Graphs' decompositions
and
resolutions of combinatorial problems

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Objectives

Objectives (of my project)

- *The first goal of this project is, starting from a graph whose tree representation is known, to solve corresponding combinatorial problems.*
  The k-coloring problem, the max clique problem or the Hamilton path problem can be explored.
- *The second purpose is to implement a graph decomposition calculator.*
- *Finally, it can be considered extensions by working on the efficiency of implementations on large-size structures, or improving the shape of displayed results.*

Contents

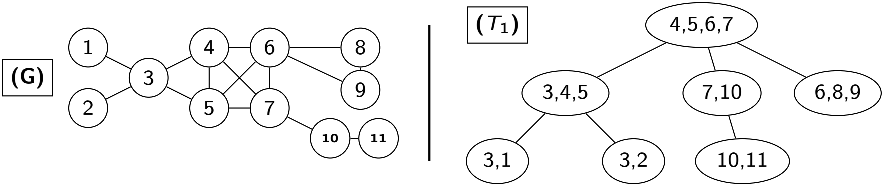
Tree decomposition of a graph

Definition (Tree decomposition of a graph)

A tree T is a tree decomposition of a graph G where its nodes are arranged satisfying the following properties :

1. If u and v are neighbors in G, then there is a bag of T containing both of them (a bag is a node of the tree).
2. For every vertex v of G, the bags of T containing v form a connected subtree

example :



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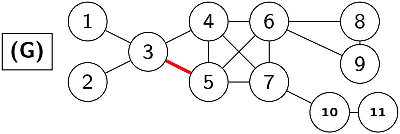

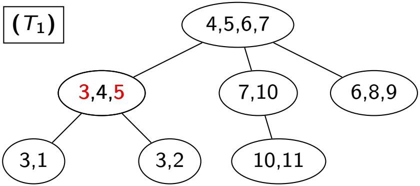

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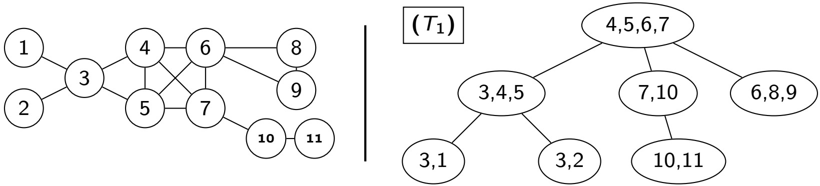
example :

(G)

(T_1)



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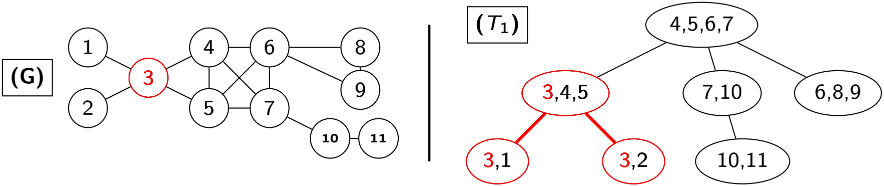

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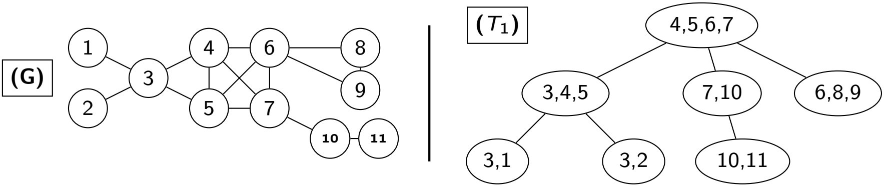

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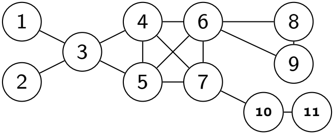

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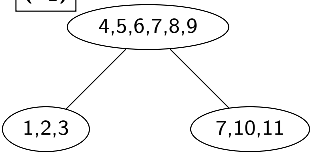

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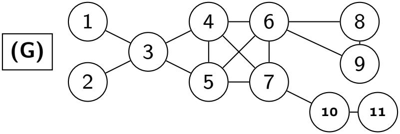
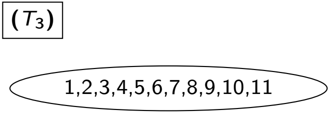


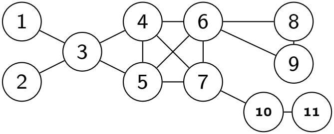


Definition (treewidth)

- The *width* of a decomposition is (largest bag size - 1).
- The *treewidth* of a graph is the lowest width of all decompositions

example :

(G)



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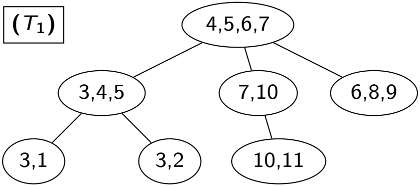

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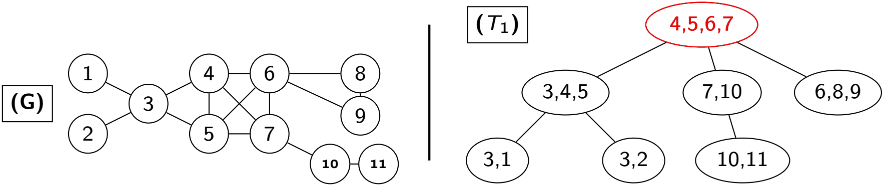

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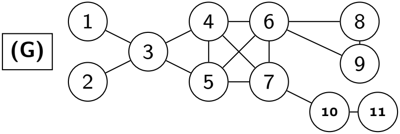

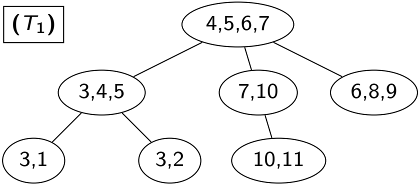

width(T_1)=4-1=3

Definition (treewidth)

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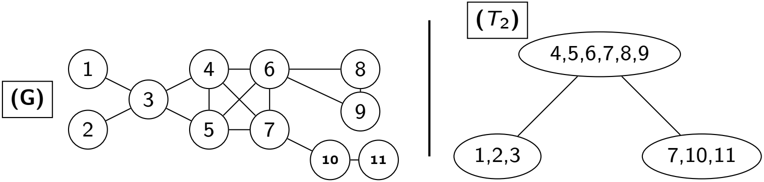
example :



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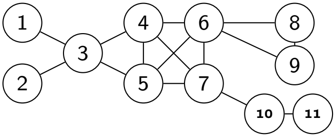
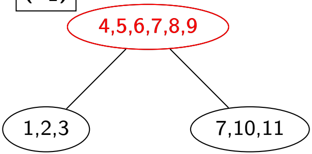


(T_2)



Treewidth

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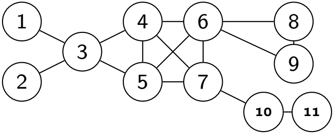
example :

(G)



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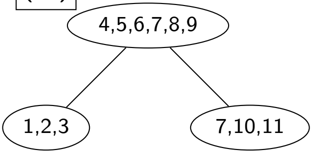

width(T_2)=6-1=5

Definition (treewidth)

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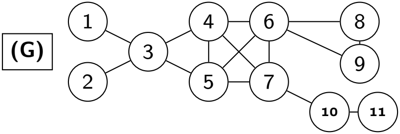
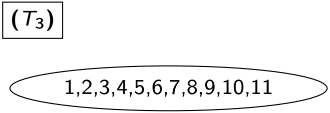
- The *treewidth* of a graph is the lowest width of all decompositions

example :





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example :

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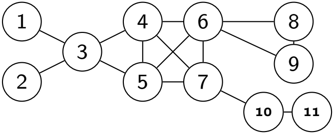

(T_3)

Definition (treewidth)

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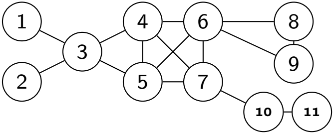
- The *treewidth* of a graph is the lowest width of all decompositions

example :

(G)



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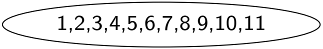
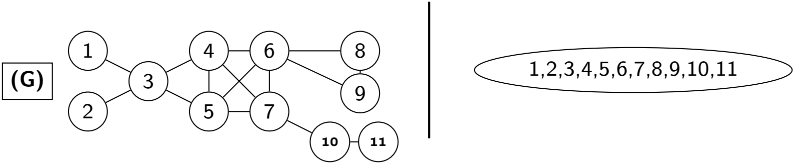


width(T_3)=11-1=10

Definition (treewidth)

- The *width* of a decomposition is (largest bag size - 1).
- The *treewidth* of a graph is the lowest width of all decompositions of this graph.

example :

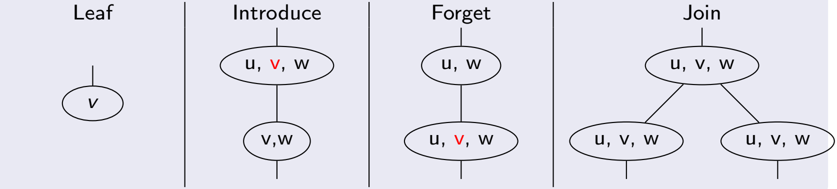


TreewidthG=Min(width(T_k)=3

Nice tree

Definition (nice tree)

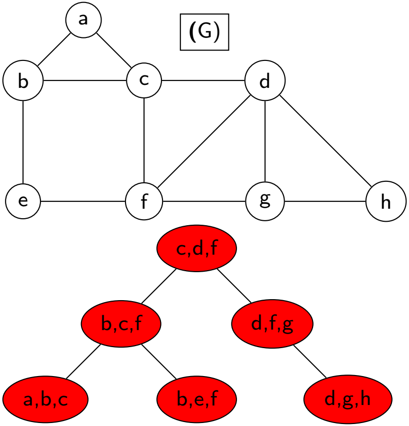
A tree decomposition is *nice* if every node is one of the following 4 types :



Remark

- *A tree decomposition can be turned into a nice tree decomposition*
- *A nice tree can be used to simplify a proof, or to find an easy program to solve a problem (as we will see later on).*

This is a tree decomposition of the graph (G).



How to obtain a nice tree ?

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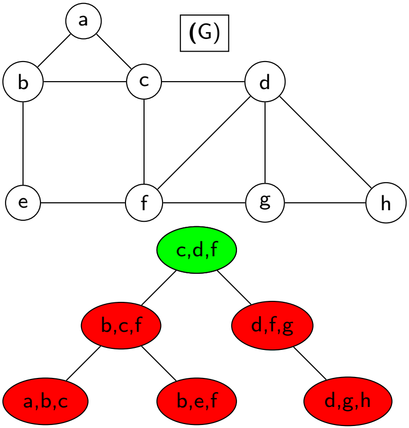

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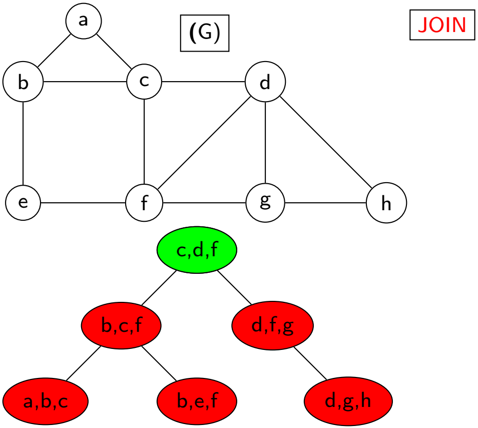

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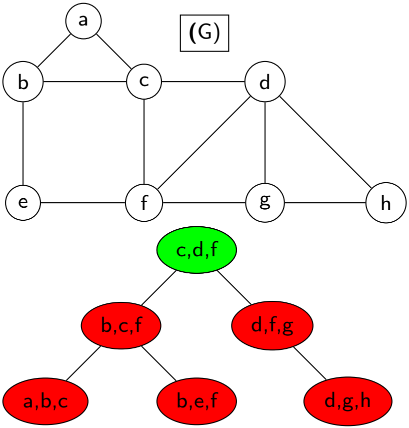

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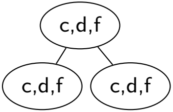

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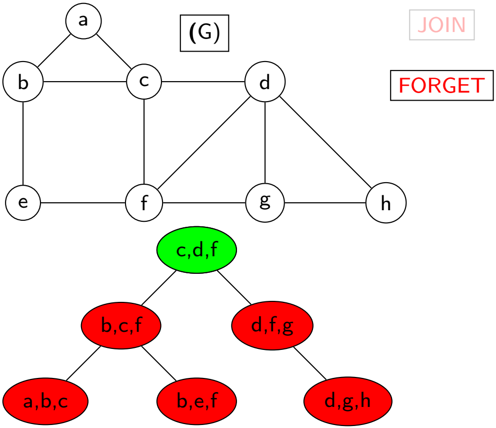

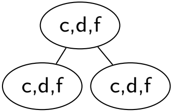

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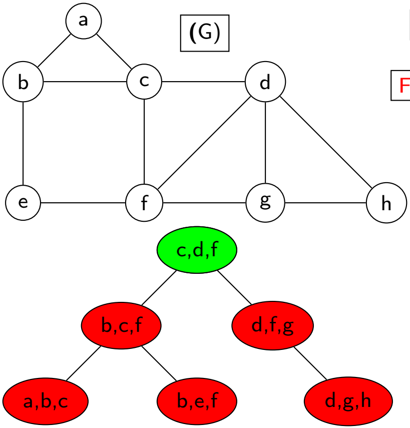

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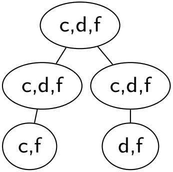

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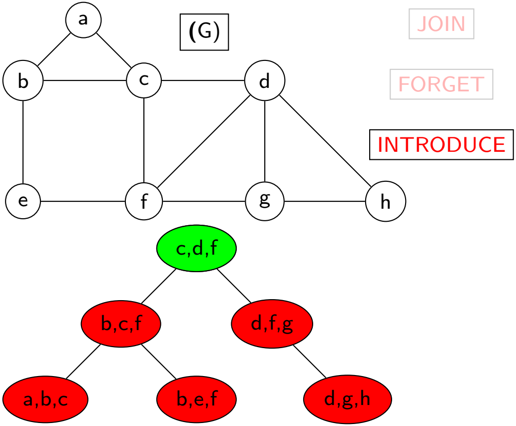

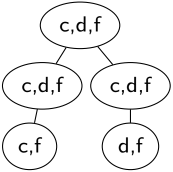

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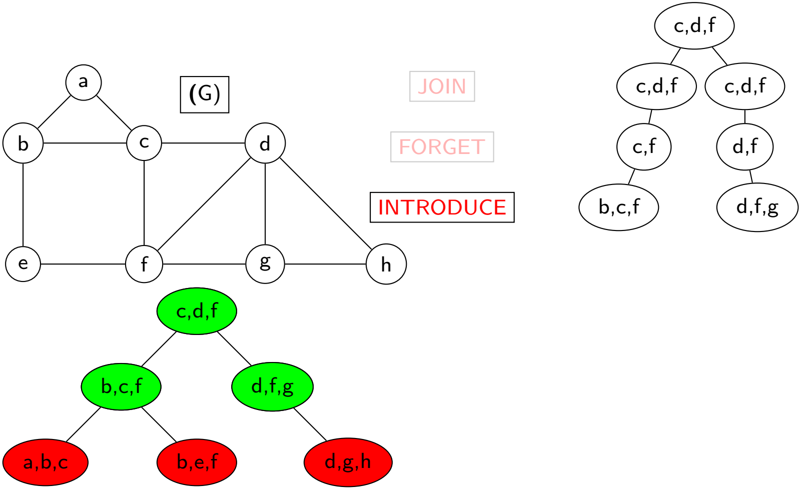

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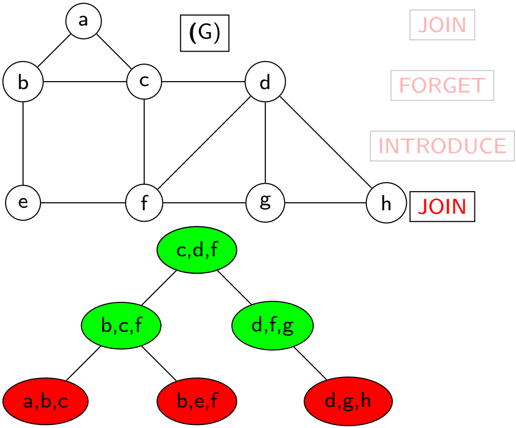

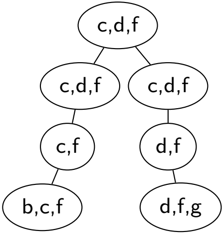

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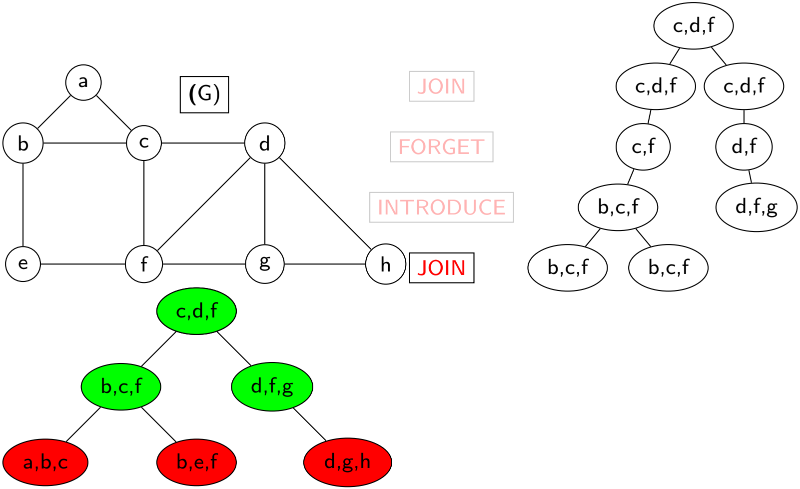

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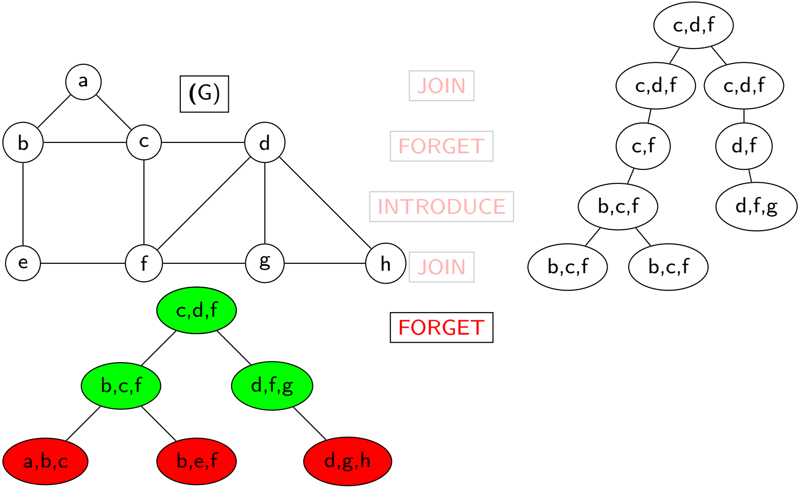

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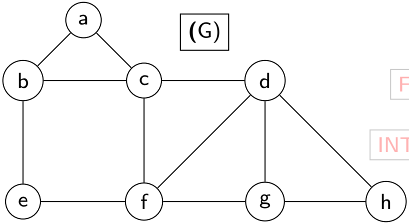

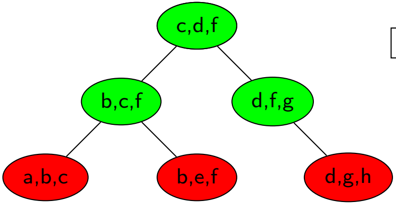

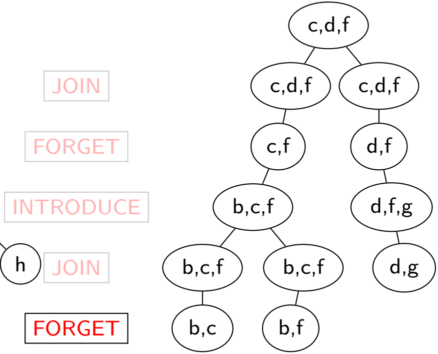

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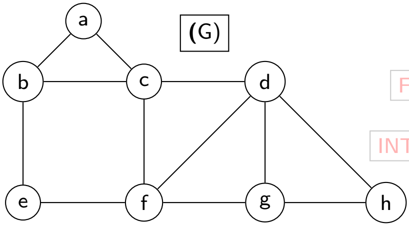

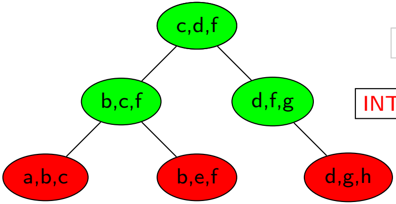

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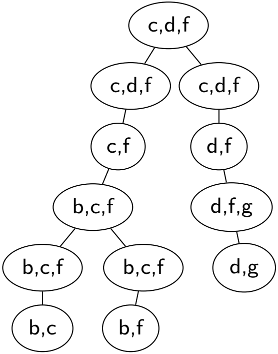

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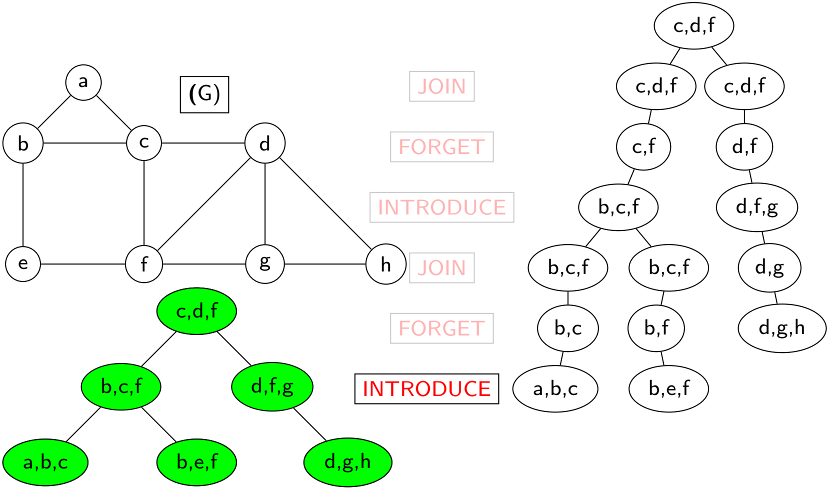

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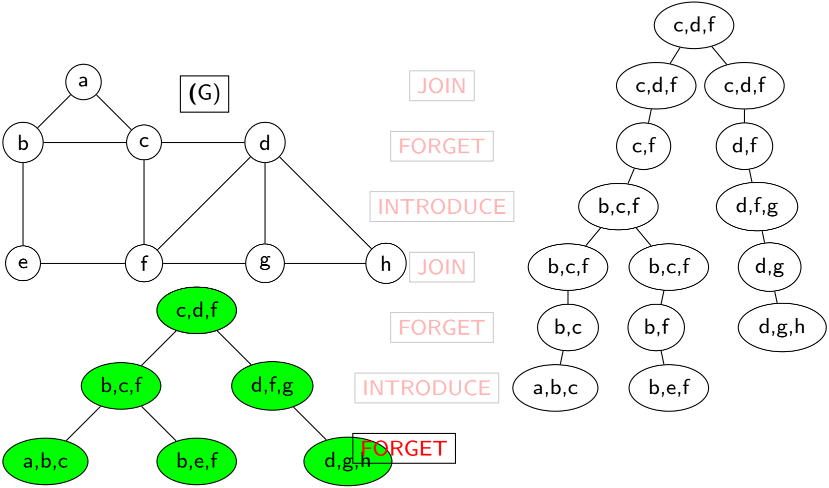

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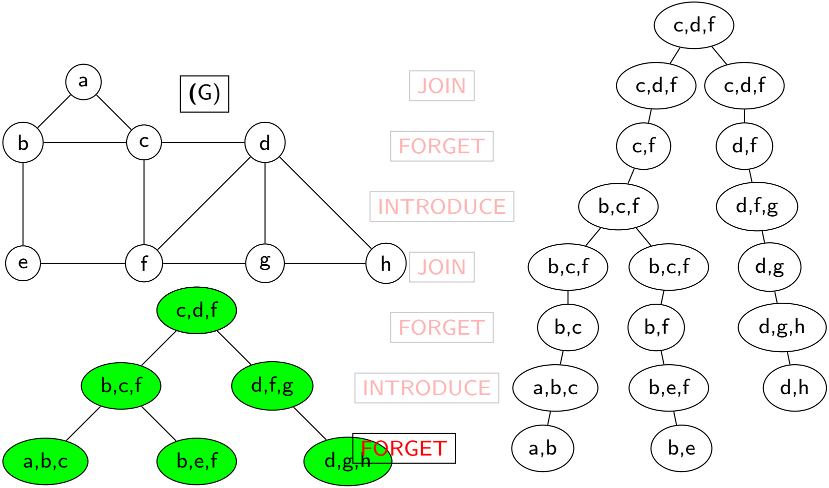

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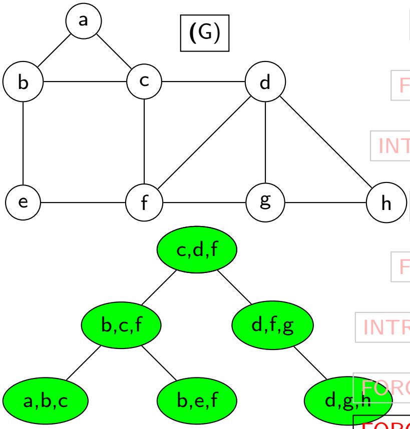

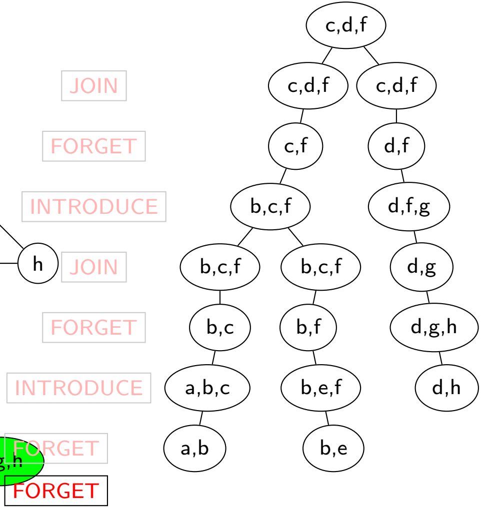

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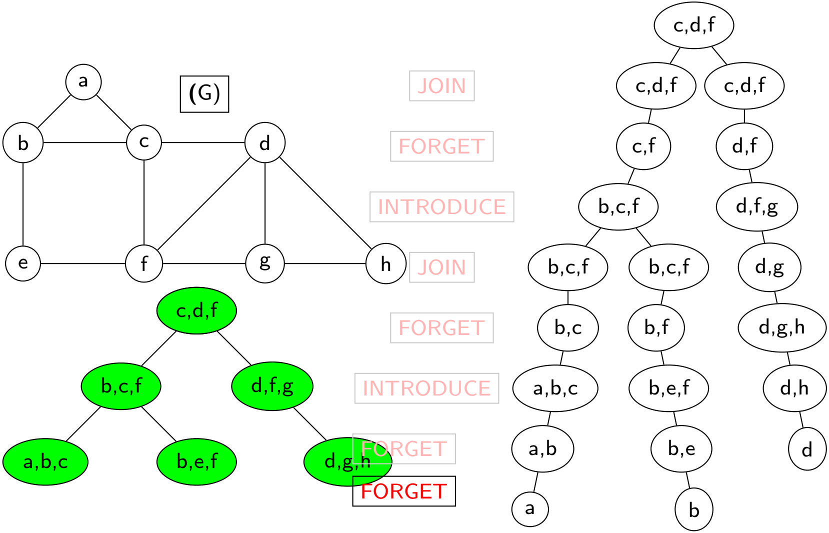

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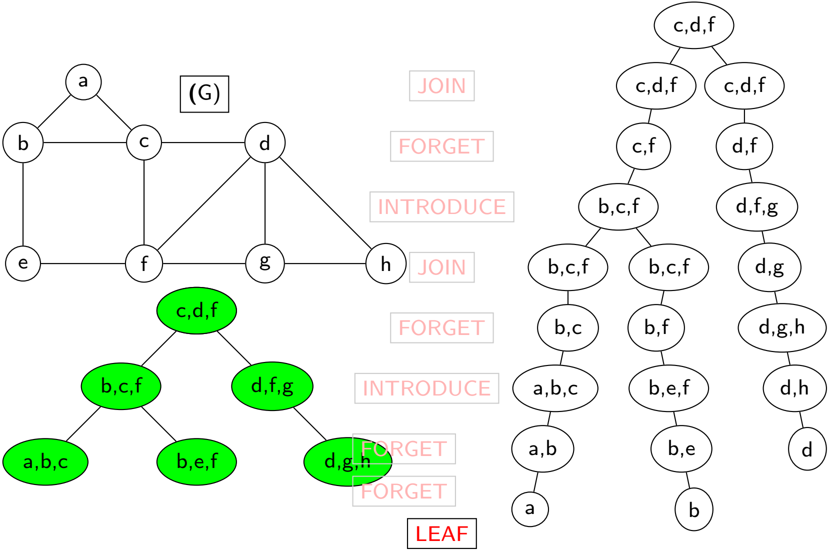

Contents

Problem (k-color)

Let (G) be a graph and k an integer. We want to know if it is possible to draw each vertice of the graph so that two neighbors have never the same color and with only k colors.
This problem is a problem of decisions problem which is NP-complet.

Solution

tree-width : *k-color is possible for a graph (G) if and only if* $k > treewidth(G)$.

nice tree : *a nice tree of (G) gives a way to find a k-coloration of (G) (if* $k > treewidth(G)$*).*

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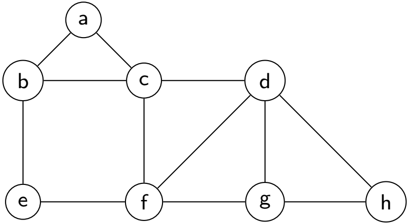
nice tree : *a nice tree of (G) gives a way to find a k-coloration of (G) (if $k >$ treewidth(G)).*

Illustration : we will use the previously nice tree to solve the problem with this graph :



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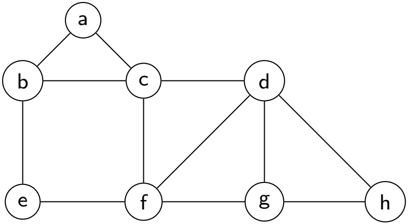

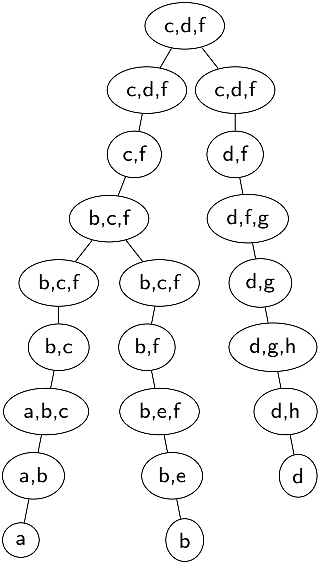

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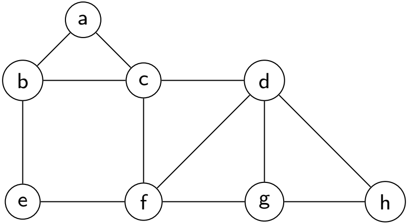

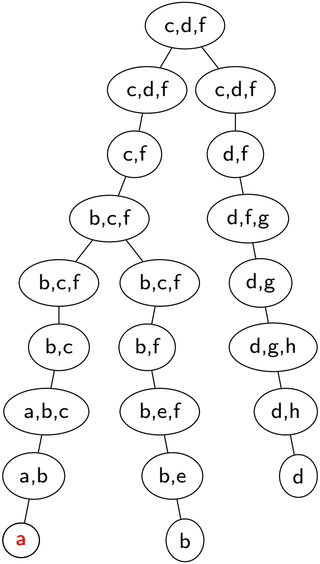

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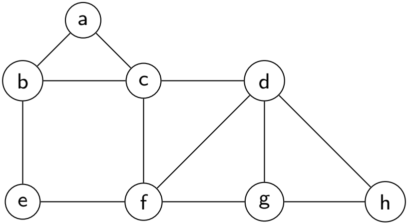

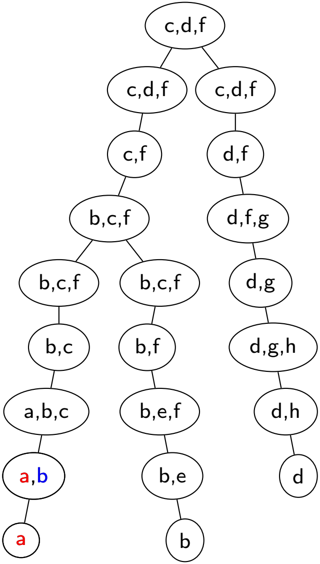

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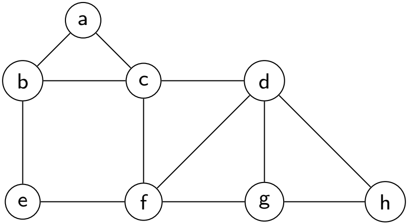

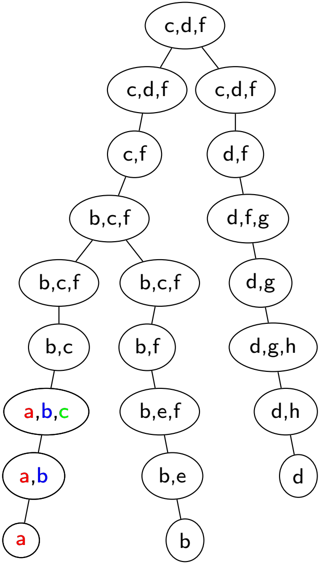

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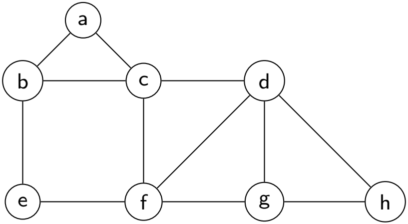

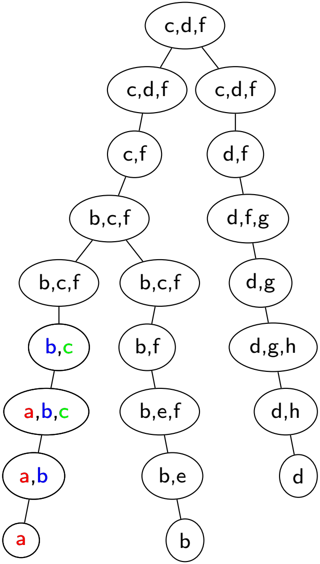

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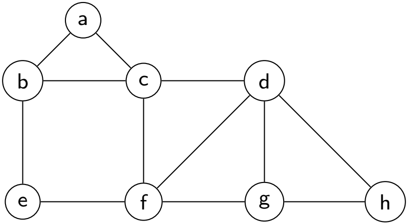

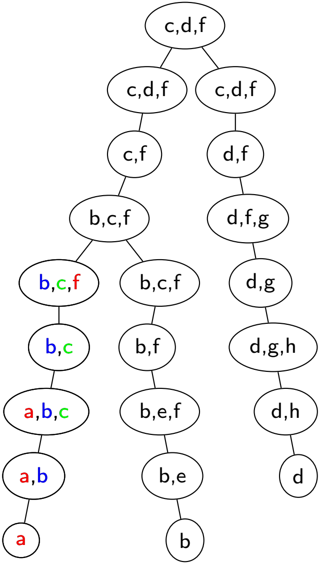

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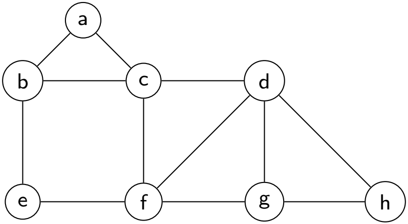

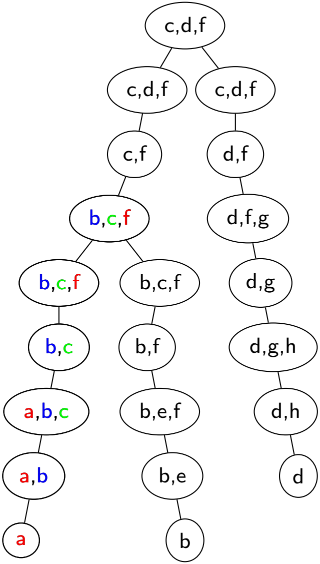

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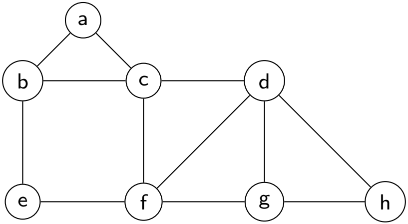

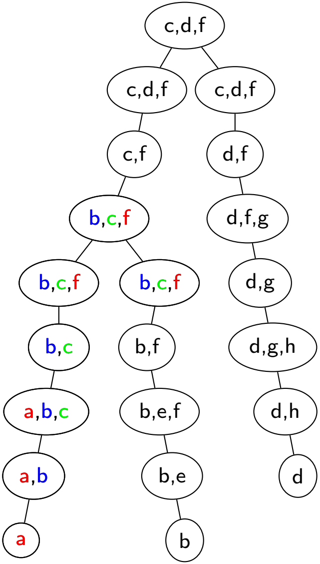

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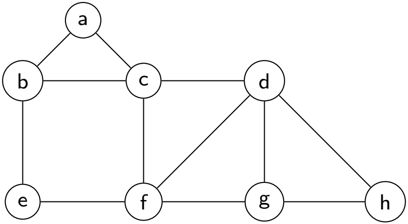

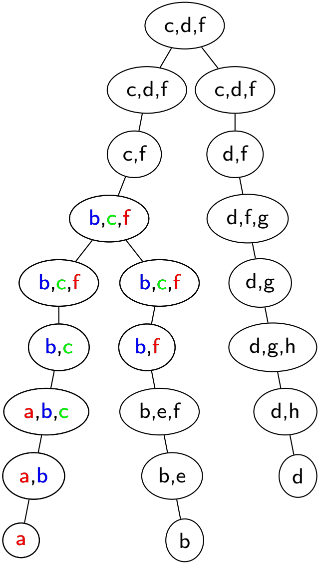

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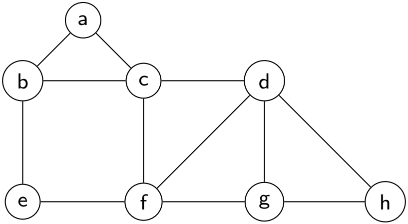

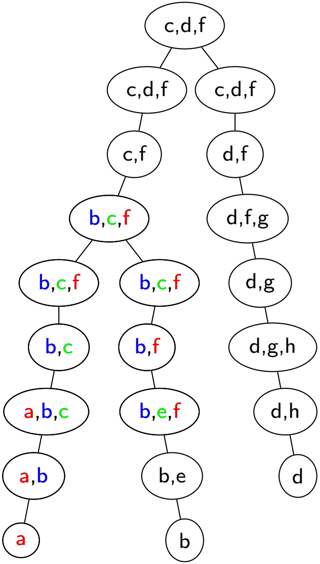

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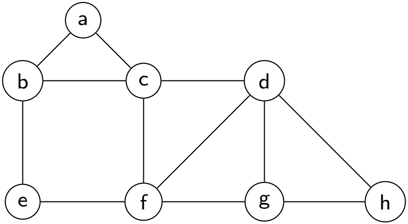

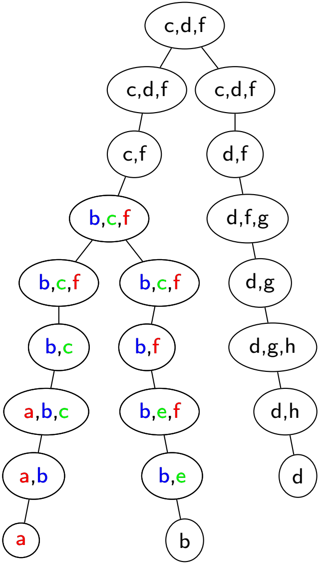

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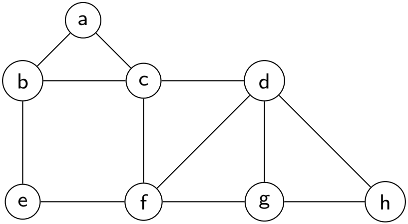

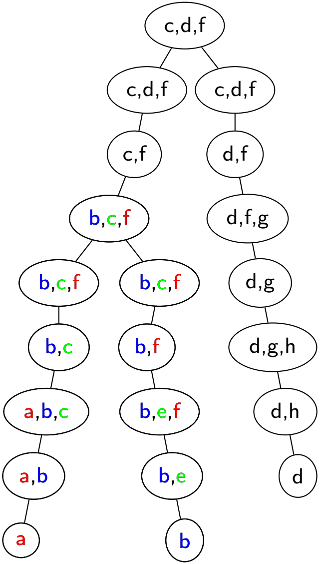

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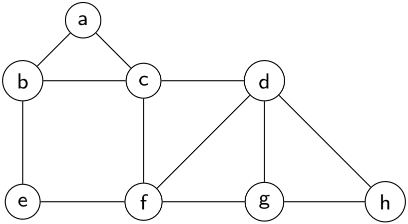

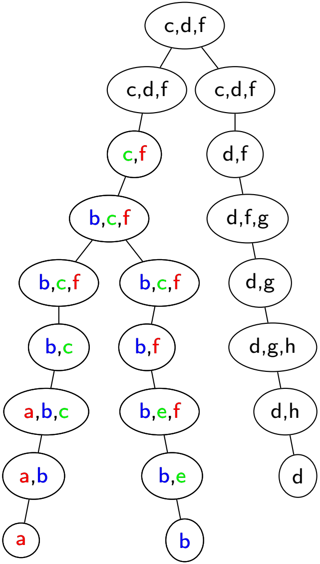

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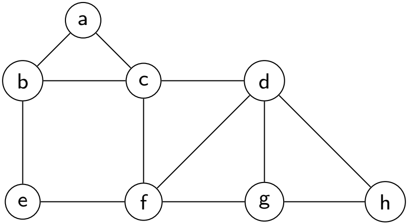

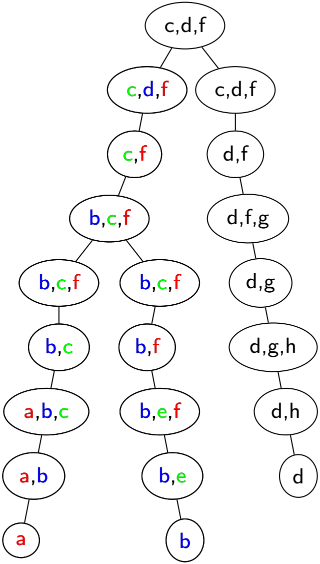

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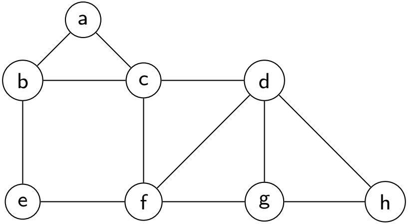

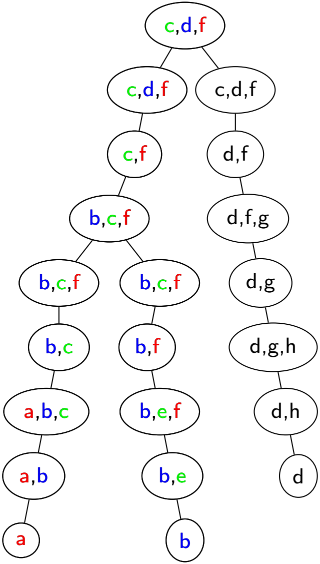

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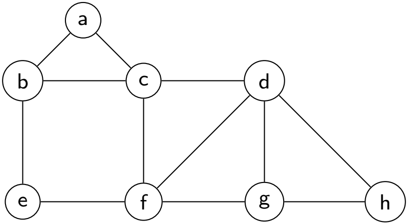

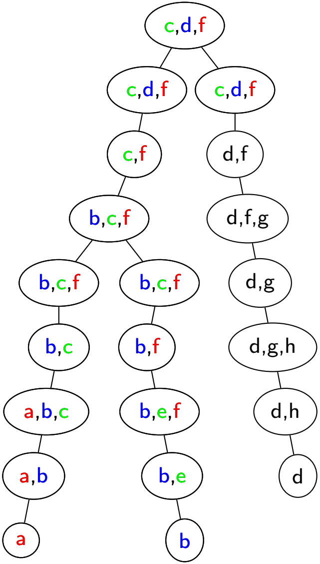

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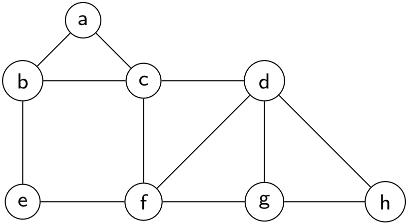

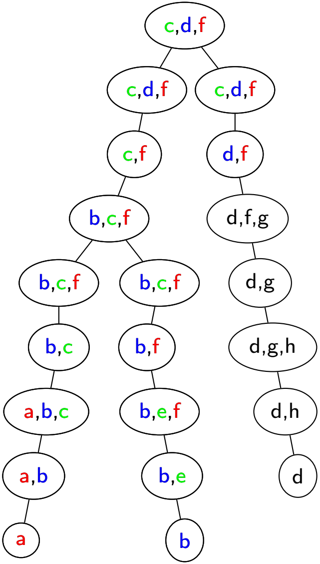

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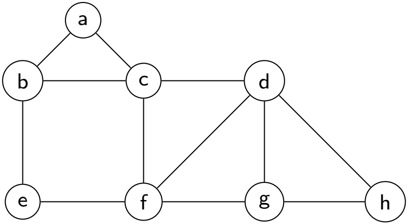

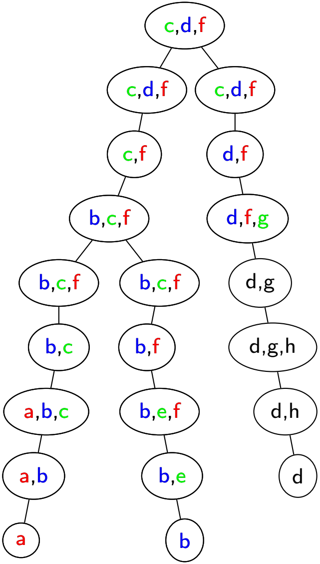

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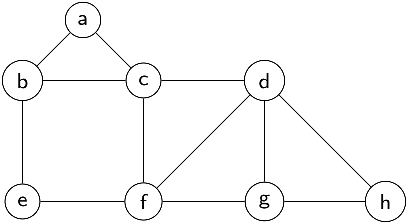

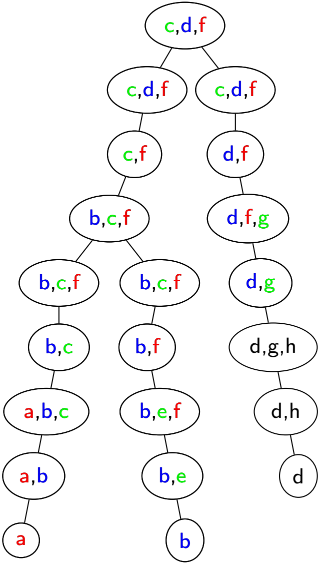

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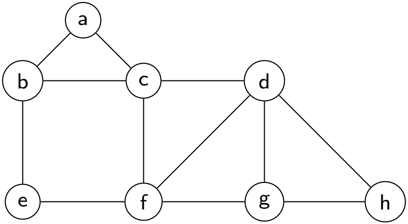

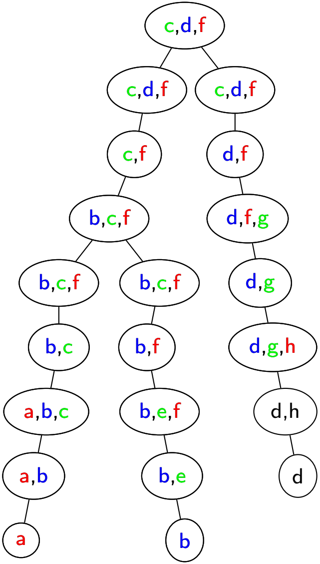

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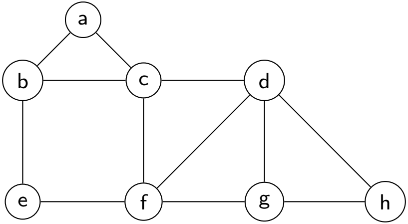

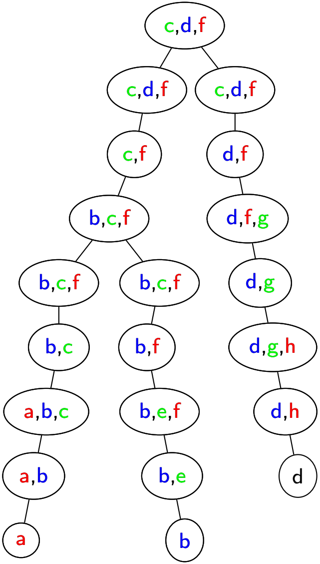

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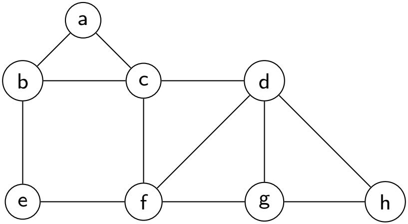

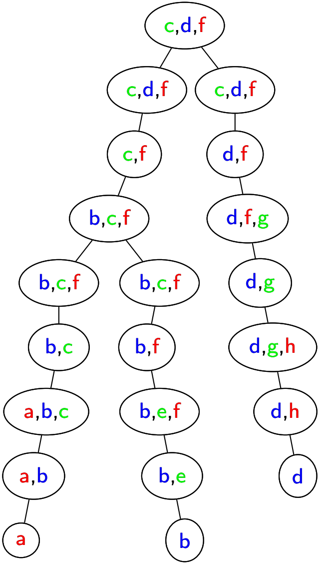

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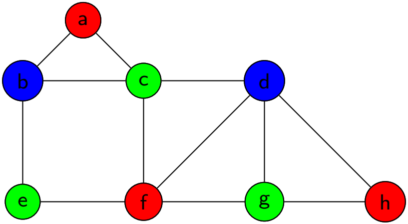

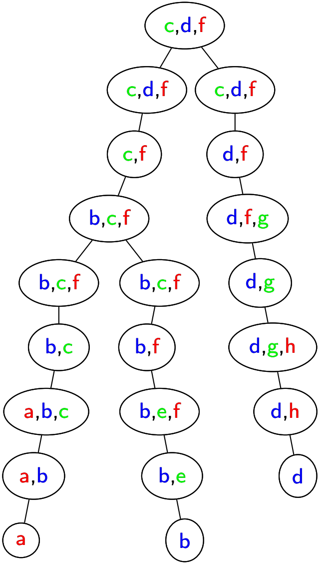

Contents

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