

**Final**  
**Algebraic Coding Theory**  
**Math 525**  
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**Problem 1:**

- (a) Find a parity-check matrix  $H$  for a cyclic Hamming code of length 15 using  $\text{GF}(2^4)$  constructed from  $1 + x + x^4$  (see Table 5.1, p. 114), where the generator polynomial is  $m_7(x)$ . Each entry of  $H$  must be expressed as a power of  $\beta$ , where  $\beta$  is the primitive element of the field, exactly as in Table 5.1.

Notice the generator polynomial:  $m_7(x)$

$$[\beta^7] \Rightarrow [\beta^{14}] \Rightarrow [\beta^{28} = \beta^{13}] \Rightarrow [\beta^{26} = \beta^{11}] \Rightarrow [\beta^{22} = \beta^7]$$

$$\begin{aligned} m_7(x) &= (x + \beta^7)(x + \beta^{14})(x + \beta^{13})(x + \beta^{11}) \\ &= \left(x^2 + (\beta^7 + \beta^{14})x + \beta^{21}\right) \left(x^2 + (\beta^{13} + \beta^{11})x + \beta^{24}\right) \\ &= \left(x^2 + \beta x + \beta^6\right) \left(x^2 + \beta^4 x + \beta^9\right) \\ &= x^4 + \left(\beta + \beta^4\right)x^3 + \left(\beta^9 + \beta^5 + \beta^6\right)x^2 + \left(\beta^{10} + \beta^{10}\right)x + \beta^{15} \\ &= x^4 + x^3 + 1 \end{aligned}$$

We get the Parity Check Matrix below:

$$H = \begin{bmatrix} 1 \\ \beta^7 \\ \beta^{14} \end{bmatrix}$$

- (b) Now suppose  $r$  is received, where  $r$  is the word of length 15 that you obtained from your last name. Find the most likely codeword transmitted.

Notice my last name is GIANG, so we get

$$r(x) = \{00100\ 10010\ 01010\}(ANG) = x^2 + x^5 + x^8 + x^{11} + x^{13}$$

Notice the following:

$$r(\beta) = \beta^2 + \beta^5 + \beta^8 + \beta^{11} + \beta^{13} = \beta^2$$

Thus we get the most likely codeword is  $\{01100\ 10010\ 01010\}$  (ING).