

Vim

| | |
|-----------------------------------|------------------------------------------------------------|
| <code>O ^ \$</code> | move to front of line, first char, end |
| <code>h j k l</code> | move left, down, up, right |
| <code>w W</code> | move forward to start of words w/wo punc |
| <code>e E</code> | move forward to end of words w/wo punc |
| <code>b B</code> | move backward by word w/wo punc |
| <code>} {</code> | move forward/backward by paragraph |
| <code>%</code> | move to matching (), {}, [] |
| <code>gg G</code> | move to top/bottom of file |
| <code>: [num] [enter]</code> | move to line number |
| <code>/ [regex]</code> | search forward for regex and move to first match |
| <code>? [regex]</code> | search backward for regex and move to last match |
| <code>n N</code> | repeat prior search in same / opposite direction |
| | |
| <code>i I</code> | insert at cursor, beginning of line |
| <code>a A</code> | append after cursor, end of line |
| <code>o O</code> | open blank line below, above current line |
| | |
| <code>c [move] cc</code> | change until move point, entire line |
| <code>d [move] dd</code> | delete until move point, entire line |
| <code>y [move] yy</code> | yank (copy) until move point, entire line |
| <code>p P</code> | put (paste) yanked content after / before cursor |
| <code>x X</code> | delete current / previous character |
| <code>r R</code> | replace a single character, multiple characters |
| <code>:s/ [old]/ [new]/g</code> | replace text on this line; g for all occurrences |
| <code>:%s/ [old]/ [new]/g</code> | replace text in file; g for all occurrences |
| | |
| <code>= [move] ==</code> | auto-indent lines |
| <code>< [move] <<</code> | unindent one step |
| <code>> [move] >></code> | indent one step |
| <code>:! [cmd]</code> | pipe file through shell command, e.g., <code>:!sort</code> |
| | |
| <code>:w :wq ZZ</code> | write, write and quit, write and quit |
| <code>:q :q!</code> | quit, quit even if unsaved changes |
| <code>:e :ls :buf [num]</code> | edit another file, list buffers, switch to buffer |
| <code>:vsplit :hsplit</code> | split window |
| <code>Ctrl-w [h j k l]</code> | switch window |
| <code>u Ctrl+r .</code> | undo, redo, repeat last action |

Nearly all commands support a prefix number, e.g., `5dd` delete 5 lines.

Beating the Averages (Paul Graham)

In the summer of 1995, my friend Robert Morris and I started a startup called Viaweb. Our plan was to write software that would let end users build online stores. What was novel about this software, at the time, was that it ran on our server, using ordinary Web pages as the interface.

A lot of people could have been having this idea at the same time, of course, but as far as I know, Viaweb was the first Web-based application. It seemed such a novel idea to us that we named the company after it: Viaweb, because our software worked via the Web, instead of running on your desktop computer.

Another unusual thing about this software was that it was written primarily in a programming language called Lisp. It was one of the first big end-user applications to be written in Lisp, which up till then had been used mostly in universities and research labs. [1]

The Secret Weapon

Eric Raymond has written an essay called "How to Become a Hacker," and in it, among other things, he tells would-be hackers what languages they should learn. He suggests starting with Python and Java, because they are easy to learn. The serious hacker will also want to learn C, in order to hack Unix, and Perl for system administration and cgi scripts. Finally, the truly serious hacker should consider learning Lisp:

Lisp is worth learning for the profound enlightenment experience you will have when you finally get it; that experience will make you a better programmer for the rest of your days, even if you never actually use Lisp itself a lot.

This is the same argument you tend to hear for learning Latin. It won't get you a job, except perhaps as a classics pro-

fessor, but it will improve your mind, and make you a better writer in languages you do want to use, like English.

But wait a minute. This metaphor doesn't stretch that far. The reason Latin won't get you a job is that no one speaks it. If you write in Latin, no one can understand you. But Lisp is a computer language, and computers speak whatever language you, the programmer, tell them to.

So if Lisp makes you a better programmer, like he says, why wouldn't you want to use it? If a painter were offered a brush that would make him a better painter, it seems to me that he would want to use it in all his paintings, wouldn't he? I'm not trying to make fun of Eric Raymond here. On the whole, his advice is good. What he says about Lisp is pretty much the conventional wisdom. But there is a contradiction in the conventional wisdom: Lisp will make you a better programmer, and yet you won't use it.

Why not? Programming languages are just tools, after all. If Lisp really does yield better programs, you should use it. And if it doesn't, then who needs it?

This is not just a theoretical question. Software is a very competitive business, prone to natural monopolies. A company that gets software written faster and better will, all other things being equal, put its competitors out of business. And when you're starting a startup, you feel this very keenly. Startups tend to be an all or nothing proposition. You either get rich, or you get nothing. In a startup, if you bet on the wrong technology, your competitors will crush you.

Robert and I both knew Lisp well, and we couldn't see any reason not to trust our instincts and go with Lisp. We knew that everyone else was writing their software in C++ or Perl. But we also knew that that didn't mean anything. If you chose technology that way, you'd be running Windows. When you choose technology, you have to ignore what other people are doing, and consider only what will work the best.

This is especially true in a startup. In a big company, you can do what all the other big companies are doing. But a startup can't do what all the other startups do. I don't think a lot of people realize this, even in startups.

The average big company grows at about ten percent a year. So if you're running a big company and you do everything the way the average big company does it, you can expect to do as well as the average big company— that is, to grow about ten percent a year.

The same thing will happen if you're running a startup, of course. If you do everything the way the average startup does it, you should expect average performance. The problem here is, average performance means that you'll go out of business. The survival rate for startups is way less than fifty percent. So if you're running a startup, you had better be doing something odd. If not, you're in trouble.

Back in 1995, we knew something that I don't think our competitors understood, and few understand even now: when you're writing software that only has to run on your own servers, you can use any language you want. When you're writing desktop software, there's a strong bias toward writing applications in the same language as the operating system. Ten years ago, writing applications meant writing applications in C. But with Web-based software, especially when you have the source code of both the language and the operating system, you can use whatever language you want.

This new freedom is a double-edged sword, however. Now that you can use any language, you have to think about which one to use. Companies that try to pretend nothing has changed risk finding that their competitors do not.

If you can use any language, which do you use? We chose Lisp. For one thing, it was obvious that rapid development would be important in this market. We were all starting from scratch, so a company that could get new features done before its competitors would have a big advantage. We knew Lisp was a really good language for writing software quickly, and server-based applications magnify the effect of rapid development, because you can release software the minute it's done.

If other companies didn't want to use Lisp, so much the better. It might give us a technological edge, and we needed all the help we could get. When we started Viaweb, we had no experience in business. We didn't know anything about marketing, or

hiring people, or raising money, or getting customers. Neither of us had ever even had what you would call a real job. The only thing we were good at was writing software. We hoped that would save us. Any advantage we could get in the software department, we would take.

So you could say that using Lisp was an experiment. Our hypothesis was that if we wrote our software in Lisp, we'd be able to get features done faster than our competitors, and also to do things in our software that they couldn't do. And because Lisp was so high-level, we wouldn't need a big development team, so our costs would be lower. If this were so, we could offer a better product for less money, and still make a profit. We would end up getting all the users, and our competitors would get none, and eventually go out of business. That was what we hoped would happen, anyway.

What were the results of this experiment? Somewhat surprisingly, it worked. We eventually had many competitors, on the order of twenty to thirty of them, but none of their software could compete with ours. We had a wysiwyg online store builder that ran on the server and yet felt like a desktop application. Our competitors had cgi scripts. And we were always far ahead of them in features. Sometimes, in desperation, competitors would try to introduce features that we didn't have. But with Lisp our development cycle was so fast that we could sometimes duplicate a new feature within a day or two of a competitor announcing it in a press release. By the time journalists covering the press release got round to calling us, we would have the new feature too.

It must have seemed to our competitors that we had some kind of secret weapon— that we were decoding their Enigma traffic or something. In fact we did have a secret weapon, but it was simpler than they realized. No one was leaking news of their features to us. We were just able to develop software faster than anyone thought possible.

When I was about nine I happened to get hold of a copy of *The Day of the Jackal*, by Frederick Forsyth. The main character is an assassin who is hired to kill the president of France. The assassin has to get past the police to get up to an apartment

that overlooks the president's route. He walks right by them, dressed up as an old man on crutches, and they never suspect him.

Our secret weapon was similar. We wrote our software in a weird AI language, with a bizarre syntax full of parentheses. For years it had annoyed me to hear Lisp described that way. But now it worked to our advantage. In business, there is nothing more valuable than a technical advantage your competitors don't understand. In business, as in war, surprise is worth as much as force.

And so, I'm a little embarrassed to say, I never said anything publicly about Lisp while we were working on Viaweb. We never mentioned it to the press, and if you searched for Lisp on our Web site, all you'd find were the titles of two books in my bio. This was no accident. A startup should give its competitors as little information as possible. If they didn't know what language our software was written in, or didn't care, I wanted to keep it that way. [2]

The people who understood our technology best were the customers. They didn't care what language Viaweb was written in either, but they noticed that it worked really well. It let them build great looking online stores literally in minutes. And so, by word of mouth mostly, we got more and more users. By the end of 1996 we had about 70 stores online. At the end of 1997 we had 500. Six months later, when Yahoo bought us, we had 1070 users. Today, as Yahoo Store, this software continues to dominate its market. It's one of the more profitable pieces of Yahoo, and the stores built with it are the foundation of Yahoo Shopping. I left Yahoo in 1999, so I don't know exactly how many users they have now, but the last I heard there were about 20,000.

The Blub Paradox

What's so great about Lisp? And if Lisp is so great, why doesn't everyone use it? These sound like rhetorical questions, but actually they have straightforward answers. Lisp is so great not because of some magic quality visible only to devotees, but be-

cause it is simply the most powerful language available. And the reason everyone doesn't use it is that programming languages are not merely technologies, but habits of mind as well, and nothing changes slower. Of course, both these answers need explaining.

I'll begin with a shockingly controversial statement: programming languages vary in power.

Few would dispute, at least, that high level languages are more powerful than machine language. Most programmers today would agree that you do not, ordinarily, want to program in machine language. Instead, you should program in a high-level language, and have a compiler translate it into machine language for you. This idea is even built into the hardware now: since the 1980s, instruction sets have been designed for compilers rather than human programmers.

Everyone knows it's a mistake to write your whole program by hand in machine language. What's less often understood is that there is a more general principle here: that if you have a choice of several languages, it is, all other things being equal, a mistake to program in anything but the most powerful one. [3]

There are many exceptions to this rule. If you're writing a program that has to work very closely with a program written in a certain language, it might be a good idea to write the new program in the same language. If you're writing a program that only has to do something very simple, like number crunching or bit manipulation, you may as well use a less abstract language, especially since it may be slightly faster. And if you're writing a short, throwaway program, you may be better off just using whatever language has the best library functions for the task. But in general, for application software, you want to be using the most powerful (reasonably efficient) language you can get, and using anything else is a mistake, of exactly the same kind, though possibly in a lesser degree, as programming in machine language.

You can see that machine language is very low level. But, at least as a kind of social convention, high-level languages are often all treated as equivalent. They're not. Technically the term "high-level language" doesn't mean anything very definite.

There's no dividing line with machine languages on one side and all the high-level languages on the other. Languages fall along a continuum [4] of abstractness, from the most powerful all the way down to machine languages, which themselves vary in power.

Consider Cobol. Cobol is a high-level language, in the sense that it gets compiled into machine language. Would anyone seriously argue that Cobol is equivalent in power to, say, Python? It's probably closer to machine language than Python.

Or how about Perl 4? Between Perl 4 and Perl 5, lexical closures got added to the language. Most Perl hackers would agree that Perl 5 is more powerful than Perl 4. But once you've admitted that, you've admitted that one high level language can be more powerful than another. And it follows inexorably that, except in special cases, you ought to use the most powerful you can get.

This idea is rarely followed to its conclusion, though. After a certain age, programmers rarely switch languages voluntarily. Whatever language people happen to be used to, they tend to consider just good enough.

Programmers get very attached to their favorite languages, and I don't want to hurt anyone's feelings, so to explain this point I'm going to use a hypothetical language called Blub. Blub falls right in the middle of the abstractness continuum. It is not the most powerful language, but it is more powerful than Cobol or machine language.

And in fact, our hypothetical Blub programmer wouldn't use either of them. Of course he wouldn't program in machine language. That's what compilers are for. And as for Cobol, he doesn't know how anyone can get anything done with it. It doesn't even have `x` (Blub feature of your choice).

As long as our hypothetical Blub programmer is looking down the power continuum, he knows he's looking down. Languages less powerful than Blub are obviously less powerful, because they're missing some feature he's used to. But when our hypothetical Blub programmer looks in the other direction, up the power continuum, he doesn't realize he's looking up. What he sees are merely weird languages. He probably considers

them about equivalent in power to Blub, but with all this other hairy stuff thrown in as well. Blub is good enough for him, because he thinks in Blub.

When we switch to the point of view of a programmer using any of the languages higher up the power continuum, however, we find that he in turn looks down upon Blub. How can you get anything done in Blub? It doesn't even have y.

By induction, the only programmers in a position to see all the differences in power between the various languages are those who understand the most powerful one. (This is probably what Eric Raymond meant about Lisp making you a better programmer.) You can't trust the opinions of the others, because of the Blub paradox: they're satisfied with whatever language they happen to use, because it dictates the way they think about programs.

I know this from my own experience, as a high school kid writing programs in Basic. That language didn't even support recursion. It's hard to imagine writing programs without using recursion, but I didn't miss it at the time. I thought in Basic. And I was a whiz at it. Master of all I surveyed.

The five languages that Eric Raymond recommends to hackers fall at various points on the power continuum. Where they fall relative to one another is a sensitive topic. What I will say is that I think Lisp is at the top. And to support this claim I'll tell you about one of the things I find missing when I look at the other four languages. How can you get anything done in them, I think, without macros? [5]

Many languages have something called a macro. But Lisp macros are unique. And believe it or not, what they do is related to the parentheses. The designers of Lisp didn't put all those parentheses in the language just to be different. To the Blub programmer, Lisp code looks weird. But those parentheses are there for a reason. They are the outward evidence of a fundamental difference between Lisp and other languages.

Lisp code is made out of Lisp data objects. And not in the trivial sense that the source files contain characters, and strings are one of the data types supported by the language. Lisp code, after it's read by the parser, is made of data structures that you

can traverse.

If you understand how compilers work, what's really going on is not so much that Lisp has a strange syntax as that Lisp has no syntax. You write programs in the parse trees that get generated within the compiler when other languages are parsed. But these parse trees are fully accessible to your programs. You can write programs that manipulate them. In Lisp, these programs are called macros. They are programs that write programs.

Programs that write programs? When would you ever want to do that? Not very often, if you think in Cobol. All the time, if you think in Lisp. It would be convenient here if I could give an example of a powerful macro, and say there! how about that? But if I did, it would just look like gibberish to someone who didn't know Lisp; there isn't room here to explain everything you'd need to know to understand what it meant. In Ansi Common Lisp I tried to move things along as fast as I could, and even so I didn't get to macros until page 160.

But I think I can give a kind of argument that might be convincing. The source code of the Viaweb editor was probably about 20-25% macros. Macros are harder to write than ordinary Lisp functions, and it's considered to be bad style to use them when they're not necessary. So every macro in that code is there because it has to be. What that means is that at least 20-25% of the code in this program is doing things that you can't easily do in any other language. However skeptical the Blub programmer might be about my claims for the mysterious powers of Lisp, this ought to make him curious. We weren't writing this code for our own amusement. We were a tiny startup, programming as hard as we could in order to put technical barriers between us and our competitors.

A suspicious person might begin to wonder if there was some correlation here. A big chunk of our code was doing things that are very hard to do in other languages. The resulting software did things our competitors' software couldn't do. Maybe there was some kind of connection. I encourage you to follow that thread. There may be more to that old man hobbling along on his crutches than meets the eye.

Aikido for Startups

But I don't expect to convince anyone (over 25) to go out and learn Lisp. The purpose of this article is not to change anyone's mind, but to reassure people already interested in using Lisp—people who know that Lisp is a powerful language, but worry because it isn't widely used. In a competitive situation, that's an advantage. Lisp's power is multiplied by the fact that your competitors don't get it.

If you think of using Lisp in a startup, you shouldn't worry that it isn't widely understood. You should hope that it stays that way. And it's likely to. It's the nature of programming languages to make most people satisfied with whatever they currently use. Computer hardware changes so much faster than personal habits that programming practice is usually ten to twenty years behind the processor. At places like MIT they were writing programs in high-level languages in the early 1960s, but many companies continued to write code in machine language well into the 1980s. I bet a lot of people continued to write machine language until the processor, like a bartender eager to close up and go home, finally kicked them out by switching to a risc instruction set.

Ordinarily technology changes fast. But programming languages are different: programming languages are not just technology, but what programmers think in. They're half technology and half religion. [6] And so the median language, meaning whatever language the median programmer uses, moves as slow as an iceberg. Garbage collection, introduced by Lisp in about 1960, is now widely considered to be a good thing. Runtime typing, ditto, is growing in popularity. Lexical closures, introduced by Lisp in the early 1970s, are now, just barely, on the radar screen. Macros, introduced by Lisp in the mid 1960s, are still *terra incognita*.

Obviously, the median language has enormous momentum. I'm not proposing that you can fight this powerful force. What I'm proposing is exactly the opposite: that, like a practitioner of Aikido, you can use it against your opponents.

If you work for a big company, this may not be easy. You

will have a hard time convincing the pointy-haired boss to let you build things in Lisp, when he has just read in the paper that some other language is poised, like Ada was twenty years ago, to take over the world. But if you work for a startup that doesn't have pointy-haired bosses yet, you can, like we did, turn the Blub paradox to your advantage: you can use technology that your competitors, glued immovably to the median language, will never be able to match.

If you ever do find yourself working for a startup, here's a handy tip for evaluating competitors. Read their job listings. Everything else on their site may be stock photos or the prose equivalent, but the job listings have to be specific about what they want, or they'll get the wrong candidates.

During the years we worked on Viaweb I read a lot of job descriptions. A new competitor seemed to emerge out of the woodwork every month or so. The first thing I would do, after checking to see if they had a live online demo, was look at their job listings. After a couple years of this I could tell which companies to worry about and which not to. The more of an IT flavor the job descriptions had, the less dangerous the company was. The safest kind were the ones that wanted Oracle experience. You never had to worry about those. You were also safe if they said they wanted C++ or Java developers. If they wanted Perl or Python programmers, that would be a bit frightening— that's starting to sound like a company where the technical side, at least, is run by real hackers. If I had ever seen a job posting looking for Lisp hackers, I would have been really worried.

Notes

[1] Viaweb at first had two parts: the editor, written in Lisp, which people used to build their sites, and the ordering system, written in C, which handled orders. The first version was mostly Lisp, because the ordering system was small. Later we added two more modules, an image generator written in C, and a back-office manager written mostly in Perl.

In January 2003, Yahoo released a new version of the editor written in C++ and Perl. It's hard to say whether the program is no longer written in Lisp, though, because to translate this program into C++ they literally had to write a Lisp interpreter: the source files of all the page-generating templates are still, as far as I know, Lisp code. (See Greenspun's Tenth Rule.)

[2] Robert Morris says that I didn't need to be secretive, because even if our competitors had known we were using Lisp, they wouldn't have understood why: "If they were that smart they'd already be programming in Lisp."

[3] All languages are equally powerful in the sense of being Turing equivalent, but that's not the sense of the word programmers care about. (No one wants to program a Turing machine.) The kind of power programmers care about may not be formally definable, but one way to explain it would be to say that it refers to features you could only get in the less powerful language by writing an interpreter for the more powerful language in it. If language A has an operator for removing spaces from strings and language B doesn't, that probably doesn't make A more powerful, because you can probably write a subroutine to do it in B. But if A supports, say, recursion, and B doesn't, that's not likely to be something you can fix by writing library functions.

[4] Note to nerds: or possibly a lattice, narrowing toward the top; it's not the shape that matters here but the idea that there is at least a partial order.

[5] It is a bit misleading to treat macros as a separate feature. In practice their usefulness is greatly enhanced by other Lisp features like lexical closures and rest parameters.

[6] As a result, comparisons of programming languages either take the form of religious wars or undergraduate textbooks so determinedly neutral that they're really works of anthropology. People who value their peace, or want tenure, avoid the topic. But the question is only half a religious one; there is something there worth studying, especially if you want to design new languages.

Linux base binaries

| | |
|--------------|--------------------------------------------------------------|
| apropos | search the manual page names and descriptions |
| autopoint | copies standard gettext infrastructure |
| b2sum | compute and check BLAKE2 message digest |
| base32 | base32 encode/decode data and print to standard output |
| base64 | base64 encode/decode data and print to standard output |
| basename | strip directory and suffix from filenames |
| bash | GNU Bourne-Again SHell |
| bashbug | report a bug in bash |
| bunzip2 | a block-sorting file compressor, v1.0.6 |
| bzcat | decompresses files to stdout |
| bzip2 | a block-sorting file compressor, v1.0.6 |
| bzip2recover | recovers data from damaged bzip2 files |
| cal | display a calendar |
| cat | concatenate files and print on the standard output |
| chage | change user password expiry information |
| chattr | change file attributes on a Linux file system |
| chcon | change file security context |
| chfn | change your finger information |
| chgrp | change group ownership |
| chmod | change file mode bits |
| chown | change file owner and group |
| chroot | run command or interactive shell with special root directory |
| chrt | manipulate the real-time attributes of a process |
| chsh | change your login shell |
| cksum | checksum and count the bytes in a file |
| cmp | compare two files byte by byte |
| col | filter reverse line feeds from input |
| colctrl | filter nroff output for CRT previewing |
| colrm | remove columns from a file |
| column | columnate lists |
| comm | compare two sorted files line by line |
| compile_et | error table compiler |
| cp | copy files and directories |
| csplit | split a file into sections determined by context lines |
| cut | remove sections from each line of files |
| date | print or set the system date and time |
| dd | convert and copy a file |
| df | report file system disk space usage |
| diff | compare files line by line |
| diff3 | compare three files line by line |
| dir | list directory contents |
| dircolors | color setup for ls |
| dirname | strip last component from file name |

| | |
|---------------|----------------------------------------------------------------|
| dmesg | print or control the kernel ring buffer |
| dnsdomainname | show DNS domain name |
| du | estimate file space usage |
| echo | display a line of text |
| edit | text editor |
| egrep | print lines matching a pattern |
| eject | eject removable media |
| env | run a program in a modified environment |
| envsubst | substitutes environment variables in shell format strings |
| ex | text editor |
| expand | convert tabs to spaces |
| expiry | check and enforce password expiration policy |
| expr | evaluate expressions |
| factor | factor numbers |
| fallocate | preallocate or deallocate space to a file |
| false | do nothing, unsuccessfully |
| fgrep | print lines matching a pattern |
| file | determine file type |
| find | search for files in a directory hierarchy |
| flock | manage locks from shell scripts |
| fmt | simple optimal text formatter |
| fold | wrap each input line to fit in specified width |
| free | Display amount of free and used memory in the system |
| ftp | File Transfer Protocol client. |
| fuser | identify processes using files or sockets |
| gawk | pattern scanning and processing language |
| getent | get entries from Name Service Switch libraries |
| getopt | parse command options (enhanced) |
| gettext | translate message |
| gettextize | install or upgrade gettext infrastructure |
| gpasswd | administer /etc/group and /etc/gshadow |
| grep | print lines matching a pattern |
| groups | display current group names |
| gunzip | compress or expand files |
| gzexe | compress executable files in place |
| gzip | compress or expand files |
| head | output the first part of files |
| hexdump | display file contents in hexadecimal, decimal, octal, or ascii |
| hostid | print the numeric identifier for the current host |
| hostname | show or set system host name |
| iconv | convert text from one character encoding to another |
| id | print real and effective user and group IDs |
| igawk | gawk with include files |
| info | read Info documents |
| init | systemd system and service manager |
| install | copy files and set attributes |
| install-info | update info/dir entries |
| ionice | set or get process I/O scheduling class and priority |

| | |
|------------|---------------------------------------------------------------------|
| ipcmk | make various IPC resources |
| ipcrm | remove certain IPC resources |
| ipcs | show information on IPC facilities |
| join | join lines of two files on a common field |
| kill | terminate a process |
| killall | kill processes by name |
| last | show a listing of last logged in users |
| lastb | show a listing of last logged in users |
| ldd | print shared object dependencies |
| less | opposite of more |
| lessecho | expand metacharacters |
| lesskey | specify key bindings for less |
| lexgrog | parse header information in man pages |
| link | call the link function to create a link to a file |
| ln | make links between files |
| locale | get locale-specific information |
| localedef | compile locale definition files |
| logger | enter messages into the system log |
| login | begin session on the system |
| logname | print user's login name |
| look | display lines beginning with a given string |
| ls | list directory contents |
| lsattr | list file attributes on a Linux second extended file system |
| lscpu | display information about the CPU architecture |
| lsipc | show information on IPC facilities currently employed in the system |
| lslogins | display information about known users in the system |
| mail | send and receive Internet mail |
| mailx | send and receive Internet mail |
| makeinfo | translate Texinfo documents |
| man | an interface to the on-line reference manuals |
| manpath | determine search path for manual pages |
| mcookie | generate magic cookies for xauth |
| md5sum | compute and check MD5 message digest |
| mesg | display (or do not display) messages from other users |
| mk_cmds | error table compiler |
| mkdir | make directories |
| mkfifo | make FIFOs (named pipes) |
| mknod | make block or character special files |
| mktemp | create a temporary file or directory |
| more | file perusal filter for crt viewing |
| mountpoint | see if a directory or file is a mountpoint |
| msgattrib | attribute matching and manipulation on message catalog |
| msgcat | combines several message catalogs |
| msgcmp | compare message catalog and template |
| msgcomm | match two message catalogs |
| msgconv | character set conversion for message catalog |
| msgen | create English message catalog |

| | |
|-------------|----------------------------------------------------------------|
| msgexec | process translations of message catalog |
| msgfilter | edit translations of message catalog |
| msgfmt | compile message catalog to binary format |
| msggrep | pattern matching on message catalog |
| msginit | initialize a message catalog |
| msgmerge | merge message catalog and template |
| msgunfmt | uncompile message catalog from binary format |
| msguniq | unify duplicate translations in message catalog |
| mtrace | interpret the malloc trace log |
| mv | move (rename) files |
| namei | follow a pathname until a terminal point is found |
| nano | Nano's ANOther editor, an enhanced free Pico clone |
| netctl | Control the netctl network profile manager |
| netctl-auto | Control automatic selection of wireless netctl profiles |
| newgidmap | set the gid mapping of a user namespace |
| newgrp | log in to a new group |
| newuidmap | set the uid mapping of a user namespace |
| ngettext | translate message and choose plural form |
| nice | run a program with modified scheduling priority |
| nl | number lines of files |
| nohup | run a command immune to hangups, with output to a non-tty |
| nproc | print the number of processing units available |
| nsenter | run program with namespaces of other processes |
| numfmt | Convert numbers from/to human-readable strings |
| od | dump files in octal and other formats |
| passwd | change user password |
| paste | merge lines of files |
| pathchk | check whether file names are valid or portable |
| pdftexi2dvi | convert Texinfo documents to DVI or PDF |
| peekfd | peek at file descriptors of running processes |
| pgrep | look up or signal processes based on name and other attributes |
| pidof | find the process ID of a running program. |
| pinky | lightweight finger |
| pgkill | look up or signal processes based on name and other attributes |
| pldd | display dynamic shared objects linked into a process |
| pmap | report memory map of a process |
| pod2texi | convert Pod to Texinfo |
| pr | convert text files for printing |
| printenv | print all or part of environment |
| printf | format and print data |
| prlimit | get and set process resource limits |
| prtstat | print statistics of a process |
| ps | report a snapshot of the current processes. |
| pstree | display a tree of processes |
| ptx | produce a permuted index of file contents |

| | |
|--------------|---------------------------------------------------------------------------|
| pwd | print name of current/working directory |
| pwdx | report current working directory of a process |
| rcp | Remote copy |
| readlink | print resolved symbolic links or canonical file names |
| realpath | print the resolved path |
| rename | rename files |
| renice | alter priority of running processes |
| rev | reverse lines characterwise |
| rlogin | Remote login |
| rm | remove files or directories |
| rmdir | remove empty directories |
| rnano | a restricted nano |
| rsh | Remote shell client |
| runcon | run command with specified security context |
| runuser | run a command with substitute user and group ID |
| script | make typescript of terminal session |
| scriptreplay | play back typescripts, using timing information |
| sdiff | side-by-side merge of file differences |
| sed | stream editor for filtering and transforming text |
| seq | print a sequence of numbers |
| setsid | run a program in a new session |
| setterm | set terminal attributes |
| sg | execute command as different group ID |
| sha1sum | compute and check SHA1 message digest |
| sha224sum | compute and check SHA224 message digest |
| sha256sum | compute and check SHA256 message digest |
| sha384sum | compute and check SHA384 message digest |
| sha512sum | compute and check SHA512 message digest |
| shred | overwrite a file to hide its contents, and optionally delete it |
| shuf | generate random permutations |
| slabtop | display kernel slab cache information in real time |
| sleep | delay for a specified amount of time |
| sort | sort lines of text files |
| split | split a file into pieces |
| sprof | read and display shared object profiling data |
| stat | display file or file system status |
| stdbuf | Run COMMAND, with modified buffering operations for its standard streams. |
| stty | change and print terminal line settings |
| su | run a command with substitute user and group ID |
| sum | checksum and count the blocks in a file |
| sync | Synchronize cached writes to persistent storage |
| systool | view system device information by bus, class, and topology |
| tac | concatenate and print files in reverse |
| tail | output the last part of files |
| talk | Talk client |
| tar | an archiving utility |
| taskset | set or retrieve a process's CPU affinity |

| | |
|-------------|-----------------------------------------------------------------------|
| tee | read from standard input and write to standard output and files |
| telnet | User interface to TELNET |
| test | check file types and compare values |
| texi2any | translate Texinfo documents |
| texi2dvi | convert Texinfo documents to DVI or PDF |
| texi2pdf | convert Texinfo documents to DVI or PDF |
| texindex | sort Texinfo index files |
| timeout | run a command with a time limit |
| tload | graphic representation of system load average |
| top | display Linux processes |
| touch | change file timestamps |
| tr | translate or delete characters |
| true | do nothing, successfully |
| truncate | shrink or extend the size of a file to the specified size |
| tsort | perform topological sort |
| tty | print the file name of the terminal connected to standard input |
| ul | do underlining |
| uname | print system information |
| unexpand | convert spaces to tabs |
| uniq | report or omit repeated lines |
| unlink | call the unlink function to remove the specified file |
| unshare | run program with some namespaces unshared from parent |
| uptime | Tell how long the system has been running. |
| usb-devices | print USB device details |
| users | print the user names of users currently logged in to the current host |
| utmpdump | dump UTMP and WTMP files in raw format |
| uuidgen | create a new UUID value |
| vdir | list directory contents |
| vedit | screen oriented (visual) display editor based on ex |
| vi | screen oriented (visual) display editor based on ex |
| view | screen oriented (visual) display editor based on ex |
| w | Show who is logged on and what they are doing. |
| wall | write a message to all users |
| watch | execute a program periodically, showing output fullscreen |
| wc | print newline, word, and byte counts for each file |
| whatis | display one-line manual page descriptions |
| whereis | locate the binary, source, and manual page files for a command |
| which | shows the full path of (shell) commands. |
| who | show who is logged on |
| whoami | print effective userid |
| write | send a message to another user |
| xargs | build and execute command lines from standard input |
| xgettext | extract gettext strings from source |
| yes | output a string repeatedly until killed |

| | |
|--------|-----------------------------------------------------------|
| zcat | compress or expand files |
| zcmp | compare compressed files |
| zdiff | compare compressed files |
| zforce | force a '.gz' extension on all gzip files |
| zgrep | search possibly compressed files for a regular expression |
| zless | file perusal filter for crt viewing of compressed text |
| zmore | file perusal filter for crt viewing of compressed text |
| znew | recompress .Z files to .gz files |

```
for pkg in `pacman -Ss | grep '(base[ ])' | \
grep -o core/[-a-z0-9]* | sed 's/core\\/'` ; \
do pacman -Ql $pkg | grep -o '/usr/bin/.*' | \
sed 's/\\usr\\bin\\/' | xargs whatis -s 1 | \
grep -v 1p | grep -v 1ssl | sed 's/(1)/'; \
done 2>/dev/null | grep -v "whatis what" | sort
```


Manifesto for Agile Software Development

We are uncovering better ways of developing software by doing it and helping others do it. Through this work we have come to value:

- **Individuals and interactions** over processes and tools
- **Working software** over comprehensive documentation
- **Customer collaboration** over contract negotiation
- **Responding to change** over following a plan

That is, while there is value in the items on the right, we value the items on the left more.

Principles behind the Agile Manifesto

We follow these principles:

Our highest priority is to satisfy the customer through early and continuous delivery of valuable software.

Welcome changing requirements, even late in development. Agile processes harness change for the customer's competitive advantage.

Deliver working software frequently, from a couple of weeks to a couple of months, with a preference to the shorter timescale.

Business people and developers must work together daily throughout the project.

Build projects around motivated individuals.
Give them the environment and support they need,
and trust them to get the job done.

The most efficient and effective method of
conveying information to and within a development
team is face-to-face conversation.

Working software is the primary measure of progress.

Agile processes promote sustainable development.
The sponsors, developers, and users should be able
to maintain a constant pace indefinitely.

Continuous attention to technical excellence
and good design enhances agility.

Simplicity—the art of maximizing the amount
of work not done—is essential.

The best architectures, requirements, and designs
emerge from self-organizing teams.

At regular intervals, the team reflects on how
to become more effective, then tunes and adjusts
its behavior accordingly.

ASCII

| 2 3 4 5 6 7 | 30 40 50 60 70 80 90 100 110 120 |
|------------------|----------------------------------|
| 0: 0 @ P ` p | 0: (2 < F P Z d n x |
| 1: ! 1 A Q a q | 1:) 3 = G Q [e o y |
| 2: " 2 B R b r | 2: * 4 > H R \ f p z |
| 3: # 3 C S c s | 3: ! + 5 ? I S] g q { |
| 4: \$ 4 D T d t | 4: " , 6 @ J T ^ h r |
| 5: % 5 E U e u | 5: # - 7 A K U _ i s } |
| 6: & 6 F V f v | 6: \$. 8 B L V ` j t ~ |
| 7: ' 7 G W g w | 7: % / 9 C M W a k u DEL |
| 8: (8 H X h x | 8: & 0 : D N X b l v |
| 9:) 9 I Y i y | 9: ' 1 ; E O Y c m w |
| A: * : J Z j z | |
| B: + ; K [k { | |
| C: , < L \ l | |
| D: - = M] m } | |
| E: . > N ^ n ~ | |
| F: / ? 0 _ o DEL | |

From the Jargon File: **ASCII**: /as'kee/, n.

[originally an acronym (American Standard Code for Information Interchange) but now merely conventional] The predominant character set encoding of present-day computers. The standard version uses 7 bits for each character, whereas most earlier codes (including early drafts of ASCII prior to June 1961) used fewer. This change allowed the inclusion of lower-case letters – a major win – but it did not provide for accented letters or any other letterforms not used in English (such as the German ß or ligature æ which is a letter in, for example, Norwegian). It could be worse, though. It could be much worse. See EBCDIC to understand how.

Computers are much pickier and less flexible about spelling than humans; thus, hackers need to be very precise when talking about characters, and have developed a considerable amount of verbal shorthand for them. Every character has one or more names – some formal, some concise, some silly. Common jargon names for ASCII characters are collected here. See also individual entries for bang, excl, open, ques, semi, shriek, splat, twiddle, and Yu-Shiang Whole Fish.

This list derives from revision 2.3 of the Usenet ASCII pronunciation guide. Single characters are listed in ASCII order; character pairs are sorted in by first member. For each character, common names are given in rough order of popularity, followed by names that are reported but rarely seen; official ANSI/CCITT names are surrounded by brokets: <>. Square brackets mark the particularly silly names introduced by INTERCAL. The abbreviations “l/r” and “o/c” stand for left/right and “open/close” respectively. Ordinary parentheticals provide some usage information.

- ! Common: bang ; pling; excl; not; shriek; ball-bat; <exclamation mark>. Rare: factorial; exclam; smash; cuss; boing; yell; wow; hey; wham; eureka; [spark-spot]; soldier, control.
- " Common: double quote; quote. Rare: literal mark; double-glitch; snakebite; <quotation marks>; <dieresis>; dirk; [rabbit-ears]; double prime.
- # Common: number sign; pound; pound sign; hash; sharp; crunch ; hex; [mesh]. Rare: grid; crosshatch; octothorpe; flash; <square>, pig-pen; tictactoe; scratchmark; thud; thump; splat .
- \$ Common: dollar; <dollar sign>. Rare: currency symbol; buck; cash; bling; string (from BASIC); escape (when used as the echo of ASCII ESC); ding; cache; [big money].
- % Common: percent; <percent sign>; mod; grapes. Rare: [double-oh-seven].
- & Common: <ampersand>; amp; amper; and, and sign. Rare: address (from C); reference (from C++); andpersand; bitand; background (from sh(1)); pretzel. [INTERCAL called this ampersand ; what could be sillier?]
- ' Common: single quote; quote; <apostrophe>. Rare: prime; glitch; tick; irk; pop; [spark]; <closing single quotation mark>; <acute accent>.
- (Common: l/r paren; l/r parenthesis; left/right; open/close; paren/thesis; o/c paren; o/c parenthesis; l/r parenthesis; l/r banana. Rare: so/already; lparen/rparen; <opening/closing parenthesis>; o/c round bracket, l/r round bracket, [wax/wane]; parenthisey/unparenthisey; l/r ear.
- * Common: star; [splat]; <asterisk>. Rare: wildcard; gear; dingle; mult; spider; aster; times; twinkle; glob.
- + Common: <plus>; add. Rare: cross; [intersection].
- , Common: <comma>. Rare: <cedilla>; [tail].

- Common: dash; <hyphen>; <minus>. Rare: [worm]; option; dak; bithorpe.
- . Common: dot; point; <period>; <decimal point>. Rare: radix point; full stop; [spot].
- / Common: slash; stroke; <slant>; forward slash. Rare: diagonal; solidus; over; slak; virgule; [slat].
- : Common: <colon>. Rare: dots; [two-spot].
- ; Common: <semicolon>; semi. Rare: weenie; [hybrid], pit-thwong.
- <> Common: <less/greater than>; bra/ket; l/r angle; l/r angle bracket; l/r broket. Rare: from/into, towards; read from/write to; suck/blow; comes-from/gozinta; in/out; crunch/zap (all from UNIX); tic/tac; [angle/right angle].
- = Common: <equals>; gets; takes. Rare: quadrathorpe; [half-mesh].
- ? Common: query; <question mark>; ques . Rare: quiz; whatmark; [what]; wildchar; huh; hook; buttonhook; hunchback.
- @ Common: at sign; at; strudel. Rare: each; vortex; whorl; [whirlpool]; cyclone; snail; ape; cat; rose; cabbage; <commercial at>.
- [] Common: l/r square bracket; l/r bracket; <opening/closing bracket>; bracket/unbracket. Rare: square/unsquare; [U turn/U turn back].
- \ Common: backslash, hack, whack; escape (from C/UNIX); reverse slash; slosh; backslant; backwhack. Rare: bash; <reverse slant>; reversed virgule; [backslat].
- ^ Common: hat; control; uparrow; caret; <circumflex>. Rare: xor sign, chevron; [shark (or shark-fin)]; to the ('to the power of'); fang; pointer (in Pascal).
- _ Common: <underline>; underscore; underbar; under. Rare: score; backarrow; skid; [flatworm].
- ` Common: backquote; left quote; left single quote; open quote; <grave accent>; grave. Rare: backprime; [backspark]; unapostrophe; birk; blugle; back tick; back glitch; push; <opening single quotation mark>; quasiquote.
- { } Common: o/c brace; l/r brace; l/r squiggly; l/r squiggly bracket/brace; l/r curly bracket/brace; <opening/closing brace>. Rare: brace/unbrace; curly/uncurly; leftit/rytit; l/r squirrelly; [em-brace/bracelet]. A balanced pair of these may be called curlies.
- | Common: bar; or; or-bar; v-bar; pipe; vertical bar. Rare: <vertical line>; gozinta; thru; pipesinta (last three from UNIX); [spike].
- ~ Common: <tilde>; squiggle; twiddle ; not. Rare: approx; wiggle; swung dash; enyay; [squiggle (sic)].

The Rise of “Worse is Better” (Richard Gabriel)

I and just about every designer of Common Lisp and CLOS has had extreme exposure to the MIT/Stanford style of design. The essence of this style can be captured by the phrase “the right thing.” To such a designer it is important to get all of the following characteristics right:

- **Simplicity** – the design must be simple, both in implementation and interface. It is more important for the interface to be simple than the implementation.
- **Correctness** – the design must be correct in all observable aspects. Incorrectness is simply not allowed.
- **Consistency** – the design must not be inconsistent. A design is allowed to be slightly less simple and less complete to avoid inconsistency. Consistency is as important as correctness.
- **Completeness** – the design must cover as many important situations as is practical. All reasonably expected cases must be covered. Simplicity is not allowed to overly reduce completeness.

I believe most people would agree that these are good characteristics. I will call the use of this philosophy of design the “MIT approach.” Common Lisp (with CLOS) and Scheme represent the MIT approach to design and implementation.

The worse-is-better philosophy is only slightly different:

- **Simplicity** – the design must be simple, both in implementation and interface. It is more important for the implementation to be simple than the interface. Simplicity is the most important consideration in a design.
- **Correctness** – the design must be correct in all observable aspects. It is slightly better to be simple than correct.

- Consistency – the design must not be overly inconsistent. Consistency can be sacrificed for simplicity in some cases, but it is better to drop those parts of the design that deal with less common circumstances than to introduce either implementational complexity or inconsistency.
- Completeness – the design must cover as many important situations as is practical. All reasonably expected cases should be covered. Completeness can be sacrificed in favor of any other quality. In fact, completeness must be sacrificed whenever implementation simplicity is jeopardized. Consistency can be sacrificed to achieve completeness if simplicity is retained; especially worthless is consistency of interface.

Early Unix and C are examples of the use of this school of design, and I will call the use of this design strategy the “New Jersey approach.” I have intentionally caricatured the worse-is-better philosophy to convince you that it is obviously a bad philosophy and that the New Jersey approach is a bad approach.

However, I believe that worse-is-better, even in its straw-man form, has better survival characteristics than the-right-thing, and that the New Jersey approach when used for software is a better approach than the MIT approach.

Let me start out by retelling a story that shows that the MIT/New-Jersey distinction is valid and that proponents of each philosophy actually believe their philosophy is better.

Two famous people, one from MIT and another from Berkeley (but working on Unix) once met to discuss operating system issues. The person from MIT was knowledgeable about ITS (the MIT AI Lab operating system) and had been reading the Unix sources. He was interested in how Unix solved the PC loser-ing problem. The PC loser-ing problem occurs when a user program invokes a system routine to perform a lengthy operation that might have significant state, such as IO buffers. If an interrupt occurs during the operation, the state of the user program must be saved. Because the invocation of the system routine is usually a single instruction, the PC of the user program does not adequately capture the state of the process. The

system routine must either back out or press forward. The right thing is to back out and restore the user program PC to the instruction that invoked the system routine so that resumption of the user program after the interrupt, for example, re-enters the system routine. It is called “PC loser-ing” because the PC is being coerced into “loser mode,” where “loser” is the affectionate name for “user” at MIT.

The MIT guy did not see any code that handled this case and asked the New Jersey guy how the problem was handled. The New Jersey guy said that the Unix folks were aware of the problem, but the solution was for the system routine to always finish, but sometimes an error code would be returned that signaled that the system routine had failed to complete its action. A correct user program, then, had to check the error code to determine whether to simply try the system routine again. The MIT guy did not like this solution because it was not the right thing.

The New Jersey guy said that the Unix solution was right because the design philosophy of Unix was simplicity and that the right thing was too complex. Besides, programmers could easily insert this extra test and loop. The MIT guy pointed out that the implementation was simple but the interface to the functionality was complex. The New Jersey guy said that the right tradeoff has been selected in Unix—namely, implementation simplicity was more important than interface simplicity.

The MIT guy then muttered that sometimes it takes a tough man to make a tender chicken, but the New Jersey guy didn’t understand (I’m not sure I do either).

Now I want to argue that worse-is-better is better. C is a programming language designed for writing Unix, and it was designed using the New Jersey approach. C is therefore a language for which it is easy to write a decent compiler, and it requires the programmer to write text that is easy for the compiler to interpret. Some have called C a fancy assembly language. Both early Unix and C compilers had simple structures, are easy to port, require few machine resources to run, and provide about 50%–80% of what you want from an operating system and programming language.

Half the computers that exist at any point are worse than median (smaller or slower). Unix and C work fine on them. The worse-is-better philosophy means that implementation simplicity has highest priority, which means Unix and C are easy to port on such machines. Therefore, one expects that if the 50% functionality Unix and C support is satisfactory, they will start to appear everywhere. And they have, haven't they?

Unix and C are the ultimate computer viruses.

A further benefit of the worse-is-better philosophy is that the programmer is conditioned to sacrifice some safety, convenience, and hassle to get good performance and modest resource use. Programs written using the New Jersey approach will work well both in small machines and large ones, and the code will be portable because it is written on top of a virus.

It is important to remember that the initial virus has to be basically good. If so, the viral spread is assured as long as it is portable. Once the virus has spread, there will be pressure to improve it, possibly by increasing its functionality closer to 90%, but users have already been conditioned to accept worse than the right thing. Therefore, the worse-is-better software first will gain acceptance, second will condition its users to expect less, and third will be improved to a point that is almost the right thing. In concrete terms, even though Lisp compilers in 1987 were about as good as C compilers, there are many more compiler experts who want to make C compilers better than want to make Lisp compilers better.

The good news is that in 1995 we will have a good operating system and programming language; the bad news is that they will be Unix and C++.

There is a final benefit to worse-is-better. Because a New Jersey language and system are not really powerful enough to build complex monolithic software, large systems must be designed to reuse components. Therefore, a tradition of integration springs up.

How does the right thing stack up? There are two basic scenarios: the "big complex system scenario" and the "diamond-like jewel" scenario.

The "big complex system" scenario goes like this:

First, the right thing needs to be designed. Then its implementation needs to be designed. Finally it is implemented. Because it is the right thing, it has nearly 100% of desired functionality, and implementation simplicity was never a concern so it takes a long time to implement. It is large and complex. It requires complex tools to use properly. The last 20% takes 80% of the effort, and so the right thing takes a long time to get out, and it only runs satisfactorily on the most sophisticated hardware.

The “diamond-like jewel” scenario goes like this:

The right thing takes forever to design, but it is quite small at every point along the way. To implement it to run fast is either impossible or beyond the capabilities of most implementors.

The two scenarios correspond to Common Lisp and Scheme.

The first scenario is also the scenario for classic artificial intelligence software.

The right thing is frequently a monolithic piece of software, but for no reason other than that the right thing is often designed monolithically. That is, this characteristic is a happenstance.

The lesson to be learned from this is that it is often undesirable to go for the right thing first. It is better to get half of the right thing available so that it spreads like a virus. Once people are hooked on it, take the time to improve it to 90% of the right thing.

A wrong lesson is to take the parable literally and to conclude that C is the right vehicle for AI software. The 50% solution has to be basically right, and in this case it isn't.

But, one can conclude only that the Lisp community needs to seriously rethink its position on Lisp design. I will say more about this later.

Emacs

Note, C- means Ctrl- and M- means Meta- (Alt- or Esc-).

| | |
|----------------|----------------------------------------------|
| C-f C-b | move forward, backward one character |
| C-n C-p | move down, up one line |
| M-f M-b | move forward, backward one word |
| C-a C-e | move to beginning, end of line |
| M-a M-e | move to beginning, end of sentence |
| M-{ M-} | move to beginning, end of paragraph |
| M-< M-> | move to beginning, end of buffer |
| C-v M-v | page down, page up |
| M-g g | go to line |
| C-d Del | delete forward, backward one character |
| M-d M-Del | delete forward, backward one word |
| C-k, M-0 C-k | delete forward, backward one line |
| M-k, C-x Del | delete forward, backward one sentence |
| C-Space | start highlighting region |
| C-w M-w | kill (cut) region, copy region |
| C-y, M-y | yank (paste) last thing killed, prior kill |
| C-s C-r | search forward, search reverse |
| C-M-s C-M-r | search regex forward, reverse |
| M-% | replace string |
| Tab C-M-\ | indent line, region |
| C-x C-f | open file |
| C-x C-s, C-x s | save buffer, save all buffers |
| C-x C-w | save buffer with new name |
| C-x C-c | quit |
| C-g | abort a command |
| M-x | run a function |
| C-h a | apropos: show commands matching some string |
| C-h k | describe a key |
| C-x 2, C-x 3 | split window (hsplit), split window (vsplit) |
| C-x 1 | close all other windows |
| C-x o | switch to other window |
| C-x b | switch buffers |
| C-x C-b | list buffers |
| C-x k | kill buffer |

Nearly all commands support a prefix number, e.g., C-u 5 C-k delete 5 lines.

Basics of the Unix Philosophy (Eric S. Raymond)

The ‘Unix philosophy’ originated with Ken Thompson’s early meditations on how to design a small but capable operating system with a clean service interface. It grew as the Unix culture learned things about how to get maximum leverage out of Thompson’s design. It absorbed lessons from many sources along the way.

The Unix philosophy is not a formal design method. It wasn’t handed down from the high fastnesses of theoretical computer science as a way to produce theoretically perfect software. Nor is it that perennial executive’s mirage, some way to magically extract innovative but reliable software on too short a deadline from unmotivated, badly managed, and underpaid programmers.

The Unix philosophy (like successful folk traditions in other engineering disciplines) is bottom-up, not top-down. It is pragmatic and grounded in experience. It is not to be found in official methods and standards, but rather in the implicit half-reflexive knowledge, the *expertise* that the Unix culture transmits. It encourages a sense of proportion and skepticism – and shows both by having a sense of (often subversive) humor.

Doug McIlroy, the inventor of Unix pipes and one of the founders of the Unix tradition, had this to say at the time:

- (i) Make each program do one thing well. To do a new job, build afresh rather than complicate old programs by adding new features.
- (ii) Expect the output of every program to become the input to another, as yet unknown, program. Don’t clutter output with extraneous information. Avoid stringently columnar or binary input formats. Don’t insist on interactive input.
- (iii) Design and build software, even operating systems, to be tried early, ideally within weeks. Don’t hesitate to throw away the clumsy parts and rebuild them.

- (iv) Use tools in preference to unskilled help to lighten a programming task, even if you have to detour to build the tools and expect to throw some of them out after you've finished using them.

He later summarized it this way (quoted in *A Quarter Century of Unix*):

This is the Unix philosophy: Write programs that do one thing and do it well. Write programs to work together. Write programs to handle text streams, because that is a universal interface.

Rob Pike, who became one of the great masters of C, offers a slightly different angle in *Notes on C Programming*:

- Rule 1. You can't tell where a program is going to spend its time. Bottlenecks occur in surprising places, so don't try to second guess and put in a speed hack until you've proven that's where the bottleneck is.
- Rule 2. Measure. Don't tune for speed until you've measured, and even then don't unless one part of the code overwhelms the rest.
- Rule 3. Fancy algorithms are slow when n is small, and n is usually small. Fancy algorithms have big constants. Until you know that n is frequently going to be big, don't get fancy. (Even if n does get big, use Rule 2 first.)
- Rule 4. Fancy algorithms are buggier than simple ones, and they're much harder to implement. Use simple algorithms as well as simple data structures.
- Rule 5. Data dominates. If you've chosen the right data structures and organized things well, the algorithms will almost always be self-evident. Data structures, not algorithms, are central to programming.
- Rule 6. There is no Rule 6.

Ken Thompson, the man who designed and implemented the first Unix, reinforced Pike's rule 4 with a gnomish maxim worthy of a Zen patriarch:

When in doubt, use brute force.

More of the Unix philosophy was implied not by what these elders said but by what they did and the example Unix itself set. Looking at the whole, we can abstract the following ideas:

1. Rule of Modularity: Write simple parts connected by clean interfaces.
2. Rule of Clarity: Clarity is better than cleverness.
3. Rule of Composition: Design programs to be connected to other programs.
4. Rule of Separation: Separate policy from mechanism; separate interfaces from engines.
5. Rule of Simplicity: Design for simplicity; add complexity only where you must.
6. Rule of Parsimony: Write a big program only when it is clear by demonstration that nothing else will do.
7. Rule of Transparency: Design for visibility to make inspection and debugging easier.
8. Rule of Robustness: Robustness is the child of transparency and simplicity.
9. Rule of Representation: Fold knowledge into data so program logic can be stupid and robust.
10. Rule of Least Surprise: In interface design, always do the least surprising thing.
11. Rule of Silence: When a program has nothing surprising to say, it should say nothing.

12. Rule of Repair: When you must fail, fail noisily and as soon as possible.
13. Rule of Economy: Programmer time is expensive; conserve it in preference to machine time.
14. Rule of Generation: Avoid hand-hacking; write programs to write programs when you can.
15. Rule of Optimization: Prototype before polishing. Get it working before you optimize it.
16. Rule of Diversity: Distrust all claims for “one true way”.
17. Rule of Extensibility: Design for the future, because it will be here sooner than you think.

If you're new to Unix, these principles are worth some meditation. Software-engineering texts recommend most of them; but most other operating systems lack the right tools and traditions to turn them into practice, so most programmers can't apply them with any consistency. They come to accept blunt tools, bad designs, overwork, and bloated code as normal – and then wonder what Unix fans are so annoyed about.

Rule of Modularity: Write simple parts connected by clean interfaces.

As Brian Kernighan once observed, “Controlling complexity is the essence of computer programming.” Debugging dominates development time, and getting a working system out the door is usually less a result of brilliant design than it is of managing not to trip over your own feet too many times.

Assemblers, compilers, flowcharting, procedural programming, structured programming, “artificial intelligence”, fourth-generation languages, object orientation, and software development methodologies without number have been touted and sold as a cure for this problem. All have failed as cures, if only because they ‘succeeded’ by escalating the normal level of program complexity to the point where (once again) human brains

could barely cope. As Fred Brooks famously observed, there is no silver bullet.

The only way to write complex software that won't fall on its face is to hold its global complexity down – to build it out of simple parts connected by well-defined interfaces, so that most problems are local and you can have some hope of upgrading a part without breaking the whole.

Rule of Clarity: Clarity is better than cleverness.

Because maintenance is so important and so expensive, write programs as if the most important communication they do is not to the computer that executes them but to the human beings who will read and maintain the source code in the future (including yourself).

In the Unix tradition, the implications of this advice go beyond just commenting your code. Good Unix practice also embraces choosing your algorithms and implementations for future maintainability. Buying a small increase in performance with a large increase in the complexity and obscurity of your technique is a bad trade – not merely because complex code is more likely to harbor bugs, but also because complex code will be harder to read for future maintainers.

Code that is graceful and clear, on the other hand, is less likely to break – and more likely to be instantly comprehended by the next person to have to change it. This is important, especially when that next person might be yourself some years down the road.

Never struggle to decipher subtle code three times. Once might be a one-shot fluke, but if you find yourself having to figure it out a second time – because the first was too long ago and you've forgotten details – it is time to comment the code so that the third time will be relatively painless.

– Henry Spencer

Rule of Composition: Design programs to be connected with other programs.

It's hard to avoid programming overcomplicated monoliths if none of your programs can talk to each other.

Unix tradition strongly encourages writing programs that read and write simple, textual, stream-oriented, device-independent formats. Under classic Unix, as many programs as possible are written as simple *filters*, which take a simple text stream on input and process it into another simple text stream on output.

Despite popular mythology, this practice is favored not because Unix programmers hate graphical user interfaces. It's because if you don't write programs that accept and emit simple text streams, it's much more difficult to hook the programs together.

Text streams are to Unix tools as messages are to objects in an object-oriented setting. The simplicity of the text-stream interface enforces the encapsulation of the tools. More elaborate forms of inter-process communication, such as remote procedure calls, show a tendency to involve programs with each others' internals too much.

To make programs composable, make them independent. A program on one end of a text stream should care as little as possible about the program on the other end. It should be made easy to replace one end with a completely different implementation without disturbing the other.

GUIs can be a very good thing. Complex binary data formats are sometimes unavoidable by any reasonable means. But before writing a GUI, it's wise to ask if the tricky interactive parts of your program can be segregated into one piece and the workhorse algorithms into another, with a simple command stream or application protocol connecting the two. Before devising a tricky binary format to pass data around, it's worth experimenting to see if you can make a simple textual format work and accept a little parsing overhead in return for being able to hack the data stream with general-purpose tools.

When a serialized, protocol-like interface is not natural for

the application, proper Unix design is to at least organize as many of the application primitives as possible into a library with a well-defined API. This opens up the possibility that the application can be called by linkage, or that multiple interfaces can be glued on it for different tasks.

Rule of Separation: Separate policy from mechanism; separate interfaces from engines.

In our discussion of what Unix gets wrong, we observed that the designers of X made a basic decision to implement “mechanism, not policy” – to make X a generic graphics engine and leave decisions about user-interface style to toolkits and other levels of the system. We justified this by pointing out that policy and mechanism tend to mutate on different timescales, with policy changing much faster than mechanism. Fashions in the look and feel of GUI toolkits may come and go, but raster operations and compositing are forever.

Thus, hardwiring policy and mechanism together has two bad effects: It makes policy rigid and harder to change in response to user requirements, and it means that trying to change policy has a strong tendency to destabilize the mechanisms.

On the other hand, by separating the two we make it possible to experiment with new policy without breaking mechanisms. We also make it much easier to write good tests for the mechanism (policy, because it ages so quickly, often does not justify the investment).

This design rule has wide application outside the GUI context. In general, it implies that we should look for ways to separate interfaces from engines.

One way to effect that separation is, for example, to write your application as a library of C service routines that are driven by an embedded scripting language, with the application flow of control written in the scripting language rather than C. A classic example of this pattern is the Emacs editor, which uses an embedded Lisp interpreter to control editing primitives written in C.

Another way is to separate your application into cooperating front-end and back-end processes communicating through a specialized application protocol over sockets. The front end implements policy; the back end, mechanism. The global complexity of the pair will often be far lower than that of a single-process monolith implementing the same functions, reducing your vulnerability to bugs and lowering life-cycle costs.

Rule of Simplicity: Design for simplicity; add complexity only where you must.

Many pressures tend to make programs more complicated (and therefore more expensive and buggy). One such pressure is technical machismo. Programmers are bright people who are (often justly) proud of their ability to handle complexity and juggle abstractions. Often they compete with their peers to see who can build the most intricate and beautiful complexities. Just as often, their ability to design outstrips their ability to implement and debug, and the result is expensive failure.

The notion of “intricate and beautiful complexities” is almost an oxymoron. Unix programmers vie with each other for “simple and beautiful” honors – a point that’s implicit in these rules, but is well worth making overt.

– Doug McIlroy

Even more often (at least in the commercial software world) excessive complexity comes from project requirements that are based on the marketing fad of the month rather than the reality of what customers want or software can actually deliver. Many a good design has been smothered under marketing’s pile of “checklist features” – features that, often, no customer will ever use. And a vicious circle operates; the competition thinks it has to compete with chrome by adding more chrome. Pretty soon, massive bloat is the industry standard and everyone is using huge, buggy programs not even their developers can love.

Either way, everybody loses in the end.

The only way to avoid these traps is to encourage a software culture that knows that small is beautiful, that actively resists bloat and complexity: an engineering tradition that puts a high value on simple solutions, that looks for ways to break program systems up into small cooperating pieces, and that reflexively fights attempts to gussy up programs with a lot of chrome (or, even worse, to design programs *around* the chrome).

That would be a culture a lot like Unix's.

Rule of Parsimony: Write a big program only when it is clear by demonstration that nothing else will do.

'Big' here has the sense both of large in volume of code and of internal complexity. Allowing programs to get large hurts maintainability. Because people are reluctant to throw away the visible product of lots of work, large programs invite over-investment in approaches that are failed or suboptimal.

Rule of Transparency: Design for visibility to make inspection and debugging easier.

Because debugging often occupies three-quarters or more of development time, work done early to ease debugging can be a very good investment. A particularly effective way to ease debugging is to design for *transparency* and *discoverability*.

A software system is *transparent* when you can look at it and immediately understand what it is doing and how. It is *discoverable* when it has facilities for monitoring and display of internal state so that your program not only functions well but can be *seen* to function well.

Designing for these qualities will have implications throughout a project. At minimum, it implies that debugging options should not be minimal afterthoughts. Rather, they should be designed in from the beginning – from the point of view that the program should be able to both demonstrate its own correctness and communicate to future developers the original developer's mental model of the problem it solves.

For a program to demonstrate its own correctness, it needs to be using input and output formats sufficiently simple so that the proper relationship between valid input and correct output is easy to check.

The objective of designing for transparency and discoverability should also encourage simple interfaces that can easily be manipulated by other programs – in particular, test and monitoring harnesses and debugging scripts.

Rule of Robustness: Robustness is the child of transparency and simplicity.

Software is said to be robust when it performs well under unexpected conditions which stress the designer's assumptions, as well as under normal conditions.

Most software is fragile and buggy because most programs are too complicated for a human brain to understand all at once. When you can't reason correctly about the guts of a program, you can't be sure it's correct, and you can't fix it if it's broken.

It follows that the way to make robust programs is to make their internals easy for human beings to reason about. There are two main ways to do that: transparency and simplicity.

For robustness, designing in tolerance for unusual or extremely bulky inputs is also important. Bearing in mind the Rule of Composition helps; input generated by other programs is notorious for stress-testing software (e.g., the original Unix C compiler reportedly needed small upgrades to cope well with Yacc output). The forms involved often seem useless to humans. For example, accepting empty lists, strings, etc., even in places where a human would seldom or never supply an empty string, avoids having to special-case such situations when generating the input mechanically.

– Henry Spencer

One very important tactic for being robust under odd inputs is to avoid having special cases in your code. Bugs often lurk in the code for handling special cases, and in the interactions among parts of the code intended to handle different special cases.

We observed above that software is *transparent* when you can look at it and immediately see what is going on. It is *simple* when what is going on is uncomplicated enough for a human brain to reason about all the potential cases without strain. The more your programs have both of these qualities, the more robust they will be.

Modularity (simple parts, clean interfaces) is a way to organize programs to make them simpler. There are other ways to fight for simplicity. Here's another one.

Rule of Representation: Fold knowledge into data, so program logic can be stupid and robust.

Even the simplest procedural logic is hard for humans to verify, but quite complex data structures are fairly easy to model and reason about. To see this, compare the expressiveness and explanatory power of a diagram of (say) a fifty-node pointer tree with a flowchart of a fifty-line program. Or, compare an array initializer expressing a conversion table with an equivalent switch statement. The difference in transparency and clarity is dramatic. See Rob Pike's Rule 5.

Data is more tractable than program logic. It follows that where you see a choice between complexity in data structures and complexity in code, choose the former. More: in evolving a design, you should actively seek ways to shift complexity from code to data.

The Unix community did not originate this insight, but a lot of Unix code displays its influence. The C language's facility at manipulating pointers, in particular, has encouraged the use of dynamically-modified reference structures at all levels of coding from the kernel upward. Simple pointer chases in such structures frequently do duties that implementations in other

languages would instead have to embody in more elaborate procedures.

Rule of Least Surprise: In interface design, always do the least surprising thing.

(This is also widely known as the Principle of Least Astonishment.)

The easiest programs to use are those that demand the least new learning from the user – or, to put it another way, the easiest programs to use are those that most effectively connect to the user's pre-existing knowledge.

Therefore, avoid gratuitous novelty and excessive cleverness in interface design. If you're writing a calculator program, '+' should always mean addition! When designing an interface, model it on the interfaces of functionally similar or analogous programs with which your users are likely to be familiar.

Pay attention to your expected audience. They may be end users, they may be other programmers, or they may be system administrators. What is least surprising can differ among these groups.

Pay attention to tradition. The Unix world has rather well-developed conventions about things like the format of configuration and run-control files, command-line switches, and the like. These traditions exist for a good reason: to tame the learning curve. Learn and use them.

The flip side of the Rule of Least Surprise is to avoid making things superficially similar but really a little bit different. This is extremely treacherous because the seeming familiarity raises false expectations. It's often better to make things distinctly different than to make them almost the same.

– Henry Spencer

Rule of Silence: When a program has nothing surprising to say, it should say nothing.

One of Unix's oldest and most persistent design rules is that when a program has nothing interesting or surprising to say, it should shut up. Well-behaved Unix programs do their jobs unobtrusively, with a minimum of fuss and bother. Silence is golden.

This “silence is golden” rule evolved originally because Unix predates video displays. On the slow printing terminals of 1969, each line of unnecessary output was a serious drain on the user's time. That constraint is gone, but excellent reasons for terseness remain.

I think that the terseness of Unix programs is a central feature of the style. When your program's output becomes another's input, it should be easy to pick out the needed bits. And for people it is a human-factors necessity – important information should not be mixed in with verbosity about internal program behavior. If all displayed information is important, important information is easy to find.

– Ken Arnold

Well-designed programs treat the user's attention and concentration as a precious and limited resource, only to be claimed when necessary.

Rule of Repair: Repair what you can – but when you must fail, fail noisily and as soon as possible.

Software should be transparent in the way that it fails, as well as in normal operation. It's best when software can cope with unexpected conditions by adapting to them, but the worst kinds of bugs are those in which the repair doesn't succeed and the problem quietly causes corruption that doesn't show up until much later.

Therefore, write your software to cope with incorrect inputs and its own execution errors as gracefully as possible. But when it cannot, make it fail in a way that makes diagnosis of the problem as easy as possible.

Consider also Postel's Prescription: "Be liberal in what you accept, and conservative in what you send." Postel was speaking of network service programs, but the underlying idea is more general. Well-designed programs cooperate with other programs by making as much sense as they can from ill-formed inputs; they either fail noisily or pass strictly clean and correct data to the next program in the chain.

However, heed also this warning:

The original HTML documents recommended "be generous in what you accept", and it has bedeviled us ever since because each browser accepts a different superset of the specifications. It is the *specifications* that should be generous, not their interpretation.

– Doug McIlroy

McIlroy adjures us to *design* for generosity rather than compensating for inadequate standards with permissive implementations. Otherwise, as he rightly points out, it's all too easy to end up in tag soup.

Rule of Economy: Programmer time is expensive; conserve it in preference to machine time.

In the early minicomputer days of Unix, this was still a fairly radical idea (machines were a great deal slower and more expensive then). Nowadays, with every development shop and most users (apart from the few modeling nuclear explosions or doing 3D movie animation) awash in cheap machine cycles, it may seem too obvious to need saying.

Somehow, though, practice doesn't seem to have quite caught up with reality. If we took this maxim really seriously throughout software development, most applications would be

written in higher-level languages like Perl, Tcl, Python, Java, Lisp and even shell – languages that ease the programmer’s burden by doing their own memory management.

And indeed this is happening within the Unix world, though outside it most applications shops still seem stuck with the old-school Unix strategy of coding in C (or C++). Later in this book we’ll discuss this strategy and its tradeoffs in detail.

One other obvious way to conserve programmer time is to teach machines how to do more of the low-level work of programming. This leads to...

Rule of Generation: Avoid hand-hacking; write programs to write programs when you can.

Human beings are notoriously bad at sweating the details. Accordingly, any kind of hand-hacking of programs is a rich source of delays and errors. The simpler and more abstracted your program specification can be, the more likely it is that the human designer will have gotten it right. Generated code (at *every* level) is almost always cheaper and more reliable than hand-hacked.

We all know this is true (it’s why we have compilers and interpreters, after all) but we often don’t think about the implications. High-level-language code that’s repetitive and mind-numbing for humans to write is just as productive a target for a code generator as machine code. It pays to use code generators when they can raise the level of abstraction – that is, when the specification language for the generator is simpler than the generated code, and the code doesn’t have to be hand-hacked afterwards.

In the Unix tradition, code generators are heavily used to automate error-prone detail work. Parser/lexer generators are the classic examples; makefile generators and GUI interface builders are newer ones.

Rule of Optimization: Prototype before polishing. Get it working before you optimize it.

The most basic argument for prototyping first is Kernighan & Plauger's; "90% of the functionality delivered now is better than 100% of it delivered never." Prototyping first may help keep you from investing far too much time for marginal gains.

For slightly different reasons, Donald Knuth (author of *The Art Of Computer Programming*, one of the field's few true classics) popularized the observation that "Premature optimization is the root of all evil." And he was right.

Rushing to optimize before the bottlenecks are known may be the only error to have ruined more designs than feature creep. From tortured code to incomprehensible data layouts, the results of obsessing about speed or memory or disk usage at the expense of transparency and simplicity are everywhere. They spawn innumerable bugs and cost millions of man-hours – often, just to get marginal gains in the use of some resource much less expensive than debugging time.

Disturbingly often, premature local optimization actually hinders global optimization (and hence reduces overall performance). A prematurely optimized portion of a design frequently interferes with changes that would have much higher payoffs across the whole design, so you end up with both inferior performance and excessively complex code.

In the Unix world there is a long-established and very explicit tradition (exemplified by Rob Pike's comments above and Ken Thompson's maxim about brute force) that says: *Prototype, then polish. Get it working before you optimize it.* Or: Make it work first, then make it work fast. 'Extreme programming' guru Kent Beck, operating in a different culture, has usefully amplified this to: "Make it run, then make it right, then make it fast."

The thrust of all these quotes is the same: get your design right with an un-optimized, slow, memory-intensive implementation before you try to tune. Then, tune systematically, looking for the places where you can buy big performance wins with the smallest possible increases in local complexity.

Prototyping is important for system design as well as optimization – it is much easier to judge whether a prototype does what you want than it is to read a long specification. I remember one development manager at Bellcore who fought against the “requirements” culture years before anybody talked about “rapid prototyping” or “agile development.” He wouldn’t issue long specifications; he’d lash together some combination of shell scripts and awk code that did roughly what was needed, tell the customers to send him some clerks for a few days, and then have the customers come in and look at their clerks using the prototype and tell him whether or not they liked it. If they did, he would say “you can have it industrial strength so-many-months from now at such-and-such cost.” His estimates tended to be accurate, but he lost out in the culture to managers who believed that requirements writers should be in control of everything.

– Mike Lesk

Using prototyping to learn which features you don’t have to implement helps optimization for performance; you don’t have to optimize what you don’t write. The most powerful optimization tool in existence may be the delete key.

One of my most productive days was throwing away 1000 lines of code.

– Ken Thompson

Rule of Diversity: Distrust all claims for “one true way.”

Even the best software tools tend to be limited by the imaginations of their designers. Nobody is smart enough to optimize for everything, nor to anticipate all the uses to which their software might be put. Designing rigid, closed software that won’t talk to the rest of the world is an unhealthy form of arrogance.

Therefore, the Unix tradition includes a healthy mistrust of “one true way” approaches to software design or implementation. It embraces multiple languages, open extensible systems, and customization hooks everywhere.

Rule of Extensibility: Design for the future, because it will be here sooner than you think.

If it is unwise to trust other people’s claims for “one true way,” it’s even more foolish to believe them about your own designs. Never assume you have the final answer. Therefore, leave room for your data formats and code to grow; otherwise, you will often find that you are locked into unwise early choices because you cannot change them while maintaining backward compatibility.

When you design protocols or file formats, make them sufficiently self-describing to be extensible. Always, always either include a version number, or compose the format from self-contained, self-describing clauses in such a way that new clauses can be readily added and old ones dropped without confusing format-reading code. Unix experience tells us that the marginal extra overhead of making data layouts self-describing is paid back a thousandfold by the ability to evolve them forward without breaking things.

When you design code, organize it so future developers will be able to plug new functions into the architecture without having to scrap and rebuild the architecture. This rule is not a license to add features you don’t yet need; it’s advice to write your code so that adding features later when you do need them is easy. Make the joints flexible, and put “If you ever need to . . .” comments in your code. You owe this grace to people who will use and maintain your code after you.

You’ll be there in the future too, maintaining code you may have half forgotten under the press of more recent projects. When you design for the future, the sanity you save may be your own.

Bash

Running commands

| | |
|----------------------------------|------------------------------------------------------------------------------------|
| <code>cmd ; cmd2</code> | run <code>cmd2</code> after <code>cmd</code> |
| <code>cmd && cmd2</code> | run <code>cmd2</code> after <code>cmd</code> if <code>cmd</code> is successful |
| <code>cmd cmd2</code> | run <code>cmd2</code> after <code>cmd</code> if <code>cmd</code> is not successful |
| <code>cmd &</code> | run <code>cmd</code> and send to background as a job |
| <code>jobs</code> | list background jobs |
| <code>fg</code> | bring recent background job back to foreground |
| <code>fg [num]</code> | bring specific job to foreground |
| <code>C-z</code> | in most programs, suspend |
| <code>bg</code> | send suspended job to background (and resume running) |
| <code>bg [num]</code> | send specific job to background |

Input tricks

| | |
|--------------------------------|-------------------------------------------------------------|
| <code>C-a C-e</code> | jump to beginning, end of line |
| <code>C-r C-s</code> | search reverse, forward in history |
| <code>C-u C-k</code> | delete before, after cursor |
| <code>C-x C-e</code> | open editor with current command line contents |
| <code>C-_</code> | undo typing |
| <code>M-.</code> | (Alt- or Esc-.) last command arg; repeat for each arg |
| <code>!!</code> | repeat last command, e.g., <code>sudo !!</code> |
| <code>\$! !*</code> | refer to last command: last arg, all args |
| <code>^foo^bar</code> | replace last command <code>foo</code> with <code>bar</code> |
| <code>{a,b,c}</code> | expand variations, e.g., <code>cp foo{.txt,.txt2}</code> |
| <code>:(){ : :& };:</code> | fork bomb |

Variables and Quoting

| | |
|------------------------------|---------------------------------------------|
| <code>foo=bar</code> | set variable |
| <code>\$foo</code> | use variable |
| <code>`cmd`</code> | run command and return its stdout |
| <code>foo=`cmd`</code> | save <code>cmd</code> 's stdout to variable |
| <code>foo=`expr 5*2`</code> | perform math |
| <code>\$0 \$1 \$2 ...</code> | script name, first arg, second arg, ... |

Redirection and Pipes

| | |
|--------------------------------|--------------------------------|
| <code>cmd < file</code> | stdin comes from file |
| <code>cmd > file</code> | save stdout to file |
| <code>cmd >> file</code> | save stdout to file, appending |
| <code>cmd cmd2</code> | cmd stdout becomes cmd2 stdin |

Conditionals and Loops

| | |
|------------------------------------------|----------------------------------------|
| <code>if [...] ; then ...; fi</code> | if-then structure |
| <code>elif [...]; then</code> | else-if snippet |
| <code>else;</code> | else snippet |
| <code>for i in list; do ...; done</code> | for loop |
| <code>-r -w -x</code> | file is readable, writable, executable |
| <code>-f -d -e</code> | file is normal file, directory, exists |
| <code>= !=</code> | string comparisons |
| <code>-eq -ne -lt -le -gt -ge</code> | numeric comparisons |
| <code>! -a -o</code> | boolean operators (not, and, or) |

Object Oriented Programming

From <http://wiki.c2.com/?ObjectOrientedProgramming>

See NygaardClassification, for the definitive definition.

Nygaard did not coin the term “Object-Oriented Programming,” AlanKay did, so I fail to see how Nygaard’s classification is “definitive”. Yes, Nygaard and Dahl’s Simula was the first language to have “objects” in it, if you ignore Dr. IvanSutherland’s SketchPad that predates it by five years, but regardless, Nygaard and Dahl did not use the term OO to describe Simula.

AlanKay and the XeroxLearningResearchGroup borrowed from many languages, including Simula and Lisp, when making Smalltalk. In Smalltalk, everything is an object, and every action is accomplished by sending messages to objects. It was for this reason Kay described his language as being “object-oriented.” Smalltalk was responsible for popularizing OO as a paradigm. Smalltalk is to OO what Lisp is to functional programming. To this day Smalltalk is considered the holy grail of OO; every single OO language is compared to it. Simula, unfortunately, is nothing but ALGOL with classes.

So if any one person’s definition of OO is definitive, it’s probably AlanKay’s. See AlanKaysDefinitionOfObjectOriented and AlanKayOnObjects.

During the course of history many authors tried to redefine the term to mean lots of things. Some of those Definitions-ForOo are discussed below.

For good or bad, the “street” ultimately is what defines existing terms, not the originators. For example, “decimation” used to mean the Roman practice of punishing armies by killing one out of every 10 soldiers. Now it means total destruction, not a tenth. Nygaard may the Roman here.

- Ah, very clever observation – if in fact you can follow up with what the street has decided the term means. Which you can’t, because the street hasn’t decided.

- *Are you suggesting we go with the originator because the street has not reached a consensus? Generally dictionaries tend to rank the different definitions. Thus, if it fits the most popular variations, then it is more OO than something that only fits a lower definition. Hmmmm. Does something that fits 2 lower definitions get considered a better fit than something that only fits the top def?*

- (1) criticizing an argument that argues against topic A is not the same thing as supporting A (might or might not be). (2) Dictionaries sometimes, but not always, give an indication as to which definitions are more common. That does **NOT** make the less common usages less correct.

Most contemporary dictionaries do not claim to tell which usages are correct; they tell which usages are or were in use.

- The “Street” does not decide the meaning of terms. That’s what the academic community is good for (and dictionaries). Otherwise the whole of human language would dissolve to ambiguity, especially if there are those who just want to forcefully use terms in speech to grab power. In any event, if there’s not a good consensus within the publications of academia, or if no leader is staking out the definitions used for and in the field, *then it should go with whoever coins it first.*
- Since we’re talking about Romans, plebeians do not write history. Nygaard may not be popular among the large masses of people monkeying around with keyboards, but he certainly is relevant where it matters. Important books like SiCp and TheoryOfObjects subscribe to Nygaard’s definition, and therefore generation of elite students are going to learn that. The books that popularize the other stuff for the consumption of the “street” are likely to be swallowed by the trash bin of history.
 - Although not an air-tight argument, I’m inclined to agree with this. Unlike street dialects, technical jar-

gon tends to have authoritative sources of judgments of correctness, as do technical issues in general. Mis-use of the technical jargon by the masses constitutes a borrowing with a shift in meaning, not a change in the original meaning. Classic example: “quantum jump” in common parlance does not mean the same thing that it does in physics. That does not make the physics meaning incorrect. It remains intact.

- *But the masses in this case are not necessarily “wrong”. Physics cannot be changed by the masses so there is a central point to compare. Software design has no objective core so far other than metrics such as run speed and code size, which most agree are only a portion of the picture. Plus, some “big names” disagree with Nygaard’s definition, not just the “masses”.*

I suggest the above be moved to DefinitionsForOo.

The heart of OOP is the Refactor ReplaceConditionalWithPolymorphism.

I don’t think this is a consensus opinion. OoBestFeaturePoll.

Consider a procedural program with several ‘switch’ statements, each switching on the same thing. Every ‘switch’ block has the same list of ‘case’ statements in it.

Topic and counter-arguments moved to SwitchStatementsSmell.

In English, the “object” of a sentence is the receiver of an action. For example, in the sentence, “Click the button”, the “button” is the object.

Unfortunately, in programming things called “objects” perform “actions”. In this sense they are more akin to the English grammar “subject”.

Correct. In English grammar the subject performs the action. The subject may receive the action when the verb is passive or intransitive. An indirect object could be argued to not receive an action, although in the whole scheme of things it does.

ObjectOrientedProgramming focuses on “objects”, so the above might be codified as:

```
button.click()
```

A contrasting example is ProceduralProgramming, which focuses on “verbs”. This yields code like this:

```
click(button)
```

This involves an ideological change as well as a syntactic change. Rather than thinking about what events happen and how they happen, we think about what entities exist and how they interact.

- An interjection here. “click” in this case is not an action being performed by **button**. Click is what someone else does to **button**, and hence **button** really is the object of the sentence. What seems missing is the subject – the entity performing the action “click” to which button is reacting. The subject seems almost always implicitly to be the object containing the code, in this case `button.click()`.

No, it's only a trivial change in syntactic sugar. There are OO systems in which the normal function calling syntax is used, e.g. CommonLispObjectSystem: (click button). In either syntax, it's perfectly clear that button is a noun, and click is a verb. There is no illusion that in `button.click()`, click is anything other than a member function, and button is an object.

No, it's much more than that. In `button.click()`, click could be one of any number of methods, on any number of buttons, while `click(button)`, click is one method for any button. The reason OO programmers use that example, it that it's a good way to get procedural programmers to see the difference, as most of us were procedural programmers at first too. It's not so much about the syntax as about creating a syntax that allows multiple virtual implementations. Procedural paradigms don't do this, object systems do, and CLOS is an object system. It has those multiple virtual implementations.

The Object-Oriented principles are emergent properties of this shifted focus.

See the Object FAQ at <http://www.cyberdyne-object-sys.com/oofaq2/>, which includes the following:

Simula was the first object-oriented language providing objects, classes, inheritance, and dynamic typing in 1967 (in addition to its AlgolSixty subset). It was intended as a conveyance of object-oriented design.

The Simula67 language [SimulaLanguage], invented by OleJohanDahl and KristenNygaard, was the first ObjectOrientedProgramming language - Simula did not invent the object oriented approach – it implemented objects and so on. The original drafts called objects ‘processes’, but Dahl decided to use the word “object” about a year before the release of Simula67 (which was in 1968 of course!). The modern terminology of classes and inheritance was all there in the language. The term Object was possibly coined by IvanSutherland before this for something very similar, and the term ObjectOriented was probably coined by AlanKay (see AlanKayOnObjects and HeDidntInventTheTerm for further discussion).

In terms of programming (as opposed to design), an OOP language is usually taken to be one which supports

- Encapsulation
- Inheritance (or delegation)
- PolyMorphism

Warning!! Warning!! This definition is considered controversial in some quarters. For instance, Encapsulation is not part of how Perl implements OO; some people might want to add classes to the mix, but that excludes prototype-based languages, like SelfLanguage, which definitely ‘feel’ OO. In particular, features commonly added and removed from that definition include:

- Object Identity
- Automated memory management (not in CeePlusPlus)

- Classes (how is this different from Encapsulation?)
- Abstract Data Types
- Mutable state

At an **absolute minimum**, “objectness” represents combining state with behavior. (Attributes/properties with methods. Or just slots as in *Self* – the distinction between attributes and methods is artificial and a Bad Thing if you ever exposed an attribute and want to change it to a method.)

See PrototypeBasedProgramming.

I'd love to try a WikiWeightedVote on the above. If two others agree here, we'll set it up.

Encapsulation: (1) attributes of an object can be seen or modified only by calling the object's methods; or (2) implementation of the object can change (within reasonable bounds) without changing the interface visible to callers.

The 2nd is the stronger definition, arguing for minimizing the number of “getter/setter” “properties” exported by the object, substituting “logical business interfaces to functionality” instead.

*The distinction I learned was between **encapsulation** as def. (2) and **data hiding** as def. (1), the more restrictive version in which encapsulation is enforced rather than advised. Encapsulation(2) is OOP-universal, as far as I can tell, as the primary abstractive technique—making the difference between, exempli gratia, a CeePlusPlus class and a CeeLanguage struct. Data hiding(1) is rejected by PythonLanguage and others. I personally don't mind it as a simplification tool as long as it isn't too hard to bust (RubyLanguage). – J. Jakes-Schauer*

Inheritance (or delegation): Definition of one object extends/changes that of another.

COM & CORBA camps would argue that this is an implementation issue not fundamental to “OO-ness.”

(However, it's a very convenient feature for code reuse, particularly in frameworks.)

Better make that COM, CORBA, and Ada. The AdaLanguage, before it supported inheritance, was generally considered to be an object-based language. This terminology was used to differentiate it from the object-oriented languages. See ObjectBasedProgramming. – DavidMcReynolds

Polymorphism: The ability to work with several different kinds of objects as if they were the same.

COM, CORBA, Java: Interfaces.

C++: “virtual” functions. (Or “interfaces” with Abstract-BaseClass-es and “pure” virtual functions (“=0” syntax).)

does this not apply to Java abstract methods, too? [It applies to all object methods in Java.]

Smalltalk: Run-time dispatching based on method name. *(Any two classes that implement a common set of method names, with compatible "meanings" are polymorphic. No inheritance relationship is needed.)*

Discussion:

I don’t believe Java Interfaces have anything to do with Polymorphism, and “Interface” in CORBA is synonymous with “Class” in other ObjectOriented languages. Any two classes that share a common parent and provide alternate implementations of a single method defined on that parent are “polymorphic”. AlanKay describes the polymorphic style as “CallBy-Desire” (as opposed to CallByValue and CallByName). – TomStambaugh

Just because you don’t believe it doesn’t mean it’s not so. Java interfaces allow for pure polymorphism (independent of inheritance). They do also have other uses, and you may never use them as a mechanism for including polymorphism in your designs, but the mechanism is there.

Absolutely right for most OO implementation languages, like C++. *But...*

In COM and Java, any two objects implementing the same interface are (for the most part) “substitutable” – you can use one in place of the other. COM has no implementation inheritance, so “sharing a common parent” (other than IUnknown) is not meaningful (unless one starts talking about a particular

implementation language, like C++). Java classes are limited to single inheritance, but may implement any number of interfaces. So Java tends to use interfaces over class inheritance for polymorphism.

Yes, CORBA “interfaces” tend to also mean “classes.” But, two CORBA servers can implement the same CORBA interface, to achieve polymorphism. – JeffGrigg

CORBA interfaces support polymorphism as well as COM or Java interfaces, and references to CORBA interfaces are just as “substitutable”. You can have as many different implementations of a CORBA interface as you want. I don’t think one-to-one correspondences between interface and implementation are any more prevalent in CORBA than they are in any other object-oriented system. (Or am I missing something here?) – KrisJohnson

I’d say that the two required features of an ObjectOrientedLanguage would be Encapsulation & Polymorphism. Inheritance is just how some languages provide Polymorphism. – JonStrayer

See “IsJavaObjectOriented?” “Irrelevant” Questions of this type, and especially on a page with the title ObjectOrientedProgramming, are always relevant. Not “Popular” with some, but worth investigation by one considering the language for use.

This exclusivist discussion does great disservice to the OO community. Really, everybody wants OO to be his side of the story.

Will anyone agree that there are several OO models, theories? And will anyone bother to give references to actual OO theory instead of going round in circles with Booch, Rumbaugh and the likes.

And by the way, the distinction between `button.click()` and `click(button)` is really naive. “Syntactic and ideological”, this gives feed to those who sustain that OO has nothing to do with ComputerScience, when in fact only a part of OO folks don’t want to depart from overly simplistic intuition about what they think OO should be. – AnonymousCoward

I propose that ProgrammingIsInTheMind, and that ObjectOrientedProgramming is one of many mental models that assists a programmer's thinking. I agree that dogmatic statements about what is and what is not OO are generally useless. I think a more useful discussion would endeavor to discover the various concepts used by programmers, then evaluate them for their costs and benefits and how they interact with each other. As I said elsewhere, I think the various concepts are better modeled on a continuous scale rather than a boolean (Is/Not OO, Good/Bad, Like/Dislike). – RobWilliams

I will add my vote to Rob's statements. As soon as you invent a label, in this case OO, you have created a church/institution. That leads to territorial squabbles. It also leads to ideology, historical revision and denial. I don't need all that baggage, I've got a job to do. – RichardHenderson.

On a continuous scale, the statements have a boolean ring, i.e: "generally useless" "concepts are better". It is hard to avoid boolean when considering use/not use. Let alone in programming constructs within OO programming such as If, ifElse, for, do, case, and While.

It might help to think in terms of what OO tries to achieve, what was the motivation for developing OO languages? Maybe that will shed some light on what OO is.

For one, I agree that programming is in the mind. Because OO is a step closer to the way humans think (taxonomy is a good example), it helps us design and write programs that are easier to understand, promote code reuse, cohesion, and decoupling, and data protection.

While hierarchical taxonomies might model the way some people think (every brain is different I would note), I do not find them very reflective of the way the real world changes over time. See LimitsOfHierarchies. I think set theory is a better classification approach than trees in most cases, but implementing it tends to lead to RelationalDatabase-like techniques, which leads us to the HolyWar of the ObjectRelationalPsychologicalMismatch.

[Top, get hierarchies off your brain, OO isn't about hier-

archies, that's just an implementation detail to help organize code.]

I was responding to the "taxonomy" statement, NOT saying that OO is all about trees.

[But that's not why we like objects. OO is about modelling the problem with "things", intelligent "things", objects. Objects, not Classes, but object instances, say for example aPerson and aCompany are easy to reason about, it's easier for most people to reason about them by imagining a conversation between them, what might those objects say to each other or do to each other... maybe aCompany.hires(aPerson), or aCompany.fires(aPerson) or aPerson.quits(aCompany), or questions like aCompany.employs(aPerson) or aPerson.worksAt(aCompany) or aCompany.isOwnedBy(aPerson) or aPerson.isCompatible(anotherPerson) or aCompany.isOwnedBy(anotherCompany)... or aCompany.employs(anotherCompany)...oops, the last two messages were used twice, but that's ok, polymorphism picks the correct method for isOwnedBy depending on whether I pass in a company or a person, same with employs, one message, but it works for several different contexts, just like real human language. If that were procedural, I'd have to remember a far greater number of language elements to accomplish the same work, and I'd have to remember where they're at and what it's the generic recordsets that I'm passing around instead of objects, which modules to include.]

- Why not model such relationships with relational modeling? Then we don't have to wade through bulky code to see the relationships, and it is more consistently implemented from developer to developer. And we can query it more easily than we could we that kind of crud hard-wired into code. CantEncapsulateLinks. – top

Why don't you use a declarative/logic programming language then? Owns(person, company). Owns(company, company). Are-compatible(person1, person2). It's funny to see you folks talk about abstraction and polymorphism while getting all excited about syntactic sugar for physical pointers.

OO is not just syntax, it is also an approach. It is possible to implement OO principles without using OO syntax. The syntax is just there to encourage the principles. For an example of OO practices without syntax, look at <http://api.drupal.org/api/file/developer/topics/oop.html/6>.

[I just explained why... Owns(person, company) forces me to remember that Owns is a valid command, if there are 300 messages in the system I must remember them all, and I simply can't remember that many commands, whereas aCompany.Owns(aPerson) forces me to remember nothing... I simply look at company and see what it can say to aPerson, or I look at person and see what it can say to aCompany. That little syntactic sugar of the . is everything, it provides scope to the set of messages an object can respond to and removes the burden of remembering those messages from me. There is a huge difference between method(arg, arg) and arg.method(arg) insofar as what the mind must comprehend and remember. It may all compile down to the same stuff, but who care's, its all 1's and 0's anyway. But at the language level that little . makes a massive difference in removing complexity from the programmers mind.]

And as ChrisDate, FabianPascal et al. say, a relational DBMS is a tool for organizing facts (i.e. declarative propositions like above) about a domain. Why would you hand-code your own incompatible and byzantine version of said tool, making it completely dependent on one particular application and architecture? (There are certainly legitimate reasons, I am just wondering what yours happen to be.)

I am not top, but I can smell industry hype when I see it.

[But with objects, one simply remembers what they can say to each other, and if we forget, we know exactly where to look, we already have objects in hand, so you just look at them. It's literally like playing with lego's when you were a kid(not sure if you did). Once you have a few basic parts in hand, then entities of the behavioral portion of the problem domain, we don't really think about the data yet, that comes later. We have all the elements to build the program and solve any problem at hand, we can use the same elements to solve any number of different

problems without having to create new elements. We can fluff out the data later as the need dictates. As behavior is added piece by piece, we add any data required to support that behavior. OO programs have an organization at the object level (not classes) that procedural programs can't. Procedural programs have procedures, OO programs have context sensitive procedures, the same procedure name may have 15 different contexts that it can work in, thus we can simplify our solution language. We could do the same with procedural by having one large procedure that has all 15 contexts in a big switch statement. While this makes adding new procedures easier, it makes adding new contexts to groups of methods damn near impossible. The OO solution allows adding entire new contexts a breeze, and we seem to find that more important and more useful. If the word "Contains" is a message/procedure to determine if one thing contains another, the procedural solution would have a big long Contains method that handled every possible case of input. The OO solution would have a Contains method for each and every different context.]

Why are you constantly bringing up the “procedural” straw-man? I explicitly mentioned declarative languages (think Haskell Language). It's hard to understand your arguments when you conflate “object-oriented” with convenient ways to define abstract data types and polymorphism. The first does not imply the second; they are different concepts.

[I'm not using is as a straw man, it doesn't matter if your procedures support polymorphism, Owns(person, company) Owns(company, company), even if polymorphic is still harder to reason about than aPerson.Owns(aCompany) or aCompany.Owns(aPerson) because I would have to know about the owns functions. Take lisp generic functions for example, very powerful, very flexible, and still, unless you know and memorize all the available functions, you're not likely to know Owns even exists. They are indispensable when you really need multiple dispatch, but in most scenarios...single dispatch is all that's needed and the messaging object.method(arg) is simply easier to reason about because I have something to look at. I can see what messages the object responds too. Generic functions

don't allow that, they make me remember it all. So, in most cases, single dispatch message oriented programming is easier, when it isn't, go for the generic function. I never claimed polymorphism was tied to abstract data types, just that they are easier to work with when tied together for most cases. Simply put...the object provides much needed scope to the list of available messages. If the Owns message only works with people and company, then it should be on either person or company, not living on it's own in a module. I love generic functions, but they aren't the answer in most scenarios. Maybe I'm missing some fundamental concept and we're just miscommunicating, I don't know Haskell, so feel free to correct me if I'm mistaken.]

*There's absolutely **no** difference between "owns aPerson aCompany" and "aPerson # owns aCompany", it is just syntax. In fact, in Haskell I just define operator # in such a way that it transforms the second form into the first. But there's not much point to it, I don't need to delude myself regarding the "naturalness" of some design methodology. This debate on how the member access operator is somehow superior to anything else under the sun is simply stupid.*

[aPerson.Contains(anAddress) is in a different context than aCompany.Contains(anAddress), but using context in this manner takes advantage of our natural social and language skills as humans. I just imagine the objects talking to each other, then it's suddenly clear what little conversation between them would solve the problem.]

[You may not agree, but most people do agree, OO is more natural, we are social animals and personifying our objects allows us to use those innate social skills to easily remember how a large number of objects interact, simply based on context. Language is based on metaphor, abstractions and context, OO provides all these using the same elements as the spoken language used to describe it. Nothing could be more natural than writing code like anOrder.add(anItem), aCustomer.add(anOrder), aCustomer.add(anAddress), anOrder.add(anAddress), add being used in several different contexts, but each one having an en-

tirely different implementation, in a separate method. Most importantly, I don't have to remember that add exists, in a procedural program I'd have to know about add, in an OO program I just ask the object, if it supports add, then I use it. Most find it easier to maintain lots of small simple procedures than to maintain a few giant procedures, and the smaller ones are safer because if you mess up one, it won't affect the rest. I'm sure you'll rebut everything I just said, but you always say you don't understand how we think and why we like OO, and I'm telling you why, I'm not arguing, and don't really need a rebuttal, I'm just saying this is how we think. It's been explained enough times by enough people that I don't think you can honestly say you don't understand how we think, you just have to say you disagree or you don't think that way... that's fine. But this is how we think and you have to understand that by now!]

It seems to me that OO is a logical progression from procedural programming. I started programming with Microsoft Professional Development System for BASIC. After a while (not long, 'cause I became a fanatic), I started designing and programming in ways that, in retrospect, I can see were object-oriented. I got a lot of satisfaction from being able to wrap up code in some high-level abstraction so I could make a procedure call that said, 'doThis(instanceData)' and not have to worry about what was involved at the nuts and bolts level. OO made it possible to wrap up the data too.

In a procedural language, my doThis(instanceData) might consist of several function calls, say doThisPart(instanceData); doThatPart(instanceData); and doSomeOtherPart(instanceData). If I wanted to mimic, so to speak, inheritance in this procedural language, I might write a second top level procedure, say doThis2(instanceData) that consisted of calls, doThis(instanceData); doStillThisOther(instanceData); and so on. One problem is that I don't have virtual functions from doThis(instanceData) that I can override, though I could further refine/break into smaller pieces, doThis(instanceData) and make calls to only the pieces I need.

OO gives us inheritance instead. Why is this approach better? – BryanHoover

See DeltaIsolation for a discussion on procedural “overrides”.

How about the paradigm of “Event-Driven Programming” (yield flow control, write only event handlers, don’t hog the flow, etc.), compared to and contrasted with “Object-Oriented Programming”?

e.g. the best O-O systems also have an event-based runtime system at their core? and the early event-based systems (e.g. Gosling’s original Andrew wm at CMU, etc.) were limited by the lack of a good O-O class hierarchy (e.g. Andrew wm rewrite leading to Andrew BE2, Insets, ATK, meanwhile paralleled with MacApp, then NeXTStep, MFC, and Java).

What exactly is EventDrivenProgramming? The only big difference I see between procedural EDP and OO EDP is the various ways to find or dispatch the various events. One can use polymorphism or MultipleDispatch, switch/case statements, or even table/structure lookups. Each group will probably champion their favorite. – ChrisKoenigsberg

General Principles or Mantra of OO (Not every one will agree with all of these)

- A thing knows by itself how to handle requests given to it
- Common behaviors can be inherited from parents instead of repeated in each child
- Objects are like a small human given a set of responsibilities
- Put behavioral “wrappers” around state or data so that class/object user does not have to know or care whether they are looking at state or behavior

Attempts at defining ObjectOrientedProgramming:

- ObjectsAreDictionaries.

- AclassIsNothingButaCyclicDependency

What is Oo?

- DefinitionsForOo
- OoIsPragmatic.

I don't know what I expected. This page looks like the beginning of a textbook on the subject. I just want to let the community know that I'm interested in OOP – FrankRobinson

Are you suggesting that this topic should be broken up into smaller topics?

There are numerous books and endless chapters on this topic, but I've found that ObjectOrientedProgramming is essentially the use of two activities: SwappingClassesAtRuntime (which some like to call PolyMorphism) and DelegationPattern. Other OO techniques, such as inheritance and encapsulation, support these two. A good ObjectOrientedProgrammer learns to use these effectively, using the tools provided by his or her LanguageOfChoice, even if it happens to be VisualBasic 6. – JoelRosenberger

Re: *"There are numerous books and endless chapters on this topic, but I've found that ObjectOrientedProgramming is essentially the use of two activities: SwappingClassesAtRuntime (which some like to call PolyMorphism) and DelegationPattern."*

One gets widely different "OO essentials" depending on who they ask, I notice.

Which is why debates on the merits of OO tend to meander endlessly. Progress is made when the individual techniques and concepts that are often rolled together in OO programming are considered separately. Then they can be decoupled (some OO languages mix the concepts together in inflexible or awkward ways), studied, and improved upon independently. It is high time for the over-marketed 'OO' phrase to fade away and make room for more clarity and progress in software development.

[That sounds very promising. Have you a pointer to such an orthogonal analysis with such clarity?]

Ah, but that question brings one back to the problem: An analysis of what? Whose notion of OO?

NobodyAgreesOnWhatOOIs

My POV (so far:) is that OOP is wrongly named, and should have been called “Model Orientated Programming”. The term “Object” is too restrictive, and gives the wrong focus when trying to understand what OOP is really all about:

An object is simply a concrete implementation of an abstract model that the programmer has in his or her mind. The methods of the object provide the ability to manipulate that model, thereby mapping that abstract model onto the implementation. Hence the “user programmer” manipulates that model (via the methods) and never has to worry about the implementation.

This idea was derived after I realized that an abstraction is simply a model (in your mind) which you then map onto something real (this provides interesting insight onto the nature of perception which isn’t relevant here). As such, OOP provides a very good “implementation” of abstraction itself.

I don't know how one verifies the objectivity or accuracy of such a statement.

You could start by trying to generalize how you deal with the real world, because the real world requires your mind to make various simplifications (which might be called abstraction)...If you define abstraction as ignoring irrelevant details, then the model is the relevant details, and mapping is the process of associating the model to the thing being abstracted.

It is then clear why the ability to be able to associate ‘internal’ state with procedures creates the essence of OOP – without such state you cannot maintain a model (which by definition always has some state). You can of course explicitly provide a (shared) state to plain old procedures, thus creating a model, but it’s a bit messy...In the case of a Singleton object, such state does not *need* to be implemented using a block of memory (such as a Struct in C or Record in Pascal), but can instead

be provided by global variables.

From such a definition it is clear that OOP's (or should that be MOP's ?:) essential characteristics does not include inheritance. You could argue that it does include encapsulation & polymorphism, but I'll leave that up to the intellectual bean-counters out there (I'm more than happy with just stating that it is model orientated and so provides a good implementation of abstraction itself).

Aside: (Single) inheritance is a cute way of organizing objects, so that you can efficiently share code & data with an incremental design strategy, but it is definitely **not** part of the basic requirements of OOP (MOP). You can imagine more complex (and much less efficient) sharing strategies (multiple inheritance being an unimaginative one). But (single) inheritance is a very easy to understand concept, which clearly doesn't conflict with using models for understanding, and indeed is commonly used in people's day-to-day understanding of the world, so that inheritance should be part of every OOP language even if it isn't actually fundamental. – [ChrisHandley](#)

(Please see [ArgumentsAgainstOop](#) for an extended explanation of what I wrote here.)

More attention should be paid to the “oriented” part of object-oriented. I see objects as just a way of wrapping manipulable values in a certain kind of high-level construct. This is a good thing, and should be done in any high-level language that can afford the speed hit. Whether a language or programming style should be oriented towards those objects is a separate issue. I think it it doesn't give enough props to the verb, but that's just my opinion.

So, by way of example, Smalltalk is object-oriented. Python is object-based. Java, less object-based, because classes aren't objects and there are those primitive types; but more object-oriented than Python, because all functions are methods. Lisp, C++, not object-oriented or object-based; they just provide objects as one of their data types.

– Brock Filer

See Also: [LearningObjectOrientedProgramming](#), [BenefitsOfOo](#), [ArgumentsAgainstOop](#), [AllObjectsAreNotCreated-Equal](#), [AlternateObjectOrientedProgrammingView](#), [ProgrammingParadigm](#)

[CategoryObjectOrientation](#) [CategoryDataOrientation](#)

tmux

Note, C- means Ctrl, e.g., C-b is Ctrl-b.

| | |
|------------------------------------------|-------------------------------|
| <code>tmux</code> | start new session |
| <code>tmux new -s [name]</code> | start new session with name |
| <code>tmux attach</code> | attach to most recent session |
| <code>tmux attach -t [name]</code> | attach to named session |
| <code>tmux ls</code> | list sessions |
| <code>tmux kill-session -t [name]</code> | kill a named session |
| | |
| <code>C-b c</code> | create window |
| <code>C-b w</code> | list windows |
| <code>C-b n</code> | next window |
| <code>C-b p</code> | previous window |
| <code>C-b ,</code> | set name for window |
| <code>C-b &</code> | kill window |
| <code>C-b d</code> | detach |
| | |
| <code>C-b %</code> | split pane (vsplit) |
| <code>C-b "</code> | split pane (hsplit) |
| <code>C-b o</code> | switch panes |
| <code>C-b x</code> | kill pane |
| <code>C-b +</code> | change pane into window |
| <code>C-b -</code> | restore pane from window |
| <code>C-b [</code> | copy and scrollbar mode |
| <code>C-b ?</code> | show key commands |

Why do my compiles take so long?

From http://www.stroustrup.com/bs_faq2.html

You may have a problem with your compiler. It may be old, you may have it installed wrongly, or your computer might be an antique. I can't help you with such problems.

However, it is more likely that the program that you are trying to compile is poorly designed, so that compiling it involves the compiler examining hundreds of header files and tens of thousands of lines of code. In principle, this can be avoided. If this problem is in your library vendor's design, there isn't much you can do (except changing to a better library/vendor), but you can structure your own code to minimize re-compilation after changes. Designs that do that are typically better, more maintainable, designs because they exhibit better separation of concerns.

Consider a classical example of an object-oriented program:

```
class Shape {
public: // interface to users of Shapes
    virtual void draw() const;
    virtual void rotate(int degrees);
    // ...
protected: // common data (for implementers of Shapes)
    Point center;
    Color col;
    // ...
};

class Circle : public Shape {
public:
    void draw() const;
    void rotate(int) { }
    // ...
protected:
    int radius;
    // ...
};

class Triangle : public Shape {
```

```

public:
    void draw() const;
    void rotate(int);
    // ...
protected:
    Point a, b, c;
    // ...
};

```

The idea is that users manipulate shapes through Shape’s public interface, and that implementers of derived classes (such as Circle and Triangle) share aspects of the implementation represented by the protected members.

There are three serious problems with this apparently simple idea:

- It is not easy to define shared aspects of the implementation that are helpful to all derived classes. For that reason, the set of protected members is likely to need changes far more often than the public interface. For example, even though “center” is arguably a valid concept for all Shapes, it is a nuisance to have to maintain a point “center” for a Triangle – for triangles, it makes more sense to calculate the center if and only if someone expresses interest in it.
- The protected members are likely to depend on “implementation” details that the users of Shapes would rather not have to depend on. For example, many (most?) code using a Shape will be logically independent of the definition of “Color”, yet the presence of Color in the definition of Shape will probably require compilation of header files defining the operating system’s notion of color.
- When something in the protected part changes, users of Shape have to recompile – even though only implementers of derived classes have access to the protected members.

Thus, the presence of “information helpful to implementers” in the base class that also acts as the interface to users is the

source of instability in the implementation, spurious recompilation of user code (when implementation information changes), and excess inclusion of header files into user code (because the “information helpful to implementers” needs those headers). This is sometimes known as the “brittle base class problem.”

The obvious solution is to omit the “information helpful to implementers” for classes that are used as interfaces to users. That is, to make interfaces, pure interfaces. That is, to represent interfaces as abstract classes:

```
class Shape {
public: // interface to users of Shapes
    virtual void draw() const = 0;
    virtual void rotate(int degrees) = 0;
    virtual Point center() const = 0;
    // ...

    // no data
};

class Circle : public Shape {
public:
    void draw() const;
    void rotate(int) { }
    Point center() const { return cent; }
    // ...
protected:
    Point cent;
    Color col;
    int radius;
    // ...
};

class Triangle : public Shape {
public:
    void draw() const;
    void rotate(int);
    Point center() const;
    // ...
protected:
    Color col;
    Point a, b, c;
    // ...
};
```

The users are now insulated from changes to implementations

of derived classes. I have seen this technique decrease build times by orders of magnitudes.

But what if there really is some information that is common to all derived classes (or simply to several derived classes)? Simply make that information a class and derive the implementation classes from that also:

```
class Shape {
public: // interface to users of Shapes
    virtual void draw() const = 0;
    virtual void rotate(int degrees) = 0;
    virtual Point center() const = 0;
    // ...

    // no data
};

struct Common {
    Color col;
    // ...
};

class Circle : public Shape, protected Common {
public:
    void draw() const;
    void rotate(int) { }
    Point center() const { return cent; }
    // ...
protected:
    Point cent;
    int radius;
};

class Triangle : public Shape, protected Common {
public:
    void draw() const;
    void rotate(int);
    Point center() const;
    // ...
protected:
    Point a, b, c;
};
```

Numbers every programmer should know

From <https://gist.github.com/hellerbarde/2843375>

| | | |
|------------------------------------|----------------|-------------|
| L1 cache reference | 0.5 ns | |
| Branch mispredict | 5 ns | |
| L2 cache reference | 7 ns | |
| Mutex lock/unlock | 25 ns | |
| Main memory reference | 100 ns | |
| Compress 1K bytes with Zippy | 3,000 ns | 3 μ s |
| Send 2K bytes over 1 Gbps network | 20,000 ns | 20 μ s |
| SSD random read | 150,000 ns | 150 μ s |
| Read 1 MB sequentially from memory | 250,000 ns | 25 μ s |
| Round trip within same datacenter | 500,000 ns | 0.5 ms |
| Read 1 MB sequentially from SSD | 1,000,000 ns | 1 ms |
| Disk seek | 10,000,000 ns | 10 ms |
| Read 1 MB sequentially from disk | 20,000,000 ns | 20 ms |
| Send packet CA→Netherlands→CA | 150,000,000 ns | 150 ms |

Multiply all values by 1 billion:

| | | |
|---------------------------|------------|----------------------------|
| L1 cache reference | 0.5 s | one heart beat |
| Branch mispredict | 5 s | yawn |
| L2 cache reference | 7 s | long yawn |
| Mutex lock/unlock | 25 s | making a coffee |
| Main memory reference | 100 s | brushing your teeth |
| Compress 1K bytes | 50 min | TV show |
| Send 2K bytes over 1 Gbps | 5.5 hr | half a work day |
| SSD random read | 1.7 days | weekend |
| Read 1 MB from memory | 2.9 days | long weekend |
| Round trip in datacenter | 5.8 days | vacation |
| Read 1 MB from SSD | 11.6 days | long vacation |
| Disk seek | 16.5 weeks | a semester in university |
| Read 1 MB from disk | 7.8 months | nearly length of pregnancy |
| Send packet CA→NL→CA | 4.8 years | time to bachelor's degree |

Editor war (Wikipedia)

Editor war is the common name for the rivalry between users of the Emacs and vi (Vim) text editors. The rivalry has become a lasting part of hacker culture and the free software community.

The Emacs vs vi debate was one of the original “holy wars” conducted on Usenet groups, with many flame wars fought between those insisting that their editor of choice is the paragon of editing perfection, and insulting the other, since at least 1985. Related battles have been fought over operating systems, programming languages, version control systems, and even source code indent style. Notably, unlike other wars (ie, UNIX vs ITS vs VMS, C vs Pascal vs Fortran), the editor war has yet to be resolved with a clear winner, and the hacker community remains split roughly 50/50.

Differences between vi and Emacs

The most important differences between vi and Emacs are presented in the following list.

Keystroke execution

vi retains each permutation of typed keys. This creates a path in the decision tree which unambiguously identifies any command.

Emacs commands are key combinations for which modifier keys are held down while other keys are pressed; a command gets executed once completely typed. This still forms a decision tree of commands, but not one of individual keystrokes. A vim-inspired Emacs package (undo-tree) provides a user interface to the tree.

Memory usage and customizability

vi is a smaller and faster program, but with less capacity for customization. The vim version of vi has evolved to provide signif-

icantly more functionality and customization than vi, making it comparable to Emacs. vi start-up time is near instantaneous for small text files, while vim is almost as fast.

Emacs executes many actions on startup, many of which may execute arbitrary user code. This makes Emacs take longer to start up (even compared to vim) and require more memory. However, it is highly customizable and includes a large number of features, as it is essentially an execution environment for a Lisp program designed for text-editing. Emacs 18 (released in 1987) introduced a server mode designed to run continuously in the background. Various instances of Emacs can then be started in client mode, attaching to this server and sharing state. Emacs client startup time is practically instantaneous as all it does is provoke the existing Emacs process to redraw the display.

User environment

vi was originally exclusively used inside of a text-mode console, offering no graphical user interface (GUI). Most modern vi derivatives, e.g. MacVim and gVim, include GUIs. However, support for proportionally spaced fonts remains absent. Also lacking is support for different sized fonts in the same document.

Emacs, while also initially designed for use on a console, grew a TUI fairly early on due to its Lisp machine heritage. X11 GUI support was added in Emacs 18, and made the default in version 19. Current Emacs GUIs include full support for proportionate spacing and font-size variation. Emacs also supports embedded images and hypertext.

Function/navigation interface

vi uses distinct editing modes. Under “insert mode”, keys insert characters into the document. Under “normal mode” (also known as “command mode”), bare keypresses execute vi commands.

Emacs uses metakey chords. Keys or key chords can be defined as prefix keys, which put Emacs into a mode where

it waits for additional key presses that constitute a key binding. Key bindings can be mode-specific, further customizing the interaction style. Emacs provides a command line accessed by M-x that can be configured to autocomplete in various ways. Emacs also provides a defalias macro, allowing alternate names for commands.

Keyboard

vi uses no <Alt> key and seldom uses the <Ctrl> key. **vi**'s keyset is mainly restricted to the alphanumeric keys, and the escape key. This is an enduring relic of its teletype heritage, but has the effect of making most of **vi**'s functionality accessible without frequent awkward finger reaches.

The expansion of one of **Emacs**' backronyms is Escape, Meta, Alt, Control, Shift, which neatly summarizes most of the modifier keys it uses, only leaving out Super. Emacs was developed on Lisp Machines with Space-cadet keyboards that were more ergonomic with respect to modifiers than modern layouts. There are multiple emacs packages, such as spacemacs or ergoemacs that replace these key combinations with ones easier to type, or customization can be done ad hoc by the user.

Language and script support

vi has rudimentary support for languages other than English. Vim is partially multilingual, with support for European, Arabic, Hebrew, and Far East Asian language support only. Notably, Indic language and script support is absent.

Emacs has full support for all Unicode-compatible writing systems and allows multiple scripts to be freely intermixed.

Benefits of Emacs

- Emacs has a non-modal interface
- One of the most ported computer programs. It runs in text mode and under graphical user interfaces on a wide variety of operating systems, including most Unix-like

systems (Linux, the various BSDs, Solaris, AIX, IRIX, macOS, etc.), MS-DOS, Microsoft Windows, AmigaOS, and OpenVMS. Unix systems, both free and proprietary, frequently provide Emacs bundled with the operating system.

- Emacs server architecture allows multiple clients to attach to the same Emacs instance and share the buffer list, kill ring, undo history and other state.
- Pervasive online help system with keybindings, functions and commands documented on the fly.
- Extensible and customizable Lisp programming language variant (Emacs Lisp), with features that include:
 - Ability to emulate vi and vim (using Evil, Viper or Vimpulse).
 - A powerful and extensible file manager (dired), integrated debugger, and a large set of development and other tools.
 - Having every command be an Emacs Lisp function enables commands to DWIM (Do What I Mean) by programmatically responding to past actions and document state. For example, a switch-or-split-window command could switch to another window if one exists, or create one if needed. This cuts down on the number of keystrokes and commands a user must remember.
 - “An OS inside an OS”. Emacs Lisp enables Emacs to be programmed far beyond editing features. Even a base install contains several dozen applications, including two web browsers, news readers, several mail agents, four IRC clients, a version of ELIZA, and a variety of games. All of these applications are available anywhere Emacs runs, with the same user interface and functionality. Starting with version 24, Emacs includes a package manager, making it easy to install additional applications includ-

ing alternate web browsers, EMMS (Emacs Multi-media System), and more. Also available are numerous packages for programming, including some targeted at specific language/library combinations or coding styles.

Benefits of vi-like editors

- Edit commands are composable
- Vi has a modal interface (which Emacs can emulate)
- Historically, vi loads faster than Emacs.
- While deeply associated with UNIX tradition, it runs on all systems that can implement the standard C library, including UNIX, Linux, AmigaOS, DOS, Windows, Mac, BeOS, OpenVMS, IRIX, AIX, HP-UX, BSD and POSIX-compliant systems.
- Extensible and customizable through VimScript or APIs for interpreted languages such as Python, Ruby, Perl, and Lua
- Ubiquitous. Essentially all Unix and Unix-like systems come with vi (or a variant) built-in. Vi (and ex, but not vim) is specified in the POSIX standard.
- System rescue environments, embedded systems (notably those with busybox) and other constrained environments often include vi, but not emacs.

Today

In the past, many small editors modeled after or derived from vi flourished. This was due to the importance of conserving memory with the comparatively minuscule amount available at the time. As computers have become more powerful, many vi clones, Vim in particular, have grown in size and code complexity. These vi variants of today, as with the old lightweight

Emacs variants, tend to have many of the perceived benefits and drawbacks of the opposing side. For example, Vim without any extensions requires about ten times the disk space required by vi, and recent versions of Vim can have more extensions and run slower than Emacs. In *The Art of Unix Programming*, Eric S. Raymond called Vim's supposed light weight when compared with Emacs "a shared myth." Moreover, with the large amounts of RAM in modern computers, both Emacs and vi are lightweight compared to large integrated development environments such as Eclipse, which tend to draw derision from Emacs and vi users alike.

Tim O'Reilly said, in 1999, that O'Reilly Media's tutorial on vi sells twice as many copies as that on Emacs (but noted that Emacs came with a free manual). Many programmers use either Emacs and vi or their various offshoots, including Linus Torvalds who uses MicroEMACS. Also in 1999, vi creator Bill Joy said that vi was "written for a world that doesn't exist anymore" and stated that Emacs was written on much more capable machines with faster displays so they could have "funny commands with the screen shimmering and all that, and meanwhile, I'm sitting at home in sort of World War II surplus housing at Berkeley with a modem and a terminal that can just barely get the cursor off the bottom line."

In addition to Emacs and vi workalikes, pico and its free and open source clone nano and other text editors such as ne often have their own third-party advocates in the editor wars, though not to the extent of Emacs and vi.

As of 2014, both Emacs and vi can lay claim to being among the longest-lived application programs of all time, as well as being the two most commonly used text editors on Linux and Unix. Many operating systems, especially Linux and BSD derivatives, bundle multiple text editors with the operating system to cater to user demand. For example, a default installation of macOS contains Emacs, ed, nano, TextEdit, and Vim.

Powers of 2

| | | | |
|----|---------------|-----------|-------------|
| 0 | 1 | 1 | 1 |
| 1 | 2 | 2 | 2 |
| 2 | 4 | 4 | 4 |
| 3 | 8 | 8 | 10 |
| 4 | 16 | 10 | 20 |
| 5 | 32 | 20 | 40 |
| 6 | 64 | 40 | 100 |
| 7 | 128 | 80 | 200 |
| 8 | 256 | 100 | 400 |
| 9 | 512 | 200 | 1000 |
| 10 | 1,024 | 400 | 2000 |
| 11 | 2,048 | 800 | 4000 |
| 12 | 4,096 | 1000 | 10000 |
| 13 | 8,192 | 2000 | 20000 |
| 14 | 16,384 | 4000 | 40000 |
| 15 | 32,768 | 8000 | 100000 |
| 16 | 65,536 | 10000 | 200000 |
| 17 | 131,072 | 20000 | 400000 |
| 18 | 262,144 | 40000 | 1000000 |
| 19 | 524,288 | 80000 | 2000000 |
| 20 | 1,048,576 | 100000 | 4000000 |
| 21 | 2,097,152 | 200000 | 10000000 |
| 22 | 4,194,304 | 400000 | 20000000 |
| 23 | 8,388,608 | 800000 | 40000000 |
| 24 | 16,777,216 | 1000000 | 100000000 |
| 25 | 33,554,432 | 2000000 | 200000000 |
| 26 | 67,108,864 | 4000000 | 400000000 |
| 27 | 134,217,728 | 8000000 | 1000000000 |
| 28 | 268,435,456 | 10000000 | 2000000000 |
| 29 | 536,870,912 | 20000000 | 4000000000 |
| 30 | 1,073,741,824 | 40000000 | 10000000000 |
| 31 | 2,147,483,648 | 80000000 | 20000000000 |
| 32 | 4,294,967,296 | 100000000 | 40000000000 |

```
base =: dyad define
syms =. ((48 + i. 10),(65 + i.0>.x-10)) { a.
idxs =. ((1+>.x^.y)#x)#:y
idxs { syms
)
b=:10 16 8
n=:33
(|:(#b),n)$n#b)base"0((n,#b)$(#b)#(2~i.n))
```


Regex

Bash globs

| | |
|----------------------------|-----------------------------------|
| <code>*</code> | match any character, zero or more |
| <code>?</code> | match any one character |
| <code>[abc]</code> | match one of a, b, c |
| <code>[a-z]</code> | match one of range of characters |
| <code>[!abc] [!a-z]</code> | inverse of [...] |

Grep regex

| | |
|----------------------------|----------------------------------------|
| <code>^ \$</code> | match start, end of line |
| <code>.</code> | match any one character |
| <code>? + *</code> | match previous 0/1, 1+, 0+ times |
| <code>[abc] [a-z]</code> | match one of list, range of characters |
| <code>[^abc] [^a-z]</code> | inverse of [...] |

Vim regex

| | |
|----------------------------|------------------------------------------|
| <code>^ \$</code> | match start, end of line |
| <code>.</code> | match any one character except newline |
| <code>\= \+ *</code> | match previous 0/1, 1+, 0+ times |
| <code>[abc] [a-z]</code> | match one of list, range of characters |
| <code>[^abc] [^a-z]</code> | inverse of [...] |
| <code>\(...\)</code> | define group |
| <code>\1 \2</code> | refer to first group, second group, etc. |
| <code>\(abc\ xyz\)</code> | match abc or xyz |

Emacs regex

| | |
|----------------------------|------------------------------------------|
| <code>^ \$</code> | match start, end of line |
| <code>.</code> | match any one character |
| <code>? + *</code> | match previous 0/1, 1+, 0+ times |
| <code>[abc] [a-z]</code> | match one of list, range of characters |
| <code>[^abc] [^a-z]</code> | inverse of [...] |
| <code>\(...\)</code> | define group |
| <code>\1 \2</code> | refer to first group, second group, etc. |
| <code>\(abc\ xyz\)</code> | match abc or xyz |

The Mythical Man-Month (Fred Brooks)

Good cooking takes time. If you are made to wait, it is to serve you better, and to please you.

– Menu of Restaurant Antoine, New Orleans

More software projects have gone awry for lack of calendar time than for all other causes combined. Why is this cause of disaster so common?

First, our techniques of estimating are poorly developed. More seriously, they reflect an unvoiced assumption which is quite untrue, i.e., that all will go well.

Second, our estimating techniques fallaciously confuse effort with progress, hiding the assumption that men and months are interchangeable.

Third, because we are uncertain of our estimates, software managers often lack the courteous stubbornness of Antoine's chef.

Fourth, schedule progress is poorly monitored. Techniques proven and routine in other engineering disciplines are considered radical innovations in software engineering.

Fifth, when schedule slippage is recognized, the natural (and traditional) response is to add manpower. Like dousing a fire with gasoline, this makes matters worse, much worse. More fire requires more gasoline, and thus begins a regenerative cycle which ends in disaster.

Schedule monitoring will be the subject of a separate essay. Let us consider other aspects of the problem in more detail.

Optimism

All programmers are optimists. Perhaps this modern sorcery especially attracts those who believe in happy endings and fairy godmothers. Perhaps the hundreds of nitty frustrations drive away all but those who habitually focus on the end goal. Perhaps it is merely that computers are young, programmers are

younger, and the young are always optimists. But however the selection process works, the result is indisputable: "This time it will surely run," or "I just found the last bug."

So the first false assumption that underlies the scheduling of systems programming is that *all will go well*, i.e., that *each task will take only as long as it "ought" to take*.

The pervasiveness of optimism among programmers deserves more than a flip analysis. Dorothy Sayers, in her excellent book, *The Mind of the Maker*, divides creative activity into three stages: the idea, the implementation, and the interaction. A book, then, or a computer, or a program comes into existence first as an ideal construct, built outside time and space, but complete in the mind of the author. It is realized in time and space, by pen, ink, and paper, or by wire, silicon, and ferrite. The creation is complete when someone reads the book, uses the computer, or runs the program, thereby interacting with the mind of the maker.

This description, which Miss Sayers uses to illuminate not only human creative activity but also the Christian doctrine of the Trinity, will help us in our present task. For the human makers of things, the incompletenesses and inconsistencies of our ideas, become clear only during implementation. Thus it is that writing, experimentation, "working out" are essential disciplines for the theoretician.

In many creative activities the medium of execution is intractable. Lumber splits; paints smear; electrical circuits ring. These physical limitations of the medium constrain the ideas that may be expressed, and they also create unexpected difficulties in the implementation.

Implementation, then, takes time and sweat both because of the physical media and because of the inadequacies of the underlying ideas. We tend to blame the physical media for most of our implementation difficulties; for the media are not "ours" in the way the ideas are, and our pride colors our judgment.

Computer programming, however, creates with an exceedingly tractable medium. The programmer builds from pure thought-stuff: concepts and very flexible representations thereof. Because the medium is tractable, we expect few diffi-

culties in implementation; hence our pervasive optimism. Because our ideas are faulty, we have bugs; hence our optimism is unjustified.

In a single task, the assumption that all will go well has a probabilistic effect on the schedule. It might indeed go as planned, for there is a probability distribution for the delay that will be encountered, and "no delay" has a finite probability. A large programming effort, however, consists of many tasks, some chained end-to-end. The probability that each will go well becomes vanishingly small.

The Man-Month

The second fallacious thought mode is expressed in the very unit of effort used in estimating and scheduling: the man-month. Cost does indeed vary as the product of the number of men and the number of months. Progress does not. *Hence the man-month as a unit for measuring the size of a job is a dangerous and deceptive myth.* It implies that men and months are interchangeable.

Men and months are interchangeable commodities only when a task can be partitioned among many workers *with no communication among them* (Fig. 1). This is true of reaping wheat or picking cotton; it is not even approximately true of systems programming.

When a task cannot be partitioned because of sequential constraints, the application of more effort has no effect on the schedule (Fig. 2). The bearing of a child takes nine months, no matter how many women are assigned. Many software tasks have this characteristic because of the sequential nature of debugging.

In tasks that can be partitioned but which require communication among the subtasks, the effort of communication must be added to the amount of work to be done. Therefore the best that can be done is somewhat poorer than an even trade of men for months (Fig. 3).

The added burden of communication is made up of two parts, training and intercommunication. Each worker must be

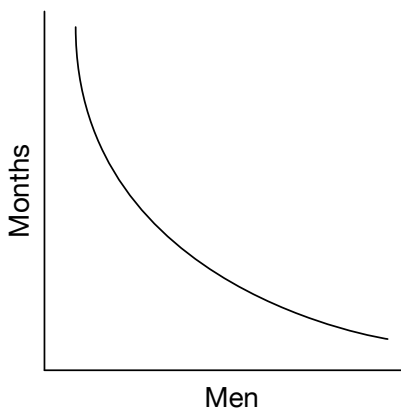


Fig. 1: Time versus number of workers – perfectly partitionable task



Fig. 2: Time versus number of workers – unpartitionable task

trained in the technology, the goals of the effort, the overall strategy, and the plan of work. This training cannot be partitioned, so this part of the added effort varies linearly with the number of workers.

Intercommunication is worse. If each part of the task must



Fig. 3: Time versus number of workers – partitionable task requiring communication

be separately coordinated with each other part, the effort increases as $n(n-1)/2$. Three workers require three times as much pairwise intercommunication as two; four require six times as much as two. If, moreover, there need to be conferences among three, four, etc., workers to resolve things jointly, matters get worse yet. The added effort of communicating may fully counteract the division of the original task and bring us to the situation of Fig. 4.

Since software construction is inherently a systems effort – an exercise in complex interrelationships – communication effort is great, and it quickly dominates the decrease in individual task time brought about by partitioning. Adding more men then lengthens, not shortens, the schedule.

Systems Test

No parts of the schedule are so thoroughly affected by sequential constraints as component debugging and system test. Furthermore, the time required depends on the number and subtlety of the errors encountered. Theoretically this number

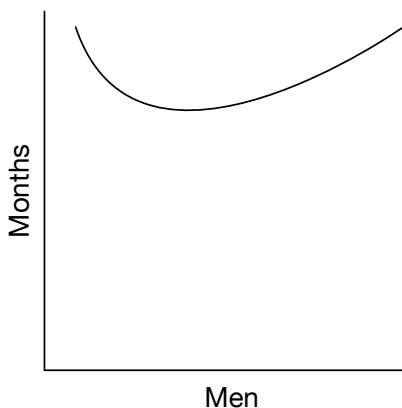


Fig. 4: Time versus number of workers – task with complex interrelationships

should be zero. Because of optimism, we usually expect the number of bugs to be smaller than it turns out to be. Therefore testing is usually the most mis-scheduled part of programming.

For some years I have been successfully using the following rule of thumb for scheduling a software task:

$\frac{1}{3}$ planning

$\frac{1}{6}$ coding

$\frac{1}{4}$ component test and early system test

$\frac{1}{4}$ system test, all components in hand.

This differs from conventional scheduling in several important ways:

1. The fraction devoted to planning is larger than normal. Even so, it is barely enough to produce a detailed and solid specification, and not enough to include research or exploration of totally new techniques.

2. The *half* of the schedule devoted to debugging of completed code is much larger than normal.
3. The part that is easy to estimate, i.e., coding, is given only one-sixth of the schedule.

In examining conventionally scheduled projects, I have found that few allowed one-half of the projected schedule for testing, but that most did indeed spend half of the actual schedule for that purpose. Many of these were on schedule until and except in system testing.

Failure to allow enough time for system test, in particular, is peculiarly disastrous. Since the delay comes at the end of the schedule, no one is aware of schedule trouble until almost the delivery date. Bad news, late and without warning, is unsettling to customers and to managers.

Furthermore, delay at this point has unusually severe financial, as well as psychological, repercussions. The project is fully staffed, and cost-per-day is maximum. More seriously, the software is to support other business effort (shipping of computers, operation of new facilities, etc.) and the secondary costs of delaying these are very high, for it is almost time for software shipment. Indeed, these secondary costs may far outweigh all others. It is therefore very important to allow enough system test time in the original schedule.

Gutless Estimating

Observe that for the programmer, as for the chef, the urgency of the patron may govern the scheduled completion of the task, but it cannot govern the actual completion. An omelette, promised in two minutes, may appear to be progressing nicely. But when it has not set in two minutes, the customer has two choices – wait or eat it raw. Software customers have had the same choices.

The cook has another choice; he can turn up the heat. The result is often an omelette nothing can save – burned in one part, raw in another.

Now I do not think software managers have less inherent courage and firmness than chefs, nor than other engineering managers. But false scheduling to match the patron's desired date is much more common in our discipline than elsewhere in engineering. It is very difficult to make a vigorous, plausible, and job-risking defense of an estimate that is derived by no quantitative method, supported by little data, and certified chiefly by the hunches of the managers.

Clearly two solutions are needed. We need to develop and publicize productivity figures, bug-incidence figures, estimating rules, and so on. The whole profession can only profit from sharing such data.

Until estimating is on a sounder basis, individual managers will need to stiffen their backbones and defend their estimates with the assurance that their poor hunches are better than wish- derived estimates.

Regenerative Schedule Disaster

What does one do when an essential software project is behind schedule? Add manpower, naturally. As Figs. 1 through 4 suggest, this may or may not help.

Let us consider an example. Suppose a task is estimated at 12 man-months and assigned to three men for four months, and that there are measurable mileposts A, B, C, D, which are scheduled to fall at the end of each month (Fig. 5).

Now suppose the first milepost is not reached until two months have elapsed (Fig. 6). What are the alternatives facing the manager?

1. Assume that the task must be done on time. Assume that only the first part of the task was misestimated, so Fig. 6 tells the story accurately. Then 9 man-months of effort remain, and two months, so $4\frac{1}{2}$ men will be needed. Add 2 men to the 3 assigned.
2. Assume that the task must be done on time. Assume that the whole estimate was uniformly low, so that Fig. 7 really describes the situation. Then 18 man-months of ef-

fort remain, and two months, so 9 men will be needed. Add 6 men to the 3 assigned.

3. Reschedule. I like the advice given by P. Fagg, an experienced hardware engineer, "Take no small slips" That is, allow enough time in the new schedule to ensure that the work can be carefully and thoroughly done, and that rescheduling will not have to be done again.
4. Trim the task. In practice this tends to happen anyway, once the team observes schedule slippage. Where the secondary costs of delay are very high, this is the only feasible action. The manager's only alternatives are to trim it formally and carefully, to reschedule, or to watch the task get silently trimmed by hasty design and incomplete testing.

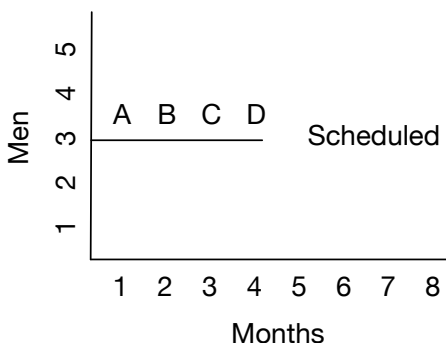


Fig. 5

In the first two cases, insisting that the unaltered task be completed in four months is disastrous. Consider the regenerative effects, for example, for the first alternative (Fig. 8). The two new men, however competent and however quickly recruited, will require training in the task by one of the experienced men. If this takes a month, *3 man-months will have been devoted to work not in the original estimate*. Furthermore, the

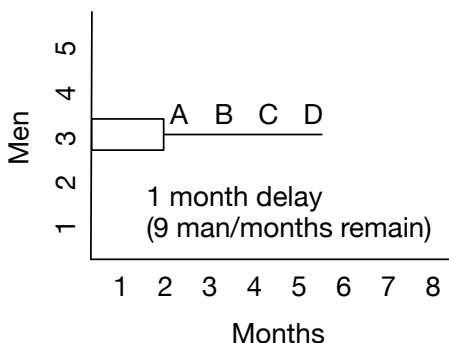


Fig. 6

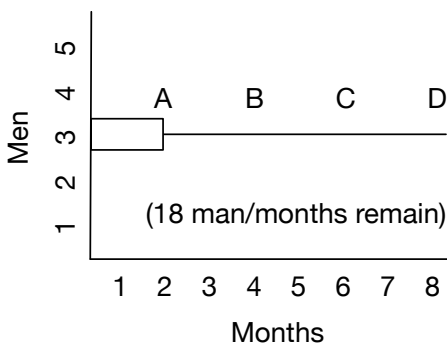


Fig. 7

task, originally partitioned three ways, must be repartitioned into five parts; hence some work already done will be lost, and system testing must be lengthened. So at the end of the third month, substantially more than 7 man-months of effort remain, and 5 trained people and one month are available. As Fig. 8 suggests, the product is just as late as if no one had been added (Fig. 6).

To hope to get done in four months, considering only training time and not repartitioning and extra systems test, would require adding 4 men, not 2, at the end of the second month.

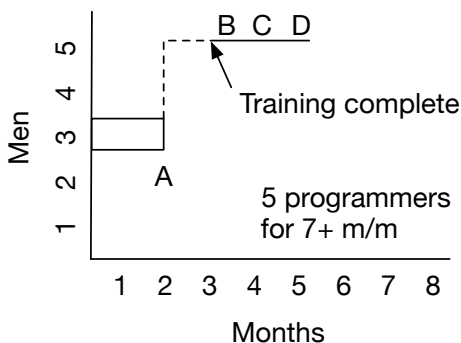


Fig. 8

To cover repartitioning and system test effects, one would have to add still other men. Now, however, one has at least a 7-man team, not a 3-man one; thus such aspects as team organization and task division are different in kind, not merely in degree.

Notice that by the end of the third month things look very black. The March 1 milestone has not been reached in spite of all the managerial effort. The temptation is very strong to repeat the cycle, adding yet more manpower. Therein lies madness.

The foregoing assumed that only the first milestone was misestimated. If on March 1 one makes the conservative assumption that the whole schedule was optimistic, as Fig. 7 depicts, one wants to add 6 men just to the original task. Calculation of the training, repartitioning, system testing effects is left as an exercise for the reader. Without a doubt, the regenerative disaster will yield a poorer product, later, than would rescheduling with the original three men, unaugmented.

Oversimplifying outrageously, we state Brooks's Law:

Adding manpower to a late software project makes it later.

This then is the demythologizing of the man-month. The number of months of a project depends upon its sequential con-

straints. The maximum number of men depends upon the number of independent subtasks. From these two quantities one can derive schedules using fewer men and more months. (The only risk is product obsolescence.) One cannot, however, get workable schedules using more men and fewer months. More software projects have gone awry for lack of calendar time than for all other causes combined.

The Story of Mel

This was posted to Usenet by its author, Ed Nather (<nather@astro.as.utexas.edu>), on May 21, 1983.

A recent article devoted to the macho side of programming made the bald and unvarnished statement:

Real Programmers write in FORTRAN.

Maybe they do now,
in this decadent era of
Lite beer, hand calculators, and “user-friendly” software
but back in the Good Old Days,
when the term “software” sounded funny
and Real Computers were made out of drums and vacuum tubes,
Real Programmers wrote in machine code.
Not FORTRAN. Not RATFOR. Not, even, assembly language.
Machine Code.
Raw, unadorned, inscrutable hexadecimal numbers.
Directly.

Lest a whole new generation of programmers
grow up in ignorance of this glorious past,
I feel duty-bound to describe,
as best I can through the generation gap,
how a Real Programmer wrote code.
I'll call him Mel,
because that was his name.

I first met Mel when I went to work for Royal McBee Computer Corp.,
a now-defunct subsidiary of the typewriter company.
The firm manufactured the LGP-30,
a small, cheap (by the standards of the day)
drum-memory computer,
and had just started to manufacture
the RPC-4000, a much-improved,
bigger, better, faster – drum-memory computer.

Cores cost too much,
and weren't here to stay, anyway.
(That's why you haven't heard of the company,
or the computer.)

I had been hired to write a FORTRAN compiler
for this new marvel and Mel was my guide to its wonders.
Mel didn't approve of compilers.

"If a program can't rewrite its own code",
he asked, "what good is it?"

Mel had written,
in hexadecimal,
the most popular computer program the company owned.
It ran on the LGP-30
and played blackjack with potential customers
at computer shows.

Its effect was always dramatic.
The LGP-30 booth was packed at every show,
and the IBM salesmen stood around
talking to each other.
Whether or not this actually sold computers
was a question we never discussed.

Mel's job was to re-write
the blackjack program for the RPC-4000.
(Port? What does that mean?)
The new computer had a one-plus-one
addressing scheme,
in which each machine instruction,
in addition to the operation code
and the address of the needed operand,
had a second address that indicated where, on the revolving drum,
the next instruction was located.

In modern parlance,
every single instruction was followed by a GO TO!
Put that in Pascal's pipe and smoke it.

Mel loved the RPC-4000

because he could optimize his code:
that is, locate instructions on the drum
so that just as one finished its job,
the next would be just arriving at the “read head”
and available for immediate execution.
There was a program to do that job,
an “optimizing assembler”,
but Mel refused to use it.

“You never know where it’s going to put things”,
he explained, “so you’d have to use separate constants”.

It was a long time before I understood that remark.
Since Mel knew the numerical value
of every operation code,
and assigned his own drum addresses,
every instruction he wrote could also be considered
a numerical constant.
He could pick up an earlier “add” instruction, say,
and multiply by it,
if it had the right numeric value.
His code was not easy for someone else to modify.

I compared Mel’s hand-optimized programs
with the same code massaged by the optimizing assembler program,
and Mel’s always ran faster.
That was because the “top-down” method of program design
hadn’t been invented yet,
and Mel wouldn’t have used it anyway.
He wrote the innermost parts of his program loops first,
so they would get first choice
of the optimum address locations on the drum.
The optimizing assembler wasn’t smart enough to do it that way.

Mel never wrote time-delay loops, either,
even when the balky Flexowriter
required a delay between output characters to work right.
He just located instructions on the drum
so each successive one was just past the read head
when it was needed;
the drum had to execute another complete revolution

to find the next instruction.
He coined an unforgettable term for this procedure.
Although “optimum” is an absolute term,
like “unique”, it became common verbal practice
to make it relative:
“not quite optimum” or “less optimum”
or “not very optimum”.
Mel called the maximum time-delay locations
the “most pessimum”.

After he finished the blackjack program
and got it to run
(“Even the initializer is optimized”,
he said proudly),
he got a Change Request from the sales department.
The program used an elegant (optimized)
random number generator
to shuffle the “cards” and deal from the “deck”,
and some of the salesmen felt it was too fair,
since sometimes the customers lost.
They wanted Mel to modify the program
so, at the setting of a sense switch on the console,
they could change the odds and let the customer win.

Mel balked.
He felt this was patently dishonest,
which it was,
and that it impinged on his personal integrity as a programmer,
which it did,
so he refused to do it.
The Head Salesman talked to Mel,
as did the Big Boss and, at the boss’s urging,
a few Fellow Programmers.
Mel finally gave in and wrote the code,
but he got the test backwards,
and, when the sense switch was turned on,
the program would cheat, winning every time.
Mel was delighted with this,
claiming his subconscious was uncontrollably ethical,
and adamantly refused to fix it.

After Mel had left the company for greener pasture\$,
the Big Boss asked me to look at the code
and see if I could find the test and reverse it.
Somewhat reluctantly, I agreed to look.
Tracking Mel's code was a real adventure.

I have often felt that programming is an art form,
whose real value can only be appreciated
by another versed in the same arcane art;
there are lovely gems and brilliant coups
hidden from human view and admiration, sometimes forever,
by the very nature of the process.
You can learn a lot about an individual
just by reading through his code,
even in hexadecimal.
Mel was, I think, an unsung genius.

Perhaps my greatest shock came
when I found an innocent loop that had no test in it.
No test. None.
Common sense said it had to be a closed loop,
where the program would circle, forever, endlessly.
Program control passed right through it, however,
and safely out the other side.
It took me two weeks to figure it out.

The RPC-4000 computer had a really modern facility
called an index register.
It allowed the programmer to write a program loop
that used an indexed instruction inside;
each time through,
the number in the index register
was added to the address of that instruction,
so it would refer
to the next datum in a series.
He had only to increment the index register
each time through.
Mel never used it.

Instead, he would pull the instruction into a machine register,
add one to its address,

and store it back.

He would then execute the modified instruction
right from the register.

The loop was written so this additional execution time
was taken into account –

just as this instruction finished,
the next one was right under the drum's read head,
ready to go.

But the loop had no test in it.

The vital clue came when I noticed
the index register bit,
the bit that lay between the address
and the operation code in the instruction word,
was turned on –

yet Mel never used the index register,
leaving it zero all the time.

When the light went on it nearly blinded me.

He had located the data he was working on
near the top of memory –
the largest locations the instructions could address –
so, after the last datum was handled,
incrementing the instruction address
would make it overflow.

The carry would add one to the
operation code, changing it to the next one in the instruction set:
a jump instruction.

Sure enough, the next program instruction was
in address location zero,
and the program went happily on its way.

I haven't kept in touch with Mel,
so I don't know if he ever gave in to the flood of
change that has washed over programming techniques
since those long-gone days.

I like to think he didn't.

In any event,

I was impressed enough that I quit looking for the
offending test,

telling the Big Boss I couldn't find it.

He didn't seem surprised.

When I left the company,
the blackjack program would still cheat
if you turned on the right sense switch,
and I think that's how it should be.
I didn't feel comfortable
hacking up the code of a Real Programmer.

[1992 postscript – the author writes: “The original submission to the net was not in free verse, nor any approximation to it – it was straight prose style, in non-justified paragraphs. In bouncing around the net it apparently got modified into the ‘free verse’ form now popular. In other words, it got hacked on the net. That seems appropriate, somehow.” The author adds that he likes the ‘free-verse’ version better than his prose original. . .]

[1999 update: Mel's last name is now known. The manual for the LGP-30 refers to “Mel Kaye of Royal McBee who did the bulk of the programming [...] of the ACT 1 system”.]

[2001: The Royal McBee LPG-30 turns out to have one other claim to fame. It turns out that meteorologist Edward Lorenz was doing weather simulations on an LGP-30 when, in 1961, he discovered the “Butterfly Effect” and computational chaos. This seems, somehow, appropriate.]

[2002: A copy of the programming manual for the LGP-30 lives at <http://ed-thelen.org/comp-hist/lgp-30-man.html>]