

 Lec_35.md

Lecture 35

Dec. 08/2020

Announcements

Final Exam

- Summative Assessment
- 35% of ECE244 Mark
- December 14, 2020
 - Starting at 9:30am ET
 - 2.5 hour duration
 - Synchronous Exam
 - Everyone takes exam at same time, regardless of time zone
- Open textbook, open notes
- No calculator, No IDE's
- No more than 50% MC type questions

Optional Lab

- 12 students completed and competed in the lab

Teaching Evaluations

Complexity Analysis

We can measure the complexity of an algorithm in time versus the input size of the problem

- Group the complexity of algorithm's into **complexity classes**

Complexity Classes

Complexity Classes:

1. $O(1)$
 - This is rare
 - Same runtime regardless of how large the size of the input is
2. $O(\log_n)$
 - Great algorithms
3. $O(n)$
 - Linear algorithms
4. $O(n \log n)$

- Somewhere between linear and n -squared

5. $O(n^2)$

- Polynomial growth

6. $O(n^3)$

7. $O(2^n)$

8. $O(n!)$

- The dreaded "factorial time"

9. $O(n^n)$

Recursive Functions

What would the **time complexity** be for a recursive function?

For example,

```
int recursive(int n){
    if(n<1) return 1;
    else return (recursive(n/2));
}
```

Notes:

1. `if(n<1) return 1;`
2. $T(n)$ = **recurrence equation**
 - C , if $n < 1$
 - $C + T(n/2)$, if $n > 1$

$T(n)$
$T(n) = C + T(n/2)$
$T(n) = C + T(n/4)$
$T(n) = C + C + T(n/8)$
$T(n) = C + C + C + T(n/16)$
$T(n) = C + C + C + \dots + T(n/n)$

Notice that this is proportional to $\log_2(n)$

- Require n divisions to reach $T(n/n)$

So `int recursive()` has **time complexity** $O(\log n)$

```
int recursive(int n){
    if(n<=1) return 1;
    else return (recursive(n-1)+recursive(n-1));
}
```

Notes:

1. `if(n<=1) return 1;`
- This is a step
2. $T(n)$ = **recurrence equation**
- C , if $n \leq 1$
- $C+2T(n-1)$, if $n > 1$

$T(n)$	
$T(n)=C+2T(n-1)$	$T(n)=C+2T(n-1)$
$T(n)=C+2(C+2T(n-2))$	$T(n)=3C+4T(n-2)$
$T(n)=C+2(C+2(C+2T(n-3)))$	$T(n)=7C+8T(n-3)$
$T(n)=\dots$	$T(n)=15C+16T(n-4)$
$T(n)=xC+yT(n-(n-1))$	$T(n)=(2^{n-1}-1)C+2^{n-1}T(1)$

Notice that in $T(n)=(2^{n-1}-1)C+2^{n-1}T(1)$, $T(1)=C$

- So the time complexity of the algorithm is $T(n)$ is proportional to $2^{n-1}C$
 - So the order of this algorithm is $O(2^n)$

```
int recursive(int n){
    if(n<=1) return 1;
    else return (recursive(n/2)+recursive(n/2));
}
```

Notes:

1. $T(n)$ =
- C , if $n \leq 1$
- $C+2T(n/2)$, if $n > 1$

$T(n)$
$T(n)=C+2T(n/2)$
$T(n)=3C+4T(n/4)$
$T(n)=7C+8T(n/8)$
$T(n)=15C+16T(n/16)$
$T(n)=(n-1)C+nT(n/n)$

Quicksort

```
int QuickSort(int left,int right,int* a){
    int p = SelectAndShuffle(left, right, a);
    if(p>left) QuickSort(left,p-1,a);
    if(p<right)QuickSort(p+1,right,a);
}
```

Notes:

1. SelectAndShuffle
- **time complexity:** $O(n)$
2. `QuickSort(left,p-1),a);` and `QuickSort(p+1,right,a)`
- $T(n) = n + T(p) + T(n-p)$
 - The **time complexity** of this *depends* on input data
 - Location of pivot with respect to data in `a`
 - Split into best, worst, and average case