

Slide #3.

- Let $r = 2$ and construct $\text{GF}(2^r)$ from $h(x) = x^2 + x + 1$. We have

$$\text{GF}(2^r) = \text{GF}(2^2) = \{0, 1, \beta, \beta^2\},$$

where β is a primitive element and $\beta^2 = \beta + 1$.

- Let $n = 3$. $\text{GF}(2^r)^n$ is a vector space over $\text{GF}(2^r)$. The elements of $\text{GF}(2^r)^n$ are listed below. There are $(2^r)^n = (2^2)^3 = 64$ of them.

$(000), (100), (\beta00), (\beta^200), (010), (110), (\beta10), (\beta^210),$

$(0\beta0), (1\beta0), (\beta\beta0), (\beta^2\beta0), (0\beta^20), (1\beta^20), (\beta\beta^20), (\beta^2\beta^20),$

$(001), (101), (\beta01), (\beta^201), (011), (111), (\beta11), (\beta^211),$

$(0\beta1), (1\beta1), (\beta\beta1), (\beta^2\beta1), (0\beta^21), (1\beta^21), (\beta\beta^21),$

$(\beta^2\beta^21), (00\beta), (10\beta), (\beta0\beta), (\beta^20\beta), (01\beta), (11\beta), (\beta1\beta),$

$(\beta^21\beta), (0\beta\beta), (1\beta\beta), (\beta\beta\beta), (\beta^2\beta\beta), (0\beta^2\beta), (1\beta^2\beta), (\beta\beta^2\beta),$

$(\beta^2\beta^2\beta), (00\beta^2), (10\beta^2), (\beta0\beta^2), (\beta^20\beta^2), (01\beta^2), (11\beta^2), (\beta1\beta^2),$

$(\beta^21\beta^2), (0\beta\beta^2), (1\beta\beta^2), (\beta\beta\beta^2), (\beta^2\beta\beta^2), (0\beta^2\beta^2), (1\beta^2\beta^2),$

$(\beta\beta^2\beta^2), (\beta^2\beta^2\beta^2).$

Slide #4. Let $r = 4$, $q = 2^r = 2^4$ and construct $\text{GF}(2^4)$ from $h(x) = x^4 + x + 1$ just as in Table 5.1, p. 114. Let, for instance,

$$\alpha_1 = 1, \alpha_2 = \beta^5, \alpha_3 = \beta^9.$$

Then

$$\begin{aligned} g(x) &= (x + \alpha_1) \cdot (x + \alpha_2) \cdot (x + \alpha_3) \\ &= (x + 1) \cdot (x + \beta^5) \cdot (x + \beta^9) \\ &= x^3 + \beta^{13}x^2 + \beta^8x + \beta^{14} \end{aligned}$$

generates a cyclic code of length $n = 2^r - 1 = 15$ over $\text{GF}(2^4)$. Note that the coefficients of $g(x)$ are not necessarily binary, that is, they do not necessarily belong to the binary field $K = \text{GF}(2) = \{0, 1\}$.

Slide #6. Derivation of a parity-check matrix for the Reed-Solomon code. Observe that

$$v = (v_0, v_1, v_2, \dots, v_{n-1}) \in RS(2^r, \delta)$$

or

$$v(x) = v_0 + v_1x + v_2x^2 + \dots + v_{n-1}x^{n-1} \in RS(2^r, \delta)$$

if and only if

$$v(\beta^{m+1}) = v(\beta^{m+2}) = \dots = v(\beta^{m+\delta-1}) = 0,$$

that is, if and only if

$$\begin{cases} v(\beta^{m+1}) = v_0 + v_1\beta^{m+1} + v_2(\beta^{m+1})^2 + \dots + v_{n-1}(\beta^{m+1})^{n-1} = 0 \\ v(\beta^{m+2}) = v_0 + v_1\beta^{m+2} + v_2(\beta^{m+2})^2 + \dots + v_{n-1}(\beta^{m+2})^{n-1} = 0 \\ \dots \\ v(\beta^{m+\delta-1}) = v_0 + v_1\beta^{m+\delta-1} + v_2(\beta^{m+\delta-1})^2 + \dots + v_{n-1}(\beta^{m+\delta-1})^{n-1} = 0. \end{cases}$$

The above system can be written as $(v_0, v_1, \dots, v_{n-1}) \cdot H = 0$ where

$$H = \begin{bmatrix} 1 & 1 & 1 \\ \beta^{m+1} & \beta^{m+2} & \beta^{m+\delta-1} \\ (\beta^{m+1})^2 & (\beta^{m+2})^2 & \dots & (\beta^{m+\delta-1})^2 \\ \vdots & \vdots & \vdots \\ (\beta^{m+1})^{n-1} & (\beta^{m+2})^{n-1} & \dots & (\beta^{m+\delta-1})^{n-1} \end{bmatrix}.$$

Slide #7. Recall:

$$\det \begin{bmatrix} ka_{11} & ka_{12} & \cdots & ka_{1n} \\ a_{21} & a_{22} & \cdots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & \cdots & a_{nn} \end{bmatrix} = k \cdot \det \begin{bmatrix} a_{11} & a_{12} & \cdots & a_{1n} \\ a_{21} & a_{22} & \cdots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & \cdots & a_{nn} \end{bmatrix}.$$

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Slide #7. Let x_1, x_2, \dots, x_n be any elements of a field (finite or not). The *Vandermonde* determinant of order n is defined as

$$V := \begin{bmatrix} 1 & x_1 & x_1^2 & \cdots & x_1^{n-1} \\ 1 & x_2 & x_2^2 & \cdots & x_2^{n-1} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & x_n & x_n^2 & \cdots & x_n^{n-1} \end{bmatrix}.$$

It is possible to show that

$$\det V = \prod_{1 \leq i < j \leq n} (x_j - x_i).$$

Property: If x_1, x_2, \dots, x_n are all distinct, then $V \neq 0$.