

Graphs' decompositions and resolutions of combinatorial problems

Stéphane Secouard
Supervised by : Florent Madeleine

Caen University - Computer science

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- 2 Tree decomposition of a graph
 - definition
 - treewidth
 - nice tree
 - Example
- 3 Application : k -color
 - the problem
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Objectives (of my project)

- *The first goal of this project is, starting from a graph whose tree representation is known, to solve corresponding combinatorial problems.
The k -coloring problem, the max clique problem or the Hamilton path problem can be explored.*
- *The second purpose is to implement a graph decomposition calculator.*
- *Finally, it can be considered extensions by working on the efficiency of implementations on large-size structures, or improving the shape of displayed results.*

Objectives (of my presentation)

Now, I am going to present :

- *important concepts to understand the goals of my project ;*
- *an example where these concepts help to solve a problem of coloration.*

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Tree decomposition of a graph

Definition (Tree decomposition of a graph)

A tree T is a tree decomposition of a graph G where its nodes are arranged satisfying the following properties :

- 1 If u and v are neighbors in G , then there is a bag of T containing both of them (a bag is a node of the tree).
- 2 For every vertex v of G , the bags of T containing v form a connected subtree

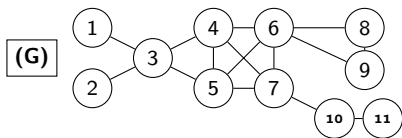
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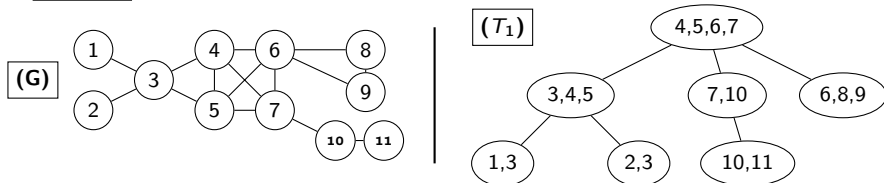
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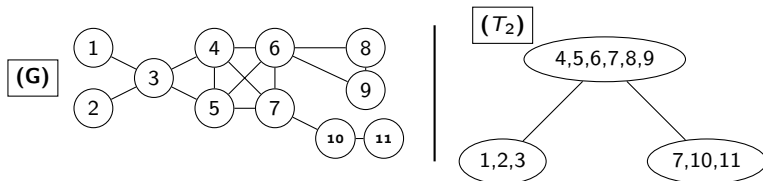
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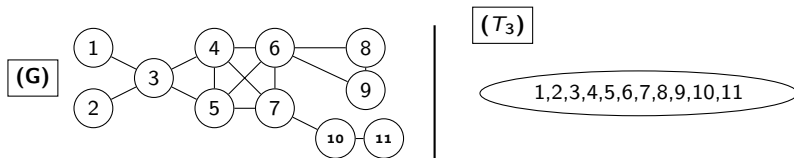
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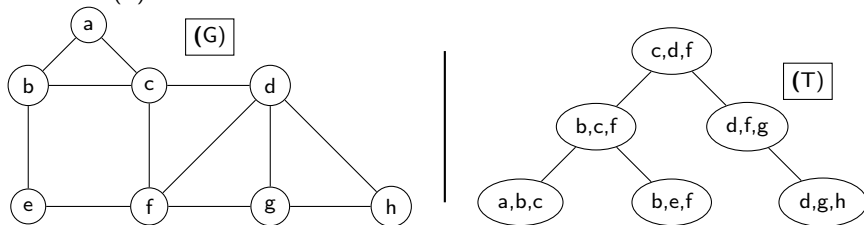
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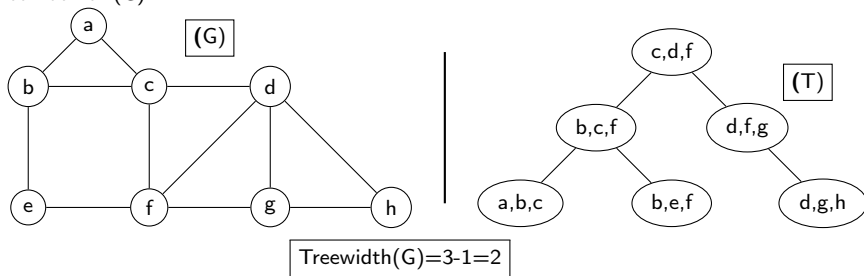
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Definition (nice tree)

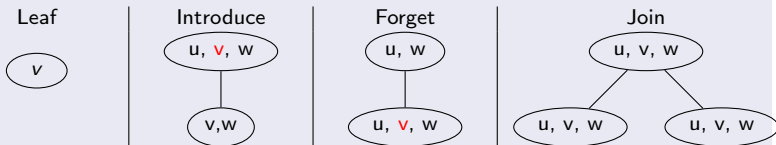
A tree decomposition is *nice* if every node x is one of the following 4 types :

Leaf : no children, $|B_x| = 1$ (B_x names a bag of the tree containing x)

Introduce : 1 child y , $B_x = B_y \cup \{v\}$ for some vertex v

Forget : 1 child y , $B_x = B_y \setminus \{v\}$ for some vertex v .

Join : 2 children y_1, y_2 with $B_x = B_{y_1} = B_{y_2}$

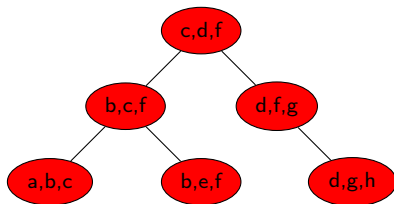


Remark

- A tree decomposition can be turned into a nice tree decomposition
- A nice tree can be used to simplify a proof, or to find an easy program to solve a problem (as we will see later on).

Example

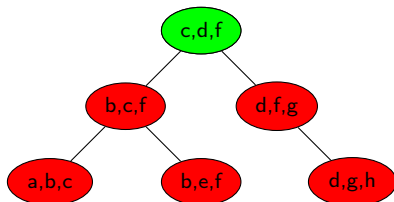
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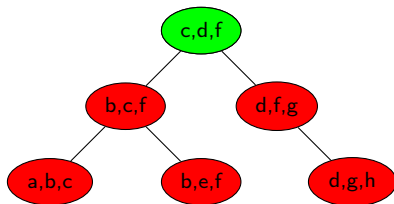


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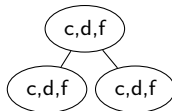
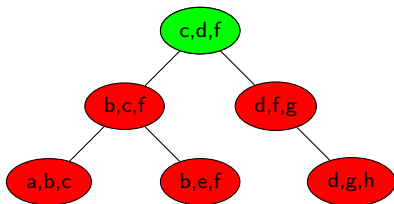
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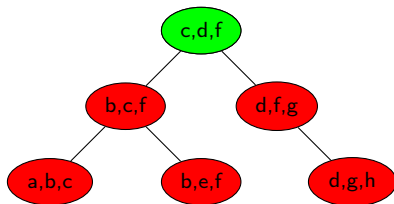
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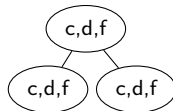
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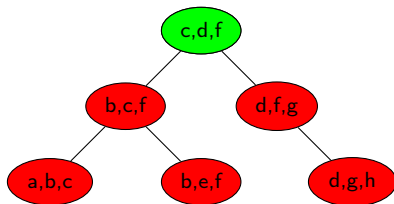
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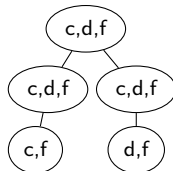
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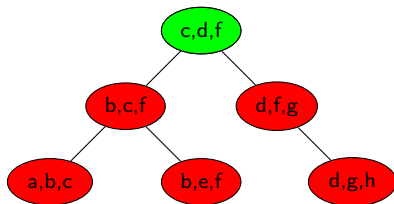
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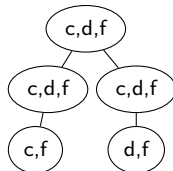
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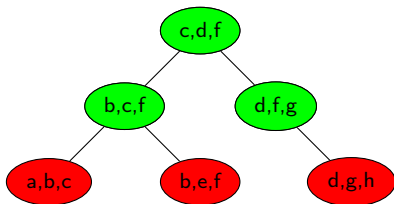
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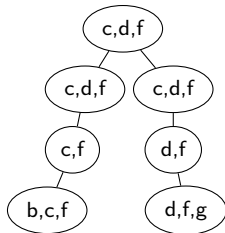
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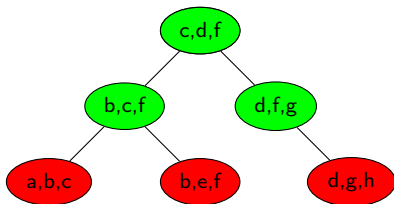
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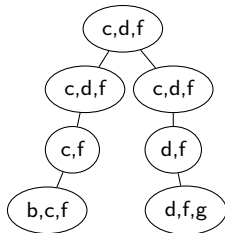


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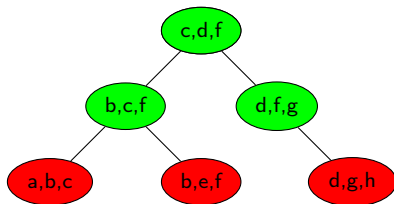
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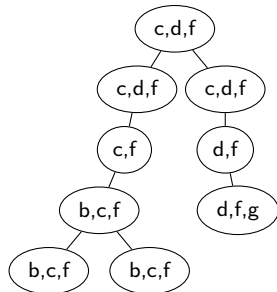


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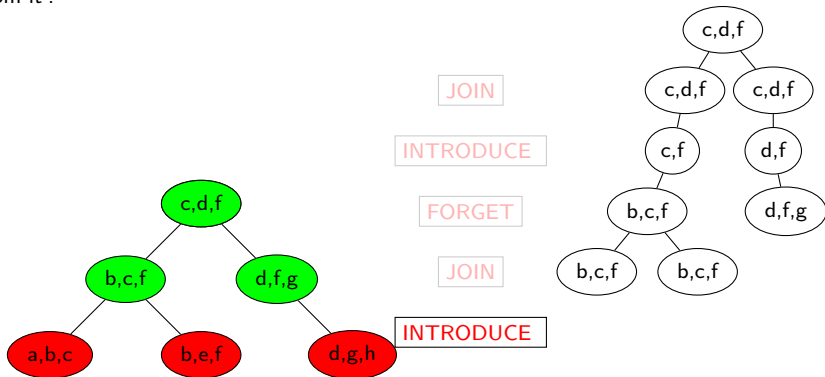
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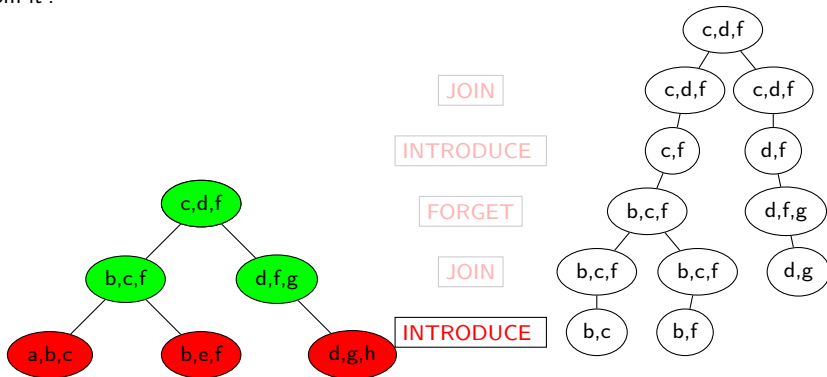
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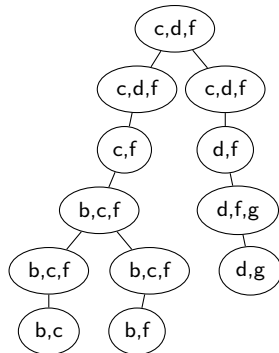
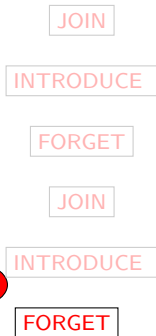
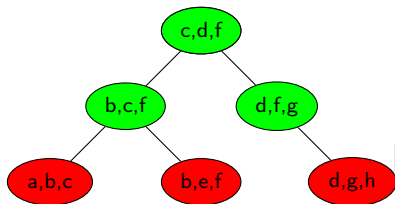
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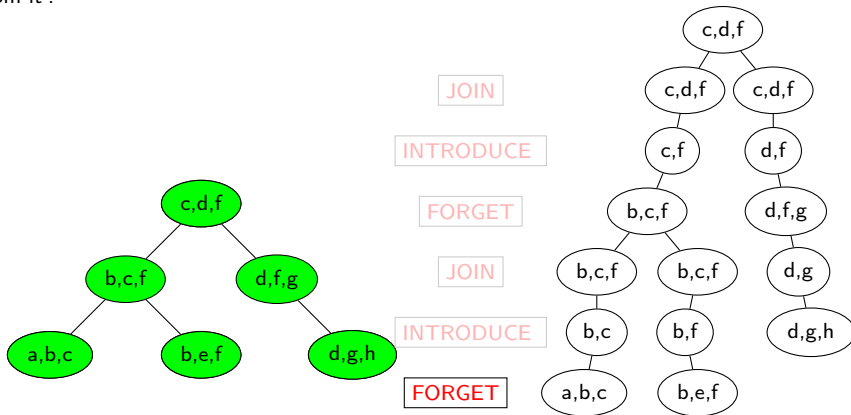
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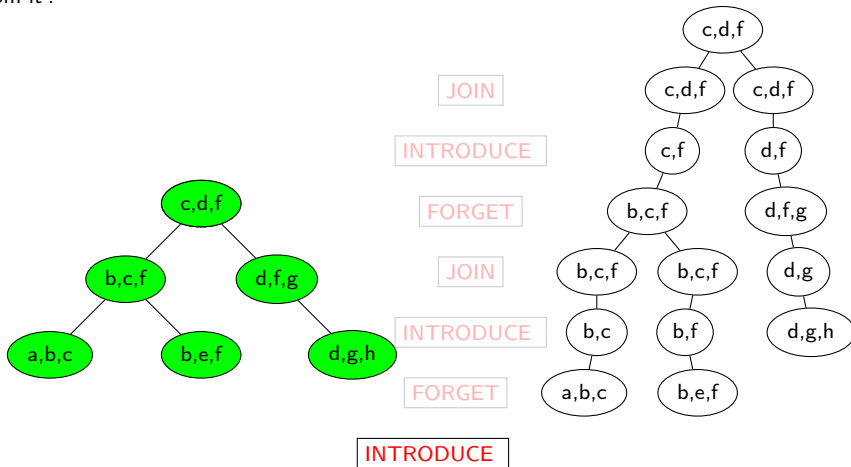
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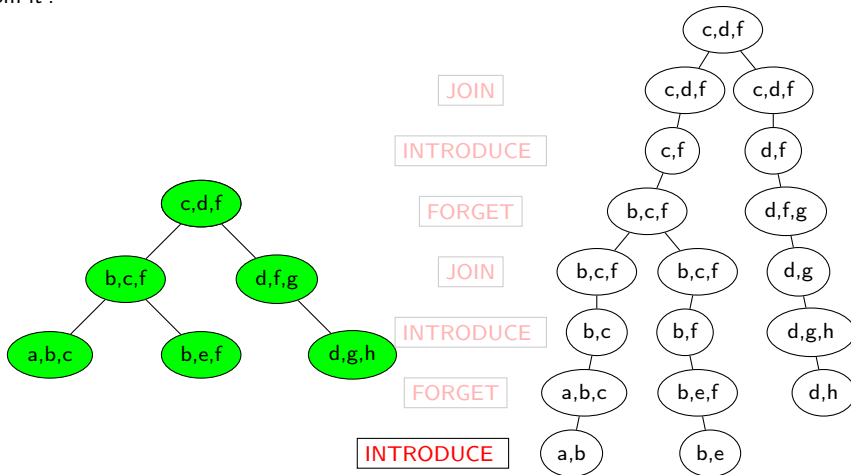
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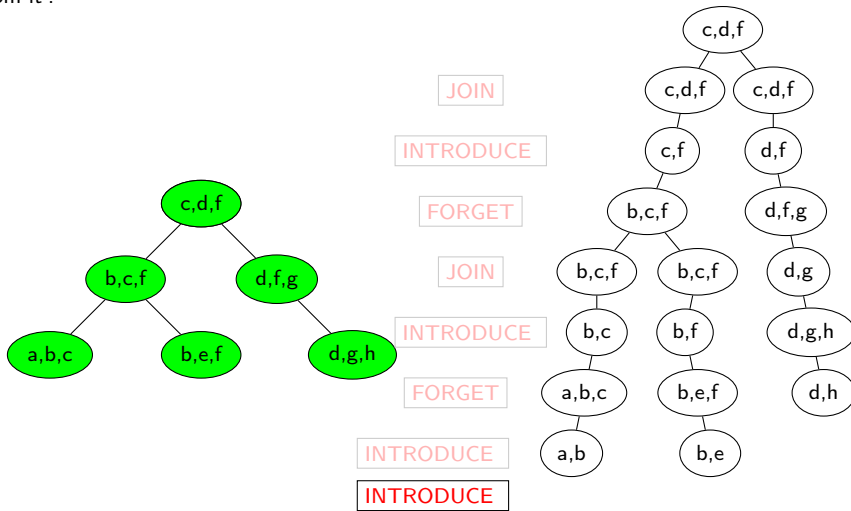
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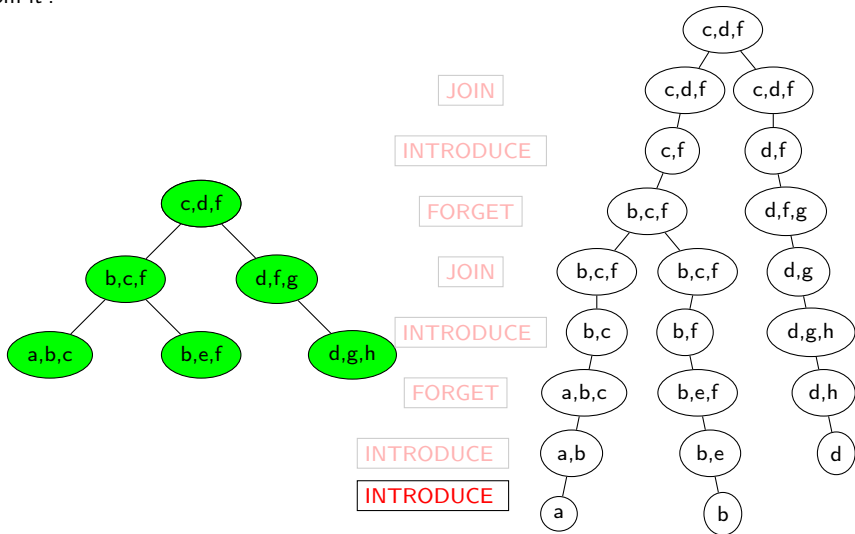
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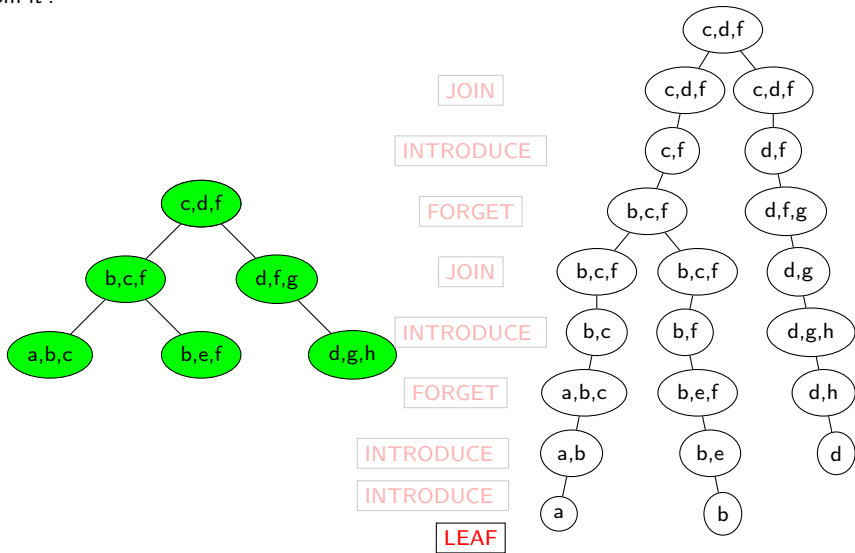
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Problem (k-color)

Let (G) be a graph and k an integer. We want to know if it is possible to draw each vertex of the graph so that two neighbors have never the same color and with only k colors.

This problem is a problem of decisions problem which is NP-complet.

Solution

tree-width : *k -color is possible for a graph (G) if and only if $k > \text{treewidth}(G)$.*

nice tree : *a nice tree of (G) gives a way to find a k -coloration of (G) (if $k > \text{treewidth}(G)$).*

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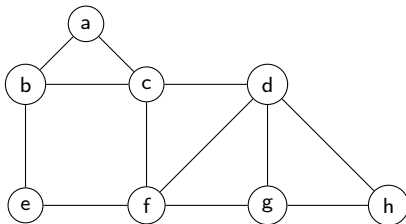
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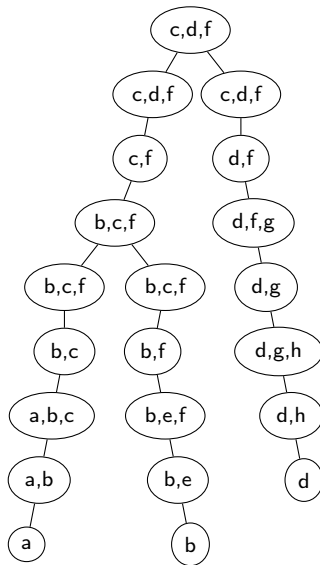
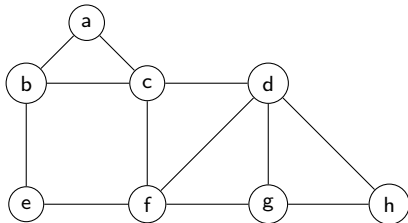
nice tree : a nice tree of (G) gives a way to find a k -coloration of (G) (if $k \geq \text{treewidth}(G)$).

Illustration : we will use the previously trees to solve the problem with this graph :



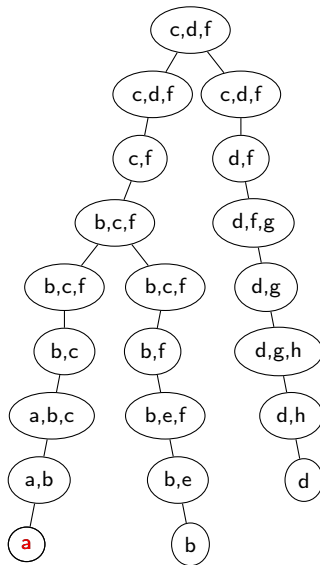
nice tree of the graph

- $\text{treewidth}(G)=2$: we can solve the problem with 3 colors
- we can fix a color for a “Leaf”
- when we meet an “Introduce” we add a color
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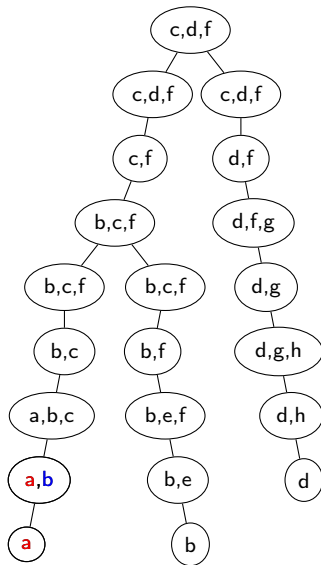
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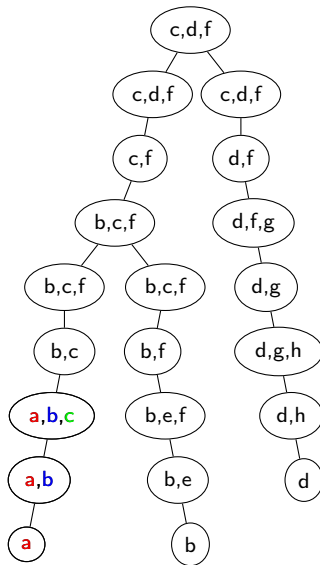
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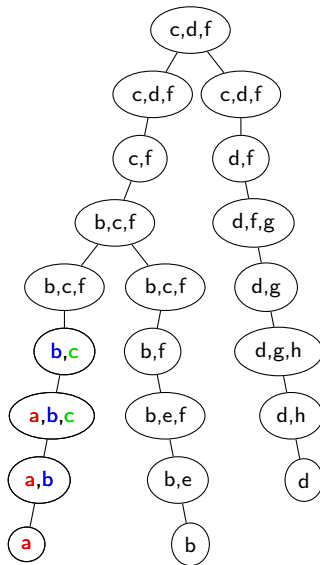
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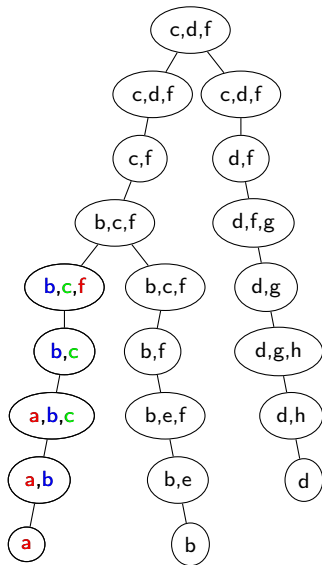
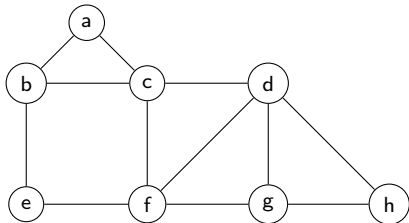
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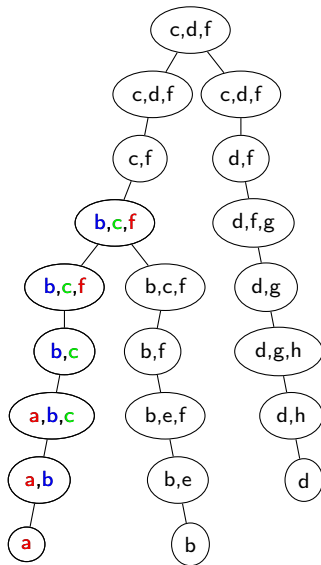
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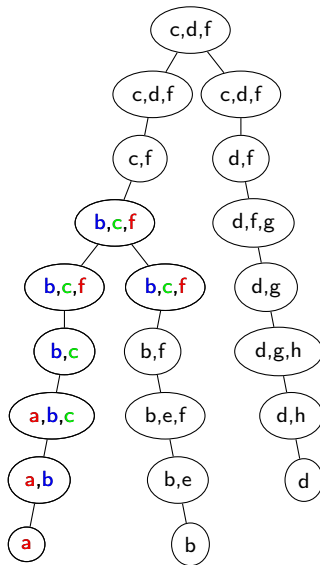
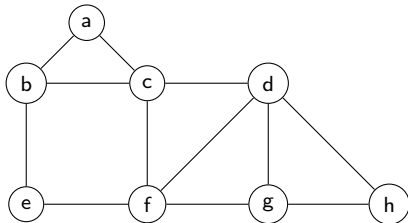
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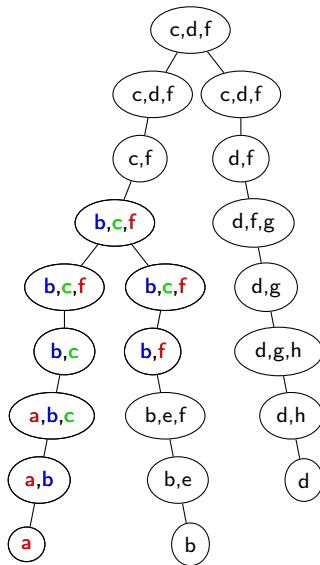
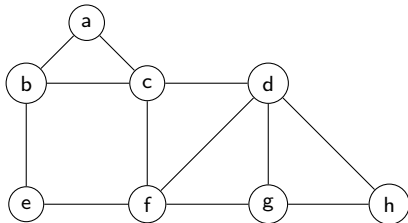
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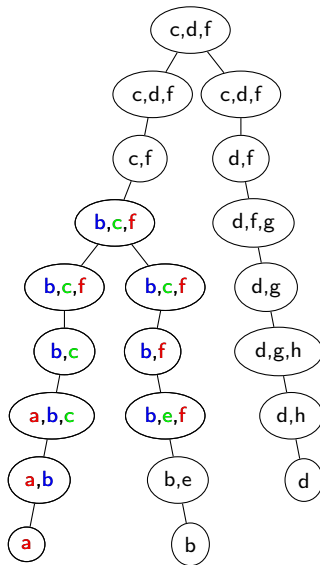
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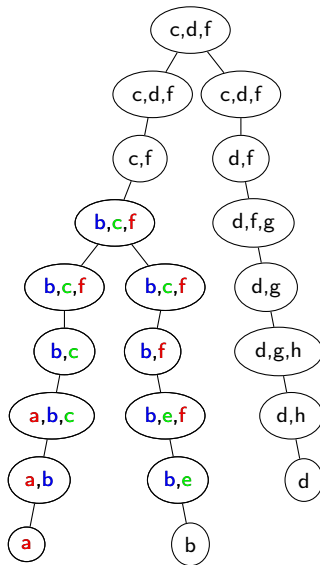
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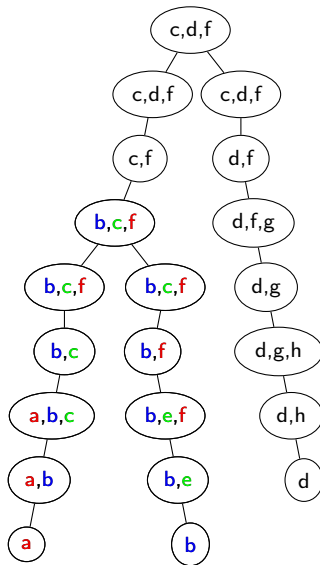
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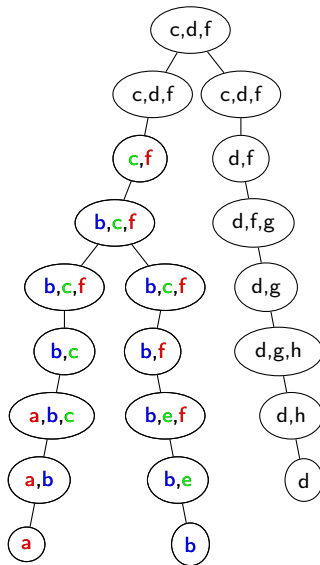
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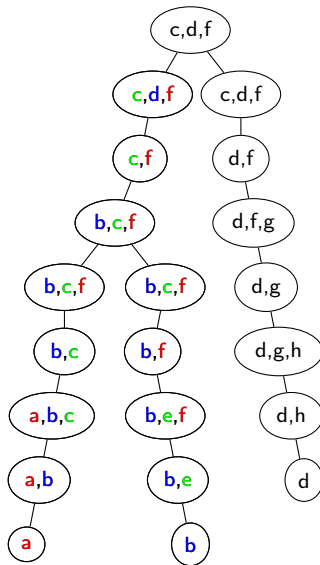
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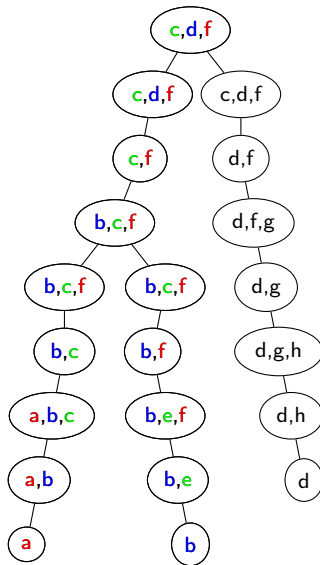
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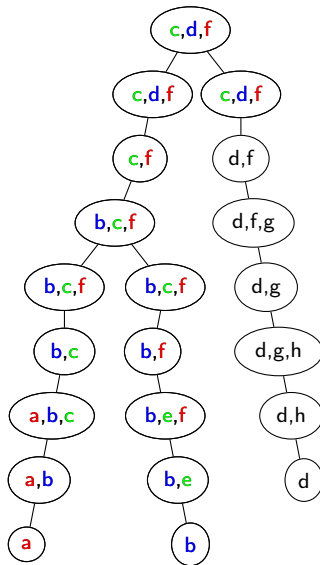
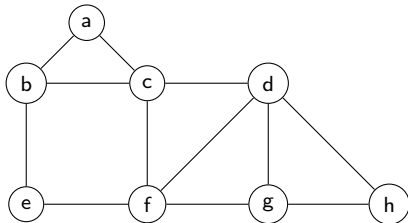
nice tree of the graph

- $\text{treewidth}(G)=2$: we can solve the problem with 3 colors
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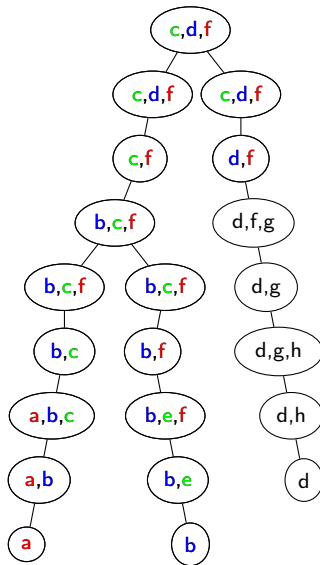
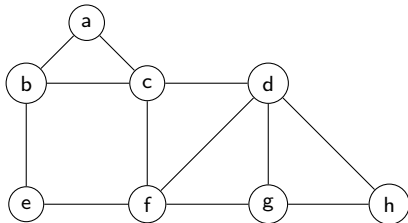
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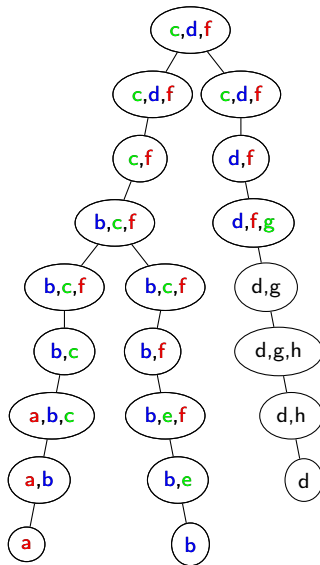
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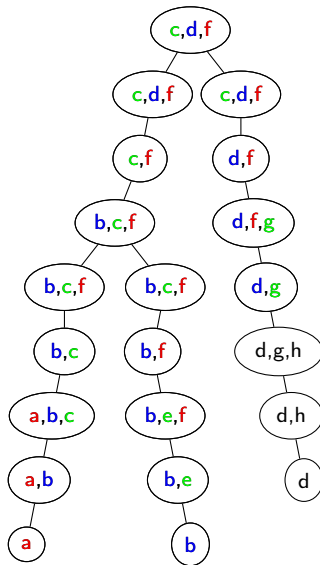
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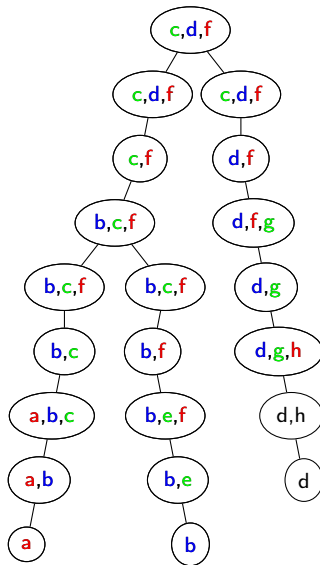
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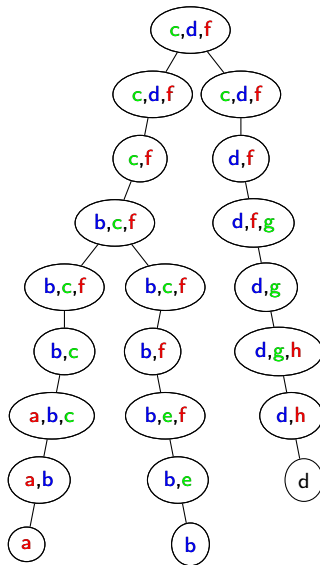
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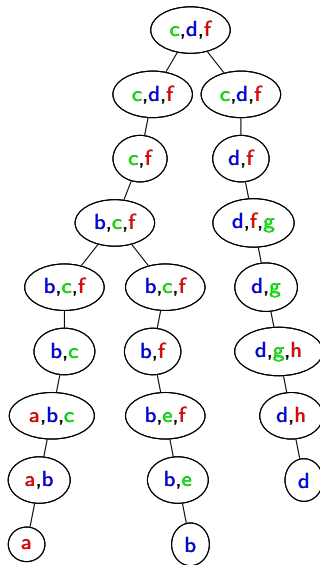
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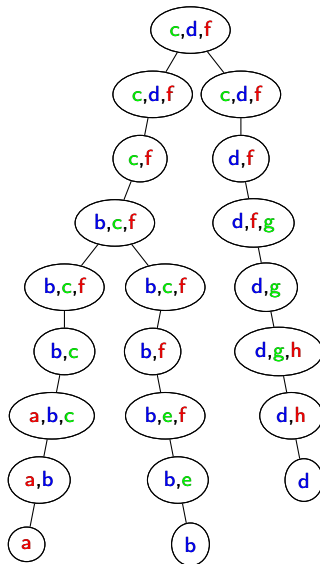
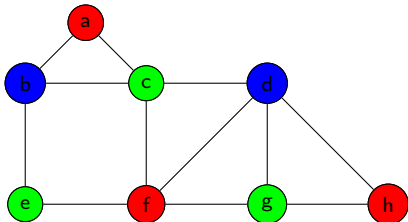
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1 Objectives

2 Tree decomposition of a graph

- definition
- treewidth
- nice tree
- Example

3 Application : k-color

- the problem
- illustration

4 Bibliography

- (1) Florent Madeleine : lesson for M2 Decim "Complexité des CSP et des requêtes"
- (2) Dániel Marx : Fixed Parameter Algorithms
- (3) wikipedia : articles of the graph section