

Graphs' decompositions and resolutions of combinatorial problems

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Tree decomposition

Definition (Graph decomposition)

A tree T is a decomposition of a graph G when its vertices are arranged satisfying the following properties :

- 1 If u and v are neighbors in G , then there is a bag of T containing both of them.
- 2 For every vertex v of G , the bags of T containing v form a connected subtree

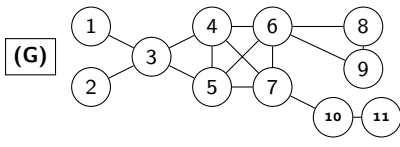
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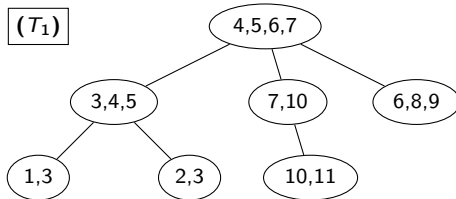
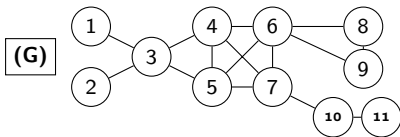
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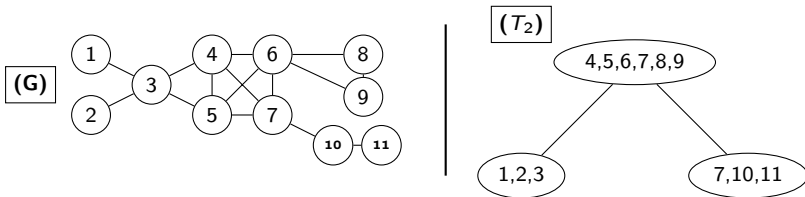
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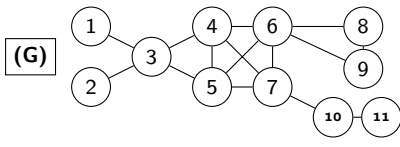
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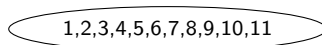
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Example :



(T_3)



Treewidth

Definition (treewidth)

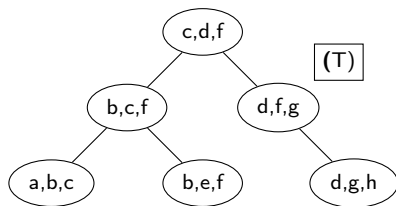
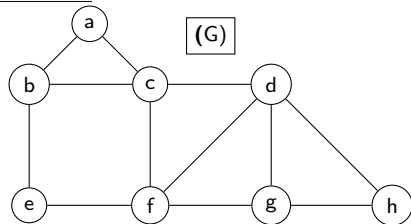
- The *width* of a decomposition is largest bag size - 1.
- The *treewidth* of a graph is the width of the best decomposition of this graph.

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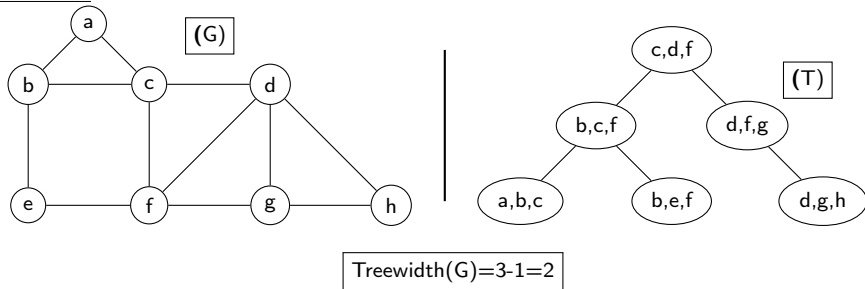


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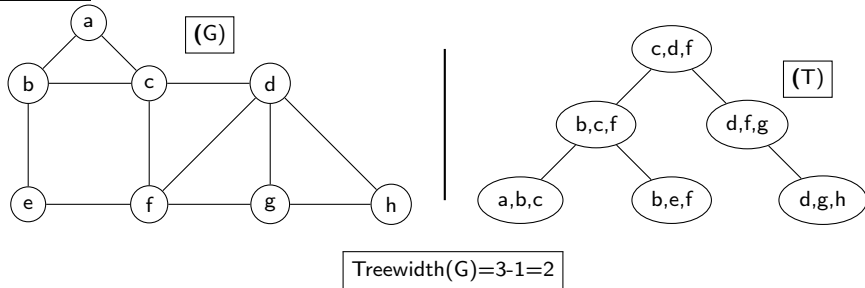


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Remark

The treewidth of a tree is 1 and if a graph have a treewidth of 1 we can claim that this graph is a forest (i.e. a collection of trees).

Nice tree

Definition (nice tree)

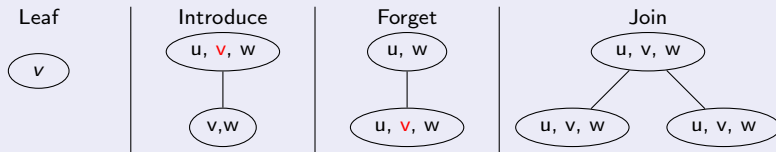
A tree decomposition is *nice* if every node x is one of the following 4 types :

Leaf : no children, $|B_x| = 1$

Introduce : 1 child y , $B_x = B_y \cup \{v\}$ for some vertex v

Forget : 1 child y , $B_x = B_y \setminus \{v\}$ for some vertex v .

Join : 2 children y_1, y_2 with $B_x = B_{y_1} = B_{y_2}$



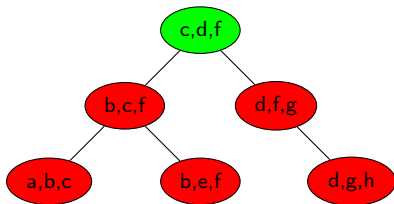
Remark

- A tree decomposition can be turned into a nice tree decomposition
- A nice tree could be very good to simplify a proof or to find an easy program to solve a problem (as we will see later).

example

This is a decomposition tree saw previously. How could we obtain a nice tree from it?

c,d,f

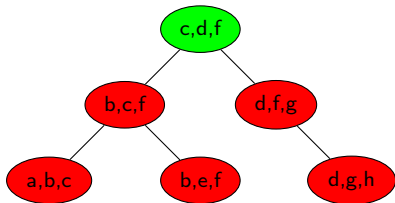


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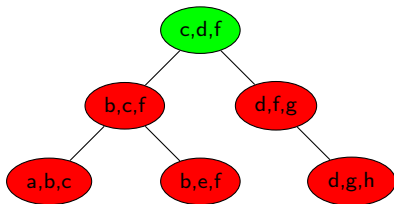
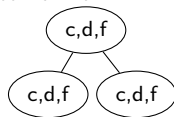
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JOIN



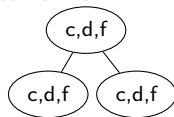
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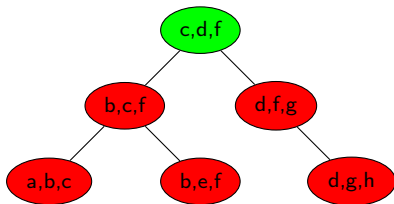


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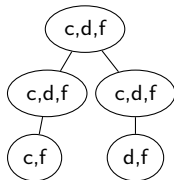
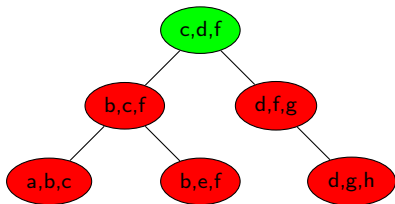


FORGET



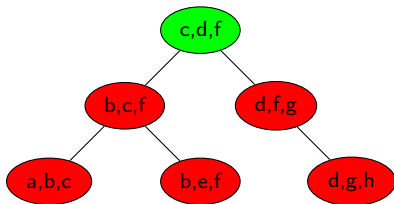
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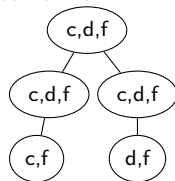


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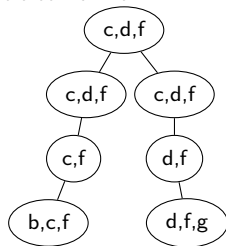
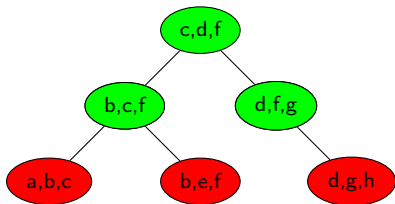


INTRODUCE



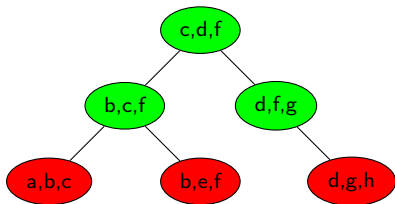
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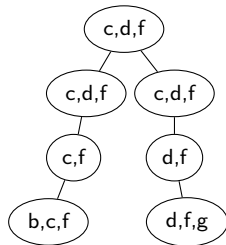


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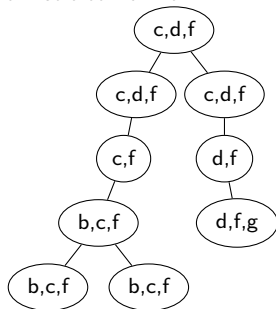
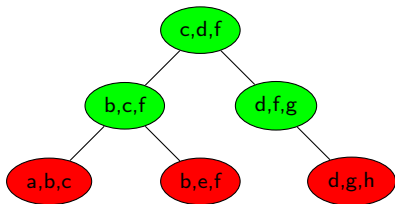


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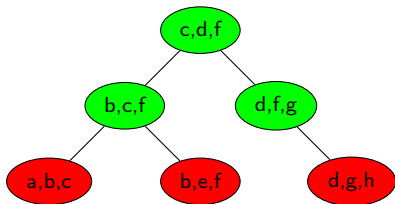
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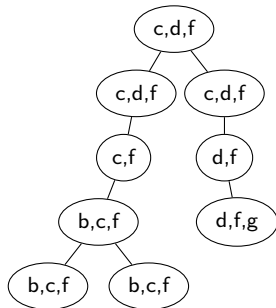


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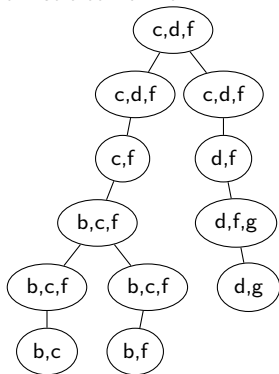
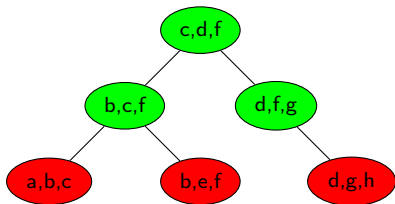


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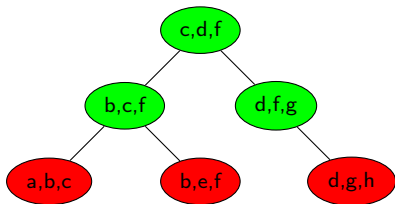
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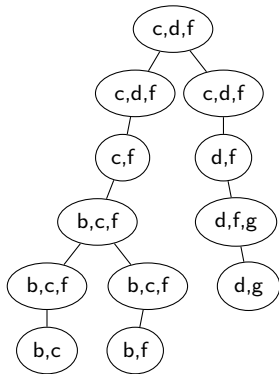


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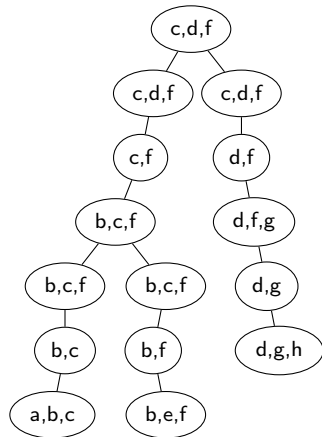
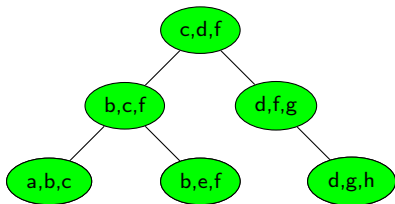


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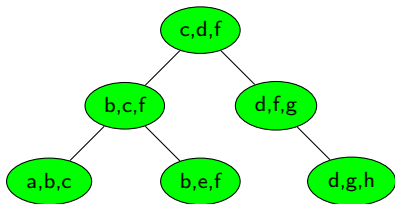
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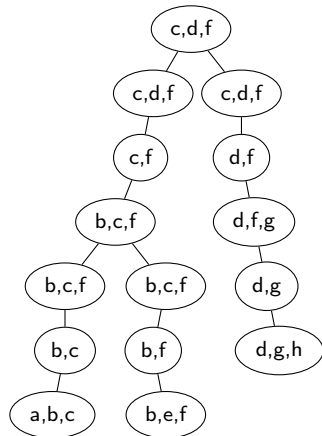


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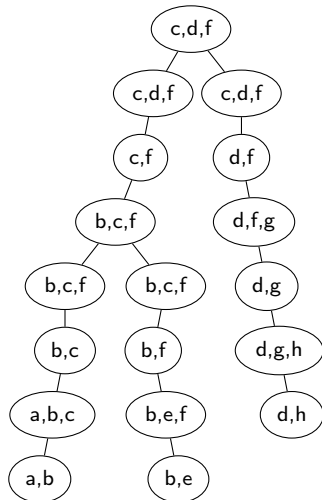
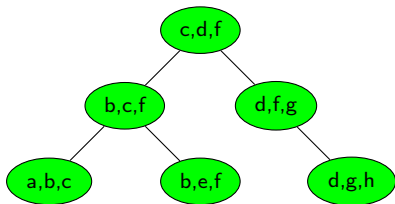


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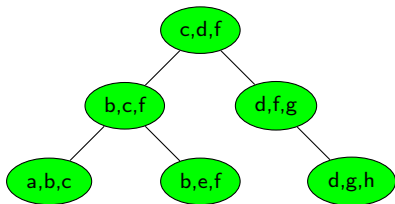
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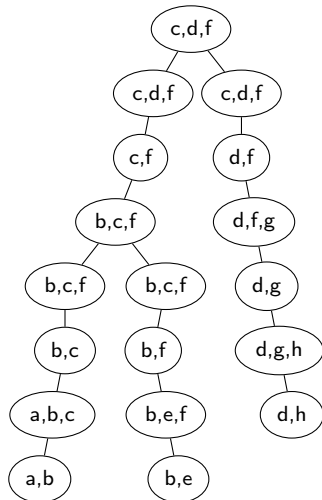


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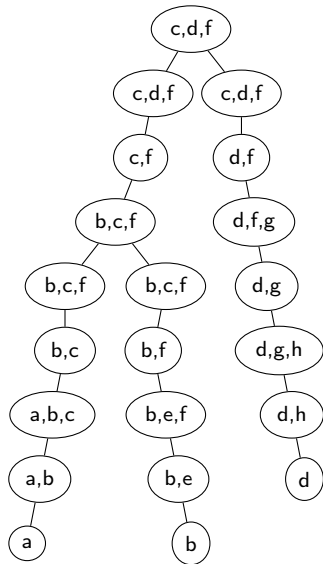
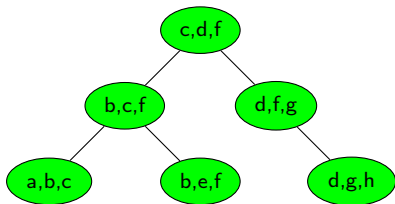


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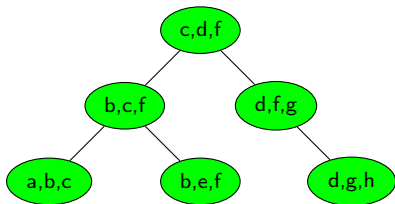
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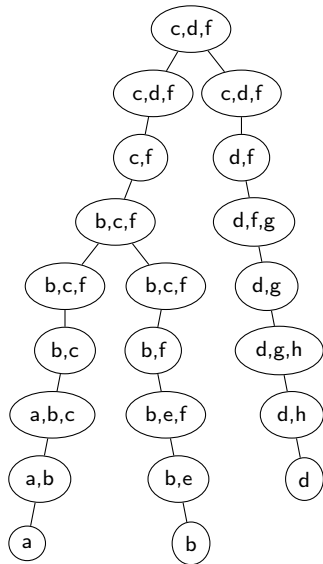


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LEAF



the problem

Problem (k-color)

problem : *Let (G) be a graph and k an integer. We want to know if it is possible to draw each vertex of the graph such that two neighbors have never the same color and with only k color.*

This problem is a problem of decision which is NP.

tree-width : *k -color is possible for a graph (G) if and only if $k \geq \text{treewidth}(G)$.*

nice tree : *a nice tree of (G) give a way to find a k -coloration of (G) (if $k \geq \text{treewidth}(G)$).*

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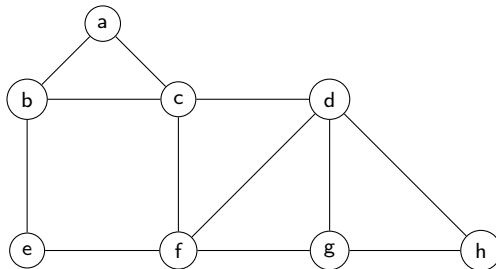
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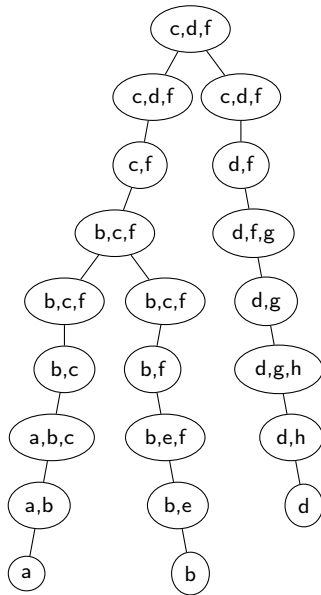
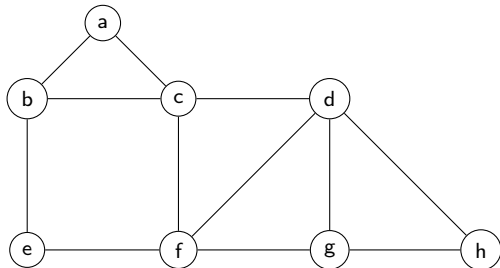
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Illustration : we will use the previously trees to solve the problem with this graph :



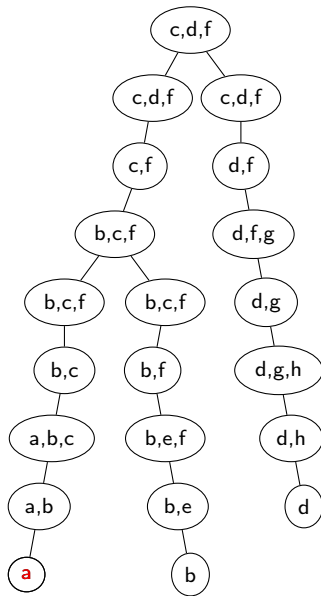
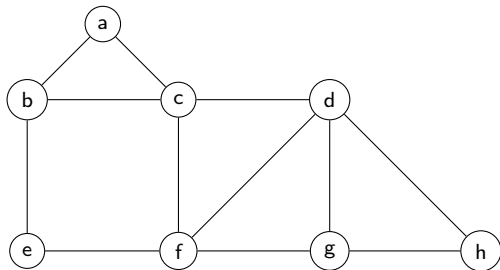
nice tree of the graph

- $\text{treewidth}(G)=3$: we can solve the problem with 3 colors
- we can fix a color for a “Leaf”
- when we meet an “Introduce” we add a color
- when we meet a “Forget” we can claim that the vertex which has disappeared won't come back (property of the decomposition) and so we can reuse the color.



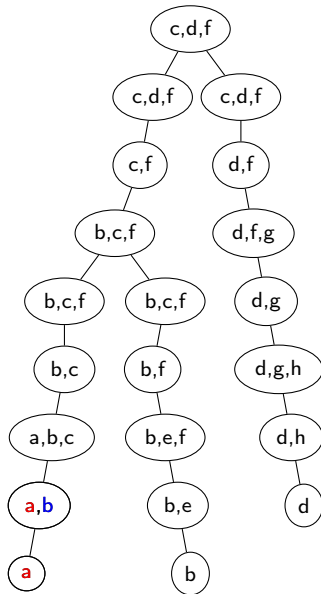
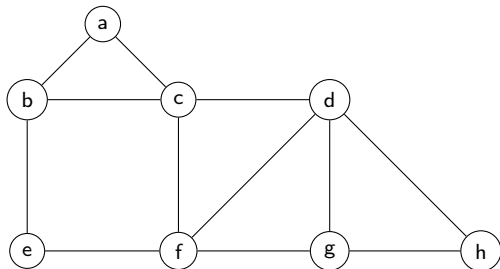
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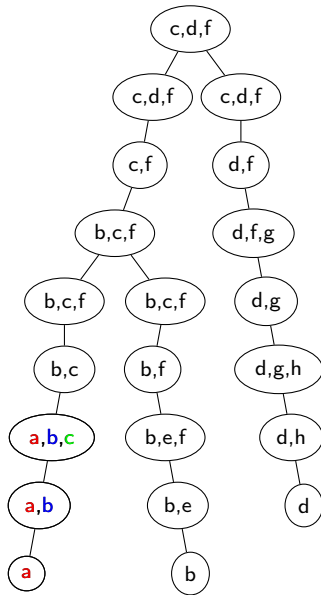
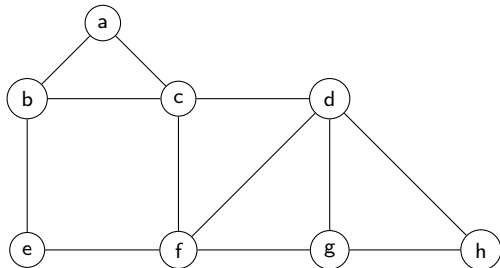
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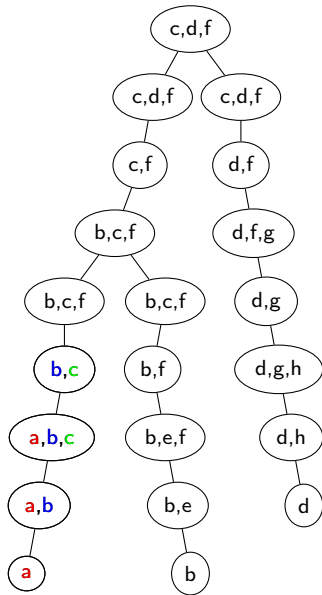
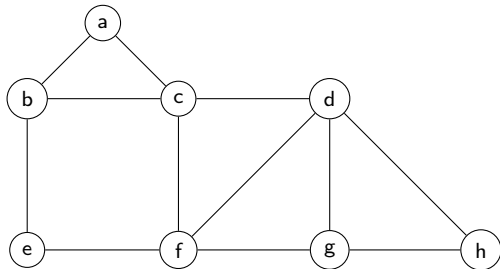
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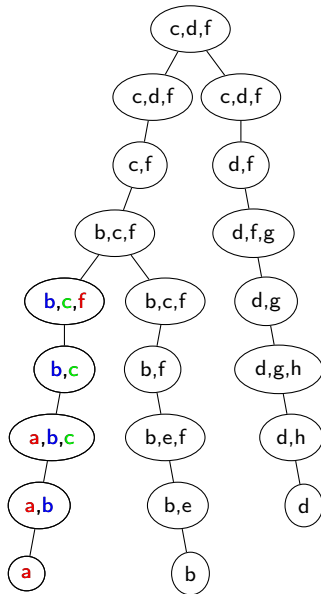
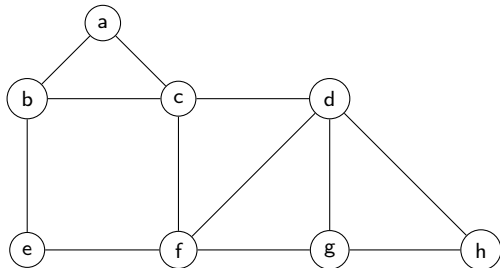
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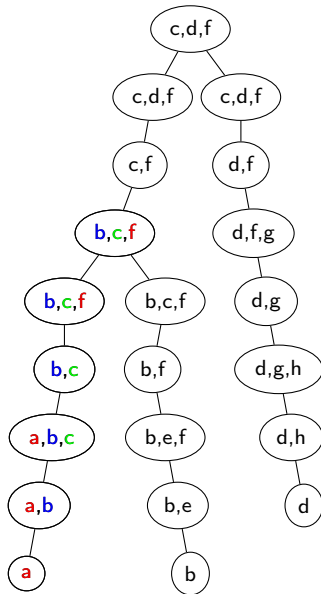
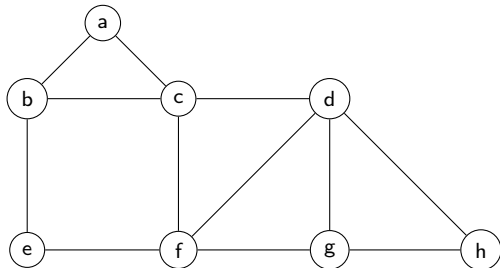
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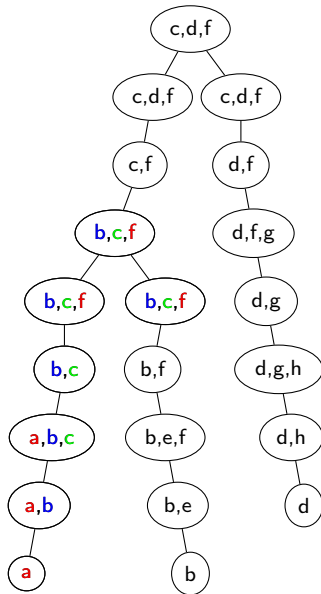
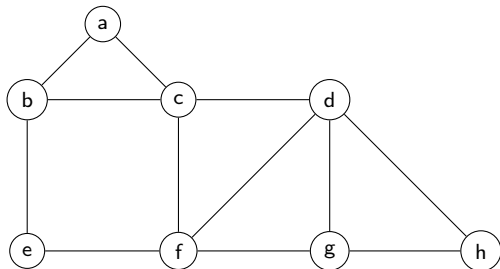
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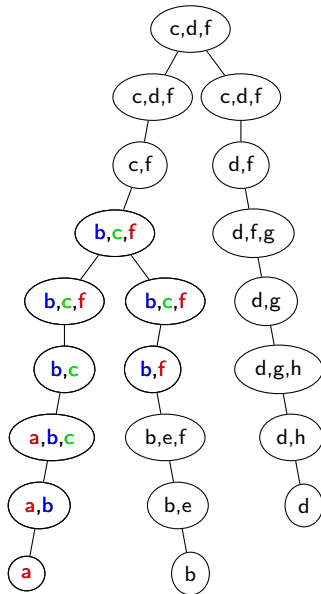
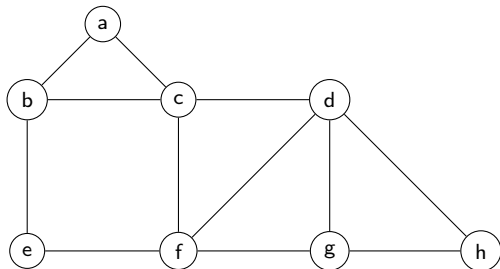
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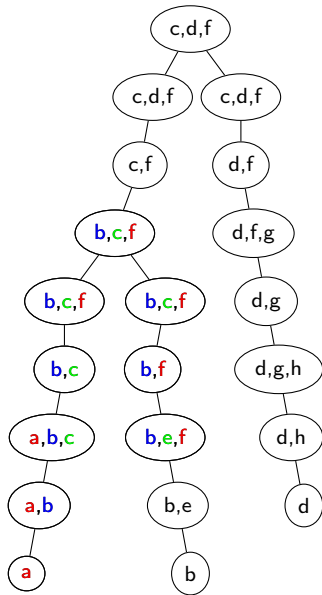
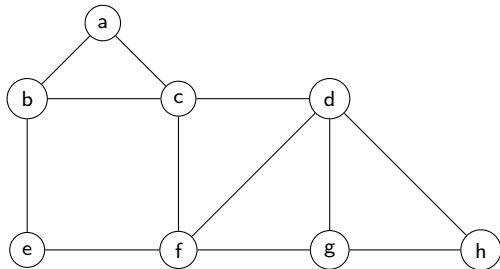
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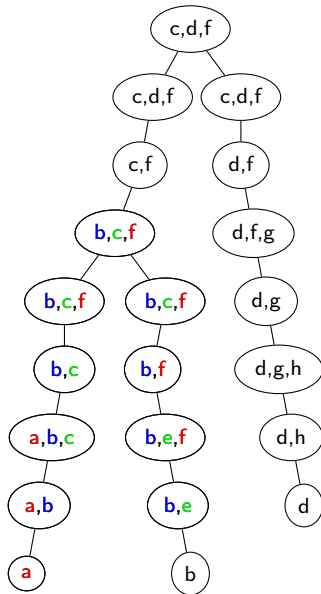
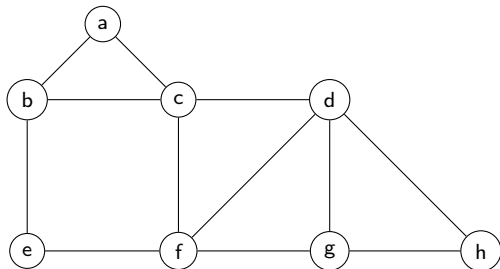
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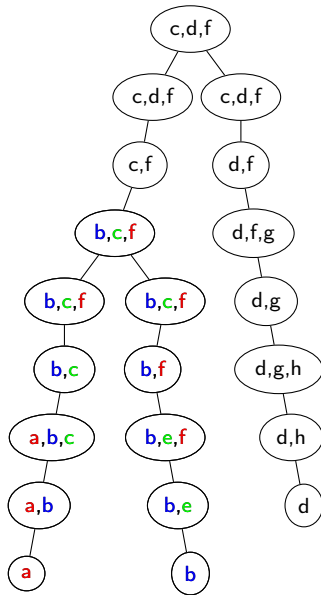
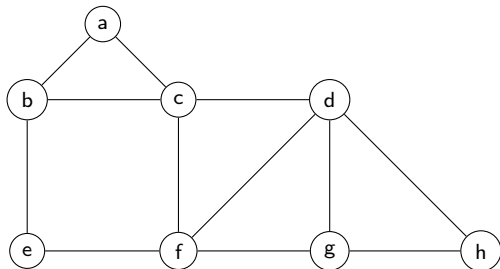
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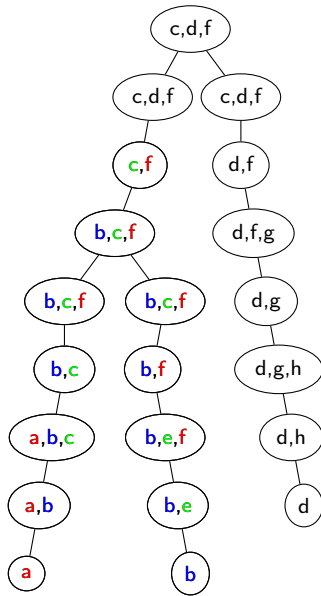
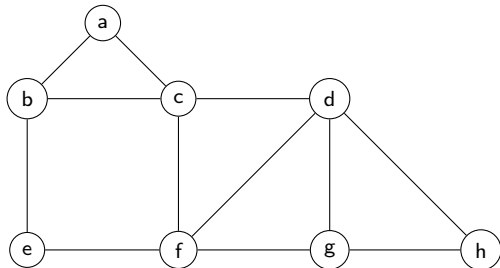
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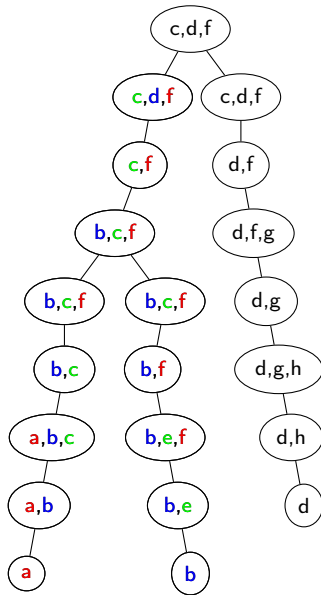
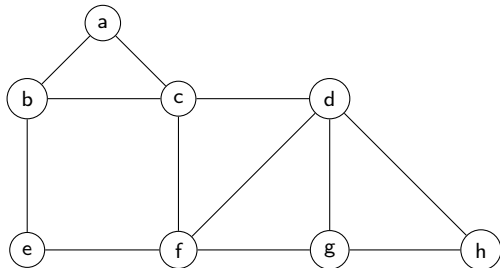
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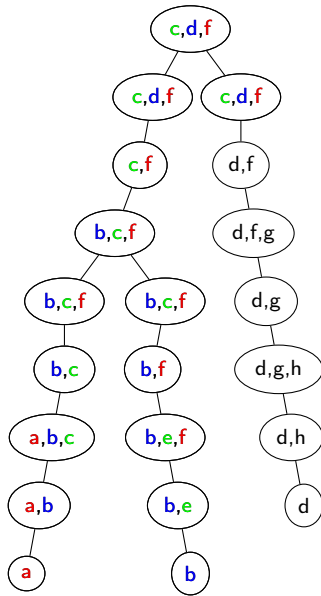
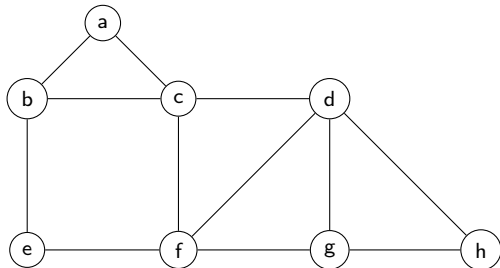
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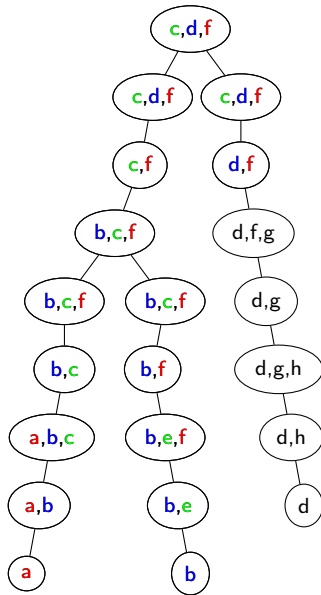
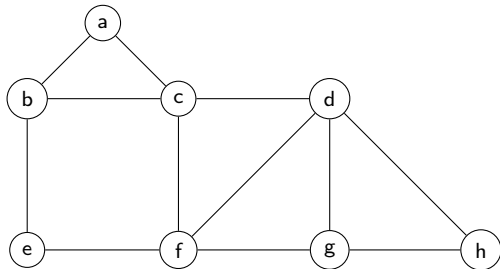
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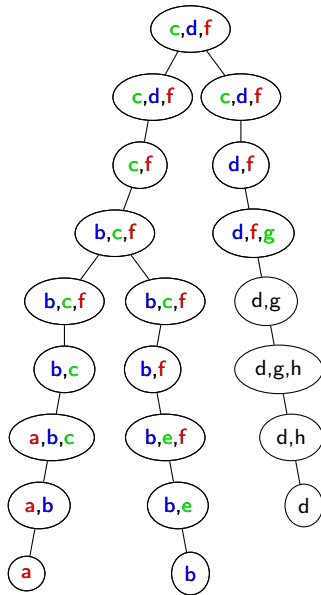
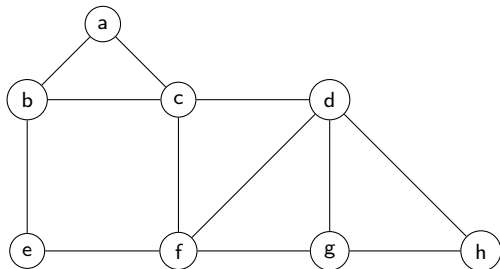
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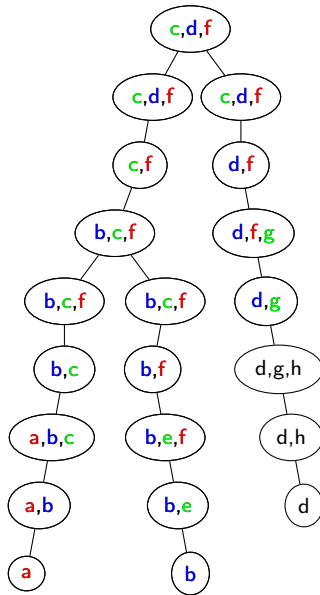
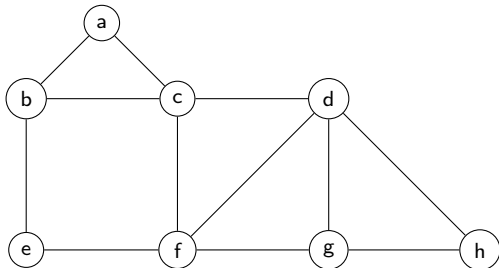
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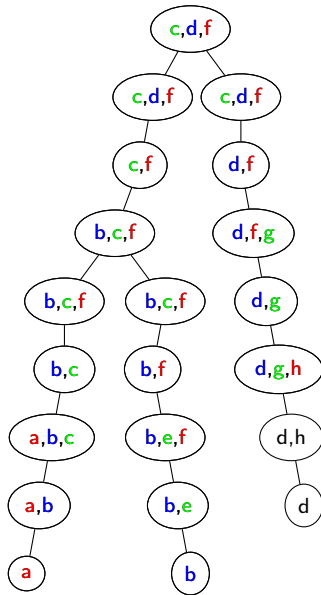
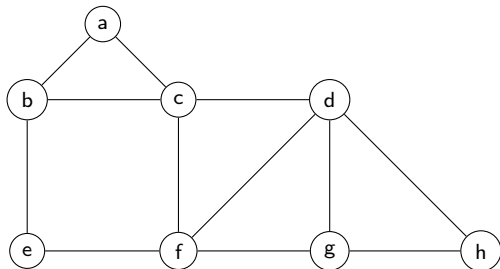
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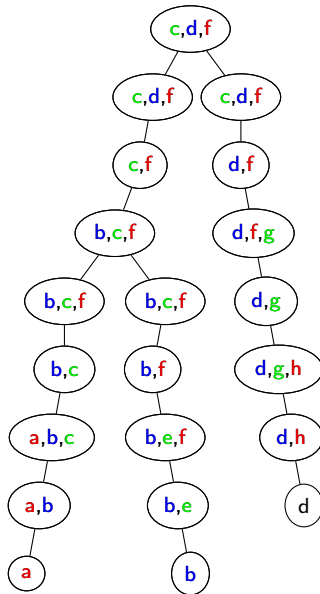
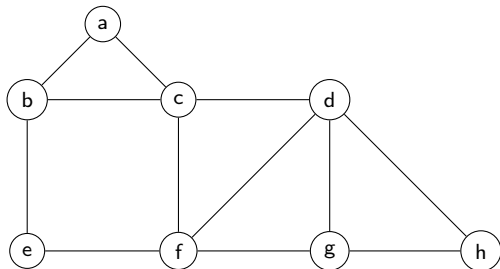
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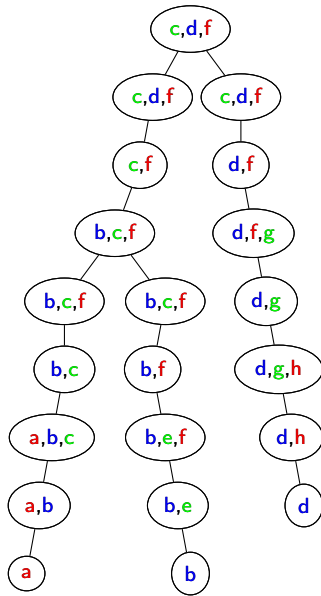
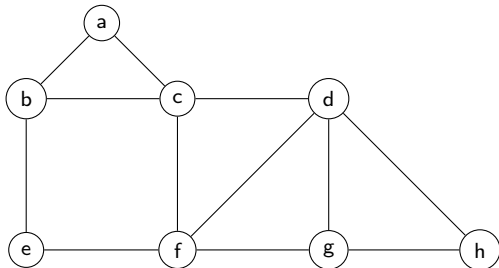
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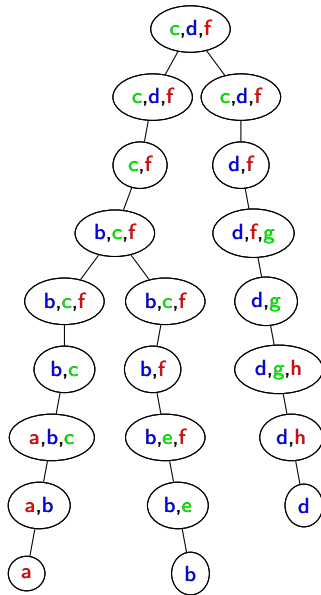
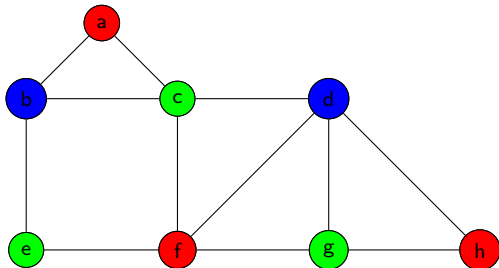
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Other applications

Mission

Bibliography