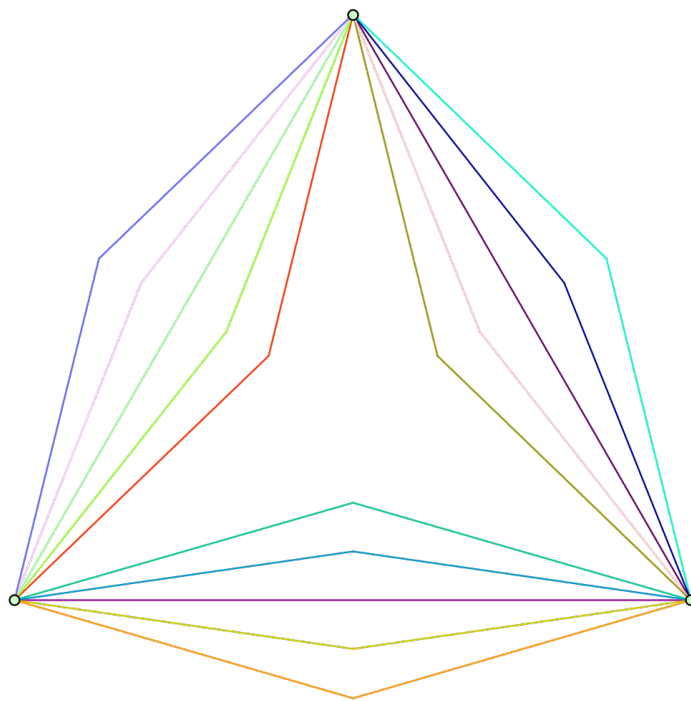


1. Give an example to show that if  $G$  is allowed to have multiple edges, then  $\chi'(G)$  may exceed  $\Delta(G) + 1$ .

**Solution:**

Figure 1: Multigraph  $G$  with  $\Delta(G) + 1 = 11$  but  $\chi'(G) = 15$



2. Without using Proposition 5.3.1, show that  $\chi'(G) = k$  for every  $k$ -regular bipartite graph  $G$ .

*Proof.* We will proceed to show that  $\chi'(G) = k$  for every  $k$ -regular bipartite graph  $G$  via induction on  $k$ . Consider the base case, where  $G$  is a 1-regular bipartite graph. Since every vertex is incident to exactly one edge, we can simply color all the edges the same color, and produce a 1 edge coloring, hence by Vizing's Theorem  $\chi'(G) = 1$ .

Suppose that  $G$  is  $n + 1$ -regular bipartite. Since  $G$  is a regular bipartite graph, it contains a one factor,  $M$ . Note that  $G - M$  is  $n$ -regular bipartite since each  $e \in M$  contributes one degree to exactly one vertex in  $A$  and one vertex in  $B$  and since  $M$  spans the vertices of  $G$ . Now by the induction hypothesis there exists an  $n$  coloring of  $G - M$ , call it  $C$ . Apply  $C$  to  $G$  and color the edges of  $M$  the  $n + 1^{th}$  color, call this coloring  $C'$ . Note that  $C'$  is a valid coloring since  $M$  is a one factor. Hence by Vizing's Theorem  $\chi'(G) = n + 1$ .  $\square$

3. Give an explicit edge-coloring to prove that the  $n$ -dimensional cube,  $Q^n$ , is Class 1.

*Proof.* Recall that the  $n$ -dimensional cube,  $Q^n$  is  $n$ -regular and therefore showing  $Q^n$  is Class 1 is equivalent to showing  $\chi'(Q^n) = n$ . Now we will proceed by induction on  $n$ . The base case is trivial since  $Q^1 = K^2$ , there is only one edge.... color it. Now consider  $Q^{n+1}$  and recall the construction of  $Q^{n+1}$  via two copies of  $Q^n$  and an independent set of edges call them  $M$  (We have discussed this before and I've proven this set is independent in HW1). So  $Q^{n+1} - M$  is two components, call them  $Q_1$  and  $Q_2$ , they are isomorphic  $Q^n$  dimensional cubes and therefore  $\chi'(Q_1) = \chi'(Q_2) = \chi'(Q^n) = n$ . Apply the same edge-coloring  $C$  on  $Q_1$  and  $Q_2$ , and color the edges of  $M$  the  $(n+1)^{th}$  color. Thus we have produced an edge-coloring of  $Q^{n+1}$  with  $n+1$  so by Vizing's Theorem  $\chi'(Q^{n+1}) = n+1$ .  $\square$

4. Prove that if  $G$  is a regular graph with a cut vertex, then  $\chi'(G) > \Delta(G)$ .

*Proof.* Suppose  $G$  is a  $k$ -regular connected graph, with  $\chi'(G) = \Delta(G) = k$ . We will proceed to show that  $G$  has no cut vertex, by proving  $G$  is 2-connected. Since  $G$  is  $k$ -regular and  $k$ -edge colorable for every  $u, v \in G$  and for all  $\alpha, \beta \in [k]$  there exists an  $\alpha, \beta$  alternating  $uv$ -path. Now let  $u, v \in G$  and consider the  $\alpha, \beta$  alternating  $uv$ -path. Regardless of the edge color entering  $v$ , there exist an edge of the other color leaving  $v$  and entering another vertex, not on the path. Since the  $\alpha, \beta$  color classes span the graph  $G$ , the path can be extended to eventually enter  $u$  via a  $\beta$  edge, forming a cycle. Hence  $G$  is 2-connected.

□

The rest is for my own good.

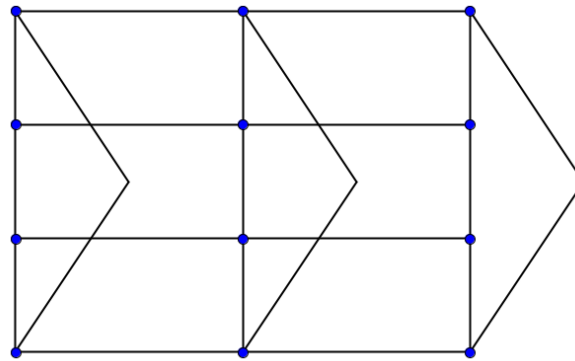
Since  $G$  is  $k$ -regular and  $k$ -edge colorable we know that for every  $\alpha \in [k]$  the set of all  $\alpha$ -edges is a one-factor by regularity, and even further for every pair of  $\alpha, \beta \in [k]$  the set of all  $\alpha$ -edges is disjoint from the set of all  $\beta$ -edges by being  $k$ -edge colorable. Thus we can conclude that the subgraph formed by all  $\alpha$ -edges and  $\beta$ -edges is connected and spans the vertices of  $G$  and clearly any path between two vertices will be alternating  $\alpha, \beta$ .

5. The **cartesian product** of two graph  $G$  and  $H$ , denoted  $G \square H$ , is the graph with vertex set  $V(G \square H) = V(G) \times V(H)$ . A pair of vertices  $(x_1, y_1), (x_2, y_2)$  are adjacent in  $G \square H$  if and only if  $x_1 = x_2$  and  $y_1 y_2 \in E(H)$  or  $x_1 x_2 \in E(G)$  and  $y_1 = y_2$ .

- (a) Draw  $P_2 \square C^4$ .

**Solution:**

Figure 2:  $P_2 \square C^4$ .



- (b) Prove that  $\Delta(G \square H) = \Delta(G) + \Delta(H)$

*Proof.* Let  $G$  and  $H$  be simple graphs, and let  $(u, v) \in V(G \square H)$ . By definition, each adjacency of  $u$  in  $G$ , and  $v$  in  $H$ , induces an adjacency of  $(u, v)$  in  $G \square H$ , so it follows that  $d((u, v)) = d(u) + d(v)$ . Clearly it follows that,  $\Delta(G \square H) = \Delta(G) + \Delta(H)$ .  $\square$

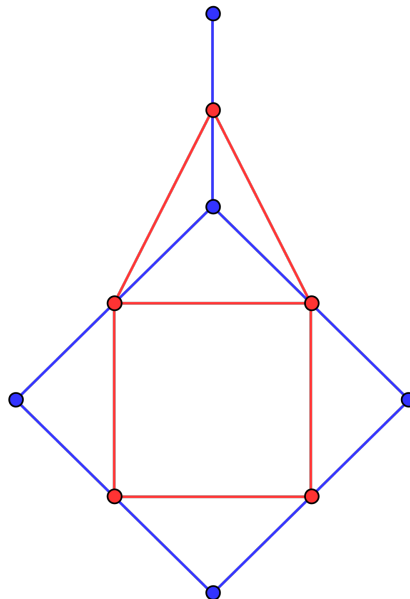
- (c) Prove that if  $\chi'(H) = \Delta(H)$ , then  $\chi'(G \square H) = \Delta(G \square H)$ .

*Proof.* Suppose  $\chi'(H) = \Delta(H)$ . Now consider  $(u, v) \in G \square H$  such that  $u \in G$  and  $v \in H$  have maximum degree. Color every copy of  $G$  in  $G \square H$  with  $\Delta(G) + 1$  colors. Now note that for every  $u \in V(G)$  there exists at least one unused color across all copies of  $u$  in  $G \square H$ , call this color  $c_u$ . For every  $u \in V(G)$  color the graph  $H$  incident to every copy of  $u$  with  $\Delta(H) - 1$  new colors and the corresponding unused color  $c_u$ . Thus we have constructed a  $\Delta(G) + 1 + \Delta(H) - 1 = \Delta(G) + \Delta(H) = \Delta(G \square H)$  coloring of  $G \square H$ .  $\square$

6. (a) Let  $G_1$  be a 5-cycle with one chord. Show that there exists a graph  $H$  such that  $L(H) = G_1$ .

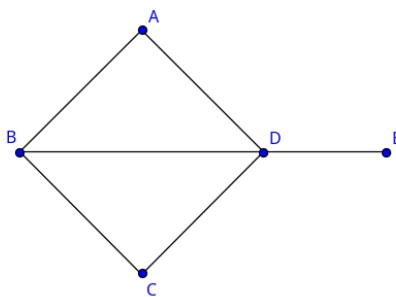
**Solution:**

Figure 3: Graph  $G_1$  in Red and  $H$  in Blue.

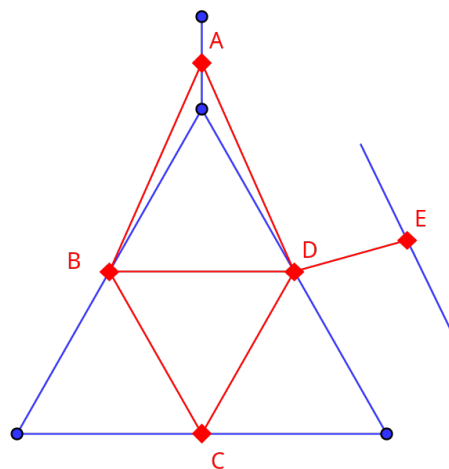


- (b) Prove that for the graph  $G_2$  (drawn below) that there does not exist any graph  $H$  such that  $L(H) = G_2$ .

Figure 4: Graph  $G_2$



*Proof.* Suppose for the sake of contradiction that there exists a graph  $H$  such that  $L(H) = G_2$ . Clearly the part of  $G_2$  which is a 4-cycle with a chord can be a line graph via a similar construction as the previous problem. However to incorporate vertex  $E$  into the line graph, would also add an edge  $CE$  or  $AE$ .

Figure 5: Graph  $G_2$  in Red and  $H$  in Blue

□