HW7: Virtual Memory and Storage Scheduling

1) Given a 3 frame VMM and the page request sequence below (it's the top row), perform the VMM **First In First Out** page replacement policy (complete the table) to determine the total number of page faults for the sequence.

F	1	2	3	4	2	4	5	1	6	3	7	5	3	2	6	6	2	4	5	3	2	1	5	1	6	3	4	2	1
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1			2	3	3	3	T	5	ì	6	δ	7	7	5	へ	Z	J	6	4	5	3	2	1	1	5	б	3	I	2
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2) Given a 4 frame VMM and the page request sequence below (it's the top row), perform the VMM **Least Recently Used** page replacement policy (complete the table) to determine the total number of page faults for the sequence.

F	1	2	3	4	2	4	5	1	6	3	7	5	3	2	6	6	2	4	5	3	2	1	5	1	6	3	4	2	1
0		2	3	4	2	4	5)	6	3	7	6	Դ	2	کا	6	7	4	5	3	~	1	5	1	6	3	ን	2	ı
1		1	2	3	4	7	7	5		9	Z	く	5	ሪ	2	2	6	7	4	6	ን	2	1	ζ	1	6	3	4	2
2			1	2	3	3	7	4	5	1	9	8	フ	7	3	3	3	8	ィ	4	5	3	2	7	5	١	6	3	4
3				ı	1	1	κ	2	4	5	7	6	6	フ	5	5	5	3	6	૨	4	5	3	3	٦	5	ļ	8	3
Hov	+ ↑ How many page faults?) (+	+				+		+	+					

How many page faults?

3) Given a 3 frame VMM and the page request sequence below (it's the top row), perform the VMM **Optimal** page replacement policy (complete the table) to determine the total number of page faults for the sequence.

F	1	2	3	4	2	4	5	1	6	3	7	5	3	2	6	6	2	4	5	3	2	1	5	1	6	3	4	2	1
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1		1	<i>1</i> 2	21	24	Zą	2,	۷,	ک ^د	づ	7 _i	ノー	?	23	ر عر	プ	7	2	3 2	ン	کر	ř-	l ₁	7	14	13	12	1.	1
2			15	14	 	اء	I,	1,4	66	ივ	જુ	3,	3,	36	6,	6	6	49	42	يمي	35	3 ₄	33	34	3,	<u>က်</u>	3:	2;	2
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