# Numerical Mathematics I, 2017/2018, Lab session 3

Deadline for discussion: Monday 14/05/18 at 5 pm

Keywords: linear algebra, direct methods, LU, pivoting, conditioning, least-squares

#### Remarks

- Make a new folder called NM1\_LAB\_3 for this lab session, save all your functions in this folder.
- Whenever a new MATLAB function is introduced, try figuring out yourself what this function does by typing help <function> in the command window.
- Make sure that you have done the preparation before starting the lab session. The answers should be worked out either by pen and paper (readable) or with any text processing software (IATEX, Word, etc.).

## 1 Preparation

## 1.1 LU factorisation and partial pivoting

- 1. Study (Textbook, Section 1.4) and (Textbook, Section 5.1 5.5) (except 5.4.1). Recall from Linear Algebra (use your Linear Algebra 1 book if needed): overdetermined, underdetermined, inverse, nonsingular, range, nullspace, eigenvalue, norm, orthogonality.
- 2. Consider the linear system of equations  $\mathbf{A}\mathbf{x} = \mathbf{b}$ , where  $\mathbf{A}$  is given by

$$\mathbf{A} = \begin{pmatrix} \epsilon & 1 \\ 1 & 1 \end{pmatrix}.$$

For the following three situations

- (i)  $\epsilon \neq 1$  and  $\mathbf{b} = (1, 2)^T$
- (ii)  $\epsilon = 1 \text{ and } \mathbf{b} = (1, 2)^T$
- (iii)  $\epsilon = 1 \text{ and } \mathbf{b} = (1, 1)^T$ ,

answer the following questions

- (a) Is **A** nonsingular?
- (b) What is the nullspace of  $\mathbf{A}$ ? And the range of  $\mathbf{A}$ ? Is  $\mathbf{b} \in \text{range}(\mathbf{A})$ ?
- (c) How many solutions are there?
- (d) Give the general form of the solution(s) if there exist any.
- 3. A matrixnorm  $\|.\|_*$  on  $\mathbb{R}^{n\times n}$  is said to be sub-multiplicative if

$$\|\mathbf{A}\mathbf{B}\|_{*} \leq \|\mathbf{A}\|_{*} \|\mathbf{B}\|_{*},$$

for all matrices  $\mathbf{A}, \mathbf{B} \in \mathbb{R}^{n \times n}$ . Show that in this case the corresponding condition number

$$K_*(\mathbf{A}) = \|\mathbf{A}\|_* \|\mathbf{A}^{-1}\|_*$$

satisfies

$$K_*(\mathbf{AB}) \le K_*(\mathbf{A})K_*(\mathbf{B}),$$

for all nonsingular matrices  $\mathbf{A}, \mathbf{B} \in \mathbb{R}^{n \times n}$ .

- 4. For a linear system of equations  $\mathbf{Ma} = \mathbf{f}$  with square nonsingular matrix  $\mathbf{M}$ , give a general upper bound for the relative error in  $\mathbf{a}$  if we perturb both  $\mathbf{f}$  and  $\mathbf{M}$  (Textbook, Section 5.5). How does this upper bound simplify if we only perturb  $\mathbf{f}$ ?
- 5. Consider again solving  $\mathbf{Ma} = \mathbf{f}$  with square nonsingular matrix  $\mathbf{M}$ . Suppose we have an approximate solution  $\hat{\mathbf{a}}$ . Let the residual  $\mathbf{r}$  be given by

$$r = f - M\hat{a}$$
.

The scaled residual norm is defined as the 2-norm of the residual  $\mathbf{r}$  divided by the 2-norm of  $\mathbf{f}$ . Is the scaled residual norm always a good measure of the relative error in  $\hat{\mathbf{a}}$ ? How does this depend on  $K_2(\mathbf{M})$ ? Hint: use the result of the previous question.

## 1.2 Conditioning of the least squares problem

- 1. Study (Textbook, Section 3.6) and (Textbook, Section 5.7).
- 2. Let  $\mathbf{A} \in \mathbb{R}^{m \times n}$  and  $\mathbf{y} \in \mathbb{R}^m$  be given with  $m \geq n$ . Show that the minimization with respect to  $\mathbf{c}$  of

$$\|\mathbf{Ac} - \mathbf{y}\|_{2}^{2}$$

leads to the normal equations

$$\mathbf{A}^T \mathbf{A} \mathbf{c} = \mathbf{A}^T \mathbf{y}. \tag{1}$$

- 3. Subject to what condition on  $\mathbf{A} \in \mathbb{R}^{m \times n}$  does there always (for any  $\mathbf{y}$ ) exist a unique solution to the normal equations?
- 4. Under the condition derived in the point above, show that the unique solution  $\mathbf{c}$  to the normal equations (1) can be found by solving the upper triangular system

$$\mathbf{Rc} = \mathbf{Q}^T \mathbf{y},\tag{2}$$

where  $\mathbf{A} = \mathbf{Q}\mathbf{R}$  is the reduced ('economy size') QR factorisation of  $\mathbf{A}$  with orthonormal  $\mathbf{Q} \in \mathbb{R}^{m \times n}$  and  $\mathbf{R} \in \mathbb{R}^{n \times n}$ . Hint: orthonormality of  $\mathbf{Q}$  implies that  $\mathbf{Q}^T \mathbf{Q} = \mathbf{I}$ .

# 2 Lab experiments

## Introduction

In this lab session you will consider the following problem: for a given set of n data points

$$(x_i, y_i), \quad i = 1, \dots, n,$$

find the coefficients **c** of the polynomial f of degree  $\leq r$ ,

$$f(x) = \sum_{i=1}^{r+1} c_j x^{r-j+1},$$

such that the error

$$E := \sum_{i=1}^{n} (f(x_i) - y_i)^2$$

is minimized. Here we assume that  $r \leq n-1$ . If we define the *Vandermonde* matrix  $\mathbf{A} = (a_{ij}) \in \mathbb{R}^{n \times (r+1)}$  by

$$a_{ij} = x_i^{r-j+1},$$

then we can write the error E as

$$E = \|\mathbf{Ac} - \mathbf{y}\|_2^2,$$

showing that the problem consists of finding the least squares solution of the overdetermined linear system

$$\mathbf{Ac} = \mathbf{y}.\tag{3}$$

In the preparation you have shown that, under certain conditions, such a problem can be uniquely solved by solving the corresponding normal equations

$$\mathbf{A}^T \mathbf{A} \mathbf{c} = \mathbf{A}^T \mathbf{y}$$

for the coefficients  $\mathbf{c}$ . In this lab session you will investigate several solution techniques for this problem.

Write a function makeVandermondeMatrix which, given the abscissae vector  $\mathbf{x}$  and degree r, generates the corresponding Vandermonde matrix  $\mathbf{A}$ . The header of this function should be

### 2.1 LU factorisation and partial pivoting

Introduction

In this section we consider the case r = n - 1 = 1. We are looking for the best linear fit through the two data points corresponding to the vectors

$$\mathbf{x} = \begin{pmatrix} \epsilon \\ 1 \end{pmatrix}, \quad \mathbf{y} = \begin{pmatrix} 1 \\ 2 \end{pmatrix}.$$

Since r = n - 1, the corresponding linear system (3) is actually not overdetermined but square, so instead of solving the normal equations we will directly solve the square system

$$\mathbf{A}_{\epsilon}\mathbf{c}_{\epsilon} = \mathbf{y},\tag{4}$$

where  $\mathbf{A}_{\epsilon} \in \mathbb{R}^{2 \times 2}$  is the Vandermonde matrix corresponding to the given data.

```
ALG. 2.1: Naive LU factorisation

Input: The matrix \mathbf{A} \in \mathbb{R}^{n \times n}
Output: LU factors \mathbf{L} and \mathbf{U}
\mathbf{L} = \mathbf{I}, \mathbf{U} = \mathbf{A}
for k = 1: n do
\begin{vmatrix} L_{ik} \leftarrow U_{ik}/U_{kk} \\ \text{for } j = k: n \text{ do} \\ | U_{ij} \leftarrow U_{ij} - L_{ik}U_{kj} \\ \text{end} \end{vmatrix}
```

#### Experiment without pivoting

We want to solve this system of equations using an LU factorisation: write a function called  $\mathtt{luNaive}$  which solves an arbitrary square linear system by using the LU factorisation technique (a pseudocode is given in Algorithm 2.1). To test that it works (for arbitrary square linear systems), you can try it out as follows: generate a random  $10 \times 10$  matrix  $\mathbf{A}$  and solution vector  $\mathbf{x}$ , compute  $\mathbf{b} = \mathbf{A}\mathbf{x}$ , and then solve for  $\mathbf{x}$  using your  $\mathtt{luNaive}$  function. Solving the upper and lower triangular subproblems may be done by using the backslash command \ (which recognizes the triangular structure of  $\mathbf{L}$  and  $\mathbf{U}$ ). The header of this function should be

```
% INPUT
2
  % A
               square matrix
3
  % b
               right hand side
  % OUTPUT
  % x
               solution such that A*x=b
  %
6
               lower triangular matrix such that A = L*U
  % U
               upper triangular matrix such that A = L*U
  function [x, L, U] = luNaive(A, b)
```

Hence the triangular factors **L** and **U** should be returned as well. Now consider solving (4) for the following values of  $\epsilon$ ,

$$\epsilon = 10^{-i}, \quad i = 1, \dots 16.$$

The resulting solutions will be denoted by  $\hat{\mathbf{c}}_{\epsilon}$ . For each  $\epsilon$  compute the following quantities:

• The relative error

$$\frac{\|\hat{\mathbf{c}}_{\epsilon} - \mathbf{c}_{\epsilon}\|_2}{\|\mathbf{c}_{\epsilon}\|_2}$$

in the obtained solution  $\hat{\mathbf{c}}_{\epsilon}$  (compute the exact solution  $\mathbf{c}_{\epsilon}$  by hand).

- The  $K_2$  condition numbers of  $\mathbf{A}_{\epsilon}$ ,  $\mathbf{L}_{\epsilon}$  and  $\mathbf{U}_{\epsilon}$
- The factorisation error

$$\|\mathbf{A}_{\epsilon} - \mathbf{L}_{\epsilon}\mathbf{U}_{\epsilon}\|_{F}$$

where  $\|\mathbf{M}\|_F$  denotes the Frobenius norm which is defined as

$$\|\mathbf{M}\|_F := \sqrt{\sum_{i,j=1}^{m,n} m_{ij}^2},$$

which in MATLAB can be computed by norm(M, 'fro').

Summarise your results in a single double-logarithmic plot (use the loglog function) where the five quantities are plotted as a function of  $\epsilon$ .

```
ALG. 2.2: LU factorisation with partial pivoting

Input: The matrix \mathbf{A} \in \mathbb{R}^{n \times n}
Output: LU factors \mathbf{L}, \mathbf{U} and the permutation matrix \mathbf{P}
\mathbf{L} = \mathbf{0}, \mathbf{U} = \mathbf{A}, \mathbf{P} = \mathbf{I}
for k = 1: n do

| Find \bar{r} such that |U_{\bar{r}k}| = \max_{r=k,...n} |U_{rk}|
Swap the \bar{r}-th row with the r-th row in \mathbf{L}, \mathbf{U} and \mathbf{P}
for i = k + 1: n do

| L_{ik} \leftarrow U_{ik}/U_{kk}
for j = k: n do

| U_{ij} \leftarrow U_{ij} - L_{ik}U_{kj}
end
end
end
L_{kk} = 1
```

Experiment using partial pivoting

Repeat the previous experiment, but now use the LU factorisation with partial pivoting such that

$$PA = LU$$
,

where **P** stands for a permutation matrix. For this you should write a function called  $\mathtt{luPivot}$  which solves an *arbitrary* square linear system using the LU factorisation technique with partial pivoting (a pseudocode is given in Algorithm 2.2). This function should also return the permutation matrix **P**. Just like you did for the previous function  $\mathtt{luNaive}$ , first test the new function on a random  $10 \times 10$  linear system. Next solve system (4) for the same values of  $\epsilon$  as before and summarise your results in a similar double-logarithmic plot.

#### Optional

Extend the functionality of your  ${\tt luNaive}$  implementation such that it can efficiently solve linear systems where the coefficient matrix is banded. The returned factors  ${\tt L}$  and  ${\tt U}$  should be  ${\tt sparse}$  matrices. Test your implementation on

```
A = \text{spdiags(kron(ones(N, 1), [1 -2 1] * N^2), [-1 0 1], N, N); b = ones(N, 1);}
```

for some N.

Determine experimentally the computational cost (that is, measure the time using tic and toc) of the resulting algorithm while varying N, and compare it to the previous implementation which did not take the bandwidth into account.

## 2.2 Conditioning of the least squares problem

Introduction

Consider the set of 21 data points given by

$$x_i = (i-1)/20, \quad y_i = x_i^8, \quad i = 1, \dots, 21,$$

which correspond to the eighth degree monomial  $f(x) = x^8$ . If we set the maximum polynomial degree equal to r = 8, the least squares solution  $\mathbf{c}$  of the overdetermined system (3) is easily seen to be given by  $\mathbf{c} = \mathbf{e}_1 := \begin{bmatrix} 1 & 0 & \dots & 0 \end{bmatrix}^T$ .

As you have shown in the preparation, solving the resulting overdetermined system can be done in at least two ways: either we solve (1) where we consider  $\mathbf{A}^T \mathbf{A}$  as the system matrix, or we use a QR factorisation of  $\mathbf{A}$  and solve (2) with  $\mathbf{R}$  as the system matrix. Note that in MATLAB one should compute the QR factorisation as [Q, R] = qr(A, 0). Inverting the system matrices can be done using MATLAB's backslash.

#### Experiment

To illustrate the difference between the two solution methods, we perturb the data points (both x and y) by adding random perturbations from the interval  $(-\epsilon, \epsilon)$ , thereby simulating measurement errors of  $\mathcal{O}(\epsilon)$ , and quantify the effect these perturbations have on the resulting solution. We consider the following values of  $\epsilon$ ,

$$\epsilon = 10^{-i}, \quad i = 1, \dots 16,$$

and denote the resulting solution by  $\hat{\mathbf{c}}_{\epsilon}$ . For each value of  $\epsilon$ , and for each of the solution methods (hence in total you should run 32 experiments), compute the following quantities:

- The relative error of the solution. We now consider the exact solution  $\mathbf{c} = \mathbf{e}_1$  which is, contrary to the previous experiment, independent of  $\epsilon$  since we now consider  $\epsilon$  as a measurement error.
- The upper bound (which you found in the preparation) for the relative error of the solution of the resulting perturbed linear system (the matrix **A** as well as the right-hand side **y** depend on  $\epsilon$ ). Note that the linear system solved is different for each of the solution methods, and so is the right-hand side.

Summarise your results in a single double-logarithmic plot, where each of the quantities is plotted as a function of  $\epsilon$ .

## 3 Discussion

### 3.1 LU factorisation and partial pivoting

- 1. What can go wrong when computing the "normal" LU factorisation of a matrix and how does partial pivoting circumvent this problem?
- 2. What is the computational cost of making an LU factorisation (for a full matrix)? What if the matrix is banded with bandwidth b?
- 3. (a) When making an LU factorisation you basically split the problem of finding  $\mathbf{x}$  such that  $\mathbf{A}\mathbf{x} = \mathbf{b}$  into two subproblems. Which subproblems?

- (b) In your numerical experiments, how do the condition numbers  $(K_2)$  of these two subproblems compare to that of the original problem? Make here a distinction between the experiments with and without partial pivoting.
- (c) Would it ever be possible that the two subproblems are well-conditioned, and the original problem is not? Hint: the matrix 2-norm  $\|.\|_2$  is sub-multiplicative.
- (d) Do your experiments confirm this?
- 4. When not using pivoting we must solve the following upper triangular subproblem

$$\begin{bmatrix} \epsilon & 1 \\ 0 & 1 - \frac{1}{\epsilon} \end{bmatrix} \mathbf{x} = \begin{bmatrix} 1 \\ 2 - \frac{1}{\epsilon} \end{bmatrix}. \tag{5}$$

(a) Verify that the solution to this subproblem is given by

$$\mathbf{x} = \begin{bmatrix} \frac{1}{1-\epsilon} \\ \frac{2\epsilon - 1}{\epsilon - 1} \end{bmatrix},$$

which is well-defined as  $\epsilon \to 0$ .

- (b) Show that the relative error in solving (5), due to the finite representation of floating point numbers, is given by  $\frac{\varepsilon_M}{\epsilon}$ .
- (c) Are your experiments in agreement with this?

## 3.2 Conditioning of the least squares problem

- 1. Consider the  $K_2$  condition numbers of the matrices  $\mathbf{A}^T \mathbf{A}$  and  $\mathbf{R}$ . How are they related? Are your experiments in agreement with this?
- 2. Are the 32 experiments that you did in Section 2.2 in agreement with the theoretical upper bounds for the relative error that you provided in Question 4 of Preparation section 1.1?