#define SYSINT unsigned int

#endif

#ifdef MAC

#define SYSINT long int

#endif

SYSINT iOpData = 0;

Condition programming can also be used as a third level of comments (above line and block comments, see Section 22.9). In the following code the whole segment between the #if and #endif directives will be omitted, regardless of the comments, as zero always is interpreted as false.

#if 0

int Square(int value)

{ return value \* value; // Square.

}

#endif

## The Standard Library

The C standard library hold around 200 functions and some macros divided into several sections.

### File Management

We can open, write to, read from, and close files with the help of file pointers. A file pointer is a pointer to a value of the FILE type; it can be considered a connection to the file. The program below reads a series of integers from the text file Input.txt and writes their squares to the file Output.txt. The function feof returns zero when there is no more value to be read from the file. Finally, we must not forget to close the file.

#include <stdlib.h>

#include <assert.h>

#include <stdio.h>

void main(void) {

FILE\* inFile = fopen("Input.txt", "r");

FILE\* outFile = fopen("Output.txt", "w");

assert(inFile != NULL);

assert(outFile != NULL);

while (!feof(inFile)) { double value;

fscanf(inFile, "%lf", &value);

fprintf(outFile, "%f\n", value \* value);

}

fclose(inFile);

fclose(outFile);

}

The text files are written in plain texts and can be viewed by the editor.

Input.txt

1

2

3

4

5

Output.txt

1

4

9

16

25

We can also read and write binary data. The program below writes the number one to ten to the file Numbers.bin and then reads the same series of values from the file. The functions write and read takes the address of the value to be read or written and the size of the value in bytes. They return the number of bytes actually read or written. When reading, we can check whether we have reached the end of the file by counting the number of read bytes; if it is zero, we have reached the end.

BinaryFile.c

#include <stdlib.h>

#include <assert.h>

#include <stdio.h>

void main(void) {

int index, value;

FILE \*outFile = fopen("Number.bin", "w"), \*inFile;

assert(outFile != NULL);

for (index = 1; index <= 10; ++index) { fwrite(&index, 1, sizeof index, outFile);

}

fclose(outFile);

inFile = fopen("Number.bin", "r");

assert(inFile != NULL);

while (fread(&value, 1, sizeof value, inFile) > 0) { printf("%d\n", value);

}

fclose(inFile);

}

The values are stored in compressed form in the binary file Numbers.bin, why they are not readable in the editor. Here is a screen dump of the file.



### Program Parameters

It possible to start the program execution with parameters. In main, the argc parameter holds the number of input strings, and argv holds the string themselves.

#include <StdIO.h>

void main(int argc, char\* argv[]) {

int index;

printf("argc: %d\n", argc);

for (index = 0; index < argc; ++index) {

printf("argv[%d]: %s\n", index, argv[index]);

}

}

The parameters are given when the program executes. Note that the value of argv[0] always is the program path.



### Environment Functions

There is a set of functions that communicates with the surrounding environment: exit quits the execution and sends an integer value to the environment; atexit states a function that is called when exit is called; getenv reads the value of an environment variable, and system executes a system command. There are two macros [EXIT\_SUCCESS](http://www.cplusplus.com/reference/clibrary/cstdlib/EXIT_SUCCESS/) and [EXIT\_FAILURE](http://www.cplusplus.com/reference/clibrary/cstdlib/EXIT_FAILURE/) that can used with the exit function.

#include <StdIO.h>

#include <stdlib.h>

void ExitFunction(void) {

printf("The Program is exiting.\n");

}

void main() {

char\* path = getenv("PATH");

system("dir \*.c");

atexit(&ExitFunction);

if (path != NULL) {

printf("Path: %s.\n", path);

exit(EXIT\_SUCCESS);

}

else {

printf("Path not found.\n");

exit(EXIT\_FAILURE);

}

}

Another way to send an exit message to the surrounding environment is to define int as main’s return type. It has the same effect as calling exit in main (except that the exit function will not be called).

#include <StdIO.h>

#include <StdLib.h>

int main() {

char\* path = getenv("PATH");

if (path != NULL) {

printf("Path: %s.\n", path);

return EXIT\_SUCCESS;

}

else {

printf("Path not found.\n");

return EXIT\_FAILURE;

}

}

### Searching and Sorting

The bsearch function performs a binary search through a sorted list of values. We need to add a pointer function parameter specifying a function comparing two values.

#include <StdLib.h>

int CompareIntegers(const int\* intPtr1, const int\* intPtr2) {

return (\*intPtr1 < \*intPtr2) ? -1 : (((\*intPtr1 > \*intPtr2) ? 1 : 0));

}

void main() {

int intArray[] = {1, 2, 3, 4, 5, 6, 7, 8, 9, 10},

arraySize = sizeof intArray / sizeof intArray[0],

value = 6;

int\* resultPtr = bsearch(&value, intArray, arraySize,

sizeof (int), &CompareIntegers);

if (resultPtr != NULL) {

int index = ((int\*) resultPtr) - intArray;

printf("Index %d.\n", index);

}

else {

printf("Not found.\n");

}

}

There is also the qsort function that performs a fast sorting.

#include <StdLib.h>

void main() {

int index, intArray[] = {7, 2, 5, 9, 3, 1, 10, 8, 4, 6},

arraySize = sizeof intArray / sizeof intArray[0],

key = 5;

qsort(intArray, arraySize, sizeof (int), &CompareIntegers);

for (index = 0; index < arraySize; ++index) {

printf("%d\n", intArray[index]);

}

}

### Functions with Variable Number of Parameters

In C, there is possible to define functions with a variable number of parameters (such as printf and scanf). One condition is that the function has at least one regular parameter in order for the va\_start macro to hook up to. The va\_arg macro reads the arguments and the va\_end macro ends the reading. Note that the macros do not give any information about the number or types of the arguments. In the following code, we use the first argument count to count the number of arguments thereafter.

#include <stdio.h>

#include <stdarg.h>

void Print(int count, ...) {

va\_list arg\_list;

va\_start(arg\_list, count);

while (count > 0) { int value = va\_arg(arg\_list, int);

printf("%d ", value);

--count;

}

printf("\n");

va\_end(arg\_list);

}

void main(void) {

Print(1, 1);

Print(2, 1, 2);

Print(3, 1, 2, 3);

}

### String Management

There are several functions dealing with strings: strcpy copies a string; strcat adds a string; strlen gives the length of a string; strcmp compares two strings; strchr and strrchr gives the first and last occurrence of a character in a string. All functions look for the finishing null character (‘\0’), strcmp returns a value less than zero if the first string is smaller than second string, a value greater than zero if the second string is larger, and zero if they are equals. Note that the comparison continues until the finishing zero character has been found, even though the character arrays may be longer. As strchr and strrchr return a pointer to the found character (or null if it was not found), we can calculate the array index by subtracting the pointers.

int FirstIndexOf(char s[], char c) {

char\* p = strchr(s, c);

return (p != NULL) ? (p - s) : -1;

}

int LastIndexOf(char s[], char c) {

char\* p = strrchr(s, c);

return (p != NULL) ? (p - s) : -1;

}

void main(void) {

char s1[] = "Hello", s2[] = "World";

char t1[20], t2[20];

strcpy(t1, s1);

strcat(t1, s2);

sprintf(t2, "%s, %s!", s1, s2);

printf("t1: \"%s\", length of t1: %d, t2: \"%s\"\n",

t1, strlen(t1), t2);

printf("Comparing t1 and t2: %d\n", strcmp(t1, t2));

printf("First Index of 'l' in t2: %d, last index of 'l' in t2: %d.\n",

FirstIndexOf(t2, 'l'), LastIndexOf(t2, 'l'));

}

It is also possible to write to a string with sprintf and read from a string with sscanf. These function come in handy when to cast between strings and numerical values.

void main() {

int i;

double x;

char s1[20], s2[20];

sprintf(s1, "%d", 123);

sprintf(s2, "%f", 123.456e-3);

printf("s1 = \"%s\", s2 = \"%s\"\n", s1, s2);

sscanf(s1, "%d", &i);

sscanf(s2, "%lf", &x);

printf("i = %d, x = %f\n", i, x);

}

The functions return the number of values written or read, which gives that we can check that the conversion worked.

if (sscanf(s2, "%lf", &x) == 1) {

printf("x = %f\n", x);

}

else {

printf("Could not cast the string \"%s\".\n", s2);

}

Moreover, we can also count the number of read character and decide whether to whole string has been read.

sscanf(s2, "%lf%n", &x, &charCount);

if (charCount == strlen(s2)) {

printf("x = %f\n", x);

}

else {

printf("Could not cast the string \"%s\".\n", s2);

}

### Memory Management

There is also a set of function for general memory management: memset sets each byte of a memory block to a given value; memcmp compares the memory blocks; memcpy and memmove copies a memory block. The difference between them is that memcpy copies the bytes of the block directly when memmove uses a buffer, which gives that memcpy is faster and memmove is safer when copying overlapping blocks.

These functions work similar to their string counterparts. However, note that they do not look for a finishing character. Instead, the sizes of the function blocks are given as a parameter.

#include <StdLib.h>

void main() {

int a1[3], a2[3];

memset(a1, 1, sizeof a1);

memcpy(a2, a1, sizeof a2);

printf("Comparing a1 and a2: %d\n", memcmp(a1, a2, sizeof a2));

}

### Mathematical Functions

C standard library includes the trigonometric cos, sin, tan, acos, asin, and atan, and atan2; the hyperbolic functions cosh, sinh, and tanh; ceil and floor that rounds a floating value to the closest higher and lower value, respectively; exp returning the natural exponent; log returning the natural logarithm; log10 returning the logarithm of ten; sqrt returning the square root, and pow returning the power of two values. All of these functions takes and returns a double value (except atan2 and pow that takes two values). The trigonometric functions deal with angles in radians.

When converting a floating value to an integral value, it will always be truncated; that is, rounded to the integer value closest to zero. One way to find the closest integer value is to add 0.5 to the double value before converting it.

printf("double %f, rounded to the closest integer value: %d\n",

x, ((int) (x + 0.5)));

### Long Jumps

Regular jumps with goto can only occur in the same function body; long jumps with the standard function setjmp and longjump can, on the other hand, occur throw call sequences; setjmp saves the state of the program execution and longjmp uses that information to restore that state, which gives the illusion of a backwards jump.

In the program below, the user inputs a number, if it does not equal zero its inverse is return. If the number does equal zero, a long jump occurs. The call to setjump returns zero when the jump is set and a non-zero value when a long jump occurs[[31]](#footnote-31).

#include <stdio.h>

#include <setjmp.h>

jmp\_buf env;

double Divide(double numerator, double denominator) {

if (denominator == 0) { longjmp(env, -1); // -1 to represents an error state.

}

return numerator / denominator;

}

double Inverse(double number) {

return Divide(1.0, number);

}

void main() {

if (setjmp(env) == 0) { double value;

printf("Value: ");

scanf("%lf", &value);

printf("Inverse: %f\n", Inverse(value));

}

else {

puts("Division by zero.");

}

}

### Time

The time function gives the number of seconds since January 1, 1970. That value can then be used to fill a tm struct with values regarding local date and time by calling the localtime function. If we call instead call gmtime, the [Greenwich Mean Time](http://sv.wikipedia.org/wiki/Greenwich_Mean_Time) (Coordinated Universal Time) will be given. Note that the tm\_year field gives the year starting from 1900.

#include <Time.h>

void main() {

time\_t t = time(NULL);

struct tm time;

char \*months[] = {"Jan", "Feb", "Mar", "Apr",

"May", "Jun", "Jul", "Aug",

"Sep", "Oct", "Nov","Dec"};

char \*days[] = {"Sun", "Mon", "Tue", "Wed",

"Thu", "Fri", "Sat"};

localtime\_s(&time, &t);

printf("%02d:%02d:%02d\n", time.tm\_hour, time.tm\_min, time.tm\_sec);

printf("%s %s %d, %d\n", days[time.tm\_wday], months[time.tm\_mon],

time.tm\_mday, 1900 + time.tm\_year);

printf("Day of year: %d.\n", time.tm\_yday);

printf("Daylight saving time: %s.\n", time.tm\_isdst ? "Yes" : "No");

}

It is also possible to measure the period between two events.

#include <StdIO.h>

#include <Time.h>

void main() {

time\_t t1 = time(NULL);

int index;

for (index = 0; index < 100000; ++index) {

printf("%d\n", index);

}

{ time\_t t2 = time(NULL);

printf("The loop took %f seconds.\n", difftime(t2, t1));

}

}

### Random Numbers

It is possible to generate a sequence of pseudo-random numbers with the rand function. In order to initialize the sequence we need to sow a random seed with the srand function. One way to generate different seeds is to use the time function. The [RAND\_MAX](http://www.cplusplus.com/reference/clibrary/cstdlib/RAND_MAX/) macro gives the largest possible random numbers; however, usually we want random numbers in a certain interval. The following code generates five sequences of ten numbers between 1 and 100.

#include <StdLib.h>

void main() {

int iIndex1, iIndex2;

for (iIndex1 = 0; iIndex1 < 5; ++iIndex1) {

srand((unsigned int) time(NULL));

for (iIndex2 = 0; iIndex2 < 10; ++iIndex2) {

int randomNumber = rand() % 100 + 1;

printf("%d\n", randomNumber);

}

printf("\n");

}

}

### Limits of Integral and Floating Types

Since the type sizes are implementation dependent, there are a set of macros defining the minimum and maximum values of the integrated and floating types. The minimum values of unsigned integral types are always zero.

#include <stdio.h>

#include <limits.h> // The integral type limit constants.

#include <float.h> // The floating type limit constants.

void main() {

printf("Minimum signed char: %d\n", SCHAR\_MIN);

printf("Maximum signed char: %d\n", SCHAR\_MAX);

printf("Minimum signed short int: %d\n", SHRT\_MIN);

printf("Maximum signed short int: %d\n", SHRT\_MAX);

printf("Minimum signed int: %d\n", INT\_MIN);

printf("Maximum signed int: %d\n", INT\_MAX);

printf("Minimum signed long int: %ld\n", LONG\_MIN);

printf("Maximum signed long int: %ld\n\n", LONG\_MAX);

printf("Maximum unsigned char: %u\n", UCHAR\_MAX);

printf("Maximum unsigned short int: %u\n", USHRT\_MAX);

printf("Maximum unsigned int: %u\n", UINT\_MAX);

printf("Maximum unsigned long int: %lu\n\n", ULONG\_MAX);

printf("Minimum float: %e\n", FLT\_MIN);

printf("Maximum float: %e\n", FLT\_MAX);

printf("Minimum double: %e\n", DBL\_MIN);

printf("Maximum double: %e\n", DBL\_MAX);

printf("Minimum long double: %le\n", LDBL\_MIN);

printf("Maximum long double: %le\n", LDBL\_MAX);

}

### Character Functions

There are a set of functions for classifying a character.

if (isdigit(c)) {

printf("%c is a digit.", c);

}

Below follows a table of the character functions. The ASCII codes are stated within brackets.

|  |  |
| --- | --- |
| Function | Checks |
| **isalpha** | Letter: a-z and A-Z. |
| **islower** | Lower case letter: a-z. |
| **isupper** | Upper case letter: A-Z. |
| **isdigit** | Digit: 0-9. |
| **Isxdigit** | Hexadecimal digit: a-f, a-f, and 0-9. |
| **Isalnum** | Alphanumeric character: letter or digit. |
| **Isspace** | White-space character: space (0x20), horizontal tab (0x09), vertical tab (0x0B), new-line (0x0A), form-feed (0x0D), or carriage return (0x0C). |
| **Iscntrl** | Control character: characters with ASCII codes 0x00 to 0x0F, inclusive, and delete (0x1F). |
| **Isprint** | Printable character: any character that is not a control character. |
| **Isgraph** | Graphical character: any printable character that is not a white-space character. |
| **ispunct** | Punctuation character: any graphical character that is not an alphanumeric character. |

There is also the tolower and toupper functions that returns the character in lower and upper case, respectively.

printf("%c in lower case: %c, and in upper case: %c.",

c, tolower(c), toupper(c));

## Further Reading

I recommend *The C Programming Language* by Brian Kernighan and Dennis Ritchie (2nd Edition, Prentice Hall, 1988). It describes ANSI C and the standard library in a short and consist way, and holds an appendix about the ANSI C standard.

If you want to learn more about C programming, I recommend *C How to Program* by P. J. [Deitel](http://www.amazon.com/s/ref=ntt_athr_dp_sr_1?_encoding=UTF8&sort=relevancerank&search-alias=books&ie=UTF8&field-author=Paul%20J.%20Deitel) (6th Edition, Prentice Hall, 2009), which which discusses in depth the theory as well as practice of programming in C, with many clear and comprehensive examples. C in a Nutshell by [Peter Prinz, Tony Crawford](http://shop.oreilly.com/product/9780596006976.do#tab_04) (O'Reilly Media, 2005) and Practical C Programming (3th edition, O'Reilly Media, 1997) are also good choices.

When you have mastered the basic of C programming, I recommend Expert C Programming by P. Van Der Linden (Prentice Hall, 1994). Do not worry; you do not have to be an expert to read it.

# A Crash Course in Java

# Foreword

The compiler is made up of several components:

#### The Type System

Every programming language has a type system, which determines the properties of its types.

#### The Symbol Table

In many languages, the source code can hold a definition of a variables that is later referred to. In that case the, the variable (its name, type, and potential value) needs to be stored in a symbol table.

#### Scanning

The scanner takes a stream of characters and put them together into a sequence of *tokens*; that is, the least meaningful parts of the source code. For instance, the characters ‘i' and ‘f’ is put together into the keyword **if**, and the characters ‘!’ and ‘=’ is put together into the operator **not\_equal**.

#### Parsing and Middle Code Generation

The syntax of a programming language is often defined by a grammar, and the parser checks whether the token sequence generated by the scanner is accepted by the grammar. Moreover, the parser does also generate middle code.

#### Middle Code Optimization

The middle code generated by the parser often holds unnecessary instructions that needs to be removed.

#### Target Code Generation

When the middle code has been generated and optimized, it is time to generate the target code. In this book, we actually generate target code for two formats: Intel x86 assembler source code and executable code in the 16-bits COM-format.

#### Register Allocation

When the target code has been generated, the registers need to be allocated. We use a graph coloring algorithm to do that.

#### Linking

When each source file has been completely compiled, they need to be linked together into a executable file format.

**The Preprocessor**

The C programming language comes with a preprocessor that needs to be.

**The Standard Library**

The C programming language also comes with a standard library.

3. I have another book proposal for you, this time about compiler technology. The book will describe the design and implementation of a C++ compiler (including the preprocessor, linker, and standard library), with all code given.

So far, I have written a C compiler and I plan to extend it to a C++ compiler in the near future. Even though C++ is a much larger language than C, it not too much extra job. Most of the job has already been done with the C compiler.

The idea of the book is basically the same as the Windows book: I explain the details and provide all source code of the compiler. The source code written in Java with CUP and JLex and generates both assembler code and executable code.

The books below are the latest editions of three classic books. They describe in detail the basic features of compiler construction and to some extent advanced features. The structure of these books is similar to mine, one can say that I have followed these books. However, while these books give many small examples, my book describes in each chapter the features necessary to be included in a C++ compiler.

Alfred V. Aho, Lam, Monica S. Lam, Ravi Sethi, Jeffrey D. Ullman. Compilers – Principles, Techniques, and Tools, 2nd edition. Prentice Hall, 2006.

Keith Cooper, Linda Torczon. Engineering a Compiler, 2nd edition. Morgan Kaufman, 2011.

Charles N. Fischer, Ron K. Cytron, Richard J. LeBlanc. Crafting a Compiler. Pearson Education, 2009.

The following books describe compiler construction in Java, C, and ML. They describe a compiler for a smaller language with excerpt from the code given in the book (the complete code is downloadable). However, these books are briefer than the books above, and describe the compiler features in less detail.

Andrew W Appel. Modern Compiler Implementation in Java, 2nd edition. Cambridge University Press, 2002.

Andrew W Appel. Modern Compiler Implementation in C, 2nd edition. Cambridge University Press, 2004.

Andrew W Appel. Modern Compiler Implementation in ML, 2nd edition. Cambridge University Press, 2004.

The following book could be said to be closely related to mine, as it also described the code for a C compiler. However, this book is more or less unreadable. It presents the code, which (in my option) is unstructured, and it does not include much text describing the code. Hopefully, my book will describe the topic in a clearer way. I also plan to describe the theory of a compiler, not just the code to implement it. Moreover, this book does not include the preprocessor, linker, or standard library, and the code is written in C, not Java.

Christopher W. Fraser, David R. Hanson. A Retargetable C Compiler : Design and Implementation. Benjamin Cummings, 1995.

This books deals with the more advanced parts of compiler construction; that is, compiler optimization. Even though I will include some optimization at the end of my book, it does not compete with this book.

Steven S. Muchnick. Advanced Compiler Design & Implementation. Morgan Kaufmann, 2003.

Here is a suggestion for the chapters:

1. Introduction. Introduces the compiler phases described in the following chapters by demonstrating a compiler for a small toy language generating MIPS-code executable in the SPIM simulator.

2. The Scanner. The scanner is a relatively small part of the compiler. Its task is to put together characters into tokens (the smallest significant parts of the source code). Examples are as key words, operators, and numerical values. The scanner is written in JLex, a lexical generator for Java, based on Lex for C.

3. The Parser. The parser, on the other hand, is a large part of the compiler. Its task is to confirm that the given tokens (generated by the scanner) agree with the syntax of the programming language, which is represented by a set of grammatical rules. The parser is defined in CUP, a syntax generator for Java, based on Yacc for C. Each rule can also be equipped with code dealing with type checking and target code generation. However, I will have try to omit as much as possible of the code in this chapter. Most of the code is made up of calls to methods defined in later chapters.

4. Declarations and the Symbol Table. C++ has a rather complicated declaration system with aggregated types such as classes, structures and arrays with a corresponding complicated syntax. All defined variables and functions are stored in the symbol table, which is a hierarchical structure matching the program structure.

5. Type Checking. C++ has a rather large set of operators with complicated rules that need to be checked.

6. Intermediate Code Generation. When the types of expressions and statements are checked, three-address-code are generated, which is a simple intermediate language used to represent the code internally. Type conversation is also included in this chapter.

7. Static and Dynamic Initialization. In C++, variables can be initialized. If the variable is static, the data shall be generated and placed in the static area of the final target code. If it is dynamic, the initialization will result in a series of assignments. One thing that complicates the issue is that it is possible to initialize hierarchical structures made up by structures and arrays with one flat list.

8. Intermediate Code Optimization. The intermediate code can be optimized. For instance, code that is never reached and assignment of variables that are never used shall be removed.

9. Target Code Generation. The target code of the C++ compiler of this book is Intel x86, which is harder to deal with than the MIPS code the first chapter. It holds a few registers and registers of different sizes overlap. Therefore, the register allocation process needs to closely keep track on which variable values that are currently stored in the registers.

10. The Preprocessor. Before the actually compilation starts, the source code has been traversed by the preprocessor that replaces macros with text, includes header files, and provides conditional programming.

11. The Linker. When the target code has finally been generated, it becomes stored into an object file. As the source code can be distributed over several files, the target code need to be merged into one executable file. This is the task of the linker, it merges the code and data area, resolve the static and extern references and generate the executable file.

12. The C Standard Library, made up by functions and macros, which will be included by the linker in the final executable file.

13. The C++ Standard Library, made up by classes, to a certain extent is based on C standard library.

The compiler is made up of several phases:

Preprocessor

Source Code

Parser

Processed Code

Syntax Tree Optimister

Syntax Tree

Middle Code Generator

Optimized Syntax Tree

Middle Code Optimizer

Middle Code List

Optimized Middle Code List

Scanner

Symbol Table

Preprocessor

Source Code

Parser

Processed Code

Syntax Tree Optimister

Syntax Tree

Middle Code Generator

Optimized Syntax Tree

Middle Code Optimizer

Middle Code List

Optimized Middle Code List

Scanner

Symbol Table

1. Then name Yacc is an abbreviation for Yet Another Compiler-Compiler. However, it is pronounced as the animal yak. Therefore, the name Bison referrers to the rather similar animal. [↑](#footnote-ref-1)
2. In other languages, which syntax includes the keyword then, it is called as the if-then-else problem. [↑](#footnote-ref-2)
3. Technically, it is possible to store null values in the map, in which case put would return null even when the key is present in the map. However, in this book, we never store null values in any map. [↑](#footnote-ref-3)
4. As we all know, goto has no place in well-structured programs. Labels and goto are included in C of historical reasons only. More recent languages have omitted goto. [↑](#footnote-ref-4)
5. This due to historical reasons. In the early versions of C, there was no void type. Instead, char was used a generic type in pointer expressions. Therefore, it was (and still is) important that a char always holds one byte. [↑](#footnote-ref-5)
6. In some books, the process of generating the low-level executable code is called the target code generation. However, in this book I have chosen the term object code generation in order not confuse it with source and targets in jump statements. That would have resulted in rather awkward variable names, like targetTargetCode. [↑](#footnote-ref-6)
7. Actually, a function’s start address and its the entry point differs in one case only: when main is defined with the argc and argv parameters and is called recursively, which is allowed in the C standard. In that case, the first part of the main function holds code for initializing argc and argv, which shall not be executed in a recursive call. [↑](#footnote-ref-7)
8. As mention in Section 17.1, a function’s start address and its the entry point differs only when main is defined with the argc and argv parameters and is called recursively. [↑](#footnote-ref-8)
9. We could add code that modifies the code in a way similar to generateJumpInfo in Section **Fel! Hittar inte referenskälla.**. However, I feel that it is too much trouble for saving a few bytes. After all, most function calls are performed by long jumps. [↑](#footnote-ref-9)
10. BCPL stands for Basic Command Programming Language and was developed at join effort at London University and Cambridge University. Later on, BCPL has jokely been said to be an acronym for the Before C Programming Language. [↑](#footnote-ref-10)
11. Actually, their program wrote "hello, world\n" without a capital letter or an exclamation mark. Their example was inherited from the internal 1974 memorandum Programming in C: A Tutorial by [Brian Kernighan](http://en.wikipedia.org/wiki/Brian_Kernighan). The oldest known instance of the usage of the words ”hello” and ”world” together occurred in Kernighan’s 1972 tutorial *Introduction to the Language* [B](http://en.wikipedia.org/wiki/B_(programming_language)), even though the syntax was different from C. [↑](#footnote-ref-11)
12. Technically, the compiler replaces every block comment with a single space character, which makes it unwise to place a comment inside a name. [↑](#footnote-ref-12)
13. Usually, the problem is solved by the following macros (see Section 22.9):

    #define BOOL int

    #define TRUE 1

    #define FALSE 0 [↑](#footnote-ref-13)
14. Technically, it is given the value that happened to be stored on the memory address. [↑](#footnote-ref-14)
15. Depending of the compiler and its warning level settings. [↑](#footnote-ref-15)
16. Actually, f represent both float or double. The reason is that a float value is always converted to double when used as input to printf (or any other function). [↑](#footnote-ref-16)
17. As we inputs pointers to the variables rather than the variables themselves when calling scanf, no conversation from float to double occur. That is why f represents float and lf represents double. That lf (long float) represents double in scanf is an emergency solution since d represents an integer in decimal notation in both printf and scanf. [↑](#footnote-ref-17)
18. A white-space character is a space (‘ ‘), horizontal tabulator (‘\t’), vertical tabulator (‘\v’), new-line (‘\n’), form-feed (‘\f’), or carriage return (‘\r’). [↑](#footnote-ref-18)
19. Technically, each enumeration value holds the type unsigned int. [↑](#footnote-ref-19)
20. However, we have no knowledge about the actual address; it may be 10,000 or anything else. [↑](#footnote-ref-20)
21. In C++, the ampersand is also used to define references variables, which are not included in C. [↑](#footnote-ref-21)
22. More accurately, the number of bytes needed to hold a value of the type, [↑](#footnote-ref-22)
23. They were added to C in order to match two specific assembler instructions. Since then, they have stuck. More recent languages, such as C++, Java, and C#, do also support increment and decrement. [↑](#footnote-ref-23)
24. Some programmers prefer to denote equal to and not equal to as comparison operators. [↑](#footnote-ref-24)
25. In C#, the break statement is mandatory. [↑](#footnote-ref-25)
26. Technically, main is a regular function. The only difference, compared to other functions, is that the linker generates code that starts the program execution by calling main. Therefore, the program must include exactly one **main** function. [↑](#footnote-ref-26)
27. In C++, it is possible to access the global variable by prefixing colons, like ::number. [↑](#footnote-ref-27)
28. There is no rational explanation, just accept it. [↑](#footnote-ref-28)
29. Note the distinction of variable definition and declaration: a definition reserves memory for the variable while a declaration is simple a notification to the compiler. [↑](#footnote-ref-29)
30. In object-oriented languages the pointer is hidden and referred to as the this pointer. [↑](#footnote-ref-30)
31. This construction performs the same task as exception handling does in object-oriented languages. [↑](#footnote-ref-31)