

2.2.

$N_f = ?$  If number of traffic flows

All-to-all traffic matrix among servers

All the servers are connected.

So, we need to consider all flows source-dest.

between the couple of servers.

In this case, we can use to calculate  $N_f$  the binomial coefficient  $\binom{a}{b}$ , where

a means all the servers, which are  $N = m^3/h$  and  $b = 2$ , because are two servers to be connected, i.e. all the possible combination of 2 among N servers.

$$N_f = \binom{N}{2} = \frac{N(N-1)}{2} = \frac{m^3}{4} \cdot \left(\frac{m^3}{h} - 1\right) \cdot \frac{1}{2} = \frac{m^3}{8} \cdot \frac{m^3 - h}{h} =$$
$$= \boxed{\frac{m^3(m^3 - h)}{32}}$$

2.3

Exact value of  $\bar{h}$  (FOR FAT-TREE)

$$\bar{h} = \frac{1}{N_f} \sum_{i=1}^{N_f} h_i =$$

Calculated sum

$$2 \underbrace{\left( \begin{array}{c} ? \\ . \\ . \end{array} \right)} + 4 \underbrace{\left( \begin{array}{c} ? \\ . \\ . \\ . \end{array} \right)} + 6 \underbrace{\left( \begin{array}{c} ? \\ . \\ . \\ . \\ . \end{array} \right)}$$

Depth search

stress test

P>D DIVING

$\binom{m}{2}^l$  with switches

$$\left(\frac{M}{2}\right)^l$$

$$2 \cdot \left[ \frac{m^3}{8} \left( \frac{m}{2} - 1 \right) \right] \Rightarrow \frac{m^3}{4} \frac{1}{2} \left( \frac{m}{2} - 1 \right)$$

$$2 \cdot \left[ \left( \frac{m}{2} - 1 \right) \left( \frac{m^3}{4} - \frac{1}{2} \right) \right] + 4 \left[ \left( \frac{m}{2} - 1 \right) \left( \frac{m^3}{8} \cdot \frac{m}{2} \right) \right] + 6 \left[ \left( \frac{m}{2} - 1 \right) \left( \frac{m^3}{8} \cdot \frac{m^2}{4} \right) \right]$$

For servers connected by some switch

2 flows for each path

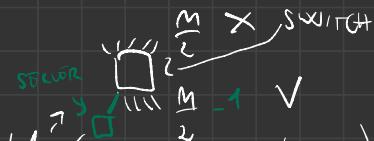
$\frac{m}{2} - 1$  ports of switch for new path ( $m/2$  ports connected to the path interested, but minus 1, bc in the case of the server)

How many different paths from A to B with no duplicates, i.e. only one path from A to B,  
so not to count the path from B to A, which is the same

$$2 \cdot \left( \frac{m}{2} - 1 \right) \left( \frac{m^3}{4} \cdot \frac{1}{2} \right)$$

For servers connected by aggregation layer switch,  
so some port

4 flows for each path



$\frac{m}{2} - 1$  ports of switch for new path (same as above)

$\frac{m^3}{4} \cdot \frac{1}{2} \cdot \frac{m}{2}$  where  $\frac{m^3}{4} \cdot \frac{1}{2}$  is the same motivation

as above, but  $m/2$  means how many aggregation layer switch are in a path

$$4 \cdot \left( \frac{m}{2} - 1 \right) \left( \frac{m^3}{8} \cdot \frac{m}{2} \right) 4 \cdot 1 \cdot 2 \cdot \frac{64}{8} = 16$$

For server connected by core layer switch

6 flows for each path

$m-1$  ports of switch for new path (for core layer switch,  $m$  ports total, minus 1, which is connected to the initial server)

$$\frac{m^3}{4} \cdot \frac{1}{2} \cdot \frac{m^2}{4}$$

$\frac{m^3}{4} \cdot \frac{1}{2}$  same as previous ones, and  $\frac{m^2}{4}$  measures the number of core layer switch

$$6 \cdot (m-1) \left( \frac{m^3}{8} \cdot \frac{m^2}{4} \right)$$

Final L

$$L = \frac{1}{N_f} \sum_{i=1}^{N_f} L_i = \frac{1}{N_f} \left\{ 2 \left[ \left( \frac{m}{2} - 1 \right) \cdot \frac{m^3}{8} \right] + 4 \left[ \left( \frac{m}{2} - 1 \right) \left( \frac{m^3}{8} \cdot \frac{m^2}{4} \right) \right] + 6 \left[ \left( m - 1 \right) \left( \frac{m^3}{8} \cdot \frac{m^2}{4} \right) \right] \right\}$$

$$\begin{aligned} & \left( \frac{m}{2} - 1 \right) \cdot \frac{m^3}{8} + \left( \frac{m}{2} - 1 \right) \frac{m^4}{16} + (m-1) \left( \frac{m^5}{32} \right) = \frac{m^4}{16} - \frac{m^3}{8} + \frac{m^8}{32} - \frac{m^6}{16} + \frac{m^6}{32} \\ & = \frac{m^6 - 4m^3}{32} = \frac{m^3(m^3 - 4)}{32} \end{aligned}$$

which is equal to  $N_f = \frac{m^3(m^3 - 4)}{32}$