

CS2510

MODERN PROGRAMMING LANGUAGES

Logic Programming 4

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Prolog Programming with Lists

- Because of their recursive definition, lists are naturally manipulated with recursive programs.
- Let's revisit the earlier loop program:
 - Enable it to build a list of values typed in at keyboard:

```
loop(L):- % a loop to create a list L
  L = [X|Xs], % List L has form [X|Xs]
  read(X), % X is read from the standard Input (keyboard)
  X \= end, % if input different from "end" then
  loop(Xs). % build the tail recursively
loop([]). % if input is "end", then List is empty.
```

```
?- loop(MyList).
|: a.
|: b.
|: end.
MyList = [a,b]
```

An argument is required to return the value of the list built.

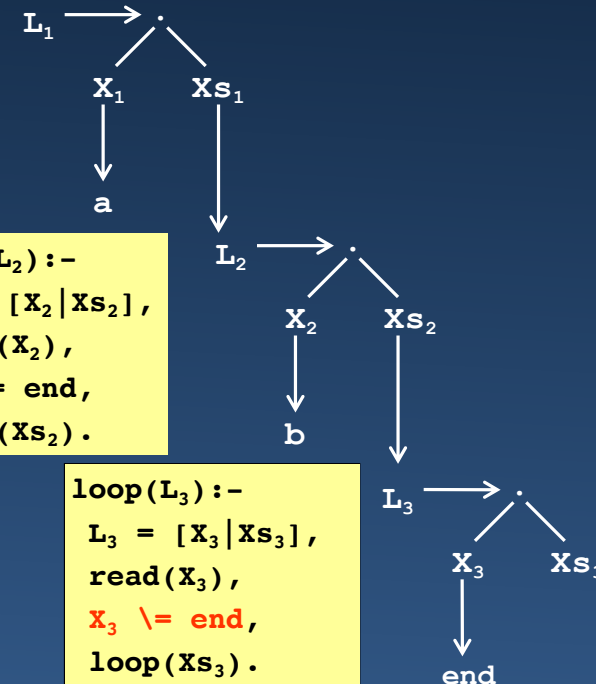
Programming with Lists

- Its execution:

```
loop(L) :-  
    L = [X|Xs],  
    read(X),  
    X \= end,  
    loop(Xs).  
loop([]).
```

```
?- loop(MyList).  
|: a.  
|: b.  
|: end.
```

```
loop(L1) :-  
    L1 = [X1|Xs1],  
    read(X1),  
    X1 \= end,  
    loop(Xs1).
```



Programming with Lists

- Its execution:

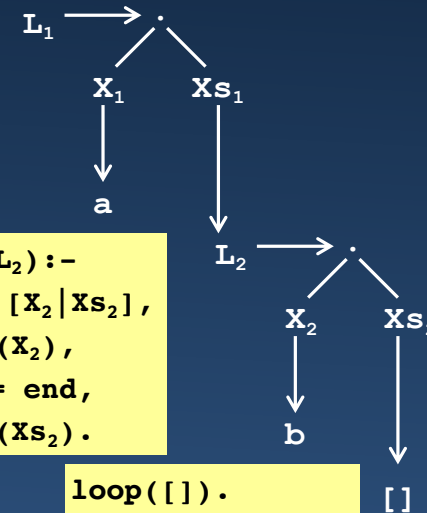
```
loop(L):-  
  L = [X|Xs],  
  read(X),  
  X \= end,  
  loop(Xs).  
loop([]).
```

```
?- loop(MyList).  
|: a.  
|: b.  
|: end.  
MyList = [a,b]
```

```
loop(L1):-  
  L1 = [X1|Xs1],  
  read(X1),  
  X1 \= end,  
  loop(Xs1).
```

```
loop(L2):-  
  L2 = [X2|Xs2],  
  read(X2),  
  X2 \= end,  
  loop(Xs2).
```

```
loop([]).
```



Programming with Lists: member ()

```
?- member(c, [a, b, c]).  
true  
?-
```

- To check if an element appears in a list:
 - Check if element is the head;
 - Otherwise, check if element appears in tail.
- This formulation is guaranteed to stop because lists are finite and the process breaks list apart.

- In *Prolog* (first formulation):

```
member(X,List):-      % X is a member of List  
    List = [Y|Ys],    % if List is of form [Y|Ys]  
    X = Y.             % and X is equal to (unifies with) Y  
member(X,List):-      % otherwise (if not as above)  
    List = [Y|Ys],    % break the list apart  
    member(X,Ys).      % and check if X appears in Ys
```

Programming with Lists: member ()

- Second formulation:

```
member(X,[Y|Ys]):-      % X is a member of [Y|Ys]
    X = Y.               % if X is equal to Y
member(X,[Y|Ys]):-      % otherwise (if not as above)
    member(X,Ys).        % check if X appears in Ys
```

- Third (final) formulation:

```
member(X,[X|_]).        % X is member if it is head of list
member(X,[_|Ys]):-      % otherwise (if not as above)
    member(X,Ys).        % check if X appears in Ys
```



Programming with Lists: `member()`

- Let's run the previous program:

```
member(X, [X|_]).  
member(X, [_|Ys]) :-  
    member(X, Ys).
```

```
?- member(3, [1,2,3]).  
true  
?-
```

The query `member(3, [1,2,3])` does not match the 1st clause!!

Prolog matches `member(3, [1,2,3])` with head of 2nd clause:

`member(X1, [_|Ys1]) :- member(X1, Ys1)`

We thus get the values for the variables: $\{X_1/3, Ys_1/[2,3]\}$

Prolog now tries to prove `member(X1, Ys1)` that is

`member(3, [2,3])`

The goal `member(3, [2,3])` does not match the 1st clause!!

Prolog matches `member(3, [2,3])` with head of 2nd clause:

`member(X2, [_|Ys2]) :- member(X2, Ys2)`

We thus get the values for the variables: $\{X_2/3, Ys_2/[3]\}$

Prolog now tries to prove `member(X2, Ys2)` that is

`member(3, [3])`

The query `member(3, [3])` matches the 1st clause!!

Programming with Lists: `member()`

- `member` code can be run in two ways:
 - To check if an element occurs in a list;
 - To obtain, the elements of a list, one at a time
- It depends on the query – the code is exactly the same!
- To obtain the elements of a list, we query:
`member(Elem, [1,2,3])`
 - We then get in `Elem`, the different elements:

```
?- member(Elem,[1,2,3]).  
Elem = 1 ? ;  
Elem = 2 ? ;  
Elem = 3 ? ;  
false  
?-
```



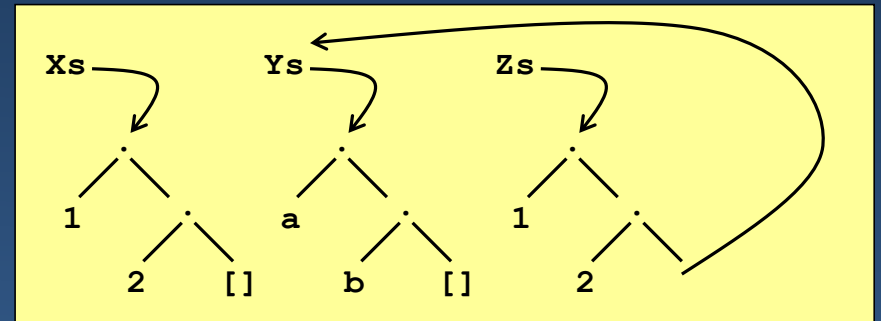
Can you trace this query?

Programming with Lists: `append()`

```
?- append([1,2], [a, b], Res).  
Res = [1,2,a,b] ?  
true  
?-
```

- To concatenate two lists **Xs** and **Ys**:
 - Create a 3rd list **Zs** by copying all elements of **Xs**;
 - When we get to the end of **Xs**, put **Ys** as the tail of **Zs**.

- Diagrammatically:



Programming with Lists: `append()`

- We need a predicate with 3 arguments:
 - 1st argument is the list `Xs`
 - 2nd argument is the list `Ys`
 - 3rd argument is the concatenation `Xs.Ys` (in this order)
- The first list `Xs` will control the loop
- As the first list `Xs` is traversed, its elements are copied onto the 3rd argument (the concatenation).
- When we get an empty list, we stop and make `Ys` the value of the 3rd argument

Programming with Lists: `append()`

- First formulation:

```
append([ ], Ys, Zs) :-           % if Xs is empty,  
    Zs = Ys.                     % then Zs is Ys  
append([X|Xs], Ys, [Z|Zs]) :-    % otherwise  
    X = Z,                       % copy X onto 3rd argument  
    append(Xs, Ys, Zs).          % carry on recursively
```

- A more compact format:

```
append([ ], Ys, Ys).             % if Xs empty, then Zs is Ys  
append([X|Xs], Ys, [X|Zs]) :-    % else copy X onto 3rd arg.  
    append(Xs, Ys, Zs).          % carry on recursively
```



SWI Prolog

Programming with Lists: append ()

- Execution:

```
append([], Ys, Ys).
append([X|Xs], Ys, [X|Zs]) :-
    append(Xs, Ys, Zs).
```

```
?- append([1,2], [a,b], Res).
Res = [1,2,a,b] ?
true
?-
```

```
append([X1|Xs1], Ys1, [X1|Zs1]) :-
    append(Xs1, Ys1, Zs1).
{X1/1, Xs1/[2], Ys1/[a,b]}
```

```
append([X2|Xs2], Ys2, [X2|Zs2]) :-
    append(Xs2, Ys2, Zs2).
{X2/2, Xs2/[], Ys2/[a,b]}
```

```
append([], Ys3, Ys3).
{Ys3/[a,b]}
```

aliasing: Res/[X₁|Zs₁]

aliasing: Zs₁/[X₂|Zs₂]

aliasing: Zs₂/Ys₃

Programming with Lists: append()

- Execution:

```
append([], Ys, Ys).
```

```
append([X|Xs], Ys, [X|Zs]) :-  
    append(Xs, Ys, Zs).
```

```
?- append([1,2], [a,b], Res).
```

```
Res = [1,2,a,b] ?
```

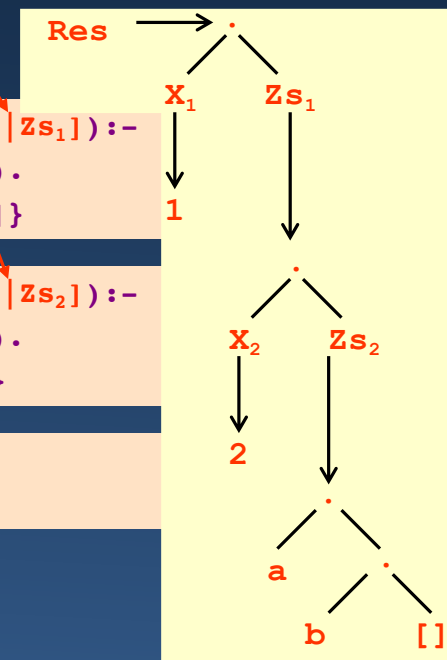
```
true
```

```
?-
```

```
append([X1|Xs1], Ys1, [X1|Zs1]) :-  
    append(Xs1, Ys1, Zs1).  
{X1/1, Xs1/[2], Ys1/[a,b]}
```

```
append([X2|Xs2], Ys2, [X2|Zs2]) :-  
    append(Xs2, Ys2, Zs2).  
{X2/2, Xs2/[], Ys2/[a,b]}
```

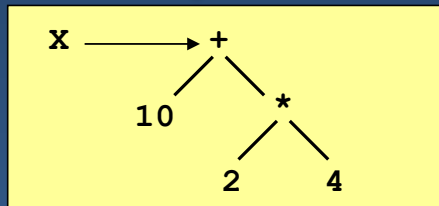
```
append([], Ys3, Ys3).  
{Ys3/[a,b]}
```



Arithmetic

- Arithmetic expressions are also trees:

```
?- x = 10 + (2 * 4).  
x = 10 + (2 * 4) ?  
true  
?-
```



- Prolog* does not treat arithmetic expressions differently, unless asked to do so.
- To evaluate arithmetic expressions, use the built-in “*is*” – as follows:

```
?- x is 10 + (2 * 4).  
x = 18 ?  
true  
?-
```

Arithmetic

- The syntax of the built-in “**is**” is:

`Variable is ArithmeticExpression`

- Examples:

`Result is (10 * 7)
NewRes is OldRes + 1`

- Variables that appear in the expression
 - Must be instantiated when expression is evaluated;
 - Otherwise execution error!

- Can delay the evaluation until all variables have value:

```
?- Exp = A + B, A = 3, B = 4, Res is Exp.  
A = 3,  
B = 4,  
Exp = 3 + 4  
Res = 7
```



SWI Prolog

Arithmetic: `length()`

```
?- length([a,b,c,d],Length).  
Length = 4 ?  
?- length([1,[2,3],4],L).  
L = 3 ?  
?-
```

- Predicate **length** of arity 2:
 - 1st argument is the list;
 - 2nd argument is the size of the list.
- Find the size (number of elements) of a list.
- Simple recursive formulation:
 - The number of elements of the empty list is zero;
 - The number of elements of a list **[X|Xs]** is one plus the number of elements of **Xs**.

Arithmetic: length ()

- First formulation:

```
length(L,S):-                % S is the length of list L
    L = [],                  % if L is an empty list
    S = 0.                   % its length S is zero
length(L,S):-                % otherwise
    L = [X|Xs],              % if L is a list [X|Xs]
    S is 1 + length(Xs,S).   % its length is 1 plus length of tail
```

CARE!

Prolog calls are **true** or **false**.

This expression does not make sense!

Arithmetic: length ()

- Second formulation:

```
length(L,S):-                % S is the length of list L
    L = [],                  % if L is an empty list
    S = 0.                   % its length S is zero
length(L,S):-                % otherwise
    L = [X|Xs],              % if L is a list [X|Xs]
    length(Xs,1 + SXs).      % get the length SXs of tail Xs
```

- Third formulation:

CARE! This is just building a structure, not computing an expression! What about S ?

```
length(L,S):-                % S is the length of list L
    L = [],                  % if L is an empty list
    S = 0.                   % its length S is zero
length(L,S):-                % otherwise
    L = [X|Xs],              % if L is a list [X|Xs]
    S is 1 + SXs,            % length S is 1 plus length of tail
    length(Xs,SXs).          % get the length SXs of tail Xs
```

CARE! When this expression is evaluated, the value of SXs won't yet exist!

Arithmetic: length()

- Fourth formulation:

```
length(L,S):-           % S is the length of list L
    L = [],             % if L is an empty list
    S = 0.              % its length S is zero
length(L,S):-           % otherwise
    L = [X|Xs],         % if L is a list [X|Xs]
    length(Xs,SXs),     % get the length SXs of its tail Xs
    S is 1 + SXs.       % add 1 to the length of its tail
```

- Simplified version:

```
length([],0).           % the empty list has length 0
length([_|Xs],S):-      % a non-empty list [_|Xs] has size S
    length(Xs,SXs),     % get the length SXs of its tail Xs
    S is 1 + SXs.       % add 1 to the length of its tail
```

Anonymous variable, used as “place holder”.

In this context, we don't care what the value of the element is – we just want to count them!



Arithmetic - Summary

- Arithmetic is via “**is**” built-in.
- Some operators:
 - **+**, **−**, *****, **/** (add, subtract, multiply, divide)
 - **//** (integer division)
 - **mod** (modulo)
- All variables on expression must have values, otherwise execution error.
- Because of this restriction, need to bear in mind the order in which *Prolog* proves/executes the body of a clause (or a query).

Failure-Driven Loops

- *Prolog* offers a built-in predicate **fail** which always fails:

```
?- fail.  
false
```

- We can use this predicate to define a failure-driven loop, an alternative to recursion:


```
?- loopFail.  
a  
b  
c  
true  
?-
```

```
p(a).  
p(b).  
p(c).  
  
loopFail:- % loopFail succeeds if  
    p(X), % we can prove p(X) and  
    write(X), % write the value of X and  
    nl, % skip a line.  
    fail. % fail and BACKTRACK!  
loopFail. % if no (more) answers, stop
```

Controlling Backtracking via Cuts (!)

- *Prolog*'s backtracking mechanism may, in some cases, lead to inefficiencies.
- We can control backtracking via the built-in “!”, called “*cut*”.
- The “!” is used as an ordinary predicate, in the body of the clause or query – it always succeeds!
- The “!” has a special property:
 - It causes the execution to commit to the solutions (proofs) of the goals to its left.
- Example:

```
p(A,B,C,D):- q(A), r(A,B), !, s(B,C), t(A,D).
```



Controlling Backtracking Using !

- Example:

```
p(a).
p(b).
p(c).
q(b,2).
q(c,3).
s(2).
s(3).
r(Y):- p(X),q(X,Y),!,s(Y).
```

```
?- r(Ans).
Ans = 2 ? ; Force backtrack...
false
?-
```

```
r(Y1):- p(X1),q(X1,Y1),!,s(Y1).
{X1/a} Fail & Backtrack...
```

```
r(Y1):- p(X1),q(X1,Y1),!,s(Y1).
{X1/b,Y1/2}
```

```
r(Y1):- p(X1),q(X1,Y1),!,s(Y1).
{X1/b,Y1/2}
```

```
r(Y1):- p(X1),q(X1,Y1),!,s(Y1).
{X1/b,Y1/2}
```

```
r(Y1):- p(X1),q(X1,Y1),!,s(Y1).
{X1/b,Y1/2} No other proof for s(Y1)
```

```
r(Y1):- p(X1),q(X1,Y1),!,s(Y1).
{X1/b,Y1/2} Cannot backtrack over "!"
```

Controlling Backtracking Using !

- Why “cut”?
 - It is a reference to the search space of a query...

